

Relations Between Redundancy Patterns of the Shannon Code and Wave Diffraction Patterns of Partially Disordered Media

Neri Merhav

Department of Electrical Engineering
Technion - Israel Institute of Technology
Haifa 32000, ISRAEL

Abstract

The average redundancy of the Shannon code, R_n , as a function of the block length n , is known to exhibit two very different types of behavior, depending on the rationality or irrationality of certain parameters of the source: It either converges to $1/2$ as n grows without bound, or it may have a non-vanishing, oscillatory, (quasi-) periodic pattern around the value $1/2$ for all large n . In this paper, we make an attempt to shed some insight into this erratic behavior of R_n , by drawing an analogy with the realm of physics of wave propagation, in particular, the elementary theory of scattering and diffraction. It turns out that there are two types of behavior of wave diffraction patterns formed by crystals, which are correspondingly analogous to the two types of patterns of R_n . When the crystal is perfect, the diffraction intensity spectrum exhibits very sharp peaks, a.k.a. Bragg peaks, at wavelengths of full constructive interference. These wavelengths correspond to the frequencies of the harmonic waves of the oscillatory mode of R_n . On the other hand, when the crystal is imperfect and there is a considerable degree of disorder in its structure, the Bragg peaks disappear, and the behavior of this mode is analogous to the one where R_n is convergent.

Index Terms: Lossless source coding, redundancy, Shannon code, scattering, diffraction, Bragg peaks, disorder.

1 Introduction

The analysis of the average redundancy of lossless codes for data compression schemes is a topic that attracted the attention of considerably many researchers throughout the history of Information Theory (cf. e.g., [1],[3],[6],[7],[8],[9],[10],[11],[12],[13] and many references therein).

In [13] Szpankowski has derived the asymptotic behavior of the average redundancy R_n , as a function of the block length n , for the Shannon code, the Huffman code, and other codes, focusing

primarily on the binary memoryless source, parametrized by p – the probability of zero. His analysis revealed a rather interesting behavior of R_n , especially in the cases of the Shannon code and the Huffman code: When $\alpha \triangleq \log_2[(1-p)/p]$ is irrational, then R_n converges to a constant (which is $1/2$ for the Shannon code and $3/2 - 1/\ln 2$ for the Huffman code) as $n \rightarrow \infty$. On the other hand, when α is rational, R_n has a non-vanishing oscillatory term of the form $\langle \beta m_0 n \rangle$, where $\beta \triangleq -\log_2(1-p)$, m_0 is the denominator of $\alpha = \ell_0/m_0$ in its representation as the ratio between two integers whose greatest common divisor is 1, and $\langle x \rangle = x - [x]$ designates the fractional part of a real number x . In several places in his paper, Szpankowski describes this behavior of R_n as “erratic” and this qualifier is, of course, understandable.

Our purpose in this paper is to make an attempt to give some insight into this erratic behavior of R_n by drawing an analogy with the physics of wave diffraction. From the theory of X-ray scattering (see, e.g., [2, Chapter 2],[14]), it is known that if the object that causes the diffraction of an incident wave is a perfect crystal, then the intensity profile of the scattered wave (as a function of the wavelength or the wave number) exhibits very sharp peaks, known as *Bragg peaks*, at wavelengths that correspond to full coherence, where the optical distance differences to all scattering elements (layers of the crystal) are exactly integer multiples of the wavelength. This continues to be the case as long as there is enough order in the medium such that all these distances are commensurable and therefore have a common divisor (common unit of length), which can serve as the fundamental wavelength. In the realm of the average redundancy analysis, this corresponds to the case where α is rational and the fundamental frequency of the oscillatory term $\langle \beta m_0 n \rangle$ of R_n is intimately related to the fundamental wavelength at which there is a Bragg peak. On the other hand, when the distances are incommensurable, perfect coherence between all scattered waves is not achieved at any wavelength and therefore no Bragg peaks are observed. This is the case of strong disorder, which in the lossless source coding problem, corresponds to the case of α irrational, where R_n is convergent.

