# On the sizes of vertex-k-maximal r-uniform hypergraphs \*

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Abstract Let H = (V, E) be a hypergraph, where V is a set of vertices and E is a set of non-empty subsets of V called edges. If all edges of H have the same cardinality r, then H is a r-uniform hypergraph; if E consists of all r-subsets of V, then H is a complete runiform hypergraph, denoted by  $K_n^r$ , where n = |V|. A hypergraph H' = (V', E') is called a subhypergraph of H = (V, E) if  $V' \subseteq V$  and  $E' \subseteq E$ . A r-uniform hypergraph H = (V, E) is vertex-k-maximal if every subhypergraph of H has vertex-connectivity at most k, but for any edge  $e \in E(K_n^r) \setminus E(H)$ , H + e contains at least one subhypergraph with vertex-connectivity at least k+1. In this paper, we first prove that for given integers n, k, r with  $k, r \ge 2$  and  $n \ge k+1$ , every vertex-k-maximal r-uniform hypergraph H of order n satisfies  $|E(H)| \ge {n \choose r} - {n-k \choose r}$ , and this lower bound is best possible. Next, we conjecture that for sufficiently large n, every vertexk-maximal r-uniform hypergraph H on n vertices satisfies  $|E(H)| \le {n \choose r} - {n-k \choose k} + {n \choose k} - 2){k \choose r}$ , where  $k, r \ge 2$  are integers. And the conjecture is verified for the case r > k.

Keywords: Vertex-connectivity; Vertex-k-maximal hypergraphs; r-uniform hypergraphs

### 1 Introduction

In this paper, we consider finite simple graphs. For graph-theoretical terminologies and notation not defined here, we follow [4]. For a graph G, we use  $\kappa(G)$  to denote the *vertex-connectivity* of G. The *complement* of a graph G is denoted by  $G^c$ . For  $X \subseteq E(G^c)$ , G + X is the graph with vertex set V(G) and edge set  $E(G) \cup X$ . We will use G + e for  $G + \{e\}$ . The *floor* of a real number x, denoted by  $\lfloor x \rfloor$ , is the greatest integer not larger than x; the *ceil* of a real number x, denoted by  $\lfloor x \rfloor$ , is the least integer greater than or equal to x. For two integers n and k, we define  $\binom{n}{k} = \frac{n!}{k!(n-k)!}$  when  $k \leq n$  and  $\binom{n}{k} = 0$  when k > n.

Matula [14] first explicitly studied the quantity  $\overline{\kappa}(G) = max\{\kappa(G') : G' \subseteq G\}$ . For a positive integer k, the graph G is *vertex-k-maximal* if  $\overline{\kappa}(G) \leq k$  but for any edge  $e \in E(G^c)$ ,  $\overline{\kappa}(G+e) > k$ . Because  $\kappa(K_n) = n-1$ , a vertex-k-maximal graph G with at most k+1 vertices must be a complete graph.

The union of two graphs  $G_1$  and  $G_2$ , denoted by  $G_1 \cup G_2$ , is the graph with vertex set  $V(G_1) \cup V(G_2)$  and edge set  $E(G_1) \cup E(G_2)$ . The join of two graphs  $G_1$  and  $G_2$ , denoted

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by  $G_1 \vee G_2$ , is the graph obtained from the union of  $G_1$  and  $G_2$  by adding all the edges that connect the vertices of  $G_1$  with  $G_2$ . Let  $G_{n,k} = ((p-1)K_k \cup K_q) \vee K_k^c$ , where  $n = pk + q \ge 2k$  $(1 \le q \le k)$  and  $(p-1)K_k$  is the union of p-1 complete graphs on k vertices. Then  $G_{n,k}$  is vertex-k-maximal and  $|E(G_{n,k})| \le \frac{3}{2}(k-\frac{1}{3})(n-k)$ , where the equality holds if n is a multiple of k. Mader [11] conjectured that, for large order of graphs, the graph  $G_{n,k}$  would in fact present the best possible upper bound for the sizes of a vertex-k-maximal graph.

**Conjecture 1.** (Mader [11]) Let  $k \ge 2$  be an integer. Then for sufficiently large n, every vertex-k-maximal graph on n vertices satisfies  $|E(G)| \le \frac{3}{2}(k - \frac{1}{3})(n - k)$ .

Some progresses towards Conjecture 1 are listed in the following.

**Theorem 1.1.** Let  $k \ge 2$  be an integer.

(i) (Mader [10], see also [11]) Conjecture 1 holds for  $k \leq 6$ .

(ii) (Mader [10], see also [11]) For sufficiently large n, every vertex-k-maximal graph G on n vertices satisfies  $|E(G)| \leq (1 + \frac{1}{\sqrt{2}})k(n-k)$ .

(iii) (Yuster [18]) If  $n \ge \frac{9k}{4}$ , then every vertex-k-maximal graph G on n vertices satisfies  $|E(G)| \le \frac{193}{120}k(n-k)$ .

