Extremal mixed metric dimension with respect to the cyclomatic number

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Abstract

In a graph G, the cardinality of the smallest ordered set of vertices that distinguishes every element of $V(G) \cup E(G)$ is called the mixed metric dimension of G, and it is denoted by $\operatorname{mdim}(G)$. In [12] it was conjectured that for a graph G with cyclomatic number c(G) it holds that $\operatorname{mdim}(G) \leq L_1(G) + 2c(G)$ where $L_1(G)$ is the number of leaves in G. It is already proven that the equality holds for all trees and more generally for graphs with edge-disjoint cycles in which every cycle has precisely one vertex of degree ≥ 3 . In this paper we determine that for every Theta graph G, the mixed metric dimension $\operatorname{mdim}(G)$ equals 3 or 4, with 4 being attained if and only if G is a balanced Theta graph. Thus, for balanced Theta graphs the above inequality is also tight. We conclude the paper by further conjecturing that there are no other graphs, besides the ones mentioned here, for which the equality $\operatorname{mdim}(G) = L_1(G) + 2c(G)$ holds.

1 Introduction

Let G be a simple connected graph with n vertices and m edges. The distance between a pair of vertices $u, v \in V(G)$ is defined as the length of the shortest path connecting u and v in G and is denoted by $d_G(u, v)$. The distance between a vertex $u \in V(G)$ and an edge $e = vw \in E(G)$ is defined by $d_G(u, e) = d_G(u, vw) = \min\{d_G(u, v), d_G(u, w)\}$. For both these distances we simply write d(u, v) and d(u, e) if no confusion arises. We say that a vertex $s \in V(G)$ distinguishes (or resolves) a pair $x, x' \in V(G) \cup E(G)$ if $d(s, x) \neq d(s, x')$. We say that a set $S \subseteq V(G)$ is a mixed metric generator if every pair $x, x' \in V(G) \cup E(G)$ is distinguished by at least one vertex from S. The cardinality of the smallest mixed metric generator is called the mixed metric dimension of G, and it is denoted by mdim(G). The notion of the mixed metric dimension is the natural generalization of the notions of the vertex metric dimension and the edge metric dimension which are defined as the cardinality of the smallest set of vertices which distinguishes all pairs of vertices and all pairs of edges respectively. The notion of vertex metric dimension for graphs was independently introduced by [3] and [13], under the names resolving sets and locating sets, respectively. Even before, this notion was introduced for the realm of metric spaces [1]. The concept of vertex metric dimension was recently extended from resolving vertices to resolving edges of a graph by Kelenc, Tratnik and Yero [5], which lead to the definition of the edge metric dimension. Finally, it was further extended to resolving mixed pairs of edges and vertices by Kelenc, Kuziak, Taranenko, and Yero [6] which resulted with the notion of the mixed metric dimension. All these variations of metric dimensions attracted interest (see [9, 11, 12, 15, 16]), while for a wider and systematic introduction of the topic metric dimension that encapsulates all three above mentioned variations, we recommend the PhD thesis of Kelenc [4].

In literature, among other questions, the mixed metric dimension of trees, unicyclic graphs and graphs with edge disjoint cycles was studied. Denoting by $L_1(G)$ the number of leaves in a graph G, we first cite the following result from [6].

Proposition 1 For every tree T, it holds

$$\mathrm{mdim}(T) = L_1(T).$$

A graph in which all cycles are pairwise edge disjoint is called a *cactus graph*. Having that in mind, the following results were proven in [12], first for unicyclic graphs and after for all cactus graphs.

Theorem 2 Let $G \neq C_n$ be a cactus graph with c cycles. Then

$$\mathrm{mdim}(G) \le L_1(G) + 2c,$$

and the upper bound is attained if and only if every cycle in G has exactly one vertex of degree ≥ 3 .

The cyclomatic number of a graph G is defined by c(G) = m - n + 1. As the number of cycles in trees and graphs with edge disjoint cycles equals the cyclomatic number, this lead the authors of [12] to make the following conjecture.

Conjecture 3 Let $G \neq C_n$ be a graph, c(G) its cyclomatic number, and $L_1(G)$ the number of leaves in G. Then

$$\mathrm{mdim}(G) \le L_1(G) + 2c(G). \tag{1}$$

Notice that Proposition 1 and Theorem 2 imply that the equality in (1) holds for all cactus graphs in which every cycle has precisely one vertex of degree ≥ 3 (this includes all trees and unicyclic graphs with precisely one vertex on the cycle with degree ≥ 3). A natural question that arises is - are there any other graphs for which the equality in (1) holds? In this paper we will try to further clarify this question.

