

Balance and Abelian Complexity of the Tribonacci word

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Abstract

G. Rauzy showed that the Tribonacci minimal subshift generated by the morphism $\tau : 0 \mapsto 01, 1 \mapsto 02$ and $2 \mapsto 0$ is measure-theoretically conjugate to an exchange of three fractal domains on a compact set in \mathbb{R}^2 , each domain being translated by the same vector modulo a lattice. In this paper we study the Abelian complexity $\rho(n)$ of the Tribonacci word \mathbf{t} which is the unique fixed point of τ . We show that $\rho(n) \in \{3, 4, 5, 6, 7\}$ for each $n \geq 1$, and that each of these five values is assumed. Our proof relies on the fact that the Tribonacci word is 2-balanced, i.e., for all factors U and V of \mathbf{t} of equal length, and for every letter $a \in \{0, 1, 2\}$, the number of occurrences of a in U and the number of occurrences of a in V differ by at most 2. While this result is announced in several papers, to the best of our knowledge no proof of this fact has ever been published. We offer two very different proofs of the 2-balance property of \mathbf{t} . The first uses the word combinatorial properties of the generating morphism, while the second exploits the spectral properties of the incidence matrix of τ .

1 Introduction

Given a finite non-empty set A , called the *alphabet*, we denote by A^* the free monoid generated by A . The identity element of A^* , called the *empty word*, will be denoted by ε . For any word $u = a_1 a_2 \cdots a_n \in A^*$, the length of u is the quantity n and is denoted by $|u|$. By convention, the length of the empty word ε is taken to be 0. For each $a \in A$, let $|u|_a$ denote the number of occurrences of the letter a in u . We denote by $A^{\mathbb{N}}$ the set of (right) infinite words

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on the alphabet A . Given an infinite word $\omega = \omega_0\omega_1\omega_2\cdots \in A^{\mathbb{N}}$, any finite word of the form $\omega_i\omega_{i+1}\cdots\omega_{i+n-1}$ (with $i \geq 0$ and $n \geq 1$) is called a *factor* of ω . Let

$$\mathcal{F}_\omega(n) = \{\omega_i\omega_{i+1}\cdots\omega_{i+n-1} \mid i \geq 0\}$$

denote the set of all factors of ω of length n , and set $p_\omega(n) = \text{Card}(\mathcal{F}_\omega(n))$. The function $p_\omega : \mathbb{N} \rightarrow \mathbb{N}$ is called the *subword complexity function* of ω . A fundamental result due to Hedlund and Morse [15] states that a word ω is ultimately periodic if and only if for some n the subword complexity $p_\omega(n) \leq n$. Words of subword complexity $p(n) = n + 1$ are called *Sturmian words*. The most well-known Sturmian word is the so-called Fibonacci word

$$\mathbf{f} = 0100101001001010010100100101001001010010100101001001010010 \cdots$$

fixed by the morphism $0 \mapsto 01$ and $1 \mapsto 0$.

Sturmian words admit various types of characterizations of geometric and combinatorial nature. We give two examples: In [16] Hedlund and Morse showed that each Sturmian word may be realized measure-theoretically by an irrational rotation on the circle. That is, every Sturmian word is obtained by coding the symbolic orbit of a point x on the circle (of circumference one) under a rotation by an irrational angle α where the circle is partitioned into two complementary intervals, one of length α and the other of length $1 - \alpha$. And conversely each such coding gives rise to a Sturmian word.

Sturmian words are equally characterized by a certain balance property. We begin with the following definition:

Definition 1.1. *An infinite word $\omega \in A^{\mathbb{N}}$ is said to be C -balanced (C a positive integer) if $||U|_a - |V|_a| \leq C$ for all factors U and V of ω of equal length, and each $a \in A$. We say ω is balanced if it is 1-balanced.*

Theorem 1.2 (Theorem 2.1.5 in [13]). *An infinite word ω is Sturmian if and only if ω is a binary aperiodic balanced word.*

In [17, 18], Rauzy showed that the regular continued fraction algorithm provides a formidable link between the arithmetic/diophantine properties of an irrational number α , the ergodic/dynamical properties of a rotation by angle α on the circle, and the combinatorial (balance) properties of Sturmian words.

A fundamental problem is to generalize and extend this rich interaction to higher dimensions, either by starting with a specified symbolic dynamical system, or by a class of words satisfying particular combinatorial conditions. In case of dimension 2, there are a couple of different dynamical systems which are natural candidates: For instance, promising results have been obtained by the third author together with Ferenczi and Holton by considering the dynamics of 3-interval exchange transformations on the unit interval [8, 9, 10, 11], and also by Arnoux, Berthé and Ito by considering two rotations on the circle [2]. However, ever since the early work of Rauzy in [19], the most natural generalization was thought to be the one stemming from a rotation on the 2-torus. In this context, the associated symbolic counterpart

is given by a class of words of subword complexity $2n + 1$, now called *Arnoux-Rauzy words*, originally introduced by Arnoux and Rauzy in [3]. In this setting, the *Tribonacci word*

$$\mathbf{t} = \lim_{n \rightarrow \infty} \tau^n(0) = 01020100102010 \dots$$

defined as the fixed point of the morphism τ

$$0 \mapsto 01 \quad 1 \mapsto 02 \quad 2 \mapsto 0,$$

is the natural analogue of the Fibonacci word \mathbf{f} . In fact, in [19] Rauzy showed that the Tribonacci minimal subshift (the shift orbit closure of \mathbf{t}) is a natural coding of a rotation on the 2-dimensional torus T^2 , i.e., is measure-theoretically conjugate to an exchange of three fractal domains on a compact set in \mathbb{R}^2 , each domain being translated by the same vector modulo a lattice. In [6], the third author together with Cassaigne and Ferenczi showed that there exist Arnoux-Rauzy words which are arbitrarily imbalanced, that is to say, which are not C -balanced for any positive integer C , and that such Arnoux-Rauzy words cannot be measure-theoretically conjugate to a rotation on the 2-torus.

In the first part of this paper, we prove:

Theorem 1.3. *The Tribonacci word \mathbf{t} is 2-balanced.*

While this result is announced in several papers (see for instance [4, 5, 6, 12, 22]), to the best of our knowledge no proof of this fact has ever been published. In this paper we give two different proofs of Theorem 1.3. The first is a proof by induction which uses the word combinatorial properties of the generating morphism τ . The second proof relies on the spectral properties of the incidence matrix and is more in the spirit of the methods developed by Adamczewski in [1].

Our hope is that one of the two proofs may be extended to establish a balance property for the general m -bonacci word ($m \geq 2$) defined as the fixed point of the morphism

$$0 \mapsto 01 \quad 1 \mapsto 02 \quad \dots \quad (m-2) \mapsto 0(m-1) \quad (m-1) \mapsto 0.$$

Computer simulations show that the 4-bonacci word is not 2-balanced. In fact, let u be the factor of the 4-bonacci word of length 3305 occurring in position 2663 (starting with 0), and v the factor of length 3305 occurring in position 9048. Then $|u|_1 = 891$ and $|v|_1 = 888$, which shows that the 4-bonacci word is not 2-balanced. It is possible that in general the m -bonacci word is $(m-1)$ -balanced as stated in [12].

