Mutual Witness Gabriel Drawings of Complete Bipartite Graphs *

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Abstract. Let Γ be a straight-line drawing of a graph and let u and v be two vertices of Γ . The Gabriel disk of u, v is the disk having u and v as antipodal points. A pair $\langle \Gamma_0, \Gamma_1 \rangle$ of vertex-disjoint straight-line drawings form a mutual witness Gabriel drawing when, for i = 0, 1, any two vertices u and v of Γ_i are adjacent if and only if their Gabriel disk does not contain any vertex of Γ_{1-i} . We characterize the pairs $\langle G_0, G_1 \rangle$ of complete bipartite graphs that admit a mutual witness Gabriel drawing. The characterization leads to a linear time testing algorithm. We also show that when at least one of the graphs in the pair $\langle G_0, G_1 \rangle$ is complete k-partite with k > 2 and all partition sets in the two graphs have size greater than one, the pair does not admit a mutual witness Gabriel drawing.

Keywords: Proximity drawings, · Gabriel drawings, · witness proximity drawings, · simultaneous drawing of two graphs.

1 Introduction

Proximity drawings, including Delaunay triangulations, rectangle of influence drawings, minimum spanning trees, and unit disk graphs, are among the most studied geometric graphs. They are commonly used as descriptors of the "shape" of a point set and are used in a variety of applications, including machine learning, pattern recognition, and computer graphics (see, e.g., [14]). They have also been used to measure the faithfulness of large graph visualizations (see, e.g., [11]).

Proximity drawings are geometric graphs in which two vertices are adjacent if and only if they are deemed close by some measure. A common approach to define the closeness of two vertices u and v uses a region of influence of u and v, which is a convex region whose shape depends only on the relative position of u with respect to v. Then we say that u and v are adjacent if and only if their region of influence does not contain some obstacle, often another vertex of the

^{*} Work partially supported by: (i) MUR, grant 20174LF3T8 AHeAD: efficient Algorithms for HArnessing networked Data", (ii) Dipartimento di Ingegneria, Universita degli Studi di Perugia, grant RICBA21LG: Algoritmi, modelli e sistemi per la rappresentazione visuale di reti.



Fig. 1. A mutual witness Gabriel drawing of two trees (Gabriel disks are dotted).

drawing. For example, a *Gabriel drawing* Γ is a proximity drawing where the region of influence u and v is the disk having u and v as antipodal points, called the Gabriel region of u and v; u and v are adjacent in Γ if and only if their Gabriel region does not contain any other vertex. See also [16] for a survey on different types of proximity regions and drawings.

An interesting generalization of proximity drawings is given in a sequence of papers by Aronov, Dulieu, and Hurtado who introduce and study witness proximity drawings and mutual witness proximity drawings [2,3,4,5]. In a witness proximity drawing the obstacles are points, called witnesses, that are suitably placed in the plane to impede the existence of edges between non-adjacent vertices; these points may or may not include some of the vertices of the drawing itself. A mutual witness proximity drawing is a pair of witness proximity drawings that are computed simultaneously and such that the vertices of one drawing are the witnesses of the other drawing. For example, Figure 1 depicts a mutual witness Gabriel drawing (MWG-drawing for short) of two trees. In the figure, the Gabriel disk of v_0, v_1 of Γ_0 includes vertex v_2 but no vertices of Γ_1 and hence v_0, v_1 are adjacent in Γ_0 ; conversely, v_1 and v_2 are not adjacent in Γ_0 because their Gabriel disk contains vertex u_1 of Γ_1 .

In this paper we characterize those pairs of complete bipartite graphs that admit an MWG-drawing. While every complete bipartite graph has a witness Gabriel drawing [3], not all pairs of complete bipartite graphs admit an MWGdrawing. To characterize the drawable pairs we also investigate some properties of MWG-drawings that go beyond complete bipartiteness. More precisely:

- We show that if $\langle \Gamma_0, \Gamma_1 \rangle$ is an MWG-drawing such that both Γ_0 and Γ_1 have diameter two, then the set of vertices of Γ_0 is linearly separable from the set of vertices of Γ_1 . This extends a result of [4], where linear separability is proved when the diameter is one, i.e. when the two graphs are complete.
- We show, perhaps surprisingly, that if $\langle G_0, G_1 \rangle$ is a pair of complete bipartite graphs that admits an MWG-drawing, then both must be planar.

- The above result let us characterize those pairs $\langle G_0, G_1 \rangle$ of complete bipartite graphs that admit an MWG-drawing and leads to a linear time testing algorithm. When the test returns that $\langle G_0, G_1 \rangle$ is drawable, an MWG-drawing can be constructed in linear time in the real RAM model.
- We show that relaxing the bipartiteness assumption does not significantly enlarge the class of representable graph pairs: We consider those pairs of complete multi-partite graphs each having all partition sets of size at least two and prove that if at least one of the graphs in the pair has more than two partition sets, then the pair does not admit an MWG-drawing.

We remark that our contribution not only fits into the rich literature devoted to proximity drawings, but it also relates to two other well studied topics in graph drawing, namely simultaneous embeddings (see, e.g., [8,20] for references) and obstacle representations (see, e.g., [1,6,9,10,12,15,17,18,19]). As in simultaneous embeddings, the coordinates of the vertices of Γ_i in a mutual witness proximity drawing are defined by taking into account the (geometric and topological) properties of Γ_{1-i} ; as in obstacle graph representations, the adjacency of the vertices Γ_i depends on whether their geometric interaction is obstructed by some external obstacles, namely the vertices of Γ_{1-i} ;. Finally, mutual witness proximity drawings are of interest in pattern recognition, where they have been used in the design of trained classifiers to convey information about the interclass structure of two sets of features (see, e.g. [13]).

2 Preliminaries

We assume familiarity with basic definitions and results of graph drawing [7]. We assume that all drawings occur in the Euclidean plane with standard x and y axes, and so concepts such as above/below a (non-vertical) line are unambiguous. Given two distinct points p and q in the plane, we denote by \overline{pq} the straight-line segment whose endpoints are p and q. Also, let a, b, c be three distinct points in the Euclidean plane, we denote by $\Delta(abc)$ the triangle whose vertices are a, b, c. Given two non-axis-parallel lines ℓ_1 and ℓ_2 intersecting at a point b, those lines divide the plane into four wedges: the top, bottom, left, and right wedges of b with respect to ℓ_1 and ℓ_2 . The top and bottom wedges lie entirely above and below the horizontal line through b, respectively; the left and right wedges lie entirely to the left and right of the vertical line through b. When the two lines are determined by providing a point (other than b) on each line, say a and c, we denote the wedges by $W_T[b, a, c], W_B[b, a, c], W_L[b, a, c], and <math>W_R[b, a, c]$ when we want to include the boundary of each wedge as part of that wedge and by $W_T(b, a, c), W_B(b, a, c), W_L(b, a, c),$ and $W_R(b, a, c)$ when we do not.

Note that exactly one of the four wedges will have both a and c on its boundary, we denote that wedge as W[b, a, c] (or W(b, a, c)). See Figure 2(a).

