# On extremal double circulant self-dual codes of lengths 90-96 

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#### Abstract

A classification of extremal double circulant self-dual codes of lengths up to 88 is known. We give a classification of extremal double circulant self-dual codes of lengths $90,92,94$ and 96 . We also classify double circulant self-dual codes with parameters $[90,45,14]$ and $[96,48,16]$. In addition, we demonstrate that no double circulant self-dual [90, 45, 14] code has an extremal self-dual neighbor, and no double circulant selfdual $[96,45,16]$ code has a self-dual neighbor with minimum weight at least 18 .


## 1 Introduction

A (binary) $[n, k]$ code $C$ is a $k$-dimensional vector subspace of $\mathbb{F}_{2}^{n}$, where $\mathbb{F}_{2}$ denotes the finite field of order 2 . All codes in this note are binary. The parameter $n$ is called the length of $C$. The weight $\mathrm{wt}(x)$ of a vector $x \in \mathbb{F}_{2}^{n}$ is the number of non-zero components of $x$. A vector of $C$ is a codeword of $C$. The minimum non-zero weight of all codewords in $C$ is called the minimum weight of $C$ and an $[n, k]$ code with minimum weight $d$ is called an $[n, k, d]$ code. The weight enumerator $W(C)$ of $C$ is given by $W(C)=\sum_{i=0}^{n} A_{i} y^{i}$ where $A_{i}$ is the number of codewords of weight $i$ in $C$. Two codes are equivalent if one can be obtained from the other by a permutation of coordinates. The dual code $C^{\perp}$ of a code $C$ of length $n$ is

[^0]defined as $C^{\perp}=\left\{x \in \mathbb{F}_{2}^{n} \mid x \cdot y=0\right.$ for all $\left.y \in C\right\}$, where $x \cdot y$ is the standard inner product. A code $C$ is called self-dual if $C=C^{\perp}$. A self-dual code $C$ is called doubly even and singly even if all codewords have weight $\equiv 0(\bmod 4)$ and if some codeword has weight $\equiv 2(\bmod 4)$, respectively.

It was shown in [15] that the minimum weight $d$ of a doubly even selfdual code of length $n$ is bounded by $d \leq 4[n / 24]+4$. We call a doubly even self-dual code meeting this upper bound extremal. The largest possible minimum weights of (singly even) self-dual codes of lengths up to 72 are given in [3, Table I]. This work was extended to lengths up to 100 in [5, Table VI] (see [6, Table 2] and [13, Table I]). According to [10], in this note, we say that a singly even self-dual code with the largest possible minimum weight given in [3, Table I] and [5, Table VI] is extremal. The largest possible minimum weight among singly even self-dual codes of lengths $90,92,94$ and 96 is $16,16,18$ and 18 , respectively. Currently, it is not known if an extremal self-dual $[90,45,16]$ code exists. There is a self-dual $[90,45,14]$ code [5]. Many extremal self-dual $[92,46,16]$ codes are known (see [4], [7], [14], [17] and [18], and references [5] and [10] in [14]). Currently, it is not known if an extremal self-dual $[94,47,18]$ code exists. There is a self-dual $[94,47,16]$ code [13]. Currently, it is not known if an extremal doubly even self-dual $[96,48,20]$ code, or an extremal singly even self-dual $[96,48,18]$ code, exists. There is a doubly even self-dual $[96,48,16]$ code (see [5]), and a singly even self-dual $[96,48,16]$ code [6].

Let $D_{p}$ and $D_{b}$ be codes with generator matrices of the form

$$
\left(\begin{array}{cc}
I_{n} & R \tag{1}
\end{array}\right)
$$

and

$$
\left(\begin{array}{ccccc} 
& 0 & 1 & \cdots & 1  \tag{2}\\
& 1 & & & \\
I_{n+1} & \vdots & & R^{\prime} & \\
& 1 & & &
\end{array}\right)
$$

respectively, where $I_{n}$ is the identity matrix of order $n$, and $R$ and $R^{\prime}$ are $n \times n$ circulant matrices. An $n \times n$ circulant matrix has the form

$$
\left(\begin{array}{ccccc}
r_{0} & r_{1} & r_{2} & \cdots & r_{n-1} \\
r_{n-1} & r_{0} & r_{1} & \cdots & r_{n-2} \\
\vdots & \vdots & \vdots & & \vdots \\
r_{1} & r_{2} & r_{3} & \cdots & r_{0}
\end{array}\right)
$$

so that each successive row is a cyclic shift of the previous one. The codes $D_{p}$ and $D_{b}$ are called pure double circulant and bordered double circulant, respectively. The two families are called double circulant codes. Many of the best known self-dual codes are double circulant codes (see [5], 8], [9, 10] and [12]). Further, constructions exist that provide double circulant selfdual codes with the largest known minimum weight (see [11] and [16]). The bordered double circulant construction provides self-dual codes only when the length is $\equiv 0(\bmod 4)$. In addition, it is known [8] that there is no bordered double circulant singly even self-dual code of length $n \equiv 0(\bmod 8)$.