We then apply Theorem 1.3 to study the so-called Abelian complexity of the Tribonacci word. Following [20], two words u and v in A^* are said to be *Abelian equivalent*, denoted $u \sim_{\text{ab}} v$, if and only if $|u|_a = |v|_a$ for all $a \in A$. For instance, the words $ababa$ and $aaabb$ are Abelian equivalent whereas they are not equivalent to $aabbb$. It is readily verified that \sim_{ab} defines an equivalence relation on A^* .

We define

$$\mathcal{F}_\omega^{\text{ab}}(n) = \mathcal{F}_\omega(n) / \sim_{\text{ab}}$$

and set $\rho_\omega = \text{Card}(\mathcal{F}_\omega^{\text{ab}}(n))$. The function $\rho = \rho_\omega : \mathbb{N} \rightarrow \mathbb{N}$ which counts the number of pairwise non Abelian equivalent factors of ω of length n is called the *Abelian complexity* or *ab-complexity* for short. Using Theorem 1.3 we compute the ab-complexity of the Tribonacci word \mathbf{t} :

Theorem 1.4. *Let \mathbf{t} denote the Tribonacci word. Then $\rho_{\mathbf{t}}(n) \in \{3, 4, 5, 6, 7\}$ for every positive integer n . Moreover, each of these five values is attained.*

We also obtain several equivalent characterizations of those values n for which $\rho(n) = 3$. As a consequence we show that $\rho(n) = 3$ for infinitely many n . Similarly we show that $\rho(n) = 7$ for infinitely many n , although the least n for which $\rho(n) = 7$ is $n = 3914$. We conclude with some open questions.

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2 Two Proofs of Theorem 1.3

2.1 Preliminaries to Proof 1 of Theorem 1.3

Let $\tau : \{0, 1, 2\} \rightarrow \{0, 1, 2\}^*$ denote the morphism

$$0 \mapsto 01 \quad 1 \mapsto 02 \quad 2 \mapsto 0,$$

and let

$$\mathbf{t} = \lim_{n \rightarrow \infty} \tau^n(0) = 01020100102010 \dots$$

denote the Tribonacci word.

For each $u \in \{0, 1, 2\}^*$, we denote by $\Psi(u)$ the *Parikh vector* associated to u , that is

$$\Psi(u) = (|u|_0, |u|_1, |u|_2).$$

For each infinite word $\omega \in \{0, 1, 2\}^{\mathbb{N}}$ let $\Psi_\omega(n)$ denote the set of Parikh vectors of factors of length n of ω :

$$\Psi_\omega(n) = \{\Psi(u) \mid u \in \mathcal{F}_\omega(n)\}.$$

Thus we have

$$\rho_\omega(n) = \text{Card}(\Psi_\omega(n)).$$

We will need the following two lemmas.

Lemma 2.1. *Let U be a non-empty factor of \mathbf{t} . Then there exists a factor u of \mathbf{t} such that $\Psi(U) = (|u| + \delta, |u|_0, |u|_1)$ for some $\delta \in \{-1, 0, 1\}$. Moreover if $|U| \geq 3$ then $|u| < |U|$.*

Proof. It is readily seen that every non-empty factor U of \mathbf{t} may be written as either $U = \tau(u)$, or $U = 0^{-1}\tau(u)$, or $U = \tau(u)0$, or $U = 0^{-1}\tau(u)0$ for some factor u of \mathbf{t} . Moreover, $|\tau(u)|_0 = |u|$, $|\tau(u)|_1 = |u|_0$, and $|\tau(u)|_2 = |u|_1$. Thus we have that

$$\Psi(U) = (|u| + \delta, |u|_0, |u|_1)$$

for some $\delta \in \{-1, 0, 1\}$. Moreover, from above it follows that $|U| \geq |\tau(u)| - 1$. Also, for all u we have $|\tau(u)| \geq |u|$, and if $|u| \geq 3$, then $|\tau(u)| \geq |u| + 2$ since either u contains at least two occurrences of 0, or u contains both an occurrence of 0 and an occurrence of 1. Finally, to see that $|u| < |U|$ whenever $|U| \geq 3$, we suppose to the contrary that $|u| \geq |U| \geq 3$; then

$$|U| \geq |\tau(u)| - 1 \geq |u| + 1 > |u|,$$

which is a contradiction. □

Note: In what follows, we will often apply the above lemma to a pair of words U and V where $|V|_0 = |U|_0 + 2$ and $|U|_i = |V|_i + 3$ for some $i \in \{1, 2\}$. Note that under these hypotheses, $|U| \geq 7$ since U contains either at least 3 occurrences of 1, or at least 3 occurrences of 2. Since every 1 and 2 occurring in \mathbf{t} is always preceded by 0, we deduce that $|U|_0 \geq 2$, and hence $|V| \geq 4$.

Lemma 2.2. *Let U and V be factors of \mathbf{t} satisfying the following two conditions:*

$$|V|_0 = |U|_0 + 2$$

$$|U|_i = |V|_i + 3$$

for some $i \in \{1, 2\}$. Then there exist factors u and v of \mathbf{t} , with $|u| < |U|$, $|v| < |V|$, $|u| \leq |v|$, and $|u|_{i-1} = |v|_{i-1} + 3$.

Proof. By the previous lemma there exist factors u, v with $|u| < |U|$, $|v| < |V|$, such that

$$\Psi(U) = (|u| + \delta_1, |u|_0, |u|_1) \quad \text{and} \quad \Psi(V) = (|v| + \delta_2, |v|_0, |v|_1)$$

for some $\delta_1, \delta_2 \in \{-1, 0, 1\}$. By hypothesis we have

$$|v| + \delta_2 = |V|_0 = |U|_0 + 2 = |u| + \delta_1 + 2,$$

whence

$$|u| = |v| + \delta_2 - \delta_1 - 2 \leq |v|.$$

Finally, the condition $|U|_i = |V|_i + 3$ for $i \in \{1, 2\}$ implies $|u|_{i-1} = |v|_{i-1} + 3$. □

In what follows, we will make use the following terminology: We say two finite words U and V on the alphabet $\{0, 1, 2\}$ are *pairwise 2-imbalanced* if $|U| = |V|$ and $||U|_i - |V|_i| \geq 3$ for some $i \in \{0, 1, 2\}$.

2.2 Proof 1 of Theorem 1.3

Suppose to the contrary that \mathbf{t} is not 2-balanced. Then there exists a shortest pair of factors U and V with $|U| = |V|$ and $|U|_i - |V|_i \geq 3$ for some $i \in \{0, 1, 2\}$. If $|U|_i - |V|_i > 3$, then by removing the last letter from each of U and V , we would obtain a shorter pair of words of equal length which are pairwise 2-imbalanced. Thus the minimality condition on $|U|$ implies that $|U|_i - |V|_i = 3$. Also, it is easily checked that all factors of length less or equal to 3 are 2-balanced whence $|U| \geq 4$. We consider three cases: $i = 0$, then $i = 2$, and finally $i = 1$.

