# Well-orders in the transfinite Japaridze algebra 

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#### Abstract

This paper studies the transfinite propositional provability logics GLP $\Lambda_{\Lambda}$ and their corresponding algebras. These logics have for each ordinal $\xi<$ $\Lambda$ a modality $\langle\alpha\rangle$. We will focus on the closed fragment of GLP $\Lambda$ (i.e., where no propositional variables occur) and worms therein. Worms are iterated consistency expressions of the form $\left\langle\xi_{n}\right\rangle \ldots\left\langle\xi_{1}\right\rangle \top$. Beklemishev has defined well-orderings $<_{\xi}$ on worms whose modalities are all at least $\xi$ and presented a calculus to compute the respective order-types.

In the current paper we present a generalization of the original $<_{\xi}$ orderings and provide a calculus for the corresponding generalized ordertypes $o_{\xi}$. Our calculus is based on so-called hyperations which are transfinite iterations of normal functions.

Finally, we give two different characterizations of those sequences of ordinals which are of the form $\left\langle o_{\xi}(A)\right\rangle_{\xi \in \mathrm{On}}$ for some worm $A$. One of these characterizations is in terms of a second kind of transfinite iteration called cohyperation.


## 1 Introduction

In this paper we study transfinite propositional provability logics $G L P_{\Lambda}$ and their corresponding algebras. For an ordinal $\Lambda$, the transfinite provability logic $G L P_{\Lambda}$ is a polymodal version of Gödel-Löb's provability logic GL where for each ordinal $\alpha<\Lambda$ the logic contains a modality [ $\alpha$ ].

These logics have been studied quite intensively lately and possess a very rich structure in various aspects. To mention just a few, it is a natural example of a logic that is not complete for its Kripke semantics but is complete for its class
of topological models [8, 10, 13]. However, for natural topologies on intervals of ordinals the completeness for these spaces is independent of ZFC giving rise to various interesting set-theoretical questions ([1, 6, 11]).

By GLP we denote the class-sized logic that extends all GLP ${ }_{\Lambda}$. In this paper we shall focus on GLP ${ }^{0}$-the closed fragment- of this class-size logic. This is the fragment that does not contain any propositional variables hence is generated by $T$, Boolean connectives and modalities only. Within $G L P^{0}$ we consider the class $\mathbb{W}$ of so-called worms. These are iterated consistency statements, that is, expressions of the form $\left\langle\alpha_{n}\right\rangle \ldots\left\langle\alpha_{1}\right\rangle \top$. By $\mathbb{W}_{\alpha}$ we denote the class of worms where each occurring modality is at least $\alpha$.

In 4. 9] it has been shown that $\mathbb{W}_{\alpha}$ can be well-ordered by defining $A<_{\alpha}$ $B: \Leftrightarrow \mathrm{GLP} \vdash B \rightarrow\langle\alpha\rangle A$. For $A \in \mathbb{W}_{\alpha}$, by $\check{o}_{\alpha}(A)$ we denote the order-type of $A$ in $\left\langle\mathbb{W}_{\alpha},<_{\alpha}\right\rangle$. It is most natural to consider these well-orders as sub-structures of the algebras that correspond to GLP which are often called Japaridze algebras.

In this paper we study the ordering $<_{\alpha}$ as an ordering on all of $\mathbb{W}$. We will see that $<_{\alpha}$ no longer defines a linear order on $\mathbb{W}$; however, we prove that it does define a well-founded relation and denote the corresponding order-types by $o_{\alpha}(A)$. We show how the $o$ order-types can be recursively reduced to the $\check{o}$ ordertypes, and in fact $o_{\xi}(A)=\check{o}_{\xi}(A)$ whenever $A \in \mathbb{W}_{\xi}$. Based on this reduction we are able to give a calculus for the ordinal sequences $\vec{o}(A):=\left\langle o_{\xi}(A)\right\rangle_{\xi \in \text { On }}$. That is, we show how to compute $\vec{o}(A)$ for a given worm $A$ and prove which ordinal sequences are attained as $\vec{o}(A)$ for some $A$. The calculus we give is based on hyperexponentials and hyperlogarithms, which are operations on ordinals related to Veblen progressions and presented in detail in [15].

Our calculus for $o_{\xi}$ is different from the $\check{o}_{\xi}$ calculus as presented in [4] in at least three essential ways. First, the definition of $o_{\xi}$ is a genuine generalization of $\check{o}_{\xi}$ and $\check{o}_{\xi}$ can be obtained a special case. Second, our presentation does not work with normal forms on worms, either in the presentation of the calculus or in any of the proofs. Finally, our calculus uses hyperexponentials whereas the calculus in 4 used Veblen functions.

It is known that the sequences $\vec{o}(A)$ can be interpreted proof-theoretically. In particular, $\mathrm{GLP}_{\omega}$ has been used to perform a $\Pi_{1}^{0}$-ordinal analysis of Peano Arithmetic (PA) and related systems (3). Meanwhile, it has been shown in [21] that there exists a close relation between Turing progressions of first-order theories and the sequences $\vec{o}(A)$. There are ongoing efforts to carry these techniques to stronger theories using transfinite provability operators [2, 12, 17, 23]

Furthermore, in [20] it is discussed how the sequences $\vec{o}(A)$ unveil important information about Kripke and other semantics for the closed fragment of $\mathrm{GLP}_{\Lambda}$ as presented in [14, 16].

Layout. Section 2 introduces the logics GLP $_{\Lambda}$ and their fragments. In Sections 3 and 4 we present the linear orders $<_{\alpha}$ and their corresponding order-types $o_{\xi}$ on substructures of the Japaridze algebra which are a central focus of this paper. We show how the computation of $o_{\xi}$ can be reduced to the computation of $o_{0}$.

In Section 5, we give a calculus to compute $o_{0}$. The calculus that we present
is actually a reduction to what we call worm enumerators $\sigma^{\xi}$. It is in Section 6 where we provide a calculus to compute the worm enumerators $\sigma^{\xi}$. Section 77reviews the notions of hyperations and cohyperations from ([15]) to show how the worm enumerators $\sigma^{\xi}$ are hyperations of ordinal exponentiation. Finally, in Section 8 we set the hyperations and cohyperations to work to obtain full characterizations of the sequences $\left\langle o_{\xi}(A)\right\rangle_{\xi \in \mathrm{On}}$.

The current paper is based on material which originally appeared in the unpublished manuscripts [18] and [20]. Portions of the latter were reported in 21].

## 2 Provability logics and the Reflection Calculus

In this section we introduce the logics $\mathrm{GLP}_{\Lambda}$ and its fragments, as well as fixing some notation.

### 2.1 The logics GLP ${ }_{\Lambda}$

The language of $\mathrm{GLP}_{\Lambda}$ is that of propositional modal logic that contains for each $\alpha<\Lambda$ a unary modal operator $[\alpha]$. In the definition below the $\alpha$ and $\beta$ range over ordinals and the $\psi$ and $\chi$ over formulas in the language of $\mathrm{GLP}_{\Lambda}$.

Definition 2.1. For $\Lambda$ an ordinal, the logic $\mathrm{GLP}_{\Lambda}$ is the propositional normal modal logic that has for each $\alpha<\Lambda$ a modality $[\alpha]$ and is axiomatized by all propositional logical tautologies together witht the following schemata:

$$
\begin{array}{ll}
{[\alpha](\chi \rightarrow \psi) \rightarrow([\alpha] \chi \rightarrow[\alpha] \psi),} & \\
{[\alpha]([\alpha] \chi \rightarrow \chi) \rightarrow[\alpha] \chi,} & \\
\langle\alpha\rangle \psi \rightarrow[\beta]\langle\alpha\rangle \psi & \text { for } \alpha<\beta, \\
{[\alpha] \psi \rightarrow[\beta] \psi} & \text { for } \alpha \leq \beta .
\end{array}
$$

The rules of inference are Modus Ponens and necessitation for each modality: $\frac{\psi}{[\alpha] \psi}$. By GLP we denote the class-size logic that has a modality $[\alpha]$ for each ordinal $\alpha$ and all the corresponding axioms and rules. The classic Gödel-Löb provability logic GL is denoted by $\mathrm{GLP}_{1}$.

### 2.2 Japaridze algebras

The relations $<_{\alpha}$ do not give proper linear orders on $\mathbb{W}_{\alpha}$, given that different worms may be equivalent in GLP and hence undistinguishable in the ordering. We remedy this by passing to the Lindenbaum algebra of GLP - that is, the quotient of the language of GLP modulo provable equivalence.

This algebra is a Japaridze algebra, as described below:
Definition 2.2 (Japaridze algebra). A Japaridze algebra is a structure

$$
\mathcal{J}=\left\langle D,\{[\alpha]\}_{\alpha<\Lambda}, \wedge, \neg, 0,1\right\rangle
$$

such that

1. $\langle D, \wedge, \neg, 0,1\rangle$ is a Boolean algebra,
2. $[\alpha] 1=1$ for all $\alpha<\Lambda$,
3. $[\alpha](x \rightarrow y) \leq[\alpha] x \rightarrow[\alpha] y$ for all $\alpha<\Lambda, x, y \in D$,
4. $[\alpha]([\alpha] x \rightarrow x) \leq[\alpha] x$ for all $\alpha<\Lambda, x \in D$,
5. $[\alpha] x \leq[\beta] x$ for all $\alpha \leq \beta<\Lambda, x \in D$ and,
6. $\langle\alpha\rangle x \leq[\beta]\langle\alpha\rangle x$ for all $\alpha<\beta<\Lambda, x \in D$,
where $\langle\alpha\rangle, \rightarrow$ are defined in the usual way.
It is in these algebras that the partial orders $<_{\alpha}$ we have described naturally reside. However, as we are mainly interested in formulas that fall within a specific fragment of our language, we will work throughout the paper in a restricted calculus.

### 2.3 The Reflection Calculus

In [12, 2, 7] Dashkov and Beklemishev introduced a calculus for reasoning about a fragment of the language of GLP and called it the Reflection Calculus (RC). (Closed) formulas of RC are built from the grammar

$$
\top|\phi \wedge \psi| \lambda \phi,
$$

where $\lambda$ is an ordinal and $\phi, \psi$ are formulas of $\mathrm{RC} ; \lambda \phi$ is interpreted as $\langle\lambda\rangle \phi$, but as RC does not contain operators of the form $[\lambda]$, the brackets become unnecessary. RC derives sequents of the form $\phi \vdash \psi$, given by the following rules and axioms:

$$
\begin{array}{lcc}
\phi \vdash \phi & \phi \vdash \top & \frac{\phi \vdash \psi \psi \vdash \chi}{\phi \vdash \chi} \\
\phi \wedge \psi \vdash \phi & \phi \wedge \psi \vdash \psi & \frac{\phi \vdash \psi \phi \vdash \chi}{\phi \vdash \psi \wedge \chi} \\
\alpha \alpha \phi \vdash \alpha \phi & \frac{\phi \vdash \psi}{\alpha \phi \vdash \alpha \psi} & \\
\beta \phi \vdash \alpha \phi & \beta \phi \wedge \alpha \psi \vdash \beta(\phi \wedge \alpha \psi) & \text { for } \alpha<\beta .
\end{array}
$$

In the context of GLP we shall sometimes denote GLP $\vdash \phi \rightarrow \psi$ by $\phi \vdash \psi$. The following is proven in [2]:

Theorem 2.3. GLP is a conservative extension of RC.
This result implies that any reasoning carried out within GLP can, in principle, be carried out within RC, and we shall use this calculus in all formal
reasoning in this paper. As such will write $\lambda \phi$ instead of $\langle\lambda\rangle \phi$, unless the brackets are needed for legibility. We will merely write $\phi \vdash \psi$ to mean " $\phi \vdash \psi$ is a theorem of RC", and for formulas of RC, we will write $\phi \equiv \psi$ if $\phi \vdash \psi$ and $\psi \vdash \phi$. The equivalence class of $\phi$ under $\equiv$ will be denoted $\bar{\phi}$. For a set of formulas $\Phi$, we denote by $\bar{\Phi}$ the set of equivalence classes of its elements.

### 2.4 Worms and the closed fragment

A closed formula in the language of GLP is simply a formula without propositional variables. In other words, closed formulas are generated by just $T$ and the Boolean and modal operators.

The closed fragment of GLP is the class of closed formulas provable in GLP and is denoted by GLP ${ }^{0}$. Within this closed fragment and the corresponding algebra, there is a particular class of privileged inhabitants/terms which are called worms. Worms are nothing more than iterated consistency statements.

Definition 2.4 (Worms, $\mathbb{W}, \mathbb{W}_{\alpha}$ ). By $\mathbb{W}$ we denote the class of worms of GLP which is inductively defined as $T \in \mathbb{W}$ and $A \in \mathbb{W} \Rightarrow\langle\alpha\rangle A \in \mathbb{W}$. Note that every worm belongs to the language of RC .

Similarly, we inductively define for each ordinal $\alpha$ the class of worms $\mathbb{W}_{\alpha}$ where all ordinals are at least $\alpha$ as $T \in \mathbb{W}_{\alpha}$ and $A \in \mathbb{W}_{\alpha} \wedge \beta \geq \alpha \Rightarrow\langle\beta\rangle A \in \mathbb{W}_{\alpha}$.

Both the closed fragment of GLP and the set of worms have been studied in [4] and [9. Worms can be conceived as the backbone of GLP ${ }^{0}$ and obtain their name from the heroic worm-battle, a variant of the Hydra battle (see [5]).

Notation 2.5. We reserve lower-case letters at the beginning of the Greek alphabet $\alpha, \beta, \gamma, \ldots$ for variables ranging over ordinals. Also $\xi$ and $\zeta$ will denote ordinals. Worms will be denoted by upper case latin letters $A, B, C, \ldots$ The Greek lower-case letters $\phi, \psi, \chi, \ldots$ will denote formulas. However, $\varphi$ shall be reserved for the Veblen enumeration function. Likewise, we reserve $\omega$ to denote the first infinite ordinal.

By $|A|$, the length of a worm $A$, we shall mean the number of modalities occurring in $A:|\top|=0$, and $|\langle\xi\rangle A|=|A|+1$. For $A$ a worm and $n$ a natural number we define the $n$-times concatenation of $A$-denoted by $A^{n}$ - as usual: $A^{0}=\top$ and $A^{n+1}=A A^{n}$. We will denote concatenation of worms by juxtaposition, defined recursively so that $\top A=A$ and $(\xi B) A=\xi(B A)$.

## 3 Ordering worms

In this section we define various natural ordering on worms and see how these orderings are related to each other.

### 3.1 Worms and consistency orderings

It is a fact of experience that natural mathematical theories can be linearly ordered in terms of consistency strength. Something similar holds for worms which motivates the next definition.
Definition $3.1\left(<,<_{\xi}\right)$. We define a relation $<_{\xi}$ on $\mathbb{W} \times \mathbb{W}$ by

$$
A<_{\xi} B: \Leftrightarrow B \vdash\langle\xi\rangle A .
$$

Instead of $<_{0}$ we shall simply write $<$.
We shall sometimes refer to the orderings $<_{\xi}$ as the $\xi$-consistency orderings, and these orderings and their corresponding order-types are the main theme of this paper. It is known (4) (9) that the class of worms is linearly ordered by $<_{0}$; that is, if $A, B$ are worms then either $A<_{0} B, A \equiv B$ or $B<_{0} A$.