(iv) (Bernshteyn and Kostochka [3]) If  $n \geq \frac{5k}{2}$ , then every vertex-k-maximal graph G on n vertices satisfies  $|E(G)| \leq \frac{19}{12}k(n-k)$ .

In [17], Xu, Lai and Tian obtained the lower bound of the sizes of vertex-k-maximal graphs.

**Theorem 1.2.** (Xu, Lai and Tian [17]) Let n, k be integers with  $n \ge k + 1 \ge 3$ . If G is a vertex-k-maximal graph on n vertices, then  $|E(G)| \ge (n-k)k + \frac{k(k-1)}{2}$ . Furthermore, this bound is best possible.

The related studies on edge-k-maximal graphs have been conducted by quite a few researchers, as seen in [7,9,12,13,15], among others. For corresponding digraph problems, see [1,8], among others.

Let H = (V, E) be a hypergraph, where V is a finite set and E is a set of non-empty subsets of V, called edges. An edge of cardinality 2 is just a graph edge. For a vertex  $u \in V$  and an edge  $e \in E$ , we say u is *incident with* e or e is *incident with* u if  $u \in e$ . If all edges of H have the same cardinality r, then H is a r-uniform hypergraph; if E consists of all r-subsets of V, then H is a *complete* r-uniform hypergraph, denoted by  $K_n^r$ , where n = |V|. For n < r, the complete r-uniform hypergraph H = (V, E), denoted by  $H^c$ , is the r-uniform hypergraph with vertex set V and edge set consisting of all r-subsets of V not in E. A hypergraph H' = (V', E') is called a *subhypergraph* of H = (V, E), denoted by  $H' \subseteq H$ , if  $V' \subseteq V$  and  $E' \subseteq E$ . For  $X \subseteq E(H^c)$ , H + X is the hypergraph with vertex set V(H) and edge set  $E(H) \cup X$ ; for  $X' \subseteq E(H)$ , H - X'is the hypergraph with vertex set V(H) and edge set  $E(H) \setminus X'$ . We use H + e for  $H + \{e\}$  and H - e' for  $H - \{e'\}$  when  $e \in E(H^c)$  and  $e' \in E(H)$ . For  $Y \subseteq V(H)$ , we use H[Y] to denote the hypergraph induced by Y, where V(H[Y]) = Y and  $E(H[Y]) = \{e \in E(H) : e \subseteq Y\}$ . H - Y is the hypergraph induced by  $V(H) \setminus Y$ . Let H be a hypergraph and  $V_1, V_2, \dots, V_l$  be subsets of V(H). An edge  $e \in E(H)$  is  $(V_1, V_2, \dots, V_l)$ -crossing if  $e \cap V_i \neq \emptyset$  for  $1 \leq i \leq l$ . If in addition,  $e \subseteq \bigcup_{i=1}^l V_i$ , then e is exact- $(V_1, V_2, \dots, V_l)$ -crossing. The set of all  $(V_1, V_2, \dots, V_l)$ -crossing edges of H is denoted by  $E_H[V_1, V_2, \dots, V_l]$ ; the set of all exact- $(V_1, V_2, \dots, V_l)$ -crossing edges of H is denoted by  $E_{H[V_1 \cup V_2 \cup \dots \cup V_l]}[V_1, V_2, \dots, V_l]$ . Let  $d_H(V_1, V_2, \dots, V_l) = |E_H[V_1, V_2, \dots, V_l]|$  and  $d_{H[V_1 \cup V_2 \cup \dots \cup V_l]}(V_1, V_2, \dots, V_l) = |E_{H[V_1 \cup V_2 \cup \dots \cup V_l]}[V_1, V_2, \dots, V_l]$ . For a vertex  $u \in V(H)$ , we call  $d_H(u) := d_H(\{u\}, V(H) \setminus \{u\})$  the degree of u in H. The minimum degree  $\delta(H)$  of H is defined as  $min\{d_H(u) : u \in V\}$ ; the maximum degree  $\Delta(H)$  of H is defined as  $max\{d_H(u) : u \in V\}$ . When  $\delta(H) = \Delta(H) = k$ , we call H k-regular.

