# ON MAGIC DISTINCT LABELLINGS OF SIMPLE GRAPHS

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ABSTRACT. A magic labelling of a graph G with magic sum s is a labelling of the edges of G by nonnegative integers such that for each vertex  $v \in V$ , the sum of labels of all edges incident to v is equal to the same number s. Stanley gave remarkable results on magic labellings, but the distinct labelling case is much more complicated. We consider the complete construction of all magic labellings of a given graph G. The idea is illustrated in detail by dealing with three regular graphs. We give combinatorial proofs. The structure result was used to enumerate the corresponding magic distinct labellings.

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#### 1. INTRODUCTION

Throughout this paper, we use standard set notations  $\mathbb{R}, \mathbb{Z}, \mathbb{N}, \mathbb{P}$  for real numbers, integers, nonnegative integers, and positive integers respectively.

Let G = (V, E) be a finite (undirected) graph with vertex set  $V = \{v_1, \ldots, v_m\}$  and edge set  $E = \{a_1, \ldots, a_n\}$ . A magic labelling of G is a labelling of the edges in E by nonnegative integers such that for each vertex  $v \in V$ , the weight wt(v) of v, defined to be the sum of labels of all edges incident to v, is equal to the same number s, called magic sum (also called index by some authors). More precisely, let  $\mu : E \mapsto \mathbb{N}$  be the labelling. Then

(1.1) 
$$wt(v_i) := \sum_{j=1, (v_i, v_j) \in E}^m \mu(v_i, v_j) = s, \qquad i = 1, 2, \dots, m.$$

A magic distinct labelling is a magic labelling whose labels are distinct. It is said to be pure if the labels are 1, 2, ..., n. These concepts were introduced by graph theorists as an analogous of magic squares, which have been objects of study for centuries.

Plenty work has been done for graph labellings. Magic labellings of simple graphs seem first introduced in [9] as vertex magic labellings. Vertex magic total labelling of a simple graph indeed corresponds to our magic labelling of the same graph with a loop attached to each vertex. For related research, see [3,5,6,7,8,11,16]. Note that "magic" may have different meaning in different context.

In the 1970s, Stanley [13] proved some remarkable facts for magic labellings:

**Theorem 1.** Let G be a finite graph and define  $h_G(s)$  to be the number of magic labellings of G of index s. There exist polynomials  $P_G(s)$  and  $Q_G(s)$  such that  $h_G(s) = P_G(s) + (-1)^s Q_G(s)$ . Moreover, if the graph obtained by removing all loops from G is bipartite, then  $Q_G(s) = 0$ , i.e.,  $h_G(s)$  is a polynomial of s.

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In terms of generating functions, the theorem asserts that  $F^G(y) = \sum_{s\geq 0} H_G(s)y^s$  is a rational function with denominator factors  $1 \pm y$ .

Though magic labellings of graphs behave nicely, magic distinct labellings of graphs behave very badly because of the "distinct" condition on the labels. For instance, for the graph  $G_4$  depicted in Figure 4, the generating function for magic labellings is  $\frac{1+y+y^2}{(1-y)^5}$ , but the generating function for magic distinct labellings is

$$\frac{72y^{12}\left(1-y\right)^2 N_4(y)}{\left(1-y^3\right)^2 \left(1-y^4\right) \left(1-y^5\right) \left(1-y^6\right) \left(1-y^7\right) \left(1-y^8\right)},$$

where  $N_4(y)$  is a polynomial of degree 19. See (3.2) and (3.3) for details.

Our starting point is a simple structure result for magic labellings, from which we are able to extract information about magic distinct labellings.

The set of magic  $\mathbb{R}$ -labellings of G is

 $S_{\mathbb{R}}(G) := \{ \boldsymbol{\alpha} = (\alpha_1, \dots, \alpha_n) \in \mathbb{R}^n : (1.1) \text{ holds for } \mu(a_k) = \alpha_k, \ 1 \le k \le n, \ s = wt(v_1) \}.$ 

Clearly,  $S_{\mathbb{R}}(G)$  is a subspace of  $\mathbb{R}^n$  and its basis can be easily computed by linear algebra. Even the structure of  $S_{\mathbb{Z}}(G) = S_{\mathbb{R}}(G) \cap \mathbb{Z}^n$  is easy: it is a finitely generated abelian group, and there are known algorithms for finding the generators. But the set of magic labellings  $S(G) = S_{\mathbb{R}}(G) \cap \mathbb{N}^n$  only forms a (commutative) monoid (semi-group with identity), and it is usually not the case that the monoid is free; that is, there exists  $\boldsymbol{\alpha}_1, \ldots, \boldsymbol{\alpha}_d \in S(G)$ , called generators, such that every  $\boldsymbol{\alpha} \in S(G)$  can be written uniquely as  $\sum_{i=1}^d k_i \boldsymbol{\alpha}_i$  where  $k_i \in \mathbb{N}$ .

We will decompose S(G) into some shifted free monoids, whose elements can be uniquely written as  $\gamma + \sum_{i=1}^{d} k_i \boldsymbol{\alpha}_i$ , where  $\gamma$  is fixed and usually in S(G), and  $\boldsymbol{\alpha}_1, \ldots, \boldsymbol{\alpha}_d$ are still called the generators. We illustrate the idea by three examples. We give combinatorial proofs.

In terms of generating functions, we define

$$F^G(\mathbf{x}) = \sum_{\boldsymbol{\alpha} \in S_{\mathbb{N}}(G)} \mathbf{x}^{\boldsymbol{\alpha}},$$

where  $\mathbf{x}^{\boldsymbol{\alpha}}$  is short for  $x_1^{\alpha_1} \cdots x_n^{\alpha_n}$ . It is known that  $F^G(\mathbf{x})$  is a rational function with denominator  $\prod_{\beta} (1-x^{\beta})$  where  $\beta$  ranges over all extreme rays of S(G). See [12, Theorem 4.6.11]. There are existing algorithms for computing  $F^G(\mathbf{x})$ , such as the Mathematica package **Omega** in [2], the Maple packages **Ell** in [17] and **CTEuclid** in [15]. But the representation of  $F^G(\mathbf{x})$  by computer is usually not ideal.

Our decomposition give rise a rational function decomposition:  $F^G(\mathbf{x}) = F_1(\mathbf{x}) + F_2(x) + \cdots$ , where each  $F_j(\mathbf{x})$  (corresponding to a shifted free monoid) is of the form

$$\frac{\mathbf{x}^{\gamma}}{(1-\mathbf{x}^{\boldsymbol{\alpha}_1})(1-\mathbf{x}^{\boldsymbol{\alpha}_2})\cdots(1-\mathbf{x}^{\boldsymbol{\alpha}_d})}$$

The paper is organized as follows. Section 1 is this introduction. In Section 2 we deal with three graphs  $G_1$ ,  $G_2$ , and  $G_3$ . We give detailed construction of their magic labellings, and combinatorial proofs. In Section 3 we introduce basic idea of MacMahon's partition analysis, outline the result for  $G_4$ , and setup basic tools for attacking magic distinct labellings of graphs. In particular, we compute the generating functions for several graphs.