Recall that $\overline{\mathbb{W}}_{\xi}$ denotes the class of worms in $\mathbb{W}_{\xi}$ modulo GLP-provable equivalence. Clearly, $<_{\xi}$ is well-defined on any of the $\overline{\mathbb{W}}_{\zeta}$ by $\bar{A}<_{\xi} \bar{B} \Leftrightarrow A<_{\xi} B$ whence we shall use the same symbol $<_{\xi}$ for both relations. The following theorem is proven in (4, (9)).
Theorem 3.2. For each ordinal $\xi$, we have that $\left\langle\overline{\mathbb{W}}_{\xi},\left\langle_{\xi}\right\rangle\right.$ is isomorphic to the class of all ordinals.

As a consequence we see that $<_{\xi}$ is irreflexive on $\mathbb{W}$. For, suppose that $A<_{\xi} A$ for some $A \in \mathbb{W}$, then $A \vdash \xi A \vdash 0 A$ contradicting the irreflexivity of $<_{0}$ on $\mathbb{W}_{0}(=\mathbb{W})$.

The existence of a minimal element and the fact that each element has a direct $<_{\xi}$ successor in $\left\langle\bar{W}_{\xi},<_{\xi}\right\rangle$ are reflected in the following easy lemma.

## Lemma 3.3.

1. $\top$ is a $<_{\xi}$-minimal element;
2. For no worms $A, B$ do we have $A<_{\xi} B<_{\xi} \xi A$.

Proof. For the first item, suppose that for some $A$ we have $A<_{\xi} \mathrm{T}$, then $\top \vdash\langle\xi\rangle A \vdash\langle\xi\rangle \top$ whence $T<_{\xi} \top$ which contradicts the irreflexivity of $<_{\xi}$.

For the second, suppose towards a contradiction that there were such a $B$. Then $\xi A \vdash \xi B \vdash \xi \xi A$ whence $\xi A<_{\xi} \xi A$ which once again contradicts the irreflexivity of $<\xi$.

The orderings $<_{\alpha}$ for any ordinal $\alpha>0$ are not linear on $\overline{\mathbb{W}}$. For example, 1 and 101 are $\alpha$ incompatible for $\alpha>0$ : Suppose $101 \vdash \alpha 1$, then $101 \vdash 11 \vdash$ $11 \wedge 01 \vdash 1101 \vdash 0101$, i.e. $101<101$ which contradicts the irreflexivity of $<$. Likewise $1 \vdash \alpha 101 \vdash 0101 \vdash 01$ yields a contradiction. Also $1 \equiv 101$ contradicts reflexivity since $1 \vdash 101 \vdash 001 \vdash 01$. Similarly, it is easy to construct infinite anti-chains -see [20] for examples- for $<_{\xi}$ when $\xi>0$, hence the $<_{\xi}$ orderings do not define a well-quasiorder on $\mathbb{W}$.

The next two lemmata are folklore and follow easily from the axioms of GLP. They shall be used repeatedly often without explicit mention in the remainder of this paper.

## Lemma 3.4.

1. Given formulas $\phi$ and $\psi$, if $\beta<\alpha$, then $(\alpha \phi \wedge \beta \psi) \equiv \alpha(\phi \wedge \beta \psi)$;
2. If $A \in \mathbb{W}_{\alpha+1}$, then $A \wedge\langle\alpha\rangle B \equiv A \alpha B$;

Proof. The left-to-right direction of the first item is an axiom of RC. For the other direction we observe that $\alpha \beta \psi \vdash \beta \psi$ in virtue of axioms $\alpha \beta \psi \vdash \beta \beta \psi$ and $\beta \beta \psi \rightarrow \beta \psi$. The second item follows directly from the first by induction on $|A|$. For more details, we refer to [9].

The next lemma tells us that in various occasions we are allowed to substitute equivalent parts into worms.

## Lemma 3.5.

1. If $A, B \in \mathbb{W}_{\alpha+1}$ and $A \equiv B$, then for any worm $C$ we have $A \alpha C \equiv B \alpha C$;
2. If $A, B, C \in \mathbb{W}$ and $B \equiv C$, then $A B \equiv A C$;
3. More generally, if $A, B, C \in \mathbb{W}$ and $B \vdash C$, then $A B \vdash A C$;
4. For $A, B \in \mathbb{W}$ we have $A B \vdash A$.

Proof. Item 1 follows directly from Lemma 3.42, Item 3 follows from an easy induction on the length of $A$ and Item 2 follows from Item 3. Also, Item 4 follows from Item 3 by taking $C=\top$.

We are not allowed to substitute just any part of a worm. For example, let us see that $1 \equiv 10$ but $11 \not \equiv 101$ : From 3.54 we see that $10 \vdash 1$. Conversely, $1 \vdash 1 \wedge 0 \vdash 10$ by Lemma 3.42 so that $1 \equiv 10$. However, if we assume that $11 \equiv 101$, then $101 \vdash 11 \vdash 11 \wedge 01 \vdash 1101 \vdash 0101$ whence $101 \vdash 0101$. But this is nothing but $101<101$ which contradicts the irreflexivity of $<$.

So, in general we are not allowed to substitute equivalent parts into the leftmost side of a worm. Lemma 3.511 gives an exception and in Corollary 5.2 we will see another exception: when $A \equiv B$, then $A 0 C \equiv B 0 C$.

### 3.2 Decomposing and manipulating worms

In studying worms, and in particular to perform inductive arguments on them it is often useful to decompose worms into smaller worms. In this subsection we will introduce such decompositions, which will appear throughout the text. We use $A \leq_{0} B$ as a shorthand for $A<_{0} B$ or $A \equiv B$. Recall that we use $|A|$ to denote the length of $A$.

Definition 3.6. Let $A$ be a worm. By $h_{\xi}(A)$ we denote the $\xi$-head of $A$. Recursively: $h_{\xi}(\top)=\top, h_{\xi}(\zeta A)=\zeta h_{\xi}(A)$ if $\zeta \geq \xi$ and $h_{\xi}(\zeta A)=\top$ if $\zeta<\xi$.

Likewise, by $r_{\xi}(A)$ we denote the $\xi$-remainder of $A: r_{\xi}(\top)=\top, r_{\xi}(\zeta A)=$ $r_{\xi}(A)$ if $\zeta \geq \xi$ and $r_{\xi}(\zeta A)=\zeta A$ if $\zeta<\xi$.

In words, $h_{\xi}(A)$ corresponds to the largest initial part (reading from left to right) of $A$ such that all symbols in $h_{\xi}(A)$ are at least $\xi$ and $r_{\xi}(A)$ is that part of $A$ that remains when removing its $\xi$-head. We thus have $A=h_{\xi}(A) r_{\xi}(A)$ for all $\xi$ and $A$.

Observe that

$$
\begin{equation*}
h_{\xi}(A) r_{\xi}(A) \equiv h_{\xi}(A) \wedge r_{\xi}(A) \tag{1}
\end{equation*}
$$

as the first symbol of $r_{\xi}(A)$ is less than $\xi$ and $h_{\xi}(A) \in \mathbb{W}_{\xi}$ (see Lemma 3.4).
As a particularly useful instance, we will write simply $h, r$ instead of $h_{1}, r_{1}$.
Definition 3.7. Given a worm $A$, we define $b(A)=B$ if $r(A)=0 B$, and $b(A)=T$ otherwise. We call $b(A)$ the body of $A$.
Lemma 3.8. Given a worm $A \neq \top$, we have that

1. $A \equiv h(A) \wedge 0 b(A)$;
2. $|b(A)|<|A|$;
3. $B<_{0} r(A)$ if and only if $\left(B \leq_{0} b(A)\right.$ and $\left.r(A) \neq \top\right)$.

Proof. We first address the first item. If $b(A) \neq \top$ then we know $A \equiv h(A) \wedge$ $r(A)=h(A) \wedge 0 b(A)$, otherwise, since $A \neq \top$, we always have $A \vdash 0$ so $A \equiv$ $h(A) \wedge 0=h(A) \wedge 0 b(A)$.

It is obvious that $b(A)$ is always shorter than $A$ so that only the last item of the lemma needs to be proven.

For the $\Leftarrow$ direction, suppose that $r(A) \neq \top$. Then, $r(A)=0 b(A)$ whence $B \leq_{0} b(A)<_{0} 0 b(A)$ and $B<_{0} r(A)$.

For the $\Rightarrow$ direction, from $B \ll_{0} r(A)$ we get that $r(A) \neq \top$ (Lemma 3.3|1) whence $B<0 b(A)$. Since $b(A)<B<0 b(A)$ is not possible (Lemma 3.3|2) we get the required $B \leq{ }_{0} b(A)$.

The following lemma tells us that if for some worm $A$ the first element is at most zero, then any worm $A^{\prime}$ equivalent to $A$ must also start with a first element that is at most zero.

Lemma 3.9. If there exists a worm $B$ such that $A \equiv r(B)$, then $h(A)=\top$.
Proof. Assuming that $A \equiv r(B)$, we clearly have $A \vdash h(A) \wedge r(B)$. If $h(A) \neq \top$ then $h(A) \vdash 1$, and thus $A \vdash 1 \wedge r(B) \vdash 1 r(B) \vdash 0 r(B)$. This contradicts the irreflexivity of $<_{0}$.

By Theorem 3.2 we knew that there is a close relation between worms and ordinals and the above lemma exhibits yet another ordinal feature: If we can write an ordinal $\alpha$ as $\alpha^{\prime}+1$, then any other way of writing $\alpha$ must necessarily end with a ' +1 ' too. This analogy will be made more precise after proving Lemma 4.7.

So far we have seen operations on worms that decompose them into different parts. Another very important manipulation on worms is a sort of translation where all modalities in a worm are shifted by a constant amount.

As we shall see in the remainder of this paper, this shift preserves a lot of structure and can even be conceived of as a functor between different spaces.

We will define a shift to the right and one to the left. In order to define the shift to the left we need to recall a very basic fact from ordinal arithmetic (see for example [24]).
Lemma 3.10. If $\zeta<\xi$ are ordinals, there exists a unique $\eta$ such that $\zeta+\eta=\xi$.
We will denote this unique $\eta$ by $-\zeta+\xi$ and it is this operation that is used to define our shift on worms to the left. We are now ready to introduce the shift to the right which is an operation $\alpha \uparrow$ that in general promotes worms to worms with higher consistency strength. As a converse operation we introduce a demoting operator $\alpha \downarrow$ which can be viewed as our shift to the left.

Definition $3.11(\alpha \uparrow$ and $\alpha \downarrow)$. Let $A$ be a worm and $\alpha$ an ordinal. By $\alpha \uparrow A$ we denote the worm that is obtained by simultaneously substituting each $\beta$ that occurs in $A$ by $\alpha+\beta$.

Likewise, if $A \in \mathbb{W}_{\alpha}$ we denote by $\alpha \downarrow A$ the worm that is obtained by replacing simultaneously each $\beta$ in $A$ by $-\alpha+\beta$.

Note that by Lemma 3.10, the operation $\alpha \downarrow$ is well-defined on $\mathbb{W}_{\alpha}$.
Lemma 3.12. For $\alpha, \beta, \gamma$ ordinals and worms $A, B$ we have:

1. $\alpha \uparrow \beta<\alpha \uparrow \gamma \Leftrightarrow \beta<\gamma$,
2. $\alpha \uparrow \beta \geq \beta$,
3. $\alpha \uparrow(\beta \uparrow A)=(\alpha+\beta) \uparrow A$,
4. $\alpha \downarrow(\beta \uparrow A)=(-\alpha+\beta) \uparrow A$, provided $\alpha \leq \beta$,
5. $\alpha \downarrow(\beta \downarrow A)=(\beta+\alpha) \downarrow A$, provided $A \in \mathbb{W}_{\beta+\alpha}$,
6. $\alpha \uparrow((\beta+\alpha) \downarrow A)=\beta \downarrow A$ for $A \in \mathbb{W}_{\beta+\alpha}$,

Proof. The first three items are trivial. It is clearly sufficient to prove items 4 - 6 only for ordinals rather than for worms. All these items have similar elementary proofs. We shall prove Item 4 as an illustration. Thus, let $\alpha \leq \beta$ and fix some ordinal $\gamma$. We see that

$$
\begin{aligned}
\alpha+(\alpha \downarrow \beta) \uparrow \gamma & =\alpha+((\alpha \downarrow \beta)+\gamma) \\
& =(\alpha+(\alpha \downarrow \beta))+\gamma \\
& =\beta+\gamma .
\end{aligned}
$$

Thus, $(\alpha \downarrow \beta) \uparrow \gamma$ is the unique ordinal $\delta$ so that $\alpha+\delta=\beta+\gamma$. In other words, $\alpha \downarrow(\beta \uparrow \gamma)=(-\alpha+\beta) \uparrow \gamma$, provided $\alpha \leq \beta$.

As announced before, the shift operators preserve important structure as is expressed in the following lemma.

Lemma 3.13. For worms $A, B \in \mathbb{W}_{\xi}$ we have

1. $A<_{\xi} B \Leftrightarrow A<B$;
2. $A<_{\xi} B \Leftrightarrow \zeta \uparrow A<_{\zeta+\xi} \zeta \uparrow B$.

Proof. The $\Rightarrow$ direction of 1 is easy. The other direction follows directly from the $\Rightarrow$ direction using irreflexivity and the fact that $<_{\xi}$ linearly orders $\overline{\mathbb{W}}_{\xi}$.

The $\Rightarrow$ direction of 2 is the consequence of a more general observation. One can easily extend the operation $\zeta \uparrow$ to any formula of RC. As the operation $\zeta \uparrow$ is order preserving on the ordinals one can easily prove by induction that any proof in RC remains a proof after applying $\zeta \uparrow$ to every formula appearing in it. Thus, if $\phi \vdash \psi$, then also $\zeta \uparrow \phi \vdash \zeta \uparrow \psi$.

The $\Leftarrow$ direction follows directly from the $\Rightarrow$ direction using irreflexivity and the fact that $<_{\xi}$ is a linear order on $\overline{\mathbb{W}}_{\xi}$.

As a consequence of this lemma we see that we can view each $\alpha \uparrow$ as an isomorphism between structures.

Lemma 3.14. The map $\alpha \uparrow$ is an isomorphism between $(\mathbb{W},<)$ and $\left(\mathbb{W}{ }_{\alpha},<_{\alpha}\right)$.
Moreover, the map $\alpha \Uparrow: \overline{\mathbb{W}} \rightarrow \overline{\mathbb{W}}_{\alpha}$ given by $\alpha \Uparrow \bar{A}=\overline{\alpha \uparrow A}$ is well-defined and also defines an isomorphism.

Proof. The first claim follows from Property 2 of Lemma 3.13. Note that by Property 6 of Lemma 3.12 we see that $\alpha \uparrow(\alpha \downarrow A)=A$ for $A \in \mathbb{W}_{\alpha}$ so that $\alpha \uparrow$ is clearly a bijection.

To check the second point it suffices to observe that if $A \equiv B$ then in view of the first claim, $\alpha \uparrow A \equiv \alpha \uparrow B$, so that the map $\alpha \Uparrow$ is well-defined.

### 3.3 Reducing the $\xi$-consistency orderings

In this subsection we shall see that any question of the form $A<_{\alpha} B$ can be reduced in various ways to simpler questions, for example, to questions of the form $A^{\prime}<B^{\prime}$.

