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# ISOMETRIC COPIES OF DIRECTED TREES IN ORIENTATIONS OF GRAPHS

TARAS BANAKH, ADAM IDZIK, OLEG PIKHURKO, IGOR PROTASOV, KRZYSZTOF PSZCZOŁA

**ABSTRACT.** The *isometric Ramsey number*  $\mathbb{R}(\vec{\mathcal{H}})$  of a family  $\vec{\mathcal{H}}$  of digraphs is the smallest number of vertices in a graph  $G$  such that any orientation of the edges of  $G$  contains every member of  $\vec{\mathcal{H}}$  in the distance-preserving way. We observe that for any finite family  $\vec{\mathcal{H}}$  of finite acyclic graphs the isometric Ramsey number  $\mathbb{R}(\vec{\mathcal{H}})$  is finite, and present upper bounds for  $\mathbb{R}(\vec{\mathcal{H}})$  in some special cases. For example, we show that the isometric Ramsey number of the family of all oriented trees with  $n$  vertices is at most  $n^{2n+o(n)}$ .

## 1. INTRODUCTION

In this paper we consider the “isometric” version of the result of Cochand and Duchet [6] who proved (generalizing a result of Rödl [11]) that for every acyclic digraph  $\vec{H}$  there exists a finite graph  $G$  such that every orientation of  $G$  contains an isomorphic copy of  $\vec{H}$ .

First we recall the necessary definitions from Graph Theory. A *graph* is a pair  $G = (V_G, E_G)$  consisting of a set  $V_G$  of *vertices* and a set  $E_G$  of two-element subsets of  $V_G$ , called the *edges* of  $G$ . By a *digraph* we will mean a pair  $\vec{G} = (V_{\vec{G}}, E_{\vec{G}})$  consisting of a set  $V_{\vec{G}}$  of vertices and a set  $E_{\vec{G}} \subset V_{\vec{G}} \times V_{\vec{G}}$  of *directed edges*, where neither loops  $(x, x)$ , nor pairs of opposite arcs  $(x, y)$  and  $(y, x)$  are allowed. By an *orientation* of a graph  $G = (V_G, E_G)$  we understand a function  $\vec{\cdot} : E_G \rightarrow V_G^2$  assigning to each edge  $e \in E_G$  an ordered pair  $\vec{e} = (a, b) \in V_G^2$  such that  $e = \{a, b\}$ . In this case the pair  $\vec{G} = (V_G, \{\vec{e}\}_{e \in E_G})$  is a digraph called an *orientation* of  $G$ .

A sequence  $(v_0, \dots, v_n)$  of distinct vertices of a graph  $G$  is called a *path* in  $G$  if for every positive  $i \leq n$  the unordered pair  $\{v_{i-1}, v_i\}$  is an edge of  $G$ . The *length* of the path  $(v_0, \dots, v_n)$  is  $n$ , that is, the number of edges. The *distance*  $d_G(x, y)$  between two vertices  $v, u$  of a graph  $G$  is the smallest length of a path in  $G$  connecting the vertices  $v$  and  $u$ . If  $u$  and  $v$  cannot be connected by a path, then we write  $d_G(x, y) = \infty$  and assume that  $\infty > n$  for all  $n \in \omega$ . A graph  $G$  is called *connected* if any two vertices  $u, v$  can be connected by a path in  $G$ . The distance in a digraph is taken with respect to the underlying undirected graph.

A sequence  $(v_0, \dots, v_n)$  of distinct vertices of a digraph  $\vec{G}$  is called a *directed path* in  $\vec{G}$  if for every positive  $i \leq n$  the ordered pair  $(v_{i-1}, v_i)$  is an edge of  $G$ . A *directed cycle* is a sequence  $(v_0, \dots, v_n)$  of distinct vertices with  $(x_i, x_{i+1})$  being a directed edge for each residue  $i$  modulo  $n + 1$ . A digraph  $\vec{G}$  is *acyclic* if it contains no directed cycles. It is well-known that each graph  $G$  admits an acyclic orientation  $\vec{G}$ : take any linear order  $\leq$  on the set  $V_G$  of vertices and for any edge  $\{u, v\} \in E_G$  put  $(u, v) \in E_{\vec{G}}$  if and only if  $u < v$ .

Following Rado’s arrow notations, for a graph  $G$  and a digraph  $\vec{H}$  we write  $G \rightarrow \vec{H}$  if for every orientation  $\vec{G}$  of  $G$  there exists an injective function  $f : V_{\vec{H}} \rightarrow V_G$  such that an ordered pair  $(u, v)$  of vertices of  $\vec{H}$  is a directed edge in  $\vec{H}$  if and only if  $(f(u), f(v))$  is a directed edge in  $\vec{G}$ . (Thus we require that  $f$  induces an isomorphism of undirected graphs and preserves all edge orientations.) If, moreover,  $d_{\vec{H}}(u, v) = d_G(f(u), f(v))$  for every pair of vertices  $u, v \in V_{\vec{H}}$ , then we write  $G \Rightarrow \vec{H}$  and say that  $f$  is an *isometric embedding* of  $\vec{H}$  in  $\vec{G}$ . Since each graph  $G$  admits an acyclic orientation, the arrow  $G \rightarrow \vec{H}$  implies that the digraph  $\vec{H}$  is acyclic.

Given a graph  $G$  and a class  $\vec{\mathcal{H}}$  of digraphs, we write  $G \rightarrow \vec{\mathcal{H}}$  (resp.  $G \Rightarrow \vec{\mathcal{H}}$ ) if for every oriented graph  $\vec{H} \in \vec{\mathcal{H}}$  we have  $G \rightarrow \vec{H}$  (resp.  $G \Rightarrow \vec{H}$ ). In this case the family  $\vec{\mathcal{H}}$  necessarily consists of acyclic digraphs. For a natural number  $n \in \mathbb{N}$  by  $\vec{\mathcal{T}}_n$  we denote the class of oriented trees on  $n$  vertices. By a *tree* we understand a connected graph without cycles. For  $n \in \mathbb{N}$ , the *directed path*  $\vec{I}_n$  is the digraph with  $V_{\vec{I}_n} = \{0, \dots, n-1\}$  and  $E_{\vec{I}_n} = \{(i-1, i) : 0 < i < n\}$ .

For a class  $\vec{\mathcal{H}}$  of digraphs let  $R(\vec{\mathcal{H}})$  (resp.  $\mathbb{R}(\vec{\mathcal{H}})$ ) be the smallest number of vertices of a graph  $G$  such that  $G \rightarrow \vec{\mathcal{H}}$  (resp.  $G \Rightarrow \vec{\mathcal{H}}$ ). If no graph  $G$  with  $G \rightarrow \vec{\mathcal{H}}$  (resp.  $G \Rightarrow \vec{\mathcal{H}}$ ) exists, then we put  $R(\vec{\mathcal{H}}) = \infty$  (resp.

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$\mathbb{R}(\vec{\mathcal{H}}) = \infty$ ). The number  $\mathbb{R}(\vec{\mathcal{H}})$  (resp.  $\mathbb{IR}(\vec{\mathcal{H}})$ ) is called the (*isometric*) *Ramsey number* of the family  $\vec{\mathcal{H}}$ . If the family  $\vec{\mathcal{H}}$  consists of a unique digraph  $\vec{H}$ , then we write  $\mathbb{R}(\vec{H})$  and  $\mathbb{IR}(\vec{H})$  instead of  $\mathbb{R}(\{\vec{H}\})$  and  $\mathbb{IR}(\{\vec{H}\})$ , respectively.

By Theorem B of Cochand and Duchet [6], for every finite acyclic digraph  $\vec{H}$ , the Ramsey number  $\mathbb{R}(\vec{H})$  is finite. This implies that for every finite family  $\vec{\mathcal{H}}$  of finite acyclic digraphs the Ramsey number  $\mathbb{R}(\vec{\mathcal{H}}) \leq \sum_{\vec{H} \in \vec{\mathcal{H}}} \mathbb{R}(\vec{H})$  is finite, too. In Section 2 we shall apply a deep Ramsey result of Dellamonica and Rödl [7] to prove that the isometric Ramsey number  $\mathbb{IR}(\vec{\mathcal{H}})$  is finite, too.