Given a hypergraph H, we define a walk in H to be an alternating sequence  $v_1, e_1, v_2, \dots, e_s$ ,  $v_{s+1}$  of vertices and edges of H such that:  $v_i \in V(H)$  for  $i = 1, \dots, s + 1$ ;  $e_i \in E(H)$  for  $i = 1, \dots, s$ ; and  $v_i, v_{i+1} \in e_i$  for  $i = 1, \dots, s$ . A path is a walk with additional restrictions that the vertices are all distinct and the edges are all distinct. A hypergraph H is connected if for every pair of vertices  $u, v \in V(H)$ , there is a path connecting u and v; otherwise H is disconnected. A component of a hypergraph H is a maximal connected subhypergraph of H. A subset  $X \subseteq V$  is called a vertex-cut of H if H - X is disconnected. We define the vertexconnectivity of H, denoted by  $\kappa(H)$ , as follows: if H had at least one vertex-cut, then  $\kappa(H)$ is the cardinality of a minimum vertex-cut of H; otherwise  $\kappa(H) = |V(H)| - 1$ . We call a hypergraph H k-vertex-connected if  $\kappa(H) \geq k$ . Let  $\overline{\kappa}(H) = max\{\kappa(H') : H' \subseteq H\}$ . For a positive integer k, the r-uniform hypergraph H is vertex-k-maximal if  $\overline{\kappa}(H) \leq k$  but for any edge  $e \in E(H^c), \overline{\kappa}(H + e) > k$ . Since  $\kappa(K_n^r) = n - r + 1$ , we note that H is complete if H is a vertex-k-maximal r-uniform hypergraph with  $n - r + 1 \leq k$ , where n = |V(H)|. The edge-kmaximal hypergraph can be defined similarly. For results on the connectivity of hypergraphs, see [2,5,6] for references.

In [16], we determined, for given integers n, k and r, the extremal sizes of an edge-k-maximal r-uniform hypergraph on n vertices.

**Theorem 1.3.** (Tian, Xu, Lai and Meng [16]) Let k and r be integers with  $k, r \ge 2$ , and let t = t(k,r) be the largest integer such that  $\binom{t-1}{r-1} \le k$ . That is, t is the integer satisfying  $\binom{t-1}{r-1} \le k < \binom{t}{r-1}$ . If H is an edge-k-maximal r-uniform hypergraph with  $n = |V(H)| \ge t$ , then

- (i)  $|E(H)| \leq {t \choose r} + (n-t)k$ , and this bound is best possible;
- (ii)  $|E(H)| \ge (n-1)k ((t-1)k \binom{t}{r})\lfloor \frac{n}{t} \rfloor$ , and this bound is best possible.

The main goal of this research is to investigate, for given integers n, k and r, the extremal sizes of a vertex-k-maximal r-uniform hypergraph on n vertices. Section 2 below is devoted to the study of some properties of vertex-k-maximal r-uniform hypergraphs. In Section 3, we give the best possible lower bound of the sizes of vertex-k-maximal r-uniform hypergraphs. We propose a conjecture on the upper bound of the sizes of vertex-k-maximal r-uniform hypergraphs and verify the conjecture for the case r > k in Section 4.

#### 2 Properties of vertex-k-maximal r-uniform hypergraphs

Combining the definition of vertex-k-maximal r-uniform hypergraph with  $\kappa(K_n^r) = n - r + 1$ , we obtain that H is isomorphic to  $K_n^r$  if H is a vertex-k-maximal r-uniform hypergraph with  $n = |V(H)| \le k + r - 1.$ 

**Lemma 2.1.** Let n, k, r be integers with  $k, r \ge 2$  and  $n \ge k + r - 1$ . If H is a vertex-k-maximal r-uniform hypergraph on n vertices, then  $\overline{\kappa}(H) = \kappa(H) = k$ .

**Proof.** Since H is vertex-k-maximal, we have  $\kappa(H) \leq \overline{\kappa}(H) \leq k$ . In order to complete the proof, we only need to show that  $\kappa(H) \geq k$ .

If n = k + r - 1, then H is complete and  $\kappa(H) = n - r + 1 = k$ . Thus, assume  $n \ge k + r$ , and so H is not complete. On the contrary, assume  $\kappa(H) < k$ . Since H is not complete, H has a vertex-cut S with  $|S| = \kappa(H) < k$ . Let  $C_1$  be a component of H - S and  $C_2 = H - (S \cup V(C_1))$ . By  $|V(C_1) \cup V(C_2)| = n - |S| \ge k + r - (k - 1) = r + 1$ , we can choose a r-subset  $e \subseteq V(C_1) \cup V(V_2)$ such that  $e \cap V(C_i) \neq \emptyset$  for i = 1, 2. Then  $e \in E(H^c)$ .

Since H is vertex-k-maximal, we have  $\overline{\kappa}(H+e) \ge k+1$ . Hence H+e contains a subhypergraph H' with  $\kappa(H') = \overline{\kappa}(H+e) \ge k+1$ . Since  $\overline{\kappa}(H) \le k$ , H' cannot be a subhypergraph of H, and so  $e \in E(H')$ . Since  $V(H') \cap V(C_i) \ne \emptyset$  for i = 1, 2, it follows that  $V(H') \cap S$  is a vertex-cut of H' - e.

