Bar 1-Visibility Drawings of 1-Planar Graphs

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Abstract. A bar 1-visibility drawing of a graph G is a drawing of G where each vertex is drawn as a horizontal line segment called a bar, each edge is drawn as a vertical line segment where the vertical line segment representing an edge must connect the horizontal line segments representing the end vertices and a vertical line segment corresponding to an edge intersects at most one bar which is not an end point of the edge. A graph Gis bar 1-visible if G has a bar 1-visibility drawing. A graph G is 1-planar if G has a drawing in a 2-dimensional plane such that an edge crosses at most one other edge. In this paper we give linear-time algorithms to find bar 1-visibility drawings of diagonal grid graphs and maximal outer 1-planar graphs. We also show that recursive quadrangle 1-planar graphs and pseudo double wheel 1-planar graphs are bar 1-visible graphs.

1 Introduction

A 1-planar drawing of a graph G is a drawing of G on a two dimensional plane where an edge can be crossed by at most another edge. A graph G is 1-planar if G has a 1-planar drawing. A straight-line drawing of a graph G is a drawing of G such that every edge of G is drawn as a straight-line segment. A right angle crossing drawing or RAC drawing is a straight-line drawing where any two crossing edges form right angles at their intersection point. A RAC graph is a graph that has a RAC drawing. A bar 1-visibility drawing of a graph G is a drawing of G where each vertex is drawn as a horizontal line segment called a bar, each edge is drawn as a vertical line segment where the vertical line segment representing an edge must connect the horizontal line segments representing the end vertices and a vertical line segment corresponding to an edge intersects at most one bar which is not an end point of the edge. A graph G is a bar 1-visibility drawing. A bar 1-visible graph and a bar 1-visibility drawing of the same graph is shown in Figures 1 (a), and (b), respectively.

Recently 1-planar graphs and RAC graphs have been extensively studied [2, 8, 10, 16]. Pach and Toth proved that 1-planar graphs with n vertices have at most 4n - 8 edges [14], whereas Didimo *et al.* showed that a RAC graph with n > 3 vertices has at most 4n - 10 edges [7]. Recognizing the both classes of graphs are NP-hard [1, 11]. Recently Eades and Liotta studied the relationship between dense RAC graphs and dense 1-planar graphs [9].

The concept of bar visibility drawing came up in the early 1980s when many new problems in visibility theory arose, originally inspired by applications dealing with determining visibilities between different electrical components in VLSI

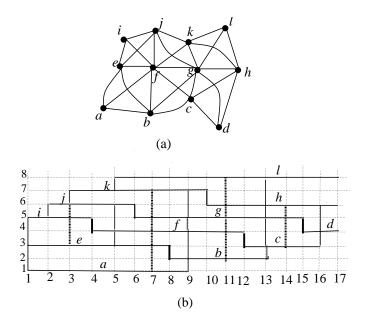


Fig. 1. (a) A bar 1-visible graph, and (b) a bar 1-visibility drawing of the same graph.

design. Other applications arise when large graphs are to be displayed in a transparent way, and in the rapidly developing field of computer graphics. A bar visibility drawing of a planar graph G is a drawing of G where each vertex is drawn as a horizontal line segment and each edge is drawn as a vertical line segment where the vertical line segment representing an edge must connect the horizontal line segments representing the end vertices. Otten and Wijk [13] have shown that every planar graph admits a visibility drawing, and Tamassia and Tollis [17] have given a linear-time algorithm for constructing a visibility drawing of a planar graph. Dean *et al.* have introduced a generalization of visibility drawing for a non-planar graph which is called bar k-visibility drawing [4]. In a bar k-visibility drawing of a graph a horizontal line corresponding to a vertex is called a *bar*, and the vertical line segment corresponding to an edge intersects at most k bars which are not end points of the edge. Thus a visibility drawing is a bar k-visibility drawing for k = 0.

In this paper we give linear-time algorithms to find bar 1-visibility drawings of diagonal grid graphs and maximal outer 1-plane graphs which are RAC drawable graphs. We also develop algorithms for finding bar 1-visibility drawings of recursive quadrangle 1-planar graphs and pseudo double wheel 1-planar graphs which are not RAC graphs.

