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THEORY OF COMPUTATION

REPORT . NO. 25

ON THE STRUCTURE OF  
FREE FINITE STATE MACHINES

BY

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1. Introduction

As explained by Birkhoff and Lipson in [1], a finite state machine  $M$  (without outputs) can be considered as an algebra with two "phyla":

$S$  = set of states,  $I$  = input alphabet

and a single operator:  $T : S \times I \rightarrow S$ , the transition function of  $M$ .

Given  $M = (S, I)$  and a pair of integers  $(m, n)$  there is an associated machine  $U_{m,n}(M)$  freely generated as an algebra by states  $t_1, \dots, t_m$  and input symbols  $e_1, \dots, e_n$  subject to the relations which hold within  $M$ . Explicitly  $U_{m,n}(M) = (\mathcal{A}, \mathcal{G})$  where

$$\mathcal{G} = \{e_1, \dots, e_n\}$$

and each state in  $\mathcal{A}$  consists of an equivalence class of expressions of the form

$$t_i w(e_1, \dots, e_n) \text{ where } 1 \leq i \leq m, w \in \mathcal{G}^*$$

and  $t_i w(e_1, \dots, e_n)$  and  $t_j v(e_1, \dots, e_n)$  are equivalent if for all pairs of maps  $\{t_1, \dots, t_m\} \xrightarrow{f} S$  and  $\{e_1, \dots, e_n\} \xrightarrow{g} I$  the relation:

$$f(t_i) w(g(e_1), \dots, g(e_n)) = f(t_j) v(g(e_1), \dots, g(e_n))$$

holds in  $M$ . The transition function then maps  $(t_i w(e_1, \dots, e_n), e_i)$  to  $t_i (w(e_1, \dots, e_n) e_i)$ .

Definition Using the notation introduced above, it will be convenient to refer to a pair of maps  $f : \{t_1, \dots, t_m\} \rightarrow S$  and  $g : \{e_1, \dots, e_n\} \rightarrow I$  as a phyla-preserving mapping from  $\{t_1, \dots, t_m, e_1, \dots, e_n\}$  to  $M$  or an interpretation of  $\{t_1, \dots, t_m, e_1, \dots, e_n\}$  in  $M$ .

The proof of the following theorem is to be found in [1].

Theorem (i)  $U_{m,n}(M)$  is a finite state machine.

(ii)  $U_{m,n}(M)$  is generated by the  $m$  states  $t_1, \dots, t_m$  and  $n$  input symbols  $e_1, \dots, e_n$ .

(iii) If  $\pi$  denotes the canonical phyla-preserving map from the set  $\{t_1, \dots, t_m, e_1, \dots, e_n\}$  to  $U_{m,n}(M)$ , and  $\theta$  is any phyla-preserving map from  $\{t_1, \dots, t_m, e_1, \dots, e_n\}$  to  $M$ , then there is an unique algebra homomorphism  $\phi : U_{m,n}(M) \rightarrow M$  such that  $\theta = \phi\pi$ .

(iv)  $U_{m,n}(M)$  is an epimorphic image of any other finite state machine having property (iii).

Definition If  $t_i w(e_1, \dots, e_n)$  and  $t_j v(e_1, \dots, e_n)$  are equivalent in  $U_{m,n}(M)$ , then

$$t_i w(e_1, \dots, e_n) = t_j v(e_1, \dots, e_n)$$

is a universal relation in  $M$ .

## 2. The case $m \geq 1$

Theorem 1: For a machine  $M = (S, I)$  to have a universal relation of the form

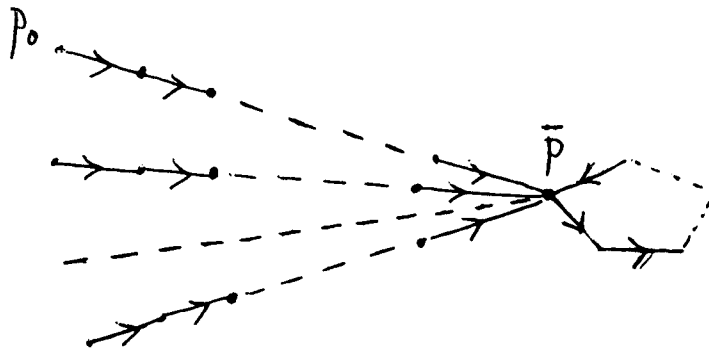
$$t_i w(e_1, \dots, e_n) = t_j v(e_1, \dots, e_n) \text{ with } i \neq j$$

it is necessary and sufficient that for each input  $\alpha$  in  $I$  there should exist a state  $t(\alpha)$  such that (a)  $t(\alpha) \cdot \alpha = t(\alpha)$

and (b) for each  $s$  in  $S$  there is a non-negative integer  $r(s)$  such that  $s \cdot \alpha^{r(s)} = t(\alpha)$

Proof: Suppose that  $M$  has a universal relation  $U$  of the form

$t_i w(e_1, \dots, e_n) = t_j v(e_1, \dots, e_n)$  for  $i \neq j$ . If  $\alpha \in I$ , there is a submachine  $M_\alpha = (S, \alpha^*)$  of  $M$ , which is a disjoint union of  $k$  machines of the following type:



Since  $U$  holds under all interpretations  $(f, g)$  for which  $g(e_i) = \alpha$  for  $1 \leq i \leq n$ , it is clear that  $k = 1$ . Moreover, taking interpretations  $(f, g)$  such that  $f(t_i) = p_0$ ,  $f(t_j) = p_0 \alpha^c$  for some non-negative integer  $c$  and  $g(e_i) = \alpha$  for  $1 \leq i \leq n$ , it follows that

$$p_0 \alpha^{\ell(w)} = p_0 \alpha^{\ell(v)+c}$$

in  $M$  for  $c = 0, 1, 2, \dots$ . This establishes that  $\overline{p\alpha} = \overline{p}$ , so that conditions (a) and (b) are satisfied with  $t(\alpha) = \overline{p}$ .

