# A Class of Linear Codes with a Few Weights

**Can Xiang<sup>1</sup>** · **Chunming Tang<sup>2</sup>** · **Keqin Feng<sup>3</sup>** 

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Abstract Linear codes have been an interesting subject of study for many years, as linear codes with few weights have applications in secrete sharing, authentication codes, association schemes, and strongly regular graphs. In this paper, a class of linear codes with a few weights over the finite field GF(p) are presented and their weight distributions are also determined, where p is an odd prime. Some of the linear codes obtained are optimal in the sense that they meet certain bounds on linear codes.

**Keywords** Linear codes · weight distribution · weight enumerator · secret sharing schemes

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## 1 Introduction

Throughout this paper, let *p* be an odd prime and let  $q = p^m$  for some positive integer *m*. An [n, k, d] code *C* over GF(*p*) is a *k*-dimensional subspace of GF(*p*)<sup>*n*</sup> with minimum (Hamming) distance *d*. Let  $A_i$  denote the number of codewords with Hamming weight *i* in a code *C* of length *n*. The weight enumerator of *C* is defined by  $1 + A_1 z + A_2 z^2 + \cdots + A_n z^n$ . The weight distribution  $(1, A_1, \ldots, A_n)$  is an important research topic in coding theory, as it contains crucial information as to estimate the error correcting capability and the probability of error detection and correction with respect to some algorithms. A code *C* is said to be a *t*-weight code if the number of nonzero  $A_i$  in the sequence  $(A_1, A_2, \cdots, A_n)$  is equal to *t*.

C. Tang

K. Feng

Department of Mathematical Sciences, Tsinghua University, Beijing, 100084, China E-mail: kfeng@math.tsinghua.edu.cn

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C. Xiang

College of Mathematics and Information Science, Guangzhou University, Guangzhou 510006, China E-mail: cxiangcxiang@hotmail.com

School of Mathematics and Information, China West Normal University, Sichuan Nanchong, 637002, China E-mail: tangchunmingmath@163.com

Let  $D = \{d_1, d_2, ..., d_n\} \subseteq GF(q)$ . Let Tr denote the trace function from GF(q) onto GF(p) throughout this paper. We define a linear code of length *n* over GF(p) by

$$\mathcal{C}_D = \{(\operatorname{Tr}(xd_1), \operatorname{Tr}(xd_2), \dots, \operatorname{Tr}(xd_n)) : x \in \operatorname{GF}(q)\},\tag{1}$$

and call *D* the *defining set* of this code  $C_D$ . By definition, the dimension of the code  $C_D$  is at most *m*.

This construction is generic in the sense that many classes of known codes could be produced by properly selecting the defining set  $D \subseteq GF(q)$ . If the defining set D is well chosen, some optimal linear codes with few weights can be obtained. This construction technique was employed in [6], [13], [14], [7], [9], [10] and [25] for obtaining linear codes with a few weights. For more details, we refer interested readers to [8,21,29,23,24,22] and the references therein.

The purpose of this paper is to construct a class of linear codes over GF(p) with a few nonzero weights using this generic construction method, and determine their weight distributions. Some of the linear codes obtained in this paper are optimal in the sense that they meet some bounds on linear codes. The linear codes with a few weights presented in this paper have applications also in secret sharing [1,4,27], authentication codes [11], combinatorial designs and graph theory [2,3], and association schemes [2], in addition to their applications in consumer electronics, communication and data storage systems.

The remainder of this paper is organized as follows. Section 2 introduces some basic notations and results of group characters, Gauss sums, exponential sums and cyclotomic fields which will be needed in subsequent sections. Section 3 presents a class of linear codes with a few weights and the proofs of their parameters are given in Section 4. Section 5 summarizes this paper.

## 2 Preliminaries

In this section, we state some notations and basic facts on group characters, Gauss sums, exponential sums and cyclotomic fields. These results will be used later in this paper.

### 2.1 Some notations fixed throughout this paper

For convenience, we adopt the following notations unless otherwise stated in this paper.

- $p^* = (-1)^{(p-1)/2}p$ .
- $\zeta_p = e^{\frac{2\pi\sqrt{-1}}{p}}$ , a primitive *p*-th root of unity.
- $L(x) = x^{p^2} + x$ .
- $\operatorname{Im}(L) = \{L(x) : x \in \operatorname{GF}(q)\}.$
- $x_b \in GF(q)$  denotes a solution of the equation  $L(x) = -b^p$  if  $b \in Im(L)$ .
- $m_p = m \mod p \in \{0, 1, ..., p-1\}$ , the least non-negative residue modulo m.
- SQ and NSQ denote the set of all squares and nonsquares in  $GF(p)^*$ , respectively.
- $\eta$  and  $\bar{\eta}$  are the quadratic characters of  $GF(q)^*$  and  $GF(p)^*$ , representing the extend these quadratic characters by letting  $\eta(0) = 0$  and  $\bar{\eta}(0) = 0$ .

#### 2.2 Group characters and Gauss sums

An *additive character* of GF(q) is a nonzero function  $\chi$  from GF(q) to the set of nonzero complex numbers such that  $\chi(x+y) = \chi(x)\chi(y)$  for any pair  $(x,y) \in GF(q)^2$ . For each  $b \in GF(q)$ , the function

$$\chi_b(c) = \zeta_p^{\text{Tr}(bc)} \quad \text{for all } c \in \text{GF}(q) \tag{2}$$

defines an additive character of GF(q). When b = 0,  $\chi_0(c) = 1$  for all  $c \in GF(q)$ , and is called the *trivial additive character* of GF(q). The character  $\chi_1$  in (2) is called the *canonical additive character* of GF(q). It is known that every additive character of GF(q) can be written as  $\chi_b(x) = \chi_1(bx)$  [20, Theorem 5.7].

The Gauss sum  $G(\eta, \chi_1)$  over GF(q) is defined by

$$G(\eta, \chi_1) = \sum_{c \in GF(q)^*} \eta(c)\chi_1(c) = \sum_{c \in GF(q)} \eta(c)\chi_1(c)$$
(3)

and the Gauss sum  $G(\bar{\eta}, \bar{\chi}_1)$  over GF(p) is defined by

$$G(\bar{\eta}, \bar{\chi}_1) = \sum_{c \in \mathrm{GF}(p)^*} \bar{\eta}(c) \bar{\chi}_1(c) = \sum_{c \in \mathrm{GF}(p)} \bar{\eta}(c) \bar{\chi}_1(c), \tag{4}$$

where  $\bar{\chi}_1$  is the canonical additive characters of GF(p).

The following three lemmas are proved in [20, Theorem 5.15 and Theorem 5.33] and [14, lemma 7], respectively.

Lemma 1 With the symbols and notations above, we have

$$G(\eta, \chi_1) = (-1)^{m-1} \sqrt{-1}^{(\frac{p-1}{2})^2 m} \sqrt{q}$$

and

$$G(\bar{\eta}, \bar{\chi}_1) = \sqrt{-1}^{(\frac{p-1}{2})^2} \sqrt{p} = \sqrt{p*}.$$

**Lemma 2** Let  $\chi$  be a nontrivial additive character of GF(q) with q odd, and let  $f(x) = a_2x^2 + a_1x + a_0 \in GF(q)[x]$  with  $a_2 \neq 0$ . Then

$$\sum_{c \in \mathrm{GF}(q)} \chi(f(c)) = \chi(a_0 - a_1^2 (4a_2)^{-1}) \eta(a_2) G(\eta, \chi)$$

**Lemma 3** If  $m \ge 2$  is even, then  $\eta(y) = 1$  for each  $y \in GF(p)^*$ . If  $m \ge 1$  is odd, then  $\eta(y) = \overline{\eta}(y)$  for each  $y \in GF(p)$ .

2.3 A type of exponential sums

For any a and b in GF(q), we define the following exponential sum

$$S(a,b) = \sum_{x \in \mathrm{GF}(q)} \chi_1\left(ax^{p+1} + bx\right)$$

in this paper. To prove our main results, we need the values of the sum S(a,b) and the help of a number of lemmas that are proved in [5, Theorem 1, Theorem 2, Theorem 2.3].

**Lemma 4** Let *m* be odd,  $f(x) = a^p x^{p^2} + ax \in F_q[x]$  and  $b \in GF(q)$ . Then f(x) is a permutation polynomial over GF(q) and

$$S(a,b) = \sqrt{p^*}^m \eta(a) \chi_1(-a x_{a,b}^{p+1}),$$

where  $p* = (-1)^{\frac{p-1}{2}}p$  and  $x_{a,b}$  is the unique solution of the equation  $f(x) = -b^p$ . Particularly,  $S(a,0) = \sqrt{p^*}^m \eta(a)$ .

**Lemma 5** Let *m* be even,  $f(x) = a^p x^{p^2} + ax \in F_q[x]$  and  $b \in GF(q)$ . There are two cases.

1. If  $a^{\frac{q-1}{p+1}} \neq (-1)^{\frac{m}{2}}$ , then f(x) is a permutation polynomial over GF(q). Let  $x_{a,b}$  be the unique solution of the equation  $f(x) = -b^p$ . Then

$$S(a,b) = (-1)^{\frac{m}{2}} p^{m/2} \chi_1(-a x_{a,b}^{p+1})$$

2. If  $a^{\frac{q-1}{p+1}} = (-1)^{\frac{m}{2}}$ , then f(x) is not a permutation polynomial over GF(q). We have S(a,b) = 0 unless the the equation  $f(x) = -b^p$  is solvable. If this equation is solvable, with solution  $x_{a,b}$  say, then

$$S(a,b) = -(-1)^{\frac{m}{2}} p^{m/2+1} \chi_1(-a x_{a,b}^{p+1}).$$

Particularly,

$$S(a,0) = \begin{cases} (-1)^{\frac{m}{2}} p^{\frac{m}{2}} & \text{if } a^{\frac{q-1}{p+1}} \neq (-1)^{\frac{m}{2}}, \\ (-1)^{\frac{m}{2}+1} p^{\frac{m}{2}+1} & \text{if } a^{\frac{q-1}{p+1}} = (-1)^{\frac{m}{2}}. \end{cases}$$

2.4 Cyclotomic fields

In this subsection, we state some basic facts on cyclotomic fields. These results will be used in the rest of this paper.

Let  $\mathbb{Z}$  be the rational integer ring and Q be the rational field. Some results on cyclotomic field  $Q(\zeta_p)$  [17] are given in the following lemma.

Lemma 6 We have the following basic facts.

