

# Conjugates of characteristic Sturmian words generated by morphisms

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## Abstract

This article is concerned with characteristic Sturmian words of slope  $\alpha$  and  $1 - \alpha$  (denoted by  $c_\alpha$  and  $c_{1-\alpha}$  respectively), where  $\alpha \in (0, 1)$  is an irrational number such that  $\alpha = [0; 1 + d_1, \overline{d_2, \dots, d_n}]$  with  $d_n \geq d_1 \geq 1$ . It is known that both  $c_\alpha$  and  $c_{1-\alpha}$  are fixed points of non-trivial (standard) morphisms  $\sigma$  and  $\hat{\sigma}$ , respectively, if and only if  $\alpha$  has a continued fraction expansion as above. Accordingly, such words  $c_\alpha$  and  $c_{1-\alpha}$  are *generated* by the respective morphisms  $\sigma$  and  $\hat{\sigma}$ . For the particular case when  $\alpha = [0; 2, \overline{r}]$  ( $r \geq 1$ ), we give a decomposition of each *conjugate* of  $c_\alpha$  (and hence  $c_{1-\alpha}$ ) into *generalized adjoining singular words*, by considering conjugates of powers of the standard morphism  $\sigma$  by which it is generated. This extends a recent result of Levé and Séébold on conjugates of the infinite Fibonacci word.

**Keywords:** combinatorics on words; characteristic Sturmian words; conjugation; Sturmian morphisms; standard morphisms; singular words.

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## 1 Introduction

In recent years, combinatorial properties of finite and infinite words have become significantly important in fields of physics, biology, mathematics, and computer science. In particular, the fascinating family of Sturmian words has been the subject of many papers (see [2, 13, 11], for example). These words, which represent the simplest family of quasi-crystals, have numerous applications in various fields of mathematics, such as symbolic dynamics, the study of continued fraction expansion, and also in some domains of physics (crystallography) and computer science (formal language theory, algorithms on words, pattern recognition).

Sturmian words are (aperiodic) infinite words with exactly  $n + 1$  distinct factors of length  $n$ , for each  $n \in \mathbb{N}$ . Since this implies that a Sturmian word has exactly two factors of length 1, then any such word is over a two-letter alphabet, say  $\mathcal{A} = \{a, b\}$ . There are many characterizations and numerous properties of Sturmian words. For a comprehensive study of the basic properties of Sturmian words, and of their transformations by morphisms, see Berstel and Séébold [2].

Here, an *infinite word*  $x$  over an alphabet  $\mathcal{A}$  is a map  $x : \mathbb{N} \rightarrow \mathcal{A}$ . For any  $i \geq 0$ , we set  $x_i = x(i)$  and write  $x = x_0x_1x_2\cdots$ , where each  $x_i \in \mathcal{A}$ . In this paper, we will utilize the following characterization of Sturmian words, which was originally proved by Morse and Hedlund [13]. An infinite word  $s$  over  $\mathcal{A} = \{a, b\}$  is Sturmian if and only if there exists an irrational  $\alpha \in (0, 1)$ , and a real number  $\rho$ , such that  $s$  is one of the following two infinite words:

$$s_{\alpha, \rho}, s'_{\alpha, \rho} : \mathbb{N} \rightarrow \mathcal{A}$$

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defined by

$$s_{\alpha,\rho}(n) = \begin{cases} a & \text{if } \lfloor (n+1)\alpha + \rho \rfloor - \lfloor n\alpha + \rho \rfloor = 0, \\ b & \text{otherwise;} \end{cases} \quad (n \geq 0)$$

$$s'_{\alpha,\rho}(n) = \begin{cases} a & \text{if } \lceil (n+1)\alpha + \rho \rceil - \lceil n\alpha + \rho \rceil = 0, \\ b & \text{otherwise.} \end{cases}$$

The irrational  $\alpha$  is called the *slope* of  $s$  and  $\rho$  is the *intercept*. If  $\rho = 0$ , we have

$$s_{\alpha,0} = ac_{\alpha} \text{ and } s'_{\alpha,0} = bc_{\alpha},$$

where  $c_{\alpha}$  is called the *characteristic Sturmian word* of slope  $\alpha$ .

The infinite Fibonacci word  $f$  is a special example of a characteristic Sturmian word of slope  $\alpha = [0; 2, \bar{1}] = (3 - \sqrt{5})/2$ , which is generated by a (standard) morphism. Wen and Wen [15] have established a factorization of the Fibonacci word into *singular words* and, in a similar fashion, Melançon [11] has proposed a generalization of singular words over a two-letter alphabet that allows for a decomposition of all the characteristic Sturmian words. More recently, Levé and Séébold [9] have obtained a generalization of Wen and Wen's 'singular' decomposition of the Fibonacci word, by establishing a similar decomposition for each *conjugate* of this infinite word into what they call *generalized singular words*. The aim of this current paper is to extend this latter result to any characteristic Sturmian word of slope  $\alpha = [0; 2, \bar{r}]$  (resp.  $1 - \alpha = [0; 1, 1, \bar{r}]$ ),  $r \geq 1$ , which we will show is generated by a particular *standard* morphism  $\sigma$  (resp.  $\hat{\sigma}$ ).