Since  $|V(C_1) \cup V(C_2)| = n - |S| \ge k + r - (k - 1) = r + 1 \ge 3$ , one of  $C_i$ , say  $C_1$ , contains at least two vertices. Let  $u_1 \in e \cap V(C_1)$ . Then  $S' = (V(H') \cap S) \cup \{u_1\}$  is a vertex-cut of H', and so we obtain

$$k+1 > |S|+1 \ge |V(H') \cap S|+1 = |S'| \ge \kappa(H') \ge k+1,$$

a contradiction.  $\Box$ 

Let *H* be a vertex-*k*-maximal *r*-uniform hypergraph with  $|V(H)| \ge k + r$ . By Lemma 2.1,  $\overline{\kappa}(H) = \kappa(H) = k$ . By  $|V(H)| \ge k + r$ , *H* is not complete, thus *H* contains vertex-cuts. Let *S* be a minimum vertex-cut of *H*, *C*<sub>1</sub> be a component of *H* - *S* and *C*<sub>2</sub> = *H* - (*S*  $\cup$  *V*(*C*<sub>1</sub>)). We call (*S*, *H*<sub>1</sub>, *H*<sub>2</sub>) a separation triple of *H*, where *H*<sub>1</sub> = *H*[*S*  $\cup$  *V*(*C*<sub>1</sub>)] and *H*<sub>2</sub> = *H*[*S*  $\cup$  *V*(*C*<sub>2</sub>)].

**Lemma 2.2.** Let n, k, r be integers with  $k, r \ge 2$  and  $n \ge k + r$ , and H be a vertex-k-maximal r-uniform hypergraph on n vertices. Assume  $(S, H_1, H_2)$  is a separation triple of H. If  $e \in E(H_1^c) \cup E(H_2^c)$ , then any subhypergraph H' of H + e with  $\kappa(H') \ge k+1$  is either a subhypergraph of  $H_1 + e$  or a subhypergraph of  $H_2 + e$ . Furthermore, if  $e \subseteq E(H_i^c) \setminus E((H[S])^c)$ , then H' is a subhypergraph of  $H_i + e$  for i = 1, 2.

**Proof.** Let  $e \in E(H_1^c) \cup E(H_2^c)$ . Since H is vertex-k-maximal, we have  $\overline{\kappa}(H+e) \geq k+1$ . Let H' be a subhypergraph of H + e with  $\kappa(H') = \overline{\kappa}(H+e) \geq k+1$ . We assume, on the contrary, that  $V(H') \cap (V(H_1) - S) \neq \emptyset$  and  $V(H') \cap (V(H_2) - S) \neq \emptyset$ . This, together with  $e \in E(H_1^c) \cup E(H_2^c)$ , implies that  $S \cap V(H')$  is a vertex-cut of H'. Hence  $k = |S| \geq |S \cap V(H')| \geq \kappa(H') \geq k+1$ , a contradiction. Therefore, we cannot have both  $V(H') \cap (V(H_1) - S) \neq \emptyset$  and  $V(H') \cap (V(H_2) - S) \neq \emptyset$ . If  $V(H') \cap (V(H_1) - S) = \emptyset$ , then H' is a subhypergraph of  $H_2 + e$ ; if  $V(H') \cap (V(H_2) - S) = \emptyset$ , then H' is a subhypergraph of  $H_1 + e$ .

If  $e \subseteq E(H_1^c) \setminus E((H[S])^c)$ , then  $V(H') \cap (V(H_1) - S) \neq \emptyset$  and  $V(H') \cap (V(H_2) - S) = \emptyset$ , thus H' is a subhypergraph of  $H_1 + e$ . Similarly, if  $e \subseteq E(H_2^c) \setminus E((H[S])^c)$ , then H' is a subhypergraph of  $H_2 + e$ .  $\Box$  **Lemma 2.3.** Let n, k, r be integers with  $k, r \ge 2$  and  $n \ge k + r$ , and H be a vertex-k-maximal r-uniform hypergraph on n vertices. Assume  $(S, H_1, H_2)$  is a separation triple of H and  $n_i = |V(H_i)|$  for i = 1, 2. Then

(i) 
$$E_{H^c}[V(H_1) - S, S, V(H_2) - S] = \emptyset$$
, and  
(ii)  $d_H(V(H_1) - S, S, V(H_2) - S) = \binom{n}{r} - \binom{n_1}{r} - \binom{n_2}{r} + \binom{k}{r} - \binom{n-k}{r} + \binom{n_1-k}{r} + \binom{n_2-k}{r}$ .

**Proof.** (i) By contradiction, assume  $E_{H^c}[V(H_1) - S, S, V(H_2) - S] \neq \emptyset$ . Let  $e \in E_{H^c}[V(H_1) - S, S, V(H_2) - S]$ . Since H is vertex-k-maximal, there is a subhypergraph H' of H + e such that  $\kappa(H') = \overline{\kappa}(H + e) \geq k + 1$ . By  $\overline{\kappa}(H) \leq k$ ,  $e \in E(H')$ . This, together with  $e \in E_{H^c}[V(H_1) - S, S, V(H_2) - S]$ , implies  $V(H') \cap S \neq \emptyset$  and  $V(H') \cap (V(H_i) - S) \neq \emptyset$  for i = 1, 2. Hence  $S \cap V(H')$  is a vertex-cut of H'. But then we obtain  $k = |S| \geq |S \cap V(H')| \geq \kappa(H') \geq k + 1$ , a contradiction. It follows  $E_{H^c}[V(H_1) - S, S, V(H_2) - S] = \emptyset$ .

