Constrained NLP via Gradient Flow Penalty Continuation: Towards Self-Tuning Robust Penalty Schemes

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Abstract

This work presents a new numerical solution approach to nonlinear constrained optimization problems based on a gradient flow reformulation. The proposed solution schemes use self-tuning penalty parameters where the updating of the penalty parameter is directly embedded in the system of ODEs used in the reformulation, and its growth rate is linked to the violation of the constraints and variable bounds. The convergence properties of these schemes are analyzed, and it is shown that they converge to a local minimum asymptotically. Numerical experiments using a set of test problems, ranging from a few to several hundred variables, show that the proposed schemes are robust and converge to feasible points and local minima. Moreover, results suggest that the GF formulations were able to find the optimal solution to problems where conventional NLP solvers fail, and in less integration steps and time compared to a previously reported GF formulation.

Keywords: gradient flow, nonlinear programming problem, convergence analysis

1 1. Introduction

Optimization problems arise in many areas of chemical engineering practice, from component 2 and systems design [1] to operation and control [2]. Due to an increasing concern of legislators 3 and the general public in environmental sustainability, optimization has been recently used to aid 4 in the design of supply chains and products considering their life cycle [3], and as a tool for the 5 design of new sustainable energy conversion systems [4, 5]. Applications encompass formulations 6 ranging from linear programming problems (LP) to mixed-integer non-linear programming problems (MINLP) and dynamic optimization problems (optimal control problems, OCP). A common 8 feature of many of these classes of problems, is that at a certain point one or several non-linear con-9 strained programming problems (NLP) need to be solved. The solution of large-scale NLP problems 10 was made possible by breakthroughs in non-linear programming during the previous decades. In 11 particular, the development of modern barrier methods [6, 7, 8], sequential quadratic programming 12 [9] and reduced gradient methods [10], led to implementations (solvers) such as IPOPT [8], SNOPT 13

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¹⁴ [9] and CONOPT [10] that can be used in user-friendly modeling and optimization environments ¹⁵ such as GAMS [11], AMPL [12] and AIMMS [13].

Most algorithms used to compute a local optimum of constrained NLP problems rely on Taylor 16 series expansions truncated after the linear or quadratic term; according to this, constraints are 17 linearized and large steps towards the local minimum are allowed. For this reason, in highly 18 nonlinear problems intermediate iterations might prove infeasible and frequent failures to converge 19 to a local optimum may arise. Alternatively to the Taylor expansion based methods, Gradient Flow 20 (GF) methods have been proposed for the solution of unconstrained and constrained nonlinear 21 programming problems. In its most simple version, the solution of an unconstrained problem 22 $\min_x f(x)$ can be obtained by solving the following set of coupled ordinary differential equations 23 (ODEs): 24

$$\frac{dx}{dt} = -\nabla_x f(x); \quad x(0) = x_0 \tag{1}$$

where $x \in \mathbb{R}^n$, $f(x) : \mathbb{R}^n \to \mathbb{R}^1$. This approach creates a smooth trajectory that might offer an advantage for highly nonlinear problems compared with the conventional optimization techniques which take finite steps along line-search directions. For the latter, finding a suitable step-size can be difficult when the optimization function is non-quadratic and has large third derivatives, resulting in a slow progress towards the solution due to the smalls steps required [14].

Another interesting feature of GF methods is the possibility of using state-of-the-art integration 30 software to find the solution of optimization problems. This approximation for the solution of 31 unconstrained problems can be traced to the work of Botsaris [15]. In the following decades, efforts 32 were made to reach a competitive level in terms of computational time and iterations compared 33 to conventional methods, with a summary found in Brown and Bartholomew-Biggs [14]. The 34 application of GF methods was further extended by introducing new formulations that were able 35 to cope with constrained nonlinear problems [16, 17, 18, 19]. The constraints of the NLP problem 36 (h(x)) are incorporated to the objective function (f(x)) with a penalty scheme in order for GF 37 methods to be employed, with one of the major issues being the updating of the penalty parameters 38 utilized. For an optimization problem with equality constraints only, Tanabe [20] proposed the 39 following Gradient Flow formulation: 40

$$\frac{dx}{dt} = -Q(x)\nabla_x f(x); \quad x(0) = x_0$$

$$Q(x) = [I - J^T(x)(J(x)J^T(x))^{-1}J(x)]$$
(2)

where $\nabla_x f(x)$ and J(x) represents the gradient of the objective function with respect to the optimization variables and the Jacobian matrix, respectively. Equation 2 is a direct generalization of the gradient projection method proposed by Rosen [21] to a differential form, which is based on projecting the search direction given by $\nabla_x f(x)$ into the subspace tangent to the active constraints. The method proposed by Tanabe [20] was further modified by Yamashita [22] and Evtushenko and Addah [17] to yield

$$\frac{dx}{dt} = -sQ(x)\nabla_x f(x) + -J(x)(J(x)J^T(x))^{-1}h(x); \quad x(0) = x_0$$
(3)

with s a positive constant. Following the work of Evtushenko and Zhadan [17], Wang et al. [23] proposed an approach to include inequality constraints and bounds where a pseudo-inverse of the square matrix $J(x)J^{T}(x)$ acts as a penalizer (equation 4), with this approach requiring a nonsingular Jacobian.

$$\frac{dx}{dt} = -[\nabla f(x) + J^T(x)(\tau h(x) - (J(x)J^T(x))^{-1}J(x)\nabla f(x))]; \quad x(0) = x_0$$
(4)

In their work, they avoid the use of slacks to account for variable bounds by using the state-space 51 transformation technique proposed by Evtushenko and Zhadan [17], otherwise the use of quadratic 52 slacks would result in singular Jacobians. As proved by the above mentioned authors, their GF 53 formulations have the property that once the equality constraints are satisfied, the trajectory of the 54 solution will stay on the manifold determined by the constraints. However, as analyzed by Brown 55 and Bartholomew-Biggs [16], the ODE system that allows following a path with those characteristics 56 needs to be solved quite accurately. This and the fact that inverses of large matrices need to be 57 calculated, produce a heavy numerical overhead. Moreover, in the formulations represented by 58 equations 2 to 4, authors do not present an approach to select the values of the parameters required 59 in their formulations (such as τ in equation 4). In practice, the value of these parameters needs 60 to be adjusted to each problem. Finally, Schropp and Singer [24] compare SQP methods and GF 61 methods for the solution of nonlinear problems from a theoretical point of view and using two case 62 studies. They concluded that SQP methods are efficient tools whereas the ODE approach may be 63 more reliable, with the ODE approach being more efficient for problems with only a small number 64 of highly nonlinear constraints. Moreover, they propose an approach combining differential and 65 algebraic equations that, according to the authors, combines efficiency and reliability. 66

In this work, a self-tuning penalty scheme is presented for the solution of constrained NLPs. The approach does not require the calculation of an inverse (or pseudo-inverse) of the Jacobian matrix, and the penalty parameters updating is directly embedded into the system of ODEs. The performance of the GF formulations presented in this work are compared to the formulation presented by Wang et al. [23] and also against several state-of-the-art NLP solvers.

72 2. New formulations using the Gradient Flow approach for solving NLP problems

73 2.1. Problem definition

⁷⁴ The minimization of the following standard constrained NLP is considered:

$$\min_{x_s} f(x_s)$$
subject to:
$$h_s(x_s) = 0$$

$$g(x_s) \le 0$$

$$r^L \le r_s \le r^U$$
(5)

where $x_s \in \mathbb{R}^{n_s}$, $f(x_s) : \mathbb{R}^{n_s} \mapsto \mathbb{R}^1$, $h_s(x_s) : \mathbb{R}^{n_s} \mapsto \mathbb{R}^{m_1}$ and $g(x_s) : \mathbb{R}^{n_s} \mapsto \mathbb{R}^{m_2}$. The subscript s stands for standard, as this problem will be converted to a penalized version were the subscripts will be dropped to simplify the notation. This problem is converted to a penalty function minimization, using a quadratic penalty scheme and standard transformations. Inequality constraints are converted to equalities via the use of squared slack variables as follows. First, inequality constraints are converted to equality constraints using the following transformation:

$$g(x_s) + w^2 = 0 \tag{6}$$

81 where $w \in \mathbb{R}^{m_2}$.

⁸² Variables bounds are transformed to equalities, by using the following equations:

$$x_s + \left(s^U\right)^2 = x_s^U \tag{7}$$

$$x_s - \left(s^L\right)^2 = x_s^L \tag{8}$$

where s^L , $s^U \in \mathbb{R}^{n_s}$. The equality constraints defined by equations 6 to 8 and the original constraints, $h_s(x)$, will be appended in the vector $h(x) : \mathbb{R}^n \mapsto \mathbb{R}^m$ with $n = 3n_s + m_2$, $m = 2n_s + m_1 + m_2$ and $x = \{x_s, s^U, s^L, w\}$. Using this new defined set of constraints and variables, the original problem posed in equation 5 can be redefined as the following (higher dimensionality) optimization problem:

$$\min_{x} f(x) \tag{9}$$

subject to:

$$h(x) = 0$$

⁸⁷ with the Lagrangian of the problem defined by:

$$L(x,\mu) = f(x) + \lambda^T h(x) \tag{10}$$

⁸⁸ A pair of points (x^*, λ^*) is said to be a stationary point of equation 10 if the following first order ⁸⁹ necessary conditions (Karush-Kuhn-Tucker conditions, KKT) are satisfied:

$$\nabla_x L(x^*, \lambda^*) = \nabla_x f(x^*) + J^T(x^*)\lambda^* = 0$$
(11)

$$\nabla_{\mu} L(x^*, \lambda^*) = h(x^*) = 0 \tag{12}$$

where $\nabla_x f(x^*)$ is the gradient vector of the objective function $(n \times 1 \text{ rows and columns})$, h(x) the vector of equality constraints $(m \times 1)$, λ the vector of Lagrange multipliers $(m \times 1)$ and J(x) the Jacobian matrix $(m \times n)$. Furthermore, the second order sufficient conditions require that for some feasible point x^* and a Lagrange multiplier vector λ^* , if

- ⁹⁴ 1. linear independence constraint qualification (LICQ) holds at x^* , and
- 2. for any vector d satisfying $J(x^*)d = 0$ holds that: $d^T \nabla_{xx} L(x^*, \lambda^*)d > 0$,

then x^* is a strict local solution of 9. In the following, it will be assumed that problem 9 show this properties.

⁹⁸ Using a penalty function, the problem defined by equation 9 can be stated as:

$$\min_{x} P_{PP}(x; M) = f(x) + \frac{1}{2} M h(x)^{T} h(x)$$
(13)

where M is a positive penalty parameter. This is a classical approach used to solve the original NLP
problem that is notorious for yielding badly conditioned unconstrained problems for conventional
NLP solvers as M is increased [25].

In the following section, a gradient flow method with a novel self-tuning scheme for updating the
 value of the penalty parameter is presented.

¹⁰⁴ 2.2. A gradient flow formulation for constrained NLP problems

Considering the approach for the solution of unconstrained NLP problems represented by equation 1, the unconstrained formulation of the originally constrained NLP (equation 9) can be written as the following set of coupled ODEs:

$$\frac{dx}{dt} = -\nabla_x P_{PP}$$

$$\frac{dx}{dt} = -\left[\nabla_x f(x) + M J(x)^T h(x)\right]; \quad x(0) = x_0$$
(14)

108 where $0 \leq t \leq +\infty$.

