# Quasiperfect Domination in Trees 

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#### Abstract

A $k$-quasiperfect dominating set $(k \geq 1)$ of a graph $G$ is a vertex subset $S$ such that every vertex not in $S$ is adjacent to at least one and at most $k$ vertices in $S$. The cardinality of a minimum $k$-quasiperfect dominating set of $G$ is denoted by $\gamma_{1 k}(G)$. Those sets were first introduced by Chellali et al. (2013) as a generalization of the perfect domination concept (which coincides with the case $k=1$ ) and allow us to construct a decreasing chain of quasiperfect dominating parameters


$$
\begin{equation*}
\gamma_{11}(G) \geq \gamma_{12}(G) \geq \ldots \geq \gamma_{1, \Delta}(G)=\gamma(G), \tag{1}
\end{equation*}
$$

in order to indicate how far is $G$ from being perfectly dominated. In this work, we study general properties, tight bounds, existence and realization results involving the parameters of the so-called $Q P$-chain (1), for trees.

Keywords: Domination, Perfect domination, Quasiperfect domination, Trees.

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## 1 Introduction

Recall that a tree is a connected acyclic graph. A leaf is a vertex of degree 1 and vertices of degree at least 2 are interior vertices. We denote by $L(T)$ the set of leaves of a tree $T$ and by $\ell(T)$ the number of leaves of $T$. A support vertex is a vertex having at least a leaf in its neighborhood and a strong support vertex is a support vertex adjacent to at least two leaves.

Given a graph $G$, a subset $S$ of its vertices is a dominating set of $G$ if every vertex $v$ not in $S$ is adjacent to at least one vertex in $S$. The domination number $\gamma(G)$ is the minimum cardinality of a dominating set of $G$, and a dominating set of cardinality $\gamma(G)$ is called a $\gamma$-code [9].

An extreme way of domination occurs when every vertex not in $S$ is adjacent to exactly one vertex in $S$. In that case, $S$ is called a perfect dominating set [2] and $\gamma_{11}(G)$, the minimum cardinality of a perfect dominating set of $G$, is the perfect domination number. A dominating set of cardinality $\gamma_{11}(G)$ is called a $\gamma_{11}$-code.

In a perfect dominating set what is gained from the point of view of accuracy is lost in size, comparing it with a dominating set. Between both notions there is a graduation of definitions: $k$-quasiperfect domination. A $k$ quasiperfect dominating set for $k \geq 1\left(\gamma_{1 k}\right.$-set for short) $[7,11]$ is a dominating set $S$ where every vertex not in $S$ is adjacent to at most $k$ vertices of $S$. Again the $k$-quasiperfect domination number $\gamma_{1 k}(G)$ is the minimum cardinality of a $\gamma_{1 k}$-set of $G$ and a $\gamma_{1 k}$-code is a $\gamma_{1 k}$-set of cardinality $\gamma_{1 k}(G)$.

Given a graph $G$ of order $n$ and maximum degree $\Delta, \gamma_{1 \Delta}$-sets are precisely dominating sets. Thus, one can construct the following chain of quasiperfect domination parameters:

$$
\begin{equation*}
n \geq \gamma_{11}(G) \geq \gamma_{12}(G) \geq \ldots \geq \gamma_{1 \Delta}(G)=\gamma(G) \tag{2}
\end{equation*}
$$

known as the quasiperfect chain of $G$, or simply the $Q P$-chain of $G$.

