Compression of Periodic Complementary Sequences and Applications¹

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Abstract

A collection of complex sequences of length v is complementary if the sum of their periodic autocorrelation function values at all non-zero shifts is constant. For a complex sequence $A = [a_0, a_1, \ldots, a_{v-1}]$ of length v = dm we define the *m*-compressed sequence $A^{(d)}$ of length d whose terms are the sums $a_i + a_{i+d} + \cdots + a_{i+(m-1)d}$. We prove that the *m*-compression of a complementary collection of sequences is also complementary. The compression procedure can be used to simplify the construction of complementary $\{\pm 1\}$ -sequences of composite length. In particular, we construct several supplementary difference sets (SDS) $(v; r, s; \lambda)$ with v even and $\lambda = (r + s) - v/2$, given here for the first time. There are 15 normalized parameter sets $(v; r, s; \lambda)$ with $v \leq 50$ for which the existence question was open. We resolve all but one of these cases.

1 Introduction

We consider the ring $\mathbf{Z}_v = \{0, 1, \dots, v-1\}$ of integers modulo a positive integer v. Let k_1, \dots, k_t be positive integers and λ an integer such that

$$\lambda(v-1) = \sum_{i=1}^{t} k_i (k_i - 1), \tag{1}$$

and let X_1, \ldots, X_t be subsets of \mathbf{Z}_v such that

$$|X_i| = k_i, \quad i \in \{1, \dots, t\}.$$
 (2)

Definition 1 We say that X_1, \ldots, X_t are supplementary difference sets (SDS) with parameters $(v; k_1, \ldots, k_t; \lambda)$, if for every nonzero element $c \in \mathbf{Z}_v$ there are exactly λ ordered pairs (a, b) such that $a - b = c \pmod{v}$ and $\{a, b\} \subseteq X_i$ for some $i \in \{1, 2, \ldots, t\}$.

These SDS are defined over the cyclic group \mathbf{Z}_v . More generally SDS can be defined over any finite abelian group, and there are also further generalizations where the group may be any finite group. However, in this paper we shall consider only the cyclic case.

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In the context of an SDS, say X_1, \ldots, X_t , with parameters $(v; k_1, \ldots, k_t; \lambda)$, we refer to the subsets X_i as the *base blocks* and we introduce an additional parameter, n, defined by:

$$n = k_1 + \dots + k_t - \lambda. \tag{3}$$

The SDS formalism is a generalization of various well-known and studied designs.

SDS consisting of just one base block, the case t = 1, are called cyclic difference sets and are denoted by $(v; k; \lambda)$. The reader is referred to the classic book [2], the more recent book [17] and the recent survey [8]. The list of $(v; k; \lambda)$ cyclic difference sets with $v \leq 50$ in a normal form can be found in section 3 of [5].

The SDS with two base blocks, the case t = 2, covers several interesting designs. First, the circulant D-optimal designs [9, 6] corresponding to the case where v is odd and n = (v - 1)/2. Second, the binary complementary pairs (periodic analogs of complementary Golay pairs) [18] correspond to the case v = 2n. Third, when the two base blocks have the same size the SDS can be used to construct some BIBD (balanced incomplete block designs) [13, 17]. We shall denote the SDS with two base blocks of size r and s by $(v; r, s; \lambda)$. Upon imposing the normalization condition $\frac{v}{2} \ge r \ge s \ge 2$, it can be seen that for $v \le 50$ there are 227 feasible parameter sets, see [5]. These parameter sets are those that satisfy the conditions (1) and (2) with $k_1 = r$, $k_2 = s$. There is a non-existence result for SDS $(v; r, s; \lambda)$ obtained in [1], which can be used to eliminate some of these 227 feasible parameter sets. In [12] the authors construct 30 new $(v; r, s; \lambda)$ and show that an additional 19 $(v; r, s; \lambda)$ do not exist. In [5] the author constructs 8 new $(v; r, s; \lambda)$. We list here the remaining 15 undecided cases from [5]. They all have v in the interval $40 < v \le 50$.

