# On the *p*-reinforcement and the complexity<sup>\*</sup>

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#### Abstract

Let G = (V, E) be a graph and p be a positive integer. A subset  $S \subseteq V$ is called a p-dominating set if each vertex not in S has at least p neighbors in S. The p-domination number  $\gamma_p(G)$  is the size of a smallest p-dominating set of G. The p-reinforcement number  $r_p(G)$  is the smallest number of edges whose addition to G results in a graph G' with  $\gamma_p(G') < \gamma_p(G)$ . In this paper, we give an original study on the p-reinforcement, determine  $r_p(G)$  for some graphs such as paths, cycles and complete t-partite graphs, and establish some upper bounds of  $r_p(G)$ . In particular, we show that the decision problem on  $r_p(G)$  is NP-hard for a general graph G and a fixed integer  $p \geq 2$ .

**Keywords:** domination, *p*-domination, *p*-reinforcement, NP-hard

AMS Subject Classification (2000): 05C69

<sup>\*</sup>The work was supported by NNSF of China (No.10711233) and the Fundamental Research Fund of NPU (No. JC201150)

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### 1 Induction

For notation and graph-theoretical terminology not defined here we follow [21]. Specifically, let G = (V, E) be an undirected graph without loops and multi-edges, where V = V(G) is the vertex-set and E = E(G) is the edge-set, where  $E \neq \emptyset$ .

For  $x \in V$ , the open neighborhood, the closed neighborhood and the degree of x are denoted by  $N_G(x) = \{y \in V : xy \in E\}$ ,  $N_G[x] = N_G(x) \cup \{x\}$  and  $deg_G(x) = |N_G(x)|$ , respectively.  $\delta(G) = \min\{deg_G(x) : x \in V\}$  and  $\Delta(G) = \max\{deg_G(x) : x \in V\}$  are the minimum degree and the maximum degree of G, respectively. For any  $X \subseteq V$ , let  $N_G[X] = \bigcup_{x \in X} N_G[x]$ .

For a subset  $D \subseteq V$ , let  $\overline{D} = V \setminus D$ . The notation  $G^c$  denotes the complement of G, that is,  $G^c$  is the graph with vertex-set V(G) and edge-set  $\{xy : xy \notin E(G) \text{ for any } x, y \in V(G)\}$ . For  $B \subseteq E(G^c)$ , we use G + B to denote the graph with vertex-set V and edge-set  $E \cup B$ . For convenience, we denote  $G + \{xy\}$  by G + xy for an  $xy \in E(G^c)$ .

A nonempty subset  $D \subseteq V$  is called a *dominating set* of G if  $|N_G(x) \cap D| \geq 1$ for each  $x \in \overline{D}$ . The *domination number*  $\gamma(G)$  of G is the minimum cardinality of all dominating sets in G. The domination is a classical concept in graph theory. The early literature on the domination with related topics is, in detail, surveyed in the two books by Haynes, Hedetniemi, and Slater [14, 15].

In 1985, Fink and Jacobson [12] introduced the concept of a generalization domination in a graph. Let p be a positive integer. A subset  $D \subseteq V$  is a p-dominating set of G if  $|N_G(x) \cap D| \ge p$  for each  $x \in \overline{D}$ . The p-domination number  $\gamma_p(G)$  is the minimum cardinality of all p-dominating sets in G. A p-dominating set with cardinality  $\gamma_p(G)$  is called a  $\gamma_p$ -set of G. For  $S, T \subseteq V$ , the set S can p-dominate T in G if  $|N_G(x) \cap S| \ge p$ for every  $x \in T \setminus S$ . Clearly, the 1-dominating set is the classical dominating set, and so  $\gamma_1(G) = \gamma(G)$ . The p-domination is investigated by many authors (see, for example, [1, 2, 4, 5, 11]). Very recently, Chellali *et al.*[4] have given an excellent survey on this topics. The following are two simple observations.

**Observation 1.1** If G is a graph with  $|V(G)| \ge p$ , then  $\gamma_p(G) \ge p$ .

**Observation 1.2** Every p-dominating set of a graph contains all vertices of degree at most p-1.

Clearly, addition of some extra edges to a graph could result in decrease of its domination number. In 1990, Kok and Mynhardt [19] first investigated this problem and proposed the concept of the reinforcement number. The *reinforcement number* r(G) of a graph G is defined as the smallest number of edges whose addition to G results in a graph G' with  $\gamma(G') < \gamma(G)$ . By convention r(G) = 0 if  $\gamma(G) = 1$ .

The reinforcement number has received much research attention (see, for example, [3, 7, 18]), and its many variations have also been well described and studied in graph

theory, including total reinforcement [16, 20], independence reinforcement [22], fractional reinforcement [6, 8] and so on. In particular, Blair *et al.* [3], Hu and Xu [17], independently, showed that the problem determining r(G) for a general graph G is NP-hard.

Motivated by the work of Kok and Mynhardt [19], in this paper, we introduce the *p*-reinforcement number, which is a natural extension of the reinforcement number. The *p*-reinforcement number  $r_p(G)$  of a graph G is the smallest number of edges of  $G^c$ that have to be added to G in order to reduce  $\gamma_p(G)$ , that is

$$r_p(G) = \min\{|B| : B \subseteq E(G^c) \text{ with } \gamma_p(G+B) < \gamma_p(G)\}.$$

It is clear that  $r_1(G) = r(G)$ . By Observation 1.1, we can also make a convention,  $r_p(G) = 0$  if  $\gamma_p(G) \leq p$ . Thus  $r_p(G)$  is well-defined for any graph G and integer  $p \geq 1$ . In this paper, we always assume  $\gamma_p(G) > p$  when we consider the *p*-reinforcement number for a graph G.

The rest of this paper is organized as follows. In Section 2 we present an equivalent parameter for calculating the *p*-reinforcement number of a graph. As its applications, we determine the values of the *p*-reinforcement numbers for special classes of graphs such as paths, cycles and complete *t*-partite graphs in Sections 3, and show that the decision problem on *p*-reinforcement is NP-hard for a general graph and a fixed integer  $p \geq 2$  in Section 4. Finally, we establish some upper bounds for the *p*-reinforcement number of a graph *G* by terms of other parameters of *G* in Section 5.