2 Preliminaries

The *(vertex) connectivity* $\kappa(G)$ of a graph G is the minimum size of a vertex cut, i.e. any subset of vertices $S \subseteq V(G)$ such that G - S is disconnected or has only one vertex. We say that a graph G is k-connected if $\kappa(G) \geq k$. As we are going to study the graphs for which the equality in (1) holds, it is useful to state the following result from [12].

Proposition 4 Let G be a 3-connected graph. Then mdim(G) < 2c(G).

This proposition implies that equality in (1) may hold only for graphs with $\kappa(G) = 1$ (beside cactus graphs in which every cycle has precisely one vertex of degree ≥ 3) and $\kappa(G) = 2$. A class of graphs with $\kappa(G) = 2$ which will be of interest to us are so called Theta graphs.

We say that a graph G is a *Theta* graph or a Θ -graph if G is a graph with two vertices u and v of degree 3 and all other vertices in G are of degree 2. We say that a Theta graph G is *balanced* if the lengths of all three paths connecting u and v differ by at most 1, otherwise we say that G is *unbalanced*. In this paper we will prove that the equality in (1) holds also for balanced Theta graphs. But, before we show that, we need to introduce the following notion which will be of use to us in the sequel.

Let G be a graph and let $S \subseteq V(G)$ be a set of vertices of a graph G. Any shortest path between two vertices from S is called a S-closed path. Let x and x' be a pair of elements from the set $V(G) \cup E(G)$. We say that a pair x and x' is enclosed by S if there is a S-closed path containing x and x'. We say that a pair x and x' is half-enclosed by S if there is a vertex $s \in S$ such that a shortest path from s to x contains x' or a shortest path from s to x' contains x.

Observation 5 Let G be a graph, let $S \subseteq V(G)$ be a set of vertices in G and let x and x' be a pair of elements from the set $V(G) \cup E(G)$. If x and x' are enclosed by S, then x and x' are distinguished by S. If x and x' are half-enclosed by S then x and x' are distinguished by S in all cases except possibly when x and x' are a pair consisting of a vertex and an edge which are incident to each other.

We say that a subgraph H of a graph G is an *isometric* subgraph, if for any two vertices $u, v \in V(H)$ it holds that $d_H(u, v) = d_G(u, v)$. The following notation for paths is used. Suppose that P is a path and $u, v \in V(P)$, then by P[u, v] we denote the subpath of P connecting vertices u and v, while by P(u, v) we denote $P[u, v] - \{u, v\}$. Notions P[u, v) and P(u, v] are also used and they denote the subpaths where only one of the end-vertices of P[u, v] is excluded.

3 Balanced Theta graphs

Notice that every Theta graph G has the cyclomatic number c(G) = 2. Also, for every Theta graph G the number of leaves equals zero, i.e. $L_1(G) = 0$. Therefore, for a Theta graph G, the equality in (1) will hold if and only if $\operatorname{mdim}(G) = 4$. In this section we will show that for balanced Theta graphs precisely that holds, i.e. $\operatorname{mdim}(G) = 4$ if and only if a Theta graph G is balanced. First we need the following lemma.

Lemma 6 Let G be a balanced Theta graph with vertices u and v of degree 3. Let $S \subseteq V(G)$ be a set of vertices in G such that |S| = 3 and S contains precisely one internal vertex from each of the three distinct paths connecting vertices u and v. Then S is not a mixed metric generator in G.

Proof. Let P_1 , P_2 and P_3 denote the three distinct paths in G connecting vertices u and v. For a pair of vertices $v_1 \in P_i$ and $v_2 \in P_j$, where $i \neq j$, we denote by $d_u(v_1, v_2)$ (resp. $d_v(v_1, v_2)$) the length of the shortest path connecting vertices v_1 and v_2 and which contains vertex u (resp. v). By C_{ij} we denote the cycle induced by paths P_i and P_j . Let $S = \{s_1, s_2, s_3\}$ be a set of vertices which contains precisely one internal vertex from each of the three distinct paths connecting vertices u and v, where the elements of S are denoted so that s_i belongs to P_i . Let P_{ij} and P'_{ij} be the two internally vertex disjoint paths connecting vertices s_i and s_j , which induce the cycle C_{ij} , denoted so that $|P_{ij}| \leq |P'_{ij}|$. If P_{12} and P_{13} do not share any other vertex besides s_1 , then either P_{12} and P_{23} share it or P_{13} and P_{23} do. Therefore, at least one pair of paths P_{ij} shares one more vertex besides vertices from S, say P_{12} and P_{13} .