Let Γ be a straight line drawing of a graph G and let u and v be two vertices of Γ (and of G). Vertices u and v may either be adjacent in G and thus \overline{uv} is an edge of Γ or u and v are not adjacent vertices, in which case \overline{uv} is a *non-edge* of



Fig. 2. (a) $W(b, a, c) = W_T(b, a, c), W_B(b, a, c), W_L(b, a, c), \text{ and } W_R(b, a, c)$ (b) If $w \in \Delta(vu_1v_1)$, at most one of $\overline{vu_1}$ and $\overline{vv_1}$ is an edge of a WG-drawing.

 Γ . For example, $\overline{v_0v_1}$ in Fig. 1 is an edge while $\overline{v_1v_2}$ is a non-edge of Γ_1 . Also, the *Gabriel disk* of p and q, denoted as D[p,q] is the disk having p and q as antipodal points; D[p,q] is a closed set.

Let V and P be two sets of distinct points in the plane. A witness Gabriel drawing (WG-drawing) with vertex set V and witness set P is a geometric graph Γ whose vertices are the points of V and such that any two vertices u and v form an edge if and only if $D[u, v] \cap P = \emptyset$. A graph G is witness Gabriel drawable (WG-drawable) if there exist two point sets V and P such that the witness Gabriel drawing with vertex set V and witness set P represents G (i.e., there is a bijection between the vertex set of G and the point set V and between the edge set of G and the point set V and between the edge set of G and the edge set of G and the edge set of F that is incidence-preserving). The following property can be proved with elementary geometric arguments (see also Figure 2(b)).

Property 1. Let Γ be a WG-drawing with witness set P and let $\overline{vu_1}$ and $\overline{vv_1}$ be two edges of Γ incident on the same vertex v. Then $\Delta(vu_1v_1) \cap P = \emptyset$.

For a pair $\langle G_0, G_1 \rangle$ of WG-drawable graphs, a mutual witness Gabriel drawing (MWG-drawing) is a pair $\langle \Gamma_0, \Gamma_1 \rangle$ of straight-line drawings such that Γ_i is a WG-drawing of G_i with witness set the vertices of Γ_{1-i} (i = 0, 1). If $\langle G_0, G_1 \rangle$ admits an MWG-drawing we say that $\langle G_0, G_1 \rangle$ is mutually witness Gabriel drawable (MWG-drawable).

Details of proofs of statements marked with '*' can be found in the appendix.

3 Linear Separability of Diameter-2 MWG-drawings

In this section we extend a result by Aronov et al. [4] about the linear separability of the MWG-drawings of complete graphs to graphs of diameter two.

Lemma 1. Let $\langle \Gamma_0, \Gamma_1 \rangle$ be an MWG-drawing such that Γ_i has diameter at most 2 (i = 0, 1). Then no segment of Γ_0 intersects any segment of Γ_1 .



Fig. 3. Illustration for the proof of Theorem 1: a non-edge $\overline{u_1v_1}$ of Γ_1 in a region \mathcal{R} bounded by portions of edges and non-edges of Γ_0 .

Proof. Note first that by Property 1, a vertex u of Γ_i cannot lie on a non-edge $\overline{u_1, v_1}$ of Γ_{1-i} since $\{u_1, v_1\}$ have at least one common neighbor v. Also no vertex of Γ_i can lie on an edge of Γ_{1-i} . Let $\overline{u_0 v_0}$ be an edge of Γ_0 and let $\overline{u_1 v_1}$ be an edge of Γ_1 . Assume that they cross and consider the quadrilateral Q whose vertices are the end-points of the two crossing edges. Since some internal angle of Q must be at least $\frac{\pi}{2}$, either $D[u_0, v_0]$ contains one of $\{u_1, v_1\}$ or $D[u_1, v_1]$ contains one of $\{u_0, v_0\}$ contradicting the fact that both $\overline{u_0 v_0}$ is an edge of Γ_0 and $\overline{u_1 v_1}$ is an edge of Γ_1 .

Let $\overline{u_0v_0}$ be an edge of Γ_0 and let $\overline{u_1v_1}$ be a non-edge of Γ_1 . Since Γ_1 has diameter at most two, there is a vertex v in Γ_1 such that both $\overline{vu_1}$ and $\overline{vv_1}$ are edges of Γ_1 . Since $\overline{u_1v_1}$ crosses $\overline{u_0v_0}$, but neither $\overline{vu_1}$ nor $\overline{vv_1}$ crosses $\overline{u_0v_0}$, we have that one of $\{u_0, v_0\}$ is a point of $\Delta(u_1, v_1, v)$. However, Γ_1 is a WG-drawing whose witness set is the set of vertices of Γ_0 and, by Property 1, no vertex of Γ_0 can be a point of $\Delta(u_1, v_1, v)$. An analogous argument applies when $\overline{u_0v_0}$ is a non-edge of Γ_0 while $\overline{u_1v_1}$ is an edge of Γ_1 .

It remains to consider the case that $\overline{u_0v_0}$ is a non-edge of Γ_0 and $\overline{u_1v_1}$ is a non-edge of Γ_1 . Let v be a vertex such that both $\overline{vu_1}$ and $\overline{vv_1}$ are edges of Γ_1 . By the previous case, neither of these two edges can cross $\overline{u_0v_0}$. It follows that one of $\{u_0, v_0\}$ is a point of $\Delta(u_1, v_1, v)$ which, by Property 1, is impossible.

We are now ready to prove the main result of this section. We denote by $CH(\Gamma)$ the convex hull of the vertex set of a drawing Γ .

Theorem 1. Let $\langle \Gamma_0, \Gamma_1 \rangle$ be an MWG-drawing such that each Γ_i has diameter 2. Then Γ_0 and Γ_1 are linearly separable.

Proof. By Lemma 1, no vertex, edge or non-edge of Γ_i intersects any vertex, edge or non-edge of Γ_{1-i} . Hence, either $CH(\Gamma_0)$ and $CH(\Gamma_1)$ are linearly separable and we are done, or one of the convex hulls – say $CH(\Gamma_1)$ – is contained in a convex region \mathcal{R} bounded by (portions of) edges and/or non-edges of Γ_0 . We



Fig. 4. Non-linearly separable MWG-drawings of: (a) a diameter two graph and a diameter three graph; (b) a diameter two graph and a diameter one graph.

prove that region \mathcal{R} cannot exist, which implies the statement. Suppose for a contradiction that Γ_1 is contained in \mathcal{R} and let $\overline{u_1v_1}$ be a non-edge of Γ_1 with $x(u_1) \leq x(v_1)$. See Figure 3 for a schematic illustration. Since u_1 and v_1 are not adjacent, there is some vertex v of Γ_0 such that $v \in D[u_1, v_1]$. Without loss of generality, assume that $\overline{u_1v_1}$ is horizontal and that v is below the line through u_1 and v_1 are points of $W_T[v, u_1, v_1]$ and Γ_1 is contained in \mathcal{R} , there is some segment \overline{pq} of the boundary of \mathcal{R} such that \overline{pq} intersects $W_T(v, u_1, v_1)$ above the line through u_1 and v_1 . Let u_0 and v_0 be the two vertices of Γ_0 such that \overline{pq} is a subset of $\overline{u_0v_0}$. For concreteness, we assume that the x-coordinates of u_0, p, q , and v_0 are such that $x(u_0) \leq x(p) \leq x(q) \leq x(v_0)$.