## 2 Preliminary

Let G be a graph with  $\gamma(G) > 1$  and  $B \subseteq E(G^c)$  with |B| = r(G) such that  $\gamma(G+B) < \gamma(G)$ . Let X be a  $\gamma$ -set of G + B. Then  $|B| \geq |V(G) \setminus N_G[X]|$ . On the other hand, given any set  $X \subseteq V(G)$ , we can always choose a subset  $B \subseteq E(G^c)$  with  $|B| = |V(G) \setminus N_G[X]|$  such that X dominates G + B. It is a simple observation that, to calculate r(G), Kok and Mynhardt [19] proposed the following parameter

$$\eta(G) = \min\{|V(G) \setminus N_G[X]| : X \subseteq V(G), |X| < \gamma(G)\},$$

$$(2.1)$$

and showed  $r(G) = \eta(G)$ . We can refine this technique to deal with the *p*-reinforcement number  $r_p(G)$ .

Let G be a graph with  $\gamma_p(G) > p$ . For any  $X \subseteq V(G)$ , let

$$X^* = \{ x \in \overline{X} : |N_G(x) \cap X| 
$$(2.2)$$$$

Let  $B \subseteq E(G^c)$  with  $|B| = r_p(G)$  such that  $\gamma_p(G + B) < \gamma_p(G)$ , and let X be a  $\gamma_p$ -set of G + B. Then

$$|B| \ge \sum_{x \in X^*} (p - |N_G(x) \cap X|).$$

On the other hand, given any set  $X \subseteq V(G)$  with  $|X| \ge p$ , we can always choose a subset  $B \subseteq E(G^c)$  with

$$|B| = \sum_{x \in X^*} (p - |N_G(x) \cap X|)$$

such that X can p-dominate G + B. Motivated by this observation, we introduce the following notations. For a subset  $X \subseteq V(G)$ ,

$$\eta_p(x, X, G) = \begin{cases} p - |N_G(x) \cap X| & \text{if } x \in X^* \\ 0 & \text{otherwise} \end{cases} \text{ for } x \in V(G), \qquad (2.3)$$

$$\eta_p(S, X, G) = \sum_{x \in S} \eta_p(x, X, G) \text{ for } S \subseteq V(G), \text{ and}$$
 (2.4)

$$\eta_p(G) = \min\{\eta_p(V(G), X, G) : |X| < \gamma_p(G)\}.$$
(2.5)

A subset  $X \subseteq V(G)$  is called an  $\eta_p$ -set of G if  $\eta_p(G) = \eta_p(V(G), X, G)$ . Clearly, for any two subsets  $S', S \subseteq V(G)$  and two subsets  $X', X \subseteq V(G)$ ,

$$\eta_p(S', X, G) \le \eta_p(S, X, G) \quad \text{if } S' \subseteq S, \\ \eta_p(S, X, G) \le \eta_p(S, X', G) \quad \text{if } |X'| \le |X|.$$

Thus, we have the following simple observation.

**Observation 2.1** If X is an  $\eta_p$ -set of a graph G, then  $|X| = \gamma_p(G) - 1$ .

The following result shows that computing  $r_p(G)$  can be referred to computing  $\eta_p(G)$  for a graph G with  $\gamma_p(G) \ge p+1$ .

**Theorem 2.2** For any graph G and positive integer p,  $r_p(G) = \eta_p(G)$  if  $\gamma_p(G) > p$ .

**Proof.** Let X be an  $\eta_p$ -set of G. Then  $|X| = \gamma_p(G) - 1$  by Observation 2.1. Let  $Y = \{y \in V(G) : \eta_p(y, X, G) > 0\}$ . Then  $Y = X^*$  is contained in  $\overline{X}$ , where  $X^*$  is defined in (2.2). Thus,  $\eta_p(G) = \eta_p(X^*, X, G)$ . We construct a new graph G' from G, for each  $y \in X^*$ , by adding  $\eta_p(y, X, G)$  edges of  $G^c$  to G joining y to  $\eta_p(y, X, G)$  vertices in X. Clearly, X is a p-dominating set of G', that is,  $\gamma_p(G') \leq |X|$ . Let B = E(G') - E(G). Then

$$\gamma_p(G) = |X| + 1 > |X| \ge \gamma_p(G') = \gamma_p(G + B),$$

which implies  $r_p(G) \leq |B|$ . It follows that

$$r_p(G) \le |B| = \sum_{y \in X^*} \eta_p(y, X, G) = \eta_p(X^*, X, G) = \eta_p(G).$$
 (2.6)

On the other hand, let B be a subset of  $E(G^c)$  such that  $|B| = r_p(G)$  and  $\gamma_p(G + B) = \gamma_p(G) - 1$ . Let G' = G + B and X' be a  $\gamma_p$ -set of G'. For every  $xy \in B$ , X' cannot p-dominate the graph G' - xy by the minimality of B. This fact means that

only one of x and y is in X'. Without loss of generality, assume  $y \in \overline{X'}$ . Since X' cannot p-dominate y in G' - xy and so in G,  $|N_G(y) \cap X'| < p$ . Let Z be all end-vertices of edges in B and  $Y = \overline{X'} \cap Z$ . Since X' is a  $\gamma_p$ -set of G',  $|N_{G'}(u) \cap X'| \ge p$  for any  $u \in \overline{X'}$ . In other words, any  $u \in \overline{X'}$  with  $|N_G(u) \cap X'| < p$  must be in Y. It follows that

$$\sum_{u \in \overline{X'}} \eta_p(u, X', G) = \sum_{y \in Y} (p - |N_G(y) \cap X'|) = |B|.$$
(2.7)

By (2.7), we immediately have that

$$\eta_p(G) \le \eta_p(V(G), X', G) = \sum_{u \in \overline{X'}} \eta_p(u, X', G) = |B| = r_p(G).$$

Combining this with (2.6), we obtain  $r_p(G) = \eta_p(G)$ , and so the theorem follows.

Note that when p = 1,  $X^*$  defined in (2.2) is  $V(G) \setminus N_G[X]$ . This fact means that  $\eta(G)$  defined in (2.1) is a special case of p = 1 in (2.5), that is,  $\eta_1(G) = \eta(G)$ . Thus, the following corollary holds immediately.