(41;15,6;6)	n=15	(43;9,4;2)	n=11	(44;19,2;8)	n=13
(45;18,2;7)	n=13	(46;21,6;10)	n=17	(47;9,5;2)	n=12
(47;12,3;3)	n=12	(47;14,2;4)	n=12	(47;15,5;5)	n=15
(48;14,3;4)	n=13	(49;10,3;2)	n=11	(49;21,4;9)	n=16
(50;8,7;2)	n=13	(50;20,4;8)	n=16	(50;22,21;18)	n=25

In this paper (see Theorem 5) we decide the existence of all of them except for the case (49; 21, 4; 9). The main tools that we use for that purpose are computational techniques based on the Power Spectral Density (PSD) test as well as the method of SDS compression. For the PSD test see section 4. The compression method is studied systematically and in full generality (for arbitrary complex sequences and in particular for SDS) in section 5. We note that techniques that essentially amount to compression have been used previously [3, 11] in the context of D-optimal designs.

Apart from the cases where there exist Golay complementary pairs of length v, the binary complementary pairs are known to exist only for v = 34 [4] and v = 50 [5, 10]. In this paper we construct such pairs of length v = 50 and v = 58 (see the section 6). The example for v = 50has different parameters, namely (50; 22, 21; 18), from the previously constructed examples with parameters (50; 25, 20; 20). The first unknown case is now for v = 68.

2 The group ring approach to SDS

In the study of SDS it is often convenient to use the group ring of the additive group \mathbf{Z}_v with coefficients in \mathbf{C} , the field of complex numbers. This group ring can be identified with the quotient ring $\mathbf{C}[x]/(x^v-1)$, where $\mathbf{C}[x]$ is the polynomial ring over \mathbf{C} in a single indeterminate x and (x^v-1) is the ideal generated by the polynomial x^v-1 . By abuse of notation, we shall consider x also as an element of this group ring, in which case we have $x^v = 1$. Then \mathbf{Z}_v is embedded in this group ring via the map which sends $i \to x^i$, $i \in \mathbf{Z}_v$. Note that because $x^v = 1$ there is a ring homomorphism (evaluation at 1) $\mathbf{C}[x]/(x^v-1) \to \mathbf{C}$ which sends $x \to 1$.

To any subset $X \subseteq \mathbf{Z}_v$ we shall associate an element of the group ring, namely $X(x) := \sum_{i \in X} x^i$. By abuse of notation, we shall denote this element also by X, except that in the case $X = \mathbf{Z}_v$ we set

$$T = T(x) = 1 + x + \dots + x^{\nu-1}.$$

It will be clear from the context which meaning of X is used.

The group ring has an involution which sends each complex number c to its complex conjugate \overline{c} and sends x to its inverse $x^{-1} = x^{v-1}$. We shall denote this involution by an asterisk, e.g., we have $x^* = x^{-1}$ and $T^* = T$. For any element X of the group ring we define its *norm* to be the element $N(X) = XX^*$. We also define the function $N_X : \mathbf{Z}_v \to \mathbf{Z}$ by declaring that $N_X(s)$ is the coefficient of x^s in N(X), i.e., we have

$$N(X) = \sum_{i=0}^{v-1} N_X(i) x^i.$$

To any complex sequence of length v, say $A = [a_0, a_1, \ldots, a_{v-1}]$ we assign the element $A(x) = \sum_i a_i x^i$ of the group ring. Obviously, the map sending A to A(x) is injective. To simplify notation, we shall write A instead of A(x) if no confusion will arise. By evaluating A(x) at 1, we obtain $A(1) = \sum_i a_i$. Similarly, we have $N(A)(1) = |A(1)|^2$.

We point out that the SDS property is equivalent to an identity in the group ring. This is stated in the next lemma whose proof is straightforward.

Lemma 1 Let X_1, \ldots, X_t be subsets of \mathbb{Z}_v and assume that Eqs. (1) and (2) hold. Then these subsets are the base blocks of an SDS with parameters $(v; k_1, \ldots, k_t; \lambda)$ if and only if

$$\sum_{i=1}^{t} N(X_i) = n + \lambda T, \tag{4}$$

where n is defined by Eq. (3).

3 Power spectral density

Many known SDS have been constructed by a computer search. In such searches it is often convenient to replace a subset $X \subseteq \mathbf{Z}_v$ with the $\{\pm 1\}$ -sequence $A = [a_0, a_1, \ldots, a_{v-1}]$ where $a_i = -1$ if and only if $i \in X$. We shall refer to A as the associated sequence of X. In that case we have

$$A(x) = T(x) - 2X(x).$$
 (5)

Let $A = [a_0, a_1, \ldots, a_{v-1}]$ be an arbitrary complex sequences of length v, and let $\omega = e^{2\pi i/v}$ be a primitive v-th root of unity.