Next, we distinguish the following three cases.

Case 1: $|P_{12}| < |P'_{12}|$ and $|P_{13}| < |P'_{13}|$. In this case let w be the neighbor of s_1 not contained in paths P_{12} and P_{13} , then s_1 and s_1w are not distinguished by S, so S is not a mixed metric generator.

Case 2: $|P_{12}| = |P'_{12}|$ and $|P_{13}| = |P'_{13}|$. In this case two edges incident to s_1 are not distinguished by S, so S cannot be a mixed metric generator.

Case 3: $|P_{12}| = |P'_{12}|$ and $|P_{13}| < |P'_{13}|$. First notice that $|P_{12}| = |P'_{12}|$ implies that the cycle C_{12} is even and the pair of vertices s_1 and s_2 is an antipodal pair on C_{12} . Since G is a balanced Theta graph, the fact that C_{12} is even implies $|P_1| = |P_2|$. Therefore, u and v are also an antipodal pair on C_{12} and it holds that

$$d(s_1, u) = d(s_2, v)$$
 and $d(s_1, v) = d(s_2, u).$ (2)

As we assumed that P_{12} and P_{13} share another vertex beside s_1 , notice that precisely one of the vertices u and v, say v, belongs to both P_{12} and P_{13} .

Claim A. If $d(s_1, v) > d(s_2, v)$, then $|P_{23}| < |P'_{23}|$.

To prove Claim A, let us assume $d(s_1, v) > d(s_2, v)$. Then (2) promptly implies $d(s_1, u) < d(s_2, u)$. Also, since P_{13} leads through the vertex v, we have

$$d(s_1, v) + d(v, s_3) \le d(s_1, u) + d(u, s_3).$$

Therefore, we obtain

$$d_v(s_2, s_3) = d(s_2, v) + d(v, s_3) < d(s_1, v) + d(v, s_3)$$

$$\leq d(s_1, u) + d(u, s_3) < d(s_2, u) + d(u, s_3) = d_u(s_2, s_3).$$

This means that P_{23} leads through v, while P'_{23} leads through u, and $|P_{23}| < |P'_{23}|$. So, the claim is proven.

In the light of Claim A, notice that either $d(s_1, v) \leq d(s_2, v)$ or $d(s_1, v) > d(s_2, v)$. If $d(s_1, v) > d(s_2, v)$, then by Claim A we have $|P_{23}| < |P'_{23}|$, so switching indices 1 and 2 reduces this case to the case $d(s_1, v) \leq d(s_2, v)$. Therefore, without loss of generality, we may assume that $d(s_1, v) \leq d(s_2, v)$. From this and (2) we immediately obtain

$$d(s_1, v) \le d(s_2, v) = d(s_1, u). \tag{3}$$

Therefore, there must exist a vertex v' on P_1 , distinct from v, such that $d(v, s_1) = d(s_1, v')$. Let a and b be vertices on P_3 such that $d(v, a) = d(v, s_1)$ and $d(u, b) = d(v, s_1)$. This situation is illustrated by Figure 1 a) and we will further consider the position of s_3 on P_3 .



Figure 1: In the proof of Lemma 6: a) the position of vertices s_1 , s_2 , v', a and b; b) the contradictory path P_{13} when $d(v, s_3) > d(v, b)$; c) the undistinguished pair s_3 and s_3w when $d(v, s_3) < d(v, a)$.

Claim B. It holds that $d(v, s_3) \leq d(v, b)$.

