Two Reflected Gray Code based orders on some restricted growth sequences

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October 16, 2018

Abstract

We consider two order relations: that induced by the *m*-ary reflected Gray code and a suffix partitioned variation of it. We show that both of them when applied to some sets of restricted growth sequences still yield Gray codes. These sets of sequences are: subexcedant or ascent sequences, restricted growth functions, and staircase words. In each case we give efficient exhaustive generating algorithms and compare the obtained results.

1 Introduction and motivations

The term 'Gray code' was taken from Frank Gray, who patented Binary Reflected Gray Code (BRGC) in 1953 [1]. The concept of BRGC is extended to Reflected Gray Code (RGC), to accommodate m-tuples (sequence), with m > 2 [2]. In these Gray codes, successive sequence differ in a single position, and by +1 or -1 in this position. More generally, if a list of sequences is such that the Hamming distance between successive sequences is upper bounded by a constant d, then the list is said a d-Gray code. So in particular, BRGC and RGC are 1-Gray codes. In addition, if the positions where the successive sequences differ are adjacent, then we say that the list is a d-adjacent Gray code.

For long time, the design of Gray codes for combinatorial classes and their corresponding generating algorithms was an ad-hoc task, that is, done case by case according to the class under consideration. Recently, general techniques which fit to large classes of combinatorial objects were developed and used. Among them are, for example, the ECO-method (initiated in Gray code context in [3], see also [4]), prefix rotations (yielding bubble languages, see [5,6] and references therein), or Reflected Gray Code based order relations; this last technique was used implicitly, for example in [7,8], and developed systematically as a general method in [4,9–13]. The results presented in this paper are in the light of this last direction. More precisely, we show that two order relations induced by Reflected Gray Code and its variation also give Gray codes for some classes of restricted growth sequences defined by means of statistics. These classes are: subexcedant or ascent sequences, restricted growth functions, and staircase words. We give efficient (CAT) generating algorithm for each obtained Gray code.

2 Preliminaries

2.1 Gray code orders

Let $G_n(m)$ be the set of length n m-ary sequences $s_1s_2...s_n$ with $s_i \in \{0, 1, ..., m-1\}$; clearly, $G_n(m)$ is the product set $\{0, 1, ..., m-1\}^n$. The Reflected Gray Code (RGC for short) for the set $G_n(m)$, denoted by $G_n(m)$, is the natural extension of the Binary Reflected Gray Code to this set. The list $G_n(m)$ is defined recursively by the following relation [2]:

$$\mathcal{G}_n(m) = \begin{cases} \epsilon & \text{if } n = 0, \\ 0\mathcal{G}_{n-1}(m), \ 1\overline{\mathcal{G}_{n-1}(m)}, \ 2\mathcal{G}_{n-1}(m), \ \dots, \ (m-1)\mathcal{G}'_{n-1}(m) & \text{if } n > 0, \end{cases}$$
(1)

where ϵ is the empty sequence, $\overline{\mathcal{G}_{n-1}(m)}$ is the reverse of $\mathcal{G}_{n-1}(m)$, and $\mathcal{G}'_{n-1}(m)$ is $\mathcal{G}_{n-1}(m)$ or $\overline{\mathcal{G}_{n-1}(m)}$ according to m is odd or even.

In $\mathcal{G}_n(m)$, two successive sequences differ in a single position and by +1 or -1 in this position. A list for a set of sequences induces an order relation to this set, and we give two order relations induced by the RGC and its variation, namely RGC order [10] and Co-RGC order.

We adopt the convention that lower case bold letters represent tuples, for example: $\mathbf{s} = s_1 s_2 \dots s_n$, $\mathbf{a} = a_1 a_2 \dots a_k$, $\mathbf{b} = b_{k+1} b_{k+2} \dots b_n$.

Definition 1. The Reflected Gray Code order \prec on $G_n(m)$ is defined as: $\mathbf{s} = s_1 s_2 \dots s_n$ is less than $\mathbf{t} = t_1 t_2 \dots t_n$, denoted by $\mathbf{s} \prec \mathbf{t}$, if either

- $\sum_{i=1}^{k-1} s_i$ is even and $s_k < t_k$, or
- $\sum_{i=1}^{k-1} s_i$ is odd and $s_k > t_k$,

where k is the leftmost position where s and t differ.

It is easy to see that $\mathcal{G}_n(m)$ defined in relation (1) lists sequences in $G_n(m)$ in \prec order. Now we give a variation of $\mathcal{G}_n(m)$. Let $s_1s_2...s_n$ be a sequence in $G_n(m)$. The complement of s_i , $1 \leq i \leq n$, is

$$(m-1-s_i),$$

and the reverse of $s_1 s_2 \dots s_n$ is

$$s_n s_{n-1} \dots s_1$$
.

Let $\widetilde{\mathcal{G}}_n(m)$ be the list obtained by transforming each sequence s in $\mathcal{G}_n(m)$ as follows:

- complementing each digit in s if m is even, or complementing only digits in odd positions if m is odd, then
- reversing the obtained sequence.

Clearly, $\widetilde{\mathcal{G}}_n(m)$ is also a Gray code for $G_n(m)$, and the sequences therein are listed in Co-Reflected Gray Code order, as defined formally below.

Definition 2. The Co-Reflected Gray Code order \prec_c on $G_n(m)$ is defined as: $s = s_1 s_2 \dots s_n$ is less than $t = t_1 t_2 \dots t_n$, denoted by $s \prec_c t$, if either

• $\sum_{i=k+1}^{n} s_i + (n-k)$ is even and $s_k > t_k$, or

• $\sum_{i=k+1}^{n} s_i + (n-k)$ is odd and $s_k < t_k$,

where k is the rightmost position where s and t differ.

Although this definition sounds somewhat arbitrary, as we will see in Section 4, it turns out that \prec_c order gives suffix partitioned Gray codes for some sets of restricted growth sequences. Obviously, the restriction of $\mathcal{G}_n(m)$ (resp. $\widetilde{\mathcal{G}}_n(m)$) to a set of sequences is simply the list of sequences in the set listed in \prec (resp. \prec_c) order.

2.2 Restricted growth sequences defined by means of statistics

Through this paper we consider sequences over non-negative integers. A *statistic* on a set of sequences is an association of an integer to each sequence in the set. For a sequence $s_1s_2...s_n$, its length minus one, numbers of ascents/levels/descents, maximal value, and last value are classical examples of statistics. They are defined as follows, see also [14]:

- $len(s_1s_2...s_n) = n-1;$
- $asc(s_1 s_2 ... s_n) = card\{i \mid 1 \le i < n \text{ and } s_i < s_{i+1}\};$
- $lev(s_1 s_2 ... s_n) = card\{i | 1 \le i < n \text{ and } s_i = s_{i+1}\};$
- $des(s_1 s_2 ... s_n) = card\{i \mid 1 \le i < n \text{ and } s_i > s_{i+1}\};$
- $m(s_1s_2...s_n) = \max\{s_1, s_2, ..., s_n\};$
- $\bullet \ \mathsf{lv}(s_1s_2\dots s_n) = s_n.$

If st is one of the statistics len, asc, m, and lv, then st satisfy the following:

$$\mathsf{st}(s_1 s_2 \dots s_n) \le n - 1,\tag{2}$$

and

if
$$s_n = \operatorname{st}(s_1 s_2 \dots s_{n-1}) + 1$$
, then $s_n = \operatorname{st}(s_1 s_2 \dots s_{n-1} s_n)$. (3)

On the contrary, the statistics lev and des do not satisfy relation (3). Accordingly, through this paper we will consider only the four statistics above. However, as we will point out, some of the results presented here are also true for arbitrary statistics satisfying relations (2) and (3).

Definition 3. For a given statistic st, an st-restricted growth sequence $s_1 s_2 \dots s_n$ is a sequence with $s_1 = 0$ and

$$0 \le s_{k+1} \le \mathsf{st}(s_1 s_2 \dots s_k) + 1 \text{ for } 1 \le k < n, \tag{4}$$

and the set of st-restricted growth sequences is the set of all sequences $s_1 s_2 \dots s_n$ satisfying relation (4).

From this definition, it follows that any prefix of an st-restricted growth sequence is also (a shorter) st-restricted growth sequence.