Case 1: $|U|_0 - |V|_0 = 3$. By Lemma 2.1 there exist factors u, v with $|u| < |U|$, $|v| < |V|$, such that

$$\Psi(U) = (|u| + \delta_1, |u|_0, |u|_1) \quad \text{and} \quad \Psi(V) = (|v| + \delta_2, |v|_0, |v|_1)$$

for some $\delta_1, \delta_2 \in \{-1, 0, 1\}$. The condition $|U|_0 = |V|_0 + 3$ implies that

$$|u| + \delta_1 = |v| + \delta_2 + 3$$

that is

$$|u| - |v| = (3 + \delta_2 - \delta_1) > 0.$$

The condition $|U| = |V|$ implies that

$$|u| + \delta_1 + |u|_0 + |u|_1 = |v| + \delta_2 + |v|_0 + |v|_1$$

or equivalently

$$2|u| + \delta_1 - |u|_2 = 2|v| + \delta_2 - |v|_2.$$

Thus

$$|u|_2 - |v|_2 = 2(|u| - |v|) + \delta_1 - \delta_2 = 2(3 + \delta_2 - \delta_1) + \delta_1 - \delta_2 = 6 + \delta_2 - \delta_1.$$

Let u' be the prefix of u of length $|u| - (3 + \delta_2 - \delta_1) = |v|$. Then, $|u'|_2 \geq |v|_2 + 3$ contradicting the minimality of $|U|$.

Case 2: $|U|_2 - |V|_2 = 3$. By Lemma 2.1 there exist factors u, v with $|u| < |U|$, $|v| < |V|$, such that

$$\Psi(U) = (|u| + \delta_1, |u|_0, |u|_1) \quad \text{and} \quad \Psi(V) = (|v| + \delta_2, |v|_0, |v|_1)$$

for some $\delta_1, \delta_2 \in \{-1, 0, 1\}$. The condition $|U|_2 = |V|_2 + 3$ implies that $|u|_1 = |v|_1 + 3$. If $|u| \leq |v|$, then u and the prefix of v of length $|u|$ are pairwise 2-imbalanced and of length less than $|U|$, contradicting the minimality of $|U|$. Thus we can suppose that

$$|u| = |v| + m \quad \text{for some } m \geq 1.$$

If $m \geq 2$, then we deduce that

$$|U|_0 = |u| + \delta_1 \geq |u| - 1 \geq |u| - (m - 1) = |v| + 1 \geq |v| + \delta_2 = |V|_0.$$

As $|U|_2 = |V|_2 + 3$, $|U|_0 \geq |V|_0$, and $|U| = |V|$, it follows that $|V|_1 \geq |U|_1 + 3$, that is $|v|_0 \geq |u|_0 + 3$. But then v and the prefix of u of length $|v|$ are pairwise 2-imbalanced and of length less than $|U|$, contradicting the minimality of $|U|$.

Thus we can suppose $m = 1$, that is $|u| = |v| + 1$. Again, if $|U|_0 \geq |V|_0$, as above we would deduce that $|v|_0 \geq |u|_0 + 3$ which would give rise to a contradiction. So we must have that $|U|_0 < |V|_0$. This gives

$$|v| + 1 + \delta_1 = |u| + \delta_1 = |U|_0 < |V|_0 = |v| + \delta_2$$

that is

$$1 + \delta_1 < \delta_2$$

which in turn implies that

$$\delta_1 = -1 \text{ and } \delta_2 = 1.$$

Thus

$$\Psi(U) = (|u| - 1, |u|_0, |u|_1) \text{ and } \Psi(V) = (|v| + 1, |v|_0, |v|_1).$$

Finally, the conditions $|U| = |V|$ together with $|u| = |v| + 1$, and $|u|_1 = |v|_1 + 3$ imply that

$$\begin{aligned} |v|_0 &= |v|_0 + (|v| + 1 + |v|_1) - (|v| + 1 + |v|_1) \\ &= |V| - (|v| + 1 + |v|_1) \\ &= |V| - (|u| + |u|_1 - 3) \\ &= |V| - (|u| + |u|_1 - 1 - 2) \\ &= |V| - (|U| - |u|_0 - 2) \\ &= |u|_0 + 2. \end{aligned}$$

In summary we have: $|u|_1 = |v|_1 + 3$ and $|v|_0 = |u|_0 + 2$. Thus we can apply Lemma 2.2 which gives a contradiction to the minimality of $|U|$.

Case 3: $|U|_1 - |V|_1 = 3$. By Lemma 2.1 there exist factors u, v with $|u| < |U|$, $|v| < |V|$, such that

$$\Psi(U) = (|u| + \delta_1, |u|_0, |u|_1) \text{ and } \Psi(V) = (|v| + \delta_2, |v|_0, |v|_1)$$

for some $\delta_1, \delta_2 \in \{-1, 0, 1\}$. The condition $|U|_1 = |V|_1 + 3$ implies that $|u|_0 = |v|_0 + 3$.

If $|u| \leq |v|$, then u and the prefix of v of length $|u|$ are pairwise 2-imbalanced and of length less than $|U|$, contradicting the minimality of $|U|$. Thus we can suppose that

$$|u| = |v| + m \text{ for some } m \geq 1.$$

Next we proceed like in Case 2. If $m \geq 2$, then we deduce that

$$|U|_0 = |u| + \delta_1 \geq |u| - 1 \geq |u| - (m - 1) = |v| + 1 \geq |v| + \delta_2 = |V|_0.$$

As $|U|_1 = |V|_1 + 3$, $|U|_0 \geq |V|_0$, and $|U| = |V|$, it follows that $|V|_2 \geq |U|_2 + 3$, that is $|v|_1 \geq |u|_1 + 3$. But then v and the prefix of u of length $|v|$ are pairwise 2-imbalanced and of length less than $|U|$, contradicting the minimality of $|U|$.

Thus we can suppose $m = 1$, that is $|u| = |v| + 1$. Again, if $|U|_0 \geq |V|_0$, as above we would deduce that $|v|_1 \geq |u|_1 + 3$ which would give rise to a contradiction. So we must have that $|U|_0 < |V|_0$. This gives

$$|v| + 1 + \delta_1 = |u| + \delta_1 = |U|_0 < |V|_0 = |v| + \delta_2$$

that is

$$1 + \delta_1 < \delta_2$$

which in turn implies that

$$\delta_1 = -1 \text{ and } \delta_2 = 1.$$

Thus

$$\Psi(U) = (|u| - 1, |u|_0, |u|_1) \text{ and } \Psi(V) = (|v| + 1, |v|_0, |v|_1).$$

The conditions $|U| = |V|$ together with $|u| = |v| + 1$, and $|u|_0 = |v|_0 + 3$ imply that

$$\begin{aligned} |v|_1 &= |v|_1 + (|v| + 1 + |v|_0) - (|v| + 1 + |v|_0) \\ &= |V| - (|v| + 1 + |v|_0) \\ &= |V| - (|u| + |u|_0 - 3) \\ &= |V| - (|u| + |u|_0 - 1 - 2) \\ &= |V| - (|U| - |u|_1 - 2) \\ &= |u|_1 + 2. \end{aligned}$$

In summary we have

$$|u| = |v| + 1 \text{ and } |v|_1 = |u|_1 + 2 \text{ and } |u|_0 = |v|_0 + 3.$$

Since $|u|, |v| \geq 3$, Lemma 2.1 implies that there exist factors u', v' with $|u'| < |u|$, $|v'| < |v|$, such that

$$\Psi(u) = (|u'| + \delta'_1, |u'|_0, |u'|_1) \text{ and } \Psi(v) = (|v'| + \delta'_2, |v'|_0, |v'|_1)$$

for some $\delta'_1, \delta'_2 \in \{-1, 0, 1\}$.

