

A novel multi-image cryptosystem based on weighted plain images and using combined chaotic maps

Ahmad Pourjabbar Kari (✉ a.pourjabbar@gmail.com)

Islamic Azad University Tehran North Branch <https://orcid.org/0000-0001-5303-3453>

Ahmad Habibizad Navin

Islamic Azad University Tabriz Branch

Amir Massoud Bidgoli

Islamic Azad University Tehran North Branch

Mirkamal Mirnia

Tabriz University: University of Tabriz

Original Research Paper

Keywords: Multimedia security, Image encryption, Information entropy, Noise attack, Chaotic maps, Henon map

Posted Date: February 3rd, 2021

DOI: <https://doi.org/10.21203/rs.3.rs-164388/v1>

License:  This work is licensed under a Creative Commons Attribution 4.0 International License.

[Read Full License](#)

Version of Record: A version of this preprint was published at Multimedia Systems on March 13th, 2021. See the published version at <https://doi.org/10.1007/s00530-021-00772-y>.

Abstract

This paper introduces a new multi-image cryptosystem based on modified Henon map and nonlinear combination of chaotic seed maps. Based on the degree of correlation between the adjacent pixels of the plain image, a unique weight is assigned to the plain image. First, the coordinates of plain images are disrupted by modified Henon map as confusion phase. In the first step of diffusion phase, the pixels content of images are changed separately by XOR operation between confused images and matrices with suitable nonlinear combination of seed maps sequences. These combination of seed maps are selected depending on the weight of plain images as well as bifurcation properties of mentioned chaotic maps. After concatenating the matrices obtained from the first step of diffusion phase, the bitwise XOR operation is applied between newly developed matrix and the other produced matrix from the chaotic sequences of the Logistic-Tent-Sine hybrid system, as second step of diffusion phase. The encrypted image is obtained after applying shift and exchange operations. The results of the implementation using graphs and histograms show that the proposed scheme, compared to some existing methods, can effectively resist common attacks and can be used as a secure method for encrypting digital images.

Full Text

This preprint is available for [download as a PDF](#).

