One-dimensional quantum cellular automata over finite, unbounded configurations.

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One-dimensional quantum cellular automata (QCA) consist in a line of identical, finite dimensional quantum systems. These evolve in discrete time steps according to a local, shift-invariant unitary evolution. By local we mean that no instantaneous long-range communication can occur. In order to define these over a Hilbert space we must restrict to a base of finite, yet unbounded configurations. We show that QCA always admit a two-layered block representation, and hence the inverse QCA is again a QCA. This is a striking result since the property does not hold for classical one-dimensional cellular automata as defined over such finite configurations. As an example we discuss a bijective cellular automata which becomes non-local as a QCA, in a rare case of reversible computation which does not admit a straightforward quantization. We argue that a whole class of bijective cellular automata should no longer be considered to be reversible in a physical sense. Note that the same two-layered block representation result applies also over infinite configurations, as was previously shown for one-dimensional systems in the more elaborate formalism of operators algebras [9]. Here the proof is made simpler and self-contained, moreover we discuss a counterexample QCA in higher dimensions.

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One-dimensional cellular automata (CA) consist in a line of cells, each of which may take one in a finite number of possible states. These evolve in discrete time steps according to a local, shift-invariant function. When defined over infinite configurations, the inverse of a bijective CA is then itself a CA, and this structural reversibility leads to a natural block decomposition of the CA. None of this holds over finite, yet possibly unbounded, configurations.

Because CA are a physics-like model of computation it seems very natural to study their quantum extensions. The flourishing research in quantum information and quantum computer science provides us with appropriate context for doing so, both in terms of the potential implementation and the theoretical framework. Right from the very birth of the field with Feynman's 1986 paper, it was hoped that QCA may prove an important path to realistic implementations of quantum computers [6] – mainly because they eliminate the need for an external, classical control and hence the principal source of decoherence. Other possible aims include providing models of distributed quantum computation, providing bridges between computer science notions and modern theoretical physics, or anything like understanding the dynamics of some quantum physical system in discrete spacetime, i.e. from an idealized viewpoint. Studying QCA rather than quantum Turing machines for instance means we bother

about the spatial structure of things [2], whether for the purpose of describing a quantum protocol, modelling a quantum physical phenomena, or again taking into account the spatial parallelism inherent to the model.

One-dimensional quantum cellular automata (QCA) consist in a line of identical, finite dimensional quantum systems. These evolve in discrete time steps according to a local, shift-invariant unitary evolution. By local we mean that information cannot be transmitted faster that a fixed number of cells per time step. Because the standard mathematical setting for quantum mechanics is the theory of Hilbert spaces, we must exhibit and work with a countable basis for our vectorial space. This is the reason why we restrict to finite, unbounded configurations. An elegant alternative to this restriction is to abandon Hilbert spaces altogether and use the more abstract mathematical setting of C^* -algebras [3] – but here we seek to make our proofs self-contained and accessible to a wider community, including those computer scientists with an interest in quantum computation. Our main result is that QCA can always be expressed as two layers of an infinitely repeating unitary gate even over such finite configurations. The existence of such a two-layered block representation implies of course that the inverse QCA is again a QCA. Our result is mainly a simplification of the same theorem over infinite configurations as expressed with operators algebra [9]. Unfortunately an example QCA disproves the theorem in further dimensions – at least in its present form.

It is a rather striking fact however that QCA admit the two-layered block representation in spite of their being defined over finite, unbounded configurations. For

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most purposes this saves us from complicated unitary tests such as [1, 4, 5]. But more importantly notice how this is clearly not akin to the classical case, where a CA may be bijective over such finite configurations, and yet not structurally reversible. In order to clarify this situation we consider a perfectly valid, bijective CA but whose inverse function is not a CA. It then turns out that its quantum version is no longer valid, as it allows superluminal signalling. Hence whilst we are used to think that any reversible computation admits a trivial quantization, this turns out not to be case in the realm of cellular automata. Curiously the non-locality of quantum states (entanglement) induces more structure upon the cellular automata – so that its evolution may remain local as an operation (no-superluminal signalling). Based upon these remarks we prove that an important, well-studied class of bijective CA may be dismissed as not physically

Outline. We reorganize a number of known mathematical results around the notion of subsystems in quantum theory (Section I). Thanks to this small theory we prove the reversibility/block structure theorem in an elementary manner (Section II). In the discussion we show why the theorem does not hold as such in further dimensions; we exhibit superluminal signalling in the XOR quantum automata, and end with a general theorem discarding all injective, non surjective CA over infinite configurations as unphysical (Section III).

I. A SMALL THEORY OF SUBSYSTEMS

Definition 1 (Algebras)

reversible.

Consider $\mathcal{A} \subseteq M_n(\mathbb{C})$. We say that \mathcal{A} is an algebra of $M_n(\mathbb{C})$ if and only if it is closed under weighting by a scalar (.), addition (+), matrix multiplication (*), adjoint (†). Moreover for any S a subset of $M_n(\mathbb{C})$, we denote by curly S its closure under the above-mentioned operations.

Note that algebras as above-defined are really just C^* -algebras over finite-dimensional systems.

Definition 2 (Subsystem algebras)

Consider \mathcal{A} an algebra of $M_n(\mathbb{C})$. We say that \mathcal{A} is a subsystem algebra of $M_n(\mathbb{C})$ if and only if there exists $p, q \in \mathbb{N} / pq = n$ and $U \in M_n(\mathbb{C}) / U^{\dagger}U = UU^{\dagger} = \mathbb{I}$ such that $U\mathcal{A}U^{\dagger} = M_p(\mathbb{C}) \otimes \mathbb{I}_q$.

Definition 3 (Center algebras)

For \mathcal{A} an algebra of $M_n(\mathbb{C})$, we note $\mathcal{C}_{\mathcal{A}} = \{A \in \mathcal{A} \mid \forall B \in \mathcal{A} \mid BA = AB\}$. $\mathcal{C}_{\mathcal{A}}$ is also an algebra of $M_n(\mathbb{C})$, which is called the center algebra of \mathcal{A} .

Theorem 1 (Characterizing one subsystem)

Let \mathcal{A} be an algebra of $M_n(\mathbb{C})$ and $\mathcal{C}_{\mathcal{A}} = \{A \in \mathcal{A} \mid \forall B \in \mathcal{A} \mid BA = AB\}$ its center algebra. Then \mathcal{A} is a subsystem algebra if and only if $\mathcal{C}_{\mathcal{A}} = \mathbb{CI}$.