As $|v|_1 = |u|_1 + 2$ we deduce that $|v'|_0 = |u'|_0 + 2$. Similarly, the condition $|u|_0 = |v|_0 + 3$ implies that

$$|u'| + \delta'_1 = |v'| + \delta'_2 + 3$$

so that

$$|v'| = |u'| + \delta'_1 - \delta'_2 - 3 \leq |u'| - 1.$$

Now if $|v'| \leq |u'| - 2$, then if v'' is any factor of the Tribonacci word of length $|u'|$ beginning in v' , we would have $|v''|_0 \geq |v'|_0 + 1 = |u'|_0 + 3$ as every factor of length 2 or greater contains at least one occurrence of the letter 0.

So we can assume that $|v'| = |u'| - 1$. Then

$$|v'| + 1 + \delta'_1 = |u'| + \delta'_1 = |u|_0 = |v|_0 + 3 = |v'| + \delta'_2 + 3$$

that is

$$\delta'_1 - \delta'_2 = 2,$$

which in turn implies that

$$\delta'_1 = 1 \text{ and } \delta'_2 = -1.$$

So

$$\begin{aligned} |u'|_1 &= |u'|_1 + (|v'| + |v'|_0) - (|v'| + |v'|_0) \\ &= |u'|_1 + (|v'| + 1 + 1 + |v'|_0 - 2) - (|v'| - 1 + |v'|_0 + 1) \\ &= |u'|_1 + (|u'| + 1 + |u'|_0) - (|v'| - 1 + |v'|_0 + 1) \\ &= |u'|_1 + (|u| - |u'|) - (|v| - |v'|_1 + 1) \\ &= |u| - (|v| - |v'|_1 + 1) \\ &= |u| - (|v| + 1) + |v'|_1 \\ &= |v'|_1 \end{aligned}$$

Finally,

$$\begin{aligned} |u'|_2 &= |u'|_2 + (|u'|_0 + |u'|_1 - 1) - (|u'|_0 + |u'|_1 - 1) \\ &= |u'| - 1 - (|u'|_0 + |u'|_1 - 1) \\ &= |v'| - (|v'|_0 - 2 + |v'|_1 - 1) \\ &= |v'|_2 + 3. \end{aligned}$$

So in summary we have

$$|v'|_0 = |u'|_0 + 2 \text{ and } |u'|_2 = |v'|_2 + 3$$

to which we can apply Lemma 2.2 to obtain the desired contradiction.

2.3 Preliminaries to Proof 2 of Theorem 1.3

For the results in this section, see [14, chapter 10].

We write $\mathbf{t} = u_0 u_1 u_2 \cdots$ with $u_i \in \{0, 1, 2\}$. The *Tribonacci numbers* T_k are defined by $T_k = |\tau^k(0)|$. Hence $T_k = T_{k-1} + T_{k-2} + T_{k-3}$ for $k \geq 3$, and

$$T_0 = 1, \quad T_1 = 2, \quad T_2 = 4, \quad T_3 = 7, \quad T_4 = 13, \dots$$

Let

$$M = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$

denote the incidence matrix of the morphism τ , i.e., $M_{ij} = |\tau(j)|_i$ for all $i, j \in \{0, 1, 2\}$. The eigenvalues of M are the roots of the polynomial $x^3 - x^2 - x - 1$. We denote them β, α , and $\bar{\alpha}$ so that β is real and α and $\bar{\alpha}$ are complex conjugates. (Given a complex number $z = a + bi$, we denote by $\bar{z} = a - bi$ the complex conjugate.) We have

$$\beta = 1.83928\dots \quad \text{and} \quad |\alpha| = |\bar{\alpha}| = 0.73735\dots$$

More information concerning the roots α and β may be found in the comments of sequences A058265 and A000073 in “The On-Line Encyclopedia of Integer Sequences”. In particular, some formulas relating α and $\bar{\alpha}$ to β are recalled.

It is well-known that the frequencies of letters in \mathbf{t} exist and

$$\text{Freq}_{\mathbf{t}}(0) = \frac{1}{\beta}, \quad \text{Freq}_{\mathbf{t}}(1) = \frac{1}{\beta^2}, \quad \text{Freq}_{\mathbf{t}}(2) = \frac{1}{\beta^3}.$$

Set

$$v_{\beta} = \begin{pmatrix} \beta^{-1} \\ \beta^{-2} \\ \beta^{-3} \end{pmatrix}, \quad v_{\alpha} = \begin{pmatrix} \alpha^{-1} \\ \alpha^{-2} \\ \alpha^{-3} \end{pmatrix}, \quad v_{\bar{\alpha}} = \begin{pmatrix} \bar{\alpha}^{-1} \\ \bar{\alpha}^{-2} \\ \bar{\alpha}^{-3} \end{pmatrix}.$$

Then $v_{\beta}, v_{\alpha}, v_{\bar{\alpha}}$ are eigenvectors of M corresponding to $\beta, \alpha, \bar{\alpha}$, respectively, normalized so that in each case the sum of coordinates equals 1. Let us denote by $a_{\beta}, a_{\alpha}, a_{\bar{\alpha}}$ the complex numbers for which

$$a_{\beta}v_{\beta} + a_{\alpha}v_{\alpha} + a_{\bar{\alpha}}v_{\bar{\alpha}} = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}.$$

The numbers a_{β}, a_{α} and $a_{\bar{\alpha}}$ are unique, and therefore

$$a_{\bar{\alpha}} = \bar{a}_{\alpha} \quad \text{and} \quad |a_{\alpha}| = |a_{\bar{\alpha}}| = 0.14135\dots$$

Denote by

$$e_1 = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}, \quad e_2 = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}, \quad e_3 = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}.$$

Then we can represent letter frequencies as $\text{Freq}_{\mathbf{t}}(i) = \langle v_{\beta}, e_i \rangle$ for each letter $i = 0, 1, 2$, where $\langle \cdot, \cdot \rangle$ denotes the Hermitian scalar product.

2.4 Proof 2 of Theorem 1.3

Every natural number N has a unique *Zeckendorff* Tribonacci representation

$$N = \sum_{k \geq 0} p_k T_k,$$

where $p_k \in \{0, 1\}$, all but finitely many p_k equal 0, and if $p_k = p_{k-1} = 1$, then $p_{k-2} = 0$.

