

A Query Efficient Non-Adaptive Long Code Test with Perfect Completeness

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Abstract

Long Code testing is a fundamental problem in the area of property testing and hardness of approximation. In this paper, we study how small the soundness s of the Long Code test with perfect completeness can be by using non-adaptive q queries. We show that $s = (2q+3)/2^q$ is possible, where q is of the form $2^k - 1$ for any integer $k > 2$.

Key words: Long Code test, dictatorship test, influence of variables, amortized query complexity

1 Introduction

Testing basic properties of Boolean functions is a well studied subject in the area of property testing, complexity theory and learning theory [19, 40]. Examples of such properties include: Linear functions, singleton functions, juntas, low-degree polynomials, and several other concise representations [12, 8, 2, 38, 3, 17, 10, 11]. In property test setting, our goal is, given oracle access to Boolean function(s) of the form $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$, to decide whether the functions satisfy a certain property, or are *far* from the property, where the definition of *farness* varies according to the property of interest and intended applications of the test. There are several parameters that characterize testers: *Completeness*, denoted by c , is the lower bound on the probability that the tester accepts functions which satisfy the property. *Soundness*, denoted by s , is the upper bound on the probability that the tester accepts functions that are far from the property. *Query complexity*, denoted by q , is the total number of oracle accesses in one execution of the tester. An *adaptive* tester means that the tester can make queries that depend on the answers of previous queries. Otherwise the tester is said to be *non-adaptive*.

The property of our interest in this paper is the Long Code. *Long Code* is a function of the form $f(x) = x_i$ for some index i . In the multi-function

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An extended abstract of this paper has appeared in Proceedings of the 14th International Workshop on Randomization and Computation (RANDOM 2010), pp. 738–751.

Long Code testing problem, the task is, given oracle access to a collection of Boolean functions $\{f_1, f_2, \dots, f_t\}$, to distinguish the following two cases: (i) there exists an index i such that $f_1(x) = f_2(x) = \dots = f_t(x) = x_i$, or (ii) for every index i and every $f_j \neq f_k$, either f_j or f_k is far from $f(x) = x_i$. The notion of Long Code was introduced by Bellare et al. [9] in the context of Probabilistically Checkable Proofs (PCPs) (also studied by Parnas et al. [38] in name of “dictatorship function”). We quickly review the use of (multi-function) Long Code test in PCPs and hardness of approximation. From now on, when we say just PCP, it means a PCP for some NP-complete language. Recall that most efficient PCP constructions today such as the celebrated work of Håstad [23] follow the paradigm of composing “outer verifier” and “inner verifier.”

An outer verifier is a certain kind of two prover one round games (or equivalently label cover problems); e.g., the combination of the PCP theorem [6, 4] and Raz’s parallel repetition theorem [39] (we often call the combination Raz verifier), or Unique Games or d -to-1 Games [29]. An inner verifier is essentially a multi-function Long Code tester. By composing outer and inner verifiers, we obtain a PCP, which often has the same parameters of the underlying Long Code tester. Therefore improving the quality of Long Code testers yields better constructions of PCPs. Here we note that criterion for soundness used in a Long Code tester/inner verifier may restrict the type of an outer verifier to ensure that they can be composed. We elaborate on this issue later.

The use of PCP constructions in hardness of approximation is quite standard. The acceptance criteria of Long Code testers/PCPs correspond to specific optimization problems and the ratio of completeness and soundness s/c gives the inapproximability ratio of the corresponding problems. For more details on the relation between Long Code tests, PCPs and hardness of approximation, see [30] for the general framework and [36] for constraint satisfaction problems (CSPs).

With the connection between Long Code testers and PCPs in mind, we sometimes compare the parameters of them at the same time without distinction. There are several interesting trade-offs between parameters as follows:

Criteria for soundness The definition of fairness is quite an important issue in the composition paradigm. Several possible criteria for rejection are; distance from any Long Codes or any juntas, quasi-randomness, and low-degree influences. Which criterion is used in Long Code test poses some restriction on an outer verifier, e.g., if we use quasi-randomness, the tester can be composed with Raz verifier, but if we use low-degree influences, it is hard to combine the tester with Raz verifier though it may be composed with Unique Games. It may also affect the lowest possible value of the soundness with other parameters fixed.

Perfect completeness vs. almost perfect completeness It is known that the soundness of PCPs with perfect completeness ($c = 1$) can be worse than that of PCPs with almost perfect completeness ($c = 1 - \epsilon$). Håstad [23]

designed a non-adaptive 3-query PCPs with almost perfect completeness and $s = 1/2$. Zwick [45] showed that any non-adaptive 3-query PCPs with perfect completeness must have soundness $s \geq 5/8$. He conjectured $s = 5/8$ is optimal. Håstad [23] showed $s = 3/4$ is possible and Khot and Saket [31] improved this to $s = 20/27$. Finally, O’Donnell and Wu [37] showed $s = 5/8$ under Khot’s d -to-1 Games conjecture. We remark that perfect completeness is important for several reasons: First, it is natural to require that testers should always accept correct objects. Second, it is required to show the inapproximability results for certain problems such as satisfiable CSPs and approximate coloring problems. Third, it is robust when we compose testers with other protocols.

Adaptive vs. non-adaptive queries It is also known that adaptive queries are more powerful than non-adaptive queries. Guruswami et al. [20] showed adaptive 3-query PCPs with $c = 1, s = 1/2$. As mentioned previously, Zwick showed non-adaptive 3-query PCPs with $c = 1$ have $s \geq 5/8$. Non-adaptivity is crucial in applications for inapproximability; we need queries to be non-adaptive to relate the parameters of PCPs to the inapproximability ratio of optimization problems.

Query complexity vs. soundness The trade-off between query complexity and soundness is also formalized as amortized query complexity, which is defined as $\frac{q}{\log_2(c/s)}$. Observe that repeating the tester t times independently, we can reduce the soundness to s^t while query complexity grows as $t \cdot q$. In this case, the amortized query complexity of the repeated tester remains the same as that of the original one. Hence this can be thought as a good and non-trivial measure for such trade-offs. Thus, constructions of testers that achieve better amortized query complexity is an interesting problem.

In this paper, we study the trade-offs between query complexity and soundness in the Long Code testing problem. Our focus is on non-adaptive testers with perfect completeness ($c = 1$).

1.1 Previous Work

The study of amortized query complexity in PCPs was initiated by Trevisan [44] and many works followed it [41, 42, 27, 26, 18]. One of the most notable earlier work is due to Samorodnitsky and Trevisan [42]: They constructed q -query non-adaptive PCPs with almost perfect completeness and $s = 2^{2\sqrt{q}}/2^q$ for infinitely many q . Later Engebretsen and Holmerin [18] improved this to $s = 2^{\sqrt{2q}}/2^q$. Håstad and Khot constructed PCPs with perfect completeness and $s = 2^{4\sqrt{q}}/2^q$. All these constructions use the distance from juntas in soundness condition, which allows them to compose Raz verifier as an outer verifier. Hence the resulting PCPs are unconditional.

Samorodnitsky and Trevisan [43] observed that the trade-offs achieved by these works are essentially tight if the soundness criterion is with respect to the distance from juntas. (Similar limitations were shown by [14, 32].) To

break the $s = 2^{\Theta(\sqrt{q})}/2^q$ barrier, they reworked the problem under the relaxed soundness criterion, i.e., low-degree influences. They proposed and analyzed the Hyper Graph Long Code test, that achieves almost perfect completeness and $s = 2q/2^q$ (or even $(q + 1)/2^q$ for infinitely many q), and combined it with Unique Games to construct conditional PCPs with the same parameters. They use a novel technique from additive combinatorics, namely, Gowers uniformity norm, in the analysis. Later Austrin and Mossel [5] improved the soundness to $s = (q + o(q))/2^q$ (or even $(q + 4)/2^q$ assuming the famous Hadamard Conjecture) for any q . Their technique is completely different from [43] and based on the Invariance Principle, which was developed by Mossel et al. [33, 34]. Similar constructions were also given in [24, 1, 21] where the emphasis is on the hardness of CSPs over “random predicates” or predicates of non-Boolean domains. However, all these works are affected by the loss of perfect completeness. Recently, Chen [16], building on the work of Samorodnitsky and Trevisan [43] and Håstad and Khot [26], showed q -query “adaptive” Long Code testers with perfect completeness and $s = q^3/2^q$ for infinitely many q .