A classification of extremal double circulant self-dual codes of lengths up to 88 was given in [9], [10] and [12]. In this note, this work is extended to length 96. Our exhaustive search shows that there is no extremal double circulant self-dual $[90,45,16]$ code. We also give a classification of double circulant self-dual $[90,45,14]$ codes. In addition, we demonstrate that every double circulant self-dual $[90,45,14]$ code has no extremal self-dual neighbor. We give a classification of extremal double circulant self-dual codes of length 92. Our exhaustive search shows that there is no double circulant self-dual $[94,47, d]$ code with $d \geq 16$. We give a classification of double circulant self-dual $[96,48, d]$ codes with $d \geq 16$. In addition, we demonstrate that every double circulant self-dual $[96,48,16]$ code has no self-dual neighbor with minimum weight at least 18 .

## 2 Double circulant self-dual [90, 45, d] codes with $d \in\{14,16\}$

Using an approach similar to that given in [9], 10] and [12], our exhaustive search found all distinct double circulant self-dual $[90,45, d]$ codes with $d \geq$ 14. This was done by considering all $45 \times 45$ orthogonal circulant matrices satisfying the condition that the weight of the first row is congruent to 1 $(\bmod 4)$ and the weight is greater than or equal to $d-1$. Since a cyclic shift of the first row of some codes defines an equivalent code, the elimination of cyclic shifts substantially reduces the number of codes which must be checked further for equivalence to complete the classification. It is useful to use the fact that self-dual codes with generator matrices of the form $\left(\begin{array}{ll}I_{45} & R\end{array}\right)$ and $\left(\begin{array}{ll}I_{45} & R^{T}\end{array}\right.$ ) are equivalent, where $R^{T}$ denotes the transpose of $R$. MAGMA [1] was employed to determine code equivalence and complete the classification.

Then we have the following results.
Proposition 1. There is no extremal double circulant self-dual code of length 90.

Proposition 2. There are 716 inequivalent double circulant self-dual [90, 45, 14] codes.

The first rows of $R$ in the generator matrices ( $\left.\begin{array}{ll}I_{45} & R\end{array}\right)$ of the 716 codes can be obtained from http://www.math.is.tohoku.ac.jp/~mharada/Paper/DCC90.txt. We verified by Magma [1] that each of the 716 codes has an automorphism group of order 90.

We determined the possible weight enumerators of self-dual [90, 45, 14] codes. For a detailed description of how this is accomplished, see [3, Theorem 5]. The possible weight enumerators of self-dual [90,45,14] codes and the shadows are as follows

$$
\begin{aligned}
& 1+(14040+a) y^{14}+(51300+3 a+8 b) y^{16}+(69920-11 a-24 b+512 c) y^{18} \\
& +(2355624-41 a-80 b-4608 c+32768 d) y^{20} \\
& +(30913560+49 a+304 b+13824 c-491520 d-2097152 e) y^{22}+\cdots, \\
& e y+(d-22 e) y^{5}+(-c-20 d+231 e) y^{9}+(b+18 c+190 d-1540 e) y^{13} \\
& +(-8 a-16 b-153 c-1140 d+7315 e) y^{17}+\cdots,
\end{aligned}
$$

respectively, where $a, b, c, d, e$ are integers. It is easy to see that the number of codewords of weights 14,16 in the code and the number of vectors of weights $1,5,9$ in the shadow uniquely determine the weight enumerator. By calculating these numbers, we verified that the 716 codes have 100 distinct weight enumerators. This was done using Magma [1]. The 100 weight enumerators have $c=d=e=0$, where $(a, b)$ are listed in Table 1. For each pair $(a, b)$, the number $N_{(a, b)}$ of codes with the weight enumerator is also listed in Table 1 .