Remark 1. If st is a statistic satisfying relations (2) and (3) above, then

1. $\max\{\operatorname{st}(s_1s_2\ldots s_n)\mid s_1s_2\ldots s_n \text{ is an st-restricted growth sequence}\}=n-1;$

2. if $s_1 s_2 \dots s_n$ is an st-restricted growth sequence, then for any $k, 1 \leq k < n, s_{k+1} = \operatorname{st}(s_1 s_2 \dots s_k) + 1$ implies $s_{k+1} = \operatorname{st}(s_1 s_2 \dots s_k s_{k+1})$.

The sets of st-restricted growth sequences, where st is one of the statistics len, asc, m, and lv, are defined below.

Definition 4.

• The set SE_n of subexcedant sequences of length n is defined as:

$$SE_n = \{s_1 s_2 \dots s_n \mid s_1 = 0 \text{ and } 0 \le s_{k+1} \le \text{len}(s_1 s_2 \dots s_k) + 1 \text{ for } 1 \le k < n\};$$

• The set A_n of ascent sequences of length n is defined as:

$$A_n = \{s_1 s_2 \dots s_n \mid s_1 = 0 \text{ and } 0 \le s_{k+1} \le \mathsf{asc}(s_1 s_2 \dots s_k) + 1 \text{ for } 1 \le k < n\};$$

• The set R_n of restricted growth functions of length n is defined as:

$$R_n = \{s_1 s_2 \dots s_n \mid s_1 = 0 \text{ and } 0 \le s_{k+1} \le \mathsf{m}(s_1 s_2 \dots s_k) + 1 \text{ for } 1 \le k < n\};$$

• The set S_n of staircase words of length n is defined as:

$$S_n = \{s_1 s_2 \dots s_n \mid s_1 = 0 \text{ and } 0 \le s_{k+1} \le \mathsf{lv}(s_1 s_2 \dots s_k) + 1 \text{ for } 1 \le k < n\}.$$

Notice that alternatively, $SE_n = \{0\} \times \{0,1\} \times \ldots \times \{0,1,\ldots,n-1\}.$

Remark 2.
$$S_n \subset R_n \subset A_n \subset SE_n \subset G_n(n)$$
.

Below we give examples to illustrate Remark 2.

Example 1.

- If s = 010145, then $s \in SE_6$, but $s \notin A_6$, $s \notin R_6$ and $s \notin S_6$.
- If s = 010103, then $s \in SE_6$, $s \in A_6$, but $s \notin R_6$ and $s \notin S_6$.
- If s = 010102, then $s \in SE_6$, $s \in A_6$, and $s \in R_6$, but $s \notin S_6$.
- If s = 010101, then $s \in SE_6$, $s \in A_6$, $s \in R_6$ and $s \in S_6$.

See Table 3 for the sets S_5 , R_5 , and A_5 listed in \prec order, and Table 4 for the same sets listed in \prec_c order.

We denote by

- \mathcal{X}_n the list for the set X_n in \prec order, and by $\widetilde{\mathcal{X}}_n$ that in \prec_c order;
- $\operatorname{succ}_X(s)$, $s \in X_n$, the successor of s in the set X_n listed in \prec order; that is, the smallest sequence in X_n larger than s with respect to \prec order;
- $\widetilde{\operatorname{succ}}_X(s)$ the counterpart of $\operatorname{succ}_X(s)$ with respect to \prec_c order;
- first(\mathcal{L}) the first sequence in the list \mathcal{L} ;
- $last(\mathcal{L})$ the last sequence in the list \mathcal{L} .

2.3 Constant Amortized Time algorithms and principle

An exhaustive generating algorithm is said to run in *constant amortized time* (*CAT* for short) if the total amount of computation is proportional to the number of generated objects. And so, a CAT algorithm can be considered an efficient algorithm.

Ruskey and van Baronaigien [15] introduced three CAT properties, and proved that if a recursive generating procedure satisfies them, then it runs in constant amortized time (see also [16]). They called this general technique to prove the efficiency of a generating algorithm as *CAT principle*, and the involved properties are:

- 1. Every call of the procedure results in the output of at least one object;
- 2. Excluding the computation done by the recursive calls, the amount of computation of any call is proportional to the *degree of the call*, that is, the number of call initiated by the current call;
- 3. The number of calls of degree one, if any, is O(N), where N is the number of generated objects.

All the generating algorithms we present in this paper satisfy these three desiderata, and so they are efficient.

3 The Reflected Gray Code order for the sets SE_n , A_n , R_n , and S_n

3.1 The bound of Hamming distance between successive sequences in the lists \mathcal{SE}_n , \mathcal{A}_n , \mathcal{R}_n , and \mathcal{S}_n

Here we will show that the Hamming distance between two successive sequences in each of the mentioned lists is upper bounded by a constant, and so the lists are Gray codes.

Without another specification, X_n generically denotes one of the sets SE_n , A_n , R_n , or S_n ; and \mathcal{X}_n denotes its corresponding list in \prec order, that is, one of the lists \mathcal{SE}_n , \mathcal{A}_n , \mathcal{R}_n , or \mathcal{S}_n . Later in this section, Theorem 1 and Proposition 2 state that, in each case, the set X_n listed in \prec order yields a Gray code.

Lemma 1. If $s = s_1 s_2 ... s_n$ and $t = t_1 t_2 ... t_n$ are two sequences in X_n with $t = \operatorname{succ}_X(s)$ and k is the leftmost position where they differ, then $s_k = t_k + 1$ or $s_k = t_k - 1$.

Proof. Let $t = \operatorname{succ}_X(s)$ and k be the leftmost position where they differ. Let us suppose that $s_k < t_k$ and $s_k \neq t_k - 1$ (the case $s_k > t_k$ and $s_k \neq t_k + 1$ being similar). It is easy to check that

$$u = s_1 s_2 \dots s_{k-1} (s_k + 1) 0 \dots 0$$

belongs to X_n , and considering the definition of \prec order relation, it follows that $s \prec u \prec t$, which is in contradiction with $t = \operatorname{succ}_X(s)$, and the statement holds.

If $\mathbf{a} = a_1 a_2 \dots a_k \in X_k$, then for any n > k, \mathbf{a} is the prefix of at least one sequence in X_n , and we denote by $\mathbf{a} \mid \mathcal{X}_n$ the sublist of \mathcal{X}_n of all sequences having the prefix \mathbf{a} . Clearly, a list in \prec order for a set of sequences is a *prefix partitioned* list (all sequences with same prefix are contiguous), and for any $\mathbf{a} \in X_k$ and n > k, it follows that $\mathbf{a} \mid \mathcal{X}_n$ is a contiguous sublist of \mathcal{X}_n .

For a given $\mathbf{a} \in X_k$, the set of all x such that $\mathbf{a}x \in X_{k+1}$ is called the *defining set of the* prefix \mathbf{a} , and obviously $\mathbf{a}x$ is also a prefix of some sequences in X_n , for any n > k. We denote by

$$\omega_X(\boldsymbol{a}) = \max\{x \mid \boldsymbol{a}x \in X_{k+1}\}\tag{5}$$

the largest value in the defining set of \boldsymbol{a} . And if we denote $\omega_X(\boldsymbol{a})$ by M, then by Remark 1 we have

$$M = \operatorname{st}(\boldsymbol{a}) + 1$$
$$= \operatorname{st}(\boldsymbol{a}M).$$

And consequently,

$$\omega_X(\mathbf{a}M) = \operatorname{st}(\mathbf{a}M) + 1
= M + 1.$$
(6)

The next proposition gives the pattern of $s \in X_n$, if $s = \text{last}(a \mid \mathcal{X}_n)$ or $s = \text{first}(a \mid \mathcal{X}_n)$.

Proposition 1. Let k < n and $\mathbf{a} = a_1 a_2 \dots a_k \in X_k$. If $\mathbf{s} = \text{last}(\mathbf{a} \mid \mathcal{X}_n)$, then the pattern of \mathbf{s} is given by:

- if $\sum_{i=1}^{k} a_i$ is odd, then $\mathbf{s} = \mathbf{a}0...0$;
- if $\sum_{i=1}^{k} a_i$ is even and M is odd, then $\mathbf{s} = \mathbf{a}M0...0$;
- if $\sum_{i=1}^k a_i$ is even and M is even, then $\mathbf{s} = \mathbf{a}M(M+1)0...0$;

where M denotes $\omega_X(\boldsymbol{a})$.

