A Low-Complexity, Full-Rate, Full-Diversity 2×2 STBC with Golden Code's Coding Gain

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Abstract—This paper presents a low-ML-decoding-complexity, full-rate, full-diversity space-time block code (STBC) for a 2 transmit antenna, 2 receive antenna multiple-input multipleoutput (MIMO) system, with coding gain equal to that of the best and well known Golden code for any QAM constellation. Recently, two codes have been proposed (by Paredes, Gershman and Alkhansari and by Sezginer and Sari), which enjoy a lower decoding complexity relative to the Golden code, but have lesser coding gain. The 2×2 STBC presented in this paper has lesser decoding complexity for non-square QAM constellations, compared with that of the Golden code, while having the same decoding complexity for square QAM constellations. Compared with the Paredes-Gershman-Alkhansari and Sezginer-Sari codes, the proposed code has the same decoding complexity for nonrectangular QAM constellations. Simulation results, which compare the codeword error rate (CER) performance, are presented.

I. INTRODUCTION

Multiple-input, multiple-output(MIMO) wireless transmission systems have been intensively studied during the last decade. The Alamouti code [1] for two transmit antennas is a novel scheme for MIMO transmission, which, due to its orthogonality properties, allows a low complexity maximumlikelihood (ML) decoder. This scheme led to the generalization of STBCs from orthogonal designs [2]. Such codes allow the transmitted symbols to be decoupled from one another and single-symbol ML decoding is achieved over quasi static Rayleigh fading channels. Even though these codes achieve the maximum diversity gain for a given number of transmit and receive antennas and for any arbitrary complex constellations. unfortunately, these codes are not full - rate, where, by a full - rate code, we mean a code that transmits at a rate of $min(n_r, n_t)$ complex symbols per channel use for an n_t transmit antenna, n_r receive antenna system.

The Golden code [3] is a full-rate, full-diversity code and has a decoding complexity of the order of M^4 , for arbitrary constellations of size M. The codes in [4] and the traceorthogonal cyclotomic code in [5] also match the Golden code. With reduction in the decoding complexity being the prime objective, two new full-rate, full-diversity codes have recently been discovered: The first code was independently discovered by Hottinen, Tirkkonen and Wichman [6] and by Paredes, Gershman and Alkhansari [7], which we call the HTW-PGA code and the second, which we call the Sezginer-Sari code, was reported in [8] by Sezginer and Sari. Both these B. Sundar Rajan Dept of ECE, Indian Institute of science Bangalore 560012, India Email:bsrajan@ece.iisc.ernet.in

codes enable simplified decoding, achieving a complexity of the order of M^3 . The first code is also shown to have the non-vanishing determinant property [7]. However, these two codes have lesser coding gain compared to the Golden code. A detailed discussion of these codes has been made in [9], wherein a comparison of the codeword error rate (CER) performance reveals that the Golden code has the best performance.

In this paper, we propose a new full-rate, full-diversity STBC for 2×2 MIMO transmission, which has low decoding complexity. The contributions of this paper may be summarized (see Table I also) as follows:

- The proposed code has the same coding gain as that of the Golden code (and hence of that in [4] and the traceorthonormal cyclotomic code) for any QAM constellation (by a QAM constellation we mean any finite subset of the integer lattice) and larger coding gain than those of the HTW-PGA code and the Sezginer-Sari code.
- Compared with the Golden code and the codes in [4] and [5], the proposed code has lesser decoding complexity for all complex constellations except for square QAM constellations in which case the complexity is the same. Compared to the HTW-PGA code and the Sezginer-Sari codes, the proposed code has the same decoding complexity for all non-rectangular QAM [Fig 3] constellations.
- The proposed code has the non-vanishing determinant property for QAM constellations and hence is Diversity-Multiplexing Gain (DMG) tradeoff optimal.