In order for this scheme to be successful, the updating of the penalty parameter needs to be embedded in a coupled way within the GF scheme and the value of the penalty parameter can be linked to the constraint norm. Thus, considering a $\kappa > 0$ the value of the penalty can be formulated as

$$\frac{dM}{dt} = \kappa \|h(x)\|_2; \quad M(0) = M_0$$
(15)

On the other hand, by considering a simple updating scheme according to the original penaltymultiplier approach (Hestenes method) [26], following the minimization at any iteration (k), if $\|h^{(k)}\| \geq \gamma \|h^{(k-1)}\|$ with $0 < \gamma < 1$, e.g. $\gamma = 0.25$, then the penalty parameter is increased by $M^{(k+1)} = \alpha \cdot M^{(k)}$. To derive a continuous variant, suitable for embedding in a GF methodology, the following algebraic steps are considered:

$$M^{(k+1)} - M^{(k)} = \kappa \cdot M^{(k)} - M^{(k)}$$
(16)

$$\Delta M^{(k)} = (\kappa - 1) \cdot M^{(k)} \tag{17}$$

with $(\kappa - 1) > 0$ and by renaming $(\kappa - 1) \rightarrow \alpha > 1$ in the limit it can be obtained that

$$\frac{dM}{dt} = \alpha M; \quad M(0) = M_0 \tag{18}$$

The scheme should allow the possibility of not increasing the penalty parameter if sufficient progress towards feasibility is made, so that the formulations in equations (15) and (18) may be combined as

$$\frac{dM}{dt} = \alpha M \|h(x)\|_2; \quad M(0) = M_0$$
(19)

Since the value of M is updated within the integration of the ODE system, and embeds a type of exponential growth for M when the norm of the constraints is large, the above scheme will limit the necessary value for α . Combining the penalty parameter differential equation and the set of coupled ODEs representing the original GF approach, the following novel scheme is obtained (scheme PP)

$$\frac{dx}{dt} = -\left[\nabla_x f(x) + MJ(x)^T h(x)\right]; \quad x(0) = x_0$$
(20a)

$$\frac{dM}{dt} = \alpha M^{\beta} \|h(x)\|_{2}; \quad M(0) = M_{0}$$
(20b)

with $0 \le t \le +\infty$. The parameter β can take any positive value, or be set to zero, however some considerations are required. Since, the term M^{β} was included to act as an acceleration parameter of the trajectory of M(t), it is convenient to analyze the behavior of the following equation

$$\frac{dM_p}{dt} = M_p^\beta; \quad M_p(0) = M_{p0} \tag{21}$$

for different values of β , where M_p represents a simplified version of the trajectory of the penalty parameter M. Clearly, this equation is a first order separable ODE. Its solution depends on the value of β according to

$$M_{p}(t) = \begin{cases} (1-\beta)^{1/(1-\beta)} \left[\frac{M_{p0}^{1-\beta}}{1-\beta} + t \right]^{1/(1-\beta)} & 0 \le \beta < 1\\ e^{t} & \beta = 1\\ (\beta-1)^{1/(1-\beta)} \left[\frac{(\beta-1)M_{p0}^{\beta-1}}{1-(\beta-1)M_{p0}^{\beta-1}t} \right]^{1/(\beta-1)} & \beta \ge 1 \end{cases}$$
(22)

¹³² When $\beta \geq 1$, the values of M_{p0} are restricted to be lower than $[(\beta - 1)t_f]^{(1-\beta)}$, where t_f is ¹³³ the final integration time, to avoid a zero value in the denominator of the equation defining $M_p(t)$. ¹³⁴ Figure 1 shows the trajectories of $M_p(t)$ for different values of β and a final integration time of 10 ¹³⁵ units. For $\beta = \{0, 0.5, 1\}$, M_{p0} was set to 1.0 while for values of β equal to 2 and 3, the value of ¹³⁶ M_{p0} was taken as $0.99 \cdot [(\beta - 1)t_f]^{(1-\beta)}$. The trajectories for $\beta = \{2, 3\}$ show smaller values when ¹³⁷ compared to β equal to zero or greater. Since the purpose of this function is to increase the rate ¹³⁸ of growth of the penalty parameter (M), only β values of 0 and 1 will be considered for further ¹³⁹ analysis as they represent two extreme cases.

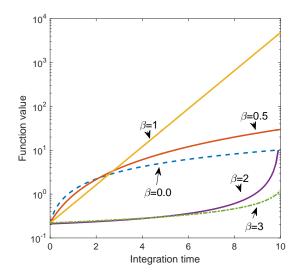


Figure 1: Trajectories of the acceleration parameter M_p as a function of time for different values of β .

¹⁴⁰ 2.2.1. Incorporating the Hestenes method to a gradient flow approach

The Hestenes multiplier scheme considers an Augmented Lagrangian for the NLP problem in equation 9 as follows [26]

$$\min_{x} P_{PM}(x; M, \lambda) = f(x) + \tilde{\lambda}^T h(x) + M \frac{1}{2} h(x)^T h(x)$$
(23)

Following minimization for a given value of $M = M^{(k)}$, the minimizer $x^*(M^{(k)})$ is obtained. The Lagrange multipliers are updated by

$$\Delta \tilde{\lambda}^{(k)} = M^{(k)} h(x^*(M^k)) \tag{24}$$

For this scheme it is known that there is a lower value of <u>M</u> for which if $M \ge \underline{M}$ it converges to the true solution of the NLP (x^*) [27]. Considering that the updating of the multipliers can be embedded also as an ODE in a GF scheme, the following formulation results (scheme PM)

$$\frac{dx}{dt} = -\left[\nabla_x f(x) + J(x)^T \left(\tilde{\lambda} + M h(x)\right)\right]; \quad x(0) = x_0$$
(25a)

$$\frac{d\lambda}{dt} = M h(x); \quad \tilde{\lambda}(0) = \tilde{\lambda}_0 \tag{25b}$$

$$\frac{dM}{dt} = \alpha M^{\beta} \|h(x)\|_{2}; \quad M(0) = M_{0}$$
(25c)

with $0 \le t \le +\infty$.

¹⁴⁹ 2.2.2. Proofs of convergence and stability analysis

C

In the following section, it will be proven that if x^* is a stationary point of the NLP defined by 9, then x^* is an equilibrium point of the PP scheme and vice versa (theorems 1 to 4). Theorems ¹⁵² 5 and 6 show that the PM scheme converges to a stationary point of the original NLP problem, ¹⁵³ and that a solution of the aforementioned problem is a stationary point of the PM scheme. Finally, ¹⁵⁴ theorems 7 and 8 establish that the equilibrium points are asymptotically stable and theorems 9 ¹⁵⁵ and 10 show that the penalty functions used in both schemes define a strictly decreasing trajectory.

Theorem 1. If x^* is the optimal solution to the unconstrained problem 13 for the penalty parameter M^{*}, then (x^*, M^*) is the equilibrium point for the system of ODEs defined by equations 20a and 20b with $\beta = 0$.

¹⁵⁹ Proof. According to the assumptions of the theorem, first order necessary conditions hold. Then,

$$\nabla_x P_{PP}(x^*; M^*) = -\frac{dx^*}{dt} = 0$$
 (26)

Furthermore, if x^* is the optimal solution of problem 13, it is also a stationary point of the NLP defined by 9 (see theorem 17.1 in Nocedal and Wright [28]). Thereby, $h(x^*) = 0$ and $\frac{dM}{dt} = \alpha \|h(x^*)\|_2 = 0$. Thus, (x^*, M^*) is the equilibrium point of the dynamic system formed by equations 20a and 20b for $\beta = 0$.

Theorem 2. If x^* is the optimal solution to the unconstrained problem 13 for the penalty parameter M^{*} and M^{*}h(x^{*}) = λ^* is the vector of Lagrange multipliers at a stationary point of the NLP defined by 9, then x^{*} is the equilibrium point of equation 20a and the left hand side of equation 20b with $\beta = 1$ is constant.

¹⁶⁸ *Proof.* According to the assumptions of the theorem,

$$\nabla_x P_{PP}(x^*; M^*) = -\frac{dx^*}{dt} = \nabla_x f(x^*) + M^* J(x^*)^T h(x^*) = \nabla_x f(x^*) + J(x^*)^T \lambda^* = 0$$
(27)

since the first order necessary conditions for problems 13 and 9 are satisfied. Now,

$$\lim_{t \to \infty} \frac{dM}{dt} = \lim_{t \to \infty} \alpha \sqrt{Mh(x)^T Mh(x)} = \alpha \lambda^{*T} \lambda^*$$
(28)

since according to theorem 17.2 in Nocedal and Wright [28], for a sufficiently large value of M^* the vector of Lagrange multipliers that satisfies the KKT conditions for the NLP defined by 9 can be approximated by $M^*h(x^*) = \lambda^*$. Thereby, x^* is the stationary point of equation 20a and the left hand side of equation 20b takes a constant value asymptotically.

Theorem 3. If (x^*, M^*) is the equilibrium point of the ODE system defined by equations 20a and 20b with $\beta = 0$, then x^* is an optimal solution of the unconstrained optimization problem defined by 13 and the NLP defined by 9.

¹⁷⁷ *Proof.* Since (x^*, M^*) is the equilibrium point of the ODE system 20a and 20b, then the first order ¹⁷⁸ necessary conditions of problem 13 are satisfied:

$$\nabla_x P_{PP}(x^*; M^*) = \nabla_x f(x^*) + M^* J(x^*)^T h(x^*) = 0$$
(29)

$$\nabla_M P_{PP}(x^*; M^*) = \frac{1}{2} h(x^*)^T h(x^*) = 0$$
(30)

¹⁷⁹ satisfaction of $\nabla_x P_{PP}(x^*; M^*) = 0$ is clear since $\nabla_x P_{PP}(x^*; M^*) = -\frac{dx^*}{dt} = 0$. Satisfaction of ¹⁸⁰ equation 30 is ensured by:

$$\frac{dM}{dt} = \alpha \|h(x^*)\|_2 = \alpha \sqrt{h(x^*)^T h(x^*)} = 0$$
(31)

The second order sufficient optimality condition requires that Hessian matrix of the penalty function defined by 13 to be positive definite:

$$\nabla_{xx} P_{PP}(x^*; M^*) = \nabla_{xx} f(x^*) + \sum_{i=1}^m M^* h_i(x^*) \nabla_{xx} h_i(x^*) + M^* J^T(x^*) J(x^*)$$
(32)

183 Replacing the Hessian matrix of problem 13:

$$\nabla_{xx} P_{PP}(x^*, M^*) = \nabla_{xx} L(x^*, \lambda^*) + M^* J^T(x^*) J(x^*)$$
(33)

Now, consider a vector d satisfying $J(x^*)d = 0$, multiplying $\nabla_{xx}P_{PP}(x^*, M^*)$ at both sides:

$$d^{T}[\nabla_{xx}L(x^{*},\lambda^{*}) + M^{*}J^{T}(x^{*})J(x^{*})]d = d^{T}\nabla_{xx}L(x^{*},\lambda^{*})d + d^{T}M^{*}J^{T}(x^{*})J(x^{*})d > 0$$
(34)

since $M^* J^T(x^*) J(x^*)$ is positive definite and $d^T \nabla_{xx} L(x^*, \lambda^*) d > 0$ by assumption of the properties of problem 9.

Theorem 4. If x^* is the equilibrium point of the ODE system defined by equation 20a and the left hand side of equation 20b with $\beta = 1$ is constant, then x^* is an optimal solution of the unconstrained optimization problem defined by 13 and the NLP defined by 9.

¹⁹⁰ Proof. Since x^* is the equilibrium point of equation 20a, equation 29 is satisfied and remains to be ¹⁹¹ proven that $\nabla_M P_{PP}(x^*, M)$ is equal to zero to show that the first order necessary conditions for ¹⁹² problem 13 are satisfied. To do so, consider that:

$$\frac{dM}{dt} = \alpha M \|h(x^*)\|_2 = \alpha \sqrt{Mh(x^*)^T Mh(x^*)}$$
(35)

Then $\lim_{t \to \infty} Mh(x^*) = \lambda^*$ and since $\lim_{t \to \infty} \frac{dM}{dt} = k$, where k is a positive real number, it follows that $\frac{dM}{dt} = k$.