### 2. Three examples

In this section, we illustrate our decomposition by considering the magic labellings of three graphs depicted in Figures 1,2 and 3. These graphs are all regular, i.e., the degree for each vertex is the same.

In what follows we will often use  $e_i$  to denote the *i*-th unit vector in  $\mathbb{R}^n$ . Then  $\boldsymbol{\alpha} = (\alpha_1, \ldots, \alpha_n) \in \mathbb{R}^n$  will be written as  $\boldsymbol{\alpha} = \sum_{i=1}^n \alpha_i e_i$ . The dimension *n* will be clear from the context.

For  $\boldsymbol{\alpha} \in S(G)$ , the magic sum  $s = s(\boldsymbol{\alpha})$  is determined by  $\boldsymbol{\alpha}$ , so it can be treated as a redundant variable. It is convenient for us to use the generating function

$$F^G(\mathbf{x}, y) = \sum_{\boldsymbol{\alpha} \in S_{\mathbb{N}}(G)} \mathbf{x}^{\boldsymbol{\alpha}} y^{s(\boldsymbol{\alpha})}.$$

Then  $F^G(\mathbf{x}) = F^G(\mathbf{x}, 1)$  and setting  $x_i = 1$  for all *i* gives the enumerating generating function

$$F^{G}(y) = F^{G}(\mathbf{1}, y) = \sum_{s \ge 0} h_{G}(s) y^{s},$$

where  $h_G(s)$  counts the number of magic labellings of G with magic sum s.

2.1. **Example 1.** Let  $G_1 = K_4$  be the complete graph with  $V = \{v_i, i = 1, 2, \dots, 4\}$ and  $E = \{a_i, i = 1, 2, \dots, 6\}$  as shown in Figure 1.

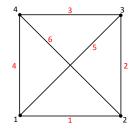


FIGURE 1. The complete graph  $G_1 = K_4$ 

Suppose the magic labelling is given by  $\mu(a_i) = x_i$ , i = 1, 2, ..., 6. Then the fulfilled equations (1.1) becomes

(2.1) 
$$\begin{cases} x_1 + x_4 + x_5 = s, \\ x_1 + x_2 + x_6 = s, \\ x_2 + x_3 + x_5 = s, \\ x_3 + x_4 + x_6 = s. \end{cases}$$

Then

$$\boldsymbol{\alpha} \in S_{\mathbb{R}}(G_1) \Leftrightarrow A(\boldsymbol{\alpha}, s)^T = 0, \ A = \begin{pmatrix} 1 & 0 & 0 & 1 & 1 & 0 & -1 \\ 1 & 1 & 0 & 0 & 0 & 1 & -1 \\ 0 & 1 & 1 & 0 & 1 & 0 & -1 \\ 0 & 0 & 1 & 1 & 0 & 1 & -1 \end{pmatrix}$$

Thus  $S_{\mathbb{R}}(G_1)$  can be identified with the null space of A, which is a subspace in  $\mathbb{R}^7$ . Since  $\operatorname{rank}(A) = 4$ ,  $\dim S_{\mathbb{R}}(G_1) = 7 - 4 = 3$ .

Indeed, we have the following result.

**Proposition 2.** For  $G_1$  as above, the magic labelling of  $G_1$  forms a free monoid with generators  $\boldsymbol{\alpha}_1 = e_2 + e_4, \boldsymbol{\alpha}_2 = e_1 + e_3, \boldsymbol{\alpha}_3 = e_5 + e_6$ .

*Proof.* It is straightforward to check that they are linearly independent and hence form a basis of  $S_{\mathbb{R}}(G_1)$ .

Let  $\boldsymbol{\alpha} = k_1 \boldsymbol{\alpha}_1 + k_2 \boldsymbol{\alpha}_2 + k_3 \boldsymbol{\alpha}_3 = k_2 e_1 + k_1 e_2 + k_2 e_3 + k_1 e_4 + k_3 e_5 + k_3 e_6 \in S_{\mathbb{R}}(G_1)$ . Then  $\boldsymbol{\alpha} \in S(G_1)$  if and only if  $k_1, k_2, k_3 \in \mathbb{N}$ . Therefore  $\boldsymbol{\alpha}_1, \boldsymbol{\alpha}_2$  and  $\boldsymbol{\alpha}_3$  freely generate  $S(G_1)$ .

Note that  $\boldsymbol{\alpha}_1, \boldsymbol{\alpha}_2, \boldsymbol{\alpha}_3$  correspond to perfect matchings of  $G_1$ .

**Corollary 3.** Let  $G_1$  be as above. Then

$$F^{G_1}(\mathbf{x}, y) = \frac{1}{(1 - yx_2x_4)(1 - yx_1x_3)(1 - yx_5x_6)}$$

Consequently,

$$F^{G_1}(y) = \frac{1}{(1-y)^3} = 1 + 3y + 6y^2 + 10y^3 + 15y^4 + 21y^5 + 28y^6 + 36y^7 + \cdots$$

2.2. **Example 2.** Let  $G_2$  be the graph with  $V = \{v_i, i = 1, 2, \dots, 6\}$  and  $E = \{a_i, i = 1, 2, \dots, 9\}$  as shown in Figure 2. Suppose the magic labelling is given by  $\mu(a_i) = x_i$ ,  $i = 1, 2, \dots, 9\}$ 

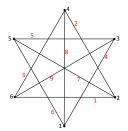


FIGURE 2. The Graph  $G_2$  with 6 vertices and 9 edges.

 $1, 2, \ldots, 9$ . Then the fulfilled equations (1.1) becomes

(2.2) 
$$\begin{cases} x_4 + x_6 + x_8 = s, \\ x_1 + x_2 + x_7 = s, \\ x_4 + x_5 + x_9 = s, \\ x_2 + x_3 + x_8 = s, \\ x_5 + x_6 + x_7 = s, \\ x_1 + x_3 + x_9 = s. \end{cases}$$

It is easy to see the following properties hold:

i) dim  $S_{\mathbb{R}}(G_2) = 4;$ 

ii)  $\beta_1 = e_1 + e_5 + e_8$ ,  $\beta_2 = e_2 + e_6 + e_9$ ,  $\beta_3 = e_3 + e_4 + e_7$ ,  $\beta_4 = e_7 + e_8 + e_9$  are linearly independent and hence form a basis of  $S_{\mathbb{R}}(G_2)$ , and they correspond to the 4 perfect matchings of  $G_2$ ;

iii)  $\beta_5 = e_1 + e_2 + e_3 + e_4 + e_5 + e_6 \in S(G_2)$  is not a nonnegative linear combination of  $\beta_1, \ldots, \beta_4$ . Indeed we have  $\beta_5 = \beta_1 + \beta_2 + \beta_3 - \beta_4$ .

