

Parallel Chip Firing Game Associated with n -cube Edges Orientations

René Ndoundam¹, Claude Tadonki², and Maurice Tchuenté¹

¹ University of Yaoundé I, Faculty of Science,
Department of Computer Science, P.O. Box. 812 Yaoundé, Cameroon

² University of Geneva, Centre Universitaire Informatique,
Département of Theoretical Computer Science 24, rue Général Dufour,
1211 Genève 4 - Switzerland

ndoundam@uycdc.uninet.cm, claudetadonki@cui.unige.ch

Abstract. We study the cycles generated by the chip firing game associated with n -cube orientations. We consider a particular class of partitions of vertices of n -cubes called *left cyclic partitions* that induce *parallel periodic evolutions*. Using this combinatorial model, we show that cycles generated by *parallel evolutions* are of even lengths from 2 to 2^n on H_n ($n \geq 1$), and of odd lengths different from 3 and ranging from 1 to $2^{n-1}-1$ on H_n ($n \geq 4$). However, the question whether there exist *parallel evolutions* with period greater than 2^n remains open.

1 Introduction

A state in the *parallel chip firing game* played on a directed graph $G = (X, A)$ is a mapping $x : V \rightarrow N$ which can be viewed as a distribution of chips onto the vertices of G . In a transition of the game, a state x is transformed into a new state by activating all nodes with more chips than its out-neighbors. The evolution is ultimately *periodic* because the total number of chips remains constant. More precisely, if $x^t, t \geq 0$, denotes the state of the system at time t , then there exists an integer q called *transient length* and another integer p called *period* or *cycle length* such that

$$x^{t+p} = x^t \text{ for } t \geq q, \text{ and } x^{t+p'} \neq x^t \text{ for } p' < p. \quad (1)$$

In this paper, we investigate the dynamics generated by the chip firing game associated with n -cube orientations and we provide a model to study possible periods generated in this particular case.

2 Recurrent Construction of Parallel Cycles

Definition 1. A partition $S_0 \cup S_1 \cup \dots \cup S_{k-1}$ of the vertices of an n -cube is called a *left cyclic partition* if the two following statements hold.

- For all i from 0 to $k-1$, every vertex of S_i has a neighbor in S_{i-1} , where index operations are performed modulo k .
- For all i from 0 to $k-1$, there is no edge between two vertices of S_i .

The model of *left cyclic partition* clearly gives a characteristic of *parallel evolutions* with unique firing within a cycle. In this paper, we investigate on possible configurations. We first present the construction of *left cyclic partitions* of even lengths.

Lemma 1. *An n -cube admits left cyclic partitions of all even lengths from 2 to 2^n .*

Proof. Let $H_n = (V, E)$ be an n -cube and let p be an even integer between 2 and 2^n . It is well known that, since p is even, there is a cycle $[x_0, x_1, \dots, x_{p-1}, x_0]$ of length p in H_n . Now, for every vertex u , let $\Gamma(u)$ denote the set of all neighbors of u in H_n . This notation is naturally extended to a set of vertices. A *left cyclic partition* of order p is obtained as follows.

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For  $i = 0, \dots, p - 1$  do
     $S_i \leftarrow \{x_i\}$ 
endfor
 $S = V - \{x_0, x_1, \dots, x_{p-1}\}$ 
while  $(S \neq \emptyset)$  do
    For  $i \leftarrow 0$  to  $p - 1$  do
         $S_{i+1} \leftarrow S_{i+1} \cup (\Gamma(S_i) \cap S)$ 
         $S \leftarrow S - (\Gamma(S_i) \cap S)$ 
    endfor
endwhile

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It is obvious that S_0, \dots, S_{p-1} is a partition of V and that every vertex in S_i has at least one neighbor in S_{i-1} . So we just need to show that two vertices of the same subset S_i cannot be neighbors. Let a and b be two vertices of S_i .

- There is a path from a to x_0 of length ℓ_1 such that $\ell_1 = i \bmod p$,
- There is a path from b to x_0 of length ℓ_2 such that $\ell_2 = i \bmod p$,

Since p is even, it follows that $\ell_1 = \ell_2 \bmod 2$. Hence, if a and b were neighbors, there would exist a cyclic path of odd length $\ell_1 + \ell_2 + 1$ joining a and b in H_n , which is not possible since H_n is a bipartite graph. This shows that two vertices of the same subset cannot be neighbors.

