# How to Solve Deadlock Situations within the Plan-Merging Paradigm for Multi-robot Cooperation<sup>\*</sup>

S. Qutub, R. Alami, F. Ingrand LAAS-CNRS 7, Avenue du Colonel Roche, 31077 Toulouse CEDEX 04 E-mail: {sam, rachid, felix}@laas.fr

### Abstract

Our motivation in proposing the Plan-Merging Paradigm as a cooperation scheme was to allow an efficient distribution of the decisions for a better reactivity to contingencies for multi-robot applications with loosely coupled tasks.

The paradigm proved to be quite efficient because it exploits the fact that most conflicts can be solved locally and because it allows a finer overlapping between plan refinement, plan coordination and execution.

However, we would like to have a scheme which distributes as much as possible the decision processes for planning and coordination while maintaining two key features:

- the coherence of the global system and the ability to detect the situations where it is not applicable
- a localized management of the planning and coordination processes with, in particularly intricate situations, a progressive transition to more global schemes which may "degrade" to a unique and centralized planning activity.

We develop in this paper the ingredients which guarantee such features as well as their consequences in terms of requirements for the task planners involved.

#### 1 Introduction

We have already presented and discussed the Plan-Merging Paradigm (PMP), a generic scheme for multirobot cooperation [3, 4, 1]. It is based on an incremental and distributed plan-merging process.

We have applied the PMP to multi-robot applications [2] with loosely coupled tasks, where each robot has a local view and a partial knowledge of the other robots activities. We showed that the PMP is quite efficient because it exploits the fact that most conflicts can be solved locally and because it allows a finer overlapping between plan refinement, plan coordination and execution.

We have also discussed the key features of this distributed cooperative scheme related to the coherence of the global system and its ability to detect situations where it is not applicable i.e. situations where it is necessary to take into account a conjunction of goals. We call such situations "Planning deadlock situations".

We present here a set of extended operators and complementary mechanisms which permit a localized management of the planning and coordination processes as well as a progressive transition to more global schemes which may even "degrade" to a unique and centralized planning activity. The result is a generic multi-robot cooperative scheme which is well suited for loosely coupled tasks but is able to treat any conflicting situation.

In section 2, we present a short discussion of related work. Section 3 describes briefly the *PMP* and explains how tasks dependencies are detected and propagated (through communication) in order to maintain the global system coherence and to detect planning deadlock situations. Section 4 introduces the notion of *Local Multi-robot Planning* that allows us to modify dynamically the distributed nature of the system in order to deal with the deadlock situations. Section 5 describes an implemented system which illustrates our approach and describes the behavior of a set of mobile robots in a very constrained environment.

### 2 Related Work

We limit our discussion here to multi-robot issues which involve the simultaneous operation of several autonomous agents, each one seeking to achieve its own task or goal. The conflicts proceed from the fact that the robots intend to use common resources simultaneously (narrow passages, crossings, devices, etc).

Many approaches have been proposed to deal with this problem especially *centralized* approaches where a central system determines a set of non-conflicting plans that solves the conflicts[11]. These approaches suffer from deficiencies in realistic applications when the number of robots becomes important. Other approaches are based on predefined *traffic rules*, which are only applicable to "route networks" modeled environment. In such cases, it is very difficult to find a set of free-deadlock rules for all the possible situations [7, 9]. Some *reactive systems* have been proposed

<sup>\*</sup>This paper has been published in the proceedings of IEEE IROS 97, Grenoble, France.

where the robots actions are the direct consequences of the information collected by the robots sensors [6] or through communication [8]. While the results of such approaches may be inefficient due to the local decision making based on sensory information, the main limitation here is that there is no guarantee of a global coherence of the system. In [5] the authors propose an idea mixing sensory information with a world discrete model to solve conflicts for a small number of robots. Finally, a *master-slave* approaches have been proposed where a robot becomes the master of the blocked robots during the conflict, resolves the conflicts and distributes the solutions [12]. In this paper, we refine this master-slave approach by mixing it with the *Plan Merging Paradigm* to generate, during a deadlock situation, a set of plans that will be validated in the global context (through a Plan-Merging Operation) before the execution phase.