To do so, we first need a reduction lemma (first published in [21]). Recall from Definition 3.6 that $h_{\xi}(A)$ is the largest initial segment of $A$ which lies in $\mathbb{W}_{\xi}$, while $r_{\xi}(A)$ is the rest/remainder of $A$ after removing $h_{\xi}(A)$.
Lemma 3.15. Let $A$ and $B$ be worms and $\xi$ an ordinal. Then, $A>_{\xi} B$ if and only if $h_{\xi}(A)>_{\xi} h_{\xi}(B)$ and $A \vdash r_{\xi}(B)$.

Proof. $(\Rightarrow)$ Assume $A>_{\xi} B$, i.e., $A \vdash\langle\xi\rangle B$. By (11), $B \equiv h_{\xi}(B) \wedge r_{\xi}(B)$ whence $A \vdash \xi B \vdash \xi\left(h_{\xi}(B) \wedge r_{\xi}(B)\right) \vdash \xi h_{\xi}(B) \wedge \xi r_{\xi}(B) \vdash \xi r_{\xi}(B) \vdash r_{\xi}(B)$, since $r_{\xi}(B)$ is either $\top$ or starts with a modality stricly below $\xi$.

It remains to show that $h_{\xi}(A)>_{\xi} h_{\xi}(B)$. Again, we write $A \equiv h_{\xi}(A) \wedge r_{\xi}(A)$. As $h_{\xi}(A), h_{\xi}(B) \in \mathbb{W}_{\xi}$ we know that either (i) $h_{\xi}(A) \equiv h_{\xi}(B)$, (ii) $h_{\xi}(B) \vdash$ $\xi h_{\xi}(A)$ or (iii) $h_{\xi}(A) \vdash \xi h_{\xi}(B)$ holds, so it suffices to discard cases (i) and (ii) under the assumption that $A \vdash \xi B$ whence $A \vdash \xi h_{\xi}(B)$.

Suppose now $h_{\xi}(A) \equiv h_{\xi}(B)$. Then, $A \equiv h_{\xi}(A) \wedge r_{\xi}(A) \vdash \xi h_{\xi}(B) \wedge r_{\xi}(A) \vdash$ $\xi h_{\xi}(A) \wedge r_{\xi}(A) \vdash \xi A$ which contradicts the irreflexivity of $<_{\xi}$.

By a similar argument, the assumption that $h_{\xi}(B) \vdash \xi h_{\xi}(A)$ contradicts the irreflexivity of $<_{\xi}$ and we conclude that $h_{\xi}(A) \vdash \xi h_{\xi}(B)$.
$(\Leftarrow)$ This is the easier direction. Assume that $h_{\xi}(A) \vdash \xi h_{\xi}(B)$ and $A \vdash$ $r_{\xi}(B)$. Then, $A \equiv h_{\xi}(A) \wedge r_{\xi}(A) \vdash \xi h_{\xi}(B) \wedge r_{\xi}(B) \vdash \xi\left(h_{\xi}(B) \wedge r_{\xi}(B)\right) \vdash \xi B$.

In the right-hand side of Lemma 3.15 we see that the first conjunct $h_{\xi}(A)>_{\xi}$ $h_{\xi}(B)$ is only referring to worms in $\mathbb{W}_{\xi}$ and their $<_{\xi}$ relations. The worm $r_{\xi}(B)$ starts with a modality strictly less than $\xi$ and thus the second conjunct $A \vdash r_{\xi}(B)$ of the lemma can be settled by calling recursively to the lemma once more. Thus, Lemma 3.15 recursively reduces the evaluation of statements of the form $A<_{\xi} B$ with $A, B \in \mathbb{W}$ to evaluation of statements of the form $A^{\prime}<_{\xi^{\prime}} B^{\prime}$ with $A^{\prime}, B^{\prime} \in \mathbb{W}_{\xi^{\prime}}$.

Moreover, we know (by Lemma 3.13) that

$$
\begin{aligned}
h_{\xi}(A)>_{\xi} h_{\xi}(B) & \Leftrightarrow h_{\xi}(A)>_{0} h_{\xi}(B) \\
& \Leftrightarrow \xi \downarrow h_{\xi}(A)>_{0} \xi \downarrow h_{\xi}(B) .
\end{aligned}
$$

Thus, Lemma 3.15 tells us that by recursion on $\xi$, any question about $B<_{\xi} A$ can be reduced to question about $B^{\prime}<_{0} A^{\prime}$.

Thus, we can reduce questions about any $<_{\xi}$ ordering to questions about the $<_{0}$ ordering. We shall now see that we reduce questions about the $<_{0}$ ordering even further in that we may restrict questions of the form $A<_{0} B$ to the case where one of $A$ or $B$ is either $\top$ or of the form $0 C$.

Before we prove this further reduction, we first need an elementary lemma that relates the notions $<, \leq$, and $\vdash$.

## Lemma 3.16.

1. $(A \leq B<C) \vee(A<B \leq C) \Rightarrow A<C$;
2. $C \leq A \Leftrightarrow C<0 A$;
3. If $B \vdash A$ then $A \leq{ }_{0} B$;
4. $A \vdash r(B) \Leftrightarrow A \geq r(B)$.

Proof. Easy and left to the reader.
Note that we cannot reverse the implication of the third item since, for example, it is easy to check that $1<01$ but $01 \nvdash 1$. We shall use our previous lemma without explicit mention in the remainder of this paper.

Lemma 3.17.

1. $A<{ }_{0} B$ if and only if one of the following occurs:
(a) $A<{ }_{0} r(B)$ or
(b) $r(A)<_{0} B$ and $h(A)<_{0} h(B)$.
2. $A \leq{ }_{0} B$ if and only if
(a) $A \leq_{0} r(B)$ or
(b) $r(A) \leq_{0} B$ and $h(A) \leq_{0} h(B)$.

Proof. We prove the first item and first focus on the $\Leftarrow$ direction omitting various 0 -subscripts. In case $A<r(B)$ we get $B \vdash h(B) \wedge r(B) \vdash r(B) \vdash 0 A$ so $A<B$.

So, now suppose that $r(A)<B$ and $h(A)<h(B)$. Clearly, $h(A)<h(B) \Leftrightarrow$ $h(A)<{ }_{1} h(B)$, and since $0 r(B) \vdash r(B)$ we get
$B \vdash h(B) \wedge B \vdash 1 h(A) \wedge B \vdash 1 h(A) \wedge 0 r(A) \vdash 1 h(A) \wedge r(A) \vdash 1 h(A) r(A) \vdash 0 A$
whence $A<B$.
For the $\Rightarrow$ direction, assume that $A<_{0} B$; let us show that if 1a fails, then 1 b holds. Clearly we have $r(A)<_{0} B$ since $B \vdash 0 A \vdash 0(h(A) \wedge r(A)) \vdash 0 r(A)$. We wish to see that $h(A)<h(B)$. Since by assumption 1a fails, we have that $r(B) \leq_{0} A$ whence $A \vdash r(B)$. Now suppose for a contradiction that $h(B) \leq h(A)$. In case $h(B) \equiv h(A)$ we get

$$
A \vdash h(A) \wedge A \vdash h(B) \wedge A \vdash h(B) \wedge r(B) \vdash B \vdash 0 A,
$$

and in case $h(B)<h(A)$ then also $h(B)<_{1} h(A)$ and we get

$$
A \vdash h(A) \wedge A \vdash 1 h(B) \wedge A \vdash 1 h(B) \wedge r(B) \vdash 1 B \vdash 10 A \vdash 00 A \vdash 0 A,
$$

contradicting the irreflexivity of $<$ so that $h(A)<h(B)$ as was to be shown.
The proof of Lemma 3.172 is similar.
Note that when asking the question $A<_{0} B$ we may always assume that one of $A$ or $B$ contains a zero, since $A<_{0} B \Leftrightarrow \alpha \downarrow A<_{0} \alpha \downarrow B$ where $\alpha$ is the smallest ordinal appearing in $A B$. Thus, indeed, by induction on $|A|+|B|$ we see that this lemma provides a reduction of questions about $<_{0}$ to questions about $<_{0}$ where one of the arguments is either $T$ or starts with a 0.

## 4 Worms and ordinals

In the previous section we introduced various orderings on the worms. In the current section we shall study the corresponding order-types.

### 4.1 Well-founded orders and order types

An important corollary to our reduction lemma, Lemma 3.15, is that the $<_{\alpha}$ orders are well-founded.

Corollary 4.1. The relation $<_{\alpha}$ on $\mathbb{W} \times \mathbb{W}$ is well-founded.

Proof. Any hypothetical infinite descending $<_{\alpha}$-chain $A_{0}>_{\alpha} A_{1}>_{\alpha} A_{2}>_{\alpha} \ldots$ in $\mathbb{W}$ yields a corresponding chain $h_{\alpha}\left(A_{0}\right)>_{\alpha} h_{\alpha}\left(A_{1}\right)>_{\alpha} h_{\alpha}\left(A_{2}\right)>_{\alpha} \ldots$ in $\mathbb{W}{ }_{\alpha}$ by Lemma 3.15. This contradicts the fact that $<_{\alpha}$ defines a well-order on $\bar{W}_{\alpha}$ (Theorem 3.2).

In virtue of this well-foundedness we can assign a $\xi$-order to each worm by bar-recursion. In this subsection we shall revisit classical theory on how this assignment can be made in the general setting of well-founded orders. In the next subsection we shall apply the general results to the orders induced by the $<_{\xi}$ orderings.

Definition 4.2. If $X$ is either a set or a class and $\prec$ is a well-founded relation on $X$, we define $o_{\prec}: X \rightarrow$ On by

$$
o_{\prec}(x)=\sup _{y \prec x}\left(o_{\prec}(y)+1\right),
$$

where $\sup \varnothing=0$.
Note that $o_{\prec}$ is well-defined due to the assumption that $\prec$ is well-founded. In the case that $\prec$ is in addition a linear order, $o_{\prec}$ is in fact an isomorphism onto an initial segment of the ordinals. To make this precise, we first need an important definition. We shall identify the ordinal $\alpha$ with the set $\{\beta \mid \beta<\alpha\}$ of ordinals $\beta$ less than alpha.

Definition 4.3. If $\langle X, \prec\rangle$ is a well-order, define

$$
\operatorname{ot}_{\prec}(X)=\sup _{x \in X}\left(o_{\prec}(x)+1\right),
$$

where possibly ot $_{\prec}(X)=\mathrm{On}$ if $X$ is a proper class.
To formulate our result we shall use the following notation

$$
X^{\prec x}=\{y \in X \mid y \prec x\} .
$$

Lemma 4.4. For $\langle X, \prec\rangle$ a well-order, we have that $o_{\prec}:\langle X, \prec\rangle \rightarrow\left\langle\operatorname{ot}_{\prec}(X),<\right\rangle$ is an isomorphism. Moreover, if $x \in X$, then $o_{\prec}(x)=$ ot $_{\prec}\left(X^{\prec x}\right)$.

The proof is easy and details can be found, e.g., in [22, 24, 19 .
In general, while $o_{\prec}$ is defined for any well-founded relation, it has much nicer behavior when $\prec$ is linear. For example, both $o_{\prec}$ and $o_{\prec}^{-1}$ are continuous as expressed by the following easy lemma whose proof we omit.
Lemma 4.5. If $\langle X, \prec\rangle$ is a well-ordered set and $S \subset X$ is bounded (in $X$ ), then $S$ has a supremum and $o_{\prec}(\sup S)=\sup o_{\prec}(S)$.

Likewise, $o_{\prec}^{-1}(\sup \Gamma)=\sup o_{\prec}^{-1}(\Gamma)$ for any set of ordinals $\Gamma$.

### 4.2 Worms and their order-types

We shall now see how the observations of the previous subsection apply to $\overline{\mathbb{W}}$. If $X=\mathbb{W}$ and $\prec=<_{\xi}$ for some ordinal $\xi$, we write $o_{\xi}$ instead of $o_{<_{\xi}}$, while if $X=\overline{\mathbb{W}}$ we will write $\mathrm{o}_{\xi}$.

Instead of $o_{0}$ we shall just write $o$. When $S$ is a set or class we shall denote by $o_{\alpha}(S)$ the image of $S$ under $o_{\alpha}$. It is easy to see that, for any $A \in \mathbb{W}$, $\mathrm{o}_{\alpha}(\bar{A})=o_{\alpha}(A)$ and moreover $\mathrm{o}_{\alpha}\left(\overline{\mathbb{W}}_{\beta}\right)=o_{\alpha}\left(\mathbb{W}_{\beta}\right)$ for any ordinals $\alpha, \beta$. We shall often refer to the $o_{\xi}$ as $\xi$-consistency order-types.

Lemma 4.6. Given $A, B \in \mathbb{W}, A<B$ if and only if $o(A)<o(B)$, and the map o: $\mathbb{W} \rightarrow$ On is surjective. Moreover, the map o : $\left\langle\overline{\mathbb{W}},<_{0}\right\rangle \rightarrow\langle\mathrm{On},<\rangle$ is an isomorphism.

Proof. Since $\left\langle\overline{\mathbb{W}},<_{0}\right\rangle$ is a well-order, we have by Lemma 4.4 that o: $\overline{\mathbb{W}} \rightarrow$ ot $_{<_{0}}(\overline{\mathbb{W}})$ is an isomorphism, and moreover since for $A, B \in \mathbb{W}$ we have that $A<_{0} B$ if and only if $\bar{A}<_{0} \bar{B}$, it also follows that $A<_{0} B$ if and only if $o(A)<o(B)$.It remains to show that $\mathrm{o}(\overline{\mathbb{W}})$ is unbounded in On. But an easy induction shows that $o(\langle\alpha\rangle) \geq \alpha$ for all $\alpha$.

Some elementary properties of the mappings $o_{\xi}$ are readily proven.

## Lemma 4.7.

1. $o_{\xi}(T)=0$;
2. $o_{\xi}(\xi A)=o_{\xi}(A)+1$;
3. $o_{\xi}\left(\xi^{n}\right)=n$.

Proof. The first item is clear (by Lemma 3.311) since $\left\{A \mid A<_{\xi} \top\right\}=\varnothing$ and $\sup \varnothing=0$.

For the second item, since $A<_{\xi} \xi A$ clearly $o_{\xi}(A)+1 \leq o_{\xi}(\xi A)$. By Lemma 3.322 we see that $\left\{B \mid B<_{\xi} \xi A\right\}=\left\{B \mid B<_{\xi} A\right\} \cup\{A\}$ so that actually $o_{\xi}(A)+1=o_{\xi}(\xi A)$.

Finally, Item 3 follows from the other two by induction on $n$.
Let us conclude this subsection by identifying the 'limit worms' and the 'successor worms'. As usual, by Succ and Lim we denote the class of successor respectively limit ordinals.

Lemma 4.8. For any worm $A$ we have

1. $o(A)=0 \Leftrightarrow A=\top$;
2. $o(A) \in \operatorname{Succ} \Leftrightarrow h(A)=\top \neq A \Leftrightarrow A=0 A^{\prime}$ for some worm $A^{\prime}$;
3. $o(A) \in \operatorname{Lim} \Leftrightarrow h(A) \neq \top \Leftrightarrow A \neq \top \& A=\sup _{B<A} B$.

