# Sparse Complementary Pairs with Additional Aperiodic ZCZ Property

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*Abstract*—This paper presents a novel class of complex-valued sparse complementary pairs (SCPs), each consisting of a number of zero values and with additional zero-correlation zone (ZCZ) property for the aperiodic autocorrelations and crosscorrelations of the two constituent sequences. Direct constructions of SCPs and their mutually-orthogonal mates based on restricted generalized Boolean functions are proposed. It is shown that such SCPs exist with arbitrary lengths and controllable sparsity levels, making them a disruptive sequence candidate for modern low-complexity, low-latency, and low-storage signal processing applications.

### I. INTRODUCTION

The concept of Golay complementary pair (GCP) [1] refers to a pair of sequences with zero aperiodic autocorrelation sum (AACS) at every non-zero time-shift. Driven by the diverse applications of GCPs in coding and telecommunication [2]– [9], there are also Golay complementary sets (GCSs) [10]– [14], each comprised of two or more constituent sequences, and Z-complementary pairs (ZCPs) with zero-correlation zone (ZCZ) sum properties [15]–[20].

With numerous research works on complementary sequence pairs/sets, the existing state-of-the-art mostly does not consider the auto- and cross-correlations among the constituent sequences. This is perhaps because of the common assumption that these constituent sequences are sent over orthogonal frequency/time channels. However, in a practical communication scenario where the transmission of one sequence interferes with the other, good inter-sequence aperiodic correlation properties are highly desirable. Among few exceptions, cross Zcomplementary pairs (C-ZCPs) were first proposed in 2020 for optimal channel training in broadband spatial modulation systems [21], where each pair is characterized by the ZCZ properties for both their AACSs and aperiodic crosscorrelation sums [22]. For Doppler-resilient radar waveforms, an efficient majorization minimization algorithm was developed in [23] for sequence pairs called quasi-orthogonal ZCPs (QO-ZCPs). In a QO-ZCP, both the AACSs and aperiodic crosscorrelation

values (rather than the aperiodic crosscorrelation sums) of the two constituent sequences have near-zero values within the zone.

The main objective of this work is to introduce a novel class of complex-valued sparse complementary pairs (SCPs), where the constituent sequences of each SCP contain a number of zeros and exhibit zero aperiodic auto- and cross-correlation zone property\*. These sparse pairs are useful in signal processing and communication/radar system design. For example, with the recent advances in sparse signal processing, the sequence sparsity can be smartly exploited for low-complexity, lowlatency, and low-storage hardware implementation [26]–[29]. For the aforementioned channel training of spatial modulation and Doppler-resilient radar sensing, such sparse pairs may be used as an alternative for C-ZCPs and QO-ZCPs, respectively.

Besides, we study the sparse complementary mate of an SCP, where the two pairs are mutually orthogonal in terms of the zero aperiodic cross-correlation sums for all the time shifts. It is shown that through a careful design, the four constituent sequences (see  $C_0, C_1, S_0, S_1$  in Fig. 1) can also enjoy zero aperiodic cross-correlation zone property. By using restricted generalized Boolean functions (RGBFs) [11], [13], the core idea of our proposed constructions is to carefully restrict certain variables of RGBFs, thus leading to flexible sequence lengths, sparsity, and ZCZ widths.

The rest of this paper is outlined as follows. Section II introduces some notations and preliminaries including the definitions of SCPs and their mates. The concept of RGBFs is introduced in Section III. Then we present direct constructions of SCPs and SCPs based on RGBFs in Section IV. Finally, we conclude our paper in Section V.

#### II. BACKGROUND AND DEFINITIONS

The following notations will be used throughout this paper:

- Z<sub>q</sub> = {0, 1, ..., q−1} denotes the set of integers modulo a positive integer q.
- $\xi = e^{2\pi\sqrt{-1/q}}$  denotes a q-th primitive root of unity.
- $(\cdot)^*$  denotes the complex conjugation.

\*Although there are ternary complementary pairs (TCPs) [24], [25] over the alphabet of  $\{+1, 0, -1\}$ , the aperiodic correlation properties of the constituent sequences have not been investigated.

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Fig. 1. Aperiodic cross-correlation properties of an SCP and its mutually orthogonal mate.