- 1. The ring of integers in  $K = Q(\zeta_p)$  is  $O_K = \mathbb{Z}[\zeta_p]$  and  $\{\zeta_p^i : 1 \le i \le p-1\}$  is an integral basis of  $O_K$ .
- 2. The field extension K/Q is Galois of degree p-1 and the Galois group  $Gal(K/Q) = \{\sigma_a : a \in (\mathbb{Z}/p\mathbb{Z})^*\}$ , where the automorphism  $\sigma_a$  of K is defined by  $\sigma_a(\zeta_p) = \zeta_p^a$ .
- {σ<sub>a</sub> : a ∈ (ℤ/pℤ)\*}, where the automorphism σ<sub>a</sub> of K is defined by σ<sub>a</sub>(ζ<sub>p</sub>) = ζ<sup>a</sup><sub>p</sub>.
  3. The field K has a unique quadratic subfield L = Q(√p\*). For 1 ≤ a ≤ p − 1, σ<sub>a</sub>(√p\*) = η(a)√p\*. Therefore, the Galois group Gal(L/Q) is {1,σ<sub>γ</sub>}, where γ is any quadratic nonresidue in GF(p).

### 3 The linear codes with a few weights

We only describe the codes and introduce their parameters in this section. The proofs of their parameters will be given in Section 4.

In this paper, the defining set D of the code  $C_D$  of (1) is given by

$$D = \{ x \in GF(q)^* : Tr(x^{p+1} - x) = 0 \}.$$
 (5)

When p = 2, the weight distribution of the code  $C_D$  was settled in [26]. In this paper, we study the code  $C_D$  for p being an odd prime.

The following three theorems are the main results of this paper.

**Theorem 1** Let  $m \ge 3$  be odd, and let D be defined in (5). Then the set  $C_D$  of (1) is an [n,m] linear code over GF(p) with at most five weights and the weight distribution in Tables 1 and 2, where

$$n = p^{m-1} - 1 + (-1)^{\frac{p-1}{2}\frac{m-1}{2}} \bar{\eta}(m_p) p^{\frac{m-1}{2}}.$$
(6)

Weight w	Multiplicity $A_w$
0	1
$(p-1)p^{m-2}$	$p^{m-1} - 1$
$(p-1)(p^{m-2}-(-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-3}{2}})$	$\frac{p-1}{2}(p^{m-2}+(-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-1}{2}})$
$(p-1)(p^{m-2}+(-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-3}{2}})$	$\frac{p-1}{2}(p^{m-2}-(-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-1}{2}})$
$(p-1)p^{m-2} + (-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-3}{2}}$	$\frac{1}{2}(p-1)^2 p^{m-2}$
$(p-1)p^{m-2} - (-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-3}{2}}$	$\frac{1}{2}(p-1)^2 p^{m-2}$

Table 1: The weight distribution of  $C_D$  of Theorem 1 when  $m_p = 0$ 

Table 2: The weight distribution of  $C_D$  of Theorem 1 when  $m_p \neq 0$ 

Weight <i>w</i>	Multiplicity $A_w$
0	1
$(p-1)p^{m-2}$	$p^{m-2} - 1 + \bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}}(p-1)p^{\frac{m-3}{2}}$
$(p-1)p^{m-2} + \bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-1}{2}}$	$p^{m-2}(p-1) - \bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}}(p-1)p^{\frac{m-3}{2}}$
$(p-1)p^{m-2} + \bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}}(p+1)p^{\frac{m-3}{2}}$	$\frac{1}{2}(p-1)(p-2)p^{\frac{m-3}{2}}(p^{\frac{m-1}{2}}-\bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}})$
$(p-1)p^{m-2} + \bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}}(p-1)p^{\frac{m-3}{2}}$	$\frac{p-1}{2}(p^{m-1}-\bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-1}{2}})$
$(p-1)p^{m-2} + \bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-3}{2}}$	$(p-1)p^{m-2} + \bar{\eta}(m_p)(-1)^{\frac{p-1}{2}\frac{m-1}{2}}(p-1)^2p^{\frac{m-3}{2}}$

*Example 1* Let (p,m) = (3,5). Then the code  $C_D$  has parameters [71,5,42] and weight enumerator  $1 + 30z^{42} + 60z^{45} + 90z^{48} + 42z^{51} + 20z^{54}$ , which is verified by a Magma program.

*Example 2* Let (p,m) = (3,9). Then the code  $C_D$  has parameters [6560,9,4320] and weight enumerator  $1 + 2268z^{4320} + 4374z^{4347} + 6560z^{4374} + 4374z^{4401} + 2106z^{4428}$ , which is verified by a Magma program.

*Remark 1* The code  $C_D$  of Theorem 1 is a five-weight linear code except in the following two cases:

- 1. When m = 3 and  $m_p = 0$ , the frequency of the weight  $(p-1)(p^{m-2} (-1)^{\frac{p-1}{2}}m^{\frac{m-3}{2}})$  of Table 1 turns out to be 0. Thus, the code  $C_D$  has four nonzero weights.
- 2. When m = 3,  $m_p \neq 0$  and  $p \equiv 2 \pmod{3}$ , the frequency of the weight  $(p-1)p^{m-2}$  of Table 2 turns out to be 0. Thus, the code  $C_D$  has only four nonzero weights.

*Example 3* Let (p,m) = (3,3). Then the code  $C_D$  has parameters [8,3,4] and weight enumerator  $1+6z^4+6z^5+8z^6+6z^7$ , which is verified by a Magma program. This code is almost optimal, since the optimal linear code has parameters [8,3,5].

*Example 4* Let (p,m) = (5,3). Then the code  $C_D$  has parameters [19,3,14] and weight enumerator  $1 + 36z^{14} + 24z^{15} + 60z^{16} + 4z^{19}$ , which is verified by a Magma program. This code is optimal.

**Theorem 2** Let  $m \ge 2$  be even and  $m \equiv 2 \pmod{4}$ , and let D be defined in (5). Then the set  $C_D$  of (1) is an [n,m] linear code over GF(p) with at most three weights and the weight distribution in Table 3 and Table 4, where

$$n = \begin{cases} p^{m-1} - 1 - (p-1)p^{\frac{1}{2}(m-1+(-1)^{\frac{m}{2}})} & \text{if } m_p = 0, \\ p^{m-1} - 1 + p^{\frac{1}{2}(m-1+(-1)^{\frac{m}{2}})} & \text{if } m_p \neq 0. \end{cases}$$
(7)

Table 3: The weight distribution of  $C_D$  of Theorem 2 when  $m_p = 0$ 

Weight w	Multiplicity A <sub>w</sub>
0	1
$(p-1)p^{m-2}$	$p^{m-2} - (p-1)p^{\frac{m}{2}-1} - 1$
$(p-1)(p^{m-2}-p^{\frac{m}{2}-1})$	$(p-1)(2p^{m-2}+p^{\frac{m}{2}-1})$
$(p-1)p^{m-2} - (p-2)p^{\frac{m}{2}-1}$	$(p-1)^2 p^{m-2}$

Table 4: The weight distribution of  $C_D$  of Theorem 2 when  $m_p \neq 0$ 

Weight w	Multiplicity $A_w$
0	1
$(p-1)p^{m-2}$	$p^{m-2} + \frac{p-1}{2}(p^{m-1} + p^{\frac{m}{2}}) - 1$
$(p-1)p^{m-2} + p^{\frac{m}{2}-1}$	$(p-1)(2p^{m-2}-p^{\frac{m}{2}-1})$
$(p-1)p^{m-2}+2p^{\frac{m}{2}-1}$	$\frac{(p-1)(p-2)}{2}(p^{m-2}-p^{\frac{m}{2}-1})$

*Example 5* Let (p,m) = (3,6). Then the code  $C_D$  has parameters [224,6,144] and weight enumerator  $1 + 342z^{144} + 324z^{153} + 62z^{162}$ , which is verified by a Magma program.

*Example 6* Let (p,m) = (5,6). Then the code  $C_D$  has parameters [3149,6,2500] and weight enumerator  $1 + 7124z^{2500} + 2525z^{4900} + 2550z^{3600}$ , which is verified by a Magma program.

*Remark* 2 The code  $C_D$  of Theorem 2 has three weights except in the case m = 2, as the frequency of the weight  $(p-1)p^{m-2} + 2p^{\frac{m}{2}-1}$  of Table 4 turns out to be 0. Hence, The code  $C_D$  of Theorem 2 is a two-weight linear code if and only if m = 2.

*Example* 7 Let (p,m) = (3,2). Then the code  $C_D$  has parameters [3,2,2] and weight enumerator  $1 + 6z^2 + 2z^3$ , which is verified by a Magma program. This code is optimal.

**Theorem 3** Let  $m \ge 6$  be even and  $m \equiv 0 \pmod{4}$ , and let *D* be defined in (5). Then the set  $C_D$  of (1) is a four-weight linear code over GF(*p*) with the parameter [n,m] and the weight distribution in Tables 5 and 6, where *n* is defined by (7).

Table 5: The weight distribution of  $C_D$  of Theorem 3 when  $m_p = 0$ 

Weight w	Multiplicity $A_w$
0	1
$p^{m-2}(p-1) - (p-1)^2 p^{\frac{m}{2}-1}$	$(p^2 - 1)p^{m-2}$
$p^{m-2}(p-1)$	$p^{m-4} - (p-1)p^{\frac{m}{2}-2} - 1$
$(p-1)p^{\frac{m}{2}}(p^{\frac{m}{2}-2}-1)$	$(p-1)(2p^{m-4}+p^{\frac{m}{2}-2})$
$(p-1)p^{m-2} - (p-2)p^{\frac{m}{2}}$	$(p-1)^2 p^{m-4}$

Table 6: The weight distribution of  $C_D$  of Theorem 3 when  $m_p \neq 0$ 

Weight w	Multiplicity $A_w$
0	1
$(p-1)(p^{\frac{m}{2}-1}+p^{m-2})$	$p^m - p^{m-2}$
$(p-1)p^{m-2}$	$p^{m-4} + \frac{p-1}{2}(p^{\frac{m}{2}-1} + p^{m-3}) - 1$
$(p-1)p^{m-2} + p^{\frac{m}{2}}$	$(p-1)(2p^{m-4}-p^{\frac{m}{2}-2})$
$(p-1)p^{m-2}+2p^{\frac{m}{2}}$	$\frac{1}{2}(p-1)(p-2)(p^{m-4}-p^{\frac{m}{2}-2})$

*Example 8* Let (p,m) = (3,8). Then the code  $C_D$  has parameters [2267,8,1458] and weight enumerator  $1 + 350z^{1458} + 5832z^{1512} + 306z^{1539} + 32z^{1620}$ , which is verified by a Magma program.

*Example 9* Let (p,m) = (5,8). Then the code  $C_D$  has parameters [78749,8,62500] and weight enumerator  $1 + 7124z^{62500} + 375000z^{63000} + 4900z^{63125} + 3600z^{63750}$ , which is verified by a Magma program.

# 4 The proofs of the main results

Our task of this section is to prove Theorems 1, 2 and 3, respectively. To this end, we shall prove a few more auxiliary results before proving the main results of this paper.

