epsilon: A tool to find a canonical basis of master integrals

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Abstract

In 2013, Henn proposed a special basis for a certain class of master integrals, which are expressible in terms of iterated integrals. In this basis, the master integrals obey a differential equation, where the right hand side is proportional to ϵ in $d = 4 - 2\epsilon$ space-time dimensions. An algorithmic approach to find such a basis was found by Lee. We present the tool epsilon, an efficient implementation of Lee's algorithm based on the Fermat computer algebra system as computational backend.

Keywords: Feynman integral; canonical basis; differential equation; Fuchsian form

Program Summary

Program Title: epsilon

Licensing provisions: GPLv3

Programming language: C++

Nature of problem: For a certain class of master integrals, a canonical basis can be found in which they fulfill a differential equation with the right hand side proportional to ϵ . In such a basis the solution of the master integrals in an ϵ -expansion becomes trivial. Unfortunately, the problem of finding a canonical basis is challenging.

Solution method: Algorithm by Lee [1]

Restrictions: The normalization step of Lee's algorithm will fail if the eigenvalues of the matrix residues are not of the form $a + b\epsilon$ with $a, b \in \mathbb{Z}$. Multi-scale problems are not supported.

[1] R. N. Lee, JHEP 1504 (2015) 108 [arXiv:1411.0911 [hep-ph]].

1 Introduction

The perturbative treatment of quantum field theories leads quite naturally to the problem of evaluating a large number of multi-loop Feynman diagrams. After a tensor reduction the Feynman diagrams can be expressed in an even larger number of scalar Feynman integrals of the form

$$\int d^d l_1 \cdots \int d^d l_L \frac{1}{D_1^{n_1} \dots D_N^{n_N}}, \qquad (1.1)$$

where L is the number of loops and $d = 4 - 2\epsilon$ the number of space-time dimensions in the context of dimensional regularization. The denominators D_i in (1.1) are usually of the form $p^2 - m^2$, where p is a linear combination of loop momenta and external momenta and m some mass.

A standard technique nowadays is the usage of integration-by-parts identities [1, 2] for the reduction of this large number of Feynman integrals to a rather small set of so-called master integrals. These identities provide linear dependences between various Feynman integrals, where the coefficients are rational functions in both the space-time dimension d and the kinematic variables of the problem.

Many methods were developed to solve these master integrals. For an overview see e.g. [3]. Among the most successful ones is the method of differential equations which is also based on integration-by-parts reductions [4–6]. Recently, significant progress was made in this method, when Henn conjectured the existence of a canonical basis for master integrals expressible in terms of iterated integrals [7]. In this basis the right hand side of the system of differential equations is proportional to $\epsilon = (4 - d)/2$. If the boundary conditions are known, the solution of the system of differential equations in an ϵ -series becomes trivial.

Two years ago, Lee proposed an algorithm to automate finding a canonical basis [8]. A first implementation for this algorithm was presented in [9, 10].

In this paper we present epsilon, a further implementation of Lee's algorithm based on the Fermat[11] computer algebra system. Our implementation utilizes the explicit dependence of the transformations used by Lee's algorithm on the kinematic variable to reduce the number of variables in intermediate steps. Another advantage of our implementation is the support of systems with singularities at complex points using Fermat's polymod capability.

In Section 2 we introduce some definitions and explain implementation details. In Section 3 the installation procedure and the usage of epsilon is described. In Section 4 we give a non-trivial example of the usage based on a real three-loop computation.

2 Implementation details

2.1 Definitions

We consider a set of N master integrals \vec{f} fulfilling an ordinary system of differential equations

$$\frac{\partial f(x,\epsilon)}{\partial x} = \mathbb{M}(x,\epsilon) \vec{f}(x,\epsilon) , \qquad (2.1)$$

where x is a kinematic variable, $\mathbb{M}(x, \epsilon)$ is an $N \times N$ -matrix and ϵ is a regulator in $d = 4 - 2\epsilon$ dimensions in the context of dimensional regularization. We restrict ourselves to the case

$$\mathbb{M}(x,\epsilon) = \sum_{x_j \in S} \sum_{k \ge 0} \frac{\mathbb{M}_k^{(x_j)}(\epsilon)}{(x-x_j)^{k+1}} + \sum_{k \ge 0} x^k \mathbb{M}_k(\epsilon) , \qquad (2.2)$$

where S is the set of all finite singularities and $\mathbb{M}_{k}^{(x_{j})}$ and $\mathbb{M}_{k}(\epsilon)$ are independent of x. In particular, singularities x_{j} depending on ϵ are forbidden. In many physically relevant cases one can use a trial and error approach to find a basis of master integrals \vec{f} fulfilling the restriction (2.2). The main strategy of our implementation is to keep the system always in the form of (2.2) since here the x-dependence is explicit.

A singularity $x_j < \infty$ has Poincaré rank p if $\mathbb{M}_p^{(x_j)} \neq 0$ and $\mathbb{M}_k^{(x_j)} = 0$ for k > p. In addition to the finite singularities, the system might also have a singularity at ∞ . The Poincaré rank p of a singularity at ∞ is defined as the Poincaré rank of the singularity at y = 0 of the system $\mathbb{M}(1/y, \epsilon)/y^2$. So (2.2) has Poincaré rank p > 0 at ∞ if $\mathbb{M}_{p-1} \neq 0$ and $\mathbb{M}_k = 0$ for $k \ge p$, and Poincaré rank p = 0 at ∞ if all $\mathbb{M}_k = 0$ and $\sum_{x_j \in S} \mathbb{M}_0^{(x_j)} = 0$, the system is not singular at ∞ .

Let p be the Poincaré rank of a singularity $x_j < \infty$, then the generalized Poincaré rank (or Moser rank) [12] of this singularity is defined as p + r/n - 1, where $r = \operatorname{rank} \mathbb{M}_p^{(x_j)}$ and n is the dimension of the system.

A system

$$\mathbb{M}(x,\epsilon) = \sum_{x_j \in S} \frac{\mathbb{M}_0^{(x_j)}(\epsilon)}{x - x_j}, \qquad (2.3)$$

where all singularities have Poincaré rank zero is called Fuchsian, and a system

$$\mathbb{M}(x,\epsilon) = \epsilon \sum_{x_j \in S} \frac{\widehat{\mathbb{M}}_0^{(x_j)}}{x - x_j}, \qquad (2.4)$$

where $\widehat{\mathbb{M}}_{0}^{(x_{j})}$ is no longer a function of ϵ , is said to be in ϵ -form. A change of basis

$$\vec{g}(x,\epsilon) = \mathbb{T}^{-1}(x,\epsilon)\vec{f}$$

modifies the system (2.1) to

$$\frac{\partial \vec{g}(x,\epsilon)}{\partial x} = \widetilde{\mathbb{M}}(x,\epsilon)\vec{g}(x,\epsilon)\,,$$

with

$$\widetilde{\mathbf{M}}(x,\epsilon) = \mathbb{T}^{-1}(x,\epsilon)\mathbb{M}(x,\epsilon)\mathbb{T}(x,\epsilon) - \mathbb{T}^{-1}(x,\epsilon)\frac{\partial}{\partial x}\mathbb{T}(x,\epsilon).$$
(2.5)

We assume the master integrals in \vec{f} to be ordered in a way that a block-triangular structure of the system is obtained (for details see e.g. [8]). We will often make use of this block-triangular structure. Therefore we write

$$\mathbf{M} = \begin{pmatrix} \mathbf{A} & 0 & 0 \\ \mathbf{B} & \mathbf{C} & 0 \\ \mathbf{D} & \mathbf{E} & \mathbf{F} \end{pmatrix} , \qquad (2.6)$$

and use the same indices as in (2.2) for the matrices $\mathbb{A}, \ldots, \mathbb{F}$ (e.g. $\mathbb{C}_{k}^{(x_{j})}(\epsilon)$). The block \mathbb{C} is called the active block as we apply Lee's algorithm to this block. As \mathbb{A} to \mathbb{F} are matrices as well, the definition of what we call the active block is more or less arbitrary as long as a block-triangular structure is obtained. But from a computational point of view a small dimension of the active block is preferable since this reduces the complexity of the resulting operations. In the following, the matrices $\mathbb{A}_{k}^{(x_{j})}, \ldots, \mathbb{F}_{k}^{(x_{j})}$ and $\mathbb{A}_{k}, \ldots, \mathbb{F}_{k}$ will be referred to as coefficient matrices.

2.2 Utilizing the explicit *x*-dependence

Lee's algorithm uses three types of transformations: balances, off-diagonal reductions and x-independent transformations.

We define balances as

$$\mathcal{B}(\mathbb{P}, x_1, x_2) = \overline{\mathbb{P}} + \frac{x - x_2}{x - x_1} \mathbb{P}, \qquad (2.7a)$$

$$\mathcal{B}(\mathbb{P}, x_1, \infty) = \overline{\mathbb{P}} + \frac{1}{x - x_1} \mathbb{P}, \qquad (2.7b)$$

$$\mathcal{B}(\mathbb{P}, \infty, x_2) = \overline{\mathbb{P}} + (x - x_2)\mathbb{P}, \qquad (2.7c)$$

where \mathbb{P} is a projector to be specified below, depending only on ϵ , $\overline{\mathbb{P}} = \mathbb{1} - \mathbb{P}$ and $x_1, x_2 < \infty$. In Lee's algorithm balances are applied to the active block \mathbb{C} in order to reduce the generalized Poincaré rank of singular points and to normalize eigenvalues of Fuchsian singularities.

Off-diagonal reductions are used to reduce the block \mathbb{B} to Fuchsian form after the blocks \mathbb{A} and \mathbb{C} were already reduced to ϵ -form. They are defined by

$$\mathcal{L}(x_1, k, \mathbb{G}) = 1 + \frac{1}{(x - x_1)^k} \mathbb{G},$$
 (2.8a)

$$\mathcal{L}(\infty, k, \mathbb{G}) = \mathbb{1} + x^k \mathbb{G}, \qquad (2.8b)$$

where

$$G = \begin{pmatrix} 0 & 0 & 0 \\ \widehat{G} & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} .$$
 (2.8c)

The block $\widehat{\mathbb{G}}$ has the same boundaries in \mathbb{G} as block \mathbb{B} in (2.6). Note that $\mathbb{G}^2 = 0$.

In both types of transformations the x-dependence is explicit. Another type of transformation which is independent of x is used in the last step of Lee's algorithm to factor out ϵ .

Our goal is to use those three types in the transformation rule (2.5) without spoiling the form (2.2) or the block-triangular structure (2.6) of the system.

