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Kenichi Morita

Theory of Reversible Computing



Kenichi Morita Professor Emeritus, Hiroshima University Hiroshima, Japan

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Preface

A reversible computing system is a "backward deterministic" system such that every state of the system has at most one predecessor. Hence, there is no pair of distinct states that go to the same state. Though its definition is so simple, it is closely related to physical reversibility. The study of reversible computing originated from an investigation of energy dissipation in reversible and irreversible computing systems. Rolf Landauer investigated the relation between reversibility in computing and reversibility in physics in his paper "Irreversibility and heat generation in the computing process" (IBM J. Res. Dev., Vol. 5, pp. 183–191, 1961). He pointed out that an irreversible logical operation inevitably causes energy dissipation in the computing system. Since then, reversible computing has been studied in relation to physical reversibility. Besides the problem of energy dissipation in computing, it is important to know how reversibility can be effectively utilized in computing. This is because future computing devices will surely be implemented directly by physical phenomena in the nano-scale level, and reversibility is one of the fundamental microscopic physical laws of Nature. For this purpose, various models of reversible computing have been proposed and investigated till now.

In this book, reversible computing is studied from the standpoint of the theory of automata and computing. We deal with various reversible computing models belonging to several different levels, which range from a microscopic level to a macroscopic one. They are reversible physical models, reversible logic elements, reversible functional modules composed of logic elements, reversible computing systems such as Turing machines, cellular automata, and others. The purpose of this book is to clarify how computation can be carried out efficiently and elegantly in these reversible computing models. We shall see that even very simple reversible systems have computational universality in spite of the constraint of reversibility. We shall also see various reversible systems in different levels are related each other, i.e., a reversible system in a higher level can be constructed out of those in a lower level. Moreover, the construction methods are often very unique and different from those in the traditional methods. Thus, these computing models as well as the designing methods will give us new insights for future computing systems.

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This book is not a comprehensive textbook on reversible computing, but describes mainly the results shown in the papers by myself and my colleagues, which were published between 1989 and 2017. In this book, readers will see how the world of reversible computing works, how reversible computing systems are constructed out of simple reversible primitives, how they are different from the traditional computing systems, and how they can be computationally universal. In fact, we shall see even very simple reversible systems have high capability of computing, and thus reversibility is not a constraint, but a useful property for computing.

This book consists of 14 chapters. Chapter 1 is an introduction. In Chaps. 2–4, reversible logic elements for constructing reversible machines are investigated. In Chaps. 5–8, reversible Turing machines, a standard model of reversible computing systems, are studied. In Chap. 9, some other reversible computing models are investigated. In Chaps. 10–14, reversible cellular automata, a spatiotemporal model of reversible dynamical systems, are studied.

There is no prerequisite knowledge to read this book besides some basics on logic, discrete mathematics and formal languages. But, it is preferable to have some knowledge on the theories of automata and computing. Fortunately, the framework of reversible computing itself is very simple. Therefore, in many cases, readers can easily understand the basic function and the structure of each such system. However, its behavior can be very complex even if the structure of the system is simple. Hence, sometimes, it becomes quite difficult to follow its behavior by using only paper and pencil. In some of these cases, readers can find files in the References that contain computer simulation results of such reversible systems.

More than 50 years have passed since Landauer's paper appeared. Thus, the history of reversible computing is relatively long. But, it is still developing, and there remain many problems to be investigated. Also, even at present, it is not so clear which results will become practically useful in the future. However, the world of reversible computing will lead readers to the new ways of thinking that cannot be found in the traditional design methodologies for computing systems. I hope the theory of reversible computing will stimulate readers' interest, and open new vistas for future computing systems.

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September 2017

Kenichi Morita

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Acronyms

CA cellular automaton
CM counter machine
CTS cyclic tag system

CTSH cyclic tag system with halting condition

ECA elementary cellular automaton

ETPCA elementary triangular partitioned cellular automaton

GoL Game of Life

ID instantaneous description

IDTM irreversible deterministic Turing machine INTM irreversible nondeterministic Turing machine

 $\mathscr{L}(\mathscr{A})$ class of languages accepted by the class of automata \mathscr{A}

LBA linear-bounded automaton MFA multi-head finite automaton

m-TS *m*-tag system

PCA partitioned cellular automaton RCA reversible cellular automaton RCM reversible counter machine

RDTM reversible deterministic Turing machine

RE rotary element

RETPCA reversible elementary triangular partitioned cellular automaton

RLEM reversible logic element with memory
RMFA reversible multi-head finite automaton
RNTM reversible nondeterministic Turing machine
RPCA reversible partitioned cellular automaton

RSM reversible sequential machine RTM reversible Turing machine SM sequential machine TM Turing machine

TM(S(n)) S(n) space-bounded Turing machine

UTM universal Turing machine

URTM universal reversible Turing machine