For the family  $\vec{\mathcal{T}}_n$  of oriented trees on  $n$  vertices Kohayakawa, Łuczak and Rödl [9] proved that  $\mathbb{R}(\vec{\mathcal{T}}_n) = O(n^4 \log n)$ . In this paper for every  $n \in \mathbb{N}$  we construct a graph  $G_n$  with  $< 2^{2^{n-1}}$  vertices such that  $G_n \Rightarrow \vec{\mathcal{T}}_n$ , showing that  $\mathbb{IR}(\vec{\mathcal{T}}_n) < 2^{2^{n-1}}$ . Using Bollobás' [3] bounds on the order of graphs of large girth and large chromatic number, we shall improve the upper bounds  $\mathbb{IR}(\vec{\mathcal{T}}_n) \leq \mathbb{IR}(\vec{\mathcal{T}}_n) < 2^{2^{n-1}}$  to  $\mathbb{IR}(\vec{\mathcal{T}}_n) = o(n^{2n})$  and  $\mathbb{IR}(\vec{\mathcal{T}}_n) = o(n^{4n})$ . In Theorem 4.5 using random graphs we improve the latter upper bound to  $\mathbb{IR}(\vec{\mathcal{T}}_n) \leq (4e + o(1))^n (n^2 \ln n)^n = n^{2n+o(n)}$ . The technique developed for the proof of Theorem 4.5 allows us to improve the upper bound  $\mathbb{R}(\vec{\mathcal{T}}_n) \leq (2500e^8 + o(1)) n^4 \ln n$  obtained by Kohayakawa, Łuczak and Rödl [9] to the upper bound  $(K + o(1)) n^4 \ln n$ , where  $K = \min_{x>1} \frac{16x^2}{1-x+x \ln x} \approx 98.8249\dots$ . In Section 5 we search for long directed paths in arbitrary orientations of graphs. In the final Section 6 we prove that every infinite graph  $G$  admits an orientation containing no directed path of infinite diameter in  $G$ . Some other results and problems related to coloring and orientations in graphs can be found in [10].

## 2. THE ISOMETRIC RAMSEY NUMBER FOR A FINITE ACYCLIC DIGRAPH

In this section we prove that each finite acyclic digraph  $\vec{H}$  has finite isometric Ramsey number  $\mathbb{IR}(\vec{H})$ . The idea of the proof of this result was suggested to the authors by Yoshiharu Kohayakawa.

**Theorem 2.1.** *For any finite acyclic digraph  $\vec{H} = (V, \vec{E})$ , the isometric Ramsey number  $\mathbb{IR}(\vec{H})$  is finite.*

*Proof.* Clearly, it is enough to prove the theorem when the graph  $\vec{H}$  is connected. Fix any vertex  $h$  of  $H$  and consider the digraph  $\vec{\Gamma}$  with

$$V_{\vec{\Gamma}} := V_{\vec{H}} \times \{0, 1\} \text{ and } \vec{E}_{\vec{\Gamma}} := \{((h, 0), (h, 1))\} \cup \{((u, 0), (v, 0)), ((v, 1), (u, 1)) : (u, v) \in E_{\vec{H}}\}.$$

Observe that the digraph  $\vec{\Gamma}$  is acyclic, connected and contains isometric copies of  $\vec{H}$  and the graph  $\vec{H}$  with the opposite orientation. Being acyclic, the graph  $\vec{\Gamma}$  admits a linear ordering  $<$  of vertices such that  $u < v$  for any directed edge  $(u, v) \in \vec{E}_{\vec{\Gamma}}$ .

By Theorem 1.8 of [7], there exists a finite graph  $G$  with a linear ordering of vertices such that for any 2-coloring of its edges there exists a monotone isometric embedding  $f : V_{\vec{\Gamma}} \rightarrow V_G$  such that the set  $\{\{f(u), f(v)\} : (u, v) \in E_{\vec{\Gamma}}\}$  is monochrome. In this case we shall say that the embedding  $f$  is *monochrome*. The monotonicity of  $f$  means that  $f$  preserves the order of vertices.

We claim that  $G \Rightarrow \vec{H}$ . Given any orientation  $\vec{G}$  of the graph  $G$ , color an edge  $\{u, v\} \in E_G$  with  $u < v$  in green if  $(u, v) \in E_{\vec{G}}$  and in red if  $(v, u) \in E_{\vec{G}}$ . By the Ramsey property of  $G$ , there exists a monochrome monotone isometric embedding  $f : V_{\vec{\Gamma}} \rightarrow V_G$ . If the color of the monochromatic set  $C = \{\{f(u), f(v)\} : (u, v) \in E_{\vec{\Gamma}}\}$  is green, then the map  $g_0 : V_{\vec{H}} \rightarrow V_G$ ,  $g_0 : v \mapsto f(v, 0)$ , is a required isometric isomorphic embedding of  $\vec{H}$  into  $\vec{G}$ . If the color of  $C$  is red, then the map  $g_1 : V_{\vec{H}} \rightarrow V_G$ ,  $g_1 : v \mapsto f(v, 1)$ , is an isometric isomorphic embedding of  $\vec{H}$  into  $\vec{G}$ . In both cases we get  $G \Rightarrow \vec{H}$ .  $\square$

**Corollary 2.2.** *Any finite family  $\vec{\mathcal{H}}$  of finite acyclic digraphs has finite isometric Ramsey number  $\mathbb{IR}(\vec{\mathcal{H}})$ .*  $\square$

**Corollary 2.3.** *For every  $n \in \mathbb{N}$  the family  $\vec{\mathcal{T}}_n$  of directed trees on  $n$  vertices has finite isometric Ramsey number  $\mathbb{IR}(\vec{\mathcal{T}}_n)$ .*  $\square$

**Remark 2.4.** The proof of [7, Theorem 1.8] proceeds by a more general induction involving amalgamation and hypergraphs, and seems to give very bad bounds on the isometric Ramsey number  $\mathbb{IR}(\vec{\mathcal{A}}_n)$  for the family  $\vec{\mathcal{A}}_n$  of all acyclic digraphs on  $n$  vertices. It would be interesting to get some reasonable upper bound on this function.

3. SIMPLE BOUNDS FOR THE ISOMETRIC RAMSEY NUMBERS  $\text{IR}(\vec{\mathcal{T}}_n)$ 

In this section we prove some simple upper bounds on the isometric Ramsey numbers  $\text{IR}(\vec{\mathcal{T}}_n)$  and  $\text{IR}(\vec{I}_n)$ . First we present a simple example of a graph witnessing that  $\text{IR}(\vec{\mathcal{T}}_n) < 2^{2^{n-1}}$ . The construction of this graph exploits rectangular products of graphs. By definition, the *rectangular product*  $G \times H$  of two graphs  $G, H$  is a graph such that  $V_{G \times H} = V_G \times V_H$  and an unordered pair  $\{(g, h), (g', h')\} \subset G \times H$  is an edge of  $G \times H$  if and only if either  $\{g, g'\} \in E_G$  and  $h = h'$  or  $g = g'$  and  $\{h, h'\} \in E_H$ . It can be shown that for any vertices  $(g, h), (g', h')$  of  $G \times H$  we get

$$d_{G \times H}((g, h), (g', h')) = d_G(g, g') + d_H(h, h').$$

For an (oriented) graph  $G$  by  $|G|$  we denote the cardinality of the set  $V_G$  of vertices of  $G$ . For a cardinal number  $m$  by  $K_m$  we denote the complete graph on  $m$  vertices.