Let  $C_0$  and  $C_1$  be two complex-valued sequences of length L given by  $C_k = (C_{k,0}, C_{k,1}, \ldots, C_{k,L-1})$ , where  $|C_{k,i}| \in \{0,1\}$  and  $C_{k,0}, C_{k,L-1} \neq 0$ . The aperiodic cross-correlation of sequences  $C_0$  and  $C_1$  at the time-shift u is defined as

$$\rho(\boldsymbol{C}_0, \boldsymbol{C}_1; u) = \begin{cases} \sum_{i=0}^{L-1-u} C_{0,i+u} C_{1,i}^*, & 0 \le u < L; \\ \sum_{i=0}^{L-1+u} C_{0,i} C_{1,i-u}^*, & -L < u < 0. \end{cases}$$
(1)

Note that  $\rho(C_0, C_1; u) = \rho^*(C_1, C_0; -u)$ . If  $C_0 = C_1$ ,  $\rho(C_0, C_0; u)$  is called the aperiodic autocorrelation of  $C_0$ , denoted as  $\rho(C_0; u)$ . In what follows, we formally define *sparse complementary pairs*.

Definition 1: Let  $(C_0, C_1)$  be a pair of complex-valued sequences with length L and N zero elements in each sequence. Suppose that the following two conditions hold:

$$C1: \rho(\boldsymbol{C}_{k}, \boldsymbol{C}_{k'}; u) = \begin{cases} L - N, & u = 0, \ k = k' \\ 0, & 0 < |u| < Z, \ k = k' \\ 0, & |u| < Z, \ k \neq k' \\ C2: \rho(\boldsymbol{C}_{0}; u) + \rho(\boldsymbol{C}_{1}; u) = \begin{cases} 0, & u \neq 0 \\ 2(L - N), & u = 0. \end{cases}$$
(2)

Then  $(C_0, C_1)$  is said to be an (L, Z, S)-SCP, where Z denotes the width of ZCZ and S = N/L represents the sparsity level.

Note that for the well-known GCP (S = 0), only C2 in the above definition is satisfied.

Definition 2: Two distinct (L, Z, S)-SCPs  $(C_0, C_1)$  and  $(S_0, S_1)$  are said to be the mate of each other if

C1: 
$$\rho(C_0, S_0; u) + \rho(C_1, S_1; u) = 0, |u| < L;$$
  
C2:  $\rho(C_k, S_{k'}; u) = 0, |u| < Z, \text{ and } k, k' = 0, 1.$ 
(3)

The cross-correlation properties of these two SCPs are depicted in Fig. 1.

#### **III. RESTRICTED GENERALIZED BOOLEAN FUNCTIONS**

In this section, we will introduce the concept of restricted generalized Boolean functions (RGBFs) [11], [13]. First, we introduce the concept of GBF. A GBF  $f: \mathbb{Z}_2^m \to \mathbb{Z}_q$  is composed of m variables  $x_1, x_2, \ldots, x_m$  where  $x_l \in \{0, 1\}$  for  $l = 1, 2, \ldots, m$ . The sequence f corresponding to the GBF f is denoted by  $f = (f_0, f_1, \ldots, f_{2^m-1})$  where  $f_i = f(i_1, i_2, \ldots, i_m)$  and  $i = \sum_{l=1}^m i_l 2^{l-1}$ . The complex-valued sequence associated with f is given by  $\xi^f = (\xi^{f_0}, \xi^{f_1}, \ldots, \xi^{f_{2^m-1}})$ .

*Example 1:* Taking q = 4, m = 3, the associated q-ary sequence f to the GBF  $f = 2x_2x_3 + x_1$  is given by f = (01010123). The complex-valued sequence is  $\xi^f = (\xi^0, \xi^1, \xi^0, \xi^1, \xi^0, \xi^1, \xi^2, \xi^3)$ .

Consider a set of t indices  $V = \{v_1, v_2, \ldots, v_t\} \subset \{1, 2, \ldots, m\}$  and a set of m - t indices  $V' = \{v'_1, v'_2, \ldots, v'_{m-t}\} = \{1, 2, \ldots, m\} \setminus V$ , where  $0 \le t < m$ . Let  $\boldsymbol{x} = (x_{v_1}, x_{v_2}, \ldots, x_{v_t})$  and  $\boldsymbol{d} = (d_1, d_2, \ldots, d_t)$  where  $d_l \in \{0, 1\}$ . A RGBF  $f|_{\boldsymbol{x}=\boldsymbol{d}}$  is defined by restricting variables  $\boldsymbol{x}$  in the f to the certain known d. For simplicity, let the associated sequence  $\boldsymbol{f}|_{\boldsymbol{x}=\boldsymbol{d}}$  be the complex-valued sequence with component equal to  $\xi^{f_i}$  if  $i_{v_\alpha} = d_\alpha$  for  $\alpha = 1, 2, \ldots, t$ , and equal to zero otherwise. Let  $\boldsymbol{F} = (F_0, F_1, \ldots, F_{2^m-1}) = \boldsymbol{f}|_{\boldsymbol{x}=\boldsymbol{d}}$ . Let  $k_0$  and  $k_1$  be the smallest and the largest integers, satisfying  $F_{k_0}, F_{k_1} \neq 0$ , respectively. It is clear that

$$k_0 = \sum_{\alpha=1}^{t} d_{\alpha} 2^{v_{\alpha}-1}$$
 and  $k_1 = k_0 + \sum_{\alpha=1}^{m-t} 2^{v'_{\alpha}-1}$ . (4)

