

# Continuants with equal values, a combinatorial approach

G Ramharter, L.Q. Zamboni

## Abstract

A regular continuant is the denominator  $K$  of a terminating regular continued fraction, interpreted as a function of the partial quotients. We regard  $K$  as a function defined on the set of all finite words on the alphabet  $1 < 2 < 3 < \dots$  with values in the positive integers. Given a word  $w = w_1 \cdots w_n$  with  $w_i \in \mathbb{N}$  we define its multiplicity  $\mu(w)$  as the number of times the value  $K(w)$  is assumed in the Abelian class  $\mathcal{X}(w)$  of all permutations of the word  $w$ . We prove that there is an infinity of different lacunary alphabets of the form  $\{b_1 < \dots < b_t < l+1 < l+2 < \dots < s\}$  with  $b_j, t, l, s \in \mathbb{N}$  and  $s$  sufficiently large such that  $\mu$  takes arbitrarily large values for words on these alphabets. The method of proof relies in part on a combinatorial characterisation of the word  $w_{\max}$  in the class  $\mathcal{X}(w)$  where  $K$  assumes its maximum.

MSC: primary 11J70; secondary 68R15, 68W05, 05A20.

Keywords: Values of continuants; Regular continued fractions; Combinatorial word problems.

**Introduction.** Given a sequence  $w = (w_1, \dots, w_n)$ , of positive  $w_i$ , let  $K(w)$  be the continuant of  $w$ , i.e., the denominator of the finite regular continued fraction  $\frac{1}{w_1 + \frac{1}{w_2 + \dots + \frac{1}{w_{n-1} + \frac{1}{w_n}}}}$ . We shall regard  $w$  as a *word* of length  $n$  over the alphabet  $\{1 < 2 < 3 < \dots\}$  and write  $w = w_1 \cdots w_n$ . Since  $K(w) = K(\bar{w})$ , where  $\bar{w} = w_n \cdots w_1$  denotes the reversal of  $w$ , we shall henceforth identify each word  $w$  with its reverse  $\bar{w}$ . Let  $\mathcal{X}(w)$  denote the Abelian class of  $w$  consisting of all permutations of  $w$ . The following problem has attracted much attention and led to a number of applications (see e.g. [1, 4, 5, 7, 8]): Let  $A = \{a_1 < \dots < a_s\}$  be a finite ordered alphabet with  $a_j \in \mathbb{N}$ . Given a word  $w = w_1 w_2 \cdots w_n$  with  $w_i \in A$ , find the arrangements  $w_{\max}, w_{\min} \in \mathcal{X}(w)$  maximizing resp. minimizing the function  $K(\cdot)$  on  $\mathcal{X}(w)$ . The first author [3] gave an explicit description of both extremal arrangements  $w_{\max}$  and  $w_{\min}$  and showed that in each case the arrangement is unique (up to reversal) and independent of the actual values of the positive integers  $a_i$ . He also investigated the analogous problem for the semi-regular continuant  $K'$  defined as the denominator of the semi-regular continued fraction  $\frac{J'}{K'} = \frac{1}{w_1 - \frac{1}{w_2 - \dots - \frac{1}{w_{n-1} - \frac{1}{w_n}}}}$  with entries  $w_i \in \{2, 3, \dots\}$ . He gave a fully combinatorial description of the minimizing arrangement  $w'_{\min}$  for  $K'(\cdot)$  on  $\mathcal{X}(w)$  and showed that the arrangement is unique (up to reversal) and independent of the actual values of the positive integers  $a_i$ . However, the determination of the maximizing arrangement  $w'_{\max}$  for the semi-regular continuant turned out to be more difficult. He showed that in the special case of a 2-digit alphabet  $\{(2 \leq) a_1 < a_2\}$ , the maximizing arrangement  $w'_{\max}$  is a Sturmian word and is independent of the values of the  $a_i$ . Recently the second author together with M. Edson and A. De Luca [8] developed an algorithm for constructing  $w'_{\max}$  over any ternary alphabet  $\{(2 \leq) a_1 < a_2 < a_3\}$ , and showed that the maximizing arrangement is independent of the choice of the digits. In contrast, they exhibited examples of words  $w = w_1 \cdots w_n$  over a 4-digit alphabet  $A = \{(2 \leq) a_1 < a_2 < a_3 < a_4\}$