**Lemma 3.1.** *Let  $\vec{\mathcal{T}}, \vec{\mathcal{T}}'$  be two families of finite oriented trees such that for every oriented tree  $\vec{T}' \in \vec{\mathcal{T}}'$  there is an oriented subtree  $\vec{T} \in \vec{\mathcal{T}}$  of  $\vec{T}'$  such that  $|\vec{T}| = |\vec{T}'| - 1$ . For any graph  $G$  with  $G \Rightarrow \vec{\mathcal{T}}$  we get  $G \times K_{|G|+1} \Rightarrow \vec{\mathcal{T}}'$ .*

*Proof.* Let  $G' = G \times K_{|G|+1}$ . To prove that  $G' \Rightarrow \vec{\mathcal{T}}'$ , take any oriented tree  $\vec{T}' \in \vec{\mathcal{T}}'$  and any orientation  $\vec{G}'$  of the graph  $G'$ . By our assumption, for the tree  $\vec{T}'$  there exists an oriented subtree  $\vec{T} \in \vec{\mathcal{T}}$  of  $\vec{T}'$  such that  $|\vec{T}| = |\vec{T}'| - 1$ . Let  $t'$  be the unique element of the set  $V_{\vec{T}'} \setminus V_{\vec{T}}$  and  $t \in V_{\vec{T}}$  be the unique vertex of  $\vec{T}$  such that  $(t', t)$  or  $(t, t')$  is an edge of  $\vec{T}'$ .

For every vertex  $u$  of the complete graph  $K_{|G|+1}$ , consider the subgraph  $G'_u = G' \times \{u\}$  of  $G'$  and its orientation  $\vec{G}'_u$ , inherited from the orientation  $\vec{G}'$  of  $G'$ . Since  $G \Rightarrow \vec{\mathcal{T}}$ , there is an isometric embedding  $f_u : \vec{T} \rightarrow \vec{G}'_u$ . By the Pigeonhole Principle, there are two distinct vertices  $u, w$  in  $K_{|G|+1}$  such that  $f_u(t) = (g, u)$  and  $f_w(t) = (g, w)$  for some vertex  $g$  of the graph  $G$ . Now look at the orientation of the edges  $\{t, t'\}$  and  $\{(g, u), (g, w)\}$  in the digraphs  $\vec{T}'$  and  $\vec{G}'$ .

If either  $(t, t') \in E_{\vec{T}'}$  and  $((g, u), (g, w)) \in E_{\vec{G}'}$ , or  $(t', t) \in E_{\vec{T}'}$  and  $((g, w), (g, u)) \in E_{\vec{G}'}$ , then we define a map  $f : \vec{T}' \rightarrow G'$  by  $f(t') = (g, w)$  and  $f|_{\vec{T}} = f_u$  and observe that  $f$  is an isometric embedding of  $\vec{T}'$  into  $\vec{G}'$ .

If either  $(t, t') \in E_{\vec{T}'}$  and  $((g, w), (g, u)) \in E_{\vec{G}'}$ , or  $(t', t) \in E_{\vec{T}'}$  and  $((g, u), (g, w)) \in E_{\vec{G}'}$ , then we define a map  $f : \vec{T}' \rightarrow G'$  by  $f(t') = (g, u)$  and  $f|_{\vec{T}} = f_w$  and observe that  $f$  is an isometric embedding of  $\vec{T}'$  into  $\vec{G}'$ .  $\square$

**Corollary 3.2.** *If for some  $n \in \mathbb{N}$  a graph  $G$  satisfies the isometric Ramsey relation  $G \Rightarrow \vec{\mathcal{T}}_n$ , then  $G \times K_{|G|+1} \Rightarrow \vec{\mathcal{T}}_{n+1}$ .*  $\square$

**Theorem 3.3.** *For every  $n \in \mathbb{N}$   $\text{IR}(\vec{\mathcal{T}}_{n+1}) \leq \text{IR}(\vec{\mathcal{T}}_n)(\text{IR}(\vec{\mathcal{T}}_n) + 1)$  and  $\text{IR}(\vec{\mathcal{T}}_n) < 2^{2^{n-1}}$ .*

*Proof.* The inequality  $\text{IR}(\vec{\mathcal{T}}_{n+1}) \leq \text{IR}(\vec{\mathcal{T}}_n)(\text{IR}(\vec{\mathcal{T}}_n) + 1)$  follows from Corollary 3.2. Indeed, for every  $n \in \omega$  we can choose a graph  $G$  with  $|G| = \text{IR}(\vec{\mathcal{T}}_n)$  vertices and  $G \Rightarrow \vec{\mathcal{T}}_n$ . By Corollary 3.2, the graph  $G' = G \times K_{|G|+1}$  satisfies the relation  $G' \Rightarrow \vec{\mathcal{T}}_{n+1}$  and hence

$$\text{IR}(\vec{\mathcal{T}}_{n+1}) \leq |G'| = |G|(|G| + 1) = \text{IR}(\vec{\mathcal{T}}_n)(\text{IR}(\vec{\mathcal{T}}_n) + 1).$$

It remains to prove that  $\text{IR}(\vec{\mathcal{T}}_n) + 1 \leq 2^{2^{n-1}}$  for  $n \in \mathbb{N}$ . For  $n = 1$  we get the equality  $\text{IR}(\vec{\mathcal{T}}_1) + 1 = 1 + 1 = 2^{2^0}$ . Assume that for some  $n \in \mathbb{N}$  we have proved that  $\text{IR}(\vec{\mathcal{T}}_n) + 1 \leq 2^{2^{n-1}}$ . Then

$$\text{IR}(\vec{\mathcal{T}}_{n+1}) + 1 \leq \text{IR}(\vec{\mathcal{T}}_n)(\text{IR}(\vec{\mathcal{T}}_n) + 1) + 1 \leq (2^{2^{n-1}} - 1)2^{2^{n-1}} + 1 = 2^{2^n} - 2^{2^{n-1}} + 1 \leq 2^{2^n}.$$

$\square$

The upper bound  $\text{IR}(\vec{\mathcal{T}}_n) < 2^{2^{n-1}}$  can be greatly improved using known upper bounds on the Erdős function  $\text{Erdős}(k, g)$ , which assigns to any positive integer numbers  $k, g$  the smallest cardinality  $|G|$  of a graph  $G$  with chromatic number  $\chi(G) \geq k$  and girth  $g(G) \geq g$ . We recall that the *girth*  $g(G)$  of a graph is the smallest cardinality of a cycle in  $G$ . If  $G$  contains no cycles, then we put  $g(G) = \infty$ . The *chromatic number*  $\chi(G)$  of a graph  $G$  is the smallest number  $k \in \mathbb{N}$  for which there exists a map  $\chi : V_G \rightarrow \{1, \dots, k\}$  such that  $\chi(x) \neq \chi(y)$  for any edge  $\{x, y\} \in E_G$ . The following bounds for the Erdős function  $\text{Erdős}(k, g)$  were proved by Erdős [8], Bollobás [3] and Spencer [12], respectively.

**Proposition 3.4.** (1) *For any  $k, g$  we get  $\text{Erdős}(k, g) \geq k^{(g-1)/2}$ ;*  
 (2) *For any  $k, g \geq 4$  we get  $\text{Erdős}(k, g) \leq \lceil h^g \rceil$  where  $h = 6(k+1) \ln(k+1)$ .*

- (3) *There exists a constant  $C$  such that for any numbers  $k, g \geq 3$  and  $m = \text{Erdős}(k, g)$  we have the inequality  $\sqrt[9]{m} \cdot \ln m < Ck$ , which implies that  $\text{Erdős}(k, g) = o(k^{g-2})$  as  $\max\{k, g\} \rightarrow \infty$ .*  $\square$

Write  $G \rightarrow \vec{H}$  if for every orientation  $\vec{G}$  of  $G$  and every  $\vec{H} \in \vec{\mathcal{H}}$  there is an injective map  $f : V_{\vec{H}} \rightarrow V_G$  such that for every directed edge  $(x, y)$  of  $H$  the pair  $(f(x), f(y))$  is a directed edge of  $\vec{G}$ . (Note that we do not require that  $f$  induces isomorphism, that is,  $G$  can have extra edges inside the set  $f(V_{\vec{H}})$ .) Another function related to  $\text{IR}(\vec{\mathcal{H}})$  is Burr's function  $\text{Burr}(\vec{\mathcal{H}})$  assigning to every family  $\vec{\mathcal{H}}$  of oriented trees the smallest number  $k$  such that  $G \rightarrow \vec{H}$  for every graph  $G$  with chromatic number  $\chi(G) \geq k$ . If such number  $k$  does not exist, then we put  $\text{Burr}(\vec{\mathcal{H}}) = \infty$ . By the Gallai-Hasse-Roy-Vitaver Theorem [13, Theorem 3.13], the chromatic number  $\chi(G)$  of a finite graph  $G$  is equal to  $\max\{n \in \mathbb{N} : G \rightarrow \vec{I}_n\}$ . This equality implies that  $\text{Burr}(\vec{I}_n) = n$  for every  $n \in \mathbb{N}$ . In [5] Burr considered the numbers  $\text{Burr}(\vec{\mathcal{T}}_n)$  and proved that  $\text{Burr}(\vec{\mathcal{T}}_n) \leq (n-1)^2$ . This upper bound was improved to the upper bound  $\text{Burr}(\vec{\mathcal{T}}_n) \leq \frac{1}{2}n^2 - \frac{1}{2}n + 1$  in [2]. According to (still unproved) Conjecture of Burr [5], the equality  $\text{Burr}(\vec{\mathcal{T}}_n) = 2n - 2$  holds for all  $n \geq 2$ .