For the converse, suppose that given input  $\alpha$  in  $I$ , there is a  $t(\alpha)$  for which conditions (a) and (b) hold. Then let  $r(\alpha) = \max_{s \in S} r(s)$ , and  $r = \max_{\alpha \in I} r(\alpha)$ . It is clear that the relation  $s\alpha^r = t\alpha^r$  holds for all  $s, t$  in  $S$  and all  $\alpha$  in  $I$ ; that is, the relation  $t_1 x_1^r = t_2 x_1^r$  holds universally in  $M$ .

### Corollary to Theorem 1:

Unless a relation of the form  $t_1 x_1^r = t_2 x_1^r$  holds universally in the machine  $M$ , the finite state machine  $U_{m,n}(M)$  is (up to isomorphism)  $m$  disjoint copies of  $U_{1,n}(M)$ .

### 3. Structure of $U_{1,n}(M)$

Definition Let  $K$  be a finite monoid generated by elements  $x_1, \dots, x_n$ .

The machine  $\mathbf{m}(K, X)$  associated with the monoid  $K$  generated by  $X$  has a set of states  $K$ , input alphabet  $X = \{x_1, \dots, x_n\}$  and transition function  $K \times X \rightarrow K$  defined by multiplication in  $K$ . The machine  $\mathbf{m}(K, X)$  will be called a monoid machine. If  $K$  is a group, then  $\mathbf{m}(K, X)$  is a group machine or Cayley diagram.

Theorem 2: (i) If  $M$  is a finite state machine, then, for  $n \geq 1$ ,  $U_{1,n}(M)$  is isomorphic to the monoid machine  $\mathcal{M}(K, X)$ , where  $K$  is the monoid freely generated by  $X = \{x_1, \dots, x_n\}$  subject to the relations:

$$w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$$

where  $tw(e_1, \dots, e_n) = tv(e_1, \dots, e_n)$

is a universal relation in  $M$ .

(ii) Let  $K$  be a finite monoid generated by  $X = \{x_1, \dots, x_n\}$

For  $\mathcal{M}(K, X)$  to be isomorphic with  $U_{1,n}(M)$  for some finite state machine  $M$ , it is necessary and sufficient that for each relation  $w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$  in  $K$  and each map  $f : \{1, 2, \dots, n\}$ , the relation  $w(x_{f(1)}, \dots, x_{f(n)}) = v(x_{f(1)}, \dots, x_{f(n)})$  also holds in  $K$ . If this condition is satisfied then  $U_{1,n}(\mathcal{M}(K, X)) \approx \mathcal{M}(K, X)$ .

(iii) For  $U_{1,n}(M)$  to be a group machine ( $n \geq 1$ ) it is necessary and sufficient that for some non-trivial  $w$  in  $\mathcal{G}^*$ , a relation of the form:

$$tw(e_1, \dots, e_n) = t$$

holds universally in  $M$ .

Definition When the necessary and sufficient conditions (stated in (ii) above) for  $\mathcal{M}(K, X)$  to be isomorphic with  $U_{1,n}(M)$  for a finite state machine  $M$  are satisfied,  $X$  is said to generate  $K$  universally or to generate a universal presentation of  $K$ .

Proof: (i) The elements of  $U_{1,n}$  are equivalence classes of expressions of the form:

$$tw(e_1, \dots, e_n)$$

where  $tw(e_1, \dots, e_n)$  and  $tv(e_1, \dots, e_n)$  are equivalent if

$tw(e_1, \dots, e_n) = tv(e_1, \dots, e_n)$  is a universal relation in  $M$ , with transition function defined by

$$(tw(e_1, \dots, e_n), e_i) \mapsto t(w(e_1, \dots, e_n)e_i)$$

The map  $tw(e_1, \dots, e_n) \mapsto w(x_1, \dots, x_n)$  then clearly induces an isomorphism  $U_{1,n}(M) \simeq K$ .

(ii) Suppose  $\mathcal{M}(K, X) \simeq U_{1,n}(M)$ . Then if the relation  $w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$  holds in  $K$  then  $tw(e_1, \dots, e_n) = tv(e_1, \dots, e_n)$  is a universal relation in  $M$ . Thus given any map  $f : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$ , the relation

$$tw(e_{f(1)}, \dots, e_{f(n)}) = tv(e_{f(1)}, \dots, e_{f(n)})$$

holds universally in  $M$ , whence  $w(x_{f(1)}, \dots, x_{f(n)}) = v(x_{f(1)}, \dots, x_{f(n)})$  in  $K$ .