Let  $L = k_1 - k_0 + 1$  and define  $f^{(L)}|_{\boldsymbol{x}=\boldsymbol{d}}$  as the truncated sequence by removing the first  $k_0$  elements and the last  $2^m - 1 - k_1$  elements. Therefore, the constructed sparse sequence is of length L. In this paper, whenever L is known from the context, we use  $f|_{\boldsymbol{x}=\boldsymbol{d}}$  to denote  $f^{(L)}|_{\boldsymbol{x}=\boldsymbol{d}}$ .

*Remark 1:* Given a sparse sequence  $f^{(L)}|_{x=d}$ , the number of non-zero elements is  $2^{m-t}$ . Hence, the sparsity level of the sequence  $f^{(L)}|_{x=d}$  is given by

$$S = \frac{L - 2^{m-t}}{L}.$$
(5)

*Example 2:* Following the same notations given in *Example 1*, let t = 1,  $x = (x_2)$ , and d = (0). The associated sequence  $f|_{x_2=0}$  to the RGBF is  $f|_{x_2=0} = (\xi^0, \xi^1, 0, 0, \xi^0, \xi^1, 0, 0)$ . From (4), we have  $k_0 = 0$  and  $k_1 = 5$  and hence L = 6. The truncated sequence is  $f^{(6)}|_{x_2=0} = (\xi^0, \xi^1, 0, 0, \xi^0, \xi^1, 0, 0, \xi^0, \xi^1, 0, 0, \xi^0, \xi^1)$  with the sparsity  $S = \frac{6-4}{6} = \frac{1}{3}$ .

## IV. PROPOSED CONSTRUCTIONS OF SCPS AND THEIR MATES

In this section, we will provide constructions of SCPs and their mates based on RGBFs.

Theorem 1: For integers m, t with  $t \leq m-1$ , let  $\pi$  be a permutation of the set  $\{1, 2, \ldots, m\}$  with  $\pi(m) > \pi(\alpha)$  for  $1 \leq \alpha \leq t$ . Let  $\boldsymbol{x} = (x_{\pi(1)}, x_{\pi(2)}, \ldots, x_{\pi(t)})$ , and binary

sequence  $d = (d_1, d_2, \dots, d_t)$ . Let us consider the RGBF given below:

$$f|_{\boldsymbol{x}=\boldsymbol{d}} = \frac{q}{2} \left( \sum_{l=1}^{t-1} d_l d_{l+1} + \sum_{l=t+1}^{m-1} x_{\pi(l)} x_{\pi(l+1)} + d_t x_{\pi(t+1)} \right) + \sum_{l=1}^{m} g_l x_l + g_0,$$
(6)

where  $g_l \in \mathbb{Z}_q$ . The pair

$$(C_0, C_1) = (f|_{x=d}, (f + (q/2)x_{\pi(t+1)})|_{x=d})$$
 (7)

is an (L, Z, S)-SCP with  $L = \sum_{\alpha=t+1}^{m} 2^{\pi(\alpha)-1} + 1$ ,  $Z = \sum_{\alpha=1}^{t} 2^{\pi(\alpha)-1} + 1$ , and  $S = \frac{L-2^{m-t}}{L}$ .

**Proof:** Without loss of generality, we consider  $d = 0_t$  where  $0_t$  is the all-zero vector of length t. Given the pair  $(C_0, C_1)$  generated from Theorem 1, they can be written as  $C_k = (C_{k,0}, C_{k,1}, \ldots, C_{k,L-1}) =$  $(f + \frac{q}{2} \cdot k \cdot \boldsymbol{x}_{\pi(t+1)})|_{\boldsymbol{x}=d}, k = 0, 1$ , where  $C_{k,i} = 0$  if  $(i_{\pi(1)}, i_{\pi(2)}, \ldots, i_{\pi(t)}) \neq (d_1, d_2, \ldots, d_t)$  or  $C_{k,i} = \xi^{c_{k,i}}$  if  $i_{\pi(\alpha)} = d_{\alpha}$  for  $\alpha = 1, 2, \ldots, t$ . We need to show that C1 and C2 in (3) hold.