## 2 Known general results

In this section, we review some results founded in the literature about quasiperfect parameters. Table 2 summarizes the values of parameters under consideration for some simple families of graphs.
Theorem 2.1 [7] If $G$ is a graph of order $n$ verifying at least one of the following conditions: (1) $\Delta(G) \geq n-3$; (2) $\Delta(G) \leq 2$; (3) $G$ is a cograph; (4) $G$ is a claw-free graph, then $\gamma_{12}(G)=\gamma(G)$.

|  | paths | cycles | cliques | stars | bicliques | wheels |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $G$ | $P_{n}$ | $C_{n}$ | $K_{n}$ | $K_{1, n-1}$ | $K_{p, n-p}$ | $W_{n}$ |
| $n$ | $n \geq 3$ | $n \geq 4$ | $n \geq 2$ | $n \geq 4$ | $2 \leq p \leq n-p$ | $n \geq 3$ |
| $\Delta(G)$ | 2 | 2 | $n-1$ | $n-1$ | $n-p$ | $n-1$ |
| $\gamma_{11}(G)$ | $\left\lceil\frac{n}{3}\right\rceil$ | $\left\lceil\frac{2 n}{3}\right\rceil-\left\lfloor\frac{n}{3}\right\rfloor$ | 1 | 1 | 2 | 1 |
| $\gamma_{12}(G)$ | $\left\lceil\frac{n}{3}\right\rceil$ | $\left\lceil\frac{n}{3}\right\rceil$ | 1 | 1 | 2 | 1 |
| $\gamma(G)$ | $\left\lceil\frac{n}{3}\right\rceil$ | $\left\lceil\frac{n}{3}\right\rceil$ | 1 | 1 | 2 | 1 |

Proposition 2.2 [3] Let $G=(V, E)$ a graph of order $n$.
(i) If $\gamma(G) \leq \Delta(G)$, then $\gamma_{1 \gamma}(G)=\ldots=\gamma_{1 \Delta}(G)=\gamma(G)$;
(ii) $\gamma_{1 \delta}(G)<n$;
(iii) $\gamma_{11}(G)=1$ if and only if $\Delta(G)=n-1$.
(iv) $\gamma_{11}(G) \leq n-\ell(G)$ where $\ell(G)$ is the number of vertices of degree one.

Theorem 2.3 [3] Let $k, n$ be positive integers such that $n \geq 6$ and $2 \leq k \leq n$. Then, there exists a graph $G$ of order $n$ such that $\Delta(G)=n-2$ and $\gamma_{11}(G)=k$.

Theorem 2.4 [3] Let $(h, k, n)$ be a triple of integers such that $2 \leq h \leq 3$, $2 \leq k \leq n$ and $n \geq 9$. Then, there exists a graph $G$ such that $|V(G)|=n$, $\Delta(G)=n-3, \gamma(G)=h$ and $\gamma_{11}(G)=k$.

Theorem 2.5 [3] Let $G$ be a graph of order $n$ and $\Delta(G)=3$, other than the bull graph. Then, $\gamma_{11}(G) \leq n-3$.

Proposition 2.6 [3] Let $G$ be either a cubic graph other than $K_{4}$, or a tree with order $n \geq 7$ and $\Delta(G)=3$. Then, $\gamma_{11}(G) \leq n-4$.

The join $G=G_{1} \vee G_{2}$ of graphs $G_{1}$ and $G_{2}$ is the graph such that $V(G)=$ $V\left(G_{1}\right) \cup V\left(G_{2}\right)$ and $E(G)=E\left(G_{1}\right) \cup E\left(G_{2}\right) \cup\left\{u v: u \in V\left(G_{1}\right), v \in V\left(G_{2}\right)\right\}$.

Theorem 2.7 [3] Let $G=G_{1} \vee G_{2}$ be a join graph of order $n$. Then,
(i) $\gamma_{11}(G)=1$ if and only if $G_{1}$ or $G_{2}$ have a universal vertex.
(ii) $\gamma_{11}(G)=2$ if and only if both $G_{1}$ and $G_{2}$ have at least an isolated vertex.
(iii) $\gamma_{11}(G)=n$ in other case.

Corollary 2.8 [3] Let $G=G_{1} \vee G_{2}$ be a connected cograph without universal vertices. Then, $\gamma_{11}(G)=2$ if both $G_{1}$ and $G_{2}$ have at least an isolated vertex, and $\gamma_{11}(G)=n$ in any other case.