Claim. $u_0 \in W_L(v, u_1, v_1)$, $v_0 \in W_R(v, u_1, v_1)$ and $\overline{u_0 v_0}$ is a non-edge of Γ_0 .

Proof of the claim: Suppose for a contradiction that a vertex of $\{u_0, v_0\}$ – say v_0 – were a point of $W_T[v, u_1, v_1]$. Since \overline{pq} intersects $W_T(v, u_1, v_1)$ above the horizontal line through u_1 and v_1 , we have that v_0 must also be above this horizontal line or else $\overline{u_0v_0}$ and $\overline{u_1v_1}$ would cross, contradicting Lemma 1. However, if v_0 is above the line through u_1 and v_1 we have that $\overline{u_1v_1}$ and $\overline{vv_0}$ cross which again contradicts Lemma 1. Therefore, $v_0 \notin W_T[v, u_1, v_1]$ and, by the same argument, $u_0 \notin W_T[v, u_1, v_1]$. Note that this argument also precludes either point of $\{u_0, v_0\}$ from being in $W_B[v, u_1, v_1]$, since, because \overline{pq} intersects $W_T[v, u_1, v_1]$, we would then have that the other point of $\{u_0, v_0\}$ lies in $W_T[v, u_1, v_1]$. Finally, observe that if u_0 and v_0 were both points of either $W_L(v, u_1, v_1)$ or $W_R(v, u_1, v_1)$, segment \overline{pq} would not intersect $W_T(v, u_1, v_1)$. It follows that $u_0 \in W_L(v, u_1, v_1)$ and $v_0 \in W_R(v, u_1, v_1)$. Note that $\angle(u_0vv_0)$ contains both u_1 and v_1 , so $\angle u_0vv_0 > \angle u_1vv_1 \ge \frac{\pi}{2}$, which implies that $\angle u_0v_1v_0 > \frac{\pi}{2}$ and so $\overline{u_0v_0}$ is a non-edge of Γ_0 . This concludes the proof of the claim.

By the claim above and by the assumption that Γ_0 has diameter two, there is some vertex z such that both $\overline{zu_0}$ and $\overline{zv_0}$ are edges of Γ_0 . Vertex z may

or may not coincide with v. If z coincides with v or if $z \in W_B[v, u_1, v_1]$, we have that $\Delta(zu_0v_0)$ contains both u_1 and v_1 and two of its sides are edges of Γ_0 , which contradicts Property 1. If $z \in W_T[v, u_1, v_1]$ and z is above the line through u_1 and v_1 , we have that \overline{vz} and $\overline{u_1v_1}$ cross, which contradicts Lemma 1. If $z \in W_T[v, u_1, v_1]$ and z is below the line through u_1 and v_1 , either $\overline{u_1v_1}$ crosses one of $\{\overline{zu_0}, \overline{zv_0}\}$ contradicting Lemma 1, or $\Delta(zu_0v_0)$ contains both u_1 and v_1 contradicting Property 1. If $z \in W_L(v, u_1, v_1)$, we consider three cases. If edge $\overline{zv_0}$ crosses $\overline{u_1v_1}$, we would violate Lemma 1. If edge $\overline{zv_0}$ is above $\overline{u_1v_1}$, then $\angle zvv_0 > \frac{\pi}{2}$ and since both u_1 and v_1 are in the interior of $\Delta(zvv_0)$, we have that u_1 and v_1 are in $D[z, v_0]$, contradicting the fact that $\overline{zv_0}$ is an edge of Γ_0 . If edge $\overline{zv_0}$ is below $\overline{u_1v_1}$, $\Delta(u_0zv_0)$ contains both u_1 and v_1 , which violates Property 1. By a symmetric argument, we have that z cannot be a point of $W_R(v, u_1, v_1)$ either. Since point z does not exist, it follows that \mathcal{R} does not exist.

Theorem 1 shows that MWG-drawings with diameter two capture useful information about the interaction of two point sets. As pointed out by both Ichino and Slansky [13] and by Aronov et al. [4], the linear separability of mutual witness proximity drawings gives useful information about the interclass structure of two set of points. It is also worth noting that if at least one of the graphs in the pair has diameter different from two, a non-linearly separable drawing may exist. For example Figure 4(a) and Figure 4(b) show MWG-drawings of graph pairs in which the diameter two property is violated for one of the two graphs.

4 MWG-drawable Complete Bipartite Graphs

In this section we exploit Theorem 1 to characterize those pairs of complete bipartite graphs that admit an MWG-drawing. In Section 4.1 we prove that any two complete bipartite graphs that form an MWG-drawable pair are planar. The complete characterization is then given in Section 4.2. In what follows we shall assume without loss of generality that the line separating a drawing Γ from its set of witnesses is horizontal and it coincides with the line y = 0, with the witnesses in the negative half-plane. The proof of the following property is trivial and therefore omitted, but Figure 5(a) and its caption illustrate it.

Property 2. Let Γ be a WG-drawing with witness set P, let \overline{uv} be a non-edge of Γ with witness $p \in P$, and let z be a vertex of Γ such that both \overline{zu} and \overline{zv} are edges of Γ . Then $z \in W(p, u, v)$.

4.1 Planarity

Let Γ be a WG-drawing; an alternating 4-cycle in Γ consists of two vertexdisjoint edges $\overline{u_0u_1}$ and $\overline{v_0v_1}$ of Γ such that $\overline{u_0v_0}$ and $\overline{u_1v_1}$ are both non-edges in the drawing. For example, Figure 5(b) shows a WG-drawing Γ whose witness set consists of points p_0 and p_1 . In the figure, $\overline{u_0u_1}$ and $\overline{v_0v_1}$ are edges of Γ while $\overline{u_0v_0}$ and $\overline{u_1v_1}$ are non-edges of Γ : these two pairs of edges and non-edges (bolder in the figure) form an alternating 4-cycle in Γ .



Fig. 5. (a) If $z \notin W(p, u, v)$, then $p \in D[z, v]$ and \overline{zv} is not an edge of Γ . (b) A WG-drawing Γ with an alternating 4-cycle highlighted in bold.

Lemma 2. Let Γ be a WG-drawing of a complete bipartite graph such that Γ is linearly separable from its witness set P and let C be an alternating 4-cycle defined on Γ . The two edges of Γ in C do not cross while the two non-edges of Γ in C do cross.

Proof. Let ℓ be the line separating Γ from its witness set. Let u_0, u_1, v_1, v_0 be the four vertices of C such that $\overline{u_0u_1}$ and $\overline{v_0v_1}$ are two edges of Γ while $\overline{u_0v_0}$ and $\overline{u_1v_1}$ are two non-edges of Γ . Since the drawing is a complete bipartite graph, $\overline{u_0v_1}$ and $\overline{v_0u_1}$ are edges of Γ . We prove that $\overline{u_0v_0}$ and $\overline{u_1v_1}$ must cross in Γ , which implies that $\overline{u_0u_1}$ and $\overline{v_0v_1}$ do not cross.

