SHADOWS AND ISOPERIMETRY UNDER THE SEQUENCE–SUBSEQUENCE RELATION

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1. Introduction and results

One of the basic results in extremal set theory was discovered in [1] and rediscovered in [2]: For a given number of k-element subsets of an n-set the shadow, that is, the set of (k-1)-element subsets contained in at least one of the specified k-element subsets, is minimal, if the k-element subsets are chosen as an initial segment in the squashed order (see [10]; called colex order in [11]), that is, a kelement subset A precedes a k-element subset B, if the largest element in $A \triangle B$ is in B. A closely related result was discovered in [3] and rediscovered in [5]: For a given number $u \in [0, 2^n]$ of arbitrary subsets of an n-set the "Hamming distance 1"-boundary is minimal for the initial segment of size u, also called in short "u-th initial segment", in the H-order (of [3]), that is, if one chooses all subsets of cardinality less than n-k (k suitable) and all remaining subsets of cardinality n-k, whose complements are in the initial segment of the squashed order.

In this paper we consider sequences and subsequences rather than sets and subsets.

The basic objects are $\mathcal{X}^n = \prod_{i=1}^n \mathcal{X}$ for $\mathcal{X} = \{0,1\}$ and $n \in \mathbb{N}$, and operations of deletion $\bigtriangledown_i, \bigtriangledown$ and of insertion $\bigtriangleup_i, \bigtriangleup$. Here \bigtriangledown_i (resp. \bigtriangleup_i) means that letter i (i = 0, 1) is deleted (resp. inserted) and \bigtriangledown (resp. \bigtriangleup) means the deletion (resp. insertion) of any letter.

So for $A \subset \mathcal{X}^n$ we get the down shadow

(1.1) $\nabla A = \{x^{n-1} \in \mathcal{X}^{n-1} : x^{n-1} \text{ is subsequence of some } a^n \in A\}$ and the *up shadow*

(1.2)
$$\triangle A = \{x^{n+1} \in \mathcal{X}^{n+1} : \text{ some subsequence of } x^{n+1} \text{ is in } A\}.$$

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(1.3)
$$\nabla A = \nabla_0 A \cup \nabla_1 A$$

and analogously

(1.4)
$$\triangle A = \triangle_0 A \cup \triangle_1 A.$$

We describe now our results.

A. Shadows for fixed level and specific letter

The ℓ -th level is the set of sequences (or words)

We consider sets $B \subset \mathcal{X}_{n-k}^n$ of cardinality $v, 0 \le v \le {n \choose k}$, and their shadows $\nabla_0 B, \Delta_1 B$ (the other shadows can be estimated similarly by symmetry).

The unique binomial representation of v is

(1.6)
$$v = \binom{a_k}{k} + \binom{a_{k-1}}{k-1} + \ldots + \binom{a_s}{s}$$

(with $a_k > a_{k-1} > \ldots > a_s \ge s \ge 1$).

Whereas Katona used in [2] and also in [6] the function F:

(1.7)
$$F(k,v) = \binom{a_k}{k-1} + \binom{a_{k-1}}{k-2} + \ldots + \binom{a_s}{s-1},$$

we introduce and need here the functions $\stackrel{\bigtriangledown}{F}$ and $\stackrel{\bigtriangleup}{F}$, which play the analogue roles for the new shadow problems:

(1.8)
$$\overrightarrow{F}(k,v) = \binom{a_k-1}{k-1} + \binom{a_{k-1}-1}{k-2} + \dots + \binom{a_s-1}{s-1}$$

 and

(1.9)
$$\stackrel{\Delta}{F}(k,v) = \binom{a_k+1}{k} + \binom{a_{k-1}+1}{k-1} + \ldots + \binom{a_s+1}{s}.$$

Theorem 1. For all $B \subset \mathcal{X}_{n-k}^n$ with |B| = v

- (i) $|\bigtriangledown_0 B| \ge \stackrel{\bigtriangledown}{F}(k, v),$ (ii) $|\bigtriangleup_1 B| \ge \stackrel{\bigtriangleup}{F}(k, v),$
- (11) $|\Delta_1 B| \ge F$ (k) and
- (iii) both bounds are optimal.

B. Shadows of arbitrary sets under deletion of any letter

For any integer $u \in [0, 2^n]$ we use the unique binomial representation

(1.10)
$$u = \binom{n}{n} + \ldots + \binom{n}{k+1} + \binom{\alpha_k}{k} + \ldots + \binom{\alpha_t}{t}$$

(with $n>\alpha_k>\ldots>\alpha_t\geq t\geq 1)$ and observe that for an initial H-order segment S with |S|=u

$$|\nabla S| = \binom{n-1}{n-1} + \binom{n-1}{n-2} + \dots + \binom{n-1}{k} + \binom{\alpha_k - 1}{k-1} + \dots + \binom{\alpha_t - 1}{t-1}$$

$$(1.11) = \overset{\nabla}{G}(n, u), \quad \text{say.}$$

Theorem 2. For every $A \subset \mathcal{X}^n$, $|\nabla A| \ge \overset{\nabla}{G}(n, |A|)$ and the bound is achieved by the |A|-th initial segment in H-order.

This result was first obtained by D. E. Daykin and T. N. Danh [8]. We are grateful to David for his dramatic story about the complexity of their (first) proof. It gave us the impetus to (quickly) find a proof with fairly lengthy calculations with binomial coefficients. Subsequently Daykin–Danh gave also another proof, which can be found in the collection [9]. Then we gave a very "short proof" in [9] based on Lemma 6 of [6] and our inequality (2.5) below. Unfortunately, as was kindly pointed out by David, the original proof of (2.5) has an error in equation (6) of [9].

C. Shadows of arbitrary sets under insertion of any letter

For u in the representation (1.10) we define

(1.12)
$$\stackrel{\triangle}{G}(n,u) = \binom{n+1}{n+1} + \binom{n+1}{n} + \ldots + \binom{n+1}{k+1} + \binom{\alpha_k+1}{k} + \ldots + \binom{\alpha_t+1}{t}.$$

Theorem 3. *For every $A \subset \mathcal{X}^n$, $|\Delta A| \ge G(n, |A|)$, and the bound is achieved by the |A|-th initial segment in H-order.

Remarks.

- 1. It must be emphasized that the *H*-order minimizes simultaneously both, the lower and the upper shadows. There is no such phenomenon in the Boolean lattice for "Kruskal-Katona"-type shadows. It has immediate consequences for isoperimetric problems.
- 2. Theorem 1 can be derived easily from Theorems 2 and 3 like Kruskal-Katona's Theorem from Harper's Theorem.