**Proposition 3.5.** *For any  $n \in \mathbb{N}$  and a subclass  $\vec{\mathcal{H}} \subset \vec{\mathcal{T}}_n$  we get the upper bound*

$$\text{IR}(\vec{\mathcal{H}}) \leq \text{Erdős}(\text{Burr}(\vec{\mathcal{H}}), 2n - 2).$$

*Proof.* Fix a graph  $G$  of cardinality  $|G| = \text{Erdős}(\text{Burr}(\vec{\mathcal{H}}), 2n - 2)$  with chromatic number  $\chi(G) \geq \text{Burr}(\vec{\mathcal{H}})$  and girth  $g(G) \geq 2n - 2$ . Let us prove that  $G \Rightarrow \vec{\mathcal{H}}$ . Take any orientation  $\vec{G}$  of  $G$  and  $\vec{H} \in \vec{\mathcal{H}}$ . Since  $G \rightarrow \vec{H}$ , there is an orientation-preserving injection  $f : \vec{H} \rightarrow \vec{G}$ . Since  $\vec{H}$  is a connected graph with at most  $n$  vertices and  $g(G) \geq 2n - 2$ , the map  $f$  is an isometric embedding. So,  $G \Rightarrow \vec{\mathcal{H}}$ .  $\square$

Combining Proposition 3.5 with known upper bounds  $\text{Burr}(\vec{I}_n) = n$  and  $\text{Burr}(\vec{\mathcal{T}}_n) \leq \frac{1}{2}n^2 - \frac{1}{2}n + 1$  we get the following upper bounds for the isometric Ramsey numbers  $\text{IR}(\vec{I}_n)$  and  $\text{IR}(\vec{\mathcal{T}}_n)$ .

**Corollary 3.6.** *For every  $n \in \mathbb{N}$  we get the upper bounds*

$$\begin{aligned} \text{IR}(\vec{I}_n) &\leq \text{Erdős}(n, 2n - 2) = o(n^{2n-4}) = o(n^{2n}) \text{ and} \\ \text{IR}(\vec{\mathcal{T}}_n) &\leq \text{Erdős}\left(\frac{1}{2}n^2 - \frac{1}{2}n + 1, 2n - 2\right) = o\left(\left(\frac{1}{2}n^2 - \frac{1}{2}n + 1\right)^{2n-4}\right) = o(n^{4n}). \end{aligned}$$

$\square$

In Theorem 4.5 we shall improve the upper bound  $o(n^{4n})$  for  $\text{IR}(\vec{\mathcal{T}}_n)$  to the upper bound  $n^{2n+o(n)}$ .

**Remark 3.7.** By Theorem 3 in [9],  $\text{R}(\vec{I}_n) \geq n^2/2$  for all  $n \in \mathbb{N}$ . This yields the lower bound

$$\frac{1}{2}n^2 \leq \text{R}(\vec{I}_n) \leq \text{IR}(\vec{I}_n) \leq \text{IR}(\vec{\mathcal{T}}_n)$$

for the isometric Ramsey numbers  $\text{IR}(\vec{I}_n)$  and  $\text{IR}(\vec{\mathcal{T}}_n)$ .

**Remark 3.8.** It can be shown that

$$\begin{aligned} \text{IR}(\vec{I}_1) &= \text{IR}(\vec{\mathcal{T}}_1) = 1 = |K_1|, \\ \text{IR}(\vec{I}_2) &= \text{IR}(\vec{\mathcal{T}}_2) = 2 = |K_2|, \\ \text{IR}(\vec{I}_3) &= 5 = |C_5|, \quad \text{IR}(\vec{\mathcal{T}}_3) = 6 = |K_2 \times K_3|, \\ \text{IR}(\vec{I}_4) &\leq 30 = |C_5 \times K_6|, \quad \text{IR}(\vec{\mathcal{T}}_4) \leq 42 = |K_2 \times K_3 \times K_7|. \end{aligned}$$

**Question 3.9.** *What is the exact value of the isometric Ramsey numbers  $\text{IR}(\vec{I}_4)$  and  $\text{IR}(\vec{\mathcal{T}}_4)$ ? Are they distinct?*

#### 4. ISOMETRIC COPIES OF DIRECTED TREES IN ORIENTATIONS OF RANDOM GRAPHS

In this section we shall apply the technique of random graphs and shall improve the upper bound  $\text{IR}(\vec{\mathcal{T}}_n) = o(n^{4n})$  established in Corollary 3.6 to the upper bound  $\text{IR}(\vec{\mathcal{T}}_n) \leq (4e + o(1))^n (n^2 \ln n)^n = n^{2n+o(n)}$ .

First we prove some technical lemmas. The first of them uses the idea of the proof of Theorem 1 in [9].

**Lemma 4.1.** *A graph  $G = (V_G, E_G)$  satisfies  $G \Rightarrow \vec{\mathcal{T}}_n$  for some  $n \in \mathbb{N}$  if there exist sequences  $(w_k)_{k=1}^{n-1}$  and  $(d_k)_{k=1}^{n-1}$  of positive real numbers such that for every  $2 \leq k < n$  the following conditions hold:*

- (1) *For every set  $S = \{s_1, \dots, s_{k-1}\} \subset V_G$  of cardinality  $k-1$  and every  $v \in V_G \setminus S$ , we have that  $|Y| \leq d_k$ , where  $Y$  consists of  $y \in V_G \setminus (S \cup \{v\})$  such that  $\{y, v\} \in E_G$  and  $\text{dist}_{G-v}(y, s_i) \leq i$  for some  $1 \leq i < k$ .*

- (2) Every set  $W \subset V_G$  of cardinality  $|W| > w_k$  spans more than  $(d_k + k - 1)w_k$  edges in  $G$ .
- (3)  $\sum_{k=1}^{n-1} w_k < |V_G|$ .

*Proof.* For a subset  $U \subset V_G$  by  $G[U]$  we denote the induced subgraph  $G[U] = (U, E[U])$  of  $G$ , where  $E[U] = \{\{u, v\} \in E_G : \{u, v\} \subset U\}$ . Also, let us write  $(G, U) \Rightarrow \vec{T}_k$ , meaning that, for every  $\vec{T} \in \vec{T}_k$ , every orientation  $\vec{G}$  of  $G$  contains a copy of  $\vec{T}$  which lies inside  $U$  and is an isometric subgraph of  $G$ .

We shall inductively prove that for every  $1 \leq k \leq n$  and every set  $U \subset V_G$  of size  $|U| > \sum_{i=1}^{k-1} w_i$ , we have  $(G, U) \Rightarrow \vec{T}_k$ . The base case  $k = 1$  is trivial. Suppose that this holds for some  $k$ . Take any  $U \subset V_G$  with  $|U| > \sum_{i=1}^k w_i$ . Take any orientation  $\vec{E}(G[U])$  of  $E(G[U])$  and any directed tree  $\vec{T} \in \vec{T}_{k+1}$ . Let  $u$  be a pendant vertex of  $\vec{T}$ . By symmetry, assume that  $(v, u)$  is an arc in  $\vec{T}$ , that is, the arc in  $\vec{T}$  goes from the unique neighbor  $v$  of  $u$  to  $u$ .