Figures

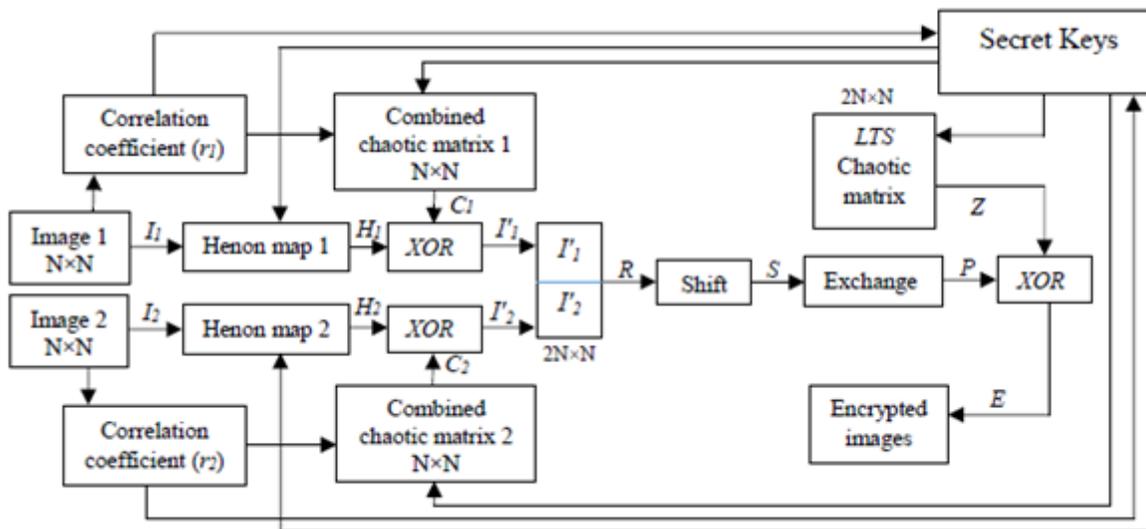


Figure 1

Block diagram of the proposed method

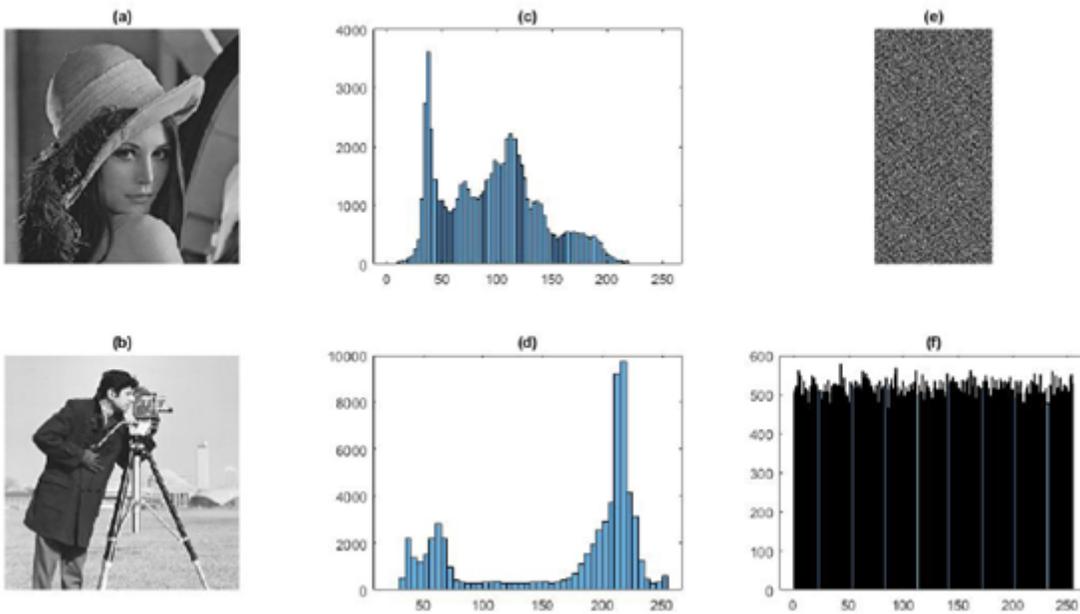


Figure 2

(a) standard 256×256 "Lena" plain image, (b) standard 256×256 "Camera man" plain image, second column (c), (d) are plain images with non-uniform histograms respectively, (e) is the 512×256 two plain image single encrypted image, (f) is the uniform histogram of double encrypted images

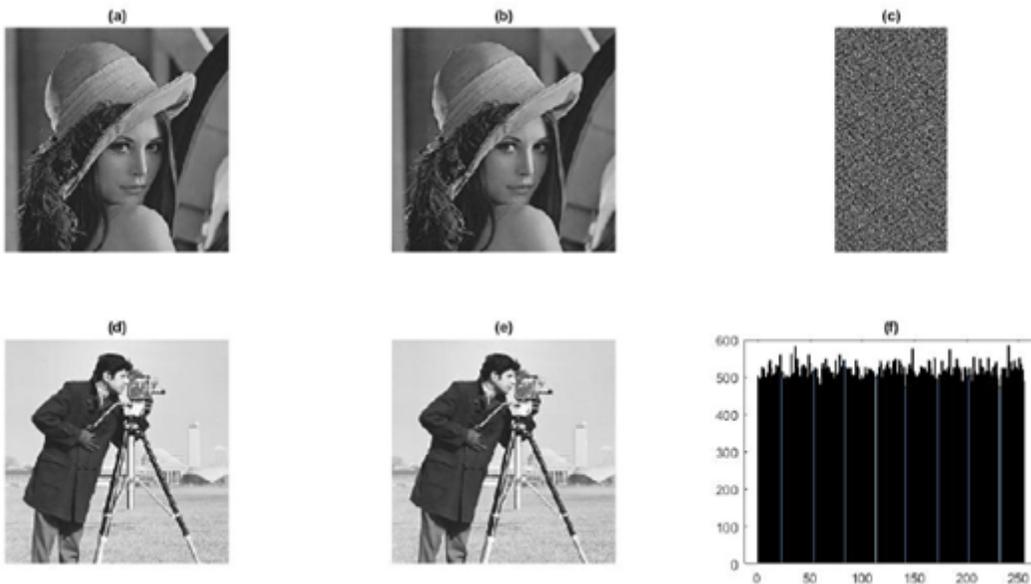


Figure 3

First column (a), (d) are two 256×256 plain images, second column (b), (e) are decrypted images respectively, (c) is the 512×256 two plain image single encrypted image, (f) is the uniform histogram of double encrypted images

