

Accepted Manuscript

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PII: S0098-1354(17)30318-6
DOI: <http://dx.doi.org/10.1016/j.compchemeng.2017.09.013>
Reference: CACE 5893

To appear in: *Computers and Chemical Engineering*

Received date: 20-6-2017
Revised date: 13-9-2017
Accepted date: 14-9-2017

Please cite this article as: Quirante, Natalia., Grossmann, Ignacio E., & Caballero, José A., Disjunctive Model for the Simultaneous Optimization and Heat Integration with Unclassified Streams and Area Estimation. *Computers and Chemical Engineering* <http://dx.doi.org/10.1016/j.compchemeng.2017.09.013>

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Disjunctive Model for the Simultaneous Optimization and Heat Integration with Unclassified Streams and Area Estimation

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Highlights

- Disjunctive formulation for simultaneous optimization and heat integration.
- It involves unclassified process streams with variable inlet/outlet temperatures.
- Extension of the model to allow area estimation assuming vertical heat transfer.
- Four examples illustrate the numerical performance of the proposed approach.
- The disjunctive formulation has excellent numerical performance.

Abstract

In this paper, we propose a disjunctive formulation for the simultaneous chemical process optimization and heat integration with unclassified process streams –streams that cannot be classified *a priori* as hot or cold streams and whose final classification depend on the process operating conditions–, variable inlet and outlet temperatures, variable flow rates, isothermal process streams, and the possibility of using different utilities.

The paper also presents an extension to allow area estimation assuming vertical heat transfer. The model takes advantage of the disjunctive formulation of the ‘max’ operator to explicitly determine all the ‘kink’ points on the hot and cold balanced composite curves and uses an implicit ordering for determining adjacent points in the balanced composite curves for area estimation.

The numerical performance of the proposed approach is illustrated with four case studies. Results show that the novel disjunctive model of the pinch location method has excellent numerical performance, even in large-scale models.

Keywords: simultaneous optimization, heat integration, variable temperatures, disjunctive model, unclassified streams.

Nomenclature

Index	
i	Hot stream
j	Cold stream
k	Unclassified process stream
m	Non-differentiable ‘kink’ point in the hot and cold composite curve and its end points
p	Process streams that are pinch candidates
s	Process stream
Sets	
$COLD$	Set of all the cold streams j
HOT	Set of all the hot streams i
ISO	Set of all the isothermal streams
M	Set of ‘kink’ points
$MCOLD$	Set of ‘kink’ points corresponding to an inlet or outlet temperature of a cold stream
$MHOT$	Set of ‘kink’ points corresponding to an inlet or outlet temperature of a hot stream
STR	Set of all the process streams s
UNC	Set of unclassified process streams k
Variables	
f^{iso}	Binary variable that takes the value ‘1’ if the isothermal stream s is a hot stream and ‘-1’ if it is a cold stream
F_i	Heat capacity flowrate of the hot stream i
f_j	Heat capacity flowrate of the cold stream j
H_m	Enthalpy value in each one of the points in the set M
m	Mass flowrate of a stream
Q^{iso}	Heat content below a pinch candidate for isothermal streams
Q_C	Heat removed by the cold utility
Q_H	Heat provided by the hot utility
Q_C^p	Cooling utilities required for each pinch candidate p
Q_H^p	Heating utilities required for each pinch candidate p
T^p	Pinch point temperature
$T_{C,k}^{in}$	Disaggregated variable for actual inlet temperature of the cold streams
$T_{C,k}^{out}$	Disaggregated variable for actual outlet temperature of the cold streams
$T_{H,k}^{in}$	Disaggregated variable for actual inlet temperature of the hot streams
$T_{H,k}^{out}$	Disaggregated variable for actual outlet temperature of the hot streams
T_i^{in}	Actual inlet temperature for the hot stream i
T_i^{out}	Actual outlet temperature for the hot stream i

t_j^{in}	Actual inlet temperature for the cold stream j
t_j^{out}	Actual outlet temperature for the cold stream j
t_m	Cold composite curve temperature of the ‘kink’ point m
T_m	Hot composite curve temperature of the ‘kink’ point m
T_s^+	Variable that will take a positive value for hot streams
T_s^-	Variable that will take a positive value for cold streams
$TS_{C,k}^{in}$	Disaggregated variable for shifted inlet temperature of the cold streams
$TS_{C,k}^{out}$	Disaggregated variable for shifted outlet temperature of the cold streams
$TS_{H,k}^{in}$	Disaggregated variable for shifted inlet temperature of the hot streams
$TS_{H,k}^{out}$	Disaggregated variable for shifted outlet temperature of the hot streams
TS_i^{in}	Shifted inlet temperature for the hot stream i
TS_i^{out}	Shifted outlet temperature for the hot stream i
ts_j^{in}	Shifted inlet temperature for the cold stream j
ts_j^{out}	Shifted outlet temperature for the cold stream j
TS_p^{in}	Pinch candidate of all the inlet temperatures of all the streams
wc	Binary variable that takes the value of ‘1’ if the stream k is classified as cold
wh	Binary variable that takes the value of ‘1’ if the stream k is classified as hot
WC	Boolean variable that takes the value of ‘True’ if the stream k is classified as cold
WH	Boolean variable that takes the value of ‘True’ if the stream k is classified as hot
$y_{m,m'}$	Binary variable that takes the value ‘1’ if the unordered enthalpy value that originally was in position m is assigned to position m' in the non-decreasing reordered enthalpies
$Y_{m,m'}$	Boolean variable that takes the value ‘True’ if the unordered enthalpy value that originally was in position m is assigned to position m' in the non-decreasing reordered enthalpies
y_s^{iso}	Binary variable that takes the value of ‘1’ if the isothermal stream is located below the pinch
Y_s^{iso}	Boolean variable that takes the value of ‘True’ if the isothermal stream is located below the pinch
λ	Specific heat associated with the change of phase
ΔT_m^{LM}	Logarithmic mean temperature
ΔT_{min}	Minimum heat recovery approach temperature

1. Introduction

One of the greatest advances in chemical process engineering was the discovery by Hohmann (1971) in his PhD thesis that it is possible to calculate the least amount of hot and cold utilities required for a process without knowing the heat exchanger network. This advance motivated the introduction of the pinch concept (Linnhoff & Flower, 1978a, 1978b; Umeda et al., 1978) and the Pinch Design Method (Linnhoff & Hindmarsh, 1983), for the design of heat exchanger networks

(HEN). Since those seminal works, hundreds of papers have been published related to heat integration.

Significant advances have been developed over the last few decades. Papoulias and Grossmann (1983) presented a mathematical programming model that takes the form of a transshipment problem that allows calculating the minimum utilities and the minimum number of matches (an alternative version that used a transportation model was presented by Cerda et al. (1983)). The first one to use the vertical heat transfer concept that allows estimating the heat transfer area without knowing the explicit design of a heat exchanger network was Jones in 1987 (Jones, 1987). However, the vertical heat transfer area assumption can be problematic if the heat transfer coefficient is significantly different for the various stream matches. A rigorous model for dealing with such a case was presented by Manousiouthakis and Martin (2004). The first automated HEN design, relying on a sequential approach –minimum utilities calculation, followed by a minimum number of heat exchangers and then the detailed network– was developed by Floudas et al. (1986). Later, Ciric and Floudas (1991), Floudas and Ciric (1989, 1990), Yee and Grossmann (1990), and Yuan et al. (1989) proposed different alternatives for the simultaneous design of the HEN, all of them based on mathematical programming approaches. Comprehensive reviews of the advances in HEN in the 20th century can be found in Gundersen and Naess (1988), Jezowski (1994a, 1994b), and Furman and Sahinidis (2002). More recent reviews can be found in Morar and Agachi (2010) and Klemeš and Kravanja (2013).

Pinch analysis has been extended to almost all branches of chemical process engineering, for example, Ahmetović presented a review of the literature for water and energy integration (Ahmetović et al., 2015; Ahmetović & Kravanja, 2013). In El-Halwagi and Manousiouthakis (1989) we can find the extension of the pinch analysis to mass exchange networks and process integration. Tan and Foo (2007) extended the pinch analysis to carbon-constrained energy sector planning. The cogeneration and total site integration can be found in Raissi (1994) and Dhole and Linnhoff (1993). Holiastos and Manousiouthakis (2002) established the theoretical basis of power and heat integration by determining the best integration between heat exchangers and heat pumps and heat engines. They showed that the optimal placement rules (Linnhoff et al., 1982; Townsend

& Linnhoff, 1983a, 1983b) of heat pumps across the pinch and heat engines entirely below or above the pinch can be violated in the optimal design. Wechsung et al. (2011) and Onishi et al. (2014b) extended the concept -they called it Work and Heat Exchanger Networks (WHEN)- and proposed superstructures for generating the optimal configuration.

One of the major limitations of the pinch technology applied to the design of heat exchanger networks is that it had to be used once the chemical process has already been designed and all the flows and temperatures fixed. However, the simultaneous design and optimization of the process and the heat integration strategy can produce larger benefits than a sequential approach (Biegler et al. (1997) presented an illustrative example).

Different alternatives have been proposed to deal with this problem. One interesting approach is the Infinite Dimensional State-space (IDEAS) approach (Drake & Manousiouthakis, 2002; Martin & Manousiouthakis, 2003; Pichardo & Manousiouthakis, 2017). It allows consideration of all process networks employing a set of unit operations. For example, a distillation column (or sequence) can be described by a distribution network, a mass exchange network and a heat exchanger network. It has the advantage that naturally results in linear (and therefore convex) optimization models.