1.2 Our Contribution

In this paper, we prove the following theorem:

Theorem 1.1. *For every $k > 2$, there exists a non-adaptive Long Code tester that makes $q = 2^k - 1$ queries, has completeness 1 and soundness $(2q + 3)/2^q$.*

1.3 Our Method

Our tester and its analysis is based on a natural but non-trivial extension of O’Donnell and Wu’s analysis for hardness of satisfiable 3-CSPs over Not-Two (NTW) predicates [37]. The most technical part of O’Donnell-Wu’s work is to bound the term of the form $\mathbf{E}[f(\mathbf{x})g(\mathbf{y})g(\mathbf{z})]$ where $\mathbf{x}, \mathbf{y}, \mathbf{z}$ are drawn from the “NTW”-distribution. We extend this to the case of the product of large number of functions with more complex distribution.

First we need to define suitable test distributions to make queries. To achieve good query complexity-soundness trade-offs, we adopt the Hyper Graph test introduced by Samorodnitsky and Trevisan [43]. The Hyper Graph test can only test linearity; for testing Long Code, we need to perturb the Hyper Graph test with some “noise” distribution.

Most often, the noise distribution is the uniform distribution. A mixture of the distribution for linearity test and uniform distribution with sufficiently small probability often yields a desired Long Code test. However, this forces the loss of perfect completeness because possible query-answers supports all over the domain, and any “yes-instance” can be rejected with small probability. To avoid this, Chen [16] introduced a noise distribution based on Håstad and Khot’s query-efficient PCPs with perfect completeness [26]. His analysis is based on Fourier analysis and Gowers norm as Samorodnitsky and Trevisan did. Although his test achieves good query-soundness trade-offs, one drawback is its

adaptivity. In this paper, we use a different noise distribution, a modification of the All-Equal noise. Previously, the All-Equal noise was used by O’donnell and Wu [36] combined with 3-bit linearity test to achieve $c = 1, s = 5/8$, non-adaptive Long Code test based on NTW predicate. The analysis is based on Fourier analysis and hypercontractivity inequality. Later the same authors [37] managed to construct PCPs with the same parameters $c = 1$ and $s = 5/8$ under Khot’s d -to-1 Games conjecture [29]. The analysis is based on Invariance Principle style argument [33, 34]. The crucial point is bounding the correlation of probability spaces.

We follow the approach of [37]. Basically our analysis becomes complicated because we have to bound terms of the form $\mathbf{E}[f_1 f_2 \cdots f_q]$ instead of a “cubic”-term, but some analysis becomes easier because we only work with functions over the same domain (or functions with “unique” constraints) rather than functions with d -to-1 constraints as in the case of [37].

1.4 Organization of Paper

In Section 2, we introduce notations and definitions needed throughout this paper. Also, we describe backgrounds on correlated spaces. In Section 3, we define our multi-function Long Code testing problem formally, then we describe our test distribution and Long Code test. In Section 4, we present completeness and soundness analysis of our tester. Finally, we notice some future work in Section 5.

2 Preliminaries

Let $[k] = \{1, 2, \dots, k\}$, $2_{\geq 1}^{[k]} = \{S \subseteq [k] : |S| \geq 1\}$ and $2_{\geq 2}^{[k]} = \{S \subseteq [k] : |S| \geq 2\}$. For a vector $x \in \{-1, 1\}^n$, x_i denotes the i th element of x . Write the all-ones vector as $\mathbf{1} = (1, 1, \dots, 1)$ and denote by $\mathbf{1}_{-i}$ a vector obtained by flipping the i th coordinate of $\mathbf{1}$ to -1 . A set $\{i\}$ is sometimes written as i . We denote by $|\mathbf{S}|$ the cardinality of a set \mathbf{S} . A *Boolean* function is a function $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$.

2.1 Fourier Expansion

Definition 2.1 (Fourier expansion). *The Fourier expansion of a function $f : \{-1, 1\}^n \rightarrow \mathbb{R}$ is*

$$f(x) = \sum_{S \subseteq [n]} \widehat{f}(S) \chi_S(x),$$

where

$$\chi_S(x) = \prod_{i \in S} x_i, \quad \widehat{f}(S) = \mathbf{E}_x [f(x) \chi_S(x)],$$

and expectation is with respect to uniform distribution.

It is easy to see that for any $S, T \subseteq [n]$, $\mathbf{E}_x [\chi_S(x) \chi_T(x)] = 1$ if $S = T$ and 0 otherwise. From this, we have

Lemma 2.2 (Parseval's theorem).

$$\mathbf{E}_x[f(x)^2] = \sum_{S \subseteq [n]} \widehat{f}(S)^2.$$

For more on Fourier expansion and Fourier analysis see, e.g., [35].

2.2 Influences

In this section we recall basic notions from Fourier analysis, influence and the Bonami-Beckner operator. We first define a notion of the influence of a coordinate on a function f :

Definition 2.3. For a function $f : \{-1, 1\}^n \rightarrow \mathbb{R}$, we define the influence of $i \in [n]$ on f to be

$$\text{Inf}_i(f) = \sum_{S \ni i} \widehat{f}(S)^2.$$

We also define the degree- w influence of $i \in [n]$ on f to be

$$\text{Inf}_i^{\leq w}(f) = \sum_{S \ni i: |S| \leq w} \widehat{f}(S)^2.$$

Next we define the cross-influence of a coordinate on a collection of functions.

Definition 2.4. Let f_1, \dots, f_t be a collection of Boolean functions. Then the cross-influence of i for them is defined as

$$\text{XInf}_i(f_1, \dots, f_t) = \max_{j \neq k} \min(\text{Inf}_i(f_j), \text{Inf}_i(f_k)).$$

Similarly,

$$\text{XInf}_i^{\leq w}(f_1, \dots, f_t) = \max_{j \neq k} \min(\text{Inf}_i^{\leq w}(f_j), \text{Inf}_i^{\leq w}(f_k)).$$

Note that XInf_i is the second highest value of $\text{Inf}_i(f_j)$ among all functions f_j .

We next recall the Bonami-Beckner operator T_ρ acting on Boolean functions:

Definition 2.5. Let $0 \leq \rho \leq 1$. The Bonami-Beckner operator T_ρ is a linear operator mapping functions $f : \{-1, 1\}^n \rightarrow \mathbb{R}$ into functions $T_\rho f : \{-1, 1\}^n \rightarrow \mathbb{R}$ via

$$(T_\rho f)(x) = \mathbf{E}[f(\mathbf{y})],$$

where in the expectation, \mathbf{y} is formed from x by setting $\mathbf{y}_i = x_i$ with probability ρ and setting \mathbf{y}_i to be a uniformly random bit with probability $1 - \rho$.

The following lemma is useful.

Lemma 2.6. (Hypercontractive inequality [13, 22]) *Let p, q, γ be real numbers such that $1 \leq p \leq q$ and $1 - \gamma \leq \sqrt{\frac{p-1}{q-1}}$, Then it holds that for all $f : \{-1, 1\}^n \rightarrow \mathbb{R}$,*

$$\|T_{1-\gamma}f\|_q \leq \|f\|_p.$$

The operator T_ρ can alternately be defined by the following formula:

Lemma 2.7.

$$T_\rho f = \sum_{S \subseteq [n]} \rho^{|S|} \widehat{f}(S) \chi_S.$$

We need the following lemmas.

Lemma 2.8. *For any function $f : \{-1, 1\}^n \rightarrow [-1, 1]$ and $0 < \gamma < 1/2$,*

$$\sum_{i=1}^n \text{Inf}_i(T_{1-\gamma}f) \leq \frac{1}{2\gamma}.$$

Lemma 2.9. *Let f be a Boolean function and $\gamma > 0$. Suppose $\text{Inf}_i(T_{1-\gamma/2}f) \geq \tau$, then there exists a constant w which only depends on γ and τ such that $\text{Inf}_i^{\leq w}(f) \geq \frac{\tau}{2}$.*

Proof. Set w such that $(1 - \frac{\gamma}{2})^w \leq \frac{\tau}{2}$, then

$$\begin{aligned} \tau &\leq \text{Inf}_i(T_{1-\gamma/2}f) = \sum_{S \ni i} (1 - \frac{\gamma}{2})^{|S|} \widehat{f}(S)^2 \\ &= \sum_{S \ni i, |S| \leq w} (1 - \frac{\gamma}{2})^{|S|} \widehat{f}(S)^2 + \sum_{S \ni i, |S| > w} (1 - \frac{\gamma}{2})^{|S|} \widehat{f}(S)^2 \\ &\leq \sum_{S \ni i, |S| \leq w} \widehat{f}(S)^2 + (1 - \frac{\gamma}{2})^w \leq \text{Inf}_i^{\leq w}(f) + \frac{\tau}{2}, \end{aligned}$$

where the second inequality is by Parseval's theorem. \square

2.3 Correlated Spaces and Noise

We now recall the definition of correlation for correlated probability spaces, as introduced by Mossel [34].

