

Expanding FL_{ew} with a Boolean connective

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Abstract We expand FL_{ew} with a unary connective whose algebraic counterpart is the operation that gives the greatest complemented element below a given argument. We prove that the expanded logic is conservative and has the Finite Model Property. We also prove that the corresponding expansion of the class of residuated lattices is an equational class.

ACKNOWLEDGEMENT

This paper was published in Soft Computing (Springer-Verlag) on 23 July 2016. The final publication is available at <http://dx.doi.org/10.1007/s00500-016-2275-y>

1 Introduction

In this paper we study the expansion of the substructural logic FL_{ew} , i.e. Full Lambek calculus with exchange and weakening, with a unary connective B whose intended algebraic semantics is as follows: given a bounded integral commutative residuated lattice (or residuated lattice for short) \mathbf{A} , Ba is the maximum, if it exists, of the Boolean elements of the universe A below a , which we call *the greatest Boolean below a* , that is,

$$Ba = \max\{b \in A : b \leq a \text{ and } b \text{ is Boolean}\}.$$

In fact, this operator is similar to the so-called Baaz-Monteiro Δ operator, very often used in the context of

mathematical fuzzy logic systems that are semilinear expansions of MTL. Baaz [2] studied it in connection with Gödel logic while Hájek [10] investigated Δ in BL logics in general, see also [8, Chapter 2] for a more general perspective. Indeed, in such a context of semilinear logics, i.e. logics that are complete with respect to a class of linearly ordered algebras, the semantics of Δ is exactly the above one for B : in a linearly-ordered MTL-algebra, $\Delta a = 1$ if $a = 1$, and $\Delta a = 0$ otherwise, since the only Boolean elements in a chain are 1 and 0; moreover, from a logical point of view, $\Delta\varphi$ represents the weakest Boolean proposition implying φ .

The operator B can be also related to the join-complement operation D , also known as dual intuitionistic negation, already considered by Skolem [20] in the context of lattices with relative meet-complement, and later independently studied by e.g. Moisil [11] and Rauszer [16] as well, the latter in the context of expansions of Heyting algebras. It turns out that the operation $\neg D$ and its iterations, where \neg is the residual negation, has also very similar properties to B , and in some classes of residuated lattices they even coincide.

In this paper we study the operator B in the context of FL_{ew} and axiomatize it. We show that the usual axiomatics of the Δ operator is actually too strong to capture the above intended semantics. In fact, the axiom

$$\Delta(\varphi \vee \psi) \rightarrow (\Delta\varphi \vee \Delta\psi)$$

is not sound for B over FL_{ew} any longer. Thus, B is a weaker operator than Δ . However, as we will see, B keeps most of the properties of Δ . In particular, the expansion of FL_{ew} with B is conservative, its corresponding class of algebras is an equational class, and has the same kind of deduction theorem as Δ . Also, B may also be interesting as $\neg B$ has a paraconsistent behaviour. On the negative side, the expansion of a semi-

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linear extension of FL_{ew} with B needs not to remain semilinear.

The paper is structured as follows. In Section 2 we overview well-known facts about residuated lattices and its Boolean elements, as well as basic facts about the logic FL_{ew} . Sections 3 and 4 contain an algebraic study of the operator B . In particular, in Section 3 we study basic properties and show, among other things, that the class \mathbb{RL}^B of residuated lattices expanded with B is an equational class and state the modalities, while in Section 4 we compare B with the mentioned Δ and with an operation using the join-complement D . Finally, in Section 5 we focus on logical aspects, introducing the logic FL_{ew}^B , i.e. the expansion of FL_{ew} with the operator B , and show that is a conservative expansion and has the Finite Model Property, and hence it is decidable. We conclude with some remarks and open problems.

We give appropriate references. However, the paper is self-contained.

2 Preliminaries

2.1 Residuated lattices and Boolean elements

In this section we recall some properties of residuated lattices as well as of their Boolean elements that we will use in the following sections.

Following [9], a bounded, integral, commutative residuated lattice, or residuated lattice for short, is an algebra $\mathbf{A} = (A; \wedge, \vee, \cdot, \rightarrow, 0, 1)$ of type $(2, 2, 2, 2, 0, 0)$ such that:

- $(A; \wedge, \vee, 0, 1)$ is a bounded lattice with $0 \leq a \leq 1$, for all $a \in A$,
- $(A; \cdot, 1)$ is a commutative monoid (i.e. \cdot is commutative, associative, with unit 1), and
- \rightarrow is the residuum of \cdot , i.e.,
 $a \cdot b \leq c$ iff $a \leq b \rightarrow c$, for all $a, b, c \in A$,

where \leq is the order given by the lattice structure. A negation operator is defined as $\neg x = x \rightarrow 0$.

The class of residuated lattices will be denoted by \mathbb{RL} . It is well known that \mathbb{RL} is an equational class and that it constitutes the algebraic semantics of the substructural logic FL_{ew} (see Section 2.2).

Example 1 In what follows we will have occasion to refer several times to the residuated lattice structure defined on the five-element lattice of Figure 1 by taking $\cdot = \wedge$ and \rightarrow its residuum. With these operations, it actually becomes a five element Gödel algebra, that is, a residuated lattice with \cdot being idempotent and satisfying the pre-linearity law $(a \rightarrow b) \vee (b \rightarrow a) = 1$.

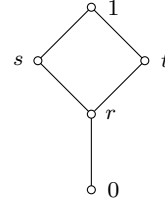


Fig. 1 A five element Gödel algebra

We omit the proof of the following well-known facts, see e.g. [9].

Lemma 1 *Let $\mathbf{A} \in \mathbb{RL}$. For any $a, b, c, d \in A$, the following properties hold:*

- (i) if $a \vee b = 1$, $a \leq c$, and $b \leq c$, then $c = 1$,
- (ii) if $a \vee b = 1$, $a \cdot c \leq d$, and $b \cdot c \leq d$, then $c \leq d$,
- (iii) if $a \leq b$, then $\neg b \leq \neg a$,
- (iv) $a \wedge \neg b \leq \neg(a \wedge b)$,
- (v) $a \cdot \neg b \leq \neg(a \rightarrow b)$,
- (vi) if $a \vee b = 1$, then $\neg a \leq b$,
- (vii) $a \leq \neg \neg a$.

Special elements in a residuated lattice are those that behave as elements in a Boolean algebra.

Definition 1 *Let $\mathbf{A} \in \mathbb{RL}$. An element a of its universe A is called Boolean or complemented iff there is an element $b \in A$ such that $a \wedge b = 0$ and $a \vee b = 1$.*

In the rest of this section we state several properties of Boolean elements that will be useful in what follows. Even if most of them are folklore, we include proofs for all of them for the sake of being self-contained.

An equivalent and simpler condition for an element to be Boolean is the following.

Lemma 2 *An element a in the universe of a residuated lattice is Boolean iff $a \vee \neg a = 1$.*

Proof \Rightarrow) Suppose there is an element b such that $a \wedge b = 0$ and $a \vee b = 1$. First, using that $a \wedge b = 0$ and $a \cdot b \leq a \wedge b$, we have that $a \cdot b = 0$. So, $b \leq a \rightarrow 0$, i.e. $b \leq \neg a$. Secondly, we have that $\neg a = \neg a \cdot 1 = \neg a \cdot (a \vee b) = (\neg a \cdot a) \vee (\neg a \cdot b) = \neg a \cdot b$. So, $\neg a = \neg a \cdot b$. As $\neg a \cdot b \leq b$, it follows that $\neg a \leq b$. So, $b = \neg a$. As we have that $a \vee b = 1$, it follows that $a \vee \neg a = 1$.

\Leftarrow) By hypothesis, we have (i) $a \vee \neg a = 1$. It is enough to see that $a \wedge \neg a = 0$. As $a \cdot \neg a = 0$, it is enough to prove that $a \wedge \neg a \leq a \cdot \neg a$. We have that $a \wedge \neg a \leq \neg a$. So, by monotonicity of \cdot , we have (ii) $a \cdot (a \wedge \neg a) \leq a \cdot \neg a$. We also have that $a \wedge \neg a \leq a$. So, again by monotonicity of \cdot , we have $\neg a \cdot (a \wedge \neg a) \leq \neg a \cdot a = a \cdot \neg a$. So, it follows (iii) $\neg a \cdot (a \wedge \neg a) \leq a \cdot \neg a$. Now, using Lemma 1(ii) with (i), (ii), and (iii), it follows that $a \wedge \neg a \leq a \cdot \neg a$. \square

Proposition 1 *Let $\mathbf{A} \in \mathbb{RL}$ and let a be a Boolean element of its universe A . Then, for all $b, c, d \in A$ the following properties hold:*

- (i) $a \wedge b = a \cdot b$,
- (ii) $a \cdot a = a$,
- (iii) $a \wedge \neg a = 0$,
- (iv) $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$,
- (v) $\neg \neg a = a$,
- (vi) $a \rightarrow b = \neg a \vee b$,
- (vii) $0 = \neg(a \vee \neg a)$,
- (viii) if $a \leq b \vee c$, $a \wedge b \leq d$, $a \wedge c \leq d$, then $a \leq d$,
- (ix) if $b \vee c = 1$, $a \wedge b \leq d$, $a \wedge c \leq d$, then $a \leq d$,
- (x) if $a \wedge b \leq c$, then $a \wedge \neg c \leq \neg b$,
- (xi) if $a \vee \neg b = 1$, then $b \leq a$.

Proof (i) Suppose that (i) $a \vee \neg a = 1$. It is enough to see that $a \wedge b \leq a \cdot b$. We have that $a \wedge b \leq b$. So, by monotonicity of \cdot , we have (ii) $a \cdot (a \wedge b) \leq a \cdot b$. We also have that $a \wedge b \leq a$. So, again by monotonicity of \cdot , we have $\neg a \cdot (a \wedge b) \leq \neg a \cdot a = 0$. So, it follows (iii) $\neg a \cdot (a \wedge b) \leq a \cdot b$. Now, using Lemma 1(ii) with (i), (ii), and (iii), it follows that $a \wedge b \leq a \cdot b$.

(ii) Using Part (i), we have $a \wedge a = a \cdot a$. Also, $a \wedge a = a$. So, $a \cdot a = a$.

(iii) Using Part(i), we have $a \wedge \neg a = a \cdot \neg a$. Also, $a \cdot \neg a = 0$. So, $a \wedge \neg a = 0$.