In a mathematical programming-based approach for the design of chemical processes, one straightforward possibility consists of extending the superstructure of the process with that of the heat exchanger network. Nevertheless, the problem rapidly becomes intractable due to the large number of variables (both continuous and integer) and equations. Despite this problem, different researchers have solved relatively complex problems following this approach (de la Cruz et al., 2014; Martelli et al., 2017; Oliva et al., 2011; Onishi et al., 2014a; Vázquez-Ojeda et al., 2013; Yee et al., 1990). To alleviate that problem, an alternative is to consider only the thermal effects (heat integration) without the design of a specific network; in other words, including in the optimization only the utilities and their nature (e.g., low, medium or high pressure steam) but not the investment costs of the heat exchangers network. The underlying idea is that energy costs have much larger impact than investment costs and can have an important effect when optimizing with the rest of the process. However, differences in the investment of two heat exchanger

networks with similar utilities and the same streams involved are not expected to be significant at least when compared with the energy effects.

Under some conditions, it is possible to solve the Pinch Tableau problem at each iteration of the optimization or explicitly include in the model the equations of the transshipment (or extended transshipment) problem (Corbetta et al., 2016). For example, some of the superstructure-based approaches for the design of chemical processes include those equations as a part of the model (Ciric & Floudas, 1991).

Gupta and Manousiouthakis (1993) addressed a similar problem in the context of mass exchanger networks with variables supply and target concentrations. They proposed a MINLP model that can be extended also to heat exchanger networks and proved some interesting results, for example, that if we minimize the utility costs with the inlet and outlet compositions of the rich and lean streams allowed to vary between lower and upper bounds, the optimal solution is always at their lower bound. Using this property, Gupta and Manousiouthakis (1996) later proposed a linear version of the model.

However, this approach relies on the concept of temperature interval. However, if inlet (outlet) temperatures can change, the number of temperature intervals and the streams present in each interval change during the optimization. Mathematically this is equivalent to introducing discontinuities and non-differentiabilities, and consequently, the complete optimization can fail. Note that fixing inlet temperatures to a bound could eventually have a large impact on the rest of the process. In some cases, it would be possible to fix some temperatures a priori, however, for the sake of generality we discuss the general case.

To overcome the problem with temperature intervals, Duran and Grossmann developed the Pinch Location Method (PLM) (Duran & Grossmann, 1986). The idea was to develop a mathematical approach that does not rely on the concept of temperature interval, and as a consequence does not suffer from the drawbacks of previous approaches. The major drawback of the original model presented by Duran and Grossmann (1986) is that in their model involves the ‘max’ operator. They proposed to use a smooth approximation to avoid the discontinuity in the derivative of the

‘max’ operator. However, the smooth approximation is nonconvex and its numerical behavior depends on parameters in the approximation function.

To avoid the non-differentiability introduced in the model of Duran and Grossmann (1986), several approaches have employed binary variables to locate pinch temperatures. In fact, Grossmann et al. (1998) presented a disjunctive formulation that explicitly takes into account the location of a stream –above, across or below– potential pinch candidate. Navarro-Amorós et al. (2013) presented an alternative MI(N)LP model that uses the concept of temperature intervals and the transshipment problem for heat integration with variable temperatures. Quirante et al. (2017) proposed a novel disjunctive model for the simultaneous optimization and heat integration of systems with variable inlet and outlet temperatures, based on the formulation of the pinch location method, modeling the ‘max’ operators by means of a disjunction. Kong et al. (2017) proposed an extension of the Navarro-Amorós et al. (2013) model for the simultaneous chemical process synthesis and heat integration considering also unclassified process streams.

A common situation that appears when the temperatures are not fixed, is that *a priori* it is not possible to decide if a process stream is a hot (it requires cooling) or a cold (it requires heating) stream (Kong et al., 2017). The objective of this paper is to extend the research made in our last work (Quirante et al., 2017) to the case in which there are unclassified process streams. The proposed model has the advantage of reducing the number of equations and binary variables compared to existing alternatives, which allows the reduction of CPU time when solving the problems.

Besides, in a chemical process, usually more than a single hot and/or cold utility are present and it is important to deal with the selection of the best set of utilities among all those available, and some streams undergo phase changes. In this paper, we will show how we can extend the pinch location method to deal with all these cases.

A drawback of the PLM is that we ignore the contribution of the area to the total cost of the heat exchanger network. Even though in most situations this is not a major problem, because as discussed above, the effect of the area can be ignored without affecting much the final solution,

this is not necessarily always the case. We will show that for medium size problems it is possible to simultaneously estimate the area of the HEN and consequently its investment cost.

The rest of the paper is structured as follows. In the two following sections, we present an overview of the pinch location method. Then, we present the disjunctive model for solving problems with unclassified streams and the extension to isothermal process streams and multiple utilities. In section 4, we present how it is possible to include area estimation in the model using the vertical heat transfer problem. In section 5, we present some case studies to illustrate the performance of the proposed approach. Finally, we provide some conclusions obtained from this work.

One final consideration. We present the Generalized Disjunctive Programming (GDP) representation of the model and then the reformulation in terms of binary variables to obtain the final MI(N)LP model. As a general «philosophy» we prefer writing a model as a GDP. Even though the MI(N)LP framework has been successfully used in many different areas, it greatly relies on the expertise of the modeler to generate models that are tractable and effective to solve. GDP models not only considers algebraic expressions but also disjunctions and logic propositions, which allows the modeler to focus on the physical description of the problem rather than on the properties of the model from a mathematical perspective. Exploiting the underlying logic structure of this representation at a higher level of abstraction can help to obtain MINLP models with tighter relaxations and, hence, develop better solution methods. However, we recognize that the selection of the best reformulation to MI(N)LP model is not always straightforward and not all researcher interested in heat integration, are necessarily familiarized with GDP, and algebraic models that automatically reformulate the problem are scarce (i.e., in JUMPS – LogMIP solvers under GAMS (GAMS Development Corporation, 2017)). So we also include, at this moment, the best reformulation as an MI(N)LP. The continuous advances in logic based algorithms and theory behind GDP could eventually find better algorithms or reformulations, but the GDP models will continue to be the same.

2. The Pinch Location Method. Overview

In a system in which all the heat flows are constant, the pinch point is always in the inlet temperature of some of the process streams. Duran and Grossmann (1986) showed that for a fixed minimum approach temperature (ΔT_{min}) between the hot and cold composite curves, if we systematically calculate all the hot and cold utilities for all the pinch candidates (all the inlet temperatures of the process streams), the correct answer corresponds to the candidate with the largest heating and cooling utilities. Mathematically this result can be written as follows:

$$Q_H = \max_{p \in STR} (Q_H^p); \quad Q_C = \max_{p \in STR} (Q_C^p); \quad (1)$$

where STR is a set of all the process streams that are pinch candidates. Q_H, Q_C are the heating and cooling utilities for a given ΔT_{min} and Q_H^p, Q_C^p are the heating and cooling utilities for each one of the pinch candidates p .

In order to take into account that the hot and cold composite curves must be separated at least by the minimum approach temperature, we must work with shifted temperatures.

Defining the following index sets:

$$\begin{aligned} HOT &= [i \mid i \text{ is a hot stream}] \quad HOT \subseteq STR \\ COLD &= [j \mid j \text{ is a cold stream}] \quad COLD \subseteq STR \end{aligned}$$

The shifted temperatures can be defined as follows:

$$\left. \begin{aligned} TS_i^{in} &= T_i^{in} - \frac{\Delta T_{min}}{2} \\ TS_i^{out} &= T_i^{out} - \frac{\Delta T_{min}}{2} \end{aligned} \right\} i \in HOT$$

$$\left. \begin{aligned} ts_j^{in} &= t_j^{in} + \frac{\Delta T_{min}}{2} \\ ts_j^{out} &= t_j^{out} + \frac{\Delta T_{min}}{2} \end{aligned} \right\} j \in COLD \quad (2)$$

where $T_i^{in}, T_i^{out}, t_j^{in}, t_j^{out}$ are the actual inlet and outlet stream process temperatures.

From a total heat balance, we obtain the following equation:

$$Q_C = Q_H + \sum_{i \in Hot} F_i T_i^{in} - T_i^{out} - \sum_{j \in Cold} f_j t_j^{out} - t_j^{in} \quad (3)$$

where F_i is the heat capacity flowrate of the hot stream i and f_j is the heat capacity flowrate of the cold stream j .

Taking into account that the pinch point divides the problem into two heat balanced parts, to calculate the hot utility requirements we need to study only the streams above the pinch and the cold utilities can be calculated from the energy balance presented in Eq.(3), or vice versa, we can calculate the cold utilities from the energy content of the streams below the pinch and the hot utilities from the energy balance.

The problem consists of determining the energy content of the streams above (below) the pinch for each of the pinch candidates. To that end, Duran and Grossmann (1986) showed that it is necessary to explicitly take into account the following three situations: The stream is above the pinch, crosses the pinch or it is below the pinch. For the case in which we study the situations of the streams below the pinch for each pinch candidate, the following equation captures the three situations:

$$Q_C^p = \sum_{j \in COLD} f_j \left[\max \left(0, T^p - ts_j^{out} \right) - \max \left(0, T^p - ts_j^{in} \right) \right] - \sum_{i \in HOT} F_i \left[\max \left(0, T^p - TS_i^{in} \right) - \max \left(0, T^p - TS_i^{out} \right) \right] \quad \forall p \in STR \quad (4)$$

where T^p is the shifted inlet temperature of all the streams.

$$T^p = \begin{cases} T_i^{in} - \frac{\Delta T_{min}}{2} & \text{if } p \text{ is a hot stream } i \\ t_j^{in} + \frac{\Delta T_{min}}{2} & \text{if } p \text{ is a cold stream } j \end{cases} \quad (5)$$

Therefore, the simultaneous optimization and heat integration model can be written as follows:

$$\begin{aligned} & \min f(x) + C_H Q_H + C_C Q_C \\ & s.t. \ h(x) = 0 \\ & \quad g(x) \leq 0 \\ & \quad Q_C \geq \sum_{j \in Cold} f_j \left[\max \left(0, T^p - ts_j^{out} \right) - \max \left(0, T^p - ts_j^{in} \right) \right] - \sum_{i \in Hot} F_i \left[\max \left(0, T^p - TS_i^{in} \right) - \max \left(0, T^p - TS_i^{out} \right) \right] \quad p \in STR \\ & \quad Q_C = Q_H + \sum_{i \in Hot} F_i (T_i^{in} - T_i^{out}) - \sum_{j \in Cold} f_j (t_j^{out} - t_j^{in}) \\ & \quad Q_C, Q_H, F_i, f_j \geq 0 \end{aligned} \quad (6)$$

where $f(x)$ refers to the effects of the rest of the process (everything but heat integration) in the objective function, $h(x)$ is the set of equations defining the process, and $g(x)$ are corresponding constraints of the process.