The remaining content of the paper is organized as follows: In Section II, the system model and the code design criteria are reviewed along with some basic definitions. The proposed STBC is described in Section III and its non-vanishing determinant property is shown in Section IV. In Section V the ML decoding complexity of the proposed code is discussed and the scheme to decode it using sphere decoding is discussed in Section VI. In Section VII, simulation results are presented to show the performance of the proposed code as well as to compare with few other known codes. Concluding remarks constitute Section VIII.

Notations: For a complex matrix X, the matrices X^T , X^H and det[X] denote the transpose, Hermitian and determinant of X, respectively. For a complex number s, $\mathcal{R}(s)$ and $\mathcal{I}(s)$

denote the real and imaginary part of s, respectively. Also, j represents $\sqrt{-1}$ and the set of all integers, all real and complex numbers are denoted by \mathbb{Z} , \mathbb{R} and \mathbb{C} , respectively. The Frobenius norm and the trace are denoted by $||.||_F$ and tr [.] respectively. The columnwise stacking operation on X is denoted by vec(X). The Kronecker product is denoted by \otimes and I_T denotes the $T \times T$ identity matrix. Given a complex vector $\mathbf{x} = [x_1, x_2, \cdots, x_n]^T$, $\tilde{\mathbf{x}}$ is defined as

$$\tilde{\mathbf{x}} \triangleq \left[\mathcal{R}\left(x_{1}\right), \mathcal{I}\left(x_{1}\right), \cdots, \mathcal{I}\left(x_{n}\right)\right]^{T}$$

and for a complex number s, the (.) operator is defined by

$$\check{\mathbf{s}} \triangleq \left[\begin{array}{cc} \mathcal{R}\left(s\right) & -\mathcal{I}\left(s\right) \\ \mathcal{I}\left(s\right) & \mathcal{R}\left(s\right) \end{array} \right]$$

The (.) operator can be extended to a complex $n \times n$ matrix by applying it to all the entries of it.

II. CODE DESIGN CRITERIA

A finite set of complex matrices is a STBC. A $n \times n$ linear STBC is obtained starting from an $n \times n$ matrix consisting of arbitrary linear combinations of k complex variables and their conjugates, and letting the variables take values from complex constellations. The rate of such a code is $\frac{k}{n}$ complex symbols per channel use. We consider Rayleigh quasi-static flat fading MIMO channel with full channel state information (CSI) at the receiver but not at the transmitter. For 2×2 MIMO transmission, we have

$$\mathbf{Y} = \mathbf{H}\mathbf{S} + \mathbf{N} \tag{1}$$

where $\mathbf{S} \in \mathbb{C}^{2 \times 2}$ is the codeword matrix, transmitted over 2 channel uses, $\mathbf{N} \in \mathbb{C}^{2 \times 2}$ is a complex white Gaussian noise matrix with i.i.d entries, i.e., $\sim \mathcal{N}_{\mathbb{C}}(0, N_0)$ and $\mathbf{H} \in \mathbb{C}^{2 \times 2}$ is the channel matrix with the entries assumed to be i.i.d circularly symmetric Gaussian random variables $\sim \mathcal{N}_{\mathbb{C}}(0, 1)$. $\mathbf{Y} \in \mathbb{C}^{2 \times 2}$ is the received matrix.

Definition 1: (Code rate) If there are k independent information symbols in the codeword which are transmitted over T channel uses, then, for an $n_t \times n_r$ MIMO system, the code rate is defined as k/T symbols per channel use. If $k = n_{min}T$, where $n_{min} = min(n_t, n_r)$, then the STBC is said to have full rate.

Considering ML decoding, the decoding metric that is to be minimized over all possible values of codewords S is given by

$$\mathbf{M}\left(\mathbf{S}\right) = \|\mathbf{Y} - \mathbf{H}\mathbf{S}\|_{F}^{2} \tag{2}$$

Definition 2: (**Decoding complexity**) The ML decoding complexity is given by the minimum number of symbols that need to be jointly decoded in minimizing the decoding metric. This can never be greater than k, in which case, the decoding complexity is said to be of the order of M^k . If the decoding complexity is lesser than M^k , the code is said to admit simplified decoding.