More concretely, the analysis of the scattered wave intensity function is based on a very simple model of disorder, which is due to Hendricks and Teller [5] (see also [4]). According to the Hendricks–Teller (HT) model, the distances between every two consecutive layers in the solid are selected independently at random from a finite set of two or more distances. In the simplest case, where there are only two possible distances d_0 and d_1 , with probabilities p and $1-p$, this random selection

process is analogous to the memoryless binary source of the data compression problem and the parameter α of this source plays a role analogous to that of the ratio d_1/d_0 . Thus, α irrational means that d_0 and d_1 are incommensurable, which is the case of strong disorder with no Bragg peaks and no oscillations in R_n . On the other hand when $\alpha = d_1/d_0$ is rational, we are in the (partially) ordered mode, as described above.

From the pure mathematical point of view, the analogy between the average redundancy problem and the diffraction problem is rooted in that at the heart of the analyzes of both problems, there is one very simple mathematical fact in common: Given a vector $(p_0, p_1, \dots, p_{M-1})$ of non-negative reals summing to unity (probabilities) and a vector $(\alpha_1, \dots, \alpha_{M-1}) \in \mathbb{R}^{M-1}$, the complex number

$$C_m = p_0 + \sum_{j=1}^{M-1} p_j e^{2\pi i m \alpha_j}, \quad i = \sqrt{-1}, \quad m = 1, 2, 3, \dots \quad (1)$$

has a modulus that obviously never exceeds unity, and $C_m = 1$ (i.e., full coherence between all M phasors) is attained for some integer values of m if and only if $\{\alpha_j\}$ are all rational. When this is the case, then $C_m = 1$ for all values of m which are integer multiples of m_0 , the first positive integer m for which $m\alpha_j$ is integer for all $1 \leq j \leq M-1$ at the same time.¹ The analogy between the Shannon code redundancy analysis and the diffraction patterns under the HT model will center around (1) and its two types of behavior depending on the rationality or irrationality of $\{\alpha_j\}$.

The remaining part of this short paper consists of two more main sections. For the sake of completeness, in Section 2, we summarize the main ingredients of the derivation in [13] (with a few shortcuts), emphasizing the use of the simple mathematical fact described in the previous paragraph. For reasons of simplicity, we focus on the Shannon code and the derivation specializes on the memoryless case. In Section 3, we bring the derivation of the diffraction patterns of the HT model, with a focus on the analogy with Section 2. We then describe in detail the mapping between the two problems under discussion. Finally, in Section 4 we summarize and conclude, with a comments on a possible extension to the Markov case.

¹The previous paragraph refers to the special case $M = 2$.

2 Average Redundancy of the Shannon Code

Throughout the remaining part of this paper, we use capital letters to designate random variables (e.g., X_i) and the corresponding lower case letters to denote specific realizations (e.g., x_i).

Consider a finite alphabet memoryless source X_1, X_2, \dots with alphabet $\mathcal{X} = \{0, 1, 2, \dots, M - 1\}$ and symbol probabilities $\{p_0, p_1, \dots, p_{M-1}\}$. The Shannon code for lossless data compression assigns to every source n -tuple $\mathbf{x} = (x_1, x_2, \dots, x_n) \in \mathcal{X}^n$ a binary codeword of length

$$\ell(\mathbf{x}) = \lceil -\log P(\mathbf{x}) \rceil = \lceil -\log \prod_{t=1}^n p_{x_t} \rceil, \quad (2)$$

where $\lceil u \rceil$ designates the smallest integer not smaller than u . The average redundancy of the Shannon code is defined as

$$R_n = \mathbf{E} \{ \ell(\mathbf{X}) \} - nH \quad (3)$$

where

$$H = - \sum_{j=0}^{M-1} p_j \log p_j \quad (4)$$

is the per-symbol entropy. The derivation of the asymptotic expression for R_n in [13] can be presented (with a few slight shortcuts and modifications) as follows. By using the Fourier series expansion of the function $\langle u \rangle$, according to