**Corollary 2.1** (Kok and Mynhardt [19])  $r(G) = \eta(G)$  if  $\gamma(G) > 1$ .

Using Observation 1.2 and Theorem 2.2, the following corollary is obvious.

**Corollary 2.2** Let  $p \ge 1$  be an integer and G be a graph with  $\gamma_p(G) > p$ . If  $\Delta(G) < p$ , then

$$r_p(G) = p - \Delta(G).$$

### 3 Some Exact Values

In this section we will use Theorem 2.2 to calculate the p-reinforcement numbers for some classes of graphs.

We first determine the *p*-reinforcement numbers for paths and cycles. Let  $P_n$  and  $C_n$  denote, respectively, a path and a cycle with *n* vertices. When p = 1, Kok and Mynhardt [19] proved that  $r(P_n) = r(C_n) = i$  if  $n = 3k + i \ge 4$ , where  $i \in \{1, 2, 3\}$ . We will give the exact values of  $r_p(P_n)$  and  $r_p(C_n)$  for  $p \ge 2$ . The following observation is simple but useful.

**Observation 3.1** For integer  $p \ge 2$ ,

$$\gamma_p(P_n) = \begin{cases} \lfloor \frac{n}{2} \rfloor + 1 & \text{if } p = 2\\ n & \text{if } p \ge 3 \end{cases} \text{ and } \gamma_p(C_n) = \begin{cases} \lceil \frac{n}{2} \rceil & \text{if } p = 2\\ n & \text{if } p \ge 3. \end{cases}$$

**Theorem 3.2** Let  $p \ge 2$  be an integer. If  $\gamma_p(P_n) > p$  then

$$r_p(P_n) = \begin{cases} 2 & \text{if } p = 2 \text{ and } n \text{ is odd} \\ 1 & \text{if } p = 2 \text{ and } n \text{ is even} \\ p - 2 & \text{if } p \ge 3. \end{cases}$$

**Proof.** Let  $P_n = x_1 x_2 \cdots x_n$  and X be an  $\eta_p$ -set of  $P_n$ . By Theorem 2.2 and  $\gamma_p(P_n) > p$ ,  $r_p(P_n) = \eta_p(P_n) = \eta_p(V(P_n), X, P_n) \ge 1$ . For  $p \ge 3$ , it is easy to see that  $r_p(P_n) = p - 2$  by Corollary 2.2. Assume that p = 2 below.

If n is even, then by Observation 3.1,  $\gamma_2(P_n) - \gamma_2(C_n) = 1$ , which implies that  $r_2(P_n) \leq 1$ . Furthermore,  $r_2(P_n) = 1$ .

If n is odd, then  $\gamma_2(P_n) = \frac{n+1}{2}$  by Observation 3.1, and so  $n \ge 5$  since  $\gamma_2(P_n) > 2$ . Let

$$X' = \bigcup_{i=1}^{\frac{n-1}{2}} \{x_{2i}\}.$$

Clearly,  $|X'| = \frac{n-1}{2} = \gamma_2(P_n) - 1$ . So

$$\eta_2(V(P_n), X, P_n) \le \eta_2(V(P_n), X', P_n) = \eta_2(x_1, X', P_n) + \eta_2(x_n, X', P_n) = 2.$$

Suppose that  $\eta_2(V(P_n), X, P_n) = 1$ . Then X can 2-dominate either  $V(P_n) \setminus \{x_1\}$  or  $V(P_n) \setminus \{x_n\}$ . In both cases, we have

$$|X| \ge \gamma_2(P_{n-1}) = \left\lfloor \frac{n-1}{2} \right\rfloor + 1 = \frac{n-1}{2} + 1,$$

which contradicts with  $|X| = \frac{n-1}{2}$ . Hence  $r_2(P_n) = \eta_2(V(P_n), X, P_n) = 2$ .

**Theorem 3.3** Let  $p \ge 2$  be an integer. If  $\gamma_p(C_n) > p$  then

$$r_p(C_n) = \begin{cases} 2 & \text{if } p = 2 \text{ and } n \text{ is odd} \\ 4 & \text{if } p = 2 \text{ and } n \text{ is even} \\ p - 2 & \text{if } p \ge 3. \end{cases}$$

**Proof.** Let  $C_n = x_1 x_2 \cdots x_n x_1$ . If  $p \ge 3$  then the result holds obviously by Corollary 2.2. In the following, we only need to calculate the values of  $r_p(C_n)$  for p = 2. Let X be an  $\eta_2$ -set of  $C_n$ . Then  $r_2(C_n) = \eta_2(C_n) = \eta_2(V(C_n), X, C_n)$  by Theorem 2.2. Note that  $n \ge 5$  since  $\gamma_2(C_n) = \lceil \frac{n}{2} \rceil > 2$ .

If n is odd, then let

$$X' = \bigcup_{i=1}^{\frac{n-1}{2}} \{x_{2i-1}\}.$$

Clearly,  $|X'| = \frac{n-1}{2} = \gamma_2(C_n) - 1$  by Observation 3.1, and  $\eta_2(V(C_n), X', C_n) = \eta_2(x_{n-1}, X', C_n) + \eta_2(x_n, X', C_n) = 2$ . So

$$r_2(C_n) = \eta_2(V(C_n), X, C_n) \le \eta_2(V(C_n), X', C_n) = 2.$$

Since X is not a 2-dominating set of  $C_n$ , there must be two adjacent vertices, denoted by  $x_i$  and  $x_{i+1}$ , of  $C_n$  not in X. This fact means that  $\eta_2(x_i, X, C_n) \ge 1$  and  $\eta_2(x_{i+1}, X, C_n) \ge 1$ . So

$$r_2(C_n) = \eta_2(V(C_n), X, P_n) \ge \eta_2(x_i, X, C_n) + \eta_2(x_{i+1}, X, C_n) \ge 2.$$

Hence  $r_2(C_n) = 2$ .