Definition 2.10. *Let (Ω, Ψ, μ) be a (finite) correlated probability space, meaning that μ is a distribution on the finite product set $\Omega \times \Psi$ and that the marginals of μ on Ω and Ψ have full support. Define the correlation $\rho(\Omega, \Psi; \mu)$ between Ω and Ψ to be*

$$\max \left\{ \mathbf{Cov}_{(\omega, \psi) \sim \mu} [f(\omega), g(\psi)] \mid f : \Omega \rightarrow \mathbb{R}, g : \Psi \rightarrow \mathbb{R}, \mathbf{Var}_{(\omega, \psi) \sim \mu} [f(\omega)] = \mathbf{Var}_{(\omega, \psi) \sim \mu} [g(\psi)] = 1 \right\}$$

It is clear that in the definition of $\rho(\Omega, \Psi; \mu)$, we can equivalently maximize $|\mathbf{E}[fg]|$ over f restricted to have $\mathbf{E}[f] = 0, \mathbf{E}[f^2] \leq 1$ under μ 's marginal on Ω ; or, over similarly restricted g (or both).

Definition 2.11 (Markov Operator). *Let (Ω, Ψ, μ) be a correlated probability space. Markov Operator associated with the probability space is the operator mapping $f \in L^p(\Psi, \mu)$ to $Uf \in L^p(\Omega, \mu)$ by*

$$(Uf)(x) = \mathbf{E}[f(Y)|X = x]$$

for $x \in \Omega$ where $(X, Y) \in (\Omega, \Psi)$ is distributed according to μ .

Lemma 2.12. (Lemma 2.9 of [34]) *Let (Ω, Ψ, μ) be two correlated spaces and let α be the minimum probability in the support of μ . Define a bipartite graph $G = (\Omega, \Psi, E)$ where $(a, b) \in \Omega \times \Psi$ satisfies $(a, b) \in E$ if $\mu(a, b) > 0$. If G is connected then*

$$\rho(\Omega, \Psi; \mu) \leq 1 - \alpha^2/2.$$

Efron-Stein Decomposition is an analogue of Fourier expansion for (non-Boolean) functions over general domains.

Definition 2.13 (Efron-Stein Decomposition). *Let $(\Omega_1, \mu_1), \dots, (\Omega_n, \mu_n)$ be discrete probability spaces $(\Omega, \mu) = \prod_{i=1}^n (\Omega_i, \mu_i)$. The Efron-Stein decomposition of $f : \Omega \rightarrow \mathbb{R}$ is given by*

$$f(x) = \sum_{S \subseteq [n]} f_S(x), \quad (1)$$

where

$$f_S(x) = \sum_{S' \subseteq S} (-1)^{|S \setminus S'|} \mathbf{E}[f(X)|(X_i)_{i \in S'} = (x_i)_{i \in S'}].$$

From the definition of $f_S(x)$, f_S depends only on $(x_i)_{i \in S}$, and it is known that for all $S \not\subseteq S'$ and $(x_i)_{i \in S'}$ it holds that $\mathbf{E}[f_S|(X_i)_{i \in S'} = (x_i)_{i \in S'}] = 0$ [34].

Proposition 2.14. (Proposition 2.11 of [34]) *Let $(\Omega_i, \Psi_i, \mu_i)$ be correlated probability spaces and let U_i be the Markov operator associated with Ω_i and Ψ_i for $1 \leq i \leq n$. Let*

$$\Omega = \prod_{i=1}^n \Omega_i, \Psi = \prod_{i=1}^n \Psi_i, \mu = \prod_{i=1}^n \mu_i, U = \bigotimes_{i=1}^n U_i.$$

Suppose $f \in L^2(\Psi)$ has Efron-Stein decomposition (1). Then the Efron-Stein decomposition of Uf satisfies

$$(Uf)_S = U(f_S).$$

Proposition 2.15. (Proposition 2.12 of [34]) *Assume the setting of Proposition 2.14 and that further for all i it holds that $\rho(\Omega_i, \Psi_i, \mu_i) \leq \rho_i$. Then for all f it holds that*

$$\|U(f_S)\|_2 \leq \left(\prod_{i \in S} \rho_i \right) \|f_S\|_2.$$

Proposition 2.16. (Proposition 2.13 of [34]) *Assume the setting of Proposition 2.14. Then*

$$\rho(\Omega, \Psi; \mu) = \max \rho(\Omega_i, \Psi_i, \mu_i).$$

3 Long Code Test

Definition 3.1. For $i \in [n]$, the i th Long Code is the function $f(x) = x_i$

Now let us define a t -function Long Code Test. Suppose we are given oracle access to a collection of Boolean functions f_1, f_2, \dots, f_t . We want to make as few queries as possible into these functions to decide if all the functions are the same Long Code, or no two functions have some common structure. More precisely, we have the following definition:

Definition 3.2. We say that a test $\mathbf{T} = \mathbf{T}^{f_1, f_2, \dots, f_t}$ is a t -functions Long Code test with completeness c and soundness s if \mathbf{T} is given oracle access to a family of t functions $f_1, f_2, \dots, f_t : \{-1, 1\}^n \rightarrow \{-1, 1\}$, such that

- if there exists some $i \in [n]$ such that for all $a \in [t]$, $f_a(x) = x_i$, then \mathbf{T} accepts with probability at least c , and
- for every $\varepsilon > 0$, there exist a constant $\tau > 0$ and a fixed positive integer w such that if \mathbf{T} accepts with probability at least $s + \varepsilon$, then there exists some $i \in [n]$ such that $\text{XInf}_i^{\leq w}(f_1, \dots, f_t) \geq \tau$.

3.1 Folding

As introduced in [9], we assume that the functions are folded. We do so by requiring our Long Code test to make queries in a special manner. Suppose the test wants to query f at point $x \in \{-1, 1\}^n$. If $x_1 = 1$, then the test queries $f(x)$ as usual. If $x_1 = -1$, then the test queries at the point $-x$ and use $-f(-x)$ as the evaluation of x . Folding ensures that $f(-x) = -f(x)$ and $\hat{f}(\emptyset) = 0$.

3.2 The Test Distribution

We now view $f_S : \{-1, 1\}^n \rightarrow \{-1, 1\}$ as a function over an n -fold product set $\mathcal{X}_S^1 \times \mathcal{X}_S^2 \times \dots \times \mathcal{X}_S^n$ where each $\mathcal{X}_S^i = \{-1, 1\}^{\{i\}}$.

For each integer $k > 2$, we construct \mathcal{T}_k as a product probability distribution over the n -fold product set

$$\prod_{i=1}^n \prod_{S \in 2_{\geq 1}^{[k]}} \mathcal{X}_S^i = \prod_{i=1}^n \mathcal{X}_1^i \times \prod_{i=1}^n \mathcal{X}_2^i \times \dots \times \prod_{i=1}^n \mathcal{X}_n^i \times \prod_{i=1}^n \mathcal{X}_{\{1,2\}}^i \times \dots \times \prod_{i=1}^n \mathcal{X}_{\{1,2,\dots,k\}}^i.$$

For each i we have a distribution \mathcal{T}_k^i on $\prod_{S \in 2_{\geq 1}^{[k]}} \mathcal{X}_S^i$ and we think of this as a “correlated space.”

The first distribution is the “Hyper Graph Test” distribution which is a slight modification of the one used by Samorodnitsky and Trevisan [43].

Definition 3.3. Define distribution \mathcal{H}_k generating $(\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_{\{1,2,\dots,k\}}) \in \prod_{S \in 2_{\geq 1}^{[k]}} \mathcal{X}_S^i$ as follows: The bits $\{\mathbf{x}_j\}_{j \in [k]}$ are independent and uniformly at random; then for each $S \in 2_{\geq 2}^{[k]}$, \mathbf{x}_S is set to be $-\prod_{j \in S} \mathbf{x}_j$.