Let $p_0 \in P$ such that $p_0 \in D[u_0, v_0]$. By Property 2, both u_1 and v_1 lie in the wedge $W(p_0, u_0, v_0) = W_T(p_0, u_0, v_0)$. Observe that p_0 cannot also be a witness for the pair u_1 and v_1 as otherwise, by Property 2, we should have that also u_0 and v_0 lie in the top wedge $W_T(p_0, u_1, v_1)$, which is impossible. So, let $p_1 \in P$ be distinct from p_0 and such that $p_1 \in D[u_1, v_1]$. If p_0 were a point in $W_T[p_1, u_1, v_1]$, p_0 would also be a point in $\triangle(v_1p_1u_1)$ and we would have $p_0 \in D[u_1, v_1]$, which we just argued is impossible (see, e.g. Figure 6). By analogous reasoning we have that p_1 cannot be a point of $W_T[p_0, u_0, v_0]$. Also, $p_1 \notin W_B[p_0, u_0, v_0]$ or else p_0 would be in $W_T[p_1, u_1, v_1]$ since $W_T(p_1, u_1, v_1)$ contains both u_0 and v_0 . It follows that either $p_1 \in W_L(p_0, u_0, v_0)$ or $p_1 \in W_R(p_0, u_0, v_0)$. In either case, $W_T(p_1, u_1, v_1)$ can contain both u_0 and v_0 only if $\overline{u_0v_0}$ and $\overline{u_1v_1}$ cross.

The following corollaries are a consequence of Lemma 2 and of Theorem 1.

Corollary 1 (*). Let G_0 and G_1 be two vertex disjoint complete bipartite graphs. If the pair $\langle G_0, G_1 \rangle$ is MWG-drawable, then both G_0 and G_1 are planar graphs.

Corollary 2 (*). Let Γ be a WG-drawing of a complete bipartite graph such that Γ is linearly separable from its witness set. Any 4-cycle formed by edges of Γ is a convex polygon.



Fig. 6. If $p_0 \in W_T[p_1, u_1, v_1]$, then $p_0 \in D[u_1, v_1]$.

4.2 Characterization

We start with two technical lemmas.

Lemma 3 (*). Let $\langle G_0, G_1 \rangle$ be an MWG-drawable pair admitting a linearly separable MWG-drawing. Then the pair $\langle G_0 \cup \{v_0\}, G_1 \rangle$ also admits a linearly separable MWG-drawing, where v_0 is a vertex not in G_0 and is adjacent to all vertices in G_0 —that is, a universal vertex of $G_0 \cup \{v_0\}$.

Figure 7(b) shows the addition of the universal vertex v_0 to the MWGdrawing of Figure 7(a). In fact, a universal vertex can be added to either drawing as long as it is positioned sufficiently far from the separating line.

Let u_0, u_1 be two points with $x(u_0) < x(u_1)$. The open vertical strip of u_0 , u_1 , denoted as $S(u_0, u_1)$, is the set of points (x, y) such that $x(u_0) < x < x(u_1)$. Assume now that u_0 and u_1 are vertices of a WG-drawing Γ such that Γ is linearly separable from its witness set by a line ℓ . Segment $\overline{u_0u_1}$ divides $S(u_0, u_1)$ into two (open) half-strips: $S_N(u_0, u_1)$ is the (near) half-strip on the same side of $\overline{u_0u_1}$ as ℓ and $S_F(u_0, u_1)$ is the other (far) half-strip. $S[u_0, u_1]$, $S_N[u_0, u_1]$, and $S_F[u_0, u_1]$ consist of $S(u_0, u_1)$, $S_N(u_0, u_1)$, and $S_F(u_0, u_1)$ along with their respective boundaries.

Lemma 4 (*). Let $\langle G_0, G_1 \rangle$ be an MWG-drawable pair admitting a linearly separable MWG-drawing. Then at least one of the pairs $\langle G_0 \cup \{v_0\}, G_1 \rangle$, $\langle G_0, G_1 \cup \{v_1\} \rangle$ also admits a linearly separable MWG-drawing, where, for $i = 0, 1, v_i$ is a vertex not in G_i and has no edges to any vertex in G_i —that is, v_i is an isolated vertex of G_i .

Figure 8(b) shows the addition of the isolated vertex v_0 to the MWG-drawing of Figure 8(a). In fact, an isolated vertex can be added to the left (right) of whichever of the two drawings has the leftmost (rightmost) vertex, as long as it



Fig. 7. (a) A linearly separable MWG-drawing $\langle \Gamma_0, \Gamma_1 \rangle$. (b)Adding a universal vertex v_0 to Γ_0 by placing it far enough from p on ℓ' .

is positioned sufficiently far enough to the left (right) of that vertex. Lemmas 3 and 4 are used in the following lemma, where we use colors to distinguish vertices in distinct partition sets.

Lemma 5. $\langle K_{1,n_0}, K_{1,n_1} \rangle$ has a MWG-drawing if $|n_0 - n_1| \leq 2$.

Proof. Observe that two independent sets whose sizes differ by at most 1, one consisting of red vertices and one consisting of blue vertices, admit a linearly separable MWG-drawing where the red vertices are above a horizonal separating line ℓ while the blue vertices are below ℓ : start with one red vertex with coordinates (0, 1) and a blue vertex with coordinates (-1, -1) and iteratively add red and blue vertices by applying the isolated vertex-addition procedure in the proof of Lemma 4. Let $G_i = K_{1,n_i}$, for i = 0, 1.

Denote by v_0 the non-leaf vertex of G_0 and by u_0 the non-leaf vertex of G_1 . Assume first that $|n_0-n_1| \leq 1$. By the previous observation, $\langle G_0 \setminus \{v_0\}, G_1 \setminus \{u_0\} \rangle$ admits a linearly separable MWG-drawing. Therefore, by Lemma 3 applied to $\langle G_0 \setminus \{v_0\}, G_1 \setminus \{u_0\} \rangle$ we have that $\langle G_0, G_1 \setminus \{u_0\} \rangle$ admits a linearly separable MWG-drawing. By Lemma 3 applied to $\langle G_0, G_1 \setminus \{u_0\} \rangle$ we have that if $|n_0-n_1| \leq 1$, the pair $\langle G_0, G_1 \rangle$ is MWG-drawable.

Consider now the case $|n_0 - n_1| = 2$ and assume that $n_0 > n_1$ (the proof when $n_1 > n_0$ is analogous). Let v_1 be a leaf of G_0 . With the same reasoning as in the previous case, $\langle G_0 \setminus \{v_0, v_1\}, G_1 \setminus \{u_0\} \rangle$ admits a linearly separable MWGdrawing that we denote as $\langle \Gamma_0 \setminus \{v_0, v_1\}, \Gamma_1 \setminus \{u_0\} \rangle$. By the technique in the proof



Fig. 8. (a) A linearly separable MWG-drawing $\langle \Gamma_0, \Gamma_1 \rangle$. (b)Adding an isolated vertex v_0 to Γ_0 .

of Lemma 3, we add the universal vertex u_0 to $\Gamma_1 \setminus \{u_0\}$ in such a way that u_0 is the rightmost vertex of the linearly separable MWG-drawing $\langle \Gamma_0 \setminus \{v_0, v_1\}, \Gamma_1 \rangle$. We now exploit the construction of Lemma 4 to add the isolated vertex v_1 to $\langle \Gamma_0 \setminus \{v_0, v_1\}, \Gamma_1 \rangle$ and obtain a linearly separable MWG-drawing $\langle \Gamma_0 \setminus \{v_0\}, \Gamma_1 \rangle$. Finally, we use Lemma 3 to construct an MWG-drawing of $\langle G_0, G_1 \rangle$ also when $|n_0 - n_1| = 2$.