## 3 Neighbors of double circulant self-dual [90, 45, 14] codes

Two self-dual codes $C$ and $C^{\prime}$ of length $n$ are said to be neighbors if $\operatorname{dim}(C \cap$ $\left.C^{\prime}\right)=n / 2-1$. We give some observations from [2] on self-dual codes constructed by neighbors. Let $C$ be a self-dual $[n, n / 2, d]$ code. Let $M$ be a

Table 1: Weight enumerators of double circulant self-dual $[90,45,14]$ codes

| $(a, b)$ | $N_{(a, b)}$ | $(a, b)$ | $N_{(a, b)}$ | $(a, b)$ | $N_{(a, b)}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $(-12555,0)$ | 1 | $(-12555,90)$ | 1 | $(-12600,180)$ | 1 |
| $(-12735,0)$ | 1 | $(-12780,0)$ | 1 | $(-12825,0)$ | 3 |
| $(-12825,90)$ | 2 | $(-12870,0)$ | 5 | $(-12870,180)$ | 1 |
| $(-12915,0)$ | 3 | $(-12915,180)$ | 1 | $(-12915,90)$ | 4 |
| $(-12960,0)$ | 6 | $(-12960,90)$ | 9 | $(-13005,0)$ | 6 |
| $(-13005,180)$ | 1 | $(-13005,90)$ | 10 | $(-13050,0)$ | 10 |
| $(-13050,180)$ | 1 | $(-13050,270)$ | 1 | $(-13050,90)$ | 8 |
| $(-13095,0)$ | 8 | $(-13095,180)$ | 4 | $(-13095,270)$ | 1 |
| $(-13095,90)$ | 6 | $(-13140,0)$ | 17 | $(-13140,180)$ | 7 |
| $(-13140,270)$ | 1 | $(-13140,90)$ | 13 | $(-13185,0)$ | 10 |
| $(-13185,180)$ | 5 | $(-13185,270)$ | 1 | $(-13185,360)$ | 1 |
| $(-13185,90)$ | 16 | $(-13230,0)$ | 6 | $(-13230,180)$ | 10 |
| $(-13230,270)$ | 4 | $(-13230,90)$ | 27 | $(-13275,0)$ | 21 |
| $(-13275,180)$ | 15 | $(-13275,270)$ | 1 | $(-13275,360)$ | 3 |
| $(-13275,90)$ | 16 | $(-13320,0)$ | 11 | $(-13320,180)$ | 20 |
| $(-13320,270)$ | 5 | $(-13320,450)$ | 1 | $(-13320,90)$ | 19 |
| $(-13365,0)$ | 13 | $(-13365,180)$ | 19 | $(-13365,270)$ | 7 |
| $(-13365,360)$ | 1 | $(-13365,90)$ | 13 | $(-13410,0)$ | 8 |
| $(-13410,180)$ | 22 | $(-13410,270)$ | 6 | $(-13410,360)$ | 2 |
| $(-13410,90)$ | 23 | $(-13455,0)$ | 11 | $(-13455,180)$ | 18 |
| $(-13455,270)$ | 13 | $(-13455,360)$ | 1 | $(-13455,90)$ | 33 |
| $(-13500,0)$ | 8 | $(-13500,180)$ | 17 | $(-13500,270)$ | 9 |
| $(-13500,360)$ | 3 | $(-13500,90)$ | 23 | $(-13545,0)$ | 4 |
| $(-13545,180)$ | 19 | $(-13545,270)$ | 8 | $(-13545,450)$ | 1 |
| $(-13545,90)$ | 18 | $(-13590,0)$ | 3 | $(-13590,180)$ | 8 |
| $(-13590,270)$ | 9 | $(-13590,360)$ | 2 | $(-13590,450)$ | 2 |
| $(-13590,90)$ | 9 | $(-13635,0)$ | 2 | $(-13635,180)$ | 7 |
| $(-13635,270)$ | 5 | $(-13635,360)$ | 2 | $(-13635,450)$ | 2 |
| $(-13635,90)$ | 6 | $(-13680,180)$ | 9 | $(-13680,270)$ | 4 |
| $(-13680,360)$ | 1 | $(-13680,90)$ | 2 | $(-13725,180)$ | 2 |
| $(-13725,270)$ | 2 | $(-13725,360)$ | 2 | $(-13725,90)$ | 2 |
| $(-13770,180)$ | 1 | $(-13770,270)$ | 2 | $(-13815,180)$ | 1 |
| $(-13815,270)$ | 2 | $(-13815,360)$ | 1 | $(-13815,90)$ | 1 |
| $(-13905,360)$ | 2 |  |  |  |  |
|  |  |  |  |  |  |

matrix whose rows are the codewords of weight $d$ in $C$. Suppose that there is a vector $x \in \mathbb{F}_{2}^{n}$ such that