As a pedagogical example we consider the transformation of block \mathbb{B} under a balance between two singularities x_1 and x_2 , i.e.

$$\mathbb{T} = \mathcal{B}(\mathbb{P}, x_1, x_2) = \overline{\mathbb{P}} + \frac{x - x_2}{x - x_1} \mathbb{P}, \quad \mathbb{T}^{-1} = \mathcal{B}(\mathbb{P}, x_2, x_1) = \overline{\mathbb{P}} + \frac{x - x_1}{x - x_2} \mathbb{P}.$$
(2.9)

Since we want to apply Lee's algorithm to the active block we can restrict the form of the projector $\mathbb P$ to

$$\mathbb{P} = \begin{pmatrix} 0 & 0 & 0 \\ 0 & \mathbb{Q} & 0 \\ 0 & 0 & 0 \end{pmatrix} , \qquad (2.10)$$

where \mathbb{Q} is a projector with the dimensions of the active block. Inserting (2.9) into (2.5)

we obtain

$$\begin{split} \widetilde{\mathbf{M}}(x,\epsilon) &= \left[\overline{\mathbb{P}} + \frac{x - x_1}{x - x_2} \mathbb{P}\right] \mathbb{M}(x,\epsilon) \left[\overline{\mathbb{P}} + \frac{x - x_2}{x - x_1} \mathbb{P}\right] + \frac{x_2 - x_1}{(x - x_1)(x - x_2)} \mathbb{P} \\ &= \mathbb{M}(x,\epsilon) - \mathbb{P}\mathbb{M}(x,\epsilon)\overline{\mathbb{P}} - \overline{\mathbb{P}}\mathbb{M}(x,\epsilon)\mathbb{P} + \frac{x - x_2}{x - x_1}\overline{\mathbb{P}}\mathbb{M}(x,\epsilon)\mathbb{P} + \frac{x - x_1}{x - x_2}\mathbb{P}\mathbb{M}(x,\epsilon)\overline{\mathbb{P}} \\ &+ \frac{x_2 - x_1}{(x - x_1)(x - x_2)}\mathbb{P} \end{split}$$

So block \mathbb{B} in (2.6) transforms as

$$\widetilde{\mathbb{B}}(x,\epsilon) = \mathbb{B}(x,\epsilon) - \mathbb{QB}(x,\epsilon) + \frac{x-x_1}{x-x_2} \mathbb{QB}(x,\epsilon) \,.$$

Inserting the form (2.2) of $\mathbb{B}(x,\epsilon)$ leads to

$$\widetilde{\mathbb{B}}(x,\epsilon) = \mathbb{B}(x,\epsilon) - \mathbb{Q}\mathbb{B}(x,\epsilon) + \sum_{x_j \in S} \sum_{k \ge 0} \frac{(x-x_1)\mathbb{Q}\mathbb{B}_k^{(x_j)}(\epsilon)}{(x-x_2)(x-x_j)^{k+1}} + \sum_{k \ge 0} \frac{(x-x_1)x^k}{x-x_2} \mathbb{Q}\mathbb{B}_k(\epsilon) \,.$$
(2.11)

Using partial fractioning and the incomplete geometric series, we can show that

$$\sum_{k=0}^{\infty} \frac{a_k}{(x-x_2)(x-x_j)^{k+1}} = \frac{1}{x-x_2} \sum_{n\geq 0} \frac{a_n}{(x_2-x_j)^{n+1}} -\sum_{k\geq 0} \frac{1}{(x-x_j)^{k+1}} \sum_{n=0}^{\infty} \frac{a_{n+k}}{(x_2-x_j)^{n+1}},$$

$$\sum_{k\geq 0} \frac{x^k}{x-x_2} a_k = \sum_{k\geq 0} x^k \sum_{n\geq 0} x_2^n a_{k+n+1} + \frac{1}{x-x_2} \sum_{n\geq 0} x_2^n a_n,$$
(2.12a)
(2.12b)

where identity (2.12a) only holds for $x_j \neq x_2$. Combining (2.11) and (2.12) yields

$$\widetilde{\mathbb{B}}(x,\epsilon) = \mathbb{B}(x,\epsilon) + \frac{x_2 - x_1}{x - x_2} \sum_{x_j \in S \setminus \{x_2\}} \sum_{n \ge 0} \frac{\mathbb{Q}\mathbb{B}_n^{(x_j)}}{(x_2 - x_j)^{n+1}} + \frac{x_2 - x_1}{x - x_2} \sum_{n \ge 0} x_2^n \mathbb{Q}\mathbb{B}_n + \sum_{k \ge 1} \frac{(x_2 - x_1)\mathbb{Q}\mathbb{B}_{k-1}^{(x_2)}(\epsilon)}{(x - x_2)^{k+1}} + \sum_{x_j \in S \setminus \{x_2\}} \sum_{k \ge 0} \frac{x_1 - x_2}{(x - x_j)^{k+1}} \sum_{n \ge 0} \frac{\mathbb{Q}\mathbb{B}_{n+k}^{(x_j)}(\epsilon)}{(x_2 - x_j)^{n+1}} + (x_2 - x_1) \sum_{k \ge 0} x_2^n \mathbb{Q}\mathbb{B}_{k+n+1}(\epsilon).$$

$$(2.13)$$

Hence, the transformation laws for the coefficient matrices can be found by comparing (2.13) with the structure of (2.2):

$$\begin{split} \widetilde{\mathbb{B}}_{0}^{(x_{2})}(\epsilon) &= \mathbb{B}_{0}^{(x_{2})}(\epsilon) + \sum_{x_{j} \in S \setminus \{x_{2}\}} \sum_{n \geq 0} \frac{x_{2} - x_{1}}{(x_{2} - x_{j})^{n+1}} \mathbb{Q} \mathbb{B}_{n}^{(x_{j})}(\epsilon) + (x_{2} - x_{1}) \sum_{n \geq 0} x_{2}^{n} \mathbb{Q} \mathbb{B}_{n}(\epsilon) \,, \\ \widetilde{\mathbb{B}}_{k>0}^{(x_{2})}(\epsilon) &= \mathbb{B}_{k}^{(x_{2})}(\epsilon) + (x_{2} - x_{1}) \mathbb{Q} \mathbb{B}_{k-1}^{(x_{2})}(\epsilon) \,, \\ \widetilde{\mathbb{B}}_{k}^{(x_{j} \neq x_{2})}(\epsilon) &= \mathbb{B}_{k}^{(x_{j})}(\epsilon) + \sum_{n \geq 0} \frac{x_{1} - x_{2}}{(x_{2} - x_{j})^{n+1}} \mathbb{Q} \mathbb{B}_{n+k}^{(x_{j})}(\epsilon) \,, \\ \widetilde{\mathbb{B}}_{k}(\epsilon) &= (x_{2} - x_{1}) \sum_{n \geq 0} x_{2}^{n} \mathbb{Q} \mathbb{B}_{k+n+1}(\epsilon) \,. \end{split}$$

An advantage of this form over the original transformation (2.9) is that now all operations are independent of x. Therefore, the underlying computer algebra system has to deal with rational functions of one less variable. The form (2.2) remains unspoiled, i.e. it is not necessary to perform a partial fraction decomposition after the transformation.

All transformations in terms of the coefficient matrices are listed in appendix A.

2.3 Overview of Lee's algorithm

Three basic steps allow Lee's algorithm[8] to transform an ordinary system of differential equations into an ϵ -form (2.4) if they are applied to the whole system:

- 1. transformation of a system into Fuchsian form,
- 2. normalization of the eigenvectors of all matrix residues,
- 3. factorization of ϵ .

In order to make use of the block-triangular structure of the system (2.6) these three steps are applied only to the active block followed by a fourth step to transform the off-diagonal block \mathbb{B} into Fuchsian form.

In this sub-section we briefly discuss all four steps. More details can be found in the original paper by Lee[8].

Fuchsification

The basic building blocks for the first part of Lee's algorithm, the transformation of the system to Fuchsian form, are the balances defined in (2.7). With the right choice of a

projector \mathbb{P} , it is possible to perform a so-called Moser reduction to strictly lower the generalized Poincaré rank of the singularity at x_1 [12]. In this discussion we restrict ourselves to the case where $x_1 < \infty$. A more general treatment is given in[8].

Let p be the Poincaré rank of the singularity at x_1 of the active block \mathbb{C} and $\{u_k^{\alpha}\}$ with $\alpha = 0, \ldots, n_k$ be the set of $n_k + 1$ right generalized eigenvectors of $\mathbb{C}_p^{(x_1)}$ belonging to a Jordan block k in the Jordan decomposition of $\mathbb{C}_p^{(x_1)}$. The Jordan blocks are ordered by their size so that $n_i \ge n_{i+1}$. If p > 0, we assume all eigenvalues of $\mathbb{C}_p^{(x_1)}$ to be zero, or else no transformation to lower the generalized Poincaré rank exists. The right generalized eigenvectors fulfill

$$\mathbb{C}_p^{(x_1)} u_k^{(0)} = 0, \quad \mathbb{C}_p^{(x_1)} u_k^{(\alpha+1)} = u_k^{(\alpha)}.$$

These relations are invariant under the transformation

$$u_k^{(\alpha)} \to u_k^{(\alpha)} + c u_l^{(\alpha)} \,, \tag{2.14}$$

where $\alpha = 0, \ldots, n_k$ and k > l. The left generalized eigenvectors $v_k^{(\alpha)}$ are related to the right generalized eigenvectors by

$$\left(v_1^{(n_1)}, \dots, v_1^{(0)}, v_2^{(n_2)}, \dots, v_2^{(0)}, \dots\right) = \left[\left(u_1^{(0)}, \dots, u_1^{(n_1)}, u_2^{(0)}, \dots, u_2^{(n_2)}, \dots\right)^{-1}\right]^{\dagger}, \quad (2.15)$$

and fulfill

$$v_k^{(0)\dagger} \mathbb{C}_p^{(x_1)} = 0, \quad v_k^{(\alpha+1)\dagger} \mathbb{C}_p^{(x_1)} = v_k^{(\alpha)\dagger}.$$

The transformation (2.14) allows us to find a basis of eigenvectors which satisfies

$$v_j^{(0)\dagger} \mathbb{C}_{p-1}^{(x_1)} u_k^{(0)} = 0, \qquad (2.16)$$

for $j \notin R$ and $k \in R \cup \{k_0\}$, where R is some set of trivial Jordan blocks $(n_i = 0 \text{ for } i \in R)$ and k_0 is a non-trivial Jordan block $(n_{k_0} > 0)$. An algorithm to find these generalized eigenvectors, together with the set R and k_0 is given in [8].

If in the definition of $\mathbb{P}(2.10)$ we use

$$\mathbb{Q} = \sum_{k \in R \cup \{k_0\}} u_k^{(0)} w_k^{\dagger}, \qquad (2.17)$$

with $w_j^{\dagger} u_k^{(0)} = \delta_{jk}$ and then apply the balance (2.7a), the generalized Poincaré rank of the singularity at x_1 is decreased. To see this, we introduce the projector

$$\mathbb{Q}_1 = \sum_{k \in R \cup \{k_0\}} u_k^{(0)} v_k^{(n_k)\dagger} \,. \tag{2.18}$$

Since $v_j^{(n_j)\dagger}u_k^{(0)} = \delta_{jk}$, the projector \mathbb{Q}_1 also is of the form (2.17). From the definitions of \mathbb{Q} and \mathbb{Q}_1 follows

$$\mathbb{Q}_1 \mathbb{Q} = \mathbb{Q}$$
, $\mathbb{Q} \mathbb{Q}_1 = \mathbb{Q}_1$, $\mathbb{C}_p^{(x_1)} \mathbb{Q} = \mathbb{C}_p^{(x_1)} \mathbb{Q}_1 = 0$

Evaluating the transformation (A.1b) at k = p leads to

$$\widetilde{\mathbb{C}}_{p}^{(x_{1})} = \overline{\mathbb{Q}}\mathbb{C}_{p}^{(x_{1})} + (x_{1} - x_{2})\overline{\mathbb{Q}}\mathbb{C}_{p-1}^{(x_{1})}\mathbb{Q}$$
$$= \left(\overline{\mathbb{Q}} + \mathbb{Q}_{1}\right)\left[\overline{\mathbb{Q}}_{1}\mathbb{C}_{p}^{(x_{1})} + (x_{1} - x_{2})\overline{\mathbb{Q}}_{1}\mathbb{C}_{p-1}^{(x_{1})}\mathbb{Q}_{1}\right]\left(\overline{\mathbb{Q}}_{1} + \mathbb{Q}\right)$$

As $(\overline{\mathbb{Q}} + \mathbb{Q}_1) = (\overline{\mathbb{Q}}_1 + \mathbb{Q})^{-1}$, the matrix rank of $\widetilde{\mathbb{C}}_p^{(x_1)}$ is given by

$$\operatorname{rank} \widetilde{\mathbb{C}}_p^{(x_1)} = \operatorname{rank} \widehat{\mathbb{C}}_p^{(x_1)}, \quad \widehat{\mathbb{C}}_p^{(x_1)} = \overline{\mathbb{Q}}_1 \mathbb{C}_p^{(x_1)} + (x_1 - x_2) \overline{\mathbb{Q}}_1 \mathbb{C}_{p-1}^{(x_1)} \mathbb{Q}_1,$$

The argument that the matrix rank (and therefore the generalized Poincaré rank) of $\widetilde{\mathbb{C}}_p^{(x_1)}$ is lower than the matrix rank of $\mathbb{C}_p^{(x_1)}$ is as follows:

- All left eigenvectors $v_j^{(0)}$ of $\mathbb{C}_p^{(x_1)}$ with $j \in R$ are left eigenvectors of $\widehat{\mathbb{C}}_p^{(x_1)}$ as $v_j^{(0)\dagger}\overline{\mathbb{Q}}_1 = 0$.
- All left eigenvectors $v_j^{(0)}$ of $\mathbb{C}_p^{(x_1)}$ with $j \notin R$ are left eigenvectors of $\widehat{\mathbb{C}}_p^{(x_1)}$ as $v_j^{(0)\dagger} \overline{\mathbb{Q}}_1 = v_j^{(0)\dagger}$ and so

$$v_j^{(0)\dagger} \widehat{\mathbb{C}}_p^{(x_1)} = (x_1 - x_2) \sum_{k \in R \cup \{k_0\}} v_j^{(0)\dagger} \mathbb{C}_{p-1}^{(x_1)} u_k^{(0)} v_k^{(n_k)\dagger} = 0,$$

where we used (2.16).