Definition 3: (Generator matrix) For any STBC S that encodes k information symbols, the *generator* matrix G is

defined by the following equation

$$\operatorname{vec}\left(\mathbf{S}\right) = \mathbf{G}\tilde{\mathbf{s}}.\tag{3}$$

where $\mathbf{s} \triangleq [s_1, s_2, \cdots, s_k]^T$ is the information symbol vector The code design criteria [12] are: (i) *Rank criterion*— To

achieve maximum diversity, the codeword difference matrix $(\mathbf{X} - \hat{\mathbf{X}})$ must be full rank for all possible pairs of codewords and the diversity gain is given by $n_t n_r$, (ii) Determinant criterion– For a full ranked STBC, the minimum determinant δ_{min} , defined as

$$\delta_{min} \triangleq \min_{\mathbf{X} \neq \hat{\mathbf{X}}} det \left[\left(\mathbf{X} - \hat{\mathbf{X}} \right) \left(\mathbf{X} - \hat{\mathbf{X}} \right)^H \right]$$
(4)

should be maximized. The coding gain is given by $(\delta_{min})^{1/n_t}$, with n_t being the number of transmit antennas.

For the 2×2 MIMO system, the target is to design a code that is full-rate, i.e transmits 2 complex symbols per channel use, has full-diversity, maximum coding gain and allows low ML decoding complexity.

III. THE PROPOSED STBC

In this section, we present our STBC for 2×2 MIMO system. The design is based on the class of codes called co-ordinate interleaved orthogonal designs (CIODs), which was studied in [11] in connection with the general class of single-symbol decodable codes and, specifically for 2 transmit antennas, is as follows.

Definition 4: The CIOD for 2 transmit antennas [11] is

$$\mathbf{X}(s_1, s_2) = \begin{bmatrix} s_{1I} + js_{2Q} & 0\\ 0 & s_{2I} + js_{1Q} \end{bmatrix}$$
(5)

where $s_i \in \mathbb{C}$, i = 1, 2 are the information symbols and s_{iI} and s_{iQ} are the in-phase (real) and quadrature-phase (imaginary) components of s_i , respectively. Notice that in order to make the above STBC full rank, the signal constellation \mathcal{A} from which the symbols s_i are chosen should be such that the real part (imaginary part, resp.) of any signal point in \mathcal{A} is not equal to the real part (imaginary part, resp.) of any other signal point in \mathcal{A} [11]. So if QAM constellations are chosen, they have to be rotated. The optimum angle of rotation has been found in [11] to be $\frac{1}{2}tan^{-1}2$ degrees and this maximizes the diversity and coding gain. We denote this angle by θ_g .

The proposed 2×2 STBC S is given by

$$\mathbf{S}(x_1, x_2, x_3, x_4) = \mathbf{X}(s_1, s_2) + e^{j\theta} \mathbf{X}(s_3, s_4) \mathbf{P}$$
(6)

where

- The four symbols s_1, s_2, s_3 and $s_4 \in \mathcal{A}$, where \mathcal{A} is a θ_g degrees rotated version of a regular QAM signal set, denoted by \mathcal{A}_q which is a finite subset of the integer lattice, and $x_1, x_2, x_3, x_4 \in \mathcal{A}_q$. To be precise, $s_i = e^{\theta_g} x_i$, i = 1, 2, 3, 4.
- **P** is a permutation matrix designed to make the STBC full rate and is given by $\mathbf{P} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$.

The choice of θ in the above expression should be such that the diversity and coding gain are maximized. A computer search was done for θ in the range [0, π/2]. The optimum value of θ was found out to be π/4.