# 4.1 Some auxiliary results

Lemma 7 With the symbols and notations above, we have

$$\sum_{c \in GF(p)^*} S(c, -c) = \begin{cases} (-1)^{\frac{p-1}{2}} \bar{\eta}(m_p) \sqrt{p^*}^{m+1} & \text{if } m \text{ is odd,} \\ -(p-1)p^{\frac{1}{2}(m+1+(-1)^{\frac{m}{2}})} & \text{if } m \text{ is even and } m_p = 0, \\ p^{\frac{1}{2}(m+1+(-1)^{\frac{m}{2}})} & \text{if } m \text{ is even and } m_p \neq 0. \end{cases}$$

*Proof* For any  $c \in GF(p)^*$ , it is easily seen that  $f(x) = c^p x^{p^2} + cx = -(-c)^p$  is solvable and  $\frac{1}{2}$  is its solution.

By definition, Lemmas 4 and 5, we have

$$\begin{split} &\sum_{c \in \mathrm{GF}(p)^*} S(c, -c) \\ &= \sum_{c \in \mathrm{GF}(p)^*} \sum_{x \in \mathrm{GF}(q)} \chi_1\left(cx^{p+1} - cx\right) \\ &= \begin{cases} \sum_{c \in \mathrm{GF}(p)^*} \sqrt{p^*} \eta(c) \chi_1\left(-\frac{c}{4}\right) & \text{if } m \text{ is odd,} \\ \sum_{c \in \mathrm{GF}(p)^*} (-1)^{\frac{m}{2}} p^{m/2} \chi_1\left(-\frac{c}{4}\right) & \text{if } m \equiv 2( \mod 4), \\ \sum_{c \in \mathrm{GF}(p)^*} -(-1)^{\frac{m}{2}} p^{m/2+1} \chi_1\left(-\frac{c}{4}\right) & \text{if } m \equiv 0( \mod 4), \\ \sum_{c \in \mathrm{GF}(p)^*} \overline{\chi}_1(c) \overline{\chi}_1\left(-\frac{c\mathrm{Tr}(1)}{4}\right) & \text{if } m \equiv 2( \mod 4), \\ -p^{m/2} \sum_{c \in \mathrm{GF}(p)^*} \overline{\chi}_1\left(\frac{c\mathrm{Tr}(1)}{4}\right) & \text{if } m \equiv 0( \mod 4), \\ -p^{m/2+1} \sum_{c \in \mathrm{GF}(p)^*} \overline{\chi}_1\left(\frac{c\mathrm{Tr}(1)}{4}\right) & \text{if } m \equiv 0( \mod 4), \\ \sqrt{p^*} \sum_{c \in \mathrm{GF}(p)^*} \overline{\chi}_1\left(\frac{c\mathrm{Tr}(1)}{4}\right) & \text{if } m \equiv 0( \mod 4), \\ \sqrt{p^*} \sum_{c \in \mathrm{GF}(p)^*} \overline{\eta}(-\frac{m_p}{4}) \overline{\eta}(-\frac{cm_p}{4}) \overline{\chi}_1\left(-\frac{cm_p}{4}\right) & \text{if } m \equiv 0 \text{ dod } m_p = 0, \\ \sqrt{p^*} \sum_{c \in \mathrm{GF}(p)^*} \overline{\eta}(-\frac{m_p}{4}) \overline{\chi}_1\left(-\frac{cm_p}{4}\right) & \text{if } m \equiv 0 \text{ dod } m_p \neq 0, \\ -p^{m/2}(p-1) & \text{if } m \equiv 2( \mod 4) \text{ and } m_p \neq 0, \\ p^{m/2} & \text{if } m \equiv 0 \text{ (mod 4) and } m_p = 0, \\ p^{m/2+1} & \text{if } m \equiv 0 \text{ (mod 4) and } m_p \neq 0, \\ -p^{m/2}(p-1) & \text{if } m \equiv 2( \mod 4) \text{ and } m_p \neq 0, \\ -p^{m/2}(p-1) & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 2( \mod 4) \text{ and } m_p = 0, \\ p^{m/2} & \text{if } m \equiv 0( \mod 4) \text{ and } m_p = 0, \\ p^{m/2+1} & \text{if } m \equiv 0( \mod 4) \text{ and } m_p = 0, \\ p^{m/2+1} & \text{if } m \equiv 0( \mod 4) \text{ and } m_p = 0. \end{cases}$$

The desired conclusion then follows from Lemma 1.

The next lemma will be employed in proving the code length.

Lemma 8 Let

$$n_0 = |\{x \in GF(q) : Tr(x^{p+1} - x) = 0\}|.$$

Then

$$n_{0} = \begin{cases} p^{m-1} + (-1)^{\frac{p-1}{2}} \bar{\eta}(m_{p}) p^{-1} \sqrt{p^{*}}^{m+1} & \text{if } m \text{ is odd,} \\ p^{m-1} - (p-1) p^{\frac{1}{2}(m-1+(-1)^{\frac{m}{2}})} & \text{if } m \text{ is even and } m_{p} = 0, \\ p^{m-1} + p^{\frac{1}{2}(m-1+(-1)^{\frac{m}{2}})} & \text{if } m \text{ is even and } m_{p} \neq 0. \end{cases}$$

*Proof* By definition, we have

$$n_{0} = \frac{1}{p} \sum_{x \in GF(q)} \sum_{y \in GF(p)} \zeta_{p}^{y(\operatorname{Tr}(x^{p+1})-x)}$$
  
=  $p^{m-1} + \frac{1}{p} \sum_{y \in GF(p)^{*}} \sum_{x \in GF(q)} \zeta_{p}^{\operatorname{Tr}(yx^{p+1}-y)}$   
=  $p^{m-1} + \frac{1}{p} \sum_{y \in GF(p)^{*}} S(y, -y).$ 

The desired conclusion then follows from Lemma 7.

From Lemma 6, the conclusion of the following lemma is straightforward and we omit their proofs.

Lemma 9 With the symbols and notations above, we have the following.

$$1. \ \sum_{y \in GF(p)^*} \sigma_y(\zeta_p^z) = \begin{cases} p-1 & \text{if } z = 0, \\ -1 & \text{if } z \neq 0. \end{cases}$$

$$2. \ \sum_{y \in GF(p)^*} \sigma_y(\sqrt{p^*}^m) = \begin{cases} 0 & \text{if } m \text{ is odd,} \\ \sqrt{p^*}^m(p-1) & \text{if } m \text{ is even} \end{cases}$$

$$3. \ \sum_{y \in GF(p)^*} \sigma_y(\sqrt{p^*} \zeta_p^z) = \begin{cases} 0 & \text{if } z = 0, \\ \bar{\eta}(z)p^* & \text{if } z \neq 0. \end{cases}$$

Particularly,

$$\sum_{\mathbf{y}\in \mathrm{GF}(p)^*} \sigma_{\mathbf{y}}(\sqrt{p^*}^m \zeta_p^{-\frac{m_p}{4}}) = \begin{cases} 0 & \text{if $m$ is odd and $m_p = 0$,} \\ (-1)^{\frac{p-1}{2}} \bar{\eta}(m_p)\sqrt{p^*}^{m+1} & \text{if $m$ is odd and $m_p \neq 0$.} \end{cases}$$

and

$$\sum_{y \in GF(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}}) = \begin{cases} p-1 & \text{if } m_p = 0, \\ -1 & \text{if } m_p \neq 0. \end{cases}$$

The following result will play an important role in proving the main results of this paper.

**Lemma 10** Let  $b \in GF(q)^*$ ,  $L(x) = x^{p^2} + x$  and

$$M = \sum_{y \in \mathrm{GF}(p)^*} \sum_{z \in \mathrm{GF}(p)^*} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\mathrm{Tr}(yx^{p+1} + (bz - y)x)}.$$

(I) If m is odd, then we have the following.

- *When*  $m_p = 0$ ,

$$M = \begin{cases} 0 & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0, \\ (-1)^{\frac{p-1}{2}} (p-1)\sqrt{p^*}^{m+1} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) = 0, \\ -(-1)^{\frac{p-1}{2}} (p-1)\sqrt{p^*}^{m+1} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{NSQ} \text{ and } \operatorname{Tr}(x_b) = 0, \\ -(-1)^{\frac{p-1}{2}}\sqrt{p^*}^{m+1} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ (-1)^{\frac{p-1}{2}}\sqrt{p^*}^{m+1} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{NSQ} \text{ and } \operatorname{Tr}(x_b) \neq 0. \end{cases}$$

- When  $m_p \in SQ$ ,

$$M = \begin{cases} (-1)^{\frac{p-1}{2}} (p-1)\sqrt{p^{*}}^{m+1} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ -(-1)^{\frac{p-1}{2}}\sqrt{p^{*}}^{m+1} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ -2 \cdot (-1)^{\frac{p-1}{2}}\sqrt{p^{*}}^{m+1} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) = 0, \\ 0 & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ}, \\ (p-2)(-1)^{\frac{p-1}{2}}\sqrt{p^{*}}^{m+1} \text{ or } -2 \cdot (-1)^{\frac{p-1}{2}}\sqrt{p^{*}}^{m+1} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) \neq 0. \end{cases}$$

- When  $m_p \in NSQ$ ,

$$M = \begin{cases} -(-1)^{\frac{p-1}{2}}(p-1)\sqrt{p^{*}}^{m+1} & \text{if } \operatorname{Tr}(x_{b}^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_{b}) = 0, \\ (-1)^{\frac{p-1}{2}}\sqrt{p^{*}}^{m+1} & \text{if } \operatorname{Tr}(x_{b}^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_{b}) \neq 0, \\ 0 & \text{if } \operatorname{Tr}(x_{b}^{p+1}) \in \operatorname{SQ}, \\ 2 \cdot (-1)^{\frac{p-1}{2}}\sqrt{p^{*}}^{m+1} & \text{if } \operatorname{Tr}(x_{b}^{p+1}) \in \operatorname{NSQ} \text{ and } \operatorname{Tr}(x_{b}) = 0, \\ -(p-2)(-1)^{\frac{p-1}{2}}\sqrt{p^{*}}^{m+1} \text{ or } 2 \cdot (-1)^{\frac{p-1}{2}}\sqrt{p^{*}}^{m+1} & \text{if } \operatorname{Tr}(x_{b}^{p+1}) \in \operatorname{NSQ} \text{ and } \operatorname{Tr}(x_{b}) \neq 0. \end{cases}$$

where  $x_b$  is the unique solution of the equation  $x^{p^2} + x = -b^p$ .

(II) If  $m \equiv 2 \pmod{4}$ , then we have the following.

- When 
$$m_p = 0$$
,

$$M = \begin{cases} -p^{\frac{m}{2}}(p-1)^2 & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{\frac{m}{2}}(p-1) & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ & \text{or } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ -p^{\frac{m}{2}} & \text{if } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } \operatorname{Tr}(x_b) \neq 0. \end{cases}$$

- When  $m_p \neq 0$ ,

$$M = \begin{cases} p^{\frac{m}{2}}(p-1) & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ -p^{\frac{m}{2}} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0 \\ & \text{or } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } A = 0, \\ p^{\frac{m}{2}}(p-1) \text{ or } -p^{\frac{m}{2}}(p+1) & \text{if } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } A \neq 0, \end{cases}$$

where  $A = -\frac{m_p}{4} + \frac{\text{Tr}(x_b)^2}{4\text{Tr}(x_b^{p+1})}$  and  $x_b$  is the unique solution of the equation  $x^{p^2} + x = -b^p$ . (III) If  $m \equiv 0 \pmod{4}$ , then we have the following.