**Proposition 4.** Let  $G_2$  be as in Figure 2. Then every  $\beta$  in  $S(G_2)$  can be uniquely written in one of the following two types.

$$B_{1}: l_{1}\boldsymbol{\beta}_{1} + l_{2}\boldsymbol{\beta}_{2} + l_{3}\boldsymbol{\beta}_{3} + l_{4}\boldsymbol{\beta}_{4}, \ l_{i} \in \mathbb{N} \ (1 \leq i \leq 4); B_{2}: \boldsymbol{\beta}_{5} + l_{1}\boldsymbol{\beta}_{1} + l_{2}\boldsymbol{\beta}_{2} + l_{3}\boldsymbol{\beta}_{3} + l_{4}\boldsymbol{\beta}_{5}, \ l_{i} \in \mathbb{N} \ (1 \leq i \leq 4).$$

*Proof.* By property ii), any element  $\boldsymbol{\beta} \in S_{\mathbb{R}}(G_2)$  can be written as  $\boldsymbol{\beta} = k_1 \boldsymbol{\beta}_1 + k_2 \boldsymbol{\beta}_2 + k_3 \boldsymbol{\beta}_3 + k_4 \boldsymbol{\beta}_4$  where  $k_i \in \mathbb{R}$ ,  $1 \leq i \leq 4$ .

Then  $\boldsymbol{\beta} = k_1 e_1 + k_2 e_2 + k_3 e_3 + k_3 e_4 + k_1 e_5 + k_2 e_6 + (k_3 + k_4) e_7 + (k_1 + k_4) e_8 + (k_2 + k_4) e_9$ . It belongs to  $S(G_2)$  if and only if  $k_1, k_2, k_3, k_1 + k_4, k_2 + k_4, k_3 + k_4 \in \mathbb{N}$ . Note that we can only deduce that  $k_4 \in \mathbb{Z}$ .

When  $k_4 \in \mathbb{N}$ , such  $\boldsymbol{\beta}$  naturally corresponds to the type  $B_1$  case (by setting  $l_i = k_i \in \mathbb{N} \ (1 \leq i \leq 4)$ ).

When  $k_4 < 0$ , i.e.,  $-k_4 \in \mathbb{P}$  we need to rewrite

$$\boldsymbol{\beta} = k_1 \boldsymbol{\beta}_1 + k_2 \boldsymbol{\beta}_2 + k_3 \boldsymbol{\beta}_3 + k_4 (\boldsymbol{\beta}_1 + \boldsymbol{\beta}_2 + \boldsymbol{\beta}_3 - \boldsymbol{\beta}_5) = \boldsymbol{\beta}_5 + (k_1 + k_4) \boldsymbol{\beta}_1 + (k_2 + k_4) \boldsymbol{\beta}_2 + (k_3 + k_4) \boldsymbol{\beta}_3 + (-k_4 - 1) \boldsymbol{\beta}_5.$$

By comparing with the type  $B_2$  case, we shall have  $l_i = k_i + k_4$  for i = 1, 2, 3 and  $l_4 = -k_4 - 1$ . The conditions on the  $k_i$ 's transform exactly to  $l_i \in \mathbb{N}$ ,  $1 \le i \le 4$ . Finally, the uniqueness follows by the linear independency of  $\beta_1, \beta_2, \beta_3, \beta_5$ .

**Corollary 5.** Let  $G_2$  be as above. Then

$$F^{G_2}(\mathbf{x}, y) = \frac{1}{(1 - yx_1x_5x_8)(1 - yx_2x_6x_9)(1 - yx_3x_4x_7)(1 - yx_7x_8x_9)} + \frac{y^2x_1x_2x_3x_4x_5x_6}{(1 - yx_1x_5x_8)(1 - yx_2x_6x_9)(1 - yx_3x_4x_7)(1 - y^2x_1x_2x_3x_4x_5x_6)}$$

Consequently,

$$F^{G_2}(y) = \frac{1+y+y^2}{(1-y)^3 (1-y^2)} = 1 + 4y + 11y^2 + 23y^3 + 42y^4 + 69y^5 + 106y^6 + \cdots$$

2.3. Example 3. Let  $G_3$  be the graph with  $V = \{v_i, i = 1, 2, \dots, 8\}$  and  $E = \{a_i, i = 1, 2, \dots, 12\}$  as shown in Figure 3. Suppose the magic labelling is given by  $\mu(a_i) =$ 

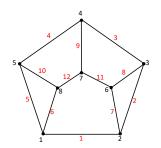


FIGURE 3. The Graph  $G_3$  with 8 vertices and 12 edges.

 $x_i, i = 1, 2, \dots, 12$ . Then the fulfilled equations (1.1) become

$$(2.3) \qquad \begin{cases} x_4 + x_5 + x_6 = s, \\ x_1 + x_2 + x_7 = s, \\ x_2 + x_3 + x_8 = s, \\ x_3 + x_4 + x_9 = s, \\ x_4 + x_5 + x_{10} = s, \\ x_7 + x_8 + x_{11} = s, \\ x_9 + x_{11} + x_{12} = s, \\ x_6 + x_{10} + x_{12} = s. \end{cases}$$

It is easy to see that dim  $S_{\mathbb{R}}(G_3) = 5$ . The structure of  $S(G_3)$  is indeed pretty complicated. Our combinatorial proof is guided by but independent of the algebraic decomposition described in Section 3.

We need the following vectors

(2.4) 
$$\begin{aligned} \gamma_1 &= e_1 + e_3 + e_{10} + e_{11}, & \gamma_2 &= e_2 + e_4 + e_6 + e_{11}, \\ \gamma_3 &= e_3 + e_5 + e_7 + e_{12}, & \gamma_4 &= e_2 + e_5 + e_6 + e_7 + e_8 + 2e_9 + e_{10}, \\ \gamma_5 &= e_2 + e_3 + e_5 + e_6 + e_7 + e_9 + e_{10} + e_{11}, & \gamma_6 &= e_1 + e_4 + e_8 + e_{12}, \\ \gamma_7 &= e_2 + e_4 + e_5 + e_6 + e_7 + e_8 + e_9 + e_{12}, & \gamma_8 &= e_1 + e_8 + e_9 + e_{10}. \end{aligned}$$

The  $\gamma_i$ 's are the extreme rays which are computed by other methods. The following relations can be easily checked.

**Proposition 6.** The  $\gamma_i(1 \leq i \leq 8)$  have the following relationship.

(2.5)  

$$\begin{array}{c|c}
(b) & \gamma_5 = \gamma_1 + \gamma_4 - \gamma_8 \\
\hline
(c_1) & \gamma_4 = -\gamma_1 + \gamma_2 + \gamma_3 - \gamma_6 + 2\gamma_8 \\
& \gamma_5 = \gamma_2 + \gamma_3 - \gamma_6 + \gamma_8 \\
\hline
(c_2) & \gamma_1 = \gamma_2 + \gamma_3 - \gamma_6 + 2\gamma_8 \\
& \gamma_5 = \gamma_2 + \gamma_3 - \gamma_6 + \gamma_8 \\
\hline
(c_3) & \gamma_1 = \gamma_2 + \gamma_3 + \gamma_4 + \gamma_6 - 2\gamma_7 \\
& \gamma_5 = \gamma_2 + \gamma_3 + \gamma_4 - \gamma_7
\end{array}$$

In order to the Theorem 9, we give the following Lemma.