D. Two isoperimetric inequalities

It has been emphasized in [7] that isoperimetric inequalities in discrete metric spaces are fundamental principles in combinatorics. The goal is to minimize the union of a specified number of spheres of constant radius. We speak of an isoperimetric inequality, if this minimum is assumed for a set of sphere-centers, which themselves form a sphere (or quasi-sphere, if numbers do not permit a sphere).

For any $A \subset \mathcal{X}^* = \bigcup_{n=0}^{\infty} \mathcal{X}^n$ and any distance d we define (the union of spheres

of radius r)

(1.13)
$$\Gamma_d^r(A) = \left\{ x^{n'} \in \mathcal{X}^* : d(x^{n'}, a^n) \le r \text{ for some } a^n \in A \right\}.$$

A prototype of a discrete isoperimetric inequality is the one discovered in [3], rediscovered in [5], and proved again in [6]. Here d equals the Hamming distance d_H and is defined on $\mathcal{X}^n \times \mathcal{X}^n$.

We recall the result. For

(1.14)
$$G(n,u) = \binom{n}{n} + \binom{n}{n-1} + \ldots + \binom{n}{k} + \binom{\alpha_k}{k-1} + \ldots + \binom{\alpha_t}{t-1}$$

and any $A \subset \mathcal{X}^n$

$$(1.15) \qquad \qquad |\Gamma_{d_H}^r(A)| \ge G(n,|A|)$$

and the bound is achieved by the |A|-th initial segment in *H*-order (this is a sphere of radius k, if $|A| = \sum_{j=0}^{k} {n \choose j}$).

^{*} A referee kindly pointed out to us that the equivalence of Theorem 2 and Theorem 3 can be derived with a theorem in "Variational principle in discrete extremal problems" by Bezrukov (Reihe Informatik Bericht tr-ri-94-152, Universität-GH-Paderborn).

We define now two distances, θ and δ , in \mathcal{X}^* . For $x^m, y^{m'} \in \mathcal{X}^*$ $\theta(x^m, y^{m'})$ counts the minimal number of insertions and deletions which transform one word into the other. $\delta(x^m, y^{m'})$ counts the minimal number of operations, if also exchanges of letters are allowed. Thus $\delta(x^m, y^{m'}) \leq \theta(x^m, y^{m'})$.

Now observe that from (1.15) and our Theorems 2 and 3, we get immediately two inequalities.

Corollary 1. For $A \subset \mathcal{X}^n$

(i) $\Gamma^{1}_{\theta}(A) \leq \overset{\nabla}{G}(n, |A|) + \overset{\Delta}{G}(n, |A|),$

(ii)
$$\Gamma^1_{\delta}(A) \leq \overset{\checkmark}{G}(n,|A|) + \overset{\leftrightarrow}{G}(n,|A|) + G(n,|A|),$$

and both bounds are achieved by the |A|-th initial segment in H-order.

Moreover, in Theorem 4 of Section 6 we have established those inequalities for every radius r. The exact formulation and the proof require a technical setup.

2. Auxiliary results A. Numerical inequalities

While working on [7] Gyula Katona drew attention to the approach of Eckhoff-Wegner [4] to prove Kruskal-Katona via the following inequality for F, defined in (1.7).

Lemma 1 (see [4]). For k > 1, $v \le v_0 + v_1$,

(2.1)
$$F(k,v) \le \max(v_0, F(k,v_1)) + F(k-1,v_0).$$

In fact, he used this idea also in his proof of the isoperimetric inequality for the Hamming space. He just had to establish the corresponding inequality for G, defined in (1.14).

Lemma 2 (Lemma 6 of [6]). If $0 \le u_1 \le u_0$ and $u \le u_0 + u_1$, then

(2.2)
$$G(n,u) \leq \max(u_0, G(n-1,u_1)) + G(n-1,u_0).$$

The discoveries in the present paper are similar inequalities for $\overset{\bigtriangledown}{F}, \overset{\bigtriangleup}{F}, \overset{\bigtriangledown}{G}$, and $\overset{\bigtriangleup}{G}$ (defined in (1.8), (1.9), (1.11), and (1.12)), which for cardinalities of shadows resp. boundaries considered describe their values for segments in the *H*-order.

We state first the inequalities for F. They are proved in the same way as those for G below.

 $\stackrel{\bigtriangledown}{F}$ -inequality: For k > 1, if $v \le v_0 + v_1$ and $v_0 < \stackrel{\bigtriangledown}{F}(k, v)$, then

(2.3)
$$\stackrel{\nabla}{F}(k,v) \leq \stackrel{\nabla}{F}(k,v_1) + \stackrel{\nabla}{F}(k-1,v_0).$$

 $\stackrel{\triangle}{F}$ -inequality: For k > 1, if $v \leq v_0 + v_1$, then

(2.4)
$$\overset{\bigtriangleup}{F}(k,v) \leq \max\left(v_0 + v_1, \overset{\bigtriangleup}{F}(k,v_1)\right) + \overset{\bigtriangleup}{F}(k-1,v_0).$$

Next we derive the inequalities for G.

 $\stackrel{\bigtriangledown}{G}$ -inequality: If $w_1 \leq w_0 < \stackrel{\bigtriangledown}{G}(n, w)$ and $w \leq w_0 + w_1$, then

(2.5)
$$\overset{\nabla}{G}(n,w) \leq \overset{\nabla}{G}(n-1,w_0) + \overset{\nabla}{G}(n-1,w_1).$$

 $\stackrel{\triangle}{G}$ -inequality: If $0 \le u_1 \le u_0, \ u \le u_0 + u_1$, then

(2.6)
$$\overset{\bigtriangleup}{G}(n,u) \leq \max\left(u_0 + u_1, \overset{\bigtriangleup}{G}(n-1,u_1)\right) + \overset{\bigtriangleup}{G}(n-1,u_0).$$

Proofs. From the definitions of the numerical functions we have

$$G(n, u) + u \stackrel{\triangle}{=} (n, u)$$
 for u as in (1.10)

and the equivalence of (2.2) and (2.6) immediately follows.

Next we show (2.5). For u as in (1.10) denote by $\ell_n(u)$ and $r_n(u)$ the smallest j with $\alpha_j > j$ and the number of i's with $\alpha_i = i$, respectively.