Explicitly, our code matrix is

 $\mathbf{S}(x_1, x_2, x_3, x_4) =$

$$\begin{bmatrix} s_{1I} + js_{2Q} & e^{j\pi/4}(s_{3I} + js_{4Q}) \\ e^{j\pi/4}(s_{4I} + js_{3Q}) & s_{2I} + js_{1Q} \end{bmatrix}$$
(7)

The minimum determinant for our code when the symbols are chosen from QAM constellations is 3.2, the same as that of the Golden code, which will be proved in the next section.

The generator matrix for our STBC, corresponding to the symbols s_i , is as follows:

It is easy to see that this generator matrix is orthonormal. In [5], it was shown that a necessary and sufficient condition for an STBC to be *Information lossless* is that its generator matrix should be unitary. Hence, our STBC has the *Information losslessness* property.

IV. NVD PROPERTY AND THE DMG OPTIMALITY

In this section it is shown that the proposed code has the non-vanishing determinant (NVD) property [3], which in conjunction with full-rateness means that our code is DMG tradeoff optimal [10].

The determinant of the codeword matrix ${\bf S}$ can be written as

$$det(\mathbf{S}) = (s_{1I} + js_{2Q})(s_{2I} + s_{1Q}) - j[(s_{3I} + js_{4Q})(s_{4I} + s_{3Q})].$$

Using $s_{iI} = (s_i + s_i^*)/2$ and $js_{iQ} = (s_i - s_i^*)/2$ in the equation above, we get,

$$4det(\mathbf{S}) = (s_1 + s_1^* + s_2 - s_2^*)(s_2 + s_2^* + s_1 - s_1^*) - j[(s_3 + s_3^* + s_4 - s_4^*)(s_4 + s_4^* + s_3 - s_3^*)] = ((s_1 + s_2) + (s_1 - s_2)^*)((s_1 + s_2) - (s_1 - s_2)^*) - j[((s_3 + s_4) + (s_3 - s_4)^*))((s_3 + s_4) - (s_3 - s_4)^*)].$$

Since $s_i = e^{j\theta_g} x_i$, i = 1, 2, 3, 4, with $s_i \in \mathcal{A}$, $x_i \in \mathcal{A}_q$, a subset of $\mathbb{Z}[i]$, defining $A \triangleq (x_1 + x_2)$, $B \triangleq (x_1 - x_2)^*$, $C \triangleq (x_3 + x_4)$ and $D \triangleq (x_3 - x_4)^*$, with A, B, C and $D \in \mathbb{Z}[i]$, we get

$$4Det(\mathbf{S}) = (e^{j\theta_g}A + e^{-j\theta_g}B)(e^{j\theta_g}A - e^{-j\theta_g}B) - j[(e^{j\theta_g}C + e^{-j\theta_g}D)(e^{j\theta_g}C - e^{-j\theta_g}D)] = e^{j2\theta_g}A^2 - e^{-j2\theta_g}B^2 - j[e^{j2\theta_g}C^2 - e^{-j2\theta_g}D^2].$$

Since $e^{j2\theta_g} = \cos(2\theta_g) + \sin(2\theta_g) = (1+2j)/\sqrt{5}$, we get $4\sqrt{5}Det(\mathbf{S}) = (1+2j)(A^2 - jC^2) - (1-2j)(B^2 - jD^2)$. (9) For the determinant of S to be 0, we must have