(iv) In a residuated lattice it holds that $a \cdot (b \vee c) = (a \cdot b) \vee (a \cdot c)$. Let a be Boolean. Then, using Part (i) three times, it follows that $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$.

(v) Suppose that (i) $a \vee \neg a = 1$. It is enough to see that $\neg \neg a \leq a$. We have that (ii) $a \cdot \neg \neg a \leq a$. Also, (iii) $\neg a \cdot \neg \neg a \leq a$, as $\neg a \cdot \neg \neg a = 0$. So, using Lemma 1(ii) with (i), (ii), and (iii), $\neg \neg a \leq a$.

(vi) It is enough to prove (i) $(\neg a \vee b) \cdot a \leq b$ and (ii) if $x \cdot a \leq b$, then $x \leq \neg a \vee b$. To see (i), note that $\neg a \cdot a \leq b$ and $b \cdot a \leq b$, whence $(\neg a \cdot a) \vee (b \cdot a) \leq b$. So, using distributivity of \cdot relative to \vee , we get (i). To see (ii), suppose $x \cdot a \leq b$. Then, $x \leq a \rightarrow b$. In order to get $x \leq \neg a \vee b$, it is enough to derive (iii) $a \rightarrow b \leq \neg a \vee b$ and use transitivity of \leq . To get (iii), let us use Lemma 1(ii). As a is Boolean, we have $a \vee \neg a = 1$. Now, $a \cdot (a \rightarrow b) \leq b \leq \neg a \vee b$. Also, $\neg a \cdot (a \rightarrow b) \leq \neg a \leq \neg a \vee b$. So, using Lemma 1(ii), we get (iii).

(vii) As we have $\neg c \leq a \vee \neg a$, for any $c \in A$, then, using Lemma 1(iii) and Part (iv), we get $\neg(a \vee \neg a) \leq \neg \neg c \leq c$.

(viii) As $a \leq b \vee c$, we have $a \leq a \wedge (b \vee c)$. Now, using Part (iv), it follows that $a \leq (a \wedge b) \vee (a \wedge c)$. Now, as $a \wedge b \leq d$ and $a \wedge c \leq d$, we have, by a basic property of \vee , that $(a \wedge b) \vee (a \wedge c) \leq d$. So, by transitivity of \leq , $a \leq d$.

(ix) As $b \vee c = 1$, reason as in Part (viii).

(x) Suppose $a \wedge b \leq c$. Then $a \cdot b \leq c$. By monotonicity of \cdot , it follows that $(a \cdot b) \cdot \neg c \leq c \cdot \neg c = 0$. Then, as \cdot is both associative and commutative, $(a \cdot \neg c) \cdot b \leq 0$. So, $a \cdot \neg c \leq \neg b$. Finally, using that a is Boolean, we get $a \wedge \neg c \leq \neg b$.

(xi) Let $a \vee \neg b = 1$. Then, $b \cdot (a \vee \neg b) = b \cdot 1 = b$. By distributivity of \cdot relative to \vee , it follows that $(b \cdot a) \vee (b \cdot \neg b) = b$. As $b \cdot \neg b = 0$, we have that $b \cdot a = b$. So, as a is Boolean, $b \wedge a = b$, i.e. $b \leq a$. \square

Lemma 3 *Let $\mathbf{A} \in \mathbb{RL}$ and let a and b be Boolean elements of A . Then,*

- (i) $a \wedge b = a \cdot b = \neg(\neg a \vee \neg b)$,
- (ii) $(a \cdot b) \vee (\neg a \cdot b) \vee (a \cdot \neg b) \vee (\neg a \cdot \neg b) = 1$.

Proof (i) Firstly, we have that $\neg a \leq \neg a \vee \neg b$. So, using Lemma 1(iii), it follows that $\neg(\neg a \vee \neg b) \leq \neg \neg a$. Now, using Proposition 1(v) and \leq -transitivity, we have that $\neg(\neg a \vee \neg b) \leq a$. Analogously, we get $\neg(\neg a \vee \neg b) \leq b$. Secondly, suppose $c \leq a$ and $c \leq b$, for $c \in A$. Then, using Lemma 1(iii) again, it follows that $\neg a \leq \neg c$ and $\neg b \leq \neg c$. So, $\neg a \vee \neg b \leq \neg c$. Then, using Lemma 1(iii) once again, $\neg \neg c \leq \neg(\neg a \vee \neg b)$. Now, using Proposition 1(v) and \leq -transitivity, we get $c \leq \neg(\neg a \vee \neg b)$.

(ii) Suppose that $a \vee \neg a = b \vee \neg b = 1$. Then, $1 = (a \vee \neg a) \cdot (b \vee \neg b) = (a \cdot b) \vee (\neg a \cdot b) \vee (a \cdot \neg b) \vee (\neg a \cdot \neg b)$. \square

Proposition 2 *Let $\mathbf{A} \in \mathbb{RL}$ and let a and b be Boolean elements of A . Then, (i) $\neg a$, (ii) $a \vee b$, (iii) $a \wedge b = a \cdot b$, (iv) $a \rightarrow b$, (v) 0 , and (vi) 1 are Boolean.*

Proof (i) Suppose $a \vee \neg a = 1$. Then, as $a \leq \neg \neg a$, we get $\neg \neg a \vee \neg a = 1$.

(ii) Use Lemma 3 (ii) and see that

$$\begin{aligned} a \cdot b &\leq a \leq a \vee b \leq (a \vee b) \vee \neg(a \vee b), \\ \neg a \cdot b &\leq b \leq a \vee b \leq (a \vee b) \vee \neg(a \vee b), \\ a \cdot \neg b &\leq a \leq a \vee b \leq (a \vee b) \vee \neg(a \vee b), \text{ and} \\ \neg a \cdot \neg b &\leq \neg a \wedge \neg b \leq \neg(a \vee b) \leq (a \vee b) \vee \neg(a \vee b). \end{aligned}$$

So, $a \vee b$ is Boolean.

(iii) Use Parts (i) and (ii), and Lemma 3(i).

(iv) Use Parts (i) and (ii), and Proposition 1(vi).

(v) Use Parts (i) and (ii), and Proposition 1(vii).

(vi) Use the definition of Boolean element and the fact that $\neg 1 = 0$. \square

From Proposition 2 it easily follows that, in any residuated lattice \mathbf{A} , the set of its Boolean elements $B(A) = \{a \in A : a \text{ is Boolean}\}$ is the domain of a subalgebra of \mathbf{A} , which is in fact a Boolean algebra. Indeed, $\mathbf{B}(\mathbf{A}) = (B(A); \wedge, \vee, \cdot, \rightarrow, 0, 1)$ is the greatest Boolean algebra contained in \mathbf{A} . $\mathbf{B}(\mathbf{A})$ is called the Boolean skeleton or the *center* of \mathbf{A} .

2.2 On the logic FL_{ew}

The logics we are interested in are extensions or expansions of the logic FL_{ew} described below.

Definition 2 *The language of FL_{ew} has four binary connectives, \wedge , \vee , \cdot , and \rightarrow , and two constants, 0 and 1.*

The axioms of FL_{ew} are:

- (1) $(\varphi \rightarrow \psi) \rightarrow ((\psi \rightarrow \gamma) \rightarrow (\varphi \rightarrow \gamma))$,
- (2) $(\gamma \rightarrow \varphi) \rightarrow ((\gamma \rightarrow \psi) \rightarrow (\gamma \rightarrow (\varphi \wedge \psi)))$,
- (3) $(\varphi \wedge \psi) \rightarrow \varphi$ and $(\varphi \wedge \psi) \rightarrow \psi$,
- (4) $\varphi \rightarrow (\varphi \vee \psi)$ and $\psi \rightarrow (\varphi \vee \psi)$,
- (5) $(\varphi \rightarrow \gamma) \rightarrow ((\psi \rightarrow \gamma) \rightarrow ((\varphi \vee \psi) \rightarrow \gamma))$,
- (6) $(\varphi \cdot \psi) \rightarrow (\psi \cdot \varphi)$,
- (7) $(\varphi \cdot \psi) \rightarrow \varphi$,
- (8) $(\varphi \rightarrow (\psi \rightarrow \gamma)) \rightarrow ((\varphi \cdot \psi) \rightarrow \gamma)$,
- (9) $((\varphi \cdot \psi) \rightarrow \gamma) \rightarrow (\varphi \rightarrow (\psi \rightarrow \gamma))$,
- (10) $0 \rightarrow \varphi$ and $\varphi \rightarrow 1$.

The only rule of FL_{ew} is modus ponens:

$$\frac{\varphi \quad \varphi \rightarrow \psi}{\psi}.$$

We define $\neg\varphi = \varphi \rightarrow 0$ and $\varphi \leftrightarrow \psi = (\varphi \rightarrow \psi) \wedge (\psi \rightarrow \varphi)$.

The following formulas and rules are derivable in FL_{ew} :

- (11) $\frac{\varphi \rightarrow \psi \quad \psi \rightarrow \gamma}{\varphi \rightarrow \gamma}$,
- (12) $\varphi \rightarrow (\psi \rightarrow \varphi)$,
- (13) $\varphi \rightarrow \varphi$,
- (14) $\varphi \rightarrow (\varphi \cdot 1)$ and $(\varphi \cdot 1) \rightarrow \varphi$,
- (15) $(\varphi \cdot (\psi \cdot \gamma)) \leftrightarrow ((\varphi \cdot \psi) \cdot \gamma)$
- (16) $(\varphi \cdot (\varphi \rightarrow \psi)) \rightarrow \psi$,
- (17) $(1 \rightarrow \varphi) \rightarrow \varphi$, $\varphi \rightarrow (1 \rightarrow \varphi)$,
- (18) $\frac{\varphi \quad \psi}{\varphi \wedge \psi}$,
- (19) $(\varphi \vee \neg\varphi) \rightarrow \neg(\varphi \wedge \neg\varphi)$.

Derivations for (11)-(19) are rather easy. Hence, they are left to the reader.

We will occasionally consider the following extensions of FL_{ew} .

