# HIGH ACCURACY SEMIDEFINITE PROGRAMMING BOUNDS FOR KISSING NUMBERS 

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#### Abstract

The kissing number in $n$-dimensional Euclidean space is the maximal number of non-overlapping unit spheres which simultaneously can touch a central unit sphere. Bachoc and Vallentin developed a method to find upper bounds for the kissing number based on semidefinite programming. This paper is a report on high accuracy calculations of these upper bounds for $n \leq 24$. The bound for $n=16$ implies a conjecture of Conway and Sloane: There is no 16-dimensional periodic sphere packing with average theta series $$
1+7680 q^{3}+4320 q^{4}+276480 q^{5}+61440 q^{6}+\cdots
$$


## 1. Introduction

In geometry, the kissing number in $n$-dimensional Euclidean space is the maximal number of non-overlapping unit spheres which simultaneously can touch a central unit sphere. The kissing number is only known in dimensions $n=1,2,3,4,8,24$, and there were many attempts to find good lower and upper bounds. We refer to Casselman [4] for the history of this problem and to Pfender, Ziegler [14], Elkies [7], and Conway, Sloane [6] for more background information on sphere packing problems.

Bachoc and Vallentin [1] develop a method (Section 2 recalls it) to find upper bounds for the kissing number based on semidefinite programming. Table 1 in Section 3, the main contribution of this paper, gives the values, i.e. the first ten significant digits, of these upper bounds for all dimensions $3, \ldots, 24$. In all cases they are the best known upper bounds. Dimension 5 is the first dimension in which the kissing number is not known. With our computation we could limit the range of possible values from $\{40, \ldots, 45\}$ to $\{40, \ldots, 44\}$. In Section 4 we show that the high accuracy computations for the upper bounds in dimension 4 result into a question about a possible approach to prove the uniqueness of the kissing configuration in 4 dimensions.

Although acquiring the data for the table is a purely computational task we think that providing this table is valuable for several reasons: The kissing number is an important constant in geometry and results can depend on good upper bounds for it. For instance in Section 5 we show that there is no periodic point set in dimension

[^0]16 with average theta series

$$
1+7680 q^{3}+4320 q^{4}+276480 q^{5}+61440 q^{6}+\cdots
$$

This proves a conjecture of Conway and Sloane [6, Chapter 7, page 190]. Furthermore, the actual computation of the table was very challenging. Bachoc and Vallentin [1] gave results for dimension $3, \ldots, 10$. However, they report on numerical difficulties which prevented them from extending their results. Now using new, more sophisticated high accuracy software and faster computers and more computation time we could overcome some of the numerical difficulties. Section 3 contains details about the computations.

## 2. Notation

In this section we set up the notation which is needed for our computation. For more information we refer to [1]. For natural numbers $d$ and $n \geq 3$ let $s_{d}(n)$ be the optimal value of the minimization problem

$$
\begin{aligned}
& \min \left\{1+\sum_{k=1}^{d} a_{k}+b_{11}+\left\langle F_{0}, S_{0}^{n}(1,1,1)\right\rangle:\right. \\
& \quad a_{1}, \ldots, a_{d} \in \mathbb{R}, \quad a_{1}, \ldots, a_{d} \geq 0 \\
& b_{11}, b_{12}, b_{22} \in \mathbb{R}, \quad\binom{b_{11} b_{12}}{b_{12} b_{22}} \text { is positive semidefinite, } \\
& F_{k} \in \mathbb{R}^{(d+1-k) \times(d+1-k), \quad F_{k} \text { is positive semidefinite, } \quad k=0, \ldots, d,} \begin{array}{l}
q, q_{1} \in \mathbb{R}[u], \operatorname{deg}\left(p+p q_{1}\right) \leq d, \quad p, p_{1} \text { sums of squares, } \\
\quad r, r_{1}, \ldots, r_{4} \in \mathbb{R}[u, v, t], \operatorname{deg}\left(r+\sum_{i=1}^{4} p_{i} r_{i}\right) \leq d, \quad r, r_{1}, \ldots, r_{4} \text { sums of squares, } \\
1+\sum_{k=1}^{d} a_{k} P_{k}^{n}(u)+2 b_{12}+b_{22}+3 \sum_{k=0}^{d}\left\langle F_{k}, S_{k}^{n}(u, u, 1)\right\rangle+q(u)+p(u) q_{1}(u)=0, \\
\left.b_{22}+\sum_{k=0}^{d}\left\langle F_{k}, S_{k}^{n}(u, v, t)\right\rangle+r(u, v, t)+\sum_{i=1}^{4} p_{i}(u, v, t) r_{i}(u, v, t)=0\right\}
\end{array} .
\end{aligned}
$$