## 3 The Plan-Merging Paradigm

Due to space limitations, we give a short presentation of the Plan-Merging Paradigm and we insist only on the situations where it does not apply. The interested reader may refer to previous papers [3, 4, 1] for a more detailed presentation of this cooperation scheme and its applications.

Let us assume that we have a set of autonomous robots and a central station which, from time to time, sends goals to robots individually. Whenever a robot  $R_i$  receives a new goal  $G_i^j$ , it elaborates an *Individual Plan*  $(IP_i^j)$  which achieves it. Each robot processes sequentially the received goals. Doing so, it incrementally appends new actions to its current plan.

However, before executing any plan step, a robot must ensure that it is valid in the multi-robot context, i.e. that there exists no other plan of another robot which may conflict with it. We call this operation *Plan Merging Operation* (PMO) and the resulting plan a *Coordinated plan* (i.e. plan valid in the current multi-robot context). Such a *Coordinated Plan* ( $CP_i$ ) consists of a sequence of actions and *execution events* to be signaled to other robots as well as *execution events* that are planned to be signaled by the other robots. Such *execution events* correspond to temporal constraints between actions involved in different coordinated plans.

At any moment, the temporal constraints between all the actions included in the union of all the coordinated plans  $(GP = \bigcup_k CP_k)$  must constitute a *di*rected acyclic graph [3, 4]. GP is a snapshot knowledge of the current global situation and its already planned evolution.

#### 3.1 The PMO and its results

When  $R_i$  receives its *j*-th goal  $G_i^j$ , it elaborates a plan  $IP_i^j$  which achieves it; then it performs a PMO in a *critical section*: it collects the coordinated plans  $CP_k$  of the robots which may interfere with  $IP_i^j$ , and builds

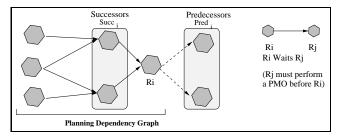


Figure 1:  $R_i$  Planning Dependency Graph  $PDG_i$  and  $Pred_i$ .

their union  $GP = \bigcup_k CP_k$ . The insertion of  $IP_i^j$  in the global plan GP, if it succeeds, adds temporal order constraints to actions in  $IP_i^j$  and transforms it into a coordinated plan  $CP_i$ . The out-coming  $CP_i$  is feasible in the current context, and does not introduce any cycle in the resulting GP.

The PMO is protected by a critical section in order to prevent other robots to perform simultaneously a modification of GP.

However a PMO may fail because the final state of at least another robot (as specified in GP) forbids  $R_i$ to insert its own plan. Let us call  $Pred_i$  (predecessors) the set of all such robots. In this case,  $R_i$  defers its PMO and waits until one of the robots in  $Pred_i$ has performed a new successful PMO which may possibly change the states preventing  $R_i$  to insert its plan. Hence, we introduce temporal order relations between robots plan-merging activities.

In addition to execution events — i.e. events elaborated by the PMO and which allow agents to synchronize their execution —, we define planning events i.e. events which occur whenever a robot performs a new successful PMO. The planning events can also be awaited for. The temporal relations between robots plan-merging activities are maintained by each robot in an additional data structure called Planning Dependency Graph  $PDG_i$  (Figure 1).

The *Planning Dependency Graph* serves to manage *PMOs* order (when necessary) as well as to detect and prevent any robot to enter a *waiting cycle*, where it would wait for itself by transitivity. We call such a situation a "Planning Deadlock Problem". The detection of deadlocks at the planning/coordinating phase permits to anticipate and avoid deadlocks during the execution phase where "backtracks" are not always possible or induce inefficient maneuvers.