Asume the contrary, i.e. $d(v, s_3) > d(v, b)$ and notice that

$$d_u(s_1, s_3) < |P_1| - d(v, s_1) + d(u, b) = |P_1|,$$

$$d_v(s_1, s_3) = d(v, s_1) + d(v, s_3) > d(v, s_1) + d(v, b) = d(v, s_1) + |P_3| - d(u, b) = |P_3|$$

Now, there are two possibilities, either $|P_3| \ge |P_1|$ or $|P_3| < |P_1|$. Assuming $|P_3| \ge |P_1|$ leads us to the conclusion that $d_v(s_1, s_3) > |P_3| \ge |P_1| = d_u(s_1, s_3)$, which means that the shortest path P_{13} from s_1 to s_3 contains u which is a contradiction (see Figure 1 b)). Assuming the other possibility $|P_3| < |P_1|$, leads us to $d_v(s_1, s_3) > |P_3| = |P_1| - 1 = d_u(s_1, s_3) - 1$, i.e. $d_v(s_1, s_3) \ge d_u(s_1, s_3)$. Therefore, either $d_v(s_1, s_3) > d_u(s_1, s_3)$ which again leads to the conclusion that P_{13} contains u which is a contradiction, or $d_v(s_1, s_3) = d_u(s_1, s_3)$ which implies $|P_{13}| = |P'_{13}|$ which contradicts the assumption of this case. Therefore, we have proven that assumption $d(v, s_3) > d(v, b)$ always leads to a contradiction, so Claim B is proven.

Claim C. If $d(v, s_3) < d(v, a)$, then S is not a mixed metric generator.

Assume $d(v, s_3) < d(v, a)$ and let w denote the neighbor of s_3 such that $d(w, v) > d(s_3, v)$ as is illustrated by Figure 1 c). Let $x = s_3$ and $x' = s_3w$ and notice that x and x' are not distinguished by s_3 . To prove that x and x' are not distinguished by s_2 either, recall that the vertex a was chosen so that $d(v, a) = d(v, s_1) = d(u, s_2)$. Having that in mind, notice that

$$d_u(w, s_2) \ge d(a, u) + d(u, s_2) = |P_3| - d(a, v) + d(u, s_2) = |P_3|,$$

$$d_v(w, s_2) \le d(a, v) + d(v, s_2) = d(a, v) + |P_2| - d(u, s_2) = |P_2|,$$

so we have

$$d_v(w, s_2) - d_u(w, s_2) \le |P_2| - |P_3| \le 1.$$

If $d_v(w, s_2) - d_u(w, s_2) \leq 0$ then the shortest path from both x and x' to s_2 leads through v, so they are not distinguished by s_2 . Otherwise, if $d_v(w, s_2) - d_u(w, s_2) = 1$, that implies $d_u(w, s_2) < d_v(w, s_2)$ and $d_v(s_3, s_2) + 1 = d_u(s_3, s_2)$, from which we further obtain

$$d(x', s_2) = d(w, s_2) = d_u(w, s_2)$$

$$d(x, s_2) = d(s_3, s_2) = d_v(s_3, s_2) = d_v(w, s_2) - 1 = d_u(w, s_2)$$

so x and x' are again not distinguished by s_2 . The inequality (3) now implies that x and x' are not distinguished by s_1 either. We conclude that x and x' are not distinguished by S, so the claim is established.

Given Claims B and C, the only remaining possibility is $d(v, a) \leq d(v, s_3) \leq d(v, b)$. Recall that $|P_1| = |P_2|$ and let $|P_3| = |P_1| + r$ where the fact that G is balanced implies that r can take only values -1, 0, 1. Let us further denote

$$q = \frac{d(v', u) + d(u, s_3) - d(v, s_3)}{2} = \frac{2|P_1| - 2d(v, s_1) + r - 2d(v, s_3)}{2}.$$
 (4)

Notice that the second expression for q in (4) implies that q is integer if and only if r = 0. We also want to use the first expression for q in (4) to derive a bound on d(v', u) from it. For that purpose notice that

$$d(u, s_3) - d(v, s_3) = |P_3| - 2d(s_3, v) = |P_1| + r - 2d(v, a) - 2d(s_3, a)$$

$$\leq |P_1| + r - 2d(v, a) = |P_1| + r - 2d(v, s_1) = d(v', u) + r$$

with equality holding if and only if $d(s_3, a) = 0$. Plugging this in (4) yields

$$q \le \frac{d(v', u) + d(v', u) + r}{2}$$

which further implies $d(v', u) \ge q - \frac{r}{2}$. Let w be the vertex from $P_1[v', u]$ such that $d(v', w) = \lfloor q \rfloor$. Notice that such a vertex w must exist on P_1 because the fact that q is integer only for r = 0 implies

$$d(v',w) = \lfloor q \rfloor \le q - \frac{r}{2} \le d(v',u).$$

Moreover, notice that in the case when r = -1 even stricter upper bound for d(v', w) holds, i.e. $d(v', w) \leq q - \frac{r}{2} - 1 \leq d(v', u) - 1$. Now, let w' be the vertex on P_2 such that $d(w', v) = \lfloor q \rfloor$.