If n is even, then  $n \ge 6$ . Deleting X and all vertices 2-dominated by X from  $C_n$ , we can obtain a result graph, denoted by H, each of whose components is a path with length at least 2. Denote all components of H by  $H_1, \dots, H_h$ , where  $h \ge 1$ . In the case that h = 1 and the length of  $H_1$  is equal to one, X can 2-dominate a subgraph of  $C_n$  that is isomorphic to  $P_{n-2}$ . By Observation 3.1,

$$|X| \ge \gamma_2(P_{n-2}) = \lfloor \frac{n-2}{2} \rfloor + 1 = \frac{n}{2},$$

which contradicts that  $|X| = \gamma_2(C_n) - 1 = \lceil \frac{n}{2} \rceil - 1 = \frac{n}{2} - 1$ . In other cases, we can find that

$$r_2(C_n) = \eta_2(V(C_n), X, C_n) \ge 4.$$

Let

$$X'' = \bigcup_{i=1}^{\frac{n}{2}-1} \{x_{2i-1}\}.$$

It is easy to check that  $|X''| = \frac{n}{2} - 1 = \gamma_2(C_n) - 1$  and  $\eta_2(V(C_n), X'', C_n) = 4$ . So

$$r_2(C_n) = \eta_2(V(C_n), X, C_n) \le \eta_2(V(C_n), X'', C_n) = 4.$$

Hence  $r_2(C_n) = 4$  and so the theorem is true.

Next we consider the *p*-reinforcement number for a complete *t*-partite graph  $K_{n_1,\dots,n_t}$ . To state our results, we need some symbols. For any subset  $X = \{n_{i_1,\dots,n_{i_r}}\}$  of  $\{n_1,\dots,n_t\}$ , define

$$|X| = r$$
 and  $f(X) = \sum_{j=1}^{r} n_{i_j}$ .

For convenience, let |X| = 0 and f(X) = 0 if  $X = \emptyset$ . let

 $\mathscr{X} = \{X : X \text{ is a subset of } \{n_1, \cdots, n_t\} \text{ with } f(X) \ge \gamma_p(G)\}$ 

and, for every  $X \in \mathscr{X}$ , define

 $f^*(X) = \max\{f(Y) : Y \text{ is a subset of } X \text{ with } |Y| = |X| - 1 \text{ and } f(Y) < p\}.$ 

**Theorem 3.4** For any integer  $p \ge 1$  and a complete t-partite graph  $G = K_{n_1,\dots,n_t}$  with  $t \ge 2$  and  $\gamma_p(G) > p$ ,

$$r_p(G) = \min\{(p - f^*(X))(f(X) - \gamma_p(G) + 1) : X \in \mathscr{X}\}.$$

**Proof.** Let  $N = \{n_1, \dots, n_t\}$  and  $V(G) = V_1 \cup \dots \cup V_t$  be the vertex-set of G such that  $|V_i| = n_i$  for each  $i = 1, \dots, t$ . Let

$$m = \min\{(p - f^*(X))(f(X) - \gamma_p(G) + 1) : X \in \mathscr{X}\}.$$

We first prove that  $r_p(G) \leq m$ . Let  $X \subseteq \mathscr{X}$  (without loss of generality, assume  $X = \{n_1, \dots, n_k, n_{k+1}\}$  for some  $0 \leq k \leq t-1$ ) such that

$$f^*(X) = n_1 + \dots + n_k$$
 and  $(p - f^*(X))(f(X) - \gamma_p(G) + 1) = m$ .

By  $X \subseteq \mathscr{X}$ , we know that  $n_{k+1} = f(X) - f^*(X) \ge \gamma_p(G) - f^*(X)$ . So we can pick a vertex-subset  $V'_{k+1}$  from  $V_{k+1}$  such that  $|V'_{k+1}| = \gamma_p(G) - f^*(X) - 1$ . Let

$$D = V_1 \cup \cdots \cup V_k \cup V'_{k+1}.$$

Clearly,  $|D| = \gamma_p(G) - 1$ . Since  $\gamma_p(G) > p$ ,  $|D| \ge p$  and so D can p-dominate  $\cup_{i=k+2}^t V_i$ . Hence by the definition of  $\eta_p(V(G), D, G)$ ,

$$\eta_p(V(G), D, G) = \eta_p(V(G) \setminus D, D, G)$$

$$= \sum_{v \in V_{k+1} \setminus V'_{k+1}} \eta_p(v, D, G) + \sum_{i=k+2}^t \eta_p(V_i, D, G)$$

$$= |V_{k+1} \setminus V'_{k+1}| (p - f^*(X)) + 0$$

$$= (p - f^*(X))[n_{k+1} - (\gamma_p(G) - f^*(X) - 1)]$$

$$= (p - f^*(X))(f(X) - \gamma_p(G) + 1)$$

$$= m.$$

By Theorem 2.2, we have  $r_p(G) = \eta_p(G) \le \eta_p(V(G), D, G) = m$ .

On the other hand, we will show that  $r_p(G) \ge m$ . For any subset M of N, we use I(M) to denote the subindex-sets of all elements in M, that is,

$$I(M) = \{i : n_i \in M\}.$$

Let S be an  $\eta_p$ -set of G and let

$$Y = \{n_i : |V_i \cap S| = |V_i| \text{ for } 1 \le i \le t\}, \text{ and} \\ A = \{n_i : 0 < |V_i \cap S| < |V_i| \text{ for } 1 \le i \le t\}.$$

Thus

$$f(Y \cup A) = f(Y) + f(A) = \sum_{i \in I(Y)} |V_i| + \sum_{i \in I(A)} |V_i| \ge |S| = \gamma_p(G) - 1$$
(3.1)

by Observation 2.1. Since  $\bigcup_{i \in I(Y)} V_i (\subseteq S)$  cannot p-dominate G,

$$f(Y) = \sum_{i \in I(Y)} n_i = |\cup_{i \in I(Y)} V_i| < p.$$
(3.2)

Hence, by (3.1) and  $\gamma_p(G) > p$ ,

$$f(A) \ge \gamma_p(G) - 1 - f(Y) > \gamma_p(G) - p - 1 \ge 0,$$

which implies that  $|A| \ge 1$ .

**Claim.** |A| = 1.

**Proof of Claim.** Suppose that  $|A| \ge 2$ . Then we can choose *i* and *j* from I(A) such that  $i \ne j$ . By the definition of *A*, we have  $0 < |V_i \cap S| < |V_i|$  and  $0 < |V_j \cap S| < |V_j|$ . Therefore, we can pick two vertices *x* and *y* from  $V_i \cap S$  and  $V_j \setminus S$ , respectively. Let

$$S' = (S \setminus \{x\}) \cup \{y\}.$$

Obviously,  $|S'| = |S| = \gamma_p(G) - 1$ ,  $|V_i \cap S'| = |V_i \cap S| - 1$  and  $|V_j \cap S'| = |V_j \cap S| + 1$ .