The rest of the paper is organized as follows. Section 2 describes some of the definitions that we have used in our paper. Sections 3 deals with linear algorithms for finding bar 1-visibility drawings of diagonal grid graphs and maximal outer 1-plane graphs and Section 4 deals with algorithms for finding bar 1-visibility drawings of recursive quadrangle 1-planar graphs and pseudo double wheel 1-planar graphs. Finally, Section 5 concludes our paper with a list of open problems.

2 Preliminaries

In this section we introduce some terminologies and definitions which will be used throughout the paper. For the graph theoretic definitions which have not been described here, see [6, 12].

A graph is *planar* if it can be embedded in the plane without edge crossing except at the vertices where the edges are incident. A *plane graph* is a planar graph with a fixed planar embedding. A plane graph divides the plane into some connected regions called the *faces*. The unbounded region is called the *outer face* and all the other faces are called the *inner faces*. The vertices on the outer face are called the *outer vertices* and all the other vertices are called *inner vertices*.

A $p \times q$ -grid graph is a graph whose vertices correspond to the grid points of a $p \times q$ -grid in the plane and edges correspond to the grid lines between two consecutive grid points. A diagonal grid graph $G_{p,q}$ is a $p \times q$ -grid graph with diagonal edges are introduced in each cell. Figure 2(a) shows a $p \times q$ -grid graph and Figure 2(b) shows a diagonal grid graph $G_{p,q}$. Let *abcd* be a cell of a diagonal grid graph as illustrated in Figure 2(c), where *a* is the bottom-left vertex, *b* is the bottom-right vertex, *c* is the up-right vertex and *d* is the up-left vertex. We call the edge (a, c) the right-diagonal edge and the edge (b, d) the left-diagonal edge of the cell *abcd*.

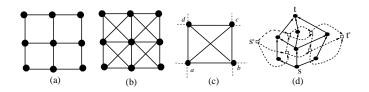


Fig. 2. (a) Grid graph, (b) diagonal grid graph, (c) one cell of a diagonal grid graph and (d) constructing the dual graph G^* from planar *st*-graph.

An acyclic digraph with a single source s and a single sink t is called an st-graph. A plane st-graph is an st-graph that is planar and embedded with vertices s and t on the boundary of the outer face. Let G be a plane st-graph and F be its set of faces. F contains two representatives of the outer face: the left outer face s^* , which is incident with the edges on the left boundary of G and the right outer face t^* , which is incident with the edges on the right boundary of G. Additionally, for each e = (u, v) we define orig(e) = u and dest(e) = v. Also, we define left(e) (respectively right(e)) to be the face to the left (respectively right) of e. We now define a dual-like graph G^* of G as follows. The vertex set of G^* is the set F of faces of G, and G^* has an edge $e^* = (f, g)$ for each edge $e \neq (st)$ of G, where f = left(e) and g = right(e). In Figure 2(d) the vertices and edges of G^* are drawn by white rectangles and dotted lines respectively.

Let G be a digraph with n vertices and m edges. A topological numbering of G is an assignment of numbers to the vertices of G, such that, for every edge (u, v) of G, the number assigned to v is greater than the one assigned to u. If the edges of

digraph G have nonnegative weights associated with them, a weighted topological numbering is a topological numbering of G, such that, for every edge (u, v) of G, the number assigned to v is greater than or equal to the number assigned to u plus the weight of (u, v). The numbering is optimal if the range of numbers assigned to the vertices is minimized.

Let G be a planar st-graph with n vertices. Two paths π_1 and π_2 of G are said to be non-intersecting if they are edge disjoint and do not cross at common vertices but they can touch at vertices. Let Π be a collection of non-intersecting paths of G. In the visibility drawing of G, for every path π of Π , if the edges of π are vertically aligned then the drawing is called *constrained visibility drawing*. The following result on constrained visibility drawing is known [3].

Lemma 1. Let G be a planar st-graph with n vertices, and let Π be a set of nonintersecting paths of G. Then one can find a constrained visibility drawing of G in O(n) time with $O(n^2)$ area, where the edges of every path π in Π are vertically aligned.

A 1-planar drawing of a graph G is a drawing of G on a two dimensional plane where an edge can be crossed by at most another edge. A graph G is 1-planar if G has a 1-planar drawing. A 1-planar graph G is optimal if no edges can be added to G without losing 1-planarity. That is, an optimal 1-planar graph of n vertices has the highest number of edges among all 1-planar graphs of n vertices. An outer-1plane graph is a topological embedding of a graph such that all vertices lie on the outer face and there is at most one crossing on each edge. An outer-1-plane graph G = (V, E) is a maximal outer-1-plane graph if for each pair u, v of vertices where (u, v) is not an edge, adding the edge (u, v) to G makes it not outer-1-plane; that is, $G' = (V, E \cup (u, v))$ is not outer-1-plane for every drawing of the edge (u, v).