It can be shown (see [14, Proposition 10.7.4 on page 510]) that for each letter $i \in \{0, 1, 2\}$,

$$|u_0 \cdots u_{N-1}|_i = a_\beta \langle v_\beta, e_i \rangle \sum_{k \geq 0} p_k \beta^k + a_\alpha \langle v_\alpha, e_i \rangle \sum_{k \geq 0} p_k \alpha^k + a_{\bar{\alpha}} \langle v_{\bar{\alpha}}, e_i \rangle \sum_{k \geq 0} p_k \bar{\alpha}^k,$$

and

$$N = a_\beta \sum_{k \geq 0} p_k \beta^k + a_\alpha \sum_{k \geq 0} p_k \alpha^k + a_{\bar{\alpha}} \sum_{k \geq 0} p_k \bar{\alpha}^k.$$

These two identities imply

$$\begin{aligned} |u_0 \cdots u_{N-1}|_i - N \langle v_\beta, e_i \rangle &= a_\alpha \langle v_\alpha - v_\beta, e_i \rangle \sum_{k \geq 0} p_k \alpha^k + a_{\bar{\alpha}} \langle v_{\bar{\alpha}} - v_\beta, e_i \rangle \sum_{k \geq 0} p_k \bar{\alpha}^k \\ &= 2 \operatorname{Re} \left(a_\alpha \langle v_\alpha - v_\beta, e_i \rangle \sum_{k \geq 0} p_k \alpha^k \right), \\ &= \sum_{k \geq 0} 2 p_k \operatorname{Re} \left(a_\alpha (\alpha^{-(i+1)} - \beta^{-(i+1)}) \alpha^k \right), \end{aligned} \tag{1}$$

where $\operatorname{Re}(\cdot)$ denotes the real part of a complex number. The second equality above holds because complex conjugation commutes with multiplication and addition.

The following lemma shows that in order to prove Theorem 1.3, it suffices to determine adequate upper and lower bounds for the above equation (1):

Lemma 2.3. *Let $x = x_0 x_1 x_2 \cdots$ be an infinite word on a finite alphabet A , and $b \in A$ a letter occurring in x . If there exist real numbers A, B and γ for which*

$$A < |x_0 x_1 \cdots x_{N-1}|_b - N\gamma < B$$

for all integers $N \geq 1$, then the word x is $\lfloor 2(B - A) \rfloor$ -balanced with respect to the letter b , i.e.,

$$\left| |U|_b - |V|_b \right| \leq \lfloor 2(B - A) \rfloor$$

for all factors U, V of equal length.

Proof. Let $i, j \geq 0$ be integers. Since

$$|x_i x_{i+1} \cdots x_{i+N-1}|_b - N\gamma = \left(|x_0 x_1 \cdots x_{i+N-1}|_b - (N+i)\gamma \right) - \left(|x_0 x_1 \cdots x_{i-1}|_b - i\gamma \right),$$

we see that

$$A - B < |x_i x_{i+1} \cdots x_{i+N-1}|_b - N\gamma < B - A. \tag{2}$$

Consequently,

$$\left| |x_i x_{i+1} \cdots x_{i+N-1}|_b - |x_j x_{j+1} \cdots x_{j+N-1}|_b \right| < (B - A) - (A - B) = 2(B - A).$$

□

We now determine good upper and lower bounds for Eq. (1). The numerical calculations in Propositions 2.4, 2.5 and 2.6 were verified independently by the first and second authors via different computer programs.

Proposition 2.4. *We have*

$$-0.6 < |u_0 u_1 \cdots u_{N-1}|_0 - \frac{N}{\beta} < 0.9.$$

Proof. By (1), we have

$$|u_0 \cdots u_{N-1}|_0 - \frac{N}{\beta} = \sum_{k \geq 0} 2p_k \operatorname{Re} \left(a_\alpha \left(\frac{1}{\alpha} - \frac{1}{\beta} \right) \alpha^k \right).$$

Set

$$S_L = \sum_{0 \leq k \leq 7} 2p_k \operatorname{Re} \left(a_\alpha \left(\frac{1}{\alpha} - \frac{1}{\beta} \right) \alpha^k \right), \quad S_R = \sum_{k \geq 8} 2p_k \operatorname{Re} \left(a_\alpha \left(\frac{1}{\alpha} - \frac{1}{\beta} \right) \alpha^k \right).$$

Then we have

$$S_L - |S_R| \leq S_L + S_R = |u_0 \cdots u_{N-1}|_0 - \frac{N}{\beta} \leq S_L + |S_R|. \quad (3)$$

Choosing the p_k so that $p_k = 1$ if and only if $\operatorname{Re} \left(a_\alpha \left(\frac{1}{\alpha} - \frac{1}{\beta} \right) \alpha^k \right)$ is positive, and similarly for negative values, it can be seen that

$$-0.42 < S_L < 0.73. \quad (4)$$

To estimate the magnitude of $|S_R|$, we compute

$$\left| \frac{1}{\alpha} - \frac{1}{\beta} \right| = 1.72457 \dots \quad (5)$$

Thus

$$|S_R| \leq 2|a_\alpha| \cdot \left| \frac{1}{\alpha} - \frac{1}{\beta} \right| \sum_{k \geq 8} |\alpha|^k = 2|a_\alpha| \cdot \left| \frac{1}{\alpha} - \frac{1}{\beta} \right| \frac{|\alpha|^8}{1 - |\alpha|} < 0.17. \quad (6)$$

The claim follows from inequalities (3), (4), and (6). \square

Proposition 2.5. *We have*

$$-0.775 < |u_0 u_1 \cdots u_{N-1}|_1 - \frac{N}{\beta^2} < 0.725.$$

Proof. Analogously to the proof of Prop. 2.4, set

$$S_L = \sum_{0 \leq k \leq 10} 2p_k \operatorname{Re} \left(a_\alpha \left(\frac{1}{\alpha^2} - \frac{1}{\beta^2} \right) \alpha^k \right), \quad S_R = \sum_{k \geq 11} 2p_k \operatorname{Re} \left(a_\alpha \left(\frac{1}{\alpha^2} - \frac{1}{\beta^2} \right) \alpha^k \right).$$

Again, by choosing the p_k appropriately, it can be verified that

$$-0.70 < S_L < 0.65.$$

To find an upper bound for $|S_R|$, we compute

$$\left| \frac{1}{\alpha^2} - \frac{1}{\beta^2} \right| = 1.96298 \dots \quad (7)$$

Therefore,

$$|S_R| \leq 2|a_\alpha| \cdot \left| \frac{1}{\alpha^2} - \frac{1}{\beta^2} \right| \sum_{k \geq 11} |\alpha|^k = 2|a_\alpha| \cdot \left| \frac{1}{\alpha^2} - \frac{1}{\beta^2} \right| \frac{|\alpha|^{11}}{1 - |\alpha|} < 0.075.$$

These inequalities imply the claim. \square

Proposition 2.6. *We have*

$$-0.88 < |u_0 u_1 \cdots u_{N-1}|_2 - \frac{N}{\beta^3} < 0.62.$$

Proof. Once again we set

$$S_L = \sum_{0 \leq k \leq 13} 2p_k \operatorname{Re} \left(a_\alpha \left(\frac{1}{\alpha^3} - \frac{1}{\beta^3} \right) \alpha^k \right), \quad S_R = \sum_{k \geq 14} 2p_k \operatorname{Re} \left(a_\alpha \left(\frac{1}{\alpha^3} - \frac{1}{\beta^3} \right) \alpha^k \right).$$

As above, it can be verified that

$$-0.8371 < S_L < 0.5764.$$

We compute

$$\left| \frac{1}{\alpha^3} - \frac{1}{\beta^3} \right| = 2.33887 \dots, \quad (8)$$

and thus

$$|S_R| \leq 2|a_\alpha| \cdot \left| \frac{1}{\alpha^3} - \frac{1}{\beta^3} \right| \sum_{k \geq 14} |\alpha|^k = 2|a_\alpha| \cdot \left| \frac{1}{\alpha^3} - \frac{1}{\beta^3} \right| \frac{|\alpha|^{14}}{1 - |\alpha|} < 0.0354.$$

The claim follows from these inequalities. \square

Propositions 2.4–2.6 together with Lemma 2.3 imply that \mathbf{t} is 2-balanced.