 $\lim_{t \to \infty} h(x^*) = 0. \text{ Thereby } \lim_{t \to \infty} \nabla_M P_{PP}(x; M) = \lim_{t \to \infty} \frac{1}{2} h(x^*)^T h(x^*) = 0. \text{ Note that } \lim_{t \to \infty} \frac{dM}{dt} = k \text{ is}$ required in order for x^* to be an equilibrium point of equation 20a and that the value of M used in this theorem is not an equilibrium value (M^*) since the growth rate $\frac{dM}{dt}$ is different than zero.

The satisfaction of the second order optimality conditions for this case is identical to Theorem 1, thus will be omitted. $\hfill \Box$

¹⁹⁹ **Theorem 5.** If x^* is the optimal solution to the unconstrained problem 23 for the penalty parameter ²⁰⁰ M^* and multiplier $\tilde{\lambda}^*$, then $(x^*, \tilde{\lambda}^*, M^*)$ is the equilibrium point for the system of ODEs defined by ²⁰¹ equations 25a to 25c.

²⁰² Proof. According to the assumptions of the theorem, first order necessary conditions hold. Then,

$$\nabla_x P_{PM}(x^*, \tilde{\lambda}^*, M^*) = -\frac{dx}{dt} = 0$$
(36)

Furthermore, if x^* is the optimal solution of problem 23, it is also a stationary point of the NLP defined by 9 (see theorem 17.5 in Nocedal and Wright [28]). Thereby, $h(x^*) = 0$ and $\frac{dM}{dt} = \frac{d\tilde{\lambda}}{dt} = 0$. Thus, $(x^*, \tilde{\lambda}^*, M^*)$ is the equilibrium point of the dynamic system formed by by equations 25a to 25c. By virtue of using the Lagrange multiplier method M goes to an equilibrium value M^* , as oppose to theorem 4. The same is true for theorem 6.

Theorem 6. If $(x^*, \tilde{\lambda}^*, M^*)$ is the equilibrium point of the ODE system defined by equations 25a to 25c, then x^* is an optimal solution of the unconstrained optimization problem defined by 23 and the NLP defined by 9.

Proof. Since $(x^*, \tilde{\lambda}^*, M^*)$ is the equilibrium point of the ODE system25a to 25c, then the first order necessary conditions of problem 23 are satisfied:

$$\nabla_x P_{PM}(x^*, \tilde{\lambda}^*, M^*) = \nabla_x f(x^*) + J(x^*)^T \left(\tilde{\lambda}^* + M h(x^*) \right) = 0$$
(37)

$$\nabla_M P_{PM}(x^*, \tilde{\lambda}^*, M^*) = \frac{1}{2} h(x^*)^T h(x^*) = 0$$
(38)

$$\nabla_{\lambda} P_{PM}(x^*, \tilde{\lambda}^*, M^*) = h(x^*) = 0$$
(39)

satisfaction of $\nabla_x P_{PM}(x^*, \tilde{\lambda}^*, M^*) = 0$ is clear since $\nabla_x P_{PM}(x^*, \tilde{\lambda}^*, M^*) = -\frac{dx^*}{dt} = 0$. Satisfaction of equations 38 and 39 are ensured by:

$$\frac{d\tilde{\lambda}^*}{dt} = M^* h(x^*) = 0 \tag{40}$$

$$\frac{dM}{dt} = \alpha M^{\beta} \left\| h(x) \right\|_{2} = 0 \tag{41}$$

The second order sufficient optimality condition requires that Hessian matrix of the penalty function defined by 23 to be positive definite:

$$\nabla_{xx} P_{PM}(x^*, \tilde{\lambda}^*, M^*) = \nabla_{xx} f(x^*) + \sum_{i=1}^{m} [\tilde{\lambda}_i^* + M^* h_i(x^*)] \nabla_{xx} h_i(x^*) + M^* J^T(x^*) J(x^*)$$
(42)

According to theorem 17.2 in Nocedal and Wright [28], for a sufficiently large value of M^* the vector of Lagrange multipliers that satisfies the KKT conditions for the NLP defined by 9 can be approximated by $\tilde{\lambda}^* + M^*h(x^*) \simeq \tilde{\lambda}^* \simeq \lambda^*$. Then, replacing the Hessian matrix of problem 9:

$$\nabla_{xx} P_{PM}(x^*; M^*) = \nabla_{xx} L(x^*, \lambda^*) + M^* J^T(x^*) J(x^*)$$

$$\tag{43}$$

which was shown to be positive definite in Theorem 1. Thereby, x^* is a strict local solution of 9.

Theorem 7. (asymptotic convergence of *PP* scheme) Assume that (x^*, λ^*) is a stationary point of equations 10. As usual, suppose LICQ holds at x^* and that for any vector d satisfying $J(x^*)d = 0$ holds that: $d^T \nabla_{xx} L(x^*, \lambda^*) d > 0$. Then, if x_0 is close enough to x^* , then $\lim_{t \to \infty} [x(t), M(t)h(t)] =$ (x^*, λ^*) . ²²⁵ *Proof.* For a system of ODEs represented by the following equation:

$$\frac{dx}{dt} = \phi(x) \tag{44}$$

the Poincaré-Lyapunov theorem (theorem 6.9 in Verhulst [29], page 191) states that if ϕ is continuously differentiable, it can be shown that x^* is an asymptotically stable point of 44 if all eigenvalues of the matrix $\nabla_x \phi(x^*)$ have negative real value. From equation 20a, $\phi(x^*, M^*) = -\nabla_x P_{PP}(x^*, M^*)$ and $\nabla_x \phi(x^*, M^*) = -\nabla_{xx} P_{PP}(x^*, M^*) = -[\nabla_{xx}^2 L(x^*, \lambda^*) + M^* J(x^*)^T J(x^*)]$. We now show that the matrix $\nabla_{xx}^2 L(x^*, \lambda^*) + M^* J(x^*)^T J(x^*)$ is positive definite. Let $\sigma_i \in \mathbb{R}$

We now show that the matrix $\nabla_{xx}^2 L(x^*, \lambda^*) + M^* J(x^*)^T J(x^*)$ is positive definite. Let $\sigma_i \in \mathbb{R}$ be an eigenvalue of $\nabla_{xx}^2 L(x^*, \lambda^*) + M^* J(x^*)^T J(x^*)$ and $z_i \in \mathbb{R}^n$ an eigenvector corresponding to σ_i . Then, σ_i and z_i satisfy

$$[\nabla_{xx}^2 L(x^*, \lambda^*) + M^* J(x^*)^T J(x^*)] z_i = \sigma_i z_i$$
(45)

²³³ Multiplying by z_i^T by the left side

$$z_i^T [\nabla_{xx}^2 L(x^*, \lambda^*) + M^* J(x^*)^T J(x^*)] z_i = \sigma_i ||z_i||_2^2$$
(46)

²³⁴ which yields,

$$\sigma_i = (z_i^T \nabla_{xx}^2 L(x^*, \lambda^*) z_i + z_i^T M^* J(x^*)^T J(x^*) z_i) / \|z_i\|_2^2$$
(47)

However, by the assumptions of the theorem $z_i^T \nabla_{xx}^2 L(x^*, \lambda^*) z_i$ is positive definite, M is positive for every t and $J(x^*)^T J(x^*)$ is also positive definite (provided LICQ holds). Moreover, $J(x^*)z_i =$ 0 if z_i satisfies the conditions of the theorem. Therefore, $\sigma_i > 0$ for all i = 1, 2, ..., n and all eigenvalues of $\nabla_x \phi(x^*)$ are strictly negative, and thus by the Poincaré-Lyapunov theorem it follows that $\lim_{t\to\infty} x(t) - x^* = 0$. Finally, since (x^*, λ^*) is a stationary point of 10, it follows from theorems 3 and 4 that:

$$\nabla_x f(x^*) + M^* J(x^*)^T h(x^*) = \nabla_x f(x^*) + J(x^*)^T \lambda^* = 0$$
(48)

and thereby $\lim_{t \to \infty} M(t)h(x) - \lambda^* = 0$

Theorem 8. (asymptotic convergence of PM scheme) Assume that (x^*, λ^*) is a stationary point of equations 10. Suppose that LICQ holds at x^* and that for any vector d satisfying $J(x^*)d = 0$, $d^T \nabla_{xx} L(x^*, \lambda^*)d > 0$, where L is the Lagrangian function for the NLP problem defined by equation 9. If x_0 is close enough to x^* , then $\lim_{t\to\infty} [x(t), \tilde{\lambda}(t) + M(t)h(t)] = (x^*, \lambda^*)$.

Proof. We start with equation 25a, defining $\phi_3(x, \tilde{\lambda}, M) = [\nabla_x f(x) + J(x)^T (\tilde{\lambda} + Mh(x))]$ and calculating $\nabla^2_{xx} \phi_3(x^*, \tilde{\lambda}^*, M^*)$:

$$\nabla_{xx}^2 \phi_3(x^*, \tilde{\lambda}^*, M^*) = -[\nabla_{xx}^2 f(x^*) + \sum_{i=1}^m [\tilde{\lambda}_i^* + M^* h_i(x^*)] \nabla_{xx} h_i(x^*) + M^* J^T(x^*) J(x^*)]$$
(49)

From the Theorem 5 under the assumptions of Theorem 8, $h(x^*) = 0$, $\tilde{\lambda}^* \simeq \lambda^*$ and,

$$\nabla_{xx}^2 \phi_3(x^*, \tilde{\lambda}^*, M^*) = -[\nabla_{xx} f(x^*) + \sum_{i=1}^m \tilde{\lambda}_i^* \nabla_{xx} h_i(x^*) + M^* J^T(x^*) J(x^*)]$$
(50)

$$\nabla_{xx}^2 \phi_3(x^*, \tilde{\lambda}^*, M^*) = -[\nabla_{xx}^2 L(x^*, \tilde{\lambda}^*) + M^* J^T(x^*) J(x^*)]$$
(51)

The right hand side of equation 51 was proven to be positive definite in Theorem 7. Recalling the definition of PM scheme (equation 25a):

$$\phi_3(x,\tilde{\lambda},M) = \nabla_x P_{PM}(x^*,\tilde{\lambda}^*,M^*) = -\frac{dx^*}{dt}$$
(52)

and using the Poincaré-Lyapunov theorem it follows that for equation 25a, $\lim_{t\to\infty} x(t) - x^* = 0$. Finally, since (x^*, λ^*) is a stationary point of 10, it follows from theorems 5 and 6 that:

$$\nabla_x f(x^*) + J(x^*)^T \left(\tilde{\lambda^*} + M^* h(x^*) \right) = \nabla_x f(x^*) + J(x^*)^T \lambda^* = 0$$
(53)

and thereby $\lim_{t \to \infty} \tilde{\lambda}(t) + M(t)h(t) - \lambda^* = 0.$

Moreover, the following theorem indicates that $P_{PP}(x, M)$ (equation 13) and $P_{PM}(x, M, \lambda)$ (equation 23) are strictly decreasing along a trajectory of x(t) that converges to x^* .

Theorem 9. Let [x(t), M(t)] be a solution trajectory of equations 25a and 20b. For a fixed $t_0 \ge 0$, if $\nabla_x P_{PP}(x(t), M(t)) \ne 0$ for all $t > t_0$, then $P_{PP}(x(t), M(t))$ is strictly decreasing with respect to $t > t_0$.

²⁵⁹ In an analogous way the following theorem is defined.

