Binary linear codes with few weights from two-to-one functions

Kangquan Li, Chunlei Li, Tor Helleseth and Longjiang Qu

Abstract

In this paper, we apply two-to-one functions over \mathbb{F}_{2^n} in two generic constructions of binary linear codes. We consider two-to-one functions in two forms: (1) generalized quadratic functions; and (2) $\left(x^{2^t} + x\right)^e$ with gcd(t, n) = 1 and $gcd(e, 2^n - 1) = 1$. Based on the study of the Walsh transforms of those functions or their related-ones, we present many classes of linear codes with few nonzero weights, including one weight, three weights, four weights and five weights. The weight distributions of the proposed codes with one weight and with three weights are determined. In addition, we discuss the minimum distance of the dual of the constructed codes and show that some of them achieve the sphere packing bound. Moreover, several examples show that some of our codes are optimal and some have the best known parameters.

Index Terms

Binary linear codes, two-to-one functions, 3-weight linear codes, constant-weight linear codes

1. INTRODUCTION

Let q be a power of a prime p, \mathbb{F}_q be the finite field of q elements and \mathbb{F}_q^* be its multiplicative group. An [n, k, d] linear code \mathcal{C} over \mathbb{F}_q is a k-dimensional subspace of \mathbb{F}_q^n with minimum (Hamming) distance d. An [n, k, d] code is sometimes usually said to be optimal when its minimum distance d achieves the maximum value with respect to the Hamming bound [19]. Given an [n, k, d] linear code \mathcal{C} over \mathbb{F}_q , the dual code of \mathcal{C} is an [n, n-k] linear code defined by $\mathcal{C}^{\perp} = \{\mathbf{x} \in \mathbb{F}_q^n : \mathbf{x} \cdot \mathbf{c} = 0, \forall \mathbf{c} \in \mathcal{C}\}$, where $\mathbf{x} \cdot \mathbf{c} = \sum_{i=1}^n x_i c_i$ is a scalar product. Let A_i denote the number of codewords with Hamming weight i in a code \mathcal{C} of length n. The weight enumerator of \mathcal{C} is defined by $1 + A_1 z + A_2 z^2 + \cdots + A_n z^n$. The sequence $(1, A_1, A_2, \cdots, A_n)$

Kangquan Li is with the College of Liberal Arts and Sciences, National University of Defense Technology, Changsha, 410073, China and is currently a visiting Ph.D. student at the Department of Informatics, University of Bergen, Bergen N-5020, Norway. Chunlei Li and Tor Helleseth are with the Department of Informatics, University of Bergen, Bergen N-5020, Norway. Longjiang Qu is with the College of Liberal Arts and Sciences, National University of Defense Technology, Changsha, 410073, China, and is also with the State Key Laboratory of Cryptology, Beijing, 100878, China. The work of Longjiang Qu was supported by the Nature Science Foundation of China (NSFC) under Grant 61722213, 11531002, National Key R&D Program of China (No.2017YFB0802000), and the Open Foundation of State Key Laboratory of Cryptology. The work of Tor Helleseth and Chunlei Li was supported by the Research Council of Norway (No. 247742/O70 and No. 311646/O70). The work of Chunlei Li was also supported in part by the Natural Science Foundation of China under Grant (No. 61771021). The work of Kangquan Li was supported by China Scholarship Council. Longjiang Qu is the corresponding author.

Emails: likangquan11@nudt.edu.cn, chunlei.li@uib.no, tor.helleseth@uib.no, ljqu_happy@hotmail.com

is called the weight distribution of C. A code C is said to be a *t*-weight code if the number of nonzero A_i in the sequence (A_1, A_2, \dots, A_n) is equal to *t*. Linear *t*-weight codes with small values of *t* have many applications, including secret sharing schemes [5, 29], authentication codes [12], association schemes [1], strongly regular graphs [2], etc. In particular, one-weight codes are known as constant-weight codes, which are closely connected to the theory of Steiner systems and designs.