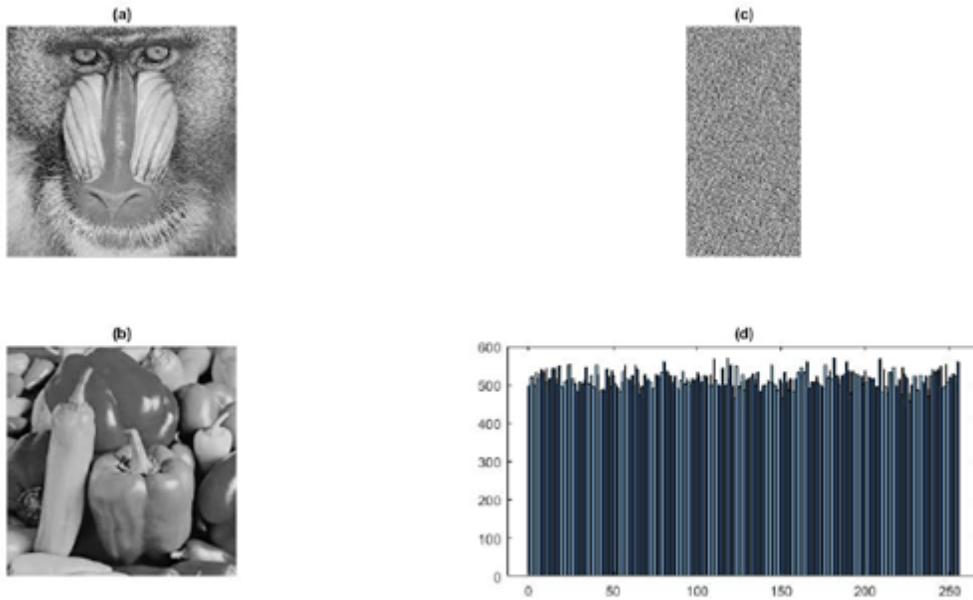


Figure 4

(a) standard 256×256 "Baboon" plain image, (b) standard 256×256 "Peppers" plain image, (c) is the 512×256 two plain image single encrypted image, (d) is the uniform histogram of encrypted images

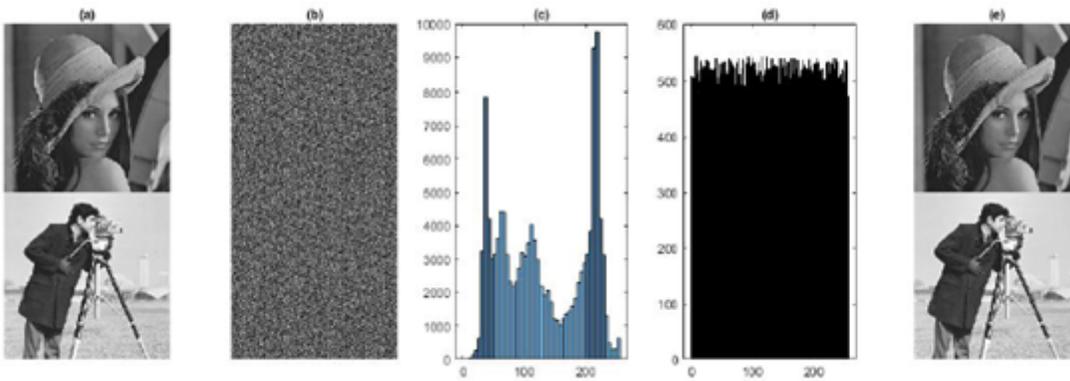


Figure 5

(a) is 512×256 merged two 256×256 plain images, (b) is the 512×256 two plain image single encrypted image, (c) is the histogram of two concatenated plain images, (d) is the uniform histogram of encrypted images, (e) is the decrypted image