- *When*  $m_p = 0$ ,

$$M = \begin{cases} 0 & \text{if } b \notin \operatorname{Im}(L), \\ -p^{\frac{m}{2}+1}(p-1)^2 & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{\frac{m}{2}+1}(p-1) & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0 \\ & \text{or } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ -p^{\frac{m}{2}+1} & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } \operatorname{Tr}(x_b) \neq 0. \end{cases}$$

- When  $m_p \neq 0$ ,

$$M = \begin{cases} 0 & \text{if } b \notin \operatorname{Im}(L), \\ p^{\frac{m}{2}+1}(p-1) & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ -p^{\frac{m}{2}+1} & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0 \\ & \text{or } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } A = 0, \\ p^{\frac{m}{2}+1}(p-1) \text{ or } -p^{\frac{m}{2}+1}(p+1) & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } A \neq 0. \end{cases}$$

where  $A = -\frac{m_p}{4} + \frac{\operatorname{Tr}(x_b)^2}{4\operatorname{Tr}(x_b^{p+1})}$  and  $x_b$  is some solution of the equation  $L(x) = -b^p$  under the condition that  $b \in \operatorname{Im}(L)$ .

Proof We have

$$M = \sum_{y \in GF(p)^{*}} \sum_{z \in GF(p)^{*}} \sum_{x \in GF(q)} \sum_{x \in GF(q)} \zeta_{p}^{\operatorname{Tr}(yx^{p+1} + (bz-y)x)}$$
  
$$= \sum_{y \in GF(p)^{*}} \sum_{z \in GF(p)^{*}} \sum_{x \in GF(q)} \zeta_{p}^{\operatorname{yTr}(x^{p+1} + (zb-1)x)}$$
  
$$= \sum_{y \in GF(p)^{*}} \sum_{z \in GF(p)^{*}} \sum_{x \in GF(q)} \zeta_{p}^{\operatorname{yTr}(x^{p+1} + (bz-1)x)}$$
  
$$= \sum_{y \in GF(p)^{*}} \sigma_{y} (\sum_{z \in GF(p)^{*}} \sum_{x \in GF(q)} \zeta_{p}^{\operatorname{Tr}(x^{p+1} + (bz-1)x)}).$$
(8)

It is easily seen that  $L(x) = x^{p^2} + x$  is a permutation polynomial over GF(q) if  $m \equiv 2$ (mod 4) or *m* is odd. Hence,  $x_b$  is the unique solution of the equation  $L(x) = -b^p$ , while  $zx_b + \frac{1}{2}$  is the unique solution of the equation  $L(x) = -(bz - 1)^p$  for  $z \in GF(p)^*$ . For  $m \equiv 0 \pmod{4}$ , although  $L(x) = x^{p^2} + x$  is not a permutation polynomial over GF(q),  $zx_b + \frac{1}{2}$  is the solution of the equation  $L(x) = -(bz - 1)^p$  for  $z \in GF(p)^*$  when

 $b \in \text{Im}(L)$  with some solution  $x_b$ .

Therefore, by Lemmas 4 and 5, Equation (8) becomes

$$M = \begin{cases} \sum_{y \in GF(p)^*} \sigma_y (\sum_{z \in GF(p)^*} \sqrt{p^*}^m \eta(1) \zeta_p^{-\operatorname{Tr}((zx_b + \frac{1}{2})^{p+1})}) & \text{if } m \text{ is odd} \\ \sum_{y \in GF(p)^*} \sigma_y (\sum_{z \in GF(p)^*} (-1)^{\frac{m}{2}} p^{\frac{m}{2}} \zeta_p^{-\operatorname{Tr}((zx_b + \frac{1}{2})^{p+1})}) & \text{if } m \equiv 2( \mod 4) \\ 0 & \text{if } m \equiv 0( \mod 4) \text{ and } b \notin \operatorname{Im}(L) \\ \sum_{y \in GF(p)^*} \sigma_y (\sum_{z \in GF(p)^*} (-1)^{\frac{m}{2}} p^{\frac{m}{2}+1} \zeta_p^{-\operatorname{Tr}((zx_b + \frac{1}{2})^{p+1})}) & \text{if } m \equiv 0( \mod 4) \text{ and } b \in \operatorname{Im}(L) \end{cases}$$
$$= \begin{cases} \sum_{y \in GF(p)^*} \sigma_y (\sqrt{p^{*m}} \zeta_p^{-\frac{m_p}{4}} \sum_{z \in GF(p)^*} \zeta_p^{-\operatorname{Tr}(x_b^{p+1})z^2 - \operatorname{Tr}(x_b)z}) & \text{if } m \text{ is odd,} \\ \sum_{y \in GF(p)^*} \sigma_y (-p^{\frac{m}{2}} \zeta_p^{-\frac{m_p}{4}} \sum_{z \in GF(p)^*} \zeta_p^{-\operatorname{Tr}(x_b^{p+1})z^2 - \operatorname{Tr}(x_b)z}) & \text{if } m \equiv 0( \mod 4), \\ 0 & \text{if } m \equiv 0( \mod 4) \text{ and } b \notin \operatorname{Im}(L), \\ \sum_{y \in GF(p)^*} \sigma_y (-p^{\frac{m}{2}+1} \zeta_p^{-\frac{m_p}{4}} \sum_{z \in GF(p)^*} \zeta_p^{-\operatorname{Tr}(x_b^{p+1})z^2 - \operatorname{Tr}(x_b)z}) & \text{if } m \equiv 0( \mod 4) \text{ and } b \notin \operatorname{Im}(L), \end{cases}$$

We distinguish the following three cases.

(I) If *m* is odd, by Lemma 2, we have

$$M = \begin{cases} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\sqrt{p^*}^m \zeta_p^{-\frac{m_p}{4}} \sum_{z \in \mathrm{GF}(p)^*} \zeta_p^{-\mathrm{Tr}(x_b)z}) & \text{if } \mathrm{Tr}(x_b^{p+1}) = 0\\ \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\sqrt{p^*}^m \zeta_p^{-\frac{m_p}{4}} (\bar{\eta}(-\mathrm{Tr}(x_b^{p+1})\zeta_p^{\frac{\mathrm{Tr}(x_b)^2}{4\mathrm{Tr}(x_b^{p+1})}} G(\bar{\eta}, \bar{\chi}_1) - 1)) & \text{if } \mathrm{Tr}(x_b^{p+1}) \neq 0\\ = \begin{cases} (p-1)A_0 & \text{if } \mathrm{Tr}(x_b^{p+1}) = 0 \text{ and } \mathrm{Tr}(x_b) = 0,\\ -A_0 & \text{if } \mathrm{Tr}(x_b^{p+1}) = 0 \text{ and } \mathrm{Tr}(x_b) \neq 0,\\ \sqrt{p^*}^{m+1} \bar{\eta}(-\mathrm{Tr}(x_b^{p+1}) \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}}) - A_0 & \text{if } \mathrm{Tr}(x_b^{p+1}) \neq 0 \text{ and } \mathrm{Tr}(x_b) = 0,\\ \sqrt{p^*}^{m+1} \bar{\eta}(-\mathrm{Tr}(x_b^{p+1}) \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-1}) - A_0 & \text{if } \mathrm{Tr}(x_b^{p+1}) \neq 0 \text{ and } \mathrm{Tr}(x_b) = 0, \end{cases}$$

where  $A_0 = \sum_{y \in GF(p)^*} \sigma_y(\sqrt{p^{*m}} \zeta_p^{-\frac{m_p}{4}})$  and  $A = -\frac{m_p}{4} + \frac{\operatorname{Tr}(x_b)^2}{4\operatorname{Tr}(x_b^{p+1})}$ . The desired conclusion in Part (I) of this lemma then follows from Lemma 9.

(II) If  $m \equiv 2 \pmod{4}$ , by Lemmas 2 and 1, we have

$$\begin{split} M &= \begin{cases} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(-p^{\frac{m}{2}} \zeta_p^{-\frac{m_p}{4}} \sum_{z \in \mathrm{GF}(p)^*} \zeta_p^{-\mathrm{Tr}(x_b)z}) & \text{if } \mathrm{Tr}(x_b^{p+1}) = 0\\ \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(-p^{\frac{m}{2}} \zeta_p^{-\frac{m_p}{4}} (\sum_{z \in \mathrm{GF}(p)} \zeta_p^{-\mathrm{Tr}(x_b^{p+1})z^2 - \mathrm{Tr}(x_b)z} - 1)) & \text{if } \mathrm{Tr}(x_b^{p+1}) \neq 0 \end{cases} \\ &= \begin{cases} -p^{\frac{m}{2}} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}} \sum_{z \in \mathrm{GF}(p)^*} \zeta_p^{-\mathrm{Tr}(x_b)z}) & \text{if } \mathrm{Tr}(x_b^{p+1}) = 0\\ -p^{\frac{m}{2}} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}} (\tilde{\eta}(-\mathrm{Tr}(x_b^{p+1})\zeta_p^{-\mathrm{Tr}(x_b)^2}) & \text{if } \mathrm{Tr}(x_b^{p+1}) = 0\\ -p^{\frac{m}{2}} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}} (\tilde{\eta}(-\mathrm{Tr}(x_b^{p+1})\zeta_p^{-\mathrm{Tr}(x_b^{p+1})} G(\tilde{\eta}, \tilde{\chi}_1) - 1)) & \text{if } \mathrm{Tr}(x_b^{p+1}) \neq 0\\ \end{cases} \\ &= \begin{cases} -p^{\frac{m}{2}} (p-1)A_1 & \text{if } \mathrm{Tr}(x_b^{p+1}) = 0 & \text{and } \mathrm{Tr}(x_b) = 0,\\ p^{\frac{m}{2}}A_1 & \text{if } \mathrm{Tr}(x_b^{p+1}) = 0 & \text{and } \mathrm{Tr}(x_b) \neq 0,\\ -p^{\frac{m}{2}} (-1)^{\frac{p-1}{2}} \tilde{\eta}(\mathrm{Tr}(x_b^{p+1})) \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\sqrt{p^*}\zeta_p^{-\frac{m_p}{4}}) + p^{\frac{m}{2}}A_1 & \text{if } \mathrm{Tr}(x_b^{p+1}) \neq 0 & \text{and } \mathrm{Tr}(x_b) = 0,\\ -p^{\frac{m}{2}} (-1)^{\frac{p-1}{2}} \tilde{\eta}(\mathrm{Tr}(x_b^{p+1})) \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\sqrt{p^*}\zeta_p^{-\frac{m_p}{4}}) + p^{\frac{m}{2}}A_1 & \text{if } \mathrm{Tr}(x_b^{p+1}) \neq 0 & \text{and } \mathrm{Tr}(x_b) = 0, \end{cases} \end{cases}$$

where  $A_1 = \sum_{y \in GF(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}})$  and  $A = -\frac{m_p}{4} + \frac{\operatorname{Tr}(x_b)^2}{4\operatorname{Tr}(x_b^{p+1})}$ . The desired conclusion in Part (II) of this lemma then follows from Lemma 9.