Proof. The argument is quite technical and its understanding not mandatory for understanding the rest of the paper. Its presentation is based on [7]. See also [3] for a proof within the setting of general C^* -algebras. $[\Rightarrow]$.

$$\mathcal{C}_{M_p(\mathbb{C})} = \mathbb{C}\mathbb{I}_p$$
, and hence $\mathcal{C}_{M_p(\mathbb{C})\otimes\mathbb{I}_q} = \mathbb{C}\mathbb{I}$.
 $[\Leftarrow].$

CONSIDER some set $P = \{P_i\}_{i=1...p}$ such that:

(i)
$$\forall i = 1 \dots p \ P_i \in \mathcal{A};$$

(*ii*) $\forall i, j = 1 \dots p P_i P_j = \delta_{ij} P_i \in \mathcal{A}.$

Moreover we take P is maximal, i.e. so that there is no set $Q = \{Q_i\}$ verifying conditions (i), (ii) and such that $\mathcal{P} \subset \mathcal{Q}$, with \mathcal{P}, \mathcal{Q} the closures of P, Q. Note that: $(iii) \sum_{i=1}^{N} P_i = \mathbb{I}$.

(*iii*) $\sum_{i=1...p} P_i = \mathbb{I}$, otherwise $\mathbb{I} - \sum_{i=1...p} P_i$ may be added to the set. FIRST we show that

$$\forall i \quad [P_i \mathcal{A} P_i = \mathbb{C} P_i]. \tag{1}$$

Intuitively this is because the contrary would allow us to refine the subspaces defined by P_i into smaller subspaces, and hence go against the fact that P is maximal. Formally consider some $M \in \mathcal{A}$ such that $P_iMP_i \not \propto P_i$. If M is proportional to a unitary let $N = M + M^{\dagger}$, else let $N = M^{\dagger}M$. In any case we have $P_iNP_i \not \propto P_i$ with N hermitian. Note that $H = P_iNP_i$ is also hermitian, and has support in the subspace P_i , hence we can write $H = \sum_k \lambda_k Q_k$ with the Q_k 's orthogonal projectors such that $\sum_k Q_k = P_i$ and the λ_k 's distinct real numbers. Any such Q_k is part of \mathcal{A} , since $Q_k = \frac{H\prod_{l \neq k}(H - \lambda_l \mathbb{I})}{\lambda_k \prod_{l \neq k}(\lambda_k - \lambda_l)}$. Consider the set $Q = P/\{P_i\} \cup_k \{Q_k\}$. It satisfies condition (i), (ii) but $\mathcal{P} \subset Q$, which is impossible.

SECOND we show that

$$\forall i, j \quad [P_i \mathcal{A} P_j \neq 0]. \tag{2}$$

Intuitively this is because the contrary would split \mathcal{A} into the direct sum of two matrix algebras, and hence go against the fact that $\mathcal{C}_{\mathcal{A}} = \mathbb{CI}$. Formally write $i \sim j$ whenever this is the case. The relation \sim is reflexive by Eq. (1), transitive by multiplicative closure of \mathcal{A} , and symmetric since

$$P_i \mathcal{A} P_j \neq 0 \Rightarrow (P_i \mathcal{A} P_j)^{\dagger} \neq 0$$
$$\Rightarrow (P_j \mathcal{A}^{\dagger} P_i) \neq 0$$
$$\Rightarrow (P_j \mathcal{A} P_i) \neq 0.$$

Say there is an equivalence class $J \subset 1 \dots p$ and let $P_I = \sum_{i \notin J} P_i$, $P_J = \sum_{j \in J} P_j$. Then

$$\begin{aligned} AP_J &= (P_I + P_J) \mathcal{A} P_J \\ &= P_J \mathcal{A} P_J \\ &= P_J \mathcal{A} \quad \text{symmetrically.} \end{aligned}$$

and so $P_J \in \mathcal{C}_A$, which is impossible.

THIRD we show that for all A, for all i, j = 1...p, if we let $M = P_i A P_j$ then

$$\exists \lambda \in \mathbb{C} \quad [M^{\dagger}M = \lambda P_i \wedge MM^{\dagger} = \lambda P_i]. \tag{3}$$

Indeed Eq. (1) gives $M^{\dagger}M = \lambda P_i$ and $MM^{\dagger} = \mu P_j$. But then $\lambda^2 P_i = M^{\dagger}MM^{\dagger}M = \mu M^{\dagger}P_jM = \mu\lambda P_i$, hence λ equals μ .

FOURTH we show that

$$\forall i, j \quad [\operatorname{Tr}(P_i) = \operatorname{Tr}(P_j) = \text{some constant } q]. \quad (4)$$

For each i, j take some $A \in \mathcal{A}$ verifying Eq. (2). Let $M = P_i A P_j$. By Eq. (3) there is a complex number λ such that we have $P_i = \lambda M M^{\dagger}$ and $P_j = \lambda M^{\dagger} M$. Then the equality follows from $\text{Tr}(\lambda M M^{\dagger}) = \text{Tr}(\lambda M^{\dagger} M)$.

FIFTH consider some unitary U which takes those $\{P_i\}_{i=1...p}$ into one-zero orthogonal diagonal matrices $I = \{\mathbf{I}_i\}_{i=1...p}$ with $\mathbf{I}_i = |i\rangle\langle i| \otimes \mathbb{I}_q$. Note that this is always possible since the $\{P_i\}_{i=1...p}$ form an orthogonal (ii), complete (iii) set of projectors of equal dimension by Eq. (4). We show that

$$\forall A \in U \mathcal{A} U^{\dagger} \,\forall i, j \quad [\mathbf{I}_{i} A \mathbf{I}_{j} = |i\rangle \langle j| \otimes A_{ij}]$$

$$\text{with } A_{ij} A_{ij}^{\dagger} = A_{ij}^{\dagger} A_{ij} \propto \mathbb{I}_{q}.$$

$$(5)$$

The first line stems from the form of I, i.e. $\mathbf{I}_i A \mathbf{I}_i$

$$= \sum_{pqkl} A_{pqkl} (|i\rangle \langle i| \otimes \mathbb{I}_q) (|p\rangle \langle q| \otimes |k\rangle \langle l|) (|j\rangle \langle j| \otimes \mathbb{I}_q)$$

$$= \sum_{kl} A_{ijkl} (|i\rangle \langle j| \otimes |k\rangle \langle l|) = |i\rangle \langle j| \otimes (\sum_{kl} A_{ijkl} |k\rangle \langle l|).$$

For the second line let $M = \mathbf{I}_i A \mathbf{I}_j$. By Eq. (3) there is a complex number λ such that we have both $\mathbf{I}_i = \lambda M M^{\dagger}$ and $\mathbf{I}_j = \lambda M^{\dagger} M$. But then

$$egin{aligned} &|i
angle\langle i|\otimes \mathbb{I}_q=\mathbf{I}_i=\ &\lambda MM^{\dagger}=\lambda(|i
angle\langle j|\otimes A_{ij})(|i
angle\langle j|\otimes A_{ij}^{\dagger})\ &=\lambda|i
angle\langle j|\otimes A_{ij}A_{ij}^{\dagger} \end{aligned}$$

and hence $\lambda A_{ij}A_{ij}^{\dagger} = \mathbb{I}_q$, and symmetrically for $\lambda A_{ij}^{\dagger}A_{ij} = \mathbb{I}_q$.