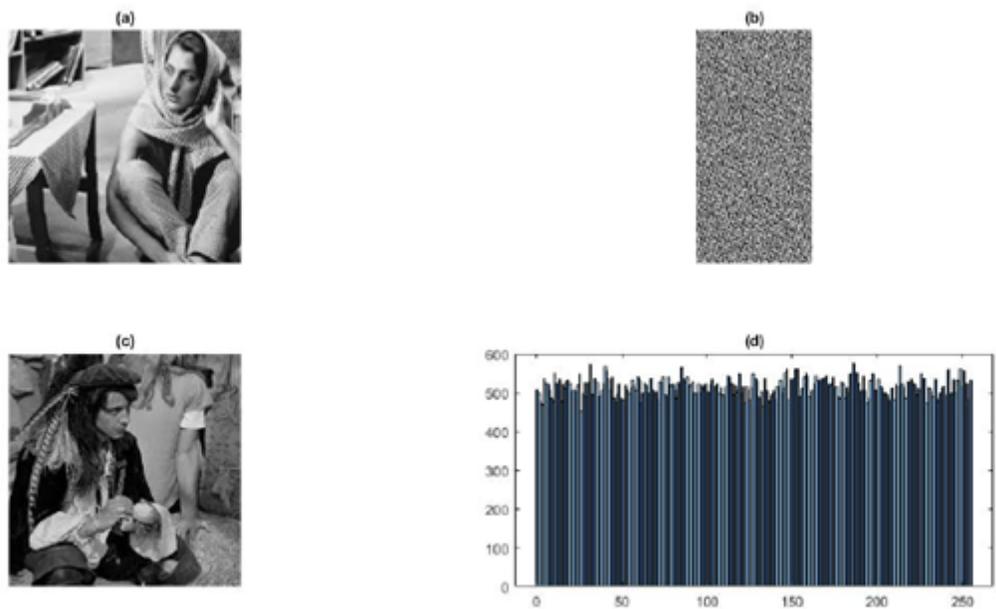


Figure 6

(a) standard 256×256 "Female" plain image, (c) standard "Male" 256×256 plain image, (b) is the 512×256 two plain image single encrypted image, (d) is the uniform histogram of encrypted images

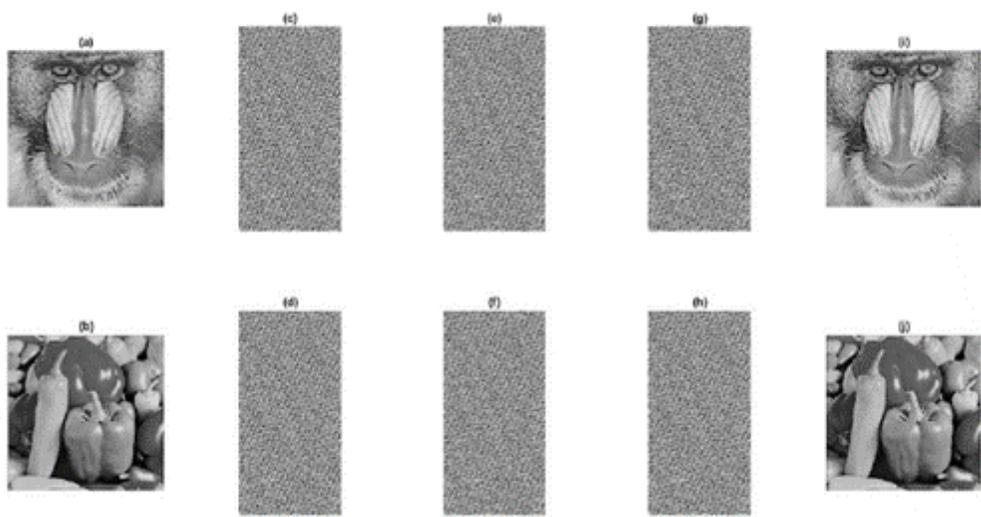


Figure 7

(a) and (b) are 256×256 plain images, (c) is the 512×256 encrypted image of two plain images with secret key $Ke_1=0.78$, (d) is the 512×256 encrypted image of two plain images with secret key $Ke_2=0.83$, (e) is the 512×256 decrypted image of two plain images with secret key $Ke_3=0.78+10^{-14}$, (f) is the 512×256 decrypted image of two plain images with secret key $Ke_4=0.83+10^{-14}$, (g) is the 512×256 decrypted image of two plain images with secret key $Ke_5=0.93$, (h) is the 512×256 decrypted image of

two plain images with secret key $Ke_4=0.63$, (i) is the decrypted image of plain image (a) with correct secret key $Ke_1=0.78$, (j) is the decrypted image of plain image (b) with correct secret key $Ke_2=0.83$