Theorem 10. Let $[x(t), M(t), \tilde{\lambda}(t)]$ be a solution trajectory of equations 25a to 25c. For a fixed $t_0 \geq 0$, if $\nabla_x P_{PM}(x(t), M(t), \tilde{\lambda}(t)) \neq 0$ for all $t > t_0$, then $P_{PM}(x(t), M(t), \tilde{\lambda}(t))$ is strictly decreasing with respect to $t > t_0$.

²⁶³ *Proof.* Here the proof of theorem 9 is presented. Proof of theorem 10 is analogous and is thus ²⁶⁴ omitted. From equation 20a, $\frac{dx}{dt} = -\phi_{PP}(x(t))$ and the trajectory of $P_{PP}(x(t))$ can be calculated ²⁶⁵ from

$$\frac{dP_{PP}(x(t))}{dt} = \frac{dP_{PP}(x(t))}{dx}\frac{dx}{dt}$$
(54)

$$\frac{dP_{PP}(x(t))}{dt} = -\|\phi_{PP}(x(t))\|_2^2 \tag{55}$$

Since $\phi_{PP}(x(t))$ is different than zero when $t > t_0$, then it can be concluded that $\frac{dP_{PP}(x(t))}{dt} < 0$. Thereby, $P_{PP}(x(t))$ is strictly decreasing with respect to $t > t_0$.

268 2.3. Comparison with previously reported GF formulations

As stated in the introductory section, the GF approach proposed in this work, unlike previously reported methods, is an infeasible path method. To analyze the reasons behind this behavior, consider the solution of the ODE system represented by equations 25a to 25c. After multiplying equation 25a by J(x) we obtain

$$J(x)\frac{dx}{dt} = \frac{dh(x)}{dt} = -J(x)[\nabla_x f(x) + J(x)^T \left(\tilde{\lambda} + Mh(x)\right)] \Rightarrow$$
$$\frac{dh(x)}{dt} + J(x)J(x)^T Mh(x) = -J(x)[\nabla_x f(x) + J(x)^T \tilde{\lambda}]$$

The last equation corresponds to a linear differential equation in h(x) with variable coefficients. Thus, defining $\nu(x(t)) = J(x(t))J(x(t))^T M$ and $\omega(x) = -J(x)[\nabla_x f(x(t)) + J(x(t))^T \tilde{\lambda}]$ and using an appropriate integration factor, the trajectory of h(x) can be implicitly expressed as

$$h(x(t)) = e^{-\int_0^t \nu(x(t))dt} \left[\int_0^t \omega(x(t)) e^{\int_0^t \nu(x(t))dt} dt + C \right]$$

Now, suppose that for some $t_0 > 0$, $h(x(t_0)) = 0$. Then,

$$e^{-\int_{0}^{t_{0}}\nu(x(t))dt} \left[\int_{0}^{t_{0}}\omega(x(t))e^{\int_{0}^{t_{0}}\nu(x(t))dt}dt + C\right] = 0 \Rightarrow$$
$$C = -\int_{0}^{t_{0}}\omega(x(t))e^{\int_{0}^{t_{0}}\nu(x(t))dt}dt$$

Therefore, for $t > t_0$:

$$h(x(t)) = e^{-\int_{t_0}^t \nu(x(t))dt} \left[\int_{t_0}^t \omega(x(t))e^{\int_{t_0}^t \nu(x(t))dt} + C \right]$$
(56)

and the only possibility for h(x(t)) to be zero for $t > t_0$ is that both C and $\omega(x(t))$ are zero for $t > t_0$, implying that also $\omega(x(t_0))$ needs to be equal to zero. These conditions can be satisfied if at time t_0 not only $h(x(t_0)) = 0$, but also $\nabla_x f(x(t)) + J(x(t))^T \tilde{\lambda} = 0$. Clearly, a point satisfying both conditions will also be the optimal solution of the problem. Thereby, unlike the formulation presented by Wang et al. [23], once the ODE system reaches a feasible point, it will generally not remain feasible.

Unlike previously proposed formulations based on the Gradient Flow approach, our GF scheme 284 does not requires the calculation of inverse matrices (such as $J(x)J^{T}(x)$, see equations 2 and 4). 285 Thereby, each integration step taken using the GF formulations presented in this work is less 286 computationally demanding compared, for example, to the GF formulation presented by Wang 287 et al. [23]. This is confirmed by the results presented in the next section. Moreover, although 288 the squared slacks used in equations 6 to 8 may cause the Jacobian to be linearly dependent, this 289 will occur at the point where the constraints are satisfied as equalities with zero slack, which in 290 turns only occurs at the solution since the proposed approach is an infeasible path method as 291 shown previously. Therefore, the use of squared slacks does not pose a problem for the GF schemes 292 proposed in this work. 293

294 2.4. Implementation

The original NLP problems (equation 5) were automatically converted to a system of differential equations (*PP*, equations 20a and 20b, and *PM*, equations 25a to 25c) using a code developed in Wolfram MathematicaTM 10.3 that takes full advantage of this Computer Algebra System (CAS).

The ODE equations are processed into a format suitable for Mathematica's built-in differential 298 equation solver NDSolve using an open package developed in MathematicaTM with flexible data 299 structures that also allows system structural analysis [30]. NDSolve options were used with default 300 values except for AccuracyGoal and PrecisionGoal which were increased to tighten constraint sat-301 isfaction when required. The option WhenEvent in NDSolve was used to reset the value of the 302 penalty parameter when it exceeds a large value (10^{12}) . Of course, this approach is only imple-303 mented for the formulation that incorporates multipliers (PM) since the formulation PP requires 304 a large penalty value to achieve convergence. For the case study Problem 6, MATLABTM ODE15s 305 was used to obtain the solution of the problem when the PM formulation was tested. MATLABTM 306 was used since it allows tailoring the execution of the integration. Specifically, for large problems 307 RAM memory usage was limited by using short integration steps, storing the solution at the final 308 integration time $(x(t_f))$ and reinitializing the integration using the stored values $(x_0 = x(t_f))$. 309

The ODE systems produced by PP and PM formulations are integrated until the merit function defined by equation 57 reaches a value below a prescribed tolerance equal to 10^{-6} , unless otherwise stated. This value was chosen to be similar to the default tolerances used by CONOPT and IPOPT, 10^{-7} and 10^{-6} , respectively.

$$Merit\,function = \|\nabla_x f(x) + J^T(x)\mu\|_2 + \|h(x)\|_2 \tag{57}$$

where μ was calculated as $\mu^* = M^* h(x^*)$ for PM formulations and as $\tilde{\lambda^*}$ for PP formulations. Finally, for the GF formulation *PP*, integration was stopped if the value of the penalty parameter *M* exceeds 10^{12} to avoid numerical problems (overflow).

The GF formulations proposed in this work are compared against the formulations reported 317 in Wang et al. [23] using two approaches: WM and WA, For problems with a small number of 318 variables (problems 1 and 2 in Section 3) a calculation procedure termed WM (Wang's Method) 319 was implemented in MathematicaTM. In this scheme, matrix $J(x)^T J(x)$ in equation 4 is symbolically 320 inverted and the resulting function is stored in MathematicaTM for its use by NDSolve. Therefore, 321 matrix $J(x)^T J(x)$ is not inverted at each integration step but the stored function is used to evaluate 322 it. Moreover, the CPU time reported for the WM scheme, as for other GF schemes, corresponds to 323 the CPU time consumed by the execution of NDSolve, thus for the WM scheme it does not include 324 the time required to calculate the inverse of $J(x)^T J(x)$ (symbolically). 325

As the number of variables increases (problems 3 to 7 in Section 3), the time required to symbolically invert the matrix $J(x)^T J(x)$ in WM approach rises to impractical levels. Therefore, in order to compare the GF approaches presented in this paper with the method proposed by Wang et al. [23] a new approach was required. This approach, termed WA (Wang's Algorithm) was also presented in Wang et al. [23] and is based in the discretization of equation 4 using the implicit backward Euler's scheme. According to this algorithmic scheme, the matrix $J(x)^T J(x)$ is not symbolically inverted but in each step $[J(x)^T J(x)]^{-1}$ is numerically calculated.

The GF schemes presented in this work and WM are compared in terms of CPU times, number of integration steps and function evaluations used by NDSolve to achieve the same merit function value. When WA and conventional NLP solvers (CONOPT, MINOS, SNOPT and IPOPT) are used, CPU times and number of iterations are compared.

337 3. Case studies

The new Gradient Flow schemes were numerically tested against five standard constrained nonlinear problems and two nonlinear problems derived from optimal control problems, all obtained from the open literature.

341 3.1. Constrained non-linear problems

342 3.1.1. Problem 1

Problem 1 corresponds to a constrained nonlinear optimization problem with a nonlinear objective function, linear constraints and variable bounds. Its optimal solution is $x = \{0, 1, 2, 0\}$ with an objective function value equal to -1.5. Problem 1 is a convex one with linear constraints; thereby, its difficulty is low and a global solution is expected, both for conventional NLP solvers and for the GF schemes presented in this work.

 $\min_{x} 1.5x_1^2 - x_1x_2 + 1.5x_2^2 + x_1 - 3x_2$

subject to:

$$-x_1 + 2x_2 + x_3 = 4 \tag{58}$$

$$x_1 + x_2 + x_4 = 1 \tag{59}$$

$$0 \le x_i \le 10; i = \{1, ..., 4\}$$

Problem 1 was solved using x = [2, 1, 3, 4] as initial values and with $\alpha = 10^6$ for scheme PM 348 with $\beta = 0$, and $\alpha = 10^3$ for scheme PM with $\beta = 1$ and for scheme PP. Table 1 shows a summary 349 of the numerical results when the integration was stopped after a merit function value equal or 350 lower than 10^{-6} was achieved for every GF formulation. Every solver tested, including this work's 351 and Wang's formulation [23], finds the optimal solution of the problem. CPU time and RAM 352 memory usage obtained using GF schemes are competitive with conventional solvers. Using an 353 identical termination criteria, the GF formulation proposed by Wang and coworkers with their 354 tuning parameter set to $\tau = 1$ [23] requires nearly 129 times more CPU time compared to the GF 355 formulations presented in this work. This difference increases as τ increases, with 23 CPU seconds 356 for $\tau = 1000$, due to the stiffness of the problem that forces the integrator to use small integration 357 steps. The quality of the solution, in terms of constraint satisfaction and objective function is 358 similar for the conventional solvers and the GF formulations. 359

or iormations	proposed i	II UIIIS WOLK	•						
	IPOPT	CONOPT	MINOS	SNOPT	$PP_{(\beta=0)}$	$PP_{(\beta=1)}$	$PM_{(\beta=0)}$	$PM_{(\beta=1)}$	WM $(\tau=1)$
CPU (s)	0.11	0.02	0.08	0.05	0.02	0.02	0.02	0.03	2.37
Memory (Mb)	3.0	2.0	3.0	3.0	1.1	1.0	1.8	2.0	11.2
Nfun	89	NA^a	4	2	1183	1384	914	1639	1172
Iterations	15	9	2	1	-	-	-	-	-
Integration steps	3 -	-	-	-	510	511	490	546	757
Obj	-1.50	-1.50	-1.50	-1.50	-1.50	-1.50	-1.50	-1.50	-1.50
	$7.1 \cdot 10^{-9}$	0	0	$2.2\cdot 10^{-16}$			$2.7\cdot 10^{-15}$	$3.4\cdot 10^{-12}$	$6.8\cdot 10^{-12}$
Merit function	NA^{b}	NA^b	NA^{b}	NA^b	$1.0 \cdot 10^{-6}$				

Table 1: Solution summary for Problem 1 including results obtained using commercial optimization solvers and the GF formulations proposed in this work.