(III) If  $m \equiv 0 \pmod{4}$ , then from Lemmas 2 and 1 we have

$$M = \begin{cases} 0 & \text{if } b \notin \operatorname{Im}(L) \\ -p^{\frac{m}{2}+1} \sum_{y \in \operatorname{GF}(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}} \sum_{z \in \operatorname{GF}(p)^*} \zeta_p^{-\operatorname{Tr}(x_b)z}) & \text{if } b \in \operatorname{Im}(L) \text{ and } \operatorname{Tr}(x_b^{p+1}) = 0 \\ -p^{\frac{m}{2}+1} \sum_{y \in \operatorname{GF}(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}} (\sum_{z \in \operatorname{GF}(p)} \zeta_p^{-\operatorname{Tr}(x_b^{p+1})z^2 - \operatorname{Tr}(x_b)z} - 1)) & \text{if } b \in \operatorname{Im}(L) \text{ and } \operatorname{Tr}(x_b^{p+1}) \neq 0 \end{cases}$$
$$= \begin{cases} 0 & \text{if } b \notin \operatorname{Im}(L), \\ -p^{\frac{m}{2}+1}(p-1)A_1 & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{\frac{m}{2}+1}A_1 & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ -p^{\frac{m}{2}+1}(-1)^{\frac{p-1}{2}} \bar{\eta}(\operatorname{Tr}(x_b^{p+1}))A + p^{\frac{m}{2}+1}A_1 & \text{if } b \in \operatorname{Im}(L) \text{ and } \operatorname{Tr}(x_b^{p+1}) \neq 0, \end{cases}$$

where  $A_1 = \sum_{y \in GF(p)^*} \sigma_y(\zeta_p^{-\frac{m_p}{4}})$  and  $A = \sum_{y \in GF(p)^*} \sigma_y(\sqrt{p^*}\zeta_p^{-\frac{m_p}{4}} + \frac{\operatorname{Tr}(x_b)^2}{4\operatorname{Tr}(x_b^{p+1})})$ . The desired conclusion in Part (III) of this lemma then follows from Lemma 9.

This completes the proof of this lemma.

In order to calculate the Hamming weight of each codeword in  $C_D$ , we need the following result.

**Lemma 11** For any  $b \in GF(q)^*$ , let

$$N(b) = |\{x \in GF(q) : Tr(x^{p+1} - x) = 0 \text{ and } Tr(bx) = 0\}|.$$

There are three cases.

(I) If m is odd, then we have the following.

• When  $m_p = 0$ ,

$$N(b) = \begin{cases} p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0, \\ (p-1)B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) = 0, \\ -(p-1)B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{NSQ} \text{ and } \operatorname{Tr}(x_b) = 0, \\ -B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{NSQ} \text{ and } \operatorname{Tr}(x_b) \neq 0, \end{cases}$$

• When  $\bar{\eta}(m_p) = 1$ ,

$$N(b) = \begin{cases} pB + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ -B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) = 0, \\ B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{NSQ}, \\ (p-1)B + p^{m-2} \text{ or } -B + p^{m-2} \text{ if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) \neq 0. \end{cases}$$

• When  $\bar{\eta}(m_p) = -1$ ,

$$N(b) = \begin{cases} -pB + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ -B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ}, \\ B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{NSQ} \text{ and } \operatorname{Tr}(x_b) = 0, \\ -(p-1)B + p^{m-2} \text{ or } B + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{NSQ} \text{ and } \operatorname{Tr}(x_b) \neq 0. \end{cases}$$

where  $B = (-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-3}{2}}$ .

(II) If  $m \equiv 2 \pmod{4}$ , then we have the following.

• When  $m_p = 0$ ,

$$N(b) = \begin{cases} -(p-1)p^{\frac{m}{2}-1} + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ & \text{or } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ -p^{\frac{m}{2}-1} + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } \operatorname{Tr}(x_b) \neq 0. \end{cases}$$

• When  $m_p \neq 0$ ,

$$N(b) = \begin{cases} p^{\frac{m}{2}-1} + p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{m-2} & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ & \text{or } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } A = 0, \\ p^{\frac{m}{2}-1} + p^{m-2} \text{ or } - p^{\frac{m}{2}-1} + p^{m-2} \text{ if } \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } A \neq 0, \end{cases}$$

where  $A = -\frac{m_p}{4} + \frac{\operatorname{Tr}(x_b)^2}{4\operatorname{Tr}(x_b^{p+1})}$  and  $x_b$  is the unique solution of the equation  $x^{p^2} + x = -b^p$ . (III) If  $m \equiv 0 \pmod{4}$ , then we have the following.

• When  $m_p = 0$ ,

$$N(b) = \begin{cases} -(p-1)p^{\frac{m}{2}-1} + p^{m-2} & \text{if } b \notin \operatorname{Im}(L), \\ -(p-1)p^{\frac{m}{2}} + p^{m-2} & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{m-2} & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ & or \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ -p^{\frac{m}{2}} + p^{m-2} & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } \operatorname{Tr}(x_b) \neq 0. \end{cases}$$

• When  $m_p \neq 0$ ,

$$N(b) = \begin{cases} p^{\frac{m}{2}-1} + p^{m-2} & \text{if } b \notin \operatorname{Im}(L), \\ p^{\frac{m}{2}} + p^{m-2} & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ p^{m-2} & \text{if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ & or \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } A = 0, \\ p^{\frac{m}{2}} + p^{m-2} \text{ or } - p^{\frac{m}{2}} + p^{m-2} \text{ if } b \in \operatorname{Im}(L), \operatorname{Tr}(x_b^{p+1}) \neq 0 \text{ and } A \neq 0. \end{cases}$$

where  $A = -\frac{m_p}{4} + \frac{\text{Tr}(x_b)^2}{4\text{Tr}(x_b^{p+1})}$  and  $x_b$  is some solution of the equation  $L(x) = x^{p^2} + x = -b^p$ when  $b \in \text{Im}(L)$ .

Proof By definition, we have

$$\begin{split} N(b) &= p^{-2} \sum_{x \in \mathrm{GF}(q)} \left( \sum_{y \in \mathrm{GF}(p)} \zeta_p^{y \operatorname{Tr}(x^{p+1}-x)} \right) \left( \sum_{z \in \mathrm{GF}(p)} \zeta_p^{z \operatorname{Tr}(bx)} \right) \\ &= p^{-2} \sum_{z \in \mathrm{GF}(p)^*} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\operatorname{Tr}(bzx)} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\operatorname{Tr}(yx^{p+1}-yx)} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sum_{z \in \mathrm{GF}(p)^*} \sum_{z \in \mathrm{GF}(p)^*} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\operatorname{Tr}(yx^{p+1}+(bz-y)x)} + p^{m-2}. \end{split}$$

Note that

$$\sum_{z\in \mathrm{GF}(p)^*}\sum_{x\in \mathrm{GF}(q)}\zeta_p^{\mathrm{Tr}(bzx)}=0.$$

The desired conclusions then follow from Lemmas 7 and 10, and the fact that

$$(-1)^{\frac{p-1}{2}}\sqrt{p^*}^{m+1} = (-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m+1}{2}}.$$

In order to calculate the frequency of each weight in  $C_D$ , we need a few more auxiliary results which are given and proved in the following three lemmas.

**Lemma 12** For any  $a \in GF(p)$ , let

$$N_a = |\{x \in GF(q) : Tr(x^{p+1}) = a\}|.$$

Then

$$N_{a} = \begin{cases} p^{m-1} & \text{if } m \text{ is odd and } a = 0, \\ p^{m-1} + p^{-1}\sqrt{p^{*}}^{m+1}\bar{\eta}(-a) & \text{if } m \text{ is odd and } a \neq 0, \\ p^{m-1} - p^{\frac{m}{2}-1}(p-1) & \text{if } m \equiv 2( \mod 4) \text{ and } a = 0, \\ p^{m-1} + p^{\frac{m}{2}-1} & \text{if } m \equiv 2( \mod 4) \text{ and } a \neq 0, \\ p^{m-1} - p^{\frac{m}{2}}(p-1) & \text{if } m \equiv 0( \mod 4) \text{ and } a = 0, \\ p^{m-1} + p^{\frac{m}{2}} & \text{if } m \equiv 0( \mod 4) \text{ and } a \neq 0. \end{cases}$$

Proof By definition, Lemmas 4 and 5, we have

$$\begin{split} N_{a} &= \frac{1}{p} \sum_{x \in \mathrm{GF}(q)} \sum_{y \in \mathrm{GF}(p)} \zeta_{p}^{\mathrm{v}(\mathrm{Tr}(x^{p+1})-a)} \\ &= p^{m-1} + \frac{1}{p} \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} \sum_{x \in \mathrm{GF}(q)} \zeta_{p}^{\mathrm{Tr}(yx^{p+1})} \\ &= p^{m-1} + \frac{1}{p} \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} S(y, 0) \\ &= \begin{cases} p^{m-1} + p^{-1} \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} \sqrt{p^{*m}} \eta(y) & \text{if } m \text{ is odd} \\ p^{m-1} + p^{-1} \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} (-p^{\frac{m}{2}}) & \text{if } m \equiv 2( \mod 4) \\ p^{m-1} + p^{-1} \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} (-p^{\frac{m}{2}+1}) & \text{if } m \equiv 0( \mod 4) \end{cases} \\ &= \begin{cases} p^{m-1} + p^{-1} \sqrt{p^{*m}} \sum_{y \in \mathrm{GF}(p)^{*}} \overline{\eta}(y) & \text{if } m \text{ is odd and } a = 0 \\ p^{m-1} + p^{-1} \sqrt{p^{*m}} \overline{\eta}(-a) \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} \overline{\eta}(-ya) & \text{if } m \text{ is odd and } a \neq 0 \\ p^{m-1} + p^{-1} (-p^{\frac{m}{2}}) \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} & \text{if } m \equiv 2( \mod 4) \\ p^{m-1} + p^{-1} \sqrt{p^{*m}} \overline{\eta}(-a) G(\overline{\eta}, \overline{\chi}_{1}) & \text{if } m \text{ is odd and } a = 0, \\ p^{m-1} + p^{-1} \sqrt{p^{*m}} \overline{\eta}(-a) G(\overline{\eta}, \overline{\chi}_{1}) & \text{if } m \text{ is odd and } a \neq 0, \\ p^{m-1} - p^{\frac{m}{2}} \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} & \text{if } m \equiv 2( \mod 4), \\ p^{m-1} - p^{\frac{m}{2}} \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} & \text{if } m \equiv 2( \mod 4), \\ p^{m-1} - p^{\frac{m}{2}} \sum_{y \in \mathrm{GF}(p)^{*}} \zeta_{p}^{-ya} & \text{if } m \equiv 0( \mod 4). \end{cases}$$

Note that  $\sum_{y \in GF(p)^*} \bar{\eta}(y) = 0$  and

$$\sum_{y \in GF(p)^*} \zeta_p^{-ya} = \begin{cases} p-1 & \text{if } a = 0, \\ -1 & \text{if } a \neq 0, \end{cases}$$

The desired conclusion then follows from Lemma 1.