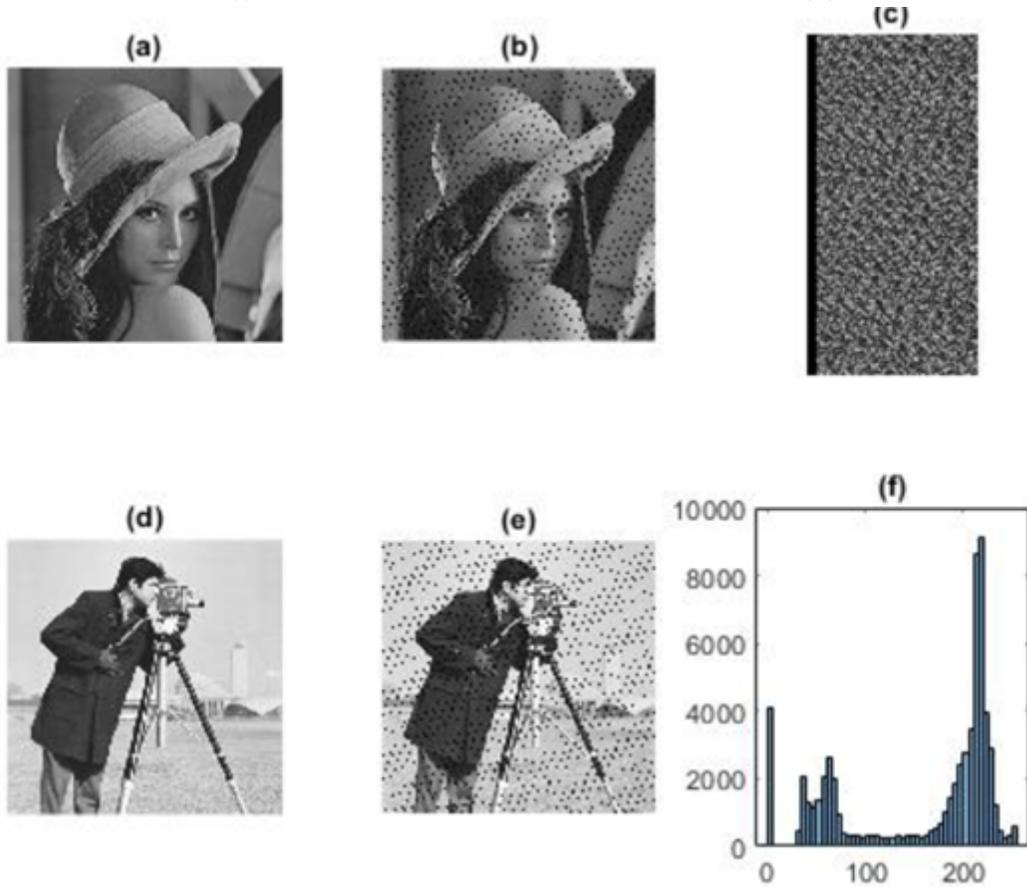


Figure 8

(a) and (d) are 256×256 plain images, (b) is the decrypted image for (a) after 1/16 cropping attack in one round, (e) is the decrypted image for (d) after 116 cropping attack in one round, (c) is the 512×256 two plain image single encrypted image after 1/16 cropping attack, (f) is the histogram of (e)