**Lemma 7.** Any  $\gamma \in S_{\mathbb{R}}(G_3)$  can be uniquely written as  $\gamma = k_1\gamma_1 + k_2\gamma_2 + k_3\gamma_3 + k_4\gamma_4 + k_5\gamma_5$ for some  $k_i \in \mathbb{R}$   $(1 \leq i \leq 5)$ . It belongs to  $S(G_3)$  if and only if the following properties hold true.

(2.6) 
$$\begin{vmatrix} k_1, & k_2, & k_3, & k_4, \\ k_1 + k_2 + k_5, & k_1 + k_3 + k_5, & k_1 + k_4 + k_5, & k_2 + k_4 + k_5, \\ k_3 + k_4 + k_5, & 2k_4 + k_5. \end{vmatrix} \in \mathbb{N}$$

Consequently,  $k_5 \in \mathbb{Z}$ .

*Proof.* The first part follows by the linear independency of  $\gamma_1, \ldots, \gamma_5$ . By writing in the  $e_i$  basis, we have

(2.7) 
$$\gamma = k_1 e_1 + (k_2 + k_4 + k_5)e_2 + (k_1 + k_3 + k_5)e_3 + k_2 e_4 + (k_3 + k_4 + k_5)e_5 + (k_2 + k_4 + k_5)e_6 + (k_3 + k_4 + k_5)e_7 + k_4 e_8 + (2k_4 + k_5)e_9 + (k_1 + k_4 + k_5)e_{10} + (k_1 + k_2 + k_5)e_{11} + k_3 e_{12}.$$

Now  $\gamma \in S(G_3)$  if and only if each coordinate belongs to N. This is the second part. The consequence  $k_5 \in \mathbb{Z}$  is obvious.

From now on, we will identify  $\boldsymbol{\gamma} \in S_{\mathbb{R}}(G_3)$  with  $(k_1, \ldots, k_5) \in \mathbb{R}^5$ . For a set S of  $(k_1, \ldots, k_5) \in \mathbb{R}^5$  and a property P, we will use P(S) to denote the subset of S satisfying P. We will use the following properties.

$$\begin{array}{ll} P: & k_i \in \mathbb{N} (1 \leqslant i \leqslant 4), k_5 \in \mathbb{Z}; \\ P_a: & k_5 \in \mathbb{N}, k_i \in \mathbb{N} (1 \leqslant i \leqslant 4); \\ P_b: & -k_5 \in \mathbb{P}, \ \min\{k_1, k_4\} \geqslant -k_5, k_i \in \mathbb{N} (1 \leqslant i \leqslant 4); \\ P_{c_1}: & -k_5 \in \mathbb{P}, \ \min\{k_1, k_4\} < -k_5, \ k_1 \geqslant k_4, k_i \in \mathbb{N} (1 \leqslant i \leqslant 4); \\ P_{c_2}: & -k_5 \in \mathbb{P}, \ \min\{k_1, k_4\} < -k_5, \ k_1 < k_4, \ 2k_1 + k_5 \geqslant 0, k_i \in \mathbb{N} (1 \leqslant i \leqslant 4); \\ P_{c_3}: & -k_5 \in \mathbb{P}, \ \min\{k_1, k_4\} < -k_5, \ k_1 < k_4, \ 2k_1 + k_5 < 0, k_i \in \mathbb{N} (1 \leqslant i \leqslant 4). \end{array}$$

It is easy to see that  $P(\mathbb{R}^5) = P_a(\mathbb{R}^5) \uplus P_b(\mathbb{R}^5) \uplus P_{c_1}(\mathbb{R}^5) \uplus P_{c_2}(\mathbb{R}^5) \uplus P_{c_3}(\mathbb{R}^5)$ , where  $\uplus$  denotes disjoint union. This implies that

$$P(S(G_3)) = P_a(S(G_3)) \uplus P_b(S(G_3)) \uplus P_{c_1}(S(G_3)) \uplus P_{c_2}(S(G_3)) \uplus P_{c_3}(S(G_3))$$

**Lemma 8.** Let *h* be any of  $a, b, c_1, c_2$  and  $c_3$ . Then  $P_h(S(G_3)) = \{(l_1, \ldots, l_5) \in \mathbb{N}^5 : l_i \text{ are given in Table 2.8}\}.$ (2.8)

Types	$ l_1 $	$l_2$	$l_3$	$ l_4 $	$l_5$
$P_a(S(G_3))$	$k_1$	$k_2$	$k_3$	$k_4$	$k_5$
$P_b(S(G_3))$	$k_1 + k_5$	$k_2$	$k_3$	$k_4 + k_5$	$-k_5 - 1$
$P_{c_1}(S(G_3))$	$k_1 - k_4$	$k_2 + k_4 + k_5$	$k_3 + k_4 + k_5$	$-k_5 - k_4 - 1$	$2k_4 + k_5$
$P_{c_2}(S(G_3))$	$k_2 + k_1 + k_5$	$k_3 + k_1 + k_5$	$k_4 - k_1 - 1$	$-k_5 - k_1 - 1$	$2k_1 + k_5$
$P_{c_3}(S(G_3))$	$k_2 + k_1 + k_5$	$k_3 + k_1 + k_5$	$k_4 + k_1 + k_5$	$k_1$	$-k_5 - 2k_1 - 1$

*Proof.* We only prove the case of  $h = c_1$ , namely,  $P_{c_1}(S(G_3)) = \{(l_1, \ldots, l_5) \in \mathbb{N}^5 : l_1 = k_1 - k_4, l_2 = k_2 + k_4 + k_5, l_3 = k_3 + k_4 + k_5, l_4 = -k_5 - k_4 - 1, l_5 = 2k_4 + k_5\}$ . Other cases are similar.

"  $\subseteq$  " If  $(k_1, k_2, k_3, k_4, k_5) \in P_{c_1}(S(G_3))$ , then we have  $-k_5 \in \mathbb{P}$ ,  $\min\{k_1, k_4\} < -k_5, k_1 \ge k_4$  and  $k_i \in \mathbb{N}(1 \le i \le 4)$ . We get  $l_1 = k_1 - k_4 \in \mathbb{N}$  and  $l_4 = -k_5 - k_4 - 1 \in \mathbb{N}$ . In addition, by Table 2.6 in Lemma 7, we get  $l_2 = k_2 + k_4 + k_5, l_3 = k_3 + k_4 + k_5$  and  $l_5 = 2k_4 + k_5$  are all in  $\mathbb{N}$ .