Theorem 2.9 [3] Let $h, k, n$ be integers such that $4 \leq n, 2 \leq h \leq k \leq n$ satisfying either $h+k \leq n$ or $3 h+k+1 \leq 2 n$. Then, there exists a claw-free graph $G$ of order $n$ such that $\gamma(G)=h$ and $\gamma_{11}(G)=k$.

The corona of a graph $G$, denoted by $\operatorname{cor}(G)$, is the graph obtained by attaching a leaf to each vertex of $G$.
Theorem 2.10 [8,10] For any graph $G$ the domination number satisfies $\gamma(G) \leq$ $n / 2$. And if $G$ is a graph of even order $n$, then $\gamma(G)=n / 2$ if and only if $G$ is the cycle of order 4 or the corona of a connected graph.

Graphs with odd order $n$ and maximum domination number $\gamma(G)=\lfloor n / 2\rfloor$ are also completely characterized in [1], as a list of six graph classes.
Proposition 2.11 [5] Let $T$ be a tree of order $n \geq 3$. Then
(i) Every $\gamma$ - code of $T$ contains all its strong support vertices.
(ii) Every $\gamma_{11}$ - code of $T$ contains all its strong support vertices.
(iii) $\gamma_{11}(T) \leq n / 2$.
(iv) $\gamma_{11}(T)=n / 2$ if and only if $\gamma(T)=n / 2$ if and only if $T=\operatorname{cor}\left(T^{\prime}\right)$ for some tree $T^{\prime}$.
A tree for which removal of all its leaves results in a path is called a caterpillar.

Proposition 2.12 [ 7 ] If $T$ is a caterpillar, then $\gamma(T)=\gamma_{12}(T)$.

## 3 Our results on Trees

Theorem 3.1[4] Let $T$ be a tree. Then, $\gamma_{1 k}(T) \leq \gamma(T)+\left\lceil\frac{\gamma(T)}{k}\right\rceil-1$, for every integer $k \in\{1, \ldots, \Delta(T)\}$.

Corollary 3.2 For every tree $T, \gamma_{11}(T) \leq 2 \gamma(T)-1$.
Remark 3.3 This bound is not true for general graphs and the difference between both parameters can be as large as desired. For example, the graph displayed in Figure 1 satisfies $\gamma(G)=2$ and $\gamma_{11}(G)=|V(G)|>2 \gamma(G)-1$.

Next, we present a realization theorem for the short chain $\gamma \leq \gamma_{11}(T) \leq$ $2 \gamma-1$. Note that, for every caterpillar $T$ of order $n \geq 3$, Proposition 2.12


Fig. 1. The pair of white vertices form a $\gamma$-code.
and Corollary 3.2 just allow two possible situations, namely, either $\gamma(T)=$ $\gamma_{11}(T) \leq n / 2$ or $\gamma(T)<\gamma_{11}(T)<n / 2$. In the following result, we show that both of them are feasible and that parameters $\gamma$ and $\gamma_{11}$ can take every possible value in each case.
Proposition 3.4 [4] Let $a, b, n$ be positive integers.
(i) If $2 \leq 2 a \leq n$, then there exists a caterpillar $T$ of order $n$ such that $\gamma(T)=\gamma_{11}(T)=a$.
(ii) If $2 \leq a<b \leq 2 a-1$ and $n>2 b$, then there exists a caterpillar $T$ of order $n$ such that $\gamma(T)=a$ and $\gamma_{11}(T)=b$.
Proposition 3.5 [4] A caterpillar $T$ satisfies $\gamma_{11}(T)=2 \gamma(T)-1$ if and only if belongs to the family shown in Figure 2.


Fig. 2. Caterpillar with $\gamma_{11}(T)=2 \gamma(T)-1$.
Let $T$ a tree with maximum degree $\Delta \geq 3$. Next theorem shows that for each inequality of the QP-chain, both possibilities, the equality and the strict inequality, are feasible.
Theorem 3.6 [4] There exists a tree with maximum degree $\Delta \geq 3$, satisfying each one of the $2^{\Delta-1}$ possible combinations of the inequalities of the $Q P$ - chain.