(*ii*) By (*i*),  $E_{H^c}[V(H_1) - S, S, V(H_2) - S] = \emptyset$ . This implies that if e is a r-subset such that  $e \cap S \neq \emptyset$  and  $e \cap (V(H_i) - S) \neq \emptyset$  for i = 1, 2, then  $e \in E(H)$ . Since the number of r-subsets contained in  $V(H_1)$  or  $V(H_2)$  is  $\binom{n_1}{r} + \binom{n_2}{r} - \binom{k}{r}$ , and the number of r-subsets exactly intersecting  $V(H_1) - S$  and  $V(H_1) - S$  is  $\binom{n-k}{r} - \binom{n_1-k}{r} - \binom{n_2-k}{r}$ , we have

$$d_H(V(H_1) - S, S, V(H_2) - S)$$
  
=  $|E_H[V(H_1) - S, S, V(H_2) - S]|$   
=  $\binom{n}{r} - \binom{n_1}{r} + \binom{n_2}{r} - \binom{k}{r} - \binom{n_{-k}}{r} - \binom{n_{1-k}}{r} - \binom{n_2-k}{r}$   
=  $\binom{n}{r} - \binom{n_1}{r} - \binom{n_2}{r} + \binom{k}{r} - \binom{n_{-k}}{r} + \binom{n_{1-k}}{r} + \binom{n_2-k}{r}.$ 

This completes the proof.  $\Box$ 

# 3 The lower bound of the sizes of vertex-k-maximal r-uniform hypergraphs

The union of two hypergraphs  $H_1$  and  $H_2$ , denoted by  $H_1 \cup H_2$ , is the hypergraph with vertex set  $V(H_1) \cup V(H_2)$  and edge set  $E(H_1) \cup E(H_2)$ . The *r*-join of two hypergraphs  $H_1$  and  $H_2$ , denoted by  $H_1 \vee_r H_2$ , is the hypergraph obtained from the union of  $H_1$  and  $H_2$  by adding all the edges with cardinality r that connect the vertices of  $H_1$  with  $H_2$ .

**Definition 1.** Let n, k, r be integers such that  $k, r \ge 2$  and  $n \ge k+1$ . We define  $H_L(n; k, r)$  to be  $K_k^r \vee_r (K_{n-k}^r)^c$ .

**Lemma 3.1.** Let n, k, r be integers such that  $k, r \geq 2$  and  $n \geq k+1$ . If  $H = H_L(n; k, r)$ , then

- (i) H is vertex-k-maximal, and
- (*ii*)  $|E(H)| = \binom{n}{r} \binom{n-k}{r}$ .

**Proof.** (i) By Definition 1, H is obtained from the union of  $K_k^r$  and  $(K_{n-k}^r)^c$  by adding all edges with cardinality r connecting  $V(K_k^r)$  with  $V((K_{n-k}^r)^c)$ .

Since  $V(K_k^r)$  is a vertex-cut of H and  $H - V(K_k^r) = (K_{n-k}^r)^c$ , there is no subhypergraph with vertex-connectivity at least k + 1, and so  $\overline{\kappa}(H) \leq k$ . If  $E(H^c) = \emptyset$ , then H is vertex-k-maximal

by the definition of vertex-k-maximal hypergraph. If  $E(H^c) \neq \emptyset$ , then for any  $e \in E(H^c)$ , e must be contained in  $V((K_{n-k}^r)^c)$ , and so  $(H+e)[V(K_k^r) \cup e]$  is isomorphic to  $K_{k+r}^r$  and  $\kappa((H+e)[V(K_k^r) \cup e]) = k+1$ . That is  $\overline{\kappa}(H+e) \geq k+1$ . Thus H is vertex-k-maximal.

(ii) holds by a direct calculation.  $\Box$ 

**Theorem 3.2.** Let n, k, r be integers such that  $k, r \ge 2$  and  $n \ge k+1$ . If H is vertex-k-maximal, then  $|E(H)| \ge {n \choose r} - {n-k \choose r}$ .

**Proof.** We will prove the theorem by induction on n. If  $n \leq k + r - 1$ , then by H is vertex-k-maximal, we have  $H \cong K_n^r$ . Thus  $|E(H)| = \binom{n}{r} = \binom{n}{r} - \binom{n-k}{r}$  by  $n - k \leq r - 1$ .