Remark 2.7. None of the inequalities in the claims of Propositions 2.4–2.6 are optimal.

3 Abelian complexity

In this section we study the Abelian complexity of the Tribonacci word and prove Theorem 1.4. We begin by recalling the following basic fact from [20]:

Fact 3.1. *If an infinite word ω has two factors u and v of the same length n for which the i th entry of the Parikh vector are p and $p+c$ respectively, then for each $\ell = 0, \dots, c$, there exists a factor u_ℓ of ω of length n such that the i th entry of $\Psi(u_\ell)$ is equal to $p + \ell$.*

Henceforth, we denote ρ the Abelian complexity of the Tribonacci word \mathbf{t} .

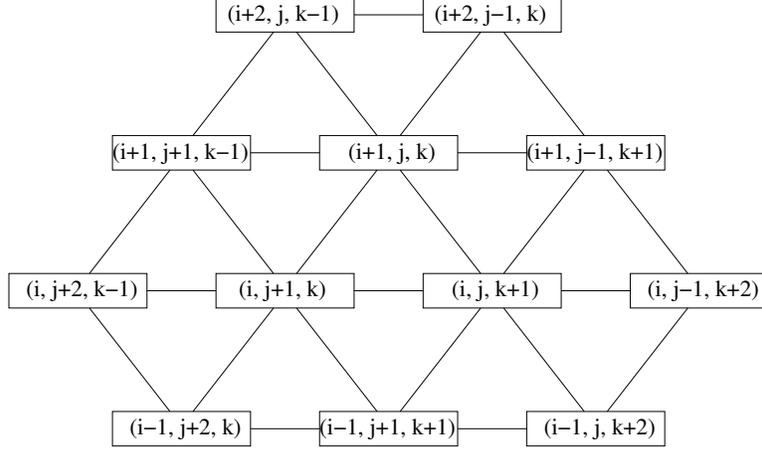


Figure 1: Links between Parikh vectors

3.1 Proof of Theorem 1.4

Let us recall that a factor u of an infinite word ω is *right special* (resp. *left special*) if there exist distinct letters a and b such that the words ua and ub (resp. au , bu) are both factors of ω . A factor which is both left and right special is called *bispecial*.

It is well-known that for every $n \geq 1$, \mathbf{t} has exactly one right special factor of length $n - 1$, and that, for this special factor that we denote $\mathbf{t}_{n-1}^<$, the three words $\mathbf{t}_{n-1}^<0$, $\mathbf{t}_{n-1}^<1$, and $\mathbf{t}_{n-1}^<2$ are each factors of \mathbf{t} of length n . Writing $\Psi(\mathbf{t}_{n-1}^<) = (i, j, k)$, where i, j, k are each non-negative integers, we define

$$\text{Central}(n) = \{(i + 1, j, k), (i, j + 1, k), (i, j, k + 1)\}.$$

We have

$$\text{Central}(n) \subseteq \Psi_{\mathbf{t}}(n). \tag{9}$$

Given a vector $\vec{v} = (\alpha, \beta, \gamma)$, set $\|\vec{v}\| = \max(|\alpha|, |\beta|, |\gamma|)$. The set of vectors \vec{v} such that $\|\vec{v} - \vec{u}\| \leq 2$ for all \vec{u} in $\text{Central}(n)$ is given by the graph in Figure 1 (whose vertices are vectors, and where two vertices \vec{u}, \vec{v} are joined by an undirected edge if and only if $\|\vec{v} - \vec{u}\| = 1$).

Since \mathbf{t} is 2-balanced, $\Psi_{\mathbf{t}}(n)$ is a subset of this set of twelve vectors. Moreover for the same reason, we should have $\|\vec{v} - \vec{u}\| \leq 2$ for all \vec{u}, \vec{v} in $\Psi_{\mathbf{t}}(n)$. This implies that the only possibility for $\Psi_{\mathbf{t}}(n)$ is to be a subset of one of the three sets delimited by a regular hexagon in Figure 1, or one of the three sets delimited by an equilateral triangle of base length 2. These sets have cardinalities 7 and 6 respectively showing that $\rho_{\omega}(n) \leq 7$. \square

Remark 3.2. By computer simulation we find that

$$(\rho(n))_{n \geq 1} = 3343444344444434444444444444345544444554444 \dots$$

In particular, the least n for which $\rho(n) = 5$ is for $n = 30$. We also found that the smallest n for which $\rho(n) = 6$ is $n = 342$, and the smallest n for which $\rho(n) = 7$ is $n = 3914$. The next four values of n for which $\rho(n) = 7$ are $n = 4063, 4841, 4990, 7199$.

3.2 More on the Abelian complexity of the Tribonacci Word

In this section, we characterize those n for which $\rho_{\mathbf{t}}(n) = 3$.

We continue to adopt the notation $\mathbf{t}_{n-1}^<$, $\Psi(\mathbf{t}_{n-1}^<)$, $\text{Central}(n)$, and $\|\vec{v}\|$ introduced in the proof of Theorem 1.4. Let $B(n)$ be the set of vectors \vec{v} for which there is exactly one \vec{u} in $\text{Central}(n)$ with $\|\vec{v} - \vec{u}\| = 2$:

$$B(n) = \{(i-1, j+1, k+1), (i+1, j-1, k+1), (i+1, j+1, k-1)\}. \quad (10)$$

Proposition 3.3. *The following are equivalent:*

- (1) For all factors v and w of \mathbf{t} of length n , we have $||v|_a - |w|_a| \leq 1$ for all $a \in \{0, 1, 2\}$.
- (2) $\rho_{\mathbf{t}}(n) = 3$.
- (3) $\Psi_{\mathbf{t}}(n) \cap B(n) = \emptyset$.
- (4) \mathbf{t} contains a bispecial special factor of length $n - 1$.
- (5) $n = 1$, or $n = \frac{1}{2}(T_m + T_{m+2} - 1)$ for $m \geq 0$, where $(T_m)_{m \geq 0} = 1, 2, 4, 7, 13, 24, \dots$ denotes the sequence of Tribonacci numbers.

Proof. Equivalence (1) \Leftrightarrow (2) is an immediate consequence of Eq. (9). To see the equivalence between items (4) and (5), we note that the bispecial factors of \mathbf{t} are precisely the palindromic prefixes of \mathbf{t} (see [21]). The lengths of these words are known to be

$$\frac{T_m + (T_{m-1} + T_m) + (T_{m-2} + T_{m-1} + T_m) - 3}{2} = \frac{T_m + T_{m+2} - 3}{2}$$

(see Corollary 3.11 in [7] or Corollary 3.3 in [23]). It follows that (4) \Leftrightarrow (5). Thus it remains to prove that items (2), (3), and (4) are equivalent. We will make use of the following lemmas.