The "planning deadlock problem" emphasizes the fact that the PMO is unable to take into account a set of goals, sent to different robots, with strong interdependencies. This limitation leads us to elaborate an extension to the PMP that allows the use of planning from a distributed to a more centralized scheme and from a local to a more global resolution.

#### 3.2 Dependency Graph Construction

This section focuses on the incremental construction of the Planning Dependency Graph  $PDG_i$  and its constraints propagation mechanism.

Each robot  $R_i$  maintains a list  $Pred_i$  and a graph  $PDG_i$ .  $Pred_i$  is the list of all of robots that block its plan-merging activity.  $PDG_i$  specifies <sup>1</sup> all the robots that depends on  $R_i$ , directly or by transitivity, for their plan-merging activities.

We call  $Succ_i$  (successors) the set of robots that are directly waiting for a *planning event* from  $R_i$  (Figure 1).

 $PDG_i$  is maintained through the following procedure:

 $\diamond$  When a robot  $R_i$  starts a *PMO*, *Pred<sub>i</sub>* is set to the empty list. After the *PMO*:

- if the PMO has succeeded:  $R_i$  signals a *planning* event to all robots in  $Succ_i$  and clears its current graph  $PDG_i$ .
- if the *PMO* has failed:  $R_i$  determines *Pred<sub>i</sub>* and checks if it induces planning dependencies which produce a cycle in *PDG<sub>i</sub>*:
  - in such a case, a *deadlock situation* is detected which means that the given goals are interdependent, they cannot be treated simply by insertion, but need to be handled by a planner that takes into account the *conjunction* of goals of all the robots involved in the cycle (see §4).
  - If the newly established planning dependencies do not introduce any cycle in  $PDG_i$ ,  $R_i$ transmits  $PDG_i$  to all robots in  $Pred_i$ .

 $\diamond$  When a robot  $R_k$  receives  $PDG_i$  from a robot  $R_i$ ,  $R_k$  adds it to its own Dependency Graph  $PDG_k$  and propagates this information to all robots in  $Pred_k$ .  $R_k$  is sure that the received  $PDG_i$  can be added to  $PDG_k$  without creating any cycle<sup>2</sup>.

### 4 Deadlock Resolution Strategy

The deadlock resolution strategy that we present is based on cooperative (not competitive) robots behavior. We assume that all robots are equipped with a multi-robot planner which can be used, when necessary, for an arbitrary number of robots.

#### 4.1 General presentation

Let us call  $DL_i$  the set of robots involved in a cycle detected by  $R_i$ . When detecting such a cycle,  $R_i$  has the necessary information in  $PDG_i$  to elaborate and validate a plan for all the robots in  $DL_i$ . Note that the blocked robots are unable to add any new executable action to their current coordinated plans  $CP_k$ . Therefore, if nothing is done, they will come to a complete stop when their plans  $CP_k$  have completed.

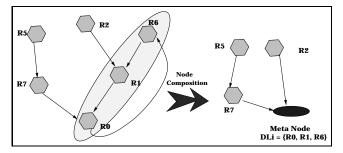


Figure 2: Node Composition in  $R_0$  Dependency Graph  $(PDG_0)$ .

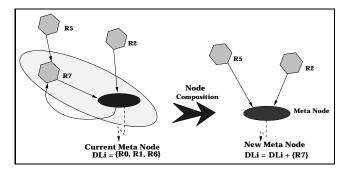


Figure 3: Forming a new Meta-Node in  $PDG_0$  by unifying the current Meta-Node and the newly detected deadlock.

To solve the deadlock, the robot  $R_i$  becomes (temporarily) the local coordinator (noted  $R_i^{LC}$ ) for all robots in  $DL_i$ . To do so, it makes use of its *Local Multi-robot Planner* that will take explicitly, in one planning operation, the conjunction of goals of all robots in  $DL_i$ . This fact will be represented in its Dependency Graph  $PDG_i^{LC}$  as a *Meta-Node* (Figure 2) which includes all robots in  $DL_i$ .