Now we assume that  $n \ge k+r$ , and that the theorem holds for smaller value of n. Since H is vertex-k-maximal and  $n \ge k+r$ , we have H is not complete. By Lemma 2.1,  $\overline{\kappa}(H) = \kappa(H) = k$ , and so H has a separation triple  $(S, H_1, H_2)$  with |S| = k. Let  $n_1 = |V(H_1)|$  and  $n_2 = |V(H_2)|$ . Then  $n_1, n_2 \ge k+1$  and  $n = n_1 + n_2 - k$ .

Since H is vertex-k-maximal, for any  $e \in E((H[S])^c)$ , there is a (k + 1)-vertex-connected subhypergraph H' of H + e. By Lemma 2.2, H' is either a subhypergraph of  $H_1 + e$  or a subhypergraph  $H_2 + e$ . Define

$$E_1 = \{e : e \in E((H[S])^c) \text{ and } \overline{\kappa}(H_1 + e) = k\}$$
$$E_2 = \{e : e \in E((H[S])^c) \text{ and } \overline{\kappa}(H_2 + e) = k\}$$

Claim. Each of the following holds.

- (i)  $E_1 \cap E_2 = \emptyset$  and  $E_1 \cup E_2 \subseteq E((H[S])^c)$ .
- (*ii*) There is a subset  $E'_1 \subseteq E_1$  such that  $H_1 + E'_1$  is vertex-k-maximal.
- (*iii*) There is a subset  $E'_2 \subseteq E_2$  such that  $H_2 + E'_2$  is vertex-k-maximal.

By the definition,  $E_1 \cup E_2 \subseteq E((H[S])^c)$ . Since H is vertex-k-maximal, we have  $E_1 \cap E_2 = \emptyset$ , and so Claim (i) holds.

Assume first that  $H_1 + E_1$  is complete. If  $n_1 \leq k + r - 1$ , then  $\overline{\kappa}(H_1 + E_1) \leq k$ , and so  $H_1 + E_1$ is vertex-k-maximal by the definition of vertex-k-maximal hypergraphs. If  $n_1 \geq k + r$ , then by  $\overline{\kappa}(H_1) \leq \overline{\kappa}(H) \leq k$  and  $\overline{\kappa}(H_1 + E_1) \geq k + 1$ , we can choose a maximum subset  $E'_1 \subseteq E_1$  such that  $\overline{\kappa}(H_1 + E'_1) \leq k$ . It follows by the maximality of  $E'_1$  and by the definition of vertex-k-maximal hypergraphs that  $H_1 + E'_1$  is vertex-k-maximal. Next, we assume  $H_1 + E_1$  is not complete. Take an arbitrary edge  $e \in E((H_1 + E_1)^c)$ . Then  $e \in E(H^c)$ , and so as H is vertex-k-maximal, H + econtains a (k + 1)-vertex-connected subhypergraph H' with  $e \in E(H')$ . If  $e \cap (V(H_1) - S) \neq \emptyset$ , then by Lemma 2.2, H' is a subhypergraph of  $H_1 + e$ . If  $e \subseteq S$ , then as  $e \notin E_1$ , we can choose H'such that H' is a subhypergraph of  $H_1 + e$ . That is,  $\overline{\kappa}(H_1 + E_1 + e) \geq k + 1$ . If  $\overline{\kappa}(H_1 + E_1) \leq k$ , then  $H_1 + E_1$  is vertex-k-maximal. If  $\overline{\kappa}(H_1 + E_1) \geq k + 1$ , then by  $\overline{\kappa}(H_1) \leq \overline{\kappa}(H) \leq k$ , we can choose a maximum subset  $E'_1 \subseteq E_1$  such that  $\overline{\kappa}(H_1 + E'_1) \leq k$ . It also follows by the maximality of  $E'_1$  and by the definition of vertex-k-maximal hypergraphs that  $H_1 + E'_1$  is vertexk-maximal. This verifies Claim (ii). By symmetry, Claim (iii) holds. Thus the proof of the Claim is complete.

By Claim (*ii*) and Claim (*iii*), there are  $E'_1 \subseteq E_1$  and  $E'_2 \subseteq E_2$  such that  $H_1 + E'_1$  and  $H_2 + E'_2$  are vertex-k-maximal. Since  $n_1, n_2 \ge k + 1$ , by induction assumption, we have  $|E(H_1 + E'_1)| \ge k + 1$ .