Note that under \mathcal{H}_k , the marginal distribution on \mathbf{x}_S for each $S \in 2_{\geq 2}^{[k]}$ is also uniformly random. The second distribution is the “Noise” distribution.

Definition 3.4. Let $p, q, r > 0$ be such that

$$\begin{aligned} p + q + (2^k - 1)r &= 1, \\ p + r &= q + (2^k - 2)r = 1/2, \\ q &= r, \end{aligned}$$

i.e., $p = (2^k - 2)/(2^{k+1} - 2)$, $q = r = 1/(2^{k+1} - 2)$. Define distribution \mathcal{N}_k generating $\mathbf{x} = (\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_{\{1,2,\dots,k\}}) \in \prod_{S \in 2_{\geq 1}^{[k]}} \mathcal{X}_S^i$ as follows: (i) With probability p , set $\mathbf{x} = \mathbf{1}$, (ii) with probability q , set $\mathbf{x} = -\mathbf{1}$, and (iii) with probability $(2^k - 1)r$, first pick $S \in 2_{\geq 1}^{[k]}$ uniformly at random, then set $\mathbf{x} = -\mathbf{1}_{-S}$, i.e., all -1 everywhere except the coordinate indexed by S is 1. Additionally, for $0 < \delta < 1$, define distribution $\mathcal{H}_{k,\delta}$ to be the mixture distribution $\mathcal{H}_{k,\delta} = (1 - \delta)\mathcal{H}_k + \delta\mathcal{N}_k$; i.e., one draws from \mathcal{H}_k with probability $1 - \delta$ and from \mathcal{N}_k with probability δ .

Again note that under \mathcal{N}_k (also $\mathcal{H}_{k,\delta}$), the marginal distribution on \mathbf{x}_S for each $S \in 2_{\geq 1}^{[k]}$ is uniformly random.

We are now ready to define the test distribution \mathcal{T}_k^i :

Definition 3.5. For each $i \in [n]$ we define \mathcal{T}_k^i to be $\mathcal{H}_{k,\delta}$, with $\delta = (\frac{\varepsilon}{9 \cdot |2_{\geq 1}^{[k]}|})^2$, where the domain of $\mathcal{H}_{k,\delta}$ is appropriately identified with the domain $\prod_{S \in 2_{\geq 1}^{[k]}} \mathcal{X}_S^i$ of \mathcal{T}_k^i . \mathcal{T}_k is the product distribution of these distributions $\mathcal{T}_k = \otimes_{i=1}^n \mathcal{T}_k^i$

Definition 3.6. For any positive integer k and $S_0 \in 2_{\geq 1}^{[k]}$, we define distribution \mathcal{I}_{k,S_0} generating $(\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_{\{1,2,\dots,k\}}) \in \prod_{S \in 2_{\geq 1}^{[k]}} \mathcal{X}_S^i$ as follows: First draw from \mathcal{H}_k ; then uniformly rerandomize \mathbf{x}_{S_0} .

Definition 3.7. For $0 < \delta < 1$, define distribution $\mathcal{I}_{k,S_0,\delta}$ to be the mixture distribution $\mathcal{I}_{k,S_0,\delta} = (1 - \delta)\mathcal{I}_{k,S_0} + \delta\mathcal{N}_k$.

3.3 A Query Efficient Long Code Test

Let $k > 2$ be an integer and $\{f_S\}_{S \in 2_{\geq 1}^{[k]}}$ be a collection of Boolean functions. We define the following Long Code test based on distribution \mathcal{T}_k :

Long Code Test \mathbf{T}_k : with oracle access to $\{f_S\}_{S \in 2_{\geq 1}^{[k]}}$

1. Pick $(\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_{\{1,2,\dots,k\}}) \in \prod_{S \in 2_{\geq 1}^{[k]}} \prod_{i=1}^n \mathcal{X}_S^i$ from \mathcal{T}_k .
2. For each $S \in 2_{\geq 1}^{[k]}$, query $f_S(\mathbf{x}_S)$.
3. Accept iff either of the followings holds:
 - (a) (Hyper Graph) For every $S \in 2_{\geq 2}^{[k]}$, $\prod_{s \in S} f_s(\mathbf{x}_s) \neq f_S(\mathbf{x}_S)$, or
 - (b) (Noise) $(f_{\{1\}}(\mathbf{x}_1), f_{\{2\}}(\mathbf{x}_2), \dots, f_{\{1,2,\dots,k\}}(\mathbf{x}_{\{1,2,\dots,k\}}))$ is either $\mathbf{1}$ or $-\mathbf{1}$ or $-\mathbf{1}_{-S}$ for some $S \in 2_{\geq 1}^{[k]}$.

Theorem 3.8 (main theorem restated). *For every integer $k > 2$, there exists a non-adaptive Long Code test \mathbf{T}_k that makes $q = 2^k - 1$ queries, has completeness 1 and soundness $(2q + 3)/2^q$.*

We briefly explain the idea behind our Long Code test.

First, the test is “approximation resistant” in the sense that random functions pass the test with probability $(2q + 3)/2^q$, i.e., the soundness of the test. Therefore, to obtain the soundness claimed above, we examine the Fourier expansion of the acceptance probability and show that all non-constant terms are sufficiently small for random functions.

Second, since our goal is to test whether functions are Long Codes or not, we have to be able to test whether functions are linear functions or not efficiently. To do so, we employ the distribution based on the Hyper Graph Test. Note that arbitrary linear functions such as parity functions pass the test using only this distribution with probability one.

Third, to test Long Code instead of linear functions, we need noise distribution to kill “large” linear functions, i.e., linear functions which depend on many input variables. Our noise distribution achieves the purpose well. If we sample from our test distribution, roughly $(1 - \delta)$ fraction of indices sample from the Hyper Graph Test distribution and δ fraction of indices sample from the Noise distribution, and large linear functions cannot pass the test with high probability.

4 Proof of Theorem 3.8

Completeness Analysis: The completeness analysis is entirely standard.

Soundness Analysis: We prove that for every $\epsilon > 0$, there exist a constant $\tau > 0$ and a positive integer w such that if for every i , $\text{XInf}_i^{\leq w}(f_1, \dots, f_{[k]}) < \tau$,

then the accepting probability is at most

$$\frac{2|2_{\geq 1}^{[k]}| + 3}{2^{|2_{\geq 1}^{[k]}|}} + \epsilon.$$

We arithmetize the probability that \mathbf{T}_k accepts. Let $\mathbf{Z}_S = -f_S(\mathbf{x}_S) \prod_{s \in S} f_s(\mathbf{x}_s)$. Then

$$\begin{aligned} \Pr[\mathbf{T}_k \text{ accepts}] &= \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in 2_{\geq 2}^{[k]}} \frac{1 + \mathbf{Z}_S}{2} + \prod_{S \in 2_{\geq 1}^{[k]}} \frac{1 - f_S(\mathbf{x}_S)}{2} + \prod_{S \in 2_{\geq 1}^{[k]}} \frac{1 + f_S(\mathbf{x}_S)}{2} \right. \\ &\quad \left. + \sum_{S' \in 2_{\geq 1}^{[k]}} \frac{1 + f_{S'}(\mathbf{x}_{S'})}{2} \prod_{S \in 2_{\geq 1}^{[k]} \setminus \{S'\}} \frac{1 - f_S(\mathbf{x}_S)}{2} \right]. \end{aligned}$$

This can be written as

$$\begin{aligned} &= \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\frac{1}{2^{|2_{\geq 2}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 2}^{[k]}} \prod_{S \in \mathbf{s}} \mathbf{Z}_S + \frac{1}{2^{|2_{\geq 1}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 1}^{[k]}} (-1)^{|\mathbf{s}|} \prod_{S \in \mathbf{s}} f_S(\mathbf{x}_S) + \frac{1}{2^{|2_{\geq 1}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 1}^{[k]}} \prod_{S \in \mathbf{s}} f_S(\mathbf{x}_S) \right. \\ &\quad \left. + \sum_{S' \in 2_{\geq 1}^{[k]}} \frac{1 + f_{S'}(\mathbf{x}_{S'})}{2^{|2_{\geq 1}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 1}^{[k]} \setminus \{S'\}} (-1)^{|\mathbf{s}|} \prod_{S \in \mathbf{s}} f_S(\mathbf{x}_S) \right] \\ &\leq \frac{1}{2^{|2_{\geq 2}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 2}^{[k]}} \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{s}} \mathbf{Z}_S \right] + \frac{2 + |2_{\geq 1}^{[k]}|}{2^{|2_{\geq 1}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 1}^{[k]}} \left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{s}} f_S(\mathbf{x}_S) \right] \right|. \end{aligned}$$