Proof. The $\Leftarrow$ direction of Item (1) is just Lemma 4.71]. For the $\Rightarrow$ direction we see that if $A \neq \top$, then $A \vdash 0 \top$ so that $o(A)>0$.

The $o(A) \in \operatorname{Succ} \Leftarrow h(A)=\top \neq A$ direction of (2) is already given in Lemma 4.72 The other direction follows from (1) in case $A=T$ so we assume that $h(A) \neq \mathrm{T}$. If $o(A) \in$ Succ, then by definition of $o$ we have $o(A)=o(B)+1$ for some $B<A$. Thus, it suffices to show that

$$
\begin{equation*}
\text { if } h(A) \neq \top \text {, and } B<A \text { then } 0 B<A \text {. } \tag{2}
\end{equation*}
$$

Indeed, $h(A) \vdash 1$, hence $A \vdash 1 \wedge 0 B \vdash 10 B \vdash 00 B$, and $0 B<A$, as claimed. The statement that $(h(A)=\top \neq A) \Leftrightarrow\left(A=0 A^{\prime}\right.$ for some worm $\left.A^{\prime}\right)$ is a mere syntactical triviality.

The $o(A) \in \operatorname{Lim} \Leftrightarrow h(A) \neq \top$ equivalence of Item (3) clearly follows from the other two items. Thus it remains to show that $o(A) \in \operatorname{Lim}$ if and only if $A \neq \top$ and $A=\sup _{B<A} B$. First assume that $o(A) \in \operatorname{Lim}$; then clearly $A \neq \top \neq h(A)$ and, by (2) if $B<A$, then $o(A)>o(B)+1=o(0 B)$ so by Lemma 4.6 $B<0 B<A$; it follows that $A=\sup _{B<A} B$, since $\{B \in \mathbb{W}: B<A\}$ can have no smaller upper bound than $A$.

Similarly, if $A \neq \top$ and $A=\sup _{B<A} B$, then given $C<A$ we have that $0 C<A$ (as we cannot have that $0 C \equiv A$ ), and thus $o(C)+1<o(A)$, which means that $o(A) \in \operatorname{Lim}$.

Thus we will say that $A$ is a successor worm whenever $A=0 A^{\prime}$ for some worm $A^{\prime}$, or a limit worm $A$ whenever $h(A) \neq \top$.

### 4.3 Reducing the $\xi$-order-types

In our reduction lemma, Lemma 3.15, we saw how questions about the $<_{\xi}$ orderings could be reduced. This reduction could be viewed as a reduction to the $<_{\zeta}$ orderings restricted to the respective $\mathbb{W}_{\zeta}$ 's as well as a reduction to the $<_{0}$ ordering. In this subsection we shall see that we have similar reductions for the order-types.

Let us temporarily introduce new orderings $\check{o}_{\xi}$ which shall turn out to be just the restriction of $o_{\xi}$ to $\mathbb{W}_{\xi}$.
Definition 4.9. For $A \in \mathbb{W}_{\xi}$ we define $\check{o}_{\xi}(A)=o_{<_{\xi} \mid \mathbb{W}_{\xi}}(A)$.
Since we know that $<_{\xi}$ linearly orders $\overline{\mathbb{W}}_{\xi}$ we have an alternative characterization of $\check{o}_{\xi}(A)$.
Lemma 4.10. Given $A \in \mathbb{W}_{\xi}$, ŏ $(A)=\operatorname{ot}_{<_{\xi}}\left\{\bar{B} \in \overline{\mathbb{W}}_{\xi} \mid B<_{\xi} A\right\}$.
Proof. By Theorem 3.2 we know that $\left\langle\overline{\mathbb{W}}_{\xi},<_{\xi}\right\rangle$ is well-ordered, thus by Lemma 4.4) $\check{o}_{\xi}: \overline{\mathbb{W}}_{\xi} \rightarrow$ ot $_{<_{\xi}}\left(\overline{\mathbb{W}}_{\xi}\right)$ is an isomorphism. It remains to check that ot $\operatorname{cog}_{\xi}\left(\overline{\mathbb{W}}_{\xi}\right)=$ On; but this readily follows by observing (by a straightforward induction) that $\check{o}_{\xi}(\langle\xi+\alpha\rangle) \geq \alpha$.

The next lemma tells us how the computation of $o_{\xi}$ can be reduced to computations of $\check{o}_{\xi}$.

Lemma 4.11. For any worm $A$ and ordinal $\xi$ we have $o_{\xi}(A)=o_{\xi} h_{\xi}(A)=$ $\check{o}_{\xi} h_{\xi}(A)$.

Proof. Recall that for a partial order $\langle X, \prec\rangle$ and $x \in X$, we defined $X^{\prec x}=$ $\{y \in X \mid y \prec x\}$. We will write $\mathbb{W}_{\xi}^{<\xi A}$ instead of $\left(\mathbb{W}_{\xi}\right)^{<\xi} A$.

We first see that

$$
\begin{equation*}
\mathbb{W}_{\xi}^{<_{\xi} h_{\xi} A}=h_{\xi} \mathbb{W}<_{\xi} A:=\left\{h_{\xi}(B) \mid B<_{\xi} A\right\} \tag{3}
\end{equation*}
$$

For the $\subseteq$ inclusion, we fix some $C \in \mathbb{W}_{\xi}$ and observe that $h_{\xi}(C)=C$. Now, if $C<_{\xi} h_{\xi}(A)$ then clearly $A \vdash h_{\xi}(A) \wedge r_{\xi}(A) \vdash h_{\xi}(A) \vdash\langle\xi\rangle C$ whence also $C<_{\xi} A$.

The other direction follows directly from Lemma 3.15 since $B<_{\xi} A$ implies $h_{\xi}(B)<_{\xi} h_{\xi} A$ and clearly $h_{\xi}(B) \in \mathbb{W}_{\xi}$.

Now that we have this equality we proceed by induction on $A$ and obtain

$$
\begin{aligned}
& o_{\xi}(A):=\sup _{B<_{\xi} A} \quad\left(o_{\xi}(B)+1\right) \\
& \mathrm{IH}_{\left(B<_{\xi} A\right)} \quad=\sup _{B<_{\xi} A} \quad\left(o_{\xi} h_{\xi}(B)+1\right) \\
& \text { by (3) } \quad=\sup _{C \in \mathbb{W}_{\xi}^{<} \xi^{h} \xi(A)}\left(o_{\xi}(C)+1\right) \\
& \mathrm{IH}_{\left(C<_{\xi} h_{\xi}(A) \leq_{\xi} A\right)} \quad=\sup _{C \in \mathbb{W}_{\xi}^{<\xi^{h}(A)}}\left(\check{o}_{\xi}\left(h_{\xi}(C)\right)+1\right) \\
& =\sup _{C \in \mathbb{W}_{\xi}^{<} \xi^{h_{\xi}(A)}}\left(\check{o}_{\xi}(C)+1\right) \\
& =\quad \check{o}_{\xi} h_{\xi}(A) \text {. }
\end{aligned}
$$

Likewise,

| $o_{\xi}(A)$ | $=\sup _{C \in \mathbb{W}_{\xi}^{<\xi^{h}(A)}}$ | $\left(o_{\xi}(C)+1\right)$ |
| :--- | :--- | :--- |
|  | $=\sup _{C \in \mathbb{W}_{\xi}{ }_{\xi} \xi^{h} h_{\xi}(A)}$ | $\left(o_{\xi}(C)+1\right)$ |
| $\left(h_{\xi} A=h_{\xi} h_{\xi} A\right)$ |  | $\sup _{C<_{\xi} h_{\xi}(A)}$ |
| by (3) | $\left(o_{\xi}\left(h_{\xi}(C)\right)+1\right)$ |  |
| $\mathrm{IH}_{\left(C<_{\xi} h_{\xi}(A) \leq_{\xi} A\right)}$ | $=\sup _{C<_{\xi} h_{\xi}(A)}$ | $\left(o_{\xi}(C)+1\right)$ |
|  | $=$ | $o_{\xi} h_{\xi}(A)$. |

Indeed, this lemma yields a reduction of questions about the orders $o_{\xi}$ defined on $\mathbb{W}$ to questions about the orders $\check{o}_{\xi}$ which are defined on $\mathbb{W}_{\xi}$. In particular, we see that $\check{o}_{\xi}$ is just the restriction of $o_{\xi}$ to $\mathbb{W}_{\xi}$.
Corollary 4.12. For $A \in \mathbb{W}_{\xi}$ we have that $o_{\xi}(A)=\check{o}_{\xi}(A)$.
Proof. Immediate from Lemma 4.11 since $h_{\xi}(A)=A$ for $A \in \mathbb{W}_{\xi}$.
We now also obtain an alternative characterization of $o_{\xi}(A)$ : If $\langle X, \prec\rangle$ is a partially ordered set (or class), a $\prec$-chain in $X$ is any subset $\mathcal{C}$ of $X$ which is linearly ordered by $\prec$. We reserve $\mathcal{C}$ to denote chains. Given $x \in X$ we will write $\mathcal{C} \prec x$ if $x$ is a strict upper bound for $\mathcal{C}$, we say that $\mathcal{C}$ is a chain below $x$.

It is straightforward to check that if $\mathcal{C} \prec x$ then ot $_{\prec}(\mathcal{C}) \leq o_{\prec}(x)$. However, it may be that $\operatorname{ot}_{\prec}(\mathcal{C})$ is always much smaller than $o_{\prec}(x)$ as is expressed in the following lemma which is folklore.

Lemma 4.13. Given an ordinal $\gamma$, there exists a partially ordered set $\langle X, \prec\rangle$ and $x \in X$ so that $o_{\prec}(x)=\gamma$ but every $\prec$-chain below $x$ is finite.

Thus it may not be the case that $o_{\prec}(x)$ is 'attained'. To be precise, given a partially ordered set $\langle X, \prec\rangle$ and $x \in X$, say that $o_{\prec}(x)$ is attained if there is a chain $\mathcal{C} \prec x$ such that ot $_{\prec}(\mathcal{C})=o_{\prec}(x)$.

Theorem 4.14. If $A \in \mathbb{W}$ and $\xi$ is any ordinal, then

$$
o_{\xi}(A)=\sup _{\mathcal{C}<\xi A} \operatorname{ot}_{\xi}(\mathcal{C})
$$

Moreover, $o_{\xi}(A)$ is attained in $\left\langle\mathbb{W}_{\xi},<\xi\right\rangle$.
Proof. To prove that $o_{\xi}(A)=\sup _{\mathcal{C}<_{\xi} A}$ ot $\xi_{\xi}(\mathcal{C})$, clearly it suffices that $o_{\xi}(A)$ is attained in $\left\langle\mathbb{W}_{\xi},<_{\xi}\right\rangle$.

Let $A \in \mathbb{W}$ and define $\mathcal{C}=\mathbb{W}_{\xi}^{<\xi h_{\xi}(A)}=\left\{B \in \mathbb{W}_{\xi}: B<_{\xi} h_{\xi}(A)\right\}$. By Theorem 3.2, $\mathbb{W}_{\xi}$ is well-ordered and hence $\mathcal{C}$ is a chain; clearly also $\mathcal{C}<_{\xi} A$. Meanwhile, by Lemma 4.11, we have that $o_{\xi}(A)=\check{o}_{\xi} h_{\xi}(A)$; however, by Lemma 4.4 $\check{o}_{\xi} h_{\xi}(A)=$ ot $_{<_{\xi}} \mathcal{C}$. We conclude that $\mathcal{C}$ is a $<_{\xi}$-chain below $A$ with $o_{\xi}(A)=$ ot $_{<\xi} \mathcal{C}$, and thus $o_{\xi}(A)$ is attained.

Just like the reduction lemma for the $<_{\xi}$ orderings yielded a reduction to $<_{0}$, by an additional lemma we shall see that Lemma 4.11 also provides a reduction of $o_{\xi}$ to $o_{0}$.

Lemma 4.15. For $A \in \mathbb{W}_{\xi}$ we have $o_{\xi}(A)=o(\xi \downarrow A)$.
Proof. By Lemma 3.12 it suffices to prove that for any worm $A \in \mathbb{W}$ we have $o_{\xi}(\xi \uparrow A)=o(A)$. Now, since $\xi \uparrow B<_{\xi} \xi \uparrow A \Leftrightarrow B<A$ (this is Lemma 3.13) and since each $B \in \mathbb{W}$ is equal to $\xi \downarrow(\xi \uparrow B)$, we have that for each worm $A$

$$
\begin{equation*}
\{B \in \mathbb{W} \mid B<A\}=\left\{\xi \downarrow C \mid C \in \mathbb{W}_{\xi} \wedge C<_{\xi} \xi \uparrow A\right\} \tag{4}
\end{equation*}
$$

We will now show by induction on $<_{0}$ that for any worm $A$ we have $o(A)=$ $\check{o}_{\xi}(\xi \uparrow A)$ :

IH

$$
\begin{array}{llll} 
& o(A) & =\sup _{B<A} & \\
\text { IH } & =\sup _{B<A} & & (B)+1) \\
\text { by (4) } & & \left.\check{o}_{\xi}(\xi \uparrow B)+1\right) \\
& & \sup _{C \in \mathbb{W}_{\xi}^{<}{ }_{\xi} \xi \uparrow A} & \\
& \left(\check{o}_{\xi}(\xi \uparrow(\xi \downarrow C))+1\right) \\
& & \sup _{C \in \mathbb{W}_{\xi}{ }_{\xi}{ }^{\xi} \xi \uparrow A} & \\
& & \left(\check{o}_{\xi}(C)+1\right) \\
& & & \check{o}_{\xi}(\xi \uparrow A) .
\end{array}
$$

We conclude our proof using Lemma 4.11 to see that

$$
o(A)=\check{o}_{\xi}(\xi \uparrow A)=o_{\xi}\left(h_{\xi}(\xi \uparrow A)\right)=o_{\xi}(\xi \uparrow A)
$$

The reduction of $o_{\xi}$ to $o$ follows by combining Lemma 4.15 and Lemma4.11.

Corollary 4.16. For any worm $A$ and ordinal $\xi$ we have $o_{\xi}(A)=o\left(\xi \downarrow h_{\xi}(A)\right)$.
Our temporary definition of $\check{o}$ will not be used further on in the paper. In previous papers on well-orders in the Japaridze algebra one used the definition $\check{o}_{\xi}$ and denoted that by $o_{\xi}$. By Corollary 4.12 we know that the definition of $o_{\xi}$ in this paper coincides with the old definition when restricted to $\mathbb{W}_{\xi}$. Thus, our notation for the new definition causes no rupture with the tradition yet merely generalizes the existing theory.

## 5 A calculus for the consistency order-types

In this section we show how to compute the order-types $o(A)$. Actually, we shall provide a calculus that reduces the computation of $o(A)$ to the computation of what we call worm enumerators. The calculus will consist of three cases: the empty worm, worms containing a zero and non-empty worms that do not contain a zero. Recall that the definitions of $h(A), b(A), r(A)$ may be found in Section 3.2

### 5.1 Worms containing a zero

For ordinals we do not have that $\xi<\zeta \Leftrightarrow \xi+\alpha<\zeta+\alpha$. However for addition on the right we do have that $\xi<\zeta \Leftrightarrow \alpha+\xi<\alpha+\zeta$. For worms we have something similar although now the left-most side of the worm is determinant:

Lemma 5.1. Given worms $A, B, C$, we have that $A 0 C<_{0} B 0 C$ if and only if $A<_{0} B$.