Let  $W$  be the set of vertices in  $U$  whose out-degree in  $G[U]$  is at most  $d_k + k - 1$ . We claim that  $|W| \leq w_k$ . Suppose not. Then  $|W| > w_k$  and Item 2 guarantees that  $W$  spans more than  $(d_k + k - 1)w_k$  edges in  $G$ , each edge contributing to out-degree of some vertex in  $W$ . Thus  $(d_k + k - 1)|W| \geq |E[W]| > (d_k + k - 1)w_k$ , which is a desired contradiction showing that  $|W| \leq w_k$ .

Thus  $U' = U \setminus W$  has size  $|U'| = |U| - |W| > (\sum_{i=1}^k w_i) - w_k = \sum_{i=1}^{k-1} w_i$ . By inductive assumption,  $U'$  has a  $G$ -isometric copy  $\vec{T}'$  of the oriented tree  $\vec{T} - u$ . Let  $\{s_1, \dots, s_{k-1}\}$  be an enumeration of the set  $S := V_{\vec{T}'} \setminus \{v\} \subset U'$  such that  $\text{dist}(s_i, v) \leq i$  for every  $i < k$ . Let  $Y$  be defined as in Item 1 with respect to  $v$  and  $\{s_1, \dots, s_{k-1}\}$ . By Item 1,  $|Y| \leq d_k$ . On the other hand, the neighbor  $v \in V_{\vec{T}'} \subset U \setminus W$  of  $u$  must have out-degree in  $U \setminus S$  greater than  $d_k + k - 1 - |S| = d_k$ . Thus there is an out-neighbor of  $v$  which is in  $U \setminus (W \cup Y)$ . Let  $u$  be mapped to this vertex. Then  $(v, u) \in \vec{E}(G[U])$  is oriented from  $v$  to  $u$ , as desired. Since  $d_{G-v}(u, s_i) > i$  for each  $i < k$ , the addition of  $u$  cannot violate the  $G$ -isometry property (since all vertices of  $\vec{T} - u$  are embedded into  $S \cup \{v\}$ ). This gives the required embedding of  $\vec{T}$  and finishes the proof.  $\square$

Our next elementary lemma yields an upper bound on the sum of a geometric progression.

**Lemma 4.2.** *For positive real numbers  $a, c$  with  $a > 1 + \frac{1}{c}$  we get  $\frac{a^n - 1}{a - 1} < (1 + c)a^{n-1}$  for every  $n \in \mathbb{N}$ .*

*Proof.* The inequality is equivalent to  $a^n - 1 < (1 + c)a^{n-1}(a - 1) = a^n - a^{n-1} + ca^{n-1}(a - 1)$  and to  $a^{n-1} - 1 < ca^{n-1}(a - 1)$ . The latter inequality follows from  $a^{n-1} < ca^{n-1}(a - 1)$ , which is equivalent to  $1 < c(a - 1)$ .  $\square$

In the proof of Lemma 4.4 we shall use the following Chernoff-type bounds; for a proof see e.g. [1, §A.1].

**Lemma 4.3** (Chernoff bounds). *Let  $X_1, \dots, X_n$  be independent random variables taking values in  $\{0, 1\}$  and let  $\mathbb{E}X$  be the expected value of their sum  $X = \sum_{i=1}^n X_i$ . Then*

$$\mathbb{P}\{X \geq C \cdot \mathbb{E}X\} \leq \left(\frac{e^{C-1}}{C}\right)^{\mathbb{E}X}, \quad \mathbb{P}\{X \geq (1+c)\mathbb{E}X\} \leq e^{-\frac{c^2}{3}\mathbb{E}X} \quad \text{and} \quad \mathbb{P}\{X \leq (1-c)\mathbb{E}X\} \leq e^{-\frac{c^2}{2}\mathbb{E}X}$$

for every  $C > 1$  and  $0 < c < 1$ .  $\square$

**Lemma 4.4.** *For positive integers  $n, N$  the inequality  $\mathbb{I}R(\vec{T}_n) \leq N$  holds if there exist real numbers  $c, p \in (0, 1)$ ,  $C \in (1, \infty)$  satisfying the following inequalities:*

- (1)  $c^2 p N > 3 \ln(3N)$ ;
- (2)  $(1 - C + C \ln C)p(1 + c)^n (pN)^{n-2} > (n - 1) \ln N + \ln(1 + c) + \ln(3)$ ;
- (3)  $c^2 C^2 (1 + c)^{2n} (pN)^{2n-4} > N \ln 2 + \ln(3n)$ ;
- (4)  $\frac{(n-1)(n-2)}{(1-c)p} + \frac{2C}{(1-c)}(n-1)(1+c)^n (pN)^{n-2} < N$ .

*Proof.* Assume that the numbers  $n, N, p, c, C$  satisfy the assumptions of the lemma. Let  $G = G(N, p)$  be a random graph on  $N$  vertices in which an edge  $\{u, v\} \subset V_G$  appears with probability  $p$ . We shall prove that with non-zero probability the random graph  $G$  has  $G \Rightarrow \vec{T}_n$ .

Let

$$\hbar := (1 + c)^n (pN)^{n-2}.$$

For every positive integer  $k < n$  let

$$d_k = C p \hbar \quad \text{and} \quad w_k = \frac{2(d_k + k - 1)}{(1 - c)p}.$$

Chernoff bound implies that any fixed vertex of  $G$  has degree  $\geq (1+c)p(N-1)$  with probability  $< e^{-\frac{c^2}{3}p(N-1)}$ . Consequently, with probability  $P_1 > 1 - Ne^{-\frac{c^2}{3}p(N-1)}$  all vertices of  $G$  have degree  $< (1+c)pN$ . The condition (1) implies that  $-\frac{c^2}{3}p(N-1) < -\ln(3N)$  and hence

$$P_1 > 1 - Ne^{-\frac{c^2}{3}p(N-1)} > 1 - Ne^{-\ln(3N)} = \frac{2}{3}.$$

For every  $k < n$ , take any pairwise distinct points  $v, s_1, \dots, s_{k-1} \in V_G$ . If the maximum degree of  $G$  is at most  $(1+c)pN$ , then for every  $i < k$  the ball  $B(s_i, i) = \{x \in V_G : \text{dist}_G(x, s_i) \leq i\}$  has cardinality

$$|B(s_i, i)| \leq \sum_{j=0}^i ((1+c)pN)^j = \frac{((1+c)pN)^{i+1} - 1}{(1+c)pN - 1} < (1+c)((1+c)pN)^i.$$

The latter strict inequality can be derived from Lemma 4.2 and the inequality  $cpN \geq c^2pN > 3\ln(3N) \geq 3$ .

By above, the set  $X$  of vertices of  $G - v$  at distance at most  $i < k$  in  $G - v$  from some  $s_i$  has size at most  $(1+c) \sum_{i=1}^{k-1} ((1+c)pN)^i = (1+c) \frac{((1+c)pN)^k - 1}{(1+c)pN - 1} < (1+c)^{k+1} (pN)^{k-1} \leq \hbar$ .

Consider the set  $Y$  of neighbors of  $v$  that fall into the set  $X$ . The definition of  $X$  does not depend on the edges incident to  $v$ , so conditioned on  $X$  (of size at most  $\hbar$ ) the size of  $Y$  is dominated by  $Y' \sim \text{Bin}(\hbar, p)$ . Chernoff bound shows that the probability that  $Y'$  is at least  $Cp\hbar = CEY'$  is at most  $(\frac{e^{C-1}}{C})^{p\hbar}$ . Since the number of possible choices of  $v, s_1, \dots, s_{k-1}$  is equal to  $\frac{N!}{(N-k)!} \leq N^k$ , with probability

$$P_2 \geq 1 - \sum_{k=1}^{n-1} N^k \left(\frac{e^{C-1}}{C}\right)^{p\hbar} = 1 - \left(\frac{e^{C-1}}{C}\right)^{p\hbar} \frac{N^n - 1}{N - 1} > 1 - (1+c)N^{n-1} \left(\frac{e^{C-1}}{C}\right)^{p\hbar}$$

the condition (1) of Lemma 4.1 is satisfied or we have a vertex of degree  $\geq (1+c)pN$ . We claim that  $P_2 > \frac{2}{3}$ . It suffices to prove that

$$\ln(1+c) + (n-1)\ln N + p\hbar(C-1 - C\ln C) < -\ln(3).$$

But this follows from condition (2).