Similar results hold for $s = \text{first}(a \mid \mathcal{X}_n)$ by replacing 'odd' by 'even', and vice versa, for the parity of $\sum_{i=1}^k a_i$.

Proof. Let $s = a_1 a_2 \dots a_k s_{k+1} \dots s_n = \text{last}(\boldsymbol{a} \mid \mathcal{X}_n)$.

If $\sum_{i=1}^k a_i$ is odd, then by considering the definition of \prec order, it follows that s_{k+1} is the smallest value in the defining set of \boldsymbol{a} , and so $s_{k+1} = 0$, and finally $\boldsymbol{s} = \boldsymbol{a}0\ldots 0$, and the first point holds. Now let us suppose that $\sum_{i=1}^k a_i$ is even. In this case s_{k+1} equals $\omega_X(\boldsymbol{a}) = M$, the largest value in the defining set of \boldsymbol{a} . When in addition M is odd, so is the summation of $\boldsymbol{a}M$, the length k+1 prefix of \boldsymbol{s} , and thus $\boldsymbol{s} = \boldsymbol{a}M0\ldots 0$, and the second point holds.

Finally, when M is even, then s_{k+2} is the largest value in the defining set of aM, which by relation (6) is M+1. In this case M+1 is odd, and thus s=aM(M+1)0...0, and the last point holds.

The proof for the case $s = \text{first}(a \mid \mathcal{X}_n)$ is similar.

By Proposition 1 above, we have the following:

Theorem 1. The lists A_n , R_n and S_n are 3-adjacent Gray codes.

Proof. Let \mathcal{X}_n be one of the lists \mathcal{A}_n , \mathcal{R}_n or \mathcal{S}_n , and $\mathbf{t} = \operatorname{succ}_X(\mathbf{s})$. Let k be the leftmost position where \mathbf{s} and \mathbf{t} differ, and let us denote by \mathbf{a} the length k prefix of \mathbf{s} and \mathbf{a}' that of \mathbf{t} ; so, $\mathbf{s} = \operatorname{last}(\mathbf{a} \mid \mathcal{X}_n)$ and $\mathbf{t} = \operatorname{first}(\mathbf{a}' \mid \mathcal{X}_n)$. If $k+3 \leq n$, then by Proposition 1, it follows that $s_{k+3} = s_{k+4} = \cdots = s_n = 0$ and $t_{k+3} = t_{k+4} = \cdots = t_n = 0$. So \mathbf{s} and \mathbf{t} differ only in position k, and possibly in position k+1 and in position k+2.

Now we show the adjacency, that is, if $k + 2 \le n$ and $s_{k+1} = t_{k+1}$ implies $s_{k+2} = t_{k+2}$. If $s_{k+1} = t_{k+1}$, by Lemma 1, it follows that the summation of the length k prefix of s and that of t have different parity, and two cases can occur:

- $s_{k+1} = t_{k+1} = 0$, and by Proposition 1, it follows that $s_{k+2} = t_{k+2} = 0$; or
- $s_{k+1} = t_{k+1} \neq 0$, and thus $s_{k+1} = t_{k+1} = \omega(\boldsymbol{a}) = \omega(\boldsymbol{a}')$. In this case, $\omega(\boldsymbol{a})$ either is odd and so $s_{k+2} = t_{k+2} = 0$, or is even and so $s_{k+2} = t_{k+2} = \omega(\boldsymbol{a}) + 1$. In both cases, $s_{k+2} = t_{k+2}$.

It is well known that the restriction of $\mathcal{G}_n(m)$ defined in relation (1) to any product space remains a 1-Gray code, see for example [13]. In particular, for $SE_n = \{0\} \times \{0,1\} \times \ldots \times \{0,1,\ldots,n-1\}$ we have the next proposition. Its proof is simply based on Lemma 1, Proposition 1, and on the additional remark: for any $\mathbf{a} \in SE_k$, k < n, it follows that $\omega_{SE}(\mathbf{a}) = k$.

Proposition 2. The list SE_n is 1-Gray code.

It is worth to mention that for any statistic st satisfying relations (2) and (3), the list in \prec order for the set of st-restricted growth sequences of length n is an at most 3-Gray code. Actually, the lists \mathcal{SE}_n , \mathcal{A}_n , \mathcal{R}_n , and \mathcal{S}_n are *circular* Gray codes, that is, the last and the first sequences in the list differ in the same way. Indeed, by the definition of \prec order, it follows that:

- first(\mathcal{X}_n) = 000 . . . 0;
- $\operatorname{last}(\mathcal{X}_n) = 010 \dots 0;$

where \mathcal{X}_n is one of the list \mathcal{SE}_n , \mathcal{A}_n , \mathcal{R}_n , or \mathcal{S}_n .

3.2 Generating algorithms for the lists SE_n , A_n , R_n , and S_n

Procedure Gen1 in Figure 1 is a general procedure generating exhaustively the list of st-restricted growth sequences, where st is a statistic satisfying relations (2) and (3). According to particular instances of the function Omega_X called by it (and so, of the statistic st), Gen1 produces specific st-restricted growth sequences, and in particular the lists \mathcal{SE}_n , \mathcal{A}_n , \mathcal{R}_n , and \mathcal{S}_n . From the length one sequence 0, Gen1 constructs recursively increasing length st-restricted growth sequences: for a given prefix $s_1s_2...s_k$ it produces all prefixes $s_1s_2...s_ki$, with i covering (in increasing or decreasing order) the defining set of $s_1s_2...s_k$; and eventually all length n st-restricted growth sequences. It has the following parameters:

- \bullet k, the position in the sequence s which is updated by the current call;
- x, belongs to the defining set of $s_1 s_2 \dots s_{k-1}$, and is the value to be assigned to s_k ;
- dir, the direction (ascending for $dir \mod 2 = 0$ and descending for $dir \mod 2 = 1$) to cover the defining set of $s_1 s_2 \ldots s_{k-1}$;
- v, the value of the statistic of the prefix $s_1 s_2 \dots s_{k-1}$ from which the value of the statistic of the current prefix $s_1 s_2 \dots s_k$ is computed. Remark that $v = \omega_X(s_1 s_2 \dots s_{k-1}) 1$.

Function Omega_X computes $\omega_X(s_1s_2...s_k)$ (see relation (5)), and the main call is Gen1(1,0,0,0). In Gen1 the amount of computation of each call is proportional with the degree of the call, and there are no degree one calls, and so it satisfies the CAT principle stated at the end of Section 2, and so it is an efficient generating algorithm. The computational tree of Gen1 producing the list \mathcal{A}_4 is given in Figure 2. Each node at level k, $1 \leq k \leq 4$, represents prefixes $s_1s_2...s_k$, and leaves sequences in \mathcal{A}_4 .

```
procedure Gen1(k, x, dir, v)
global: n, s;
s_k := x;
\quad \text{if } k=n
then Print s;
else u := \text{Omega}_X(v, k);
     if dir \mod 2 = 0
      then for i := 0 to u do
                Gen1(k+1, i, i, u-1);
           enddo
      else for i := u downto 0 do
                Gen1(k+1, i, i+1, u-1);
           enddo
endif
end.
                    (a)
```

```
\overline{\text{(i)}} function Omega_SE(w,p)
   return p;
   end.
(ii) function Omega\_A(w, p)
    if p > 1 and s_p > s_{p-1}
    then return w+2;
    else return w+1;
    endif
    end
(iii) function {\tt Omega\_R}(w,p)
    if p > 1 and s_p > w
    then return w+2;
    else return w+1;
    endif
    end.
(iv)function Omega_S(w,p)
   return s_p + 1;
                (b)
```

Figure 1: (a) Algorithm Gen1, generating the list \mathcal{X}_n ; (b) Particular function Omega_X called by Gen1, and returning the value for $\omega_X(s_1s_2...s_k)$, if X_n is one of the sets: (i) SE_n , (ii) A_n , (iii) R_n , and (iv) S_n .

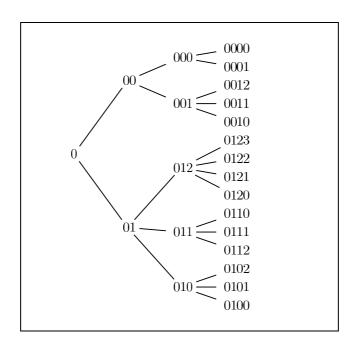


Figure 2: The tree induced by the initial call Gen1(1,0,0,0) for n=4 and generating the list \mathcal{A}_4 .