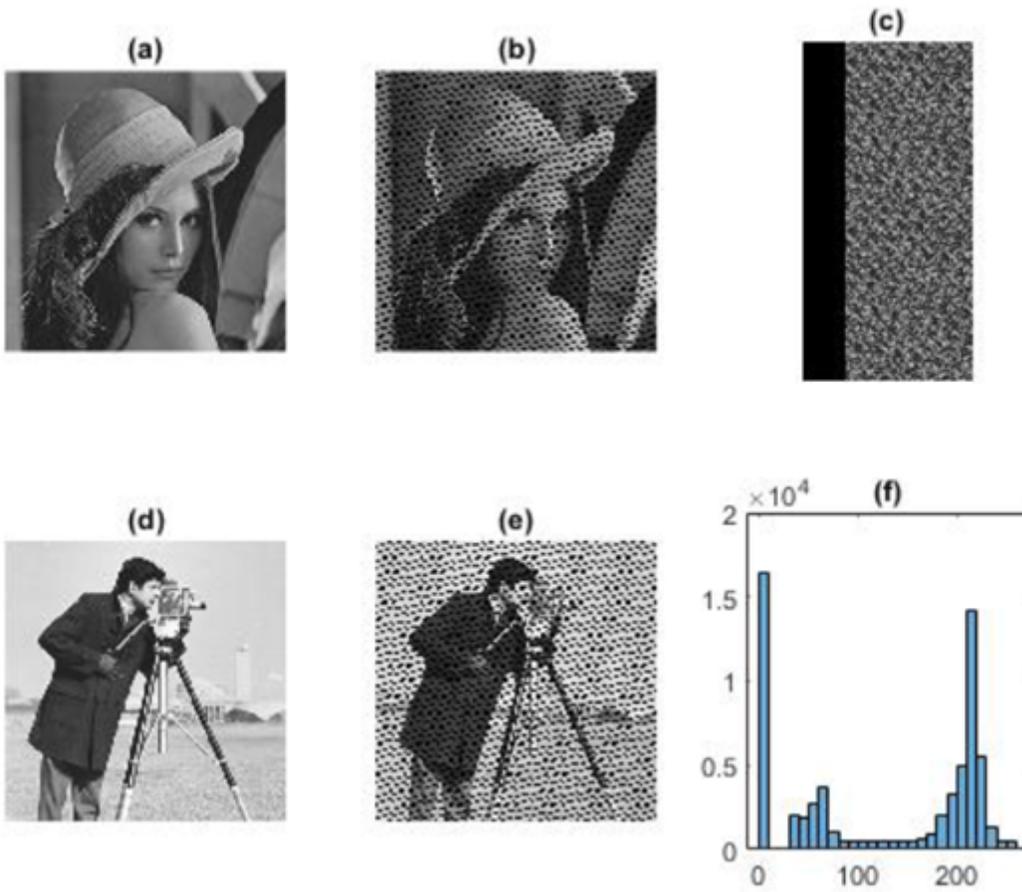


Figure 9

a) and (d) are 256×256 plain images, (b) is the decrypted image for (a) after 1/4 cropping attack, (e) is the decrypted image for (d) after 1/4 cropping attack, (c) is the 512×256 two plain image single encrypted image after 1/4 cropping attack, (f) is the histogram of (e)