In the first part, we start with checking C1, i.e.,

$$\rho(\mathbf{C}_k, \mathbf{C}_{k'}; u) = \sum_{i=0}^{L-1-u} C_{k,i+u} C_{k',i}^*$$
$$= \begin{cases} 0, & 0 < u < Z, \ k = k' \\ 0, & 0 \le u < Z, \ k \ne k' \end{cases}$$

where k, k' = 0, 1. Let j = i + u and also let  $(j_1, j_2, \ldots, j_m)$ and  $(i_1, i_2, \ldots, i_m)$  be binary representations of the integers  $j = \sum_{l=1}^{m} j_l 2^{l-1}$  and  $i = \sum_{l=1}^{m} i_l 2^{l-1}$ , respectively. Note that  $C_{k,j}C_{k',i}^* = 0$  if  $(j_{\pi(1)}, j_{\pi(2)}, \ldots, j_{\pi(t)})$ or  $(i_{\pi(1)}, i_{\pi(2)}, \ldots, i_{\pi(t)})$  is not equal to  $(d_1, d_2, \ldots, d_t)$ . Hence, throughout the proof, we only need to focus on the multiplications of non-zero elements, i.e.,  $C_{k,j}C_{k',i}^* \neq 0$ implying  $(j_{\pi(1)}, j_{\pi(2)}, \ldots, j_{\pi(t)}) = (i_{\pi(1)}, i_{\pi(2)}, \ldots, i_{\pi(t)}) =$  $(d_1, d_2, \ldots, d_t)$ . Then three cases are considered below.

*Case 1 (i):* k = k'. In this case, we consider the aperiodic autocorrelation of  $C_k$ . For  $i_{\pi(\alpha)} = j_{\pi(\alpha)}$ ,  $1 \le \alpha \le t$ , we show next that  $i_{\pi(m)} = j_{\pi(m)}$  holds. Let us assume that  $i_{\pi(m)} \ne j_{\pi(m)}$ . Denote the set  $W = \{w_0, w_1, w_2, \ldots, w_\gamma\}$ such that  $\pi(m) = \pi(w_0) < \pi(w_1) < \pi(w_2) < \cdots < \pi(w_\gamma)$ where  $0 \le \gamma \le m - t - 1$ . Assume that  $\kappa$  is the largest integer satisfying  $i_{\pi(w_{\kappa})} \ne j_{\pi(w_{\kappa})}$  and  $A = \{1, 2, \ldots, m\} \setminus$  $\{1, \ldots, t, w_{\kappa+1}, \ldots, w_{\gamma}\}$ . That is to say, A contains all the indices  $\alpha$  for which  $i_{\pi(\alpha)} \ne j_{\pi(\alpha)}$ . Since u = j - i > 0, we have  $j_{\pi(w_{\kappa})} = 1$  and  $i_{\pi(w_{\kappa})} = 0$ . Then, we obtain

$$u = j - i = 2^{\pi(w_{\kappa}) - 1} + \sum_{\alpha \in A \setminus \{w_{\kappa}\}} (j_{\pi(\alpha)} - i_{\pi(\alpha)}) 2^{\pi(\alpha) - 1}$$
  
$$\geq 2^{\pi(w_{\kappa}) - 1} - \sum_{\alpha \in A \setminus \{w_{\kappa}\}} 2^{\pi(\alpha) - 1} = \sum_{\alpha = 1}^{t} 2^{\pi(\alpha) - 1} + 1,$$
(8)

since  $\sum_{\alpha=1}^{t} 2^{\pi(\alpha)-1} + \sum_{\alpha \in A \setminus \{w_{\kappa}\}} 2^{\pi(\alpha)-1} = 2^{\pi(w_{\kappa})-1} - 1$ . Obviously, this contradicts the assumption that  $u < Z = \sum_{\alpha=1}^{t} 2^{\pi(\alpha)-1} + 1$ . Thus, we have  $i_{\pi(m)} = j_{\pi(m)}$ . For simplicity, we denote  $\sigma(\alpha) = \pi(m+1-\alpha)$ . Then we have  $i_{\sigma(1)} = j_{\sigma(1)}$ . Assume that  $\beta$  is the smallest integers such that  $i_{\sigma(\beta)} \neq j_{\sigma(\beta)}$ . Let i' and j' be integers different from i and j, respectively, in the position  $\sigma(\beta - 1)$ , i.e.,  $i'_{\sigma(\beta-1)} = 1 - i_{\sigma(\beta-1)}$  and  $j'_{\sigma(\beta-1)} = 1 - j_{\sigma(\beta-1)}$ . We can obtain

$$i', j' \le \sum_{\alpha=t+1}^{m} 2^{\pi(\alpha)-1} - 2^{\sigma(\beta-1)-1} + 1 < L.$$