Next, we prove that with probability  $> \frac{2}{3}$  the condition (2) of Lemma 4.1 holds. Take any positive  $k < n$  and put  $\bar{w}_k = \min\{m \in \mathbb{N} : w_k < m\}$ . For any fixed set  $W \subset V_G$  of cardinality  $|W| = \bar{w}_k$ , the number of edges it spans is  $\text{Bin}(\binom{\bar{w}_k}{2}, p)$ . By Chernoff bound, the probability that it is less than  $(1-c)p\binom{\bar{w}_k}{2}$  is less than  $e^{-\frac{1}{2}c^2p\binom{\bar{w}_k}{2}}$ . The probability  $P_{3,k}$  that some set  $W \subset V_G$  of cardinality  $|W| = \bar{w}_k$  spans less than  $(1-c)p\binom{\bar{w}_k}{2}$  edges is  $P_{3,k} < \binom{N}{\bar{w}_k} e^{-\frac{1}{2}c^2p\binom{\bar{w}_k}{2}} < 2N e^{-\frac{1}{4}c^2p\bar{w}_k(\bar{w}_k+1)}$ . We claim that  $P_{3,k} < \frac{1}{3n}$  which will follow as soon as we show that  $N \ln 2 - \frac{1}{4}c^2p\bar{w}_k(\bar{w}_k+1) < -\ln(3n)$ . For this it suffices to check that  $\frac{1}{4}c^2p\bar{w}_k(\bar{w}_k+1) > N \ln 2 + \ln(3n)$ .

This follows from the chain of the inequalities

$$\frac{1}{4}c^2\bar{w}_k(\bar{w}_k+1) > \frac{1}{4}c^2w_k^2 > c^2C^2\hbar^2 = c^2C^2(1+c)^{2n}(pN)^{2n-4} > N \ln 2 + \ln(3n),$$

the last inequality postulated in (3). Therefore,  $P_{3,k} < \frac{1}{3n}$  and the probability  $P_3$  that for every  $k < n$  every set  $W \subset V[G]$  of cardinality  $|W| > w_k$  spans at least

$$(1-c)p\binom{\bar{w}_k}{2} > (1-c)pw_k(w_k+1)/2 = (d_k+k-1)(w_k+1) > (d_k+k-1)w_k$$

edges is  $> 1 - \sum_{k=1}^{n-1} P_{3,k} > 1 - \frac{n-1}{3n} > \frac{2}{3}$ . So, with probability  $> \frac{2}{3}$  the condition (2) of Lemma 4.1 holds.

Since  $(1-P_1) + (1-P_2) + (1-P_3) < 1$ , there is a non-zero probability that the random graph  $G = G(N, p)$  satisfies the conditions (1) and (2) of Lemma 4.1.

It remains to show that the condition (3) of Lemma 4.1 holds, too. For this observe that

$$\begin{aligned} \sum_{k=1}^{n-1} w_k &= \sum_{k=1}^{n-1} \frac{2(Cp\hbar + k - 1)}{(1-c)p} = \frac{2}{(1-c)p} \sum_{k=1}^{n-1} (k-1) + \frac{2C}{1-c}(n-1)\hbar = \\ &= \frac{(n-1)(n-2)}{(1-c)p} + \frac{2C}{1-c}(n-1)(1+c)^n(pN)^{n-2} < N. \end{aligned}$$

The last inequality follows from the condition (4) of the Lemma.

Now it is legal to apply Lemma 4.1 and conclude that  $G \Rightarrow \vec{T}_n$  and hence  $\mathbb{R}(\vec{T}_n) \leq |G| = N$ .  $\square$

Now we are able to prove the promised upper bound  $\mathbb{R}(\vec{T}_n) \leq (4e + o(1))^n (n^2 \ln n)^n = n^{2n+o(n)}$ .

**Theorem 4.5.** *For every  $\varepsilon \in (0, 1)$  there is  $n_\varepsilon \in \mathbb{N}$  such that  $\text{IR}(\vec{T}_n) \leq (4e(1 + \varepsilon) n^2 \ln n)^n$  for all  $n \geq n_\varepsilon$ .*

*Proof.* Choose any positive  $\delta, c \in (0, 1)$  such that

$$(1 + \delta)(1 + c) < 1 + \varepsilon \quad \text{and} \quad 4(1 + \delta) \frac{1 - c}{2 + c} > 2 + \delta.$$

For every  $n \in \mathbb{N}$  let  $N$  be the smallest integer number, which is greater than

$$\frac{(2+c)e^n}{1-c}(n-1)(1+c)^n(4(1+\delta)n^2 \ln n)^{n-2}$$

and let

$$p := \frac{4(1+\delta)n^2 \ln n}{N}.$$

So,  $N > \frac{(2+c)e^n}{1-c}(n-1)(1+c)^n(pN)^{n-2} \geq N - 1$ . It is easy to see that

$$N = o((4e(1 + \varepsilon) n^2 \ln n)^n)$$

and for  $C = e^n$  the conditions (1),(3),(4) of Lemma 4.4 hold for all sufficiently large  $n$ . To verify the condition (2), observe that

$$\begin{aligned} (1 - C + C \ln C)p(1 + c)^n(pN)^{n-2} &\geq (1 - e^n + e^n \ln e^n)p \frac{(N-1)(1-c)}{(2+c)e^n(n-1)} = \\ &\frac{1 + e^n(n-1)}{e^n(n-1)} \frac{1-c}{2+c} \frac{N-1}{N} pN = \left(1 + \frac{1}{e^n(n-1)}\right) \frac{N-1}{N} \frac{1-c}{2+c} 4(1+\delta)n^2 \ln n > \\ &> \left(1 + \frac{1}{e^n(n-1)}\right) \frac{N-1}{N} (2+\delta)n^2 \ln n = (2+\delta + o(1))n^2 \ln n. \end{aligned}$$

On the other hand,  $(n-1) \ln N + \ln(1+c) + \ln 3 = (2 + o(1))n^2 \ln n$ . So, the condition (2) holds for large  $n$ . Applying Lemma 4.4, we conclude that

$$\text{IR}(\vec{T}_n) \leq N \leq (4e(1 + \varepsilon) n^2 \ln n)^n$$

for all sufficiently large  $n$ . □

By Corollary 3.6 and Theorem 4.5,  $\text{IR}(\vec{I}_n) = o(n^{2n})$  and  $\text{IR}(\vec{T}_n) \leq n^{2n+o(n)}$ .

**Question 4.6.** *What is the growth rate of the sequence  $\text{IR}(\vec{T}_n)$ ? Is  $\text{IR}(\vec{T}_n) = n^{o(n)}$ ?*

The technique developed for the proof of Theorem 4.5 allows us to improve the upper bound

$$\text{R}(\vec{T}_n) \leq (4(5e^2)^4 + o(1))n^4 \ln n,$$

obtained by Kohayakawa, Luczak and Rödl in (the proof of) Theorem 1 of [9], and replace the constant  $4(5e^2)^4 = 2500e^8 \approx 7452395.96...$  by the a much smaller constant  $K \approx 98.82...$ .

**Theorem 4.7.** *Let  $K := \min_{x>1} \frac{16x^2}{1-x+x \ln x} \approx 98.8249...$  For any positive  $\varepsilon > 0$  there exists  $n_\varepsilon \in \mathbb{N}$  such that  $\text{R}(\vec{T}_n) < (K + \varepsilon)n^4 \ln n$  for all  $n \geq n_\varepsilon$ . Consequently,  $\text{R}(\vec{T}_n) < 99n^4 \ln n$  for all sufficiently large  $n$ .*

*Proof.* We indicate which changes should be made in the proof of Theorem 4.5 to obtain Theorem 4.7.

In the condition (1) of Lemma 4.1 the inequality  $d_{G-v}(y, s_i) \leq i$  should be replaced by  $d_{G-v}(y, s_i) \leq 1$ .