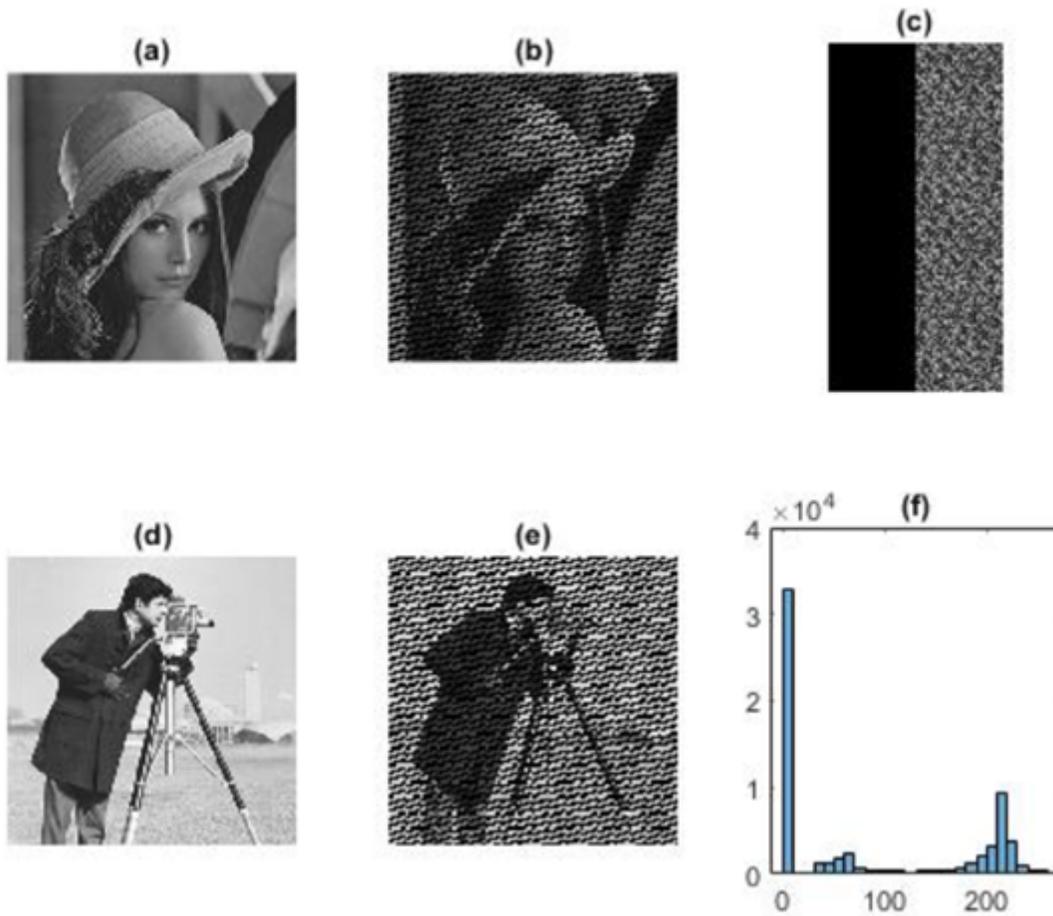


Figure 10

(a) and (d) are 256×256 plain images, (b) is the decrypted image for (a) after 1/2 cropping attack, (e) is the decrypted image for (d) after 1/2 cropping attack, (c) is the 512×256 two plain image single encrypted image after 1/2 cropping attack, (f) is the histogram of (e)



Figure 11

a) and (c) are 256×256 plain images, (b) is the decrypted image for (a) after 5% salt and pepper attack, (d) is the decrypted image for (c) after 5% salt and pepper attack