The Discrete Fourier Transform (DFT) of the sequence A is the function $\mathbf{Z}_v \to \mathbf{C}$ defined by the formula

$$DFT_A(s) = \sum_{j=0}^{\nu-1} a_j \omega^{js}.$$

We often identify a function on \mathbf{Z}_v with the sequence of its values, e.g., we write

$$DFT_A = B = [b_0, b_1, \dots, b_{v-1}], \text{ where } b_j = DFT_A(j).$$

This convention is used throughout the paper.

The *Power Spectral Density* (PSD) of the same sequence A is the function $\mathbf{Z}_v \to \mathbf{R}$ defined by the formula

$$\operatorname{PSD}_A(s) = |\operatorname{DFT}_A(s)|^2$$
.

The *Periodic Autocorrelation Function* (PAF) of A is defined as

$$\operatorname{PAF}_A(s) = \sum_{j=0}^{v-1} a_{j+s} \bar{a}_j.$$

We shall refer to the argument s as the *shift* variable.

It is straightforward to verify that for any complex sequence $A = [a_0, a_1, \ldots, a_{v-1}]$ and the corresponding element $A = A(x) = a_0 + a_1x + \cdots + a_{v-1}x^{v-1}$ of the group ring we have

$$N_A = \mathrm{PAF}_A.$$
 (6)

In general we have $N_A(-s) = \overline{N_A(s)}$, and for a real sequence A we have $N_A(-s) = N_A(s)$. Another important fact is the following classical theorem (see e.g. [14]).

Theorem 1 (Wiener–Khinchin) For any complex sequence $A = [a_0, a_1, \ldots, a_{v-1}]$, or equivalently any element $A = a_0 + a_1x + \cdots + a_{v-1}x^{v-1}$ of the group ring, we have

$$PSD_A = DFT(PAF_A).$$

By using Eq. (6) this can be rewritten as follows:

$$PSD_A(s) = \sum_{r=0}^{\nu-1} N_A(r)\omega^{rs}.$$
(7)

If the sequence A is real, then also

$$PSD_A(s) = \sum_{j=0}^{v-1} N_A(j) \cos \frac{2\pi j s}{v}.$$
 (8)

4 Complementary sequences

The associated sequences of the base blocks of an SDS have an important property known as complementarity. Let us begin with the general definition of this property.

Definition 2 Let A_1, \ldots, A_t be complex sequences of length v. We say that these sequences are complementary if

$$\sum_{i=1}^{i} \operatorname{PAF}_{A_i} = [\alpha_0, \underbrace{\alpha, \dots, \alpha}_{v-1 \text{ terms}}]$$

for some α_0 and α (the PAF-constants).

In the next theorem we show that a finite collection of complex sequences has constant sum of PAF values at nonzero shifts if and only if the same is true for the PSD values, and we find simple formulae expressing the PSD-constants in terms of the PAF-constants. (This theorem is a generalization of [7, Theorem 2].)

Theorem 2 Let A_1, \ldots, A_t be complex sequences of length v. These sequences are complementary, *i.e.*,

$$\sum_{i=1}^{t} \operatorname{PAF}_{A_i} = [\alpha_0, \underbrace{\alpha, \dots, \alpha}_{v-1 \text{ terms}}]$$
(9)

if and only if

$$\sum_{i=1}^{t} \operatorname{PSD}_{A_i} = [\beta_0, \underbrace{\beta, \dots, \beta}_{v-1 \text{ terms}}].$$
(10)

The PAF-constants α_0 and α and the PSD-constants β_0 and β are related as follows:

$$\beta_0 = \alpha_0 + (v - 1)\alpha, \quad \beta = \alpha_0 - \alpha. \tag{11}$$

Proof Note that (9) can be written as:

$$\sum_{i=1}^{t} PAF(A_i) = (\alpha_0 - \alpha) \cdot [1, 0, \dots, 0] + \alpha \cdot [1, \dots, 1].$$

By applying the DFT and by using the Wiener–Khinchin theorem and the fact that DFT is a linear operator, this implies that

$$\sum_{i=1}^{t} PSD(A_i) = (\alpha_0 - \alpha) \cdot DFT[1, 0, \dots, 0] + \alpha \cdot DFT[1, \dots, 1]$$
$$= (\alpha_0 - \alpha) \cdot [1, \dots, 1] + \alpha \cdot [v, 0, \dots, 0]$$
$$= [\beta_0, \beta, \dots, \beta]$$