In order to simultaenously prove both implications, we will instead prove the equivalent claim that if $A, B, C, D$ are worms such that $h(C)=h(D)=\top$ (i.e., $C, D$ begin with a zero or are $T$ ), then $A C<_{0} B C$ implies that $A D<_{0} B D$. We work by induction on $|A|+|B|$. Observe that if $A_{1}, A_{0}$ are worms with $h\left(A_{0}\right)=\top$, then $h\left(A_{1} A_{0}\right)=h\left(A_{1}\right)$ while $r\left(A_{1} A_{0}\right)=r\left(A_{1}\right) A_{0}$.

If $A C<_{0} B C$, then by Lemma 3.17 we have that either (i) $A C<_{0} r(B C)$, or (ii) $r(A C)<_{0} B C$ and $h(A C)<_{0} h(B C)$. If (i) holds, since $r(B C)=r(B) C$, we have that $C \leq_{0} A C<_{0} r(B) C$, hence by irreflexivity $r(B) \neq \top$ (or else $C<_{0} C$ ) and thus $r(B)=0 b(B)$ and $b(B C)=b(B) C$. Thus by Lemma 3.8. $A C \leq_{0} b(B C)=b(B) C$. By linearity this is equivalent to $b(B) C \nless_{0} A C$, so that by our induction hypothesis $b(B) D \nless_{0} A D$, i.e. $A D \leq_{0} b(B) D$. It follows that $r(B) D=0 b(B) D>_{0} b(B) D \geq_{0} A D$, and thus $B D=h(B) r(B) D \vdash$ $h(B) 0 A D \vdash 0 A D$, so that $A D<_{0} B D$.

If (ii) holds, we first claim that $B D \vdash r(A) D$. If $r(A)=\top$ this is obvious, otherwise we note that $b(A) C<_{0} A C<_{0} B C$, so by the induction hypothesis $b(A) D<_{0} B D$ and $B D \vdash 0 b(A) D=r(A) D$. Thus, since $h(A)=h(A C)<_{0}$ $h(B C)=h(B)$, we also have $h(A)<_{1} h(B)$ and get $B D \equiv h(B) \wedge B D \vdash$ $1 h(A) \wedge r(A) D \equiv 1 h(A) r(A) D \vdash 0 A D$, and $A D<_{0} B D$.

We have seen at the end of Lemma 3.5 that on the left side of a worm, one is not allowed to replace a part by any equivalent part. The next corollary tells us that if we have a zero, then that allows us such a substitution and as such the zero functions as a sort of buffer.

Corollary 5.2. For worms $A, B$ and $C$ we have that $A \equiv B$ if and only if $A 0 C \equiv B 0 C$.

Proof. Immediate from Lemma 5.1.
The following is an analogue of Lemma 3.10 it says that, for worms, we have a form of subtraction. In this case, however, it becomes "right subtraction".

Lemma 5.3. $A<{ }_{0} B$ if and only if there exists $C$ such that $B \equiv C 0 A$.
Proof. One direction is trivial: if $B \equiv C 0 A$, then clearly $B \vdash 0 A$.
So assume that $A<_{0} B$. We shall prove by induction on $|B|$ that if $A<_{0} B$, then we can find a worm $C$ so that $B \equiv C 0 A$.

We consider two cases depending on $A<_{0} r(B)$ or $r(B) \leq_{0} A$.
In case $A<_{0} r(B)$ we must have $r(B) \neq \top$, so $r(B)=0 b(B)$, whence $A \leq{ }_{0} b(B)$. If indeed $A \equiv b(B)$ then we have $B=h(B) 0 b(B) \equiv h(B) 0 A$. Otherwise, $A \ll_{0} b(B)$ so by induction hypothesis $b(B) \equiv C^{\prime} 0 A$ for some worm $C^{\prime}$. We set $C=h(B) 0 C^{\prime}$ and readily see that $B \equiv C 0 A$.

In case $A \geq_{0} r(B)$ we claim that $B \equiv h(B) 0 A$. Indeed since $A<_{0} B$, $B \vdash h(B) \wedge 0 A \equiv h(B) 0 A$, while since $A \geq_{0} r(B)$, we also have $h(B) 0 A \vdash$ $h(B) \wedge 0 A \vdash h(B) \wedge 0 r(B) \vdash h(B) \wedge r(B) \equiv B$.

The above lemmas suggest that concatenations of the form $A 0 B$ behave much like addition; the following result makes this precise.

Lemma 5.4. Given worms $A, B, o(A 0 B)=o(B)+1+o(A)$.
Proof. By induction on $<_{0}$. First let us show that $o(A 0 B) \leq o(B)+1+o(A)$. Suppose that $C<_{0} A 0 B$. If $C \leq_{0} B$, then we already have $o(C) \leq o(B)<$ $o(B)+1+o(A)$. If $C>_{0} B$, then by Lemma 5.3 we have that $C \equiv D 0 B$ for some $D$, and thus by the induction hypothesis we have $o(C)=o(B)+1+o(D)$. But by Lemma5.1, $D<_{0} A$ so $o(D)<_{0} o(A)$ and thus $o(B)+1+o(D)<o(B)+1+o(A)$.

We now will show that $o(B)+1+o(A) \leq o(A 0 B)$. Note that if $A=\top$, then $o(A 0 B)=o(0 B)=o(B)+1$ by Lemma 4.7. Thus, we assume that $A \not \equiv \top$ so that we can choose $\xi<o(A)$. Then, since $o$ is an isomorphism between $\left\langle\overline{\mathbb{W}},<_{0}\right\rangle$ and the ordinals (Lemma 4.6), there is a worm $C<_{0} A$ with $o(C)=\xi$. By Lemma 5.1 we have that $C 0 B<_{0} A 0 B$, and by induction $o(C 0 B)=o(B)+$ $1+o(C)=o(B)+1+\xi$. Since $\xi<o(A)$ was arbitrary, we conclude that $o(A 0 B) \geq o(B)+1+o(A)$.

In this subsection we have dealt with worms that do contain a zero and could recursively compute their order-types. We shall reduce worms that do not contain a zero to worms that do contain a zero via the $\xi \uparrow$ and $\xi \downarrow$ mappings.

### 5.2 A calculus using the worm enumerators $\sigma^{\alpha}$

A key role in the larger calculus is reserved for the functions $\sigma^{\alpha}$ that enumerate the <-orders of worms in $\mathbb{W}_{\alpha}$ in increasing order. We shall prove sufficiently many structural properties of these functions $\sigma^{\alpha}$ so that we end up with a recursive calculus to compute them.

Moreover, it shall turn out that the functions $\sigma^{\alpha}$ can be viewed as transfinite iterations of a certain ordinal exponentiation that we shall call hyperexponential functions and which we shall discuss in Section 7.1

Definition 5.5 (Worm enumerators $\sigma^{\alpha}$ ). We define $\sigma^{\alpha}$ to be the function that enumerates o $\left(\mathbb{W}_{\alpha}\right)$ in increasing order.

We shall first see how a calculus for $o$ can be reduced to a calculus for these functions $\sigma^{\alpha}$. The following nice lemma characterizes $\sigma^{\alpha}$ as a conjugate of the map $\alpha \Uparrow$ on worms.

Lemma 5.6. $o\left(\mathbb{W}_{\alpha}\right)=\mathrm{o}\left(\overline{\mathbb{W}}_{\alpha}\right)$ is enumerated in increasing order by $\mathrm{o} \circ \alpha \Uparrow \circ \mathrm{o}^{-1}$, that is,

$$
\sigma^{\alpha}=\mathrm{o} \circ \alpha \Uparrow \circ \mathrm{o}^{-1}
$$

Proof. In the proof we shall explicitly write $<_{0}$ for the ordering on worms and $<$ for the ordering on ordinals. Lemma 4.6 told us that o : $\left\langle\overline{\mathbb{W}},<_{0}\right\rangle \cong\langle\mathrm{On},<\rangle$. Thus by Lemma 3.1311 we see that for $A, B \in \mathbb{W}_{\alpha}$

$$
\begin{equation*}
A<_{\alpha} B \Leftrightarrow A<_{0} B \Leftrightarrow o(A)<o(B) . \tag{5}
\end{equation*}
$$

If we combine this with the fact that $o\left(\mathbb{W}_{\alpha}\right)$ is an unbounded class of ordinals, we see that an order-preserving enumeration of $o\left(\mathbb{W}_{\alpha}\right)$ is nothing but the unique isomorphism between $\langle\mathrm{On},<\rangle$ and $\left\langle o\left(\mathbb{W}_{\alpha}\right),<_{\alpha}\right\rangle$.

We can reformulate (5) as o: $\left\langle\overline{\mathbb{W}}_{\alpha},<_{\alpha}\right\rangle \cong\left\langle o\left(\mathbb{W}_{\alpha}\right),<\right\rangle$. We also have by Lemma 3.14 that $\alpha \Uparrow:\left\langle\overline{\mathbb{W}},<_{0}\right\rangle \cong\left\langle\overline{\mathbb{W}}_{\alpha},<_{\alpha}\right\rangle$. Once more using the fact that $\mathrm{o}^{-1}:\langle\mathrm{On},<\rangle \cong\left\langle\overline{\mathbb{W}},<_{0}\right\rangle$, we see by composing these three isomorphisms that $\circ \circ \alpha \Uparrow \circ \mathrm{o}^{-1}:\langle\mathrm{On},<\rangle \cong\left\langle o\left(\mathbb{W}_{\alpha}\right),<\right\rangle$, whence $\sigma^{\alpha}=\mathrm{o} \circ \alpha \Uparrow \circ \mathrm{o}^{-1}$.

So seeing $\alpha \Uparrow$ as an action of the ordinals on $\overline{\mathbb{W}}$, and $\sigma^{\alpha}$ as an action of the ordinals on the ordinals, the above tells us that the two actions are isomorphic. Let us draw a nice corollary from our lemma.

Corollary 5.7. For any worm $A, o(\alpha \uparrow A)=\sigma^{\alpha} o(A)$
Proof. We have that

$$
\mathrm{o}(\alpha \uparrow A)=\mathrm{o}(\overline{\alpha \uparrow A})=\mathrm{o}(\alpha \Uparrow \bar{A})=\mathrm{o}\left(\alpha \Uparrow \mathrm{o}^{-1}(o(A))\right)=\sigma^{\alpha} o(A) .
$$

We may recast this by stating that the following diagram commutes:


With this we now obtain a complete calculus for computing $o$ and $o_{\alpha}$ once we know how to compute the functions $\sigma^{\alpha}$.
Theorem 5.8. Given worms $A, B$ and an ordinal $\xi$,

1. $o(T)=0$;
2. $o(A 0 B)=o(B)+1+o(A)$;
3. $o(\xi \uparrow A)=\sigma^{\xi} o(A)$.

The calculus in this theorem looks efficient and elegant but we seem to be running in circles here. To compute $o$ we need to know how to compute the worm-enumerators $\sigma^{\xi}$. The $\sigma^{\xi}$ in their turn are defined in terms of $o$. In the next section we shall see how we can break out of this circle and provide a stand-alone calculus for our worm-enumerators.

## 6 Computing the worm enumerators $\sigma^{\xi}$

In this section we shall see how the worm enumerators $\sigma^{\alpha}$ can be computed. We shall provide a recursive calculus in Theorem 6.11

### 6.1 Worm enumerators: basic properties

The first step in characterizing the worm enumerators we get almost for free and consists of determining the ordinal function $\sigma^{0}$.
Lemma 6.1. The function $\sigma^{0}$ is the identity function on the ordinals.
Proof. Recall that by definition, $\sigma^{0}$ enumerates $o\left(\mathbb{W}_{0}\right)$ in increasing order. Since o defines an isomorphism between $\overline{\mathbb{W}}$ and On, we see that $o\left(\mathbb{W}_{0}\right)=o(\mathbb{W})=$ On. Evidently, On is enumerated by the identity function, so that $\sigma^{0}(\alpha)=\alpha$ for each ordinal $\alpha$.

As a second step in characterizing the worm enumerators $\sigma^{\alpha}$, we shall prove that for each ordinal $\alpha$, the corresponding $\sigma^{\alpha}$ is a normal (both increasing and continuous) function. Since each $\sigma^{\alpha}$ by definition enumerates a class of ordinals in increasing order, it is clear that each $\sigma^{\alpha}$ is increasing.

So, next we need to see that each $\sigma^{\alpha}$ is continuous, that is, that computing $\sigma^{\alpha}$ commutes with taking suprema: if $\Delta$ is a set of ordinals then $\sigma^{\alpha}(\sup \Delta)=$ $\sup \sigma^{\alpha} \Delta$. Since we know that $\sigma^{\alpha}=0 \circ \alpha \Uparrow \circ \mathrm{o}^{-1}$ (this is Lemma 5.6), it suffices
to prove that taking suprema commutes with all of $\mathrm{o}, \mathrm{o}^{-1}$, and $\Uparrow$. For o and $\mathrm{o}^{-1}$ this has been established in Lemma 4.5 it remains to show that $\alpha \uparrow$ can be viewed as continuous on $\mathbb{W}$. We shall state continuity for 'limit worms'. Recall from Lemma 4.8 that $o(A) \in \operatorname{Lim}$ iff $h(A) \neq \top$ in which case $A=\sup _{B<A} B$.

Lemma 6.2. Given a worm $A$ with $h(A) \neq \top$ and an ordinal $\alpha>0$, then

$$
\begin{aligned}
& \alpha \uparrow A=\alpha \uparrow\left(\sup _{B<A} B\right)=\sup _{B<A} \alpha \uparrow B \quad \text { and } \\
& \alpha \Uparrow A=\sup _{B<A} \alpha \Uparrow B
\end{aligned}
$$

Proof. Clearly, the second item follows from the first one. In Lemma 4.8 it is proven that $A=\sup _{B<A} B$ whenever $h(A) \neq \top$ so we will concentrate on showing $\alpha \uparrow A=\sup _{B<A} \alpha \uparrow B$.

From $B<A$ we get by Lemma 3.13 that $\alpha \uparrow B<{ }_{\alpha} \alpha \uparrow A$ whence $\alpha \uparrow B<$ $\alpha \uparrow A$ so that $\sup _{B<A} \alpha \uparrow B \leq \alpha \uparrow A$ is clear.

In order to show that $\alpha \uparrow A \leq \sup _{B<A} \alpha \uparrow B$ it suffices to prove that for any $C<\alpha \uparrow A$ we can find $B<A$ such that $C<\alpha \uparrow B$. This we prove by induction on $|C|$, with the base case $(C=\top)$ being trivial so that we may write $C=\beta \uparrow\left(C_{1} 0 C_{0}\right)$. Moreover we shall pick $C_{0}$ and $C_{1}$ in the unique way so that $C_{1} \in \mathbb{W}_{1}$ (this includes the case $C_{1}=\top$ ).