Conversely, suppose that if  $w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$  in  $K$  and  $f$  is a map  $\{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$ , then  $w(x_{f(1)}, \dots, x_{f(n)}) = v(x_{f(1)}, \dots, x_{f(n)})$ . It follows that the relation  $tw(e_1, \dots, e_n) = tv(e_1, \dots, e_n)$  holds universally in  $\mathcal{M}(K, X)$ . Conversely if  $tw(e_1, \dots, e_n) = tv(e_1, \dots, e_n)$  is a universal relation in  $\mathcal{M}(K, X)$  then certainly  $w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$  in  $K$  (interpreting  $t$  as  $1$ , and  $e_i$  as  $x_i$  for  $i = 1, 2, \dots, n$ ). The isomorphism

$$U_{1,n}(\mathcal{M}(K, X)) \simeq \mathcal{M}(K, X)$$

follows from (i).

(iii) Let  $x_1, \dots, x_n$  generate  $U_{1,n}(M)$  freely subject to the relations:

$$w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$$

where  $tw(e_1, \dots, e_n) = tv(e_1, \dots, e_n)$

is a universal relation in  $M$ . Since  $x_1^r = 1$  for some  $r \geq 1$ , the relation  $te_i^r = t$  must hold universally in  $M$  for some  $r$ .

Conversely, suppose  $tw(e_1, \dots, e_n) = t$  holds universally in  $M$ , with  $w$  non-trivial. Then given  $f : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$  the relation  $w(x_{f(1)}, \dots, x_{f(n)}) = 1$  holds in  $U_{1,n}(M)$ . In particular,  $w(x_i, \dots, x_i) = 1$  for each  $i$ , which proves the existence of  $x_i^{-1}$  for each  $i$ , as  $w$  is non-trivial.

Definition Let  $M = (S, I)$  be a finite state machine, and let  $F(S)$

denote the semigroup of mappings  $S \rightarrow S$  under composition. For each  $\alpha$  in  $I$ , let  $T(\alpha)$  be the map  $S \rightarrow S$  in  $F(S)$ . The map  $T$  extends naturally to a semigroup homomorphism  $I^* \rightarrow F(S)$ . The image of this homomorphism is the syntactic monoid  $\mathcal{S}(M)$  of  $M$ .

Lemma: For each  $n > 1$ ,  $U_{1,n}(M)$  and  $U_{1,n}(\mathcal{M}(\mathcal{S}(M), T(I)))$  are isomorphic.

Proof: Suppose that  $sw(\alpha_1, \dots, \alpha_n) = sv(\alpha_1, \dots, \alpha_n)$  for all  $s$  in  $S$  and all  $\alpha_i$  in  $I$ . Then  $w(T(\alpha_1), \dots, T(\alpha_n))$  and  $v(T(\alpha_1), \dots, T(\alpha_n))$  represent the same element of  $\mathcal{S}(M)$ , so that  $fw(T(\alpha_1), \dots, T(\alpha_n)) = fv(T(\alpha_1), \dots, T(\alpha_n))$  for all  $f$  in  $\mathcal{S}(M)$  and all  $\alpha_i$  in  $I$ .

Conversely, if  $fw(T(\alpha_1), \dots, T(\alpha_n)) = fv(T(\alpha_1), \dots, T(\alpha_n))$  for all  $f$  in  $\mathcal{S}(M)$  and all  $\alpha_i$  in  $I$ , then  $w(T(\alpha_1), \dots, T(\alpha_n)) = v(T(\alpha_1), \dots, T(\alpha_n))$  in  $\mathcal{S}(M)$ . Thus  $sw(\alpha_1, \dots, \alpha_n) = sv(\alpha_1, \dots, \alpha_n)$  in  $M$  for all  $s$  in  $S$  and all  $\alpha_i$  in  $I$ .

This proves the required isomorphism.

It is evident that a universal relation of the form  $tw(e_1, \dots, e_m) = tv(e_1, \dots, e_m)$  holds in a monoid machine  $\mathcal{M}(K, X)$  if and only if  $w(x_1, \dots, x_m) = v(x_1, \dots, x_m)$  for all  $x_i$  in  $X$ . This result will be used in the proof of the next theorem, which describes a simple method for constructing  $U_{1,m}(M)$  when  $M$  is a monoid machine.

Theorem 3: Let  $K$  be a finite monoid generated by  $X = \{x_1, \dots, x_m\}$ .

Let  $X^*$  be the set of rows of the  $n$  by  $m^n$  array whose columns are the elements of  $X^n$ . Then  $X^*$  has  $n$  elements  $X_1, \dots, X_n$ , which generate a submonoid  $K^*$  of  $K^{m^n}$ , and  $U_{1,n}(\mathcal{M}(K, X))$  and  $\mathcal{M}(K^*, X^*)$  are isomorphic.

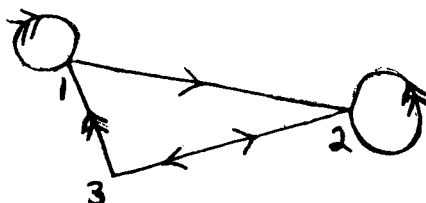
Proof: Suppose that  $w(y_1, \dots, y_n) = v(y_1, \dots, y_n)$  in  $K$  for all interpretations of  $y_1, \dots, y_n$  in  $X$ . Then the identity  $w(y_1, \dots, y_n) = v(y_1, \dots, y_n)$  necessarily holds in  $K^*$ .

Conversely  $w(X_1, \dots, X_n) = v(X_1, \dots, X_n)$  in  $K^*$  entails  $w(y_1, \dots, y_n) = v(y_1, \dots, y_n)$  for all interpretations of  $y_1, \dots, y_n$  in  $X$ , each interpretation corresponding to a projection of the identity  $w(X_1, \dots, X_n) = v(X_1, \dots, X_n)$  onto a single component.