Finally, we present the general form of the QP-chain in the case of $k$-ary trees, that has just two different terms.
Proposition 3.7 [4] Let $T=T(k, h)$ the full $k$-ary tree of order $n=\frac{k^{h+1}-1}{k-1}$, where all leaves are at distance $h-1$ from the root, with $k \geq 2, h \geq 3$. Then

$$
n-\ell(T)=\gamma_{11}(T)=\gamma_{12}(T)=\ldots=\gamma_{1, k-1}(T)>\gamma_{1, k}(T)=\gamma_{1, k+1}(T)=\gamma(T)
$$

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## Appendix

## Proof of Theorem 3.1

Remark 1 Let $T$ be a tree and $S$ a dominating set. Then, since $T$ has no cyles, every vertex not in $S$ has at most one neighbor at each connected component of the subgraph $T[S]$.
Remark 2 Let $T$ be a tree and $S$ a dominating set such that the subgraph $T[S]$ has at most $k$ connected components. Then, $S$ is a $\gamma_{1 k}$-set.

Let $S$ be a $\gamma$-code of $T$. If $S$ is also a $\gamma_{1 k}$-set, then the inequality stated in the theorem holds.

Suppose on the contrary that $S$ is not a $\gamma_{1 k}$-set.
We construct a $\gamma_{1 k}$-set $S^{*}$ containing $S$ and satisfying the inequality stated in the theorem. Let $r$ be the number of connected components of the subgraph induced by $S$, denoted by $T[S]$. Then, $\gamma(T) \geq r$ and, by Remark 2, $r>k$.

Consider a vertex $x_{0} \in V(T) \backslash S$ with at least $k+1$ neighbors in $S$ and let $S_{1}=S \cup\left\{x_{0}\right\}$. By Remark 1, all the neighbors of $x_{0}$ in $S$ lie in different connected components of $T[S]$, therefore $S_{1}$ is a dominating set inducing a subgraph $T\left[S_{1}\right]$ with at most $r-k$ connected components. If $S_{1}$ is a $\gamma_{1 k}$-set, let $S^{*}=S_{1}$.

Otherwise, consider a vertex $x_{1} \in V(T) \backslash S_{1}$ having at least $k+1$ neighbors in $S_{1}$ and let $S_{2}=S_{1} \cup\left\{x_{1}\right\}$. By Remark 1, all the neighbors of $x_{1}$ in $S_{1}$ lie in different connected components of $T\left[S_{1}\right]$, therefore $S_{2}$ is a dominating set inducing a subgraph $T\left[S_{2}\right]$ with at most $(r-k)-k=r-2 k$ connected components. If $S_{2}$ is a $\gamma_{1 k}$-set, let $S^{*}=S_{2}$.

Otherwise, we repeat this procedure until we obtain a $\gamma_{1 k}$-set. Observe that this procedure will end since the number of connected components induced by the sets $S_{1}, S_{2}, \ldots$ is strictly decreasing. Moreover, since $T\left[S_{i}\right]$ has at most $r-i k$ connected components, by Remark 2, $S_{i}$ is a $\gamma_{1 k}$-set whenever $r-i k \leq$ $k$. Therefore, the number of steps needed in order to obtain that $S_{i}$ is a $\gamma_{1 k}$-set, is at most $i=\left\lceil\frac{r-k}{k}\right\rceil$.

Let $S^{*}=S_{j}$ be a $\gamma_{1 k}$-set obtained in this way, where $j \leq\left\lceil\frac{r-k}{k}\right\rceil$. Then,

$$
\gamma_{1 k}(T) \leq\left|S^{*}\right|=|S|+j \leq \gamma(T)+\left\lceil\frac{r-k}{k}\right\rceil \leq \gamma(T)+\left\lceil\frac{\gamma(T)-k}{k}\right\rceil=\gamma(T)+\left\lceil\frac{\gamma(T)}{k}\right\rceil-1 .
$$

## Proof of Proposition 3.4

(i) Consider the caterpillar obtained by attaching a leaf to each of the first $a-1$ vertices of a path of order $a$ and $n-2 a+1 \geq 1$ leaves to the last vertex of the path (see Figure 3). Then the vertices of the path is both a $\gamma$-code and a $\gamma_{11}$-code, and $\gamma(T)=\gamma_{11}(T)=a$.