If $\beta<\alpha$, then clearly both $C$ and $\alpha \uparrow A$ belong to $\mathbb{W}_{\beta}$. Using Lemmas 3.12 and 3.13, we have $C_{1} 0 C_{0}=\beta \downarrow C<\beta \downarrow(\alpha \uparrow A)=(-\beta+\alpha) \uparrow A$. Since $C_{0}, C_{1} \leq C_{1} 0 C_{0}$, by the induction hypothesis for $\left|C_{0}\right|,\left|C_{1}\right|<|C|$ there are $B_{0}, B_{1}<A$ with $C_{i}<(-\beta+\alpha) \uparrow B_{i}$. Taking $B=\max \left\{B_{0}, B_{1}\right\}$ we see that for $i=0,1,(-\beta+\alpha) \uparrow B \geq(-\beta+\alpha) \uparrow B_{i}>C_{i}$. Since we chose $C_{1}$ with $C_{1} \in \mathbb{W}_{1}$, we also have $(-\beta+\alpha) \uparrow B>_{1} C_{1}$ whence $(-\beta+\alpha) \uparrow B \vdash$ $1 C_{1} \wedge 0 C_{0} \vdash 1 C_{1} 0 C_{0} \vdash 0 C_{1} 0 C_{0}$, i.e., $C_{1} 0 C_{0}<(-\beta+\alpha) \uparrow B$. It follows that $C=\beta \uparrow\left(C_{1} 0 C_{0}\right)<\beta \uparrow((-\beta+\alpha) \uparrow B)=\alpha \uparrow B$, as was to be proven.

If $\beta \geq \alpha$ we get that $C, \alpha \uparrow A \in \mathbb{W}_{\alpha}$, so that from $C<\alpha \uparrow A$ we obtain $\alpha \downarrow$ $C<\alpha \downarrow(\alpha \uparrow A)=A$. Since $A$ is a limit, we have that $\alpha \downarrow C<0(\alpha \downarrow C)<A$, thus $C=\alpha \uparrow(\alpha \downarrow C)<\alpha \uparrow(0(\alpha \downarrow C))$, as was to be shown.

Now that we have established that $\alpha \uparrow$ can be viewed as continuous on $\mathbb{W}$ we can prove that the $\sigma^{\alpha}$ are continuous ordinal functions.
Lemma 6.3. Each $\sigma^{\alpha}$ is a normal function.
Proof. It is clear that $\sigma^{\alpha}$ is increasing so we only need to see that $\sigma^{\alpha}$ is continuous for each $\alpha$. For the continuity of $\sigma^{\alpha}$ we reason as follows. Let $\lambda \in \operatorname{Lim}$ so that for some worm $A$ with $h(A) \neq \top$ we have $\lambda=o(A)$.

$$
\begin{array}{rlllll}
\sigma^{\alpha}(\lambda) & = & \sigma^{\alpha} o(A) & = & \mathrm{o}(\alpha \Uparrow) \mathrm{o}^{-1} o(A) & = \\
& = & \mathrm{o}(\alpha \Uparrow) \sup _{B<A} \bar{B} & = & \mathrm{o} \sup _{B<A}(\alpha \Uparrow) \bar{B} & = \\
\sup _{B<A} \mathrm{o}(\alpha \Uparrow) \bar{B} \\
& =\sup _{B<A} \mathrm{o}(\alpha \Uparrow) \mathrm{o}^{-1} o(B) & = & \sup _{B<A} \sigma^{\alpha} o(B) & = & \sup _{\beta<\lambda} \sigma^{\alpha} \beta,
\end{array}
$$

where we use Lemmas 4.5 and 6.2 to commute sup with o and $\alpha \Uparrow$, respectively.

As a next step in characterizing the $\sigma^{\alpha}$ functions we shall set out to determine $\sigma^{1}$. We first look at the $<_{1}$-first non-trivial worm in $\mathbb{W}_{1}$. It is not hard to prove by elementary methods that $1=\sup _{n \in \omega} 0^{n}$. In this sense, $\left\langle 0^{n}\right\rangle_{n \in \omega}$ is a natural fundamental sequence for the worm 1. Since we know that $o\left(o^{n}\right)=n$ we see that $o(1)=o\left(\sup _{n \in \omega} 0^{n}\right)=\sup _{n \in \omega} o\left(0^{n}\right)=\omega$.

In a similar fashion we see that for the $<_{1}$-second non-trivial worm in $\mathbb{W}_{1}$ which is 11 , can prove that $11=\sup _{n \in \omega}(10)^{n} 1$ so that $\left\langle(10)^{n} 1\right\rangle_{n \in \omega}$ can be conceived as a natural fundamental sequence for the worm 11. Using that $o(1)=\omega$, by repeatedly applying Lemma 5.4 we know that $o\left((10)^{n} 1\right)=\omega \cdot(n+1)$ so that $o(11)=o\left(\sup _{n \in \omega}(10)^{n} 1\right)=\sup _{n \in \omega} o\left((10)^{n} 1\right)=\sup _{n \in \omega} \omega \cdot(n+1)=\omega^{2}$.

The following two lemmas are inspired by these examples and establish that limit worms admit uniform fundamental sequences and that there is essentially a power $\omega$ difference between $o(A)$ and $o(1 \uparrow A)$.

Lemma 6.4. Let $B \in \mathbb{W}_{1}$. For $A=1 B$, we have that $A=\sup _{n<\omega} A[n]$, where $A[0]=B$ and $A[n+1]=B 0 A[n]$.

Proof. That $A>A[n]$ for all $n$ follows by induction; the base case is easy since $A=1 B \vdash 0 B$. For the induction step,

$$
A \stackrel{\mathrm{IH}}{\vdash} 1 B \wedge 0 A[n] \equiv 1 B 0 A[n] \vdash 0 B 0 A[n]=0 A[n+1] .
$$

Meanwhile, we prove by induction on the length of $C$ that if $C<A$ then $C<A[n]$ for some $n$. We assume $C \neq \top$, otherwise the claim is trivial. Thus, from $b(C)<C<A$, the fact that $|b(C)|<|C|$ and the induction hypothesis we obtain $b(C)<A[n]$ for some $n$ so that also $0 A[n] \vdash 0 b(C)$.

Then, from $h(C) \leq C<_{1} A$ we obtain $h(C) \leq B$ whence $h(C) \leq_{1} B$ and $B \vdash 1 h(C)$. Thus,

$$
A[n+1]=B 0 A[n] \equiv B \wedge 0 A[n] \vdash 1 h(C) \wedge 0 b(C) \equiv 1 h(C) 0 b(C) \vdash 0 h(C) 0 b(C)
$$

whence $C<A[n+1]$.
With the use of this lemma, we can now establish a relation between $o(A)$ and $o(1 \uparrow A)$, as to determine $\sigma^{1}$ :

## Lemma 6.5.

1. Given a worm $A$, we have that $o(1 \uparrow A)=-1+\omega^{o(A)}$;
2. For each ordinal $\xi$ we have $\sigma^{1} \xi=-1+\omega^{\xi}$.

Proof. We first prove 1 by induction on $<_{0}$. For the base case, $A=1 \uparrow A=\top$, we verify that $o(1 \uparrow A)=0=-1+\omega^{0}=-1+\omega^{o(A)}$.

If $A$ is a limit worm, i.e. $h(A) \neq \top$, the claim follows from Lemmas 4.5 and 6.2 since

$$
-1+\omega^{o(A)}=\sup _{B<A}\left(-1+\omega^{o(B)}\right) \stackrel{\mathrm{IH}}{=} \sup _{B<A} o(1 \uparrow B)=o\left(1 \uparrow \sup _{B<A} B\right)=o(1 \uparrow A)
$$

If $A=0 B$, then by Lemma 6.4 we have that $1 \uparrow A=\sup _{n<\omega} B_{n}$, where $B_{0}:=1 \uparrow B$ and $B_{n+1}:=(1 \uparrow B) 0 B_{n}$. Therefore $o(1 \uparrow A)=\sup _{n<\omega} o\left(B_{n}\right)$.

By an easy subsidiary induction on $n$ we now see that $o\left(B_{n}\right)=-1+\omega^{o(B)}$. $(n+1)$. For $n=0$ this is just the induction hypothesis of the main induction. For $n+1$ we apply Lemma 5.4 to obtain

$$
o\left(B_{n+1}\right)=o\left((1 \uparrow B) 0 B_{n}\right)=o\left(B_{n}\right)+1+o((1 \uparrow B))
$$

By the subsidiary induction we have that $o\left(B_{n}\right)=-1+\omega^{o(B)} \cdot(n+1)$ and by the main induction we have that $o((1 \uparrow B))=-1+\omega^{o(B)}$ so that
$o\left(B_{n}\right)+1+o((1 \uparrow B))=-1+\omega^{o(B)} \cdot(n+1)+1+-1+\omega^{o(B)}=-1+\omega^{o(B)} \cdot(n+2)$
as was to be shown. We now conclude the argument by observing that for $A=0 B$ we have
$o(1 \uparrow A)=\sup _{n<\omega} o\left(1 \uparrow B_{n}\right)=\sup _{n<\omega}-1+\omega^{o(B)} \cdot(n+1)=-1+\omega^{o(B)+1}=-1+\omega^{o(0 B)}$ whence $o(1 \uparrow A)=-1+\omega^{o(A)}$, as claimed.

To see Item 2, choose an ordinal $\xi$ and a worm $A$ such that $o(A)=\xi$. By Theorem 5.833, we have that $\sigma^{1} \xi=\sigma^{1} o(A)=o(1 \uparrow A)$; but by the previous item this is equal to $-1+\omega^{o(A)}=-1+\omega^{\xi}$, as desired.

Note that $-1+\omega^{\xi}=\omega^{\xi}$ whenever $\xi \neq 0$. At this point we may give a convenient breakdown of $o(A)$ in terms of its head, rest and body.

Corollary 6.6. If $A$ is any worm, then

$$
o(A)=o(r(A))+o(h(A))=o(b(A))+\omega^{o(1 \downarrow h(A))} .
$$

Proof. Immediate from Theorem 5.8 and Lemma 6.5
With a lemma similar to Lemma 6.5 we can now characterize $\sigma^{2}$ since $2 \uparrow=$ $1 \uparrow \circ 1 \uparrow$ and we just need to iterate Lemma 6.5. The more general fact that $(\alpha+\beta) \uparrow=\alpha \uparrow \circ \beta \uparrow$ is reflected in the following lemma.

## Lemma 6.7.

1. $\sigma^{0}=\mathrm{id}$;
2. $\sigma^{1}=e$ where $e(\xi)=-1+\omega^{\xi}$;
3. $\sigma^{\alpha+\beta}=\sigma^{\alpha} \circ \sigma^{\beta}$.

Proof. The first item is Lemma 6.1 and the second item is Lemma 6.5.
For the last item we see that

$$
\begin{aligned}
\sigma^{\alpha} \circ \sigma^{\beta} & =\mathrm{O} \circ \alpha \Uparrow \circ \mathrm{o}^{-1} \circ \mathrm{\circ} \circ \beta \Uparrow \circ \mathrm{o}^{-1}=\mathrm{O} \circ \alpha \Uparrow \circ \beta \Uparrow \circ \mathrm{o}^{-1} \\
& =\mathrm{O} \circ(\alpha+\beta) \Uparrow \circ \mathrm{o}^{-1}=\sigma^{\alpha+\beta} .
\end{aligned}
$$

Clearly this lemma, together with Theorem 5.8 completely determines the order-types for worms that only use natural numbers as ordinals.
Example 6.8. $o(2103)=o(3)+1+o(21)=\sigma^{3} o(0)+1+o(21)=\sigma^{3}(1)+1+$
$o(21)=\sigma^{1} \circ \sigma^{1} \circ \sigma^{1}(1)+1+o(21)=\omega^{\omega^{\omega}}+1+o(21)=\omega^{\omega^{\omega}}+1+\sigma^{1} o(10)=$
$\omega^{\omega^{\omega}}+1+\sigma^{1} o(10 \top)=\omega^{\omega^{\omega}}+1+\sigma^{1}(o(\top)+1+o(1))=\omega^{\omega^{\omega}}+1+\sigma^{1} \omega=$
$\omega^{\omega^{\omega}}+1+\omega^{\omega}=\omega^{\omega^{\omega}}+\omega^{\omega}$.

### 6.2 A recursive calculus for the worm enumerators

It is evident that Lemma 6.7 says nothing about the behavior of $\sigma^{\alpha}$ for additively indecomposable $\alpha$. To deal with those ordinals we have the following lemmas.

Lemma 6.9. If $\lambda$ is infinite and additively indecomposable then

$$
\lambda \uparrow(0 A)=\sup _{\eta<\lambda} \eta(\lambda \uparrow A)
$$

Proof. It is evident that $\sup _{\eta<\lambda} \eta(\lambda \uparrow A) \leq \lambda(\lambda \uparrow A)=\lambda \uparrow(0 A)$ so we shall show the other inequality by proving that for each $B<\lambda \uparrow(0 A)$ there is some $\eta<\lambda$ so that $B \leq \eta(\lambda \uparrow A)$. We distinguish two cases.

First assume that $B=\lambda \uparrow B^{\prime}$. Then, $B^{\prime}<0 A$ and thus $B^{\prime} \leq A$ so that $0\left(\lambda \uparrow B^{\prime}\right) \leq 0(\lambda \uparrow A)<\lambda \uparrow(0 A)$.

Otherwise, there are $\gamma<\lambda$, a worm $B_{1} \in \mathbb{W}_{1}$ and $B_{0} \in \mathbb{W}$ such that $B=\gamma \uparrow$ $\left(B_{1} 0 B_{0}\right)$. By induction on length, there are $\eta_{0}, \eta_{1}<\lambda$ such that $B_{i}<\eta_{i}(\lambda \uparrow A)$. Letting $\eta^{\prime}=1+\max \left\{\eta_{0}, \eta_{1}\right\}$ (so that $\eta^{\prime}>0$ ) we see that $B_{1}<\eta^{\prime}(\lambda \uparrow A)$ whence $B_{1}<_{1} \eta^{\prime}(\lambda \uparrow A)$. Thus, $\eta^{\prime}(\lambda \uparrow A) \vdash 1 B_{1} \wedge 0 B_{0} \vdash 1 B_{1} 0 B_{0} \vdash 0 B_{1} 0 B_{0}$ whence $B_{1} 0 B_{0}<\eta^{\prime}(\lambda \uparrow A)$. Since $\lambda$ is additively indecomposable we see that $\gamma+\lambda=\lambda$ and $\eta=\gamma+\eta^{\prime}<\lambda$, while $B=\gamma \uparrow\left(B_{1} 0 B_{0}\right)<\eta(\lambda \uparrow A)$, as needed.

For any ordinal $\lambda$ we have that $\sigma^{\lambda}(0)=o(T)=0$. Moreover, since $\sigma^{\lambda}$ is continuous, we can compute $\sigma^{\lambda}$ on limit ordinals if we have computed the values for all smaller ordinals. Thus, we only need to study the behavior of $\sigma^{\lambda}$ on successor ordinals for which we have the next lemma.