$$(1+2j)(A^2 - jC^2) = (1-2j)(B^2 - jD^2)$$

$$\Rightarrow (1+2j)^2(A^2 - jC^2) = 5(B^2 - jD^2).$$

The above can be written as

=

$$A_1^2 - jC_1^2 = 5(B^2 - jD^2)$$
⁽¹⁰⁾

where $A_1 = (1+2j)A, C_1 = (1+2j)C$ and clearly $A_1, C_1 \in \mathbb{Z}[i]$. It has been shown in [4] that (10) holds only when $A_1 = B = C_1 = D = 0$, i.e., only when $x_1 = x_2 = x_3 = x_4 = 0$. This means that the determinant of the codeword difference matrix is 0 only when the codeword difference matrix is itself 0. So, for any distinct pair of codewords, the codeword difference matrix is always full rank for any constellation which is a subset of $\mathbb{Z}[i]$. Also, the minimum value of the modulus of R.H.S of (9) can be seen to be 4. So, $|Det(\mathbf{S})| \ge 1/\sqrt{5}$. In particular, when the constellation chosen is the standard OAM constellation, the difference between any two signal points is a multiple of 2. Hence, for such constellations, $|Det(\mathbf{S}\cdot\mathbf{S}')| \ge 4/\sqrt{5}$, where S and S' are distinct codewords. The minimum determinant is consequently 16/5. This means that the proposed codes has the non-vanishing determinant (NVD) property [3]. In [10], it was shown that full-rate codes which satisfy the non-vanishing determinant property achieve the optimal DMG tradeoff. So, our proposed STBC is DMG tradeoff optimal.

V. DECODING COMPLEXITY

The decoding complexity of the proposed code is of the order of M^3 . This is due to the fact that conditionally given the symbols x_3 and x_4 , the symbols x_1 and x_2 can be decoded independently. This can be proved as follows. Writing the STBC in terms of its weight matrices/dispersion matrices A_i , $i = 1, 2, \dots, 8$, [11], we have

$$\mathbf{S} = \sum_{m=1}^{4} \underbrace{x_{mI}A_{2m-1} + x_{mQ}A_{2m}}_{\mathbf{T}_m} = S_1 + S_2$$

where

$$S_1 = \sum_{m=1}^{2} x_{mI} A_{2m-1} + x_{mQ} A_{2m}$$

and

$$S_2 = \sum_{m=3}^{4} x_{mI} A_{2m-1} + x_{mQ} A_{2m}.$$

For our code, we have

$$A_{1} = \begin{bmatrix} \cos\theta_{g} & 0\\ 0 & j\sin\theta_{g} \end{bmatrix}; \quad A_{2} = \begin{bmatrix} -\sin\theta_{g} & 0\\ 0 & j\cos\theta_{g} \end{bmatrix}$$
$$A_{3} = \begin{bmatrix} j\sin\theta_{g} & 0\\ 0 & \cos\theta_{g} \end{bmatrix}; \quad A_{4} = \begin{bmatrix} j\cos\theta_{g} & 0\\ 0 & -\sin\theta_{g} \end{bmatrix}$$
$$A_{5} = e^{j\pi/4} \begin{bmatrix} 0 & \cos\theta_{g}\\ j\sin\theta_{g} & 0 \end{bmatrix}$$
$$A_{6} = e^{j\pi/4} \begin{bmatrix} 0 & -\sin\theta_{g}\\ j\cos\theta_{g} & 0 \end{bmatrix}$$

$$A_{7} = e^{j\pi/4} \begin{bmatrix} 0 & jsin\theta_{g} \\ cos\theta_{g} & 0 \end{bmatrix}$$
$$A_{8} = e^{j\pi/4} \begin{bmatrix} 0 & jcos\theta_{g} \\ -sin\theta_{g} & 0 \end{bmatrix}.$$

The ML decoding metric in (2) can be written as

$$M(S) = tr \left[(Y - HS) (Y - HS)^{H} \right]$$

= $tr \left[(Y - HS_{1} - HS_{2}) (Y - HS_{1} - HS_{2})^{H} \right]$
= $tr \left[(Y - HS_{1}) (Y - HS_{1})^{H} \right]$
 $- tr \left[HS_{2} (Y - HS_{1})^{H} \right]$
 $- tr \left[(Y - HS_{1}) (HS_{2})^{H} \right]$
 $+ tr \left[HS_{2} (HS_{2})^{H} \right].$

It can be verified that the following hold true for $l, m \in [1, 4]$

$$A_m A_l^H + A_l A_m^H = 0 \left\{ \begin{array}{ll} \forall l \neq m, m+1, & \text{if m is odd} \\ \forall l \neq m, m-1, & \text{if m is even.} \end{array} \right.$$