We obtain the bounds

$$\frac{1}{2^{|2_{\geq 2}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 2}^{[k]}} \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{s}} \mathbf{Z}_S \right] = \frac{1}{2^{|2_{\geq 2}^{[k]}|}} + \frac{1}{2^{|2_{\geq 2}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 2}^{[k]}: \mathbf{s} \neq \emptyset} \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{s}} \mathbf{Z}_S \right] \leq \frac{1}{2^{|2_{\geq 2}^{[k]}|}} + \epsilon/3$$

and

$$\begin{aligned} \frac{2 + |2_{\geq 1}^{[k]}|}{2^{|2_{\geq 1}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 1}^{[k]}} \left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{s}} f_S(\mathbf{x}_S) \right] \right| &= \frac{2 + |2_{\geq 1}^{[k]}|}{2^{|2_{\geq 1}^{[k]}|}} + \frac{2 + |2_{\geq 1}^{[k]}|}{2^{|2_{\geq 1}^{[k]}|}} \sum_{\mathbf{s} \subseteq 2_{\geq 1}^{[k]}: \mathbf{s} \neq \emptyset} \left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{s}} f_S(\mathbf{x}_S) \right] \right| \\ &\leq \frac{2 + |2_{\geq 1}^{[k]}|}{2^{|2_{\geq 1}^{[k]}|}} + 2\epsilon/3 \end{aligned}$$

using the following lemma.

Lemma 4.1. For any $\varepsilon > 0$, there exists τ, δ, w such that if $\text{XInf}_i^{\leq w}(f_1, \dots, f_{[k]}) < \tau$ for any i then, for any $\mathbf{S} \subseteq 2_{\geq 1}^{[k]}$, $\mathbf{S} \neq \emptyset$,

$$\left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{S}} f_S(\mathbf{x}_S) \right] \right| \leq \frac{\varepsilon}{3 \cdot |2_{\geq 1}^{[k]}|}.$$

Note that the term $\mathbf{E}_{\mathcal{H}_{k,\delta}^n} [\prod_{S \in \mathbf{S}} \mathbf{Z}_S]$ is also represented as $\pm \mathbf{E}_{\mathcal{H}_{k,\delta}^n} [\prod_{S \in \mathbf{S}'} f_S(\mathbf{x}_S)]$ since each f_S is a Boolean function.

4.1 Proof of Lemma 4.1

First note that for any $\mathbf{S} \subseteq 2_{\geq 1}^{[k]}$, $|\mathbf{S}| \leq 2$, the LHS becomes at most δ since the marginal distribution of $\mathcal{H}_{k,\delta}^n$ on $\prod_{S \in \mathbf{S}} \prod_{i=1}^n \mathcal{X}_S^i$ is uniformly random and functions are folded. We need to prove Lemma 4.1 for $\mathbf{S} \subseteq 2_{\geq 1}^{[k]}$, $|\mathbf{S}| \geq 3$ and actually we only prove Lemma 4.1 for $\mathbf{S} = 2_{\geq 1}^{[k]}$; any choice of \mathbf{S} can be handled in almost the same way. It is easy to see that following lemmas with the triangle inequality concludes the proof of Lemma 4.1. Note that in Definition 3.5, we set δ to satisfy $\max\{\delta, 3\sqrt{\delta}\} \leq \frac{\varepsilon}{3 \cdot |2_{\geq 1}^{[k]}|}$.

Lemma 4.2. By taking $\gamma > 0$ small enough as a function of δ, k ,

$$\left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{S}} f_S(\mathbf{x}_S) - \prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] \right| \leq \sqrt{\delta}.$$

Lemma 4.3. For any $S_0 \in \mathbf{S}$, $|\mathbf{S}| \geq 3$, γ, δ , there exist τ, w such that if $\text{XInf}_i^{\leq w}(f_1, \dots, f_{[k]}) < \tau$ for any i then,

$$\left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] - \mathbf{E}_{\mathcal{I}_{k,S_0,\delta}^n} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] \right| \leq \sqrt{\delta}.$$

Lemma 4.4. For any $S_0 \in \mathbf{S}, \gamma, \delta$, it holds that

$$\left| \mathbf{E}_{\mathcal{I}_{k,S_0,\delta}^n} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] \right| \leq \sqrt{\delta}.$$

Proofs of Lemmas 4.2, 4.3 and 4.4 are given in Subsections 4.2, 4.3 and 4.4, respectively.

4.2 Proof of Lemma 4.2

We define a lexicographical order $<_{\text{lex}}$ between elements of $2_{\geq 1}^{[k]}$ as

$$1 <_{\text{lex}} 2 <_{\text{lex}} \dots <_{\text{lex}} n <_{\text{lex}} \{1, 2\} <_{\text{lex}} \{1, 3\} <_{\text{lex}} \dots <_{\text{lex}} \{1, 2, \dots, k\},$$

and let $\mathbf{S}_l \subseteq 2_{\geq 1}^{[k]}$ be the first l elements according to $<_{\text{lex}}$. It is easy to see that the following lemma with the triangle inequality completes the proof of Lemma 4.2.

Lemma 4.5. *By taking $\gamma > 0$ small enough as a function of δ, k , we ensure, for any $l \in [|2_{\geq 1}^{[k]}|]$,*

$$\left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{S}_l} f_S(\mathbf{x}_S) \cdot \prod_{S \in \mathbf{S} \setminus \mathbf{S}_l} T_{1-\gamma} f_S(\mathbf{x}_S) - \prod_{S \in \mathbf{S}_{l-1}} f_S(\mathbf{x}_S) \cdot \prod_{S \in \mathbf{S} \setminus \mathbf{S}_{l-1}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] \right| \leq \sqrt{\delta} / 2^{|2_{\geq 1}^{[k]}|}.$$

Proof. We only prove the case for $l = |2_{\geq 1}^{[k]}|$. Other cases can be handled in almost the same way by changing the role of $f_{[k]}$ and $\prod_{i=1}^n \mathcal{X}_{[k]}^i$ to f_S and $\prod_{i=1}^n \mathcal{X}_S^i$ for each $S \in 2_{\geq 1}^{[k]}$. (Notice that we can do so by Lemma 4.6 that ensures the correlation is strictly less than 1. This is not the case in the proof of Theorem C.1 in [37]; they need some additional analysis for some choice of l , i.e., the proof of Lemma C.5.) We have

$$\begin{aligned} & \left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in 2_{\geq 1}^{[k]}} f_S(\mathbf{x}_S) - \prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} f_S(\mathbf{x}_S) \cdot T_{1-\gamma} f_{[k]}(\mathbf{x}_{[k]}) \right] \right| \\ &= \left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} f_S(\mathbf{x}_S) \cdot (id - T_{1-\gamma}) f_{[k]}(\mathbf{x}_{[k]}) \right] \right| \\ &= \left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} f_S(\mathbf{x}_S) \cdot U(id - T_{1-\gamma}) f_{[k]}(\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_{\{2,3,\dots,k\}}) \right] \right|, \quad (2) \end{aligned}$$

where U is the Markov operator for the correlated probability space $(\prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} (\prod_{i=1}^n \mathcal{X}_S^i), \prod_{i=1}^n \mathcal{X}_{[k]}^i; \mathcal{H}_{k,\delta}^n)$, which maps functions on the latter space $\prod_{i=1}^n \mathcal{X}_{[k]}^i$ into functions on the former space $\prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} (\prod_{i=1}^n \mathcal{X}_S^i)$.