3 Bar 1-Visibility Drawings of 1-Planar RAC Graphs

Some interesting labeling properties of diagonal grid graphs have been studied recently by Selvaraju and Pricilla [15]. Recently Dehkordi showed that outer-1-plane graphs are RAC graphs [5]. In this section we give linear-time algorithms for finding bar 1-visibility drawings of diagonal grid graphs and maximal outer 1-plane graphs which are RAC drawable graphs. In Section 3.1 we develop an algorithm for finding a bar 1-visibility drawing of a diagonal grid graph and in Section 3.2 we develop an algorithm for finding a bar 1-visibility drawing of a maximal outer 1-plane graph.

3.1 Diagonal Grid Graphs

In this Section we prove the following theorem.

Theorem 1. A bar 1-visibility drawing of a diagonal grid graph can be drawn in linear-time.

Proof. Let $G_{p,q}$ be a diagonal grid graph as illustrated in Figure 3(a). Clearly $G_{p,q}$ is RAC drawable and 1-planar. We will prove that $G_{p,q}$ is bar 1-visible by

constructing a bar 1-visibility drawing of $G_{p,q}$. We first obtain a graph G from $G_{p,q}$ by deleting the left-diagonal edge of each cell. Clearly G is a plane graph as illustrated in Figure 3(b). Let v_{ij} , $1 \leq i \leq p$ and $1 \leq j \leq q$, be the vertex corresponding to the grid point on the *i*th row and *j*th column of the $p \times q$ grid. We now assign a number $Y(v_{ij})$ to each vertex v_{ij} as follows. We set $Y(v_{1,j}) = j$ for $1 \leq j \leq q$, and for $1 < i \leq p$ we set $Y(v_{ij}) = Y(v_{(i-1)j}) + 2$. We now construct a directed graph from G by assigning direction to each edge from lower number to higher number. We add a vertex s below the first row and a vertex t above the last row. We also add directed edges $(s, v_{1,j})$ and $(v_{p,j}, t)$ for $1 \leq j \leq q$. Let G_{st} be the resulting digraph as illustrated in Figure 3(c). From the construction one can observe that G_{st} is an *st*-graph. We now construct a visibility drawing of G_{st} as follows [17].

We first construct G_{st}^* of G_{st} and assign unit weights to the edges of G_{st}^* and compute an optimal weighted topological numbering X of G_{st}^* as illustrated in Figure 3(d). We then draw each vertex v as a horizontal line segment $\Gamma(v)$ at y-coordinate Y(v) and between x-coordinate X(left(v)) and X(right(v) - 1). We call X(left(v)) the start of $\Gamma(v)$ and X(right(v) - 1) the end of $\Gamma(v)$. For each edge e, we draw the vertical line segment $\Gamma(e)$ at x-coordinate X(left(e)), between y-coordinate Y(orig(e)) and Y(dest(e)). Let u be the upper-left vertex and v be the bottom-right vertex of a cell in a diagonal grid graph. Then one can observe from the drawing algorithm that X(right(v) - 1) - X(left(u)) = 2.

We now obtain a bar 1-visibility drawing of G from the visibility drawing Γ of G_{st} as follows. We first delete $\Gamma(s)$ and $\Gamma(t)$ from the drawing together with the drawings of the edges incident to s and t. The visibility drawing Γ of G_{st} is illustrated in Figure 3(e). We insert one vertical grid line (column) between the two consecutive vertical grid lines i and j if i = left(e) and j = right(e) for some right diagonal edge e by expanding the drawing towards +x direction. We perform this insertion operation for every i, j. After this insertion operation the difference of x-coordinate of end of $\Gamma(v)$ and start of $\Gamma(u)$ will be three, where u is the upper-left vertex and v is the bottom-right vertex of a cell in a diagonal grid graph. We thus extend the end of $\Gamma(v)$ by 2 unit in +x-direction and the end of $\Gamma(u)$ by one unit to the -x-direction. We can place the deleted left diagonal edges in the vertical segment which will be placed between starting point of the horizontal bar corresponding to bottom-right vertex and end point of the extended horizontal segment corresponding to up-left vertex in each cell. By extending these bars, the right diagonal edge in each cell crosses horizontal bar corresponding to up-left vertex in the drawing. Since all edges including left diagonal edges can be placed at end point and start point of the horizontal bars then only right diagonal edges always pass through one horizontal bar corresponding to the vertex, the drawing becomes a bar 1-visibility drawing. The bar 1-visibility drawing of $G_{p,q}$ is illustrated in Figure 3(f).