Lemma 6. Let Γ be a WG-drawing of a graph such that Γ is linearly separable from its witness set P. If \overline{uv} is an edge of Γ and $z \in S_F[u, v]$ is a vertex of Γ , then both \overline{uz} and \overline{vz} are edges of Γ .

Proof. Consider, w.l.o.g., the segment \overline{uz} . If it is not an edge of Γ , then it must have a witness in $S_N(u, z)$. But any such point will also be in D[u, v], contradicting the fact that \overline{uv} is an edge of Γ .

Lemma 7 (*). Let Γ be a WG-drawing of a graph such that Γ is linearly separable from its witness set P. Let u_0, u_1, v_0, v_1 be such that u_0, v_0, u_1, v_1 induce a C_4 in Γ . We have that: (i) v_0 and v_1 are in opposite half-planes with respect to the line through u_0, u_1 , and (ii) one of $\{v_0, v_1\}$ is a point of $S_N(u_0, u_1)$ and the other is not in $S[u_0, u_1]$.

By means of Lemma 7 we can restrict the set of complete bipartite graph pairs that are MWG-drawable.

Lemma 8 (*). Let Γ be a WG-drawing of a complete bipartite graph such that Γ is linearly separable from its witness set. Then Γ does not have $K_{2,3}$ as a subgraph.

We now characterize the MWG-drawable pairs of complete bipartite graphs. We recall that Aronov et al. prove that every complete bipartite graph admits a WG-drawing (Theorem 5 of [3]). The following theorem can be regarded as an analog of the result by Aronov et al. in the context of MWG-drawings.



Fig. 9. An MWG-drawing of $K_{2,2}$ and of an independent set of size three.

Theorem 2. Let $\langle G_0, G_1 \rangle$ be a pair of complete bipartite graphs such that G_i has n_i vertices. The pair $\langle G_0, G_1 \rangle$ admits an MWG-drawing if and only if, for $i = 0, 1, G_i$ is either K_{1,n_i-1} or $K_{2,2}$ and $|n_0 - n_1| \leq 2$.

Proof. By Theorem 1, any MWG-drawing $\langle \Gamma_0, \Gamma_1 \rangle$ of G_0 and G_1 is linearly separable, so any witnes for a non-edge \overline{uv} in Γ_i must lie in S(u, v). By Corollary 1 both G_0 and G_1 must be planar. By Lemma 8 each of the two graphs is either $K_{2,2}$ or a star (i.e. K_{1,n_i-1} , i = 0, 1). Together, these imply that the difference in the cardinalities of the vertex sets in the two graphs is at most two.

If $G_0 = K_{1,n_0-1}$ and $G_1 = K_{1,n_1-1}$, the theorem follows by Lemma 5. If $G_0 = K_{2,2}$ and $G_1 = K_{2,2}$ the pair $\langle G_0, G_1 \rangle$ has an MWG-drawing as shown, for example, in Figure 7(a). By removing one of the bottom-most vertices of Γ_1 in Figure 7(a) we obtain an MWG-drawing of $\langle K_{2,2}, K_{1,2} \rangle$ and by removing both the bottom-most vertices of Γ_1 in Figure 7(a) we obtain an MWGdrawing of $\langle K_{2,2}, K_{1,1} \rangle$. To complete the proof we have to show that $\langle K_{2,2}, K_{1,3} \rangle$, $\langle K_{2,2}, K_{1,4} \rangle$, and $\langle K_{2,2}, K_{1,5} \rangle$ are also MWG-drawable pairs. To this end refer to Figure 9 that shows an MWG-drawing $\langle \Gamma_0, \Gamma_1 \rangle$ where Γ_0 is $K_{2,2}$ while Γ_1 is an independent set consisting of three vertices. By applying Lemma 3 we can add a universal vertex to Γ_1 , thus obtaining an MWG-drawing of $\langle K_{2,2}, K_{1,3} \rangle$. In order to construct MWG-drawings of $\langle K_{2,2}, K_{1,4} \rangle$ and of $\langle K_{2,2}, K_{1,5} \rangle$, notice that in Figure 9 v_0 is the leftmost vertex and v_1 is the rightmost vertex of $\langle \Gamma_0, \Gamma_1 \rangle$. By Lemma 4 we can add either one isolated vertex or two isolated vertices to Γ_1 . In the former case we obtain an MWG-drawing of $K_{2,2}$ and of an independent set of size four which can be extended to an MWG-drawing of $\langle K_{2,2}, K_{1,4} \rangle$ by means of Lemma 3. In the latter case, we again use Lemma 3 to add a universal vertex to the drawing of the independent set of size five and obtain an MWG-drawing of $\langle K_{2,2}, K_{1,5} \rangle$.

The following theorem is a consequence of Theorem 2 and of the constructive arguments of Lemma 5 $\,$



Fig. 10. $v_4 \in S_N(u_2, u'_2)$ implies $W_L(u_2, v_2, v_{2+k}) \cap W_B(v_4, u_2, u'_2) = \emptyset$.

Theorem 3. Let $\langle G_0, G_1 \rangle$ be a pair of complete bipartite graphs such that G_0 has n_0 vertices and G_1 has n_1 vertices. There exists an $O(n_0+n_1)$ -time algorithm that tests whether $\langle G_0, G_1 \rangle$ admits an MWG-drawing. In the affirmative case, there exists an $O(n_0 + n_1)$ -time algorithms to compute an MWG-drawing of $\langle G_0, G_1 \rangle$ in the real RAM model of computation.

5 MWG-drawable Complete k-partite graphs

Aronov et al. also showed that there exists a complete multipartite graph, namely $K_{3,3,3,3}$, which does not admit a WG-drawing (Theorem 15 of [3]). We extend this result in the context of MWG-drawings by proving the following result.

Theorem 4 (*). Let $\langle G_0, G_1 \rangle$ be a pair of complete multi-partite graphs such that for each of the graphs every partition set has size at least two. The pair is mutually Gabriel drawable if and only if it is $\langle K_{2,2}, K_{2,2} \rangle$.

Proof (Sketch). Let Γ be a WG-drawing of a complete k-partite graph, with $k \geq 2$ such that Γ is linearly separable from its witness set P by a separating line ℓ . Note that any induced subgraph G' of G admits a WG-drawing with witness set P, which can be derived from Γ by removing the vertices not in G'. By this observation, Theorem 1, and Lemma 8 we conclude that if $\langle G_0, G_1 \rangle$ is a pair of complete MWG-drawable multi-partite graphs, then neither G_0 nor G_1 can have $K_{2,3}$ as a subgraph. Therefore we can assume that all partition sets in each of the two graphs have size exactly two. Refer to Figure 10.