 $R_i$  becomes a local coordinator  $R_i^{LC}$  whose responsibility is to:

- 1. Find a multi-robot solution  $(Sol_i^{LC})$ , if it exists, to the conjunction of goals. This solution is represented by a *lattice* whose nodes are high level actions to be performed to break the cycle and whose arcs are "deadlock synchronization events" between these actions.
- 2. Try to insert  $Sol_i^{LC}$  in the current multi-robot context, i.e. the set of current plans  $CP_k$  of the robots which are not involved in  $DL_i$ : 2.1. If the insertion succeeds,  $R_i^{LC}$  sends to each robot in  $DL_i$  its plan and waits for an acknowledgment. If all the blocked robots accept these plans, the coordinator  $R_i^{LC}$  gives the permission to start the execution. The deadlock cycle is broken. Each robot in  $DL_i$  recovers its "planning and plan-merging" autonomy.

2.2. If the insertion fails, this means that the final state of at least one robot (not included in

 $<sup>^1\</sup>mathrm{A}$  node in  $PDG_i$  represents a set of robots, their current states and their goals.

<sup>&</sup>lt;sup>2</sup>If such cycle existed,  $R_i$  would have discovered it.

 $DL_i$ ) forbids  $R_i^{LC}$  to insert  $Sol_i^{LC}$ .  $R_i^{LC}$  determines the set of such robots  $Pred_i^{LC}$ . It then verifies that  $Pred_i^{LC}$  does not create a cycle when inserted  $PDG_i^{LC}$ :

- 2.2.1 if no cycle is created,  $R_i^{LC}$  defers its PMO, transmits  $PDG_i^{LC}$  to all robots in  $Pred_i^{LC}$ and waits until one of them has performed a new successful PMO;
- 2.2.2 if a new deadlock  $DL_i^{LC}$  is detected,  $R_i^{LC}$  generates a new *Meta-Node* containing the union of  $DL_i$  and  $DL_i^{LC}$ . It then restarts the same process (Figure 3), acting as a co-ordinator of a greater set of robots.

#### 4.2 Deadlock Automata

We describe here the finite state automata (Figure 4, 5) that defines the behavior of the robots during a deadlock situation.

Figure 4 describes the behavior of a robot<sup>3</sup>  $R_i$ when it detects a deadlock  $DL_i$ .  $R_i$  sends a message  $(Deadlock-give-info R_i)$  to all robots in  $DL_i$  (state 0) and waits for replies. If it receives $(Cycle(DL_k))$  from  $R_k \in DL_i$  (state 3), this means that  $R_k$  is the coordinator of another deadlock  $DL_k$ ; a new Meta-node is created containing the union of  $DL_i \leftarrow DL_i \cup DL_k$ . When all the expected replies are received,  $R_i$  becomes the local coordinator  $R_i^{LC}$  of all robots in  $DL_i$ . It invokes its Local Multi-robot Planner in order to find a plan  $Sol_i^{LC}$  for the conjunction of goals of  $DL_i$  (state 2); then it distributes  $Sol_i^{LC}$  to the robots in  $DL_i$  and waits for an acknowledgment (state 4).

Figure 5 describes the behavior of a  $R_j$ , involved in a cycle  $DL_i$  when it receives a message (*Deadlock-give-info*  $R_i$ ) from a coordinator  $R_i$ . There are three possible cases:

- 1.  $R_j$  is the coordinator of another deadlock  $DL_j$ (state 4) : it transmits the message  $(Cycle(DL_j))$ to  $R_i$ ; This message contains all the necessary information concerning  $DL_j$ .  $DL_j$  will be merged with  $DL_i$  (see also state 3 in Figure 4).
- 2.  $R_j$  participates in another cycle  $DL_k$  whose coordinator is  $R_k$   $(R_j \neq R_k)$ (state 3) : it transmits the message (*Deadlock-give-info*  $R_i$ ) to  $R_k$ .
- 3.  $R_j$  does not participate to any cycle (state 2) : it sends its current state and goal to  $R_i$  and waits for a plan from it.