$$\begin{aligned} \binom{n_1}{r} - \binom{n_1 - k}{r} & \text{and } |E(H_2 + E'_2)| \ge \binom{n_2}{r} - \binom{n_2 - k}{r}. \text{ By Claim } (i) \text{ and the definition of } (H[S])^c, \text{ we} \\ & \text{have } |E'_1| + |E'_2| + |E(H[S])| \le |E_1| + |E_2| + |E(H[S])| \le |E((H[S])^c)| + |E(H[S])| = \binom{k}{r}. \text{ Thus } \\ & |E(H)| = |E(H_1)| + |E(H_2)| - |E(H[S])| + |E_H[V(H_1) - S, S, V(H_2) - S]| \\ & = |E(H_1 + E'_1)| - |E'_1| + |E(H_2 + E'_2)| - |E'_2| - |E(H[S])| + |E_H[V(H_1) - S, S, V(H_2) - S]| \\ & \ge \binom{n_1}{r} - \binom{n_1 - k}{r} + \binom{n_2}{r} - \binom{n_2 - k}{r} - \binom{k}{r} \\ & + \binom{n}{r} - \binom{n_1}{r} - \binom{n_2}{r} + \binom{k}{r} - \binom{n_1 - k}{r} + \binom{n_1 - k}{r} + \binom{n_2 - k}{r} \text{ (By Lemma 2.3)} \\ & = \binom{n}{r} - \binom{n - k}{r}. \end{aligned}$$

This proves Theorem 3.2.  $\Box$ 

By Lemma 3.1, the lower bound of the sizes of vertex-k-maximal hypergraphs given in Theorem 3.2 is best possible. If r = 2, then a r-uniform hypergraph H is just a graph. Thus Theorem 1.2 is a corollary of Theorem 3.2.

**Corollary 3.3.** (Xu, Lai and Tian [17]) Let n, k be integers with  $n \ge k+1 \ge 3$ . If G is a vertexk-maximal graph on n vertices, then  $|E(G)| \ge {\binom{n}{2}} - {\binom{n-k}{2}} = (n-k)k + \frac{k(k-1)}{2}$ . Furthermore, this bound is best possible.

## 4 The upper bound of the sizes of vertex-k-maximal r-uniform hypergraphs

**Definition 2.** Let n, k, r be integers such that  $k, r \ge 2$  and  $n \ge 2k$ . Assume n = pk + q, where p, q are integers and  $1 \le q \le k$ . We define  $H_U(n; k, r)$  to be  $((p-1)K_k^r \cup K_q^r) \lor_r (K_k^r)^c$ , where  $(p-1)K_k^r$  is the union of p-1 complete r-uniform hypergraphs on k vertices.

**Lemma 4.1.** Let n, k, r be integers such that  $k, r \ge 2$  and  $n \ge 2k$ . If  $H = H_U(n; k, r)$ , then

- (i) H is vertex-k-maximal, and
- (ii)  $|E(H)| \leq {n \choose r} {n-k \choose k} + (\frac{n}{k} 2){k \choose r}$ , where the equality holds if n is a multiple of k.

**Proof.** (i) By Definition 2,  $H = ((p-1)K_k^r \cup K_q^r) \vee_r (K_k^r)^c$ . Denote the p-1 complete r-uniform hypergraphs on k vertices by  $K_k^r(1), \cdots, K_k^r(p-1)$ . Let  $H_0 = H[V((K_k^r)^c)], H_p = H[V(K_q^r)]$  and  $H_i = H[V(K_k^r(i))]$  for  $1 \le i \le p-1$ . Then  $H = H_0 \vee_r (H_1 \cup \cdots \cup H_p)$ .

Since  $V(H_0)$  is a vertex-cut of size k and every component of  $H - V(H_0)$  has at most k vertices. It follows that H contains no (k + 1)-vertex-connected subhypergraphs, and so  $\overline{\kappa}(H) \leq k$ . If  $E(H^c) = \emptyset$ , then H is vertex-k-maximal by the definition of vertex-k-maximal hypergraphs. Thus we assume  $E(H^c) \neq \emptyset$  in the following. Let  $e \in E(H^c)$ . If  $e \subseteq V(H_0)$ , then  $H' = H[V(H_1) \cup e]$  is isomorphic to  $K_{k+r}^r$ , and so  $\kappa(H') = k + 1$ . If  $e \subseteq V(H_1) \cup \cdots \cup V(H_p)$ , let e be exact- $(V(H_{i1}), \cdots, V(H_{is}))$ -crossing. We will prove that  $H'' = H[V(H_0) \cup V(H_{i1}) \cup \cdots \cup (H_{is})] + e$  is (k + 1)-vertex-connected. It suffices to prove that H'' - S is connected for any  $S \subseteq V(H'')$  with |S| = k. If  $S = V(H_0)$ , then, by e is exact- $(V(H_{i1}) \cup \cdots \cup V(H_{is}))$ -crossing, H'' - S is connected. So assume  $V'_0 = V(H_0) \setminus S \neq \emptyset$ . Let  $V'_1 = (V(H_{i1}) \cup \cdots \cup V(H_{is})) \setminus S$ . Then H'' - S is isomorphic to  $H[V'_0] \lor_r H[V'_1]$  if  $S \cap e \neq \emptyset$ ; and H'' - S is connected. Thus  $\overline{\kappa}(H + e) \geq k + 1$  for any  $e \in E(H^c)$ , and so H is vertex-k-maximal.