We now turn to the construction of *left cyclic partitions* of odd lengths.

Lemma 2. *If S_0, S_1, S_2 is a left cyclic partition of $H_n, n \geq 2$, then every vertex of S_i has at least two neighbors in S_{i-1} for $i = 0, 1, 2$.*

Proof. Because of symmetry considerations, we can assume that $i = 2$. So let x be a vertex of S_2 . From the definition of *left cyclic partitions*,

- x has a neighbor $x \oplus e_j$ in S_1 , where \oplus is the XOR operator and e_j is a vector of the canonical basis.
- similarly, $x \oplus e_j$ has a neighbor $x \oplus e_j \oplus e_k$ in S_0 .

Now consider the vertex $x \oplus e_k$.

- It is a neighbor of x , hence it does not belong to S_2 .
- It is a neighbor of $x \oplus e_j \oplus e_k$, hence it does not belong to S_0 .

It then follows that $x \oplus e_k$ belongs to S_1 , hence x admits two neighbors $x \oplus e_j$ and $x \oplus e_k$ which are both in S_1 .

Lemma 3. *If H_n , $n \geq 3$ admits a left cyclic partition of order 3, then H_{n-1} admits a left cyclic partition of order 3.*

Proof. Obvious.

Proposition 1. *n -cubes do not admit left cyclic partitions of order 3.*

Proof. From lemma 2, if a hypercube H admits a left-cyclic partition of order 3, then $|H| \geq 6$, which is not the case for H_2 . By application of lemma 3, we deduce that no n -cube, $n \geq 3$ admits a left cyclic partition of order 3.

Proposition 2. *If S_0, \dots, S_{p-1} is a left cyclic partition of odd order p of H_n , then $p \leq 2^{n-1} - 1$.*

Proof. We just have to show that in such a case, $|S_i| \geq 2$ for $i = 0, \dots, p-1$. Indeed, starting from a vertex $a_{p-1} \in S_{p-1}$, we construct a chain $[a_{p-1}, a_{p-2}, \dots, a_0, b_{p-1}, b_{p-2}, \dots, b_0]$ such that $a_i, b_i \in S_i$ for $i = 0, \dots, p-1$. It is clear that $a_i \neq b_i, i = 0, \dots, p-1$, otherwise we would have displayed a closed path of odd length in H_n which is not possible.

Lemma 4. *If H_n admits a left cyclic partition of order p , then H_{n+1} admits left cyclic partition of order p .*

Proof. If S_0, \dots, S_{p-1} is a left cyclic partition of order p in H_n , then $1S_i \cup 0S_{i-1}, i = 0, \dots, p-1$ is a also left cyclic partition of order p in H_{n+1} .

Lemma 5. *If H_n admits a left cyclic partition of odd order p , $p \geq 5$ then H_{n+1} admits a left cyclic partition of order $2p-1$. Moreover, if $p \geq 7$, then H_{n+1} admits a left cyclic partition of order $2p-3$.*

Proof. Let S_0, S_1, \dots, S_{p-1} be a left cyclic partition of odd order p .

- Case $p \geq 5$

The following sequence is a left cyclic partition of order $2p-1$ in H_{n+1} .

$0S_0, 1S_0 \cup 0S_1, 1S_1, 1S_2, 0S_2, 0S_3, 1S_3, \dots, 1S_{2i}, 0S_{2i}, 0S_{2i+1}, 1S_{2i+1}, \dots, 1S_{p-3}, 0S_{p-3}, 0S_{p-2}, 1S_{p-2}, 1S_{p-1}, 0S_{p-1}$.

- Case $p \geq 7$

A left cyclic partition of order $2p-3$ in H_{n+1} is obtained from the left cyclic partition exhibited in the case $p \geq 5$ by replacing the subsequence $1S_2, 0S_2, 0S_3, 1S_3, 1S_4, 0S_4, 0S_5, 1S_5$ by $1S_2, 0S_2 \cup 1S_3, 0S_3, 0S_4, 1S_4 \cup 0S_5, 1S_5$.

Lemma 6. *H_4 admits left cyclic partitions of orders 5 and 7.*

Proof.