4 The Co-Reflected Gray Code order for the sets SE_n , A_n , R_n , and S_n

In this section we will consider, as in the previous one, the sets SE_n , A_n , R_n and S_n , but listed in \prec_c order. Our main goal is to prove that the obtained lists are Gray codes as well, and to develop generating algorithms for these lists. Recall that X_n generically denotes one of the sets SE_n , A_n , R_n , or S_n ; and let $\widetilde{\mathcal{X}}_n$ denote their corresponding list in \prec_c order, that are, $\widetilde{\mathcal{SE}}_n$, $\widetilde{\mathcal{A}}_n$, $\widetilde{\mathcal{R}}_n$, or $\widetilde{\mathcal{S}}_n$. Clearly, a set of sequences listed in \prec_c order is a *suffix partitioned* list, that is, all sequences with same suffix are contiguous, and such are the lists we consider here.

For a set X_n and a sequence $\mathbf{b} = b_{k+1}b_{k+2} \dots b_n$, we call \mathbf{b} an admissible suffix in X_n if there exists at least a sequence in X_n having suffix \mathbf{b} . For example, 124 is an admissible suffix in A_6 , because there are sequences in A_6 ending with 124, namely 012124 and 010124. On the other hand, 224 is not an admissible suffix in A_6 ; indeed, there is no length 6 ascent sequence ending with 224.

We denote by $\widetilde{\mathcal{X}}_n \mid \boldsymbol{b}$ the sublist of $\widetilde{\mathcal{X}}_n$ of all sequences having suffix \boldsymbol{b} , and clearly, $\widetilde{\mathcal{X}}_n \mid \boldsymbol{b}$ is a contiguous sublist of $\widetilde{\mathcal{X}}_n$. The set of all x such that $x\boldsymbol{b}$ is also an admissible suffix in X_n is called the *defining set of the suffix* \boldsymbol{b} .

For \prec order discussed in Section 3, the characterization of prefixes is straightforward: $a_1a_2...a_k$ is the prefix of some sequences in X_n , n > k, if and only if $a_1a_2...a_k$ is in X_k . And the defining set of the prefix $a_1a_2...a_k$ is $\{0,1,...,\mathsf{st}(a_1a_2...a_k)+1\}$. In the case of \prec_c order, it turns out that similar notions are more complicated: for example, 13 is an admissible suffix in A_5 , but 13 is not in A_2 ; and the defining set of the suffix 13 is $\{0,2\}$, because 013 and 213 are both admissible suffixes in A_5 , but 113 is not. See Table 4 for the set A_5 listed in \prec_c order.

4.1 Suffix expansion of sequences in the sets SE_n , A_n , R_n , and S_n

For a suffix partitioned list, we need to build st-restricted growth sequences under consideration from right to left, i.e., by expanding their suffix. For this purpose, we need the notions defined below

Definition 5. Let $b = b_{k+1}b_{k+2} \dots b_n$, $1 \le k < n$, an admissible suffix in X_n .

• $\alpha_X(b)$ is the set of all elements in the defining set of the suffix b. Formally:

$$\alpha_X(\boldsymbol{b}) = \{x \mid x\boldsymbol{b} \text{ is an admissible suffix in } X_n\},\$$

and for the empty suffix ϵ , $\alpha_X(\epsilon) = \{0, 1, \dots, n-1\}$.

• $\mu_X(\mathbf{b})$ is the minimum required value of the statistic defining the set X_n , and provided by a length (k+1) prefix of a sequence in X_n having suffix \mathbf{b} . Formally:

$$\mu_X(\mathbf{b}) = \min\{\mathsf{st}(s_1 s_2 \dots s_k b_{k+1}) \mid s_1 s_2 \dots s_k \mathbf{b} \in X_n\}.$$

Notice that $\mu_X(x\mathbf{b}) \in {\{\mu_X(\mathbf{b}) - 1, \mu_X(\mathbf{b}), x\}}$ for $x \in \alpha_X(\mathbf{b})$.

Remark 3. Let st be one of statistics asc, m or lv, and $s = s_1 s_2 \dots s_n$ be an st-restricted growth sequence. If there is a k < n such that $s_{k+1} = k$, then $s_i = i - 1$ for all $i, 1 \le i \le k$.

Proof. If $s_k < k-1$, then in each case for st, $\mathsf{st}(s_1 s_2 \dots s_k) < k-1$, which is in contradiction with $s_{k+1} = k$, and so $s_k = k-1$. Similarly, $s_{k-1} = k-2, \dots, s_2 = 1$, and $s_1 = 0$.

Under the conditions in the previous remark, $s_{k+1} = k$ imposes that all values at the left of k+1 in s are uniquely determined. As we will see later, in the induced tree of the generating algorithm, all descendants of a node with $s_{k+1} = k$ have degree one, and we will eliminate the obtained degree-one path in order not to alter the algorithm efficiency.

It is routine to check the following propositions. (Actually, Proposition 3 is a consequence of Remark 1.)

Proposition 3. Let X_n be one of the sets SE_n , A_n , R_n , or S_n . If $\mathbf{b} = b_{k+1}b_{k+2} \dots b_n$, $1 \le k < n$, is an admissible suffix in X_n , then $b_{k+1} \le \mu_X(\mathbf{b})$.

Proposition 4. Let Y_n be one of the sets A_n , R_n , or S_n . If $\mathbf{b} = b_{k+1}b_{k+2} \dots b_n$, $1 \le k < n$, is an admissible suffix in Y_n , then

- 1 if $\mathbf{b} = b_n$, that is, a length one admissible suffix, then $\mu_Y(\mathbf{b}) = b_n$;
- $2 \mu_Y(\mathbf{b}) = k \text{ if and only if } b_{k+1} = k;$
- 3 if $x\mathbf{b}$ is also an admissible suffix in Y_n (i.e., $x \in \alpha_Y(\mathbf{b})$) and $x \geq b_{k+1}$, then

$$\mu_Y(x\boldsymbol{b}) = \max\{x, \mu_Y(\boldsymbol{b})\}.$$

The following propositions give the values for $\alpha_X(\mathbf{b})$ and $\mu_X(x\mathbf{b})$, if X_n is one of the sets SE_n , A_n , R_n , or S_n . We do not provide the proofs for Propositions 5, 6, 11, and 12, because they are obviously based on the definition of the corresponding sequences.

Proposition 5. Let $b = \epsilon$ or $b = b_{k+1}b_{k+2}...b_n$ be an admissible suffix in SE_n . Then

$$\alpha_{SE}(\boldsymbol{b}) = \left\{ \begin{array}{ll} \{0,1,,\ldots,n-1\} & \text{if } \boldsymbol{b} = \epsilon, \\ \{0,1,\ldots,k-1\} & \text{otherwise.} \end{array} \right.$$

Proposition 6. Let $\mathbf{b} = b_{k+1}b_{k+2}\dots b_n$ be an admissible suffix in SE_n and $x \in \alpha_{SE}(\mathbf{b})$. Then

$$\mu_{SE}(x\boldsymbol{b}) = \mu_{SE}(\boldsymbol{b}) - 1.$$

Obviously, for a length one suffix $b = b_n$, it follows that $\mu_{SE}(b) = n - 1$.

Example 2. If $\mathbf{b} = \epsilon$, and n = 10, then $\alpha_{SE}(\mathbf{b}) = \{0, 1, \dots, 9\}$; and for $\mathbf{b} = b_{10}$, $b_{10} \in \{0, 1, \dots, 9\}$, it follows that $\mu_{SE}(x\mathbf{b}) = 9 - 1 = 8$, for all $x \in \alpha_{SE}(\mathbf{b})$.

Proposition 7. Let $\mathbf{b} = \epsilon$ or $\mathbf{b} = b_{k+1}b_{k+2} \dots b_n$ be an admissible suffix in A_n . Then

$$\alpha_{A}(\boldsymbol{b}) = \begin{cases} \{0, 1, \dots, n-1\} & \text{if } \boldsymbol{b} = \epsilon, \\ \{k-1\} & \text{if } \mu_{A}(\boldsymbol{b}) = k, \text{ or } \mu_{A}(\boldsymbol{b}) = k-1 \text{ and } b_{k+1} = 0, \\ \{0, 1, \dots, b_{k+1} - 1\} \cup \{k-1\} & \text{if } \mu_{A}(\boldsymbol{b}) = k-1 \text{ and } 0 < b_{k+1} < k, \\ \{0, 1, \dots, k-1\} & \text{if } \mu_{A}(\boldsymbol{b}) < k-1. \end{cases}$$

Proof. If $b = \epsilon$, the result is obvious.