• The vector $v_{k_0}^{(n_{k_0})}$ which is not a left eigenvector of $\mathbb{C}_p^{(x_1)}$ is an additional left eigenvector of $\widehat{\mathbb{C}}_p^{(x_1)}$ as $v_{k_0}^{(n_{k_0})\dagger}\overline{\mathbb{Q}}_1 = 0$.

So $\widehat{\mathbb{C}}_p^{(x_1)}$ has one eigenvector more than $\mathbb{C}_p^{(x_1)}$ and therefore has a lower matrix rank.

Unfortunately, the balance (2.7a) might also increase the Poincaré rank at x_2 . Therefore, we have a closer look at (A.1d) evaluated at k = q + 1, where q is the Poincaré rank at the singularity x_2 :

$$\widetilde{\mathbb{C}}_{q+1}^{(x_2)} = (x_2 - x_1) \mathbb{Q} \mathbb{C}_q^{(x_2)} \overline{\mathbb{Q}} \,.$$

If this expression vanishes, the Poincaré rank at x_2 is not increased. This is the case if the vectors w_k^{\dagger} in (2.17) span a left-invariant space of $\mathbb{C}_q^{(x_2)}$.

In Lee's algorithm the above steps are used to decrease the Poincaré rank of a singularity x_1 as long as it is possible to find a projector (2.17) with w_k^{\dagger} spanning a left-invariant space of some $\mathbb{C}_q^{(x_2)}$, where $x_2 \neq x_1$ is some other singular point. If no such projector exists, a regular point y is chosen and the projector \mathbb{Q}_1 defined by (2.18) is used in the balance. This of course creates a new (apparent) Fuchsian singularity at y. This way it is possible to reduce the system to an equivalent one with all singularities having Poincaré rank zero. Fortunately, the next step, the normalization of the eigenvalues, removes all apparent singularities introduced in this step.

Normalization

The second step of Lee's algorithm is the normalization of the eigenvalues of the matrix residues. Again, the main ingredients are the balances defined in (2.7). We assume all singularities to be Fuchsian now and the eigenvalues of the matrix residues to be of the form $a + b\epsilon$ with $a \in \mathbb{Z}$. This is a necessary condition for the existence of a transformation normalizing all eigenvalues. If this condition cannot be fulfilled a redefinition of the kinematic variable might be helpful. A normalized eigenvalue is an eigenvalue proportional to ϵ (i.e. with a = 0). The matrix residue at ∞ is equal to $-\sum_{x_j \in S} \mathbb{C}_0^{(x_j)}$. Therefore, the sum of all matrix residues (including the residue at ∞) vanishes and hence the sum of all eigenvalues of all matrix residues must vanish as well.

For brevity's sake, we will restrict our discussion to finite singularities $x_1, x_2 < \infty$; the generalized treatment can again be found in[8].

Let $\{u_k^{(\alpha)}\}$ with $\alpha = 0, \ldots, n_k$ be the set of $n_k + 1$ right generalized eigenvectors of $\mathbb{C}_0^{(x_1)}$ belonging to the Jordan block k:

$$\mathbb{C}_{0}^{(x_{1})}u_{k}^{(0)} = \lambda_{k}u_{k}^{(0)}, \quad \mathbb{C}_{0}^{(x_{1})}u_{k}^{(\alpha+1)} = \lambda_{k}u_{k}^{(\alpha+1)} + u_{k}^{(\alpha)}$$

As in the fuch sification step, we define left generalized eigenvectors $v_k^{(\alpha)}$ via (2.15) which satisfy

$$v_k^{(0)\dagger} \mathbb{C}_0^{(x_1)} = \lambda_k v_k^{(0)\dagger}, \quad v_k^{(\alpha+1)\dagger} \mathbb{C}_0^{(x_1)} = \lambda_k v_k^{(\alpha+1)\dagger} + v_k^{(\alpha)\dagger}$$

A balance (2.7a) with a right choice of a projector \mathbb{P} can now be used to shift one eigenvalue of the matrix residue at x_1 up by one and/or one eigenvalue of the matrix residue at x_2 down by one. Let us consider a projector \mathbb{P} defined by (2.10) with

$$\mathbb{Q} = u_k^{(0)} w^\dagger \,, \tag{2.19}$$

where $w^{\dagger}u_k^{(0)} = 1$. Also an additional projector

$$\mathbb{Q}_1 = u_k^{(0)} v_k^{(n_k)}$$

is useful for the discussion. The following relations involving these two projectors hold:

$$\mathbb{Q}\mathbb{Q}_1 = \mathbb{Q}_1, \quad \mathbb{Q}_1\mathbb{Q} = \mathbb{Q}, \quad \mathbb{C}_0^{(x_1)}\mathbb{Q} = \lambda_k\mathbb{Q}, \quad \mathbb{C}_0^{(x_1)}\mathbb{Q}_1 = \lambda_k\mathbb{Q}_1$$

Let us now consider the transformation of $\mathbb{C}_{0}^{(x_{1})}$ under a balance (2.7a) as given in (A.1a):

$$\begin{split} \widetilde{\mathbb{C}}_{0}^{(x_{1})} &= \mathbb{C}_{0}^{(x_{1})} - \mathbb{Q}\mathbb{C}_{0}^{(x_{1})}\overline{\mathbb{Q}} + \sum_{x_{j} \in S \setminus \{x_{1}\}} \frac{x_{1} - x_{2}}{x_{1} - x_{j}} \overline{\mathbb{Q}}\mathbb{C}_{0}^{(x_{j})}\mathbb{Q} + \mathbb{Q} \\ &= (\overline{\mathbb{Q}} + \mathbb{Q}_{1})\widehat{\mathbb{C}}_{0}^{(x_{1})}(\overline{\mathbb{Q}}_{1} + \mathbb{Q}) \,, \end{split}$$

where

$$\widehat{\mathbb{C}}_{0}^{(x_{1})} = \mathbb{C}_{0}^{(x_{1})} - \mathbb{Q}_{1}\mathbb{C}_{0}^{(x_{1})}\overline{\mathbb{Q}}_{1} + \sum_{x_{j}\in S\setminus\{x_{1}\}}\frac{x_{1}-x_{2}}{x_{1}-x_{j}}\overline{\mathbb{Q}}_{1}\mathbb{C}_{0}^{(x_{j})}\mathbb{Q}_{1} + \mathbb{Q}_{1}$$

Because of $(\overline{\mathbb{Q}} + \mathbb{Q}_1) = (\overline{\mathbb{Q}}_1 + \mathbb{Q})^{-1}$, $\widetilde{\mathbb{C}}_0^{(x_1)}$ and $\widehat{\mathbb{C}}_0^{(x_1)}$ are related by a similarity transformation and thus have the same eigenvalues. To evaluate the eigenvalues of $\widetilde{\mathbb{C}}_0^{(x_1)}$ it is therefore sufficient to analyze the eigenvalues of $\widehat{\mathbb{C}}_0^{(x_1)}$ which are much simpler to determine. We consider $\widehat{\mathbb{C}}_0^{(x_1)}$ in the basis of the generalized eigenvectors of $\mathbb{C}_0^{(x_1)}$. In this basis $\mathbb{C}_0^{(x_1)}$ is in Jordan normal form. The second term $-\mathbb{Q}_1 \mathbb{C}_0^{(x_1)} \overline{\mathbb{Q}}_1$ removes all elements from the row corresponding to $u_k^{(0)}$ but the diagonal one containing the eigenvalue to $u_k^{(0)}$. The last term $+\mathbb{Q}_1$ increases the diagonal element by one. The terms proportional to $\overline{\mathbb{Q}}_1 \mathbb{C}_0^{(x_2)} \mathbb{Q}_1$ contribute to the non-diagonal elements of the column corresponding to $u_k^{(0)}$ (or $v_k^{(n_k)\dagger}$). Hence, the transformations can be summarized as



where * stands for contributions from the $\overline{\mathbb{Q}}_1 \mathbb{C}_0^{(x_j)} \mathbb{Q}_1$ terms. Calculating the characteristic polynomial by means of a Laplace expansion along the row corresponding to $u_k^{(0)}$ leads to the conclusion that all eigenvalues but one stay the same; only one eigenvalue λ_k is changed to $\lambda_k + 1$.

In the same way, a projector

$$\mathbb{Q}' = w v_k^{(0)\dagger} \tag{2.20}$$

with a left eigenvector $v_k^{(0)\dagger}$ of $\mathbb{C}_0^{(x_2)}$ and $v_k^{(0)\dagger}w = 1$ shifts one eigenvalue λ_k of $\mathbb{C}_0^{(x_2)}$ down by one.

A balance with a projector (2.19) could spoil the Fuchsian form of the system if it increases the Poincaré rank at any singularity. This is in principle possible at x_2 . Evaluating (A.1d) at k = 1 leads to

$$\widetilde{\mathbb{C}}_1^{(x_2)} = (x_2 - x_1) \mathbb{Q} \mathbb{C}_0^{(x_2)} \overline{\mathbb{Q}} \,.$$

This vanishes if w^{\dagger} is a left eigenvector of $\mathbb{C}_{0}^{(x_{2})}$. In that case not only the Fuchsian form of the system is preserved but we also arrive at a projector of the form (2.20). So the projector of choice is

$$Q = u_k^{(0)} v_l^{(0)\dagger} , \qquad (2.21)$$

where $u_k^{(0)}$ is a right eigenvector of $\mathbb{C}_0^{(x_1)}$, $v_l^{(0)\dagger}$ is a left eigenvector of $\mathbb{C}_0^{(x_2)}$ and $v_l^{(0)\dagger}u_k^{(0)} = 1$. This projector, used in balance (2.7a), increases one eigenvalue λ_k of $\mathbb{C}_0^{(x_1)}$ by one and decreases one eigenvalue μ_l of $\mathbb{C}_0^{(x_2)}$ by one while conserving the Fuchsian form of the system.

In order to utilize the considerations above in an algorithmic approach, one first selects a singularity x_0 as 'fallback'. Then balances with projectors of the form (2.21) are used to 'mutually balance' eigenvalues between two singularities. Certainly, the eigenvalue to be increased should be negative and the eigenvalue to be decreased should be positive (for $\epsilon = 0$).