Let

(2.7)
$$\overline{u}(n-1) \stackrel{\Delta}{=} u - \stackrel{\nabla}{G}(n,u) = \binom{n-1}{n-1} + \dots + \binom{n-1}{k+1} + \binom{\alpha_k - 1}{k} + \dots + \binom{\alpha_{\ell_n(u)} - 1}{\ell_n(u)}.$$

By (1.11) and (1.14)

$$\stackrel{\nabla}{G}(n,u) = \binom{n-1}{n-1} + \ldots + \binom{n-1}{k} + \binom{\alpha_k - 1}{k-1} + \ldots + \binom{\alpha_{\ell_n(u)} - 1}{\ell_n(u) - 1} + r_n(u)$$

(2.8)
$$= G(n-1,\overline{u}(n-1)) + r_n(u).$$

Moreover by the binomial coefficient representation

(2.9)
$$u+1 = \begin{cases} \binom{n}{n} + \ldots + \binom{n}{k+1} + \binom{\alpha_k}{k} + \ldots + \binom{\ell_n(u)}{\ell_n(u)-1} & \text{if } \alpha_t = t = 1 \\ \binom{n}{n} + \ldots + \binom{n}{k+1} + \binom{\alpha_k}{k} + \ldots + \binom{\alpha_t}{t} + \binom{t-1}{t-1} & \text{otherwise} \end{cases}$$

(1.11) implies

(2.10)
$$\overset{\nabla}{G}(n, u+1) = \begin{cases} \overset{\nabla}{G}(n, u) & \text{if } \alpha_t = t = 1\\ \overset{\nabla}{G}(n, u) + 1 & \text{otherwise.} \end{cases}$$

By the definition of binomial coefficient representation, $\overline{u}(n-1)$ in (2.7) is non-decreasing in u for fixed n (c.f. (2.9)).

For w_1 , w_0 and w with

$$(2.11) w_1 \le w_0 < \overset{\bigtriangledown}{G} (n, u)$$

 and

$$(2.12) w \le w_0 + w_1,$$

we let $w^* = w$, if $r_n(w) = 0$, and otherwise

$$w^* = \binom{n}{n} + \ldots + \binom{n}{k'+1} + \binom{\beta_{k'}}{k'} + \ldots + \binom{\beta_{\ell_n(w)}}{\ell_n(w)} + \binom{\ell_n(w)}{\ell_n(w)-1}$$

$$(2.13) = w - r_n(w) + \ell_n(w),$$

if the representation of w is

$$w = \binom{n}{n} + \dots + \binom{n}{k'+1} + \binom{\beta_{k'}}{k'} + \dots + \binom{\beta_s}{s}$$
$$= \binom{n}{n} + \dots + \binom{n}{k'+1} + \binom{\beta_{k'}}{k'} + \dots + \binom{\beta_{\ell_n(w)}}{\ell_n(w)} + \binom{\ell_n(w)-1}{\ell_n(w)-1} + \dots + \binom{\ell_n(w)-r_n(w)}{\ell_n(w)-r_n(w)}.$$

Write

(2.14)
$$w_0^* = w_0 + (w^* - w)$$
 and $w_1^* = w_1$.

Then by the definitions of w^* , w_0^* , w_1^* , and (2.10) (used repeatedly),

(2.15)
$$r_n(w^*) = 0,$$

(2.16)
$$\stackrel{\nabla}{G}(n, w^*) = \stackrel{\nabla}{G}(n, w) + (w^* - w) - 1, \text{ if } w^* \neq w,$$

 and

(2.17)
$$\overset{\nabla}{G}(n-1,w_0^*) \leq \overset{\nabla}{G}(n-1,w_0) + (w^*-w) - \tau(w_0^*),$$

where $\tau(w_0^*) = 1$, if $r_{n-1}(w_0^*) = 0$ and $w_0^* \neq w_0$, and $\tau(w_0^*) = 0$ otherwise. So, by (2.11), (2.14), and (2.16)

(2.18)
$$w_1^* \le w_0^* \le \overset{\bigtriangledown}{G}(n, w^*),$$

which with (2.7), (2.8) and (2.15) yields

$$\overline{w}_{0}^{*}(n-2) + G(n-2, \overline{w}_{0}^{*}(n-2)) = w_{0}^{*} - \overleftarrow{G}(n-1, w_{0}^{*}) + G(n-2, \overline{w}_{0}^{*}(n-2))$$
$$= w_{0}^{*} - r_{n-1}(w_{0}^{*}) \le w_{0}^{*} \le \overleftarrow{G}(n, w^{*})$$
$$= G(n-1, \overline{w}^{*}(n-1)).$$

Moreover, by the first inequality in (2.18) and the monotonicity of $\overline{u}(n-1)$ (as a function of u),

(2.20)
$$\overline{w}_1^*(n-2) \le \overline{w}_0^*(n-2).$$

Now we assume that (2.5) does not hold and derive a contradiction. With (2.12) we obtain

(2.21)
$$w - \overset{\nabla}{G}(n,w) < w_0 - \overset{\nabla}{G}(n-1,w_0) + w_1 - \overset{\nabla}{G}(n-1,w_1).$$

When $w^* \neq w$ then by (2.7) and (2.16) the LHS of (2.22) is $w - \overleftarrow{G}(n, w^*) + (w^* - w) - 1 = \overline{w}^*(n-1) - 1$ and by (2.7), (2.14) and (2.17) the RHS of (2.22) is not bigger than $w_0 - \overleftarrow{G}(n-1, w_0^*) + (w^* - w) - \tau(w_0^*) + \overline{w}_1^*(n-2) \leq \overline{w}_0^*(n-2) + \overline{w}_1^*(n-2)$. Thus we have

(2.22)
$$\overline{w}^*(n-1) \le \overline{w}^*_0(n-2) + \overline{w}^*_1(n-2)$$

By our notation in (2.7), (2.21) certainly implies (2.22), when $w^* = w$ (so $w^*_0 = w_0$).

Finally, with (2.19), (2.20), and (2.22) we obtain from Lemma 2,

$$(2.23) \qquad G\left(n-1,\overline{w}^*(n-1)\right) \leq G\left(n-2,\overline{w}_0^*(n-2)\right) + G\left(n-2,\overline{w}_1^*(n-2)\right).$$

This implies (2.5) (a contradiction to our assumption), because by (2.8), (2.15), and (2.16) the LHS of (2.23) is

$$\stackrel{
abla}{G}(n,w^*) = egin{cases} \stackrel{
abla}{G}(n,w) + w^* - w - 1 & ext{if } w
eq w^*, \ \stackrel{
abla}{G}(n,w) & ext{if } w = w^* ext{ (note } r_n(w) = 0) \end{cases}$$

and by (2.8), (2.14), and (2.17) the RHS of (2.23) is

$$\begin{split} & \left\{ \begin{matrix} \nabla G \left(n-1, w_0^* \right) + \overleftarrow{G} \left(n-1, w_1^* \right) - \left(r_{n-1} (w_0^*) + r_{n-1} (w_1^*) \right) \\ \leq & \left\{ \begin{matrix} \nabla G \left(n-1, w_0 \right) + \overleftarrow{G} \left(n-1, w_1 \right) + w^* - w - \tau (w_0^*) - \left(r_{n-1} (w_0^*) + r_{n-1} (w_1^*) \right) \\ \leq & \left\{ \begin{matrix} \nabla G \left(n-1, w_0 \right) + \overleftarrow{G} \left(n-1, w_1 \right) + \left\{ \begin{matrix} w^* - w - 1 & \text{if } w \neq w^* \\ 0 & \text{if } w = w^*, \end{matrix} \right\} \end{matrix} \right\} \end{split}$$