Known linear codes with good properties are constructed largely by two generic approaches [5, 9, 11]. By the first approach, linear codes over \mathbb{F}_q are defined based on a function f from \mathbb{F}_{q^n} to itself as

$$\overline{\mathcal{C}}_f = \{ (\operatorname{Tr}_n(af(x) + bx))_{x \in \mathbb{F}_{q^n}}, \ a, b \in \mathbb{F}_{q^n} \}$$

or

$$\mathcal{C}_f = \{ (\operatorname{Tr}_n(af(x) + bx))_{x \in \mathbb{F}_{q^n}}, \ a, \ b \in \mathbb{F}_{q^n} \}$$

when f(0) = 0, where *n* is a positive integer and Tr_n is the trace function from \mathbb{F}_{q^n} to \mathbb{F}_q . This generic construction has a long history and its importance is supported by Delsartes Theorem [7]. It also gives a coding-theory characterisation of APN monomials, almost bent functions, and bent functions [4], and cross-correlation between *m*-sequences and their decimations [14]. The second generic construction, introduced by Ding and Niederreiter, is described in terms of defining sets [11]. More specifically, it takes a subset $D = \{d_1, d_2, \dots, d_\ell\} \subseteq \mathbb{F}_{q^n}$ and defines a linear code of length ℓ over \mathbb{F}_q as

$$\mathcal{C}_D = \{(\mathrm{Tr}_n(xd_1), \mathrm{Tr}_n(xd_2), \cdots, \mathrm{Tr}_n(xd_\ell)) : x \in \mathbb{F}_{q^n}\},\$$

where D is called the defining set of the code C_D . When the defining set D is properly chosen, the code C_D can have good or optimal properties. This construction is generic in the sense that many classes of known codes could be produced by selecting proper defining sets D. In recent years, researchers have proposed new families of linear codes with few weights by considering defining sets derived from the support and image of certain functions over \mathbb{F}_{q^n} , see [8, 9, 13, 17, 27, 30] etc. Interested readers may refer to a recent survey by Li and Mesnager in [23] and references therein for good or optimal linear codes constructed from these two generic approaches.

Nonlinear functions over finite fields play important roles in cryptography, combinatorics, designs and sequence design. In coding theory, they have been employed in the above two generic constructions, resulting in a number of linear codes with good or optimal properties. Recently Mesnager and Qu in [25] made a systematic study of two-to-one functions over arbitrary finite fields, motivated by their close connection to special important primitives in symmetric cryptography. Later, Li et al. further developed the study of two-to-one functions over finite fields with characteristic 2 and proposed some two-to-one trinomials and quadrinomials [22].

The research of linear codes from two-to-one functions, to our best knowledge, began in [8, 10] when o-monomials and APN functions are discussed in the context. In this paper we will provide a more comprehensive study of two-to-one functions in constructing binary linear codes with few weights. Two

forms of two-to-one functions from \mathbb{F}_{2^n} to itself are considered. The first form is the generalized quadratic polynomial f(x), for which there exists a positive integer e with $gcd(e, 2^n - 1) = 1$ such that $f(x^e)$ is a quadratic function. The second form is the function $(x^{2^t} + x)^e$ with gcd(t, n) = 1 and $gcd(e, 2^n - 1) = 1$. Among the generalized quadratic polynomials, of particular interest are those with large ranks because they can produce linear codes with few weights. Hence some two-to-one functions in [22] and two newly constructed two-to-one polynomials are considered. As a result, we obtain many classes of 1-weight (a.k.a constant-weight), 3-weight, 5-weight binary linear codes by the two generic constructions. For the second form $(x^{2^t} + x)^e$, we give an interesting connection between the weight distribution of linear codes in second construction and the Walsh spectrum of the Boolean function $Tr_n(x^e)$. The connection enables us to derive many classes of 3-weight, 4-weight and 5-weight binary linear codes. By applying the Pless power moments, the weight distribution of the proposed constant-weight and 3-weight linear codes in this paper. In the end, based on the experiment results, we propose some open problems for the linear codes.

The remainder of this paper is organized as follows. Section 2 introduces mathematic foundations and auxiliary results. Section 3 first recalls some known two-to-one functions in [22] and then investigates the parameters of binary linear codes constructed from those two-to-one functions. In Section 4, we construct two new classes of two-to-one functions and propose 3-weight linear codes from them. In Section 5, we discuss the properties of linear codes from the two-to-one functions of the form $(x^{2^t} + x)^e$. Finally, the concluding remark on our work is given in Section 6.

2. PRELIMINARIES

This section presents basic notations, definitions and auxiliary results for the subsequent sections. Throughout this paper, we will restrict our discussion to finite fields with characteristic 2.