#### 4. Illustrative Examples

##### Example 1:

Let  $M$  be the machine having three states, and input alphabet  $\{a, b\}$ , as indicated below:



(This machine is considered by Birkhoff and Lipson in [1]).

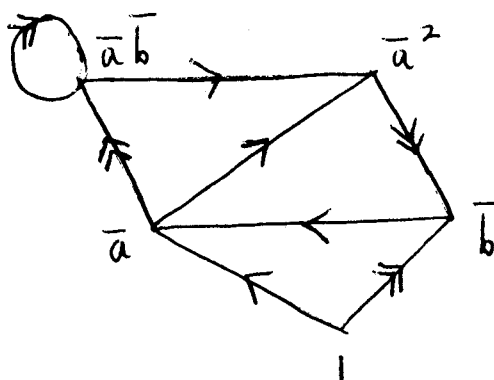
For this machine  $M$ ,  $\Delta(M)$  is the subsemigroup of maps  $\{1, 2, 3\}$  generated by  $a, b$  where

$$\bar{a}(1) = 2, \quad \bar{a}(2) = 3, \quad \bar{a}(3) = 2$$

$$\text{and } \bar{b}(1) = 1, \quad \bar{b}(2) = 2, \quad \bar{b}(3) = 1$$

The syntactic monoid then consists of five maps viz.  $1, \bar{a}, \bar{b}, \bar{a}^2, \bar{a}\bar{b}$ , and the additional relations  $\bar{a}^3 = \bar{a}, \bar{b}^2 = \bar{b}, \bar{b}.\bar{a} = \bar{a}, \bar{a}^2\bar{b} = \bar{b}, \bar{a}\bar{b}\bar{a} = \bar{a}^2$  hold.

The machine  $m = m(\Delta(M), \{\bar{a}, \bar{b}\})$  is:



The free machine  $U_{1,2}^{(M)} \approx U_{1,2}(M)$  is now the semigroup machine associated with the subsemigroup of  $\mathcal{S}(M)^4$  generated by

$$A = (\bar{a}, \bar{a}, \bar{b}, \bar{b}) \text{ and } B = (\bar{a}, \bar{b}, \bar{a}, \bar{b})$$

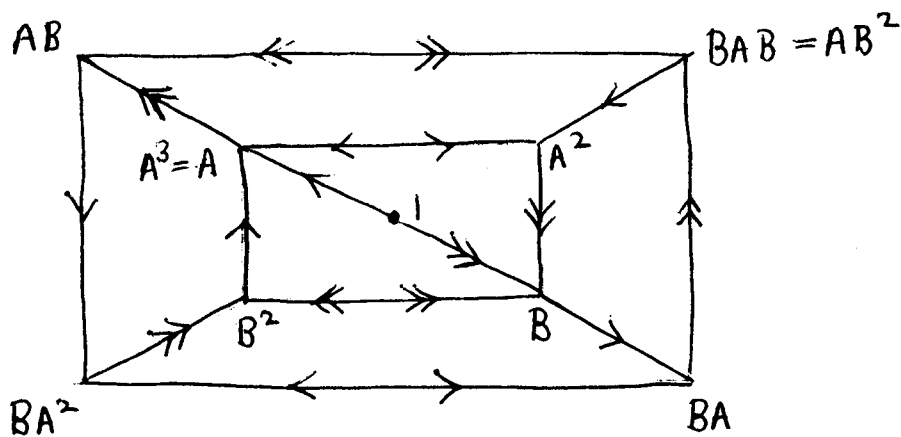
It has 9 elements viz.:

$$1, A, B, A^2 = (\bar{a}^2, \bar{a}^2, \bar{b}, \bar{b}), B^2 = (\bar{a}^2, \bar{b}, \bar{a}^2, \bar{b}),$$

$$AB = (\bar{a}^2, \bar{a}\bar{b}, \bar{a}, \bar{b}), BAB = (\bar{a}, \bar{a}\bar{b}, \bar{a}^2, \bar{b})$$

$$BA = (\bar{a}^2, \bar{a}, \bar{a}\bar{b}, \bar{b}) \text{ and } BA^2 = (\bar{a}, \bar{a}^2, \bar{a}\bar{b}, \bar{b})$$

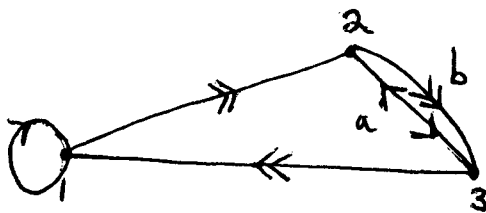
The resulting semigroup machine is:



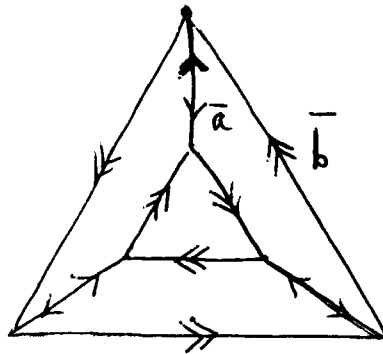
(Note that there is an error in the representation of  $U_{1,n}(M)$  given by Birkhoff and Lipson in [1], and that a similar error occurs in [2]. The relation  $AB^2 = BA^2$  does not hold universally in  $M$  as the diagrams in [1] and [2] suggest).