Definition 3 *Consider the following axiomatic extensions of FL_{ew} :*

- *Intuitionistic logic IL is FL_{ew} plus the axiom*
(Contr) $\varphi \rightarrow (\varphi \cdot \varphi)$.
- *The logic MTL is FL_{ew} plus the axiom*
(Prel) $(\varphi \rightarrow \psi) \vee (\psi \rightarrow \varphi)$.
- *SMTL logic is MTL plus the axiom*
(PC) $\neg(\varphi \wedge \neg\varphi)$

- *WNM logic is MTL plus the axiom*
(WNM) $\neg(\varphi \& \psi) \vee (\varphi \wedge \psi \rightarrow \varphi \& \psi)$
- *NM logic is WNM plus the axiom*
(Inv) $\neg\neg\varphi \rightarrow \varphi$
- *BL logic is MTL plus the axiom*
(Div) $(\varphi \wedge \psi) \rightarrow (\varphi \& (\varphi \rightarrow \psi))$
- *Product logic is BL plus (PC) and the axiom*
(C) $\neg\varphi \vee ((\varphi \rightarrow \varphi \& \psi) \rightarrow \psi)$
- *Gödel logic G is BL plus (Contr)*
- *Lukasiewicz logic L is BL plus (Inv).*

Remark 1 *Note that Gödel logic G arises also as MTL plus (Contr) or as IL plus (Prel).*

In MTL and its extensions we may define $\phi \vee \psi := ((\phi \rightarrow \psi) \rightarrow \psi) \wedge ((\psi \rightarrow \phi) \rightarrow \phi)$. In BL and its extensions we may define $\varphi \wedge \psi := \varphi \& (\varphi \rightarrow \psi)$. Moreover, in IL the formula $(\phi \cdot \psi) \leftrightarrow (\phi \wedge \psi)$ is derivable, i.e. connectives \wedge and \cdot coincide in IL.

All these logics are algebraizable, and hence they are strongly complete with respect to their corresponding classes of algebras. Namely, FL_{ew} is complete with respect to the variety \mathbb{RL} of residuated lattices, MTL is complete with respect to the variety of pre-linear residuated lattices (MTL-algebras), and IL is complete with respect to the variety of contractive residuated lattices (Heyting algebras). Moreover, all axiomatic extensions of MTL are semilinear logics, that is, they are strongly complete with respect to the corresponding class of linearly ordered algebras. For instance, Gödel logic is complete with respect to the class of linearly ordered Heyting algebras, or Gödel chains.

3 Residuated lattices enriched with B

As explained in the previous section, the set of Boolean elements of a residuated lattice \mathbf{A} forms a Boolean algebra denoted *the center or Boolean skeleton of \mathbf{A}* . Cignoli and Monteiro considered Boolean elements in Łukasiewicz algebras in [6] and [7]. However, as far as we know, the operator defining the greatest Boolean element below, i.e. the operator B studied in this paper, has not yet been studied in the general context of residuated lattices. One relevant exception is the paper [18], where Reyes and Zolfaghari define modal operators \Box and \Diamond in the context of Bi-Heyting algebras that are shown to correspond respectively to the greatest and the smallest complemented element below and above, respectively. Thus, the \Box operation coincides with B . In the cited paper, using dual negation (or join-complement) D , always in the context of Bi-Heyting algebras, the authors also study a family of

modal operators \Box_n and \Diamond_n , in a similar way to the one we shall employ in Section 4.2.

We will be considering residuated lattices \mathbf{A} enriched with a unary operation B such that, for all $a \in A$, Ba is the greatest Boolean element below a , as defined in the Introduction. It is clear that B can be characterized by the following three conditions, for a, b in A :

- (BE1) $Ba \leq a$,
- (BE2) $Ba \vee \neg Ba = 1$,
- (BI) if $b \leq a$ and $b \vee \neg b = 1$, then $b \leq Ba$.

The class of residuated lattices with B will be denoted by \mathbb{RL}^B . Namely, an \mathbb{RL}^B -algebra is an algebra $\mathbf{A} = (A; \wedge, \vee, \cdot, \rightarrow, B, 0, 1)$ such that $(A; \wedge, \vee, \cdot, \rightarrow, 0, 1)$ is a residuated lattice and B satisfies the above three conditions.

First of all, note that B is new, that is, B is not expressible by a $\{\wedge, \vee, \cdot, \rightarrow, 0\}$ -term. Indeed, for instance, in the Gödel algebra $\mathbf{G}_2 \times \mathbf{G}_3$ (the direct product of the two-element Boolean algebra with universe $\{0, 1\}$ and the three-element Gödel algebra with universe $\{0, \frac{1}{2}, 1\}$) we have, for any $\{\wedge, \vee, \cdot, \rightarrow, 0\}$ -term t , that $ta \in \{0, a, 1\}$, where $a = (1, \frac{1}{2})$ is the join reducible coatom, while $Ba = (1, 0)$ is the join-irreducible atom, which does not belong to $\{0, a, 1\}$.

In the next proposition we see that all operations remain independent.

Proposition 3 *The set of operators $\{\wedge, \vee, \cdot, \rightarrow, B, 0\}$ is independent.*

Proof To see that \wedge is independent of the rest take the distributive lattice in Figure 1 and define the monoidal operation \cdot as $s \cdot t = 0$, $s \cdot s = s$, and $t \cdot t = t$, for coatoms s and t . This operation has a corresponding residuum \rightarrow . Since the only Boolean elements are 0 and 1, the operator B is defined as $B1 = 1$ and $Ba = 0$, for all $a \neq 1$. Then, note that the set $S = \{0, s, t, 1\}$ is closed for $\vee, \cdot, \rightarrow, 0$ and B , but $s \wedge t \notin S$.

To see that \vee is independent of the rest take the algebra that results from inverting the lattice order in the algebra of Example 1 and note that the set S with bottom, both atoms a_1 and a_2 , and top is closed for $\wedge, \cdot, \rightarrow, 0$, and B , but $a_1 \vee a_2 \notin S$.

To see that \cdot is independent of the rest take the four-element chain $0 < a < b < 1$ where $a \cdot b = a \cdot a = b \cdot b = a$, for the atom a and the coatom b , and note that the set $S = \{0, b, 1\}$ is closed for $\wedge, \vee, \rightarrow, 0$, and B , but $b \cdot b \notin S$.

To see that \rightarrow is independent of the rest take the algebra of Example 1 and note that the set $S = \{0, r, s, 1\}$ is closed for $\wedge, \vee, \neg, 0$, and B , but $s \rightarrow r \notin S$.

To see that 0 is independent of the rest take the two element Boolean algebra and note that the set $S = \{1\}$ is closed for $\wedge, \vee, \cdot, \rightarrow$, and B , but $0 \notin S$.

The independence of B has already been considered. \square

Lemma 4 *Let $\mathbf{A} \in \mathbb{RL}^B$ and let $a \in A$. Then,*

- (i) $Ba = a$ iff a is Boolean,
- (ii) $Ba = 1$ iff $a = 1$,
- (iii) $BBa = Ba$.

Proof (i) Suppose $Ba = a$. Using (BE2) it follows that $a \vee \neg a = 1$. For the other conditional, suppose $a \vee \neg a = 1$. Then, as $a \leq a$, using (BI) it follows that $a \leq Ba$. The other inequality follows by (BE2).

(ii) Suppose $Ba = 1$. Using (BE1), it follows that $1 \leq a$, i.e. $a = 1$. For the other conditional, suppose $a = 1$. Then, $a \leq 1$. Using (BI) and the fact that 1 is Boolean (see Proposition 2(vi)), it follows that $1 \leq Ba$.

(iii) Considering (BE1), it is enough to see that $Ba \leq BBa$, which follows using (BI) and (BE2). \square

We also have the following properties.

Lemma 5 *Let $\mathbf{A} \in \mathbb{RL}^B$ and let $a, b \in A$. Then,*

- (i) B -Monotonicity: if $a \leq b$, then $Ba \leq Bb$,
- (ii) $B(a \wedge b) = Ba \wedge Bb$,
- (iii) $B(a \wedge b) \leq a \cdot b$,
- (iv) $B(a \cdot b) = B(a \wedge b)$,
- (v) $B(a \cdot b) = Ba \cdot Bb$,
- (vi) $Ba \vee Bb \leq B(a \vee b)$,
- (vii) $B(a \rightarrow b) \leq Ba \rightarrow Bb$,
- (viii) $B0 = 0$,
- (ix) $B\neg a \leq \neg Ba$.

Proof (i) Suppose $a \leq b$. Using (BI), it is enough to have $Ba \leq b$ and $Ba \vee \neg Ba = 1$. Now, the former follows by (BE1) and the hypothesis, and the latter is (BE2).

(ii) $B(a \wedge b) \leq Ba \wedge Bb$ follows from $a \wedge b \leq a, b$ using B -monotonicity. The other inequality follows using (BI), (BE1), and (iii) in Proposition 2.

(iii) By (i) in Proposition 1 and part (ii) we have $B(a \wedge b) = Ba \wedge Bb = Ba \cdot Bb$. The goal follows using $Ba \leq a$, $Bb \leq b$, and monotonicity of \cdot .

(iv) From $a \cdot b \leq a \wedge b$ by (i), we get $B(a \cdot b) \leq B(a \wedge b)$. For the other inequality, using (BI), it is enough to have $B(a \wedge b) \leq a \cdot b$ and $B(a \wedge b)$ Boolean. Now, the former is (iii) and the latter follows from (BE2).

(v) As Ba is Boolean, by (i) of Proposition 1, we have $Ba \wedge Bb = Ba \cdot Bb$. Moreover, by (ii), $B(a \wedge b) = Ba \wedge Bb$. We get our goal using (iv).

(vi) It follows using (i) (B -Monotonicity).

(vii) As $a \rightarrow b \leq a \rightarrow b$, we have $(a \rightarrow b) \cdot a \leq b$. Then, by B -monotonicity, $B((a \rightarrow b) \cdot a) \leq Bb$. So, using (v), $B(a \rightarrow b) \cdot Ba \leq Bb$. So, $B(a \rightarrow b) \leq Ba \rightarrow Bb$.

(viii) It follows because 0 is Boolean.

(ix) It follows from (vii), (viii), and $\neg a = a \rightarrow 0$. \square

Regarding the inequalities in the previous lemma, that is, (iii), (vi), (vii), and (ix), their reciprocals do not hold. Indeed, inequality $a \cdot b \leq B(a \wedge b)$ fails in the three-element Gödel algebra \mathbf{G}_3 taking the top and the middle element. Inequality $B(a \vee b) \leq Ba \vee Bb$ fails in the algebra of Example 1 taking a and b to be the coatoms s and t . Also, inequality $\neg Ba \leq B\neg a$ fails in \mathbf{G}_3 , taking a to be the middle element. So, also inequality $Ba \rightarrow Bb \leq B(a \rightarrow b)$ fails.