This paper is organized in the following manner. In Section 2, after recalling some combinatorial notions used in the study of words and morphisms (§2.1), we will consider right conjugation of standard morphisms (§2.2). We shall then discuss characteristic Sturmian words  $c_{\alpha}$  and a 'singular' decomposition of such words, which we will later generalize to each conjugate of  $c_{\alpha}$  for particular  $\alpha$ . In the section to follow (§3), we describe all irrationals  $\alpha \in (0, 1)$  such that  $c_{\alpha}$  is generated by a morphism, and subsequently obtain generalizations of Levé and Séébold's [9] results (on conjugates of the Fibonacci word) for  $c_{\alpha}$  and  $c_{1-\alpha}$  with  $\alpha = [0; 2, \bar{r}]$ .

## 2 Preliminaries

### 2.1 Words and Morphisms

Any of the following terminology that is not further clarified can be found in either [10] or [2], which give more detailed presentations.

Let  $\mathcal{A}$  be a finite set of symbols that we shall call an *alphabet*, the elements of which are called *letters*. A (finite) *word* is an element of the *free monoid*  $\mathcal{A}^*$  generated by  $\mathcal{A}$ , in the sense of concatenation. The identity  $\varepsilon$  of  $\mathcal{A}^*$  is called the *empty word*, and the *free semi-group*, denoted by  $\mathcal{A}^+$ , is defined by  $\mathcal{A}^+ := \mathcal{A}^* \setminus \{\varepsilon\}$ . We denote by  $\mathcal{A}^{\omega}$  the set of all infinite words over  $\mathcal{A}$ , and define  $\mathcal{A}^{\infty} := \mathcal{A}^* \cup \mathcal{A}^{\omega}$ .

A finite word  $w$  is a *factor* of  $x \in \mathcal{A}^{\infty}$  if  $x = uwv$  for some  $u \in \mathcal{A}^*$  and  $v \in \mathcal{A}^{\infty}$ . Furthermore,  $w$  is called a *prefix* (resp. *suffix*) of  $x$  if  $u = \varepsilon$  (resp.  $v = \varepsilon$ ). An infinite word  $z \in \mathcal{A}^{\omega}$  is called a *suffix* of  $x \in \mathcal{A}^{\omega}$  if there is a word  $w \in \mathcal{A}^*$  such that  $x = wz$ . The *length*  $|w|$  of a finite word  $w$  is defined to be the number of letters it contains. (Note that  $|\varepsilon| = 0$ .)

The *inverse* of  $w \in \mathcal{A}^*$ , written  $w^{-1}$ , is defined by  $ww^{-1} = w^{-1}w = \varepsilon$ . It must be emphasized that this is merely notation, i.e. for  $u, v, w \in \mathcal{A}^*$ , the words  $u^{-1}w$  and  $wv^{-1}$  are defined only if  $u$  (resp.  $v$ ) is a prefix (resp. suffix) of  $w$ . Also note that if  $w = uv$  then  $wv^{-1} = u$  and  $u^{-1}w = v$ , and if  $x = wx'$ , where  $w \in \mathcal{A}^*$  and  $x' \in \mathcal{A}^{\omega}$ , then  $w^{-1}x = x'$ .

Two words  $w, z \in \mathcal{A}^*$  are said to be *conjugate* if there exist words  $u, v$  such that  $w = uv$  and  $z = vu$ . If  $|u| = k$ , then  $z$  is called the  $k$ -th conjugate of  $w$ . This notion extends to infinite words as follows. For  $k \in \mathbb{N}$ , the  $k$ -th *conjugate* of an infinite word  $x$  over  $\mathcal{A}$  is the infinite word  $x'$  such that  $x = ux'$ , where  $u \in \mathcal{A}^*$  and  $|u| = k$ .

A *morphism* on  $\mathcal{A}$  is a map  $\psi : \mathcal{A}^* \longrightarrow \mathcal{A}^*$  such that  $\psi(uv) = \psi(u)\psi(v)$  for all  $u, v \in \mathcal{A}^*$ . It is uniquely determined by its image on the alphabet  $\mathcal{A}$ . If  $\psi(c) = cw$ , for some letter  $c \in \mathcal{A}$  and some  $w \in \mathcal{A}^+$ , then  $\psi$  is said to be *prolongable* on  $c$ . In this case, the word  $\psi^n(c)$  is a proper prefix of the word  $\psi^{n+1}(c)$  for each  $n \in \mathbb{N}$ , and the sequence  $(\psi^n(c))_{n \geq 0}$  converges to a unique infinite word

$$x = \lim_{n \rightarrow \infty} \psi^n(c) = \psi^\omega(c).$$

An infinite word  $x$  is *generated by a morphism* if  $x = \psi^\omega(c)$  for some letter  $c$  and some morphism  $\psi$ .

In what follows, it is assumed that all words are over the two-letter alphabet  $\mathcal{A} = \{a, b\}$ .

## 2.2 Conjugation of Standard Morphisms

Define on  $\mathcal{A}$  the following three morphisms

$$E : \begin{array}{l} a \mapsto b \\ b \mapsto a \end{array}, \quad \varphi : \begin{array}{l} a \mapsto ab \\ b \mapsto a \end{array}, \quad \tilde{\varphi} : \begin{array}{l} a \mapsto ba \\ b \mapsto a \end{array}.$$

A morphism  $\psi$  is *Sturmian* if and only if  $\psi \in \{E, \varphi, \tilde{\varphi}\}^*$ , i.e. if and only if it is a composition of  $E$ ,  $\varphi$ , and  $\tilde{\varphi}$  in any number and order (see [12]). Furthermore, a morphism  $\psi$  is *standard* if and only if  $\psi \in \{E, \varphi\}^*$  (see [6]). Note that a morphism is *non-trivial* if it is neither  $E$  nor  $Id_{\mathcal{A}}$  (the identity morphism).