Then we have

$$\begin{aligned} c_{k,i'} &- c_{k,i} \\ &= \frac{q}{2} (i_{\sigma(\beta-2)} i'_{\sigma(\beta-1)} - i_{\sigma(\beta-2)} i_{\sigma(\beta-1)} \\ &+ i'_{\sigma(\beta-1)} i_{\sigma(\beta)} - i_{\sigma(\beta-1)} i_{\sigma(\beta)}) \\ &+ g_{\sigma(\beta-1)} i'_{\sigma(\beta-1)} - g_{\sigma(\beta-1)} i_{\sigma(\beta-1)} \\ &\equiv \frac{q}{2} (i_{\sigma(\beta-2)} + i_{\sigma(\beta)}) + g_{\sigma(\beta-1)} (1 - 2i_{\sigma(\beta-1)}) \pmod{q}. \end{aligned}$$

Since  $i_{\sigma(\beta-1)} = j_{\sigma(\beta-1)}$  and  $i_{\sigma(\beta-2)} = j_{\sigma(\beta-2)}$ , we possess

$$c_{k,j} - c_{k,i} - c_{k,j'} + c_{k,i'} \equiv \frac{q}{2}(i_{\sigma(\beta)} - j_{\sigma(\beta)}) \equiv \frac{q}{2},$$
 (9)

implying  $\xi^{c_{k,j}-c_{k,i}} + \xi^{c_{k,j'}-c_{k,i'}} = 0.$ 

*Case 1 (ii):*  $k \neq k'$  and  $u \neq 0$ . In this case, we show that  $\rho(C_0, C_1; u) = 0$  for  $0 < u < Z_A$ . Based on *Case 1 (i)*, we know  $i_{\pi(m)} = j_{\pi(m)}$  and let  $\pi(\alpha) = \sigma(m + 1 - \alpha)$ . Assume that  $\beta$  is the smallest integer satisfying  $i_{\sigma(\beta)} \neq j_{\sigma(\beta)}$ . Let *i'* and *j'* be integers distinct from *i* and *j*, respectively, in the position  $\sigma(\beta - 1)$ . Following the similar arguments as given above, we can obtain  $\xi^{c_{0,j}-c_{1,i}} + \xi^{c_{0,j'}-c_{1,i'}} = 0$ .

Case 1 (iii):  $k \neq k'$  and  $u \neq 0$ . In this case, we prove  $\rho(\mathbf{C}_0, \mathbf{C}_1; 0) = 0$ . From (7), we have  $\rho(\mathbf{C}_0, \mathbf{C}_1; 0) = \sum_{i=0}^{L-1} C_{0,i}C_{1,i}^*$  where  $C_{0,i}C_{1,i}^* = \xi^{\frac{q}{2}i_{\pi(t+1)}}$  for  $i_{\pi(\alpha)} = d_{\alpha}$ ,  $1 \leq \alpha \leq t$ , and  $i_{\pi(t+1)}$  is the  $\pi(t+1)$ -th bit of the binary representation of *i*. According to *Remark 1*, each sequence has  $2^{m-t}$  non-zero elements so that we can find  $2^{m-t-1}$ pairs fulfilling  $\xi^{c_{0,i}-c_{1,i}} = \xi^{q/2} = -1$  and another  $2^{m-t-1}$ pairs such that  $\xi^{c_{0,i}-c_{1,i}} = \xi^0 = 1$ . Therefore, we have  $\rho(\mathbf{C}_0, \mathbf{C}_1; 0) = 0$ .

From Case 1 (i) to Case 1 (iii), we see C1 of (2) holds.

Now we consider condition C2. That is to say, we need to show that

$$\rho(\boldsymbol{C}_0; u) + \rho(\boldsymbol{C}_1; u) = 0$$

for 0 < u < L. Taking  $i_{\pi(\alpha)} = j_{\pi(\alpha)} = d_{\alpha}$  for  $\alpha = 1, 2, ..., t$ , the following two cases are considered.

*Case 2 (i):* Suppose that  $i_{\pi(t+1)} \neq j_{\pi(t+1)}$ . We can obtain

$$c_{0,j} - c_{0,i} - c_{1,j} + c_{1,i} = \frac{q}{2} (i_{\pi(t+1)} - j_{\pi(t+1)}) \equiv \frac{q}{2} \pmod{q}$$

meaning that  $\xi^{c_{0,j}-c_{0,i}} + \xi^{c_{1,j}-c_{1,i}} = 0.$ 

Case 2 (ii): Considering  $i_{\pi(t+1)} = j_{\pi(t+1)}$ , we assume  $\beta$  is the smallest integer which fulfills  $i_{\pi(\beta)} \neq j_{\pi(\beta)}$ . Likewise, let (8) i' and j' be integers different from i, j, respectively, in the