(*ii*) By a direct calculation, we have  $|E(H)| \leq {\binom{n}{r}} - {\binom{n-k}{r}} + {\binom{n}{k}} - 2){\binom{k}{r}}$ , where the equality holds if n is a multiple of k.  $\Box$ 

Motivated by Conjecture 1, we propose the following conjecture for vertex-k-maximal r-uniform hypergraphs.

**Conjecture 2.** Let k, r be integers with  $k, r \ge 2$ . Then for sufficiently large n, every vertex-k-maximal r-uniform hypergraph H on n vertices satisfies  $|E(H)| \le {n \choose r} - {n-k \choose r} + {n \choose k} - 2){k \choose r}$ .

The following theorem confirms Conjecture 2 for the case k < r.

**Theorem 4.2.** Let n, k, r be integers such that  $k, r \ge 2$  and  $n \ge 2k$ . If k < r, then every vertexk-maximal r-uniform hypergraph H on n vertices satisfies  $|E(H)| \le {n \choose r} - {n-k \choose r} + (\frac{n}{k} - 2){k \choose r} = {n \choose r} - {n-k \choose r}$ .

**Proof.** We will prove the theorem by induction on n. If  $n \leq k + r - 1$ , then by H is vertex-k-maximal, we have  $H \cong K_n^r$ . Thus  $|E(H)| = \binom{n}{r} = \binom{n}{r} - \binom{n-k}{r}$  by  $n - k \leq r - 1$ .

Now we assume that  $n \ge k + r$ , and that the theorem holds for smaller value of n. Since H is vertex-k-maximal and  $n \ge k + r$ , we have H is not complete. Let S be a minimum vertex-cut of H. By Lemma 2.1, |S| = k. Let  $C_1$  be a minimum component of H - S and  $C_2 = H - (V(C_1) \cup S)$ . Assume  $H_1 = H[V(C_1) \cup S]$  and  $H_2 = H[V(C_2) \cup S]$ . Since k < r, we have  $E((H[S])^c) = \emptyset$ , and so  $H_1$  and  $H_2$  are both vertex-k-maximal by Lemma 2.2. Let  $n_1 = |V(H_1)|$  and  $n_2 = |V(H_2)|$ . Then  $n = n_1 + n_2 - k$  and  $k + 1 \le n_1 \le n_2$ . We consider two cases in the following.

**Case 1.**  $|V(C_1)| = 1$ .

By  $|V(C_1)| = 1$ , we obtain that  $n_2 = n - 1 \ge k + r - 1 \ge 2k$ . Since  $H_2$  is vertex-k-maximal, by induction assumption, we have  $|E(H_2)| \le {\binom{n-1}{r}} - {\binom{n-k-1}{r}}$ . Thus

$$|E(H)| = |E(H_1)| + |E(H_2)| - |E(H[S])| + |E_H[V(H_1) - S, S, V(H_2) - S]|$$
  

$$\leq \binom{k}{r-1} + \binom{n-1}{r} - \binom{n-k-1}{r-1} + \binom{n-1}{r-1} - \binom{k}{r-1} - \binom{n-k-1}{r-1}$$
  

$$= \binom{n}{r} - \binom{n-k}{r}.$$

**Case 2.**  $|V(C_1)| \ge 2$ .

By  $|V(C_1)| \ge 2$ , we obtain that  $C_1$  contains edges, and so  $|V(C_1)| \ge r$ . Thus  $n_2 \ge n_1 \ge k+r \ge 2k+1$ . Since both  $H_1$  and  $H_2$  are vertex-k-maximal, by induction assumption, we have  $|E(H_i)| \le {n_i \choose r} - {n_i-k \choose r}$  for i = 1, 2. Thus

$$\begin{aligned} |E(H)| &= |E(H_1)| + |E(H_2)| - |E(H[S])| + |E_H[V(H_1) - S, S, V(H_2) - S]| \\ &\leq \binom{n_1}{r} - \binom{n_1 - k}{r} + \binom{n_2}{r} - \binom{n_2 - k}{r^2} \\ &+ \binom{n}{r} - \binom{n_1}{r} - \binom{n_2}{r} + \binom{k}{r} - \binom{n-k}{r} + \binom{n_1 - k}{r} + \binom{n_2 - k}{r^2}$$
(By Lemma 2.3)  
$$&= \binom{n}{r} - \binom{n-k}{r}. \end{aligned}$$

This completes the proof.  $\Box$ 

Combining Theorem 3.2 with Theorem 4.2, we have the following corollary.

**Corollary 4.3.** Let n, k, r be integers such that  $k, r \ge 2$  and  $n \ge 2k$ . If k < r, then every vertex-k-maximal r-uniform hypergraph H on n vertices satisfies  $|E(H)| = \binom{n}{r} - \binom{n-k}{r}$ .

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