for which the maximizing arrangement for  $K'(\cdot)$  is not unique and depends on the actual values of the positive integers  $a_1$  through  $a_4$ . In the course of these investigations the following problem came up: given an alphabet  $A$  of positive integers, we say that a word  $w$  on  $A$  has multiplicity  $\mu = \mu(w)$  if the value  $K(w)$  occurs precisely  $\mu$  times in the multi-set  $\{K(x) : x \in \mathcal{X}(w)\}$ . The multiplicity  $\mu'(w)$  is defined analogously for the semi-regular continuant  $K'(w)$ . Thus each Abelian class  $\mathcal{X}(w)$  is split into subclasses of equally valued words. Question: is it true that  $\mu$  can take arbitrarily large values for infinitely many alphabets and is there a combinatorial proof of this? Our aim here is to give a positive answer to this question in the case of regular continued fractions.

**Theorem.** Fix positive integers  $1 \leq t \leq l < s$ ,  $b_1 < \dots < b_t \leq l$  and let  $A$  be an ordered alphabet of the form  $\{b_1 < \dots < b_t < l + 1 < \dots < s\}$ . Then for all  $s$  sufficiently large, there exists an infinite sequence of words  $w_k$  over  $A$  with multiplicities  $\mu(w_k) \rightarrow \infty$  as  $k \rightarrow \infty$ .

It should be noted that for fixed  $s$  one obtains the largest possible alphabet  $A' = \{1 < 2 < \dots < s\}$  by choosing  $b_1 = t = l = 1$  ( $< s$ ). Our proof makes use of the combinatorial structure of  $w_{\max}$  found by the first author in [3].

**Preliminaries.** We introduce some notation. Let  $w = w_1 \cdots w_n$  be a word of length  $n \geq 2$  with  $w_j \in \mathbb{N}$  ( $j = 1, \dots, n$ ). The regular continuant of  $w$  has a matrix representation

$$K(w_1) = w_1 \quad \text{and} \quad K(w) = \det \begin{pmatrix} w_1 & -1 & 0 & \cdots & 0 \\ 1 & w_2 & -1 & \ddots & \vdots \\ 0 & 1 & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & w_{n-1} & -1 \\ 0 & \cdots & 0 & 1 & w_n \end{pmatrix}, \quad n \geq 2$$

It can also be defined recursively by  $K(\{\}) = 1$  ( $\{\}$  = empty word),  $K(w_1) = w_1$  and  $K(w_1 \cdots w_j) = w_j K(w_1 \cdots w_{j-1}) + K(w_1 \cdots w_{j-2})$  for  $j \geq 2$ . For each  $1 \leq k \leq m \leq n$  we set  $w_{k,m} := w_k \cdots w_m$  and  $W := W_{1,n}$ ,  $W_{k,m} := K(w_{k,m})$ . The following fundamental formula goes back to the late 19<sup>th</sup> century and can be found in Perron [2], p.11, (4):  $(W =) W_{1,n} = W_{1,j} W_{j+1,n} + W_{1,j-1} W_{j+2,n}$  ( $j \in \{1, \dots, n-1\}$ ). From this we infer the simple but useful inequality

$$W_{1,n} < 2W_{1,j} W_{j+1,n}. \quad (1)$$

Let  $A = \{a_1 < \dots < a_s\} \subset \mathbb{N}$ . We consider a word  $w = w_1 \cdots w_n := a_1^{p_1} \cdots a_s^{p_s}$  of length  $n$  with Parikh vector  $\mathbf{p} = (p_1, \dots, p_s)$  with  $p_1 + \dots + p_s = n$  where

$$a^r = \underbrace{aa \cdots a}_{r\text{-times}}$$

denotes a sequence of  $r$  equal elements  $a$ . Let  $\mathcal{X} = \mathcal{X}(A, \mathbf{p})$  denote the set of all permutations of  $w$  where we identify each word  $v$  with its reverse  $\bar{v}$ . Let  $N(A, \mathbf{p})$  denote the cardinality of  $\mathcal{X}$ . Then,  $N(A, \mathbf{p}) \geq \frac{n!}{2^{p_1! \cdots p_s!}}$ . We put  $W_{\max} = W_{\max}(A, \mathbf{p}) := \max\{K(v) : v \in \mathcal{X}\}$ . It was shown in [3] (see (3), p. 190) that  $W_{\max}$  is uniquely attained (up to reversal) by the arrangement

$$a_s L_{s-1} a_{s-2} L_{s-3} \cdots a_1^{p_1} \cdots a_{s-3} L_{s-2} a_{s-1} L_s \quad (2)$$

where  $L_i = a_i^{p_i-1}$ . Let  $P = P(A, \mathbf{p}) = \#\{K(v) : v \in \mathcal{X}\}$ .