FINALLY consider some unitary $V = \sum_{i} |i\rangle \langle i| \otimes A_{1i}$, where $U^{\dagger}AU$ is some matrix verifying Eq. (2), and rescaled so that Eq. (5) makes A_{1j} it unitary. We show that

$$\forall M \in VU\mathcal{A}U^{\dagger}V^{\dagger} \forall i, j \quad [\mathbf{I}_{i}M\mathbf{I}_{j} = |i\rangle\langle j| \otimes \lambda \mathbb{I}_{q}]$$

with λ a complex number.

For a better understanding of V notice that

$$\begin{split} V &= \sum_{i} (|i\rangle \langle i| \otimes A_{1i}) \\ &= \sum_{i} (|i\rangle \langle 1| \otimes \mathbb{I}_{q}) (|1\rangle \langle i| \otimes A_{1i}) \\ &= \sum_{i} (|i\rangle \langle 1| \otimes \mathbb{I}_{q}) \mathbf{I}_{1} A \mathbf{I}_{i} \end{split}$$

Consider $B = V^{\dagger}MV$. It belongs to $U\mathcal{A}U^{\dagger}$ and by Eq. (5) it is of the form $B = \sum_{ij} |i\rangle\langle j| \otimes B_{ij}$. Now $\mathbf{I}_i M \mathbf{I}_j$

$$= \mathbf{I}_{i} V B V^{\dagger} \mathbf{I}_{j}$$

$$= (|i\rangle\langle 1| \otimes \mathbb{I}_{q}) \mathbf{I}_{1} A \mathbf{I}_{i} B \mathbf{I}_{j} A^{\dagger} \mathbf{I}_{1} (|1\rangle\langle j| \otimes \mathbb{I}_{q})$$

$$= (|i\rangle\langle 1| \otimes \mathbb{I}_{q}) \lambda \mathbf{I}_{1} (|1\rangle\langle j| \otimes \mathbb{I}_{q})$$

$$= \lambda (|i\rangle\langle 1| \otimes \mathbb{I}_{q}) (|1\rangle\langle 1| \otimes \mathbb{I}_{q}) (|1\rangle\langle j| \otimes \mathbb{I}_{q})$$

$$= \lambda |i\rangle\langle j| \otimes \mathbb{I}_{q}$$

Theorem 2 (Characterizing several subsystems)

Let \mathcal{A} and \mathcal{B} be commuting algebras of $M_n(\mathbb{C})$ such that $\mathcal{AB} = M_n(\mathbb{C})$. Then there exists a unitary matrix U such that, $U\mathcal{A}U^{\dagger}$ is $M_p(\mathbb{C}) \otimes \mathbb{I}_q$ and $U\mathcal{B}U^{\dagger}$ is $\mathbb{I}_p \otimes M_q(\mathbb{C})$, with pq = n.

Proof.

First, let us note that $\mathcal{C}_{\mathcal{A}}$ includes \mathbb{CI}_n . Next, the elements of $\mathcal{C}_{\mathcal{A}}$ commute by definition with all matrices in \mathcal{A} , but also with all matrices in \mathcal{B} , since \mathcal{A} and \mathcal{B} commute. Therefore, as $\mathcal{AB} = M_n(\mathbb{C})$, $\mathcal{C}_{\mathcal{A}}$ is equal to \mathbb{CI}_n . Thus, according to proposition 1, it is a subsystem algebra. For simplicity matters, and without loss of generality, we will assume that \mathcal{A} is actually equal to $M_p(C) \otimes \mathbb{I}_q$ for some p and q such that pq = n. Now for the same reasons \mathcal{B} is also a subsystem algebra. Because it commutes with \mathcal{A} it must act on a disjoint subsystem as \mathcal{A} . And since together they generate $M_n(\mathbb{C})$, there is no other choice but to have \mathcal{B} actually equal to $\mathbb{I}_p \otimes M_q(\mathbb{C})$. \Box

Definition 4 (Restriction Algebras)

Consider \mathcal{A} an algebra of $M_p(\mathbb{C}) \otimes M_q(\mathbb{C}) \otimes M_r(\mathbb{C})$. For A an element of \mathcal{A} , we write $A|_1$ for the matrix $Tr_{02}(A)$. Similarly so we call $\mathcal{A}|_1$ the restriction of \mathcal{A} to the middle subsystem, i.e. the algebra generated by the matrices of the set $\{Tr_{02}(A) | A \in \mathcal{A}\}.$

Lemma 1 (Restriction of commuting algebras)

Consider \mathcal{A} an algebra of $M_p(\mathbb{C}) \otimes M_q(\mathbb{C}) \otimes \mathbb{I}_r$ and \mathcal{B} an algebra of $\mathbb{I}_p \otimes M_q(\mathbb{C}) \otimes M_r(\mathbb{C})$. Say \mathcal{A} and \mathcal{B} commute. Then so do $\mathcal{A}|_1$ and $\mathcal{B}|_1$.

Proof.

In the particular case where \mathcal{A} and \mathcal{B} have only subsystem 1 in common we have

$$\forall A \in \mathcal{A}, B \in \mathcal{B} \quad pr.\mathrm{Tr}_{02}(AB) = \mathrm{Tr}_{02}(A)\mathrm{Tr}_{02}(B). \quad (6)$$

Indeed take $A = \sum_{i} \alpha_i . (\sigma_i \otimes \tau_i \otimes \mathbb{I})$ and $B = \sum_{j} \beta_j . (\mathbb{I} \otimes \mu_j \otimes \nu_j)$. We have

$$pr.\operatorname{Tr}_{02}(AB) = \operatorname{Tr}_{02}(\sum_{ij} pr\alpha_i\beta_j.(\sigma_i \otimes \tau_i\mu_j \otimes \nu_j))$$
$$= (\sum_i p\alpha_i.\operatorname{Tr}(\sigma_i).\tau_i)(\sum_j r\beta_j.\operatorname{Tr}(\nu_j).\mu_j)$$
$$= \operatorname{Tr}_{02}(A)\operatorname{Tr}_{02}(B).$$