Let n be a positive integer. For $m \mid n$, let $\operatorname{Tr}_{n/m}(\cdot)$ denote the relative trace function from \mathbb{F}_{2^n} onto \mathbb{F}_{2^m} , i.e., $\operatorname{Tr}_{n/m}(x) = x + x^{2^m} + \cdots + x^{2^{\binom{n}{m}-1}m}$ for any $x \in \mathbb{F}_{2^n}$. Particularly, when m = 1, we use $\operatorname{Tr}_n(\cdot)$ to denote the absolute trace function from \mathbb{F}_{2^n} onto \mathbb{F}_2 . For any set E, we denote by |E| the cardinality of E.

A. Binary codes from two-to-one functions

Let f(x) be a mapping from \mathbb{F}_{2^n} to itself with f(0) = 0. Recall that the Walsh transform of f(x) at $(a,b) \in \mathbb{F}_{2^n} \times \mathbb{F}_{2^n}$ is given by

$$W_f(a,b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(ax+bf(x))}.$$
 (1)

Here we add the case that b = 0 in the definition for convenience.

In the first generic construction, the binary linear code from f is given by

$$\mathcal{C}_f = \left\{ \mathbf{c}_{a,b} = (\operatorname{Tr}_n(ax + bf(x)))_{x \in \mathbb{F}_{2^n}}, \ a, b \in \mathbb{F}_{2^n} \right\}.$$
(2)

Note that with the restriction f(0) = 0, we have $\operatorname{Tr}_n(ax+bf(x)) = 0$ for any $a, b \in \mathbb{F}_{2^n}$ when x = 0. Hence the code \mathcal{C}_f is commonly considered in the literature over the code $\overline{\mathcal{C}}_f$ in the first generic construction. In this case the linear code \mathcal{C}_f has length $2^n - 1$ and dimension at most 2n. Furthermore, in order to determine the dimension of \mathcal{C}_f , we only need to compute the number of $a, b \in \mathbb{F}_{2^n}$ such that $\operatorname{Tr}_n(ax+bf(x)) = 0$ for any $x \in \mathbb{F}_{2^n}$ since the code is linear. Equivalently, the dimension of \mathcal{C}_f is equal to $2n - d_{K_1}$, where d_{K_1} is the dimension of the \mathbb{F}_2 -vector space

$$K_1 = \left\{ (a,b) \in \mathbb{F}_{2^n}^2 : \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(ax+bf(x))} = 2^n \right\}.$$
(3)

For a codeword $\mathbf{c}_{a,b}$, its Hamming weight is given by

$$wt(\mathbf{c}_{a,b}) = 2^{n} - 1 - \# \{ x \in \mathbb{F}_{2^{n}}^{*} : \operatorname{Tr}_{n}(ax + bf(x)) = 0 \}$$

$$= 2^{n-1} - \frac{1}{2} \sum_{x \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n}(ax + bf(x))}$$

$$= 2^{n-1} - \frac{1}{2} W_{f}(a, b).$$
(4)

Therefore, the weight distribution of C can be directly derived from the Walsh spectrum of f(x):

$$\left\{ \left\{ W_f(a,b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(ax+bf(x))} : a, b \in \mathbb{F}_{2^n} \right\} \right\}.$$

Namely, if a value of $W_f(a, b)$ occurs X times in the Walsh spectrum of f, then there are $X/2^{d_{K_1}}$ codewords in C_f with Hamming weight $2^{n-1} - \frac{1}{2}W_f(a, b)$. In particular, when the Walsh transforms of f take only three values v_1, v_2, v_3 for all $a \in \mathbb{F}_{2^n}$ and $b \in \mathbb{F}_{2^n}^*$, the value distribution of $W_f(a, b)$ can be calculated by solving the following equations derived from the first three power moment identities:

$$\begin{cases} \sum_{a \in \mathbb{F}_{2^n}} \sum_{b \in \mathbb{F}_{2^n}} W_f^0(a, b) &= 2^{2n} = 2^{d_{K_1}} + X_1 + X_2 + X_3, \\ \sum_{a \in \mathbb{F}_{2^n}} \sum_{b \in \mathbb{F}_{2^n}} W_f(a, b) &= 2^{2n} = 2^{d_{K_1}} \cdot 2^n + v_1 X_1 + v_2 X_2 + v_3 X_3, \\ \sum_{a \in \mathbb{F}_{2^n}} \sum_{b \in \mathbb{F}_{2^n}} W_f^2(a, b) &= 2^{3n} = 2^{d_{K_1}} \cdot 2^{2n} + v_1^2 X_1 + v_2^2 X_2 + v_3^2 X_3, \end{cases}$$
(5)

where X_i is the occurrences of $W_f(a, b) = v_i$'s, i = 1, 2, 3 in the Walsh spectrum of f. Then the weight distribution of C_f can be determined accordingly.

