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Recovery Guarantees for Rank Aware Pursuits

Jeffrey D. Blanchard and Mike E. Davies

Abstract—This paper considers sufficient conditions for sparse recovery in the sparse multiple measurement vector (MMV) problem for some recently proposed rank aware greedy algorithms. Specifically we consider the compressed sensing framework with Gaussian random measurement matrices and show that the rank of the measurement matrix in the noiseless sparse MMV problem allows such algorithms to reduce the effect of the $\log n$ term that is present in traditional OMP recovery.

Index Terms—Multiple Measurement Vectors, Greedy Algorithm, Orthogonal Matching Pursuit, rank.

I. INTRODUCTION

Sparse signal representations provide a general signal model that make it possible to solve many ill-posed problems such as source separation, denoising and most recently compressed sensing [1] by exploiting the additional sparsity constraint. While the general problem of finding the sparsest $\mathbf{x} \in \mathbb{R}^n$ given an observation vector $\mathbf{y} = \Phi\mathbf{x}$, $\mathbf{y} \in \mathbb{R}^m$ is known to be NP-hard [2] a number of suboptimal strategies have been shown to be able to recover k -sparse signals, \mathbf{x} , when $m \sim Ck \log(n/k)$ for some constant C , if Φ is chosen judiciously.

An interesting extension of the sparse recovery problem is the sparse multiple measurement vector (MMV) problem, $\mathbf{Y} = \Phi\mathbf{X}$, $\mathbf{Y} \in \mathbb{R}^{m \times l}$, $\mathbf{X} \in \mathbb{R}^{n \times l}$, which has also received much attention, e.g. [3], [4], [5]. Initially the algorithms proposed for this problem were straightforward extensions of existing single measurement vector (SMV) solutions. However, most of these are unable to exploit the additional information available through the rank of \mathbf{Y} . In contrast, some new greedy algorithms for joint sparse recovery have been proposed [6], [7], [8] based around the MUSIC (Multiple Signal Classification) algorithm [9] from array signal processing which provides optimal recovery in the maximal rank scenario $r = k$ [10].

The aim of this paper is to analyse the recovery performance of two Rank Aware algorithms when the observation matrix does not have maximal rank, $\text{rank}(\mathbf{Y}) < k$. Our approach follows the recovery analysis of [11] where it was shown that, in the noiseless case, Orthogonal Matching Pursuit (OMP) can recover k -sparse vectors from $m \gtrsim Ck \log n$ Gaussian measurements with high probability. We extend this analysis to the MMV sparse recovery problem and show joint k -sparse matrices, \mathbf{X} , can be recovered from $m \gtrsim Ck(\frac{1}{r} \log n + 1)$

MMVs using a rank aware algorithm.¹ This implies that the $\log n$ penalty term observed for OMP recovery can be essentially removed with very modest values of rank, $r \gtrsim \log n$.

II. NOTATION AND PROBLEM FORMULATION

We define the support of a collection of vectors $\mathbf{X} = [\mathbf{x}_1, \dots, \mathbf{x}_l]$ as the union over all the individual supports: $\text{supp}(\mathbf{X}) := \bigcup_i \text{supp}(\mathbf{x}_i)$. A matrix \mathbf{X} is called k joint sparse if $|\text{supp}(\mathbf{X})| \leq k$. We make use of the subscript notation Φ_Ω to denote a submatrix composed of the columns of Φ that are indexed in the set Ω , while the notation $\mathbf{X}_{\Omega,:}$ denotes a row-wise submatrix composed of the rows of \mathbf{X} indexed by Ω . Thus denoting by $|\Omega|$ the cardinality of Ω , the matrix $\mathbf{X}_{\Omega,:}$ is $|\Omega|$ -sparse.

We can now formally define the sparse MMV problem. Consider the observation matrix $\mathbf{Y} = \Phi\mathbf{X}$, $\mathbf{Y} \in \mathbb{R}^{m \times l}$ where $\Phi \in \mathbb{R}^{m \times n}$ with $m < n$ is the dictionary matrix and $\mathbf{X} \in \mathbb{R}^{n \times l}$ is assumed to be jointly k -sparse. The task is then to recover \mathbf{X} from \mathbf{Y} given Φ . We will further assume that $\text{rank}(\mathbf{Y}) = r$ and without loss of generality that $r = l$.