Now we consider the quantity inside the expectation in (2) to be a product of two functions on $\prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} (\prod_{i=1}^n \mathcal{X}_S^i)$, namely, $F = \prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} f_S$ and $G = U(id - T_{1-\gamma}) f_{[k]}$. We take the Efron-Stein decomposition of these two functions with respect to the (product) distribution $\prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} (\prod_{i=1}^n \mathcal{X}_S^i)$. Then by the orthogonality of the Efron-Stein decomposition and the Cauchy-Schwartz inequality,

$$\begin{aligned} (2) &= \left| \sum_{V \subseteq [n]} \mathbf{E}_{\mathcal{H}_{k,\delta}^n} [F_V(\mathbf{x}_1, \dots, \mathbf{x}_{\{2,\dots,k\}}) \cdot G_V(\mathbf{x}_1, \dots, \mathbf{x}_{\{2,\dots,k\}})] \right| \\ &\leq \sqrt{\sum_{V \subseteq [n]} \|F_V\|_2^2} \sqrt{\sum_{V \subseteq [n]} \|G_V\|_2^2}. \quad (3) \end{aligned}$$

where the 2-norms $\|\cdot\|_2$ are with respect to $\mathcal{H}_{k,\delta}^n$'s marginal distribution on $\prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} (\prod_{i=1}^n \mathcal{X}_S^i)$. By orthogonality again, the quantity $\sqrt{\sum_{V \subseteq [n]} \|F_V\|_2^2}$ is just $\|F\|_2$, which is precisely 1 because F 's range is $\{-1, 1\}$. Hence we have

$$(3) \leq \sqrt{\sum_{V \subseteq [n]} \|G_V\|_2^2}. \quad (4)$$

From Lemma 2.14, $G_V = UG'_V$ where $G' = (id - T_{1-\gamma})f_{[k]}$. Here the Efron-Stein decomposition is with respect to $\mathcal{H}_{k,\delta}^n$'s marginal distribution on $\prod_{i=1}^n \mathcal{X}_{[k]}^i$, i.e., the uniform distribution. It is also easy to check that this Efron-Stein decomposition of $f_{[k]}$ has

$$(f_{[k]})_V = \widehat{f}_{[k]}(V)\chi_V.$$

It follows that applying the Bonami-Beckner operator $T_{1-\gamma}$ to $f_{[k]}$ also commutes with taking the Efron-Stein decomposition. Hence we have $G_V = UG'_V = U(id - T_{1-\gamma})(f_{[k]})_V$. Substituting this into (4) yields

$$(4) = \sqrt{\sum_{V \subseteq [n]} \|U(id - T_{1-\gamma})(f_{[k]})_V\|_2^2}. \quad (5)$$

From Lemma 4.6, which is proved later, the correlation

$$\rho \left(\prod_{S \in 2_{\geq 1}^{[k]} \setminus [k]} \mathcal{X}_S^i, \mathcal{X}_{[k]}^i; \mathcal{H}_{k,\delta} \right) \leq \rho_0 = 1 - \alpha^2/2$$

for every $i \in [n]$, where $\alpha = \delta/(2^{k+1}-2)$; Applying Proposition 2.15, we conclude that for each $V \subseteq [n]$,

$$\|U(id - T_{1-\gamma})(f_{[k]})_V\|_2 \leq \rho_0^{|V|} \|(id - T_{1-\gamma})(f_{[k]})_V\|_2, \quad (6)$$

where the 2-norm on the right is with respect to the uniform distribution on $\{-1, 1\}^n$.

Next, observe that

$$\|(id - T_{1-\gamma})(f_{[k]})_V\|_2^2 = (1 - (1 - \gamma)^{|V|})^2 \widehat{f}_{[k]}(V)^2.$$

Substituting (6) and then into (5), we determine

$$(5) \leq \sqrt{\sum_{V \subseteq [n]} \rho_0^{2|V|} (1 - (1 - \gamma)^{|V|})^2 \widehat{f}_{[k]}(V)^2}. \quad (7)$$

We now bound

$$\rho_0^{|V|} (1 - (1 - \gamma)^{|V|}) \leq \max\{\rho_0^{|V|}, (1 - (1 - \gamma)^{|V|})\} \leq \sqrt{\delta}/2^{2^{\lfloor k/2 \rfloor}}.$$

Set $\beta = \sqrt{\delta}/2^{2^{\lfloor k/2 \rfloor}}$ and $\gamma = \beta \log \rho_0 / \log \beta$. Then, if $|V| \geq \log \beta / \log \rho_0$, we have $\rho_0^{|V|} \leq \beta$. If $|V| < \log \beta / \log \rho_0$, $(1 - (1 - \gamma)^{|V|}) \leq \gamma|V| \leq \beta$.

Doing so, we get

$$(7) \leq \sqrt{\sum_{V \subseteq [n]} (\sqrt{\delta}/2^{2^{\lfloor k/2 \rfloor}})^2 \widehat{f}_{[k]}(V)^2} = (\sqrt{\delta}/2^{2^{\lfloor k/2 \rfloor}}) \sqrt{\mathbf{E}[f_{[k]}^2]} = \sqrt{\delta}/2^{2^{\lfloor k/2 \rfloor}}$$

where the first equality is by Parseval's theorem. \square

Now, we give Lemma 4.6.

Lemma 4.6. *For any $i \in [n]$ and $S' \in 2_{\geq 1}^{[k]}$,*

$$\rho \left(\prod_{S \in 2_{\geq 1}^{[k]} \setminus \{S'\}} \mathcal{X}_S^i, \mathcal{X}_{S'}^i; \mathcal{H}_{k,\delta} \right) \leq 1 - \alpha^2/2,$$

where $\alpha = \delta/(2^{k+1} - 2)$.

Proof. We use Lemma 2.12 to show that

$$\rho \left(\prod_{S \in 2_{\geq 1}^{[k]} \setminus \{S'\}} \mathcal{X}_S^i, \mathcal{X}_{S'}^i; \mathcal{H}_{k,\delta} \right) \leq 1 - \alpha^2/2 \quad (8)$$

for each i . First observe that for any $(a, b) \in \prod_{S \in 2_{\geq 1}^{[k]} \setminus \{S'\}} \mathcal{X}_S^i \times \mathcal{X}_{S'}^i$, either $b = 1$ or $b = -1$ holds. Next observe that vertices $(\mathbf{1}), (-\mathbf{1}) \in \mathcal{X}_{S'}^i$ are connected because there are edges $(-\mathbf{1}, 1), (-\mathbf{1}, -1)$ due to the definition of the distribution \mathcal{N}_k . Thus we can guarantee that the bipartite graph is connected. Easy calculation of p, q, r shows that the probability of the smallest atom is $\alpha = \delta/(2^{k+1} - 2)$. \square

4.3 Proof of Lemma 4.3

In this section, we prove the following.

Lemma 4.7 (Restatement of Lemma 4.3). *For any $S_0 \in \mathbf{S}, |\mathbf{S}| \geq 3, \gamma, \delta$, there exist τ, w such that if $\text{XInf}_i^{\leq w}(f_1, \dots, f_{[k]}) < \tau$ for any i then,*

$$\left| \mathbf{E}_{\mathcal{H}_{k,\delta}^n} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] - \mathbf{E}_{\mathcal{I}_{k,S_0,\delta}^n} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] \right| \leq \sqrt{\delta}.$$

From Proposition 2.9, by setting w as a function of γ and τ , we can assume $\text{XInf}_i(T_{1-\gamma/2} f_1, \dots, T_{1-\gamma/2} f_{[k]}) < 2\tau$ for all i where γ is the value chosen in Section 4.2. To prove Lemma 4.3 for the case $|\mathbf{S}| \geq 3$, we use an inductive proof. We replace the distribution from $\mathcal{H}_{k,\delta}^n = \otimes_{i=1}^n \mathcal{H}_{k,\delta}$ to $\mathcal{I}_{k,\delta}^n = \otimes_{i=1}^n \mathcal{I}_{k,S_0,\delta}$ one component at a time. We will show the following lemma.

Lemma 4.8. For any $l \in [n]$,

$$\left| \mathbf{E}_{\substack{\otimes_{i=1}^{l-1} \mathcal{I}_{k, S_0, \delta} \\ \times \otimes_{i=l}^n \mathcal{H}_{k, \delta}}} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] - \mathbf{E}_{\substack{\otimes_{i=1}^l \mathcal{I}_{k, S_0, \delta} \\ \times \otimes_{i=l+1}^n \mathcal{H}_{k, \delta}}} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] \right| \leq \Delta_l, \quad (9)$$

where

$$\Delta_l = (2\tau)^{\gamma/2|\mathbf{S}|} 2^{|\mathbf{S}|} \sum_{S \in \mathbf{S}} \text{Inf}_l(T_{1-\gamma/2} f_S).$$

We first give the proof Lemma 4.3.