 $^a\mathrm{Not}$ reported by GAMS CONOPT. $^b\mathrm{Default}$ values in GAMS were used

The effect of varying the value of α in the *PP*, *PM* and Wang's (increasing τ) formulations 360 [23] is shown in Figure 2, indicating that the satisfaction of the constraints is largely independent 361 of the value of parameters α and τ for this problem (where only linear constraints are present). 362 However, as the value of these parameters increases the system of differential equations becomes 363 stiff, requiring more integration steps (and computational time) to achieve the same norm of the 364 constraints. This is evident in Wang's formulation (WM) where the merit function value oscillates 365 for $\tau = 10^6$. Since in PP and PM formulations the value of the penalty parameter M is tied to 366 the violation of the constraints, it is not necessary to use large values of α , especially with the PM 367 formulation where multipliers are also used. 368

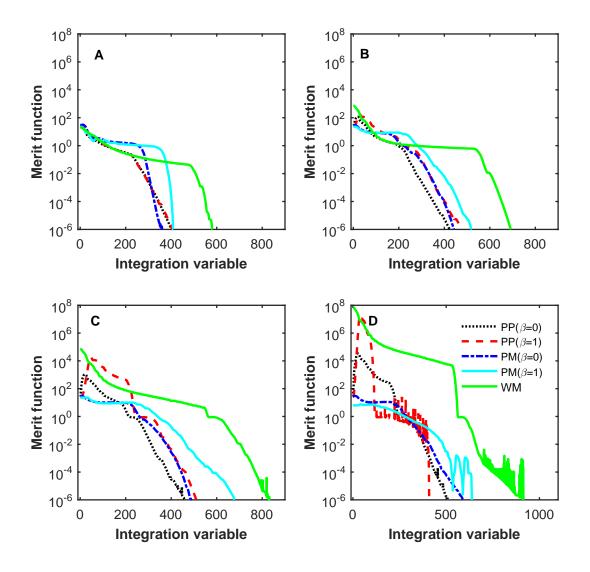


Figure 2: Norm of the constraints vector as the integration proceeds for Problem 1 and GF formulations using a penalty parameter scheme (*PP*) or a penalty and multipliers scheme (*PM*). For panels A, B, C and D the value of α is 10^{-2} , 1, 10^3 and 10^6 respectively.

369 3.1.2. Problem 2

This problem has been taken from Example 6.8 in Biegler [25], and it was designed to challenge Newton-based interior point methods when starting from an infeasible point. The problem has only one solution at $x^* = [1, 0, 0.5]$. $\min x_1$

subject to:

$$x_1^2 - x_2 - 1 = 0$$

 $x_1 - x_3 - 0.5 = 0$

$$-10 \le x_1 \le 10$$

$$0 \le x_i \le 10; \ i = \{2, 3\}$$

The infeasible initial point for this problem is x = [-2, 3, 1] and for the GF formulations an 373 α value of 10⁶, for PM with $\beta = 0$, and 10³ for the rest of the formulations was chosen. Table 2 374 shows the numerical results for this problem where the merit function satisfaction was set to 10^{-6} 375 for GF based formulations (including the one proposed by Wang et al. [23]). As expected, IPOPT 376 fails in finding the optimal solution, but CONOPT, MINOS and SNOPT provide the solution. 377 The solutions provided by GF schemes are optimal, show excellent constraint satisfaction and are 378 obtained using less RAM memory and CPU time compared to the conventional NLP solvers. On 379 the other hand, the GF formulation proposed by Wang et al. [23] reaches the limit of iterations 380 (10^5) without achieving the solution of the problem when $\tau = 1$, finds the solution for $\tau = 1000$ and 381 fails again for a larger value of this parameter. This result is indicative of the importance of the 382 selection of the τ value for NLP problems with nonlinear constraints, unlike the situation shown 383 for Problem 1, where τ values only affect the stiffness of the problem. The independence to the α 384 value shown by the algorithms presented in this paper is a consequence of the self-tuning behavior 385 of the penalty parameter in the GF schemes presented in this work. On the contrary, Wang et al. 386 [23] reports linear rate of convergence for small values of the penalty parameter and quadratic rate 387 of convergence for large values. Thereby, in their approach the value of the penalty parameter 388 needs to be tuned carefully for each problem to obtain a quadratic convergence and to avoid a stiff 389 problem that prevents finding a solution. 390

391 3.1.3. Problem 3

Problem 3 corresponds to the flowsheet optimization problem of the Williams-Otto process [31], 392 adapted from Biegler [25]. The process flowsheet is shown in Figure 3 where also the reaction 393 network for the synthesis of P (main product), E (by-product) and G (waste product) is presented. 394 Two feed streams with pure A and B components (streams F_A and F_B) are fed to a stirred tank 395 reactor whose operating temperature, T, is subject to optimization. The effluent stream is cooled 396 and sent to a centrifuge to separate G (in stream F_G). The clarified stream is fed to a column 397 separator to recover P where 90% of the product P is recovered in the column's top stream. This 398 stream is separated into purge (F_{purge}) and a recycled stream (F_R) that is recycled to the reactor. 399 The optimization problem is represented by equations 62 to 74. Variable bounds, initial values, 400 optimal values and the values obtained using formulation PM with $\beta = 0$ are shown in Table 3. 401

(61)

	IPOPT	IPOPT CONOPT MINOS	MINOS	SNOPT	PPian	PP(a-1)	PM	PM(a-1)	WM $(\tau = 1)$	SNOPT $PP_{\ell^{o}}$, $PP_{\ell^{o}}$, $PM_{\ell^{o}}$, $PM_{\ell^{o}}$, WM ($\tau = 1$) WM ($\tau = 10^3$) WM ($\tau = 10^7$)	$WM (\tau = 10^7)$
CPU (s)	0.39	0.18	0.28	0.23	0.02	(1=0)	0.02	0.018	65.3	0.60	0.2
Memory (Mb) 1.3	1.3	0.7	2.1	1.2	1.4	0.9	1.7	2.0	96.9	2.2	1.7
Nfun	89	NA	36	16	1239	873	936	873	225694	4131	2504
Iterations	13	7	7	°	,	ı	·	ı	ı	ı	ı
ntegration steps	ı	ı	ı	ı	754	507	554	623	10^{5}	1912	1162
Obj	INF^b	1.00	1.00	1.00	1.00	1.00		1.00	0.883	1.00	0.885
$\ y(x) \ $	1.5	0	$3.3\cdot 10^{-9}$	$2.4\cdot 10^{-9}$	$3.3 \cdot 10^{-9} \ 2.4 \cdot 10^{-9} \ 6.9 \cdot 10^{-7} \ 1.2 \cdot 10^{-7} \ 1.4 \cdot 10^{-9} \ 6.3 \cdot 10^{-12}$	$1.2\cdot 10^{-7}$	$1.4\cdot 10^{-9}$	$6.3\cdot 10^{-12}$	0.221	$9.2\cdot 10^{-10}$	0.221
Merit Function NA ^b	NA^{b}	NA^{b}	NA^{b}	NA^{b}	$NA^{b} = 1.0 \cdot 10^{-6} \ 1.0 \cdot 10^{-6} \ 1.0 \cdot 10^{-6} \ 1.0 \cdot 10^{-6}$	$1.0 \cdot 10^{-6}$	$1.0\cdot 10^{-6}$	$1.0 \cdot 10^{-6}$	0.377	$1.0\cdot 10^{-6}$	$3.9\cdot10^{10}$

2.
Problem
for
summary
Solution
2:
(1)

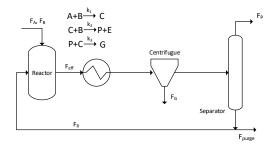


Figure 3: Flowsheet for Problem 3, the Williams-Otto process.

$$\min_{x} - \frac{2207 \cdot F_P + 50 \cdot F_{purge} - 168 \cdot F_A - 252 \cdot F_B - 2.22 \cdot F_{eff}^{sum} - 84 \cdot F_G - 60 \cdot V \cdot \rho}{6 \cdot \rho \cdot V}$$

subject to:

$$k_1 - 5.9755 \cdot 10^9 e^{-\frac{120}{T}} = 0 \tag{63}$$

$$k_2 - 2.5962 \cdot 10^{12} e^{-\frac{150}{T}} = 0 \tag{64}$$

$$k_3 - 9.6283 \cdot 10^{15} e^{-\frac{200}{T}} = 0 \tag{65}$$

$$F_{eff}^P - 0.1F_{eff}^E - F_P = 0 ag{66}$$

$$F_A + F_B - F_G - F_P - F_{purge} = 0$$

$$-k_1 F^A \epsilon_F F^B \epsilon_F V \rho \qquad F_{purge} F^A \epsilon_F$$

$$(67)$$

$$\frac{-\kappa_1 \Gamma_{eff} \Gamma_{eff} \nu \rho}{(F_{eff}^{sum})^2} - \frac{\Gamma_{purge} \Gamma_{eff}}{F_{eff}^{sum} - F_G - F_P} + F_A = 0$$
(68)

$$\frac{(-k_1 F_{eff}^A F_{eff}^B - k_2 F_{eff}^B F_{eff}^C) V \rho}{(F_{off}^{sum})^2} - \frac{F_{purge} F_{eff}^B}{F_{off}^{sum} - F_G - F_P} + F_B = 0$$
(69)

$$\frac{((2k_1F_{eff}^A - k_2F_{eff}^C)F_{eff}^B - k_3F_{eff}^CF_{eff}^P)V\rho}{(F_{eff}^{sum})^2} - \frac{F_{purge}F_{eff}^C}{F_{eff}^{sum} - F_G - F_P} = 0$$
(70)

$$\frac{2k_2 F_{eff}^B F_{eff}^C V \rho}{(F_{eff}^{sum})^2} - \frac{F_{purge} F_{eff}^E}{F_{eff}^{sum} - F_G - F_P} = 0$$
(71)

$$\frac{(k_2 F_{eff}^B - 0.5 k_3 F_{eff}^P) F_{eff}^C V \rho}{(F_{eff}^{sum})^2} - \frac{F_{purge}(F_{eff}^P - F_P)}{F_{eff}^{sum} - F_G - F_P} - F_P = 0$$
(72)

$$\frac{-1.5k_3F_{eff}^CF_{eff}^PV\rho}{(F_{eff}^{sum})^2} - F_G = 0$$
(73)

$$F_{eff}^{A} + F_{eff}^{B} + F_{eff}^{C} + F_{eff}^{E} + F_{eff}^{P} + F_{G} - F_{eff}^{sum} = 0$$
(74)

Table 3: Bounds, initial conditions, optimal solution and solution obtained using formulation $PM(\beta = 0)$ for Problem 3.