Theorem 2. A diagonal grid graph $G_{p,q}$ admits a bar 1-visibility drawing Γ on a grid of size $(q+2p-2) \times (3(p+q)-3)$. Furthermore, Γ can be found in linear-time.

Proof. Let $G_{p,q}$ be a diagonal grid graph. We first obtain a graph G from $G_{p,q}$ by deleting the left-diagonal edge of each cell. Then according to proof of theorem 1, we can get visibility drawing of the graph G. The height of the drawing is

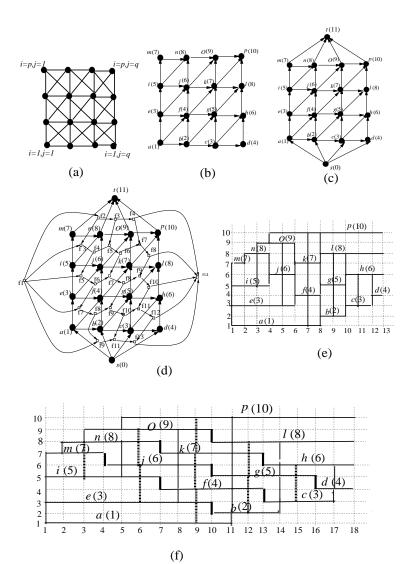


Fig. 3. (a) Diagonal grid graph $G_{p,q}$, (b) graph G with numbering of vertices and (c) the digraph G_{st} , (d) Dual graph G_{st}^* , (e) bar visibility drawing and (f) bar 1-visibility drawing of diagonal grid graph $G_{p,q}$.

 $Y(v_{pq}) = q + 2p - 2$. The width of the drawing is the length of the longest path in G_{st}^* . Since each cell of G contains two faces, the longest path of G_{st}^* is at most 2(p+q). We now compute the size of bar 1-visibility drawing. One can observe that at most (p-1) + (q-1) - 1 = p + q - 3 new columns are inserted for constructing bar 1-visibility drawing of $G_{p,q}$ from the visibility drawing of G. Thus the width of the bar 1-visibility drawing is (3(p+q)-3). Thus we can obtain a bar 1-visibility drawing on a grid of size $(q+2p-2) \times (3(p+q)-3)$.

A diagonal grid graph G has n = pq vertices. If p > q then the area of the bar 1-visibility drawing is $O(p^2)$ and if p < q then the area is $O(q^2)$. When $p \cong q$ then the area of the bar 1-visibility drawing is O(pq) = O(n). The bar 1-visibility drawing obtained by our algorithm is "compact" in a sense that there is at least one line segment for every vertical and horizontal grid line except the last vertical grid line.

3.2 Maximal Outer 1-Planar Graphs

In this Section we give an algorithm for obtaining a bar 1-visibility drawing of a maximal outer 1-plane graph. This problem has an interesting correlation with a constrained visibility drawing of a planar *st*-graph. To describe the algorithm, we have need some definitions.

Let G be a maximal outer 1-plane graph. In a maximal outer 1-plane graph, each edge crossing is surrounded by a cycle of length 4 and no edge in this cycle has a crossing [5]. We call such a cycle of length four a *quadrangle*. An edge e of an outer-1-plane graph is called a *crossing edge* if e has a crossing; otherwise e is called a *non-crossing edge*. In the Figure 4(a), a, b, c, d and b, c, e, f are quadrangles. In quadrangle a, b, c, d, edges (a, c) and (b, d) are crossing edges and edges (a, b), (b, c), (c, d), (d, a) are non-crossing edges. The crossing edges of a quadrangle are called *diagonal* of the quadrangle.

We call a labeling of vertices of a maximal outer 1-plane graph by integers 1 to n (where n is the number of vertices in the graph) a *diagonal labeling* if vertices of each quadrangle got numbers in such a way that two ends of a diagonal of a quadrangle get the lowest and highest numbers among the numbers assigned to the four vertices of the quadrangle.