B. A calculus of iterative applications for
$$\overset{\bigtriangledown}{G}, \overset{\bigtriangleup}{G}$$
, and G

We present here a rather technical result (Lemma 4 below), which is needed only for the proof of Theorem 4. Recall that for $u, 1 \le u \le 2^n$,

$$u = \binom{n}{n} + \ldots + \binom{n}{k+1} + \binom{\alpha_k}{k} + \binom{\alpha_{k-1}}{k-1} + \ldots + \binom{\alpha_t}{t},$$

$$\overrightarrow{G}(n, u) = \binom{n-1}{n-1} + \ldots + \binom{n-1}{k+1} + \binom{n-1}{k} + \binom{\alpha_k}{k-1} + \ldots + \binom{\alpha_t-1}{t-1},$$

$$G(n, u) = \binom{n}{n} + \ldots + \binom{n}{k+1} + \binom{n}{k} + \binom{\alpha_k}{k-1} + \ldots + \binom{\alpha_t}{t-1},$$

and

$$\overset{\Delta}{G}(n,u) = \binom{n+1}{n+1} + \ldots + \binom{n+1}{k+1} + \binom{\alpha_k+1}{k} + \ldots + \binom{\alpha_t+1}{t}.$$

All these functions are increasing in u and they transform binomial representations into binomial representations. This makes it easy to apply them repeatedly.

We notice that the representation of $\overset{\nabla}{G}(n,u)$ may be not unique, due to the appearance of the term $\begin{pmatrix} 0\\0 \end{pmatrix}$. However, it causes no difficulties to apply the functions, because both representations (if they exist) always give the same result, when $\overset{\nabla}{G}$, G or $\overset{\Delta}{G}$ are applied. More specifically, the non-uniqueness happens only when $\alpha_t = t = 1$ in (1.10), and with the notation $\ell_n(u) = \ell$ (say) in the proof of (2.5),

$$\overset{\bigtriangledown}{G}(n,u) = \binom{n-1}{n-1} + \ldots + \binom{n-1}{k+1} + \binom{n-1}{k} + \binom{\alpha_k - 1}{k-1} + \ldots \\ + \binom{\alpha_\ell - 1}{\ell-1} + \binom{\ell-2}{\ell-2} + \ldots + \binom{1}{1} + \binom{0}{0}$$

$$= \binom{n-1}{n-1} + \ldots + \binom{n-1}{k} + \binom{\alpha_k - 1}{k-1} + \ldots + \binom{\alpha_\ell - 1}{\ell-1} + \binom{\ell-1}{\ell-2} \stackrel{\triangle}{=} v \text{ say.}$$

For the first representation of v

$$\begin{split} & \stackrel{\nabla}{G}(n-1,v) = \\ & \begin{pmatrix} n-2\\ n-2 \end{pmatrix} + \ldots + \begin{pmatrix} n-2\\ k-1 \end{pmatrix} + \begin{pmatrix} \alpha_k - 2\\ k-2 \end{pmatrix} + \ldots + \begin{pmatrix} \alpha_\ell - 2\\ \ell-2 \end{pmatrix} + \begin{pmatrix} \ell-3\\ \ell-3 \end{pmatrix} + \ldots + \begin{pmatrix} 0\\ 0 \end{pmatrix}, \\ & G(n-1,v) = \\ & \begin{pmatrix} n-1\\ n-1 \end{pmatrix} + \ldots + \begin{pmatrix} n-1\\ k-1 \end{pmatrix} + \begin{pmatrix} \alpha_k - 1\\ k-2 \end{pmatrix} + \ldots + \begin{pmatrix} \alpha_\ell - 1\\ \ell-2 \end{pmatrix} + \begin{pmatrix} \ell-2\\ \ell-3 \end{pmatrix} + \ldots + \begin{pmatrix} 1\\ 0 \end{pmatrix} \\ & \text{and,} \\ & \stackrel{\Delta}{\frown} \\ & \downarrow \\ &$$

$$\overset{\Delta}{G}(n-1,v) = \binom{n}{n} + \ldots + \binom{n}{k} + \binom{\alpha_k}{k-1} + \ldots + \binom{\alpha_\ell}{\ell-1} + \binom{\ell-1}{\ell-2} + \ldots + \binom{2}{1} + \binom{1}{0},$$

and for the second representation of v,

$$\overrightarrow{G}(n-1,v) = \binom{n-2}{n-2} + \dots + \binom{n-2}{k-1} + \binom{\alpha_k-2}{k-2} + \dots + \binom{\alpha_\ell-2}{\ell-2} + \binom{\ell-2}{\ell-3},$$

$$G(n-1,v) = \binom{n-1}{n-1} + \dots + \binom{n-1}{k-1} + \binom{\alpha_k-1}{k-2} + \dots + \binom{\alpha_\ell-1}{\ell-2} + \binom{\ell-1}{\ell-3},$$

and

$$\overset{\Delta}{G}(n-1,v) = \binom{n}{n} + \ldots + \binom{n}{k} + \binom{\alpha_k}{k-1} + \ldots + \binom{\alpha_\ell}{\ell-1} + \binom{\ell}{\ell-2},$$

They really have the same values.

For two functions $\phi, \psi: \mathbb{N} \to \mathbb{N}$ we write $\psi(\phi(\cdot))$ as $\psi \circ \phi(\cdot)$ and thus we can define

(2.24)
$$\overrightarrow{G}^{\circ p}(n,\cdot) = \overrightarrow{G}(n-p+1,\cdot) \circ \overrightarrow{G}(n-p+2,\cdot) \circ \ldots \circ \overrightarrow{G}(n,\cdot),$$

(2.25)
$$G^{\circ q}(n,\cdot) = G(n,\cdot) \circ G(n,\cdot) \circ \ldots \circ G(n,\cdot),$$

and

(2.26)
$$\overset{\triangle^{\circ s}}{G}(n,\cdot) \stackrel{\triangle}{=} \stackrel{(n+s-1,\cdot)\circ}{G}(n+s-2,\cdot)\circ\ldots\circ \overset{(\Delta)}{G}(n,\cdot)$$

with p, q, and s factors, respectively. We can also define $\overset{\bigtriangledown}{G}^{\circ p}(n+s,\cdot) \circ \overset{\bigtriangleup}{G}^{\circ s}(n,\cdot), \ G^{\circ q} \circ \overset{\bigtriangleup}{G}^{\circ p}$ etc. Directly from the definitions the functions in (2.24) - (2.26) can be calculated.