III. GREEDY MMV ALGORITHMS

Despite the fact that the rank of the observation matrix \mathbf{Y} can be exploited to improve recovery performance, to date most popular techniques have ignored this fact and have been shown to be “rank-blind” [6]. In contrast, a discrete version of the MUSIC algorithm [10], [13] is able to recover \mathbf{X} from \mathbf{Y} under mild conditions on Φ whenever $m \geq k + 1$ if we are in the maximal rank case, i.e. $\text{rank}(\mathbf{Y}) = k$.

While MUSIC provides guaranteed recovery for the MMV problem in the maximal rank case there are no performance guarantees for when $\text{rank}(\mathbf{X}) < k$ and empirically MUSIC does not perform well in this scenario. This motivated a number of works [6], [7], [8] to investigate the possibility of an algorithm that in some way interpolates between a classical greedy algorithm for the SMV problem and MUSIC when $\text{rank}(\mathbf{X}) = k$.

The approach proposed in [7], [8] was to use a greedy selection algorithm to find the first $t = k - r$ coefficients. The remaining components can then be found by applying MUSIC to an augmented data matrix $[\mathbf{Y}, \Phi_\Lambda^{(t)}]$ which under identifiability assumptions will span the range of Φ_Λ .

In [6] two “rank aware” (RA) algorithms were presented. In RA-OMP the greedy selection step was modified to measure the distance of the columns of Φ from the subspace spanned by the residual matrix at iteration t by measuring the correlation of columns of Φ with an orthonormal basis of the residual

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¹These results were previously announced at the “SMALL” workshop, London, January 2011 [12].

Algorithm 1 RA-OMP / RA-ORMP

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1: initialization:  $\mathbf{R}^{(0)} = \mathbf{Y}, \mathbf{X}^{(0)} = \mathbf{0}, \Omega^0 = \emptyset$ 
2: for  $j = 1; j := j + 1$  until stopping criterion do
3:   Calculate  $\mathbf{U}^{(j-1)} = \text{ortho}(\mathbf{R}^{(j-1)})$ ,
4:   if (RA-OMP) then
5:      $i^{(j)} = \arg \max_i \|\varphi_i^T \mathbf{U}^{(j-1)}\|_2$ 
6:   else if (RA-ORMP) then
7:      $i^{(j)} = \arg \max_{i \notin \Omega^{(j-1)}} \|\varphi_i^T \mathbf{U}^{(j-1)}\|_2 / \|\mathbf{P}_{\Omega^{(j-1)}}^\perp \varphi_i\|_2$ 
8:   end if
9:    $\Omega^{(j)} = \Omega^{(j-1)} \cup i^{(j)}$ 
10:   $\mathbf{X}_{\Omega^{(j)},:}^{(j)} = \Phi_{\Omega^{(j)}}^\dagger \mathbf{Y}$ 
11:   $\mathbf{R}^{(j)} = \mathbf{Y} - \Phi \mathbf{X}^{(j)}$ 
12: end for

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matrix: $\mathbf{U}^{(j-1)} = \text{ortho}(\mathbf{R}^{(j-1)})$.² However, the recovery performance was shown to deteriorate even in the maximal rank scenario as the algorithm selected more coefficients. To compensate for this, the column normalization used in Order Recursive Matching Pursuit (ORMP) was included. Specifically at the start of the t th iteration, if we have a selected support set $\Omega^{(t)}$, a new column Φ_i is then chosen based upon the following selection rule:

$$i^{(t)} = \arg \max_i \frac{\|\varphi_i^T \mathbf{U}^{(t)}\|_2}{\|\mathbf{P}_{\Omega^{(t)}}^\perp \varphi_i\|_2}, \quad (1)$$

where $\mathbf{P}_{\Omega^{(t)}}^\perp$ denotes the orthogonal projection onto the null space of $\Phi_{\Omega^{(t)}}$. The righthand side of (1) measures the distance of the *normalized* vector $\mathbf{P}_{\Omega^{(t)}}^\perp \varphi_i / \|\mathbf{P}_{\Omega^{(t)}}^\perp \varphi_i\|$ from the subspace spanned by $\mathbf{U}^{(t)}$. This ensures that correct selection is maintained at each iteration in the maximal rank scenario. The full description of the RA-OMP and RA-ORMP are summarized in Algorithm 1.