For $\boldsymbol{b} \neq \epsilon$, let $x \in \alpha_A(\boldsymbol{b})$.

If $\mu_A(\mathbf{b}) = k$, by Proposition 4 point 2, $b_{k+1} = k$ and by Remark 3 we have x = k - 1.

If $\mu_A(\mathbf{b}) = k - 1$ and $b_{k+1} = 0$, then $\operatorname{\mathsf{asc}}(xb_{k+1}) = 0$, and so $\mu_A(x\mathbf{b}) = \mu_A(\mathbf{b}) = k - 1$, and again by Proposition 4 point 2 we have x = k - 1.

If $\mu_A(\mathbf{b}) = k - 1$ and $0 < b_{k+1} < k$, then there are two possibilities for $\mu_A(x\mathbf{b})$:

- $\mu_A(xb) = \mu_A(b) = k 1$, if $asc(xb_{k+1}) = 0$, and as above x = k 1;
- $\mu_A(xb) = \mu_A(b) 1 = k 2$, if $\operatorname{asc}(xb_{k+1}) = 1$. In this case $x \in \{0, 1, \dots, b_{k+1} 1\}$.

If $\mu_A(\mathbf{b}) < k-1$ (and consequently $0 \le b_{k+1} < k$), then there are two possibilities for $\mu_A(x\mathbf{b})$:

- $\mu_A(xb) = \mu_A(b)$, if $\operatorname{asc}(xb_{k+1}) = 0$, and we have $x \in \{b_{k+1}, b_{k+1} + 1, \dots, k-1\}$;
- $\mu_A(xb) = \mu_A(b) 1$, if $asc(xb_{k+1}) = 1$, and we have $x \in \{0, 1, \dots, b_{k+1} 1\}$.

Proposition 8. Let $\mathbf{b} = b_{k+1}b_{k+2}\dots b_n$ be an admissible suffix in A_n and $x \in \alpha_A(\mathbf{b})$. Then

$$\mu_A(x\boldsymbol{b}) = \begin{cases} x & \text{if } x \ge \mu_A(\boldsymbol{b}), \\ \mu_A(\boldsymbol{b}) & \text{if } b_{k+1} \le x < \mu_A(\boldsymbol{b}), \\ \mu_A(\boldsymbol{b}) - 1 & \text{if } x < b_{k+1}. \end{cases}$$

Proof. If $x \ge \mu_A(\mathbf{b})$, by Proposition 3 it follows that $x \ge b_{k+1}$, and by Proposition 4 point 3, that $\mu_A(x\mathbf{b}) = \max\{x, \mu_A(\mathbf{b})\} = x$.

If $b_{k+1} \le x < \mu_A(\boldsymbol{b})$, then, again by Proposition 4 point 3, it follows that $\mu_A(x\boldsymbol{b}) = \max\{x, \mu_A(\boldsymbol{b})\} = \mu_A(\boldsymbol{b})$.

If
$$x < b_{k+1}$$
, then $asc(xb_{k+1}) = 1$, so $\mu_A(xb) = \mu_A(b) - 1$.

Example 3. Let k = 5, n = 9, and $\mathbf{b} = b_6b_7b_8b_9 = 2050$ be an admissible suffix in A_9 . Clearly, $\mu_A(\mathbf{b})$, the minimum number of ascents in a prefix $s_1s_2...s_5b_6$ such that $s_1s_2...s_5\mathbf{b} \in A_9$, is 4. In this case, denoting s_5 by x, we have

- the set $\alpha_A(b)$ of all possible values for x is $\{0, 1, \dots, b_{k+1} 1\} \cup \{k 1\} = \{0, 1\} \cup \{4\}$.
- $\mu_A(x\mathbf{b}) = \mu_A(\mathbf{b}) 1 = 4 1 = 3$, if $x \in \{0, 1\}$; or $\mu_A(x\mathbf{b}) = \mu_A(\mathbf{b}) = 4$, if x = 4.

Proposition 9. Let $b = \epsilon$ or $b = b_{k+1}b_{k+2}...b_n$ be an admissible suffix in R_n . Then

$$\alpha_R(\boldsymbol{b}) = \begin{cases} \{0, 1, \dots, n-1\} & \text{if} \quad \boldsymbol{b} = \epsilon, \\ \{k-1\} & \text{if} \quad \mu_R(\boldsymbol{b}) = k, \text{ or } \mu_R(\boldsymbol{b}) = k-1 \text{ and } b_{k+1} < k-1, \\ \{0, 1, \dots, k-1\} & \text{if} \quad \mu_R(\boldsymbol{b}) = k-1 \text{ and } b_{k+1} = k-1, \text{ or } \mu_R(\boldsymbol{b}) < k-1. \end{cases}$$

Proof. If $b = \epsilon$, the result is obvious.

For $\boldsymbol{b} \neq \epsilon$, let $x \in \alpha_R(\boldsymbol{b})$.

If $\mu_R(\mathbf{b}) = k$, by Proposition 4 point 2, $b_{k+1} = k$ and by Remark 3 we have x = k - 1.

If $\mu_R(\boldsymbol{b}) = k-1$ and $b_{k+1} < k-1$, then the maximal value of the statistic **m** (defining the set R_n) of a length k+1 prefix ending with $b_{k+1} < k-1$ is k-1, and it is reached when x = k-1. If $\mu_R(\boldsymbol{b}) = k-1$ and $b_{k+1} = k-1$, then there are two possibilities for $\mu_R(x\boldsymbol{b})$:

- $\mu_R(x\mathbf{b}) = \mu_R(\mathbf{b}) = k 1$, and as above, this implies x = k 1;
- $\mu_R(xb) = \mu_R(b) 1 = k 2$, which implies $x \in \{0, 1, \dots, k 2\}$.

Finally, if $\mu_R(\mathbf{b}) < k-1$, then x can be any value in $\{0, 1, \dots, k-1\}$.

Proposition 10. Let $\mathbf{b} = b_{k+1}b_{k+2}\dots b_n$ be an admissible suffix in R_n and $x \in \alpha_R(\mathbf{b})$. Then

$$\mu_R(x\mathbf{b}) = \begin{cases} x & \text{if } x \ge \mu_R(\mathbf{b}), \\ \mu_R(\mathbf{b}) & \text{if } b_{k+1} \le x < \mu_R(\mathbf{b}) \text{ or } x < b_{k+1} < \mu_R(\mathbf{b}), \\ \mu_R(\mathbf{b}) - 1 & \text{if } x < b_{k+1} = \mu_R(\mathbf{b}). \end{cases}$$

Proof. The case $x \geq \mu_R(\mathbf{b})$ is analogous with the similar case in Proposition 8.

The next case is equivalent with $x < \mu_R(\mathbf{b})$ and $b_{k+1} < \mu_R(\mathbf{b})$, and since R_n corresponds to the statistic m, the result holds.

Finally, if
$$x < b_{k+1} = \mu_R(\boldsymbol{b})$$
, then $\mu_R(\boldsymbol{b}) = \mu_R(x\boldsymbol{b}) + 1$, and so $\mu_R(x\boldsymbol{b}) = \mu_R(\boldsymbol{b}) - 1$.

Example 4. Let k = 4, n = 7, and $\mathbf{b} = b_5b_6b_7 = 241$ be an admissible suffix in R_7 . It follows that $b_{k+1} = 2$, $\mu_R(\mathbf{b}) = 3$ and

- $\alpha_R(\mathbf{b}) = \{k-1\} = \{3\};$
- $\mu_R(xb) = x = 3$.

Proposition 11. Let $b = \epsilon$ or $b = b_{k+1}b_{k+2}...b_n$ be an admissible suffix in S_n . Then

$$\alpha_S(\mathbf{b}) = \begin{cases} \{0, 1, \dots, n-1\} & \text{if } \mathbf{b} = \epsilon, \\ \{k-1\} & \text{if } b_{k+1} = k, \\ \{C, C+1, \dots, k-1\} & \text{if } 0 \le b_{k+1} \le k-1, \end{cases}$$

where $C = \max\{0, b_{k+1} - 1\}$.