- A left cyclic partition of order 5 in H_4 is the following :

$\{0000, 1101\}, \{0001, 1100, 0010, 1111\}, \{0110, 1011\}, \{0100, 0111, 1001, 1010\}, \{0011, 0101, 1000, 1110\}$.

- A left cyclic partition of order 7 in H_4 is the following :

$\{0000, 1101\}, \{0001, 1100\}, \{0011, 1110\}, \{0010, 1111\}, \{0110, 1011\}, \{0100, 0111, 1001, 1010\}, \{0101, 1000\}$.

Lemma 7. H_n , $n \geq 4$, admits a left cyclic partition of order $2^{n-1} - 1$.

Proof. Consider the sequence $\{u_i; 0 \leq i \leq 2^{n-1} - 1\}$, defined by $u_i = \text{bin}(i) \oplus \text{bin}(i/2)$, where $\text{bin}(x)$ is the n -position binary representation of the integer x , and symbol $/$ denotes integer division. It can be easily checked that this sequence corresponds to a *hamiltonian cycle* in H_{n-1} . Now, let us denote $v_i = u_i \oplus 1 \oplus 2^{n-2}$ (i.e. v_i is obtained from u_i by changing the first and last bits) and $N = 2^n$. It is also easy to check that $\{v_i; 0 \leq i \leq 2^{n-1} - 1\}$ is a *hamiltonian cycle* of H_{n-1} . Now, observe that $0u_i \oplus 1v_i = 2^{n-1} \oplus (u_i \oplus v_i) = 2^{n-1} \oplus 1 \oplus 2^{n-2}$. Hence, $0u_i$ and $1v_i$ are not neighbors in the hypercube H_n . On the other hand, $u_{N-4} = 100\dots010$, $u_{N-2} = 10\dots01$, $u_{N-1} = 10\dots0$ and $v_0 = u_0 \oplus 1 \oplus 2^{n-2} = 10\dots01 = u_{N-2}$. Hence, $0u_{N-4}, 0u_{N-1}, 0u_{N-2}, 1v_0$ is a chain of H_n . Moreover, $v_{N-4} = 0\dots011$, $v_{N-2} = 0\dots0 = u_0$ and $v_{N-1} = 0\dots01 = u_1$. Hence $1v_{N-4}, 1v_{N-1}, 1v_{N-2}, 0u_0$ is a chain of H_n . Hence, by considering the two chains and the two previous *hamiltonian cycles*, we see that the partition $\{0u_0, 1v_0\}, \{0u_{N-4}, 1v_{N-4}\}, \{0u_{N-3}, 0u_{N-1}, 1v_{N-3}, 1v_{N-1}\}, \{0u_{N-2}, 1v_{N-2}\}$ is a *left cyclic partition*.

Proposition 3. H_n , $n \geq 4$, admits left cyclic partitions of all odd orders from 5 to $2^{n-1} - 1$.

Proof. We proceed by induction. For $n = 4$ the result follows from lemma 6. Assuming that the result holds for $n \geq 4$, let us consider an $(n+1)$ -cube together with an odd integer $p \in [5, 2^n - 1]$.

- Case 1 : $5 \leq p \leq 2^{n-1} - 1$. The result follows from the induction hypothesis and lemma 4.
- Case 2 : $2^{n-1} - 1 < p < 2^n - 1$. There is an odd integer q , $7 < q < 2^{n-1} - 1$, such that $p = 2q - 1$ or $p = 2q - 3$. The result follows from the induction hypothesis and lemma 5.
- Case 3 : $p = 2^n - 1$. The result follows from lemma 7.

Theorem 1. The cycles generated by the parallel chip firing game associated with n -cube orientations, $n \geq 4$, are of even lengths from 2 to 2^n , and of odd lengths different from 3 and ranging from 1 to $2^{n-1} - 1$.

3 Conclusion

In this paper, we have given, in the particular case of *parallel evolutions*, a range of possible periods which can be generated by the chip firing game associated with n -cube orientations. The question of the existence of cycles with length greater than 2^n remains to be clarified. We suggest, for the case of n -cubes, to consider a recurrent approach based on the sub-evolutions induced in every face of the n -cube by a *block sequential evolution*.

References

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