Now $\mathcal{A}|_1$ is generated by $\{\operatorname{Tr}_{02}(A) | A \in \mathcal{A}\}$, and $\mathcal{B}|_1$ is generated by $\{\operatorname{Tr}_{02}(B) | B \in \mathcal{B}\}$. Since commutation is preserved by $*, +, \alpha$. and \dagger all we need to check is that the generating elements commute. Consider $A|_1$ an element of $\mathcal{A}|_1$ and take A such that $A|_1 = \operatorname{Tr}_{02}(A)$. Similarly take $B|_1$ and B such that $B|_1 = \operatorname{Tr}_{02}(B)$. We have $A|_1B|_1 = \operatorname{Tr}_{02}(A)\operatorname{Tr}_{02}(B) = pr.\operatorname{Tr}_{02}(AB) =$ $pr.\operatorname{Tr}_{02}(BA) = \operatorname{Tr}_{02}(B)\operatorname{Tr}_{02}(A) = B|_1A|_1$. \Box

Lemma 2 (Restriction of generating algebras)

Consider \mathcal{A} an algebra of $M_p(\mathbb{C}) \otimes M_q(\mathbb{C}) \otimes \mathbb{I}_r$ and \mathcal{B} an algebra of $\mathbb{I}_p \otimes M_q(\mathbb{C}) \otimes M_r(\mathbb{C})$. Say $\mathcal{AB}|_1 = M_p(\mathbb{C})$. Hence we have that $\mathcal{A}|_1\mathcal{B}|_1 = M_p(\mathbb{C})$.

Proof.

 $\mathcal{AB}|_1$ is generated by $\{\operatorname{Tr}_{02}(AB) | A \in \mathcal{A}, B \in \mathcal{B}\}$. However by Eq. (6) this is the same as $\{\operatorname{Tr}_{02}(A)\operatorname{Tr}_{02}(B) | A \in \mathcal{A}, B \in \mathcal{B}\}$, which generates $\mathcal{A}|_1\mathcal{B}|_1$. \Box

Lemma 3 (Duality)

Let \mathcal{H}_0 and \mathcal{H}_1 be Hilbert spaces, with \mathcal{H}_0 of dimension p. Let A, ρ, ρ' denote some elements of $\mathcal{L}(\mathcal{H}_0 \otimes \mathcal{H}_1)$ with ρ, ρ' having partial traces $\rho|_0, \rho'|_0$ over \mathcal{H}_0 . We then have that $A \in M_p(\mathbb{C}) \otimes \mathbb{I}$ is equivalent to

$$\forall \rho, \rho' \quad [\rho|_0 = \rho'|_0 \Rightarrow Tr(A\rho) = Tr(A\rho')].$$

Moreover we have that $\rho|_0 = \rho'|_0$ is equivalent to

$$\forall A \in M_p(\mathbb{C}) \otimes \mathbb{I} \quad [Tr(A\rho) = Tr(A\rho')].$$

Proof.

Physically the first part of the lemma says that "a measurement is local if and only if it depends only upon the reduced density matrices".

[⇒]. Suppose that $A = A_0 \otimes \mathbb{I}_1$. In this case we have $\operatorname{Tr}(A\rho) = \operatorname{Tr}(A_0\rho|_0)$. Assuming $\rho|_0 = \rho'|_0$ yields $\operatorname{Tr}(A\rho) = \operatorname{Tr}(A_0\rho|_0) = \operatorname{Tr}(A_0\rho'|_0) = \operatorname{Tr}(A\rho')$.

$$[\Leftarrow]. \text{ Let's write } A = \sum_{i,j} |i\rangle \langle j| \otimes B_{ij}, \text{ with } |i\rangle \text{ and } |j\rangle$$

ranging over some unitary basis of \mathcal{H}_0 . If A is not of the form $A_0 \otimes \mathrm{Id}_1$, then for some i and j, B_{ij} is not a multiple of the identity. Then there exist unit vectors $|x\rangle$ and $|y\rangle$ of \mathcal{H}_1 such that $\langle x|B_{ij}|x\rangle \neq \langle y|B_{ij}|y\rangle$. In other words, $\mathrm{Tr}(B_{ij}|x\rangle\langle x|) \neq \mathrm{Tr}(B_{ij}|y\rangle\langle y|)$. If we now consider $\rho = |j\rangle\langle i| \otimes |x\rangle\langle x|$ and $\rho' = |j\rangle\langle i| \otimes |y\rangle\langle y|$, we get what we wanted, i.e. $\rho|_0 = \rho'|_0$ but $\mathrm{Tr}(A\rho) \neq \mathrm{Tr}(A\rho')$.

Physically the second part of the lemma says that "two reduced density matrices are the same if and only if their density matrices cannot be distinguished by a local measurement".

 $[\Rightarrow]$. This ' $[\Rightarrow]$ ' is actually exactly the same as the first one, so we have already proved it.

 $\begin{array}{ll} [\Leftarrow]. & \text{Supposing Tr} \left(A\rho\right) = \text{Tr} \left(A\rho'\right) \text{ for } A = |j\rangle\langle i| \otimes \mathbb{I} \\ \text{yields } \rho|_{0ij} = \text{Tr} \left(|j\rangle\langle i|\rho|_0\right) = \text{Tr} \left(A\rho\right) = \text{Tr} \left(A\rho'\right) = \\ \text{Tr} \left(|j\rangle\langle i|\rho'|_0\right) = \rho'|_{0ij}. \end{array} \end{array}$ Because we can do this for all ij we have $\rho|_0 = \rho'|_0.$

II. BLOCK STRUCTURE

We will now introduce the basic definitions of onedimensional QCA.

In what follows Σ will be a fixed finite set of symbols (i.e. 'the alphabet', describing the possible basic states each cell may take) and q is a symbol such that $q \notin \Sigma$, which will be known as 'the quiescent symbol', which represents an empty cells. We write $q\Sigma = \{q\} \cup \Sigma$ for short.

Definition 5 (finite configurations)

A (finite) configuration c of over $q\Sigma$ is a function c : $\mathbb{Z} \longrightarrow q\Sigma$, with $i \longmapsto c(i) = c_i$, such that there exists a (possibly empty) interval I verifying $i \in I \Rightarrow c_i \in q\Sigma$ and $i \notin I \Rightarrow c_i = q$. The set of all finite configurations over $\{q\} \cup \Sigma$ will be denoted C_f .

Whilst configurations hold the basic states of an entire line of cells, and hence denote the possible basic states of the entire QCA, the global state of a QCA may well turn out to be a superposition of these. The following definition works because C_f is a countably infinite set.