In the proof of Lemma 4.4 the constant  $h$  should be redefined as  $h := (1+c)(n-2)pN$  and the conditions (1)–(4) of Lemma 4.4 should be changed to the conditions:

- (1')  $c^2 pN > 3 \ln(3N)$ ;
- (2')  $(1 - C + C \ln C)(1 + c)(n - 2)p^2 N > (n - 1) \ln N + \ln(1 + c) + \ln(3)$ ;
- (3')  $(cC(1 + c)(n - 2)pN)^2 > N \ln 2 + \ln(3n)$ ;
- (4')  $\frac{n(n-1)}{(1-c)p} + \frac{2C(1+c)}{(1-c)}(n-1)(n-2)pN < N$ .

Now we are able to prove Theorem 4.7. Let  $C \approx 4.92155...$  be the unique real number in  $(1, \infty)$  such that

$$\frac{16C^2}{1 - C + C \ln C} = K := \min_{x>1} \frac{16x^2}{1 - x + x \ln x} \approx 98.8249...^1$$

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<sup>1</sup>The approximate values of  $C$  and  $K$  were found by the online WolframAlpha computational knowledge engine at [www.wolframalpha.com](http://www.wolframalpha.com)



Given any  $\varepsilon > 0$ , choose real numbers  $\delta, c \in (0, 1)$  such that  $K\delta < \varepsilon$  and

$$4(1 + \delta) \frac{(1 - c)^2}{(1 + c)^3} > 4 + \delta.$$

For every  $n \in \mathbb{N}$  let  $p := \frac{1-c}{2C(1+c)^2 n^2}$  and let  $N$  be the smallest integer, which is greater than  $K(1 + \delta)n^4 \ln n$ . It is easy to see that  $N = o((K + \varepsilon)n^4 \ln n)$  and the conditions (1'), (3') and (4') are satisfied for all sufficiently large  $n$ . To see that (2') holds, observe that

$$\begin{aligned} (1 - C + C \ln C)(1 + c)(n - 2)p^2 N &\geq \frac{(1 - C + C \ln C)(1 + c)(1 - c)^2}{(2C(1 + c)^2 n^2)^2} (n - 2)K(1 + \delta)n^4 \ln n = \\ &= \frac{1 - C + C \ln C}{C^2} \frac{(1 - c)^2}{4(1 + c)^3} (n - 2)K(1 + \delta) \ln n = (1 + \delta)K \frac{16}{K} \frac{(1 - c)^2}{4(1 + c)^3} (n - 2) \ln n > \\ &> (4 + \delta)(n - 2) \ln n = (4 + \delta + o(1))n \ln n. \end{aligned}$$

On the other hand,

$$(n - 1) \ln N + \ln(1 + c) + \ln 3 \leq (n - 1) \ln(1 + K(1 + \delta)n^4 \ln n) + \ln(1 + c) + \ln 3 = (4 + o(1))n \ln n,$$

so for large  $n$  the condition (2') is satisfied, too.

Applying the modified version of Lemma 4.4, we get

$$R(\vec{T}_n) \leq N \leq (K + \varepsilon)n^4 \ln n$$

for all sufficiently large numbers  $n$ . □

## 5. LONG DIRECTED PATHS IN ORIENTATIONS OF A GRAPH

By the Gallai-Hasse-Roy-Vitaver Theorem [13, Theorem 3.13], each finite graph  $G$  has chromatic number

$$\chi(G) = \max\{n \in \mathbb{N} : G \rightarrow \vec{I}_n\},$$

where the symbol  $G \rightarrow \vec{I}_n$  means that each orientation of  $G$  contains a simple directed path of length  $n$ . Having in mind this characterization, for every graph  $G$  consider the numbers

$$\bar{\chi}_I(G) = \sup\{n \in \mathbb{N} : G \Rightarrow \vec{I}_n\}, \quad \bar{\chi}_T(G) = \sup\{n \in \mathbb{N} : G \Rightarrow \vec{T}_n\},$$

and observe that  $\bar{\chi}_T(G) \leq \bar{\chi}_I(G) \leq \chi(G)$  and

$$\bar{\chi}_I(G) \leq \sup\{\text{diam}(G') + 1 : G' \text{ is a connected component of } G\}.$$

Observe that  $\text{IR}(\vec{I}_n)$  (resp.  $\text{IR}(\vec{T}_n)$ ) is equal to the smallest cardinality  $|G|$  of a graph  $G$  with  $\bar{\chi}_I(G) \geq n$  (resp.  $\bar{\chi}_T(G) \geq n$ ). So, the characteristics  $\bar{\chi}_I$  and  $\bar{\chi}_T$  determine the isometric Ramsey numbers  $\text{IR}(\vec{I}_n)$  and  $\text{IR}(\vec{T}_n)$ .

We shall show that a graph  $G$  has  $\bar{\chi}_I(G) \leq 2$  if and only if  $G$  is a comparability graph. We recall that a graph  $G$  is called a *comparability graph* if  $G$  admits a *transitive* orientation  $\vec{G}$  (that is, for any directed edges  $(x, y)$  and  $(y, z)$  of  $\vec{G}$  the pair  $(x, z)$  is a directed edge of  $\vec{G}$ ); equivalently, the set  $V_G$  of vertices of  $G$  admits a partial order such that a pair  $\{u, v\}$  of distinct vertices of  $G$  is an edge of  $G$  if and only if  $u$  and  $v$  are comparable in the partial order. By the results of Ghouila-Houri and of Gilmore and Hoffman (see [4, Theorem 6.1.1]), comparability graphs can be characterized as graphs  $G$  whose every cycle of odd length has a triangular chord (more precisely, for every  $(2n + 3)$ -cycle on  $(v_0, \dots, v_{2n+2})$  with  $n \geq 1$ , there is a residue  $i$  modulo  $2n + 3$  such that  $\{v_i, v_{i+2}\} \in E_G$ ). More information on comparability graphs can be found in Chapter 6 of the survey [4].

**Proposition 5.1.** *A graph  $G$  has  $\bar{\chi}_I(G) \leq 2$  if and only if  $G$  is a comparability graph.*

*Proof.* If  $G$  is comparability graph, then  $G$  has a transitive orientation  $\vec{G}$ . It follows that for any directed path  $(v_0, v_1, v_2)$  in  $\vec{G}$  the pair  $(v_0, v_2)$  is an edge of  $\vec{G}$  and hence  $d_G(v_0, v_2) \leq 1$ . This means that  $G \not\Rightarrow \vec{I}_3$  and hence  $\bar{\chi}_I(G) \leq 2$ .

If  $G$  is not a comparability graph, then  $G$  contains an odd cycle  $C$  without a triangular chord. It is easy to see that any orientation  $\vec{C}$  of the cycle  $C$  contains a directed path  $(v_0, v_1, v_2)$ . Since  $C$  has no triangular chords,  $d_G(v_0, v_2) = 2$ , which means that  $\{v_0, v_1, v_2\}$  is an isometric copy of  $\vec{I}_3$  in  $\vec{C}$  and in  $G$ . Therefore,  $\bar{\chi}_I(G) \geq 3$ . □

**Problem 5.2.** *Characterize graphs  $G$  with  $\bar{\chi}_I(G) \leq 3$  ( $\bar{\chi}_I(G) \leq n$  for  $n \geq 4$ ).*

**Problem 5.3.** *Characterize graphs  $G$  with  $\bar{\chi}_T(G) \leq 2$  ( $\bar{\chi}_T(G) \leq n$  for  $n \geq 3$ ).*

**Remark 5.4.** Any cycle  $C$  of odd length  $n \geq 5$  satisfies  $\bar{\chi}_I(C) = 3$  and  $\bar{\chi}_T(C) = 2$ .

Now we prove a weak 3-space property for the number  $\bar{\chi}_I(G)$ . By a *weak homomorphism*  $f : G \rightarrow H$  of graphs  $G, H$  we understand a function  $f : V_G \rightarrow V_H$  such that for every edge  $\{u, v\}$  of  $G$  we have either  $f(u) = f(v)$  or  $\{f(u), f(v)\}$  is an edge of  $H$ . For a weak homomorphism  $f : G \rightarrow H$  and vertex  $y$  of  $H$  the preimage  $f^{-1}(y)$  is a graph with the set of edges  $\{\{u, v\} \in E_G : f(u) = y = f(v)\}$ .