Variable	x^L	x^U	x_0	x^*	$x_{PM(\beta=0)}^*$	Variable	x^L	x^U	x_0	x^*	$x_{PM(\beta=0)}^*$
F_{eff}^{sum}	0.0	1000	52	366.369	370.523	F_{purge}	0.0	100	0.1	35.910	36.126
	0.0	100	10	46.907	43.675	V	0.03	0.1	0.06	0.03	0.03
	0.0	500	30	145.444	149.517	F_A	0	100	30	13.357	13.170
ejj	0.0	100	3	7.692	6.989		0.0	100	20	30.442	31.071
F_{eff}^E	0.0	1000	3	144.033	147.479	T	2	6.8	5.8	6.744	6.782
F_{eff}^P	0.0	100	5	19.115	19.511	k_1	0.0	200	6.2	111.7	123.6
F_P	0	4.763	0.5	4.712	4.763	k_2	0	1000	15.2	567.6	643.8
F_G	0.0	100	1	3.178	3.352	k_3	0	1500	10.2	1268.2	1500

Computational results obtained using conventional NLP solvers and the GF formulations pre-402 sented in this work are shown in Table 4. Commercial solver CONOPT achieves a feasible local 403 optimum, while the commercial solver MINOS reports the problem as infeasible. On the other 404 hand, GF formulations achieve the demanded value for the merit function ($< 10^{-6}$), although only 405 local minima solutions are attained. However, this is not unexpected since the GF formulations 406 do not incorporate provisions to achieve global optima. Using the algorithmic version of the for-407 mulation proposed by Wang and coworkers (WA), an algorithm based on the use of the implicit 408 Euler method, with a small integration step of 0.01, the algorithm reaches a merit function value of 409 $2.75 \cdot 10^{25}$ in four iterations producing numerical errors . The solution of this problem using Wang's 410 method (WM) implemented in NDSolve (as in Problems 1 and 2) was also not possible, since this 411 requires the calculation of a 44×44 inverse matrix with symbolic entries which proves extremely 412 time consuming. 413

Table 4: Solution summary for Problem 3.

		2.1.2						
	IPOPT	CONOPT	MINOS	SNOPT	$PP_{(\beta=0)}$	$PP_{(\beta=1)}$	$PM_{(\beta=0)}$	$PM_{(\beta=1)}$
CPU (s)	0.4	0.25	0.26	0.61	54.8	97.8	0.55	8.8
Memory (Mb)	0.4	1.8	2.1	1.6	31.5	39.6	21.2	22.7
Nfun	152	NA	385	4406	12692		936	7866
Iterations	33	71	42	233	-	-	-	-
Integration steps	-	-	-	-	4751	6490	1151	2507
Obj	-121.1	10.0	INF^{b}	-121.1	-118.9	-118.9	-120.2	-71.3
$\ h(x)\ $	$6.4\cdot 10^{-7}$	0	2.9	$1.1\cdot 10^{-12}$	$8\cdot 10^{-6}$	$1.9\cdot 10^{-6}$	$1.6\cdot 10^{-9}$	$3.7 \cdot 10^{-11}$
Merit Function	$\mathbf{N}\mathbf{A}^{c}$	$\mathbf{N}\mathbf{A}^{c}$	$\mathbf{N}\mathbf{A}^{c}$	\mathbf{NA}^{c}	$1.0\cdot 10^{-6}$	$1.0\cdot 10^{-6}$	$1.0\cdot 10^{-6}$	$1.0\cdot 10^{-6}$

^aNot reported by GAMS CONOPT. ^bSolver reports an infeasible solution. ^cDefault values in GAMS were used

Using the solution obtained with $PM(\beta = 1)$ scheme as a starting point, CONOPT declares the initial point as feasible after two iterations and finds the optimal solution in 18 iterations. Starting from the solution given by $PM(\beta = 0)$, CONOPT reports a feasible solution in two iterations and an optimal solution in 19 iterations. Although, MINOS fails in finding a feasible solution from the starting point given in Table 3, it reports an optimal solution after 12 iterations starting from the solution provided by $PM(\beta = 1)$ and in 10 iterations when starting from $PM(\beta = 0)$.

Therefore, the GF formulations presented in this paper result useful as an initialization method for this highly infeasible problem. It is important to stress that the commercial solvers only fail in the initial point reported in Table 3 and in other highly infeasible starting points, while they are able to solve the problem to optimality for most initial points.

424 3.1.4. Problem 4

This problem corresponds to the simplified alkylation process presented in Berna et al. [32], including 14 continuous variables, 1 linear constraint, 5 nonlinear constraints and a nonlinear objective function. Bounds for variables are shown in Table 5.

$$\min_{x} -6.3x_4x_7 + 5.04x_1 + 0.35x_2 + x_3 + 3.36x_5 \tag{75}$$

subject to:

$$x_4 - (x_1 + x_5)/1.22 = 0 \tag{76}$$

$$0.98x_3 - x_6(\frac{x_4x_9}{100} + x_3) = 0 \tag{77}$$

$$10x_2 + x_5 - x_1 x_8 = 0 \tag{78}$$

$$x_4 x_{11} - x_1 (1.12 + 0.13167 x_8 - 0.0067 x_8^2) = 0$$

$$0.8635 + \frac{(1.098x_8 - 0.038x_8^2)}{100} + 0.325(x_6 - 0.89) - x_7x_{12} = 0$$
(79)

$$35.82 - 22.2x_{10} - x_9 x_{13} = 0 \tag{80}$$

$$-1.33 + 3x_7 - x_{10}x_{14} = 0 \tag{81}$$

Table 5: Bo<u>unds and initial conditions for</u> Problem 4.

	x^L	x^U		x^L	x^U
x_1	$\begin{array}{c} 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \end{array}$	2	x_8	3.0	12.0
x_2	0.0	1.6	x_9	1.2	4.0
x_3	0.0	1.2	x_{10}	1.45	1.62
x_4	0.0	5	$ x_{11} $	0.99	1.01
x_5	0.0	2	x_{12}	0.99	1.01
x_6	$\begin{array}{c} 0.85\\ 0.90 \end{array}$	0.93	x_{13}	0.90	1.11
x_7	0.90	0.95	x_{14}	0.99	1.01

Problem 4 can be solved to optimality by every conventional NLP solver, and also by all the GF formulations proposed in this work (see Table 6) with demanded values of the merit function below 10⁻⁵ and 10⁻⁶ for PP and PM formulations. CPU time are smaller and RAM memory usage are generally larger for GF formulations, except when compared to IPOPT.. As shown in Table 6, the algorithm proposed by Wang and coworkers fails in achieving a solution, even for a small value of the integration step.

		Tai	ble 0. Soluti	on summa	I Y 101 F 100	nem 4.			
	IPOPT	CONOPT	MINOS	SNOPT	$PP_{(\beta=0)}$	$PP_{(\beta=1)}$	$PM_{(\beta=0)}$	$PP_{(\beta=1)}$	WA $(h = 10)$
CPU (s)	0.3	0.29	0.29	0.25	0.45	0.13	0.15	0.25	1612.08
Memory (Mb)	22.9	1.7	2.1	1.6	6.5	5.6	13.5	11.5	0.2
Nfun	79	NA^{a}	316	192		1981	1955	2785	-
Iterations	15	19	13	22	-	-	-	-	10^{4}
Integration steps	-	-	-	-	1081	835	1109	1070	-
Obj	-1.765	-1.765	-1.765	-1.765	-1.765	-1.765	-1.765	-1.765	$-9.7\cdot10^{6}$
$\ h(x)\ $	$8.3\cdot 10^{-9}$	$4.7\cdot 10^{-9}$	$2.5\cdot 10^{-11}$	$4.3\cdot 10^{-8}$	$6.3\cdot 10^{-5}$	$8.2\cdot 10^{-7}$	$1.9\cdot 10^{-9}$	$6.4\cdot 10^{-12}$	$5.6\cdot 10^6$
Merit Function	NA^{b}	NA^{b}	NA^{b}	NA^{b}	$1.0\cdot 10^{-5}$	$1.0 \cdot 10^{-5}$	$1.0\cdot 10^{-6}$	$1.0 \cdot 10^{-6}$	$1.3 \cdot 10^7$

Table 6: Solution summary for Problem 4.

^aNot reported by GAMS CONOPT. ^bDefault values in GAMS were used

434 3.1.5. Problem 5

Problem 5 was adapted from Wang et al. [23], where only seven variables were considered. In this work, the problem was modified to accommodate an arbitrary number of variables (n_v) while maintaining its qualification as a constrained concave programming problem. The problem has multiple local minima and a global minimum of -1.0.

$$\min_{x} - \sum_{i=1}^{n_{v}} x_{i}^{2} \tag{82}$$

subject to:

$$\sum_{i=1}^{n_v} x_i - 1 + slack = 0 \tag{83}$$

$$0 \le x_1 \le 0.8 \tag{84}$$

$$slack \ge 0$$
 (85)

$$0 \le x_i \le 1; i = \{2, ..., n_v\}$$

For all solvers, the initial point was taken as $x_i = 0.5$, $i = \{1, ..., n_v\}$. Table 7 shows the solution 439 of the problem for 50 to 600 variables using conventional NLP solvers and the GF formulations 440 introduced in this work when integration was stopped after the merit function achieves values lower 441 than 10^{-6} . Using the algorithm presented by Wang et al. [23], only problems with 50 and 100 442 variables were solved in less than 3600 CPU seconds. Every conventional solver reports a feasible 443 solution, however, all fail to find the global minimum as they are all local solvers. Although the GF 444 formulations presented in this work find the globally optimal solution of this problem with multiple 445 local minima, there is no theoretical reason to support this behavior. Moreover, using a different 446 starting point the commercial solvers also achieve the global solution for this problem. 447

The CPU time required by GF formulations is competitive with the commercial solvers for 50 and 100 variables but for a large number of variables, commercial solvers find a local optima for the problem using less CPU time and RAM memory. Considering that the commercial solvers represent the state of the art, rely on extensive preprocessing of the problem to achieve an efficient solution ⁴⁵² and are coded on a faster platform, this is an expected result.

We point the reader's attention to the fact that the GF formulation proposed in this work are implemented using Mathematica NDSolve, a general purpose algebraic solver. This explains the increase in RAM memory usage, as NDSolve stores the solution of the problem as polynomial splines. Clearly, an algorithmic implementation of the GF formulations, where the trajectories are not stored, will consume less RAM memory. Despite the difference in performance as the number of variables increase, the GF formulations presented in this work achieve optimal solutions with sharp constraint satisfaction.

⁴⁶⁰ Ås shown in Table 7, the algorithmic implementation of the GF formulation proposed by Wang ⁴⁶¹ and coworkers also achieves an optimal solution, however the CPU times required are several hun-⁴⁶² dred times larger.