"  $\supseteq$  " If  $(l_1, l_2, l_3, l_4, l_5) \in \{(l_1, \dots, l_5) \in \mathbb{N}^5 : l_1 = k_1 - k_4, l_2 = k_2 + k_4 + k_5, l_3 = k_3 + k_4 + k_5, l_4 = -k_5 - k_4 - 1, l_5 = 2k_4 + k_5\}$ , then we get  $l_1 = k_1 - k_4 \in \mathbb{N}$  and  $l_4 = -k_5 - k_4 - 1 \in \mathbb{N}$ . So  $k_1 \ge k_4$  and  $\min\{k_1, k_4\} = k_4$ . Also, by inversely solving the

 $k_i$ 's from the  $l_i$ 's, we obtain  $k_5 = -2l_4 - l_5 - 2 < 0$ . Hence,  $-k_5 \in \mathbb{P}$  and  $\min\{k_1, k_4\} = k_4 < -k_5$ . Thus,  $(k_1, k_2, k_3, k_4, k_5) \in P_{c_1}(S(G_3))$ .

Now we are ready to state and prove our result.

**Theorem 9.** Let  $G_3$  be the graph in Figure 3. Then any  $\gamma \in S(G_3)$  can be uniquely written in one of the following five types.

 $\begin{array}{l} T_{a}: l_{1}\boldsymbol{\gamma}_{1} + l_{2}\boldsymbol{\gamma}_{2} + l_{3}\boldsymbol{\gamma}_{3} + l_{4}\boldsymbol{\gamma}_{4} + l_{5}\boldsymbol{\gamma}_{5}; \\ T_{b}: \boldsymbol{\gamma}_{8} + l_{1}\boldsymbol{\gamma}_{1} + l_{2}\boldsymbol{\gamma}_{2} + l_{3}\boldsymbol{\gamma}_{3} + l_{4}\boldsymbol{\gamma}_{4} + l_{5}\boldsymbol{\gamma}_{8}; \\ T_{c_{1}}: \boldsymbol{\gamma}_{6} + l_{1}\boldsymbol{\gamma}_{1} + l_{2}\boldsymbol{\gamma}_{2} + l_{3}\boldsymbol{\gamma}_{3} + l_{4}\boldsymbol{\gamma}_{6} + l_{5}\boldsymbol{\gamma}_{8}; \\ T_{c_{2}}: \boldsymbol{\gamma}_{4} + \boldsymbol{\gamma}_{6} + l_{1}\boldsymbol{\gamma}_{2} + l_{2}\boldsymbol{\gamma}_{3} + l_{3}\boldsymbol{\gamma}_{4} + l_{4}\boldsymbol{\gamma}_{6} + l_{5}\boldsymbol{\gamma}_{8}; \\ T_{c_{3}}: \boldsymbol{\gamma}_{7} + l_{1}\boldsymbol{\gamma}_{2} + l_{2}\boldsymbol{\gamma}_{3} + l_{3}\boldsymbol{\gamma}_{4} + l_{4}\boldsymbol{\gamma}_{6} + l_{5}\boldsymbol{\gamma}_{7}, \end{array}$ 

where  $\gamma_j$ ,  $1 \leq j \leq 8$  are given by Equation (2.4) and  $l_i \in \mathbb{N}$ ,  $1 \leq i \leq 5$ .

*Proof.* Let  $\gamma \in S(G_3)$  be written as

(2.9) 
$$\boldsymbol{\gamma} = k_1 \boldsymbol{\gamma}_1 + k_2 \boldsymbol{\gamma}_2 + k_3 \boldsymbol{\gamma}_3 + k_4 \boldsymbol{\gamma}_4 + k_5 \boldsymbol{\gamma}_5, \qquad k_1, \dots, k_5 \in \mathbb{R}.$$

Then  $\boldsymbol{\gamma} \in P_h(S(G_3))$  for  $h \in \{a, b, c_1, c_2, c_3\}$ .

We can use Table 2.5 to rewrite  $\gamma$ . The results are given in the following table.

Types	$\gamma$
$P_a(S(G_3))$	$k_1 \boldsymbol{\gamma}_1 + k_2 \boldsymbol{\gamma}_2 + k_3 \boldsymbol{\gamma}_3 + k_4 \boldsymbol{\gamma}_4 + k_5 \boldsymbol{\gamma}_5$
$P_b(S(G_3))$	$ig  oldsymbol{\gamma}_8 + (k_1+k_5)oldsymbol{\gamma}_1 + k_2oldsymbol{\gamma}_2 + k_3oldsymbol{\gamma}_3 + (k_4+k_5)oldsymbol{\gamma}_4 + (-k_5-1)oldsymbol{\gamma}_8$
$P_{c_1}(S(G_3))$	$\gamma_6 + (k_1 - k_4)\gamma_1 + (k_2 + k_4 + k_5)\gamma_2 + (k_3 + k_4 + k_5)\gamma_3$
	$+(-k_5-k_4-1)oldsymbol{\gamma}_6+(2k_4+k_5)oldsymbol{\gamma}_8$
$P_{c_2}(S(G_3))$	$\begin{vmatrix} \boldsymbol{\gamma}_4 + \boldsymbol{\gamma}_6 + (k_2 + k_1 + k_5)\boldsymbol{\gamma}_2 + (k_3 + k_1 + k_5)\boldsymbol{\gamma}_3 + (k_4 - k_1 - 1)\boldsymbol{\gamma}_4 \\ + (-k_5 - k_1 - 1)\boldsymbol{\gamma}_6 + (2k_1 + k_5)\boldsymbol{\gamma}_8 \end{vmatrix}$
$P_{c_3}(S(G_3))$	$\begin{vmatrix} \boldsymbol{\gamma}_{7} + (k_{2} + k_{1} + k_{5})\boldsymbol{\gamma}_{2} + (k_{3} + k_{1} + k_{5})\boldsymbol{\gamma}_{3} + (k_{4} + k_{1} + k_{5})\boldsymbol{\gamma}_{4} + k_{1}\boldsymbol{\gamma}_{6} \\ + (-k_{5} - 2k_{1} - 1)\boldsymbol{\gamma}_{7} \end{vmatrix}$

By Lemma 8,  $P_h(S(G_3))$  is transformed exactly to type  $T_h$  for each h.