Though B may not exist for every element in a residuated lattice, B exists in every finite residuated lattice.

Proposition 4 *Let $\mathbf{A} \in \mathbf{RL}$ be finite. Then, B exists in A .*

Proof In a finite residuated lattice \mathbf{A} , for any $a \in A$, we have $Ba = \bigvee \{b \in A : b \leq a \text{ and } b \vee \neg b = 1\}$. It is enough to see that if $b_1 \leq a$, $b_1 \vee \neg b_1 = 1$, $b_2 \leq a$, and $b_2 \vee \neg b_2 = 1$, then (i) $b_1 \vee b_2 \leq a$ and (ii) $(b_1 \vee b_2) \vee \neg(b_1 \vee b_2) = 1$. Now, (i) follows immediately and (ii) follows using (ii) in Proposition 2. \square

On the other hand, there are infinite residuated lattices where B does not exist. Indeed, we have the following example due to Franco Montagna (see [1]).

Proposition 5 *There is an (infinite) Gödel algebra \mathbf{A} and $a \in A$ such that Ba does not exist, i.e. where B does not exist.*

Proof Let $[0, \frac{1}{2}, 1]_G$ be the three-element Gödel algebra. Let us consider

$$\begin{aligned} A_1 &= \{a \in ([0, \frac{1}{2}, 1]_G)^{\mathbb{N}} \text{ such that } \{i \in \mathbb{N} : a_i = 0\} \text{ is finite}\}, \\ A_2 &= \{a \in ([0, \frac{1}{2}, 1]_G)^{\mathbb{N}} \text{ such that } \{i \in \mathbb{N} : a_i \neq 0\} \text{ is finite}\}, \text{ and} \\ A &= A_1 \cup A_2. \end{aligned}$$

The set A is the domain of a subalgebra of $([0, \frac{1}{2}, 1]_G)^{\mathbb{N}}$. Indeed, if $a, b \in A_1$, then $a \wedge b \in A_1$ and $a \rightarrow b \in A_1$, if $a, b \in A_2$, then $a \wedge b \in A_2$ and $a \rightarrow b \in A_1$, if $a \in A_1$ and $b \in A_2$, then $a \wedge b \in A_2$ and $a \rightarrow b \in A_2$, and if $a \in A_2$ and $b \in A_1$, then $a \wedge b \in A_2$ and $a \rightarrow b \in A_1$. Also, $0 \in A_2$. So, A is the domain of a subalgebra \mathbf{A} of $([0, \frac{1}{2}, 1]_G)^{\mathbb{N}}$.

Now, take a to be such that $a_i = 1$ if i is even and $a_i = \frac{1}{2}$ if i is odd. Next, consider the set $\{b \in A : b \leq a \text{ and } b \text{ is Boolean}\}$. It consists of all elements b such that $b_i = 0$ for all odd i and for all but finitely many even i , and $b_i = 1$ otherwise. It can be seen that this set has no maximum in A . \square

Actually, Montagna's example of Proposition 5 can be generalized as follows.

Proposition 6 *Let \mathbb{V} be a variety of MTL-algebras such that there is a linearly ordered algebra $\mathbf{A} \in \mathbb{V}$ with a proper filter F (i.e. $\{1\} \subsetneq F \subsetneq A$), that is, such that \mathbf{A} is not simple. Then, \mathbb{V} contains an infinite algebra where B does not exist.*

Proof Let $\mathbf{D} \in \mathbb{V}$ be a chain and F be a filter of \mathbf{A} satisfying the hypothesis of the proposition. Let us define $F^\neg = \{x \in D \mid \exists y \in F, x \leq \neg y\}$ and let $C = F \cup F^\neg$. It is easy to check that C is the domain of a subalgebra of \mathbf{D} . Finally define the following sets:

$$\begin{aligned} A_1 &= \{a \in C^{\mathbb{N}} \text{ such that } \{i \in \mathbb{N} : a_i \in F\} \text{ is finite}\}, \\ A_2 &= \{a \in C^{\mathbb{N}} \text{ such that } \{i \in \mathbb{N} : a_i \in F^\neg\} \text{ is finite}\}, \\ A &= A_1 \cup A_2. \end{aligned}$$

One can check that again A is the domain of a subalgebra of $\mathbf{C}^{\mathbb{N}}$, taking into account that if $x \in F$ and $y \in F^\neg$, then $x \wedge y, x * y, x \rightarrow y \in F^\neg$, and if $x, y \in F^\neg$, then $x \rightarrow y \in F$.

Thus, A is a subalgebra and taking an element a such that $a_i = 1$ if i is even and $a_i = b$, for a given $b \in F \setminus \{1\}$, then the same argument as in Montagna's example proves that Ba does not exist. \square

For readers familiar with the main systems of mathematical fuzzy logic and their algebraic semantics (see [8]), we provide the following corollary with further examples of subvarieties of residuated lattices containing algebras where B does not exist.

Corollary 1 *In the following varieties of MTL-algebras, there is an infinite algebra where B does not exist:*

- the variety generated by any continuous t -norm,
- the varieties generated by either the NM t -norm or a WNM t -norm.¹

Proof In all these varieties there is an algebra \mathbf{A} satisfying the conditions of Proposition 6.

If the t -norm is either a Gödel, Product, or a WNM t -norm (including NM), then take as \mathbf{A} the standard chain and as F the positive elements respect to \neg , i.e., the elements such that $\neg x \leq x$. If the t -norm is Lukasiewicz, then take \mathbf{A} as the Chang algebra and F as the set of its positive elements. Finally, if the continuous t -norm is a proper ordinal sum, then take \mathbf{A} as the standard chain and $F = [a, 1]$, where $a \in (0, 1)$ is the end point of a component. It is clear that in all cases F is a proper filter and thus Proposition 6 applies. \square

¹ Actually, this could be generalized in the following sense. In [14] Noguera proves that the variety generated by simple n -contractive MTL-chains is the variety of S_n -MTL algebras, i.e. MTL-algebras satisfying the law $x \vee \neg x^{n-1} = 1$. Therefore, any variety of n -contractive MTL-algebras that are not S_n -MTL has a chain with a proper filter. In particular, this is the case for the varieties of WNM and NM-algebras, since they are 3-contractives and are not S_3 -MTL.

Some papers (e.g. [4]) consider the notion of *compatible operation*. Operation B is not compatible, that is, the congruences of \mathbb{RL} and \mathbb{RL}^B are not the same. To see this, take the three-element Heyting or Gödel algebra \mathbf{G}_3 with universe $\{0, \frac{1}{2}, 1\}$ and the equivalence relation given by $\theta = \{(0, 0), (\frac{1}{2}, \frac{1}{2}), (\frac{1}{2}, 1), (1, \frac{1}{2}), (1, 1)\}$. It holds that θ is a \mathbb{RL} -congruence, but not an \mathbb{RL}^B -congruence, as $B\frac{1}{2} = 0$ and $B1 = 1$.

Taking *modality* to mean a finite combination of the unary operators \neg and B , the next statement shows how many different modalities there are in \mathbb{RL}^B .

Proposition 7 *In \mathbb{RL}^B there are nine different modalities. They may be ordered as follows: on the one hand, the positive modalities $B \leq Id \leq \neg\neg \leq \neg B\neg$ (with also $B \leq B\neg\neg \leq \neg\neg$) and, on the other hand, the negative modalities $B\neg \leq \neg \leq \neg B\neg\neg \leq \neg B$. See Figure 2.*

Proof The inequalities are immediate. The reverse inequalities can be seen not to be the case by considering either the only atom in the three-element Gödel algebra \mathbf{G}_3 or any of the two non-comparable elements of the Heyting algebra obtained by adding a top element to the Boolean algebra of 4 elements. There are no other modalities, because if we apply operations \neg and B to the given nine modalities, we do not get anything new, as $\neg\neg B = B$, $BB = B$, and $B\neg B = \neg B$. \square

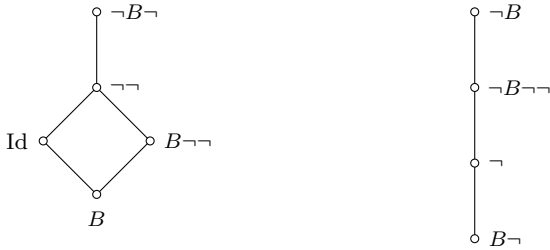


Fig. 2 The positive and negative modalities of \neg and B

3.1 An equational class

It is natural to inquire whether the class \mathbb{RL}^B is in fact an equational class. To this end, we start focusing our attention on the following equations, using $x \preccurlyeq y$ as an abbreviation for $x \vee y \approx y$:

$$(BI1) \quad Bx \preccurlyeq B(x \vee y),$$

$$(BI2) \quad B1 \approx 1.$$

Lemma 6 *Equations (BI1) and (BI2) hold in \mathbb{RL}^B .*

Proof The given equations follow immediately from lemmas 4(ii) and 5(vi), respectively. \square

We are also interested in the equation

$$(BI3) \quad B(x \vee \neg x) \preccurlyeq Bx \vee \neg x,$$

but it is not easy to see that it holds in \mathbb{RL}^B . Towards this goal, we state and prove the following result.

Lemma 7 *In \mathbb{RL}^B the following hold:*

- (i) $(B(x \vee \neg x) \wedge x) \vee \neg(B(x \vee \neg x) \wedge x) \approx 1$,
- (ii) $B(x \vee \neg x) \wedge x \preccurlyeq Bx$,
- (iii) $B(x \vee \neg x) \wedge \neg Bx \preccurlyeq \neg x$.

Proof (i) Using Lemma 1(i), and using T for the left hand side of the given equation, it is enough to get

- (iv) $B(x \vee \neg x) \vee \neg B(x \vee \neg x) \approx 1$,
- (v) $B(x \vee \neg x) \preccurlyeq T$, and
- (vi) $\neg B(x \vee \neg x) \preccurlyeq T$.

Part (iv) is immediate because of (BE2).

To see (v), using Proposition 1(viii), note that we have that $B(x \vee \neg x) \preccurlyeq x \vee \neg x$, (immediate using (BE1)), $B(x \vee \neg x) \wedge x \preccurlyeq T$ (also immediate), and $B(x \vee \neg x) \wedge \neg x \preccurlyeq T$, which follows from $B(x \vee \neg x) \wedge \neg x \preccurlyeq \neg(B(x \vee \neg x) \wedge x)$, which holds because of Lemma 1(iv).