$$
\begin{equation*}
M x^{T}=\mathbf{1}^{T} \tag{3}
\end{equation*}
$$

where $\mathbf{1}$ is the all-one vector. Set $C_{0}=\langle x\rangle^{\perp} \cap C$, where $\langle x\rangle$ denotes the code generated by $x$. Then $C_{0}$ is a subcode of index 2 in $C$. If the weight of $x$ is even, then we have two self-dual neighbors $\left\langle C_{0}, x\right\rangle$ and $\left\langle C_{0}, x+y\right\rangle$ of $C$ for some $y \in C \backslash C_{0}$, which do not contain any codewords of weight $d$ in $C$, where $\langle C, x\rangle=C \cup(x+C)$. When $C$ has a self-dual $\left[n, n / 2, d^{\prime}\right]$ neighbor $C^{\prime}$ with $d^{\prime} \geq d+2$, (3) has a solution $x$ and we can obtain $C^{\prime}$ in this way. If $\operatorname{rank} M<\operatorname{rank}\left(M \mathbf{1}^{T}\right)$, then $C$ has no self-dual $\left[n, n / 2, d^{\prime}\right]$ neighbor $C^{\prime}$ with $d^{\prime} \geq d+2$. If $\operatorname{rank} M=t$, then we have at most $2 \times 2^{n / 2-t}$ self-dual neighbors of $C$. Furthermore, if the subcode generated by the codewords of weight $d$ in $C$ contains 1, then $C$ has exactly $2 \times 2^{n / 2-t}$ self-dual neighbors. When $C$ has a self-dual $\left[n, n / 2, d^{\prime}\right]$ neighbor $C^{\prime}$ with $d^{\prime} \geq d+2$, (3) has a solution $x$ and we can obtain $C^{\prime}$ in this way.

We verified by Magma [1] that

$$
\left(\operatorname{rank} M, \operatorname{rank}\left(M \mathbf{1}^{T}\right)\right)=(43,43),
$$

for one of the 716 double circulant self-dual $[90,45,14]$ codes and

$$
\left(\operatorname{rank} M, \operatorname{rank}\left(M \mathbf{1}^{T}\right)\right)=(45,45)
$$

for the remaining 715 codes. In addition, using the above method, we verified by Magma [1] that the self-dual neighbors constructed by the above argument have minimum weight at most 14 . Hence, we have the following result.

Proposition 3. No double circulant self-dual [90, 45, 14] code has an extremal self-dual neighbor of length 90.

It is still an open problem whether an extremal self-dual code of length 90 exists.

## 4 Extremal double circulant self-dual codes of length 92

Using a method similar to that given in Section2, our exhaustive search found all distinct extremal pure and bordered double circulant self-dual codes of length 92 . Then we have the following results.

Proposition 4. There is no extremal pure double circulant self-dual code of length 92.

Remark 5. Alfred Wassermann in a private communication indicated that there is no extremal pure double circulant self-dual code of length 92, which provides an independent confirmation of our results.

Proposition 6. There are 158 inequivalent extremal bordered double circulant self-dual codes of length 92.

We denote the 158 inequivalent extremal bordered double circulant selfdual codes of length 92 by $B_{92, i}(i=1,2, \ldots, 158)$. For the codes $B_{92, i}$ ( $i=1,2, \ldots, 158$ ), the first rows $r$ of $R^{\prime}$ in (22) are listed in Table 3) In the table, the rows are written in octal using $0=(000), 1=(001), \ldots, 6=(110)$ and $7=(111)$. We verified by MAGMA [1] that $B_{92, i}$ has an automorphism group of order 90 for $i=1,2, \ldots, 158$.

The possible weight enumerators of extremal self-dual codes of length 92 are given in [5] as follows

$$
\begin{aligned}
W_{92,1}= & 1+(4 \beta+4692) y^{16}+(174800-8 \beta+256 \alpha) y^{18} \\
& +(-2048 \alpha+2425488-52 \beta) y^{20}+\cdots, \\
W_{92,2}= & 1+(4 \beta+4692) y^{16}+(174800-8 \beta+256 \alpha) y^{18} \\
& +(-2048 \alpha+2441872-52 \beta) y^{20}+\cdots, \\
W_{92,3}= & 1+(4 \beta+4692) y^{16}+(121296-8 \beta) y^{18} \\
& +(3213968-52 \beta) y^{20}+\cdots,
\end{aligned}
$$

where $\alpha, \beta$ are integers. By calculating the numbers of codewords of weights $16,18,20$ in the codes, we verified that $B_{92, i}$ has weight enumerator $W_{92,3}$, where $i$ and $\beta$ in $W_{92,3}$ are listed in Table 2,

## 5 Double circulant self-dual [94, 47, d] codes with $d \geq 16$

As mentioned in Section 1, it is currently not known if an extremal self-dual code of length 94 exists. There is a self-dual [94, 47, 16] code [13].