If no such balance exists, the eigenvalues will be balanced with the 'fallback singularity' x_0 regardless of the sign of the eigenvalue at x_0 . This normalizes the eigenvalues at all singularities but x_0 . To normalize even the eigenvalues at x_0 , we balance one unnormalized eigenvalue with some regular point creating a new apparent singularity there and restart the algorithm. Hopefully, the unnormalized eigenvalue at this new apparent singularity can now be mutually balanced with another unnormalized eigenvalue at x_0 .

ϵ -Factorization

In the next step we find an x-independent transformation $\mathbb{T}(\epsilon)$ to factor out ϵ . For an x-independent \mathbb{T} , (2.5) becomes a similarity transformation and does not change the eigenvalues of the system. This makes it necessary for the eigenvalues to be proportional to ϵ , i.e. the normalization step before must have been successful. We use the fact that

 $\mathbb{T}^{-1}(\epsilon)[\mathbb{M}_0^{(x_j)}(\epsilon)/\epsilon]\mathbb{T}(\epsilon)$ should be independent of ϵ , so that the equation

$$\mathbb{T}^{-1}(\epsilon)\frac{\mathbb{M}_{0}^{(x_{j})}(\epsilon)}{\epsilon}\mathbb{T}(\epsilon) = \mathbb{T}^{-1}(\mu)\frac{\mathbb{M}_{0}^{(x_{j})}(\mu)}{\mu}\mathbb{T}(\mu)$$

holds for all $x_j \in S$. Multiplying this equation from the left by $\mathbb{T}(\epsilon)$ and from the right by $\mathbb{T}^{-1}(\mu)$, leads to

$$\frac{\mathbb{M}_{0}^{(x_{j})}(\epsilon)}{\epsilon}\mathbb{T}(\epsilon,\mu) = \mathbb{T}(\epsilon,\mu)\frac{\mathbb{M}_{0}^{(x_{j})}(\mu)}{\mu}, \qquad (2.22)$$

where $\mathbb{T}(\epsilon, \mu) = \mathbb{T}(\epsilon)\mathbb{T}^{-1}(\mu)$. This linear system can be solved e.g. with Gaussian elimination and the constants should be fixed so that $\mathbb{T}(\epsilon, \mu_0)$ is an invertible matrix, where μ_0 is some arbitrary number. The transformation $\mathbb{T}(\epsilon, \mu_0)$ can now be used to factor out ϵ .

Fuchsification of off-diagonal blocks

Since the definition of the active block is somewhat arbitrary, we should in principle be able to transform all diagonal blocks to ϵ -form by redefining the active block and applying the three steps described above. However, we still need to reduce the off-diagonal block \mathbb{B} to Fuchsian form. Again we will restrict the discussion to finite singularities $x_1 < \infty$ and assume the block \mathbb{A} and \mathbb{C} to be already in ϵ -form.

Let p be the Poincaré rank of the off-diagonal block \mathbb{B} at the singularity x_1 , i.e. $\mathbb{B}_p^{(x_1)} \neq 0$ and $\mathbb{B}_k^{(x_1)} = 0$ for k > p. The behavior of $\mathbb{B}_p^{(x_1)}$ under a transformation (2.8a) with k = pis given by (A.3a):

$$\widetilde{\mathbb{B}}_p^{(x_1)} = \mathbb{B}_p^{(x_1)} + \mathbb{C}_0^{(x_1)} \widehat{\mathbb{G}} - \widehat{\mathbb{G}} \mathbb{A}_0^{(x_1)} + p \widehat{\mathbb{G}} .$$

In order to decrease the Poincaré rank, $\widehat{\mathbb{G}}$ has to be determined such that $\widetilde{\mathbb{B}}_p^{(x_1)}$ vanishes. This linear system of equations can be solved e.g. with Gaussian elimination.

Hence, with transformations of the form (2.8) it is possible to reduce all singularities of the off-diagonal block \mathbb{B} to Fuchsian form.

3 Usage

3.1 Installation guide on Linux systems

Ensure that the dependencies

- Fermat (≥ 6.0) [11],
- GiNaC ($\geq 1.6.2$) [13] (for epsilon-prepare only)

are installed.

As a next step, libFermat has to be installed. libFermat, which was developed in connection with epsilon, is a C++ library designed to communicate with Fermat. Nevertheless, we decided to publish libFermat in a separate repository since it might be useful elsewhere. Internally, the communication is done with PStreams[14] which is included in the package. The source code of the most recent version of libFermat can be obtained via github using

```
git clone https://github.com/mprausa/libFermat.git
```

This will create a directory libFermat/ and clone the library into that location. Now inside this directory, run

```
cmake -DCMAKE_INSTALL_PREFIX=/path/to/install .
make
make install
```

where /path/to/install is your desired installation directory and defaults to /usr/local on a typical Linux system. The library is installed into the sub-directory lib and the header files into the sub-directory include of /path/to/install. If your choice is a global directory you will require root access for the last step make install. Remember to include the sub-directory lib of /path/to/install into the LD_LIBRARY_PATH environment variable if you are using a non-standard directory.

The next step is to install epsilon and epsilon-prepare. In principle the procedure is the same as for the installation of libFermat. First, obtain the most recent version of the source code with

```
git clone https://github.com/mprausa/epsilon.git
```

then change into the newly created directory epsilon/ and run

```
cmake -DCMAKE_INSTALL_PREFIX=/path/to/install .
make
make install
```

It is recommended to use the same /path/to/install as for libFermat, else the cmake step might require additional options to find libFermat. The programs epsilon and epsilon-prepare are installed into the sub-directory bin of the installation path. As before, make install might need root access depending on the installation prefix. For a non-standard installation prefix, the environment variable PATH should be adjusted to include the sub-directory bin of /path/to/install so that the programs epsilon and epsilon-prepare can be found by the shell.

It is also possible to build epsilon and epsilon-prepare individually. This can be done by changing into the corresponding sub-directory and running cmake and make from within there.

The epsilon-repository also offers a Mathematica package EpsilonTools.m found in the sub-directory mma/. Run

./install.sh

from within this sub-directory to install EpsilonTools.m into the Applications/ directory of your Mathematica installation.

3.2 Input/Output format

epsilon uses its own file format to represent a system of differential equations of the form (2.2), where every line represents one coefficient matrix. A line starts with either $[A[x_j,k]:' \text{ or } B[k]:' \text{ followed by a matrix.}$ The matrix is stored as a list of rows, where each row is itself a list of matrix elements. Lists are enclosed in curly-braces and list entries are separated by commas. A line starting with $[A[x_j,k]:' (B[k]:')]$ represents a matrix $M_k^{(x_j)}(M_k)$ in (2.2).

The name of the symbol used to represent ϵ is fixed to ep.

An example of an input file for epsilon is given in section 4.

3.3 Usage of epsilon-prepare

The tool epsilon-prepare is used to convert a matrix in Mathematica format (a list of lists) into epsilon input format (see section 3.2). If, for example, the file containing a

symbol	polymod	expression
i	$i^2 + 1$	i
rN	$\mathbf{r}N^2 - \mathbf{r}N + (1+N)/4$	$\frac{1+i\sqrt{N}}{2}$
qN	$qN^2 - qN + (1-N)/4$	$\frac{1+\sqrt{N}}{2}$
$\mathtt{sqrt}N$	${\tt sqrt}N^2-N$	\sqrt{N}
${\tt isqrt}N$	${\tt isqrt}N^2+N$	$i\sqrt{N}$

Table 1: Additional symbols used by epsilon-prepare to represent zeros of quadratic polynomials. N is a positive integer.

matrix in Mathematica format is called matrix.m, then the command

```
epsilon-prepare matrix.m > matrix.dat
```

is used to create a file matrix.dat in epsilon format. The matrix elements of the input matrix are expected to be rational functions in ϵ and x (represented by the symbols ep and x, respectively). The command epsilon-prepare performs a partial fraction decomposition over complex numbers in the variable x of the matrix elements and results in an expression of the form of (2.2).

To perform a partial fractioning of a rational function, the zeros of its denominator have to be determined. Therefore, epsilon-prepare applies GiNaC's polynomial factorization algorithm to the denominator in order to factorize it into polynomials that are irreducible over the integers. In a second step, the zeros of all factors are found individually.

This step will fail, if the considered irreducible polynomial has a degree larger than two. An error is thrown as well if the system is not in form (2.2), i.e. a zero depends on the parameter ϵ .

The additional symbols listed in table 1 might be introduced by epsilon-prepare. See section 3.5 for how to make epsilon accept them.

3.4 Usage of epsilon

The general command syntax for the tool epsilon is

```
epsilon [OPTIONS] JOBS...
```

The path to the Fermat binary can be set inside the environment variable FERMAT. If this variable is not set explicitly, epsilon will look for a binary fer64 inside the directories defined in the environment variable PATH. Options are set once at the start of epsilon.

Jobs are processed one by one in the same order they are defined on the command line. Some jobs will perform transformations to the system. These transformations are stored in a so-called internal transformation queue in RAM and can also be written to an external file.

Options:

• --verbose: Enable verbose output.

This option prints out all communication between libFermat and Fermat during a regular run. This is useful as a debug tool.

• --timings: Enable timings.

This option prints the elapsed real time after every job and the total time of the complete run.

• --symbols symbols: Adjoin additional symbols to Fermat.

This option adds the specified symbols to Fermat. The symbols defined in this option are the very first variables adjoint to Fermat followed by the symbol ep and further internally used symbols. *symbols* has to be a comma-separated list.

• --echelon-fermat: Use the Redrowech function in Fermat to solve linear systems.

At various points in the code linear systems of equations have to be solved. Use this option to choose the **Redrowech** function over our own implementation of Gaussian elimination.

Jobs:

• --fermat *file*: Execute Fermat commands.

This job reads *file* line by line and sends all non-empty lines to Fermat.

• --load file start end: Load system.

This job loads *file* and activates the block {*start*, *end*}. Hereby *file* is expected to be in the format specified in section 3.2

• --write *file*:

Write system.

This job writes the system of differential equations to *file* using the format specified in section 3.2.

• --queue *file*: Set external transformation queue.

This job enables an external transformation queue. An external transformation queue is a file containing all transformations already performed by epsilon during a run. This is particularly useful in connection with the options --load-queue and --replay to restore an aborted run to the state after the last successful transformation.

• --load-queue *filename*: Load transformation queue.

This job loads an external transformation queue from *filename* into an internal transformation queue stored in RAM. This job does *not* apply the transformations stored in the file to the system.

• --replay: Apply internal transformation queue.

This job 'replays' the internal transformation queue, i.e. the transformations in the queue are applied to the system one by one. It should only be used immediately after --load-queue.

• --export *file*: Export transformation matrix.

This job computes a transformation matrix out of the transformations inside the internal transformation queue. The matrix is written in Mathematica format to file.

• --block start end: Activate a block.

This job activates the block {*start*, *end*}.

• --fuchsify: Transform active block into Fuchsian form.

This job reduces the active block to Fuchsian form (2.3). See section 2.3 for details.

• --normalize: Normalize eigenvalues.

This job normalizes the eigenvalues of all residue matrices making them proportional to ϵ . See section 2.3 for details.

• --factorep:

Factor out ϵ (auto detect μ).

This job transforms the active block into ϵ -form (2.4) using the method described in section 2.3. The variable μ in (2.22) is left as an unknown and will be fixed only *after* the system is solved to ensure that the transformation is invertible.

• --factorep-at mu: Factor out ϵ (with predefined μ).

This job transforms the active block into ϵ -form (2.4) using the method described in section 2.3. In this variant, the variable μ in (2.22) is set to **mu** before the system is solved. This is faster than **--factorep** because **epsilon** has to deal with one less variable in the polynomials. Unfortunately, an unlucky choice of **mu** can hit a pole in the matrix elements of the system or one can end up with an uninvertible transformation. In both cases an error is thrown.

• --left-fuchsify: Fuchsify off-diagonal block.

This job is used to transform the block left of the active block (block \mathbb{B} in (2.6)) to Fuchsian form. See section 2.3 for details.