In the second construction, let $D(f) = \{f(x) : x \in \mathbb{F}_{2^n}\} \setminus \{0\} = \{d_1, d_2, \dots, d_\ell\}$ and define the binary linear code

$$\mathcal{C}_{D(f)} = \left\{ \mathbf{c}_b = (\operatorname{Tr}_n(bd_1), \operatorname{Tr}_n(bd_2), \dots, \operatorname{Tr}_n(bd_\ell)) : b \in \mathbb{F}_{2^n} \right\}.$$
(6)

It is clear that the code $C_{D(f)}$ has length $\ell = |D(f)|$ and dimension at most n. Furthermore, in order to determine the dimension of $C_{D(f)}$, we need to compute the number of $b \in \mathbb{F}_{2^n}$ such that $\operatorname{Tr}_n(bf(x)) = 0$

for any $x \in \mathbb{F}_{2^n}$ since the code is linear. Equivalently, the dimension of $\mathcal{C}_{D(f)}$ is equal to $n - d_{K_2}$, where d_{K_1} is the dimension of the \mathbb{F}_2 -vector space

$$K_{2} = \left\{ b \in \mathbb{F}_{2^{n}} : \sum_{x \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n}(bf(x))} = 2^{n} \right\}.$$
(7)

For any $b \in \mathbb{F}_{2^n}$, the Hamming weight of a codeword \mathbf{c}_b in $\mathcal{C}_{D(f)}$ is given by

wt(
$$\mathbf{c}_b$$
) = $|\{1 \le i \le \ell : \operatorname{Tr}_n(bd_i) = 1\}| = \frac{1}{2} \left(|D_f| - \sum_{d \in D(f)} (-1)^{\operatorname{Tr}_n(bd)} \right).$

From the above formula, one sees that the weight distribution of the linear code $C_{D(f)}$ is essentially the value distribution of a partial exponential sum, which is generally intractable if f is not properly chosen.

Suppose that f(x) is a two-to-one mapping from \mathbb{F}_{2^n} to itself, which means that $|f^{-1}(a)| = 2$ for any $a \in \text{Im}(f)$. Then the linear code $\mathcal{C}_{D(f)}$ has length $|D_f| = 2^{n-1} - 1$ and the Hamming weight of its codeword is given by

$$\operatorname{wt}(\mathbf{c}_{b}) = \frac{1}{2} \left(|D_{f}| - \frac{1}{2} \sum_{x \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n}(bf(x))} + 1 \right) = 2^{n-2} - \frac{1}{4} \sum_{x \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n}(bf(x))}.$$
(8)

From (7) and (8), one sees that the dimension and the weight distribution of $C_{D(f)}$ heavily depend on the value of

$$W_f(b) \triangleq W_f(0,b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(bf(x))}, \quad b \in \mathbb{F}_{2^n}.$$
(9)

In particular, if $W_f(b)$ takes only three values v_1, v_2 and v_3 for $b \in \mathbb{F}_{2^n}^*$, then the code $\mathcal{C}_{D(f)}$ has three nonzero weights, namely, $w_i = 2^{n-2} - v_i/4$ for i = 1, 2, 3. Denote by A_i the number of codewords with weight w_i in $\mathcal{C}_{D(f)}$. Note that the dual of $\mathcal{C}_{D(f)}$ has Hamming weight no less than 3 since $\operatorname{Tr}_n(xd) = 0$ holds for all $x \in \mathbb{F}_{2^n}$ if and only if d = 0. The first three Pless Power Moments [19, p. 260] leads to the following system of equations:

$$\begin{cases}
A_1 + A_2 + A_3 = 2^n - 1 \\
w_1 A_1 + w_2 A_2 + w_3 A_3 = \ell 2^{n-1} \\
w_1^2 A_1 + w_2^2 A_2 + w_3^2 A_3 = \ell(\ell+1)2^{n-2},
\end{cases}$$
(10)

where $\ell = 2^{n-1} - 1$. Therefore, the weight distribution of $C_{D(f)}$ can be determined from the above system of equations when it is shown to have only three nonzero weights.