The proof proceeds by first showing that $CH(\Gamma)$ is a *convex terrain* with respect to ℓ ; that is for each vertex v on the boundary of $CH(\Gamma)$, the segment from v to ℓ perpendicular to ℓ does not intersect $CH(\Gamma)$. Using this property we can order the vertices of Γ by increasing x-coordinate and show that the *i*-th partite set consists of vertices v_i, v_{i+k} . Let p_i be a witness for vertices v_i and v_{k+i} and let p_j be a witness for vertices v_j and v_{k+j} . Thirdly, we show that if i < j then $p_i \in W_L(p_j, v_j, v_{j+k})$; if i > j then $p_i \in W_R(p_j, v_j, v_{j+k})$. Consider now an MWG-drawing $\langle \Gamma_0, \Gamma_1 \rangle$ where Γ_0 is in the upper half-plane with respect to the separating line ℓ and it has at least three partition sets. Let u_1, u_2 and u_3 be three vertices of Γ_1 that act as witnesses for partition sets $\{v_1, v_{1+k}\}$, $\{v_2, v_{2+k}\}$, and $\{v_3, v_{3+k}\}$, respectively. We have $u_1 \in W_L(u_2, v_2, v_{2+k})$ and $u_3 \in$ $W_R(u_2, v_2, v_{2+k})$. Let u'_2 be the vertex of Γ_1 that is in the same partition set as u_2 and assume that u'_2 is to the right of u_2 (the proof in the other case being symmetric). Let v be a vertex of Γ_0 that is a witness of u_2 and u'_2 (i.e. $v \in D[u_2, u'_2]$).

By Property 2 either $v \in \{v_2, v_{2+k}\}$ or $v \in W_T(u_2, v_2, v_{2+k})$, hence $v \in W_T[u_2, v_2, v_{2+k}]$. Again by Property 2, all vertices of Γ_1 must lie in $W_B[v, u_2, u'_2]$. Because $v \in W_T[u_2, v_2, v_{2+k}]$, $W_B[v, u_2, u'_2]$ is disjoint from $W_L(u_2, v_2, v_{2+k})$. But $u_1 \in W_L(u_2, v_2, v_{2+k})$, a contradiction.

6 Open Problems

The results of this paper naturally suggest many interesting open problems. For example: (i) Can one give a complete characterization of those pairs of complete multipartite graphs that admit an MWG-drawing extending Theorem 4 by taking into account graphs some of whose partition sets have size one? It is not hard to see that the ideas of Lemmas 3 and 5 can be used to construct MWG-drawings of graph pairs of the form $\langle K_{1,\dots,1,n_0}, K_{1,\dots,1,n_1} \rangle$ as long as the number of partition sets of size one in the two graphs differ by at most two. However, this may not be a complete characterization. (ii) Which other pairs of diameter-2 graphs admit an MWG-drawing? (iii) Which pairs of (not necessarily complete) bipartite graphs admit an MWG-drawing? (iv) Finally, it would be interesting to study mutual witness drawings for other proximity regions.

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— Appendix

Proofs omitted from Section 4

Corollary 1 (*). Let G_0 and G_1 be two vertex disjoint complete bipartite graphs. If the pair $\langle G_0, G_1 \rangle$ is MWG-drawable, then both G_0 and G_1 are planar graphs.

Proof. Let $\langle \Gamma_0, \Gamma_1 \rangle$ be an MWG-drawing of $\langle G_0, G_1 \rangle$. By Theorem 1, Γ_0 and Γ_1 are linearly separable. Any two vertex-disjoint edges of Γ_i induce an alternating 4-cycle and thus, by Lemma 2, they cannot cross. And, of course, two edges that share a vertex cannot cross, so both graphs are planar.

Corollary 2 (*). Let Γ be a WG-drawing of a complete bipartite graph such that Γ is linearly separable from its witness set. Any 4-cycle formed by edges of Γ is a convex polygon.

Proof. By the same argument as in Corollary 1, no two edges of Γ can cross. Let C be a 4-cycle formed ny four edges $\overline{u_0v_0}, \overline{v_0u_1}, \overline{u_1v_1}, \overline{v_1u_0}$ of Γ . Since any the two pairs of non adjacent edges of C form an alternating 4-cyle with the two non-edges $\overline{u_0u_1}$ and $\overline{v_0v_1}$, by Lemma 2 we have that non-edges $\overline{u_0u_1}$ and $\overline{v_0v_1}$ cross in Γ , which implies that C is a convex polygon.

Lemma 3 (*). Let $\langle G_0, G_1 \rangle$ be an MWG-drawable pair admitting a linearly separable MWG-drawing. Then the pair $\langle G_0 \cup \{v_0\}, G_1 \rangle$ also admits a linearly separable MWG-drawing, where v_0 is a vertex not in G_0 and is adjacent to all vertices in G_0 —that is, a universal vertex of $G_0 \cup \{v_0\}$.

Proof. Let $\langle \Gamma_0, \Gamma_1 \rangle$ be a linearly separable MWG-drawing of $\langle G_0, G_1 \rangle$ and let ℓ be the horizontal line separating Γ_0 from Γ_1 ; see, for example, Figure 7(a). Assume that Γ_0 is in the top half-plane defined by ℓ and let ℓ' be a vertical line not passing through any vertices of Γ_0 . Let p be the intersection point of ℓ with ℓ' . For each vertex $v \in \Gamma_0$ let ℓ_v be the line passing through v and orthogonal to segment \overline{pv} . Let p_v be the intersection point of ℓ_v with ℓ' and let p' the topmost of such intersection points defined by considering all vertices of Γ_0 . Also, let p''be the topmost intersection point between D[u, u'] and ℓ' , taken over all edges $\overline{uu'}$ of Γ_1 . Let p_{\max} be any point of ℓ' above both p' and p''. A linearly separable MWG-drawing of $\langle G_0 \cup \{v_0\}, G_1 \rangle$ such that v_0 is a universal vertex of G_0 can be obtained from $\langle \Gamma_0, \Gamma_1 \rangle$ by adding v_0 at any point along ℓ' above p_{max} . See, for example, Figure 7(b). By construction, we have that: (i) for any vertex $v \in \Gamma_0$, we have $D[v, v_0] \cap \ell = \emptyset$ which implies that v_0 is adjacent to all vertices of Γ_0 ; (ii) for any pair uu' of vertices of Γ_1 , the Gabriel disk D[u, u'] contains a vertex of $\Gamma_0 \cup \{v_0\}$ if and only if it contains a vertex of Γ_0 ; (iii) since no vertex has been added or deleted in Γ_1 , all edges and non-edges of Γ_0 are maintained in $\langle \Gamma_0 \cup \{v_0\}, \Gamma_1 \rangle.$

Lemma 4 (*). Let $\langle G_0, G_1 \rangle$ be an MWG-drawable pair admitting a linearly separable MWG-drawing. Then at least one of the pairs $\langle G_0 \cup \{v_0\}, G_1 \rangle$, $\langle G_0, G_1 \cup \{v_0\}, G_2 \rangle$, $\langle G_0, G_2 \cup \{v_0\}, G_2 \rangle$, $\langle G_0, G_2 \cup \{v_0\}, G_2 \rangle$, $\langle G_0, G_2 \cup \{v_0\}, G$

 $\{v_1\}$ also admits a linearly separable MWG-drawing, where, for $i = 0, 1, v_i$ is a vertex not in G_i and has no edges to any vertex in G_i —that is, v_i is an isolated vertex of G_i .