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Lemma 3. With the convention $\binom{k}{\ell} = 0$ for $\ell < 0$

(2.27)
$$\overset{\nabla}{G}^{\circ p}(n,u) = \binom{n-p}{n-p} + \ldots + \binom{n-p}{k+1-p} + \binom{\alpha_k-p}{k-p} + \ldots + \binom{\alpha_t-p}{t-p},$$

(2.28)
$$G^{\circ q}(n,u) = \binom{n}{n} + \ldots + \binom{n}{k+1-q} + \binom{\alpha_k}{k-q} + \ldots + \binom{\alpha_t}{t-q},$$

and

(2.29)
$$\overset{\Delta^{\circ s}}{G}(n,u) = \binom{n+s}{n+s} + \ldots + \binom{n+s}{k+1} + \binom{\alpha_k+s}{k} + \ldots + \binom{\alpha_t+s}{t}.$$

Here (2.28) is well-known from the isoperimetric theorem in the Hamming space.

Another important property of G-type functions is the commutativity of the \circ -operation:

$$\overrightarrow{G} \circ G(n, u) = G \circ \overrightarrow{G}(n, u) =$$

$$(2.30) \qquad \binom{n-1}{n-1} + \ldots + \binom{n-1}{k} + \binom{n-1}{k-1} + \binom{\alpha_k - 1}{k-2} + \ldots + \binom{\alpha_t - 1}{t-2},$$

$$\overrightarrow{\nabla} \quad \overrightarrow{\Delta} \quad \overrightarrow{\nabla} \quad \overrightarrow{\Delta} \quad \overrightarrow{\nabla}$$

(2.31)
$$\begin{array}{c} \stackrel{\Delta}{G} \circ \stackrel{\Delta}{G} (n, u) \stackrel{\Delta}{=} \stackrel{\Delta}{G} \circ \stackrel{\nabla}{G} (n, u) = \\ \begin{pmatrix} n \\ n \end{pmatrix} + \ldots + \begin{pmatrix} n \\ k+1 \end{pmatrix} + \begin{pmatrix} n \\ k \end{pmatrix} + \begin{pmatrix} \alpha_k \\ k-1 \end{pmatrix} + \ldots + \begin{pmatrix} \alpha_t \\ t-1 \end{pmatrix},$$

 and

$$G \circ \overset{\triangle}{G}(n,u) = \overset{\triangle}{G} \circ G(n,u) =$$

$$(2.32) \qquad \binom{n+1}{n+1} + \ldots + \binom{n+1}{k+1} + \binom{n+1}{k} + \binom{\alpha_k+1}{k-1} + \ldots + \binom{\alpha_t+1}{t-1}.$$

Applying (2.27) - (2.29) and (2.30) - (2.32) repeatedly or by calculation we establish general rules.

Lemma 4. We have

$$\begin{split} & \stackrel{\bigtriangledown}{G}^{\circ p} \circ G^{\circ q} \circ \stackrel{\bigtriangleup}{G}^{\circ s} (n, u) = \stackrel{\bigtriangledown}{G}^{\circ p} \circ \stackrel{\bigtriangleup}{G}^{\circ s} \circ G^{\circ q}(n, u) \\ & = G^{\circ q} \circ \stackrel{\bigtriangledown}{G}^{\circ p} \circ \stackrel{\bigtriangleup}{G}^{\circ s} (n, u) = G^{\circ q} \circ \stackrel{\bigtriangleup}{G}^{\circ s} \circ \stackrel{\bigtriangledown}{G}^{\circ p} (n, u) \\ & = \stackrel{\bigtriangleup}{G}^{\circ s} \circ \stackrel{\bigtriangledown}{G}^{\circ p} \circ G^{\circ q}(n, u) = \stackrel{\bigtriangleup}{G}^{\circ s} \circ G^{\circ q} \circ \stackrel{\bigtriangledown}{G}^{\circ r} (n, u) \\ \end{split}$$

$$= \binom{n+s-p}{n+s-p} + \binom{n+s-p}{n+s-p-1} + \dots + \binom{n+s-p}{k+1-p-q}$$

$$(2.33) \qquad \qquad + \binom{\alpha_k+s-p}{k-p-q} + \dots + \binom{\alpha_t+s-p}{t-p-q}$$

for u as in (1.10), $0 \leq p, q, s$.

3. Proof of Theorem 1

Denote an initial segment in squashed order (see [10]) over \mathcal{X}_k^n by S and write \overline{S} for the set of complements of the members of S. Thus $\overline{S} \subset \mathcal{X}_{n-k}^n$ and $|\overline{S}| = |S| = v$, say. We speak here about the complementary squashed order or in short about the CS-order.

We consider first $\nabla_0 \overline{S}$ and $\triangle_1 \overline{S}$.

Lemma 5. For the initial segment \overline{S} defined above

(i) $\nabla_0 \overline{S}$ is the $\overset{\vee}{F}(k,v)$ -th initial segment in the CS-order on \mathcal{X}_{n-k}^{n-1} and

(ii)
$$\triangle_1 \overline{S}$$
 is the $\overset{\triangle}{F}(k,v)$ -th initial segment in the CS-order on $\mathcal{X}_{n+1-k}^{n+1}$.

Proof. (i) We use the expansion (1.6) for v and look at any $s^n \in \overline{S}$:

 $s_{t_i} = 0$ for $i = 1, 2, \dots, k$ and $1 \le t_1 < t_2 < \dots < t_k \le n$.

By the definition of the CS-order there must be a j such that for all $i \in (j,k]$ $t_i = a_i + 1$ and for all $i \leq j$ $t_i \leq a_j$. Now suppose that we delete for some index ℓ s_{t_ℓ} . We can assume that $s_{t_\ell-1}=1$, because otherwise we can delete $S_{t_\ell-1}$ and get the same subsequence. Let s'^{n-1} be the resulting subsequence, $t'_i = t_i$ for $i < \ell$ and $t'_{i-1} = t_i$ for $i > \ell$.

Choose now $j' = \max(\ell, j)$ and notice that for $i \le j'-1$, $t'_i \le a_{j'}-1$, for i > j'-1 $t'_{i-1} = a_i = (a_i - 1) + 1$, and for all $i s'_{t'} = 0$.