• --dyson file order type format: Generate Dyson operator.

This job writes a Dyson operator $U(x, x_0)$ for the active block up to order order in ϵ to *file*. The active block has to be in ϵ -form. The Dyson operator fulfills

$$\frac{\partial}{\partial x}U(x,x_0) = \epsilon \sum_{x_j \in S} \frac{\widehat{\mathbb{C}}_0^{(x_j)}}{x - x_j} U(x,x_0), \quad U(x_0,x_0) = \mathbb{1}.$$

The option *type* specifies the type of multiple polylogarithms in the output and can be set to GPL, HPL or HPLalt for Goncharov polylogarithms[15] or harmonic polylogarithms in the "a"- or "m"-notation[16, 17], respectively. *format* should be mma or form to specify Mathematica or FORM[18] as output format.

3.5 Using field extensions

The Fermat computer algebra system works with multivariate polynomials over the ground ring \mathbb{Z} and the corresponding quotient field, the rational functions over \mathbb{Z} . Fortunately, Fermat offers a way to extend the ground ring by setting so-called polymods.

A polymod $p(\xi)$ is a univariate polynomial in ξ , where ξ is one of the variables adjoint to Fermat. This instates a new quotient ring $\mathbb{Z}[\xi]/\langle p(\xi) \rangle$ as the ground ring, forcing

the condition $p(\xi) = 0$ onto the variable ξ . In other words every polynomial $q(\xi) \in \mathbb{Z}[\xi]$ encountered by Fermat is replaced immediately by the remainder of a polynomial division $q(\xi) \div p(\xi)$. For more details see the Fermat manual[11].

The polymod $i^2 + 1$ for example leads to the quotient ring $\mathbb{Z}[i]/\langle i^2 + 1 \rangle$ which is equivalent to the Gaussian integers, a field extension of the integers by a number *i* with $i^2 = -1$. Also other complex numbers can be represented in Fermat using polymods. The complex numbers introduced via epsilon-prepare can be represented in Fermat using the polymods listed in table 1.

The syntax to set a polymod in Fermat is &(P=polymod,1), where the variable of the polymod must be the first symbol adjoint to Fermat which does not have a polymod assigned yet. This Fermat command should be stored in a text file and can be read in by epsilon with the --fermat job.

For internal reasons Fermat runs very slow if more than one polymod is assigned to it. Therefore, in complicated cases with more than one complex number, it is useful to set a polymod only for the most frequent variable appearing in the block to work with next. As the order of symbols once set cannot be changed inside Fermat, the only way to change the 'active' polymod is by saving the system with --write and reload it into a new session after the symbols are adjoined in a different order. In our tests we were able to solve huge systems with up to three complex numbers applying this strategy.

3.6 Usage of EpsilonTools.m

The Mathematica package EpsilonTools.m provides functions which help to set up epsilon and to work with the epsilon input/output file format (see section 3.2). It is not essential in order to run epsilon. After a successful installation, the package can be loaded into a Mathematica session with

```
<<EpsilonTools'
```

EpsilonTools.m provides three functions:

• EpsilonSymRules[expression]

This function scans *expression* for symbols of the form given in table 1 and compiles a list of rules for these symbols to their corresponding Mathematica expressions, e.g.

```
{r3->(1+I*Sqrt[3])/2, q5->(1+Sqrt[5])/2, i->I}
```

• EpsilonRead[file]

This function reads *file* into Mathematica where *file* is in the format described in section 3.2.

Options:

- ReplaceSymbols (default: True)

If this option is set to True, all symbols introduced by epsilon-prepare will be replaced by their corresponding Mathematica expression.

- CheckFuchsian (default: False)

If this option is set to True, the function returns **\$Failed** if the system in *file* is not in Fuchsian form.

- CheckEpsilon (default: False)

If this option is set to True, the function returns Failed if the system in *file* is not in ϵ -form.

• EpsilonBlocks[M]

This function scans for a block-triangular structure of the matrix M and returns a list of the boundaries of all diagonal blocks. The returned boundaries are in the form $\{start, end\}$. The values start and end can be used in the --block job of epsilon.

4 A physical example

As an example, we consider a set of three-loop master integrals $\{I_j\}$ with $j = 1, \ldots, 9$ in $d = 4 - 2\epsilon$ dimensions with internal lines of mass one or zero. A graphical representation of the master integrals except I_3 is contained in fig. 4.1. Further we define I_3 as

$$I_3 = \int d^d p \int d^d l \int d^d k \, \frac{(k-l)^2}{[(p+q_1)^2 - 1][l^2 - 1][(p-l+k)^2 - 1][(l-p+q_2)^2]}$$

The kinematics is given by

$$q_1^2 = q_2^2 = 0$$
, $q_1 \cdot q_2 = -\frac{(1-x)^2}{2x}$.

The vector $\vec{f} = (I_1, \ldots, I_9)$ obeys a differential equation

$$\frac{\partial}{\partial x}\vec{f} = \mathbb{M}(x,\epsilon)\vec{f},$$



Figure 4.1: Three-loop master integrals to be solved with epsilon. The not pictured integral I_3 has the same topology as I_2 but with an additional numerator. The thick (thin) lines are massive (massless). The thin external lines carry the momenta q_1 and q_2 , while the double line carries the momentum $q_1 + q_2$.

where the 9 \times 9-matrix $\mathbb{M}(x, \epsilon)$ has the structure

The * represents any non-zero entry and we point out the block-triangular structure of the system. First, this matrix should be stored in a Mathematica compatible file matrix.m. The next step is to convert this file to the format described in section 3.2 via

```
epsilon-prepare matrix.m > matrix.dat
```

The generated file matrix.dat should read

A[r3,0]:	{{0,0,0,0,0,0,0,0,0},{0,0,0,0,0,0,0,0,0,	
A[-1,0]:	{{0,0,0,0,0,0,0,0,0},{1-ep,5-6*ep,-6+	
A[-1,1]:	$\{\{0,0,0,0,0,0,0,0,0,0\},\{0,0,0,0,0,0,0,0,$	
A[-1,2]:	$\{\{0,0,0,0,0,0,0,0,0,0\},\{0,0,0,0,0,0,0,0,$	
A[1-r3,0]:	$\{\{0,0,0,0,0,0,0,0,0,0\},\{0,0,0,0,0,0,0,0,$	
A[1,0]:	{{0,0,0,0,0,0,0,0,0},{-1+ep,-11+10*ep	
A[1,1]:	$\{\{0,0,0,0,0,0,0,0,0,0\},\{0,0,0,0,0,0,0,0,$	
A[1,2]:	{{0,0,0,0,0,0,0,0,0},{0,0,0,0,0,0,0,0,0,	
A[1,3]:	{{0,0,0,0,0,0,0,0,0},{0,0,0,0,0,0,0,0,0,	
A[1,4]:	{{0,0,0,0,0,0,0,0,0},{0,0,0,0,0,0,0,0,0,	

The arguments of A and B reveal singularities at $\{r_3, -1, 1 - r_3, 1, 0, \infty\}$ with Poincaré ranks $\{0, 2, 0, 4, 1, 1\}$, respectively. The symbol r_3 introduced by epsilon-prepare is a root of the polynomial $r_3^2 - r_3 + 1$ and is given by $r_3 = (1 + i\sqrt{3})/2$ (see table 1). Hence, we need a file enable.r3.fer to set a polynomial for Fermat to mod out with containing

 $\&(P=r3^2-r3+1,1)$

To understand the origin of r_3 we consider for example the partial fraction decomposition over the complex of $[\mathbb{M}(x, \epsilon)]_{51}$:

$$\begin{split} [\mathbb{M}(x,\epsilon)]_{51} &= \frac{(\epsilon-1)^2(1+x)\left(1+4x+x^2+\epsilon\left(1-10x+x^2\right)\right)}{4\epsilon(2\epsilon-1)(x-1)^3\left(x^2-x+1\right)} \\ &= \frac{1}{x-r_3} \left\{ \frac{5-9\epsilon^3+23\epsilon^2-19\epsilon}{4(2\epsilon^2-\epsilon)} \right\} + \frac{1}{x-(1-r_3)} \left\{ \frac{5-9\epsilon^3+23\epsilon^2-19\epsilon}{4(2\epsilon^2-\epsilon)} \right\} \\ &\quad + \frac{1}{x-1} \left\{ \frac{-5+9\epsilon^3-23\epsilon^2+19\epsilon}{2(2\epsilon^2-\epsilon)} \right\} + \frac{1}{(x-1)^2} \left\{ \frac{3-4\epsilon^3+11\epsilon^2-10\epsilon}{2(2\epsilon^2-\epsilon)} \right\} \\ &\quad + \frac{1}{(x-1)^3} \left\{ \frac{3-4\epsilon^3+11\epsilon^2-10\epsilon}{2\epsilon^2-\epsilon} \right\} \,. \end{split}$$

The possible blocks to run epsilon with can be either read off from the matrix (4.1) or determined by the function EpsilonBlocks of the Mathematica package EpsilonTools.m. In this case, the possible blocks are

$$\{1,1\},\{2,3\},\{4,4\},\{5,9\}.$$

Now we can run epsilon with the command

```
epsilon --timings --symbols r3 --load matrix.dat 1 1 \
    --queue out.queue \
    --fermat enable.r3.fer \
    --fuchsify --normalize --factorep-at -1 \
    --block 2 3 \
    --fuchsify --normalize --factorep-at -1 --left-fuchsify \
    --block 4 4 \
    --fuchsify --normalize --factorep-at -1 --left-fuchsify \
    --block 5 9 \
    --fuchsify --normalize --factorep-at 1 --left-fuchsify \
```

--block 1 9 \
--factorep-at -1 \
--write epsilon.dat --export transformation.m

The --load job in the first line loads the file matrix.dat with an active block $\{1, 1\}$; the --block jobs are used to change the active block. In all blocks except for block $\{5, 9\}$ we factor out ϵ at $\mu = -1$. For this block $\mu = -1$ would lead to a singular system so we choose $\mu = 1$ instead. The final system is written into the file epsilon.dat and the corresponding transformation matrix into transformation.m. The whole reduction process takes about three minutes on an Intel Core i5-3320M CPU with 2.60 GHz. The package EpsilonTools.m offers the function EpsilonRead for reading the file epsilon.dat into Mathematica. The result is in ϵ -form (2.4), with