Proof of Lemma 4.3. From Lemma 4.8,

$$\begin{aligned} & \left| \mathbf{E}_{\mathcal{H}_{k, \delta}^n} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] - \mathbf{E}_{\mathcal{I}_{k, \delta}^n} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] \right| \\ & \leq (2\tau)^{\gamma/2|\mathbf{S}|} 2^{|\mathbf{S}|} \sum_{i \in [n]} \sum_{S \in \mathbf{S}} \text{Inf}_i(T_{1-\gamma/2} f_S) \\ & \leq |\mathbf{S}| (2\tau)^{\gamma/2|\mathbf{S}|} 2^{|\mathbf{S}|} / \gamma. \end{aligned}$$

We use Lemma 2.8 for the last inequality. By choosing $\tau = (\gamma\sqrt{\delta}/(|\mathbf{S}|2^{|\mathbf{S}|}))^{O(|\mathbf{S}|/\gamma)}$, the lemma follows. \square

In what follows, we show the proof of Lemma 4.8. We only show the proof for $l = 1$. The case for $l = 2, \dots, n$ is the same. Actually, we only use the fact that the unchanged product distribution $\mathcal{H}'_{k, \delta} = \bigotimes_{i=2}^n \mathcal{H}_{k, \delta}$ (when $l = 1$) have the property that its marginals on \mathbf{x}_S for $S \in \mathbf{S}$ are uniform. Here, We identify a set $S \subseteq [k]$ with a vector over \mathbb{F}_2^k , i.e., the i th element of the vector corresponding to S is 1 iff S contains i . The following lemma holds.

Lemma 4.9. Let $\mathfrak{X}_{\mathbf{S}}$ be the set of all $\mathbf{T} \subseteq \mathbf{S}$ satisfying the property that $S_0 \in \mathbf{T}$ and $\sum_{S \in \mathbf{T}} S = 0$ over \mathbb{F}_2^k . For any functions $f_S : \{-1, 1\} \rightarrow \mathbb{R}$,

$$\mathbf{E}_{\mathcal{H}_{k, \delta}} \left[\prod_{S \in \mathbf{S}} f_S(\mathbf{x}_S) \right] - \mathbf{E}_{\mathcal{I}_{k, S_0, \delta}} \left[\prod_{S \in \mathbf{S}} f_S(\mathbf{x}_S) \right] = -(1-\delta) \sum_{\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}} (-1)^{|\mathbf{T} \setminus [k]|} \prod_{S \in \mathbf{T}} \widehat{f}_S(\{1\}).$$

Proof. Taking Fourier expansion of each function in the left hand side, we get

$$\sum_{\{U_S\}_{S \in \mathbf{S}}} \prod_{S \in \mathbf{S}} \widehat{f}_S(U_S) \left(\mathbf{E}_{\mathcal{H}_{k, \delta}} \left[\prod_{S \in \mathbf{S}} \chi_{U_S}(\mathbf{x}_S) \right] - \mathbf{E}_{\mathcal{I}_{k, S_0, \delta}} \left[\prod_{S \in \mathbf{S}} \chi_{U_S}(\mathbf{x}_S) \right] \right).$$

From $\mathcal{H}_{k, \delta} = (1-\delta)\mathcal{H}_k + \delta\mathcal{N}_k$ and $\mathcal{I}_{k, S_0, \delta} = (1-\delta)\mathcal{I}_{k, S_0} + \delta\mathcal{N}_k$, it follows that

$$\mathbf{E}_{\mathcal{H}_{k, \delta}} \left[\prod_{S \in \mathbf{S}} \chi_{U_S}(\mathbf{x}_S) \right] - \mathbf{E}_{\mathcal{I}_{k, S_0, \delta}} \left[\prod_{S \in \mathbf{S}} \chi_{U_S}(\mathbf{x}_S) \right] = (1-\delta) \left(\mathbf{E}_{\mathcal{H}_k} \left[\prod_{S \in \mathbf{S}} \chi_{U_S}(\mathbf{x}_S) \right] - \mathbf{E}_{\mathcal{I}_{k, S_0}} \left[\prod_{S \in \mathbf{S}} \chi_{U_S}(\mathbf{x}_S) \right] \right).$$

Notice that \mathcal{H}_k and \mathcal{I}_{k,S_0} have the same marginal distribution except \mathcal{X}_{S_0} . Thus, to make the expression above non zero, \mathbf{S} must contain S_0 , and U_{S_0} must be $\{1\}$. In such a case,

$$\mathbf{E}_{\mathcal{I}_{k,S_0}} \left[\chi_{\{1\}}(\mathbf{x}_{S_0}) \prod_{S \in \mathbf{S} \setminus S_0} \chi_{U_S}(\mathbf{x}_S) \right] = \mathbf{E}_{\mathcal{I}_{k,S_0}} [\chi_{\{1\}}(\mathbf{x}_{S_0})] \mathbf{E}_{\mathcal{I}_{k,S_0}} \left[\prod_{S \in \mathbf{S} \setminus S_0} \chi_{U_S}(\mathbf{x}_S) \right] = 0.$$

It is not hard to see that $\mathbf{E}_{\mathcal{H}_k} [\prod_{S \in \mathbf{S}} \chi_{U_S}(\mathbf{x}_S)]$ becomes non zero iff there exists $\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}$ such that $U_S = \{1\}$ for $S \in \mathbf{T}$ and $U_S = \emptyset$ for $S \in \mathbf{S} \setminus \mathbf{T}$, and then the value is $-1^{|\mathbf{T} \setminus [k]|}$. \square

Lemma 4.10. *For any function $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$, $p \geq 3$ and $\gamma > 0$ satisfying $2 + 2\gamma \leq p$,*

$$\|T_{1-\gamma}f\|_p \leq \|T_{1-\gamma/2}f\|_2^{(2+\gamma)/p}.$$

Proof. It holds that

$$\|T_{1-\gamma}f\|_p = \mathbf{E}[\|T_{1-\gamma}f\|^p]^{1/p} \leq \mathbf{E}[\|T_{1-\gamma}f\|^{2+2\gamma}]^{1/p} = \|T_{1-\gamma}f\|_{2+2\gamma}^{(2+\gamma)/p}.$$

Here, we use $|T_{1-\gamma}f| \leq 1$. Since $1 - \gamma \leq \frac{1}{\sqrt{1+2\gamma}}$, applying the hypercontractive inequality (Lemma 2.6) to $\|T_{1-\gamma}f\|_{2+2\gamma}^{(2+\gamma)/p}$, we have

$$\|T_{1-\gamma}f\|_p \leq \|f\|_2^{(2+\gamma)/p}.$$

Using this,

$$\|T_{1-\gamma}f\|_p \leq \|T_{1-\gamma/2}T_{1-\gamma/2}f\|_p \leq \|T_{1-\gamma/2}f\|_2^{(2+\gamma)/p}.$$

\square

Proof of Lemma 4.8. Let $x = (x_1, x')$ where $x' = (x_2, \dots, x_n)$. We divide the Fourier expansion of $f_S(x)$ according to its dependence to x_1 , i.e., $f_S(x) = F_{S,\emptyset}(x') + x_1 F_{S,1}(x')$. Then, we note that $T_{1-\gamma}f_S(x) = T_{1-\gamma}F_{S,\emptyset}(x') + (1 - \gamma)x_1 T_{1-\gamma}F_{S,1}(x')$.