For example, Figure 4(a) shows an input maximal outer 1-plane graph G and Figure 4(b) shows a diagonal labeling of G.

We can introduce diagonal labeling on the maximal outer 1-plane graph. We have the following lemma.

Lemma 2. Every maximal outer 1-plane graph admits diagonal labeling.

Proof. Let G be a maximal outer 1-plane graph. We will prove that G has a diagonal labeling. Let v be a vertex in G. Then assign 1 to v. After that we will give next numbers to those vertices which are incident to non-crossing edges from v in counterclockwise order. Then assign numbers to the vertices which are incident to crossing edges from v. We then consider the vertex labeled by 2 and assign next numbers to the vertices in the same way. Since the labeling has done always in increasing order and the vertices incident to crossing edges are labeled later, so the

diagonal of a quadrangle got the highest and lowest numbers among the numbers of four vertices of the quadrangle. Figure 4(b) illustrates a diagonal labeling of the maximal outer 1-plane graph in Figure 4(a).

Using this diagonal labeling, we can construct a bar 1-visibility drawing of a maximal outer 1-plane graph as mentioned in the following theorem.

Theorem 3. A bar 1-visibility drawing of a maximal outer 1-plane graph can be drawn in linear-time.

Proof. Let G be a maximal outer 1-plane graph as illustrated in Figure 4(a). By Lemma 2, G has a diagonal labeling as illustrated in Figure 4(b). We first give direction to every edge from lower number to higher number as illustrated in Figure 4(b). We next identify the crossing edge containing the highest number and the lowest number in each quadrangle. We next construct a planar graph G'from G by passing the crossing edge containing the highest number and the lowest number through the vertex which are right side of the edge in each quadrangle as illustrated in Figure 4(c). The crossing edge which is passed through the vertex is identified as a non-intersecting path. Since we pass all the crossing edges through the vertices which are right side of the edges, these satisfies the conditions of nonintersecting paths stated in the result on constrained visibility drawing [3]. In the graph illustrated in Figure 4(c), more than one sink vertices are found. We next construct planar st-graph by adding dummy edges between the sink vertices to the highest labeled sink vertex as illustrated in Figure 4(d). Since the graph is outer planar, the obtained graph remains planar after adding dummy edges. We next construct a constrained visibility drawing of this planar st-graph according to Lemma 1 [3]. All edges can be placed at end point and start point of the horizontal bars. Since crossing edges containing the highest number and the lowest number for all quadrangles always pass through one horizontal bar corresponding to the vertex, the drawing becomes a bar 1-visibility drawing. Since every step of the algorithm can be done in linear-time, a bar 1-visibility drawing of a maximal outer 1-plane graph can be drawn in linear-time [3]. The bar 1-visibility drawing of G is illustrated in Figure 4(e). П

4 Bar 1-Visibility Drawings of 1-Planar Non-RAC Graphs

In the previous section, we showed that diagonal grid graphs and maximal outer 1-planar graphs are two classes of 1-planar graphs. Recently Eades and Liotta showed that every maximally dense RAC graph is 1-planar [9]; on the other hand, they introduced a class of 1-planar graphs which is not RAC drawable. In this section we call the class recursive quadrangle 1-planar graphs. Suzuki studied the combinatorial properties of the optimal 1-planar graphs having 4n - 8 edges [16] which we will define as "pseudo double wheel 1-planar graphs". Since a RAC graph with n > 3 vertices has at most 4n - 10 edges [7], pseudo double wheel 1-planar graphs are not RAC graphs. In this section we show that recursive quadrangle 1planar graphs and pseudo double wheel 1-planar graphs are bar 1-visible graphs. In Section 4.1 we develop an algorithm for finding a bar 1-visibility drawing of a recursive quadrangle 1-planar graph and in Section 4.2 we develop an algorithm for finding a bar 1-visibility drawing of a pseudo double wheel 1-planar graph.