In the next section the recovery guarantees for RA-ORMP and RA-OMP+SA-MUSIC (using RA-OMP to select the first $k-r$ indices followed by the subspace augmented SA-MUSIC of [7], [8]) are examined and shown to exploit the rank of \mathbf{Y} very effectively.

RA-OMP+SA-MUSIC is very similar to the approach considered in [7]. However in [7] only a single orthogonalization of \mathbf{Y} was performed followed by simultaneous OMP (SOMP). [8] considered SOMP+SA-MUSIC but without an initial orthogonalization.³ Both theoretical and empirical recovery performance of SOMP+SA-MUSIC is limited due to the “rank-blind” property of SOMP [6].

IV. SPARSE MMV RECOVERY BY RA-OMP

Correct selection by the RA-OMP algorithm at the j th iteration is characterized by the following quantity.

²In practice, significant complications arise in the noisy case where the estimation of the rank and signal subspace is non-trivial. For an analysis of related algorithms in the noisy case, see [14].

³While writing up this work we became aware of an updated version of [8] where the authors have switched to considering the RA-OMP+SA-MUSIC proposed here instead of SOMP+SA-MUSIC. The updated version also analyses the presence of noise in an asymptotic setting allowing the problem size to tend to ∞ .

Definition 1 (Greedy Selection Ratio for RA-OMP). In iteration j of RA-OMP, let $\Lambda = \text{supp}(\mathbf{X})$ and define the *greedy selection ratio for RA-OMP* as

$$\rho(j) = \frac{\max_{i' \in \Lambda^c} \|\varphi_{i'}^T \mathbf{U}^{(j)}\|_2}{\max_{i \in \Lambda} \|\varphi_i^T \mathbf{U}^{(j)}\|_2}. \quad (2)$$

The following observation is obvious.

Lemma 1. In iteration j , RA-OMP will correctly select an atom $\varphi_j, j \in \Lambda = \text{supp}(\mathbf{X})$ if and only if $\rho(j) < 1$.

To bound $\rho(j)$ we introduce the following lemmas.

Lemma 2. Suppose $\mathbf{U} \in \mathbb{R}^{m \times r}$ with columns $\mathbf{U}_i \in \text{range}(\Phi_\Lambda)$ and $\|\mathbf{U}_i\|_2 = 1$ for all i . Let $\alpha = \sigma_{\min}(\Phi_\Lambda)$ be the smallest singular value of Φ_Λ with $|\Lambda| = k < m$. Then

$$\max_{i \in \Lambda} \|\varphi_i^T \mathbf{U}\|_2 \geq \alpha \sqrt{\frac{r}{k}}. \quad (3)$$

Proof: We can write

$$\|\varphi_i^T \mathbf{U}\|_2^2 = \|\mathbf{e}_i^T \Phi_\Lambda^T \mathbf{U}\|_2^2 \quad (4)$$

where \mathbf{e}_i is the standard (dirac) basis in \mathbb{R}^k .

Define $\tilde{\mathbf{U}} = \Phi_\Lambda^T \mathbf{U}$. Since \mathbf{U}_i is in the range of Φ_Λ we have the following bound:

$$\|\tilde{\mathbf{U}}_i\|_2^2 \geq \sigma_{\min}(\Phi_\Lambda)^2 \|\mathbf{U}_i\|_2^2 = \alpha^2 \quad (5)$$

Hence

$$\begin{aligned} \max_{i \in \Lambda} \|\varphi_i^T \mathbf{U}\|_2^2 &= \max_{i \in \Lambda} \|\mathbf{e}_i^T \tilde{\mathbf{U}}\|_2^2 = \max_{i \in \Lambda} \sum_{l=1}^r [\tilde{\mathbf{U}}_{i,l}^2] \\ &\geq \text{mean}_i [\tilde{\mathbf{U}}_{i,l}^2] = \frac{1}{k} \sum_{i=1}^k \sum_{l=1}^r [\tilde{\mathbf{U}}_{i,l}^2] \quad (6) \\ &= \frac{1}{k} \sum_{l=1}^r \|\tilde{\mathbf{U}}_l\|_2^2 \geq \frac{r}{k} \alpha^2 \end{aligned}$$

where we have bounded the maximum by the mean and then swapped the order of summation. \square