**Proof of the Theorem.** Our first goal is to describe how to specify the last digit  $s$  ( $\geq 2$ ) in an alphabet  $A : \{b_1 < \dots < b_t < l+1 < \dots < s\}$ . We consider 'equipartitioned' words

$$w = w_1 \dots w_n := b_1^m \dots b_t^m (l+1)^m \dots s^m.$$

corresponding to the Parikh vector  $\mathbf{p} = (m, m, \dots, m)$  in which each digit of  $A$  occurs precisely  $m$ -times in  $w$ . We will give a lower bound for  $s$  (see (7) below). To this end, we introduce the quantities  $Q_{r,m-1} := K(r^{m-1})$  ( $r \in 1, 2, \dots$ ). They are the elements of the  $r$ -th generalised Fibonacci sequence which is determined by the recursion  $Q_{r,0} := 1$ ,  $Q_{r,1} := K(r) = r$ ,  $Q_{r,j+1} := rQ_{r,j} + Q_{r,j-1}$  ( $j = 1, 2, \dots$ ).

**Claim:**  $Q_{r,j-1} < (r+1)^j$  for each fixed  $r \geq 1$  and all  $j \geq 1$ .

To prove the claim, we proceed by induction on  $j$ : This is obviously true for  $j = 1$  and  $j = 2$ . Then by the induction hypothesis

$$\begin{aligned} Q_{r,j-1} &= rQ_{r,j-2} + Q_{r,j-3} < r(r+1)^{j-1} + (r+1)^{j-2} \\ &= (r+1)^{j-2}(r(r+1) + 1) < (r+1)^{j-2}(r+1)^2 \\ &= (r+1)^j. \end{aligned}$$

In order to obtain an upper bound for the number  $P(A, \mathbf{p})$ , it suffices to consider words over the largest allowed  $s$ -digit alphabet  $A' : \{1 < \dots < s\}$ ,  $b_1 = t = l = 1$  ( $< s$ ), with Parikh vector  $\mathbf{p}' = (\underbrace{m, m, \dots, m}_{s\text{-times}})$ . Clearly

$$P(A, \mathbf{p}) \leq W_{\max}(A, \mathbf{p}) \leq W_{\max}(A', \mathbf{p}')$$

and by (2)

$$w_{\max}(A', \mathbf{p}') = s \cdot (s-1)^{m-1} \cdot (s-2) \dots 1 \cdot 1^{m-1} \dots (s-2)^{m-1} \cdot (s-1) \cdot s^{m-1}. \quad (3)$$

By iteration of (1) applied to the decomposition in (3) we obtain the inequalities

$$\begin{aligned} W_{\max}(A', \mathbf{p}') &= K(w_{\max}(A', \mathbf{p}')) \\ &< 2^{2s} s \cdot (s-1) \dots 3 \cdot 2 \prod_{j=1}^s K(j^{m-1}) \\ &= 2^{2s} s! \prod_{j=1}^s Q_{j,m-1} \\ &< 2^{2s} s! \prod_{j=1}^s (j+1)^m \\ &= 2^{2s} s!((s+1)!)^m \end{aligned}$$

and hence

$$P(A, \mathbf{p}) < 2^{2s} s! ((s+1)!)^m. \quad (4)$$

For each  $s \geq 2$  we define  $m_0 = m_0(s)$  to be the smallest positive integer such that

$$2^{2s} s! \leq \left( \frac{100}{99} \right)^{m_0}.$$

Then

$$P(A, \mathbf{p}) < \left( \frac{100}{99} ((s+1)!) \right)^m \quad \text{for all } m \geq m_0(s). \quad (5)$$