### 4.3 Discussion

The problem discussed in this paper is a typical problem in distributed multi-robot applications where each robot does not have a global view of the world. To solve it, we have accepted to reduce momentarily the "distribution level" of a part of the system at some very particular instants to increase its ability to treat

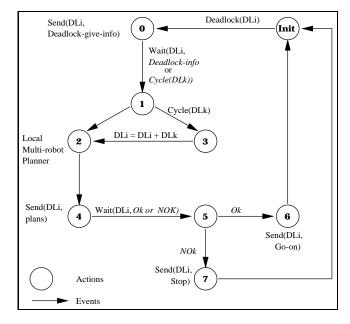


Figure 4: The coordinator finite state automate.

intricate situations. When a subset of robots enters in a deadlock cycle, the robot that detects the cycle becomes the cycle coordinator whose responsibility is to find a solution to the problem. If the coordinator finds a solution, it tries to validate it in the global plan GP, constructed from the coordinated plans of the non-blocked robots, by a PMO. If the insertion succeeds, the coordinator distributes the solution to the concerned robots and the system returns to its initial distributed state (Figure 6).

If the insertion fails and produces a new cycle, the coordinator recursively applies the same algorithm to the current *Meta-Node* and to the detected cycle to create a new *Meta-Node*. So, we may imagine some very complicated situations where the *Meta-Node* starts to grow up and does not stop until the inclusion of the whole system (all robots). In such situation, our completely distributed system tends to a completely centralized system (Figure 6).

Note also that we may have, in parallel, many deadlocks which do not interfere and which are solved independently. At the same time, we may have other cycles that group and un-group dynamically depending on the context.

#### 5 Examples

In order to illustrate our approach, we have implemented a generalized PMP that takes into account a conjunction of goals characterizing a deadlock situation. The application involves a large fleet of autonomous mobile robots [4, 1]. While the overall system allows to operate a large number of robots in a route-network environment, we will limit ourselves here to intricate situations that may happen from time

 $<sup>{}^3</sup>R_i$  acts here for itself or as a local coordinator for a set of robots determined in a former step

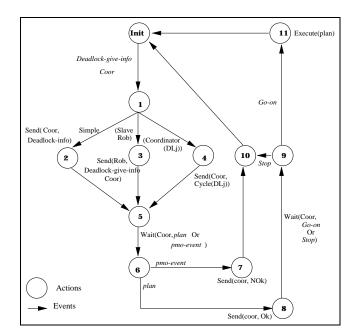


Figure 5: The other robots finite state automate.

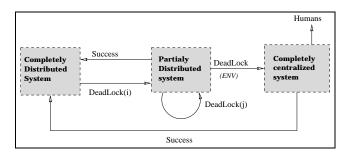


Figure 6: The evolution of the global system to a more or less distributed system is function of the situations complexity.

to time.

Each robot  $R_i$  is equipped with a multi-robot planner which can be used for an arbitrary number of robots. Such a planner allows to plan and synchronize paths in an environment described as a graph of spatial entities (cells & stations) using an A\* algorithm.

Note that we could have used also a multi-robot motion planner. However, even though it would allow to solve intricate situations without a pre-structuring of the environment into a discrete set of places, such planners can hardly be used when the number of robots is greater than 3[10].

#### 5.1A Simple Deadlock Situation

This example treats a simple deadlock  $DL_3$  involving two robots  $(R_3, R_6)$  where the goal of each robot is the initial position of the other one (Figure 7).

Robot Initial-State Goal-State

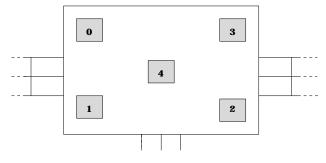


Figure 7: A part of the testbed environment: an area with 5 stations.