TABLE ISCPs of Lengths up to 35

Length	15	17	18	19	21	22	23	24	25	27	28	29	30	31	33	34	35
m	4	5	5	5	5	5	5	5	5	5	5	5	5	5	6	6	6
Restricting variables	$x_1$	$\begin{array}{c} x_1, x_2 \\ x_3, x_4 \end{array}$	$egin{array}{c} x_2, x_3 \ x_4 \end{array}$	$egin{array}{c} x_1, x_3 \ x_4 \end{array}$	$egin{array}{c} x_1, x_2 \ x_4 \end{array}$	$x_2$ $x_4$	$\begin{array}{c} x_1 \\ x_4 \end{array}$	$x_4$	$egin{array}{c} x_1, x_2 \ x_3 \end{array}$	$\begin{array}{c} x_1 \\ x_3 \end{array}$	$x_3$	$\begin{array}{c} x_1 \\ x_2 \end{array}$	$x_2$	$x_1$	$egin{array}{c} x_1, x_2 \ x_3, x_4 \ x_5 \end{array}$	$egin{array}{l} x_2, x_3 \ x_4, x_5 \end{array}$	$egin{array}{c} x_1, x_3 \ x_4, x_5 \end{array}$
ZCZ width	2	16	15	14	12	11	10	9	8	6	5	4	3	2	32	31	30
Sparsity	$\frac{7}{15}$	$\frac{15}{17}$	$\frac{14}{18}$	$\frac{15}{19}$	$\frac{17}{21}$	$\frac{14}{22}$	$\frac{15}{23}$	$\frac{8}{24}$	$\frac{21}{25}$	$\frac{19}{27}$	$\frac{12}{28}$	$\frac{21}{29}$	$\frac{24}{30}$	$\frac{15}{31}$	$\frac{31}{33}$	$\frac{30}{34}$	$\frac{31}{35}$

position  $\pi(\beta - 1)$ . Following the similar arguments as given in *Case 1 (i)*, we obtain

$$\xi^{c_{0,j}-c_{0,i}}+\xi^{c_{0,j'}-c_{0,i'}}+\xi^{c_{1,j}-c_{1,i}}+\xi^{c_{1,j'}-c_{1,i'}}=0,$$

which completes the proof.

Remark 2: If we consider t = 0 in Theorem 1, the pair  $(C_0, C_1)$  is reduced to a GCP of length  $2^m$  which is the so-called Golay-Davis-Jedweb (GDJ) pair [2].

*Remark 3:* According to *Theorem 1*, the sequence length and the ZCZ width satisfy

$$L + Z = \sum_{\alpha=t+1}^{m} 2^{\pi(\alpha)-1} + 1 + \sum_{\alpha=1}^{t} 2^{\pi(\alpha)-1} + 1 = 2^{m} + 1.$$

Since  $\{\pi(1), \pi(2), \ldots, \pi(m)\} = \{1, 2, \ldots, m\}$  with  $\pi(m) > \pi(\alpha)$  for  $1 \le \alpha \le t$ , by taking different values of t and choices of  $\pi$ , we can flexibly generate SCPs with arbitrary length in the range  $[2^{m-1} + 1, 2^m]$ . In addition, we point out that there is a one-to-one correspondence between the sequence length L and t, permutation  $\pi$ .

*Remark 4:* The sparsity S is controlled by the sequence length and the number of restricted variables t, where the sequence length is associated with the restricted variables of RGBFs as stated in *Remark 3*.

In Table I, we list some examples of the proposed SCPs of lengths from 15 to 35 except for the lengths 16, 20, 26, 32 for binary GCPs, where the corresponding ZCZ widths and sparsities are given as well. It can be observed that by restricting different variables in *Theorem 1*, flexible lengths of SCPs are obtained.

*Example 3:* Taking q = 4, m = 5, t = 2, and  $\pi = (1, 3, 2, 4, 5)$ , let  $\boldsymbol{x} = (x_1, x_3)$  and  $\boldsymbol{d} = (d_1, d_2) = (0, 0)$ . The RGBF is

$$f|_{\boldsymbol{x}=\boldsymbol{d}} = 2 \cdot (x_2 x_4 + x_4 x_5 + d_1 d_2 + d_2 x_2) + 3x_2.$$

According to Theorem 1, the pair

$$(C_0, C_1) = (f|_{\boldsymbol{x}=\boldsymbol{d}}, (f+2\boldsymbol{x}_2)|_{\boldsymbol{x}=\boldsymbol{d}})$$

is a (27, 6, 19/27)-SCP where

$$\boldsymbol{C}_0 = (\xi^0 0 \xi^3 00000 \xi^0 0 \xi^1 00000 \xi^0 0 \xi^3 00000 \xi^2 0 \xi^3)$$

and

$$\boldsymbol{C}_1 = (\xi^0 0 \xi^1 00000 \xi^0 0 \xi^3 00000 \xi^0 0 \xi^1 00000 \xi^2 0 \xi^1).$$

In Fig. 2, we can see  $C_0$  and  $C_1$  have zero AACS for any non-zero time-shift. Besides, in Fig. 3, the ZCZ width of 6 of the aperiodic autocorrelations of  $C_0$  and cross-correlations of  $C_0$  and  $C_1$  are illustrated.