Using a method similar to that given in Section 2, our exhaustive search found no double circulant self-dual $[94,46, d]$ code with $d \geq 16$. Then we have the following result.

Table 2: Weight enumerators of $B_{92, i}(i=1,2, \ldots, 158)$

| $\beta$ | $i$ |
| :---: | :--- |
| 1527 | 81,140 |
| 1572 | $57,64,102$ |
| 1617 | $52,55,77,106$ |
| 1662 | $6,14,18,33,56,87,121,151$ |
| 1707 | $3,36,38,50,69,99,101,109,111,123,143$ |
| 1752 | $5,21,40,49,63,116,125,127$ |
| 1797 | $15,27,32,46,89,95,105,138,141,147,152,153,156$ |
| 1842 | $1,8,10,17,22,66,72,85,90,97,108$ |
| 1887 | $13,26,39,41,44,48,58,62,74,84,91,103,110,112,113,119,130$, |
|  | $136,139,154,155$ |
| 1932 | $16,51,80,98,131,134$ |
| 1977 | $2,23,24,47,53,59,61,86,120,126$ |
| 2022 | $28,31,37,60,67,79,82,88,92,114,117,118,128,137,144,146,150$ |
| 2067 | $4,7,19,35,100,107,157,158$ |
| 2112 | $20,65,76,94,96,104,129,132,142,145$ |
| 2157 | $11,34,45,68,70,133$ |
| 2202 | $25,54,71,75,83,149$ |
| 2247 | $30,42,115$ |
| 2292 | $12,29,43,93,122$ |
| 2337 | $73,124,135$ |
| 2382 | 78 |
| 2427 | 148 |
| 2607 | 9 |

Proposition 7. There is no double circulant self-dual code of length 94 and minimum weight $d \geq 16$.

## 6 Double circulant self-dual [96,48, $d]$ codes with $d \geq 16$

As described in Section 1 it is currently not known if an extremal doubly even self-dual $[96,48,20]$ code exists, or if an extremal singly even self-dual $[96,48,18]$ code exists. There is a doubly even self-dual $[96,48,16]$ code [5], and a singly even self-dual $[96,48,16]$ code [6].

Using a method similar to that given in Section 2, our exhaustive search found all distinct pure double circulant self-dual $[96,48, d]$ codes with $d \geq 16$