To see (vi), note that $B(x \vee \neg x) \wedge x \leq B(x \vee \neg x)$. So, using Lemma 1(iii), it follows that $\neg B(x \vee \neg x) \preccurlyeq \neg(B(x \vee \neg x) \wedge x)$. And so, $\neg B(x \vee \neg x) \preccurlyeq T$.

(ii) Use (BI), Part (i), and $B(x \vee \neg x) \wedge x \preccurlyeq x$.

(iii) Use Part (ii) and Proposition 1(x). \square

Proposition 8 *The equation (BI3) holds in \mathbb{RL}^B .*

Proof Using Proposition 1(ix), it is enough to check the following three conditions:

- (i) $Bx \vee \neg Bx \approx 1$,
- (ii) $B(x \vee \neg x) \wedge Bx \preccurlyeq Bx \vee \neg x$, and
- (iii) $B(x \vee \neg x) \wedge \neg Bx \preccurlyeq Bx \vee \neg x$.

Now, (i) is immediate due to (BE2) and (ii) is also immediate as $B(x \vee \neg x) \wedge Bx \preccurlyeq Bx$. Regarding (iii), it follows from Lemma 7(iii). \square

Remark 2 *Arguing as in the proof of Proposition 8 and noting that $B(x \vee \neg x) \wedge \neg Bx \preccurlyeq B\neg x$ follows using (BI) from Lemma 7(iii) and the fact that the term $B(x \vee \neg x) \wedge \neg Bx$ is Boolean, it may be seen that also the inequality $B(x \vee \neg x) \preccurlyeq Bx \vee B\neg x$ holds in \mathbb{RL}^B .*

Lemma 8 *B is monotone just using equations.*

Proof Suppose $x \vee y \approx y$. Then, (i) $B(x \vee y) \approx By$. Now, using (BI1), we have (ii) $Bx \vee B(x \vee y) \approx B(x \vee y)$. From (i) and (ii) we get $Bx \vee By \approx By$. \square

The following theorem answers positively the question posed above.

Theorem 1 \mathbb{RL}^B is an equational class. An equational basis relative to \mathbb{RL} is the following set of equations:

- (BE1) $Bx \preceq x$,
- (BE2) $Bx \vee \neg Bx \approx 1$,
- (BI1) $Bx \preceq B(x \vee y)$,
- (BI2) $B1 \approx 1$,
- (BI3) $B(x \vee \neg x) \preceq Bx \vee \neg x$.

Proof It is enough to prove (BI) using the given equations. Suppose (i) $b \leq a$ and (ii) $b \vee \neg b = 1$. From (i), using Lemma 8, it follows (iii) $Bb \leq Ba$. From (ii), using (BI2), it follows $B(b \vee \neg b) = B1 = 1$, which, using (BI3), implies that $Bb \vee \neg b = 1$, which, using (iii) gives $Ba \vee \neg b = 1$, which, using Proposition 1(xi), implies $b \leq Ba$. \square

Remark 3 Note that, as expected, the just given proof of (BI) only uses equations (BI1), (BI2), and (BI3).

It is also natural to inquire whether the given equations are independent.

Proposition 9 The set $\{(\text{BE1}), (\text{BE2}), (\text{BI1}), (\text{BI2}), (\text{BI3})\}$ is independent.

Proof To see that (BE1) is independent of the rest, take the three-element Gödel algebra \mathbf{G}_3 and define $B0 = 0$, and $Ba = 1$, if a is not 0.

To see that (BE2) is independent of the rest, take the four-element Gödel chain \mathbf{G}_4 and define $B0 = 0$, $B1 = 1$, $Ba = 0$, for the only atom a of the chain, and $Bc = c$, for the remaining element c .

To see that (BI1) is independent of the rest, take the Gödel algebra $\mathbf{G}_3 \times \mathbf{G}_3$. If a is any of the four Boolean elements, then put $Ba = a$, else put $Ba = 0$.

To see that (BI2) is independent of the rest, take again \mathbf{G}_3 , but now define $Ba = 0$, for every a .

Finally, to see that (BI3) is independent of the rest, take the four-element Boolean $\mathbf{G}_2 \times \mathbf{G}_2$ algebra and define, for any a , if $a = 1$, then $Ba = 1$, else $Ba = 0$. \square

3.2 Subdirectly irreducible \mathbb{RL}^B -algebras

In this section we show the subdirectly irreducible members of \mathbb{RL}^B are those whose Boolean elements are only the top and bottom elements.

Definition 4 Let $\mathbf{A} \in \mathbb{RL}^B$. A set F contained in A is said to be a \mathbb{RL}^B -filter iff for all $a, b \in A$ it satisfies

- (1) $1 \in F$,
- (2) if $a \in F$ and $a \leq b$, then $b \in F$,
- (3) if $a, b \in F$, then $a \cdot b \in F$,
- (4) if $a \in F$, then $Ba \in F$.

Proposition 10 Let $\mathbf{A} \in \mathbb{RL}^B$. The lattice of \mathbb{RL}^B -congruences is isomorphic to the set of \mathbb{RL}^B -filters. Indeed, let $f : \text{Con}(\mathbf{A}) \rightarrow \text{Fil}(\mathbf{A})$ be defined by: if \equiv is a \mathbb{RL}^B -congruence, then $f(\equiv)$ is the \mathbb{RL}^B -filter $F_\equiv = \{a \in A : a \equiv 1\}$. Then, the function f is an isomorphism such that if F is a \mathbb{RL}^B -filter, then $f^{-1}(F)$ is the \mathbb{RL}^B -congruence \equiv_F defined by $a \equiv_F b$ iff $a \rightarrow b, b \rightarrow a \in F$.

Proof It is obvious that F_\equiv is a \mathbb{RL}^B -filter. In order to prove that \equiv_F is a congruence we need to prove that if $a \equiv_F b$, then $Ba \equiv_F Bb$, since the other conditions are known to be true for any residuated lattice. So, suppose $a \rightarrow b$ and $b \rightarrow a \in F$. Then, by the fourth condition in the definition of filter, $B(a \rightarrow b) \in F$. Now, using (vi) in Lemma 5 and the second condition in the definition of filter, it follows that $Ba \rightarrow Bb \in F$. Analogously, we obtain that $Bb \rightarrow Ba \in F$. Finally, it is also obvious that $f^{-1} \circ f = \text{Id}$. \square

Now we can characterize a family of \mathbb{RL}^B -filters.

Proposition 11 Let $\mathbf{A} \in \mathbb{RL}^B$. If $a \in B(A)$, then $F_a = [a, 1] = \{x \in A : a \leq x \leq 1\}$ is a \mathbb{RL}^B -filter.

Proof It is obvious that F_a satisfies the first two conditions of a \mathbb{RL}^B -filter. The third is an easy consequence of the fact that if $a \in B(A)$, then $a * x = a \wedge x$ and thus if $x, y \in F_a$, then $a = a \wedge y \leq x * y$ and thus $x * y \in F_a$. Finally, if $x \in F$, then $a = Ba \leq Bx$. \square

From now on, F_a denotes the principal filter defined by $a \in B(A)$.

In order to characterize the subdirectly irreducible \mathbb{RL}^B -algebras, we will use the result of [22, Theorem 97]: an algebra \mathbf{A} is subdirectly reducible iff there exists a family of non-trivial congruences σ_i such that their intersection is the identity. In our case, this means that \mathbf{A} is subdirectly irreducible iff there is a unique coatom in the lattice of \mathbb{RL}^B -congruences of \mathbf{A} .

Proposition 12 Let $\mathbf{A} \in \mathbb{RL}^B$. Then, \mathbf{A} is subdirectly irreducible iff $B(A) = \{0, 1\}$.

Proof Observe first that if F is a \mathbb{RL}^B -filter of A , then F contains a Boolean element a (by the third condition of \mathbb{RL}^B -filter) and, thus, F contains F_a . So, to obtain the intersection of the non-trivial \mathbb{RL}^B -filters of \mathbf{A} it is enough to compute the intersection of the filters F_a . However, this intersection is not the identity iff there exists a unique Boolean element a such that a is a coatom of $B(A)$. So, being $B(A)$ a Boolean algebra, this implies that $B(A) = \{0, 1\}$. \square

4 Comparing B with other operations

Operation B is strongly related to other operations considered in the literature, e.g. the Monteiro-Baaz Δ and an operation defined with the join-complement D . In this section we study these relationships.

4.1 Comparing B with Δ

The operation Δ was already considered by Monteiro in his paper on symmetric Heyting algebras in 1980 (see [13]). Monteiro considered the same definitions of possibility and necessity operations given by Moisil in [11] (see p. 67 in [13]). However, instead of using Moisil's notation, Monteiro used ∇ and Δ , respectively. When so doing, he did not explicitly mention Moisil. However, many works by Moisil appear in the list of references of [13], including [11] and [12]. Monteiro also considered the Δ operator in the setting of linear symmetric Heyting algebras and studied the properties of Δ in the totally linear case (see [13, Ch. 5, Sect. 3]). In 1996, independently, Baaz in [2] considered an expansion of Gödel logic with a connective he also called Δ satisfying certain axioms and the rule $\varphi/\Delta\varphi$. Although he did not cite Monteiro, the proposed axioms are equivalent to the properties that Monteiro proved for his Δ operator in the framework of totally linear symmetric Heyting algebras. Baaz also provided a deduction theorem using Δ : $\Gamma, \varphi \vdash \psi$ iff $\Gamma \vdash \Delta\varphi \rightarrow \psi$. In 1998, Hájek considered Baaz's Δ in the context of BL-algebras and BL logic (see pp. 57-61 [10]). He gave for Δ exactly the same axioms as Baaz presented in [2] for Gödel logic. He observed that all Δ axioms make it behave like a necessity operator, with the exception of the axiom $\Delta(\varphi \vee \psi) \rightarrow (\Delta\varphi \vee \Delta\psi)$, that is characteristic of possibility operations (see Remark 2.4.7 of [10]). The Δ operation has also been studied in the more general context of Mathematical fuzzy logic, see several chapters in the handbook [8]. More recently, in [1] the authors study, among other things, the expansion of FL_{ew} with the Δ operator and show that it is conservative.