Note that G is a complete t-partite graph. For any  $v \in V(G)$ , we can easily find the value of  $\eta_p(v, S', G) - \eta_p(v, S, G)$  by the definitions of  $\eta_p(v, S', G)$  and  $\eta_p(v, S, G)$ as follows:

$$\eta_p(v, S', G) - \eta_p(v, S, G) = \begin{cases} (p - |S| + |V_i \cap S| - 1) - 0 & \text{if } v = x \\ -1 & \text{if } v \in V_i \setminus S \\ 0 - (p - |S| + |V_j \cap S|) & \text{if } v = y \\ 1 & \text{if } v \in (V_j \setminus S) \setminus \{y\} \\ 0 & \text{otherwise.} \end{cases}$$

Since S is an  $\eta_p$ -set of G and |S'| = |S|, we have

$$\begin{array}{ll} 0 &\leq & \eta_p(V(G), S', G) - \eta_p(V(G), S, G) \\ &= & \sum_{v \in V(G)} (\eta_p(v, S', G) - \eta_p(v, S, G)) \\ &= & (p - |S| + |V_i \cap S| - 1) - |V_i \setminus S| - (p - |S| + |V_j \cap S|) + |(V_j \setminus S) \setminus \{y\}| \\ &= & (|V_i \cap S| - |V_i \setminus S|) - (|V_j \cap S| - |V_j \setminus S|) - 2. \end{array}$$

This means that

$$(|V_i \cap S| - |V_i \setminus S|) \ge (|V_j \cap S| - |V_j \setminus S|) + 2$$

However, by the symmetry of  $V_i$  and  $V_j$ , we can also obtain

$$(|V_j \cap S| - |V_j \setminus S|) \ge (|V_i \cap S| - |V_i \setminus S|) + 2$$

by applying the similar discussion. This is a contradiction, and so the claim holds.  $\Box$ 

By **Claim**, we can assume that  $I(A) = \{h\}$ . From the definitions of Y and A, we have  $|Y \cup A| = |Y| + 1$  and

$$f(Y \cup A) = \sum_{i \in I(Y)} |V_i| + |V_h| \ge \sum_{i \in I(Y)} |V_i| + (|V_h \cap S| + 1) = |S| + 1 = \gamma_p(G).$$

It follows that  $Y \cup A \in \mathscr{X}$ . Thus, by (3.2) and the definition of  $f^*(Y \cup A)$ , we have  $f(Y) \leq f^*(Y \cup A)$ . Since  $\gamma_p(G) > p$ ,  $|S| = \gamma_p(G) - 1 \geq p$ , and so S p-dominates

 $V(G) \setminus (\bigcup_{i \in I(Y \cup A)} V_i)$ . Therefore, by Theorem 2.2,

$$\begin{aligned} r_p(G) &= \eta_p(G) = \eta_p(V(G), S, G) &= \eta_p(V(G) \setminus S, S, G) \\ &= \sum_{v \in V_h \setminus S} \eta_p(v, S, G) \\ &= (p - f(Y)) |V_h \setminus S| \\ &= (p - f(Y)) [|V_h| - (|S| - f(Y))] \\ &= (p - f(Y)) (f(Y \cup A) - \gamma_p(G) + 1) \\ &\geq (p - f^*(Y \cup A)) (f(Y \cup A) - \gamma_p(G) + 1) \\ &\geq m. \end{aligned}$$

This completes the proof of the theorem.

For example, let  $G = K_{2,2,10,17}$  and p = 11. Then  $\gamma_{11}(G) = 12$ , and so

$$\mathscr{X} = \{\{17\}, \{2, 10\}, \{2, 17\}, \{10, 17\}, \{2, 2, 10\}, \{2, 2, 17\}, \{2, 10, 17\}, \{2, 2, 10, 17\}\}.$$

By Theorem 3.4, for any  $X \in \mathscr{X}$ , we have that

$$f^*(X) = \begin{cases} 0 & \text{if } X = \{17\}, \{2, 10, 17\} \text{ or } \{2, 2, 10, 17\}; \\ 2 & \text{if } X = \{2, 17\}; \\ 4 & \text{if } X = \{2, 2, 10\} \text{ or } \{2, 2, 17\}; \\ 10 & \text{if } X = \{2, 10\} \text{ or } \{10, 17\}. \end{cases}$$

Hence

$$r_{11}(G) = \min\{(11 - f^*(X))(f(X) - \gamma_{11}(G) + 1) : X \in \mathscr{X}\}\$$
  
=  $\min\{(11 - f^*(X))(f(X) - 11) : X \in \mathscr{X}\}\$   
=  $(11 - f^*(\{2, 10\}))(f(\{2, 10\}) - 11)\$   
= 1.

### 4 Complexity

Blair et al. [3], Hu and Xu [17], independently, showed that the 1-reinforcement problem is NP-hard. Thus, for any positive integer p, the p-reinforcement problem is also NP-hard since the 1-reinforcement is a sub-problem of the p-reinforcement problem.

For each fixed p, p-dominating set is polynomial-time computable (see Downey and Fellows [9, 10] for definitions and discussion). However, the p-reinforcement number problem is hard even for specific values of the parameters. In this section, we will consider the following decision problem.

#### *p*-Reinforcement

Instance: A graph  $G, p \ (\geq 2)$  is a fixed integer.

Question: Is  $r_p(G) \leq 1$ ?

We will prove that *p*-Reinforcement  $(p \ge 2)$  is also NP-hard by describing a polynomial transformation from the following NP-hard problem (see [13]).

#### 3-Satisfiability (3SAT)

Instance: A set  $U = \{u_1, \dots, u_n\}$  of variables and a collection  $\mathscr{C} = \{C_1, \dots, C_m\}$ of clauses over U such that  $|C_i| = 3$  for  $i = 1, 2, \dots, m$ .

Furthermore, every literal is used in at least one clause.

Question: Is there a satisfying truth assignment for C?