Fig. 3. $T$ has order $n, \gamma(T)=\gamma_{11}(T)=a$.
(ii) Note that $\gamma(T)=1$ implies $\gamma_{11}(T)=1$, so if both parameter do not agree them $\gamma(T) \geq 2$.

Using that $1 \leq b-a \leq a-1$, let $P$ be the path of order $b$ with consecutive vertices labeled with

$$
u_{1}, v_{1}, \ldots, u_{b-a}, v_{b-a}, u_{b-a+1}, u_{b-a+2}, \ldots, u_{a}
$$

and consider the caterpillar obtained by attaching two leaves to each of the vertices $u_{1}, u_{2}, \ldots, u_{b-a}$, one leaf to each of the vertices $u_{b-a+2}, u_{b-a+3}, \ldots, u_{a}$ and $n-2 b+1$ leaves to vertex $u_{b-a+1}$ (see Figure 4). Since $n-2 b+$ $1 \geq 2$ we obtain that $\left\{u_{1}, u_{2}, \ldots, u_{a}\right\}$ is a $\gamma$-code with $a$ vertices and $\left\{u_{1}, u_{2}, \ldots, u_{a}\right\} \cup\left\{v_{1}, \ldots, v_{b-a}\right\}$ is a $\gamma_{11}$-code with $b$ vertices.


Fig. 4. $T$ has order $n>2 b, a=\gamma(T)<\gamma_{11}(T)=b \leq 2 a-1$.

## Proof of Theorem 3.6

Remark 1 If $u$ is a vertex of a graph $G$ with at least $d$ leaves in its neighborhood, then $u$ is in every $\gamma_{1, h}$-code, for any $h \in\{1, \ldots, d-1\}$.
Remark 2 If $G$ is a graph with maximum degree $\Delta$ and $u$ is a vertex with at least $\Delta-1$ leaves in its neighborhood, then $u$ is in every $\gamma_{1, h}$-code, for any $h \in\{1, \ldots, \Delta-2\}$.
Remark 3 Let $T$ be a tree with maximum degree $\Delta$ and $s$ support vertices. Then $\gamma_{1, \Delta}(T)=\gamma(T) \geq s$.

Let $\Delta \geq 3$. For all $i \in\{1, \ldots, \Delta-1\}$, we write $\circledast_{i}$ for the symbol ' $=$ ' or ' $>$ ' in $\gamma_{1, i}(T) \geq \gamma_{1, i+1}(T)$.
(i) Case 1. If $\circledast_{i}$ is ' $=$ ' for all $i \in\{1, \ldots, \Delta-2\}$. We distinguish two subcases.
(a) Case 1.1. If $\circledast_{\Delta-1}$ is ' $=$ '. The complete bipartite graph $T=K_{1, \Delta}$ is a tree with maximum degree $\Delta$ satisfying:

$$
\gamma_{11}(T)=\gamma_{12}(T)=\ldots=\gamma_{1, \Delta-1}(T)=\gamma_{1, \Delta}(T)=\gamma(T)=1
$$

(b) Case 1.2. If $\circledast_{\Delta-1}$ is ' $>$ '. We consider the following tree $T$ with maximum degree $\Delta$ : let $u$ be a vertex of degree $\Delta$ adjacent to vertices $x_{1}, x_{2}, \ldots, x_{\Delta}$, and attach $\Delta-1$ leaves to each $x_{i}, 1 \leq i \leq \Delta$. Then, we easily derive from Remark 2 that $\left\{x_{1}, \ldots, x_{\Delta}\right\}$ is a $\gamma$-code and $\left\{u, x_{1}, \ldots, x_{\Delta}\right\}$ is a $\gamma_{1, i}$-code for any $i$ such that $i<\Delta$. Therefore, $T$ satisfies
$\Delta+1=\gamma_{11}(T)=\gamma_{12}(T)=\ldots=\gamma_{1, \Delta-1}(T)>\gamma_{1, \Delta}(T)=\gamma(T)=\Delta$.