Claim D. It holds that $d_u(w, s_3) < d_v(w, s_3)$.

Notice that

$$d_u(w, s_3) = d(w, u) + d(u, s_3) = d(v', u) - d(w, v') + d(u, s_3),$$

$$d_v(w, s_3) = d(w, v') + 2(v, s_1) + d(v, s_3).$$

Therefore, we have

$$d_u(w, s_3) - d_v(w, s_3) = d(v', u) - 2d(w, v') + d(u, s_3) - 2(v, s_1) - d(v, s_3).$$

Now the fact that $d(v', w) = \lfloor q \rfloor$ and the definition of q further imply

$$d_u(w, s_3) - d_v(w, s_3) = -2 \lfloor q \rfloor - 2(v, s_1) + 2q < 0$$

which concludes the proof of Claim D.

Claim E. It holds that $d_v(w', s_3) < d_u(w', s_3)$.

Notice that

$$d_v(w', s_3) = d(w', v) + d(v, s_3) = \lfloor q \rfloor + d(v, s_3), d_u(w', s_3) = d(w', u) + d(u, s_3) = |P_2| - \lfloor q \rfloor + d(u, s_3).$$

Therefore, having in mind that $|P_2| = |P_1|$ we further have

$$d_v(w', s_3) - d_u(w', s_3) = 2 \lfloor q \rfloor + d(v, s_3) - |P_1| - d(u, s_3) = 2 \lfloor q \rfloor + d(v, s_3) - d(v', u) - 2d(v, s_1) - d(u, s_3)$$

from which, given the definition of q, we obtain

$$d_v(w', s_3) - d_u(w', s_3) = 2 \lfloor q \rfloor - 2d(v, s_1) - 2q < 0$$

which proves Claim E.

Now we distinguish the following three subcases with respect to the value of r.

Subcase 3.a: r = -1. Recall that in this case $d(v', w) \leq d(v', u) - 1$, which implies $w \neq u$, so there exists a neighbor z of the vertex w on P_1 , which is further from s_1 than w. Also, let z' be the neighbor of w' on P_2 which is further from v than w'. We want to prove that the edges x = wz and x' = w'z' are not distinguished by S, see Figure 2 a) for illustration. First note that

$$d(x, s_1) - d(x', s_1) = \lfloor q \rfloor + d(v', s_1) - (\lfloor q \rfloor + d(v, s_1)) = 0$$

which implies that s_1 does not distinguish x and x'. Since s_1 and s_2 are an antipodal pair of vertices on the even cycle C_{12} , if s_1 does not distinguish x and x' then s_2 does not distinguish them either. The only remaining possibility is for s_3 to distinguish x and x', but Claims D and E here imply

$$d(x, s_3) - d(x', s_3) = d(v', u) - \lfloor q \rfloor - 1 + d(u, s_3) - (\lfloor q \rfloor + d(v, s_3)),$$

where the definition of q further implies

$$d(x, s_3) - d(x', s_3) = -2 \lfloor q \rfloor - 1 + 2q = 0$$

from which we conclude that s_3 does not distinguish x and x' either. Therefore, x and x' are not distinguished by S, which means that S cannot be a mixed metric generator.



Figure 2: With the proof of Lemma 6, the undistinguished pair x and x' when $d(v, a) \le d(v, s_3) \le d(v, b)$ and a) r = -1; b) r = 0; c) r = 1.

Subcase 3.b: r = 0. In this subcase vertices x = w and x' = w' are not distinguished by S (see Figure 2 b)), which is shown similarly as in Subcase 2.a, so S cannot be a mixed metric generator.

Subcase 3.c: r = 1. In this subcase notice that w = u if and only if $d(s_3, a) = 0$, i.e. $s_3 = a$. When $s_3 = a$ then $x = s_3$ and $x' = s_3 z$ are not distinguished by S, so we may assume $s_3 \neq a$ which implies $w \neq u$, so there is a neighbor z of w on P_1 further from s_1 than w. Now, let z' be the neighbor of w' on P_2 which is further from v than w'. Then x = wz and x' = w'z' are not distinguished by S (see Figure 2 c)), which is shown by a similar calculation as in Subcase 3.a.