Lemma 3.4. *Let $n \geq 1$, and let v be a factor of \mathbf{t} . Put $\Psi(\mathbf{t}_{n-1}^<) = (i, j, k)$, and $\Psi(v) = (i', j', k')$. Suppose either*

- $|v| \geq n$ and either $i' \leq i - 1$ or $j' \leq j - 1$

or

- $|v| \leq n$, and either $i' \geq i + 2$ or $j' \geq j + 2$.

Then $\rho(n) > 3$.

Proof. The first condition implies the existence of a factor v' of \mathbf{t} of length n with either $|v'|_0 \leq i - 1$ or $|v'|_1 \leq j - 1$. By (9) it follows that the factors of \mathbf{t} of length n are not balanced. Similarly, the second condition implies the existence of a factor v' of \mathbf{t} of length n with either $|v'|_0 \geq i + 2$ or $|v'|_1 \geq j + 2$. Again by (9) it follows that the factors of \mathbf{t} of length n are not balanced. \square

For each $n \geq 1$, let $\phi(n) = |\tau(\mathbf{t}_{n-1}^{\leq})0| + 1$.

Lemma 3.5. *If $\rho(n) = 3$, then $\rho(\phi(n)) = 3$.*

Proof. Assume $\rho(n) = 3$, and set $N = \phi(n)$. Let $\Psi(\mathbf{t}_{n-1}^{\leq}) = (i, j, k)$, and $\mathbf{t}_{N-1}^{\leq} = (I, J, K)$. Since $\mathbf{t}_{N-1}^{\leq} = \tau(\mathbf{t}_{n-1}^{\leq})0$, it follows that

$$I = n; \quad J = i; \quad K = j. \quad (11)$$

Now we have

$$\Psi_{\mathbf{t}}(n) = \{(i+1, j, k), (i, j+1, k), (i, j, k+1)\}$$

and

$$\{(I+1, J, K), (I, J+1, K), (I, J, K+1)\} \subseteq \Psi_{\mathbf{t}}(N)$$

and we must show that in fact we have equality in the previous containment. By Theorem 1.3, the only other possible Parikh vectors of factors of \mathbf{t} of length N are one of the following six vectors

$$\begin{pmatrix} I-1 \\ J \\ K+2 \end{pmatrix}; \begin{pmatrix} I-1 \\ J+2 \\ K \end{pmatrix}; \begin{pmatrix} I \\ J-1 \\ K+2 \end{pmatrix}; \begin{pmatrix} I \\ J+2 \\ K-1 \end{pmatrix}; \begin{pmatrix} I+2 \\ J-1 \\ K \end{pmatrix}; \begin{pmatrix} I+2 \\ J \\ K-1 \end{pmatrix}$$

or one of the following three vectors

$$\begin{pmatrix} I+1 \\ J+1 \\ K-1 \end{pmatrix}; \begin{pmatrix} I+1 \\ J-1 \\ K+1 \end{pmatrix}; \begin{pmatrix} I-1 \\ J+1 \\ K+1 \end{pmatrix}.$$

Using Lemma 3.4 we will show that if any one of these nine vectors belong to $\Psi_{\mathbf{t}}(n)$, then $\rho(n) > 3$, a contradiction. Hence $\rho(N) = 3$. We proceed one vector at a time.

◇ Suppose $(I-1, J, K+2) \in \Psi_{\mathbf{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathbf{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I-1 &= |v| + \delta \\ J &= i' \\ K+2 &= j'. \end{aligned}$$

It follows from (11) that $|v| = n-1-\delta \leq n$, and $j' = j+2$ which contradicts Lemma 3.4.

◇ Suppose $(I-1, J+2, K) \in \Psi_{\mathbf{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathbf{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I-1 &= |v| + \delta \\ J+2 &= i' \\ K &= j'. \end{aligned}$$

It follows from (11) that $|v| = n-1-\delta \leq n$, and $i' = i+2$ which contradicts Lemma 3.4.

◇ Suppose $(I, J - 1, K + 2) \in \Psi_{\mathfrak{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathfrak{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I &= |v| + \delta \\ J - 1 &= i' \\ K + 2 &= j'. \end{aligned}$$

It follows from (11) that $|v| = n - \delta \leq n + 1$, $i' = i - 1$, and $j' = j + 2$ which contradicts Lemma 3.4.

◇ Suppose $(I, J + 2, K - 1) \in \Psi_{\mathfrak{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathfrak{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I &= |v| + \delta \\ J + 2 &= i' \\ K - 1 &= j'. \end{aligned}$$

It follows from (11) that $|v| = n - \delta \leq n + 1$, $i' = i + 2$, and $j' = j - 1$ which contradicts Lemma 3.4.

◇ Suppose $(I + 2, J - 1, K) \in \Psi_{\mathfrak{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathfrak{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I + 2 &= |v| + \delta \\ J - 1 &= i' \\ K &= j'. \end{aligned}$$

It follows from (11) that $|v| = n + 2 - \delta \geq n + 1$, and $i' = i - 1$ which contradicts Lemma 3.4.

◇ Suppose $(I + 2, J, K - 1) \in \Psi_{\mathfrak{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathfrak{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I + 2 &= |v| + \delta \\ J &= i' \\ K - 1 &= j'. \end{aligned}$$

It follows from (11) that $|v| = n + 2 - \delta \geq n + 1$, and $j' = j - 1$ which contradicts Lemma 3.4.

This concludes the first six cases. Now we consider the last three:

◇ Suppose $(I + 1, J + 1, K - 1) \in \Psi_{\mathfrak{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathfrak{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I + 1 &= |v| + \delta \\ J + 1 &= i' \\ K - 1 &= j'. \end{aligned}$$

It follows from (11) that $|v| = n + 1 - \delta \geq n$, and $j' = j - 1$ which contradicts Lemma 3.4.
 \diamond Suppose $(I + 1, J - 1, K + 1) \in \Psi_{\mathbf{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathbf{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I + 1 &= |v| + \delta \\ J - 1 &= i' \\ K + 1 &= j'. \end{aligned}$$

It follows from (11) that $|v| = n + 1 - \delta \geq n$, and $i' = i - 1$ which contradicts Lemma 3.4.

\diamond Suppose $(I - 1, J + 1, K + 1) \in \Psi_{\mathbf{t}}(N)$. By Lemma 2.1 there exists a factor v of \mathbf{t} with $\Psi(v) = (i', j', k')$ and $\delta \in \{-1, 0, 1\}$ such that

$$\begin{aligned} I - 1 &= |v| + \delta \\ J + 1 &= i' \\ K + 1 &= j'. \end{aligned}$$

It follows from (11) that $|v| = i' + j' + k' = n - 1 - \delta = i + j + k - \delta$, $i' = i + 1$, and $j' = j + 1$. Thus we have $k' = k - 2 - \delta$. If $\delta = -1$, then $(i + 1, j + 1, k - 1) \in \Psi_{\mathbf{t}}(n)$ contradicting that $\rho(n) = 3$. If $\delta = 0$, then $|v| = n - 1$ and $k' = k - 2$. Thus any factor of \mathbf{t} of length n beginning in v will have at most $k - 1$ many 2's, contradicting that the balance between any two factors of length n is at most one (there is a factor of length n with Parikh vector $(i, j, k + 1)$). If $\delta = 1$, then $|v| = n - 2$ and $k' = k - 3$. Thus any factor of length n beginning in v will have at most $k - 2$ many 2's, contradicting Theorem 1.3.