In MTL, Δ can always be defined over chains, namely as $\Delta 1 = 1$ and $\Delta x = 0$ for all $x \neq 1$, and thus, Δ and B over MTL-chains coincide. But there are (non-linearly) MTL-algebras where Δ does not exist. Nevertheless, this is not a problem because MTL is semilinear, and the semantics of Δ over chains is clear. However, there is not a clear semantical interpretation of the axioms of Δ in the general context of residuated lattices.

In the context of a residuated lattice, the operator Δ is introduced e.g. in [1] by the same equations as in MTL or BL (cf. [10, p. 58]):

- ($\Delta E1$) $\Delta x \preceq x$,
- ($\Delta E2$) $\Delta x \vee \neg \Delta x \approx 1$,
- ($\Delta I1$) $\Delta(x \vee y) \preceq \Delta x \vee \Delta y$,
- ($\Delta I2$) $\Delta 1 \approx 1$,
- ($\Delta I3$) $\Delta x \preceq \Delta \Delta x$,
- ($\Delta I4$) $\Delta(x \rightarrow y) \preceq \Delta x \rightarrow \Delta y$,

where, again, $x \preceq y$ abbreviates $x \vee y \approx y$. Note that ($\Delta I3$) may be derived from the rest: it is enough to check that an operator satisfying the rest of the equations, satisfies all the equations in Theorem 1, and hence the quasi-equation (**BI**) as well; then use (iii) of Lemma 4. Also, regarding their defining equations, the only difference between Δ and B is that Δ satisfies $\Delta(x \vee y) \preceq \Delta x \vee \Delta y$, whereas B only satisfies the particular case $y = \neg x$, that is, B only satisfies $B(x \vee \neg x) \preceq Bx \vee B\neg x$, as stated in Remark 2.

We will denote by \mathbb{RL}^Δ the class of residuated lattices expanded with Δ .

It will be useful to bear in mind the following fact.

Lemma 9 *Let $\mathbf{A} \in \mathbb{RL}^\Delta$ and $a \in A$. Then, $\Delta a = a$ iff a is Boolean.*

Proof Supposing $\Delta a = a$, using ($\Delta E2$) it follows that a is Boolean. On the other hand, suppose a is Boolean. Considering ($\Delta E1$), it is enough to prove that $a \leq \Delta a$. By Lemma 1(ii), it is enough in turn to prove $\Delta a \vee \Delta \neg a = 1$, $a \cdot \Delta a \leq \Delta a$, and $a \cdot \Delta \neg a \leq \Delta a$. The first condition holds using ($\Delta I1$) and ($\Delta I2$), since a is Boolean. The second condition is immediate. For the third, observe that $a \cdot \Delta \neg a \leq a \cdot \neg a = 0 \leq \Delta a$.

Actually, Δ is somewhat stronger than B in the following sense.

Proposition 13 *Let $\mathbf{A} \in \mathbb{RL}$. If Δ exists in \mathbf{A} , then so does B , with $B = \Delta$.*

Proof Considering Theorem 1, all we have to see is that Δ satisfies the equational basis given for B . This is immediate excepting (**BI1**). Let us see that the equation $\Delta x \preceq \Delta(x \vee y)$ also holds. As we have $x \rightarrow (x \vee y) \approx 1$, using ($\Delta I2$) and ($\Delta I4$) we get $1 \preceq \Delta x \rightarrow \Delta(x \vee y)$, which gives $\Delta x \preceq \Delta(x \vee y)$. \square

On the other hand, we have the following result.

Proposition 14 *There exist finite residuated lattices where B exists, but Δ does not.*

Proof Using Proposition 4, it follows that B exists in the Gödel algebra of Example 1, as the algebra is finite. Now, take its coatoms s and t . To see that Δ does not exist, note that ($\Delta E1$) and ($\Delta E2$) imply that $\Delta s = \Delta t = 0$. So, $\Delta s \vee \Delta t = 0$. However, $\Delta(s \vee t) = 1$, due to ($\Delta I2$). Then, ($\Delta I1$) is not satisfied. \square

Example 1 makes clear the basic difference between Δ and B , when we define them over MTL-algebras. It is well known that MTL^Δ , the expansion of MTL with Δ , is semilinear, i.e. each algebra of the variety is a subdirect product of linearly ordered MTL^Δ -algebras. Moreover, we have seen that Δ and B coincide over chains. Thus, Example 1 proves that MTL^B , the expansion of MTL with B , is not semilinear. In fact, this was already clear from Proposition 12, since there exist subdirectly irreducible MTL^B -algebras (like the one defined in Example 1) that are not linearly ordered.

4.2 Comparing B with an operation using the join-complement

The join-complement operation D has a long history. In 1919, Skolem considered lattices expanded with both meet and join relative complements (see §2 of [20] or pp. 77-85 of [21]). He just worked from an algebraic point of view. He noted that existence of both top 1 and bottom 0 is implied. Also, he briefly considered the meet and join-complements, for which, for an arbitrary argument a , he used the notations $\frac{0}{a}$ and $1 - a$, respectively.

In 1942, Moisil defined possibility as $\neg\neg$ and necessity as DD in a logical setting where he had both intuitionistic negation \neg and its dual D (see §4 of [11] or p. 365 in [12]). He did not mention Skolem. In 1949, Ribenboim proved that distributive lattices with D form an equational class (see [19]). In fact, the meet is not needed, as the class with join and join-complement D is already an equational class. In 1974, Rauszer, mainly considering algebraic aspects, studied a logic with conjunction, disjunction, conditional, and its dual (see [16]). She also included both intuitionistic negation \neg and its dual D , though these can be easily defined. Her axiomatization included the expected axioms plus the rules *modus ponens* and $\varphi/\neg D\varphi$. She also provided a deduction theorem using $(\neg D)^n$. She neither mentioned Skolem nor Moisil.

In the context of a join semi-lattice \mathbf{A} , it is possible to postulate the existence of the join-complement $Da = \min\{b \in A : \text{for all } c \in A, c \leq a \vee b\}$, for $a \in A$. This is equivalent to the following two conditions:

- (DI) $b \leq a \vee Da$, for all $a, b \in A$,
- (DE) for any $a, b \in A$, if for all $c \in A$, $c \leq a \vee b$, then $Da \leq b$.

In a join semi-lattice the existence of D implies the existence of both top $1 = a \vee Da$, for any a , and bottom $0 = D(a \vee Da)$, for any a . Moreover, D can be equationally characterized by the following three equations, where, again, we use $x \preceq y$ as an abbreviation for $x \vee y \approx y$:

- (DI) $y \preceq x \vee Dx$,
- (DE1) $D(x \vee Dx) \preceq y$,
- (DE2) $Dy \preceq x \vee D(x \vee y)$.

In what follows, \mathbb{RL}^D will denote the class of residuated lattices expanded with an operation D satisfying these equations. Obviously, by definition, \mathbb{RL}^D is an equational class. Notice that in a residuated lattice \mathbf{A} , having in the signature the symbols 0 and 1 for the bottom and top elements, respectively, the above definition of D can be simplified to $Da = \min\{b \in A : a \vee b = 1\}$, and the condition (DE) simplifies to be

- (DE') for any $a, b \in A$, if $a \vee b = 1$, then $Da \leq b$.

Moreover the equations (DI) and (DE1) can also be simplified to:

- (DI') $x \vee Dx \approx 1$,
- (DE1') $D1 \approx 0$.

Remark 4 Note that, while $x \preceq \neg\neg x$ holds in \mathbb{RL} , from (DE') and (DI') it follows that $DDx \preceq x$ holds in \mathbb{RL}^D . Note also that in a Heyting algebra D is the dual of \neg , since in that case \neg coincides with the meet complement.

As in the case of B , D may not exist in some residuated lattices, but it always exists in the finite ones.

Proposition 15 Let \mathbf{A} be a finite residuated lattice. Then D exists in A .

Proof It is enough to prove that $\bigwedge\{b \in A : a \vee b = 1\}$ exists in A . For that, it is enough to see that if $a \vee b_1 = 1$ and $a \vee b_2 = 1$, then $a \vee (b_1 \wedge b_2) = 1$. Now, from the antecedent it follows that $(a \vee b_1) \cdot (a \vee b_2) = 1$ and using twice the distributive law of \cdot with respect to \vee , we have that $(a \cdot a) \vee (a \cdot b_2) \vee (b_1 \cdot a) \vee (b_1 \cdot b_2) = 1$. Any subterm t of the left-hand term is such that $t \leq a \vee (b_1 \wedge b_2)$. \square

Following [16] and [18], we consider now the compound operation $\neg D$ and its relation to B . First, let us state the following fact.

Lemma 10 Let $\mathbf{A} \in \mathbb{RL}^D$ and $a, b \in A$. Then,

- (i) $\neg Da \leq a$,
- (ii) if $a \leq b$, then $Db \leq Da$ and $\neg Da \leq \neg Db$.

Proof (i) follows from $a \vee Da = 1$ using Lemma 1(vi). (ii) Assume $a \leq b$. Then, $1 = a \vee Da \leq b \vee Da$. Hence, by (DE) we have $Db \leq Da$. Now, apply Lemma 1(iii) to get $\neg Da \leq \neg Db$. \square

In [5, Section 5] the authors prove a result about iterations of the operation $\neg D$ in the context of meet-complemented distributive lattices with D . Once trivially adapted to \mathbb{RL}^D , it is the following fact.

Proposition 16 (i) For any natural $n > 0$, let \mathbf{RL}^{Dn} be the subvariety of \mathbf{RL}^D defined by adding to those of \mathbf{RL}^D the following equation:

$$(\neg D)^{n+1}x \approx (\neg D)^n x.$$

Then, the sequence of varieties $\mathbf{RL}^{D1} \subset \mathbf{RL}^{D2} \subset \dots \subset \mathbf{RL}^{Dn} \subset \dots$ is strictly increasing.

(ii) There are algebras of \mathbf{RL}^D where none of the equations given in (i) hold.

Next, consider the following example of a Heyting algebra where B exists but D does not.

Example 2 B exists in the Heyting algebra $1 + (\mathbb{N} \times \mathbb{N})^\partial$ of Figure 3, where $(\mathbb{N} \times \mathbb{N})^\partial$ is obtained ‘turning upside down’ the partial order $\mathbb{N} \times \mathbb{N}$. In that Heyting algebra, $Ba = 1$ if $a = 1$ else $Ba = 0$. However, D does not exist for the elements $(0, n)$ and $(n, 0)$, with non-zero n .