We will use Lemma 6 to prove the exact value of the mixed metric dimension of balanced Theta graphs.

Theorem 7 If G is a balanced Theta graph, then mdim(G) = 4.

Proof. Let u and v be the two vertices in G of degree 3 and let w and z be two neighbors of v. We want to prove that $S = \{u, v, w, z\}$ is a mixed metric generator. Let x and x' be a pair of elements from $V(G) \cup E(G)$. Let P_1 , P_2 and P_3 be the three distinct paths

connecting vertices u and v in G. If x and x' belong to a same path P_i then they are certainly distinguished by $u \in S$ or by $v \in S$. Assume therefore that x belongs to P_i and x' belongs to P_j where $i \neq j$. If x and x' are distinguished by u or v, then the proof is done, so we may assume that they are not distinguished by either u or v. If C_{ij} is even, this implies x and x' are a pair of vertices or a pair of edges, while in the case of odd cycle C_{ij} the pair x and x' must be a mixed pair consisting of a vertex and an edge. Notice that at least one of the paths P_i and P_j must contain at least one more vertex from Sbesides u and v, say P_i contains w. Then w certainly distinguishes x and x', so we proved that S is a mixed metric generator in G which implies mdim $(G) \leq 4$.

To conclude the proof, we still need to prove that any set S with |S| < 4 cannot be a mixed metric generator. Asume first that both u and v are contained in S, then |S| < 4 implies that two of the paths P_1 , P_2 and P_3 do not share an internal vertex with S, say P_i and P_j . Let u_i and u_j be the neighbors of vertex u on paths P_i and P_j respectively. If P_i and P_j are of the same length, then u_i and u_j are not distinguished by S, so S cannot be a mixed metric generator. If P_i and P_j are not of the same length, say $|P_i| < |P_j|$, then the fact that G is balanced implies $|P_i| = |P_j| - 1$, but this further implies u and uu_j are not distinguished by S.

Assume now that precisely one of the vertices u and v is contained in S, say u. This implies that there is a path P_i in G which does not share any other vertex with S besides u. Denote by u_i the neighbor of u on P_i . Since G is a balanced Theta graph it holds that $d(u_i, w) \ge d(u, w)$ for every internal vertex w from the path P_j , $j \ne i$. Therefore, u and uu_i are not distinguished by S.

Finaly, assume that neither u nor v are contained in S. If there is a path P_i which does not share an internal vertex with S, then u and uu_i are obviously not distinguished by S. Assume therefore that each path P_i shares at least one internal vertex with S, which together with the fact |S| < 4 further implies |S| = 3 and each P_i shares precisely one internal vertex with S. But then Lemma 6 implies S cannot be a mixed metric generator and the proof is established.

Since c(G) = 2 and $L_1(G) = 0$ holds for any Theta graph G, Theorem 7 immediately yields the following corollary.

Corollary 8 For a balanced Theta graph G, it holds that $mdim(G) = L_1(G) + 2c(G)$.

This result implies that Conjecture 3 holds for balanced Theta graphs, moreover it holds with equality in (1).

4 Unbalanced Theta graphs

To complete the results we will now prove that Conjecture 3 holds also for unbalanced Theta graphs, but for them the equality in (1) does not hold.

Lemma 9 Let G be an unbalanced Theta graph, then $mdim(G) \leq 3$.

Proof. Let u and v be the two vertices of degree 3 in G and let P_1 , P_2 and P_3 be three distinct paths in G connecting vertices u and v, where without loss of generality we may

assume that $|P_1| \leq |P_2| \leq |P_3|$. Since G is an unbalanced Theta graph, it follows that $|P_3| - |P_1| \geq 2$. By C_{ij} we denote the cycle induced by paths P_i and P_j . Let us now consider the cycle C_{13} and vertices u and v which belong to it. If the cycle C_{13} is even, then each of the two vertices u and v has precisely one antipodal vertex on C_{13} . On the other hand, if C_{13} is an odd cycle, then vertices u and v each have precisely two antipodal vertex.