Definition 6 (superpositions of configurations)

Let \mathcal{H}_{C_f} be the Hilbert space of configurations, defined as follows. To each finite configuration c is associated a unit vector $|c\rangle$, such that the family $(|c\rangle)_{c\in C_f}$ is an orthonormal basis of \mathcal{H}_{C_f} . A superposition of configurations is then a unit vector in \mathcal{H}_{C_f} .

Definition 7 (Unitarity)

A linear operator $G : \mathcal{H}_{\mathcal{C}_f} \longrightarrow \mathcal{H}_{\mathcal{C}_f}$ is unitary if and only if $\{G|c\rangle | c \in \mathcal{C}_f\}$ is an orthonormal basis of $\mathcal{H}_{\mathcal{C}_f}$.

Definition 8 (Shift-invariance)

Consider the shift operation which takes configuration $c = \ldots c_{i-1}c_ic_{i+1}\ldots$ to $c' = \ldots c'_{i-1}c'_ic'_{i+1}\ldots$ where for all $i c'_i = c_{i+1}$. Let $\sigma : \mathcal{H}_{\mathcal{C}_f} \longrightarrow \mathcal{H}_{\mathcal{C}_f}$ be its linear extension to superpositions of configurations. A linear operator $G : \mathcal{H}_{\mathcal{C}_f} \longrightarrow \mathcal{H}_{\mathcal{C}_f}$ is said to be shift invariant if and only if $G\sigma = \sigma G$.

Definition 9 (Locality)

A linear operator $G : \mathcal{H}_{\mathcal{C}_f} \longrightarrow \mathcal{H}_{\mathcal{C}_f}$ is said to be local with radius $\frac{1}{2}$ if and only if for any ρ, ρ' two states over $\mathcal{H}_{\mathcal{C}_f}$, and for any $i \in \mathbb{Z}$, we have

$$\rho|_{i,i+1} = \rho'|_{i,i+1} \quad \Rightarrow G\rho G^{\dagger}|_i = G\rho' G^{\dagger}|_i. \tag{7}$$

In the classical case, the definition would be that the letter to be read in some given cell i at time t+1 depends only the state of the cells i and i+1 at time t. This seemingly restrictive definition of locality is known in the classical case as a $\frac{1}{2}$ -neighborhood cellular automaton. This is because the most natural way to represent such

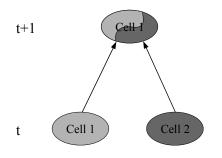


FIG. 1: A $\frac{1}{2}$ -neighborhood CA

an automaton is to shift the cells by $\frac{1}{2}$ at each step, so that the state of a cell depends on the state of the two cells under it, as shown in figure 1. This definition of locality is actually not so restrictive, since by grouping cells into 'supercells' one can construct a $\frac{1}{2}$ -neighborhood CA simulating the first one. The same thing can easily be done for QCA, so that this definition of locality is essentially done without loss of generality. Transposed to a quantum setting, we get the above definition: to know the state of cell number *i*, we only need to know the state of cells *i* and *i* + 1 before the evolution.

We are now set to give the formal definition of onedimensional quantum cellular automata.

Definition 10 (QCA)

A one-dimensional quantum cellular automaton (QCA) is an operator $G : \mathcal{H}_{\mathcal{C}_f} \longrightarrow \mathcal{H}_{\mathcal{C}_f}$ which is unitary, shift-invariant and local.

The next theorem provides us with another characterization of locality, more helpful in the proofs. But more importantly it entails structural reversibility, i.e. the fact that the inverse function of a QCA is also a QCA. Actually this theorem works for *n*-dimensional QCA as well as one.

Theorem 3 (Structural reversibility)

Let G be a unitary operator of \mathcal{H}_{C_f} and \mathcal{N} a finite subset of \mathbb{Z} . The two properties are equivalent:

- (i) For every states ρ and ρ' over the finite configurations, if $\rho|_{\mathcal{N}} = \rho'|_{\mathcal{N}}$ then $(G\rho G^{\dagger})|_0 = (G\rho' G^{\dagger})|_0$.
- (ii) For every operator A localized on cell 0, then $G^{\dagger}AG$ is localized on the cells in \mathcal{N} .
- (iii) For every states ρ and ρ' over the finite configurations, if $\rho|_{-\mathcal{N}} = \rho'|_{-\mathcal{N}}$ then $(G^{\dagger}\rho G)|_{0} = (G^{\dagger}\rho' G)|_{0}.$
- (iv) For every operator A localized on cell 0, then GAG^{\dagger} is localized on the cells in $-\mathcal{N}$.

When G satisfies these properties, we say that G is local at 0 with neighbourhood \mathcal{N} .

Proof.

 $[(i) \Rightarrow (ii)]$. Suppose (i) and let A be an operator acting on cell 0. For every states ρ and ρ' such that $\rho|_{\mathcal{N}} = \rho'|_{\mathcal{N}}$, we have Tr $(AG\rho G^{\dagger}) =$ Tr $(AG\rho' G^{\dagger})$, using lemma 3 and our hypothesis that $(G\rho G^{\dagger})|_0 = (G\rho' G^{\dagger})|_0$. We thus get Tr $(G^{\dagger}AG\rho) =$ Tr $(G^{\dagger}AG\rho')$. Since this is true of every ρ and ρ' such that $\rho|_{\mathcal{N}} = \rho'|_{\mathcal{N}}$, this means, again according to lemma 3, that $G^{\dagger}AG$ is localized on the cells in \mathcal{N} . $[(ii) \Rightarrow (i)]$. Suppose (ii) and $\rho|_{\mathcal{N}} = \rho'|_{\mathcal{N}}$. Then, for every operator B localized on the cells in \mathcal{N} , lemma 3 gives Tr $(B\rho) =$ Tr $(B\rho')$, so for every operator A localized on cell 0, we get:

$$\begin{aligned} \operatorname{Tr} \left(A G \rho G^{\dagger} \right) &= \operatorname{Tr} \left(G^{\dagger} A G \rho \right) \\ &= \operatorname{Tr} \left(G^{\dagger} A G \rho' \right) \\ &= \operatorname{Tr} \left(A G \rho' G^{\dagger} \right) \end{aligned}$$

Again by lemma 3, this means $(G\rho G^{\dagger})|_{0} = (G\rho' G^{\dagger})|_{0}$. $[(ii) \Rightarrow (iv)]$. Suppose (ii) and let A be an operator acting on cell 0. Consider some operator M acting on a cell i which does not belong to $-\mathcal{N}$. According to our hypothesis we know that $G^{\dagger}MG$ does not act upon cell 0, and hence it commutes with A. But $AB \mapsto GAG^{\dagger}GBG^{\dagger} = GABG^{\dagger}$ is a morphism, hence $GG^{\dagger}MGG^{\dagger} = M$ also commutes with GAG^{\dagger} . Because M can be chosen amongst to full matrix algebra $M_d(\mathbb{C})$ of cell i, this entails that GAG^{\dagger} must be the identity upon this cell. The same can be said of any cell outside $-\mathcal{N}$.