					Ta	ble	7:	Sol	utic	on s	um	ma	ry f	for	Pro	ble	m 5	i .					
$n_v = 600$	0.19	2.8	599	0.0000	0					76.0	432.3	3567	1780	-1.0000	$1.5\cdot 10^{-9}$	$1.0\cdot 10^{-6}$				ı			
$n_v = 300 \ n_v = 600$	0.18	2.3	299	0.0000	0					21.6	225.4	3647	1814	-1.0000	$1.8\cdot 10^{-9}$	$6.9\cdot 10^{-7}$				ı			
$n_v = 100$	0.17	2.2	66	0.0000	0					2.11	64.4	2785	1501	-1.0000	$1.6\cdot 10^{-8}\ 8.7\cdot 10^{-9}\ 1.8\cdot 10^{-9}\ 1.5\cdot 10^{-9}$	$1.0\cdot 10^{-6}\ 1.0\cdot 10^{-6}\ 6.9\cdot 10^{-7}\ 1.0\cdot 10^{-6}$	1594.4	1.0	ı	63	-1.000	$3.1 \cdot 10^{-7}$ $4.6 \cdot 10^{-7}$	$1.0 \cdot 10^{-6}$ $1.0 \cdot 10^{-6}$
$n_{v} = 50$	0.18	0.8	49	0.0000	0					0.72	31.6	2682	1379	-1.0000	$1.6\cdot 10^{-8}$	$1.0\cdot 10^{-6}$	155.9	0.5	ı	48	0.0000	$3.1\cdot 10^{-7}$	$1.0\cdot 10^{-6}$
	MINOS, $CPU(s)$	Memory (Mb)	Nfun	Obj	$\ y(x) \ $					$PP(\beta = 1), CPU(s)$	Memory (Mb)	Nfun	Iterations	Obj	$\ y(x) \ $	Merit Function	WA(h = 1000), CPU(s)	Memory (Mb)	Nfun	Iterations	Obj	$\ y(x) \ $	Merit Function
$n_v = 600$	2.3	30.3	74	$-1.7 \cdot 10^{-3}$	$5.4\cdot 10^{-9}$	0.31	3.4	-0.5000	0	49.5	327.4	2989	1305	-1.0000	$3.5\cdot 10^{-7}$	$1.0\cdot 10^{-6}$	93.4	577.01	2714	1364	-1.0000	$2.7\cdot 10^{-13}$	$1.0\cdot 10^{-6}$
$n_v = 300$	0.34	13.1	69	$-3.4\cdot10^{-3}$	$2.1\cdot 10^{-10}$	0.27	2.4	-0.5000	0	18.8	194.3	3809	1571	-1.0000	$3.1\cdot 10^{-8}$	$1.0\cdot 10^{-6}$	19.2	287.1	2719	1347	-1.0000	$5.1\cdot10^{-12}$	$1.0 \cdot 10^{-6}$
$n_v = 100$	0.33	30.4	62	-0.6404	$3.0\cdot10^{-11}$	0.29	1.9	-0.5000	0	1.8	52.3	2965	1266	-1.0000	$9.3\cdot 10^{-8}$	$1.0\cdot 10^{-6}$	1.8	78.5	2189	1077	-1.0000	$2.7\cdot 10^{-12}$	$1.0\cdot 10^{-6}$
$n_v = 50$	0.32	30.9	104	-0.6408	$2.6\cdot 10^{-10}$	0.24	1.8	-0.5000	0	0.55	21.1	2489	1058	-1.0000	$1.8\cdot 10^{-6}$	$1.0\cdot 10^{-6}$	0.54	34.5	1829	204	-1.0000	$3.8\cdot 10^{-12}$	$1.0\cdot 10^{-6}$
	IPOPT, $CPU(s)$	Memory (Mb)	Nfun	Obj	$\ y(x) \ $	CONOPT, CPU(s)	Memory (Mb)	Obj	$\ v(x)\ $	$PP(\beta = 0), CPU(s)$	Memory (Mb)	Nfun	Iterations	Obj	$\ v(x) \ $	Merit Function	$PM(\beta = 0), CPU(s)$	Memory (Mb)	Nfun	Iterations	Obj	$\ v(x) \ $	Merit Function

Table 7: Solution summary for Problem 5 . $\circ \circ$

Figure 4 shows the value of the constraint as integration proceeds for increasing values of α 463

(panels A to C) and the values of the objective function (panel D). As it can be seen, constraint 464 satisfaction when the maximum number of integration steps was limit to 2000 units depends on 465 the value of α , at least for formulations $PP(\beta = 0)$ and $PP(\beta = 1)$. For example, Panel A shows 466 that formulation $PP(\beta = 0)$ achieves a constraint satisfaction in the order of 10^{-2} at the end of 467 integration, requiring nearly 2000 integration steeps. However, this does not mean that formulation 468 $PP(\beta = 0)$ is unable to produce sharp constraint satisfaction. In fact, as shown in Table 7, a 469 constraint satisfaction of $9.3 \cdot 10^{-8}$ can be attained using a longer integration time (controlled by 470 demanding a value of the merit function smaller than a certain threshold, 10^{-6} for this problem) 471

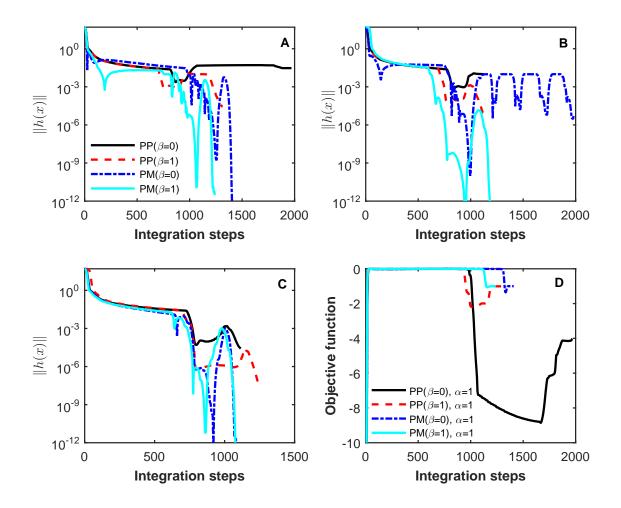


Figure 4: Constraint satisfaction for Problem 5 (with $n_v = 100$) for $\alpha = 1$ (panel A), $\alpha = 10$ (panel B) and $\alpha = 1000$ (panel C). Panel D shows the value of the objective function for the Gradient Flow formulations with $\alpha = 1$. Integration time was set to 500 units.

As shown in Figure 4, once a feasible point is achieved (for example $||h(x)|| = 10^{-12}$ in panel C for formulation PM) the trajectory of the ODE system does not remain feasible, unlike the Gradient ⁴⁷⁴ Flow approach presented by Wang et al. [23] for the solution of NLP problems (see Section 2.3).

475 3.2. Optimal control problems as NLP problems

476 3.2.1. Problem 6

Problem 6 corresponds to the determination of the optimal acceleration along time such that the total travel time is minimized for a car, subject to a path constraint (speed should be less than 10 units), final point constraint (distance should be equal to 300 units ($y_2 = 300$)), final velocity should be zero ($y_1 = 0$) and bounds for acceleration. The optimal control problem is defined by

$$\min_{t_f,u} t_f \tag{86}$$

subject to:

$$y'_1(t) = u(t)$$
 (87)

$$y_2'(t) = y_1(t)$$
 (88)
 $y_2(t) \le 10$ (89)

$$y_1(t) \le 10 \tag{89} y_1(t_f) = 0 \tag{90}$$

$$y_2(t_f) = 300$$
 (91)

$$-2 \le u(t) \le 1$$

$$0 \le t_f \le 50 \tag{92}$$

⁴⁸¹ Using the Euler backward difference formula, the optimal control problem can be written as a ⁴⁸² finite dimensional NLP problem:

$$\min_{t_f,u} t_f \tag{93}$$

subject to:

$$y_{1,i} - y_{1,i-1} - \left(\frac{t_f}{n_h}\right) u_i = 0 \tag{94}$$

$$y_{2,i} - y_{2,i-1} - \left(\frac{t_f}{n_h}\right) y_{1,i} = 0 \tag{95}$$

$$y_{1,0} = 0 (96) y_{2,0} = 0$$

$$y_{1,n_h} = 0$$
 (97)

$$y_{2,n_h} = 300$$
 (98)

$$0 \le y_i \le 10; \ i = \{0, 1, ..., n_h\}$$
(99)

$$-2 \le u_i \le 1; \ i = \{0, 1, ..., n_h\}$$

$$(100)$$

where n_h is the number of integration elements. For *PP* it was necessary to reduce the α value to 10, while the value of this parameter for formulation PP was maintained in 10³.

Results are presented in Table 8 showing that every conventional NLP solver achieved an optimal solution for 5 and 20 intervals. Formulation PP ($\beta = 1$) was unable to achieve the optimal solution as the integration terminated when the penalty parameter value exceeded 10^{12} . On the other hand, formulation *PM* achieves feasibility and locally optimal solutions for 5 and 20 intervals in a fraction of the time required by the *PP* formulation. Still, the time consumed by the *PM* formulation to attain the solution of the problem is larger compared to the one required by CONOPT or other commercial NLP solvers. Using the algorithm presented by Wang et al. [23], WA with $h = 10^{-5}$ and $\tau = 1000$ for $n_h = 5$, no solution was attained after 1000 iterations and 7492 CPU seconds, reaching a merit function value equal to $7.9 \cdot 10^5$ with very slow progress towards constraints satisfaction.

	$n_h = 5$	$n_h = 20$		$n_h = 5$	$n_h = 20$
IPOPT, CPU (s)	0.8	1.1	CONOPT	0.7	0.7
Memory (Mb)	22.3	29.2	Memory (Mb)	0.6	2.1
Iters	14	16	Iters	20	63
Obj	39.56	37.62	Obj	39.64	37.62
$\ h(x)\ $	$1.0 \cdot 10^{-11}$	$3.1\cdot10^{-11}$	$\ h(x)\ $	0	0.00
MINOS, CPU (s)	0.9	0.46	SNOPT	0.5	0.7
Memory (Mb)	2.0	2.1	Memory (Mb)	0.7	0.9
Iters	51	88	Iters	23	109
Obj	39.56	37.62	Obj	36.56	37.62
$\ h(x)\ $	$2.8\cdot10^{-14}$	$3.2\cdot 10^{-7}$	$\ h(x)\ $	$1.6\cdot 10^{-7}$	$9.1\cdot 10^{-8}$
$PP(\beta = 0), \text{CPU (s)}$	0.77	272.6	$PP(\beta = 1)$	2.06	310.1
Memory (Mb)	72.3	344.4	Memory (Mb)	10.4	425.4
Nfun	17648	17205	Nfun	14572	26905
Integration steps	8935	10195	Integration steps	7340	13551
Obj	49.97		Obj	50	50.1
$\ h(x)\ $	$2.1\cdot 10^{-6}$	$3.5\cdot 10^{-7}$	$\ h(x)\ $	$2.4\cdot 10^{-8}$	0.60
Merit function	$1\cdot 10^{-6}$	$9.93\cdot 10^{-7}$	Merit Function	$8.9\cdot 10^{-7}$	168.1^{*}
$PM(\beta = 0), CPU (s)$	0.64	12.1	$PM(\beta = 1, \alpha = 1)$	1.6	15.1
Memory (Mb)	28.0	39.2	Memory (Mb)	76.4	13.0
Nfun	2897	16760	Nfun	8527	20683
Integration steps	1662	7970	Integration steps	4637	10431
Obj	50	37.62	Obj	50	37.62
$\ h(x)\ $	$2.1\cdot 10^{-6}$	$9.3\cdot 10^{-6}$	$\ h(x)\ $	$3.4\cdot10^{-13}$	
Merit function	$1.0\cdot 10^{-6}$	$1.0\cdot 10^{-5}$	Merit function	$1.0\cdot 10^{-6}$	$1.0\cdot 10^{-5}$
	* Integra	tion stops	when $M > 10^{12}$.		

Table 8: Solution summary for Problem 6. The set of ODEs generated by the schemes $PM(\beta = 1)$ and $PM(\beta = 0)$ were solved in MATLAB.

^{*} Integration stops when $M > 10^{12}$.

In Table 8 the system of differential equations $PM(\beta = 1)$ reaches the demanded value for the 494 merit function after 10431 integration steps reporting a local optimum. However, if the integration 495 is allowed to continue and smaller values of α are used, a feasible point is achieved after 2500 496 integration steps, and after several integration steps where the value of the objective function 497 remains fixed, it starts decreasing towards the optimal value. When the norm of the constraint 498 vector is plotted against the number of integration steps for $PM(\beta = 1)$ and $n_h = 20$, a behavior 499 similar to Problem 5 is observed (compare figures 4 and 5). Interestingly, smaller α values require 500 fewer integration steps to move from the plateau value of 50 units of time to the globally optimum 501 value reported by the commercial solvers. This can be explained by the stiffness caused by large α 502

503 values.