**Theorem 4.1** For a fixed integer  $p \ge 2$ , p-Reinforcement is NP-hard.

**Proof.** Let  $U = \{u_1, \ldots, u_n\}$  and  $\mathscr{C} = \{C_1, \ldots, C_m\}$  be an arbitrary instance I of **3SAT**. We will show the NP-hardness of *p*-Reinforcement by reducing **3SAT** to it in polynomial time. To this aim, we construct a graph G as follows:

- **a.** For each variable  $u_i \in U$ , associate a graph  $H_i$ , where  $H_i$  can be obtained from a complete graph  $K_{2p+2}$  with vertex-set  $\{u_i, \overline{u}_i\} \cup (\bigcup_{j=1}^p \{v_{i_j}, \overline{v}_{i_j}\})$  by deleting the edge-subset  $\bigcup_{i=1}^{p-1} \{u_i \overline{v}_{i_j}, \overline{u}_i v_{i_j}\}$ ;
- **b.** For each clause  $C_j \in \mathscr{C}$ , create a single vertex  $c_j$  and join  $c_j$  to the vertex  $u_i$  (resp.  $\overline{u}_i$ ) in  $H_i$  if and only if the literal  $u_i$  (resp.  $\overline{u}_i$ ) appears in clause  $C_j$  for any  $i \in \{1, \ldots, n\}$ ;
- **c.** Add a complete graph  $T \cong K_p$  and join all of its vertices to each  $c_j$ .

For convenience, let  $X_i = \bigcup_{j=1}^p \{v_{i_j}\}$  and  $\overline{X}_i = \bigcup_{j=1}^p \{\overline{v}_{i_j}\}$ . Then  $V(H_i) = \{u_i, \overline{u}_i\} \cup X_i \cup \overline{X}_i$ . Use  $H_0$  to denote the induced subgraph by  $\{c_1, \cdots, c_m\} \cup V(T)$ .

It is clear that the construction of G can be accomplished in polynomial time. To complete the proof of the theorem, we only need to prove that  $\mathscr{C}$  is satisfiable if and only if  $r_p(G) = 1$ . We first prove the following two claims.

**Claim 1.** Let D be a  $\gamma_p$ -set of G. Then |D| = p(n+1), moreover,  $|V(H_i) \cap D| = p$ and  $|\{u_i, \overline{u}_i\} \cap D| \leq 1$  for each  $i \in \{1, 2, \ldots, n\}$ .

**Proof of Claim 1.** Suppose there is some  $i \in \{1, 2, \dots, n\}$  such that  $|V(H_i) \cap D| < p$ . Then there must be a vertex, say x, of  $V(H_i) \setminus D$  such that  $N_G(x) \subseteq V(H_i)$ . And so  $|N_G(x) \cap D| \leq |V(H_i) \cap D| < p$ , which contradicts that D is a  $\gamma_p$ -set of G. Thus  $|V(H_i) \cap D| \geq p$  for each  $i \in \{0, 1, \dots, n\}$ , and so

$$\gamma_p(G) = |D| = \sum_{i=0}^n |V(H_i) \cap D| \ge p(n+1).$$
 (4.1)

On the other hand, let

$$D' = \bigcup_{i=1}^{n} [(X_i - \{v_{i_p}\}) \cup \{\overline{u}_i\}] \cup V(T).$$

Clearly, |D'| = p(n+1) and D' is a p-dominating set of G. Hence by (4.1),

$$p(n+1) \le \sum_{i=0}^{n} |V(H_i) \cap D| = \gamma_p(G) \le |D'| = p(n+1),$$

which implies that  $\gamma_p(G) = p(n+1)$  and  $|V(H_i) \cap D| = p$  for each  $0 \leq i \leq n$ . Furthermore, if  $|\{u_i, \overline{u}_i\} \cap D| = 2$  then  $|(X_i \cup \overline{X}_i) \cap D| = p - 2$ . So we can choose a vertex from  $X_i \cup \overline{X}_i$  that is not *p*-dominated by *D*. This is impossible since *D* is a  $\gamma_p$ -set of *G*, and so  $|\{u_i, \overline{u}_i\} \cap D| \leq 1$ . The claim holds.  $\Box$ 

**Claim 2.** If there is an edge  $e = xy \in G^c$  such that  $\gamma_p(G+e) < \gamma_p(G)$ , then any  $\gamma_p$ -set  $D_e$  of G + e satisfies the following properties.

- (i)  $|V(H_i) \cap D_e| = p$  and  $|\{u_i, \overline{u}_i\} \cap D_e| \leq 1$  for each  $i \in \{1, \cdots, n\}$ ;
- (ii)  $\{c_1, \cdots, c_m\} \cap D_e = \emptyset$ , and so  $|V(T) \cap D_e| = p 1$ ;
- (iii) One of x and y belongs to  $V(T) \setminus D_e$  and the other belongs to  $H \cap D_e$ , where  $H = \bigcup_{i=1}^n V(H_i)$ .

**Proof of Claim 2.** Because  $D_e$  is a  $\gamma_p$ -set of G + e and  $\gamma_p(G + e) < \gamma_p(G)$ , one of x and y is not in  $D_e$  but the other is in  $D_e$ . Without loss of generality, say  $x \notin D_e$  and  $y \in D_e$ . It is clear that  $|N_G(x) \cap D_e| = p - 1$ . Since vertex x is the unique vertex not be p-dominated by  $D_e$ , we have

$$\eta_p(V(G), D_e, G) = \eta_p(x, D_e, G) = p - (p - 1) = 1.$$
(4.2)

Let

$$D = D_e \cup \{x\}.$$

Then D is a p-dominating set of G and  $|D| = |D_e| + 1 = \gamma_p(G + e) + 1 \le \gamma_p(G)$ . That is, D is a  $\gamma_p$ -set of G. By Claim 1,

$$|V(H_i) \cap D| = p \text{ for each } i = 0, 1, \cdots, n,$$

$$(4.3)$$

and  $|\{u_i, \overline{u}_i\} \cap D_e| \le |\{u_i, \overline{u}_i\} \cap D| \le 1$  for  $1 \le i \le n$ .