Here $P_{k}^{n}$ is the normalized Jacobi polynomial of degree $k$ with $P_{k}^{n}(1)=1$ and parameters $((n-3) / 2,(n-3) / 2)$. In general, Jacobi polynomials with parameters $(\alpha, \beta)$ are orthogonal polynomials for the measure $(1-u)^{\alpha}(1+u)^{\beta} d u$ on the interval $[-1,1]$. Before we can define the matrices $S_{k}^{n}$ we first define the entry $(i, j)$ with $i, j \geq 0$ of the (infinite) matrix $Y_{k}^{n}$ containing polynomials in the variables $u, v, w$ by

$$
\begin{aligned}
& \left(Y_{k}^{n}\right)_{i, j}(u, v, t)=u^{i} v^{j} \\
& \quad\left(\left(1-u^{2}\right)\left(1-v^{2}\right)\right)^{k / 2} P_{k}^{n-1}\left(\frac{t-u v}{\sqrt{\left(1-u^{2}\right)\left(1-v^{2}\right)}}\right)
\end{aligned}
$$

Then we get $S_{k}^{n}$ by symmetrization: $S_{k}^{n}=\frac{1}{6} \sum_{\sigma} \sigma Y_{k}^{n}$, where $\sigma$ runs through all permutations of the variables $u, v, t$ which acts on the matrix coefficients in the
obvious way. The polynomials $p, p_{1}, \ldots, p_{4}$ are given by

$$
\begin{aligned}
& p(u)=-(u+1)(u+1 / 2) \\
& p_{1}(u, v, t)=p(u), \quad p_{2}(u, v, t)=p(v), \quad p_{3}(u, v, t)=p(t), \\
& p_{4}(u, v, t)=1+2 u v t-u^{2}-v^{2}-t^{2} .
\end{aligned}
$$

By $\langle A, B\rangle$ we denote the inner product between symmetric matrices trace $(A B)$.
In 11 it is shown that this minimization problem is a semidefinite program and that every upper bound on $s_{d}(n)$ provides an upper bound for the kissing number in dimension $n$. Clearly, the numbers $s_{d}(n)$ form a monotonic decreasing sequence in $d$.