Since $\mu_S(\mathbf{b}) = b_{k+1}$, the next result follows:

Proposition 12. Let $\mathbf{b} = b_{k+1}b_{k+2}\dots b_n$ be an admissible suffix in S_n and $x \in \alpha_S(\mathbf{b})$. Then

$$\mu_S(x\boldsymbol{b}) = x.$$

Example 5. Let k = 6, n = 9, and $\mathbf{b} = b_7 b_8 b_9 = 457$ be an admissible suffix in S_9 . So we have $\mu_S(\mathbf{b}) = 4$, $C = \max\{0, b_{k+1} - 1\} = \max\{0, 3\} = 3$, and

- $\alpha_S(\mathbf{b}) = \{C, C+1, \dots, k-1\} = \{3, 4, 5\};$
- $\mu_S(xb) = x$, where $x \in \{3, 4, 5\}$.

4.2 The bound of Hamming distance between successive sequences in the lists $\widetilde{\mathcal{SE}}_n$, $\widetilde{\mathcal{A}}_n$, $\widetilde{\mathcal{R}}_n$, and $\widetilde{\mathcal{S}}_n$

Now we show that the Hamming distance between two successive sequences in the mentioned lists is upper bounded by a constant, which implies that the lists are Gray codes. These results are embodied in Theorems 2, 3 and 4.

4.2.1 The list $\widetilde{\mathcal{SE}}_n$

Theorem 2. The list $\widetilde{\mathcal{SE}}_n$ is 1-Gray code.

Proof. The result follows from the fact that the restriction of the 1-Gray code list $\mathcal{G}_n(n)$ to any product space remains a 1-Gray code (see [13]), in particular to the set

$$V_n = \vartheta_1 \times \vartheta_2 \times \ldots \times \vartheta_n$$

where

- $\vartheta_i = \{0, 1, ..., n-i\}$, if n is odd and i is even, or
- $\vartheta_i = \{i-1, i, ..., n-1\}$, if n is even, or n and i are both odd.

Then by applying to each sequence s in the list V_n the two transforms mentioned before Definition 2, namely:

• complementing each digit in s if n is even, or only digits in odd positions if n is odd, then

• reversing the obtained sequence,

the desired 1-Gray code for the set SE_n in \prec_c order is obtained.

4.2.2 The lists $\widetilde{\mathcal{A}}_n$ and $\widetilde{\mathcal{R}}_n$

The next proposition describes the pattern of $s = \operatorname{last}(\widetilde{\mathcal{Y}}_n \mid \boldsymbol{b})$ and $s = \operatorname{first}(\widetilde{\mathcal{Y}}_n \mid \boldsymbol{b})$, where $\widetilde{\mathcal{Y}}_n$ is one of the lists $\widetilde{\mathcal{A}}_n$ or $\widetilde{\mathcal{R}}_n$.

Proposition 13. Let Y_n be one of the sets A_n or R_n , and $\mathbf{b} = b_{k+1}b_{k+2}\dots b_n$ be an admissible suffix in Y_n . If $\mathbf{s} = \text{last}(\widetilde{\mathcal{Y}}_n \mid \mathbf{b})$ or $\mathbf{s} = \text{first}(\widetilde{\mathcal{Y}}_n \mid \mathbf{b})$, then \mathbf{s} has one of the following patterns:

- s = 012...(k-2)(k-1)b, or
- s = 012...(k-2)0b.

Proof. Let $\mathbf{s} = s_1 s_2 \dots s_k b_{k+1} b_{k+2} \dots b_n$. Since $\mathbf{s} = \operatorname{last}(\widetilde{\mathcal{Y}}_n \mid \mathbf{b})$ or $s = \operatorname{first}(\widetilde{\mathcal{Y}}_n \mid \mathbf{b})$, according to $\alpha_Y(\mathbf{b})$ given in Propositions 7 and 9, it follows that $s_k \in \{0, k-1\}$. In other words, s_k is either the smallest or the largest value in $\alpha_Y(\mathbf{b})$.

If $s_k = k - 1$, then by Remark 3 we have $\mathbf{s} = 012 \dots (k - 2)(k - 1)\mathbf{b}$.

If $s_k = 0$, then considering the definition of \prec_c order we have either

- $s = \text{first}(\widetilde{\mathcal{Y}}_n \mid b)$ and $\sum_{i=k+1}^n b_i + (n-k)$ is odd, or
- $s = \text{last}(\widetilde{\mathcal{Y}}_n \mid b)$ and $\sum_{i=k+1}^n b_i + (n-k)$ is even.

For the first case, again by the definition of \prec_c order, it follows that s_{k-1} must be the largest value in $\alpha_Y(0\mathbf{b})$, and so $s_{k-1} = k-2$, and by Remark 3, $\mathbf{s} = 012...(k-2)0\mathbf{b}$. Similarly, the same result is obtained for the second case.

A direct consequence of the previous proposition is the next theorem.

Theorem 3. The lists $\widetilde{\mathcal{A}}_n$ and $\widetilde{\mathcal{R}}_n$ are 2-adjacent Gray codes.

Proof. Let $s, t \in Y_n$, with $t = \widetilde{\operatorname{succ}}_Y(s)$. If k+1 is the rightmost position where s and t differ, then there are admissible suffixes $b = b_{k+1}b_{k+2}\dots b_n$ and $b' = b'_{k+1}b_{k+2}\dots b_n$ in Y_n such that $s = \operatorname{last}(\widetilde{\mathcal{Y}}_n | b)$ and $t = \operatorname{first}(\widetilde{\mathcal{Y}}_n | b')$.

By Proposition 13, s has pattern

$$012...(k-2)(k-1)\mathbf{b}$$
, or $012...(k-2)0\mathbf{b}$;

and t has pattern

$$012...(k-2)(k-1)\mathbf{b}'$$
, or $012...(k-2)0\mathbf{b}'$.

And in any case, s and t differ in position k+1 and possibly in position k.

4.2.3 The list $\widetilde{\mathcal{S}}_n$

The next proposition gives the pattern of $\operatorname{last}(\widetilde{\mathcal{S}}_n \mid \boldsymbol{b})$ and $\operatorname{first}(\widetilde{\mathcal{S}}_n \mid \boldsymbol{b})$ for an admissible suffix \boldsymbol{b} in S_n .

Proposition 14. Let $\mathbf{b} = b_{k+1}b_{k+2} \dots b_n$ be an admissible suffix in S_n . If $\mathbf{s} = \operatorname{last}(\widetilde{S}_n \mid \mathbf{b})$, then the pattern of \mathbf{s} is given by:

• if $b_{k+1} = k$ or $\sum_{i=k+1}^{n} b_i + (n-k)$ is odd, then

$$s = 012 \dots (k-2)(k-1)b;$$

• if $b_{k+1} < k$ and $\sum_{i=k+1}^{n} b_i + (n-k)$ is even, and either $b_{k+1} = 0$ or b_{k+1} is odd, then

$$s = 012...(k-2)(\max\{0, b_{k+1} - 1\})b;$$

• if $b_{k+1} < k$ and $\sum_{i=k+1}^{n} b_i + (n-k)$ is even, and $b_{k+1} > 0$ is even, then

$$s = 012...(k-3)(b_{k+1}-2)(b_{k+1}-1)\boldsymbol{b}.$$

Similar results hold for $s = \operatorname{first}(\widetilde{S}_n \mid b)$ by replacing 'odd' by 'even', and vice versa, for the parity of $\sum_{i=k+1}^n b_i + (n-k)$.

Proof. Let $\mathbf{s} = s_1 s_2 \dots s_k b_{k+1} \dots b_n = \operatorname{last}(\widetilde{\mathcal{S}}_n \mid \mathbf{b}).$

If $b_{k+1} = k$ or $\sum_{i=k+1}^{n} b_i + (n-k)$ is odd, then s_k is the largest value in $\alpha_S(\boldsymbol{b})$, so $s_k = k-1$, and by Remark 3, $s_i = i-1$ for $1 \le i \le k$. So the first case holds.

If $\sum_{i=k+1}^{n} b_i + (n-k)$ is even and $b_{k+1} < k$, then s_k is the smallest value in $\alpha_S(\boldsymbol{b})$, namely $\max\{0, b_{k+1} - 1\}$, which is even if $b_{k+1} = 0$ or b_{k+1} is odd. Thus, by the definition of \prec_c order, s_{k-1} is the largest value in $\alpha_S(s_k \boldsymbol{b})$, which is k-2, and by Remark 3, the second case holds.