Suppose  $\psi$  and  $\xi$  are morphisms on  $\mathcal{A}$ . If there exists a word  $u$  such that

$$\psi(w)u = u\xi(w) \quad \text{for all words } w \in \mathcal{A}^*,$$

then  $\xi$  is called the  $|u|$ -th *right conjugate* of  $\psi$ , denoted  $\psi_{|u|}$ .

It has been shown by Séébold [14] that the number of distinct right conjugates of a standard morphism  $\psi$  is  $|\psi(ab)| - 1$ ; namely, the morphisms  $\psi_0$  to  $\psi_{|\psi(ab)|-2}$ . The following useful lemma is proved in [9].

**Lemma 2.1.** *Suppose the infinite word  $x$  is generated by the standard morphism  $\psi$ . Let  $k \in \mathbb{N}$  with  $0 \leq k \leq |\psi(ab)| - 2$ , and let  $v$  denote the prefix of  $x$  of length  $k$ . Then  $\psi_k$  is such that  $\psi_k(x) = v^{-1}x$ .  $\square$*

Thus, if  $\psi$  is a standard morphism that generates an infinite word  $x$ , one deduces from Lemma 2.1 that the result of applying  $\psi_k$  to  $x$  simply consists of deleting the first  $k$  letters of  $x$ , i.e.  $\psi_k(x)$  is the  $k$ -th conjugate of  $x$ .

## 2.3 Characteristic Sturmian Words $c_\alpha$ and Singular Words

Note that every irrational  $\alpha \in (0, 1)$  has a unique continued fraction expansion

$$\alpha = [0; a_1, a_2, a_3, \dots] = \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \dots}}},$$

where each  $a_i$  is a positive integer. If the sequence  $(a_i)_{i \geq 1}$  is eventually periodic, with  $a_i = a_{i+m}$  for all  $i \geq n$ , we use the notation  $\alpha = [0; a_1, a_2, \dots, a_{n-1}, \overline{a_n, a_{n+1}, \dots, a_{n+m-1}}]$ . The  $n$ -th *convergent* of  $\alpha$  is defined by

$$\frac{p_n}{q_n} = [0; a_1, a_2, \dots, a_n] \quad \text{for all } n \geq 1,$$

where the sequences  $(p_n)_{n \geq 0}$  and  $(q_n)_{n \geq 0}$  are given by

$$\begin{aligned} p_0 &= 0, & p_1 &= 1, & p_n &= a_n p_{n-1} + p_{n-2}, & n &\geq 2; \\ q_0 &= 1, & q_1 &= a_1, & q_n &= a_n q_{n-1} + q_{n-2}, & n &\geq 2. \end{aligned}$$

Suppose  $\alpha = [0; 1 + d_1, d_2, d_3, \dots]$ , with  $d_1 \geq 0$  and all other  $d_n > 0$ . To the *directive sequence*  $(d_1, d_2, d_3, \dots)$ , we associate a sequence  $(s_n)_{n \geq -1}$  of words defined by

$$s_{-1} = b, \quad s_0 = a, \quad s_n = s_{n-1}^{d_n} s_{n-2}; \quad n \geq 1.$$

Such a sequence of words is called a *standard sequence*, and we have

$$|s_n| = q_n \quad \text{for all } n \geq 0.$$

Note that  $ab$  is a suffix of  $s_{2n-1}$  and  $ba$  is a suffix of  $s_{2n}$ , for all  $n \geq 1$ .

Standard sequences are related to characteristic Sturmian words in the following way. Observe that, for any  $n \geq 0$ ,  $s_n$  is a prefix of  $s_{n+1}$ , which gives obvious meaning to  $\lim_{n \rightarrow \infty} s_n$  as an infinite word. In fact, one can prove [7, 3] that each  $s_n$  is a prefix of  $c_\alpha$ , and we have

$$c_\alpha = \lim_{n \rightarrow \infty} s_n.$$

### 2.3.1 Singular Decomposition of $c_\alpha$

Melançon [11] (also see [4]) has proposed a generalization of Wen and Wen's [15] singular factors of the Fibonacci word to the case of any characteristic Sturmian word  $c_\alpha$  and, in doing so, has established a decomposition of  $c_\alpha$  into *adjoining singular words*, as shown below.

For  $c_\alpha$  such that  $\alpha = [0; 1 + d_1, d_2, d_3, \dots]$  with  $d_1 \geq 1$ , Melançon [11] introduced the singular words  $w_n$  of  $c_\alpha$  defined by

$$w_n = \begin{cases} as_nb^{-1} & \text{if } n \text{ is odd,} \\ bs_na^{-1} & \text{otherwise,} \end{cases}$$

for  $n \geq 1$ , with  $w_{-2} = \varepsilon, w_{-1} = a, w_0 = b$ . Furthermore, the following words  $v_n$  are also defined in [11]. For all  $n \geq -1$ ,

$$v_n = \begin{cases} as_{n+1}^{d_{n+2}-1} s_n b^{-1} & \text{if } n \text{ is odd,} \\ bs_{n+1}^{d_{n+2}-1} s_n a^{-1} & \text{otherwise.} \end{cases}$$

Clearly, the word  $v_n$  differs from  $w_{n+2}$  by a factor  $s_{n+1}$ , and it can be proved that all  $v_n$  and  $w_n$  are palindromes (i.e. words that read the same backwards as forwards). We shall call  $v_n$  the *n-th adjoining singular word* of  $c_\alpha$ , and set  $v_{-2} = \varepsilon$ . In terms of the singular and adjoining singular words of  $c_\alpha$ , the following generalization of Wen and Wen's [15] singular decomposition of the Fibonacci word has been established.