On the other hand, we have the following lower bound for the number of different words in  $\mathcal{X}(w)$ :

$$N(A, \mathbf{p}) \geq \frac{((s-l+t)m)!}{2(m!)^{s-l+t}} \quad (6)$$

Based on the condition (7) below, we will later make a choice of  $s = s'(t, l)$  depending on the parameters  $t, l$ . We apply the estimates provided by Sterling's formula to the factorial terms occurring in relations (5) and (6) to obtain

$$\begin{aligned} (P(A, \mathbf{p}))^{1/m} &< \frac{100}{99} (s+1)! < \frac{100}{99} \frac{12}{11} e^{-(s+1)} (s+1)^{s+1} \sqrt{2\pi(s+1)}. \\ (((s-l+t)m)!)^{1/m} &> e^{-(s-l+t)} ((s-l+t)m)^{s-l+t} \sqrt{2\pi(s-l+t)m}^{1/m}. \\ (2(m!)^{s-l+t})^{1/m} &< e^{-(s-l+t)} m^{s-l+t} \left( 2 \frac{12}{11} (\sqrt{2\pi m})^{s-l+t} \right)^{1/m}. \end{aligned}$$

When we put the right hand sides of the last two inequalities together, the terms  $e^{s-l+t}$  and  $m^{s-l+t}$  cancel out, and if we keep the parameters  $t, l$  fixed for the moment, the terms of the form  $\sqrt{\cdot}^{1/m}$  tend to 1 as  $m \rightarrow \infty$ . Letting  $m \rightarrow \infty$  we get

$$\begin{aligned} \lim_{m \rightarrow \infty} \left( \frac{N(A, \mathbf{p})}{P(A, \mathbf{p})} \right)^{1/m} &\geq \frac{99}{100} \frac{11}{12} \frac{e^{s+1} (s-l+t)^{s-l+t}}{\sqrt{2\pi(s+1)} (s+1)^{s+1}} = \\ &= \frac{363}{400} \frac{e^{s+1} (s+1-l+t-1)^{s+1-l+t-1}}{\sqrt{2\pi(s+1)} (s+1)^{s+1}} \\ &= \frac{363}{400} \frac{e^{s+1}}{\sqrt{2\pi(s+1)} (s-l+t)^{l-t+1}} \left( 1 - \frac{l-t+1}{s+1} \right)^{s+1}. \end{aligned}$$

For fixed  $t, l$  ( $l-t \geq 1$ ) the function  $f(t, l, s) = \left( 1 - \frac{l-t+1}{s+1} \right)^{s+1}$  in the variable  $s$  is strictly increasing on the interval  $[l-t+1, \infty)$  with  $f(t, l, s) \nearrow e^{-(l-t)-1}$  as  $s \rightarrow \infty$ . We define  $s_0$  to be the lowest integer such that  $f(t, l, s_0) \geq \frac{1}{2} e^{-(l-t)-1}$ . Then

$$\lim_{m \rightarrow \infty} \left( \frac{N(A, \mathbf{p})}{P(A, \mathbf{p})} \right)^{1/m} \geq \frac{363}{400} \frac{e^{s+1}}{\sqrt{2\pi(s+1)} (s-l+t)^{l-t+1}} \frac{1}{2} e^{-(l-t)-1} =: H(t, l, s)$$

for all  $s \geq s_0$ . Obviously there exists some sufficiently large  $s' = s'(t, l) \geq s_0$  such that

$$H(t, l, s') > 1. \quad (7)$$

Therefore the right hand side of

$$\left( \frac{N(A, \mathbf{p})}{P(A, \mathbf{p})} \right) > (H(t, l, s'))^m \quad (8)$$