Lemma 3. If $\Phi \in \mathbb{R}^{m \times n}$ with entries draw i.i.d. from $\mathcal{N}(0, m^{-1})$, $\Lambda \subset \{1, \dots, n\}$ is an index set with $|\Lambda| = k$, and $\mathbf{U} \in \mathbb{R}^{m \times r}$ is a matrix with orthonormal columns, $\text{rank}(\mathbf{U}) = r$ and with $\text{span}(\mathbf{U}) \subset \text{span}(\Phi_\Lambda)$, then

$$\mathbb{P}\{\max_{i \in \Lambda^c} \|\varphi_i^T \mathbf{U}\|_2^2 < \mu^2\} \geq 1 - (n - k)e^{-(m\mu^2 - 2r)/4}. \quad (7)$$

Proof: Let $z = \varphi_i^T \mathbf{U}$ then $z \in \mathbb{R}^r$ and for $i \notin \Lambda$ the entries in z follow the normal distribution $\mathcal{N}(0, m^{-1})$. We can now use the Laplace transform method [15] to bound $\|z\|_2$.

$$\begin{aligned} \mathbb{P}\{\|z\|_2^2 \geq \mu^2\} &\leq e^{-\lambda \mu^2} \mathbb{E}\{e^{\lambda \|z\|_2^2}\} \\ &= e^{-\lambda \mu^2 + \frac{r}{2} \ln\left(\frac{m}{m-2\lambda}\right)} \quad (8) \end{aligned}$$

for any $\lambda > 0$. Selecting $\lambda = m/4$ gives

$$\mathbb{P}(\|z\|_2^2 \geq \mu^2) \leq e^{-(m\mu^2 - 2r)/4}. \quad (9)$$

Applying the union bound completes the result. \square

We will also require that the residual matrix, $\mathbf{R}^{(j)}$, retains generic rank (equivalent to the row non-degeneracy condition

in [7]) which is given by:

Lemma 4. *Let $\Phi \in \mathbb{R}^{m \times n}$ with entries draw i.i.d. from $\mathcal{N}(0, m^{-1})$ and \mathbf{X} be a joint sparse matrix with support $\Lambda \subset \{1, \dots, n\}$, $|\Lambda| = k$, such that $\mathbf{X}_{\Lambda, \cdot}$ is in general position. If $\Omega^{(j)} \subset \Lambda$ for $j \leq k - r$, then $\text{rank}(\mathbf{R}^{(j)}) = r$.*

Proof: Note that:

$$\mathbf{R}^{(j)} = \mathbf{P}_{\Omega^{(j)}}^\perp \mathbf{Y} = \mathbf{P}_{\Omega^{(j)}}^\perp \Phi_{\Lambda - \Omega^{(j)}} \mathbf{X}_{\Lambda - \Omega^{(j)}, \cdot} \quad (10)$$

Since $\mathbf{X}_{\Lambda, \cdot}$ is in general position $\text{rank}(\mathbf{X}_{\Lambda - \Omega^{(j)}, \cdot}) = \min\{r, k - j\} = r$, and since Φ is i.i.d. Gaussian then $\mathbf{P}_{\Omega^{(j)}}^\perp \Phi_{\Lambda - \Omega^{(j)}}$ will have maximal rank with probability 1. Therefore $\text{rank}(\mathbf{R}^{(j)}) = r$. \square

We can combine the above lemmata to give:

Lemma 5. *Suppose that after $j < k - r$ iterations of RA-OMP, $\Omega^{(j)} \subset \Lambda = \text{supp}(\mathbf{X})$ with $|\Lambda| = k$. Define $\mathbf{R}^{(j)} = \mathbf{Y} - \Phi \mathbf{X}^{(j)}$ and assume \mathbf{X} is in general position and $\Phi \in \mathbb{R}^{m \times n}$ with entries drawn i.i.d. from $\mathcal{N}(0, m^{-1})$. Then, for $\alpha = \sigma_{\min}(\Phi_\Lambda)$,*

$$\mathbb{P}\{\Omega^{(j+1)} \subset \Lambda\} \geq 1 - (n - k)e^{-(m\alpha^2 \frac{k}{r} - 2r)/4}. \quad (11)$$