$$\langle u \rangle = \frac{1}{2} - \sum_{m \neq 0} a_m e^{2\pi i m u}, \quad a_m = \frac{1}{2\pi i m}, \quad (5)$$

we have the following:

$$\begin{aligned} R_n &= \mathbf{E} \{ \lceil -\log P(\mathbf{X}) \rceil + \log P(\mathbf{X}) \} \\ &= 1 - \mathbf{E} \{ -\log P(\mathbf{X}) - \lceil -\log P(\mathbf{X}) \rceil \} \\ &= 1 - \mathbf{E} \langle -\log P(\mathbf{X}) \rangle \\ &= 1 - \mathbf{E} \left\{ \frac{1}{2} - \sum_{m \neq 0} a_m \exp[-2\pi i m \log P(\mathbf{X})] \right\} \\ &= \frac{1}{2} + \sum_{m \neq 0} a_m \mathbf{E} \{ \exp[-2\pi i m \log P(\mathbf{X})] \} \\ &= \frac{1}{2} + \sum_{m \neq 0} a_m \sum_{\mathbf{x} \in \mathcal{X}^n} \left(\prod_{t=1}^n p_{x_t} \right) \cdot \exp \left[-2\pi i m \sum_t \log p_{x_t} \right] \end{aligned}$$

$$\begin{aligned}
&= \frac{1}{2} + \sum_{m \neq 0} a_m \sum_{\mathbf{x} \in \mathcal{X}^n} \prod_{t=1}^n (p_{x_t} \exp[-2\pi i m \log p_{x_t}]) \\
&= \frac{1}{2} + \sum_{m \neq 0} a_m \prod_{t=1}^n \left(\sum_{x_t=0}^{M-1} p_{x_t} \exp[-2\pi i m \log p_{x_t}] \right) \\
&= \frac{1}{2} + \sum_{m \neq 0} a_m \left(\sum_{j=0}^{M-1} p_j \exp[-2\pi i m \log p_j] \right)^n \\
&= \frac{1}{2} + \sum_{m \neq 0} a_m e^{-2\pi i m n \log p_0} \left[p_0 + \sum_{j=1}^{M-1} p_j \exp\{2\pi i m \log(p_0/p_j)\} \right]^n. \tag{6}
\end{aligned}$$

Denoting $\alpha_j = \log(p_0/p_j)$, $j = 1, 2, \dots, M-1$, the expression in the square brackets is exactly C_m as was defined in (1). The behavior of R_n for large n is then as follows. If $\{\alpha_j\}$ are not all rational, then $|C_m| < 1$ for all m , and so, $\lim_{n \rightarrow \infty} C_m^n = 0$, which causes the entire summation over m to vanish for large n . In this case, $R_n \rightarrow 1/2$ as $n \rightarrow \infty$. On the other hand, if $\{\alpha_j\}$ are all rational, then there exists an integer m such that $m\alpha_j$ are all integers. Let m_0 be the smallest positive integer with this property. Then all other integers with the same property are integer multiples of m_0 . Consequently, $\lim_{n \rightarrow \infty} C_m^n = 1$ whenever m is an integer multiple of m_0 and $\lim_{n \rightarrow \infty} C_m^n = 0$ otherwise. Thus, denoting $\beta = -\log p_0$, we now have for large n ,

$$\begin{aligned}
R_n &\approx \frac{1}{2} + \sum_{k \neq 0} a_{km_0} e^{2\pi i k m_0 n \beta} \\
&= \frac{1}{2} + \frac{1}{m_0} \sum_{k \neq 0} a_k e^{2\pi i k m_0 n \beta} \\
&= \frac{1}{2} + \frac{1}{m_0} \left(\frac{1}{2} - \langle \beta m_0 n \rangle \right), \tag{7}
\end{aligned}$$

where the second line holds since a_m is inversely proportional to m (see (5) above) and in the third line we used again (5) with $u = \beta m_0 n$. As can easily be seen from the second line of (7), for large n , the sequence R_n is harmonic with a fundamental frequency $\omega_0 = 2\pi m_0 \beta$. In other words, the Fourier transform of $\{R_n\}$ contains Dirac delta functions at integer multiples of ω_0 (modulo 2π). We will see later on that these spectral spikes are analogous to the Bragg peaks of the HT model.