Lemma 6.10. Let $\lambda$ be an additively indecomposable limit ordinal. We have that

$$
\sigma^{\lambda}(\beta+1)=\sup _{\eta<\lambda} \sigma^{\eta}\left(\sigma^{\lambda}(\beta)+1\right)
$$

Proof. Pick $B^{\prime}$ so that $o\left(0 B^{\prime}\right)=\beta+1$ and let $B:=\lambda \uparrow B^{\prime}$ so that by Corollary 5.7 we obtain

$$
\begin{equation*}
o(B)=o\left(\lambda \uparrow B^{\prime}\right)=\sigma^{\lambda} o\left(B^{\prime}\right)=\sigma^{\lambda} \beta \tag{6}
\end{equation*}
$$

Moreover, as $\lambda$ is additively indecomposable, we see that $-\eta+\lambda=\lambda$ for any $\eta<\lambda$. In particular we get that

$$
\begin{equation*}
\eta \downarrow B=B \quad \text { for any } \eta<\lambda \tag{7}
\end{equation*}
$$

By Lemma 6.9, $\lambda B=\sup _{\eta<\lambda} \eta B$, so that

$$
\begin{equation*}
o(\lambda B)=\sup _{\eta<\lambda} o(\eta B) \tag{8}
\end{equation*}
$$

Recall that $o\left(0 B^{\prime}\right)=\beta+1$ so that we can reason

$$
\begin{aligned}
\sigma^{\lambda}(\beta+1) & =\sigma^{\lambda}\left(o\left(0 B^{\prime}\right)\right) & & \\
& =o(\lambda B) & & \text { by Corollary } 5.7 \\
& =\sup _{\eta<\lambda} o(\eta B) & & \text { by (8) } \\
& =\sup _{\eta<\lambda} \sigma^{\eta} o(0(\eta \downarrow B)) & & \text { by Corollary } 5.7 \\
& =\sup _{\eta<\lambda} \sigma^{\eta}(o(\eta \downarrow B)+1) & & \text { by Lemma 4.7 } \\
& =\sup _{\eta<\lambda} \sigma^{\eta}(o(B)+1) & & \text { by (7) } \\
& =\sup _{\eta<\lambda} \sigma^{\eta}\left(\sigma^{\lambda}(\beta)+1\right) & & \text { by (6). }
\end{aligned}
$$

Now that we have proved this lemma we finally have fully determined all functions $\sigma^{\alpha}$.

Theorem 6.11. For ordinals $\alpha$ and $\beta$, the values $\sigma^{\alpha}(\beta)$ are determined by the following recursion.

1. $\sigma^{\alpha} 0=0$ for all $\alpha \in \mathrm{On}$;
2. $\sigma^{1}=e$ with $e(\xi)=-1+\omega^{\xi}$;
3. $\sigma^{\alpha+\beta}=\sigma^{\alpha} \sigma^{\beta}$;
4. $\sigma^{\alpha}(\lambda)=\sup _{\beta<\lambda} \sigma^{\alpha}(\beta)$ for limit ordinals $\lambda$;
5. $\sigma^{\lambda}(\beta+1)=\sup _{\eta<\lambda} \sigma^{\eta}\left(\sigma^{\lambda}(\beta)+1\right)$ for $\lambda$ an additively indecomposable limit ordinal.

It is clear that this theorem embodies a full calculus. Let us see a simple example.

Example 6.12. $\sigma^{\omega} 1=\varepsilon_{0}$ so that $o(\langle\omega\rangle \top)=\varepsilon_{0}$ with $\varepsilon_{0}:=\sup \left\{\omega, \omega^{\omega}, \omega^{\omega^{\omega}}, \ldots\right\}$.
Proof. By definition, $\sigma^{1}(1)=e 1=\omega$. Consequently, $\sigma^{2}(1)=\sigma^{1} \sigma^{1}(1)=\omega^{\omega}$ and likewise $\sigma^{3}(1)=\omega^{\omega^{\omega}}$, etc. Thus, $\sigma^{\omega} 1=\sup _{\eta<\omega} \sigma^{\eta}\left(\sigma^{\omega}(0)+1\right)=\sup _{\eta<\omega} \sigma^{\eta}(1)=$ $\varepsilon_{0}$.

It should not come as a surprise that $\sigma^{\omega}(1)$ is a fixpoint of $e$ and something more general holds. If $\lambda$ is additively indecomposable, then $\eta+\lambda=\lambda$ for all $\eta<\lambda$. Thus

$$
\sigma^{\lambda}(\xi)=\sigma^{\eta+\lambda}(\xi)=\sigma^{\eta} \sigma^{\lambda}(\xi)
$$

so that for each $\xi$ we see that $\sigma^{\lambda}(\xi)$ is a fixpoint of $\sigma^{\eta}$ for each $\eta<\lambda$; this is very similar to what happens with the Veblen hierarchy. In Theorem 7.5 the relation between the worm enumerators and the one-placed Veblen functions is made precise.

## 7 Hyperations and Cohyperations

A main theme of this paper is how to compute given a worm $A$ and ordinal $\xi$ its $\xi$-consistency order-type $o_{\xi}(A)$. In Corollary 4.16 we have reduced the $o_{\xi}$ ordertypes to the plain o order-type. Subsequently, in Theorem 5.8 we provided a calculus for $o$ in terms of the so-called worm-enumerators $\sigma^{\alpha}$. Finally, in Theorem 6.11 we worked out a recursive calculus for computing the wormenumerators thereby completing all ingredients needed to compute any ordertype $o_{\xi}(A)$.

In the final section of this paper we wish to characterize what, given a worm $A$, the sequence $\left\langle o_{\xi}(A)\right\rangle_{\xi \in \text { On }}$ can look like. It shall turn out that to give a smooth characterization of these sequences, we need certain well-behaved left-inverses of our worm enumerator functions. These inverses can be computed within the general framework of what the authors call hyperations and co-hyperations.

Hyperations are a kind of transfinite iteration of certain ordinal functions and were introduced and studied in full generality by the authors in [15]. In this section we shall briefly state -without proof- the main properties of hyperations and the related cohyperations that we need in the remainder of this paper and refer to [15] for further background. Moreover, we shall prove that the wormenumerators $\sigma^{\alpha}$ are the hyperation of a special form of ordinal exponentiation. For definitions and basic properties of ordinals, we refer the reader to [22, 24].

### 7.1 Hyperations

As mentioned before, hyperation is a form of transfinite iteration of normal functions. It is based on the additivity of finite iterations, that is $f^{m+n}=f^{m} \circ f^{n}$ generalizing this to the transfinite setting. Let us first recall the definition of a normal function.

We call a function on the ordinals $f$ increasing if $\alpha<\beta$ implies $f(\alpha)<$ $f(\beta)$. An ordinal function is called continuous if $\sup _{\zeta<\xi} f(\zeta)=f(\xi)$ for all limit ordinals $\xi$. Functions which are both increasing and continuous are called normal.

Definition 7.1 (Weak hyperation). A weak hyperation of a normal funcion $f$ is a family of normal functions $\left\langle g^{\xi}\right\rangle_{\xi \in \mathrm{On}}$ such that

1. $g^{0} \xi=\xi$ for all $\xi$,
2. $g^{1}=f$,
3. $g^{\xi+\zeta}=g^{\xi} g^{\zeta}$.

Par abuse de langage we will often write just $g^{\xi}$ instead of $\left\langle g^{\xi}\right\rangle_{\xi \in \mathrm{On}}$. In Lemma 6.7 we have proven that the family of worm enumerators $\sigma^{\xi}$ is a weak hyperation of the function $e$ defined as $\xi \mapsto-1+\omega^{\xi}$.

Weak hyperations are not unique. However, if we impose a minimality condition, we can prove that there is a unique minimal hyperation.

Definition 7.2 (Hyperation). A weak hyperation $g^{\xi}$ of $f$ is minimal if it has the property that, whenever $h^{\xi}$ is a weak hyperation of $f$ and $\xi, \zeta$ are ordinals, then $g^{\xi} \zeta \leq h^{\xi} \zeta$.

If $f$ has a (unique) minimal weak hyperation, we call it the hyperation of $f$ and denote it $f^{\xi}$.

We shall now prove that the worm enumerators $\sigma^{\xi}$ are the hyperation of the function $e$.
Theorem 7.3. $\sigma^{\alpha}$ is the hyperation of the function $e: \xi \mapsto-1+\omega^{\xi}$.
Proof. The properties 1.-3. of Lemma 6.7 express that the $\sigma^{\alpha}$ are a weak hyperation of $e$. To see that it is the unique hyperation we only need to check for minimality.

So, suppose that $\left\{f^{\alpha}\right\}_{\alpha \in \text { Ord }}$ is a collection of normal functions such that 1. -3 . holds. By induction on $\alpha$ we shall see that $\sigma^{\alpha}(\beta) \leq f^{\alpha}(\beta)$.

For $\alpha=0$ and $\alpha=1$ this is obvious and for additively decomposable ordinals we see that $\sigma^{\alpha+\beta}=\sigma^{\alpha} \sigma^{\beta} \leq_{\mathrm{IH}} f^{\alpha} f^{\beta}=f^{\alpha+\beta}$.

So, let $\alpha$ be an indecomposable limit ordinal. We proceed by an auxiliary induction on $\beta$ to show that $\sigma^{\alpha}(\beta) \leq f^{\alpha}(\beta)$ which clearly holds for $\beta=0$. As both $f^{\alpha}$ and $\sigma^{\alpha}$ are continuous, we only need to consider successor ordinals in which case, by Lemma 6.10 we see that

$$
\begin{equation*}
\sigma^{\alpha}(\beta+1)=\sup _{\alpha^{\prime}<\alpha} \sigma^{\alpha^{\prime}}\left(\sigma^{\alpha}(\beta)+1\right) \leq_{\mathrm{H}} \sup _{\alpha^{\prime}<\alpha} f^{\alpha^{\prime}}\left(f^{\alpha}(\beta)+1\right) \tag{9}
\end{equation*}
$$

As for $\alpha^{\prime}<\alpha$ we have $\alpha^{\prime}+\alpha=\alpha$, by Property 3 , we see that $f^{\alpha}(\beta+1)=$ $\sup _{\alpha^{\prime}<\alpha} f^{\alpha^{\prime}} f^{\alpha}(\beta+1)$. But, as $f^{\alpha}$ is monotone we also see that $f^{\alpha}(\beta+1) \geq$ $f^{\alpha}(\beta)+1$ whence by monotonicity of all of the $f^{\alpha^{\prime}}$ we see that

$$
f^{\alpha}(\beta+1)=\sup _{\alpha^{\prime}<\alpha} f^{\alpha^{\prime}} f^{\alpha}(\beta+1) \geq \sup _{\alpha^{\prime}<\alpha} f^{\alpha^{\prime}}\left(f^{\alpha}(\beta)+1\right)
$$

We combine this with (9) to conclude that $\sigma^{\alpha}(\beta+1) \leq f^{\alpha}(\beta+1)$.
Since we have proven that $e^{\xi}=\sigma^{\xi}$, from now on we shall only use the notation based on $e$. We call the family $e^{\xi}$ hyperexponentials.

Hyperations in general allow for an explicit recursive definition very much in the style of Theorem 6.11. Moreover, there turns out to be a close connection between hyperations and Veblen progressions:

Definition 7.4. For $f$ a normal function, the Veblen progression based on $f$ is denoted by $\left\langle f_{\xi}\right\rangle_{\xi \in \mathrm{On}}$ and defined by $f_{0}:=f$ and for $\xi>0, f_{\xi}$ is the normal function that enumerates in increasing order the ordinals which are simultaneous fixpoints for all the $f_{\zeta}$ for $\zeta<\xi$.

The Veblen progression based on $\varphi(\xi):=\omega^{\xi}$ are the well-known one-place Veblen functions $\varphi_{\xi}$. Note that $\varphi(\xi)=e(\xi)$ for $\xi \neq 0$. In 15 it is proven that hyperations can be seen as a natural refinement of Veblen progressions.

Theorem 7.5. For $f$ a normal function we have that $f^{\omega^{\xi}}=f_{\xi}$.
It should be clear that this theorem provides an easy link between the hyperexponentials and the earlier known Veblen functions. In particular, if we write some ordinal $\xi$ in its unique Cantor Normal Form with base $\omega$ so that for some $n \geq 0$ we have $\xi=\omega^{\xi_{1}}+\ldots+\omega^{\xi_{n}}$ and $\xi_{i} \geq \xi_{i+1}$ for all $i<n$, then for $\alpha>0$ and $\xi \notin \operatorname{Lim}$ (for $\xi \in \operatorname{Lim}$ we should replace the $\alpha$ on the right side by $-1+\alpha$ ):

$$
e^{\xi}(\alpha)=\varphi_{\xi_{1}} \circ \ldots \circ \varphi_{\xi_{n}}(\alpha)
$$

### 7.2 Cohyperations

We shall see that in order to relate the different $\xi$-consistency order-types $o_{\xi}$ to each other we shall need left inverses to hyperexponentials. Hyperations are injective and hence invertible on the left; however, a left-inverse of a hyperation is typically not a hyperation, but a different form of transfinite iteration we call cohyperation.

Instead of transfinitely iterating normal functions we shall consider initial functions. We will say a function $f$ is initial if, whenever $I$ is an initial segment (i.e., of the form $[0, \beta)$ for some $\beta$ ), then $f(I)$ is an initial segment. It is easy to see that $f \xi \leq \xi$ for initial functions $f$.

Definition 7.6 (Cohyperation). A weak cohyperation of an initial function $f$ is a family of initial functions $\left\langle g^{\xi}\right\rangle_{\xi \in \mathrm{On}}$ such that

1. $g^{0} \xi=\xi$ for all $\xi$,
2. $g^{1}=f$,
3. $g^{\xi+\zeta}=g^{\zeta} g^{\xi}$.

If $g$ is maximal in the sense that $g^{\xi} \zeta \geq h^{\xi} \zeta$ for every weak cohyperation $h$ of $f$ and all ordinals $\xi, \zeta$, we say $g$ is the cohyperation of $f$ and write $f^{\xi}=g^{\xi}$.

Both hyperations and cohyperations are denoted using the superscript; however, this does not lead to a clash in notation as the only function that is both normal and initial is the identity. In [15], a general recursive scheme to compute actual cohyperations is given very much in the spirit of Theorem 6.11.

Let $f$ be a normal function. Then, $g$ is a left adjoint for $f$ if, for all ordinals $\alpha, \beta$,

1. if $\alpha=f(\beta)$, then $g(\alpha)=\beta$ and
2. if $\alpha<f(\beta)$, then $g(\alpha)<\beta$.

Left-adjoints are natural left-inverses and cohyperating them yields leftadjoints to the corresponding hyperations in a uniform way:
Theorem 7.7. Given a normal function $f$ with left adjoint $g$ and ordinals $\xi \leq \zeta$ and $\alpha, g^{\xi} f^{\zeta}=f^{-\xi+\zeta}$ and $g^{\zeta} f^{\xi}=g^{-\xi+\zeta}$.

We shall need left-adjoints to our hyperexponentials. In order to formulate them, let us first recall some basic properties of the ordinals.

Lemma 7.8. Given any ordinal $\eta>0$, there exist ordinals $\alpha, \beta$, where $\beta$ is uniquely determined, such that $\eta=\alpha+\omega^{\beta}$. We will denote this unique $\beta$ by $\ell \eta$ and define $\ell 0=0$.

The following theorem is proven in [15].
Theorem 7.9. The function $\ell$ is a left adjoint to $e$, and thus $\ell^{\xi}$ is left adjoint to $e^{\xi}$ for all $\xi$.

We will refer to the functions $\ell^{\xi}$ as hyperlogarithms.

### 7.3 Exact sequences

A nice feature of cohyperations is that, in a sense, they need only be defined locally. To make this precise, we introduce the notion of an exact sequence.
Definition 7.10. Let $g^{\xi}$ be a cohyperation, and $f: \Lambda \rightarrow \Theta$ be an ordinal function.