Therefore the resulting subsequence s'^{n-1} falls into the $\overset{\bigtriangledown}{F}(k,v)$ -th initial segment in CS-order.

Conversely, given a sequence s'^{n-1} in \mathcal{X}_{n-k}^{n-1} and in the $\overset{\nabla}{F}(k,v)$ -th initial segment the forgoing argument provides a way to find an s^n in the v-th initial segment from which s'^{n-1} is obtainable by deleting a 0.

(ii) Use again the s^n described above and let s''^{n+1} be obtained by inserting a 1 before $s_{t_{\rho''}}$, $t''_i = t_i$ for $i < \ell''$ and $t''_i = t_i + 1$ for $i \ge \ell''$.

Then $s_{t_i''}'=0$ for all i and for $i \le j'' t_i'' \le a_i+1$; for $i > j'' t_i''=a_i+2=(a_i+1)+1$, if we choose $j''=\max(j,\ell-1)$.

Clearly, such an s''^{n-1} is in the $\stackrel{\triangle}{F}(k,v)$ -th initial segment in the CS-order. The same argument gives also the reverse implication.

Proof of Theorem 1 (i) and (ii) by induction on n.

The cases n = 1, 2 are done by simple inspection. For any ℓ , $m, j, C \subset \mathcal{X}^{\ell}$, $D \subset \mathcal{X}^m$, and $E \subset \mathcal{X}^j$ let

(3.1)
$$C_i = \{(c_1, \ldots, c_{\ell-1}) : (c_1, \ldots, c_{\ell-1}, i) \in C\} (\subset \mathcal{X}^{\ell-1}),$$

(3.2)
$$D * i = \{(d_1, \ldots, d_m, i) : (d_1, \ldots, d_m) \in D\} (\subset \mathcal{X}^{m+1}),$$

 and

(3.3)
$$\hat{E}_i = \{(e_1, \dots, e_j) : e_j = i \text{ and } (e_1, \dots, e_j) \in E\} (\subset \mathcal{X}^j)$$

for i = 0, 1.

(i) for n > 2. Since $B_0 \subset \bigtriangledown_0 B, (\bigtriangledown_0 B_i) * i \subset \bigtriangledown_0 B(i = 0, 1)$ and $(\bigtriangledown_0 B_0) * 0 \cap (\bigtriangledown_0 B_1) * 1 = \emptyset$, either $|\bigtriangledown_0 B| \ge |B_0| \ge \stackrel{\frown}{F}(k, |B|)$ or by (2.3) and induction hypothesis (IH) $|\bigtriangledown_0 B| \ge |\bigtriangledown_0 B_0| + |\bigtriangledown_0 B_1| \ge \stackrel{*}{F}(k-1, |B_0|) + \stackrel{\frown}{F}(k, |B_1|) \ge \stackrel{\frown}{F}(k, |B|)$, where (*) is justified by $B_0 \subset \mathcal{X}_{n-k}^{n-1}$, and $B_1 \subset \mathcal{X}_{n-k-1}^{n-1}$.

(ii) for n > 2.

Recall the definition of the operator " \wedge " in (3.3).

Considering $\triangle_1 B = (\widehat{\triangle_1 B_1})_1 \cup (\widehat{\triangle_1 B})_0$, $(\widehat{\triangle_1 B})_0 = (\triangle_1 B_0) * 0$, $B * 1 \subset (\widehat{\triangle_1 B})_1$ and $(\triangle_1 B_1) * 1 \subset (\widehat{\triangle_1 B})_1$, by (2.4) and IH, $|\triangle_1 B| \ge \max(|B| | \triangle_2 B_1|) + |\triangle_2 B_0| \ge |A|$

$$|\Delta_1 B| \ge \max(|B|, |\Delta_1 B_1|) + |\Delta_1 B_0| \ge \\ \max(|B|, \stackrel{\Delta}{F}(k, |B_1|)) + \stackrel{\Delta}{F}(k-1, |B_0|) \ge \stackrel{\Delta}{F}(k, |B|)$$

(iii) follows by Lemma 5.

4. Proof of Theorem 2

Lemma 6. For the initial segment S in the H-order $\triangle S$ equals the $\stackrel{\triangle}{G}(n,|S|)$ -th initial segment in the H-order, and ∇S equals the $\overset{\vee}{G}(n,|S|)$ -th initial segment in H-order.

Proof. By the definitions of the two orders and direct inspection, we first get, that for some k and m, and the m-th initial segment S' (of level n-k) in the CS-order

(5.1)
$$S = \begin{pmatrix} n - (k+1) \\ \bigcup_{\ell=0} \mathcal{X}_{\ell}^{n} \end{pmatrix} \cup S',$$

(5.2)
$$\Delta S = \left(\bigcup_{\ell=0}^{n-k} \mathcal{X}_{\ell}^{n+1}\right) \cup \Delta_1 S',$$

and

(5.3)
$$\nabla S = \begin{pmatrix} n - (k+1) \\ \bigcup_{\ell=0} \mathcal{X}^{n-1} \end{pmatrix} \cup \nabla_0 S'.$$

The rest of the proof follows from Lemma 5.

Proof of Theorem 2 by induction on n. For n=2 the statement is readily verified. From the IH for n-1 we proceed to n.

Next observe that, by convention (3.1) and (3.2), $\bigcup_{i=0}^{1} (\nabla A_i) * i \subset \nabla A, \bigcap_{i=0}^{1} (\nabla A_i) *$ $i = \emptyset$ and that therefore

$$|\bigtriangledown A| \ge \sum_{i=0}^{1} |\bigtriangledown A_i| \ge \sum_{i=0}^{1} \overset{\bigtriangledown}{G} (n-1, |A_i|)$$
 (by the IH).

According to the ∇ -inequality this can be lower bounded with the desired $\stackrel{\bigtriangledown}{G}(n,|A|),$ if $|A_0|,|A_1|$ $\stackrel{\bigtriangledown}{<}\stackrel{\bigtriangledown}{G}(n,|A|).$ Otherwise we have for some i $|A_i| = \max(|A_0|, |A_1|) \ge \overset{\nabla}{G}(n, |A|)$ and we are done again, because $\bigtriangledown A \supset A_i$.

The achievability follows from Lemma 6.

5. Proof of Theorem 3

The proof goes in exactly the same way as the proof of Theorem 1, (ii) (and the " \triangle_1 " part of (iii)), except that here we use (2.6), Lemma 6 and the observations: $\triangle A = (\widehat{\triangle A})_1 \cup (\widehat{\triangle A})_0, (\triangle A_i) * i \subset (\widehat{\triangle A})_i$ and $A * i \subset (\widehat{\triangle A})_i$ (for i = 0, 1).