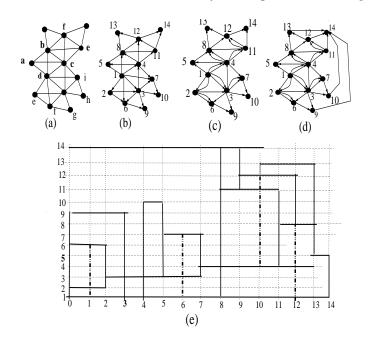


Fig. 4. (a) A maximal outer 1-plane graph G, (b) a diagonal labeling of G, (c) planar graph G', (d) planar *st*-graph and (e) bar 1-visibility drawing of G

4.1 Recursive Quadrangle 1-Planar Graphs

The class recursive quadrangle 1-planar graph is defined recursively, as follows. G_0 is a 1-planar graph of eight vertices . G_0 has an 1-planar drawing as illustrated in Figure 5(a) where the cycle *abcd* is drawn as the outer rectangle and none of the four edges on the outer rectangle has a crossing. Let *abcd* be the outer rectangle of G_i , $i \ge 0$. Graph G_{i+1} is obtained from G_i by adding a new outer rectangle a'b'c'd' and 16 new edges as described in Figure 5(c) and (e), where the four edges on the the outer face do not have any crossing.

Let G be a 1-planar graph and x be a vertex of G and (x, y) be an edge of G. We denote by $\Gamma(G)$, $\Gamma(x)$, $\Gamma(x, y)$ a bar 1-visibility drawing of G, the drawing of a vertex x as a horizontal bar in $\Gamma(G)$, and the drawing of a an edge (x, y) as a vertical line segment in $\Gamma(G)$, respectively. For a vertex x in G, let a and b be the x-coordinates of the two ends of $\Gamma(x)$ such that a < b. We call a and b the *left end* and the *right end* of $\Gamma(x)$, respectively. We now have the following theorem.

Theorem 4. Every recursive quadrangle 1-planar graph G_i , $i \ge 0$ is a bar 1-visible graph.

Proof. We prove the claim by induction on *i*. Let *abcd* be the outer rectangle of G_0 , as illustrated in Figure 5(a). Then G_0 has a bar 1-visibility drawing, as illustrated in Figure 5(b) where (i) $\Gamma(a)$ is the bottommost bar, $\Gamma(c)$ is the topmost bar, $\Gamma(d)$ is the second bottommost bar and $\Gamma(b)$ is the second topmost bar; (ii) starts of $\Gamma(c)$ and $\Gamma(d)$ have the smallest x-coordinate of the drawing and starts of $\Gamma(a)$ and $\Gamma(b)$ have the second smallest x-coordinate of the drawing; and (iii) $\Gamma(a, b)$

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crosses $\Gamma(d)$. We assume that i > 0 and G_j , for j < i, has a bar visibility drawing satisfying (i)-(iii) above. We now show that G_i has a bar 1-visibility drawing satisfying (i)-(iii) above. Let a'b'c'd' be the outer rectangle of G_i . We obtain a graph G_{i-1} by deleting the four vertices on the outer rectangle of G_i . Let *abcd* be the outer rectangle of G_{i-1} . By induction hypothesis, G_{i-1} has a bar 1-visibility drawing satisfying (i)-(iii) as illustrated in Figure 5(d). We now obtain a bar 1visibility drawing of G_i from the visibility drawing of G_{i-1} as follows. Extend the left ends of $\Gamma(a)$, $\Gamma(d)$ and $\Gamma(b)$ by four, four and two units, respectively, to the -x-direction. Extend the right ends of $\Gamma(c)$ and $\Gamma(b)$ by five unit each to the +xdirection. Draw $\Gamma(a')$, $\Gamma(d')$, $\Gamma(b')$ and $\Gamma(c')$ as the bottommost, 2nd bottommost, 2nd topmost and topmost bars outside $\Gamma(G_{i-1})$ and draw the new edges as vertical line segments such that (i)-(iii) are satisfied in $\Gamma(G_i)$, as illustrated in Fig. 5(f).

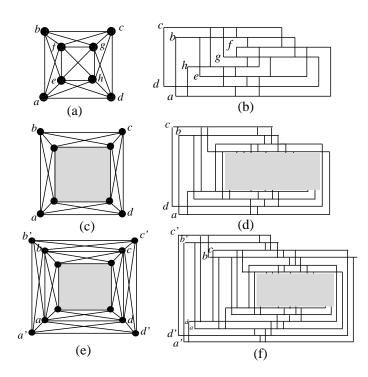


Fig. 5. Illustration for the proof of Theorem 4

By adding two more edges to recursive quadrangle 1-planar graph G of 4p, $p \ge 2$ vertices, we can obtain an optimal 1-planar graph G' of 4p, $p \ge 2$ vertices. From the bar 1-visibility drawing of G we can obtain a bar 1-visibility drawing of G' by adding the drawing of additional two edges as vertical line segment, as illustrated in Figure 6. Thus the following theorem holds.