Proof: Let $\mathbf{U}^{(j)} = \text{ortho}(\mathbf{R}^{(j)})$ then from Lem. 4 we have $\text{rank}(\mathbf{U}^{(j)}) = \text{rank}(\mathbf{R}^{(j)}) = \text{rank}(\mathbf{X}) = r$. Now, Lem. 1 and the assumption that $\Omega^{(j)} \subset \Lambda$ allow us to rewrite the probability statement as

$$\mathbb{P}\{\Omega^{(j+1)} \subset \Lambda\} = \mathbb{P}\{\rho(j+1) < 1\}.$$

Lemmas 2 and 3 with $\mu^2 = \alpha^2 r/k$ combine to show that

$$\begin{aligned} \mathbb{P}\{\rho(j+1) < 1\} &= \mathbb{P}\left\{\max_{j' \in \Lambda^c} \|\varphi_{j'}^T \mathbf{U}\|_2 < \max_{j \in \Lambda} \|\varphi_j^T \mathbf{U}\|_2\right\} \\ &\geq \mathbb{P}\left\{\max_{j' \in \Lambda^c} \|\varphi_{j'}^T \mathbf{U}\|_2^2 < \frac{\alpha^2 r}{k}\right\} \\ &\geq 1 - (n - k)e^{-(m\alpha^2 \frac{r}{k} - 2r)/4}. \end{aligned}$$

\square

We can now state our main theorem for RA-OMP.

Theorem 6 (RA-OMP + SA-MUSIC recovery). *Assume $\mathbf{X} \in \mathbb{R}^{n \times r}$, $\text{supp}(\mathbf{X}) = \Lambda$, $|\Lambda| = k > r$ with \mathbf{X}_Λ in general position and let Φ be a random matrix, independent of \mathbf{X} , with i.i.d. entries $\Phi_{i,j} \sim \mathcal{N}(0, m^{-1})$. Then, for some C and with probability greater than $1 - \delta$, RA-OMP + SA-MUSIC will recover \mathbf{X} from $\mathbf{Y} = \Phi \mathbf{X}$ if:*

$$m \geq Ck\left(\frac{1}{r} \log(n/\sqrt{\delta}) + 1\right). \quad (12)$$

Proof: It is sufficient to bound the probability of making $q \leq k - r$ successive correct selections after which SA-MUSIC [7] is guaranteed to recover the remaining coefficients [7], [8]. Suppose $\sigma_{\min}(\Phi_\Lambda) = \alpha$, then

$$\begin{aligned} \mathbb{P}\{\max_{t \leq q} \rho(t) < 1\} &\geq \prod_{t \leq q} \mathbb{P}\{\rho(t) < 1\} \\ &\geq \left[1 - (n - k)e^{-(m\alpha^2 r/k - 2r)/4}\right]^q \quad (13) \\ &\geq 1 - q(n - k)e^{-(m\alpha^2 r/k - 2r)/4}. \end{aligned}$$

Now recall [11] that $\mathbb{P}\{\sigma_{\min}(\Phi_\Lambda) > 0.5\} \geq 1 - e^{-cm}$ for

some $c > 0$. Hence:

$$\begin{aligned} \mathbb{P}\{\max_{t \leq q} \rho(t) < 1\} &\geq \left(1 - n^2 e^{-(m\alpha^2 r/k - 2r)/4}\right) \left(1 - e^{-cm}\right) \\ &\geq 1 - n^2 e^{-(mr/k - 8r)/16} - e^{-cm} \\ &\geq 1 - n^2 e^{-C(mr/k - 8r)} \end{aligned} \quad (14)$$

where we have used the fact that $q(n - k) < n^2$. Now choosing $\delta \geq n^2 e^{-C(mr/k - 8r)}$ and rearranging gives (12). \square

V. SPARSE MMV RECOVERY BY RA-ORMP

As RA-OMP and RA-ORMP only differ in the selection step we can use similar arguments to above, with additional control of the normalization term $\|P_{\Omega^{(j)}}^\perp \varphi_i\|_2$ given by the following.