#### Example 2:

Let  $N$  be the machine with 3 states, and input alphabet  $\{a, b\}$ , as represented below:



In this case,  $\mathcal{A}(N)$  is the subgroup of the semigroup of maps  $\{1,2,3\}$  consisting of all permutations, with generators  $\bar{a} = (23)$ ,  $\bar{b} = (123)$ . The machine  $m(\mathcal{A}(N), \{a,b\})$  is then a Cayley diagram for the symmetric group  $S_3$  viz:



By the previous results,  $U_{1,2}(M) \approx U_{1,2}(m)$  is the group machine associated with the subgroup of  $S_3^4$  generated by  $A = (\bar{a}, \bar{a}, \bar{b}, \bar{b})$  and  $B = (\bar{a}, \bar{b}, \bar{a}, \bar{b})$  generated by  $A = (\bar{a}, \bar{a}, \bar{b}, \bar{b})$  and  $B = (\bar{a}, \bar{b}, \bar{a}, \bar{b})$ . Since  $(AB)^2 = (1, 1, 1, \bar{b})$ , and it can be shown that  $(\bar{a}, \bar{a}, \bar{b})$  and  $(\bar{a}, \bar{b}, \bar{a})$  generate the subgroup of  $C_1 \times S_3^2$  consisting of triples  $(c, p, q)$  such that  $pq$  and  $c$  are permutations of the same parity, it follows that  $U_{1,2}(M)$  is (up to isomorphism) the group machine  $m(G, X)$  where  $G = S_3 \times S_3 \times C_3$  and  $X = \{(\bar{a}, \bar{b}, \bar{b}), (\bar{b}, \bar{a}, \bar{b})\}$ . This result would be difficult to obtain by the direct method for computing  $U_{1,2}(M)$  described in [1].

## 5. Monoids with a universal presentation

It has been shown in the previous section that every free finite state machine is of the form  $m(S, X)$  where  $S$  is a finite monoid, and  $X = \{x_1, \dots, x_n\}$  is a set of generators for  $S$ . In this section, necessary conditions on  $S$  and  $X$  for  $m(S, X)$  to be a free finite state machine are described. For instance, it is evident that if  $X$  generates  $S$  universally then the set of relations between  $x_1, \dots, x_n$  holding in  $S$  is invariant under any permutation of  $\{1, 2, \dots, n\}$ . This fact is relevant for the interpretation of the results and proofs of this section.

Lemma 3: If the elements  $x_1, \dots, x_n$  generate the finite monoid  $S$  universally, then either  $x_1 = x_2 = \dots = x_n$  or  $x_1, x_2, \dots, x_n$  are pairwise distinct and generate  $S$  irredundantly.

Proof: Suppose that  $x_1 = w(x_2, \dots, x_n)$ . Let  $f$  be the map  $\{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$  such that  $f(1) = 2$  and  $f(i) = i$  for  $i > 2$ . Then  $x_{f(1)} = x_2 = w(x_{f(2)}, \dots, x_{f(n)}) = w(x_2, \dots, x_n) = x_1$ , whence  $x_1 = x_2 = \dots = x_n$ .

Notation: Let  $P$  be a partition of  $\{1, 2, \dots, n\}$ . The rank of  $P$  (the number of blocks in the partition  $P$ ) will be denoted by  $\rho(P)$ .

Theorem 4: Let  $S$  be a finite monoid generated universally by distinct generators  $x_1, \dots, x_n$ . For each partition  $P$  let  $E(P)$  be the smallest congruence on  $S$  such that  $x_i$  and  $x_j$  are congruent for all  $(i, j) \in P$ . Then

- (i) the map  $E$  is a join-preserving bijection from the lattice of partitions of  $\{1, 2, \dots, n\}$  (ordered by  $P < Q$  if  $P$  is a refinement of  $Q$ ), to the congruence lattice of  $S$ .
- (ii) the quotient  $S/E(P)$  is isomorphic with the subsemigroup  $S_{\rho(P)}$  of  $S$  generated by  $x_1, \dots, x_{\rho(P)}$ .

Proof: (ii) Let  $F$  be the monoid freely generated by  $e_1, \dots, e_n$ .

Define  $p : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, \rho(P)\}$  by setting  $p(i) =$  smallest integer in the block of  $P$  which contains  $i$ . There is a unique monoid homomorphism  $\phi : F \rightarrow G$  such that  $\phi(e_i) = x_{p(i)}$  for  $i = 1, 2, \dots, n$ . Since  $w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$  implies  $w(x_{p(1)}, \dots, x_{p(n)}) = v(x_{p(1)}, \dots, x_{p(n)})$  there is a monoid homomorphism  $\phi' : G \rightarrow G$ , induced by  $\phi$ , such that  $\phi'(w(x_1, \dots, x_n)) = w(x_{p(1)}, \dots, x_{p(n)})$ . Consider the equivalence relation  $E$  on  $S$  defined by  $x \equiv y$  if and only if there exist  $w$  and  $v$  in  $F$  such that  $w(x_1, \dots, x_n) = x$ ,  $v(x_1, \dots, x_n) = y$  and  $w(x_{p(1)}, \dots, x_{p(n)}) = v(x_{p(1)}, \dots, x_{p(n)})$ . Clearly  $(i, j) \in P$  implies  $(x_i, x_j) \in E$ , whilst it is easy to show that  $\text{Ker } \phi' = E \subseteq E(P)$ . Since  $E(P)$  is the smallest congruence in which  $x_i$  and  $x_j$  are equivalent whenever  $(i, j) \in P$ , it

follows that  $\text{Ker } \phi' = E(P)$ . Thus  $S/E(P)$  is isomorphic to  $\text{Im } \phi'$ , the submonoid of  $S$  generated by  $\{x_{p(1)}, \dots, x_{p(n)}\}$ , and this set comprises  $\rho(P)$  distinct elements.