can be made arbitrarily large by letting  $m \rightarrow \infty$ . We call an  $(s' - l + t)$ -digit alphabet  $A = \{(1 \leq) b_1 < \dots < b_t < \dots < s'\}$  *admissible* if  $s' = s'(t, l)$  fulfills condition (7). We consider the word  $u(A, \mathbf{p}_1) = (b_1)^{m_1} \dots (b_t)^{m_1} (l+1)^{m_1} \dots (s')^{m_1}$  of length  $n = (s' - l + t)m_1$  with Parikh vector  $\mathbf{p}_1 = ((m_1)^{s'-l+t})$  where we choose  $m_1 \geq m_0$  such that  $\left( \frac{N(A, \mathbf{p}_1)}{P(A, \mathbf{p}_1)} \right) > (H(t, l, s'))^{m_1}$ . The multi-set  $\mathcal{X}_1 = \mathcal{X}(A, \mathbf{p}_1)$  is made up of the  $N(A, \mathbf{p}_1) = \#\mathcal{X}_1$  permuted arrangements of  $u$ . There exists at least one word  $w_1 \in \mathcal{X}_1$  with multiplicity  $\mu \geq 2$  because otherwise we would have  $N(A, \mathbf{p}_1) = P(A, \mathbf{p}_1)$  which contradicts (8) with  $m = m_1$ . Let  $\mu_1 (\geq 2)$  be the maximal multiplicity attained by words  $w \in \mathcal{X}_1$ . Next choose  $m_2 > m_1(s')$  such that  $H(t, l, s')^{m_2} > \mu_1$ . We claim that at least one word  $w_2$  from  $\mathcal{X}_2 = \mathcal{X}(A, \mathbf{p}_2)$ ,  $\mathbf{p}_2 = ((m_2)^{s'-l+t})$  has multiplicity  $\mu > \mu_1$ . Otherwise we would have  $N(A, \mathbf{p}_2) \leq \mu_1 P(A, \mathbf{p}_2)$  which contradicts (8) with  $m = m_2$ . Next let  $\mu_2 (\geq \mu_1)$  be the maximal multiplicity attained by words  $w \in \mathcal{X}_2$ . Proceeding with this construction step by step we end up with a sequence of words  $w_k$  on  $A$  with multiplicities  $\mu_k \rightarrow \infty$  as  $k \rightarrow \infty$ . The construction can be carried out for infinitely many different admissible alphabets. This completes the proof of the Theorem.  $\square$

The question remains largely unsolved in the case of semi-regular continuants though it seems certain that the behavior is quite similar to the regular case.

There is some evidence supporting the following

**Conjecture.** Given any ordered alphabet  $A = \{a_1 < \dots < a_s\}$  ( $a_j \in \mathbb{N}$ ,  $s \geq 2$ ), let  $\mu \geq 2$  be a positive integer. Then there exist infinitely many words on  $A$  whose multiplicity is precisely  $\mu$ . The problem appears to require a difficult investigation into the values of continuants. Most likely our theorem and the conjecture also hold for continuants of semi-regular continued fractions. Unfortunately no higher-dimensional analogue of the theorem is available at present for  $s \geq 4$  due to the fact that very little is known about the maximizing arrangements  $w'_{\max}$  for  $s \geq 4$  (see [8]).

## References

- [1] C. Baxa, Extremal values of continuants and transcendence of certain continued fractions, *Adv. Appl. Math.* 32 (2004), 754–790.
- [2] O. Perron, *Die Lehre von den Kettenbrüchen*, Bd. 1, Vieweg&Teubner, 1977.
- [3] G. Ramharter, Extremal values of continuants, *Proc. Amer. Math. Soc.* 89 (2) (1983), 189–201.
- [4] G. Ramharter, Some metrical properties of continued fractions, *Mathematica* 30 (2) (1983), 117–132.
- [5] G. Ramharter, Maximal continuants and the Fine-Wilf theorem, *J. Comb. Th. A* 111 (2005), 59–77.
- [6] G. Ramharter, Maximal continuants and periodicity, *Integers* (2006-01), (2005), 10 p.
- [7] J. Shallit, J. Sorensen, Analysis of a left shift binary GCD algorithm, *J. Symb. Comput.* 17 (6) (1994), 473–486.
- [8] A. De Luca, M. Edson, L.Q. Zamboni, Extremal values of semi-regular continuants and codings of interval exchange transformations, preprint 2021, arXiv:2105.00496.

---

Gerhard Ramharter  
Dept. of Convex and Discrete Geometry  
Technische Universität Wien  
Wiedner Hauptstrasse 8-10  
A-1040 Vienna, Austria  
e-mail: Gerhard Ramharter <gerhard.ramharter@tuwien.ac.at>

Luca Q. Zamboni  
Institut Camille Jordan  
Université Claude Bernard Lyon 1  
43 boulevard du 11 novembre 1918  
F-69622 Villeurbanne Cedex  
e-mail: zamboni <zamboni@math.univ-lyon1.fr>