 $[(iv) \Rightarrow (ii)], [(iii) \Rightarrow (iv)], [(iii) \Leftarrow (iv)]$ are symmetrical to $[(ii) \Rightarrow (iv)], [(i) \Rightarrow (ii)], [(ii) \Leftarrow (i)]$ just by interchanging the roles of G and G^{\dagger} . Now this is done we proceed to prove the structure theorem for QCA over finite, unbounded configurations. This is a simplification of [9]. The basic idea of the proof is that in a cell at time t we can separate what information will be sent to the left at time t+1 and which information will be sent to the right at time t+1. But first of all we shall need two lemmas. These are better understood by referring to Figure 2.

Lemma 4 Let \mathcal{A} be the image of the algebra of the cell 1 under the global evolution G. It is localized upon cells 0 and 1, and we call $\mathcal{A}|_1$ the restriction of \mathcal{A} to cell 1.

Let \mathcal{B} be the image of the algebra of the cell 2 under the global evolution G. It is localized upon cells 1 and 2, and we call $\mathcal{B}|_1$ the restriction of \mathcal{B} to cell 1.

There exists a unitary U acting upon cell 1 such that $U\mathcal{A}|_0 U^{\dagger}$ is of the form $M_p(\mathbb{C}) \otimes \mathbb{I}_q$ and $U\mathcal{B}|_1 U^{\dagger}$ is of the form $\mathbb{I}_p \otimes M_q(\mathbb{C})$, with pq = d.

Proof.

 \mathcal{A} and \mathcal{B} are indeed localized as stated due to the locality of G and a straightforward application of lemma 3 with $\mathcal{N} = \{0, 1\}$, which we can apply at position 1 and 2 by shift-invariance.

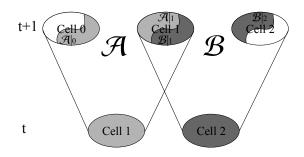


FIG. 2: Definitions of the algebras for the proof of the structure theorem.

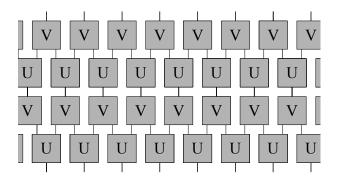


FIG. 3: QCA with two-layered block representation (U, V). Each line represents a cell, which is a quantum system. Each square represents a unitary U/V which gets applied upon the quantum systems. Time flows upwards.

 \mathcal{A} and \mathcal{B} commute because they are the image of two commuting algebras, those of Cell 1 and 2, via a morphism $AB \mapsto GAG^{\dagger}GBG = GABG^{\dagger}$.

Moreover by lemma 3 the antecedents of the operators localized in cell 1 are all localized in cells 1 and 2. Plus they all have an antecedent because G is surjective. Hence $\mathcal{AB}|_1$ is the entire cell algebra of cell 1, i.e. $M_d(\mathbb{C})$.

So now we can apply Theorem 2 and the result follows. \square

Lemma 5

Let \mathcal{B} be the image of the algebra of the cell 2 under the global evolution G. It is localized upon cells 1 and 2, and we call $\mathcal{B}|_1$ the restriction of \mathcal{B} to cell 1 and $\mathcal{B}|_2$ the restriction of \mathcal{B} to cell 2. We have that $\mathcal{B} = \mathcal{B}|_1 \otimes \mathcal{B}|_2$.

Proof.

We know that \mathcal{B} is isometric to $M_d(\mathbb{C})$ and we know that $\mathcal{B}|_1 \otimes \mathcal{B}|_2 \subset \mathcal{B}$. But then by the previous lemma applied upon cell 1 we also know that $\mathcal{B}|_1$ is isometric to $M_q(\mathbb{C})$ and if we apply it to cell 2 then we have that $\mathcal{B}|_2$ is isometric to $M_p(\mathbb{C})$. Hence the inclusion is an equality. \Box

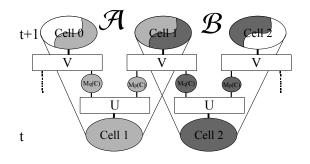


FIG. 4: Zooming into the two-layered block representation. The unitary interactions U and V are alternated repeatedly as shown.

Theorem 4 (Structure theorem)

Any QCA G is of the form described by Figures 3 and 4.

Proof.

Let \mathcal{A} and \mathcal{B} be respectively the images of the algebra of the cells 1 and 2 under the global evolution G. By virtue of lemma 4 we know that \mathcal{A} and \mathcal{B} are respectively isometric to $M_p(\mathbb{C}) \otimes \operatorname{Id}_q$ and $\operatorname{Id}_p \otimes M_q(\mathbb{C})$; let V^{\dagger} be the unitary transformation over \mathbb{C}^d which accomplishes this separation. From lemma 5, we know that $(V^{\dagger} \otimes V^{\dagger})G$ maps the algebra of one cell into $\operatorname{Id}_p \otimes M_q(\mathbb{C}) \otimes M_p(\mathbb{C}) \otimes \operatorname{Id}_q$, so we can choose a unitary operator U over \mathbb{C}^d which realizes this mapping by conjugation. By shift-invariance, the same V and U will do for every position in the line. Therefore $G = (\bigotimes V) (\bigotimes U)$ as in Fig. 4. The rest of the proof serves only to give a formal meaning to these infinite tensor products as unitary operators over $\mathcal{H}_{\mathcal{C}_f}$.

Let us consider $|q\rangle\langle q| \in M_d(\mathbb{C})$. Its image by V^{\dagger} , i.e. $V|q\rangle\langle q|V^{\dagger}$, is some one-dimensional projector in $M_p(\mathbb{C}) \otimes M_q(\mathbb{C})$. Now consider the state corresponding to the quiescent state on every cells. It is invariant by G, so this $V|q\rangle\langle q|V^{\dagger}$ has to be separable in $M_p(\mathbb{C}) \otimes M_q(\mathbb{C})$, because after applying independent U transformations on each side we get the everywhere quiescent state, which is unentangled. This means that $V|q\rangle\langle q|V^{\dagger}$ can be written as $|q_1\rangle\langle q_1| \otimes |q_2\rangle\langle q_2|$, where $|q_1\rangle$ and $|q_2\rangle$ are respectively unit vectors of \mathbb{C}^p and \mathbb{C}^q . So we can assume that V maps $|q_1\rangle|q_2\rangle$ to $|q\rangle$. Moreover we know $U^{\dagger}(|q_2\rangle\langle q_2| \otimes |q_1\rangle\langle q_1|) U$ must be equal to $|q\rangle\langle q|$, so we can assume that U maps $|q\rangle$ to $|q_2\rangle|q_1\rangle$.