3. The Pinch Location Method with Unclassified Process Streams

In this section, we present a disjunctive model for the simultaneous optimization and heat integration that also takes into account the possibility of including unclassified process streams. These streams can behave as hot or cold streams depending on the operating conditions of the rest of the process and, therefore, cannot be classified *a priori*. The model is based on the pinch location in which the ‘max’ operators are replaced by disjunctions following the procedure presented by Quirante et al. (2017).

To formally introduce the model let us define the following index sets:

$$\begin{aligned} STR &= [s \mid s \text{ is a process stream}] \\ HOT &= [i \mid i \text{ is a hot stream}] \quad HOT \subseteq STR \\ COLD &= [j \mid j \text{ is a cold stream}] \quad COLD \subseteq STR \\ UNC &= [k \mid k \text{ is an unclassified stream}] \quad UNC \subseteq STR \end{aligned}$$

Note that $HOT \cup COLD \cup UNC = STR$

Classification constraints

Here we follow the approach presented by Kong et al. (2017).

$$\begin{aligned} T_s^{in} - T_s^{out} &= T_s^+ - T_s^- \quad s \in STR \\ T_s^- &= 0 \quad \forall s \in HOT \\ T_s^+ &= 0 \quad \forall s \in COLD \end{aligned} \tag{7}$$

In Eq.(7), we have introduced the variables T_s^+ , T_s^- . The first one (T_s^+) will take a positive value for hot streams, and the second one (T_s^-) for the cold streams. The correct classification of the unclassified streams can be forced by the following disjunction:

$$\left[\begin{array}{c} WH_k \\ T_k^+ \geq 0 \\ T_k^- = 0 \end{array} \right] \vee \left[\begin{array}{c} WC_s \\ T_k^- \geq 0 \\ T_k^+ = 0 \end{array} \right] \quad \forall k \in UNC \tag{8}$$

where WH and WC are Boolean variables that take the value of “True” if the stream ‘ k ’ is classified as hot or cold respectively. This disjunction can be reformulated in terms of binary variables using the hull reformulation (Trespalcios & Grossmann, 2014):

$$\left. \begin{aligned} wh_k + wc_k &= 1 \\ T_k^+ &\leq \overline{T_k^+} \cdot wh_k \\ T_k^- &\leq \overline{T_k^-} \cdot wc_k \end{aligned} \right\} \forall k \in UNC \quad (9)$$

where wh and wc are now binary variables that take the value 1 if the stream is classified as hot or cold respectively, and 0 otherwise.

Definition of shifted temperatures

For the hot and cold streams, shifted temperatures are equivalent to those presented in Eq.(2) (we rewrite them here for the sake of clarity):

$$\left. \begin{aligned} TS_i^{in} &= T_i^{in} - \frac{\Delta T_{min}}{2} \\ TS_i^{out} &= T_i^{out} - \frac{\Delta T_{min}}{2} \end{aligned} \right\} i \in HOT$$

$$\left. \begin{aligned} TS_j^{in} &= t_j^{in} + \frac{\Delta T_{min}}{2} \\ TS_j^{out} &= t_j^{out} + \frac{\Delta T_{min}}{2} \end{aligned} \right\} j \in COLD \quad (10)$$

The correct displacement of the unclassified streams can be forced with the following disjunction:

$$\left[\begin{array}{c} WH_k \\ TS_k^{in} = T_k^{in} - \frac{\Delta T_{min}}{2} \\ TS_k^{out} = T_k^{out} - \frac{\Delta T_{min}}{2} \end{array} \right] \vee \left[\begin{array}{c} WC_k \\ TS_k^{in} = T_k^{in} + \frac{\Delta T_{min}}{2} \\ TS_k^{out} = T_k^{out} + \frac{\Delta T_{min}}{2} \end{array} \right] \quad \forall k \in UNC \quad (11)$$

The previous disjunction can be written in terms of binary variables using the hull reformulation as follows:

$$\left. \begin{aligned}
 & wh_k + wc_k = 1 \\
 & TS_k^{in} = TS_{H,k}^{in} + TS_{C,k}^{in} \quad TS_k^{out} = TS_{H,k}^{out} + TS_{C,k}^{out} \\
 & T_k^{in} = T_{H,k}^{in} + T_{C,k}^{in} \quad T_k^{out} = T_{H,k}^{out} + T_{C,k}^{out} \\
 & TS_{H,k}^{in} = T_{H,k}^{in} - \frac{\Delta T_{min}}{2} wh_k \quad TS_{H,k}^{out} = T_{H,k}^{out} - \frac{\Delta T_{min}}{2} wh_k \\
 & TS_{C,k}^{in} = T_{C,k}^{in} + \frac{\Delta T_{min}}{2} wc_k \quad TS_{C,k}^{out} = T_{C,k}^{out} + \frac{\Delta T_{min}}{2} wc_k \\
 & T_{H,k}^{in} \leq \overline{T_{H,k}^{in}} \cdot wh_k \quad T_{H,k}^{out} \leq \overline{T_{H,k}^{out}} \cdot wh_k \\
 & TS_{H,k}^{in} \leq \overline{TS_{H,k}^{in}} \cdot wh_k \quad TS_{H,k}^{out} \leq \overline{TS_{H,k}^{out}} \cdot wh_k \\
 & T_{C,k}^{in} \leq \overline{T_{C,k}^{in}} \cdot wc_k \quad T_{C,k}^{out} \leq \overline{T_{C,k}^{out}} \cdot wc_k \\
 & TS_{C,k}^{in} \leq \overline{TS_{C,k}^{in}} \cdot wc_k \quad TS_{C,k}^{out} \leq \overline{TS_{C,k}^{out}} \cdot wc_k
 \end{aligned} \right\} \forall k \in UNC \quad (12)$$

The new variables $TS_{H,k}^{in}$, $TS_{C,k}^{in}$, $TS_{H,k}^{out}$, $TS_{C,k}^{out}$, $T_{H,k}^{in}$, $T_{C,k}^{in}$, $T_{H,k}^{out}$, $T_{C,k}^{out}$ in Eq.(12) correspond to the disaggregated variables needed in the hull reformulation.

Pinch Candidates

As previously commented, the pinch candidates are all the inlet temperatures of all the streams.

For clarity in notation, we introduce the variable T^p .

$$T^p = TS_p^{in} \quad p \in STR \quad (13)$$

Minimum utilities.

In order to calculate the utilities, we must introduce the unclassified streams in the Pinch Location Method. To that end, let us rearrange the Eq.(4) as follows:

$$\begin{aligned}
 Q_C^p = & \sum_{j \in COLD} f_j \left[\max \ 0, T^p - ts_j^{out} \right] + \sum_{i \in HOT} F_i \left[\max \ 0, T^p - TS_i^{out} \right] - \\
 & \sum_{j \in COLD} f_j \left[\max \ 0, T^p - ts_j^{in} \right] - \sum_{i \in HOT} F_i \left[\max \ 0, T^p - TS_i^{in} \right] \quad \forall p \in STR \quad (14)
 \end{aligned}$$

In previous equation, the ‘max’ terms related to the output temperatures on the right side of the equation are additive and those related to the input temperatures have a negative sign. The introduction of the unclassified streams is then straightforward.

$$\begin{aligned}
 Q_C^p = & \sum_{j \in COLD} f_j \left[\max \ 0, T^p - ts_j^{out} \right] + \sum_{i \in HOT} F_i \left[\max \ 0, T^p - TS_i^{out} \right] + \sum_{k \in UNC} F_k \left[\max \ 0, T^p - TS_k^{out} \right] - \\
 & \sum_{j \in COLD} f_j \left[\max \ 0, T^p - ts_j^{in} \right] - \sum_{i \in HOT} F_i \left[\max \ 0, T^p - TS_i^{in} \right] - \sum_{k \in UNC} F_k \left[\max \ 0, T^p - TS_k^{in} \right] \quad (15) \\
 & \forall p \in STR
 \end{aligned}$$

In Eq.(15), we introduce the summation over hot, cold and unclassified streams (the complete set of process streams). Therefore, it is not necessary to maintain the differentiation between hot, cold or unclassified streams, and Eq.(15) can be written in the more compact form using a single index for all the process streams.

$$Q_C^p = \sum_{s \in STR} F_s \left[\max \left(0, T^p - TS_s^{out} \right) - \max \left(0, T^p - TS_s^{in} \right) \right] \quad \forall p \in STR \quad (16)$$

The ‘max’ operator has the drawback that it is nondifferentiable and, therefore, cannot be directly included in an optimization model. In the original paper, Duran and Grossmann (1986) try to overcome that problem by using a smooth approximation. The major problem with this approach is that these kind of smooth approximations are nonconvex and they depend on parameters that must be adjusted to accurately approximate the ‘max’ operator, and at the same time avoid numerical conditioning problems (Balakrishna & Biegler, 1992).

