# Mokhov A, Khomenko V, Sokolov D, Yakovlev A. Opportunistic Merge <br> Element. In: 21st IEEE International Symposium on Asynchronous Circuits and Systems (ASYNC'15). 2015, Mountain View, Silicon Valley, California, USA: IEEE Computing Society Press. 

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DOI link to article:
http://dx.doi.org/10.1109/ASYNC.2015.25

Date deposited:

21/07/2015 and skills

# Opportunistic <br> Merge Element 

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## Merge Element


Purpose: merge independent requests
Example: count the total number of requests
Property: requests are never lost, $\mathbf{I}_{\mathbf{1}}+\mathbf{I}_{\mathbf{2}}=\mathbf{0}$
Requires arbitration

- between requests
- better outside the critical path


## Opportunistic Merge Element

input channel $1\left\{\begin{array}{ll}\text { A1 } \\ \text { a1 } & r \\ r 2 & a-\end{array}\right\}$ output channel
input channel $2\left\{\begin{array}{l}\text { a }\end{array}\right.$

Purpose: merge independent requests, bundling closely arriving requests together

Example: respond to an alarm (two sensors)
Property: $\max \left(I_{1}, I_{2}\right) \leq 0 \leq I_{1}+I_{2}$

## OMs in the real world



## Conceptual specification



Merge

## Conceptual specification



OM

## Conceptual specification (unrolled)



## Conceptual specification (unrolled)



## Decomposing the bundle



## OM with bundle

## Decomposing the bundle

Problem: decomposed specification cannot be synthesised due to irreducible state encoding (CSC) conflicts between $s_{1}$ and $s_{4}$, and between $\mathbf{s}_{\mathbf{2}}$ and $\mathbf{s}_{\mathbf{3}}$


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## Is this a dead end?

Decomposing the bundle $\{\mathrm{a} 1, \mathrm{a} 2\}$ is highly nontrivial:

- Output-determinacy violations
- Non-commutativity of inputs
- Irreducible CSC conflicts


## ...then a miracle occurs...

1 THINK YOU SHOULD BE MORE SPECIFIC HERE IN STEP TWO


## STG specification



## STG specification



## CSC resolution (MPSAT)



## CSC resolution (MPSAT)



## Deadlock free

## No hazards

## Synthesisable

Fast response: no metastability on the critical path


## Synthesised circuit (MPSAT)



## Simplified (hacked up) circuit



New optimisation technique: fairness-based optimisation

## Simplified (hacked up) circuit



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Scenario 1: acknowledgement a wins the arbitration

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Scenario 2: request $r 2$ wins the arbitration

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Scenario 3: sequential bundling of requests

## Simplified (hacked up) circuit



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Scenario 3: sequential bundling of requests

## Simplified (hacked up) circuit



Fair mutexes do not permit sequential bundling

## Scaling to more inputs



## Scaling to more inputs

 m1

## Scaling to more inputs



## Conclusion

- New reusable asynchronous component - surprisingly difficult for just 3 handshakes!
- Fast implementation - no metastability on critical path
- Discovered fairness-based optimisation
- Scalable
- Formally verified using Workcraft and Versify
- To be integrated into a real multiphase buck
- Challenge for asynchronous community: Design OM in a non-monolithic way
(how to design it without a miracle?)


## Thank you!

Opportunistic bundling of questions is encouraged (fairness assumption on the session chair to prevent sequential bundling) ()