**Lemma 13** For any  $a \in GF(p)$ , let

$$N_{(a,0)} = |\{x \in GF(q) : Tr(x^{p+1}) = a \text{ and } Tr(x) = 0\}|$$

Then, for  $m_p = 0$ , we have

$$N_{(a,0)} = \begin{cases} p^{m-2} & \text{if } m \text{ is odd } and \ a = 0, \\ p^{m-2} + p^{-1}\sqrt{p^*}^{m+1}\bar{\eta}(-a) & \text{if } m \text{ is odd } and \ a \neq 0, \\ p^{m-2} - p^{\frac{m}{2}-1}(p-1) & \text{if } m \equiv 2( \mod 4) \text{ and } a = 0, \\ p^{m-2} + p^{\frac{m}{2}-1} & \text{if } m \equiv 2( \mod 4) \text{ and } a \neq 0, \\ p^{m-2} - p^{\frac{m}{2}}(p-1) & \text{if } m \equiv 0( \mod 4) \text{ and } a \neq 0, \\ p^{m-2} + p^{\frac{m}{2}} & \text{if } m \equiv 0( \mod 4) \text{ and } a = 0, \\ p^{m-2} + p^{\frac{m}{2}} & \text{if } m \equiv 0( \mod 4) \text{ and } a \neq 0. \end{cases}$$

and for  $m_p \neq 0$ , we have

$$N_{(a,0)} = \begin{cases} p^{m-2} + p^{-2}\bar{\eta}(-m_p)\sqrt{p^{*}}^{m+1}(p-1) & \text{if } m \text{ is odd and } a = 0, \\ p^{m-2} - p^{-2}\bar{\eta}(-m_p)\sqrt{p^{*}}^{m+1} & \text{if } m \text{ is odd and } a \neq 0, \\ p^{m-2} & \text{if } m \text{ is even and } a = 0, \\ p^{m-2} - \bar{\eta}(m_p)\bar{\eta}(a)(-1)^{\frac{p-1}{2}}p^{\frac{m}{2}-1} & \text{if } m \equiv 2( \mod 4) \text{ and } a \neq 0, \\ p^{m-2} - \bar{\eta}(m_p)\bar{\eta}(a)(-1)^{\frac{p-1}{2}}p^{\frac{m}{2}} & \text{if } m \equiv 0( \mod 4) \text{ and } a \neq 0. \end{cases}$$

Proof By definition, we have

$$\begin{split} N_{(a,0)} &= p^{-2} \sum_{x \in \mathrm{GF}(q)} \left( \sum_{y \in \mathrm{GF}(p)} \zeta_p^{y(\mathrm{Tr}(x^{p+1})-a)} \right) \left( \sum_{z \in \mathrm{GF}(p)} \zeta_p^{z\mathrm{Tr}(x)} \right) \\ &= p^{-2} \sum_{z \in \mathrm{GF}(p)} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\mathrm{Tr}(zx)} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sum_{z \in \mathrm{GF}(p)} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\mathrm{Tr}(yx^{p+1}+zx-ya)} \\ &= p^{-2} (q + \sum_{z \in \mathrm{GF}(p)^*} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\mathrm{Tr}(zx)}) + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \left( \zeta_p^{-ya} \sum_{z \in \mathrm{GF}(p)} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{y(\mathrm{Tr}(x^{p+1}+zx))} \right). \end{split}$$

Note that

$$\sum_{z \in \mathrm{GF}(p)^*} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\mathrm{Tr}(zx)} = 0.$$

Therefore,

$$N_{(a,0)} = p^{m-2} + p^{-2} \sum_{y \in GF(p)^*} \left( \zeta_p^{-ya} \sum_{z \in GF(p)} \sum_{x \in GF(q)} \zeta_p^{y(\operatorname{Tr}(x^{p+1} + zx))} \right)$$
  
=  $p^{m-2} + p^{-2} \sum_{y \in GF(p)^*} \sigma_y \left( \zeta_p^{-a} \sum_{z \in GF(p)} \sum_{x \in GF(q)} \zeta_p^{\operatorname{Tr}(x^{p+1} + zx)} \right).$  (9)

It is clear that  $-\frac{1}{2}z$  is the solution of the equation  $x^{p^2} + x = -z^p$  for any  $z \in GF(p)$ . Hence, by Lemmas 4 and 5, Equation (9) becomes

$$\begin{split} N_{(a,0)} &= p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y \left( \zeta_p^{-a} \sum_{z \in \mathrm{GF}(p)} \sum_{x \in \mathrm{GF}(q)} \zeta_p^{\mathrm{Tr}(x^{p+1} + zx)} \right) \\ &= \begin{cases} p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y (\zeta_p^{-a} \sum_{z \in \mathrm{GF}(p)} \sqrt{p^{*m}} \zeta_p^{-\frac{m_p}{4} z^2}) & \text{if } m \text{ is odd,} \\ p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y (\zeta_p^{-a} \sum_{z \in \mathrm{GF}(p)} - p^{\frac{m}{2}} \zeta_p^{-\frac{m_p}{4} z^2}) & \text{if } m \equiv 2( \text{ mod } 4), (10) \\ p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y (\zeta_p^{-a} \sum_{z \in \mathrm{GF}(p)} - p^{\frac{m}{2} + 1} \zeta_p^{-\frac{m_p}{4} z^2}) & \text{if } m \equiv 0( \text{ mod } 4). \end{cases}$$

We distinguish the following two cases.

1. When  $m_p = 0$ , Equation (10) becomes

$$\begin{split} N_{(a,0)} &= \begin{cases} p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-a} \sum_{z \in \mathrm{GF}(p)} \sqrt{p^*}^m) & \text{if } m \text{ is odd,} \\ p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-a} \sum_{z \in \mathrm{GF}(p)} - p^{\frac{m}{2}}) & \text{if } m \equiv 2( \bmod 4), \\ p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-a} \sum_{z \in \mathrm{GF}(p)} - p^{\frac{m}{2}+1}) & \text{if } m \equiv 0( \bmod 4), \end{cases} \\ &= \begin{cases} p^{m-2} + p^{-1} \sqrt{p^*}^{m-1} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-a} \sqrt{p^*}) & \text{if } m \text{ is odd,} \\ p^{m-2} - p^{\frac{m}{2}-1} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-a}) & \text{if } m \equiv 2( \bmod 4), \\ p^{m-2} - p^{\frac{m}{2}} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-a}) & \text{if } m \equiv 0( \bmod 4). \end{cases} \end{split}$$

2. When  $m_p \neq 0$ , from Lemmas 2 and 1, we can easily get

$$\sum_{z\in\mathrm{GF}(p)}\zeta_p^{-\frac{m_p}{4}z^2}=\bar{\eta}(-m_p)\sqrt{p^*}.$$

Thus, Equation (10) becomes

$$\begin{split} N_{(a,0)} &= \begin{cases} p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y \left( \zeta_p^{-a} \bar{\eta}(-m_p) \sqrt{p^{*}}^{m+1} \right) & \text{if } m \text{ is odd,} \\ p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y \left( \zeta_p^{-a}(-p^{\frac{m}{2}}) \bar{\eta}(-m_p) \sqrt{p^*} \right) & \text{if } m \equiv 2 \pmod{4}, \\ p^{m-2} + p^{-2} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y \left( \zeta_p^{-a}(-p^{\frac{m}{2}+1}) \bar{\eta}(-m_p) \sqrt{p^*} \right) & \text{if } m \equiv 0 \pmod{4}. \end{cases} \\ &= \begin{cases} p^{m-2} + p^{-2} \bar{\eta}(-m_p) \sqrt{p^{*m+1}} \sum_{y \in \mathrm{GF}(p)^*} \sigma_y \left( \zeta_p^{-a} \right) & \text{if } m \text{ is odd,} \\ p^{m-2} - p^{\frac{m}{2}-2} \bar{\eta}(-m_p) \sum_{y \in \mathrm{GF}(p)^*} \sigma_y \left( \zeta_p^{-a} \sqrt{p^*} \right) & \text{if } m \equiv 2 \pmod{4}, \\ p^{m-2} - p^{\frac{m}{2}-1} \bar{\eta}(-m_p) \sum_{y \in \mathrm{GF}(p)^*} \sigma_y \left( \zeta_p^{-a} \sqrt{p^*} \right) & \text{if } m \equiv 0 \pmod{4}. \end{cases} \end{split}$$

The desired conclusions then follow from the facts that

$$\sum_{y\in \mathrm{GF}(p)^*} \sigma_y(\zeta_p^{-a}) = \begin{cases} p-1 & \text{if } a=0,\\ -1 & \text{if } a\neq 0, \end{cases}$$

and

$$\sum_{\mathbf{y}\in \mathrm{GF}(p)^*} \sigma_{\mathbf{y}}(\zeta_p^{-a}\sqrt{p^*}) = \begin{cases} 0 & \text{if } a=0 \ ,\\ \left(\frac{-a}{p}\right)p^* & \text{if } a\neq 0. \end{cases}$$

**Lemma 14** Suppose that  $m_p \neq 0$ . Let

$$\bar{N}_0 = |\{x \in \mathrm{GF}(q) : \mathrm{Tr}(x^{p+1}) - \frac{1}{m_p}\mathrm{Tr}(x)^2 = 0\}|.$$

Then

$$\bar{N}_0 = \begin{cases} p^{m-1} + p^{-1}\bar{\eta}(-m_p)\sqrt{p^{*}}^{m+1}(p-1) & \text{if } m \text{ is odd,} \\ p^{m-1} & \text{if } m \text{ is even.} \end{cases}$$

Proof By definition, we have

$$\bar{N}_{0} = \frac{1}{p} \sum_{x \in GF(q)} \sum_{y \in GF(p)} \zeta_{p}^{y(\operatorname{Tr}(x^{p+1}) - \frac{1}{m_{p}}\operatorname{Tr}(x)^{2})} \\
= p^{m-1} + \frac{1}{p} \sum_{y \in GF(p)^{*}} \sum_{x \in GF(q)} \zeta_{p}^{y(\operatorname{Tr}(x^{p+1}) - \frac{1}{m_{p}}\operatorname{Tr}(x)^{2})} \\
= p^{m-1} + \frac{1}{p} \sum_{y \in GF(p)^{*}} \sigma_{y} (\sum_{x \in GF(q)} \zeta_{p}^{\operatorname{Tr}(x^{p+1}) - \frac{1}{m_{p}}\operatorname{Tr}(x)^{2}}).$$
(11)