## 3. Bounds for kissing numbers

| $n$ | best lower bound known | best upper bound previously known | SDP bound |
| :---: | :---: | :---: | :---: |
| 3 | 12 |  | $\begin{aligned} & \hline s_{11}(3)=12.42167009 \ldots \\ & s_{12}(3)=12.40203212 \ldots \\ & s_{13}(3)=12.39266509 \ldots \\ & s_{14}(3)=12.38180947 \ldots \end{aligned}$ |
| 4 | 24 | (11) Musin, 2008 | $\begin{aligned} & s_{11}(4)=24.10550859 \ldots \\ & s_{12}(4)=24.09098111 \ldots \\ & s_{13}(4)=24.07519774 \ldots \\ & s_{14}(4)=24.06628391 \ldots \end{aligned}$ |
| 5 | 40 | [1] Bachoc, Vallentin, 2008 | $\begin{aligned} & s_{11}(5)=45.06107293 \ldots \\ & s_{12}(5)=45.02353644 \ldots \\ & s_{13}(5)=45.00650838 \ldots \\ & s_{14}(5)=\underline{44} .99899685 \ldots \end{aligned}$ |
| 6 | 72 | (1) Bachoc, $\frac{78}{\text { Vallentin, } 2008}$ | $\begin{aligned} & s_{11}(6)=78.58344077 \ldots \\ & s_{12}(6)=78.35518719 \ldots \\ & s_{13}(6)=78.29404232 \ldots \\ & s_{14}(6)=78.24061272 \ldots \end{aligned}$ |
| 7 | 126 | [1] Bachoc, Vallentin, 2008 | $\begin{aligned} & s_{11}(7)=134.8824614 \ldots \\ & s_{12}(7)=134.7319671 \ldots \\ & s_{13}(7)=134.5730609 \ldots \\ & s_{14}(7)=134.4488169 \ldots \end{aligned}$ |
| 8 | 240 | (12] Odlyzko, Sloane, 1979 (9) Levenshtein, 1979 | $s_{11}(8)=240.0000000 \ldots$ |
| 9 | 306 | [1] Bachoc, Vallentin, 2008 | $\begin{aligned} & s_{11}(9)=365.3229274 \ldots \\ & s_{12}(9)=364.7282746 \ldots \\ & s_{13}(9)=364.3980087 \ldots \\ & s_{14}(9)=364.0919287 \ldots \end{aligned}$ |
| 10 | 500 | (1) Bachoc, Vallentin, 2008 | $\begin{aligned} & s_{11}(10)=558.1442813 \ldots \\ & s_{12}(10)=556.2840736 \ldots \\ & s_{13}(10)=555.2399024 \ldots \\ & s_{14}(10)=554.5075418 \ldots \end{aligned}$ |
| 11 | 582 | [12] Odlyzko, Sloane, 1979 | $\begin{aligned} & \hline s_{11}(11)=878.6158044 \ldots \\ & s_{12}(11)=873.3790094 \ldots \\ & s_{13}(11)=871.9718533 \ldots \\ & s_{14}(11)=\underline{870.8831157 \ldots} \end{aligned}$ |


| 12 | 840 | $1416$ <br> (12] Odlyzko, Sloane, 1979 | $\begin{aligned} & s_{11}(12)=1364.683810 \ldots \\ & s_{12}(12)=1362.200297 \ldots \\ & s_{13}(12)=1359.283834 \ldots \\ & s_{14}(12)=1357.889300 \ldots \end{aligned}$ |
| :---: | :---: | :---: | :---: |
| 13 | 1130 | (12) Odlyzko, Sloane, 1979 | $\begin{aligned} & s_{11}(13)=2089.116331 \ldots \\ & s_{12}(13)=2080.631518 \ldots \\ & s_{13}(13)=2073.074796 \ldots \\ & s_{14}(13)=\underline{2069.587585 \ldots} \end{aligned}$ |
| 14 | 1582 | 3492 [12] Odlyzko, Sloane, 1979 | $\begin{aligned} & s_{11}(14)=3224.950751 \ldots \\ & s_{12}(14)=3202.448902 \ldots \\ & s_{13}(14)=3189.127644 \ldots \\ & s_{14}(14)=3183.133169 \ldots \\ & \hline \end{aligned}$ |
| 15 | 2564 | 5431 (12] Odlyzko, Sloane, 1979 | $\begin{aligned} & s_{11}(15)=4949.650431 \ldots \\ & s_{12}(15)=4893.479446 \ldots \\ & s_{13}(15)=4876.037229 \ldots \\ & s_{14}(15)=4866.245659 \ldots \end{aligned}$ |
| 16 | 4320 | 13 Pfender, 2007 | $\begin{aligned} & s_{11}(16)=7515.952644 \ldots \\ & s_{12}(16)=7432.720718 \ldots \\ & s_{13}(16)=7374.093742 \ldots \\ & s_{14}(16)=7355.809036 \ldots \end{aligned}$ |
| 17 | 5346 | 12210 (13) Pfender, 2007 | $\begin{aligned} & s_{11}(17)=11568.41674 \ldots \\ & s_{12}(17)=11333.84265 \ldots \\ & s_{13}(17)=11128.26227 \ldots \\ & s_{14}(17)=11072.37543 \ldots \end{aligned}$ |
| 18 | 7398 | 17877 [12] Odlyzko, Sloane, 1979 | $\begin{aligned} & s_{11}(18)=17473.48016 \ldots \\ & s_{12}(18)=17034.32488 \ldots \\ & s_{13}(18)=16686.28908 \ldots \\ & s_{14}(18)=16572.26478 \ldots \end{aligned}$ |
| 19 | 10668 | [3] Boyvalenkov, 1994 | $\begin{aligned} & s_{11}(19)=26397.34794 \ldots \\ & s_{12}(19)=25636.98958 \ldots \\ & s_{13}(19)=25029.87432 \ldots \\ & s_{14}(19)=\underline{24812.30254 \ldots} \end{aligned}$ |
| 20 | 17400 | 37974 [12] Odlyzko, Sloane, 1979 | $\begin{aligned} & s_{11}(20)=39045.32761 \ldots \\ & s_{12}(20)=37844.10380 \ldots \\ & s_{13}(20)=37067.18966 \ldots \\ & s_{14}(20)=\underline{36764} .40138 \ldots \\ & \hline \end{aligned}$ |
| 21 | 27720 | 56851 [3] Boyvalenkov, 1994 | $\begin{gathered} s_{11}(21)=58087.03857 \ldots \\ s_{12}(21)=56079.21685 \ldots \\ s_{13}(21)=55170.03449 \ldots \\ s_{14}(21)=54584.76757 \ldots \\ \hline \end{gathered}$ |
| 22 | 49896 | 86537 (12) Odlyzko, Sloane, 1979 | $\begin{aligned} & s_{11}(22)=87209.06261 \ldots \\ & s_{12}(22)=84922.09101 \ldots \\ & s_{13}(22)=84117.92103 \ldots \\ & s_{14}(22)=82340.08003 \ldots \end{aligned}$ |
| 23 | 93150 | 128095 [3] Boyvalenkov, 1994 | $\begin{aligned} & s_{11}(23)=128360.7969 \ldots \\ & s_{12}(23)=127323.7095 \ldots \\ & s_{13}(23)=125978.7655 \ldots \\ & s_{14}(23)=\underline{124416.9796 \ldots} \end{aligned}$ |
| 24 | 196560 | 196560 <br> [12] Odlyzko, Sloane, 1979 <br> (9) Levenshtein, 1979 | $s_{11}(24)=196560.0000 \ldots$ |