Fig. 2. Sum of aperiodic autocorrelations of  $C_0$  and  $C_1$  in *Example 3*.



Fig. 3. Aperiodic autocorrelations of  $C_0$  and cross-correlations of  $C_0$  and  $C_1$  in *Example 3*.

Next, we present a mate construction of an SCP based on *Theorem 1*.

Theorem 2: Let  $\pi(m-1) > \pi(\alpha)$  for  $1 \le \alpha \le t$ , and  $0 \le t \le m-2$  in Theorem 1, and  $(C_0, C_1)$  be the SCP given in (7). Then the pair  $(S_0, S_1)$  given by

$$\left(\left(\boldsymbol{f} + \frac{q}{2}\boldsymbol{x}_{\pi(m)}\right)\Big|_{\boldsymbol{x}=\boldsymbol{d}}, \left(\boldsymbol{f} + \frac{q}{2}\boldsymbol{x}_{\pi(t+1)} + \frac{q}{2}\boldsymbol{x}_{\pi(m)}\right)\Big|_{\boldsymbol{x}=\boldsymbol{d}}\right)$$

is a sparse mate of  $(C_0, C_1)$ .

*Proof:* For the two (L, Z, S)-SCPs  $(C_0, C_1)$  and  $(S_0, S_1)$ , we show that (3) holds. Likewise, we only consider the condition  $i_{\pi(\alpha)} = j_{\pi(\alpha)} = d_{\alpha}$ , where  $1 \leq \alpha \leq t$ . In the first part, we prove

$$\rho(C_k, S_{k'}; u) = 0, \ k, k' = 0, 1$$

for  $0 \leq u < Z$ . Following similar notations as given in the proof of *Theorem 1*, let  $\mathbf{S}_{k'} = (S_{k',0}, S_{k',1}, \dots, S_{k',L-1})$ where  $S_{k',i} = \xi^{s_{k',i}}$  if  $(i_{\pi(1)}, i_{\pi(2)}, \dots, i_{\pi(t)}) = \mathbf{d}$  and  $S_{k',i} = 0$  otherwise. We consider two cases below.

*Case 1 (i):* Given u > 0,  $0 \le t \le m - 2$ , we already have  $i_{\pi(m)} = j_{\pi(m)}$ . In this case, we next show that  $i_{\pi(m-1)} = j_{\pi(m-1)}$ . To proceed, let us assume that  $i_{\pi(m-1)} \ne j_{\pi(m-1)}$ . Denote the set  $\hat{W} = \{\hat{w}_0, \hat{w}_1, \dots, \hat{w}_\gamma\}$  such that  $\pi(m-1) = \pi(\hat{w}_0) < \pi(\hat{w}_1) < \pi(\hat{w}_2) < \dots < \pi(\hat{w}_\gamma)$  where  $0 \le \gamma \le m - t - 1$ . Assume that  $\lambda$  is the largest integer satisfying  $i_{\pi(\hat{w}_\lambda)} \ne j_{\pi(\hat{w}_\lambda)}$  and  $\hat{A} = \{1, 2, \dots, m\} \setminus \{1, \dots, t, \hat{w}_{\lambda+1}, \dots, \hat{w}_\gamma\}$ . If  $\pi(\hat{w}_\lambda) < \pi(m)$ , we have

$$u = j - i = 2^{\pi(\hat{w}_{\lambda}) - 1} + \sum_{\alpha \in \hat{A} \setminus \{\hat{w}_{\lambda}\}}^{m} (j_{\pi(\alpha)} - i_{\pi(\alpha)}) 2^{\pi(\alpha) - 1}$$
$$\geq 2^{\pi(\hat{w}_{\lambda}) - 1} - \sum_{\alpha \in \hat{A} \setminus \{\hat{w}_{\lambda}\}}^{m} 2^{\pi(\alpha) - 1} = \sum_{\alpha = 1}^{t} 2^{\pi(\alpha) - 1} + 1$$

which violates the assumption  $u < \sum_{\alpha=1}^{t} 2^{\pi(\alpha)-1} + 1$ . Similar results can be obtained if  $\pi(\hat{w}_{\lambda}) > \pi(m)$ . Hence, we have  $i_{\pi(m-1)} = j_{\pi(m-1)}$ . Similarly we denote  $\pi(\alpha) = \sigma(m+1-\alpha)$  and assume  $\beta$  is the smallest integers such that  $i_{\sigma(\beta)} \neq j_{\sigma(\beta)}$  where  $\beta > 2$ . Then let i' and j' be integers different from i and j, respectively, in the position  $\sigma(\beta-1)$ . Similar results can be obtained as provided in *Case 1 (i)* of the proof of *Theorem 1*. Hence, we possess  $\xi^{c_{k,j}-s_{k',i}} + \xi^{c_{k,j'}-s_{k',i'}} = 0$ .