Hence (10) holds as well as (11). Conversely, assume that (10) holds. By applying the inverse DFT we obtain

$$\sum_{i=1}^{t} PAF(A_i) = (\beta_0 - \beta) \cdot DFT^{-1}[1, 0, \dots, 0] + \beta \cdot DFT^{-1}[1, \dots, 1]$$

= $\frac{\beta_0 - \beta}{v}[1, \dots, 1] + \beta[1, 0, \dots, 0]$
= $[\alpha_0, \alpha, \dots, \alpha]$

By solving the equations (11) we obtain that

$$\alpha_0 = \frac{\beta_0 - \beta}{v} + \beta, \quad \alpha = \frac{\beta_0 - \beta}{v}.$$

Note that the PSD values are always nonnegative. Hence, if A_1, \ldots, A_t are complementary complex sequences of length v with the PSD-constants β_0 and β , then for $r = 1, \ldots, t$ we have

$$PSD_{A_r}(s) \le \beta, \quad s = 1, 2, \dots, v - 1.$$
 (12)

We shall refer to this inequality as the *PSD-test*.

In the following proposition we show that the associated sequences of the base blocks of an SDS are complementary sequences and we compute the PAF and PSD-constants. In a special case these constants were computed in [7, Example 3].

Proposition 1 Let A_1, \ldots, A_t be the $\{\pm 1\}$ -sequences associated to the base blocks X_1, \ldots, X_t of an SDS with parameters $(v; k_1, \ldots, k_t; \lambda)$. Then

$$\sum_{i=1}^{t} N(A_i) = 4n + (tv - 4n)T.$$
(13)

Moreover, the sequences A_1, \ldots, A_t are complementary with PAF-constants

$$\alpha_0 = tv, \quad \alpha = tv - 4n, \tag{14}$$

and PSD-constants

$$\beta_0 = tv, \quad \beta = 4n. \tag{15}$$

Proof Note that $x^i T = T$ for each *i*, and so we have $T^2 = vT$ and $TX_i = TX_i^* = k_i T$. By Eq. (5) we have $N(A_i) = (T - 2X_i)(T - 2X_i^*)$ which gives

$$N(A_i) = (v - 4k_i)T + 4N(X_i).$$
(16)

Summing over all i = 1, ..., t and by using Eq. (4), we obtain Eq. (13). Using Eq. (6), we have

$$N(A_i) = \sum_{j=0}^{v-1} \operatorname{PAF}_{A_i}(j) x^j.$$

By summing these equations over i = 1, ..., t and by using (13), we deduce that

$$\sum_{i=1}^{t} \operatorname{PAF}_{A_i}(j) = \begin{cases} tv - 4n, & j \neq 0, \\ tv, & j = 0. \end{cases}$$

Hence, the equations (14) hold, and by using Theorem 2 we deduce that the equations (15) also hold.

When searching for SDS, we can discard the trial base block X_r if the associated sequence A_r fails the PSD-test: $PSD_{A_r}(s) > \beta$ for some $s \neq 0$. In this case we can restate the PSD-test as follows.

Lemma 2 Under the hypotheses of Proposition 1, for r = 1, ..., t we have

$$\sum_{j=1}^{\nu-1} N_{X_r}(j) \cos \frac{2\pi j s}{\nu} \le n - k_r, \quad s \in \{1, \dots, \nu - 1\}.$$
(17)

Proof Recall that the DFT of a constant function vanishes for all shifts s except for s = 0. Hence, by using Eqs. (6) and (16), we obtain that $DFT(PAF_{A_r})(s) = 4 \cdot DFT(N_{X_r})(s)$. Thus, we must have

$$4\left(k_r + \sum_{j=1}^{v-1} N_{X_r}(j)\cos\frac{2\pi js}{v}\right) \le \beta = 4n, \quad s \ne 0,$$

i.e., (17) holds.

5 Compression of sequences

If we have a collection of complementary sequences of length v = dm, then we can compress them to obtain complementary sequences of length d. We refer to the ratio v/d = m as the *compression factor*. Here is the precise definition.

Definition 3 Let $A = [a_0, a_1, \ldots, a_{v-1}]$ be a complex sequence of length v = dm and set

$$a_j^{(d)} = a_j + a_{j+d} + \ldots + a_{j+(m-1)d}, \quad j = 0, \ldots, d-1.$$
 (18)

Then we say that the sequence $A^{(d)} = [a_0^{(d)}, a_1^{(d)}, \dots, a_{d-1}^{(d)}]$ is the m-compression of A.