In 1998, Grossmann et al. (1998) proposed a disjunctive formulation for calculating the energy content of a stream above (below) the pinch (Q_H^p, Q_C^p) that explicitly takes into account for each pinch candidate the three alternatives: the stream is above the pinch, the stream crosses it, or it is below the pinch. This disjunctive model was reformulated as an MI(N)LP model using a big-M approach. If the heat flows of all the streams are constant –which is a good approximation in most cases– the resulting model is linear and can be easily added to any process model.

Quirante et al. (2017) presented an alternative disjunctive model in which they deal directly with the ‘max’ operator:

$$\begin{aligned} \phi = \max[0, c^T x] \\ \underline{x} \leq x \leq \bar{x} \end{aligned} \Rightarrow \begin{aligned} & \left[\begin{array}{c} Y \\ c^T x \geq 0 \\ \phi = c^T x \\ \underline{x} \leq x \leq \bar{x} \end{array} \right] \bigvee \left[\begin{array}{c} \neg Y \\ c^T x \leq 0 \\ \phi = 0 \\ \underline{x} \leq x \leq \bar{x} \end{array} \right] \\ & Y \in \{True, False\} \end{aligned} \quad (17)$$

Quirante et al. (2017) showed that the hull reformulation of the disjunction of Eq.(17) can be written as follows:

$$\begin{aligned}
\phi &= c^T x + s \\
y\phi^{LO} &\leq \phi \leq y\phi^{UP} \\
(1-y)s^{LO} &\leq s \leq (1-y)s^{UP} \\
s &\geq 0; \quad \phi \geq 0
\end{aligned} \tag{18}$$

They also showed that previous reformulation requires a smaller number of binary variables and equations and has better relaxation gap than the disjunctive model presented by Grossmann et al. (1998). In Appendix A, the interested reader can find a derivation of the previous formulation as well as tight bounds for ϕ and s . In this paper, we have followed this approach.

The complete disjunctive model for the simultaneous optimization and heat integration considering unclassified streams can be written as follows:

$$\begin{aligned}
 & \min f(x) + C_H Q_H + C_C Q_C \\
 & \text{s.t. } \left. \begin{aligned} h(x) &= 0 \\ g(x) &\leq 0 \end{aligned} \right\} \text{Process Constraints} \\
 & \left. \begin{aligned} T_s^{in} - T_s^{out} &= T_s^+ - T_s^- \quad s \in STR \\ T_s^- &= 0 \quad \forall s \in HOT \\ T_s^+ &= 0 \quad \forall s \in COLD \\ \left[\begin{array}{c} WH_k \\ T_k^+ \geq 0 \\ T_k^- = 0 \end{array} \right] \vee \left[\begin{array}{c} WC_k \\ T_k^- \geq 0 \\ T_k^+ = 0 \end{array} \right] & \quad \forall k \in UNC \end{aligned} \right\} \text{Classification Constraints} \\
 & \left. \begin{aligned} TS_i^{in} &= T_i^{in} - \frac{\Delta T_{min}}{2} \\ TS_i^{out} &= T_i^{out} - \frac{\Delta T_{min}}{2} \end{aligned} \right\} i \in HOT \\
 & \left. \begin{aligned} TS_j^{in} &= t_j^{in} + \frac{\Delta T_{min}}{2} \\ TS_j^{out} &= t_j^{out} + \frac{\Delta T_{min}}{2} \end{aligned} \right\} j \in COLD \\
 & \left. \left[\begin{array}{c} WH_k \\ TS_k^{in} = T_k^{in} - \frac{\Delta T_{min}}{2} \\ TS_k^{out} = T_k^{out} - \frac{\Delta T_{min}}{2} \end{array} \right] \vee \left[\begin{array}{c} WC_k \\ TS_k^{in} = T_k^{in} + \frac{\Delta T_{min}}{2} \\ TS_k^{out} = T_k^{out} + \frac{\Delta T_{min}}{2} \end{array} \right] \quad \forall k \in UNC \right\} \text{Shifted Temperatures} \\
 & T^p = TS_p^{in} \quad p \in STR \quad \} \text{Pinch Candidates} \\
 & \left. \begin{aligned} \phi_{s,p}^{out} &= T^p - TS_s^{out} + S_{s,p}^{out} \\ \left[\begin{array}{c} Y_{s,p}^{out} \\ S_{s,p}^{out} = 0 \end{array} \right] \vee \left[\begin{array}{c} \neg Y_{s,p}^{out} \\ \phi_{s,p}^{out} = 0 \end{array} \right] \\ \phi_{s,p}^{in} &= T^p - TS_s^{in} + S_{s,p}^{in} \\ \left[\begin{array}{c} Y_{s,p}^{in} \\ S_{s,p}^{in} = 0 \end{array} \right] \vee \left[\begin{array}{c} \neg Y_{s,p}^{in} \\ \phi_{s,p}^{in} = 0 \end{array} \right] \end{aligned} \right\} p, s \in STR \quad \} \text{max operator disjunctive formulation} \\
 & Q_C \geq \sum_{s \in STR} F_s [\phi_{s,p}^{out} - \phi_{s,p}^{in}] \quad \forall p \in STR \quad \} \text{Streams heat content below the pinch} \\
 & Q_C = Q_H + \sum_{s \in STR} F_s (T_s^+ - T_s^-) \quad \} \text{Overall heat balance} \\
 & Q_C, Q_H, F_s, T_s^+, T_s^-, \phi_{s,p}^{out}, \phi_{s,p}^{in}, S_{s,p}^{out}, S_{s,p}^{in} \geq 0 \\
 & WH_k, WC_k, Y_{s,p}^{out}, Y_{s,p}^{in} \in (True, False) \\
 & \underline{T_s^{in}} \leq T_s^{in} \leq \overline{T_s^{in}}; \underline{T_s^{out}} \leq T_s^{out} \leq \overline{T_s^{out}} \\
 & \underline{\phi_{s,p}^{in}} \leq \phi_{s,p}^{in} \leq \overline{\phi_{s,p}^{in}}; \underline{\phi_{s,p}^{out}} \leq \phi_{s,p}^{out} \leq \overline{\phi_{s,p}^{out}} \\
 & \underline{S_{s,p}^{out}} \leq S_{s,p}^{out} \leq \overline{S_{s,p}^{out}}; \underline{S_{s,p}^{in}} \leq S_{s,p}^{in} \leq \overline{S_{s,p}^{in}}
 \end{aligned} \tag{19}$$

Note that the previous model is linear if the heat flows (F) are constant.

4. Extension to Isothermal Streams and Multiple Utilities

In the case of an isothermal process stream (for example, a pure component that undergoes a phase change at constant pressure), we cannot use Eq.(16) because all terms cancel each other. However, the heat content below a pinch candidate can be easily calculated by the following disjunction:

$$\left[\begin{array}{c} Y_{s,p}^{Iso} \\ T_s^{Iso} \leq T^p \\ Q_{C,s}^{Iso} = m_s \lambda_s \end{array} \right] \vee \left[\begin{array}{c} \neg Y_{s,p}^{Iso} \\ T_s^{Iso} > T^p \\ Q_{C,s}^{Iso} = 0 \end{array} \right] \quad \forall s \in ISO, p \in STR \quad (20)$$

where ISO is an index set that makes reference to the isothermal streams ($ISO \subseteq STR$). λ is the specific heat for the change of phase and m the flowrate. Y_s^{Iso} is a Boolean variable that takes the value of ‘True’ if the isothermal stream is located below the pinch and ‘False’ otherwise.

The hull reformulation of the previous disjunction can be written as follows:

$$\left. \begin{array}{l} Q_{C,s,p}^{Iso} = m_s \lambda_s y_{s,p}^{Iso} \\ T^p - T_s^{Iso} \leq \frac{T^p - T_s^{Iso}}{y_{s,p}^{Iso}} y_{s,p}^{Iso} \end{array} \right\} s \in ISO, p \in STR \quad (21)$$

When there are isothermal streams, the heat content of the streams below the pinch must be modified as follows:

$$Q_C \geq \sum_{s \in STR|_{STR \neq ISO}} F_s [\phi_{s,p}^{out} - \phi_{s,p}^{in}] + \sum_{s \in ISO} Q_{C,s,p}^{ISO} \quad \forall p \in STR \quad (22)$$

Note that the isothermal stream can also be either a hot or a cold stream and we must take this fact into account in the overall heat balance. Using the parameter f^{Iso} that takes value ‘1’ if the isothermal stream s is a hot stream and “-1” if it is a cold stream, the energy balance becomes:

$$Q_C = Q_H + \sum_{s \in STR|_{STR \neq ISO}} F_s T_s^+ - T_s^- + \sum_{s \in ISO} f^{Iso} m_s \lambda_s \quad (23)$$

The handling of multiple utilities is straightforward. In the case of the utilities, we know the inlet and outlet temperatures, but the heat flowrate is unknown. However, from the point of view of modeling, the extra utilities are completely equivalent to process streams, except for the fact that we must include their costs in the objective function.

Note that if for the utilities, the inlet and outlet temperatures are constant, the model continues to be linear.

4.1. Area Estimation

In most of the chemical processes, the energy savings have an important economic (and environmental) impact. While the investment costs can eventually be also important, as a general rule, we would not expect important differences in investment costs between two different heat exchanger network designs for the same process in comparison with the energy impact. As a consequence, the simultaneous optimization of the process and the energy integration with a posteriori design of the heat exchanger network guarantees a good design. However, in some cases (e.g, expensive materials) the estimation of the area (and therefore of the cost) together with the energy savings could be of interest.