Suppose that there exists some  $i \in \{1, \dots, n\}$  such that  $|V(H_i) \cap D_e| \neq p$ . Then by (4.3),  $x \in V(H_i)$  and  $|V(H_i) \cap D_e| = p-1$ . Thus every vertex in  $(X_i \cup \overline{X_i}) \setminus (D_e \cup \{x\})$  is dominated by at most p-1 vertices of  $D_e$ . Hence by  $|X_i \cup \overline{X_i}| = 2p$ ,

$$\eta_p(V(G), D_e, G) \ge \eta_p(X_i \cup \overline{X}_i, D_e, G) \ge |(X_i \cup \overline{X}_i) \setminus D_e| - 1 \ge 2p - (p - 1) - 1 > 1,$$

which contradicts with (4.2). Hence (i) holds.

Suppose that there is some  $j \in \{1, \dots, m\}$  such that  $c_j \in D_e$ . By (i) and (4.3),  $x \in V(H_0)$  and so  $|V(H_0) \cap D_e| = |V(H_0) \cap D| - 1 = p - 1$ . Hence  $|V(T) \cap D_e| \le p - 2$  by  $V(H_0) = \{c_1, \dots, c_m\} \cup V(T)$ . Since each vertex of  $T \ (\cong K_p)$  has exact p - 1 neighbors in  $D_e$ ,

$$\eta_p(V(G), D_e, G) \ge \eta_p(V(T), D_e, G) = |V(T) \setminus D_e| = p - |V(T) \cap D_e| \ge 2.$$

This contradicts with (4.2). Thus  $\{c_1, \dots, c_m\} \cap D_e = \emptyset$ , and so  $|V(T) \cap D_e| = |V(H_0) \cap D_e| = p - 1$ . Hence (*ii*) holds.

By (*ii*), *T* has a unique vertex, say *z*, not in  $D_e$ . From  $|N_G(z) \cap D_e| = |V(H_0) \cap D_e| = p - 1$ , the vertex *z* is not *p*-dominated by  $D_e$ . However, *x* is the unique vertex not be *p*-dominated by  $D_e$  in *G* by (4.2). Thus z = x, and so  $x = z \in V(T) \setminus D_e$ . By the construction of *G* and  $xy \in G^c$ , it is clear that  $y \in (\bigcup_{i=1}^n V(H_i)) \cap D_e$ . Hence (*iii*) holds.  $\Box$ 

We now show that  $\mathscr{C}$  is satisfiable if and only if  $r_p(G) = 1$ .

If  $\mathscr{C}$  is satisfiable, then  $\mathscr{C}$  has a satisfying truth assignment  $t : U \to \{T, F\}$ . According to this satisfying assignment, we can choose a subset S from V(G) as follows:

$$S = S_0 \cup S_1 \cup \cdots \cup S_n,$$

where  $S_0$  consists of p-1 vertices of T and

$$S_i = \begin{cases} u_i \cup (\overline{X}_i - \{\overline{v}_{i_p}\}) & \text{if } t(u_i) = T \\ \overline{u}_i \cup (X_i - \{v_{i_p}\}) & \text{if } t(u_i) = F \end{cases} \text{ for each } i \in \{1, \cdots, n\}.$$

It can be verified easily that  $|S| = p(n+1) - 1 = \gamma_p(G) - 1$  and  $\bigcup_{i=1}^n V(H_i)$  can be p-dominated by S. Since t is a satisfying true assignment for  $\mathscr{C}$ , each clause  $C_j \in \mathscr{C}$  contains at least one true literal. That is, the corresponding vertex  $c_j$  has at least one neighbor in  $\{u_1, \bar{u}_1 \cdots, u_n, \bar{u}_n\} \cap S$  by the definitions of G and S, and so every  $c_j \in \{c_1, \cdots, c_m\}$  has at least p neighbors in S since  $S_0 \subseteq N_G(c_j)$ . Note that the unique vertex in  $V(T) \setminus S_0$  has exact p-1 neighbors in S. By Theorem 2.2 and  $|S| = \gamma_p(G) - 1$ ,

$$r_p(G) = \eta_p(G) \le \eta_p(V(G), S, G) = \eta_p(V(T) \setminus S_0, S, G) = p - (p - 1) = 1.$$

Furthermore, we have  $r_p(G) = 1$  since  $\gamma_p(G) > p$  by Claim 1.

Conversely, assume  $r_p(G) = 1$ . That is, there exists an edge e = xy in  $G^c$  such that  $\gamma_p(G+e) < \gamma_p(G)$ . Let  $D_e$  be a  $\gamma_p$ -set of G + e. Define  $t: U \to \{T, F\}$  by

$$t(u_i) = \begin{cases} T & \text{if vertex } u_i \in D_e \\ F & \text{if vertex } u_i \notin D_e \end{cases} \text{ for } i = 1, \cdots, n.$$

$$(4.4)$$

We will show that t is a satisfying truth assignment for  $\mathscr{C}$ . Let  $C_j$  be an arbitrary clause in  $\mathscr{C}$ . By (ii) and (iii) of Claim 2, the corresponding vertex  $c_j$  is not in  $D_e$  and  $|N_G(c_j) \cap D_e| \ge p$  since  $c_j \notin \{x, y\}$ . Then there must be some  $i \in \{1, \dots, n\}$  such that

$$|\{u_i, \overline{u}_i\} \cap N_G(c_j) \cap D_e| = 1, \tag{4.5}$$

since T contains exact p-1 vertices of  $D_e$  by (i) and (ii) of Claim 2. If  $u_i \in N_G(c_j) \cap D_e$ , then  $u_i \in C_j$  and  $t(u_i) = T$  by the construction of G and (4.4). If  $\overline{u_i} \in N_G(c_j) \cap D_e$ , then the literal  $\overline{u_i}$  belongs to  $C_j$  by the construction of G. Note that  $u_i \notin D_e$  from  $\overline{u_i} \in D_e$  and (i) of Claim 2. This means that  $t(u_i) = F$  by (4.4). Hence  $t(\overline{u_i}) = T$ . The arbitrariness of  $C_j$  with  $1 \leq j \leq m$  shows that all the clauses in  $\mathscr{C}$  is satisfied by t. That is,  $\mathscr{C}$  is satisfiable.

The theorem follows.