For the last case, as above, $s_k = \max\{0, b_{k+1} - 1\}$, and considering $b_{k+1} > 0$ and even, it follows that $s_k = b_{k+1} - 1$ is odd. Thus s_{k-1} is the minimal value in $\alpha_S(s_k \mathbf{b})$, that is $b_{k+1} - 2$, which in turns is even, and the last case holds.

The proof for the case $s = \operatorname{first}(\widetilde{\mathcal{S}}_n \mid b)$ is similar.

Theorem 4. The list $\widetilde{\mathcal{S}}_n$ is 3-adjacent Gray codes.

Proof. Let $\mathbf{t} = \widetilde{\operatorname{succ}}_S(\mathbf{s})$, and k+1 be the rightmost position where \mathbf{s} and \mathbf{t} differ. Let us denote by \mathbf{b} the length (n-k) suffix of \mathbf{s} and \mathbf{b}' that of \mathbf{t} ; so, $\mathbf{s} = \operatorname{last}(\mathbf{b} \mid \mathcal{S}_n)$ and $\mathbf{t} = \operatorname{first}(\mathbf{b}' \mid \mathcal{S}_n)$. It follows by Proposition 14, that $s_i = t_i = i-1$ for all $i \leq k-2$, and so the other differences possibly occur in position k and in position k-1.

Considering all valid combinations for s and t as given in Proposition 14, the proof of the adjacency is routine, and based on the following: $s_k \neq t_k$ if and only if $s_{k-1} \neq t_{k-1}$. It follows that s and t differ in one position, or three positions which are adjacent.

In addition, the lists $\widetilde{\mathcal{A}}_n$, $\widetilde{\mathcal{R}}_n$, and $\widetilde{\mathcal{S}}_n$, are circular Gray codes. This is a consequence of the following remarks based on Propositions 13 and 14:

- first($\widetilde{\mathcal{Y}}_n$) = 012...(n-2)(n-1);
- $\operatorname{last}(\widetilde{\mathcal{Y}}_n) = 012\dots(n-2)0;$

where $\widetilde{\mathcal{Y}}_n$ is one of the lists $\widetilde{\mathcal{A}}_n$, $\widetilde{\mathcal{R}}_n$, or $\widetilde{\mathcal{S}}_n$.

4.3 Generating algorithm for $\widetilde{\mathcal{SE}}_n$, $\widetilde{\mathcal{A}}_n$, $\widetilde{\mathcal{R}}_n$, and $\widetilde{\mathcal{S}}_n$

Here we explain algorithm Gen2 in Figure 3 which generates suffix partitioned Gray codes for restricted growth sequences; according to particular instances of the functions called by it, Gen2 produces the list $\widetilde{\mathcal{SE}}_n$, $\widetilde{\mathcal{A}}_n$, $\widetilde{\mathcal{R}}_n$, or $\widetilde{\mathcal{S}}_n$. Actually, for convenience, Gen2 produces length (n+1) sequences $\mathbf{s} = s_1 s_2 \dots s_{n+1}$ with $s_{n+1} = 0$, and so, neglecting the last value in each sequence \mathbf{s} the desired list is obtained. Notice that with this dummy value for s_{n+1} we have $\mu_X(s_k s_{k+1} \dots s_n) = \mu_X(s_k s_{k+1} \dots s_n 0)$, for $k \leq n$, and similarly for α_X .

In Gen2, the sequence s is a global variable, and initialized by 01...(n-1)0, which is the first length n sequence in \prec_c order, followed by a 0; and the main call is Gen2(n+1,0,0,0). Procedure Gen2 has the following parameters (the first three of them are similar with those of procedure Gen1):

- \bullet k, the position in the sequence s which is updated by the current call;
- x, the value to be assigned to s_k ;
- dir, gives the direction in which s_{k-1} covers $\alpha_X(s_k s_{k+1} \dots s_n 0)$, the defining set of the current suffix;
- v, the value of $\mu_X(s_{k+1}...s_n0)$.

The functions called by Gen2 are given in Figures 4 and 5. They are principally based on the evaluation of α_X and μ_X for the current suffix of s, and are:

- Mu_X(k, x, v) returns the value of $\mu_X(s_k s_{k+1} \dots s_n 0)$, with $x = s_k$.
- IsDegreeOne_X(k) stops the recursive calls when $\alpha_X(s_ks_{k+1}\dots s_n0)$ has only one element, namely k-2. In this case, by Remark 3 the sequence is uniquely determined by the current suffix, and this prevents Gen2 to produce degree one calls. In addition, IsDegreeOne_X(k) sets appropriately d-1 values at the left of s_k , where d is the upper bound of the Hamming distance in the list (changes at the left of s_1 are considered with no effect). This can be considered as a Path Elimination Technique [16].

- Lowest_X(k) is called when IsDegreeOne_X returns false, and gives the lowest value in $\alpha_X(s_ks_{k+1}\dots s_n0)$.
- SecLargest_X(k, u), is called when IsDegreeOne_X returns false, and gives the second largest value in $\alpha_X(s_k s_{k+1} \dots s_n 0)$ (the largest value being always k-2).

By this construction, algorithm Gen2 has no degree one calls and it satisfies the CAT principle. Figure 6 shows the tree induced by the algorithm when generates $\widetilde{\mathcal{A}}_4$.

```
procedure Gen2(k, x, dir, v)
global n, s;
s_k := x;
u := \mathbf{Mu} X(k, x, v);
if IsDegreeOne\_X(k, u)
then Print s;
else c:=Lowest_X(k);
      d:=SecLargest\_X(k,u);
     \text{if } dir \!\!\!\mod 2 = 1
      then for i := c to d do
                 Gen2(k-1,i,i,u);
            enddo
      endif
      Gen2(k-1, k-2, (k-1-dir) \mod 2, u);
      \text{if } dir \!\!\!\mod 2 = 0
      then for i := d downto c do
                 Gen2(k-1, i, i+1, u);
            enddo
      endif
endif
end.
```

Figure 3: Algorithm Gen2, generating the list $\widetilde{\mathcal{X}}_n$.

```
function Mu\_A(m, i, w)
                                        function Mu_R(m, i, w)
  \texttt{if} \ i \geq w
                                           \text{if } i \geq w
  then return i;
                                          then return i;
  else if i \geq s_{m+1}
                                           else if s_{m+1} < w
        then return w;
                                                then return w;
        else return w-1;
                                                 else return w-1;
        endif
                                                 endif
  endif
                                           endif
end.
                                        end.
function IsDegreeOne_A(m, v)
                                        function IsDegreeOne_\mathbb{R}(m,v)
  if v=m-1 or
                                           \quad \text{if } v=m-1 \text{ or } \\
      (v=m-2 \text{ and } s_m=0)
                                              (v = m - 2 \text{ and } s_m < m - 2)
  then s_{m-1} := m-2;
                                          then s_{m-1} := m-2;
        return true;
                                                return true;
  else return false;
                                          else return false;
  endif
                                           endif
end.
                                        end.
function Lowest_A(m)
                                        function Lowest_\mathbb{R}(m)
  return 0;
                                          return 0;
function SecLargest_A(m, w)
                                        function SecLargest_\mathbb{R}(m, w)
  if w=m-2 and s_m>0
                                          return m-3;
     and s_m < m-1
                                        end.
  then return s_m-1;
  else return m-3;
  endif
end.
                (a)
                                                         (b)
```

Figure 4: Particular functions called by Gen2, generating the lists: (a) $\widetilde{\mathcal{A}}_n$, and (b) $\widetilde{\mathcal{R}}_n$.

```
function Mu\_S(m, i, w)
 return i;
end.
function IsDegreeOne_S(m, v)
  if s_m = m-1
  then s_{m-1} := m-2;
       s_{m-2} := m-3;
       return true;
  else return false;
  endif
end.
function Lowest_S(m)
  if s_m>1 and s_m\leq m-2
 then return s_m-1;
  else return 0;
  endif
end.
function SecLargest_S(m, w)
 return m-3;
end.
```

```
function Mu_SE(m,i,w)

if m=n

then return n-1;

else return w-1;

end.

function IsDegreeOne_SE(m,v)

if m=2

then return true;

else return false;

end.

function Lowest_SE(m)

return 0;

end.

function SecLargest_SE(m,w)

return m-3;

end.
```

Figure 5: Particular functions called by Gen2, generating the lists: (a) $\widetilde{\mathcal{S}}_n$, and (b) $\widetilde{\mathcal{SE}}_n$.