Case 1 (ii): In this case, we show

$$\rho(\mathbf{C}_k, \mathbf{S}_{k'}; 0) = \sum_{i=0}^{L-1} C_{k,i} S_{k',i} = 0,$$

where  $C_{k,i}S_{k',i} = \xi^{c_{k,i}-s_{k',i}}$  and

$$c_{k,i} - s_{k',i} = \frac{q}{2} \left( (k - k')i_{\pi(t+1)} - i_{\pi(m)} \right)$$
(10)

for  $i_{\pi(\alpha)} = d_{\alpha}$ ,  $1 \le \alpha \le t$ . Clearly, (10) can be regarded as the linear combination of the terms  $i_{\pi(t+1)}$  and  $i_{\pi(m)}$ . Hence,  $\sum_{i=0}^{L-1} C_{k,i} S_{k',i} = 0$ .

From *Case 1 (i)* and *Case 1 (ii)*, we can see that the crosscorrelation of sequences  $C_k$  and  $S_{k'}$  is zero within the ZCZ Z. Then in the second part, we need to show

$$\rho(C_0, S_0; u) + \rho(C_1, S_1; u) = 0.$$

Three cases are considered below.

*Case 2 (i):* For u > 0, if  $\pi(t+1) \neq \pi(t+1)$ , we follow a similar derivation in *Case 2 (i)* of the proof of *Theorem 1* and have  $\xi^{c_{0,j}-s_{0,i}} + \xi^{c_{1,j}-s_{1,i}} = 0$ .

*Case 2 (ii):* For u > 0, assume  $i_{\pi(t+1)} = j_{\pi(t+1)}$ . Let  $\beta$  be the smallest integers such that  $i_{\pi(\beta)} \neq j_{\pi(\beta)}$ . Similar results can be obtained as given in *Case 2 (ii)* of the proof of *Theorem 1*. We have  $\xi^{c_{0,j}-s_{0,i}} + \xi^{c_{0,j'}-s_{0,i'}} + \xi^{c_{1,j'}-s_{1,i'}} = 0$ .

*Case 2 (iii):* Finally, it only suffices to validate that  $\rho(\mathbf{C}_0, \mathbf{S}_0; 0) + \rho(\mathbf{C}_1, \mathbf{S}_1; 0) = 0$ . By recalling the *Case 1 (ii)* of this proof, we complete the proof.

*Example 4:* Let us follow the same notations given in *Example 3.* According to *Theorem 2*, we let  $\pi(4) = 4$  and obtain the pair  $(S_0, S_1) = ((f + 2x_5)|_{x=d}, (f + 2x_2 + 2x_5)|_{x=d})$  given by

$$\boldsymbol{S}_0 = (\xi^0 0 \xi^3 00000 \xi^0 0 \xi^1 00000 \xi^2 0 \xi^1 00000 \xi^0 0 \xi^1)$$

and

$$\boldsymbol{S}_1 = (\xi^0 0 \xi^1 0 0 0 0 \xi^0 0 \xi^3 0 0 0 0 \xi^2 0 \xi^3 0 0 0 0 \xi^0 0 \xi^3).$$

In Fig. 4, we illustrate the aperiodic cross-correlations of  $C_0$  and  $S_0$  and cross-correlations of  $C_0$  and  $S_1$ . The ZCZ width is 6. Besides, it can be validated that they have zero cross-correlation sums for every time-shift.



Fig. 4. Aperiodic cross-correlations of  $C_0$  and  $S_0$  and cross-correlations of  $C_0$  and  $S_1$  in *Example 4*.

#### V. CONCLUSION

In this paper, we have introduced a novel class of sequence pairs called SCPs in which every SCP has zero AACS at every non-zero time-shift and the aperiodic autocorrelations and cross-correlations of the two constituent sequences are zero within a zone. Thanks to the use of RGBFs, we have introduced a direct construction of SCPs with arbitrary lengths and controllable sparsity levels. We have further constructed the mutually orthogonal mate of an SCP by considering crosscorrelation properties between the two distinct SCPs.

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