6. General isoperimetric theorems

We use now the calculus of iterative applications of $\overset{\bigtriangledown}{G}$, $\overset{\bigtriangleup}{G}$, and G described in Section 2 B.

Fortunately our Theorems 2, 3 and Harper's Theorem ([3]) establish the Inheritance property for the operations ∇ , \triangle , and $\Gamma^1_{d_H}$ (recall definition (1.13)). In the sequel, we abbreviate $\Gamma^1_{d_H}$ as Γ_{d_H} and as Γ . If S is an initial segment in Horder, then so are ∇S , $\triangle S$, and $\Gamma_{d_H}S$. This enables us to apply these theorems repeatedly. Formally, we introduce

(6.1)
$$\nabla^{\ell} A = \nabla \big(\bigtriangledown \dots \bigtriangledown (\bigtriangledown A) \dots \big),$$

(6.2)
$$\Delta^{\ell} A = \Delta \big(\Delta \dots \Delta (\Delta A) \dots \big),$$

and

(6.3)
$$\Gamma_{d_H}^{\ell} A = \Gamma(\Gamma \dots \Gamma(\Gamma A) \dots)$$
$$= \left\{ x^n \in \mathcal{X}^n : d_H(x^n, a^n) \le \ell \text{ for some } a^n \in A \right\}$$

and state the results.

Proposition 1. For every $A \subset \mathcal{X}^n$, |A| = u

(i)
$$|\nabla^{\ell}A| \ge \overset{\nabla^{\circ\ell}}{G}(n,u)$$

(ii) $|\Delta^{\ell}A| \ge \overset{\Delta^{\circ\ell}}{G}(n,u)$
(iii) $|\Gamma^{\ell}_{d_H}A| \ge G^{\circ\ell}(n,u)$

and all these bounds are achieved by the u-th initial segment in H-order.

Now we turn to the distances θ and δ in order to generalize Corollary 1. Here operations are combined and the commutative law for the numerical functions (Lemma 4 in Section 2) is needed.

Fortunately this commutative law holds also for the operations ∇ , \triangle , and Γ ! Indeed, using the short notation

$$\nabla\{x^n\} = \nabla x^n, \ \Delta\{x^n\} = \Delta x^n, \ \Gamma\{x^n\} = \Gamma x^n,$$

we see that

(6.4)
$$\nabla \{ \Delta x^n \} = \Delta \{ \nabla x^n \}, \ \Gamma \{ \Delta x^n \} = \Delta \{ \Gamma x^n \}, \ \nabla \{ \Gamma x^n \} = \Gamma \{ \nabla x^n \}.$$

Therefore the commutative law holds for every $A \subset \mathcal{X}^n$:

(6.5)
$$\nabla(\triangle A) = \triangle(\nabla A), \ \Gamma(\triangle A) = \triangle(\Gamma A), \ \nabla(\Gamma A) = \Gamma(\nabla A).$$

Moreover, it is clear that for every $A \subset \mathcal{X}^n$

(6.6)
$$\Gamma^{\ell'}A \subset \nabla^{\ell}(\Delta^{\ell}A) = \Delta^{\ell}(\nabla^{\ell}A) \text{ for } \ell \leq n.$$

Here strict inclusion can occur:

(6.7)
$$\Gamma(1,0) = \{(0,0), (1,0), (1,1)\} \neq \mathcal{X}^2 = \bigtriangledown (\bigtriangleup(1,0)).$$

However, strict inclusion does not occur, if S is an initial segment in H-order.

Proposition 2. If S is an initial segment in H-order, |S| = u, then

(i)
$$|\Delta^{\ell}(\nabla^{\ell}S)| = |\nabla^{\ell}(\Delta^{\ell}S)| = G^{\circ \ell} \circ G^{\circ \ell} \circ G^{\circ \ell}(n,u) = G^{\circ \ell}(n,u) = |\Gamma^{\ell}S|$$

and
(ii) $\Delta^{\ell}(\nabla^{\ell}S) = \nabla^{\ell}(\Delta^{\ell}S) = \Gamma^{\ell}S.$

Proof. For (i) the first equalities are justified by (6.6) and Proposition 1 and the last equality is (the easy) part of Harper's Theorem. The remaining equality follows from Lemma 4 with the choices $p = s = \ell$, q = 0 and p = s = 0, $q = \ell$, respectively: both quantities equal $\binom{n}{n} + \ldots + \binom{n}{k+1-\ell} + \binom{\alpha_k}{k-\ell} + \ldots + \binom{\alpha_t}{t-\ell}$. Notice that (i) and (6.6) imply (ii).

Now we consider arbitrary sets $A \subset \mathcal{X}^n$ and the distances θ , δ .

Proposition 3. For any $A \subset \mathcal{X}^n$, r > 0 and any ℓ_i, ℓ'_i (i = 1, 2) with $\ell_2 - \ell_1 = \ell'_2 - \ell'_1$ and $\ell_2 < \ell'_2$

(i)
$$\nabla^{\ell_2}(\triangle^{\ell_1}A) \subset \nabla^{\ell'_2}(\triangle^{\ell'_1}A)$$

and
(ii) $\Gamma^r A = \prod_{l=1}^r \nabla^{\lfloor (r+\ell)/2 \rfloor}(\triangle^{\lfloor (r+\ell)/2 \rfloor})$

(ii)
$$\Gamma_{\theta}^{r} A = \bigcup_{\ell=-r} \nabla^{\lfloor (r+\ell)/2 \rfloor} (\triangle^{\lfloor (r-\ell)/2 \rfloor} A)$$
$$= \bigcup_{\ell=0}^{r-1} \left[\left(\nabla^{\ell} (\triangle^{r-\ell} A) \right) \cup \left(\nabla^{\ell} (\triangle^{r-1-\ell} A) \right) \right] \cup \nabla^{r} A,$$

where by convention $\triangle^0 A = \nabla^0 A = A$.

Proof. Obviously, for all ℓ ,

and therefore by the commutative law (6.5)

$$\nabla^{\ell_2}(\triangle^{\ell_1}A) \subset \nabla^{\ell_2}\left(\triangle^{\ell_1}\left(\nabla^{\ell_2'-\ell_2}(\triangle^{\ell_2'-\ell_2}A)\right)\right) = \\\nabla^{\ell_2}\left(\triangle^{\ell_1}\left(\nabla^{\ell_2'-\ell_2}(\triangle^{\ell_1'-\ell_1}A)\right) = \nabla^{\ell_2'}(\triangle^{\ell_1'}A)\right),$$

and thus (i) is verified.