**Proposition 5.5.** *If  $f : G \rightarrow H$  is a weak homomorphism of finite graphs, then*

$$\bar{\chi}_I(G) \leq \max \left\{ \sum_{y \in F} \bar{\chi}_I(f^{-1}(y)) : F \subset V_H, |F| \leq \chi(H) \right\}.$$

*Proof.* By definition of the chromatic number  $\chi(H)$ , there exists a coloring  $c : V_H \rightarrow \{1, \dots, \chi(H)\}$  of the graph  $H$  such that for every edge  $\{u, v\}$  of  $G$  the colors  $c(u)$  and  $c(v)$  are distinct. For every  $y \in H$  choose an orientation  $\vec{G}_y$  of the graph  $G_y = f^{-1}(y)$  such that  $\vec{G}_y \not\cong \vec{I}_k$  for  $k = \bar{\chi}_I(G_y) + 1$ . Let  $\vec{G}$  be the orientation of the graph  $G$  such that for an edge  $\{u, v\}$  of  $G$  the ordered pair  $(u, v)$  is an edge of  $\vec{G}$  if and only if either  $c(f(u)) < c(f(v))$  or  $(u, v)$  is an edge of  $\vec{G}_y$  for some  $y \in H$ .

We claim that the digraph  $\vec{G}$  contains no isometric copy of the graph  $\vec{I}_{m+1}$ , where

$$m = \max \left\{ \sum_{y \in F} \bar{\chi}_I(G_y) : F \subset V_H, |F| \leq \chi(H) \right\}.$$

Suppose on the contrary that  $\vec{G}$  contains a directed path  $(v_0, \dots, v_m)$  such that  $d_G(v_0, v_m) = m$ . It follows that  $(c(f(v_0)), \dots, c(f(v_m)))$  is a non-decreasing sequence of numbers in the interval  $\{1, \dots, \chi(H)\}$ . Consequently, for every number  $i$  in the set  $C = \{c(f(v_0)), \dots, c(f(v_m))\}$  the set  $J_i = \{j \in \{0, \dots, m\} : c(f(v_j)) = i\}$  coincides with some subinterval  $[a_i, b_i]$  of  $\{0, \dots, m\}$  and the set  $\{f(v_j) : j \in [a_i, b_i]\}$  is a singleton  $\{y_i\}$  for some vertex  $y_i \in H$ . It follows that  $(v_{a_i}, \dots, v_{b_i})$  is a directed path isometric to  $\vec{I}_{|[a_i, b_i]|}$  in the graph  $G_{y_i}$  and hence  $|[a_i, b_i]| \leq \bar{\chi}_I(G_{y_i})$ . The choice of the orientation  $\vec{G}$  guarantees that the set  $F = \{y_i : i \in C\}$  has cardinality  $|F| = |C| \leq \chi(H)$ . Then

$$m + 1 = |[0, m]| = \sum_{i \in C} |[a_i, b_i]| \leq \sum_{i \in C} \bar{\chi}_I(G_{y_i}) = \sum_{y \in F} \bar{\chi}_I(G_y) \leq m,$$

which is a desired contradiction.  $\square$

## 6. INFINITE DIRECTED PATHS IN ORIENTATIONS OF GRAPHS

Now we discuss the problem of existence of infinite directed paths in orientations of graphs. Consider the infinite digraphs  $\vec{I}_\omega$  and  $\vec{I}_{-\omega}$  with  $V_{\vec{I}_\omega} = \omega = V_{\vec{I}_{-\omega}}$ ,  $E_{\vec{I}_\omega} = \{(i, i+1) : i \in \omega\}$ , and  $E_{\vec{I}_{-\omega}} = \{(i+1, i) : i \in \omega\}$ .

First, observe that Theorem 3.3 implies the following:

**Corollary 6.1.** *There exists a countable graph  $G$  such that  $G \Rightarrow \vec{I}_n$  for every  $n \in \mathbb{N}$ .*  $\square$

On the other hand, we shall prove that each graph  $G$  admits an orientation containing no isometric copy of the digraphs  $\vec{I}_\omega$  or  $\vec{I}_{-\omega}$  and, more generally, no directed paths of infinite diameter in  $G$ . (For a subset  $A \subset V_G$  of a graph  $G$  its *diameter* is defined as  $\text{diam}(A) = \sup\{d_G(u, v) : u, v \in A\} \in \omega \cup \{\infty\}$ .)

A sequence  $(v_n)_{n \in \omega} \in V_G^\omega$  of distinct vertices of a graph  $G$  is called an  $\omega$ -*path* in  $G$  if for every  $n \in \omega$  the pair  $\{v_n, v_{n+1}\}$  is an edge of  $G$ . An  $\omega$ -path  $(v_n)_{n \in \omega}$  in a graph  $G$  is called  $\vec{\omega}$ -*directed* (resp.  $\overleftarrow{\omega}$ -*directed*) in an orientation  $\vec{G}$  of  $G$  if for every  $n \in \omega$  the pair  $(v_n, v_{n+1})$  (resp.  $(v_{n+1}, v_n)$ ) is a directed edge of  $\vec{G}$ . An  $\omega$ -path in  $G$  is called *directed* in an orientation  $\vec{G}$  of  $G$  if it is either  $\vec{\omega}$ -directed or  $\overleftarrow{\omega}$ -directed.

The Ramsey Theorem implies that every orientation of the complete countable graph  $K_\omega$  contains  $\vec{I}_\omega$  or  $\vec{I}_{-\omega}$ . On the other hand, we have the following result.

**Theorem 6.2.** *Every graph  $G$  has an orientation  $\vec{G}$  containing no directed  $\omega$ -paths of infinite diameter in  $G$ . This implies that  $G \not\Rightarrow \vec{I}_\omega$  and  $G \not\Rightarrow \vec{I}_{-\omega}$ .*

*Proof.* Without loss of generality, the graph  $G$  is connected. Fix any vertex  $o$  in  $G$  and for every vertex  $v$  of  $G$  let  $\|v\|$  be the smallest length of a path linking the vertices  $v$  and  $o$ . Choose an orientation  $\vec{G}$  of  $G$  such that for any edge  $\{u, v\}$  in  $G$  with  $\|v\| = \|u\| + 1$  the pair  $(u, v)$  is an edge of  $\vec{G}$  if  $\|u\|$  is even and  $(v, u)$  is an edge of  $\vec{G}$  if  $\|u\|$  is odd.

We claim that the orientation  $\vec{G}$  contains no directed  $\omega$ -paths of infinite diameter. To derive a contradiction, assume that  $(v_n)_{n \in \omega}$  is a directed  $\omega$ -path of infinite diameter. Fix any even number  $n \in \omega$  such that  $\|v_0\| < n$ . Since the  $\omega$ -path  $(v_n)_{n \in \omega}$  has infinite diameter, there exists a number  $k \in \omega$  such that  $\|v_k\| \geq n$ . We can assume that  $k$  is the smallest number with this property. Taking into account that  $\|v_n\| - \|v_{n+1}\| \leq 1$  for all  $n \in \omega$ , we conclude that  $\|v_k\| = n > \|v_0\|$  and  $\|v_{k-1}\| = n - 1$ , and hence  $(v_{k-1}, v_k)$  is an edge of  $\vec{G}$ . Let also  $m$  be the smallest number such that  $\|v_m\| \geq n + 1$ . For this number we get  $\|v_m\| = n + 1$ ,  $\|v_{m-1}\| = n$  and hence  $(v_m, v_{m-1})$  is a directed edge of  $\vec{G}$ . Since both pairs  $(v_{k-1}, v_k)$  and  $(v_m, v_{m-1})$  are directed edges of the oriented graph  $\vec{G}$ , the  $\omega$ -path  $(v_n)_{n \in \omega}$  is not directed in  $\vec{G}$ . Since the graphs  $\vec{I}_\omega$  and  $\vec{I}_{-\omega}$  have infinite diameters, the digraph  $\vec{G}$  does not contain isometric copies of  $\vec{I}_\omega$  or  $\vec{I}_{-\omega}$ .  $\square$

**Remark 6.3.** Theorem 6.2 implies that every locally finite graph  $G$  admits an orientation containing no directed  $\omega$ -paths.

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