The above discussion shows that for a two-to-one mapping f, the parameters of the linear codes C_f in (2) and $C_{D(f)}$ in (6) depend on the investigation of the Walsh transform of f. In addition, it is clear that the number of nonzero weights in $C_{D(f)}$ is no more than that of C_f . Therefore, we will focus on the two-to-one functions of which the Walsh transforms have few different values.

At the end of this subsection, we consider the minimum distance of dual codes of C_f in (2) and $C_{D(f)}$ in

(6) in the following theorem.

Theorem 1. Let f be a two-to-one mapping over \mathbb{F}_{2^n} , C_f and $C_{D(f)}$ be defined as in (2) and (6), respectively. Let C_f^{\perp} and $C_{D(f)}^{\perp}$ be the dual code of C_f and $C_{D(f)}$, respectively. Let d_{K_1} and d_{K_2} be defined as in (3) and (7), respectively. Then

- (1) C_f^{\perp} is a $[2^n 1, 2^n 1 2n + d_{K_1}]$ binary code with the minimum distance d_f^{\perp} satisfying $3 \le d_f^{\perp} \le 6$. Particularly, if $d_{K_1} \ge 2$, $3 \le d_f^{\perp} \le 4$. Moreover, $d_f^{\perp} = 3$ if and only if there exist two distinct elements $x_1, x_2 \in \mathbb{F}_{2^n}^*$ such that $f(x_1) + f(x_2) + f(x_1 + x_2) = 0$.
- (2) $C_{D(f)}^{\perp}$ is a $[2^{n-1}-1, 2^{n-1}-1-n+d_{K_2}]$ binary code with the minimum distance $d_{D(f)}^{\perp}$ satisfying $3 \leq d_{D(f)}^{\perp} \leq 4$. Particularly, when $d_{K_2} = 1$, the equality of the sphere packing bound can be achieved. Moreover, $d_{D(f)}^{\perp} = 3$ if and only if there exist three distinct elements $x_1, x_2, x_3 \in \mathbb{F}_{2^n}^*$ such that $f(x_i) \neq f(x_j)$ for $1 \leq i < j \leq 3$ and $f(x_1) + f(x_2) + f(x_3) = 0$.

Proof. From the above discussion, we know that the linear code C_f (resp. $C_{D(f)}$) has length $2^n - 1$ (resp. $2^{n-1}-1$) and dimension $2n - d_{K_1}$ (resp. $n - d_{K_2}$). Then according to the definition, the length and dimension of C_f^{\perp} and $C_{D(f)}^{\perp}$ can be trivially determined. Thus it suffices to consider the minimum distances of the dual codes.

(1) For C_f^{\perp} , it is clear that $d_f^{\perp} \neq 1$. If C_f^{\perp} has a codeword of weight two, then there exist two distinct elements $x_1, x_2 \in \mathbb{F}_{2^n}^*$ such that $x_1 + x_2 = 0$ and $f(x_1) + f(x_2) = 0$, which is impossible. Thus $d_f^{\perp} \geq 3$. We now prove that $d_f^{\perp} \leq 6$. Suppose that $d_f^{\perp} \geq 7$. We would then have

$$\sum_{i=0}^{3} {\binom{2^{n}-1}{i}} (2-1)^{i}$$

$$= 1+2^{n}-1+(2^{n}-1)\left(2^{n-1}-1\right)+\frac{(2^{n}-1)\left(2^{n}-3\right)\left(2^{n-1}-1\right)}{3}$$

$$= 2^{2n-1}-2^{n-1}+1+\frac{(2^{n}-1)\left(2^{n}-3\right)\left(2^{n-1}-1\right)}{3}$$

$$> 2^{2n-d_{K_{1}}}.$$

which is contrary to the sphere packing bound. Thus $d_f^{\perp} \leq 6$. In particular, obviously, if $d_{K_1} \geq 2$, $d_f^{\perp} \leq 4$. Moreover, according to the definition, C_f^{\perp} has a codeword of weight three if and only if there are pairwise distinct elements $x_1, x_2, x_3 \in \mathbb{F}_{2^n}^*$ such that

$$\begin{cases} x_1 + x_2 + x_3 = 0\\ f(x_1) + f(x_2) + f(x_3) = 0 \end{cases}$$

i.e., there exist two distinct elements $x_1, x_2 \in \mathbb{F}_{2^n}^*$ such that $f(x_1) + f(x_2) + f(x_1 + x_2) = 0$.