## 5 Upper Bounds

For a graph G and p = 1, Kok and Mynhardt [19] provided an upper bound for r(G) in terms of the smallest private neighborhood of a vertex in some  $\gamma$ -set of G. Let  $X \subseteq V(G)$  and  $x \in X$ . The private neighborhood of x with respect to X is defined as the set

$$PN(x, X, G) = N_G[x] \setminus N_G[X \setminus \{x\}].$$
(5.1)

Set

$$\mu(X,G) = \min\{|PN(x,X,G)| : x \in X\}$$

and

$$\mu(G) = \min\{\mu(X, G) : X \text{ is a } \gamma \text{-set of } G\}.$$
(5.2)

Using this parameter, Kok and Mynhardt [19] showed that  $r(G) \leq \mu(G)$  if  $\gamma(G) \geq 2$  with equality if  $\gamma(G) = 1$ . We generalize this result to any positive integer p.

In order to state our results, we need some notations. Let  $X \subseteq V(G)$  and  $x \in X$ . A vertex  $y \in \overline{X}$  is called a *p*-private neighbor of x with respect to X if  $xy \in E(G)$  and  $|N_G(y) \cap X| = p$ . The *p*-private neighborhood of x with respect to X is defined as

$$PN_p(x, X, G) = \{y : y \text{ is a } p \text{-private neighbor of } x \text{ with respect to } X\}.$$
(5.3)

Let

$$\mu_p(x, X, G) = |PN_p(x, X, G)| + \max\{0, p - |N_G(x) \cap X|\},$$
(5.4)

$$\mu_p(X,G) = \min\{\mu_p(x,X,G) : x \in X\}, \text{ and}$$
 (5.5)

$$\mu_p(G) = \min\{\mu_p(X, G) : X \text{ is a } \gamma_p \text{-set of } G\}.$$
(5.6)

**Theorem 5.1** For any graph G and positive integer p,

$$r_p(G) \le \mu_p(G)$$

with equality if  $r_p(G) = 1$ .

**Proof.** If  $\gamma_p(G) \leq p$ , then  $r_p(G) = 0 \leq \mu_p(G)$  by our convention. Assume that  $\gamma_p(G) \geq p+1$  below. Let X be a  $\gamma_p$ -set of G and  $x \in X$  such that

$$\mu_p(G) = \mu_p(X, G) = \mu_p(x, X, G).$$

Since  $|X| = \gamma_p(G) \ge p + 1$ , we can choose a vertex, say  $u_y$ , from  $X \setminus N_G(y)$  for each  $y \in PN_p(x, X, G)$ , and a subset X' with  $|X'| = \max\{0, p - |N_G(x) \cap X|\}$  from  $X \setminus N_G[x]$ . Let

$$G' = G + \{yu_y : y \in PN_p(x, X, G)\} + \{xv : v \in X'\}.$$

Obviously,  $X \setminus \{x\}$  is a p-dominating set of G', which implies that

$$r_p(G) \le |PN_p(x, X, G)| + |X'| = \mu_p(x, X, G) = \mu_p(G).$$

Assume  $r_p(G) = 1$ . Then  $\gamma_p(G) \ge p + 1$  and there exists an edge  $xy \in E(G^c)$  such that  $\gamma_p(G + xy) = \gamma_p(G) - 1$ . Let G' = G + xy and X be a  $\gamma_p$ -set of G'. Without loss of generality, assume that  $x \in X$  and  $y \in \overline{X}$ . Clearly, y is a p-private neighbor of x with respect to X in G and  $X \cup \{y\}$  is a  $\gamma_p$ -set of G, which implies

$$PN_p(y, X \cup \{y\}, G) = \emptyset$$
 and  $p - |N_G(y) \cap (X \cup \{y\})| = 1$ 

that is,  $\mu_p(y, X \cup \{y\}, G) = 1$ . It follows that

$$r_p(G) \le \mu_p(G) \le \mu_p(X \cup \{y\}, G) \le \mu_p(y, X \cup \{y\}, G) = 1$$

Thus,  $r_p(G) = \mu_p(G) = 1$ . The theorem follows.

Note that  $|PN_p(x, X, G)| \leq deg_G(x)$  for any  $X \subseteq V(G)$  and  $x \in X$ . By Theorem 5.1, we obtain the following corollary immediately.

**Corollary 5.1** For any graph G with maximum degree  $\Delta(G)$  and positive integer p,  $r_p(G) \leq \Delta(G) + p$ .

**Corollary 5.2** Let p be a positive integer and G be a graph with minimum degree  $\delta(G)$ . If  $\delta(G) < p$ , then  $r_p(G) \leq \delta(G) + p$ .

**Proof.** Let X be a  $\gamma_p$ -set of G and  $x \in V(G)$  with degree  $\delta(G)$ . Since  $deg_G(x) = \delta(G) < p, x \in X$  by Observation 1.2. Note that  $|PN_p(x, X, G)| \le deg_G(x) = \delta(G)$  and  $p - |N_G(x) \cap X| \le p$ . By Theorem 5.1,

$$r_p(G) \leq \mu_p(G)$$
  

$$\leq \mu_p(x, X, G)$$
  

$$= |PN_p(x, X, G)| + \max\{0, p - |N_G(x) \cap X|\}$$
  

$$\leq \delta(G) + p.$$

The corollary follows.

Consider p = 1. Let  $X \subseteq V(G)$  and  $x \in X$ . If x is not an isolated vertex of the induced subgraph G[X], then PN(x, X, G) defined in (5.1) does not contain x and  $\max\{0, 1 - |N_G(x) \cap X|\} = 0$  in (5.4). Otherwise, PN(x, X, G) contains x and  $\max\{0, 1 - |N_G(x) \cap X|\} = 1$ . Notice that  $PN_1(x, X, G)$  defined in (5.3) does not contain x. Hence, by (5.5),

$$\mu_1(x, X, G) = PN_1(x, X, G) + \max\{0, 1 - |N_G(x) \cap X|\} = |PN(x, X, G)|.$$

This fact means that  $\mu(G)$  defined in (5.2) is a special case of p = 1 in (5.6), that is,  $\mu_1(G) = \mu(G)$ . Thus, by Theorem 5.1, the following corollary holds immediately.

**Corollary 5.3** (Kok and Mynhardt [19]) For any graph G with  $\gamma(G) \ge 2$ ,  $r(G) \le \mu(G)$ , with equality if r(G) = 1.

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