The area estimation can be done assuming a vertical heat transfer between the hot and cold balanced composite curves (Jones, 1987) (Smith, 2016). However, if the heat transfer coefficients of different streams are significantly different this assumption could reach to poor results. Manousiouthakis and Martin (2004) proposed an alternative model that explicitly take into account the possibility of non-vertical heat transfer between two streams. In any case, the location of the area intervals, the implicit re-ordering and the calculation of all ‘kink’ points (see next paragraphs) is common in both methods and the approach by Manousiouthakis and Martin (2004) requires more integer variables. As will be commented later, that area calculation is the part in the model with worse numerical behavior, so for simplicity, we present the model with the vertical heat transfer assumption.

To that end, let us define the new index sets:

$$\begin{aligned}
 M &= [\text{m} \mid \text{m is a non-differentiable (kink) point in the hot and cold composite curve} \\
 &\quad \text{and its end points}] \mid K = 2 \mid HOT \mid + \mid COLD \mid \\
 MHOT &= [\text{the ‘kink’ point m corresponds to an inlet or outlet temperature of a hot} \\
 &\quad \text{stream}] \\
 MCOLD &= [\text{the ‘kink’ point m corresponds to an inlet or outlet temperature of a cold} \\
 &\quad \text{stream}]
 \end{aligned}$$

According to Watson and Barton (2016) and Watson et al. (2015), if we denote as H_m the enthalpy value in each one of the points in the set M , we can create a set of triples (H_m, T_m, t_m) ordered by non-decreasing enthalpy values. T_m makes reference to the hot composite curve temperature and t_m to the cold composite curve temperature, both at H_m .

Two adjacent pairs of triples define a zone for the vertical heat transfer area between the hot and cold composite curves:

$$UA_m = \frac{H_{m+1} - H_m}{\Delta T_m^{ML}} \quad m \in M \mid_{m \neq |M|} \quad (24)$$

The difficulty is to calculate all the triples from an arbitrarily ordered set of hot and cold streams in which inlet and outlet temperatures are also unknown. Watson et al. (2015) and Watson and Barton (2016) showed that the enthalpy values for each of the ‘kink’ points can be calculated by the following expressions:

$$\begin{aligned} H_m &= \sum_{i \in HOT} F_i \left[\max \left(0, T^L - T_i^{out} \right) - \max \left(0, T^L - T_i^{in} \right) \right]; \quad T^L \in T_i^{in} \cup T_i^{out} : m \in MHOT \\ H_m &= \sum_{j \in COLD} f_j \left[\max \left(0, t^L - t_j^{in} \right) - \max \left(0, t^L - t_j^{out} \right) \right]; \quad t^L \in t_j^{in} \cup t_j^{out} : m \in MCOLD \end{aligned} \quad (25)$$

With the previous equations, we can calculate all the enthalpy values and the corresponding temperatures of the ‘kink’ points for the hot and cold balanced composite curves. However, we still need to calculate the temperature values of hot streams for the ‘kink’ points of the cold composite curve and the temperatures of cold streams for the ‘kink’ points of the hot composite curve. In other words, there is one unknown temperature in each triple:

$$\begin{aligned} (H_m, T_m, ?) \quad m \in MHOT \\ (H_m, ?, t_m) \quad m \in MCOLD \end{aligned} \quad (26)$$

Watson et al. (2015) and Watson and Barton (2016) showed that if we know the enthalpies, the following expressions allow calculating the unknown temperatures (T_m, t_m):

$$\begin{aligned} H_m &= \sum_{i \in HOT} F_i \left[\max \left(0, T_m - T_i^{out} \right) - \max \left(0, T^L - T_i^{in} \right) \right]; \quad m \in MCOLD \\ H_m &= \sum_{j \in COLD} f_j \left[\max \left(0, t_m - t_j^{in} \right) - \max \left(0, t^L - t_j^{out} \right) \right]; \quad m \in MHOT \end{aligned} \quad (27)$$

At this point, it is worth mentioning that in Eq.(27) the ‘max’ operator can be formulated as a disjunction following the procedure presented by Quirante et al. (2017). Note also that the terms

in Eq.(27) in which T^L (or t^L) correspond to inlet temperatures have already been included in the model because these temperatures are also the pinch candidates (T^p) in Eq.(16).

Unfortunately, the values of enthalpy (H_m), and therefore the temperatures of the hot and cold balanced composite curves, are unordered. To calculate the area, we must know which triple is adjacent to each other. This can be done using the following disjunctive model:

$$\begin{aligned} \bigvee_{m \in M} \left[\begin{array}{l} Y_{m,m'} \\ H_{m'}^{Ord} = H_m \\ T_{m'}^{Ord} = T_m \\ t_{m'}^{Ord} = t_m \end{array} \right] & \forall m' \in M \\ \bigvee_{m' \in M} Y_{m,m'} & \forall m \in M \\ \left. \begin{array}{l} H_m^{Ord} \leq H_{j+1}^{Ord} \\ T_m^{Ord} \leq T_{m+1}^{Ord} \\ t_m^{Ord} \leq t_{m+1}^{Ord} \end{array} \right\} & \forall m \in M \mid m \neq |M| \end{aligned} \quad (28)$$

where the Boolean variable $Y_{m,m'}$ takes the value ‘True’ if the unordered enthalpy value that originally was in position m is assigned to position m' in the non-decreasing reordered enthalpies and ‘False’ otherwise. The subscript ‘ord’ makes reference to the ordered variables.

The disjunctions in Eq.(28) can be reformulated as a linear problem in terms of binary variables using either a big-M or a convex hull reformulation (Trespalcios & Grossmann, 2014). However, in this case, numerical tests have shown that the Big-M have better numerical performance since in the convex hull reformulation a large number of new variables is not compensated by the improved relaxation.

An estimation of the area can be obtained from:

$$UA = \sum_{m \in M \mid m \neq |M|} \frac{H_{m+1}^{ord} - H_m^{ord}}{\Delta T_m^{LM}} \quad (29)$$

where ΔT_m^{LM} is the logarithmic mean temperature in the interval formed by two consecutive triples. To avoid eventual numerical problems when the difference of temperatures is the same at both ends of the interval, we substitute the logarithmic mean temperature by Chen’s approximation (Chen, 1987).

$$\Delta T_m^{LM} \approx \left[\theta_m \theta_{m+1} \frac{\theta_m + \theta_{m+1}}{2} \right]^{\frac{1}{3}} \quad \forall m \in M \mid_{m \neq |M|} \quad (30)$$

where:

$$\theta_m = T_m - t_m \quad \forall m \in M \quad (31)$$

Then the final model is given by all the equations of the pinch location method and Eq.(25) and Eqs.(27)-(31).

The previous model allows the simultaneous optimization and heat integration considering the effect of the investment in the heat exchanger network. Not only the energy, ordering equations and the inherent non-convexities in the model constrain it into small or medium size problems, but the complexity of the problem depends also on the bounds on inlet and outlet temperatures and on the number of ‘real’ alternatives for ordering temperatures and enthalpies.

It is possible to increase the numerical performance by fixing *a priori* some $Y_{m,m'}$ variables. In other words, a point m in the balanced hot/cold composite curve cannot be assigned to any m' position. It is constrained to a subset of m' positions depending on the bounds of its inlet/outlet temperatures and the bounds of the inlet/outlet temperatures of the rest of streams. For example, if all the inlet/outlet temperatures are fixed, all $Y_{m,m'}$ variables can be fixed *a priori*, and if all bounds of the inlet/outlet temperatures are equal, we *a priori* cannot fix any $Y_{m,m'}$.

Navarro-Amorós et al. (2013) and Kong et al. (2017) proposed, in the context of implicitly ordering, a set of values in a mathematical programming model algorithms that allow reducing the ordering alternatives. These algorithms can also be used for this particular problem.

Alternatively, it is also possible to reduce the reordering alternatives by solving a sequence of MILP problems. Note that if the heat flow values of the process streams are constant and the inlet and outlet temperatures of the utilities are fixed, all the reformulations in terms of binary variables of the equations of pinch location method, and the equation for interpolation and reordering in the area estimation, are linear. Therefore, if we search for the highest (lowest) position in which the point m could be reordered in the nondecreasing sequence of enthalpy values, we can fix to ‘0’

those values of the binary $y_{m,m'}$ outside of those limits. This can be achieved by solving, for each point m , the following MILP:

$$\begin{aligned} Y_m^{LO} &= \arg \min (m' Y_{m,m'}) / Y_m^{UP} = \arg \max (m' Y_{m,m'}) \\ \text{s.t. } & \text{Eqs: } 16, 25, 27, 28 \\ & \underline{T_s^{in}} \leq T_s^{in} \leq \overline{T_s^{in}} \\ & \underline{T_s^{out}} \leq T_s^{out} \leq \overline{T_s^{out}} \end{aligned} \quad (32)$$

Then:

$$\begin{aligned} y_{m,m'} &= 0 \quad \forall m' < Y_m^{LO}, \forall m \in M \\ y_{m,m'} &= 0 \quad \forall m' > Y_m^{UP}, \forall m \in M \\ y_{m,m'} &= 1 \quad \forall m' / Y_m^{LO} = Y_m^{UP}, \forall m \in M \end{aligned} \quad (33)$$

where m' makes reference to the position that the point m' occupies in the ordered set M .

If the heat flow values are not constant, then we still can solve the problem of Eq.(32) by using the corresponding upper/lower bounds for the heat flows.