Proof. Let $\langle \Gamma_0, \Gamma_1 \rangle$ be a linearly separable MWG-drawing of $\langle G_0, G_1 \rangle$, let ℓ be the horizontal line separating Γ_0 from Γ_1 , and assume that Γ_0 is in the top halfplane defined by ℓ . Let u be the rightmost vertex of the MWG-drawing. If ubelongs to Γ_1 we add an isolated vertex to Γ_0 , else we add an isolated vertex to Γ_1 . For concreteness, we describe the construction in the case that u is a vertex of Γ_1 , the other case being symmetric. See, for example, Figure 8(a). For each vertex $v \in \Gamma_0$ let ℓ_v be the line through u orthogonal to segment \overline{vu} . Let p_v be the intersection point of ℓ_v with ℓ and let p_{\max} be the rightmost of such intersection points defined by considering all vertices of Γ_0 and ℓ_{\max} be the line from the set $\{\ell_v : v \in \Gamma_0\}$ intersecting ℓ at p_{\max} . See, for example, Figure 8(b). Add a vertex v_0 to the right of ℓ_{\max} and above ℓ . We now show that $\langle \Gamma_0 \cup \{v_0\}, \Gamma_1 \rangle$ is an MWG-drawing.

Since v_0 is to the right of any vertex of Γ_1 , v_0 is in the top half-plane and Γ_1 is in the bottom half-plane, by Property 3 we have that for any pair of vertices u_i, u_j of Γ_1 , v_0 is not a point of $D[u_i, u_j]$. It follows that for any pair u_i, u_j of vertices of Γ_1 , the Gabriel disk $D[u_i, u_j]$ contains a vertex of $\Gamma_0 \cup \{v_0\}$ if and only if it contains a vertex of Γ_0 . Also, for any vertex $v \in \Gamma_0$, we have $\angle vuv_0 > \frac{\pi}{2}$ which implies that $u \in D[v, v_0]$ and v_0 is not adjacent to any vertex of Γ_0 . Finally, since no vertex has been added or deleted in Γ_1 , all edges and non-edges of Γ_0 are maintained in $\langle \Gamma_0 \cup \{v_0\}, \Gamma_1 \rangle$. We conclude that if $\langle \Gamma_0, \Gamma_1 \rangle$ is a linearly separable MWG-drawing, also $\langle \Gamma_0 \cup \{v_0\}, \Gamma_1 \rangle$ is a linearly separable MWG-drawing.

The following property can be proved with elementary geometric arguments.

Property 3. Let $\langle \Gamma_0, \Gamma_1 \rangle$ be a linearly separable MWG-drawing of $\langle G_0, G_1 \rangle$, let ℓ be the (horizontal) line separating Γ_0 from Γ_1 , let $i \in \{0, 1\}$, let u_0, u_1 be a pair of non-adjacent vertices of Γ_i with $x(u_0) < x(u_1)$, and let z be a vertex of Γ_{1-i} such that $z \in D[u_0, u_1]$. We have that $z \in S_N(u_0, u_1)$.

Lemma 7 (*). Let Γ be a WG-drawing of a graph such that Γ is linearly separable from its witness set P. Let u_0, u_1, v_0, v_1 be such that u_0, v_0, u_1, v_1 induce a C_4 in Γ . We have that: (i) v_0 and v_1 are in opposite half-planes with respect to the line through u_0, u_1 , and (ii) one of $\{v_0, v_1\}$ is a point of $S_N(u_0, u_1)$ and the other is not in $S[u_0, u_1]$.

Proof. Let λ be the line through u_0 and u_1 . If both v_0 and v_1 lie in the same half-plane defined by λ we have that one of $\overline{v_0u_0}, \overline{v_0u_1}$ crosses one of $\overline{v_1u_0}, \overline{v_1u_1}$, which is impossible by Lemma 2. Hence, Property (i) holds. Let p be a witness point such that $p \in D[u_0, u_1]$; by Property 3 $p \in S_N(u_0, u_1)$. Let v_0 be the vertex in the same half-plane as p with respect to λ . Vertex v_0 must also be a point of $S_N(u_0, u_1)$ or else we would have that either $\angle u_0 pv_0 > \frac{\pi}{2}$ or $\angle u_1 pv_0 > \frac{\pi}{2}$, which is impossible because both $\overline{v_0u_0}$ and $\overline{v_0u_1}$ are edges of Γ .

Now suppose $v_1 \in S[u_0, u_1]$. Then either $v_1 \in S[u_0, v_0]$ or $v_1 \in S[v_0, u_1]$. Assume that $v_1 \in S[u_0, v_0]$ (the other case is analogous). Then, by Lemma 6, $\overline{v_0v_1}$ is an edge of Γ_0 , a contradiction. Thus $v_1 \notin S[u_0, u_1]$ and hence Property (ii) holds.

Lemma 8 (*). Let Γ be a WG-drawing of a complete bipartite graph such that Γ is linearly separable from its witness set. Then Γ does not have $K_{2,3}$ as a subgraph.

Proof. Suppose Γ contains $K_{2,3}$ as a subgraph. Let u_0, u_1, u_2, v_0, v_1 be the vertices of $K_{2,3}$ such that u_0, u_1, u_2 are in the same partition set and assume $x(u_0) \leq x(u_1) \leq x(u_2)$. Consider the 4-cycle C_0 with vertices v_0, u_0, v_1, u_1 and the 4-cycle C_1 with vertices v_0, u_1, v_1, u_2 . By Corollary 2 Γ is such that both C_0 and C_1 are convex quadrilaterals, which implies that v_0, u_1 , and v_1 are vertically aligned points. However, by Lemma 7 one of $\{v_0, v_1\}$ -say v_0 -must be a point of $S_N(u_0, u_1)$ that is $x(u_0) < x(v_0) < x(u_1)$, a contradiction.

Proofs omitted from Section 5

Observe that if Γ is a WG-drawing of a graph G with witness set P, then any induced subgraph G' of G admits a WG-drawing with witness set P, which can be derived from Γ by removing the vertices not in G'. Also, by Theorem 1, if G_0 and G_1 are MWG-drawable, their drawings must be linearly separable. By this observation and Lemma 8 we conclude that if $\langle G_0, G_1 \rangle$ is a pair of complete MWG-drawable multi-partite graphs, then neither G_0 nor G_1 can have $K_{2,3}$ as a subgraph. It follows that in order to prove Theorem 4 we can assume that all partition sets in each of the two graphs have size exactly two.

To prove Theorem 4 we introduce some terminology and a technical result from Aronov et al [3]. Let ℓ be a line and let Π be a convex polygon in one of the half-planes of ℓ . We say that Π is a *convex terrain with respect to* ℓ if for each vertex v on the boundary of Π , the segment from v to ℓ perpendicular to ℓ does not intersect Π . To be consistent with [3], we associate each partition set of Γ with a distinct color. We shall also sometimes say that two vertices are in the same color class to mean that they belong to the same partition set.