Let us denote antipodal vertices of u on C_{13} by a_u and a'_u (where we assume $a_u = a'_u$ when C_{13} is an even cycle) and by a_v and a'_v the two antipodals of the vertex v. We may asume that the pair of antipodal vertices a_u and a'_u are denoted so that $d(v, a_u) \leq d(v, a'_u)$. Similarly, we denote vertices a_v and a'_v so that the inequality $d(u, a_v) \leq d(u, a'_v)$ holds. From $|P_3| - |P_1| \geq 2$ it follows that all of the vertices a_u, a'_u, a_v, a'_v belong to P_3 , where a_u and a'_u are distinct from v and, similarly a_v and a'_v are distinct from u. Further, let w be a vertex from the path P_2 such that distance from w to vertices u and v differs by at most one (i.e. w is the middle or "almost middle" vertex of the path P_2). Finally, we define the set $S = \{a_u, a_v, w\}$ for which we will prove that it is a mixed metric generator in G. All these vertices and the set S are illustrated by Figure 3.



Figure 3: In the proof of Lemma 9, the position of vertices w, a_u , a'_u , a_v , a'_v and the mixed metric generator $S = \{w, a_u, a_v\}$ in the case when both P_2 and C_{13} are: a) of even length, b) of odd length.

In order to prove that S is a mixed metric generator, let x and x' be a pair of elements from the set $V(G) \cup E(G)$. We distinguish the following three cases regarding the position of x and x' in G.

Case 1: Both x and x' belong to C_{13} . Notice that there are two subpaths of C_{13} which connect vertices a_u and a_v , one of them is $P_3[a_u, a_v]$, the other is $C_{13} - P_3(a_u, a_v)$. If x and x' belong to different subpaths of C_{13} connecting a_u and a_v , since a_u and a_v is not an antipodal pair on C_{13} , it follows that x and x' are distinguished by a_u or a_v . If both x and x' belong to $P_3[a_u, a_v]$, they are distinguished by a_u or a_v according to Observation 5. The only remaining possibility is that both x and x' belong to $C_{13} - P_3(a_u, a_v)$. Here we distinguish several further possibilities. If x or x' is an internal vertex or an edge from P_1 , then x and x' are distinguished by a_u or a_v . Similarly, if one from the pair x and x' belongs to $P_3[u, a_v]$ and the other to $P_3[v, a_u]$, then again x and x' are distinguished by a_u or a_v . Now, if both x and x' belong to $P_3[u, a_v]$ they are enclosed by a_v and w, so they are distinguished by S according to Observation 5. Finally, if both x and x' belong to $P_3[v, a_u]$, they are enclosed by a_u and w, so Observation 5 again implies x and x' are distinguished by S.

Case 2: Both x and x' belong to P_2 . Notice that if both x and x' belong to $P_2[u, w]$, or they both belong to $P_2[v, w]$, then x and x' are enclosed by a_v and w in the first case and by a_u and w in the second case, either way they are S-enclosed and therefore distinguished by S according to Observation 5. The only remaining possibility is when they belong to different sides of w, say x belongs to $P_2[u, w]$ and x' belongs to $P_2[v, w]$, with both x and x' being distinct from w. In this case if x and x' are not distinguished by w, this implies that $|d(x, u) - d(x', v)| \leq 2$ where without loss of generality we may assume $d(x, u) \leq d(x', v)$, i.e. $0 \leq d(x', v) - d(x, u) = \Delta$ where $\Delta \leq 2$. If the shortest path from x to a_u leads through w, then x and x' are distinguished by a_u according to Observation 5. The similar argument holds when the shortest path from x' to a_v leads through P_1 . Notice that if C_{13} is even, then there is a shortest path which leads both through P_1 . Hence, we have

$$d(x, a_u) = d(x, u) + |P_1| + d(v, a_u) = d(x', v) - \Delta + |P_1| + d(v, a_u) = d(x', a_u) - \Delta + |P_1|.$$

Therefore, x and x' are not distinguished by a_u only when $|P_1| = \Delta$. But in that case we have

$$d(x', a_v) = d(x', v) + |P_1| + d(u, a_v) = d(x, u) + \Delta + |P_1| + d(u, a_v) = d(x, a_v) + 2|P_1| > d(x, a_v),$$

so x and x' are distinguished by $a_v \in S$ and we are finished.