Table 3: First rows of $R^{\prime}$ in (2) for $B_{92, i}(i=1,2, \ldots, 158)$

| 2 | $r$ | $i$ | $r$ | $i$ | $r$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 045722771307000 | 2 | 046354263735000 | 3 | 054130272607000 |
| 4 | 102155447541000 | 5 | 104521145473000 | 6 | 110545607071000 |
| 7 | 111222513531000 | 8 | 115060555603000 | 9 | 126644436177000 |
| 10 | 130023052373000 | 11 | 130115321127000 | 12 | 141232724213000 |
| 13 | 165172671757000 | 14 | 206201322771000 | 15 | 207022470467000 |
| 16 | 221254213777000 | 17 | 233413676413000 | 18 | 236265461527000 |
| 19 | 243271301677000 | 20 | 246165167155000 | 21 | 263645737137000 |
| 22 | 265543417117000 | 23 | 271737137037000 | 24 | 304151364577000 |
| 25 | 306307103017000 | 26 | 330141216433000 | 27 | 341136314537000 |
| 28 | 367621175177000 | 29 | 406233754353000 | 30 | 407777023131000 |
| 31 | 436453744513000 | 32 | 442536745265000 | 33 | 453136757327000 |
| 34 | 506762231273000 | 35 | 522423127177000 | 36 | 533555410563000 |
| 37 | 536456502533000 | 38 | 545721744703000 | 39 | 552732211353000 |
| 40 | 577604234513000 | 41 | 605303705657000 | 42 | 616171605527000 |
| 43 | 616352763673000 | 44 | 645661771573000 | 45 | 652610723547000 |
| 46 | 041045754474740 | 47 | 041074234673640 | 48 | 043135457542240 |
| 49 | 043250274476540 | 50 | 043574372741740 | 51 | 043603253362640 |
| 52 | 044104457046640 | 53 | 044226336701740 | 54 | 044671546702540 |
| 55 | 044736177354640 | 56 | 046241222776640 | 57 | 047663617660740 |
| 58 | 050576621023740 | 59 | 050751623626240 | 60 | 052354251553140 |
| 61 | 053047243053740 | 62 | 053452447124740 | 63 | 054716632374740 |
| 64 | 055727742734140 | 65 | 057165135375140 | 66 | 060736704331140 |
| 67 | 061065253646540 | 68 | 063057644302740 | 69 | 063731763152340 |
| 70 | 065235232367740 | 71 | 066764121737540 | 72 | 072237363775740 |
| 73 | 101742440560540 | 74 | 103257370547740 | 75 | 104571361141740 |
| 76 | 105644361251740 | 77 | 107752466111140 | 78 | 112175633752540 |
| 79 | 112453746103640 | 80 | 113343466667340 | 81 | 114740315712340 |
| 82 | 115256766234740 | 83 | 115303767255340 | 84 | 115331337561640 |
| 85 | 116226336753640 | 86 | 116277617462540 | 87 | 123320123173340 |
| 88 | 123663146657540 | 89 | 123763453707140 | 90 | 127654632533640 |
| 91 | 127737665370740 | 92 | 131553671516340 | 93 | 132721322740340 |
| 94 | 134537632035740 | 95 | 136446677675740 | 96 | 140352750117540 |
| 97 | 141147475510340 | 98 | 142376737627740 | 99 | 145756370077140 |
| 100 | 146514734137740 | 101 | 146542652434540 | 102 | 147307747376740 |
| 103 | 151056130545740 | 104 | 153556753442340 | 105 | 153743314476340 |
| 106 | 154356132761740 | 107 | 155237453760340 | 108 | 162603763561740 |
| 109 | 162755377664740 | 110 | 163131275323740 | 111 | 203315731776440 |
| 112 | 204336305345340 | 113 | 205721743736640 | 114 | 206127347363740 |
| 115 | 216236374745540 | 116 | 216712337262740 | 117 | 216776346322340 |
| 118 | 223601644714740 | 119 | 225103675656740 | 120 | 225756665264340 |
| 121 | 227266265646740 | 122 | 227656146713640 | 123 | 231235753751440 |
| 124 | 233466766660640 | 125 | 233475224764740 | 126 | 235161677207340 |
| 127 | 236612727317440 | 128 | 236657266701540 | 129 | 244353765463340 |
| 130 | 246531765347240 | 131 | 257573767412740 | 132 | 261574375164340 |
| 133 | 265576361776640 | 134 | 272764241653740 | 135 | 275373174710140 |
| 136 | 277333517335540 | 137 | 277444625674540 | 138 | 277464773666340 |
| 139 | 277757446333340 | 140 | 306156645077740 | 141 | 306533337306340 |
| 142 | 310773753761740 | 143 | 313675510752340 | 144 | 317071170770740 |
| 145 | 324761561777740 | 146 | 357234774374740 | 147 | 463637676050640 |
| 148 | 512662622756740 | 149 | 516703753357740 | 150 | 535715465737340 |
| 151 | 537761713627540 | 152 | 547776766153140 | 153 | 626053773755740 |
| 154 | 656707543160740 | 155 | 043557136776166 | 156 | 047172572571772 |
| 157 | 051150777762736 | 158 | 066775577156766 |  |  |

and all distinct bordered double circulant doubly even self-dual $[96,48, d]$ codes with $d \geq 16$. Then we have the following result.

Proposition 8. There is no extremal double circulant doubly even self-dual code of length 96. There is no extremal double circulant singly even self-dual code of length 96.

Remark 9. Alfred Wassermann in a private communication indicated that there is no double circulant self-dual code of length 96 and minimum weight $d \geq 18$, which provides an independent confirmation of our results.

Proposition 10. There are 49 inequivalent pure double circulant singly even self-dual $[96,48,16]$ codes. There are 4565 inequivalent pure double circulant doubly even self-dual $[96,48,16]$ codes. There are 1532 inequivalent bordered double circulant doubly even self-dual $[96,48,16]$ codes.