Then, we say $f$ is g-exact if, given ordinals $\xi, \zeta$ with $\xi+\zeta<\Lambda, f(\xi+\zeta)=$ $g^{\zeta} f(\xi)$.

A $g$-exact function $f$ describes the values of $g^{\xi} f(0)$. However, for $f$ to be $g$-exact, we need only check a fairly weak condition:

Lemma 7.11. The following are equivalent:

1. $f$ is $g$-exact
2. for every ordinal $\xi, f(\xi)=g^{\xi} f(0)$
3. for every ordinal $\zeta>0$ there is $\xi<\zeta$ such that $f(\zeta)=g^{-\xi+\zeta} f(\xi)$.

Example 7.12. By $e^{\alpha} \beta \cdot \gamma$ we denote $\left(e^{\alpha} \beta\right) \cdot \gamma$. Then, the following sequence whose initial sub-sequence of non-zero elements is of length $\omega \cdot 2+2$ is $\ell$-exact:

$$
\left\langle e^{\omega^{2}} 1+e^{\omega}\left(e^{\omega+1} 1 \cdot 2\right), e^{\omega}\left(e^{\omega+1} 1 \cdot 2\right), \ldots, e^{\omega+1} 1 \cdot 2, e^{\omega+1} 1, \ldots, \omega, 1,0 \ldots\right\rangle
$$

Proof. By Lemma 7.11 and Theorem 7.9.
In the light of Theorem 7.5 we can reformulate this example in terms of the better-known Veblen functions. Thus, using the usual convention that $\varphi_{1}(\alpha)$ is denoted by $\varepsilon_{\alpha}$ we can rewrite our example as:

$$
\left\langle\varphi_{2}(1)+\varepsilon_{\varepsilon_{\omega}+\varepsilon_{\omega}}, \varepsilon_{\varepsilon_{\omega}+\varepsilon_{\omega}}, \ldots, \varepsilon_{\omega}+\varepsilon_{\omega}, \varepsilon_{\omega}, \ldots, \omega, 1,0 \ldots\right\rangle
$$

## 8 Consistency Sequences

Given a worm $A$, we define its consistency sequence to be the sequence $\left\langle o_{\xi}(A)\right\rangle_{\xi \in \mathrm{On}}$. In this section we give a full characterization of consistency sequences. That is, we will determine which sequences $\left\langle\alpha_{\xi}\right\rangle_{\xi \in O_{n}}$ are of the form $\vec{o}(A)$ for some worm $A$. Moreover, given $o_{0}(A)$, we will compute $o_{\xi}(A)$ for all $\xi>0$, even when $A$ itself is not explicitly given.

It is intuitively clear that for constant $A$, the function $o_{\xi}(A)$ is weakly decreasing in $\xi$ as is expressed in the following lemma.

Lemma 8.1. For $\xi<\zeta$ we have that $o_{\xi}(A) \geq o_{\zeta}(A)$.
Proof. By induction on $o_{\xi}(A)$ we see that
$o_{\xi}(A):=\sup _{B<\xi A}\left(o_{\xi}(B)+1\right) \geq \geq_{\mathrm{IH}} \sup _{B<\xi A}\left(o_{\zeta}(B)+1\right) \geq \sup _{B<{ }_{\zeta} A}\left(o_{\zeta}(B)+1\right)=o_{\zeta}(A)$.
Note that we have the last inequality since for $\xi<\zeta$ we have $B<\zeta$ implies $B<_{\xi} A$.

We will use the notation $\vec{o}(A)$ for the sequence $\left\langle o_{\xi}(A)\right\rangle_{\xi \in \mathrm{O}_{\mathrm{n}}}$; that is,

$$
\vec{o}(A):=\left\langle o_{0}(A), o_{1}(A), \ldots, o_{\omega}(A), o_{\omega+1}(A), \ldots\right\rangle
$$

By Lemma 8.1, $\vec{o}(A)$ is a weakly decreasing sequence of ordinals. In particular, we have that $\left\{o_{\xi}(A) \mid \xi \in \mathrm{On}\right\}$ is a finite set for any worm $A$. Moreover, we see that any consistency sequence eventually hits zero.

Corollary 8.2. Given a worm $A \neq \top$, we may write $A=\xi B$ for a unique ordinal $\xi$ and worm $B$. Then, given an arbitrary ordinal $\zeta$, we have that $o_{\zeta}(A)=$ 0 if and only if $\zeta>\xi$.

Proof. If $A=\xi B$ then clearly $h_{\xi}(A) \neq \top$, so that by Lemma 4.11) $o_{\xi}(A)=$ $o_{\xi} h_{\xi}(A) \neq 0$. On the other hand, for $\zeta>\xi, h_{\zeta}(A)=\top$ whence $o_{\zeta}(A)=0$.

Consistency sequences are of interest of their own. Moreover, they shed light on (Kripke) semantics for the closed fragment of GLP. Also, they admit a proof-theoretic interpretation as explained in 20.

### 8.1 A local characterization

We shall first provide a local characterization of the consistency sequences in that we relate the values in the sequence to its neighbors. To this end, let us first compute $o_{\xi+1}(A)$ in terms of $o_{\xi}(A)$. Recall that $\ell \alpha$ denotes the unique $\beta$ such that $\alpha=\alpha^{\prime}+\omega^{\beta}$ for $\alpha>0$, while $\ell 0=0$. The following lemma will be useful:

Lemma 8.3. Given an ordinal $\xi$ and a worm $A, o_{\xi+1} h_{\xi+1}(A)=\ell_{o} h_{\xi}(A)$.

Proof. Let $B=\xi \downarrow h_{\xi}(A)$, so that from Corollary 6.6 we have

$$
o_{\xi} h_{\xi}(A)=o\left(\xi \downarrow h_{\xi}(A)\right)=o(B)=o b(B)+\omega^{o(1 \downarrow h(B))},
$$

and thus $\ell_{o \xi} h_{\xi}(A)=o(1 \downarrow h(B))$. But $h(B)=\xi \downarrow h_{\xi+1}(A)$, so

$$
o(1 \downarrow h(B))=o\left((\xi+1) \downarrow h_{\xi+1}(A)\right)=o_{\xi+1} h_{\xi+1}(A),
$$

where the last equality is an instance of Corollary 5.7
Now we are ready to describe the relation between successive coordinates of the $\vec{o}(A)$ sequence.
Theorem 8.4. Given an ordinal $\xi$ and a worm $A, o_{\xi+1}(A)=\ell_{o \xi}(A)$
Proof. We have that $o_{\xi+1}(A)=o_{\xi+1} h_{\xi+1}(A)=\ell_{\xi} h_{\xi}(A)=\ell_{o_{\xi}}(A)$, where the second equality uses Lemma 8.3 and the others Lemma 4.11 .

Theorem 8.4 tells us what the relation between successor coordinates of $\vec{o}(A)$ is. We may also infer from it when successor coordinates are different; if $o_{\xi}(A)$ is a fixed point of $\zeta \mapsto-1+\omega^{\zeta}$ then $o_{\xi}(A)=o_{\xi+1}(A)$.

Next we shall determine what happens at limit steps in the consistency sequences. Since we know that for a given worm $A$, the set $\left\{o_{\xi}(A) \mid \xi \in \mathrm{On}\right\}$ is finite it is clear that for limit $\xi$, the value $o_{\xi}(A)$ can only be non-zero, if at some point before $\xi$, the sequence $\vec{o}(A)$ had stabilized. We shall now compute the relation between this stabilized value and the limit value. Here, our functions $e^{\xi}$ come back into play:

Theorem 8.5. Let $\zeta \in \operatorname{Lim}$ then, for $\theta$ large enough we have that

$$
o_{\theta}(A)=e^{-\theta+\zeta} o_{\zeta}(A)=e^{\omega^{\ell \zeta}} o_{\zeta}(A)=e_{\ell \zeta}\left(o_{\zeta}(A)\right)
$$

Proof. That $e^{\omega^{\ell \zeta}} o_{\zeta}(A)=e_{\ell \zeta}\left(o_{\zeta}(A)\right)$ is just Theorem 7.5 so we focus on the first equalities. Since $\zeta \in \operatorname{Lim}$, for $\xi$ large enough we have that $h_{\zeta}(A)=h_{\xi}(A)$, and more generally,

$$
\begin{equation*}
h_{\zeta}(A)=h_{\theta}(A) \tag{10}
\end{equation*}
$$

whenever $\theta \in[\xi, \zeta]$. Moreover, writing $\zeta=\zeta^{\prime}+\omega^{\ell \zeta}$, we may without loss of generality choose $\xi \geq \zeta^{\prime}$ so that $-\theta+\zeta=\omega^{l \zeta}$ for all $\theta \in[\xi, \zeta)$. For the sake of abbreviating, let $\delta=-\theta+\zeta=\omega^{\ell \zeta}$. By definition

$$
\begin{equation*}
\theta+\delta=\zeta \tag{11}
\end{equation*}
$$

Now we can prove our theorem starting with an application of Lemma 4.11

$$
\begin{aligned}
o_{\theta}(A) & =o_{\theta} h_{\theta}(A) & & \\
& =o_{\theta} h_{\zeta}(A) & & \text { By (10) } \\
& =o\left(\theta \downarrow h_{\zeta}(A)\right) & & \text { Corollary 4.16 } \\
& =o\left(\delta \uparrow\left((\theta+\delta) \downarrow h_{\zeta}(A)\right)\right) & & \text { Lemma } 3.126 \\
& =o\left(\delta \uparrow\left(\zeta \downarrow h_{\zeta}(A)\right)\right) & & \text { By (11) } \\
& =e^{\delta} o\left(\zeta \downarrow h_{\zeta}(A)\right) & & \text { Theorem 5.83 } \\
& =e^{\delta} o_{\zeta}(A) & & \text { Corollary 4.16 }
\end{aligned}
$$

Theorem 8.4 and Theorem 8.5 have been presented in a somewhat different guise in [21]. Note that, indeed, these theorems provide a local characterization of $\vec{o}(A)$ given any worm $A$ : start by computing $o_{0}(A)$ and compute how $o_{\xi}(A)$ changes for increasing values of $A$. Moreover, in 21] a characterization is given for those values $\xi$ where $o_{\xi}(A)$ changes value.

### 8.2 From local to global

The computations we have presented give the value of $o_{\xi}(A)$ from $o_{\zeta}(A)$ for $\zeta<\xi$ provided $\zeta$ is large enough. As such, we have only characterized them locally. In the next subsection we will give a global characterization of $\vec{o}(A)$, so that all values may be computed directly from $o_{0}(A)$.

In our local characterization we have distinguished two cases: successor coordinates and limit coordinates. It turns out that one can conceive both successor and limit steps as one of the same kind. For the successor case when $\zeta=\xi+1$ we saw that

$$
\begin{equation*}
o_{\zeta}(A)=o_{\xi+1}(A)=\ell_{o \xi}(A)=\ell^{-\xi+(\xi+1)} o_{\xi}(A)=\ell^{-\xi+\zeta_{o \xi}}(A) \tag{12}
\end{equation*}
$$

For limit steps, we say $o^{\xi}(A)=e^{-\xi+\zeta_{o \zeta}}(A)$ for $\xi<\zeta$ large enough. By Lemma 7.9, $\ell^{\alpha}$ is a left-inverse to $e^{\alpha}$ for all $\alpha$. Then our characterization for limit coordinates becomes

$$
\begin{equation*}
o_{\zeta}(A)=\ell^{-\xi+\zeta} o_{\xi}(A) \quad \text { for } \xi<\zeta \text { large enough. } \tag{13}
\end{equation*}
$$

Written in this way, we see that (12) and (13) actually are the same. Moreover, as we shall see, Lemma 7.11 will allow us to drop the requirement "for $\xi<\zeta$ large enough" giving rise to our desired global characterization. Let us unify the results obtained so far by describing the sequences $\vec{o}(A)$ using hyperexponentials and hyperlogarithms.

### 8.3 A global characterization

Theorem 8.6. If $A$ is any worm, $\vec{o}(A)$ is the unique $\ell$-exact sequence with $o_{0}(A)=o(A)$.

Proof. In view of Lemma 7.11 it suffices to show that, given any ordinal $\zeta$, there is $\xi<\zeta$ such that $o_{\zeta}(A)=\ell^{-\xi+\zeta} o_{\xi}(A)$.

If $\zeta$ is a successor ordinal, write $\zeta=\xi+1$. Then, by Theorem 8.4, we have that $o_{\zeta}(A)=\ell o_{\xi}(A)$.

Meanwhile, if $\zeta$ is a limit ordinal, we know from Lemma 8.5 that, for $\xi<\zeta$ large enough, $o_{\xi}(A)=e^{-\xi+\zeta} o_{\zeta}(A)$. Applying $\ell^{-\xi+\zeta}$ on both sides and using Theorem 7.7 we see that $\ell^{-\xi+\zeta_{o \xi}}(A)=o_{\zeta}(A)$. Thus we can use Lemma 7.11 to see that $\vec{o}(A)$ is $\ell$-exact, so that, for all $\xi, o_{\xi}(A)=\ell^{\xi} o_{0}(A)$, as claimed.

Example 8.7. For $A=\langle\omega \cdot 2+1\rangle\langle\omega\rangle\langle\omega \cdot 2+1\rangle\langle 0\rangle\left\langle\omega^{2}\right\rangle \top$ we have that
$\vec{o}(A)=\left\langle e^{\omega^{2}} 1+e^{\omega}\left(e^{\omega+1} 1 \cdot 2\right), e^{\omega}\left(e^{\omega+1} 1 \cdot 2\right), \ldots, e^{\omega+1} 1 \cdot 2, e^{\omega+1} 1, \ldots, \omega, 1,0 \ldots\right\rangle$.
Proof. Theorem6.11 yields $o(A)=e^{\omega^{2}} 1+e^{\omega}\left(e^{\omega+1} 1 \cdot 2\right)$. Then, the result follows from Theorem 8.6 above and Example 7.12 .

Recall that we can recast our example in terms of the more familiar Veblen functions as

$$
\vec{o}(A)=\left\langle\varphi_{2}(1)+\varepsilon_{\varepsilon_{\omega}+\varepsilon_{\omega}}, \varepsilon_{\varepsilon_{\omega}+\varepsilon_{\omega}}, \ldots, \varepsilon_{\omega}+\varepsilon_{\omega}, \varepsilon_{\omega}, \ldots, \omega, 1,0 \ldots\right\rangle
$$

Hyperexponentials give us lower bounds on $\ell$-exact sequences. The value of $o_{\xi}(A)$ fully determines the values of $o_{\zeta}(A)$ for $\zeta>\xi$ but not vice versa. However for $\zeta>\xi$ we do have a lower-bound on $o_{\xi}(A)$ :
Theorem 8.8. Given a worm $A$ and ordinals $\xi, \zeta, o_{\xi}(A) \geq e^{\zeta} o_{\xi+\zeta}(A)$.
Proof. Towards a contradiction, assume that there is a worm $A$ and ordinals $\xi<\zeta$ such that $o_{\xi}(A)<e^{-\xi+\zeta_{o \zeta}}(A)$. Then, by Theorem 7.9, $\ell^{-\xi+\zeta_{o \xi}}(A)<o_{\zeta}(A)$. But this is impossible by Theorem 8.6 given that $\ell^{-\xi+\zeta_{o \xi}}(A)=o_{\zeta}(A)$.

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