Since the Fourier expansion of  $\zeta_p^{-\frac{1}{m_p}X^2}$  can be expressed as

$$\zeta_p^{-\frac{1}{m_p}X^2} = \sum_{z \in \mathrm{GF}(p)} a_z \zeta_p^{zX}$$

for any  $X \in GF(p)$ , we have

$$a_{z} = \frac{1}{p} \sum_{X \in \mathrm{GF}(p)} \zeta_{p}^{-\frac{1}{m_{p}}X^{2} - zX}$$
(12)

for any  $z \in \operatorname{GF}(p)$  and

$$\zeta_p^{-\frac{1}{m_p}\operatorname{Tr}(x)^2} = \sum_{z \in \operatorname{GF}(p)} a_z \zeta_p^{z\operatorname{Tr}(x)}.$$

Therefore, Equation (11) becomes

$$\bar{N}_{0} = p^{m-1} + \frac{1}{p} \sum_{y \in GF(p)^{*}} \sigma_{y} \left( \sum_{z \in GF(p)} a_{z} \sum_{x \in GF(q)} \zeta_{p}^{\operatorname{Tr}(x^{p+1}) + z\operatorname{Tr}(x)} \right) \\
= \begin{cases} p^{m-1} + p^{-1} \sum_{y \in GF(p)^{*}} \sigma_{y} \left( \sqrt{p^{*}}^{m} \sum_{z \in GF(p)} a_{z} \zeta_{p}^{-\frac{m_{p}}{4} z^{2}} \right) & \text{if } m \text{ is odd }, \\ p^{m-1} + p^{-1} \sum_{y \in GF(p)^{*}} \sigma_{y} \left( -p^{\frac{m}{2}} \sum_{z \in GF(p)} a_{z} \zeta_{p}^{-\frac{m_{p}}{4} z^{2}} \right) & \text{if } m \equiv 2 \pmod{4}, \\ p^{m-1} + p^{-1} \sum_{y \in GF(p)^{*}} \sigma_{y} \left( -p^{\frac{m+1}{2}} \sum_{z \in GF(p)} a_{z} \zeta_{p}^{-\frac{m_{p}}{4} z^{2}} \right) & \text{if } m \equiv 0 \pmod{4}. \end{cases}$$
(13)

However, by Equation (12) and Lemma 2 we obtain

$$\sum_{z \in GF(p)} a_z \zeta_p^{-\frac{m_p}{4}z^2} = \sum_{z \in GF(p)} \frac{1}{p} \sum_{X \in GF(p)} \zeta_p^{-\frac{1}{m_p}X^2 - zX - \frac{m_p}{4}z^2}$$
$$= \sum_{z \in GF(p)} \frac{1}{p} \bar{\eta}(-\frac{1}{m_p}) G(\bar{\eta}, \bar{\chi}_1)$$
$$= \bar{\eta}(-m_p) \sqrt{p^*}.$$

Hence, Equation (13) becomes

The desired conclusions then follow from the fact that

$$\begin{cases} \sigma_{y}(1) = 1, \\ \sum_{y \in \mathrm{GF}(p)^{*}} \sigma_{y}(\sqrt{p^{*}}) = 0. \end{cases}$$

4.2 The proof of Theorems 1, 2 and 3

By definition, the code length of  $C_D$  is  $n = |D| = n_0 - 1$ , where  $n_0$  was defined in Lemma 8. This means that Equations (6) and (7) follow.

For each  $b \in GF(q)^*$ , define

$$\mathbf{c}_b = (\mathrm{Tr}(bd_1), \mathrm{Tr}(bd_2), \dots, \mathrm{Tr}(bd_n)), \tag{14}$$

where  $d_1, d_2, \ldots, d_n$  are the elements of *D*. Then the Hamming weight wt( $\mathbf{c}_b$ ) of  $\mathbf{c}_b$  is

$$\operatorname{wt}(\mathbf{c}_b) = n_0 - N(b), \tag{15}$$

where  $n_0$  and N(b) were defined before. By Lemmas 8 and 11, we have wt( $\mathbf{c}_b$ ) =  $n_0 - N(b) > 0$  for each  $b \in GF(q)^*$ . This means that the code  $C_D$  has q distinct codewords. Hence, the dimension of the code  $C_D$  is m.

Next we shall prove the multiplicities  $A_{w_i}$  of the codewords with weight  $w_i$  in  $C_D$ , and will distinguish the following three cases.

1. The case that *m* is odd.

It follows from Lemma 4 that  $L(x) = x^{p^2} + x$  is a permutation polynomial over GF(q). Thus the unique solution  $x_b$  of  $L(x) = -b^p$  runs through  $GF(q)^*$  when *b* runs through  $GF(q)^*$ .

- When  $m_p = 0$ , the desired conclusion of Table 1 follows from Equation (15) and Lemmas 8, 11, 12 and 13.
- When m<sub>p</sub> ≠ 0, we only give the proof for the case η
   (m<sub>p</sub>) = 1 and omit the proof for the case η
   (m<sub>p</sub>) = −1 whose proof is similar. Suppose that m<sub>2</sub> = 1 and η
   (m<sub>p</sub>) = 1, then from Lemmas 8 and 11 we obtain

$$\begin{aligned} \mathrm{wt}(\mathbf{c}_b) &= n_0 - N(b) \\ &= \begin{cases} B_1 & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) = 0, \\ B_1 + Bp & \text{if } \operatorname{Tr}(x_b^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_b) \neq 0, \\ B_1 + B(p+1) & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) = 0, \\ B_1 + B(p-1) & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{NSQ}, \\ B_1 + B \text{ or } B_1 + B(p+1) & \text{if } \operatorname{Tr}(x_b^{p+1}) \in \operatorname{SQ} \text{ and } \operatorname{Tr}(x_b) \neq 0, \end{cases}$$

where  $B_1 = p^{m-2}(p-1)$  and  $B = (-1)^{\frac{p-1}{2}\frac{m-1}{2}}p^{\frac{m-3}{2}}$ . Therefore, the weight wt( $\mathbf{c}_b$ ) satisfies

$$wt(\mathbf{c}_b) \in \{B_1, B_1 + Bp, B_1 + B(p+1), B_1 + B(p-1), B_1 + B\}$$

for each  $b \in GF(q)^*$ . Define

$$w_1 = B_1,$$
  

$$w_2 = B_1 + B_p,$$
  

$$w_3 = B_1 + B(p+1),$$
  

$$w_4 = B_1 + B(p-1),$$
  

$$w_5 = B_1 + B.$$

We now determine the number  $A_{w_i}$  of the codewords with weight  $w_i$  in  $C_D$ . By Lemmas 12 and 13, we can directly determine

$$\begin{cases} A_{w_1} = N_{(0,0)} - 1 = p^{m-2} - 1 + (p-1)B, \\ A_{w_2} = (N_0 - 1) - (N_{(0,0)} - 1) = B_1 - (p-1)B, \\ A_{w_4} = \frac{p-1}{2}N_a = \frac{p-1}{2}(p^{m-1} - Bp), \end{cases}$$
(16)

where *a* is a nonsquare in  $GF(p)^*$ . Since  $0 \notin D$ , the minimum distance  $d^{\perp}$  of the dual code  $C_D^{\perp}$  of  $C_D$  cannot be 1. This means that the minimum weight of the dual code  $C_D^{\perp}$  is at least 2. The first two Pless Power Moments [16, p.260] lead to the following system of equations:

$$\begin{cases} A_{w_1} + A_{w_2} + A_{w_3} + A_{w_4} + A_{w_5} = p^m - 1, \\ w_1 A_{w_1} + w_2 A_{w_2} + w_3 A_{w_3} + w_4 A_{w_4} + w_5 A_{w_5} = (p - 1)np^{m-1}, \end{cases}$$
(17)

where the code length  $n = n_0 - 1 = p^{m-1} - 1 + pB$ . Solving the system of equations in (17) yields

$$\begin{cases} A_{w_3} = \frac{1}{2}(p-1)(p-2)p^{\frac{m-3}{2}}(p^{\frac{m-1}{2}} - (-1)^{\frac{p-1}{2}\frac{m-1}{2}}), \\ A_{w_5} = (p-1)p^{m-2} + (-1)^{\frac{p-1}{2}\frac{m-1}{2}}(p-1)^2p^{\frac{m-3}{2}}. \end{cases}$$

This completes the proof of the weight distribution of Table 2.

# 2. The case that $m \equiv 2 \pmod{4}$ .

It follows from Lemma 5 that  $L(x) = x^{p^2} + x$  is a permutation polynomial over GF(q). Similar to the analysis for the case *m* being odd, we have the following.

- When  $m_p = 0$ , it is clear that the desired conclusion of Table 3 follows from Equation (15) and Lemmas 8, 11, 12 and 13.
- When  $m_p \neq 0$ , by Lemmas 8 and 11 we have

$$\begin{aligned} \operatorname{wt}(\mathbf{c}_{b}) &= n_{0} - N(b) \\ &= \begin{cases} (p-1)p^{m-2} & \text{if } \operatorname{Tr}(x_{b}^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_{b}) = 0, \\ (p-1)p^{m-2} + p^{\frac{m}{2}-1} & \text{if } \operatorname{Tr}(x_{b}^{p+1}) = 0 \text{ and } \operatorname{Tr}(x_{b}) \neq 0, \\ \text{or } \operatorname{Tr}(x_{b}^{p+1}) \neq 0 \text{ and } A = 0, \\ (p-1)p^{m-2} \text{ or } (p-1)p^{m-2} + 2p^{\frac{m}{2}-1} & \text{if } \operatorname{Tr}(x_{b}^{p+1}) \neq 0 \text{ and } A \neq 0, \end{cases} \end{aligned}$$

where  $A = -\frac{m_p}{4} + \frac{\text{Tr}(x_b)^2}{4\text{Tr}(x_b^{p+1})}$ . Therefore, the weight wt( $\mathbf{c}_b$ ) satisfies

wt(
$$\mathbf{c}_b$$
)  $\in \{(p-1)p^{m-2}, (p-1)p^{m-2} + p^{\frac{m}{2}-1}, (p-1)p^{m-2} + 2p^{\frac{m}{2}-1}\}$ 

for each  $b \in GF(q)^*$ . Define

$$\begin{split} & w_1 = (p-1)p^{m-2}, \\ & w_2 = (p-1)p^{m-2} + p^{\frac{m}{2}-1}, \\ & w_3 = (p-1)p^{m-2} + 2p^{\frac{m}{2}-1}. \end{split}$$

We now determine the number  $A_{w_i}$  of the codewords with weight  $w_i$  in  $C_D$ . Note that

$$\begin{split} &|\{b \in \mathrm{GF}(q) : \mathrm{Tr}(x_b^{p+1}) \neq 0 \text{ and } A = 0\}| \\ &= |\{b \in \mathrm{GF}(q) : A = 0\}| - |\{b \in \mathrm{GF}(q) : \mathrm{Tr}(x_b^{p+1}) = 0 \text{ and } A \neq 0\}| \\ &= |\{x \in \mathrm{GF}(q) : \mathrm{Tr}(x^{p+1}) - \frac{1}{m_p} \mathrm{Tr}(x)^2 = 0\}| - |\{x \in \mathrm{GF}(q) : \mathrm{Tr}(x^{p+1}) = 0 \text{ and } \mathrm{Tr}(x) = 0\}|, \end{split}$$

as  $L(x) = x^{p^2} + x$  is a permutation polynomial over GF(q) for this case  $m \equiv 2 \pmod{4}$ . Thus from Equation (18) and Lemmas 12, 13 and 14, we can directly obtain

$$A_{w_2} = (N_0 - N_{(0,0)}) + (\bar{N}_0 - N_{(0,0)}) = (p-1)(2p^{m-2} - p^{\frac{m}{2}-1}).$$

Since the minimum weight of the dual code  $C_D^{\perp}$  is at least 2, from the first two Pless Power Moments [16, p.260] we can compute  $A_{w_1}$  and  $A_{w_3}$ . This completes the proof of the weight distribution of Table 4.