Table 4: First rows of $R$ in (1) for $C_{96, i}(i=1,2, \ldots, 49)$

| $i$ | $r$ | $(a, b, c, d)$ | $i$ | $r$ | $(a, b, c, d)$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 5532465545470000 | $(9798,0,0,0)$ | 2 | 1011116717627400 | $(10050,0,0,0)$ |
| 3 | 2117213667133520 | $(10164,0,0,0)$ | 4 | 5160450553527400 | $(10416,0,0,0)$ |
| 5 | 0411642402747400 | $(10422,0,0,0)$ | 6 | 1110737636054400 | $(10434,0,0,0)$ |
| 7 | 2730315332407400 | $(10566,0,0,0)$ | 8 | 4127775466731720 | $(10566,0,0,0)$ |
| 9 | 5247741422235400 | $(10740,0,0,0)$ | 10 | 1104701417751460 | $(10854,0,0,0)$ |
| 11 | 1334257665167760 | $(10980,0,0,0)$ | 12 | 1551523722207400 | $(11154,0,0,0)$ |
| 13 | 1072513135756620 | $(11364,0,0,0)$ | 14 | 5302720447547400 | $(11508,0,0,0)$ |
| 15 | 1115027566566720 | $(11820,0,0,0)$ | 16 | 1176414666173320 | $(12108,0,0,0)$ |
| 17 | 1252510325477400 | $(12180,0,0,0)$ | 18 | 0707334570645560 | $(9618,-48,0,0)$ |
| 19 | 0536450432504760 | $(10326,-48,0,0)$ | 20 | 2727311567536720 | $(10326,-48,0,0)$ |
| 21 | 4260735067342400 | $(10422,-48,0,0)$ | 22 | 5720417224633400 | $(10434,-48,0,0)$ |
| 23 | 0465637224357620 | $(10566,-48,0,0)$ | 24 | 0447671345066400 | $(11124,-48,0,0)$ |
| 25 | 0644667174474660 | $(11844,-48,0,0)$ | 26 | 1667375134475360 | $(11994,-48,0,0)$ |
| 27 | 1233543431133400 | $(10458,-96,0,0)$ | 28 | 5772526161347400 | $(10806,-96,0,0)$ |
| 29 | 4072735065262400 | $(11190,-96,0,0)$ | 30 | 1077057777245360 | $(11634,-96,0,0)$ |
| 31 | 1473646640067400 | $(11670,-96,0,0)$ | 32 | 7242336777667400 | $(11748,-96,0,0)$ |
| 33 | 0411474700534400 | $(11796,-96,0,0)$ | 34 | 7005761177137400 | $(11940,-96,0,0)$ |
| 35 | 0420777500236160 | $(11940,-96,0,0)$ | 36 | 0407113175431520 | $(11952,-96,0,0)$ |
| 37 | 0603237114035560 | $(12390,-96,0,0)$ | 38 | 2577350620527260 | $(12852,-96,0,0)$ |
| 39 | 0670641356075760 | $(10818,-144,0,0)$ | 40 | 1462237456233660 | $(11616,-144,0,0)$ |
| 41 | 5176656144756400 | $(12090,-144,0,0)$ | 42 | 0463117521602660 | $(12132,-144,0,0)$ |
| 43 | 0411766016336120 | $(12198,-144,0,0)$ | 44 | 1237215132353660 | $(12384,-144,0,0)$ |
| 45 | 2271227740255400 | $(12690,-144,0,0)$ | 46 | 2114462227575400 | $(13050,-144,0,0)$ |
| 47 | 1176617376233720 | $(13218,-144,0,0)$ | 48 | 0477501403733400 | $(12282,-192,0,0)$ |
| 49 | 5764337750370000 | $(14124,-288,0,0)$ |  |  |  |

We denote the 49 inequivalent pure double circulant singly even self-dual $[96,48,16]$ codes by $C_{96, i}(i=1,2, \ldots, 49)$. For these codes, the first rows $r$
of $R$ in (1) are listed in Table 4. In the table, the rows are written in octal using $0=(000), 1=(001), \ldots, 6=(110)$ and $7=(111)$.

The possible weight enumerators of singly even self-dual $[96,48, d]$ codes with $d \geq 16$ and their shadows (see [3] for the definition) are

$$
\begin{aligned}
& 1+(-5814+a) y^{16}+(97280+64 b) y^{18}+(1694208-16 a-384 b+4096 c) y^{20} \\
& +(18969600+192 b-49152 c-262144 d) y^{22} \\
& +(184315200+120 a+3328 b+237568 c+4718592 d) y^{24}+\cdots, \\
& d y^{4}+(c-22 d) y^{8}+(-b-20 c+231 d) y^{12}+(a+18 b+190 c-1540 d) y^{16} \\
& +(3231744-16 a-153 b-1140 c+7315 d) y^{20} \\
& +(369664000+120 a+816 b+4845 c-26334 d) y^{24}+\cdots,
\end{aligned}
$$

respectively, where $a, b, c, d$ are integers. It is easy to see that the number of codewords of weight 16 in the code and the numbers of vectors of weights $4,8,12$ in the shadow uniquely determine the weight enumerator. By calculating these numbers, we determined the weight enumerators of the codes $C_{96, i}$. This was done using Magma [1]. We display in Table $4(a, b, c, d)$ for the weight enumerators of the codes $C_{96, i}$. We verified by Magma [1] that each of the 49 codes has an automorphism group of order 96 .