Theorem 5. Every optimal 1-planar graph obtained from recursive quadrangle 1-planar graph by adding two edges is a bar 1-visible graph.

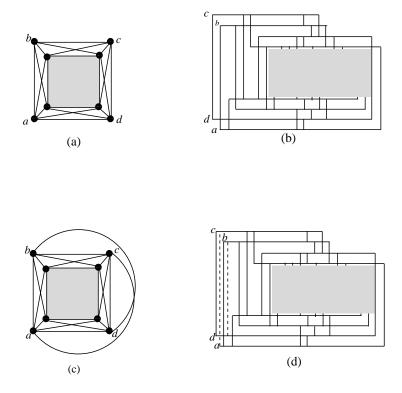


Fig. 6. (a) A recursive quadrangle 1-planar graph G, (b) a bar 1-visibility drawing of G, (c) an optimal 1-planar graph G' obtained from G by adding two edges and (d) a bar 1-visibility drawing of G'.

4.2 Pseudo Double Wheel 1-Planar Graphs

Let C be a cycle $v_1u_1v_2u_2...v_nu_n$ of even number of vertices embedded on a plane. Let x and y be two vertices outside and inside of C, respectively. We add x with u_i and y with v_i for i = 1...n. Let H be the resulting plane graph, as illustrated in Figure 7(a). We add a pair of crossing edges to each face of H, as illustrated in Figure 7(b). The resulting graph is an optimal 1-planar graph as introduced by Suzuki [16]. We call this optimal 1-planar graph even pseudo double wheel 1-planar graph.

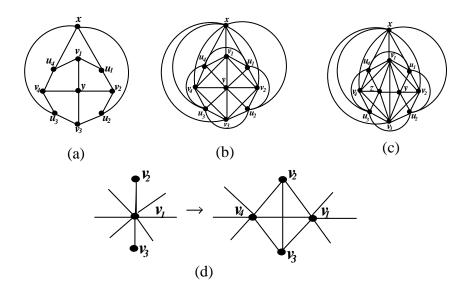


Fig. 7. (a) Pseudo double wheel, (b) even pseudo double wheel 1-planar graph G, (c) odd pseudo double wheel 1-planar graph G' and (d) Q_v splitting.

 Q_v splitting is an expansion operation at v_1 defined as follows: (i) Identify v_2 and v_3 such that there are v_1v_2 and v_1v_3 edges but no v_2v_3 edge. (ii) Split $v_1v_2v_3$ path (iii) Rename one copy of v_1 as v_4 . (iv) Join v_1 , v_4 and v_2 , v_3 . Splitting is illustrated in Figure 7(d). An optimal 1-planar graph obtained from even pseudo double wheel 1-planar graph by one splitting operation is called *odd pseudo double wheel 1-planar graph* as illustrated in Figure 7(c). We now have the following theorem.

Theorem 6. Every pseudo double wheel 1-planar graph is a bar 1-visible graph.

Proof. We will prove every even pseudo double wheel 1-planar graph and odd pseudo double wheel 1-planar graph is a bar 1-visible graph.

Let G be an even pseudo double wheel 1-planar graph. We will prove that G is a bar 1-visible graph by constructing a bar 1-visibility drawing of G. We first draw n bars where 1st topmost bar is y, 2nd topmost bar is x and other n-2 bars starting from 3rd topmost bar are $v_1, u_1, v_2, u_2...v_n, u_n$ respectively. We next draw n-2 vertical lines from $v_1, u_1, v_2, u_2...v_n, u_n$ to y each crossing x and another

n-2 vertical lines from v_1 , u_1 , v_2 , $u_2...v_n$, u_n to x without crossing any bar. We next join each of the v_2 , u_2 , v_3 , $u_3...v_{n-1}$, u_{n-1} bars to 4 bars 1 and 2 unit up and below itself. We next consider v_1 , u_1 , v_n and u_n . We next join v_1 to u_1 , v_1 to v_n , v_1 to u_n and v_n to u_n by vertical lines without crossing any bar. At last we draw a vertical line from u_1 to u_n crossing v_n . Since each vertical line crosses at most one bar, the drawing becomes bar 1-visibility drawing. The bar 1-visibility drawing of G is illustrated in Figure 8(a).