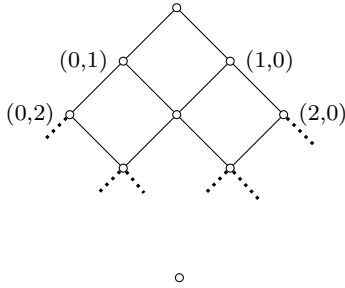


Fig. 3 The residuated lattice $1 + (\mathbb{N} \times \mathbb{N})^\partial$, where B exists but D does not

There are also \mathbf{RL} -algebras where D exists, but B does not (see the end of Section 2 of [5] for an example of a Heyting algebra). Note that in Franco Montagna’s example, neither B nor D exist.

In the following case, existence of D implies existence of B .

Proposition 17 Let $\mathbf{A} \in \mathbf{RL}$ and $\neg a \vee \neg \neg a = 1$, for all $a \in A$. Then, if D exists in \mathbf{A} , then B also exists in \mathbf{A} , with $B = \neg D$.

Proof Let $a \in A$. Then, $\neg Da$ exists in A . We have to see (i) $\neg Da \leq a$, (ii) $\neg Da \vee \neg \neg Da = 1$, and (iii) if $b \leq a$ and $b \vee \neg b = 1$, then $b \leq \neg Da$. Now, (i) holds as seen in Lemma 10(i) and (ii) follows from the hypothesis that $\neg a \vee \neg \neg a = 1$, for any $a \in A$. To see (iii), suppose (iv) $b \leq a$ and (v) $b \vee \neg b = 1$. Due to Lemma 10(ii), we have that (iv) implies $Da \leq Db$ and, using (DE), (v) implies $Db \leq \neg b$. So, by \leq -transitivity it follows that $Da \leq \neg b$. Then, in a residuated lattice we have $b \leq \neg Da$. \square

Remark 5 Given the conditions of Proposition 17, taking any of the coatoms in the algebra of Example 1, it is easy to see that $Dx \approx \neg Bx$ does not hold. Also, the reciprocal of Proposition 17 is not the case, as the algebra in Example 2 satisfies the equation $\neg x \vee \neg \neg x \approx 1$ and B exists in that algebra, but D does not exist.

Taking into account De Morgan laws valid in any MTL-algebra², we can easily obtain the following consequence of the previous proposition.

Corollary 2 Let \mathbf{A} be a SMTL-algebra, i.e. an MTL-algebra such that for all $a \in A$, $a \wedge \neg a = 0$. Then, for any $a \in A$, if Da exists, then so does Ba , and $Ba = \neg Da$.

Lemma 11 Let $\mathbf{A} \in \mathbf{RL}^D$ and $a \in A$. Then, the following are equivalent:

- (i) a is Boolean,
- (ii) $\neg Da = a$,
- (iii) $Da = \neg a$.

Proof (i) \Rightarrow (ii) Suppose $a \vee \neg a = 1$. Then, using (DE), $Da \leq \neg a$. Then, $a \leq \neg Da$. Now, by Lemma 10(i) we have $\neg Da \leq a$. So, $\neg Da = a$.

(ii) \Rightarrow (iii) Suppose $\neg Da = a$. Then, $\neg \neg Da = \neg a$. As $Da \leq \neg \neg Da$, we have that $Da \leq \neg a$. Now, by (DI), $a \vee Da = 1$. So, also $\neg a \leq Da$. Then, $Da = \neg a$.

(iii) \Rightarrow (i) Suppose $Da = \neg a$. As using (DI) we have $a \vee Da = 1$, it follows that $a \vee \neg a = 1$. \square

As a direct consequence we have the following fact.

Corollary 3 Let \mathbf{A} be a residuated lattice where both B and D exist. Then, $Da \leq \neg Ba$, for all $a \in A$. Equivalently, $Ba \leq \neg Da$, for all $a \in A$.

Proof Let $a \in A$. We have that $Ba \leq a$. Hence, using Lemma 10(ii), $Da \leq DBa$. Now, since Ba is Boolean, using Lemma 11 it follows that $Da \leq \neg Ba$. \square

Remark 6 The equality $B \approx \neg D$ does not hold. Indeed, consider the join-irreducible coatom c in the Heyting algebra in Figure 4, where $0 = Bc < \neg Dc$.

Lemma 12 Let \mathbf{A} be a \mathbf{RL}^D -algebra and $a, b \in A$. We have that if $b \leq a$ and $b \vee \neg b = 1$, then $b \leq \neg Da$.

Proof Suppose $b \leq a$. Then, $\neg Db \leq \neg Da$. Now, using the hypothesis $b \vee \neg b = 1$ and Lemma 11, we have that $\neg Db = b$. So, $b \leq \neg Da$. \square

² Notice that not all instances of De Morgan laws are valid in the variety of MTL-algebras, for instance the equations $\neg(x \wedge y) \approx \neg x \vee \neg y$ and $\neg(x \vee y) \approx \neg x \wedge \neg y$ are valid, but $x \wedge y \approx \neg(\neg x \vee \neg y)$ is not.

Proof Considering Propositions 13 and 21, it is enough to prove that $\Delta = DD$, which follows from $\neg\Delta\neg\Delta = \Delta$. As for any $a \in A$, Δa is Boolean due to **(ΔE2)**, using Lemma 9 it is enough to check that $\neg\neg\Delta = \Delta$, which follows again from **(ΔE2)** and Lemma 1(vi). \square

5 The logic FL_{ew}^B

In this section we introduce an expansion of FL_{ew} with a unary connective B , whose intended algebraic semantics is the variety of \mathbb{RL}^B -algebras studied in Section 3.

Indeed, we define FL_{ew}^B as the expansion of FL_{ew} with the following axiom schemas:

- (B1)** $B\varphi \rightarrow \varphi$,
- (B2)** $B\varphi \vee \neg B\varphi$,
- (B3)** $B(\varphi \vee \neg\varphi) \rightarrow (B\varphi \vee \neg\varphi)$,
- (B4)** $B(\varphi \rightarrow \psi) \rightarrow (B\varphi \rightarrow B\psi)$.

and the following additional rule:

- (B)** From φ derive $B\varphi$.

We denote (finitary) derivability in FL_{ew}^B by \vdash .

Note that we have the following facts.

Lemma 14 (i) $\vdash B\varphi \rightarrow BB\varphi$ and

(ii) $B\varphi \rightarrow \psi \vdash B\varphi \rightarrow B\psi$.

Proof For (i) check the following derivation:

- | | |
|---|---------------------------|
| 1. $B\varphi \vee \neg B\varphi$ | (B2) |
| 2. $B(B\varphi \vee \neg B\varphi)$ | 1, rule (B) |
| 3. $BB\varphi \vee \neg B\varphi$ | (B3) , 2, mp |
| 4. $BB\varphi \rightarrow (B\varphi \rightarrow BB\varphi)$ | FL_{ew} |
| 5. $\neg B\varphi \rightarrow (B\varphi \rightarrow BB\varphi)$ | FL_{ew} |
| 6. $B\varphi \rightarrow BB\varphi$ | 3, 4, 5, FL_{ew} |

(ii) follows easily using (i). \square

Clearly, FL_{ew}^B is a *Rasiowa implicative* logic (cf. [17]). Then, it follows that it is algebraizable in the sense of Blok and Pigozzi [3]. It is also straightforward to check that the variety \mathbb{RL}^B is its equivalent algebraic semantics. Algebraizability immediately implies strong completeness of FL_{ew}^B with respect to \mathbb{RL}^B .

Theorem 2 For every set $\Gamma \cup \{\varphi\}$ of formulas, $\Gamma \vdash \varphi$ iff for every $\mathbf{A} \in \mathbb{RL}^B$ and every \mathbf{A} -evaluation e , $e(\varphi) = 1$, whenever $e[\Gamma] \subseteq \{1\}$.

In FL_{ew}^B the usual form of the deduction theorem does not hold. Indeed, $\varphi \vdash B\varphi$, but $\not\vdash \varphi \rightarrow B\varphi$. Which fails in the three-element Gödel algebra $[0, \frac{1}{2}, 1]_G$, where $B1 = 1$ and $B\frac{1}{2} = B0 = 0$: for any evaluation e in this algebra, if $e(\varphi) = 1$, then $e(B\varphi) = 1$, but for $e(\varphi) = \frac{1}{2}$ we have $e(B\varphi) = 0$, and thus $e(\varphi \rightarrow B\varphi) = 0$.

Actually, FL_{ew}^B enjoys the same form of deduction theorem holding for logics with the Δ operator (cf. [2, Proposition 2.2]).

Theorem 3 $\Gamma, \varphi \vdash \psi$ iff $\Gamma \vdash B\varphi \rightarrow \psi$.

Proof \Rightarrow) We prove by induction on every formula χ_i ($1 \leq i \leq n$) of the given derivation of ψ from $\Gamma \cup \{\varphi\}$ that $\Gamma \vdash B\varphi \rightarrow \chi_i$. If $\chi_i = \varphi$, then the result follows due to axiom schema **(B1)**. If χ_i belongs to Γ or is an instance of an axiom, then the result follows using *modus ponens* and the derivability of the schema $\chi_i \rightarrow (B\varphi \rightarrow \chi_i)$. If χ_i comes by application of *modus ponens* on previous formulas in the derivation, then the result follows, because from $B\varphi \rightarrow \chi_k$ and $B\varphi \rightarrow (\chi_k \rightarrow \chi_i)$ we may derive $(B\varphi \& B\varphi) \rightarrow (\chi_k \& (\chi_k \rightarrow \chi_i))$ and then also $B\varphi \rightarrow \chi_i$, using transitivity of \rightarrow applied to the derivable formulas $B\varphi \rightarrow (B\varphi \& B\varphi)$ and $(\chi_k \& (\chi_k \rightarrow \chi_i)) \rightarrow \chi_i$. Finally, if $\chi_i = B\chi_k$ comes using rule **(B)** from formula χ_k , then from $B\varphi \rightarrow \chi_k$ we may derive $B\varphi \rightarrow B\chi_k$ using Lemma 14(ii).