**Corollary 10.** Let  $G_3$  be as above. Then we have the decomposition

$$F^{G_3}(\mathbf{x}, y) = \sum_{\boldsymbol{\gamma} \in S(G_3)} \mathbf{x}^{\boldsymbol{\gamma}} y^{s(\boldsymbol{\gamma})} = \sum_{i=1}^5 F_i^{G_3}(\mathbf{x}, y), \quad \text{where}$$

$$\begin{split} F_1^{G_3}(\mathbf{x}, y) &= \frac{1}{(1 - \mathbf{x}^{\gamma_1} y)(1 - \mathbf{x}^{\gamma_2} y)(1 - \mathbf{x}^{\gamma_3} y)(1 - \mathbf{x}^{\gamma_4} y^2)(1 - \mathbf{x}^{\gamma_5} y^2)};\\ F_2^{G_3}(\mathbf{x}, y) &= \frac{\mathbf{x}^{\gamma_8} y}{(1 - \mathbf{x}^{\gamma_8} y)(1 - \mathbf{x}^{\gamma_2} y)(1 - \mathbf{x}^{\gamma_3} y)(1 - \mathbf{x}^{\gamma_1} y)(1 - \mathbf{x}^{\gamma_4} y^2)}; \end{split}$$

$$\begin{split} F_3^{G_3}(\mathbf{x},y) &= \frac{\mathbf{x}^{\gamma_6} y}{(1-\mathbf{x}^{\gamma_6} y)(1-\mathbf{x}^{\gamma_2} y)(1-\mathbf{x}^{\gamma_3} y)(1-\mathbf{x}^{\gamma_1} y)};\\ F_4^{G_3}(\mathbf{x},y) &= \frac{\mathbf{x}^{(\gamma_4+\gamma_6)} y^3}{(1-\mathbf{x}^{\gamma_3} y)(1-\mathbf{x}^{\gamma_2} y)(1-\mathbf{x}^{\gamma_8} y)(1-\mathbf{x}^{\gamma_6} y)(1-\mathbf{x}^{\gamma_4} y^2)};\\ F_5^{G_3}(\mathbf{x},y) &= \frac{\mathbf{x}^{\gamma_7} y^2}{(1-\mathbf{x}^{\gamma_6} y)(1-\mathbf{x}^{\gamma_2} y)(1-\mathbf{x}^{\gamma_3} y)(1-\mathbf{x}^{\gamma_4} y^2)(1-\mathbf{x}^{\gamma_7} y^2)}. \end{split}$$

Consequently,

$$F^{G_3}(y) = \frac{1+2y+4y^2+2y^3+y^4}{(1-y)^3(1-y^2)^2} = 1+5y+18y^2+46y^3+101y^4+193y^5+\cdots$$

### 3. MACMAHON'S PARTITION ANALYSIS AND MAGIC DISTINCT LABELLINGS

We first introduce the basic idea of MacMahon's partition analysis and discuss possible applications of our results.

3.1. MacMahon's Partition Analysis. MacMahon's partition analysis was introduced by MacMahon in [10], and has been restudied by Andrews and his coauthors in a series of papers starting with [1]. The Mathematica package Omega was developed in [2]. The main idea of MacMahon's partition analysis is to replace linear constraints by using new variables and MacMahon's Omega operators on formal series:

$$\Omega \sum_{k=-\infty}^{\infty} c_i \lambda^i = c_0, \qquad \Omega \sum_{k=-\infty}^{\infty} c_i \lambda^i = \sum_{k=0}^{\infty} c_i.$$

MacMahon's operators always acting on the  $\lambda$  variables, which will be clear from the context. We explain to how to compute  $F^{G_1}(\mathbf{x}, y)$  in Example 1. By (2.1), we have

$$F^{G_1}(\mathbf{x}, y) = \sum_{(\boldsymbol{\alpha}, s) \in \mathbb{N}^7} x_1^{\alpha_1} \cdots x_6^{\alpha_6} y^s \underset{=}{\Omega} \lambda_1^{\alpha_1 + \alpha_4 + \alpha_5 - s} \lambda_2^{\alpha_1 + \alpha_2 + \alpha_6 - s} \lambda_3^{\alpha_2 + \alpha_3 + \alpha_5 - s} \lambda_4^{\alpha_3 + \alpha_4 + \alpha_6 - s}$$
$$= \underset{=}{\Omega} \frac{\lambda_1 \lambda_2 \lambda_3 \lambda_4}{D_1(\mathbf{x}, y)}, \text{ where}$$

$$D_1(\mathbf{x}, y) = (1 - \lambda_1 \lambda_4 x_1) (1 - \lambda_1 \lambda_2 x_2) (1 - \lambda_2 \lambda_3 x_3) (1 - \lambda_3 \lambda_4 x_4) (1 - \lambda_2 \lambda_4 x_5) (1 - \lambda_1 \lambda_3 x_6) (-y + \lambda_1 \lambda_2 \lambda_3 \lambda_4).$$

Eliminating the  $\lambda$ 's will give a representation of  $F^{G_1}(\mathbf{x}, y)$ . The whole theory relies on unique series expansion of rational functions. See [17] for the field of iterated Laurent series and the partial fraction algorithm implemented by the Maple package Ell. The maple package CTEuclid in [15] is better in most situations.

The normal form of  $F^{G_1}(\mathbf{x}, y)$  (by Maple) already has combinatorial meaning. The normal form of  $F^{G_2}(\mathbf{x}, y)$  is

$$F^{G_2} = \frac{1 - y^3 x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_9}{\left(1 - y x_7 x_8 x_9\right) \left(1 - y x_2 x_6 x_9\right) \left(1 - y x_1 x_5 x_8\right) \left(1 - y x_3 x_4 x_7\right) \left(1 - y^2 x_1 x_2 x_3 x_4 x_5 x_6\right)},$$

which can be easily decomposed by inspection.

For the graph  $G_3$ , CTEuclid gives an expression of  $F^{G_3}(\mathbf{x}, y)$  quickly, but the normal form of  $F^{G_3}(\mathbf{x}, y)$  is

$$F^{G_3}(\mathbf{x}, y) = \frac{N(\mathbf{x}, y)}{(1 - yx^{\gamma_1})(1 - yx^{\gamma_2})(1 - yx^{\gamma_3})(1 - y^2x^{\gamma_4})(1 - y^2x^{\gamma_5})(1 - yx^{\gamma_6})(1 - y^2x^{\gamma_7})(1 - yx^{\gamma_8})},$$

where

$$\begin{split} N(\mathbf{x},y) =& 1 - y^3 x_1 x_2 x_4 x_5 x_6 x_7 x_8^2 x_9^2 x_{10} x_{12} - y^3 x_1 x_2 x_3 x_5 x_6 x_7 x_8 x_9^2 x_{10}^2 x_{11} \\ &\quad - 2 y^3 x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_9 x_{10} x_{11} x_{12} - y^4 x_2^2 x_3 x_4 x_5^2 x_6^2 x_7^2 x_8 x_9^2 x_{10} x_{11} x_{12} \\ &\quad + y^4 x_1^2 x_2 x_3 x_4 x_5 x_6 x_7 x_8^2 x_9^2 x_{10}^2 x_{11} x_{12} + 2 y^5 x_1 x_2^2 x_3 x_4 x_5^2 x_6^2 x_7^2 x_8^2 x_9^3 x_{10}^2 x_{11} x_{12} \\ &\quad + y^5 x_1 x_2^2 x_3 x_4^2 x_5^2 x_6^2 x_7^2 x_8^2 x_9^2 x_{10} x_{11} x_{12}^2 + y^5 x_1 x_2^2 x_3^2 x_4 x_5^2 x_6^2 x_7^2 x_8 x_9^2 x_{10}^2 x_{11}^2 x_{12} \\ &\quad - y^8 x_1^2 x_2^3 x_3^2 x_4^2 x_5^3 x_6^3 x_7^3 x_8^3 x_9^4 x_{10}^3 x_{11}^2 x_{12}^2, \end{split}$$

is polynomial of 10 terms. It is not clear how to decompose  $F^{G_3}(\mathbf{x}, y)$  as a sum of simple rational functions. We guessed such a decomposition (in Corollary 10) by certain criterion. The verification of the formula by computer is easy.