**Theorem 2.2.** [11]  $c_\alpha = \prod_{j=-1}^{\infty} (v_{2j} w_{2j+1})^{d_{2j+3}} = \prod_{j=-1}^{\infty} v_j$ . □

In the next section, for the case  $\alpha = [0; 2, \overline{r}]$ , we will generalize this factorization of  $c_\alpha$  (and hence  $c_{1-\alpha}$ ), by showing that, for each prefix  $v$  of  $c_\alpha$ ,  $v^{-1}c_\alpha$  can be decomposed into *generalized adjoining singular words*. Such a result has already been established (by Levé and Séébold [9]) for the case of the Fibonacci word  $f = c_{(3-\sqrt{5})/2}$ , where  $\frac{3-\sqrt{5}}{2} = [0; 2, \overline{1}]$ .

## 3 Decompositions of Conjugates of $c_\alpha$

### 3.1 Characteristic Sturmian Words Generated by Morphisms

Here, we describe all irrationals  $\alpha \in (0, 1)$  such that the characteristic Sturmian word  $c_\alpha$  is generated by a morphism. In order to do this, we must first define a special set of irrational numbers. A *Sturm number* (see [2]) is an irrational number  $\alpha \in (0, 1)$  that has a continued fraction expansion of one of the following types:

- (i)  $\alpha = [0; 1 + d_1, \overline{d_2, \dots, d_n}] < \frac{1}{2}$  with  $d_n \geq d_1 \geq 1$ ;

(ii)  $\alpha = [0; 1, d_1, \overline{d_2, \dots, d_n}] > \frac{1}{2}$  with  $d_n \geq d_1$ .

Observe that if  $\alpha = [0; 1 + d_1, \overline{d_2, \dots, d_n}]$  with  $d_n \geq d_1 \geq 1$ , then

$$1 - \alpha = \frac{1}{1 + \alpha/(1 - \alpha)} = [0; 1, d_1, \overline{d_2, \dots, d_n}].$$

Hence,  $\alpha$  has an expansion of type (i) if and only if  $1 - \alpha$  has an expansion of type (ii). Accordingly,  $\alpha$  is a Sturm number if and only if  $1 - \alpha$  is a Sturm number.

In what follows, we will always assume (unless otherwise stated) that  $\alpha$  is a Sturm number of type (i). Also, we shall denote the standard sequence of  $c_\alpha$  (resp.  $c_{1-\alpha}$ ) by  $(s_n)_{n \geq -1}$  (resp.  $(\hat{s}_n)_{n \geq -1}$ ). Clearly, we have  $\hat{s}_1 = \hat{s}_{-1} = b$  since  $1 - \alpha = [0; 1, d_1, \overline{d_2, \dots, d_n}]$ . Consequently,  $c_{1-\alpha}$  is obtained from  $c_\alpha$  by exchanging all letters  $a$  and  $b$  in  $c_\alpha$ , i.e.  $c_{1-\alpha} = E(c_\alpha)$ . Indeed, it is easily checked that

$$\hat{s}_n = E(s_{n-1}) \quad \text{for all } n \geq 0.$$

Hence,

$$E(c_\alpha) = E\left(\lim_{n \rightarrow \infty} s_n\right) = \lim_{n \rightarrow \infty} E(s_n) = \lim_{n \rightarrow \infty} \hat{s}_{n+1} = c_{1-\alpha}.$$

Therefore, we can restrict our attention to characteristic Sturmian words  $c_\alpha$  such that  $\alpha$  is a Sturm number of type (i). Later, an analogue of the main result of this paper (Theorem 3.7) will be deduced for  $c_{1-\alpha}$ .

We say that a morphism  $\psi$  *fixes* an infinite word  $x$  if  $\psi(x) = x$ , in which case  $x$  is called a *fixed point* of  $\psi$ . The following result describes all irrationals  $\alpha \in (0, 1)$  such that  $c_\alpha$  is a fixed point of a non-trivial morphism.

**Theorem 3.1.** [5, 8, 1] *Let  $\alpha \in (0, 1)$  be irrational. Then  $c_\alpha$  is a fixed point of a non-trivial morphism  $\sigma$  if and only if  $\alpha$  is a Sturm number. In particular, if  $\alpha = [0; 1 + d_1, \overline{d_2, \dots, d_n}]$  with  $d_n \geq d_1 \geq 1$ , then  $c_\alpha$  is the fixed point of any power of the morphism*

$$\sigma : \begin{array}{ll} a & \mapsto s_{n-1} \\ b & \mapsto s_{n-1}^{d_n - d_1} s_{n-2} \end{array}.$$

Further,  $c_{1-\alpha}$  is a fixed point of any power of the morphism

$$\hat{\sigma} : \begin{array}{ll} a & \mapsto \hat{s}_n^{d_n - d_1} \hat{s}_{n-1} \\ b & \mapsto \hat{s}_n \end{array}.$$

□

Note that, for any  $m \in \mathbb{N}$ , both  $\sigma^m$  and  $\hat{\sigma}^m$  are standard morphisms. In fact, it was shown by Crisp et al. [5] that

$$\begin{aligned} \sigma &= (\varphi E)^{d_1} E(\varphi E)^{d_2} E \dots (\varphi E)^{d_{n-1}} E(\varphi E)^{d_n - d_1}; \quad \text{and} \\ \hat{\sigma} &= E(\varphi E)^{d_1} E(\varphi E)^{d_2} E \dots (\varphi E)^{d_{n-1}} E(\varphi E)^{d_n - d_1} E. \end{aligned}$$

Observe that  $\hat{\sigma} = E\sigma E$ .