At this point, a technical comment is in order. At first glance, it may seem that the above approximate expression of R_n is asymmetric with respect to permutations of the alphabet, because β was defined as $-\log p_0$ and the choice of the symbol $x = 0$ as having a special role in the last line of (6) was completely arbitrary (we could have chosen, of course, any other symbol j as well).

However, note that $\langle \beta m_0 n \rangle = \langle -m_0 n \log p_0 \rangle$ is identical to $\langle -m_0 n \log p_j \rangle$ for all $j = 1, \dots, M - 1$ because in the rational case considered above, the numbers $\{-m_0 n \log p_j\}_{j=0}^{M-1}$ differ from each other by integers, and therefore their fractional parts are all the same. Thus, the above expression of R_n is, in fact, invariant to permutations of the alphabet.

3 Diffraction Patterns of the HT Model

The simplest way to think of the HT model is as a one-dimensional model of an alloy, which is characterized by a sequence of mass points, positioned along the real line at random locations Z_0, Z_1, \dots, Z_{n-1} . The ensemble of the HT model is defined in terms of the spacings $\Delta_j \triangleq Z_j - Z_{j-1}$, $j = 1, 2, \dots, n - 1$, which are $n - 1$ i.i.d. random variables taking on values in a finite set $\{d_0, d_1, \dots, d_{M-1}\}$ with probabilities p_0, p_1, \dots, p_{M-1} , respectively (thus, Z_0, Z_1, \dots is a random walk). Each point Z_i contributes a scattered wave described by the phasor e^{-iqZ_j} , where in the one-dimensional setting considered here, q can be understood as the wave number, that is, $q = 2\pi/\lambda$, where λ is the wavelength. Assuming the same amplitudes at all points, the superposition of all these contributions is then the sum $U(q) = \sum_j e^{-iqZ_j}$, which can be interpreted as the Fourier transform of the function $u(z) = \sum_j \delta(z - Z_j)$. The overall intensity of this superposition of waves is designated by the *structure function* [2, Chapter 2]

$$I(q) = \mathbf{E}\{|U(q)|^2\} = \mathbf{E}\left\{\sum_{k,\ell} e^{iq(Z_k - Z_\ell)}\right\} = \sum_{k,\ell} \mathbf{E}\{e^{iq(Z_k - Z_\ell)}\}, \quad (8)$$

where the expectation is with respect to the random variables $\{Z_j\}$.

The derivation of $I(q)$ is fairly simple (see, e.g., [4]) and it is brought here for the sake of completeness.

$$\begin{aligned} I(q) &= \sum_{k,\ell} \mathbf{E}\{e^{iq(Z_k - Z_\ell)}\} \\ &= n + \sum_{k>\ell} \mathbf{E}\{e^{iq(Z_k - Z_\ell)}\} + \sum_{k<\ell} \mathbf{E}\{e^{iq(Z_k - Z_\ell)}\} \\ &= n + \sum_{k>\ell} \mathbf{E}\{e^{iq(Z_k - Z_\ell)}\} + \sum_{k>\ell} \mathbf{E}\{e^{-iq(Z_k - Z_\ell)}\} \\ &\triangleq n + I_0(q) + I_0^*(q) \end{aligned} \quad (9)$$

where $I_0(q)$ is defined as the second term of the third line and $I_0^*(q)$ is the complex conjugate of