We should mention that for some complicated graphs G, Maple will stuck when normal  $F^{G}(\mathbf{x}, y)$ .

We conclude the subsection by reporting the following result. Let  $G_4$  be given in Figure 4, with 6 vertices and 9 edges.

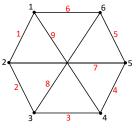


FIGURE 4. Regular graph  $G_4$  with 6 vertices and 9 edges.

Then

 $(3.1) F^{G_4}(\mathbf{x}, y) = \frac{yx_2x_4x_6}{(1 - yx_1x_4x_8)(1 - yx_2x_4x_6)(1 - yx_2x_5x_9)(1 - yx_3x_6x_7)(1 - yx_1x_3x_5)} \\ + \frac{1}{(1 - yx_1x_4x_8)(1 - yx_2x_5x_9)(1 - yx_3x_6x_7)(1 - yx_7x_8x_9)(1 - yx_1x_3x_5)} \\ + \frac{y^2x_2x_4x_6x_7x_8x_9}{(1 - yx_7x_8x_9)(1 - yx_3x_6x_7)(1 - yx_2x_5x_9)(1 - yx_2x_4x_6)(1 - yx_1x_4x_8)}.$ 

(3.2) 
$$F^{G_4}(y) = \frac{1+y+y^2}{(1-y)^5} = 1+6y+21y^2+55y^3+120y^4+231y^5+406y^6+\cdots$$

It is not hard to give a combinatorial proof using similar ideas.

3.2. Magic Labellings. The complete generating function  $F^G(\mathbf{x}, y)$  encodes almost all information of S(G).

Let  $S^*(G)$  be the set of magic distinct labellings of G. In some literature, the technique condition "positive" is added because it is possible that any  $\boldsymbol{\alpha} \in S(G)$  must have label 0 on some edges. For instance, if G is given by Figure 5 (a), then it is easy to check that S(G) only contains the all 0 labelling; if G is given by Figure 5 (b), then for any  $\boldsymbol{\alpha} \in S(G)$  the labels of 2 and 3 must be 0. In deed, we have  $F^{G_{5(b)}}(\mathbf{x}, y) = \frac{1}{1-yx_1x_4}$ .

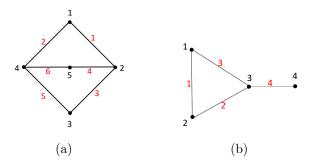


FIGURE 5. Graph  $G_{5(a)}$  with 5 vertices and 6 edges and graph  $G_{5(b)}$  with 4 vertices and 4 edges.

For a *d* regular graph *G*, the all 1 labelling **1** is magic with magic sum *d*. Thus  $S_{\mathbb{P}}(G) = \mathbf{1} + S(G) = \{\mathbf{1} + \boldsymbol{\alpha} : \boldsymbol{\alpha} \in S(G)\}$ . In this sense, the positive condition makes no difference for magic labellings of a regular graph.

The graphs  $G_1, G_2, G_3$  in our examples are all regular graphs. It is an accident that none of them have magic distinct labellings. Indeed, they do not have magic distinct  $\mathbb{R}$ -labellings. To see this for  $G_3$ , any  $\gamma \in S_{\mathbb{R}}(G_3)$  can be written as in (2.7) for some  $k_1, \ldots, k_5 \in \mathbb{R}$ . Then the  $a_2$  and  $a_6$  labels are the same. The situation for the other two graphs are similar: look at the  $a_1$  and  $a_3$  labels for  $G_1$ , and the  $a_1$  and  $a_5$  labels for  $G_2$ .

In general, the structure of  $S^*(G)$  is pretty complicated. It is obtained by slicing out all  $\boldsymbol{\alpha} \in S(G)$  in the  $\binom{n}{2}$  hyper planes  $\alpha_i = \alpha_j$ ,  $1 \leq i < j \leq n$ . Using inclusion and exclusion principle will be too expensive since that will involve  $2^{\binom{n}{2}}$  cases. It is possible to obtain the generating function

$$E^G(\mathbf{x}, y) = \sum_{\boldsymbol{\alpha} \in S^*(G)} \mathbf{x}^{\boldsymbol{\alpha}} y^{s(G)}$$

by MacMahon's partition analysis.

Here we introduce two operators that can be realized by MacMahon's partition analysis. If  $A(x,y) = \sum_{i,j\geq 0} a_{i,j}x^iy^j$  is a formal power series in x and y. Then the diagonal operator defined by

$$\operatorname{diag}_{x,y} A(x,y) = \sum_{i \ge 0} a_{i,i} x^i y^i = \sum_{i,j \ge 0, \ i-j=0} a_{i,j} x^i y^j$$

can be realized by MacMahon's Omega (linear) operator. We have

$$\operatorname{diag}_{x,y} A(x,y) = \sum_{i,j\geq 0} a_{i,j} x^i y^j \underset{=}{\Omega} \lambda^{i-j} = \underset{i,j\geq 0}{\Omega} \sum_{i,j\geq 0} a_{i,j} (\lambda x)^i (y/\lambda)^j = \underset{=}{\Omega} A(\lambda x, y/\lambda).$$

Similarly, if we define

$$\operatorname{diag}_{x>y} A(x,y) = \sum_{i>j\ge 0} a_{i,j} x^i y^j = \sum_{i,j\ge 0, \ i-j-1\ge 0} a_{i,j} x^i y^j$$

then it can be realized by

$$\operatorname{diag}_{x>y} A(x,y) = \sum_{i,j\geq 0} a_{i,j} x^i y^j \underset{\geq}{\Omega} \lambda^{i-j-1} = \underset{\geq}{\Omega} \sum_{i,j\geq 0} a_{i,j} \lambda^{-1} (\lambda x)^i (y/\lambda)^j = \underset{\geq}{\Omega} \lambda^{-1} A(\lambda x, y/\lambda) A(\lambda x, y/\lambda)$$

We have

$$E^{G}(\mathbf{x}, y) = \prod_{1 \le i < j \le n} (1 - \operatorname{diag}_{x_i, x_j}) F^{G}(\mathbf{x}, y),$$

whose expansion is just the inclusion and exclusion result. We can normal the result after each application of  $1 - \text{diag}_{x_i,x_j}$ , provided that the result would not explode, i.e., the numerator becomes too large for Maple to handel. Note that the computation highly relies on the order of the operators. For instance,  $(1 - \text{diag}_{x_2,x_6})F^{G_3}(\mathbf{x}, y) = 0$ . But the result quickly explodes for some orders.