We denote the 4565 inequivalent pure double circulant doubly even selfdual $[96,48,16]$ codes by $P_{96, i}(i=1,2, \ldots, 4565)$. The first rows $r$ of $R$ in (1) can be obtained fromhttp://www.math.is.tohoku.ac.jp/~mharada/Paper/DCCp96.txt. By the Gleason theorem (see [15]), the possible weight enumerators of doubly even self-dual $[96,48, d]$ codes with $d \geq 16$ are

$$
\begin{aligned}
& 1+a y^{16}+(3217056-16 a) y^{20}+(369844880+120 a) y^{24} \\
& +(18642839520-560 a) y^{28}+(422069980215+1820 a) y^{32}+\cdots
\end{aligned}
$$

where $a$ is an integer with $0 \leq a \leq 201066$. By calculating the number of codewords of weight 16, we verified that the 4565 codes have 614 distinct weight enumerators. The numbers $a$ in the weight enumerators are listed in Table 5. We verified by Magma [1] that 4530, 34 and 1 of the 4565 codes have automorphism groups of orders 96, 192 and 89280, respectively. For the unique code with an automorphism group of order 89280, the first row $r$ of $R$ in (1) is
(010001101111001001011111101110100100101111110000).

Table 5: Weight enumerators of pure double circulant doubly even self-dual [96, 48, 16] codes


We denote the 1532 inequivalent bordered double circulant doubly even self-dual $[96,48,16]$ codes by $B_{96, i}(i=1,2, \ldots, 1532)$. The first rows $r$ of $R^{\prime}$ in (2) can be obtained from http://www.math.is.tohoku.ac.jp/~mharada/Paper/DCCb96.txt. By calculating the numbers of codewords of weight 16 , we verified that the 1532 codes have 25 distinct weight enumerators. For each $a$, the number $N_{a}$ of codes with the weight enumerator is listed in Table 6. We verified by MAGMA [1] that the 1532 codes have automorphism groups of order 94.

Table 6: Weight enumerators of bordered double circulant doubly even selfdual $[96,48,16]$ codes

| $a$ | $N_{a}$ | $a$ | $N_{a}$ | $a$ | $N_{a}$ | $a$ | $N_{a}$ | $a$ | $N_{a}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 6204 | 1 | 6768 | 2 | 7050 | 5 | 7332 | 8 | 7614 | 17 |
| 7896 | 30 | 8178 | 72 | 8460 | 82 | 8742 | 116 | 9024 | 141 |
| 9306 | 157 | 9588 | 197 | 9870 | 141 | 10152 | 160 | 10434 | 130 |
| 10716 | 92 | 10998 | 74 | 11280 | 32 | 11562 | 31 | 11844 | 20 |
| 12126 | 12 | 12408 | 4 | 12690 | 5 | 12972 | 1 | 13254 | 2 |

## 7 Neighbors of double circulant self-dual [96, 48, 16] codes

Let $C$ be a double circulant self-dual $[96,48,16]$ code. Let $M$ be a matrix whose rows are the codewords of weight 16 in $C$. We verified by Magma [1] that

$$
\left(\operatorname{rank} M, \operatorname{rank}\left(M \mathbf{1}^{T}\right)\right)=(47,48),
$$

for $C=C_{96, i}(i=1,2, \ldots, 49)$ and

$$
\left(\operatorname{rank} M, \operatorname{rank}\left(M \mathbf{1}^{T}\right)\right)=(48,49)
$$

for $C=P_{96, i}(i=1,2, \ldots, 4565)$ and $C=B_{96, i}(i=1,2, \ldots, 1532)$. By the method given in Section 3, we have the following results.

Proposition 11. No double circulant self-dual $[96,48,16]$ code has a selfdual $[96,48, d]$ neighbor with $d \geq 18$.

It is still an open problem whether an extremal doubly even self-dual code or an extremal singly even self-dual code of length 96 exists.

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