From [11], we obtain

$$tr\left[(Y - HS_1)(Y - HS_1)^H\right] = \sum_{m=1}^2 ||Y - HT_m||_F^2 - tr\left(YY^H\right)$$

and hence,

$$M(S) = \sum_{m=1}^{2} \|Y - HT_m\|_F^2 - tr(YY^H) + tr[HS_2(HS_1)^H] + tr[HS_1(HS_2)^H] - tr[HS_2Y^H] - tr[Y(HS_2)^H] + tr[HS_2(HS_2)^H] = \sum_{m=1}^{2} \|Y - HT_m\|_F^2 + \sum_{m=1}^{2} tr[HS_2(HT_m)^H] + \sum_{m=1}^{2} tr[HT_m(HS_2)^H] + \|Y - HS_2\|_F^2 - 2tr(YY^H).$$

Hence, when S_2 is given, i.e., symbols x_3 and x_4 are given, the ML metric can be decomposed as

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$$M(S) = \sum_{m=1}^{2} M(x_m) + M_c$$
(11)

with $M_c = ||Y - HS_2||_F^2 - 2tr(YY^H)$ and $M(x_m)$ being a function of symbol x_m alone. Thus decoding can be done as follows: choose the pair (x_3, x_4) and then, in parallel, decode x_1 and x_2 so as to minimize the ML decoding metric. With this approach, there are $2M^3$ values of the decoding metric that need to be computed in the worst case. So, the decoding complexity is of the order of M^3 .

VI. SIMPLIFIED DECODING USING SPHERE DECODER

In this section, it is shows that sphere decoding can be used to achieve the decoding complexity of M^3 . It can be shown that (1) can be written as

$$\widetilde{vec}(\mathbf{Y}) = \mathbf{H}_{eq}\tilde{\mathbf{s}} + \widetilde{vec}(\mathbf{N})$$
(12)

where $\mathbf{H}_{eq} \in \mathbb{R}^{8 \times 8}$ is given by

$$\mathbf{H}_{eq} = \left(\mathbf{I}_2 \otimes \check{\mathbf{H}}\right) \mathbf{G} \tag{13}$$

with $\mathbf{G} \in \mathbb{R}^{8 \times 8}$ being the generator matrix as in (8) and

$$\tilde{\mathbf{s}} \triangleq [\mathcal{R}(s_1), \mathcal{I}(s_1), \cdots, \mathcal{R}(s_4), \mathcal{I}(s_4)]^T$$

with $s_i, i = 1, \dots, 4$ drawn from \mathcal{A} , which is a rotation of the regular QAM constellation \mathcal{A}_q . Let

$$\mathbf{x}_q \triangleq [x_1, x_2, x_3, x_4]^T$$

Then,

$$\tilde{\mathbf{s}} = \mathbf{F}\tilde{\mathbf{x}}_q.$$

where $\mathbf{F} \in \mathbb{R}^{8 \times 8}$ is $diag[\mathbf{J}, \mathbf{J}, \mathbf{J}, \mathbf{J}]$ with \mathbf{J} being a rotation matrix and is defined as follows

$$\mathbf{J} \triangleq \begin{bmatrix} \cos(\theta_g) & -\sin(\theta_g) \\ \sin(\theta_g) & \cos(\theta_g) \end{bmatrix}.$$

So, (12) can be written as

$$\widetilde{vec}(\mathbf{Y}) = \mathbf{H}'_{eq} \tilde{\mathbf{x}_q} + \widetilde{vec}(\mathbf{N})$$
(14)

where $\mathbf{H}'_{eq} = \mathbf{H}_{eq}\mathbf{F}$. Using this equivalent model, the ML decoding metric can be written as

$$\mathbf{M}\left(\tilde{\mathbf{x}_{q}}\right) = \|\widetilde{vec}\left(\mathbf{Y}\right) - \mathbf{H}_{eq}'\tilde{\mathbf{x}_{q}}\|^{2}$$
(15)