Another way is to use the natural decomposition

$$S^*(G) = \biguplus_{\pi \in \mathfrak{S}_n} S^{\pi}(G),$$

where  $\pi$  ranges over all permutations of  $\{1, 2, ..., n\}$ , and  $S^{\pi}(G)$  consists of all  $\boldsymbol{\alpha} = (\alpha_1, ..., \alpha_n) \in S(G)$  compatible with  $\pi$ , i.e.,  $\alpha_{\pi_1} > \alpha_{\pi_2} > \cdots > \alpha_{\pi_n}$ .

The generating function  $E^{G,\pi}(\mathbf{x}, y)$  of  $S^{\pi}(G)$  can be extracted from  $F^{G}(\mathbf{x}, y)$  by applying MacMahon's Omega operator. We have

$$E^{G,\pi}(\mathbf{x},y) = \operatorname{diag}_{x_{\pi_1} > x_{\pi_2}} \cdots \operatorname{diag}_{x_{\pi_{n-1}} > x_{\pi_n}} F^G(\mathbf{x},y).$$

For  $F^{G_4}(\mathbf{x}, y)$  with the combinatorial decomposition as  $F_1 + F_2 + F_3$  as in (3.1), and for each  $F_i$ , i = 1, 2, 3, we can extract

$$E_i^{\pi} = \operatorname{diag}_{x_{\pi_1} > x_{\pi_2}} \cdots \operatorname{diag}_{x_{\pi_8} > x_{\pi_9}} F_i(\mathbf{x}, y).$$

for any particular  $\pi \in \mathfrak{S}_9$ . Only  $432 = 2^4 3^3$  out of 362880 = 9! permutations give nonvarnishing results. And the three sets of permutations do not overlap. Each result are simple rational functions with numerator either a monomial or a binomial. For instance,

$$E_1^{134568279} = \frac{x_4^7 x_2^3 x_6^5 y^{15} x_1^{10} x_3^8 x_5^6 x_8^4 x_7^2 \left(1 - x_4^7 x_2^3 x_6^5 y^{15} x_1^{10} x_3^8 x_5^6 x_8^4 x_7^2\right)}{D(\mathbf{x}, y)},$$

where

$$D(\mathbf{x}, y) = \left(1 - y^3 x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_9\right) \left(1 - y^4 x_1^3 x_2 x_3^2 x_4^2 x_5^2 x_6 x_8\right) \\ \times \left(1 - y^5 x_1^3 x_2 x_3^3 x_4^2 x_5^2 x_6^2 x_7 x_8\right) \left(1 - y^6 x_1^4 x_2 x_3^3 x_4^3 x_5^2 x_6^2 x_7 x_8^2\right) \\ \times \left(1 - y^7 x_1^5 x_2 x_3^4 x_4^3 x_5^3 x_6^2 x_7 x_8^2\right) \left(1 - y^8 x_1^5 x_2^2 x_3^4 x_4^4 x_5^3 x_6^3 x_7 x_8^2\right).$$

There are total of  $2^4 3^4 = 1296$  permutations  $\pi$  such that  $E^{G_4,\pi}(\mathbf{x}, y) \neq 0$ . As a consequence, we have

$$(3.3) \quad E^{G_4}(y) = \frac{72y^{12} (1-y)^2 N_4(y)}{(1-y^3)^2 (1-y^4) (1-y^5) (1-y^6) (1-y^7) (1-y^8)} \\ = 72(y^{12}+2y^{13}+4y^{14}+8y^{15}+12y^{16}+20y^{17}+29y^{18}+42y^{19}+54y^{20}+\cdots)$$

where  $N_4(y)$  is given by

$$\frac{1+4\,y+11\,y^2+24\,y^3+44\,y^4+73\,y^5+109\,y^6+152\,y^7+192\,y^8+233\,y^9+258\,y^{10}}{+274\,y^{11}+268\,y^{12}+249\,y^{13}+207\,y^{14}+166\,y^{15}+117\,y^{16}+79\,y^{17}+41\,y^{18}+18\,y^{19}+18\,y^{1$$

The 72 in the numerator seems a surprise. Observe that the symmetry group of  $G_4$  is the Dihedral group  $D_6$  which is of cardinality 12. Thus we have the following result.

**Corollary 11.** Let  $\tilde{h}(s)$  be the number of magic distinct labellings of  $G_4$  of magic sum s up to isomorphism under the Dihedral group  $D_6$ . Then it is divisible by 6 for all s.

*Proof.* Clearly we have

$$E^{G_4}(y)/12 = \sum_{s \ge 0} \widetilde{h}(s)y^s.$$

The corollary then follows by the formula (3.3).

Corollary 11 needs a combinatorial proof. We list in Figure 6 all non-isomorphic magic labellings of G4 with minimum magic sum s = 12.

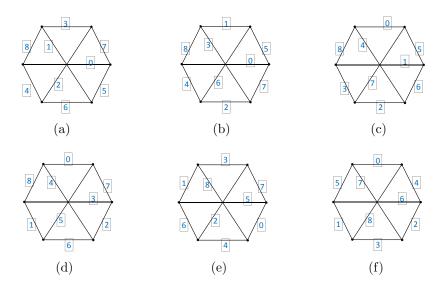


FIGURE 6. All non-isomorphic magic labellings of  $G_4$  with minimum magic sum s = 12. The edges are labeled by the blue boxed numbers.

## 4. Concluding Remark

We have studied the complete construction of magic labelling of graphs S(G). Our aim is to decompose S(G) into some shifted free monoids. We have achieved this for four graphs, and give combinatorial proofs of the decompositions.

In general, there are algorithms to compute the generating functions  $F^G(\mathbf{x}, y)$ . Then the decomposition corresponds to algebraic decomposition of  $F^G(\mathbf{x}, y)$ . Such a decomposition seems easier to attack, and it is a guide for combinatorial proofs.

Our approach to magic distinct labellings is by using MacMahon's partition analysis, especially the Maple package CTEuclid. The package extracts constant term of an Elliott-rational function, i.e., a rational function whose denominator is a product of binomials. The number of binomials in the denominator affects the performance of Maple significantly. This is why we prefer a good decomposition of  $F^G(\mathbf{x}, y)$ .

Magic distinct labellings of the cube was studied in [18], where the cube has 8 vertices and 12 edges. The generating function is more complicated than that of  $G_4$ . It seems that the more edges the graph have, the more complex the generating function is.

Another direction is to restrict the number of vertices. In an upcoming paper, we will report the results for all graphs with 5 vertices.

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