Table 1. New upper bounds for the kissing number (best known values are underlined).

Finding the solution of the semidefinite program defined in Section 2 is a computational challenge. It turns out that the major obstacle is numerical instability and not the problem size. When $d$ is fixed, then the size of the input matrices stays constant with $n$; when $n$ is fixed, then it grows rather moderately with $d$.

There is a number of available software packages for solving semidefinite programs. Mittelmann compares many existing packages in 10. For our purpose first order, gradient-based methods such as SDPLR are far too inaccurate, and second order, primal-dual interior point methods are more suitable. Here increasingly ill-conditioned linear systems have to be solved even if the underlying problem is well-conditioned. This happens in the final phase of the algorithm when one approaches an optimal solution. Our problems are not well-conditioned and even the most robust solver SeDuMi which uses partial quadruple arithmetic in the final phase does not produce reliable results for $d>10$.

We thus had to fall back on the implementation SDPA-GMP [8] which is much slower but much more accurate than other software packages because it uses the GNU Multiple Precision Arithmetic Library. We worked with 200-300 binary digits and relative stopping criteria of $10^{-30}$. The ten significant digits listed in the table are thus guaranteed to be correct. One problem was the convergence. Even with the control parameter settings recommended by the authors of SDPA-GMP for "slow but stable" computations, the algorithm failed to converge in several instances. However, we found parameter choices which worked for all cases: We varied the parameter lambdaStar between 100 and 10000 depending on the case while the other parameters could be chosen at or near the values recommended for "slow but stable" performance.

The computations were done on Intel Core 2 platforms with one and two Quad processors. The computation time varied between five and ten weeks per case for $d=12$. An accuracy of 128 bits in SDPA-GMP did yield sufficient accuracy but did not yield a reduction in computing time.

After the computations for the cases $d=11$ and $d=12$ were finished new 128bit versions (quadruple precision) of SDPA and CSDP became available; partly with our assistance. These new versions do not rely on the GNU Multiple Precision Arithmetic Library. So the computation time for the cases $d=13$ and $d=14$ were reasonable: from one week to two and a half weeks.