			1		0		C		0	0		0	0	0	0	0	\		
$\widehat{\mathbf{n}}(-1)$	($\frac{1}{5}$	_	18		$\frac{36}{5}$	0		0	0	0	0	0					
			$\frac{1}{3}$	-	30		12	0		0	0	0	0	0					
			0		D		0	0		0	0	0	0	0					
		77	90185	-4068	887355	133	803239	-371	5319	-21	323	1	-121	26	3				
		47	43936	316	2624	26	3552	6325	248	58	58	1	58	11	6				
INI ⁰		_	1	222	799291	-5578	897095	133	303239	-8545	52337	-483	7429	23	-2783	604	19	,	
				75	902976	1265	0496	21	08416	12650	0496	116	116	2	116	23	2		
			·	-193	66519759	71280	791995	-287	0982795	672844	12709	38031	-584953	-1811	219131	-476	293		
			1	220	1186304	3668	64384	611	44064	36686	4384	3364	3364	58	3364	672	28		
				90	725693	299	6225	-27	685119	-4829	99147	-273	4199	19	-1573	341	.9		
				37	951488	89	088	10	54208	6325	248	58	58	15	58	11	6		
				-48	8778543	1264	40695	-32	235923	18576	3595	105	-1615	_5	605	-13	15		
			/	37	951488	632	5248	10	54208	6325	248	58	58	-5	58	11	₆ /		
,	/	0		0		0		0		ſ)	0	0	0		0	,	、	
	1		$^{-2}$		0		0		0			, ,	0	0	0		0		1
		-	5		26		-12		0		()	0	0	0		0		
		-	$\frac{-7}{9}$		50		-23		0		()	0	0	0		0		
	L		0		0		0		0		()	0	0	0		0		
		-303	3217	97	3565507	75	-1198	49	19535	387	111	37	-51841	-1225	3513	35	-1893	319	
$\widehat{\mathbf{M}}^{(0)}$ _		711	590	4	237196	68	2059		47439	36	55	68	1392	192	278	4	1113	36	
$\mathbf{IW}_{0} = \begin{bmatrix} -68613\\ 113854\\ 4967800\\ 330177\\ -32347\\ 56927\\ 158320\\ 569272 \end{bmatrix}$	-686	1355	525	2235183	885	291233	07	529852	429	419	699	-1585715	-37835	11472	213	-4875	557	,	
	1138	3544	64	379514	88	105420)8	37951_{-}	488	445	544	11136	1536	2227	2	8908	38		
	49678	3009	745	-2298805	59165	1437349	057	-300314	02873	-257	17127	91504967	2219903	-6570	5785	274163	3633		
	3301	7794	156	1100593	152	305720	32	1100593	3152	1291	776	322944	44544	6458	88	25835	552		
	-323	4724	451	-219383	5945	504564	29	344565	875	262	093	-1020493	$\frac{-24181}{}$	7468	67	-3104	155		
	569	2723	32	189757	44	52710	4	18975'	744	222	272	5568	768	1113	86	4454	14		
	1583	3205	29	4020933	<u>895</u>	-14262	031	-13722	7105	-10	1759	404543	9527	-290	113	13136	557	/	
	2723	32	189757	44	52710	4	18975'	744	222	272	5568	768	1113	86	4454	44	/		

	1	0	0	0	0	0	0	0	0	0	\	
$\widehat{\mathbf{M}}_{0}^{(1)} =$	1	-10	24	0	0	0	0	0	0	١	1	
	5	10	5	0	0	0	0	0	0			
	4	-20	10	0	0	0	0	0	0			
	$\frac{-}{9}$	-20	10	0	0	0	0	0	0			
	0	0	0	0	0	0	0	0	0			
	41108207	-636997435	21452971	28883609	673	-1484	-13	1651	1785			
	14231808	9487872	790656	18975744	116	87	4	174	232			
	733595729	22172065	-1318339	-334258861	-1507	5221	167	-14597	24933		,	
	227708928	4743936	109056	75902976	464	87	16	696	928			
	-41044567181	2404770185	-2068713589	20579631337	194887	-331405	-11707	952145	-1183713	3		
	_	6603558912	137574144	91716096	2201186304	13456	2523	464	20184	26912	-	
		391067287	50935825	-71070457	-281285603	-1053	7870	121	-11539	17115		
		113854464	1185984	1581312	37951488	232	87	8	348	464		
		-188522477	-28164515	18097199	43744885	-237	-1634	-23	1877	-7157		
		113854464	4743936	1581312	37951488	232	87	8	348	464	/	
		0	0	0	0	0	0	0	0	0	\	
		0	0	0	0	0	0	0	0	0	1	
		0	0	0	0	0	0	0	0	0		
		0	0	0	0	0	0	0	0	0		
		-119849	1797735	-838943	-2996225	-2425	1355	97	-1053	5839		
		889488	65888	49416	1185984	696	58	24	116	1392	İ	
$\widehat{\mathbf{n}_{I}}(r_{3})$		-3715319	55729785	-26007233	-92882975	-75175	42005	3007	-32643	181009		
$IM_0 =$	=	56927232	4216832	3162624	75902976	44544	3712	1536	7424	89088		,
		-26486629	397299435	-185406403	-662165725	-535925	299455	21437	-232713	1290419		
		1650889728	122288128	91716096	2201186304	1291776	107648	44544	215296	2583552		
		-2037433	30561495	-14262031	-50935825	-41225	23035	1649	-17901	99263		
		28463616	2108416	1581312	37951488	22272	1856	768	3712	44544		
		2277131	-34156965	15939917	56928275	46075	-25745	-1843	20007	-110941		
		28463616	2108416	1581312	37951488	22272	1856	768	3712	44544	/	
		`								,		
	0	0	0	0	0	0	0	0	0			
		0	0	0	0	0	0	0	0	0	1	
		0	0	0	0	0	0	0	0	0		
		0	0	0	0	0	0	0	0	0		
		-119849	1797735	-838943	-2996225	-2425	1355	97	-1053	5839	I	
	889488	65888	49416	1185984	696	58	24	116	1392			
$\widehat{\mathbf{M}}^{(1-r_3)}$		-3715319	55729785	-26007233	-92882975	-75175	42005	3007	-32643	181009		
TNT 0	_	56927232	4216832	3162624	75902976	44544	3712	1536	7424	89088		•
	-26486629	397299435	-185406403	-662165725	-535925	299455	21437	-232713	$\frac{1290419}{1290419}$			
	1650889728	122288128	91716096	2201186304	1291776	107648	44544	215296	2583552			
	-2037433	30561495	-14262031	-50935825	-41225	23035	1649	-17901	99263	ļ		
	28463616	2108416	1581312	37951488	22272	1856	768	3712	44544			
	2277131	-34156965	15939917	56928275	46075	-25745	-1843	20007	-110941	1		
	28463616	2108416	1581312	37951488	22272	1856	768	3712	44544	/		

5 Summary

In this paper we presented epsilon. The tool epsilon is an efficient implementation of an algorithm proposed by R.N. Lee to reduce a system of ordinary differential equations with rational coefficients to a canonical form, where the right hand side is proportional to ϵ .

In physically relevant situations, the small parameter ϵ usually is a regulator in dimensional regularization (e.g. in $d = 4 - 2\epsilon$ dimensions). We showed its applicability in

a three-loop example and demonstrated the possibility to reduce systems with complex singular points.

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The author wants to thank Robert Harlander for useful discussions and comments on the manuscript. This work was supported by BMBF contract 05H15PACC1. The computing resources were granted by RWTH Aachen University under project rwth0119. The Feynman diagrams in this article have been drawn with JaxoDraw [19] based on Axodraw [20].

A Transformations

A.1 Balances

The main building blocks of Lee's algorithm are balances. In this section we describe how they act on a system in the form (2.2) which is also assumed to be in block-triangular form (2.6).

The projector \mathbb{P} is determined by its impact on the active block \mathbb{C} . Hence, it has to be of the form

$$\mathbb{P} = \begin{pmatrix} 0 & 0 & 0 \\ 0 & \mathbb{Q} & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

Naturally, this also modifies the blocks \mathbb{B} and \mathbb{E} . In this appendix we omit writing down the ϵ -dependencies explicitly.

First, we consider a balance between two singularities $x_1, x_2 < \infty$ (2.7a):

$$\mathbb{T} = \mathcal{B}(\mathbb{P}, x_1, x_2), \quad \mathbb{T}^{-1} = \mathcal{B}(\mathbb{P}, x_2, x_1).$$

The transformation (2.5) can now be written in terms of the coefficient matrices.

We find for the active block

$$\widetilde{\mathbb{C}}_{0}^{(x_{1})} = \mathbb{C}_{0}^{(x_{1})} - \sum_{n \ge 0} \frac{1}{(x_{2} - x_{1})^{n}} \mathbb{Q} \mathbb{C}_{n}^{(x_{1})} \overline{\mathbb{Q}} + \sum_{x_{j} \in S \setminus \{x_{1}\}} \sum_{n \ge 0} \frac{x_{1} - x_{2}}{(x_{1} - x_{j})^{n+1}} \overline{\mathbb{Q}} \mathbb{C}_{n}^{(x_{j})} \mathbb{Q} + (x_{1} - x_{2}) \sum_{n \ge 0} \frac{x_{1}^{n} \overline{\mathbb{Q}} \mathbb{C}_{n} \mathbb{Q} + \mathbb{Q}, \qquad (A.1a)$$

$$\widetilde{\mathbb{C}}_{k>0}^{(x_1)} = \mathbb{C}_k^{(x_1)} + (x_1 - x_2)\overline{\mathbb{Q}}\mathbb{C}_{k-1}^{(x_1)}\mathbb{Q} - \sum_{n\geq 0} \frac{1}{(x_2 - x_1)^n} \mathbb{Q}\mathbb{C}_{n+k}^{(x_1)}\overline{\mathbb{Q}}, \qquad (A.1b)$$

$$\widetilde{\mathbb{C}}_{0}^{(x_{2})} = \mathbb{C}_{0}^{(x_{2})} - \sum_{n \ge 0} \frac{1}{(x_{1} - x_{2})^{n}} \overline{\mathbb{Q}} \mathbb{C}_{n}^{(x_{2})} \mathbb{Q} + \sum_{x_{j} \in S \setminus \{x_{2}\}} \sum_{n \ge 0} \frac{x_{2} - x_{1}}{(x_{2} - x_{j})^{n+1}} \mathbb{Q} \mathbb{C}_{n}^{(x_{j})} \overline{\mathbb{Q}}$$
(A.1c)

$$+ (x_2 - x_1) \sum_{n \ge 0} x_2^n \mathbb{Q} \mathbb{C}_n \mathbb{Q} - \mathbb{Q},$$

$$\widetilde{\mathbb{C}}_{k>0}^{(x_2)} = \mathbb{C}_k^{(x_2)} + (x_2 - x_1) \mathbb{Q} \mathbb{C}_{k-1}^{(x_2)} \overline{\mathbb{Q}} - \sum_{n \ge 0} \frac{1}{(x_1 - x_2)^n} \overline{\mathbb{Q}} \mathbb{C}_{n+k}^{(x_2)} \mathbb{Q},$$
 (A.1d)

$$\widetilde{\mathbb{C}}_{k}^{(x_{j}\neq x_{1},x_{2})} = \mathbb{C}_{k}^{(x_{j})} + \sum_{n\geq 0} \frac{x_{2}-x_{1}}{(x_{1}-x_{j})^{n+1}} \overline{\mathbb{Q}} \mathbb{C}_{n+k}^{(x_{j})} \mathbb{Q} + \sum_{n\geq 0} \frac{x_{1}-x_{2}}{(x_{2}-x_{j})^{n+1}} \mathbb{Q} \mathbb{C}_{n+k}^{(x_{j})} \overline{\mathbb{Q}}, \qquad (A.1e)$$

$$\widetilde{\mathbb{C}}_{k} = \mathbb{C}_{k} + (x_{1} - x_{2}) \sum_{n \ge 0} \left\{ x_{1}^{n} \overline{\mathbb{Q}} \mathbb{C}_{k+n+1} \mathbb{Q} - x_{2}^{n} \mathbb{Q} \mathbb{C}_{k+n+1} \overline{\mathbb{Q}} \right\} , \qquad (A.1f)$$