Lemma 9. [3] Let Γ be a WG-drawing of a complete k-partite graph, with $k \geq 2$. Then $CH(\Gamma)$ contains at least two vertices of each color class on its boundary.

Lemma 10. Let Γ be a WG-drawing of a complete k-partite graph, with $k \geq 2$ such that Γ is linearly separable from its witness set by a separating line ℓ and such that every color class of Γ has size two. Then $CH(\Gamma)$ is a convex terrain with respect to ℓ .

Proof. As usual, we assume that ℓ is a horizontal line. Note that by Lemma 9, every vertex of Γ is on the boundary of $CH(\Gamma)$. Label the vertices as v_1, \ldots, v_{2k}



Fig. 11. $p_1 \in W_L(p_2, v_2, v_5)$ and $p_3 \in W_R(p_2, v_2, v_5)$.

in order of increasing x-coordinate (we can assume that the vertices have distinct x-coordinates). We first show that, for i = 1, ..., k, v_i and v_{k+i} have the same color.

Let $u_1 = v_1$ and let u'_1 be the other vertex having the same color as u_1 ; call this color c_1 . Now let w_1 and w'_1 be two vertices of the same color (but different from c_1). By Lemma 7, exactly one of w_1, w'_1 is in $S(u_1, u'_1)$; in fact it is in $S_N(u_1, u'_1)$. Thus, $S_N(u_1, u'_1)$ contains exactly one vertex of every color other than c_1 . Thus $u'_1 = v_{k+1}$ and v_1, \ldots, v_k all have distinct colors c_1, \ldots, c_k .

Now let $u_2 = v_2$ and let u'_2 be the other vertex having the same color as u_2 ; so u_2 has color c_2 . By the same argument, $S(u_2, u'_2)$ (in fact $S_N(u_2, u'_2)$) contains exactly one vertex of each of the other colors and so u'_2 must be v_{k+2} , and, similarly, for $i = 1, \ldots k$, we have that $v_i = v_{k+i}$, and that $v_{i+1}, \ldots, v_{k+i-1}$ are in $S_N(v_i, v_{k+i})$.

Thus v_1, \ldots, v_{2k}, v_1 is a convex polygon and so every diagonal v_i, v_j is contained within $CH(\Gamma)$ and thus, for each v_i , the vertical segment from v_i to ℓ does not intersect $CH(\Gamma)$, making $CH(\Gamma)$ a convex terrain with respect to ℓ .

An example of the configuration described in the statement of Lemma 10 is shown in Figure 11. In the following lemma we assume, without loss of generality, that for every $p \in P$, p is contained in the Gabriel disk of some pair of vertices of Γ . Indeed, if P contained a point p which is not in any Gabriel disk of any pair of non-adjacent vertices of Γ , then $P \setminus \{p\}$ would also be a witness set for Γ .

Lemma 11. Let Γ be a WG-drawing of a complete k-partite graph, with $k \geq 2$ such that Γ is linearly separable from its witness set P by a separating line ℓ and such that every color class of Γ has size two. Let p_i be a witness for vertices v_i and v_{k+i} and let p_j be a witness for vertices v_j and v_{k+j} . If i < j then $p_i \in W_L(p_j, v_j, v_{j+k})$; if i > j then $p_i \in W_R(p_j, v_j, v_{j+k})$. *Proof.* We assume, w.l.o.g, the separating line is horizontal and that Γ is in the upper half-plane (see for example Figure 11). We begin with an observation. Let $\{v_i, v_{i+k}\}$ be one of the color classes of Γ and let $p_i \in P$ be in $D[v_i, v_{i+k}]$. Since every other vertex v of Γ is adjacent to both v_i and v_{i+k} , by Property 2 it follows that v must lie in $W_T(p_i, v_i, v_{i+k})$ and that no $p \in P$ can be a witness for two different color classes.

It is also the case that if $\{v_i, v_{i+k}\}$ is a color class distinct from $\{v_j, v_{j+k}\}$, then no witness p_i for $\{v_i, v_{i+k}\}$ can be contained in $W_T(p_j, v_j, v_{j+k})$, because then p_i would be contained in $\triangle(p_j, v_j, v_{j+k})$, which would imply that

$$\angle v_j, p_i, v_{j+k} > \angle v_j, p_j, v_{j+k} \ge \frac{\pi}{2},$$

which makes p_i a witness for $\{v_j, v_{j+k}\}$ as well as for $\{v_i, v_{i+k}\}$.

Similarly, p_i cannot be contained in $W_B(p_j, v_j, v_{j+k})$ since then, in order for $W_T(p_i, v_i, v_{i+k})$ to contain both of $\{v_j, v_{j+k}\}, W_T(p_i, v_i, v_{i+k})$ would also need to contain p_j , which we just showed cannot occur. So, either $p_i \in W_L(p_j, v_j, v_{j+k})$ or $p_i \in W_R(p_j, v_j, v_{j+k})$.

Assume that i < j (the proof when i > j is analogous). Since, by Lemma 10, the vertices of Γ form a convex terrain with respect to the separating line, we have that p_i is in $W_L(p_j, v_j, v_{j+k})$, as otherwise $W_T(p_i, v_i, v_{i+k})$ would not contain v_j .

We are now ready to prove Theorem 4. As noted at the beginning of this section, it suffices to consider graphs all of whose partitions sets have size two.

Theorem 4 (*). Let $\langle G_0, G_1 \rangle$ be a pair of complete multi-partite graphs such that for each of the graphs every partition set has size at least two. The pair is mutually Gabriel drawable if and only if it is $\langle K_{2,2}, K_{2,2} \rangle$.

Proof. Assume that Γ_0 is above the separating line and that it has at least three color classes. Let u_1, u_2 and u_3 be three vertices of Γ_1 that act as witnesses for color classes $\{v_1, v_{1+k}\}, \{v_2, v_{2+k}\}, \text{ and } \{v_3, v_{3+k}\}, \text{ respectively. By Lemma 11}, we have <math>u_1 \in W_L(u_2, v_2, v_{2+k})$ and $u_3 \in W_R(u_2, v_2, v_{2+k})$. Let u'_2 be the vertex of Γ_1 that is in the same color class as u_2 and assume that u'_2 is to the right of u_2 (the proof in the other case being symmetric). Let v be a vertex of Γ_0 that is a witness of u_2 and u'_2 (i.e. $v \in D[u_2, u'_2]$). See for example Figure 10 where vertex v is v_4 . Vertex v is above the separating line and, by Property 3, $v \in S_N(u_2, u'_2)$. Also either v coincides with one of $\{v_2, v_{2+k}\}$ or, by Property 2 $v \in W_T(u_2, v_2, v_{2+k})$. In either case, we have that $W_L(u_2, v_2, v_{2+k}) \cap W_B(v, u_2, u'_2) = \emptyset$. However, $W_B(v, u_2, u'_2)$ must, by Property 2, contain all vertices of Γ_1 , a contradiction since $u_1 \in W_L(u_2, v_2, v_{2+k})$.