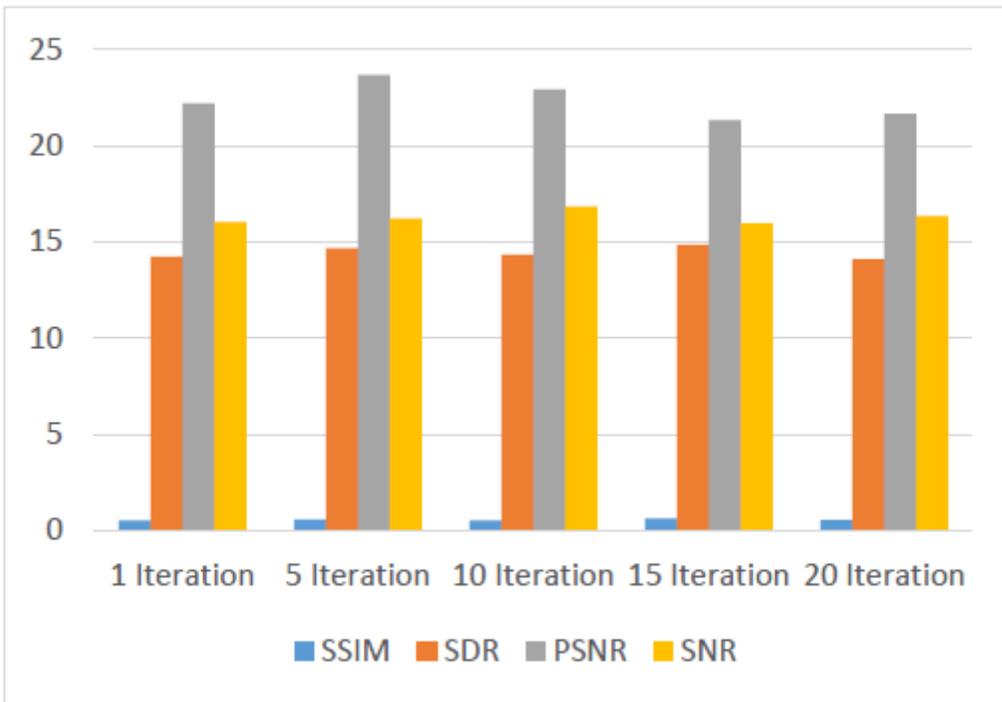


Figure 12

Evaluating parameters for various iterations, and 1/16 cropping attack

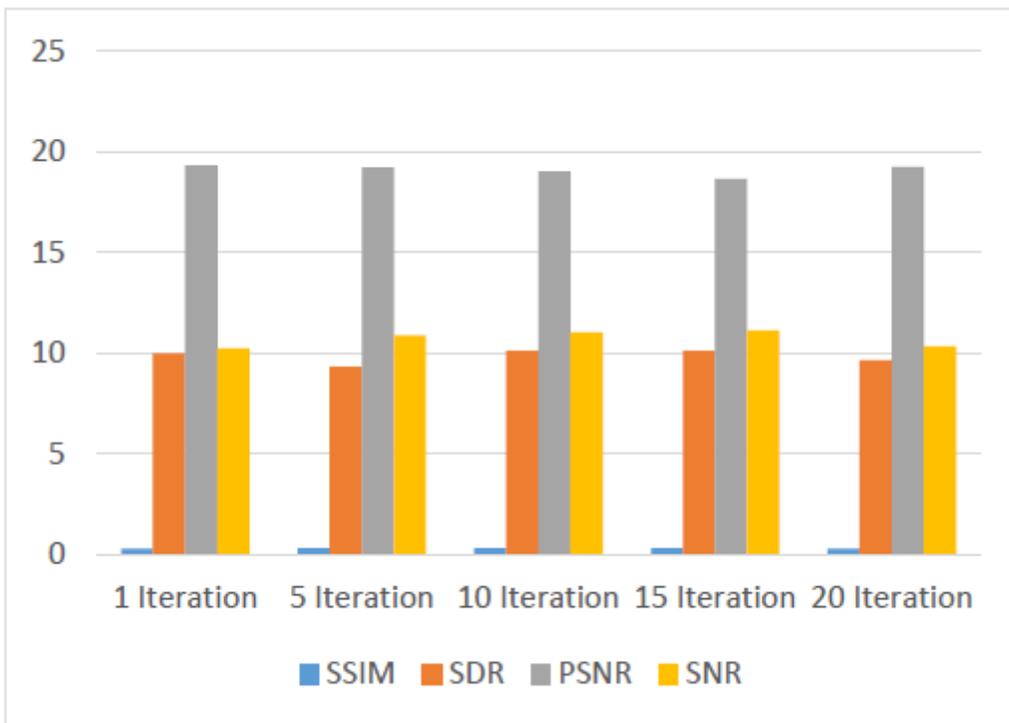


Figure 13

Evaluating parameters for various iterations, and 1/4 cropping attack

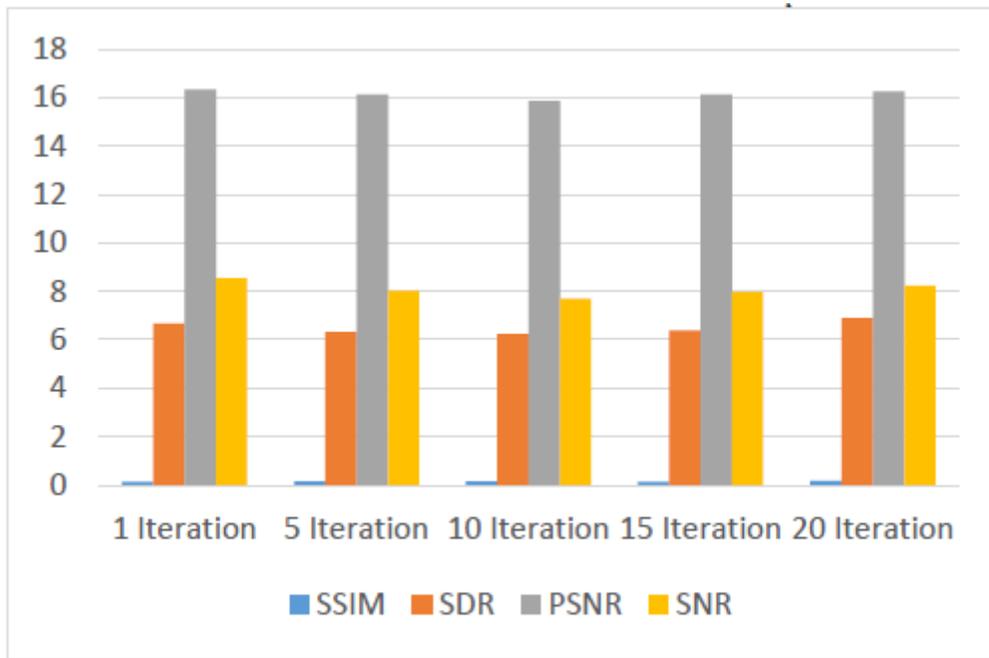


Figure 14

Evaluating parameters for various iterations, and 1/2 cropping attack