Lemma 7. *If $\Phi \in \mathbb{R}^{m \times n}$ with i.i.d. entries $\Phi_{i,j} \sim \mathcal{N}(0, m^{-1})$, $\Lambda \subset \{1, \dots, n\}$, $|\Lambda| = k$, and $\Omega^{(j)} \subset \Lambda$, $|\Omega^{(j)}| = j$, then for $i \notin \Omega^{(j)}$ and $m \geq 2j$ we have:*

$$\mathbb{P}\left\{\|P_{\Omega^{(j)}}^\perp \varphi_i\|_2^2 \geq \frac{1}{4}\right\} \geq 1 - e^{-m/32} \quad (15)$$

Proof: Since $P_{\Omega^{(j)}}^\perp$ and φ_i are independent, $z := P_{\Omega^{(j)}}^\perp \varphi_i$ is a Gaussian random vector within the $m - j$ dimensional subspace $\text{Null}(\Phi_{\Omega^{(j)}}^T)$ whose entries have variance m^{-1} . Hence we can use the following concentration of measure bound [15]:

$$\mathbb{P}\left\{\|z\|_2^2 \geq (1 - \epsilon) \frac{m - j}{m}\right\} \geq 1 - e^{-\epsilon^2(m - j)/4} \quad (16)$$

Selecting $\epsilon = 1/2$ and noting that by assumption $m - j \geq m/2$ gives the required result. \square

Lemma 8. *If $\Phi \in \mathbb{R}^{m \times n}$ with i.i.d. entries $\Phi_{i,j} \sim \mathcal{N}(0, m^{-1})$, $\Lambda \subset \{1, \dots, n\}$, $|\Lambda| = k$, and $\Omega^{(j)} \subset \Lambda$, $|\Omega^{(j)}| = j$, then for $i \notin \Omega^{(j)}$ and $m \geq 2j$ we have:*

$$\mathbb{P}\left\{\|P_{\Omega^{(j)}}^\perp \varphi_i\|_2^2 \leq 2\right\} \geq 1 - e^{-m/16} \quad (17)$$

Proof: Standard Gaussian bounds [15] for φ_i give:

$$\mathbb{P}\left\{\|\varphi_i\|_2^2 \leq (1 - \epsilon)^{-1}\right\} \geq 1 - e^{-\epsilon^2 m/4} \quad (18)$$

Selecting $\epsilon = 1/2$ and using $\|P_{\Omega^{(j)}}^\perp \varphi_i\|_2 < \|\varphi_i\|_2$ gives the required result. \square

Theorem 9 (RA-ORMP recovery). *Assume $\mathbf{X} \in \mathbb{R}^{n \times r}$, $\text{supp}(\mathbf{X}) = \Lambda$, $|\Lambda| = k > r$ with \mathbf{X}_Λ in general position and let Φ be a random matrix, independent of \mathbf{X} , with i.i.d. entries $\Phi_{i,j} \sim \mathcal{N}(0, m^{-1})$. Then, for some C and with probability greater than $1 - \delta$, RA-ORMP will recover \mathbf{X} from $\mathbf{Y} = \Phi \mathbf{X}$ if m satisfies (12).*

Proof: We first note that if $\Omega^{(j)} \subset \Lambda$, $j < k - r$ and $\alpha = \sigma_{\min}(\Phi_\Lambda)$, then

$$\begin{aligned} &\mathbb{P}\{\Omega^{(j+1)} \subset \Lambda\} \\ &= \mathbb{P}\left\{\max_{i' \in \Lambda^c} \frac{\|\varphi_{i'}^T \mathbf{U}^{(j)}\|_2^2}{\|P_{\Omega^{(j)}}^\perp \varphi_{i'}\|_2^2} \leq \max_{i \in \Lambda} \frac{\|\varphi_i^T \mathbf{U}^{(j)}\|_2^2}{\|P_{\Omega^{(j)}}^\perp \varphi_i\|_2^2}\right\} \end{aligned} \quad (19)$$

Now using the bounds from Lemma 3 with $\mu^2 = \alpha^2 r/8k$ and

Lemmas 7 and 8 along with the union bound gives:

$$\begin{aligned} \mathbb{P}\{\Omega^{(j+1)} \subset \Lambda\} &\geq (1 - (n - k)e^{-(m\alpha^2 \frac{r}{k} - 2r)/4}) \\ &\times (1 - (n - k)e^{-m/32})(1 - me^{-m/16}) \\ &\geq 1 - (n - k) \left(e^{-(m\alpha^2 \frac{r}{k} - 2r)/4} + e^{-m/32} + e^{-m/16} \right) \\ &\geq 1 - (n - k)e^{-C(m\alpha^2 \frac{r}{k} - 2r)} \end{aligned} \quad (20)$$

for an appropriately chosen C . To complete the proof for the correct selection of $\Omega^{(j)}$ for all $j < k - r$ we can again apply the union bound and remove the dependence on α as in (13) and (14) above. We leave the details to the reader.