Now, Séébold [14] proved that a standard morphism  $\psi$  generates an infinite (characteristic Sturmian) word if and only if

$$\psi \in \{\varphi, E\varphi, \varphi E, E\varphi E\}^+ \setminus (\{E\varphi\}^+ \cup \{\varphi E\}^+).$$

Here, we prove that a characteristic Sturmian word  $c_\gamma$  is generated by a (standard) morphism if and only if  $\gamma$  is a Sturm number. Specifically, we prove that  $c_\alpha = \sigma^\omega(a)$  and  $c_{1-\alpha} = \hat{\sigma}^\omega(b) = (E\sigma E)^\omega(b)$ . These are direct results of the following lemma and corollary.

**Lemma 3.2.** *For any  $k \in \mathbb{N}$ ,  $\sigma(s_k) = s_{k+(n-1)}$ . Consequently, if  $k \in \mathbb{N}$  is fixed, then*

$$\sigma^m(s_k) = s_{k+m(n-1)} \quad \text{for all } m \geq 0.$$

*Proof.* Mathematical induction. □

**Corollary 3.3.** *For any integer  $m \geq 1$ ,  $\sigma^m(a) = s_{m(n-1)}$  and  $\sigma^m(b) = s_{m(n-1)}^{d_n - d_1} s_{m(n-1)-1}$ .* □

As an immediate consequence of the above corollary, we have the following result.

**Corollary 3.4.** *Let  $\alpha = [0; 1 + d_1, \overline{d_2, \dots, d_n}]$  with  $d_n \geq d_1 \geq 1$ . Then*

- (i)  $c_\alpha = \lim_{m \rightarrow \infty} \sigma^m(a) = \sigma^\omega(a)$ ;
- (ii)  $c_{1-\alpha} = \lim_{m \rightarrow \infty} \hat{\sigma}^m(b) = \hat{\sigma}^\omega(b)$ , where  $\hat{\sigma} = E\sigma E$ .

*Proof.* The fact that  $c_\alpha = \lim_{m \rightarrow \infty} \sigma^m(a)$  follows from Corollary 3.3 since  $\sigma^m(a) = s_{m(n-1)}$ , for any integer  $m \geq 1$ . Moreover, we know that  $c_{1-\alpha} = E(c_\alpha)$ , so that (ii) is obtained by realizing

$$c_{1-\alpha} = E \left( \lim_{m \rightarrow \infty} \sigma^m(a) \right) = \lim_{m \rightarrow \infty} E\sigma^m E(b) = \lim_{m \rightarrow \infty} (E\sigma E)^m(b).$$

□

In light of Corollary 3.4, one has that if  $\gamma$  is a Sturm number, then  $c_\gamma$  is generated by a (standard) morphism. The converse is trivially true in view of Theorem 3.1.

By considering Melançon's factorization of  $c_\alpha$  into adjoining singular words (Theorem 2.2), we will now extend Levé and Séébold's result (Theorem 4.6 in [9]) to the case  $\alpha = [0; 2, \bar{r}]$ .

### 3.2 The case $\alpha = [0; 2, \bar{r}]$

Now, if  $\alpha = [0; 2, \bar{r}]$ , then for each  $m \in \mathbb{N}$ ,

$$v_m = \begin{cases} as_{m+1}^{r-1} s_m b^{-1} & \text{if } m \text{ is odd,} \\ bs_{m+1}^{r-1} s_m a^{-1} & \text{otherwise,} \end{cases}$$

and  $v_{-1} = as_0^0 s_{-1} b^{-1} = a$ . Observe that, for any integer  $m \geq 0$ ,  $|\sigma^m(ab)| = |\sigma(s_1)| = |s_{m+1}| = q_{m+1}$ . Furthermore, using Corollary 3.3, it is easily checked that, for any  $m \geq -1$ ,

$$v_m = \begin{cases} a\sigma^{m+1}(b)b^{-1} & \text{if } m \text{ is odd,} \\ b\sigma^{m+1}(b)a^{-1} & \text{otherwise,} \end{cases}$$

since  $\sigma^{m+1}(b) = s_{m+1}^{r-1} s_m$ , for all  $m \in \mathbb{N}$ . Also note that  $|v_m| = q_{m+2} - q_{m+1}$ , for all  $m \geq -1$ .

Whence, we have the following special case of Lemma 2.1.

**Lemma 3.5.** *Suppose  $\alpha = [0; 2, \bar{r}]$  and let  $k, m \in \mathbb{N}$  be such that  $0 \leq k \leq q_{m+1} - 2$ . If  $v$  denotes the prefix of length  $k$  of  $c_\alpha (= (\sigma^m)^\omega(a))$ , then  $(\sigma^m)_k(c_\alpha) = v^{-1}c_\alpha$ .* □

The next lemma shows how to remove a prefix from the 'singular' decomposition of  $c_\alpha$  (cf. Proposition 4.5 in [9]).