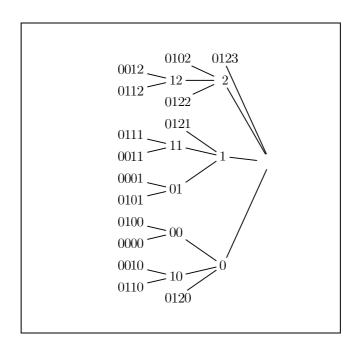


Figure 6: The tree induced by the call Gen2(5,0,0,0), generating the list $\widetilde{\mathcal{A}}_4$.

5 Final remarks

We conclude this paper by comparing for each of the sets SE_n , A_n , R_n and S_n the prefix partitioned Gray codes induced by \prec order and the suffix partitioned one induced by \prec order. This can be done by comparing the Hamming distance between all pairs of successive sequences either in the worst case or in average.

Table 1 summarizes Theorems 1, 3, and 4, and Proposition 2, and gives the upper bound of the Hamming distance (that is, the worst case Hamming distance) for the two order relations. It shows that for the sets SE_n and S_n these relations have same performances, and for the sets A_n and R_n , \prec_c order induces more restrictive Gray codes.

Set	\prec (RGC)	$\prec_c \text{(Co-RGC)}$
	order	order
$\overline{SE_n}$	1	1
A_n	3	2
R_n	3	2
S_n	3	3

Table 1: The bound of the Hamming distance between two successive sequences in \prec and \prec_c orders.

For a list \mathcal{L} of sequences, the average Hamming distance is defined as

$$\frac{\sum d(\boldsymbol{s}, \boldsymbol{t})}{N-1},$$

where the summation is taken over all s in \mathcal{L} , except its last element, t is the successor of s in \mathcal{L} , d is the Hamming distance, and N the number of sequences in \mathcal{L} .

Surprisingly, despite \prec_c order has same or better performances in terms of worst case Hamming distance, if we consider the average Hamming distance, numerical evidences show that \prec order is 'more optimal' than \prec_c order on A_n , R_n $(n \ge 5)$, and S_n $(n \ge 6)$. And this phenomenon strengthens for large n; see Table 2.

	_	(RGC	C) orde	er	\prec_c (Co-RGC) order					
n	\mathcal{SE}_n	\mathcal{A}_n	\mathcal{R}_n	\mathcal{S}_n	$\widetilde{\mathcal{SE}}_n$	$\widetilde{\mathcal{A}}_n$	$\widetilde{\mathcal{R}}_n$	$\widetilde{{\mathcal S}}_n$		
4	1	1.21	1.21	1.31	1	1.14	1.14	1.15		
5	1	1.13	1.12	1.29	1	1.19	1.18	1.24		
6	1	1.09	1.07	1.27	1	1.23	1.20	1.31		
7	1	1.06	1.06	1.26	1	1.25	1.22	1.35		
8	1	1.04	1.04	1.25	1	1.26	1.23	1.37		
9	1	1.03	1.03	1.24	1	1.28	1.24	1.39		
10	1	1.02	1.03	1.23	1	1.28	1.24	1.41		

Table 2: The average Hamming distance for \prec order and \prec_c order.

Algorithmically, \prec_c order has the advantage that its corresponding generating algorithm, Gen2, is more appropriate to be parallelized than its \prec order counterpart, Gen1. Indeed, the main call of Gen2 produces n recursive calls (compare to two recursive calls produced by the

main call of Gen1), and so we can have more parallelized computations; and this is more suitable for large n. See Figure 2 and 6 for examples of computational trees.

Finally, it will be of interest to explore order relation based Gray codes for restricted growth sequences defined by statistics other than those considered in this paper. In this vein we suggest the following conjecture, checked by computer for $n \leq 10$, and concerning descent sequences (defined similarly with ascent sequences in Section 2).

Conjecture 1. The set of length n descent sequences listed in \prec_c order is a 4-adjacent Gray code.

Appendix

Sequence	S_5	R_5	A_5	Sequence	S_5	R_5	A_5	Sequence	S_5	R_5	A_5
00000	√	√	√	01233	√	√	√	01111	√	√	√
00001	✓	\checkmark	✓	01234	\checkmark	✓	\checkmark	01110	✓	✓	\checkmark
00012	✓	\checkmark	\checkmark	01223	\checkmark	✓	✓	01120	✓	✓	\checkmark
00011	✓	\checkmark	\checkmark	01222	\checkmark	✓	✓	01121	✓	✓	\checkmark
00010	✓	\checkmark	✓	01221	\checkmark	✓	\checkmark	01122	✓	✓	\checkmark
00123	✓	\checkmark	✓	01220	\checkmark	✓	\checkmark	01123	✓	✓	\checkmark
00122	✓	\checkmark	✓	01210	\checkmark	✓	\checkmark	01023		✓	\checkmark
00121	✓	\checkmark	✓	01211	\checkmark	✓	\checkmark	01022		✓	\checkmark
00120	✓	\checkmark	\checkmark	01212	\checkmark	✓	✓	01021		✓	✓
00110	✓	\checkmark	\checkmark	01213		✓	\checkmark	01020		✓	\checkmark
00111	✓	\checkmark	✓	01203		✓	\checkmark	01010	✓	✓	\checkmark
00112	✓	\checkmark	✓	01202		✓	\checkmark	01011	✓	✓	\checkmark
00102		\checkmark	✓	01201	\checkmark	✓	\checkmark	01012	✓	✓	\checkmark
00101	✓	\checkmark	✓	01200	\checkmark	✓	\checkmark	01013			\checkmark
00100	✓	\checkmark	✓	01100	\checkmark	✓	\checkmark	01002		✓	\checkmark
01230	✓	\checkmark	✓	01101	✓	✓	\checkmark	01001	✓	✓	\checkmark
01231	✓	\checkmark	✓	01102		✓	\checkmark	01000	✓	✓	✓
01232	√	√	✓	01112	√	✓	✓				

Table 3: The sets $S_5,\,R_5,\,$ and A_5 listed in \prec order.

Sequence	S_5	R_5	A_5	Sequence	S_5	R_5	A_5	Sequence	S_5	R_5	A_5
01234	√	√	√	01222	√	√	√	01101	√	√	√
01233	✓	✓	✓	01122	✓	✓	✓	01201	\checkmark	✓	✓
01023		✓	✓	00122	✓	✓	✓	01200	\checkmark	✓	✓
00123	✓	✓	✓	01022		✓	✓	01100	\checkmark	✓	✓
01123	✓	✓	✓	01232	\checkmark	✓	\checkmark	00100	\checkmark	✓	✓
01223	✓	✓	✓	01231	✓	✓	✓	00000	\checkmark	✓	✓
01213		✓	✓	01021		✓	✓	01000	\checkmark	✓	✓
01013			✓	00121	\checkmark	✓	\checkmark	01010	\checkmark	✓	✓
01203		✓	✓	01121	\checkmark	✓	\checkmark	00010	\checkmark	✓	✓
01202		✓	✓	01221	\checkmark	✓	\checkmark	00110	\checkmark	✓	✓
01102		✓	✓	01211	\checkmark	✓	\checkmark	01110	\checkmark	✓	✓
00102		✓	✓	01111	\checkmark	✓	\checkmark	01210	\checkmark	✓	✓
01002		✓	✓	00111	\checkmark	✓	\checkmark	01220	\checkmark	✓	✓
01012	✓	✓	✓	00011	\checkmark	✓	\checkmark	01120	\checkmark	✓	✓
00012	✓	✓	✓	01011	\checkmark	✓	\checkmark	00120	\checkmark	✓	✓
00112	✓	✓	✓	01001	\checkmark	✓	✓	01020		✓	✓
01112	✓	✓	✓	00001	\checkmark	✓	\checkmark	01230	\checkmark	✓	✓
01212	✓	\checkmark	\checkmark	00101	\checkmark	√	√				

Table 4: The sets $S_5,\,R_5,\,{\rm and}~A_5$ listed in \prec_c order.

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