Again by (6.5) and the definition of distance θ

(6.9)
$$\Gamma_{\theta}^{r}A = \bigcup_{r_{1}+r_{2} \leq r} (\nabla^{r_{2}}(\triangle^{r_{1}}A)).$$

Thus by (i) and (6.9)

$$\Gamma_{\theta}^{r} = \bigcup_{\ell=-r}^{r} \bigcup_{\substack{r_{1}+r_{2} \leq r \\ r_{2}-r_{1}=\ell}} \left(\nabla^{r_{2}}(\triangle^{r_{1}}A) \right) = \bigcup_{\ell=-r}^{r} \left(\nabla^{\lfloor (r+\ell)/2 \rfloor}(\triangle^{\lfloor (r-\ell)/2 \rfloor}A) \right)$$
$$= \bigcup_{\ell=0}^{r-1} \left[\left(\nabla^{\ell}(\triangle^{r-\ell}A) \right) \cup \left(\nabla^{\ell}(\triangle^{r-1-\ell}A) \right) \right] \cup \nabla^{r}A.$$

We are now ready to state and prove the main result.

Theorem 4. For all $A \subset \mathcal{X}^n$ and $r \ge 0$

(i)
$$|\Gamma_{\theta}^{r}A| \ge \sum_{\ell=-r}^{r} \overset{\nabla^{\circ}\left\lfloor \frac{r+\ell}{2} \right\rfloor}{G} \overset{\bigtriangleup}{\circ} \overset{\Box}{G} \overset{\left\lfloor \frac{r-\ell}{2} \right\rfloor}{(n,|A|)}$$

and

(ii)
$$|\Gamma_{\delta}^{r}A| \geq \sum_{\ell=1}^{r} \left[\overset{\nabla^{\circ\ell}}{G} \circ G^{\circ(r-\ell)}(n,|A|) + \overset{\Delta^{\circ\ell}}{G} \circ G^{\circ(r-\ell)}(n,|A|) \right] + G^{\circ r}(n,|A|),$$

where $G^{\circ 0}(n, u) = u$, and both bounds are achieved by the |A|-th initial segment in *H*-order.

Proof. By our definitions for $0 \le \ell_i$ (i=1,2) and $n-\ell_2+\ell_1 \ge 0$

(6.10)
$$\nabla^{\ell_2} \left(\triangle^{\ell_1} (\Gamma^{\ell_0} A) \right) \subset \mathcal{X}^{n-\ell_2+\ell_1}.$$

(Here Γ^{ℓ_0} is only used for proving (ii).)

Therefore also

(6.11)
$$\nabla^{\ell_2} \left(\Delta^{\ell_1}(\Gamma^{\ell_0} A) \right) \cap \nabla^{\ell'_2} \left(\Delta^{\ell'_1}(\Gamma^{\ell'_0} A) \right) = \emptyset, \quad \text{if} \quad \ell_1 - \ell_2 \neq \ell'_1 - \ell'_2$$

and (i) as well as its optimality immediately follows from Proposition 3, (6.11) and Proposition 1 (applied twice).

(ii) Similarly to (6.9), we have also

(6.12)
$$\Gamma^r_{\delta}A = \bigcup_{r_1+r_2+r_0 \le r} \left(\bigtriangledown^{r_2} \left(\bigtriangleup^{r_1}(\Gamma^{r_0}_{d_H}A) \right) \right)$$

and therefore

(6.13)
$$\Gamma_{\delta}^{r}A \supset \bigcup_{\ell=1}^{r} \left(\nabla^{\ell} (\Gamma_{d_{H}}^{r-\ell}A) \right) \cup \left(\triangle^{\ell} (\Gamma_{d_{H}}^{r-\ell}A) \right) \cup \Gamma_{d_{H}}^{r}A.$$

Hence (ii) follows from (6.11), (6.13) and Proposition 1 (applied twice).

Finally, we have to show that the |A|-th initial segment in H-order S achieves equality.

By Proposition 2 (ii), Proposition 3 (i), (6.12) and (6.11), and by the monotonicity of Δ^t , ∇^t , $\Gamma^t_{d_H}$ in the sets it suffices to show that for all parameters $-r \leq \ell \leq r, \ \ell_1 + \ell_2 + \ell_0 = r, \ \ell_2 - \ell_1 = \ell$, and $\ell_i \geq 0$ for i = 0, 1, 2

$$\nabla^{\ell_2} \left(\triangle^{\ell_1} (\Gamma^{\ell_0} S) \right) \subset \begin{cases} \nabla^{\ell} (\Gamma^{r-\ell} S), & \text{if } \ell_2 > \ell_1 \\ \triangle^{|\ell|} (\Gamma^{r-|\ell|} S), & \text{if } \ell_2 < \ell_1 \\ \Gamma^r S, & \text{if } \ell_2 = \ell_1. \end{cases}$$

Let us abbreviate $\nabla^{\ell_2} (\triangle^{\ell_1}(\Gamma^{\ell_0}S)) = L.$

Using Proposition 2 (ii) and Proposition 3 (i) we show the desired inclusions. Case $\ell_2 > \ell_1$.

$$L = \bigtriangledown^{\ell_2 - \ell_1} \left(\bigtriangledown^{\ell_1} \left(\bigtriangleup^{\ell_1} (\Gamma^{\ell_0} S) \right) \right)$$

= $\bigtriangledown^{\ell} \left(\Gamma^{\ell_1} (\Gamma^{\ell_0} S) \right) = \bigtriangledown^{\ell} (\Gamma^{\ell_1 + \ell_0} S) = \bigtriangledown^{\ell} (\Gamma^{r - \ell_2} S)$
 $\subset \bigtriangledown^{\ell} (\Gamma^{r - \ell} S) \text{ (as } \ell_2 > \ell_1 \ge 0, r - \ell_2 \le r - \ell).$

Case $\ell_2 < \ell_1$.

$$\begin{split} L &= \Delta^{\ell_1 - \ell_2} \left(\Delta^{\ell_2} \left(\nabla^{\ell_2} (\Gamma^{\ell_0} S) \right) \right) \\ &= \Delta^{|\ell|} (\Gamma^{\ell_0 + \ell_2} S) = \Delta^{|\ell|} (\Gamma^{r - \ell_1} S) \\ &\subset \Delta^{|\ell|} (\Gamma^{r - |\ell|} S) \quad (\text{as} \quad 0 \le \ell_2 < \ell_1, r - \ell_1 \le r + \ell = r - |\ell|), \end{split}$$

Case $\ell_2 = \ell_1$.

$$L = \nabla^{\ell_1} \left(\triangle^{\ell_1}(\Gamma^{\ell_0} S) \right) = \Gamma^{\ell_0 + \ell_1} S \subset \Gamma^r S.$$

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