**Lemma 3.6.** *Suppose  $\alpha = [0; 2, \bar{r}]$  and let  $k, m \in \mathbb{N}$  be such that  $k = q_{m+1} - p$  with  $2 \leq p \leq q_{m+1} - q_m + 1$ . Then*

$$(\sigma^m)_k(c_\alpha) = u^{-1}v_{m-1} \prod_{j=m}^{\infty} v_j,$$

where  $u$  is the prefix of  $v_{m-1}$  of length  $|u| = q_{m+1} - q_m + 1 - p$ .

*Proof.* We have  $q_m - 1 \leq k \leq q_{m+1} - 2$ . Thus, if  $k = 0$ , then we must have  $m = 0$ . Now,  $k = q_1 - p = 0$  implies  $p = q_1 = 2$ , and therefore,  $|u| = q_1 - q_0 + 1 - 2 = 0$ . Further, from Theorem 2.2, we have

$$(\sigma^0)_0(c_\alpha) = c_\alpha = \prod_{j=-1}^{\infty} v_j,$$

Suppose  $k \geq 1$ , then  $m \geq 1$  and, in this case, observe that

From Lemma 3.5, we know that  $(\sigma^m)_k(c_\alpha)$  is the word obtained from  $c_\alpha$  by removing its prefix of length  $k$ , i.e. of length at least  $\sum_{j=-1}^{m-2} (q_{j+2} - q_{j+1})$ . Whence, since  $|v_j| = q_{j+2} - q_{j+1}$  for any  $j \geq -1$ , then  $(\sigma^m)_k(c_\alpha)$  is obtained from  $c_\alpha = \prod_{j=-1}^{\infty} v_j$  by first removing the prefix  $v_{-1}v_0v_1 \cdots v_{m-2}$ . Then, from the remaining infinite word  $v_{m-1} \prod_{j=m}^{\infty} v_j$ , we remove the prefix  $u$  of  $v_{m-1}$  of length



[illegible]
$$\begin{aligned} c_\alpha &= (\sigma^0)_0(c_\alpha) = v_{-1}v_0v_1v_2v_3 \cdots \\ a^{-1}c_\alpha &= (\sigma^1)_1(c_\alpha) = v_0v_1v_2v_3 \cdots \\ (ab)^{-1}c_\alpha &= (\sigma^1)_2(c_\alpha) = b^{-1}v_0v_1v_2v_3 \cdots \\ (aba)^{-1}c_\alpha &= (\sigma^1)_3(c_\alpha) = (ba)^{-1}v_0v_1v_2v_3 \cdots \\ (abab)^{-1}c_\alpha &= (\sigma^1)_4(c_\alpha) = (bab)^{-1}v_0v_1v_2v_3 \cdots \\ (ababa)^{-1}c_\alpha &= (\sigma^1)_5(c_\alpha) = (baba)^{-1}v_0v_1v_2v_3 \cdots \\ (ababab)^{-1}c_\alpha &= (\sigma^2)_6(c_\alpha) = v_1v_2v_3 \cdots, \text{ etc.} \end{aligned}$$

**Theorem 3.7.** *Let  $k, m \in \mathbb{N}$  be such that  $k = q_{m+1} - p$  with  $2 \leq p \leq q_{m+1} - q_m + 1$ . Then*

*Proof.* It follows immediately from Lemma 3.6 that

where  $u$  is the prefix of  $v_{m-1}$  such that  $|u| = q_{m+1} - q_m + 1 - p$ . Thus, if  $u = \varepsilon$ , then

and so the result holds since  $|u| = 0 = q_{m+1} - q_m + 1 - p$  implies  $q_{m+1} - q_m - p = -1$ .

Observe that, by definition of the adjoining singular word  $v_m$ , there exist letters  $x, y \in \mathcal{A}$  ( $x \neq y$ ) such that, for any integer  $p \geq -1$ ,

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Hence,

$$\begin{aligned}
\prod_{j=m-1}^{\infty} v_j &= (x\sigma^m(b)y^{-1})(y\sigma^{m+1}(b)x^{-1})(x\sigma^{m+2}(b)y^{-1})\dots \\
&= x\sigma^m(b)\sigma^{m+1}(b)\sigma^{m+2}(b)\dots \\
&= x \prod_{j=m-1}^{\infty} \sigma^{j+1}(b),
\end{aligned}$$

and therefore,  $(\sigma^m)_k(c_\alpha) = u^{-1}x \prod_{j=m-1}^{\infty} \sigma^{j+1}(b)$ . If  $u \neq \varepsilon$ , then there exists a word  $\hat{u}$  such that  $u^{-1}x = \hat{u}^{-1}$ , which implies that  $\hat{u} = x^{-1}u$  with  $|\hat{u}| = q_{m+1} - q_m - p$ .