and for the off-diagonal blocks

$$\begin{split} \widetilde{\mathbb{B}}_{0}^{(x_{2})} &= \mathbb{B}_{0}^{(x_{2})} + \sum_{x_{j} \in S \setminus \{x_{2}\}} \sum_{n \geq 0} \frac{x_{2} - x_{1}}{(x_{2} - x_{j})^{n+1}} \mathbb{Q}\mathbb{B}_{n}^{(x_{j})} + (x_{2} - x_{1}) \sum_{n \geq 0} x_{2}^{n} \mathbb{Q}\mathbb{B}_{n} \,, \\ \widetilde{\mathbb{B}}_{k}^{(x_{2})} &= \mathbb{B}_{k}^{(x_{2})} + (x_{2} - x_{1}) \mathbb{Q}\mathbb{B}_{k-1}^{(x_{2})} \,, \\ \widetilde{\mathbb{B}}_{k}^{(x_{j} \neq x_{2})} &= \mathbb{B}_{k}^{(x_{j})} + \sum_{n \geq 0} \frac{x_{1} - x_{2}}{(x_{2} - x_{j})^{n+1}} \mathbb{Q}\mathbb{B}_{n+k}^{(x_{j})} \,, \\ \widetilde{\mathbb{B}}_{k}^{(x_{j})} &= \mathbb{E}_{k}^{(x_{1})} + \sum_{n \geq 0} \frac{x_{1} - x_{2}}{(x_{2} - x_{j})^{n+1}} \mathbb{E}_{n}^{(x_{j})}\mathbb{Q} + (x_{1} - x_{2}) \sum_{n \geq 0} x_{1}^{n}\mathbb{E}_{n}\mathbb{Q} \,, \\ \widetilde{\mathbb{E}}_{k}^{(x_{1})} &= \mathbb{E}_{k}^{(x_{1})} + (x_{1} - x_{2})\mathbb{E}_{k-1}^{(x_{1})}\mathbb{Q} \,, \\ \widetilde{\mathbb{E}}_{k}^{(x_{j} \neq x_{1})} &= \mathbb{E}_{k}^{(x_{j})} + \sum_{n \geq 0} \frac{x_{2} - x_{1}}{(x_{1} - x_{j})^{n+1}} \mathbb{E}_{n+k}^{(x_{j})}\mathbb{Q} \,, \\ \widetilde{\mathbb{E}}_{k}^{(x_{j} \neq x_{1})} &= \mathbb{E}_{k}^{(x_{j})} + \sum_{n \geq 0} \frac{x_{2} - x_{1}}{(x_{1} - x_{j})^{n+1}} \mathbb{E}_{n+k}^{(x_{j})}\mathbb{Q} \,, \\ \widetilde{\mathbb{E}}_{k}^{(x_{j} \neq x_{1})} &= \mathbb{E}_{k}^{(x_{j})} + \sum_{n \geq 0} \frac{x_{2} - x_{1}}{(x_{1} - x_{j})^{n+1}} \mathbb{E}_{n+k}^{(x_{j})}\mathbb{Q} \,. \end{split}$$

All other blocks are unaffected.

Next, we consider the case $x_1 < \infty$, $x_2 = \infty$, i.e.

$$\mathbb{T} = \mathcal{B}(\mathbb{P}, x_1, \infty), \quad \mathbb{T}^{-1} = \mathcal{B}(\mathbb{P}, \infty, x_1).$$

In that case, the balances are given by (2.7b) and (2.7c). Here the active block transforms

$$\begin{split} \widetilde{\mathbf{C}}_{0}^{(x_{1})} &= \mathbf{C}_{0}^{(x_{1})} - \overline{\mathbf{Q}} \mathbf{C}_{0}^{(x_{1})} \mathbf{Q} - \mathbf{Q} \mathbf{C}_{0}^{(x_{1})} \overline{\mathbf{Q}} + \mathbf{Q} \mathbf{C}_{1}^{(x_{1})} \overline{\mathbf{Q}} \\ &+ \sum_{x_{j} \in S \setminus \{x_{1}\}} \sum_{n=0}^{\infty} \frac{1}{(x_{1} - x_{j})^{n+1}} \overline{\mathbf{Q}} \mathbf{C}_{n}^{(x_{j})} \mathbf{Q} + \sum_{n=0}^{\infty} x_{1}^{n} \overline{\mathbf{Q}} \mathbf{C}_{n} \mathbf{Q} + \mathbf{Q} , \\ \widetilde{\mathbf{C}}_{k>0}^{(x_{1})} &= \mathbf{C}_{k}^{(x_{1})} - \overline{\mathbf{Q}} \mathbf{C}_{k}^{(x_{1})} \mathbf{Q} - \mathbf{Q} \mathbf{C}_{k}^{(x_{1})} \overline{\mathbf{Q}} + \mathbf{Q} \mathbf{C}_{k+1}^{(x_{1})} \overline{\mathbf{Q}} + \overline{\mathbf{Q}} \mathbf{C}_{k-1}^{(x_{1})} \mathbf{Q} , \\ \widetilde{\mathbf{C}}_{k}^{(x_{j} \neq x_{1})} &= \mathbf{C}_{k}^{(x_{j})} - \overline{\mathbf{Q}} \mathbf{C}_{k}^{(x_{j})} \mathbf{Q} - \mathbf{Q} \mathbf{C}_{k}^{(x_{j})} \overline{\mathbf{Q}} + \mathbf{Q} \mathbf{C}_{k+1}^{(x_{j})} \overline{\mathbf{Q}} + (x_{j} - x_{1}) \mathbf{Q} \mathbf{C}_{k}^{(x_{j})} \overline{\mathbf{Q}} \\ &- \sum_{n=0}^{\infty} \frac{1}{(x_{1} - x_{j})^{n+1}} \overline{\mathbf{Q}} \mathbf{C}_{n+k}^{(x_{j})} \mathbf{Q} , \\ \widetilde{\mathbf{C}}_{0} &= \mathbf{C}_{0} - \overline{\mathbf{Q}} \mathbf{C}_{0} \mathbf{Q} - (1 + x_{1}) \mathbf{Q} \mathbf{C}_{0} \overline{\mathbf{Q}} + \sum_{x_{j} \in S} \mathbf{Q} \mathbf{C}_{0}^{(x_{j})} \overline{\mathbf{Q}} + \sum_{n=0}^{\infty} x_{1}^{n} \overline{\mathbf{Q}} \mathbf{C}_{n+1} \mathbf{Q} , \\ \widetilde{\mathbf{C}}_{k>0} &= \mathbf{C}_{k} - \overline{\mathbf{Q}} \mathbf{C}_{k} \mathbf{Q} - (1 + x_{1}) \mathbf{Q} \mathbf{C}_{k} \overline{\mathbf{Q}} + \mathbf{Q} \mathbf{C}_{k-1} \overline{\mathbf{Q}} + \sum_{n=0}^{\infty} x_{1}^{n} \overline{\mathbf{Q}} \mathbf{C}_{k+n+1} \mathbf{Q} , \end{split}$$

and the blocks $\mathbb B$ and $\mathbb E$ as

$$\begin{split} \widetilde{\mathbf{B}}_{k}^{(x_{j})} &= \mathbf{B}_{k}^{(x_{j})} - \mathbf{Q}\mathbf{B}_{k}^{(x_{j})} + \mathbf{Q}\mathbf{B}_{k+1}^{(x_{j})} + (x_{j} - x_{1})\mathbf{Q}\mathbf{B}_{k}^{(x_{j})}, \\ \widetilde{\mathbf{B}}_{0} &= \mathbf{B}_{0} - (1 + x_{1})\mathbf{Q}\mathbf{B}_{0} + \sum_{x_{j} \in S} \mathbf{Q}\mathbf{B}_{0}^{(x_{j})}, \\ \widetilde{\mathbf{B}}_{k>0} &= \mathbf{B}_{k} - (1 + x_{1})\mathbf{Q}\mathbf{B}_{k} + \mathbf{Q}\mathbf{B}_{k-1}, \\ \widetilde{\mathbf{E}}_{0}^{(x_{1})} &= \mathbf{E}_{0}^{(x_{1})} - \mathbf{E}_{0}^{(x_{1})}\mathbf{Q} + \sum_{x_{j} \in S \setminus \{x_{1}\}} \sum_{n \geq 0} \frac{1}{(x_{1} - x_{j})^{n+1}} \mathbf{E}_{n}^{(x_{j})}\mathbf{Q} + \sum_{n \geq 0} x_{1}^{n}\mathbf{E}_{n}\mathbf{Q}, \\ \widetilde{\mathbf{E}}_{k>0}^{(x_{1})} &= \mathbf{E}_{k}^{(x_{1})} - \mathbf{E}_{k}^{(x_{1})}\mathbf{Q} + \mathbf{E}_{k-1}^{(x_{1})}\mathbf{Q}, \\ \widetilde{\mathbf{E}}_{k}^{(x_{j} \neq x_{1})} &= \mathbf{E}_{k}^{(x_{j})} - \mathbf{E}_{k}^{(x_{j})}\mathbf{Q} - \sum_{n \geq 0} \frac{1}{(x_{1} - x_{j})^{n+1}} \mathbf{E}_{n+k}^{(x_{j})}\mathbf{Q}, \\ \widetilde{\mathbf{E}}_{k} &= \mathbf{E}_{k} - \mathbf{E}_{k}\mathbf{Q} + \sum_{n \geq 0} x_{1}^{n}\mathbf{E}_{k+n+1}\mathbf{Q}. \end{split}$$

Finally, we consider the case $x_1 = \infty, x_2 < \infty$, i.e.

$$\mathbb{T} = \mathcal{B}(\mathbb{P}, \infty, x_2), \quad \mathbb{T}^{-1} = \mathcal{B}(\mathbb{P}, x_2, \infty).$$

as

This is similar to the previous case. The active block transforms to

$$\begin{split} \widetilde{\mathbf{C}}_{0}^{(x_{2})} &= \mathbf{C}_{0}^{(x_{2})} - \overline{\mathbf{Q}} \mathbf{C}_{0}^{(x_{2})} \mathbf{Q} - \mathbf{Q} \mathbf{C}_{0}^{(x_{2})} \overline{\mathbf{Q}} + \overline{\mathbf{Q}} \mathbf{C}_{1}^{(x_{2})} \mathbf{Q} \\ &+ \sum_{x_{j} \in S \setminus \{x_{2}\}} \sum_{n \geq 0} \frac{1}{(x_{2} - x_{j})^{n+1}} \mathbf{Q} \mathbf{C}_{n}^{(x_{j})} \overline{\mathbf{Q}} + \sum_{n \geq 0} x_{2}^{n} \mathbf{Q} \mathbf{C}_{n} \overline{\mathbf{Q}} - \mathbf{Q} \,, \\ \widetilde{\mathbf{C}}_{k>0}^{(x_{2})} &= \mathbf{C}_{k}^{(x_{2})} - \overline{\mathbf{Q}} \mathbf{C}_{k}^{(x_{2})} \mathbf{Q} - \mathbf{Q} \mathbf{C}_{k}^{(x_{2})} \overline{\mathbf{Q}} + \mathbf{Q} \mathbf{C}_{k-1}^{(x_{2})} \overline{\mathbf{Q}} + \overline{\mathbf{Q}} \mathbf{C}_{k+1}^{(x_{2})} \mathbf{Q} \,, \\ \widetilde{\mathbf{C}}_{k}^{(x_{j} \neq x_{2})} &= \mathbf{C}_{k}^{(x_{j})} - \overline{\mathbf{Q}} \mathbf{C}_{k}^{(x_{j})} \mathbf{Q} - \mathbf{Q} \mathbf{C}_{k}^{(x_{j})} \overline{\mathbf{Q}} + \overline{\mathbf{Q}} \mathbf{C}_{k+1}^{(x_{j})} \mathbf{Q} + (x_{j} - x_{2}) \overline{\mathbf{Q}} \mathbf{C}_{k}^{(x_{j})} \mathbf{Q} \\ &- \sum_{n \geq 0} \frac{1}{(x_{2} - x_{j})^{n+1}} \mathbf{Q} \mathbf{C}_{n+k}^{(x_{j})} \overline{\mathbf{Q}} \,, \\ \widetilde{\mathbf{C}}_{0} &= \mathbf{C}_{0} - \overline{\mathbf{Q}} \mathbf{C}_{0} \mathbf{Q} - \mathbf{Q} \mathbf{C}_{0} \overline{\mathbf{Q}} + \sum_{n \geq 0} x_{2}^{n} \mathbf{Q} \mathbf{C}_{n+1} \overline{\mathbf{Q}} - x_{2} \overline{\mathbf{Q}} \mathbf{C}_{0} \mathbf{Q} + \sum_{x_{j} \in S} \overline{\mathbf{Q}} \mathbf{C}_{0}^{(x_{j})} \mathbf{Q} \,, \\ \widetilde{\mathbf{C}}_{k>0} &= \mathbf{C}_{k} - \overline{\mathbf{Q}} \mathbf{C}_{k} \mathbf{Q} - \mathbf{Q} \mathbf{C}_{k} \overline{\mathbf{Q}} + \sum_{n \geq 0} x_{2}^{n} \mathbf{Q} \mathbf{C}_{k+n+1} \overline{\mathbf{Q}} + \overline{\mathbf{Q}} \mathbf{C}_{k-1} \mathbf{Q} - x_{2} \overline{\mathbf{Q}} \mathbf{C}_{k} \mathbf{Q} \,, \end{split}$$