(i) If  $P$  is a refinement of  $Q$ , then certainly  $E(P) \subseteq E(Q)$ . Moreover  $P < Q$  ensures  $\rho(P) > \rho(Q)$ , so  $S/E(P)$  and  $S/E(Q)$  are non-isomorphic by (ii) and the previous lemma. Thus  $P < Q$  entails  $E(P) \subsetneq E(Q)$ .

Now  $E(P \vee Q)$  is the congruence generated by the relations  $x_i = x_j$  for  $(i, j) \in P \vee Q$ . Since  $P \vee Q$  is the smallest equivalence relation which contains both  $P$  and  $Q$ , it follows that  $E(P \vee Q) = E(P \cup Q) = E(P) \vee E(Q)$ , showing that  $E$  is a join-preserving map.

Suppose that  $E(P) = E(Q)$ . Then  $E(P) = E(P) \vee E(Q) = E(P \vee Q)$ . Since  $P \vee Q \geq P$ , this implies  $Q \leq P$ . Similarly  $P \leq Q$ , so that  $E$  is a bijective map.

Note: The map  $E$  is not in general a lattice homomorphism.

Let  $S$  be universally generated by  $x_1, \dots, x_n$ , and suppose that  $x_1$  (and thus each generator) has stem of length  $c$  and period  $t$ .

Suppose that  $w(e_1, \dots, e_n)$  and  $v(e_1, \dots, e_n)$  are elements of length  $\ell(w)$  and  $\ell(v)$  respectively in  $F$ , the monoid freely generated by  $e_1, \dots, e_n$ . If  $w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$  in  $S$ , then  $x_1^{\ell(w)} = w(x_1, \dots, x_1) = v(x_1, \dots, x_1) = x_1^{\ell(v)}$  whence either

- (i)  $\ell(w) = \ell(v) < c$
- or (ii)  $\min(\ell(w), \ell(v)) \geq c$   
and  $\ell(w) \equiv \ell(v) \pmod{t}$

Given an element  $x$  in  $S$ , it is then consistent to define the length of  $x$  as the unique number  $\ell(x)$  such that if  $w(x_1, \dots, x_n) = x$  then  $\ell(x) \equiv \ell(w) \pmod{t}$  and  $\ell(x) < c + t$ .

Corollary: Let  $U$  be the partition of  $\{1,2,\dots,n\}$  consisting of a single block. Then  $(x,y) \in E(U)$  if and only if  $l(x) = l(y)$ .

Proof: Let  $w(x_1, \dots, x_n) = x$  and  $v(x_1, \dots, x_n) = y$ . Then  $l(x) = l(y)$  if and only if  $w(x_1, \dots, x_n) = v(x_1, \dots, x_n)$ , and this is equivalent to  $(x,y) \in E(U)$  as observed in the proof of part (ii) of the theorem.

## 6. Algebras with a universal presentation

Necessary conditions for a finite monoid to possess a universal presentation have already been described. In this section, stronger conditions are derived for special varieties of monoid.

Theorem 5: Let  $S$  be an upper semilattice with least element 0 (i.e. a monoid  $(S, \vee)$  in which the binary operation  $\vee$  is commutative and idempotent and 0 is the identity element).

The generators  $x_1, \dots, x_n$  of  $S$  generate a universal presentation of  $S$  if and only if either  $x_1 = x_2 = \dots = x_n$  or  $x_1, \dots, x_n$  freely generate  $S$  as an upper semilattice with zero element.

Proof: The sufficiency of the stated conditions is clear. Accordingly, it suffices to show that if a relation of the form

$$(U) \quad \bigvee_{i \in A} x_i = \bigvee_{i \in B} x_i \quad A, B \subseteq \{1, 2, \dots, n\}$$

holds in  $S$  then either  $x_1 = x_2 = \dots = x_n$  or  $A = B$ .

Assume without loss of generality that  $A \neq \emptyset$ . Then if  $B = \emptyset$  the relation (U) is of the form

$$\bigvee_{i \in A} x_i = 0$$

whence  $x_i = 0$  for all  $i$  in  $A$ , and  $x_1 = x_2 = \dots = x_n = 0$ .

Suppose  $A, B$  both non-empty, and let  $I = A \cap B$ ,  $\bar{A} = A \setminus I$  and  $\bar{B} = B \setminus I$ . If  $\bar{A} = \bar{B} = \emptyset$  then  $A = B$ . Otherwise assume without loss of generality that  $\bar{A} \neq \emptyset$  and let  $f : \{1, 2, \dots, n\} \rightarrow \bar{A}$  be such that

$$f(i) = \begin{cases} 1 & \text{if } i \in \bar{A} \\ 2 & \text{otherwise} \end{cases} \quad \text{Since } (U) \text{ is a universal relation,}$$

$$\bigvee_{i \in A} x_{f(i)} = \bigvee_{i \in B} x_{f(i)} \quad \text{also holds in } S. \quad \text{If } I = \emptyset, \text{ this entails } x_1 = x_2,$$

whence  $x_1 = x_2 = \dots = x_n$ . If  $I \neq \emptyset$ , then  $x_2 \vee x_1 = x_2$  whence (by universality and commutativity)

$$x_1 = x_1 \vee x_2 = x_2 \vee x_1 = x_2.$$

Theorem 6: Let  $G$  be a finite Abelian group. The elements  $g_1, \dots, g_n$  of  $G$  are the generators of a universal presentation if and only if for some  $t$  and some  $d$  dividing  $t$ , the group  $G$  is freely generated by  $g_1, \dots, g_n$  subject to the relations:

$$\left. \begin{array}{ll} \forall i & g_i^t = 1 \\ \forall i, j & g_i g_j = g_j g_i \\ \forall i, j & g_i^d = g_j^d \end{array} \right\} (*)$$