For any  $n \in \mathbb{N}$ ,  $\sigma^{n+1}(b) = s_{n+1}^{r-1}s_n$  is a prefix of  $\sigma^{n+2}(b) = s_{n+2}^{r-1}s_{n+1} = (s_{n+1}^{r-1}s_n)^{r-1}s_{n+1}$ , and  $|\sigma^{n+1}(b)| = |v_n| = q_{n+2} - q_{n+1}$ . Whence, for any integer  $r \geq m-1$ , we have  $\sigma^{r+1}(b) = \hat{u}u_{r+1}$ , for some  $u_{r+1} \in \mathcal{A}^*$  with  $|u_{r+1}| \geq p$ . Indeed, for  $r \geq m-1$ , we have

$$\begin{aligned}
|u_{r+1}| &= |\sigma^{r+1}(b)| - |\hat{u}| = q_{r+2} - q_{r+1} - q_{m+1} + q_m + p \\
&= (q_{r+2} - q_{m+1}) - (q_{r+1} - q_m) + p \geq p.
\end{aligned}$$

Consequently, by definition of right conjugation of morphisms,

$$(\sigma^{r+1})_{q_{m+1}-q_m-p}(b) = u_{r+1}\hat{u}.$$

From the above observations, we therefore find

$$\begin{aligned}
(\sigma^m)_k(c_\alpha) &= u^{-1}x \prod_{j=m-1}^{\infty} \sigma^{j+1}(b) \\
&= \hat{u}^{-1} \prod_{j=m-1}^{\infty} \sigma^{j+1}(b) \\
&= \hat{u}^{-1}\sigma^m(b) \prod_{j=m}^{\infty} \sigma^{j+1}(b) \\
&= \hat{u}^{-1}\hat{u}u_m \prod_{j=m}^{\infty} \sigma^{j+1}(b) \\
&= u_m \prod_{j=m}^{\infty} \sigma^{j+1}(b) \\
&= u_m\hat{u}\hat{u}^{-1} \prod_{j=m}^{\infty} \sigma^{j+1}(b) \\
&= (\sigma^m)_{q_{m+1}-q_m-p}(b)\hat{u}^{-1} \prod_{j=m}^{\infty} \sigma^{j+1}(b) \\
&= \dots \\
&= \prod_{j=m-1}^{\infty} (\sigma^{j+1})_{q_{m+1}-q_m-p}(b).
\end{aligned}$$

□

So if  $v$  is a prefix of  $c_\alpha$ , then  $v^{-1}c_\alpha$  can be obtained by concatenating all the words  $[(\sigma^{j+1})_i(b)]_{j \geq l}$ , where  $i$  and  $l$  are integers depending only on  $|v|$ . The characteristic Sturmian word  $c_\alpha$  is the special case when  $i = l = -1$ , so that

$$c_\alpha = \prod_{j=-1}^{\infty} (\sigma^{j+1})_{-1}(b),$$

where Melançon's adjoining singular words,  $v_j$ , are all the words  $(\sigma^{j+1})_{-1}(b)$ ,  $j \geq -1$ .



Recall that the Fibonacci word  $f$  is the characteristic Sturmian word  $c_\alpha$  such that  $\alpha = [0; 2, \overline{1}] = (3 - \sqrt{5})/2$ . In this case, one has  $\sigma = \varphi$ , i.e.

$$f = \varphi^\omega(a) = abaababaabaababaababaababaabaab \dots$$

If we set  $\varphi^{-1}(a) = b$ , then for any integer  $n \geq -1$ ,  $\varphi^n(a) = \varphi^{n+1}(b)$  with  $|\varphi^n(a)| = F_n = |\varphi^{n+1}(b)|$ , where  $F_n$  is the  $n$ -th *Fibonacci number* defined by

$$F_{-1} = F_0 = 1, \quad F_n = F_{n-1} + F_{n-2}; \quad n \geq 1.$$

Note that  $q_m = F_m$  for every  $m \in \mathbb{N}$ , and

$$F_{m+1} - F_m = F_{m+1} - F_m = F_{m-1} \quad \text{for all } m \in \mathbb{N}.$$

Hence, it is deduced from Theorem 3.7 that if  $k, m \in \mathbb{N}$  are such that  $k = F_{m+1} - p$  with  $2 \leq p \leq F_{m-1} + 1$ , then

$$(\varphi^m)_k(f) = \prod_{j=m-1}^{\infty} (\varphi^j)_{F_{m-1}-p}(a),$$

which is Levé and Séébold’s result (Theorem 4.6 in [9]).

As an example, we list some decompositions of conjugates of the characteristic Sturmian word  $c_\alpha$  for  $\alpha = [0; \overline{2}] = \sqrt{2} - 1$ .