$R_6$	$Station_1$	$Station_0$
$R_3$	$Station_0$	$Station_1$

 $R_6$  fails in its PMO and decides to wait for a *plan*ning event from  $Pred_6 = \{R_3\}$ . It propagates this information to  $R_3$  which adds  $R_6$  to  $PDG_3$ .  $R_3$  performs a PMO, detects a cycle and becomes the coordinator for  $DL_3 = \{R_3, R_6\}$ .  $R_3$  invokes its local multi-robot planner and finds a solution to the given conflict. The solution (trivially) uses  $station_4$  as a buffer. finally,  $R_3$ performs a PMO in order to insert such a plan. No other robot is present in the area; the PMO succeeds.

#### 5.2Two Independent Deadlock Cycles

To increase the complexity of the situation, let us create a second deadlock  $DL_1 = \{R_1, R_0\}$  and assume that the coordinator  $R_1$  elaborates a plan for  $R_1$  and  $R_0$  which uses the same station (*Station*<sub>4</sub>) as a buffer. (Figure 7).

$\operatorname{Robot}$	Initial-State	Goal-State
$R_0$	$Station_3$	$Station_2$
$R_1$	$Station_2$	$Station_3$

Two independent cycles  $DL_3 = \{R_3, R_6\}$  and  $DL_1 = \{R_1, R_0\}$  are created and use the same buffer station  $(Station_4)^4$ , two coordinators work these two deadlocks and two independent "lattice solutions" are found. The resultant lattices are coordinated with each other and with the other robots' coordinated plans by a PMO.

The validation of these two lattices in the global context imposes a new synchronization event between  $R_1$ and  $R_6$  concerning the occupation of Station<sub>4</sub> (Figure 8)

#### 5.3Two Incompatible Deadlock Cycles

This example treats the case where the system switches to a centralized system when many deadlocks emerge requiring one global centralized planning activity.

The initial and final states are given below (Figure 7) :

<sup>&</sup>lt;sup>4</sup>This station minimizes the distance criteria given to the planner.

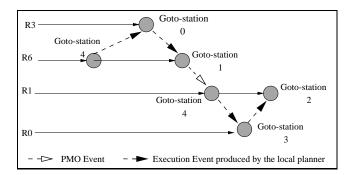


Figure 8: The set of the totally ordered actions that solves  $DL_1$  and  $DL_3$ .

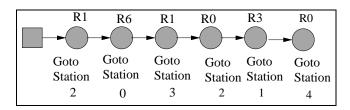


Figure 9: The set of synchronized actions that solves the conflict after the aggregation of two deadlocks  $DL_1 = \{R_1, R_6\}$  and  $DL_3 = \{R_3, R_0\}$ .

$\operatorname{Robot}$	Initial-State	Goal-State
$R_6$	$Station_3$	$Station_0$
$R_3$	$Station_4$	$Station_1$
$R_0$	$Station_1$	$Station_4$
$R_1$	$Station_0$	$Station_3$

Two coordinators  $R_3$  and  $R_1$  try to break independently two detected deadlocks  $(DL_1 = \{R_1, R_6\} \& DL_3 = \{R_3, R_0\})$ . Each coordinator produces a solution  $Sol_i^{LC}$  that resolves locally the given deadlock. However,  $Sol_3^{LC}$  and  $Sol_1^{LC}$  cannot be merged together without introducing a cycle in GP. Therefore,  $R_3^{LC}$  takes the "control" of the situation, informs the blocked robots  $(R_0, R_1, R_6)$  that it became the new coordinator for this new deadlock  $(DL_3 \Leftarrow DL_3 \cup DL_1)$ . The two *Meta-nodes* are unified to produce one *Metanode* representing the four conflicting robots. Finally,  $R_3^{LC}$  produces a solution plan for the overall task (Figure 9).

Note that, in this example, the whole distributed system switched to a totally centralized one where the generated solution is thus valid in the multi-robot context without a PMO.