Let us now show that the complementarity is preserved by the compression process. At the same time we shall compute the PAF and PSD-constants of the compressed sequences.

Theorem 3 Let $A_i = [a_{i0}, a_{i1}, \ldots, a_{i,v-1}]$, $i = 1, \ldots, t$, be complementary complex sequences of length v = dm with the PAF-constants α_0, α . Then the corresponding m-compressed sequences $A_i^{(d)} = [a_{i0}^{(d)}, \ldots, a_{i,d-1}^{(d)}]$, $i = 1, \ldots, t$, are complementary with PAF-constants:

$$\alpha_0^{(d)} = \alpha_0 + (m-1)\alpha, \quad \alpha^{(d)} = m\alpha.$$
 (19)

The PSD-constants of the original and compressed sequences are the same.

Proof By using (18) we compute $PAF_{A_i^{(d)}}(k)$ for arbitrary shift k:

$$PAF_{A_{i}^{(d)}}(k) = \sum_{\substack{j=0\\m-1}}^{d-1} a_{i,j+k}^{(d)} \overline{a_{i,j}^{(d)}} \\ = \sum_{r,s=0}^{m-1} \sum_{j=0}^{d-1} a_{i,j+k+rd} \overline{a_{i,j+sd}}.$$

By applying the substitution $r \to r + s$, we obtain

$$\operatorname{PAF}_{A_i^{(d)}}(k) = \sum_{\substack{r,s=0\\m-1}}^{m-1} \sum_{j=0}^{d-1} a_{i,j+sd+k+rd} \overline{a_{i,j+sd}}$$
$$= \sum_{r=0}^{m-1} \operatorname{PAF}_{A_i}(k+rd)$$

By summing over i from 1 to t, we obtain

$$\sum_{i=1}^{t} \mathrm{PAF}_{A_{i}^{(d)}}(k) = \sum_{r=0}^{m-1} \sum_{i=1}^{t} \mathrm{PAF}_{A_{i}}(k+rd).$$

Since $\operatorname{PAF}_{A_i}(k+rd)$ is equal to α_0 if k=r=0 and is equal to α otherwise, we conclude that the sequences $A_1^{(d)}, \ldots, A_t^{(d)}$ are complementary and that their PAF-constants are given by Eqs. (19). The assertion about the PSD-constants follows from Theorem 2.

By applying Theorem 3 directly to SDS with parameters $(v; k_1, \ldots, k_t; \lambda)$, we obtain the following result.

Theorem 4 Let A_1, \ldots, A_t be the sequences associated to the base blocks of an SDS with parameters $(v; k_1, \ldots, k_t; \lambda)$. Suppose that v = dm and let $A_1^{(d)}, \ldots, A_t^{(d)}$ be the corresponding m-compressed sequences. Then the PAF-constants of the compressed sequences are given by

$$\alpha_0^{(d)} = m(tv - 4n) + 4n, \quad \alpha^{(d)} = m(tv - 4n), \tag{20}$$

where n is defined by Eq. (3).

Proof By Proposition 1, the sequences A_1, \ldots, A_t are complementary with PAF-constants given by Eqs. (14). It remains to apply Theorem 3.

Remark 1 Theorem 4 generalizes [6, Theorem 1], which deals with the special case: SDS $(v; r, s; \lambda)$ with $n = r + s - \lambda = (v - 1)/2$. (These SDS are used to construct circulant D-optimal matrices.)

Remark 2 Theorem 4 has two useful corollaries for

- m = 2, in which case $|a_{ij}^{(d)}| \in \{0, 2\}$,
- m = 3, in which case $|a_{ij}^{(d)}| \in \{1, 3\}$.

An easy counting argument allows us to compute the number of $a_{ij}^{(d)}$ that assume a specific absolute value.

Corollary 1 With the hypotheses and notation of Theorem 4, let m = 2 and denote by ν_0, ν_2 the number of $a_{ij}^{(d)}$ terms in all 2-compressed sequences $A_i^{(d)}$ that have absolute value equal to 0, 2 respectively. Then $\nu_0 = n$ and $\nu_2 = td - n$.

Proof Since m = 2, from Eq. (20) we have $\alpha_0^{(d)} = 4(td - n)$. Therefore $0 \cdot \nu_0 + 4 \cdot \nu_2 = 4(td - n)$, i.e., $\nu_2 = td - n$. As $\nu_0 + \nu_2 = td$, we have $\nu_0 = n$.