$I_0(q)$. Now,

$$\begin{aligned}
I_0(q) &= \sum_{k>\ell} \mathbf{E}\{e^{iq(Z_k - Z_\ell)}\} \\
&= \sum_{k>\ell} \mathbf{E}\left\{\exp\left[iq \sum_{s=\ell+1}^k \Delta_s\right]\right\} \\
&= \sum_{k>\ell} \mathbf{E}\left\{\prod_{s=\ell+1}^k \exp[iq\Delta_s]\right\} \\
&= \sum_{k>\ell} \left[\sum_{j=0}^{M-1} p_j e^{iqd_j}\right]^{k-\ell} \\
&= \sum_{r=1}^{n-1} (n-r)[C(q)]^r, \tag{10}
\end{aligned}$$

where we have denoted

$$C(q) = \sum_{j=0}^{M-1} p_j e^{iqd_j}. \tag{11}$$

For n large, whenever $|C(q)| < 1$, the last expression is dominated by the term $n \sum_{r=1}^{\infty} [C(q)]^r = nC(q)/[1 - C(q)]$, which together with the two other terms of (9), yields

$$I(q) \approx n \left(1 + \frac{C(q)}{1 - C(q)} + \frac{C^*(q)}{1 - C^*(q)}\right) = n \cdot \frac{1 - |C(q)|^2}{|1 - C(q)|^2}, \tag{12}$$

or equivalently,

$$\hat{I}(q) = \lim_{n \rightarrow \infty} \frac{I(q)}{n} = \frac{1 - |C(q)|^2}{|1 - C(q)|^2}. \tag{13}$$

If there are values of q for which $|C(q)| = 1$, yet $C(q) \neq 1$, then the geometric series diverges at these points, but these are only points of removable discontinuity in $\hat{I}(q)$ because for every other point, arbitrarily close to such a discontinuity point, again $|C(q)| < 1$, and the geometric series converges. The real problematic points, if any, are those where $C(q) = 1$ if they exist. For $C(q) = 1$, we have to re-derive the expression of $I(q)$ separately, which is very simple as $I(q)$ is just the sum of n^2 1's, namely, $I(q) = n^2$. In other words, the intensity scales quadratically rather than linearly with n , which means that these are extremely high peaks in $I(q)$, namely, the Bragg peaks.

For $C(q)$ to take the value 1 for some q , the products qd_j must all be integer multiples of 2π . Suppose that q is such that $qd_0 = 2\pi m$ for some integer m , i.e., $q = q_m \triangleq 2\pi m/d_0$, in which case we shall denote $C(q_m)$ by C_m , as before. In this case,

$$C_m = p_0 + \sum_{j=1}^{M-1} p_j e^{2\pi i m d_j / d_0}. \tag{14}$$

But this is again exactly the expression in (1), this time with $\alpha_j = d_j/d_0$, which as mentioned earlier, may assume the value 1, for some integer values of m , if and only if $\alpha_j = d_j/d_0$ are all rational, or equivalently, d_0, d_1, \dots, d_{M-1} are commensurable. When this is the case, then as before, there exists an integer m for which md_j/d_0 are all integers simultaneously. Analogously to the derivation in Section 2, let m_0 be the smallest integer with this property. Then, the Bragg peaks appear at wave-numbers q_{km_0} , $k = 1, 2, \dots$, which correspond to wavelengths λ_0/k , where $\lambda_0 = d_0/m_0$.

The analogy between the two settings is now clear: The memoryless source of Section 2 is parallel to the random selection process in the HT model. The parameters $\alpha_j = \log(p_0/p_j)$ of the source are analogous to distance ratios d_j/d_0 , $j = 1, 2, \dots, M - 1$. Their rationality/irrationality dictates the mode of behavior in both problems. The integer parameter m_0 is then defined in both settings in the very same way. The partially ordered mode in the diffraction model is parallel to the oscillatory mode of R_n in the data compression problem, and the Bragg peaks at all harmonics of the fundamental wave-number $q_{m_0} = 2\pi m_0/d_0$ correspond to all harmonics of the fundamental frequency $\omega_0 = 2\pi\beta m_0$ in the oscillatory component of R_n . In other words, the parameter β is conjugate, in this sense, to $1/d_0$.