For the selection of the remaining coefficients we note that for $j \geq k - r$ the original RA-ORMP task is equivalent to solving $\mathbf{R}^{(j)} = [\mathbf{P}_{\Omega^{(j)}}^\perp \Phi_{\Lambda - \Omega^{(j)}}] \mathbf{X}_{\Lambda - \Omega^{(j)} \cdot}$ using RA-ORMP. However by our assumptions on \mathbf{X} , $\text{rank}(\mathbf{X}_{\Lambda - \Omega^{(j)} \cdot}) = r = |\Lambda - \Omega^{(j)}|$ therefore we are in the maximal rank scenario and from [6] we have guaranteed recovery by RA-ORMP. \square

VI. NUMERICAL RESULTS

Here we demonstrate empirically that the $(\log n)/r$ term in our recovery result appears to accurately capture the effect of rank on the recovery performance. To this end we performed a number of experiments using Gaussian random matrices for both Φ and \mathbf{X}_Λ . The parameters m and k were held fixed, first at $m = 3k/2 = 30$ and then with $m = 2k = 40$. We then varied the number of channels of \mathbf{X} from $r = 1, \dots, 15$ and n in powers of 2 from 64 to 4096. For each set of $\{k, r, m, n\}$ we performed 100 trials and calculated the empirical probability of recovery. The included simulations are restricted to the noiseless case; the algorithms demonstrate robustness to noise similar to that observed in [6], [7], [8].

Figure 1 shows the recovery plots for the recovery algorithms RA-OMP + MUSIC and RA-ORMP. In each case the “phase transition” appears to exhibit an approximate linear dependency between r and $\log n$ as highlighted by the red lines (note the lines are constrained to pass through the origin).

VII. CONCLUSION

Our theoretical results predict that the rank of the coefficient matrix in the noiseless sparse MMV recovery problem can be successfully exploited in RA-OMP+SA-MUSIC and RA-ORMP to enable joint sparse recovery when $m \gtrsim Ck((\log n)/r + 1)$ using a Gaussian measurement matrix, although no attempt has been made to optimize the constant C . This removes the $\log n$ penalty term that is observed in OMP when there is only a modest number of multiple measurement vectors $r \gtrsim \log n$. Numerical experiments suggest that this form may reasonably characterize the recovery behaviour in practice. Empirically the RA-ORMP algorithm appears to perform slightly better than RA-OMP+SA-MUSIC. However this comes with an additional computational expense.

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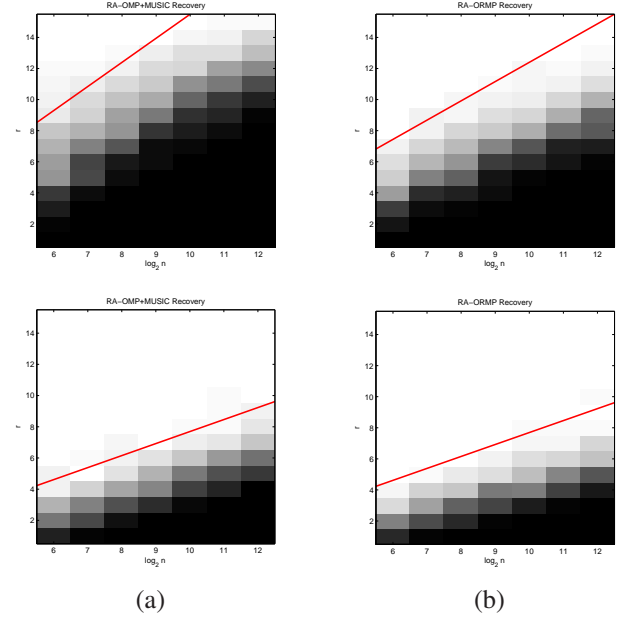


Fig. 1. Sparse MMV recovery plots showing the “phase transitions” for RA-OMP+SA-MUSIC (a) and RA-ORMP (b) with $m = 3k/2 = 30$ (top) and $m = 2k = 40$ (bottom) while varying the size of the dictionary $n = 64, 128, \dots, 4096$ and number of channels, $r = 1, 2, \dots, 15$. The red line indicates a linear relation between r and $\log n$.

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