Let G' be an odd pseudo double wheel 1-planar graph of n vertices. We will prove that G' is a bar 1-visible graph by constructing a bar 1-visibility drawing of G'. We first find v_i and v_j on C such that v_i and v_j have degree 8. We next draw n bars where 1st topmost bar is z, 2nd topmost bar is y, 3rd topmost bar is x and other n-3 bars starting from 4th topmost bar are v_1 , u_1 , v_2 , $u_2...v_n$, u_n respectively. We next draw n-3 vertical lines from v_1 , u_1 , v_2 , $u_2...v_n$, u_n to x without crossing any bar. Then join y to $v_i, u_i, v_{i+1}, u_{i+1}...v_j$ by vertical lines. These lines cross bar x. We next draw vertical lines from z to $v_j, u_j, v_{j+1}, u_{j+1}...v_n$, $u_n, v_1, u_1...v_j$ crossing bar x. We next join each of the $v_2, u_2, v_3, u_3...v_{n-1}, u_{n-1}$ bars to 4 bars 1 and 2 unit up and below itself. Now we consider v_1, u_1, v_n and u_n . We next join v_1 to u_1, v_1 to v_n, v_1 to u_n and v_n to u_n, y to z by vertical lines without crossing any bar. We next draw a vertical line from u_1 to u_n crossing v_n . At last we join v_i to v_j . Since each vertical line crosses at most one bar, the drawing becomes bar 1-visibility drawing. The bar 1-visibility drawing of G' is illustrated in Figure 8(b).



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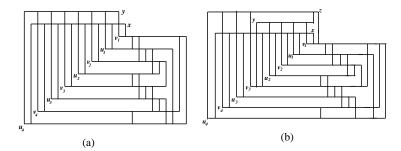


Fig. 8. (a) Bar 1-visibility drawing of even pseudo double wheel 1-planar graph and (b) bar 1-visibility drawing of odd pseudo double wheel 1-planar graph.

5 Conclusion

In this paper we give linear-time algorithms to find bar 1-visibility drawings of diagonal grid graphs and maximal outer 1-planar graphs which are RAC drawable graphs. We also developed algorithms for finding bar 1-visibility drawings of recursive quadrangle 1-planar graphs and pseudo double wheel 1-planar graphs which are not RAC drawable graphs.

Pach and Toth [14] proved that a 1-planar graph can have at most 4n-8 edges whereas Dean *et al.* [4] showed that a bar 1-visible graph can have at most 6n-20edges. The bound for bar 1-visible graph is tight since for each $n \ge 8$ there exists a bar 1-visible graph with exactly 6n-20 edges [4]. Thus not all bar 1-visible graphs are 1-planar graphs. This can be well illustrated by the following example. A bar 1-visibility drawing of an optimal 1-planar graph G of eight vertices in Figure 9(a), is shown in Figure 9(b), but Figure 9(c) shows a bar 1-visibility drawing of a graph of eight vertices which has more edges. Suzuki [16] proved that every optimal 1planar graph can be obtained from a even pseudo double wheel 1-planar graph by a sequence of Q_v splittings and " Q_4 additions". We were able to construct bar 1-visibility drawing of every 1-planar graph that we considered as an example, but yet to find a formal proof. We thus conjecture as follows.

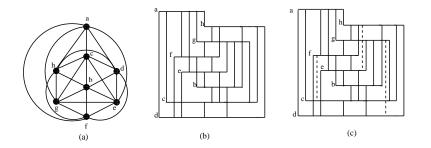


Fig. 9. (a) An optimal 1-planar graph G, (b) a bar 1-visibility drawing of G and (c) a bar 1-visibility drawing of G' obtained by adding some edges to G.

Conjecture 1. Every 1-planar graph is a bar 1-visible graph.

Several interesting open problems have come out from this works.

- 1. Recognition of both RAC graphs and 1-planar graphs are NP-complete problems. It is interesting to know the complexity of recognizing a bar 1-visible graph. Finding a complete characterization of bar 1-visible graph is also an interesting open problem.
- 2. How to find a 1-planar embedding of a 1-planar graph?
- 3. Can we find a complete characterization of bar k-visibility drawing?

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