This concludes the proof of the lemma. \square

Lemma 3.6. *If $\Psi_{\mathbf{t}}(n) \cap B(n) \neq \emptyset$, then $\Psi_{\mathbf{t}}(\phi(n)) \cap B(\phi(n)) \neq \emptyset$.*

Proof. Since $\Psi(\mathbf{t}_{n-1}^<) = (i, j, k)$, we have that $\phi(n) = n + i + j + 1$. Following (10), if $\Psi(v) = (i - 1, j + 1, k + 1)$ for some factor v of \mathbf{t} of length n , then $\Psi(\tau(v)0) = (n + 1, i - 1, j + 1) \in \Psi_{\mathbf{t}}(\phi(n)) \cap B(\phi(n))$. If $\Psi(v) = (i + 1, j - 1, k + 1)$ for some factor v of \mathbf{t} of length n , then $\Psi(\tau(v)0) = (n + 1, i + 1, j - 1) \in \Psi_{\mathbf{t}}(\phi(n)) \cap B(\phi(n))$. Finally if $\Psi(v) = (i + 1, j + 1, k - 1)$ for some factor v of \mathbf{t} of length n , then $\Psi(0^{-1}\tau(v)) = (n - 1, i + 1, j + 1) \in \Psi_{\mathbf{t}}(\phi(n)) \cap B(\phi(n))$. \square

Lemma 3.7. *If $\Psi_{\mathbf{t}}(n) \cap B(n) = \emptyset$, then $\mathbf{t}_{n-1}^<$ is a bispecial factor of \mathbf{t} .*

Proof. We proceed by induction on n . The result is clear for $n = 1, 2$. Let $N > 2$, and suppose the result is true for all $n < N$. Assume $\Psi_{\mathbf{t}}(N) \cap B(N) = \emptyset$. Set $\Psi(\mathbf{t}_{N-1}^<) = (I, J, K)$. Since $\mathbf{t}_{N-1}^<$ is right special, it ends to 0. We now show that $\mathbf{t}_{N-1}^<$ also begins in 0. Suppose $\mathbf{t}_{N-1}^<$ begins in 1. Then, $1^{-1}\mathbf{t}_{N-1}^<20$ is a factor of \mathbf{t} of length N , and $\Psi(1^{-1}\mathbf{t}_{N-1}^<20) = (I + 1, J - 1, K + 1) \in \Psi_{\mathbf{t}}(N) \cap B(N)$ contrary to our assumption. Similarly, assume $\mathbf{t}_{N-1}^<$ begins in 2. Then, $2^{-1}\mathbf{t}_{N-1}^<10$ is a factor of \mathbf{t} of length N , and $\Psi(2^{-1}\mathbf{t}_{N-1}^<10) = (I + 1, J + 1, K - 1) \in \Psi_{\mathbf{t}}(N) \cap B(N)$ contrary to our assumption.

Having established that $\mathbf{t}_{N-1}^<$ begins in 0, we can write $\mathbf{t}_{N-1}^< = \tau(u)0$ for some factor u of \mathbf{t} . Also, as $\mathbf{t}_{N-1}^<$ is right special, so is u and hence $u = \mathbf{t}_{n-1}^<$ for some n , and so $N = \phi(n)$. It follows from the previous lemma that $\Psi_{\mathbf{t}}(n) \cap B(n) = \emptyset$. Hence by induction hypothesis we have that $\mathbf{t}_{n-1}^<$ is bispecial, and hence so is $\tau(\mathbf{t}_{n-1}^<)0 = \mathbf{t}_{N-1}^<$. \square

We are now ready to establish the remaining equivalences in 3.3: We will show that (3) \Rightarrow (4) \Rightarrow (2) \Rightarrow (3). The previous lemma states precisely that (3) \Rightarrow (4). That (2) \Rightarrow (3) is clear from (9) and (10). Finally to see that (4) \Rightarrow (2), we proceed by induction on n . The result is clear for $n = 1, 2$. Let $N > 2$ and suppose the result is true for all $n < N$. Assume $\mathbf{t}_{N-1}^<$ is bispecial. Then as $\mathbf{t}_{N-1}^<$ begins and ends in 0, we can write $\mathbf{t}_{N-1}^< = \tau(u)0$ for some factor u of \mathbf{t} . Moreover as $\mathbf{t}_{N-1}^<$ is bispecial, so is u . It follows that $u = \mathbf{t}_{n-1}^<$ for some n , and hence $N = \phi(n)$. By induction hypothesis we have that $\rho(n) = 3$. It now follows from Lemma 3.5 that $\rho(N) = 3$ as required. This completes our proof of Proposition 3.3. \square

As an immediate consequence of Proposition 3.3, we have that $\rho(n) = 3$ for infinitely many values of n . We next show that the same is true for $\rho(n) = 7$.

Proposition 3.8. *The Abelian complexity of the Tribonacci word attains the value 7 infinitely often.*

Proof. Note that \mathbf{t} begins with the square $(0102010)^2 = \tau^3(00)$ and so with the squares $\tau^{n+3}(00)$. Consider any integer $n \geq 0$ for which all factors of length $m = 3914$ occur in $\tau^{n+3}(0)$ (we recall that a computer computation showed that $\rho(3914) = 7$). For any factor y of length m there exists a conjugate x of $\tau^{n+3}(0)$ such that xy is a factor of $\tau^{n+3}(00)$. Since $\Psi(x) = \Psi(\tau^{n+3}(0))$ and since $\Psi(xy) = \Psi(x) + \Psi(y)$, we deduce that $\text{Card}\Psi_{\mathbf{t}}(|\tau^{n+3}(0)| + m) \geq \text{Card}\Psi_{\mathbf{t}}(m) = 7$. Since the Abelian complexity of the Tribonacci word is bounded by 7, we get $\text{Card}\Psi_{\mathbf{t}}(|\tau^{n+3}(0)| + m) = 7$. \square

4 Conclusion

Various aspects of the Abelian complexity of the Tribonacci word remain a mystery. For instance, it is surprising to us that the value $\rho_{\mathbf{t}}(n) = 7$ does not occur until $n = 3914$, but then re-occurs relatively shortly thereafter. As another example, it is verified that for all $n \leq 184$, if U and V are factors of \mathbf{t} of length n , with U a prefix of \mathbf{t} , then $||U|_a - |V|_a| \leq 1$, for all $a \in \{0, 1, 2\}$. But then this property fails for $n = 185$. We list a few related open questions:

Open problem 1. Does the Abelian complexity of the Tribonacci word attain each value in $\{4, 5, 6\}$ infinitely often?

Open problem 2. For each value $m \in \{4, 5, 6, 7\}$, characterize those n for which $\rho(n) = m$.

Another problem concerns the m -bonacci word, the generalization of the Tribonacci word (and of the celebrated Fibonacci word) to an alphabet on m letters. This word is defined

as the fixed point of the morphism $\tau_m : \{0, \dots, m - 1\}^* \rightarrow \{0, \dots, m - 1\}^*$ defined by $\tau_m(i) = 0(i + 1)$ for $0 \leq i \leq m - 2$ and $\tau(m - 1) = 0$.

Open problem 3. Prove or disprove that the m -bonacci word is $(m - 1)$ -balanced.

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