Fig. 5. Trees illustrating Case 1. of Theorem 3.6.
(ii) Case 2. If $\circledast_{i}$ is ' $>$ ' for some $i \in\{1, \ldots, \Delta-2\}$.

If $\Delta=3$, consider the graphs showed in Figure 6. The tree $T$ on the left side satisfies $6=\gamma_{11}(T)>\gamma_{12}(T)=\gamma_{1,3}(T)=\gamma(T)=4$, since support vertices form a $\gamma$-code (and also a $\gamma_{12}$-code and a $\gamma_{13}$-code), and
all vertices but the leaves form a $\gamma_{11}$-code. The tree $T$ on the right side satisfies $\gamma_{11}(T)=18>\gamma_{12}(T)=12>\gamma_{1,3}(T)=\gamma(T)=11$, since support vertices together with vertex $u$ form a $\gamma$-code (and also a $\gamma_{13^{-}}$ code), support vertices together with vertices $u$ and $v$ form a $\gamma_{12}$-code, and all vertices but the leaves form a $\gamma_{11}$-code.

$\gamma_{11}>\gamma_{12}=\gamma_{13}$
$\gamma_{11}>\gamma_{12}>\gamma_{13}$
Fig. 6. Trees illustrating Case 2 of Theorem 3.6 when $\Delta=3$.
Now suppose $\Delta \geq 4$. Let

$$
\left\{i_{1}, i_{2}, \ldots, i_{k}\right\}=\left\{j: \gamma_{1, j}(T)>\gamma_{1, j+1}(T), j \leq \Delta-2\right\}
$$

where $k \geq 1$ by hypotheses, and assume $1 \leq i_{1}<\ldots<i_{k} \leq \Delta-2$. We distinguish two subcases.


Fig. 7. Trees illustrating Case 2.1 (above) and Case 2.2 (bottom).
(a) Case 2.1. If $\circledast \Delta-1$ is ' $=$ '.

Consider a path $P$ of length $k+2$ with consecutive vertices labeled $u_{i_{1}}, \ldots, u_{i_{k}}, v, w$. Attach $i_{j}$ new vertices to $u_{i_{j}}$ and $\Delta-1$ leaves to each one of those new vertices. Attach also $\Delta-2$ leaves to vertex $v$.

For each vertex $x$ of the path $P$, let $N^{\prime}(x)$ be the set of vertices of $N(x)$ not belonging to the path $P$. Let $A=\cup_{j=1}^{k} N^{\prime}\left(u_{i_{j}}\right)$.

It is not hard to verify that $A \cup\{v\}$ is a $\gamma$-code of $T$, and also a $\gamma_{1, \Delta-1}$-code. Moreover, $A \cup\{v\} \cup\left\{u_{i_{j}}: h \leq j \leq k\right\}$ is a $\gamma_{1 i}$-code if $i_{h-1}<i \leq i_{h}$.
(b) Case 2.2. If $\circledast_{\Delta-1}$ is ' $>$ '.

Consider the tree constructed in case 2.1 and attach $\Delta-1$ new vertices to $w$ and $\Delta-1$ leaves to each one of those new vertices.

With the same notations as in Case 2.1, it is easy to verify that $A \cup\{v\} \cup N^{\prime}(w)$ is a $\gamma$-code of $T$ and $A \cup\{v, w\} \cup N^{\prime}(w)$ is a $\gamma_{1, \Delta-1^{-}}$ code. Moreover, $A \cup\{v, w\} \cup N^{\prime}(w) \cup\left\{u_{i_{j}}: h \leq j \leq k\right\}$ is a $\gamma_{1 i}$-code if $i_{h-1}<i \leq i_{h}$.