5. Case Studies

In this paper, we present four case studies to illustrate and discuss the performance of the PLM with unclassified streams, multiple utilities isothermal streams and area estimation. As commented above, the area estimation is constrained to medium size problems, therefore, in the first three examples that deals with a large number of process streams, we consider only the heat integration and in the fourth example, we introduce the area (investment) cost estimation.

The first example integrated unclassified multiple utilities and isothermal process streams. In the second example, we introduce a large-scale problem and we show the excellent numerical performance of the proposed approach. To study the performance of the proposed approach without the interference of external factors, these two first examples deal only with heat integration without taking into account the rest of the process, but in the third one, we simultaneously consider the process synthesis and heat integration. Finally, in the last example, we introduce the area estimation and illustrate the effect of the preprocessing in the numerical behavior of the model.

Problem calculations were carried out in GAMS (McCarl et al., 2016), using BARON (Sahinidis, 1996) as a solver. The computations were performed in a computer with a 3.60 GHz Intel® Core™ i7 Processor and 8 GB of RAM under Windows 10.

5.1. Case Study 1

The first example includes seven process streams: two hot streams, two cold streams, and three unclassified streams. All relevant data for this first case study is in Table 1.

Table 1. Data for case study 1.

Stream	Type	Inlet T (°C)	Outlet T (°C)	FCp (MW/°C)
1	Hot	400 – 440	110 – 130	1
2	Hot (isothermal)	340 – 380	340 – 380	100
3	Cold	160 – 180	415 – 425	3 – 4
4	Cold	100 – 120	250 – 260	3 – 4
5	Unclassified	130 – 240	150 – 300	1
6	Unclassified	180 – 430	210 – 300	2
7	Unclassified	30 – 100	40 – 300	1
				Cost (\$/kW year)
Hot Utility 1		500	500	80
Hot Utility 2		380	380	60
Cold Utility		20	30	20

$\Delta T_{\min} = 20\text{ }^{\circ}\text{C}$

We consider that stream 2 is an isothermal stream, while the other streams are non-isothermal.

We assume that a second hot utility is available at 380 °C with a unit cost of \$60/kW·year.

The objective function consists of minimizing the utility costs. The results obtained and some relevant parameters for the case study are presented in Table 2 and Table 3.

Table 2. Stream temperatures, flow rates, and heat loads for the optimal solution of case study 1.

Stream	Type	Inlet T (°C)	Outlet T (°C)	FCp (MW/°C)
1	Hot	440.0	130.0	1
2	Hot	341.0	341.0	100
3	Cold	180.0	415.0	3
4	Cold	120.0	250.0	3
5	Hot	240.0	150.0	1
6	Hot	430.0	210.0	2
7	Cold	30.1	40.1	1
				Q (MW)
Hot Utility 1		500	500	5
Hot Utility 2		380	380	160
Cold Utility		20	30	0

The optimal solution was \$10.0 million/year. Unclassified streams 5 and 6 are classified as hot streams, while unclassified stream 7 is defined as a cold stream. After heat integration, the process requires 165 MW of heating duty, which is satisfied by the hot utility (5 MW) and the intermediate hot utility (160 MW), and no cooling is required. It is worth remarking the model is solved very efficiently in a fraction second of CPU time.

Table 3. Computational statistics and solution of case study 1.

No equations	602.000
No variables	435.000
No binary variables	46.000
CPU time (s) ^a	0.436
Optimal solution (MM\$/y)	10.000

5.2. Case Study 2

In the second example, we apply the methodology to a large-scale problem. This second example includes 17 process streams: six hot streams, seven cold streams, and four unclassified streams. Temperature and flowrate bounds are shown in Table 4. This problem was originally proposed by Kong et al. (2017). We use it as a means to validate the model –as far as we know, the work by Kong et al. (2017) is the only one that deals with unclassified stream– and show the performance of the proposed approach.

Table 4. Stream specifications for case study 2.

Stream	Type	Inlet T (°C)	Outlet T (°C)	FCp (MW/°C)
1	Hot	400 – 440	110 – 130	1
2	Hot (isothermal)	340 – 380	340 – 380	100
3	Cold	160 – 180	415 – 425	3 – 4
4	Cold	100 – 120	250 – 260	3 – 4
5	Unclassified	130 – 240	150 – 300	1
6	Unclassified	180 – 430	210 – 300	2
7	Hot	280	140	1.5 – 2
8	Hot	355	190 – 200	1.1 – 1.3
9	Cold	360 – 410	411	3.3 – 4
10	Cold	230	320	3 – 3.5
11	Cold	390	460	0.9
12	Unclassified	150 – 160	120 – 180	3
13	Hot	220	170 – 180	0.5 – 1
14	Cold	300	400 – 408	1.6
15	Cold	170	440 – 450	3.5
16	Hot	480	440 – 460	1.8
17	Unclassified	170 – 190	180	3.2 – 4
Cost (\$/kW year)				
Hot Utility		500	500	80
Hot Utility		380	380	60
Cold Utility		20	30	20

 $\Delta T_{\min} = 20\text{ }^{\circ}\text{C}$

In this second case, stream 2 is a hot isothermal stream, while the rest of streams are not isothermal. We have also two hot utilities. The heatflow rate of some of the streams is not constant with becomes the problem in nonlinear and non-convex due to the bilinear term that appears in energy balances. Under these conditions, the resulting problem is an MINLP that is solved to global optimality using the deterministic global solver BARON (Sahinidis, 1996).

The objective function consists of minimizing the utility cost. The results obtained and some relevant parameters for the case study are presented in Table 5 and Table 6.

Table 5. Stream temperatures, flow rates, and heat loads for the optimal solution of case study 2.

Stream	Type	Inlet T (°C)	Outlet T (°C)	FCp (MW/°C)
1	Hot	440	130	1
2	Hot	380	380	100
3	Cold	180	415	3
4	Cold	100	250	3
5	Hot	240	190	1
6	Hot	430	210	2
7	Hot	280	140	2
8	Hot	355	190	1.3
9	Cold	410	411	3.3
10	Cold	230	320	3
11	Cold	390	460	0.9
12	Hot	150	140	3
13	Hot	220	170	1
14	Cold	300	400	1.6
15	Cold	170	440	3.5
16	Hot	480	440	1.8
17	Hot	180	180	3.6
				Q (MW)
Hot Utility 1		500	500	190.3
Hot Utility 2		380	380	859.5
Cold Utility		20	30	0.0

Table 6. Computational statistics and solution of case study 2.

	Present work	Kong et al. (2017)
No equations	3901.000	9714.000
No variables	2681.000	5801.000
No binary variables	163.000	2083.000
CPU time (s) ^a	2.95	13275.000
Heating requirements (MW)		
Hot utility	190.300	190.300
Intermediate hot utility	859.500	859.500
Cooling requirements (MW)	0.000	0.000
Optimal solution (MM\$/y)	66.79	66.800

^a Intel Core i7-4790 3.60 GHz, using BARON 14.4.0 for MINLP.

The optimal solution achieved with our model is \$66.79 million/year. After heat integration, the process requires 1049.8 MW of heating duty, which is satisfied by the hot utility (190.3 MW) and the intermediate hot utility (859.5 MW), and no cooling is required.

The results show that the number of continuous and binary variables and the total number of equations is much lower in the proposed model in comparison to the MINLP model developed by Kong et al. (2017).

The model is solved in around three seconds of CPU time. Although the model in this work and the one presented by Kong et al. (2017) have been solved in different computers and therefore we cannot do a direct comparison, the four orders of magnitude reduction in CPU time and the lower number of variables (specially the number of binaries) and constraints show the potential applicability of the new approach.

5.3. Case Study 3

The following case study corresponds to an example of simultaneous process synthesis and heat integration. This case study is adapted from the work by Kong et al. (2017). Unfortunately, in the original paper some data are missing and consequently, both models cannot be compared.

The superstructure for the chemical process is shown in Fig. 1.

< Insert Fig. 1 >

Fig. 1. Superstructure for the chemical process considered in case study 3.

Four components (A, B, C, and D) are taken into account in the process. The raw materials (components A and B) are used to produce the intermediate product C (Eq.(34)). The reaction can be carried out in two alternative isothermal continuous stirred-tank reactors (CSTR1 and CSTR2) that work with different conditions.



The outlet stream from the reactor is sent to a flash unit in order to separate unreacted A and B from intermediate C. Unreacted A and B are separated at the top and recycled, while C is separated at the bottom. Pure component C is sent to another isothermal stirred-tan reactor (CSTR3) to produce final product D. This second reaction Eq.(35) is assumed to be an equilibrium reaction, and the equilibrium constant (K_c) is a function of the reactor temperature.



$$K_c = K_c^o \cdot \exp\left(-\frac{\Delta H^o}{R} \left(\frac{1}{T(K)} - \frac{1}{298}\right)\right) \quad (36)$$

where $K_c^o = 0.4$ is the equilibrium constant at standard state (298 K, 1 bar), $\Delta H^o = 8 \text{ kJ/mol}$ is the heat of reaction at standard state, R is the universal gas constant, and T is the temperature (in Kelvin) of the CSTR3.

For simplicity, we assume ideal behavior:

$$K_c = \frac{[D]}{[C]} \quad (37)$$

where [C] and [D] are the concentration of component C and D in stream 13, respectively.

Reactor CSTR3 requires heating because the reaction is assumed endothermic. Finally, unreacted C is separated from D in one of the alternative separation technologies before recycled back to CSTR3. Table 7 summarizes the unit specifications for the superstructure.

Table 7. Unit specifications for the superstructure of case study 3.