\Leftarrow) To the derivation given by the hypothesis add a step with φ . In the next step put $B\varphi$, which follows from the previous formula using rule **(B)**. Finally, derive ψ using *modus ponens*. \square

Thanks to this B -deduction theorem, the logic FL_{ew}^B has the following property: if we expand FL_{ew}^B with any further rule $\varphi_1, \dots, \varphi_n / \varphi$, then it is possible to dispose of the rule just adding the axiom $(B\varphi_1 \wedge \dots \wedge B\varphi_n) \rightarrow \varphi$. This property also holds for the logics FL_{ew}^A and FL_{ew}^D .

Proposition 22 FL_{ew}^B is a conservative expansion of FL_{ew} .

Proof Use Proposition 4 and the Finite Model Property of FL_{ew} (see [15]). \square

One could analogously define the expansion of MTL with the pre-linearity axiom $(\varphi \rightarrow \psi) \vee (\psi \rightarrow \varphi)$ with B , with the same additional axioms and rule, yielding the logic MTL^B , which is algebraizable and strongly complete with respect to the variety MTL^B of MTL^B -algebras. However, unlike the case of expansion with Δ , MTL^B is not a semilinear logic, that is, it is not complete with respect to the class of MTL^B -chains. The reason is that the \vee -form of rule **(B)**, “from $\psi \vee \varphi$ derive $\psi \vee B\varphi$ ”, is not derivable in MTL^B . Indeed, taking the coatoms s and t in the Gödel algebra of Example 1, it is clear that $s \vee t = 1$, while $s \vee Bt = s \vee 0 = s$.

As a final result, we can show that FL_{ew}^B inherits from FL_{ew} the Finite Model Property (FMP). Before proving this, we introduce some preliminary notation.

For a logic $L \in \{\text{FL}_{ew}, \text{FL}_{ew}^B\}$, let us denote by $Fm(L, Var)$ the set of L -formulas built from a set Var of propositional variables. Now let us define the enlarged set of propositional variables $Var^* = Var \cup \{“B\varphi” \mid B\varphi \in Fm(\text{FL}_{ew}^B, Var)\}$, where “ $B\varphi$ ” is intended to denote a fresh propositional variable, one

for each formula $B\varphi \in Fm(\mathbb{FL}_{ew}^B, Var)$. Then, we can define a one-to-one translation of every formula $\varphi \in Fm(\mathbb{FL}_{ew}^B, Var)$ into a formula $\varphi^* \in Fm(\mathbb{FL}_{ew}, Var^*)$, by just inductively defining:

- $0^* = 0$,
- if $\varphi = p \in Var$, then $\varphi^* = p$,
- if $\varphi = B\psi$, then $\varphi^* = "B\psi"$,
- if $\varphi = \psi \odot \chi$, then $\varphi^* = \psi^* \odot \chi^*$, for $\odot \in \{\wedge, \vee, \&, \rightarrow\}$.

If Γ is a set of formulas, we write $\Gamma^* = \{\varphi^* \mid \varphi \in \Gamma\}$. Note that for any $\psi \in Fm(\mathbb{FL}_{ew}, Var^*)$, there is a formula $\varphi \in Fm(\mathbb{FL}_{ew}^B, Var)$ such that $\varphi^* = \psi$.

Moreover, we need the following result that will allow us to reduce proofs in \mathbb{FL}_{ew}^B to proofs in \mathbb{FL}_{ew} .

Lemma 15 *Let T be the set of all instances of axioms of \mathbb{FL}_{ew}^B . For each set $\Gamma \cup \{\varphi\} \subseteq Fm(\mathbb{FL}_{ew}^B, Var)$, it holds that*

$$\Gamma \vdash_{\mathbb{FL}_{ew}^B} \varphi \text{ iff } \Gamma^* \cup Cg^* \cup T^* \vdash_{\mathbb{FL}_{ew}} \varphi^*,$$

where $Cg = \{B\varphi \leftrightarrow B\psi \mid \Gamma \vdash_{\mathbb{FL}_{ew}^B} \varphi \leftrightarrow \psi\}$.

The proof is quite straightforward and analogous to those of similar results that can be found in the literature in slightly different contexts.

Theorem 4 \mathbb{FL}_{ew}^B enjoys the FMP, that is, if $\Gamma \not\vdash_{\mathbb{FL}_{ew}^B} \varphi$, then there is a finite $\mathbf{A} \in \mathbb{RL}^B$ and an \mathbf{A} -evaluation e such that $e(\Gamma) = 1$ and $e(\varphi) < 1$.

Proof If $\Gamma \not\vdash_{\mathbb{FL}_{ew}^B} \varphi$, by Lemma 15, it holds that $\Gamma^* \cup Cg^* \cup T^* \not\vdash_{\mathbb{FL}_{ew}} \varphi^*$, and by strong completeness and FMP of \mathbb{FL}_{ew} , there is a finite algebra $\mathbf{C} \in \mathbb{RL}$ and \mathbf{C} -evaluation v such that $v(\Gamma^* \cup Cg^* \cup T^*) = 1$ and $v(\varphi^*) < 1$. Then, the result will follow from the following facts:

Claim 1: $G = \{v("B\varphi") \mid B\varphi \in Fm(\mathbb{FL}_{ew}^B, Var)\}$ is a set of Boolean elements of \mathbf{C} . “You were a little grave,”

Proof of the claim: It is enough to check that $v((B\varphi)^*) \vee \neg v((B\varphi)^*) = v((B\varphi)^* \vee \neg(B\varphi)^*) = v((B\varphi \vee \neg B\varphi)^*) = 1$, where the latter holds because $B\varphi \vee \neg B\varphi$ is the axiom (B2) of \mathbb{FL}_{ew}^B . \dashv

Claim 2: Let \mathbf{A} be the \mathbb{RL} -algebra generated by the set $X = \{v(\varphi) \mid \varphi \in Fm(\mathbb{FL}_{ew}, Var^*)\}$, which is finite since \mathbf{A} is a subalgebra of \mathbf{C} . Then, B exists in \mathbf{A} and $B(\mathbf{A}) = G$. Therefore, \mathbf{A} is indeed an \mathbb{RL}^B -algebra.

Proof of the claim: That \mathbf{A} is finite is obvious, and thus, by Proposition 5, B exists. On the other hand, the elements of G keep being Boolean in \mathbf{A} . Hence, the only missing thing to check is that any Boolean element of \mathbf{A} already belongs to G . This is also clear since Boolean elements are closed by propositional combinations with connectives. \dashv

Claim 3: Let us define the \mathbf{A} -evaluation (taking \mathbf{A} as \mathbb{RL}^B -algebra) $e : Var \rightarrow A$ defined by $e(p) = v(p)$. Then, for any φ , $e(\varphi) = v(\varphi^*)$, in particular, $e(B\varphi) = v("B\varphi")$.

Proof of the claim: We prove that $e(\varphi) = v(\varphi^*)$ by structural induction.

- if φ is a propositional variable, it holds by construction
 - if $\varphi = \psi \odot \chi$ for $\odot \in \{\wedge, \vee, \&, \rightarrow\}$, by induction hypothesis we have $e(\psi) = v(\psi^*)$ and $e(\chi) = v(\chi^*)$, and hence $e(\varphi) = e(\psi \odot \chi) = e(\psi) \odot e(\chi) = v(\psi^*) \odot v(\chi^*) = v(\psi^* \odot \chi^*) = v((\psi \odot \chi)^*) = v(\varphi^*)$.
 - If $\varphi = B\psi$, then we have to prove that $v("B\psi") = B(e(\psi))$, the latter being equal to $e(B\psi)$ by definition. Therefore, we have to prove in turn that the three defining conditions (BE1), (BE2), and (BI) are satisfied by $v("B\psi") = v((B\psi)^*)$ to be the greatest Boolean below $e(\psi)$, assuming by induction that $v(\psi^*) = e(\psi)$.
- (BE1) Since $B\psi \rightarrow \psi$ is axiom (BE1) of \mathbb{FL}_{ew}^B , we have that $1 = v((B\psi \rightarrow \psi)^*) = v((B\psi)^* \rightarrow v(\psi^*)) = v((B\psi)^* \rightarrow e(\psi))$. Hence, $v((B\psi)^*) \leq e(\psi)$.
- (BE2) is clear from Claim 1.
- (BI) We have to check that if $b \in B(A)$ is such that $b \leq e(\psi) = v(\psi^*)$, then $b \leq v((B\psi)^*)$. If $b \in B(A)$, by construction of \mathbf{A} , then there exists a formula χ such that $b = v((B\chi)^*)$. On the other hand, by (ii) of Lemma 12, we know that $B\chi \rightarrow \psi, B\chi \vee \neg B\chi \vdash B\chi \rightarrow B\psi$. Thus, we also know that if $v((B\chi)^*) \leq v(\psi^*)$ and $v((B\chi)^*) \vee \neg v((B\chi)^*) = 1$, then $v((B\chi)^*) \leq v((B\psi)^*)$. Now, the two conditions are satisfied, hence we have $b = v((B\chi)^*) \leq v((B\psi)^*)$.

This closes the proof of Claim 3. \dashv

Finally, from these claims it readily follows that $e(\Gamma) = v(\Gamma^*) = 1$ and $e(\varphi) = v(\varphi^*) < 1$, as required. \square

6 Conclusions and dedication

In this paper we have considered the expansion of \mathbb{FL}_{ew} with the operator B , that in algebraic terms provides the greatest Boolean below a given element of a residuated lattice. Among other things, we have axiomatized it and shown that the resulting logic is a conservative expansion enjoying the Finite Model Property. The axioms for B turn out to be very close to those of the Monteiro-Baaz Δ operator, in fact only one axiom is a weaker version of the one for Δ . Even if the properties are very similar, that small difference causes, e.g. that in the context of MTL, the expansion with B is not any

longer a semilinear logic, in contrast to the expansion with Δ .

As a matter of fact, we have chosen this topic for our humble contribution to honour the memory of our beloved and late friend Franco Montagna, because it was suggested by Franco to the first author during the preparation of their joint manuscript [1], together with Amidei, where they study the expansion of FL_{ew} and other substructural logics with Δ .

Acknowledgments and Compliance with Ethical Standards

The authors are thankful to an anonymous reviewer for his/her comments that have helped to improve the final layout of this paper. The authors have been funded by the EU H2020-MSCA-RISE-2015 project 689176-SYSMICS. Esteve and Godo have been also funded by the FEDER/MINECO Spanish project TIN2015-71799-C2-1-P. The authors declare that they have no conflict of interest. This article does not contain any studies with human participants or animals performed by any of the authors.

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