We can now give a meaning to the infinite product of unitary operators. For each n we consider the operator $\left(\bigotimes_{[-n,n]}U\right)$ where the U's are only applied on the portion [-n,n] of the line. The action of $\left(\bigotimes U\right)$ is simply the limit of its images by $\left(\bigotimes_{[-n,n]}U\right)$, when n goes to infinity. That U maps $|q\rangle$ to $|q_2\rangle|q_1\rangle$ insures that this limit does exist. Indeed, for every finite configuration c, the sequence $\left(\bigotimes_{[-n,n]}U\right)|c\rangle$ will be ultimately constant, due to the quiescent boundaries. \Box Note that this structure could be further simplified if we

III. QUANTIZATIONS AND CONSEQUENCES

The structure theorem for QCA departs in several important ways from the classical situation, giving rise to a number of apparent paradoxes. We begin this section by discussing some of these concerns in turns. Each of them is introduced via an example, which we then use to derive further consequences or draw the limits of the structure theorem.

Bijective CA and superluminal signalling.

First of all, it is a well-known fact that not all bijective CA are structurally reversible. The XOR CA is a standard example of that.

Definition 11 (XOR CA)

were to allow ancillary cells [1].

Let C_f be the set of finite configurations over the alphabet $q\Sigma = \{q, 0, 1\}$. For all x, y in $q\Sigma$ Let $\delta(qx) = x$, $\delta(xq) = q$, and $\delta(xy) = x \oplus y$ otherwise. We call $F : C_f \longrightarrow C_f$ the function mapping $c = \ldots c_{i-1}c_ic_{i+1}\ldots$ to $c' = \ldots \delta(c_{i-1}c_i)\delta(c_ic_{i+1})\ldots$

The XOR CA is clearly shift-invariant, and local in the sense that the state of a cell at t + 1 only depends from its state and that of its right neighbour at t. It is also bijective. Indeed for any $c' = \ldots qqc'_kc'_{k+1}\ldots$ with c'_k the first non quiescent cell, we have $c_k = q$, $c_{k+1} = c'_k$, and thereon for $l \ge k + 1$ we have either $c_{l+1} = c_l \oplus c'_l$ if $c'_l \ne q$, or once again $c_{l+1} = q$ otherwise, etc. In other words the antecedent always exists (surjectivity) and is uniquely derived (injectivity) from left till right. But the XOR CA is not structurally reversible. Indeed for some $c' = \ldots 000000000\ldots$ we cannot know whether the antecedent of this large zone of zeroes is another large zone of zeroes or a large zone of ones – unless we deduce this from the left border as was previously described...but the left border may lie arbitrary far.

So classically there are bijective CA whose inverse is not a CA, and thus who do not admit any *n*-layered block representation at all. Yet surely, just by defining F over \mathcal{H}_{C_f} by linear extension (e.g. $F(\alpha.|\ldots01\ldots) + \beta.|\ldots11\ldots\rangle) = \alpha.F|\ldots01\ldots\rangle + \beta.F|\ldots11\ldots\rangle$) we ought to have a QCA, together with its block representation, hence the apparent paradox.

In order to lift this concern let us look at the properties of this quantized $F : \mathcal{H}_{\mathcal{C}_f} \longrightarrow \mathcal{H}_{\mathcal{C}_f}$. It is indeed unitary as a linear extension of a bijective function, and it is shift-invariant for the same reason. Yet counterintuitively it is non-local. Indeed consider configurations $c_{\pm} = 1/\sqrt{2}.|\ldots qq\rangle(|00\ldots 00\rangle \pm |11\ldots 11\rangle)|qq\ldots\rangle$. We have $Fc_{\pm} = |\ldots qq00\ldots 0\rangle|\pm\rangle|qq\ldots\rangle$, where we have used the usual notation $|\pm\rangle = 1/\sqrt{2}.(|0\rangle \pm |1\rangle)$. Let *i* be the position of this last non quiescent cell. Clearly $(Fc_{\pm})|_i = |\pm\rangle\langle\pm|$ is not just a function of $c|_{i,i+1} =$ $(|0q\rangle\langle 0q| + |1q\rangle\langle 1q|)/2$, but instead depends upon this t the que

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global \pm phase. Another way to put it is that the quantized XOR may be used to transmit information faster than light. Say the first non quiescent cell is with Alice in Paris and the last non quiescent cell is with Bob in New York. Just by applying a phase gate Z upon her cell Alice can change c_+ into c_- at time t, leading to a perfectly measurable change from $|+\rangle$ to $|-\rangle$ for Bob. Again another way to say it is that operators localized upon cell 1 are not taken to operators localized upon cells 0 and 1, as was the case for QCA. For instance take $\mathbb{I} \otimes Z \otimes \mathbb{I} |calized upon cell 1$. This is taken to $F(\mathbb{I} \otimes Z \otimes \mathbb{I})F^{\dagger}$. But this operation is not localized upon cells 0 and 1, as it takes $|\ldots qq00\ldots 0\rangle|+\rangle|qq\ldots\rangle$ to $|\ldots qq00\ldots 0\rangle|-\rangle|qq\ldots\rangle$, whatever the position i of the varying $|\pm\rangle$.

Now let us take a step back. If a CA is not structurally reversible, there is no chance that its QCA will be. Moreover according the current state of modern physics, quantum mechanics is the theory for describing all closed systems. Therefore we reach the following proposition.

Proposition 1 (Class B is not locally quantizable) The class of bijective but not structurally reversible CAupon finite configurations is known to coincide with the class of surjective but non injective CA upon infinite configurations, or again the class of bijective CA upon finite configurations but not upon infinite configuration. Call it B.

The quantization of a class B automata is not local. It cannot be implemented by a series of finite closed quantum systems.

As far as CA are concerned this result removes much of the motivation of several papers which focus upon class B. As regards QCA the structure theorem removes much of the motivation of the papers [1, 4, 5], which contain unitary decision procedures for possibly non-structurally reversible QCA.

Faster quantum signalling.

Second, it is a well-known fact that there exists some radius 1/2, structurally reversible CA, whose inverse is also of radius 1/2, and yet which do not admit a two-layered block representation unless the cells are grouped into supercells. The Toffoli CA is a good example of that.