Corollary 2 With the hypotheses and notation of Theorem 4, let m = 3 and denote by ν_1, ν_3 the number of $a_{ij}^{(d)}$ terms in the 3-compressed sequences $A_i^{(d)}$ that have absolute value equal to 1, 3 respectively. Then $\nu_1 = n$ and $\nu_3 = td - n$.

Proof Since m = 3, from Eq. (20) we have $\alpha_0^{(d)} = 9td - 8n$. Therefore $1 \cdot \nu_1 + 9 \cdot \nu_3 = 9td - 8n$. As $\nu_1 + \nu_3 = td$, we have $\nu_1 = n$ and $\nu_3 = td - n$.

6 Computational results for SDS with two base blocks

6.1 Necklaces, bracelets and charmed bracelets

When compressing the pair A, B of periodic complementary $\{\pm 1\}$ sequences of length v = dm, with compression factor m, we obtain the pair $A^{(d)}, B^{(d)}$ of periodic complementary sequences of length d whose elements are integers belonging to the set $\{m, m - 2, \ldots, 2 - m, -m\}$. The first equation of (15) shows that the sum of squares of the elements of $A^{(d)}$ and $B^{(d)}$ together is equal to 2v. In our applications we only used the compression factors m = 2 and m = 3. In these two cases the Corollaries 1 and 2 determine the total number of squares 0^2 , 2^2 and 1^2 , 3^2 , respectively. There are several ways to distribute these squares over the sequences $A^{(d)}$ and $B^{(d)}$, and we treat each of them separately.

We shall use combinatorial objects known as necklaces, bracelets and charmed bracelets [15, 16] Roughly speaking the *necklaces* are cyclic arrangements of v objects where we do not distinguish the arrangements obtained by cyclic shifts. If we enlarge the equivalence classes by allowing also that the arrangement be reversed, then we refer to these equivalence classes as *bracelets*. We can further enlarge these equivalence classes by moving, for each i, the object in position i to the position $si \pmod{v}$ where s is a fixed integer relatively prime to v. We shall refer to this operation as a *multiplication* by s. Note that the multiplication by s = -1 is the same as the reversal. We refer to these enlarged equivalence classes as *charmed racelets*. We emphasize that the multiplication operation does change the PAF of a sequence, but when performed simultaneously on a pair of complementary sequences it preserves the complementarity property.

The sequences A and B, as well as the compressed sequences $A^{(d)}$ and $B^{(d)}$, can be viewed as being cyclic arrangements of length v and d, respectively. The cyclic shifts and the reversal do not change the periodic autocorrelation of the sequence, and so they can be applied separately on A and B without destroying the complementarity property. However, the multiplication by $s \neq -1 \pmod{v}$ has to be performed simultaneously, with the same s, on both A and B as otherwise the complementarity property may be destroyed.

In our searches using the compression method we first construct the compressed sequences $A^{(d)}$ and $B^{(d)}$. They must have the specified row sums and pass the PSD test. We choose the first sequence to be a representative of a charmed bracelet while the second sequence can be chosen only as a representative of an ordinary bracelet. Then we proceed to select the pairs $A^{(d)}$, $B^{(d)}$ which are complementary. If there are no such pairs this means that the SDS that we are looking for do not exist. Otherwise we examine each of the complementary pairs $A^{(d)}$, $B^{(d)}$ and try to lift them to obtain a complementary pair A, B.

6.2 The range $v \le 50$

We shall prove the following theorem resolving the existence of all but one of the 15 open cases from the list in the Introduction. The only remaining undecided case is now (49; 21, 4; 9).

Theorem 5

- There do not exist SDS with following parameters: (41; 15, 6; 6) (43; 9, 4; 2) (44; 19, 2; 8) (45; 18, 2; 7) (46; 21, 6; 10) (47; 9, 5; 2) (47; 12, 3; 3) (47; 14, 2; 4) (47; 15, 5; 5) (48; 14, 3; 4) (49; 10, 3; 2) (50; 8, 7; 2) (50; 20, 4; 8)
- There exist SDS with the following parameters: (50; 22, 21; 18)

Proof We proceed by examining the 14 $(v; r, s; \lambda)$ cases one by one.