#### 6 Conclusion

The effectiveness of the Plan-merging paradigm has already been discussed and illustrated through the implementation of a system involving up to 30 simulated mobile robots. It has also been implemented on a set of 3 real robots in a laboratory environment. The Plan-merging paradigm is a well suited paradigm to multi-robot applications with looselycoupled tasks. However, even if an application is designed to ease robots interaction, one cannot guarantee in the general case that tightly-coupled tasks will never happen. For example, the robots may find themselves in intricate situations simply because of an unknown obstacle placed in a critical place. This is why it is important to design a system which is able to efficiently exploit the tasks decoupling, but which is also able to detect and solve transient "puzzle-like" situations.

We have presented here a set of extended operators and associated mechanisms which allow not only to detect but also to solve situations where the robots goals are tightly coupled. This extension is done for the sake of completeness. The operators permit a coherent management of the distributed planning and coordination processes as well as a progressive transition to more global schemes which may even "degrade" to a unique and centralized planning activity.

#### References

- Luis Aguilar, Rachid Alami, Sara Fleury, Matthieu Herrb, Félix Ingrand, and Frédéric Robert. Ten autonomous mobile robots (and even more) in a route network like environment. In *IEEE International Workshop on Intelligent Robots and Systems (IROS '95), Pittsburgh, (Pennsylvania USA)*, 1995.
- [2] R. Alami, S. Fleury, M. Herrb, F. Ingrand, and F. Robert. Multi Robot Cooperation in the Martha Project. *IEEE Robotics and Automation* Magazine, Robotics and Automation in Europe : Projects funded by the Commission of the European Union, ?(?), ? 1997. To be published in 1997. Also available as LAAS/CNRS Technical Note 96392.
- [3] R. Alami, F. Robert, F. F. Ingrand, and S. Suzuki. A paradigm for plan-merging and its use for multi-robot cooperation. In *IEEE International Conference on Systems, Man, and Cybernetics, San Antonio, Texas (USA)*, 1994.
- [4] R. Alami, F. Robert, F. F. Ingrand, and S. Suzuki. Multi-robot cooperation through incremental plan-merging. In *IEEE Interna*tional Conference on Robotics and Automation, Nagoaya, (Japan), 1995.
- [5] K. Azarm and G. Schmidt. A decentralized approach for the conflict-free motion of multiple mobile robots. *iros-96*, 3, November 1996.
- [6] Harinaraya and Lumelsky. Sensored-based motion planning for multiple mobile robots in an uncertain environment. In *IEEE International* Workshop on Intelligent Robots and Systems

(IROS '94), Munich, (Germany), pages 1485-1492, 1994.

- S. Kato, S. Nishiyama, and J. Takeno. Coordinating mobile robots by applying traffic rules. In *IEEE International Conference on Intelligent Robots and Systems (IROS '92), Raleigh (North Carolina, USA)*, pages 1535-1541, July 1992.
- [8] L.E. Parker. Heterogeneous multi-robot cooperation. Technical Report AITR-1465, MIT, 1994.
- [9] Y. Shoham and M. Tennenholtz. On social laws for artificial societies: Off-line design. Artificial Intelligence, 734, 1995.
- [10] P. Svestka and M.H. Overmars. Coordinated Motion Planning for Multiple Car-Like Robots using Probabilistic Roadmaps. In *IEEE International Conference on Robotics and Automation*, *Nagoaya*, (Japan), 1995.
- [11] Warren. Multiple-robots path coordination using artificial potentiel fields. In IEEE International Conference on Robotics and Automation, Cincinnati, (USA), pages 500-505, 1990.
- [12] S. Yuta and S.Premvuti. Coordinating autonomous and centralized decision making to achieve cooperative behaviors between multiple mobile robots. In *IEEE International Conference* on Intelligent Robots and Systems (IROS '92), Raleigh (North Carolina, USA), pages 1566-1574, July 1992.