## 4. Question about the optimality of the $D_{4}$ Root system

Looking at the values $s_{d}(4)$ in Table 1 one is led to the following question:
Question 4.1. Is $\lim _{d \rightarrow \infty} s_{d}(4)=24$ ?
If the answer to this question is yes (which at the moment appears unlikely because we computed $s_{15}(4)=23.06274835 \ldots$ ), then it would have two noteworthy consequences about optimality properties of the root system $D_{4}$.

The root system $D_{4}$ defines (up to orthogonal transformations) a point configuration on the unit sphere $S^{3}=\left\{x \in \mathbb{R}^{4}: x \cdot x=1\right\}$ consisting of 24 points; it is the same point configuration as the one coming from the vertices of the regular 24 cell. This point configuration has the property that the spherical distance of every two distinct points is at least arccos $1 / 2$. Hence, these points can be the maximal 24 touching points of unit spheres kissing the central unit sphere $S^{3}$.

If $\lim _{d \rightarrow \infty} s_{d}(4)=24$, then this would prove that the root system $D_{4}$ is the unique optimal point configuration of cardinality 24 . Here optimality means that one cannot distribute 24 points on $S^{3}$ so that the minimal spherical distance between two distinct points exceeds arccos $1 / 2$. Thus, the root system $D_{4}$ would be characterized by its kissing property. This is generally believed to be true but so far no proof could be given.

Another consequence would be that there is no universally optimal point configuration of 24 points in $S^{3}$ as conjectured by Cohn, Conway, Elkies, Kumar [5]. Universally optimal point configurations minimize every absolutely monotonic potential function. The conjecture would follow if the answer to our question is yes: Every universally optimal point configuration is automatically optimal and Cohn, Conway, Elkies, Kumar [5] show that the root system $D_{4}$ is not universally optimal.

## 5. Nonexistence of a sphere packing

Our new upper bound of 7355 for the kissing number in dimension 16 implies that there is no periodic point set in dimension 16 whose average theta series equals

$$
\begin{equation*}
1+7680 q^{3}+4320 q^{4}+276480 q^{5}+61440 q^{6}+\cdots \tag{1}
\end{equation*}
$$

This settles a conjecture of Conway and Sloane [6, Chapter 7, page 190]. In this section we explain this result. We refer to Conway, Sloane [6, Elkies 7], and to Bowert [2] for more information.

An $n$-dimensional periodic point set $\Lambda$ is a finite union of translates of an $n$ dimensional lattice, i.e. one can write $\Lambda$ as $\Lambda=\left(A \mathbb{Z}^{n}+v_{1}\right) \cup \ldots \cup\left(A \mathbb{Z}^{n}+v_{N}\right)$, with $v_{1}, \ldots, v_{N} \in \mathbb{R}^{n}$, and $A: \mathbb{R}^{n} \rightarrow \mathbb{R}^{n}$ is a linear isomorphism. The average theta series of $\Lambda$ is

$$
\Theta_{\Lambda}(z)=\frac{1}{N} \sum_{i=1}^{N} \sum_{j=1}^{N} \sum_{v \in \mathbb{Z}^{n}} q^{\left\|A v-v_{i}+v_{j}\right\|^{2}}, \text { with } q=e^{\pi i z}
$$

This is a holomorphic function defined on the complex upper half-plane. A holomorphic function $f$ which is defined on the complex upper half-plane, which is meromorphic for $z \rightarrow i \infty$, and which satisfies the transformation laws

$$
f(-1 / z)=z^{8} f(z), \text { and } f(z+2)=f(z) \text { for all } z \in \mathbb{C} \text { with } \Im z>0
$$

is called a modular forms of weight 8 for the Hecke group $G(2)$. The expression (11) defines the unique modular form of weight 8 for the Hecke group $G(2)$ which starts off with $1+0 q^{1}+0 q^{2}$. It is also called an extremal modular form, see Scharlau and Schulze-Pillot [15]

If there would be a 16-dimensional periodic point set whose average theta series coincides with (11) then this periodic point set would define the sphere centers of a sphere packing with extraordinary high density (see [6, Chapter 7, page 190]). At the same time the existence of such a periodic point set would show that the kissing number in dimension 16 is at least 7680 . This is not the case.

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