Case 3: x belongs to P_2 and x' belongs to C_{13} . Recall that w is the middle (or "almost" middle) vertex of the path P_2 , where without the loss of generality we may assume $d(w, u) \ge d(w, v)$. Denote by u_i the neighbor of u on the path P_i and by v_i the neighbor of v on the path P_i . Now, when P_2 is of even length, then x and x' are not distinguished by w only if $x \in \{u, v\}$ and $x' \in \{u, v, uu_1, uu_3, vv_1, vv_3\}$. But, since both u and v belong also to C_{13} , this means x belongs also to C_{13} , so this case reduces to Case 1. Assume, therefore, that the length of P_2 is odd. Recall that in this case d(w, u) > d(w, v). Notice that x and x' are not distinguished by w only if $x \in \{u, vv_1, vv_3\}$. Also, notice that there certainly exists a shortest path from x to a_u which leads through P_3 , then C_{13} is even and there is also a shortest path which leads through P_1 , so the claim holds). Finally, notice that $d(x, a_u) > d(v, a_u)$. On the other hand it obviously holds that $d(x', a_u) \le d(v, a_u)$. We conclude that $d(x, a_u) > d(x', a_u)$, which implies that x and x' are distinguished by S.

We have established that any x and x' are distinguished by S, so S is a mixed metric generator. Since |S| = 3, this implies $\operatorname{mdim}(G) \leq 3$.

Theorem 10 If G is an unbalanced Theta graph, then mdim(G) = 3.

Proof. Given the result from Lemma 9, it is sufficient to prove that a set $S \subseteq V(G)$ such that |S| = 2 cannot be a mixed metric generator. We use the same notation as before, i.e. P_1, P_2 and P_3 are the three paths connecting vertices u and v of degree 3 in G denoted so that $|P_1| \leq |P_2| \leq |P_3|$. By C_{ij} we denote the cycle induced by paths P_i and P_j . Let us further denote by s_1 and s_2 the only pair of elements from S.

Assume first that both s_1 and s_2 belong to C_{12} . As the mixed metric dimension of any cycle equals three, it follows that S is not a mixed metric generator in C_{12} . Since C_{12} is an isometric subgraph of G it follows that S cannot be a mixed metric generator in G either. The similar argument holds when both s_1 and s_2 belong to C_{13} . The only remaining possibility is that s_1 belongs to P_2 and s_2 belongs to P_3 . Notice that both s_1 and s_2 in this case must be of degree two, otherwise the pair s_1 and s_2 would belong to C_{12} or C_{13} and the case would reduce to the already proven cases. Let w and z be two neighbors of s_1 . There are only two possibilities with respect to the distances from w and z to s_2 , it is either $d(w, s_2) = d(z, s_2)$ or $d(w, s_2) \neq d(z, s_2)$. If $d(w, s_2) = d(z, s_2)$, then ws_2 and zs_2 are not distinguished by S. On the other hand, if $d(w, s_2) \neq d(z, s_2)$ we may, without the loss of generality, assume that $d(w, s_2) < d(z, s_2)$. But then s_2 and s_2z are not distinguished by S and we are finished.

As for an unbalanced Theta graph G we have $L_1(G) = 0$ and c(G) = 2, we immediately obtain the following result.

Corollary 11 For an unbalanced Theta graph G it holds that $mdim(G) < L_1(G) + 2c(G)$.

We conclude from this result that Conjecture 3 holds for the class of unbalanced Theta graphs with the strict inequality in (1).

5 Concluding remarks

In [12] it was conjectured that $\operatorname{mdim}(G) \leq L_1(G) + 2c(G)$ for all graphs, where c(G) is the cyclomatic number and $L_1(G)$ the number of leaves in a graph G (see Conjecture 3). In this paper we focused our interest on graphs for which the conjecture holds with equality. It was already proven in literature that the equality holds for all trees, even more for all cactus graphs in which every cycle has precisely one vertex of degree ≥ 3 . We wanted to find other graphs for which the equality holds. By Proposition 4 we know that the equality can hold only for graphs with $\kappa(G) = 1$ or $\kappa(G) = 2$. Since all cactus graphs (except a cycle graph) have vertex conectivity equal to 1, this means there were no known graphs with $\kappa(G) = 2$ for which the equality holds. In this paper we found such a family, i.e. balanced Theta graphs, for which the equality also holds. We further proved that for unbalanced Theta graphs Conjecture 3 also holds, but with strict inequality in (1). The natural question that arises is: Are there any other graphs for which the equality holds? Our investigation of the question leads us to the opinion that there are not, i.e. having in mind that trees are a subclass of cactus graphs we state the following formal conjecture.

Conjecture 12 For a graph G it holds that $mdim(G) = L_1(G) + 2c(G)$ if and only if G is a cactus graph in which every cycle has precisely one vertex of degree ≥ 3 or a balanced Theta graph.

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