Using Lemma 4.9, the left hand side of (9) is equal to

$$\begin{aligned} & \left| \left[\mathbf{E}_{\mathcal{H}'_{k,\delta}} \left[\mathbf{E}_{\mathcal{H}_{k,\delta}} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma}f_S(\mathbf{x}) \right] - \mathbf{E}_{\mathcal{I}_{k,S_0,\delta}} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma}f_S(\mathbf{x}) \right] \right] \right] \right| \\ &= \left| (1 - \delta) \sum_{\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}} (1 - \gamma)^{|\mathbf{T}|} \mathbf{E}_{\mathcal{H}'_{k,\delta}} \left[\prod_{S \in \mathbf{T}} T_{1-\gamma}F_{S,1}(\mathbf{x}'_S) \right] \right| \\ &\leq (1 - \delta) \sum_{\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}} (1 - \gamma)^{|\mathbf{T}|} \mathbf{E}_{\mathcal{H}'_{k,\delta}} \left[\left| \prod_{S \in \mathbf{T}} T_{1-\gamma}F_{S,1}(\mathbf{x}'_S) \right| \right] \\ &\leq (1 - \delta) \sum_{\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}} (1 - \gamma)^{|\mathbf{T}|} \prod_{S \in \mathbf{T}} \|T_{1-\gamma}F_{S,1}\|_{|\mathbf{T}|}. \end{aligned} \quad (10)$$

The last step uses the Hölder's inequality and the fact that the marginal of $\mathcal{H}'_{k,\delta}$ on \mathbf{x}'_S is uniform for each S . Using Lemma 4.10,

$$\prod_{S \in \mathbf{T}} \|T_{1-\gamma} F_{S,1}\|_{|\mathbf{T}|} \leq \prod_{S \in \mathbf{T}} \|T_{1-\gamma/2} F_{S,1}\|_2^{(2+\gamma)/|\mathbf{T}|}$$

Here,

$$\|T_{1-\gamma/2} F_{S,1}\|_2^2 = \sum_{Q \ni 1} (1-\gamma/2)^{2|Q|} \widehat{f}_S(Q)^2 \leq \text{Inf}_1(T_{1-\gamma/2} f_S).$$

Thus, (10) is upper bounded by

$$\sum_{\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}} \prod_{S \in \mathbf{T}} \text{Inf}_1(T_{1-\gamma/2} f_S)^{(2+\gamma)/2|\mathbf{T}|}. \quad (11)$$

Since $|\mathbf{T}| \geq 2$ for all $\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}$, we have $\min_{S \in \mathbf{T}} \text{Inf}_1(T_{1-\gamma/2} f_S) \leq 2\tau$. Then,

$$\begin{aligned} (11) &\leq \sum_{\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}} (2\tau)^{\gamma/2|\mathbf{T}|} \prod_{S \in \mathbf{T}} \text{Inf}_1(T_{1-\gamma/2} f_S)^{1/|\mathbf{T}|} \\ &\leq (2\tau)^{\gamma/2|\mathbf{S}|} \sum_{\mathbf{T} \in \mathfrak{X}_{\mathbf{S}}} \sum_{S \in \mathbf{T}} \text{Inf}_1(T_{1-\gamma/2} f_S) \quad (\text{using the AM-GM inequality}) \\ &\leq (2\tau)^{\gamma/2|\mathbf{S}|} 2^{|\mathbf{S}|} \sum_{S \in \mathbf{S}} \text{Inf}_1(T_{1-\gamma/2} f_S) \quad (\text{since } |\mathfrak{X}_{\mathbf{S}}| \leq 2^{|\mathbf{S}|}). \end{aligned}$$

□

4.4 Proof of Lemma 4.4

In this section, we prove the following.

Lemma 4.11 (Restatement of Lemma 4.4). *For any $S_0 \in \mathbf{S}$, γ, δ , it holds that*

$$\left| \mathbf{E}_{\mathcal{I}_{k,S_0,\delta}^n} \left[\prod_{S \in \mathbf{S}} T_{1-\gamma} f_S(\mathbf{x}_S) \right] \right| \leq \sqrt{\delta}.$$

We need the following lemma.

Lemma 4.12. *For any $i \in [n]$ and $S_0 \in \mathbf{S}$, it holds that*

$$\rho(\mathcal{X}_{S_0}^i, \prod_{S \in \mathbf{S} \setminus \{S_0\}} \mathcal{X}_S^i; \mathcal{I}_{k,S_0,\delta}) \leq \sqrt{\delta}.$$

Proof. Let $f : \mathcal{X}_{S_0}^i \rightarrow \mathbb{R}$ be any function with $\mathbf{E}_{\mathcal{I}_{k,S_0,\delta}}[f] = 0$, $\mathbf{E}_{\mathcal{I}_{k,S_0,\delta}}[f^2] \leq 1$. Also, let $g : \prod_{S \in \mathbf{S} \setminus \{S_0\}} \mathcal{X}_S^i \rightarrow \mathbb{R}$ be any functions with $\mathbf{E}_{\mathcal{I}_{k,S_0,\delta}}[g] = 0$, $\mathbf{E}_{\mathcal{I}_{k,S_0,\delta}}[g^2] \leq 1$. From the latter property we get,

$$1 \geq \mathbf{E}_{\mathcal{I}_{k,S_0,\delta}}[g^2] = (1-\delta) \mathbf{E}_{\mathcal{I}_{k,S_0}}[g^2] + \delta \mathbf{E}_{\mathcal{N}_k}[g^2] \geq \delta \mathbf{E}_{\mathcal{N}_k}[g^2]. \quad (12)$$

We now observe that

$$\begin{aligned} \mathbf{E}_{\mathcal{I}_{k,S_0,\delta}} [fg] &= (1-\delta) \mathbf{E}_{\mathcal{I}_{k,S_0}} [fg] + \delta \mathbf{E}_{\mathcal{N}_k} [fg] \\ &\leq (1-\delta) \mathbf{E}_{\mathcal{I}_{k,S_0}} [f] \mathbf{E}_{\mathcal{I}_{k,S_0}} [g] + \delta \sqrt{\mathbf{E}_{\mathcal{N}_k} [f^2]} \sqrt{\mathbf{E}_{\mathcal{N}_k} [g^2]}. \end{aligned}$$

We have equality for the first term because the marginals of $\mathcal{I}_{k,S_0,\delta}$ on $\mathcal{X}_{S_0}^i$ and $\prod_{S \in \mathbf{S} \setminus \{S_0\}} \mathcal{X}_S^i$ are independent. We used Cauchy-Schwarz on the second term. Since $\mathbf{E}_{\mathcal{I}_{k,S_0}} [f] = 0$ and From (12),

$$\delta \sqrt{\mathbf{E}_{\mathcal{N}_k} [f^2]} \sqrt{\mathbf{E}_{\mathcal{N}_k} [g^2]} \leq \delta \sqrt{1} \sqrt{1/\delta}.$$

Hence,

$$\mathbf{E}_{\mathcal{I}_{k,S_0,\delta}} [fg] \leq \sqrt{\delta}.$$

This completes the proof. \square

Proof of Lemma 4.4. From Lemma 4.12 and Proposition 2.16,

$$\rho(\mathcal{X}_{S_0}, \prod_{S \in \mathbf{S} \setminus \{S_0\}} \mathcal{X}_S; \mathcal{I}_{k,S_0,\delta}^n) \leq \sqrt{\delta}.$$

Since $\mathbf{E}_{\mathcal{I}_{k,S_0,\delta}^n} [T_{1-\gamma} f_S]$ has mean 0 and variance at most 1 for any S , the lemma immediately follows. \square

5 Discussion

Unfortunately, our Long Code test does not immediately imply a new PCP characterization of NP. A Long Code test without consistency checks is most easily composed with Unique Games [29] as an outer verifier in PCP constructions. However, since Unique Games cannot have perfect completeness, an obvious approach that combines our tester and Unique Games does not imply a new PCP construction. A variant of Unique Games, which is called d -to-1 Games [29], is conjectured to have perfect completeness. Though it seems hard to combine our tester and d -to-1 Games, we hope that it is possible and we obtain a new conditional PCP construction.

Recent work After the extended abstract of this paper was published, several papers have given new unconditional PCP constructions under the standard assumption that $P \neq NP$ instead of Khot's Unique Games conjecture or d -to-1 Games conjecture. In [25], Håstad constructed non-adaptive 3-query PCPs with perfect completeness and soundness $s = 5/8$, removing Khot's d -to-1 conjecture in [37]. In [15], Chan obtained non-adaptive q -query PCPs with almost perfect

completeness and soundness $s = 2q/2^q$, removing Khot’s Unique Games conjecture in [43]. In [28], Huang showed non-adaptive q -query PCPs with perfect completeness and soundness $s = 2^{O((q \log q)^{1/3})}/2^q$, improving the previous best result due to [26]. The techniques introduced by these papers will probably be useful for constructing PCPs with perfect completeness based on our Long Code test.

Acknowledgments

We would like to thank three reviewers for their careful reading and comments on our paper. The comments helped us a lot in improving the paper.

ST is supported in part by MEXT KAKENHI Grant Number 24106003 and JSPS KAKENHI Grant Numbers 23700015, 25240002. YY is supported in part by JSPS Grant-in-Aid for Research Activity Start-up (24800082), MEXT Grant-in-Aid for Scientific Research on Innovative Areas (24106003), and JST, ERATO, Kawarabayashi Large Graph Project.

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