- 3. The case that  $m \equiv 0 \pmod{4}$ .
  - If  $m_p = 0$ , then by Lemmas 8 and 11 we have

$$\begin{split} \mathrm{wt}(\mathbf{c}_b) &= n_0 - N(b) \\ &= \begin{cases} B_1 - (p-1)^2 p^{\frac{m}{2}-1} & \text{if } b \notin \mathrm{Im}(L), \\ B_1 & \text{if } b \in \mathrm{Im}(L), \, \mathrm{Tr}(x_b^{p+1}) = 0 \text{ and } \mathrm{Tr}(x_b) = 0, \\ B_1 - (p-1) p^{\frac{m}{2}} & \text{if } b \in \mathrm{Im}(L), \, \mathrm{Tr}(x_b^{p+1}) = 0 \text{ and } \mathrm{Tr}(x_b) \neq 0 \\ & \text{or } \mathrm{Tr}(x_b^{p+1}) \neq 0 \text{ and } \mathrm{Tr}(x_b) = 0, \\ B_1 - (p-2) p^{\frac{m}{2}} & \text{if } b \in \mathrm{Im}(L), \, \mathrm{Tr}(x_b^{p+1}) \neq 0 \text{ and } \mathrm{Tr}(x_b) \neq 0, \end{split}$$

where  $B_1 = p^{m-2}(p-1)$ .

Recall that  $L(x) = x^{p^2} + x$ . If  $b \in \text{Im}(L)$ , then this means that  $L(x) = -b^p$  is solvable. We point out the following facts:

- If  $m \equiv 0 \pmod{4}$ , then L(x) is not a permutation polynomial over GF(q). However, for any  $b \in GF(q)$ , if  $L(x) = -b^p$  is solvable in this case  $m \equiv 0 \pmod{4}$ , then it has  $p^2$  solutions.
- If  $m \equiv 0 \pmod{4}$ , then the number of *b* is  $p^{m-2}$  such that  $L(x) = -b^p$  is solvable when *b* runs through GF(q).

Therefore, we have

$$\operatorname{wt}(\mathbf{c}_{b}) = \begin{cases} p^{m-2}(p-1) - (p-1)^{2} p^{\frac{m}{2}-1} & \operatorname{occurs} p^{m} - p^{m-2} \text{ times,} \\ p^{m-2}(p-1) & \operatorname{occurs} p^{-2} N_{(0,0)} - 1 \text{ times,} \\ (p-1) p^{\frac{m}{2}} (p^{\frac{m}{2}-2} - 1) & \operatorname{occurs} (N_{0} + p^{m-1} - 2N_{(0,0)}) p^{-2} \text{ times,} \\ (p-1) p^{m-2} - (p-2) p^{\frac{m}{2}} & \operatorname{occurs} (p^{m} - p^{m-1} - N_{0} + N_{(0,0)}) p^{-2} \text{ times,} \end{cases}$$

when *b* runs through  $GF(q)^*$ , where  $N_0$  and  $N_{(0,0)}$  were defined before. The desired conclusion of Table 5 then follows from Lemmas 12 and 13.

• If  $m_p \neq 0$ , then from Lemmas 8 and 11 we get

$$\begin{split} \mathrm{wt}(\mathbf{c}_b) &= n_0 - N(b) \\ &= \begin{cases} (p-1)p^{\frac{m}{2}-1} + B_1 & \text{if } b \notin \mathrm{Im}(L), \\ B_1 & \text{if } b \in \mathrm{Im}(L), \, \mathrm{Tr}(x_b^{p+1}) = 0 \text{ and } \mathrm{Tr}(x_b) = 0, \\ B_1 + p^{\frac{m}{2}} & \text{if } b \in \mathrm{Im}(L), \, \mathrm{Tr}(x_b^{p+1}) = 0 \text{ and } \mathrm{Tr}(x_b) \neq 0 \ (19) \\ & \text{or } \mathrm{Tr}(x_b^{p+1}) \neq 0 \text{ and } A = 0, \\ B_1 \text{ or } B_1 + 2p^{\frac{m}{2}} & \text{if } b \in \mathrm{Im}(L), \, \mathrm{Tr}(x_b^{p+1}) \neq 0 \text{ and } A \neq 0, \end{cases}$$

,

where  $B_1 = p^{m-2}(p-1)$ . Therefore, the weight wt( $\mathbf{c}_b$ ) satisfies

wt(
$$\mathbf{c}_b$$
)  $\in \{(p-1)p^{\frac{m}{2}-1} + B_1, B_1, B_1 + p^{\frac{m}{2}}, B_1 + 2p^{\frac{m}{2}}\}$ 

for each  $b \in GF(q)^*$ . Define

$$w_1 = (p-1)p^{\frac{m}{2}-1} + B_1$$
  

$$w_2 = B_1,$$
  

$$w_3 = B_1 + p^{\frac{m}{2}},$$
  

$$w_4 = B_1 + 2p^{\frac{m}{2}}.$$

We now determine the number  $A_{w_i}$  of the codewords with weight  $w_i$  in  $C_D$ . It is clear that  $A_{w_i} = p^m - p^{m-2}.$ 

$$\begin{array}{l} A_{w_3} = p^{-2} (N_0 - N_{(0,0)} + \bar{N}_0), \\ = (p-1) (2 p^{m-4} - p^{\frac{m}{2}-2}). \end{array}$$

Since the minimum weight of the dual code  $C_D^{\perp}$  is at least 2, the first two Pless Power Moments [16, p.260] lead to the following system of equations:

$$\begin{cases} A_{w_1} + A_{w_2} + A_{w_3} + A_{w_4} = p^m - 1, \\ w_1 A_{w_1} + w_2 A_{w_2} + w_3 A_{w_3} + w_4 A_{w_4} = (p-1)np^{m-1}, \end{cases}$$
(20)

where  $n = n_0 - 1 = p^{m-1} + p^{\frac{m}{2}} - 1$ . Solving the system of equations in (20) gives

$$\begin{cases} A_{w_2} = p^{m-4} + \frac{p-1}{2}(p^{\frac{m}{2}-1} + p^{m-3}) - 1, \\ A_{w_4} = \frac{1}{2}(p-1)(p-2)(p^{m-4} - p^{\frac{m}{2}-2}). \end{cases}$$

This completes the proof of the weight distribution of Table 6.

Summarizing all the conclusions above completes the proofs of Theorems 1, 2 and 3.

# 5 Concluding remarks

In this paper, we presented a class of linear codes  $C_D$  with a few weights and completely determined their weight distributions. The result showed that they have at most five weights. Particularly, the codes  $C_D$  presented in this paper have three weights when  $m \ge 6$  and  $m \equiv 2 \pmod{4}$ . Many classes of linear codes with a few weights were constructed (see, for example, [14, 29, 23, 2, 15, 18, 4, 28, 12]). When *m* is odd, the code  $C_D$  presented in this paper has the same parameter and weight distribution as that of  $C_D$  presented in [19]. But they have different defining set *D*.

Any linear code over GF(p) can be used to obtain secret sharing schemes [4,27]. In order to obtain secret sharing schemes with interesting access structures, one would like to have linear codes C such that  $w_{\min}/w_{\max} > \frac{p-1}{p}$  [27], where  $w_{\min}$  and  $w_{\max}$  denote the minimum and maximum nonzero weight of the linear code.

When  $m \ge 5$  is odd and  $m_p = 0$ , the code  $C_D$  of Section 3 satisfies that

$$\frac{w_{\min}}{w_{\max}} = \frac{p^{m-2} - p^{(m-3)/2}}{p^{m-2} + p^{(m-3)/2}} > \frac{p-1}{p}.$$

When  $m \ge 5$  is odd and  $m_p \ne 0$ , the code  $C_D$  of Section 3 satisfies that

$$\frac{w_{\min}}{w_{\max}} = \frac{(p-1)p^{m-2} - (p+1)p^{(m-3)/2}}{(p-1)p^{m-2} + (p+1)p^{(m-3)/2}} > \frac{p-1}{p}.$$

When  $m \ge 6$ ,  $m \equiv 2 \pmod{4}$  and  $m_p = 0$ , the code  $\mathcal{C}_D$  of Section 3 satisfies that

$$\frac{w_{\min}}{w_{\max}} = \frac{p^{m-2} - p^{m/2-1}}{p^{m-2}} > \frac{p-1}{p}.$$

When  $m \ge 6$ ,  $m \equiv 2 \pmod{4}$  and  $m_p \ne 0$ , the code  $C_D$  of Section 3 satisfies that

$$\frac{w_{\min}}{w_{\max}} = \frac{(p-1)p^{m-2}}{(p-1)p^{m-2} + 2p^{m/2-1}} > \frac{p-1}{p}.$$

When  $m \ge 6$ ,  $m \equiv 0 \pmod{4}$  and  $m_p = 0$ , the code  $C_D$  of Section 3 satisfies that

$$\frac{w_{\min}}{w_{\max}} = \frac{p^{m-2} - p^{m/2}}{p^{m-2}} > \frac{p-1}{p}.$$

When  $m \ge 6$ ,  $m \equiv 0 \pmod{4}$  and  $m_p \ne 0$ , the code  $\mathcal{C}_D$  of Section 3 satisfies that

$$\frac{w_{\min}}{w_{\max}} = \frac{(p-1)p^{m-2}}{(p-1)p^{m-2} + 2p^{m/2}} > \frac{p-1}{p}.$$

Hence, the linear codes  $C_D$  presented in this paper satisfy the condition that  $w_{\min}/w_{\max} > \frac{p-1}{p}$ , and can be employed to obtain secret sharing schemes with interesting access structures [1,4,14,27].

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