	$v_{-1}$	$v_0$	$v_1$	$v_2$	$v_3$	$\cdots$
$m = 0$	$c_\alpha = (\sigma^0)_{-1}(b)$	$(\sigma^1)_{-1}(b)$	$(\sigma^2)_{-1}(b)$	$(\sigma^3)_{-1}(b)$	$(\sigma^4)_{-1}(b)$	$\cdots$
$m = 1$	$a^{-1}c_\alpha =$	$(\sigma^1)_{-1}(b)$	$(\sigma^2)_{-1}(b)$	$(\sigma^3)_{-1}(b)$	$(\sigma^4)_{-1}(b)$	$\cdots$
	$(ab)^{-1}c_\alpha =$	$(\sigma^1)_0(b)$	$(\sigma^2)_0(b)$	$(\sigma^3)_0(b)$	$(\sigma^4)_0(b)$	$\cdots$
	$(aba)^{-1}c_\alpha =$	$(\sigma^1)_1(b)$	$(\sigma^2)_1(b)$	$(\sigma^3)_1(b)$	$(\sigma^4)_1(b)$	$\cdots$
$m = 2$	$(abab)^{-1}c_\alpha =$		$(\sigma^2)_{-1}(b)$	$(\sigma^3)_{-1}(b)$	$(\sigma^4)_{-1}(b)$	$\cdots$
	$(ababa)^{-1}c_\alpha =$		$(\sigma^2)_0(b)$	$(\sigma^3)_0(b)$	$(\sigma^4)_0(b)$	$\cdots$
	$(ababaa)^{-1}c_\alpha =$		$(\sigma^2)_1(b)$	$(\sigma^3)_1(b)$	$(\sigma^4)_1(b)$	$\cdots$
	$(ababaab)^{-1}c_\alpha =$		$(\sigma^2)_2(b)$	$(\sigma^3)_2(b)$	$(\sigma^4)_2(b)$	$\cdots$
	$(ababaaba)^{-1}c_\alpha =$		$(\sigma^2)_3(b)$	$(\sigma^3)_3(b)$	$(\sigma^4)_3(b)$	$\cdots$
	$(ababaabab)^{-1}c_\alpha =$		$(\sigma^2)_4(b)$	$(\sigma^3)_4(b)$	$(\sigma^4)_4(b)$	$\cdots$
	$(ababaababa)^{-1}c_\alpha =$		$(\sigma^2)_5(b)$	$(\sigma^3)_5(b)$	$(\sigma^4)_5(b)$	$\cdots$
$m = 3$	$(ababaababaa)^{-1}c_\alpha =$			$(\sigma^3)_{-1}(b)$	$(\sigma^4)_{-1}(b)$	$\cdots$

### 3.3 The case $1 - \alpha = [0; 1, 1, \overline{r}]$

Now, if  $\alpha = [0; 2, \overline{7}]$ , then  $1 - \alpha = [0; 1, 1, \overline{7}]$ . By observing that  $c_{1-\alpha} = \hat{\sigma}^\omega(b)$ , where  $\hat{\sigma} = E\sigma E$ , it is clear that the following theorem is an immediate consequence of Theorem 3.7 and the lemma below.

**Lemma 3.8.** *For any standard morphism  $\psi$ ,*

$$(E\psi E)_k = E\psi_k E; \quad 0 \leq k \leq |\psi(ab)| - 2.$$

*Proof.* Let  $w \in \mathcal{A}^*$ , with  $|w| = k$ , be such that  $\psi(z)w = w\psi_k(z)$  for all  $z \in \mathcal{A}$ . Then, for some  $z \in \mathcal{A}$ ,

$$E\psi E(z)E(w) = E(\psi E(z)w) = E(\psi(y)w), \quad \text{for some } y \in \mathcal{A}, y \neq z.$$

Therefore, for  $z \in \mathcal{A}$  and  $0 \leq k \leq |\psi(ab)| - 2$ , we have

$$E\psi E(z)E(w) = E(\psi(y)w) = E(w\psi_k(y)) = E(w)E\psi_k E(z).$$

Thus, there exists a word of length  $k$ , namely  $w' = E(w)$ , such that for any  $u \in \mathcal{A}^*$ ,

$$E\psi E(u)w' = w'E\psi_k E(u).$$

□

**Theorem 3.9.** *Let  $k, m \in \mathbb{N}$  be such that  $k = q_{m+1} - p$  with  $2 \leq p \leq q_{m+1} - q_m + 1$ . Then*

$$(\hat{\sigma}^m)_k(c_{1-\alpha}) = \prod_{j=m-1}^{\infty} (\hat{\sigma}^{j+1})_{q_{m+1}-q_m-p}(a).$$

□

## 4 Concluding Remarks

Note that, by Corollary 3.3, for any integer  $m \geq 1$ ,

$$\sigma^m(b) = s_{m(n-1)}^{d_n-d_1} s_{m(n-1)-1},$$

where  $\sigma$  is the standard morphism that generates  $c_\alpha$ , for  $\alpha = [0; 1 + d_1, \overline{d_2, \dots, d_n}]$  with  $d_n \geq d_1 \geq 1$ . Now, by the periodicity of the continued fraction expansion of  $\alpha$ ,  $d_i = d_{i+n-1}$ , for all  $i \geq 2$ . Hence, it is easily deduced that

$$d_n = d_{m(n-1)+1}, \text{ for any integer } m \geq 1,$$

and one may write

$$\sigma^m(b) = s_{m(n-1)}^{d_{m(n-1)+1}-d_1} s_{m(n-1)-1}.$$

Whence, if  $d_1 = 1$ , then for each  $m \geq 1$ ,

$$v_{m(n-1)-1} = \begin{cases} a\sigma^m(b)b^{-1} & \text{if } m \text{ is even,} \\ b\sigma^m(b)a^{-1} & \text{otherwise.} \end{cases}$$

Accordingly, we do not have a ‘nice’ expression for each  $v_k$  ( $k \in \mathbb{N}$ ) in terms of a power of  $\sigma$  unless  $n = 2$ , i.e. unless  $\alpha = [0; 2, \overline{r}]$  for some  $r \in \mathbb{N}^+$ . Therefore, we cannot establish an extension of Theorem 3.7 to the case of  $c_\alpha$  (nor  $c_{1-\alpha}$ ) with  $\alpha$  (as above) having  $d_1 \geq 2$  and  $n \geq 3$ .

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