Proof: The group  $G$  with free presentation on generators  $g_1, \dots, g_n$  subject to the relations  $(*)$  is universally presented on generators  $g_1, \dots, g_n$ , since the set of relations  $(*)$  is closed under the application of any function  $f : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$  to the indices of the  $g_i$ 's.

Conversely, suppose that  $g_1, \dots, g_n$  universally generate  $G$  and have common order  $t$ . Then  $G$  has a free presentation on  $g_1, \dots, g_n$  with relations

$$\begin{array}{ll} \forall i & g_i^t = 1 \\ \forall i, j & g_i g_j = g_j g_i \end{array}$$

and other relations of the form  $\prod_{i=1}^n g_i^{r_i} = 1$ . Since  $g_i^t = g_j^t = 1$ ,

there is a least number  $d$  such that  $g_i^d = g_j^d$  (for some, whence all

pairs of indices  $(i,j)$ ). As  $g_i^t = g_j^t$  and  $g_i^d = g_j^d$  ensure  $g_i^{\text{HCF}(d,t)} = g_j^{\text{HCF}(d,t)}$ , it must be that  $d$  divides  $t$ . If the relation  $\prod_{i=1}^n g_i^{r_i} = 1$  holds in  $G$ , then  $\prod_{i=1}^n g_{f(i)}^{r_i} = 1$  for all maps  $f : \{1,2,\dots,n\}$ . In particular, if  $r = \sum_{i=1}^n r_i$  then  $g_1^r = 1$ , whence  $t$  divides  $r$ . Moreover,  $g_1^{r_1} g_2^{r-r_1} = 1$ , showing that  $g_1^{r_1} = g_2^{r_1}$  and thus that  $d$  divides  $r_1$ . By symmetry,  $d$  divides  $r_i$  for each  $i$ , so that the relation  $\prod_{i=1}^n g_i^{r_i} = 1$  is a consequence of the set of relations  $g_i^d = g_j^d$  for all  $i$  and  $j$ .

Cor.1:

$G$  is an Abelian group universally generated by elements  $g_1, \dots, g_n$  of order  $t$  if and only if  $G$  and  $C_t \times C_d^{n-1} \cong \langle \alpha \rangle \times \langle \beta \rangle^{n-1}$  are isomorphic via the mapping  $\phi$  such that  $\phi(g_1) = (\alpha, 1, \dots, 1)$  and  $\phi(g_i) = (\alpha, 1, \dots, 1, \beta, 1, \dots, 1)$    
  $i$ th component  
 for  $i = 2, 3, \dots, n$ .

Proof: It is not difficult to show that the group freely generated by  $g_1, \dots, g_n$  subject to the relations (\*) is indeed isomorphic to  $C_t \times C_d^{n-1}$  via the mapping  $\phi$ .

## 7. Groups with a universal presentation

Necessary and sufficient conditions for a finite Abelian group to have a universal presentation are given in Theorem 6. The results and examples in this section relate to the harder (and unresolved) problem of determining which finite non-Abelian groups admit a universal presentation.

The following result is a corollary to Theorem 6:

Cor 2 to Theorem 6:

Suppose that  $g_1, \dots, g_n$  generate a universal presentation for the finite group  $G$ . Let  $G'$  be the commutator subgroup of  $G$ . The images  $\bar{g}_1, \dots, \bar{g}_n$  of  $g_1, \dots, g_n$  generate a universal presentation of  $G/G'$ . In particular,  $G/G'$  is isomorphic with  $C_t \times C_d^{n-1}$  for some positive integers  $k$  and  $d$  where  $d$  divides  $k$ .

Proof: The elements of  $G'$  are products of commutators. Thus if

$w(g_1, \dots, g_n) \in G'$  and  $f : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$  is any map, then

$w(g_{f(1)}, \dots, g_{f(n)}) \in G'$ . That is, the relations imposed upon  $g_1, \dots, g_n$  by taking the quotient by  $G'$  hold universally in  $G/G'$ . By Theorem 6,  $G/G'$  (being a finite Abelian group) is isomorphic with some  $C_t \times C_d^{n-1}$ .

The next result is the analogue for groups of Theorem 4.

Theorem 7: Let  $G$  be a finite group generated universally by distinct generators  $g_1, \dots, g_n$ . For each partition  $P$ , let  $N(P)$  be the normal subgroup of  $G$  generated by all elements of the form  $g_i g_j^{-1}$  such that  $(i, j) \in P$ .

Then

(i) the map  $N$  is a join-preserving bijection from the lattice of partitions of  $\{1, 2, \dots, n\}$  (ordered by refinement) to the lattice of normal subgroups of  $G$ .