Lemma 3.8 Let $T$ be a tree of order $n \geq k+1(k \geq 2)$ with all interior vertices of degree at least $k+1$, except at most one vertex of degree $k$, then $\gamma_{1, k-1}(T)=n-\ell(T)$.

Proof. Notice that $V(T) \backslash L(T)$ is a $\gamma_{1, k-1}$-set for all $k \geq 2$. Suppose that $S$ is a $\gamma_{1, k-1}$-code such that $S \neq V(T) \backslash L(T)$. If $V(T) \backslash L(T) \subset S$, then $|S|>|V(T) \backslash L(T)|$ which is a contradiction. Therefore, there exists a vertex $u_{0} \in V(T) \backslash L(T)$ such that $u_{0} \notin S$. Consider the connected component $T_{0}$ of $u_{0}$ in $T \backslash S$. Notice that $T_{0}$ is a tree of order $n_{0} \geq 1$. If $T_{0}$ has only the vertex $u_{0} \notin L(T)$, then $u_{0}$ is adjacent to at least $k$ vertices of $S$, which is a contradiction. If $T_{0}$ has at least two vertices, $T_{0}$ has at least two leaves in $T_{0}$. Observe that a leaf $w$ of $T_{0}$ can not be a leaf of $T$, otherwise the only neighbor of $w$ is not in $S$, contradicting the fact that $S$ is a dominating set. Therefore, $T_{0}$ has a leaf $w_{0}$ that is a vertex of degree al least $k+1$, implying that $\geq k$ neighbors of $w_{0}$ are in $S$, which is again a contradiction.

## Proof of Proposition 3.7

The set of interior vertices of a tree is a $\gamma_{1, i}$-set for any $i \geq 1$. Therefore, by Lemma 3.8, $n-\ell(T)=\gamma_{11}(T)=\gamma_{12}(T)=\ldots=\gamma_{1, k-1}(T)$. On the other hand, for any $h \geq 3$ consider the set $S$ described as follows:

$$
\begin{aligned}
& S=\bigcup_{0 \leq i \leq r-1} L_{2+3 i}, \text { if } h=3 r, r \geq 1 ; \\
& S=\{z\} \cup \bigcup_{1 \leq i \leq r} L_{3 i}, \text { where } z \in L_{2}, \text { if } h=3 r+1, r \geq 1 ; \\
& S=\bigcup_{0 \leq i \leq r} L_{1+3 i}, \text { if } h=3 r+2, r \geq 1 .
\end{aligned}
$$

Notice that $S$ contains exactly the vertices of one of each three consecutive levels, taking into account that $S$ must contain the strong support vertices, i.e., the vertices of level $h-1$, and in the case $h=3 r+1$ we have to add a vertex $z$ of level 2 to dominate the root (see in Figure 8 an illustration of case $k=2$ ).


Fig. 8. If we add new groups of three levels in each case, being black vertices those of the middle level, the set of black vertices is a dominating code of $T(2, h), h \geq 3$.

By construction, it is obvious that $S$ is a $\gamma_{1, k}$-set and a $\gamma_{1, k+1}$-set, since a vertex not in $S$ has at most $k$ neighbors in $S$. We claim that $S$ is a dominating code and consequently a $\gamma_{1, k}$-code and a $\gamma_{1, k+1}$-code. Let $S$ be a dominating code of $T(k, h), k \geq 2, h \geq 3$. We know that $S$ contains all its strong support vertices, $L_{h-1}$, and these vertices dominate vertices of levels $h, h-1$ and $h-2$. So, we may assume that $S$ does not contain any vertex of level $h-2$, otherwise we can change a vertex $x \in S \cap L_{h-2}$ by its neighbor in level $h-3$ obtaining also a dominating code. Therefore, $S$ is obtained by adding a dominating code of the tree $T(k, h-3)$. Reasoning recursively, we deduce that $S$ is a dominating code.


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