and the blocks $\mathbb B$ and $\mathbb E$ transform as

$$\begin{split} \widetilde{\mathbb{B}}_{0}^{(x_{2})} &= \mathbb{B}_{0}^{(x_{2})} - \mathbb{Q}\mathbb{B}_{0}^{(x_{2})} + \sum_{x_{j} \in S \setminus \{x_{2}\}} \sum_{n \geq 0} \frac{1}{(x_{2} - x_{j})^{n+1}} \mathbb{Q}\mathbb{B}_{n}^{(x_{j})} + \sum_{n \geq 0} x_{2}^{n} \mathbb{Q}\mathbb{B}_{n}, \\ \widetilde{\mathbb{B}}_{k>0}^{(x_{2})} &= \mathbb{B}_{k}^{(x_{2})} - \mathbb{Q}\mathbb{B}_{k}^{(x_{2})} + \mathbb{Q}\mathbb{B}_{k-1}^{(x_{2})}, \\ \widetilde{\mathbb{B}}_{k}^{(x_{j} \neq x_{2})} &= \mathbb{B}_{k}^{(x_{j})} - \mathbb{Q}\mathbb{B}_{k}^{(x_{j})} - \sum_{n \geq 0} \frac{1}{(x_{2} - x_{j})^{n+1}} \mathbb{Q}\mathbb{B}_{n+k}^{(x_{j})}, \\ \widetilde{\mathbb{B}}_{k} &= \mathbb{B}_{k} - \mathbb{Q}\mathbb{B}_{k} + \sum_{n \geq 0} x_{2}^{n} \mathbb{Q}\mathbb{B}_{k+n+1}, \\ \widetilde{\mathbb{B}}_{k} &= \mathbb{E}_{k}^{(x_{j})} - \mathbb{E}_{k}^{(x_{j})} \mathbb{Q} + \mathbb{E}_{k+1}^{(x_{j})} \mathbb{Q} + (x_{j} - x_{2})\mathbb{E}_{k}^{(x_{j})} \mathbb{Q}, \\ \widetilde{\mathbb{E}}_{0} &= \mathbb{E}_{0} - (1 + x_{2})\mathbb{E}_{0}\mathbb{Q} + \sum_{x_{j} \in S} \mathbb{E}_{0}^{(x_{j})}\mathbb{Q}, \\ \widetilde{\mathbb{E}}_{k>0} &= \mathbb{E}_{k} - (1 + x_{2})\mathbb{E}_{k}\mathbb{Q} + \mathbb{E}_{k-1}\mathbb{Q}. \end{split}$$

A.2 Fuchsification of off-diagonal blocks

In this appendix we consider the transformation of the off-diagonal block \mathbb{B} to Fuchsian form. We assume that the diagonal blocks \mathbb{A} and \mathbb{C} are already in ϵ -form. The transformation needed has the form

$$\mathbb{T} = \mathcal{L}(x_1, k, \mathbb{G}), \quad \mathbb{T}^{-1} = \mathcal{L}(x_1, k, -\mathbb{G}), \qquad (A.2)$$

where we used the definitions (2.8).

In addition to block $\mathbb B$ only block $\mathbb D$ is influenced by this transformation, i.e. the blocks \mathbbm{A} and \mathbbm{C} are unaffected.

The transformation (2.5), with T defined as in (A.2) translates to rules for the coefficient matrices. For $x_1 < \infty$ we find

$$\widetilde{\mathbb{B}}_{k}^{(x_{1})} = \mathbb{B}_{k}^{(x_{1})} + \mathbb{C}_{0}^{(x_{1})}\widehat{\mathbb{G}} - \widehat{\mathbb{G}}\mathbb{A}_{0}^{(x_{1})} + k\widehat{\mathbb{G}}, \qquad (A.3a)$$

$$\widetilde{\mathbb{B}}_{n< k}^{(x_1)} = \mathbb{B}_n^{(x_1)} - \sum_{x_j \in S \setminus \{x_1\}} \frac{\mathbb{C}_0^{(x_j)} \mathbb{G} - \mathbb{G} \mathbb{A}_0^{(x_j)}}{(x_j - x_1)^{k-n}}, \qquad (A.3b)$$

$$\widetilde{\mathbb{B}}_{n>k}^{(x_1)} = \mathbb{B}_n^{(x_1)}, \qquad (A.3c)$$

$$\widetilde{\mathbb{B}}_{0}^{(x_{j}\neq x_{1})} = \mathbb{B}_{0}^{(x_{j})} + \frac{\mathbb{C}_{0}^{(x_{j})}\mathbb{G} - \mathbb{G}A_{0}^{(x_{j})}}{(x_{j} - x_{1})^{k}}, \qquad (A.3d)$$

$$\widetilde{\mathbf{B}}_{n>0}^{(x_j \neq x_1)} = \mathbf{B}_n^{(x_j)},$$

$$\widetilde{\mathbf{B}}_n = \mathbf{B}_n,$$
(A.3e)
(A.3f)

$$\mathbf{B}_n = \mathbf{B}_n \,, \tag{A.3f}$$

and

$$\begin{split} \widetilde{\mathbb{D}}_{n < k}^{(x_{1})} &= \mathbb{D}_{n}^{(x_{1})} - \sum_{x_{j} \in S \setminus \{x_{1}\}} \sum_{i \geq 0} (-1)^{i} \frac{\binom{k+i-n-1}{i}}{(x_{j} - x_{1})^{k+i-n}} \mathbb{E}_{i}^{(x_{j})} \widehat{\mathbb{G}} \\ &+ \sum_{i \geq 0} x_{1}^{i} \binom{i+k-n-1}{i} \mathbb{E}_{i+k-n-1} \widehat{\mathbb{G}} , \\ \widetilde{\mathbb{D}}_{n \geq k}^{(x_{1})} &= \mathbb{D}_{n}^{(x_{1})} + \mathbb{E}_{n-k}^{(x_{1})} \widehat{\mathbb{G}} , \\ \widetilde{\mathbb{D}}_{n}^{(x_{j \neq 1})} &= \mathbb{D}_{n}^{(x_{j})} + (-1)^{k} \sum_{i \geq 0} \frac{\binom{k+i-1}{i}}{(x_{1} - x_{j})^{k+i}} \mathbb{E}_{n+i}^{(x_{j})} \widehat{\mathbb{G}} , \\ \widetilde{\mathbb{D}}_{n}^{(x_{j \neq 1})} &= \mathbb{D}_{n} + \sum_{m \geq 0} (-1)^{m} \sum_{i \geq 0} x_{1}^{m+i} \binom{n+m}{n} \binom{i+n+m+k}{i} \mathbb{E}_{i+n+m+k} \widehat{\mathbb{G}} . \end{split}$$

The case $x_1 = \infty$ leads to

$$\widetilde{\mathbb{B}}_{k-1} = \mathbb{B}_{k-1} + \sum_{x_j \in S} \left[\mathbb{C}_0^{(x_j)} \mathbb{G} - \widehat{\mathbb{G}} \mathbb{A}^{(x_j)} \right] - k \mathbb{G} ,$$

$$\widetilde{\mathbb{B}}_{n < k-1} = \mathbb{B}_n + \sum_{x_j \in S} x_j^{k-n-1} \left[\mathbb{C}_0^{(x_j)} \widehat{\mathbb{G}} - \widehat{\mathbb{G}} \mathbb{A}_0^{(x_j)} \right] ,$$

$$\widetilde{\mathbb{B}}_{n > k-1} = \mathbb{B}_n ,$$

$$\widetilde{\mathbb{B}}_0^{(x_j)} = \mathbb{B}_0^{(x_j)} + x_j^k \left[\mathbb{C}_0^{(x_j)} \widehat{\mathbb{G}} - \widehat{\mathbb{G}} \mathbb{A}_0^{(x_j)} \right] ,$$

$$\widetilde{\mathbb{B}}_{n > 0}^{(x_j)} = \mathbb{B}_n^{(x_j)} ,$$

and

$$\begin{split} \widetilde{\mathbb{D}}_{n}^{(x_{j})} &= \mathbb{D}_{n}^{(x_{j})} + \sum_{i=0}^{k} x_{j}^{i} \binom{k}{k-i} \mathbb{E}_{n+k-i}^{(x_{j})} \widehat{\mathbb{G}} ,\\ \widetilde{\mathbb{D}}_{n < k} &= \mathbb{D}_{n} + \sum_{x_{j} \in S} \sum_{m=0}^{k-n-1} \sum_{i=0}^{m} (-1)^{k-n-m-1} x_{j}^{k-n-i-1} \binom{k-m-1}{n} \binom{k}{k+i-m} \mathbb{E}_{i}^{(x_{j})} \widehat{\mathbb{G}} ,\\ \widetilde{\mathbb{D}}_{n \geq k} &= \mathbb{D}_{n} + \mathbb{E}_{n-k} \widehat{\mathbb{G}} . \end{split}$$

A.3 ϵ -Factorization

The transformation required in the ϵ -factorization is x-independent. Hence, every coefficient matrix in (2.2) transforms the same and the transformation rule (2.5) becomes a similarity transformation

$$\widetilde{\mathbb{M}}(x,\epsilon) = \mathbb{T}^{-1}(\epsilon)\mathbb{M}(x,\epsilon)\mathbb{T}(\epsilon) \,.$$

The matrices $\mathbb{T}(\epsilon)$ and $\mathbb{T}^{-1}(\epsilon)$ have the form

$$\mathbb{T}(\epsilon) = \begin{pmatrix} \mathbb{1} & 0 & 0 \\ 0 & \widehat{\mathbb{T}}(\epsilon) & 0 \\ 0 & 0 & 1 \end{pmatrix}, \quad \mathbb{T}^{-1}(\epsilon) = \begin{pmatrix} \mathbb{1} & 0 & 0 \\ 0 & \widehat{\mathbb{T}}^{-1}(\epsilon) & 0 \\ 0 & 0 & 1 \end{pmatrix},$$

with $\widehat{\mathbb{T}}(\epsilon)$ and $\widehat{\mathbb{T}}^{-1}(\epsilon)$ corresponding to block \mathbb{C} . Using the block-triangular structure (2.6) yields

$$\begin{pmatrix} \widetilde{\mathbb{A}}(x,\epsilon) & 0 & 0\\ \widetilde{\mathbb{B}}(x,\epsilon) & \widetilde{\mathbb{C}}(x,\epsilon) & 0\\ \widetilde{\mathbb{D}}(x,\epsilon) & \widetilde{\mathbb{E}}(x,\epsilon) & \widetilde{\mathbb{F}}(x,\epsilon) \end{pmatrix} = \begin{pmatrix} \mathbb{A}(x,\epsilon) & 0 & 0\\ \mathbb{T}^{-1}(\epsilon)\mathbb{B}(x,\epsilon) & \mathbb{T}^{-1}(\epsilon)\mathbb{C}(x,\epsilon)\mathbb{T}(\epsilon) & 0\\ \mathbb{D}(x,\epsilon) & \mathbb{E}(x,\epsilon)\mathbb{T}(\epsilon) & \mathbb{F}(x,\epsilon) \end{pmatrix} .$$

Thus, besides block \mathbb{C} , block \mathbb{B} and block \mathbb{E} are influenced as well.

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