4 Conclusion

In this short paper, we have made an attempt to provide some insight into the erratic behavior of the redundancy pattern of the Shannon code for lossless data compression. The insight we propose is rooted in the physical point of view, where the two modes of the behavior of the redundancy patterns are respectively analogous to partial order and complete disorder of a wave diffraction medium, which dictates the existence or non-existence of Bragg peaks pertaining to perfectly constructive interference. It is hoped that this physical insight contributes to the intuitive understanding of the redundancy of the Shannon code and perhaps other codes as well.

Finally, we comment that the above analyses are, in principle, generalizable to the finite-state Markov case (and indeed, Markov models have been proposed in the diffraction setting too [5],[14]). When it comes to the Markov case, then both in the data compression problem and in the HT model, the role played by high powers of C_m is essentially replaced by high powers of state transition probability matrix whose entries are weighted by the appropriate complex exponentials

(which depend on m). What matters then are the eigenvalues of this matrix. More concretely, it is not difficult to see that the spectral radius, in both settings, never exceeds unity. In the data compression problem, the critical behavior is dictated by the existence or non-existence of integer values $\{m\}$ for which the spectral radius is exactly 1. When such values of m exist, then R_n has an oscillatory behavior. In the diffraction problem, the distinction between the two types of behavior is dictated by the existence of values of m for which one of the eigenvalues is exactly equal to one.

References

- [1] R. M. Capocelli and A. De Santis, "On the redundancy of optimal codes with limited word length," *IEEE Trans. Inform. Theory*, vol. IT-38, no. 2, pp. 439–445, March 1992.
- [2] P. M. Chaikin and T. C. Lubensky, *Principles of Condensed Matter Physics*, Cambridge University Press, Cambridge England, 1995.
- [3] R. G. Gallager, "Variations on the theme by Huffman," *IEEE Trans. Inform. Theory*, vol. IT-24, no. 6, pp. 668–674, November 1978.
- [4] A. Garg and D. Levine, "Speckle in the diffraction patterns of Hendricks–Teller and icosahedral glass models," *Physical Review Letters*, vol. 60, no. 21, pp. 2160–2163, 23 May 1988.
- [5] S. Hendricks and E. Teller, "X-ray interference in partially ordered layer lattices," *The Journal of Chemical Physics*, vol. 10, no. 3, pp. 147–167, March 1942.
- [6] D. A. Huffman, "A method for the construction of minimum-redundancy codes," *Proc. IRE*, vol. 40, pp. 1098–1101, 1952.
- [7] P. Jackuet and W. Szpankowski, "Asymptotic behavior of the Lempel–Ziv parsing scheme and digital search trees," *Theoretical Computer Science*, vol. 144, pp. 161–197, 1995.
- [8] R. E. Krichevskii, "The relation between redundancy coding and the reliability of information from a source," *Problems of Information Transmission (IPPI)*, vol. 4, no. 3, pp. 37–45, 1968.
- [9] G. Louchard and W. Szpankowski, "Average redundancy of the Lempel–Ziv code," *IEEE Trans. Inform. Theory*, vol. 43, no. 1, pp. 2–8, January 1997.

- [10] J. Rissanen, “Complexity of strings in the class of Markov sources,” *IEEE Trans. Inform. Theory*, vol. IT-32, no. 4, pp. 526–532, July 1986.
- [11] S. A. Savari, “Variable-to-fixed length codes for predictable sources,” *Proc. Data Compression Conference (DCC)*, Snowbird, UT, pp. 481–490, 1998.
- [12] S. A. Savari and R. G. Gallager, “Generalized Tunstall codes for sources with memory,” *IEEE Trans. Inform. Theory*, vol. 43, no. 2, pp. 658–668, March 1997.
- [13] W. Szpankowski, “Asymptotic average redundancy of Huffman (and other) block codes,” *IEEE Trans. Inform. Theory*, vol. 46, no. 7, pp. 2434–2443, November 2000.
- [14] T. R. Welberry, “Diffuse X-ray scattering and models of disorder,” *Rep. Prog. Phys.*, vol. 48, pp. 1543–1593, 1985.