(ii) the quotient  $G/N(P)$  is isomorphic with the subgroup  $G_{\rho(P)}$  of  $G$  generated by  $g_1, \dots, g_{\rho(P)}$ , and  $G$  is isomorphic to a semi-direct product of  $G_{\rho(P)}$  and  $N(P)$ .

Proof: It suffices to show that  $G \approx G_{\rho(P)} * N(P)$ ; the other results are interpretations of Theorem 4.

For  $i = 1, 2, \dots, n$  let  $q(i)$  be the least integer such that  $(i, q(i)) \in P$ . Let  $\theta : G \rightarrow G$  be the group homomorphism such that  $\theta(g_i) = g_{q(i)}$  (c.f. proof of Theorem 4 (ii)). Then  $\text{Ker}\theta = N(P)$ , and  $\text{Im}\theta = \langle g_{q(1)}, \dots, g_{q(n)} \rangle$ . Note that  $g_i = g_{q(i)} (g_{q(i)}^{-1} g_i) \in \text{Im}\theta \cdot \text{Ker}\theta$

so that  $G = \text{Im } \theta \cdot \text{Ker } \theta$ . Moreover, if  $g \in \text{Ker } \theta \cap \text{Im } \theta$ , then  $\theta(g) = g = 1$ .

Thus  $G = \text{Im } \theta * \text{Ker } \theta \approx C_{\rho(P)} * N(P)$

Corollary 1: If  $G$  has a universal presentation by generators  $g_1, \dots, g_n$  of common order  $t$  then the elements of  $G$  of length 0 form a normal subgroup  $N$  of  $G$ , and  $G \approx C_t * N$ .

Proof: See the corollary to Theorem 4, and apply Theorem 7 (ii).

Corollary 2: If  $k \geq 4$ , then the symmetric group  $S_k$  has no universal presentation.

Proof: Suppose  $g_1, \dots, g_n$  are permutations generating  $S_k$  universally, and let  $g_1, \dots, g_n$  have common order  $t$ . Since  $S_k \approx C_t * N$  for some normal subgroup  $N$ , it must be that  $N = A_k$  and  $t = 2$ . On the other hand, in view of Theorem 7 (i),  $n \leq 2$ . But, if a group is generated by two elements of order 2 it is dihedral (see [4] p.49 Ex.1 ).

## 8. Examples of groups with universal presentations

### Example 1:

As suggested by the proof of the previous corollary, the dihedral group  $D_n$  of order  $2n$  has a universal presentation by two generators of order 2, viz.  $\langle x, y \mid x^2 = y^2 = (xy)^n = 1 \rangle$ . In particular  $S_3 (\approx D_3)$  is universally generated by a pair of transpositions.

### Example 2:

Every finite Burnside group  $B(t, n)$  (which is generated by  $n$  elements  $x_1, \dots, x_n$  subject to relations  $g^t = 1$  for every  $g$  in  $B(t, n)$ ) is universally generated by its canonical generating set.

The Burnside group  $B(3, 3)$  of order 2187 illustrates that the map  $N$  in Theorem 7 (and likewise the map  $E$  in Theorem 4) is not in general a lattice homomorphism. As described in [3], every element of  $B(3, 3)$  has a unique representation of the form:

$$x_1^{a_1} x_2^{a_2} x_3^{a_3} (x_1, x_2)^{b_3} (x_1, x_3)^{b_2} (x_2, x_3)^{b_1} (x_1, x_2, x_3)^c$$

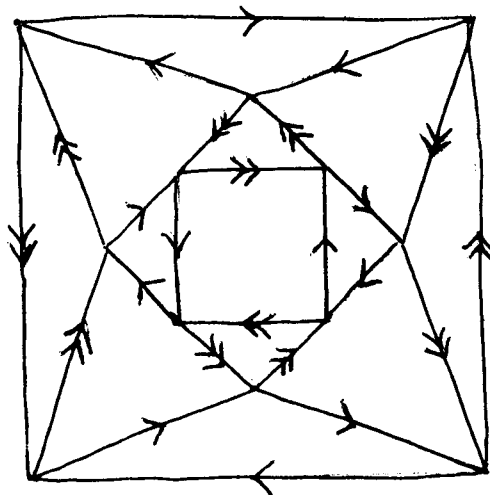
where  $0 \leq a_i, b_i, c \leq 2$ . Let  $P$  be the partition (12)(3) and  $Q$  the partition (1)(23). The partition  $P \wedge Q$  is (1)(2)(3) whence  $N(P \wedge Q) = \{1\}$ . But  $(x_1, x_2)^2 (x_1, x_3) (x_2, x_3)^2 \in N(P) \cap N(Q)$  (it reduces to 1 under adjunction of the relation  $x_1 = x_2$  or  $x_2 = x_3$ ) whence  $N(P) \cap N(Q) \neq N(P \wedge Q)$ .

### Example 3:

The group  $A_4$  is universally generated by  $x$  and  $y$  subject to the relations:

$$x^3 = y^3 = (xy)^3 = (xy^2)^2 = 1$$

The Cayley diagram associated with this presentation is:



### Example 4:

The group  $G$  of order 56 generated by  $x$  and  $y$  subject to the relations

$$x^2 y x y^3 = y^2 x y x^3 = 1$$

is universally generated by  $x$  and  $y$ , which are elements of order 7 (see [4] p.60). The semi-direct product decomposition of  $G$  referred to in Corollary 2 to Theorem 7 exhibits  $G$  as  $C_7 * C_2^3$ .

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