Reactors	RXN	Temperature (°C)	Conversion ^a	Unit cost pre-factor ^b , k (\$/kmol ^{0.6} ·year ^{0.4})
CSTR1	$A + B \rightarrow C$	227	0.9	0.90
CSTR2	$A + B \rightarrow C$	127	0.8	0.85
CSTR3	$C \leftrightarrow D$	57 – 127	variable	1.00
Separators	Top/Bottom			
SEP1	AB/C	157		1.00
SEP2	C/D ^c	107		1.10
SEP3	C/D	67		1.10
SEP4	C/D	87		0.90
SEP5	C/D	77		0.80

^a The conversion is with respect to the limiting component B.

^b Cost pre-factor relates the total molar flow at the inlet to the annualized cost: $k = k_i (F_i^T)^{0.6}$

^c The split fractions in SEP2 are 0.6 and 0 for component C and D, respectively. The remaining separations are assumed sharp.

It is assumed that the feed stream (stream 1) flow rates are 2 kmol/s of A and 1 kmol/s of B, with a raw material cost of \$0.02/kmol A and \$0.01/kmol B, respectively. We are selling the final product D at a price of \$0.17/kmol. The objective is to maximize the profit, which takes into account the revenue, cost of raw materials, unit capital cost, and utility cost.

The case study contains four process streams that require heating or cooling (streams 2, 7, 12, and 13) which are unknown *a priori*, one process stream that requires cooling (stream 15) and two isothermal streams that represent the heat duties of SEP1 and CSTR3.

We assume that a hot utility is available at 500 °C with a unit cost of \$80/kW·year, and the cold utility enters at 20 °C and exists at 30 °C with a cost of \$20/kW·year. All the problems were solved for a minimum heat recovery temperature (ΔT_{min}) of 20 °C.

The resulting model consists of 651 variables (95 binary variables) and 917 equations. It was solved in 390 seconds using BARON as a solver, with an optimal objective function of \$2.829 million/year. CSTR1 is selected for the first reaction, where the reaction takes place at 227 °C with a 0.9 conversion of reactant B. Intermediate C is converted to D in CSTR3 at 115.14 °C. Finally, the product D is sent to SEP5, where is separated at a rate of 0.930 kmol/s. The optimal stream conditions are shown in Table 8 and the optimal solution for streams in the heat integration is shown in Table 9. After heat integration, the process requires 81.006 MW of heating utility and 2.505 MW of cooling water.

Table 8. Optimal solution for streams in the chemical process.

Stream	Component molar flow rates (kmol/s)			
	A	B	C	D
1	2.000	1.000	-	-
2	2.512	1.033	-	-
3	2.512	1.033	-	-
4	-	-	-	-
5	1.582	0.103	0.930	-
6	-	-	-	-
7	1.582	0.103	0.930	-
8	1.582	0.103	-	-
9	1.070	0.070	-	-
10	0.512	0.033	-	-
11	-	-	0.930	-
12	-	-	1.951	-
13	-	-	1.021	0.930
14	-	-	-	-
15	-	-	-	-
16	-	-	-	-
17	-	-	-	-
18	-	-	-	-
19	-	-	-	-
20	-	-	-	-
21	-	-	-	-
22	-	-	1.021	0.930
23	-	-	-	0.930
24	-	-	1.021	-
25	-	-	-	-
26	-	-	1.021	-
27	-	-	-	0.930

Table 9. Stream specifications for the superstructure of case study 3.

Stream	T_{in} (°C)	T_{out} (°C)	FC_p (MW/°C)	Type
2	47.00	227.00	4.050	Cold
7	227.00	157.00	143.835	Hot
12	115.14	127.00	303.707	Cold
13	127.00	77.00	4.910	Hot
15	107.00	67.00	0.000	Hot
SEP1	157.00	157.00	706.309	Cold
CSTR3	127.00	127.00	5356.113	Cold

The optimal superstructure obtained through the simultaneous optimization and heat integration is shown in Fig. 2.

< Insert Fig. 2 >

Fig. 2. Optimal superstructure for the chemical process of case study 3.

5.4. Case Study 4

In this last case study, we introduce the equations for area estimation together with those of the pinch location method. However, as commented in previous sections, the numerical performance of the model is very dependent on the number of process streams and on the bounds of the inlet and outlet temperatures. As a general rule, the model is constrained to medium size problems mainly due to the inefficient behavior of the implicit reordering equations. In any case, it could be useful in models in which the investment is as important as energy savings.

It is worth noting that although for large-scale problems we cannot ensure a globally optimal solution, it is always possible to get a good solution even though we cannot prove it is the best one.

Table 10 shows the data for this problem. Costs of utilities were obtained from Turton et al. (2013). The investment costs were also correlated from shell and tube heat exchangers also from Turton et al. (2013) and updated to 2017 using the Chemical Engineering Plant Cost Index (CEPCI).

Table 10. Data for case study 4.

Stream	Type	Inlet T (°C)	Outlet T (°C)	FCp (MW/°C)
1	Hot	230 – 260	30 – 50	0.15
2	Hot	135 – 155	110 – 125	0.50
3	Hot	80 – 100	20 – 30	0.25
4	Hot	110 – 120	80 – 100	0.30
5	Cold	10 – 40	170 – 190	0.20
6	Cold	90 – 110	180 – 225	0.30
7	Cold	125 – 160	225 – 235	0.15
8	Cold	130 – 150	200 – 240	0.40
				Cost (k\$/MW year)
Hot Utility		250	250	408.96
Cold Utility		10	20	10.19
$\Delta T_{\min} = 10\text{ }^{\circ}\text{C}$				
$U = 0.002\text{ MW/m}^2\text{ }^{\circ}\text{C}$				
$\text{Area Cost (k\$/year)} = 47.65 + 0.7313\text{ Area (m}^2\text{)}$				

The objective in this problem consists of minimizing the Total Annualized Cost (TAC). We will use the following objective:

$$TAC(k\$ / y) = 408.96 Q_{Hot} + 10.19 Q_{Cold} + 0.7313 Area + 47.65 \quad (38)$$

The problem of determining which minimum utility consumption are required can be very efficiently solved by using the PLM. This problem was solved in 0.06 seconds of CPU time. The minimum hot utility consumption was 49.5 MW and the minimum cold utility consumption was 5 MW. In these conditions, it is possible to estimate the area of the heat exchanger network using the vertical heat transfer approach either by solving the MINLP model in which we fix all the temperatures or using the classical approach using a spreadsheet or even manually (Smith, 2016). If we fix all the inlet and outlet temperatures to the values obtained when solved the PLM method, BARON (Sahinidis, 1996) is able of solving this MINLP problem in less than 10 seconds of CPU time. The area estimation yields 2013 m², with a total annualized cost of 21814 k\$/year. This relatively short CPU time shows that the proposed interpolation approach, using the ‘max’ operator for calculating the missing points in each triple is very efficient. Table 11 shows the optimal results with the *a posteriori* area estimation.

Table 11. Solution for case study 4.

Stream	Type	A posteriori area estimation		Simultaneous area estimation	
		Inlet T (°C)	Outlet T (°C)	Inlet T (°C)	Outlet T (°C)
1	Hot	260	50	260	50
2	Hot	155	120.5	155	120.91
3	Hot	80	30	80	30
4	Hot	110	100	118.62	100
5	Cold	10	170	10	170
6	Cold	90	180	100	180
7	Cold	160	225	160	225
8	Cold	150	250	150	250
Hot Utility (MW)		49.5		49.5	
Cold Utility (MW)		5		10.37	
Area (m ²)		1213		1712	
TAC (k\$/year)		21814		21649	

However, if we include the equations of area estimation, without any pretreatment, the solver BARON is not even able of finding a feasible solution in 500 s of CPU time.

If we solve the pretreatment MILPs, then we can significantly reduce the number of alternatives to be considered in the implicit reordering (see Table 12). Even though BARON is not able to guarantee the global optimal solution in 500 s of CPU time, we get a good solution with just a relative gap of 5.3%.

The obtained solution shows just a marginal improvement in TAC (21649 k\$/year) around a 0.8 %, which is in agreement with the assumption that, in general, neglecting the effect of area cost in the preliminary design of a heat exchanger network does not significantly affect the final result. The area is reduced from 2013 to 1712 m², (~ 15 %) but this reduction is only around 1 % of the TAC, and we must take into account also the fact that the cold utility consumption increases from 5 to 10.3 MW.

Table 12. Feasible intervals for each inlet/outlet temperatures after pretreatment.

Stream		Type	Lower Position**	Upper position**
1	Inlet	Hot	12	15
1	Outlet	Hot	1	5
2	Inlet	Hot	10	13
2	Outlet	Hot	8	10
3	Inlet	Hot	4	8
3	Outlet	Hot	1	4
4	Inlet	Hot	8	10
4	Outlet	Hot	4	8
Hot Utility	Inlet	Hot	16	20
Hot Utility	Outlet	Hot	13	16
5	Inlet	Cold	1	7
5	Outlet	Cold	14	18
6	Inlet	Cold	6	11
6	Outlet	Cold	16	19
7	Inlet	Cold	11	15
7	Outlet	Cold	17	19
8	Inlet	Cold	11	15
8	Outlet	Cold	19	20
Cold Utility	Inlet	Cold	1	4
Cold Utility	Outlet	Cold	3	7

**It refers to the lower/upper position that the inlet/outlet enthalpy point of a given stream could be placed when ordered in non-decreasing enthalpy values.

Table 13. Computational statistics and solution of case study 4.

PLM with area estimation	
N° equations	6439
No variables	2917
No binary variables	234(1124) ^(a)
Pre-processing time (s)	23.9
CPU time (s)	500 ^(b)
Best solution found	21648.9

(a) Number of binary variables after the preprocessing and before (into brackets) preprocessing.

(b) Fixed a maximum CPU time in 500 s.