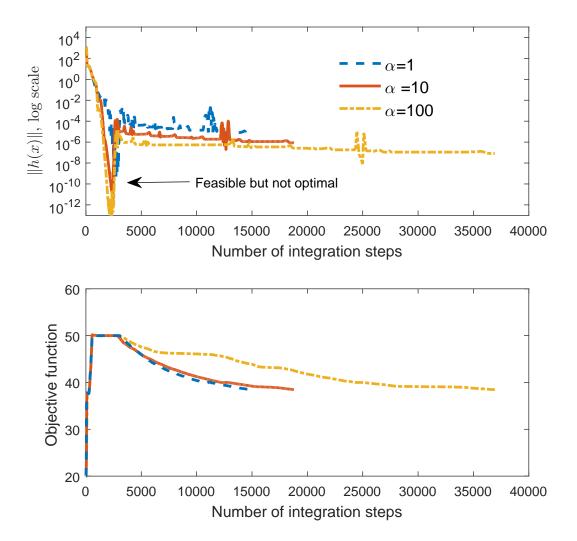


Figure 5: Norm of the constraints vector in Problem 6 (with $n_h = 20$) as the integration proceeds (as number of integration steps) for increasing values of α . Integration time was set to 500 units.

504 3.2.2. Problem 7

Problem 7 corresponds to the maximization of the harvested amount of a biological resource, provided this resource grows as time passes. The growth rate is assumed to be proportional to the amount of the resource present at a given time. The optimal control problem is defined by

$$\min_{u} - \int_{0}^{t_f} \sqrt{u(t)} dt \tag{101}$$

subject to:

$$x'(t) = (\varphi - 1)x(t) - u(t)$$
(102)
$$x(0) = x(t_{*})$$
(103)

$$\begin{aligned} x(0) &= x(t_f) \tag{103}\\ x(t) &> 0 \tag{104} \end{aligned}$$

$$\begin{aligned} u(t) \ge 0 \tag{104} \\ u(105) \end{aligned}$$

$$0 \le t \le t_f \tag{106}$$

$$x(0) = x_0 \tag{107}$$

where
$$x(t)$$
 is a scalar function representing the amount of a biological resource, $(\varphi - 1) > 0$ is the
growth rate and $u(t)$ is the amount of the resource extracted at a given time. Moreover, the problem
demands that the amount of resource at the final time be equal to its initial amount. By using a
discretization in time of 1 unit, the problem can be rewritten as the following finite-dimensional
NLP problem [33]:

$$\min_{u_k} - \sum_{k=0}^{N-1} \sqrt{u_k}$$
(108)

subject to:

$$x_{k+1} = \varphi x_k - u_k \tag{109}$$

$$x_0 = x_N \tag{110}$$

$$0 \le u_k \le u_k^U; \ k = \{0, 1, ..., N\}$$
(111)

$$0 \le x_k \le x_k^U; \ k = \{0, 1, ..., N\}$$
(112)

The optimal value of the NLP problem represented by equations 108 to 112 is known analytically to be [33] given by

512

$$J^* = \sqrt{\frac{x_{ini}(\varphi^N - 1)^2}{\varphi^{N-1}(\varphi - 1)}}$$
(113)

513

The optimal control problem was solved for N with values of 10, 20, 30 and 50, $x_{ini} = 1$ and $\varphi = 1.1$. For $PP(\beta = 0)$ and $PM(\beta = 0)$, α was set as 10^4 while for $PM(\beta = 0)$ and $PM(\beta = 1)$, α was set as 10^2 . The systems of ODEs were integrated until the merit function value was less than 10^{-6} . The upper bounds for the amount of resource (x_k^U) and control variable (u_k^U) need to be increased as N increases. Hence, for values of $N = \{10, 30, 50\}$, $x_k^U = \{2, 5, 30\}$ and $u_k^U = \{1, 2, 10\}$ for every k.

Table 9 shows the computational results for this case study. Solvers MINOS and SNOPT report infeasible solutions for every value of N, despite the problems having a moderate number of variables $(2 \cdot N)$. CONOPT and IPOPT achieve the optimal solution values, -3.282 for N = 10, -13.060for N = 30 and -35.629 for N = 50. GF formulations achieve between 98.2% to 100% of the

optimal solution with excellent constraint satisfaction, especially for $PM(\beta = 0)$ and $PM(\beta = 1)$ 524 formulations. 525

	Table 9: S	olution sum	mary for Problem 7.		
	N = 10	N = 30		N = 10	N = 30
IPOPT, CPU(s)	0.32	0.30	CONOPT	0.23	0.25
Memory (Mb)	30.8	31.0	Memory (Mb)	1.8	1.9
Nfun	54	74	Iters	NA^{a}	NA
Obj	-3.282	-13.060	Obj	-3.282	-13.060
$\ h(x)\ $	$1.0\cdot10^{-11}$	$1.0\cdot 10^{-11}$	$\ h(x)\ $	$2.2\cdot 10^{-16}$	$3.6\cdot 10^{-8}$
$PP(\beta = 0), CPU(s)$	0.23	0.81	$PP(\beta = 1)$	0.14	0.66
Memory (Mb)	9.0	28.7	Memory (Mb)	8.2	34.9
Nfun	1680	2398	Nfun	1801	3060
Integration steps	853	1075	Integration steps	902	1360
Obj	-3.281	-13.061	Obj	-3.282	-13.061
$\ h(x)\ $	$9.4\cdot 10^{-7}$	$4.8\cdot 10^{-7}$	$\ h(x)\ $	$8.4 \cdot 10^{-7}$	$3.4\cdot10^{-7}$
Merit function	$1.0\cdot 10^{-6}$	$1.0\cdot 10^{-6}$	Merit function	$1.0\cdot 10^{-6}$	$1.0\cdot 10^{-6}$
$PM(\beta = 0), CPU(s)$	0.078	0.55	$PM(\beta = 1)$	0.12	0.68
Memory (Mb)	7.3	45.2	Memory (Mb)	12.9	58.3
Nfun	678	2083	Nfun	1293	2701
Integration steps	363	956	Integration steps	684	1207
Obj	-3.282	-13.060	Obj	-3.282	-13.060
$\ h(x)\ $	$6.2\cdot 10^{-9}$	$1.8\cdot 10^{-9}$	$\ h(x)\ $	$8.3\cdot 10^{-10}$	$1.7\cdot 10^{-10}$
Merit function	$1.0 \cdot 10^{-6}$	$1.0 \cdot 10^{-6}$	Merit function	$1.0\cdot 10^{-6}$	$1.0 \cdot 10^{-6}$

^aNot reported by GAMS CONOPT. ^bSolver reports an infeasible solution

An attempt to use the algorithm presented by Wang et al. [23], WA, to solve this problem with 526 N = 10 was undertaken. After a systematic search of τ and integration step values (h), the values 527 producing the best solution were $\tau = 1000$ and h = 0.001. The algorithm achieves a merit function 528 value of 0.66, objective function value of -5.2 and norm of the constraint vector equal to $1.2 \cdot 10^{-4}$ 529 in 1797 CPU seconds and 3000 iterations. 530

Figure 6 shows the progress of merit function and constraint satisfaction as well as the obtained 531 solution for the amount of resource and the control variable. Panels C and D show the trajectory 532 of the state and control variables. Trajectories produced using formulation $PM(\beta = 1)$ are slightly 533 different compared to those obtained using CONOPT or IPOPT, which explains the near optimal 534 objective function value calculated for $PM(\beta = 1)$ and N = 30 in Table 9. 535

Summarizing our results, the case studies show that the GF approach presented in this work 536 is faster, in terms of CPU time, compared to a previously reported GF formulation [23], which is 537 representative of a family of formulations used to transform NLP problems to a system of ODEs, 538 and requires fewer function evaluations and integration steps (WM implementation versus PP and 539 PM schemes). This family of formulations makes extensive use of the inverse matrix $[J(x)J^{T}(x)]^{-1}$, 540 introducing a computational overhead that it is eliminated in our formulations thanks to a simpler 541 approach to penalizing functions. Moreover, in the GF formulation reported by Wang et al. [23] a 542 penalty parameter τ has to be tuned for each problem, and for the algorithmic version of the GF 543 scheme reported by the authors only τ has to be tuned for each problem but also the integration 544 step of the backward Euler's scheme (h). Results obtained for Problem 2 show that in WM method, 545

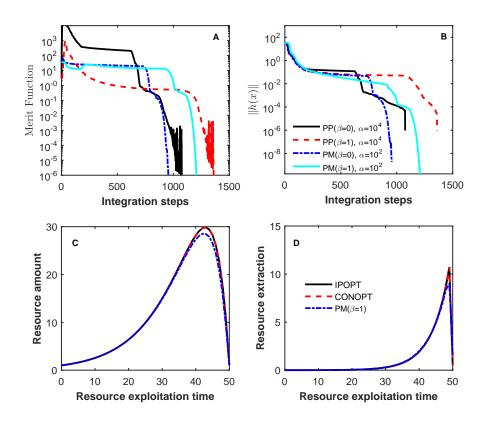


Figure 6: Progress towards the solution as integration proceeds for Problem 7.Merit function for Problem 7 (with N = 30, panel A), norm of the constraints vector value (panel B), Panels C and D show the trajectory of the state variables (x(t), representing the amount of resource) and control variable (u(t), amount of resource) extracted at a given time) for N = 30.

selection of τ plays a critical role, while low values produce slow converge, too high values produce 546 a very stiff system of ODEs. For problem 4, 5, 6 and 7 WA scheme fails to reach a solution for 547 the problem. For these problems, the reported solutions correspond to the best solution achieved 548 after testing several combination of h and τ values, being the best ones obtained for low values of 549 h and high values of τ , however, the low h values produce a slow convergence, a result that can be 550 expected from the analysis presented by Wang et al. [23], who claims a linear rate of convergence 551 for small h values. The GF schemes presented in this work successfully avoid this issues by using a 552 self-tuning approach of the penalty parameters. 553

In our approach, the solution of the problem is approached from the exterior of the feasible set at every step, except when the optimal solution is attained. This in turns allow the use of conventional slacks to bound variables. Compared to the state-of-the-art NLP solvers used in this work as benchmark, the proposed GF formulation remains competitive for problems with less than one hundred variables, being especially useful for problems with highly non-linear constraints, either as a solution or an initialization method.

560 4. Conclusions

This work presents a novel Gradient Flow formulation for the solution of nonlinear optimization 561 problems with equality and inequality constraints. The proposed schemes were shown theoretically 562 to converge (asymptotically) to a local minimum of the original problem under conventional as-563 sumptions on the objective function and constraints. These formulations were theoretically and 564 numerically compared to other reported Gradient Flow formulations for the solution of nonlin-565 ear constrained optimization problems, showing that the proposed formulations outperforms the 566 reference method in terms of computational time and the size of problems that can be solved. More-567 over, the self-tuning nature of the proposed approach reduces the numerical problems introduced 568 by increasing the value of the penalty parameter. 569

Numerical experiments using a set of seven specially selected problems, ranging from 3 to 600 variables, show that the proposed schemes are robust and converge to feasible points and local minima, irrespective of the choice of the value of parameter α used in the formulations, due to the self-tuning properties of the penalty parameters introduced in this work.

Moreover, results suggest that, for the set of problems analyzed, the GF formulations were able to find the optimal solution to problems where conventional NLP solvers fail. Primarily this is due to the fact that constraints are not linearized at intermediate iterations, with the solution being approached from the exterior of the feasible set.

As shown by the computational experiments, the GF formulations presented in this work achieve feasibility for problems with difficult nonlinear constraints. The combined multiplier and penalty approach in formulation PM provides solutions in shorter times and with sharp constraint satisfaction for almost every problem compared to the other GF formulations presented and compared in this work.

Most likely, if a customized integrator were used to solve the ODE systems produced by this and the other GF formulations, solution times and the number of function evaluations and iterations will be significantly reduced.

Future work includes an algorithmic implementation of the formulations presented in this work using a customized integrator and the exploration of the possibility of exploiting the special structure of the linear constraints and bounds as to reduce the ODE system size, for example incorporating the ideas presented in Schropp and Singer [24] and Shikhman and Stein [34], that might allow the decoupling of the original variables from the variables used to enforce bounds, thereby reducing the size of the dynamical system.

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