Definition 12 (Toffoli CA) Let C_f be the set of finite configurations over the alphabet $\{00, 01, 01, 11\}$, with 00 now taken as the quiescent symbol. For all ab and cd taken in the alphabet let $\delta(abcd) = (b \oplus a.c)c$. We call $F: C_f \longrightarrow C_f$ the function mapping $c = \ldots c_{i-1}c_ic_{i+1}\ldots$ to $c' = \ldots \delta(c_{i-1}c_i)\delta(c_ic_{i+1})\ldots$ This is best described by Figure 5.

The Toffoli CA is clearly shift-invariant, and of radius 1/2. Let us check that its inverse is also of radius 1/2. For instance say we seek to retrieve (c,d). c is easy of course. By shift-invariance retrieving d is like retrieving b. But since we have a and d in cleartext we can easily

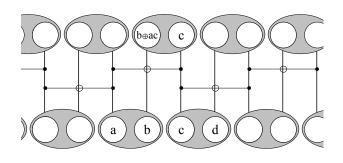


FIG. 5: The Toffoli CA.

substract a.d from $b \oplus a.d$. Now why does it not have a two-layered block representation without cell grouping? Remember the toffoli gate is the controlled-controlled-NOT gate. Here b is NOTed depending upon a and c, which pass through unchanged, same for d with the left and right neighbouring subcells, etc. So actually the Toffoli CA is just two layers of the toffoli gate, as we have shown in Figure 5. But we know that the toffoli gate cannot be obtained from two bit gates in classical reversible electronics, hence there cannot be a two-layered block representation without cell grouping.

So classically there exists some structurally reversible CA, of radius 1/2, whose inverse is also of radius 1/2, but do not admit a two-layered block representation without cell grouping. Yet surely, just by defining F over $\mathcal{H}_{\mathcal{C}_f}$ by linear extension we ought to have a QCA, together with its block representation, and that construction does not need any cell grouping, hence again the apparent paradox.

Again in order to lift this concern let us look at the properties of this quantized $F : \mathcal{H}_{\mathcal{C}_f} \longrightarrow \mathcal{H}_{\mathcal{C}_f}$. It is indeed unitary and shift-invariant of course. This time it is also local, but counter-intuitively it turns out not to be of radius 1/2. Indeed from the formulation in terms of Toffoli gates as in Figure 5 one can show that the radius is 3/2 in a quantum mechanical setting. For instance one can check that putting $|+\rangle$ in the *a*-subcell, $|-\rangle$ in the *b*-subcell, and either $|0\rangle$ or $|1\rangle$ in the *c*-subcell of Fig. 5 at time *t* will yield either $|+\rangle$ or $|-\rangle$ in the *a*-subcell at time t + 1.

Once more let us take a step back. The Toffoli CA is yet another case where exploiting quantum superpositions of configurations enables us to have information flowing faster than in the classical setting, just like for the XOR CA. But unlike the XOR CA, the speed of information remains bounded in the Toffoli CA, and so up to cell grouping it can still be considered a QCA. Therefore we reach the following proposition.

Proposition 2 (Quantum information flows faster) Let $F : C_f \longrightarrow C_f$ be a CA and $F : \mathcal{H}_{C_f} \longrightarrow \mathcal{H}_{C_f}$ the corresponding QCA, as obtained by linear extension of F. Information may flow faster in the the quantized

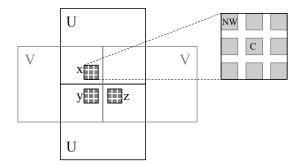


FIG. 6: The Kari CA and the choice of x, y, z cells.

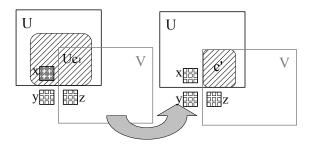


FIG. 7: Uc_1 and c'.

version of F.

This result is certainly intriguing, and one may wonder whether it might contain the seed of a novel development quantum information theory, as opposed to its classical counterpart.

No-go for n-dimensions.

Finally, it is again well-known that in two-dimensions there exists some structurally reversible CA which do not admit a two-layered block representation, even after a cell-grouping. The standard example is that of Kari [8]:

Definition 13 (Kari CA) Let C_f be the set of finite configurations over the alphabet $\{0,1\}^9$, with 0^9 is now taken as the quiescent symbol. So each cell is made of 8 bits, one for each cardinal direction (North, North-East...) plus one bit in the center, as in Figure 6. At each time step, the North bit of a cell undergoes a NOT only if the cell lying North has center bit equal to 1, the North-East bit of a cell undergoes a NOT only if the cell lying North-East has center bit equal to 1, and so on. Call F this CA.

This is clearly shift-invariant, local and structurally reversible. Informally the proof that the Kari CA does not admit a a two-layered block representation, even if we group cells into supercells, runs as follows [8]:

- Suppose F admits a decomposition into U and V blocks;

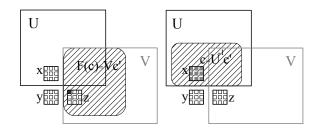


FIG. 8: F(c) and c: a contradiction.

- Consider cells x, y, z such that they are all neighbours; x, y are in the same V-block but not z; x and y, z are not in the same U-block, as in Figure 6;

- Consider c_1 the configuration with 1 as the center bit of x and zero everywhere. Run the *U*-blocks, the hatched zone left of Figure 7 represents those cells which may not be all zeros;

- Consider the configuration obtained from the previous Uc_1 by putting to zero anything which lies not the V-block of z, as in Figure 7. Call it c', and let c be defined as the antecedent of c under a run of the U-blocks;

- On the one hand we know that F(c) is obtained from a run of the V-blocks from c'. In the left V-block there are just zeroes so F(c) has only zeroes, in particular the North bit of the y cell is 0. But the right V-block is ex-

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actly the same as that of Uc_1 , and so we know that the North-West bit of the y cell of F(c) is 1, as in the left of Figure 8:

- On the other hand since c' has zeroes everywhere but in the above U-block, the same is true of c, as shown right of the Figure 8. A consequence of this is that the North bit of the y cell must be equal to the North-West bit of the z cells, as they are both obtained from the same function ofs the center bit of the x cell;

- Hence the contradiction.

Now by defining F over $\mathcal{H}_{\mathcal{C}_f}$ by linear extension we have a QCA, and the proof applies in a very similar fashion, the contradiction being that the state of the North qubit of the y cell must be equal to the state of the North-West bit of the z cells, whether these are mixed states or not. Hence we have a counterexample to the higherdimensional case of the Theorem in [9]. We reach the following proposition.

Proposition 3 (No-go for *n***-dimensions)** There exists some 2-dimensional QCA which do not admit a twolayered block representation.

Understanding the structure of the *n*-dimensional QCA is clearly the main challenge that remains ahead of us. Acknowledgements

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