On obtaining the QR decomposition of \mathbf{H}'_{eq} , we get $\mathbf{H}'_{eq} = \mathbf{QR}$, where $\mathbf{Q} \in \mathbb{R}^{8 \times 8}$ is an orthonormal matrix and $\mathbf{R} \in \mathbb{R}^{8 \times 8}$ is an upper triangular matrix. The ML decoding metric now can be written as

$$\mathbf{M}(\tilde{\mathbf{x}_q}) = \|\mathbf{Q}^T \widetilde{\mathbf{vec}}(\mathbf{Y}) - \mathbf{R}\tilde{\mathbf{x}_q}\|^2$$
(16)

If $\mathbf{H}'_{eq} \triangleq [\mathbf{h}_1 \ \mathbf{h}_2 \cdots \mathbf{h}_8]$, where $\mathbf{h}_i, i = 1, 2, \cdots, 8$ are column vectors, then \mathbf{Q} and \mathbf{R} have the general form obtained by Gram - Schmidt process as shown below

$$\mathbf{Q} = [\mathbf{q}_1 \ \mathbf{q}_2 \ \mathbf{q}_3 \cdots \mathbf{q}_8]$$

where $\mathbf{q}_i, i = 1, 2, \cdots, 8$ are column vectors, and

$$\mathbf{R} = \begin{bmatrix} \|\mathbf{r}_1\| & \langle \mathbf{h}_2, \mathbf{q}_1 \rangle & \langle \mathbf{h}_3, \mathbf{q}_1 \rangle & \dots & \langle \mathbf{h}_8, \mathbf{q}_1 \rangle \\ 0 & \|\mathbf{r}_2\| & \langle \mathbf{h}_3, \mathbf{q}_2 \rangle & \dots & \langle \mathbf{h}_8, \mathbf{q}_2 \rangle \\ 0 & 0 & \|\mathbf{r}_3\| & \dots & \langle \mathbf{h}_8, \mathbf{q}_3 \rangle \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & \|\mathbf{r}_8\| \end{bmatrix}$$

where $\mathbf{r}_1 = \mathbf{h}_1$, $\mathbf{q}_1 = \frac{\mathbf{r}_1}{\|\mathbf{r}_1\|}$, $\mathbf{r}_i = \mathbf{h}_i - \sum_{j=1}^{i-1} \langle \mathbf{h}_i, \mathbf{q}_j \rangle \mathbf{q}_j$, $\mathbf{q}_i = \frac{\mathbf{r}_i}{\|\mathbf{r}_i\|}$, $i = 2, \cdots, 8$.

It can be shown by direct computation that \mathbf{R} has the following structure

- ~	~	\cap	0	~	~	~	~ -
u	u	0	0	u	u	u	a
0	a	0	0	a	a	a	a
0	0	a	a	a	a	a	a
0	0	0	a	a	a	a	a
0	0	0	0	a	a	a	a
0	0	0	0	0	a	a	a
0	0	0	0	0	0	a	a
0	0	0	0	0	0	0	a
-							

where a stands for a possibly non-zero entry.

The structure of the matrix **R** allows us to perform a 4 dimensional real sphere decoding (SD) [13] to find the partial vector $[\mathcal{R}(x_3), \mathcal{I}(x_3), \mathcal{R}(x_4), \mathcal{I}(x_4)]^T$ and hence obtain the symbols x_3 and x_4 . Having found these, x_1 and x_2 can be decoded independently. Observe that the real and imaginary parts of symbol x_1 are entangled with one another because of constellation rotation but are independent of the real and imaginary parts of x_2 when x_3 and x_4 are conditionally given. found the partial Having vector $[\mathcal{R}(x_3), \mathcal{I}(x_3), \mathcal{R}(x_4), \mathcal{I}(x_4)]^T$, we proceed to find the rest of the symbols as follows. We do two parallel 2 dimensional real search to decode the symbols x_1 and x_2 . So, overall, the worst case decoding complexity of the proposed STBC is $2M^3$. This is due to the fact that

- 1) A 4 dimensional real SD requires M^2 metric computations in the worst possible case.
- 2) Two parallel 2 dimensional real SD require 2M metric computations in the worst case.