6. Conclusions

We have proposed a disjunctive model for the simultaneous process optimization and heat integration of systems that include variable temperatures, streams that cannot be classified as hot or cold streams *a priori*, and whose classification as hot or cold stream depends on the operating conditions, isothermal streams and multiple utilities. The idea underlying the proposed approach is that the energy-related costs have a much larger impact than investment cost. The model is based on the disjunctive approach of the pinch location method proposed by Quirante et al. (2017) and the treatment of the unclassified streams presented by Kong et al. (2017).

The proposed formulation has proved to be numerically very efficient. The total number of variables and equations is lower than alternative formulations for dealing with the same problem proposed by Navarro-Amorós et al. (2013) for problems without unclassified streams, or the extension proposed by Kong et al. (2017) that also considers unclassified streams, and the CPU time is reduced by 3-4 orders of magnitude.

The model has also been extended to allow estimating the area of the heat exchanger network. Following the assumption of vertical heat transfer, it is possible to obtain an area estimation with a small error –usually, lower than 10 %– (Smith, 2016), showing that the heat transfer coefficients are not significantly different. To that end, it is necessary to calculate for each ‘kink’ point in the hot and cold balanced composite curves (all the inlet and outlet temperatures) the triples (H_m, T_m, t_m) and order those triples by non-decreasing enthalpy values. The first part (calculate the triples) can be efficiently performed using the approach presented by Watson et al. (2015) and Watson and Barton (2016) that relies also on the ‘max’ operator, and therefore can be efficiently reformulated as a disjunction following the procedure presented by Quirante et al. (2017). The advantage of this approach, in particular, the part related to the interpolation, is that for constant heat flow values it preserves the linearity and it has shown to be numerically efficient. The second part –determining adjacent triples– requires an implicit ordering that significantly complicates the model. However, in some situations, it is possible to reduce the combinatorial difficulties related to the ordering by reducing *a priori* the ‘positions’ that a given point can reach when sorted.

The proposed disjunctive model with unclassified and isothermal process streams and multiple utilities has proved to be robust and numerically very efficient in large-scale problems.

Even though the model is NP-Hard, it has been proved to be very effective in large-scale problems (in Quirante et al. (2017) a problem with 20 hot streams and 20 cold streams was solved in less than a second). The inclusion of unclassified process streams, isothermal streams, and multiple utilities does not affect the nice problem performance. For example, case study 2 with 14 process streams, 4 of them unclassified streams and 1 isothermal stream, two hot utilities and 1 cold utility was solved to global optimality in less than three seconds of CPU time.

The performance of the model extended with the area estimation depends on the problem characteristics –how large are the bounds on the inlet and outlet temperatures and the degree of overlap between those bounds–. However, even in the case where we cannot prove global optimality, we are able to obtain good solutions with a relatively small gap for medium size problems.

Acknowledgments

The authors gratefully acknowledge the financial support by the Ministry of Economy, Industry, and Competitiveness of Spain (CTQ2016-77968-C3-02-P, AEI/FEDER, UE), and Call 2013 National Sub-Program for Training, Grants for pre-doctoral contracts for doctoral training (BES-2013-064791).

Appendix A. Hull Reformulation of the $\langle \max[0, \mathbf{c}^T \mathbf{x}] \rangle$ operator

In this Appendix A, we present the disjunctive reformulation of the ‘max’ operator following the approach presented by Quirante et al. (2017). We also show how to obtain good lower and upper bounds for the variables.

Consider the following expression:

$$\begin{aligned} \phi &= \max(0, \mathbf{c}^T \mathbf{x}) \\ \mathbf{x}^{LO} &\leq \mathbf{x} \leq \mathbf{x}^{UP} \end{aligned} \tag{A.1}$$

where c is a vector of known coefficients and x is a vector of variables. An equivalent disjunctive formulation of Eq.(A.1) is as follows:

$$\begin{aligned} &\left(\begin{array}{c} Y \\ \mathbf{c}^T \mathbf{x} \geq 0 \\ \phi = \mathbf{c}^T \mathbf{x} \\ \mathbf{x}^{LO} \leq \mathbf{x} \leq \mathbf{x}^{UP} \end{array} \right) \bigvee \left(\begin{array}{c} \neg Y \\ \mathbf{c}^T \mathbf{x} \leq 0 \\ \phi = 0 \\ \mathbf{x}^{LO} \leq \mathbf{x} \leq \mathbf{x}^{UP} \end{array} \right) \\ &Y \in [True, False] \end{aligned} \tag{A.2}$$

In disjunction from Eq.(A.2), if the Boolean variable takes the value ‘True’ the first term is enforced and ϕ must be positive, otherwise, it is equal to zero. The hull reformulation of Eq.(A.2) is as follows:

$$\begin{aligned} \mathbf{x} &= \mathbf{x}_1 + \mathbf{x}_2 \\ \phi &= \phi_1 + \phi_2 \\ \mathbf{c}^T \mathbf{x}_1 &\geq 0 & \mathbf{c}^T \mathbf{x}_2 &\leq 0 \\ \phi_1 &= \mathbf{c}^T \mathbf{x}_1 & \phi_2 &= 0 \\ y \cdot \mathbf{x}^{LO} &\leq \mathbf{x}_1 \leq y \cdot \mathbf{x}^{UP} & (1-y) \cdot \mathbf{x}^{LO} &\leq \mathbf{x}_2 \leq (1-y) \cdot \mathbf{x}^{UP} \\ y &\in (0, 1) \end{aligned} \tag{A.3}$$

where the superscripts *LO* and *UP* make reference to the lower and upper bounds respectively.

The model in Eq.(A.3) introduces new variables and equations. However, Quirante et al. (2017) showed that this formulation can be simplified taking into account that variable ϕ_2 is fixed to zero and it does not have much sense to add a new variable and then fix it to zero. Therefore, it can be removed.

The particular values of variables x_2 are not relevant to the problem because they are not used in the model. It is possible to lump the term $\mathbf{c}^T \mathbf{x}_2$ in a single variable as follows:

$$\mathbf{x} = \mathbf{x}_1 + \mathbf{x}_2 \rightarrow \mathbf{c}^T \mathbf{x} = \mathbf{c}^T \mathbf{x}_1 + \mathbf{c}^T \mathbf{x}_2 \rightarrow \mathbf{c}^T \mathbf{x} = \mathbf{c}^T \mathbf{x}_1 - s \rightarrow \mathbf{c}^T \mathbf{x}_1 = \mathbf{c}^T \mathbf{x} + s \tag{A.4}$$

Consequently, we can rewrite the Hull reformulation in terms of the original x variables and the new variable s :

$$\begin{aligned} \phi &= \mathbf{c}^T \mathbf{x} + s \\ y\phi^{LO} &\leq \phi \leq y\phi^{UP} \\ (1-y)s^{LO} &\leq s \leq (1-y)s^{UP} \\ s &\geq 0; \phi \geq 0; y \in (0,1) \end{aligned} \tag{A.5}$$

Good lower and upper bounds for ϕ and s variables can be obtained from the bounds of x and c values.

It is worth noting that Eq.(A.5) can be obtained directly from the hull reformulation of the disjunctive reformulation of the ‘max’ operator formulated as an optimization problem with complementarity constraints (Biegler, 2010).

$$\phi = \max(0, c^T x) \Rightarrow \begin{cases} \phi = c^T x + s \\ 0 \leq \phi \perp s \geq 0 \end{cases} \Rightarrow \begin{cases} \phi = c^T x + s \\ \begin{bmatrix} Y \\ s = 0 \end{bmatrix} \vee \begin{bmatrix} -Y \\ \phi = 0 \end{bmatrix} \\ s \geq 0; \quad \phi \geq 0 \end{cases} \quad (\text{A.6})$$

Note that the hull reformulation of the disjunction in Eq.(A.6) is the set of equations shown in Eq.(A.5).

As an example consider one of the terms that appear in the PLM:

$$\begin{aligned} \phi_{j,p} &= \max(0, t_j^{\text{out}} - T^p) \\ \phi_{j,p} &= t_j^{\text{out}} - T^p + s_{j,p}^{\text{out}} \\ y_{j,p}^{\text{out}} \phi_{j,p}^{\text{out}} \text{ }^{LO} &\leq \phi_{j,p}^{\text{out}} \leq y_{j,p}^{\text{out}} \phi_{j,p}^{\text{out}} \text{ }^{UP} \\ 1 - y_{j,p}^{\text{out}} s_{j,p}^{\text{out}} \text{ }^{LO} &\leq s_{j,p}^{\text{out}} \leq 1 - y_{j,p}^{\text{out}} s_{j,p}^{\text{out}} \text{ }^{UP} \\ s_{j,p}^{\text{out}} &\geq 0; \phi_{j,p}^{\text{out}} \geq 0 \end{aligned} \quad (\text{A.7})$$

The upper and lower bounds can be inferred from the bounds on temperatures as follows:

$$\begin{aligned} \phi_{j,p}^{\text{out}} \text{ }^{LO} &= \max(0, t_j^{\text{out}} \text{ }^{LO} - T^p \text{ }^{UP}) \\ \phi_{j,p}^{\text{out}} \text{ }^{UP} &= \max(0, t_j^{\text{out}} \text{ }^{UP} - T^p \text{ }^{LO}) \\ s_{j,p}^{\text{out}} \text{ }^{LO} &= \max(0, T^p \text{ }^{LO} - t_j^{\text{out}} \text{ }^{UP}) \\ s_{j,p}^{\text{out}} \text{ }^{UP} &= \max(0, T^p \text{ }^{UP} - t_j^{\text{out}} \text{ }^{LO}) \end{aligned} \quad (\text{A.8})$$

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Figure captions

Fig. 1. Superstructure for the chemical process considered in case study 3.

Fig. 2. Optimal superstructure for the chemical process of case study 3.