- 1. For the case (41; 15, 6; 6) the application of the PSD test with the constant value 60 gave 1040 normalized candidate A-sequences and 13104 normalized candidate B-sequences. Subsequently, there was no match found between these two sets of candidate sequences.
- 2. For the case (43; 9, 4; 2) there are no A-sequences that pass the PSD test with the constant value 44, therefore we conclude immediately that such SDS do not exist.
- 3. For the case (44; 19, 2; 8) there are no A-sequences that pass the PSD test with the constant value 52, therefore we conclude immediately that such SDS do not exist.
- 4. (45; 18, 2; 7) there are no A-sequences that pass the PSD test with the constant value 52, therefore we conclude immediately that such SDS do not exist.
- 5. For the case (46; 21, 6; 10) we used 2-compression to prove that there do not exist such SDS. Here are the details of this computation. First note that in this case we have $n = r + s \lambda = 21 + 6 10 = 17$. Taking m = 2, d = 23 in Theorem 4, we see that we need to find two sequences of length 23 with elements from $\{-2, 0, +2\}$ such that their PSD values add up to $4n = 4 \cdot 17 = 68$. Since we know from corollary 1 that the total number of 0 elements in the two sequences is equal to $\nu_2 = 29$, we deduce that there are four cases to consider, using ordinary and charmed bracelets [15]. For each one of these four cases: (1) we compute the charmed bracelets that give 2-compressed sequences of length 23 all of whose PSD values are smaller than 68 and (2) we compute the ordinary bracelets that give 2-compressed sequences are smaller than 68. We summarize the results in the following tables.

A-sequences						
Case	charmed bracelets	charmed bracelets passing PSD				
1	2,116,296	85				
2	475,020	$2,\!009$				
3	54,264	4,552				
4	$3,\!015$	1,442				
B-sequences						
	B-s	sequences				
Case		sequences ordinary bracelets passing PSD				
Case 1		1				
	ordinary bracelets	ordinary bracelets passing PSD				
1	ordinary bracelets 2,277	ordinary bracelets passing PSD 1,749				

Notice that case 4 can be eliminated at this stage, since there are no ordinary bracelets that pass the PSD test. The next step is for each of the remaining three cases to look for pairs of sequences that have constant PAF and we summarize the results in the following table (where the symbol \rightarrow indicates removal of sequences with duplicate PAF)

Case	A-sequences	B-sequences	# of pairs
1	$85 \rightarrow 84$	$1,\!749 \rightarrow 1,\!716$	39
2	$2{,}009 \rightarrow 1{,}970$	$1{,}419 \rightarrow 1{,}419$	34
3	$4{,}552 \rightarrow 4{,}497$	$22 \rightarrow 22$	0

Notice that case 3 can be eliminated at this stage, since there are no pairs of matching sequences at all. Subsequently we used the 39 + 34 = 73 matching pairs of sequences of length 23 with elements from $\{-2, 0, +2\}$ to construct all the corresponding pairs of uncompressed sequences of length 46 and check whether they form an SDS with the required parameters. We did not find any such SDS.

- 6. For the case (47; 9, 5; 2) there are no A-sequences that pass the PSD test with the constant value 48, therefore we conclude immediately that such SDS do not exist.
- 7. For the case (47; 12, 3; 3) there are no A-sequences that pass the PSD test with the constant value 48, therefore we conclude immediately that such SDS do not exist.
- 8. For the case (47; 14, 2; 4) there are no A-sequences that pass the PSD test with the constant value 48, therefore we conclude immediately that such SDS do not exist.
- 9. For the case (47; 15, 5; 5) there are no A-sequences that pass the PSD test with the constant value 60, therefore we conclude immediately that such SDS do not exist.
- 10. For the case (48; 14, 3; 4) there are no A-sequences that pass the PSD test with the constant value 52, therefore we conclude immediately that such SDS do not exist.
- 11. For the case (49; 10, 3; 2) there are no A-sequences that pass the PSD test with the constant value 44, therefore we conclude immediately that such SDS do not exist.
- 12. For the case (50; 8, 7; 2) the application of the PSD test with the constant value 52 gave 1130 normalized candidate A-sequences and 2910 normalized candidate B-sequences. Subsequently, there was no match found between these two sets of candidate sequences.
- 13. For the case (50; 20, 4; 8) we used 2-compression to prove that there do not exist such SDS.
- 14. We give four non-equivalent examples of SDS (50; 22, 21; 18):