This decoding complexity is the same as that achieved by the HTW-PGA code and the Sezginer-Sari code.

Though it has not been mentioned anywhere to the best of our knowledge, the ML decoding complexity of the Golden code, Dayal-Varanasi code and the trace-orthogonal cyclotomic code is also $2M^3$ for square QAM constellations. This follows from the structure of the **R** matrices for these codes which are counterparts of the one in (17). The **R** matrices of these codes are similar in structure and as shown below:

$$\mathbf{R} = \begin{bmatrix} a & 0 & a & 0 & a & a & a & a \\ 0 & a & 0 & a & a & a & a & a \\ 0 & 0 & a & 0 & a & a & a & a \\ 0 & 0 & 0 & a & a & a & a & a \\ 0 & 0 & 0 & 0 & a & a & a & a \\ 0 & 0 & 0 & 0 & 0 & a & a & a \\ 0 & 0 & 0 & 0 & 0 & 0 & a & a \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & a \end{bmatrix}$$

Table I presents the comparison of the known full-rate, fulldiversity 2×2 codes in terms of their ML decoding complexity and the coding gain.

VII. SIMULATION RESULTS

Fig 1 shows the codeword error performance plots for the Golden code, the proposed STBC and the HTW-PGA code for the 4-QAM constellation. The performance of the proposed



Fig. 1. CER PERFORMANCE FOR 4-QAM

code is the same as that of the Golden code. The HTW-PGA code performs slightly worse due to its lower coding gain. Fig 2, which is a plot of the CER performance for 16-QAM, also highlights these aspects. Table I gives a comparison between the well known full-rate, full-diversity codes for 2×2 MIMO.

VIII. CONCLUDING REMARKS

In this paper, we have presented a full-rate STBC for 2×2 MIMO systems which matches the best known codes for such systems in terms of error performance, while at the same time, enjoys simplified-decoding complexity that the codes presented in [7] and [8] do. Recently, a Rate-1 STBC, based on scaled repetition and rotation of the Alamouti code, was proposed [15]. This code was shown to have a hard-decision performance which was only slightly worse than that of the Golden code for a spectral efficiency of 4b/s/Hz, but the complexity was significantly lower.

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	Min det	Ν	ML Decoding complexity			
Code	for QAM	Square QAM	Rectangular QAM	Non-rectangular		
			$M = M_1 \times M_2$	QAM		
Yo-Wornell[14]	0.8000	$2M^{3}$	$M^2(M_1^2 + M_2^2)$	M^4		
Dayal-Varanasi code[4]	3.2000	$2M^{3}$	$M^2(M_1^2 + M_2^2)$	M^4		
Golden code [3]	3.2000	$2M^{3}$	$M^2(M_1^2 + M_2^2)$	M^4		
Trace-orthonormal cyclotomic code [5]	3.2000	$2M^{3}$	$M^2(M_1^2 + M_2^2)$	M^4		
HTW-PGA code [7]	2.2857	$4M^2\sqrt{M}$	$2M^2(M_1 + M_2)$	$2M^3$		
Sezginer-Sari code [8]	2.0000	$4M^2\sqrt{M}$	$2M^2(M_1 + M_2)$	$2M^3$		
The proposed code	3.2000	$2M^3$	$2M^3$	$2M^3$		

TABLE I

COMPARISION BETWEEN THE MINIMUM DETERMINANT AND DECODING COMPLEXITY OF SOME WELL KNOWN FULL-RATE 2×2 STBCs



Fig. 2. CER PERFORMANCE FOR 16-QAM

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Fig. 3. AN EXAMPLE OF A NON-RECTANGULAR QAM CONSTELLATION

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