 $\{0, 1, 2, 3, 6, 7, 9, 13, 14, 16, 18, 20, 22, 23, 26, 27, 30, 35, 37, 41, 45, 46\}$ $\{0, 1, 2, 3, 4, 5, 6, 8, 11, 12, 14, 17, 20, 22, 29, 30, 32, 37, 38, 39, 42\}$

 $\{0, 1, 2, 3, 4, 6, 7, 8, 9, 14, 16, 18, 20, 21, 25, 31, 32, 35, 36, 42, 44, 45\}$ $\{0, 1, 2, 4, 5, 8, 9, 10, 12, 14, 18, 21, 23, 24, 27, 29, 32, 34, 35, 39, 42\}$

 $\{0, 1, 2, 3, 5, 8, 9, 11, 14, 15, 19, 21, 24, 25, 29, 30, 32, 36, 38, 39, 41, 43\}$ $\{0, 1, 3, 5, 6, 7, 8, 9, 10, 13, 16, 18, 20, 21, 24, 25, 31, 32, 33, 37, 41\}$ $\{0, 2, 3, 4, 6, 9, 10, 12, 13, 17, 19, 20, 24, 25, 28, 29, 30, 33, 38, 39, 41, 47\}$ $\{0, 1, 3, 5, 6, 7, 8, 10, 12, 13, 14, 17, 20, 22, 24, 28, 32, 37, 38, 39, 40\}$

All four examples are in the canonical form defined in [5] and since they are different, this implies that they are non-equivalent. Note that since $2v - 4n = 2 \cdot 50 - 4 \cdot 25 = 0$, these SDS give rise to $\{\pm 1\}$ sequences of length 50 with PAF zero, because of (14).

6.3 The range v > 50

In this section we consider the SDS $(v; r, s; \lambda)$ with v = 2n, where $n = r + s - \lambda$. In particular v is even. Recall from Proposition 1 that these SDS give rise to binary periodic complementary pairs, i.e., pairs of $\{\pm 1\}$ sequences of length v with PAF constant $\alpha = 0$, because of (14).

Since 52 is a Golay number, there exist Golay pairs of length 52. At the same time they are also binary periodic complementary pairs. Since the Diophantine equations $x^2 + y^2 = 2 \cdot 54$ and $x^2 + y^2 = 2 \cdot 56$ do not have solutions, there are no binary periodic complementary pairs of lengths 54 and 56. Therefore, the first interesting value of v > 50 is v = 58, since this value is not excluded by the restriction of Arasu and Xiang [1]. In this paper we find 4 non-equivalent examples of SDS (58; 27, 24; 22), for the first time. These SDS give 4 respective pairs of binary sequences of length 58 with zero periodic autocorrelation function.

We give four non-equivalent examples of SDS (58; 27, 24; 22):

 $\{0, 1, 2, 3, 4, 7, 8, 10, 11, 12, 13, 16, 18, 20, 24, 26, 29, 31, 32, 33, 36, 38, 43, 46, 47, 50, 53\}$ $\{0, 1, 2, 3, 7, 8, 10, 11, 12, 13, 16, 17, 21, 22, 24, 27, 30, 34, 41, 42, 43, 45, 47, 49\}$

 $\{0, 1, 2, 3, 5, 6, 7, 9, 11, 12, 14, 15, 17, 19, 23, 24, 25, 26, 29, 32, 33, 39, 40, 43, 45, 48, 52\}$ $\{0, 1, 2, 3, 4, 5, 9, 11, 14, 15, 16, 18, 22, 26, 27, 31, 32, 34, 37, 39, 41, 42, 45, 51\}$

 $\{0, 1, 2, 3, 5, 8, 9, 11, 12, 13, 14, 18, 19, 21, 24, 25, 27, 29, 32, 34, 35, 39, 41, 43, 44, 48, 49\}$ $\{0, 2, 3, 4, 6, 8, 10, 13, 16, 17, 19, 20, 21, 25, 28, 29, 32, 33, 34, 39, 40, 41, 43, 46\}$

 $\{0, 2, 3, 4, 6, 7, 8, 10, 11, 14, 16, 17, 18, 20, 23, 25, 26, 28, 31, 32, 36, 37, 38, 41, 42, 47, 49\}$ $\{0, 1, 2, 3, 5, 8, 9, 10, 12, 16, 17, 18, 22, 25, 28, 30, 35, 37, 41, 44, 45, 46, 48, 49\}$

All four examples are in the canonical form defined in [5] and since they are different, this implies that they are non-equivalent.

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