(2) For $\mathcal{C}_{D(f)}^{\perp}$, it is also clear that $d_{D(f)}^{\perp} \neq 1$. Moreover, if $\mathcal{C}_{D(f)}^{\perp}$ has a codeword of weight two, then there exist two distinct elements $d_1, d_2 \in \text{Im}(f)$ such that $d_1 + d_2 = 0$, which is contrary. Thus $d_{D(f)}^{\perp} \geq 3$. We next show that $d_f^{\perp} \leq 4$. Suppose that $d_f^{\perp} \geq 5$. We would then have

$$\sum_{i=0}^{2} {\binom{2^{n-1}-1}{i}} (2-1)^{i}$$

= $1+2^{n-1}-1+(2^{n-1}-1)(2^{n-2}-1)$
= $2^{2n-3}-2^{n-2}+1$
> $2^{n-d_{K_2}}$,

which is contrary to the sphere packing bound. Thus $d_f^{\perp} \leq 4$. Particularly, when $d_{K_2} = 1$, the equality of the sphere packing bound can be achieved, namely,

$$\sum_{i=0}^{1} \binom{2^{n-1}-1}{i} (2-1)^{i} = 2^{n-1} = 2^{n-d_{\kappa_2}}.$$

Moreover, from the definition, $\mathcal{C}_{D(f)}^{\perp}$ has a codeword of weight three if and only if there are three distinct elements $d_1, d_2, d_3 \in \text{Im}(f)$ such that $d_1+d_2+d_3=0$, i.e., there exist three distinct elements $x_1, x_2, x_3 \in \mathbb{F}_{2^n}^*$ such that $f(x_i) \neq f(x_j)$ for $1 \leq i < j \leq 3$ and $f(x_1) + f(x_2) + f(x_3) = 0$.

We need to recall some useful results on the Walsh transforms of quadratic functions.

B. Quadratic functions and Walsh transforms

Let Q(x) be a quadratic function from \mathbb{F}_{2^n} to itself, i.e., it has algebraic degree 2. Let $\varphi(x) = \text{Tr}_n(Q(x))$ and define its associated bilinear mapping as

$$B_{\varphi}(x,y) = \varphi(x+y) + \varphi(x) + \varphi(y).$$

The set

$$V_{\varphi} = \{y \in \mathbb{F}_{2^n} \, : \, \varphi(x+y) + \varphi(x) + \varphi(y) = 0 \text{ for any } x \in \mathbb{F}_{2^n}\}$$

is an \mathbb{F}_2 -vector space and is known as the kernel of the bilinear mapping $B_{\varphi}(x, y)$. The rank of $\varphi(x)$ is defined by

$$\operatorname{Rank}(\varphi(x)) = n - \dim_{\mathbb{F}_2}(V_{\varphi}).$$

Observe that

$$\left(\sum_{x\in\mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(Q(x))}\right)^2 = \sum_{x\in\mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(Q(x))} \sum_{y\in\mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(Q(y))}$$
$$= \sum_{x,y\in\mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(Q(x+y)+Q(y))}$$
$$= \sum_{y\in\mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(Q(y))} \sum_{x\in\mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(Q(x+y)+Q(x)+Q(y))}$$
$$= 2^n \sum_{y\in V_{\varphi}} (-1)^{\operatorname{Tr}_n(Q(y))}.$$
(11)

By the definition of the kernel V_{φ} , it is readily seen that the function $\varphi(y) = \text{Tr}_n(Q(y))$ is linear over V_{φ} . Then one has

$$\sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(Q(x))} = \begin{cases} \pm 2^{\frac{n+d}{2}}, & \text{if } \varphi(x) = 0 \text{ for all } x \in V_{\varphi}, \\ 0, & \text{otherwise,} \end{cases}$$
(12)

where d is the dimension of V_{φ} over \mathbb{F}_2 .

For a quadratic function f from \mathbb{F}_{2^n} to itself, the bilinear mapping of $\operatorname{Tr}_n(bf(x) + ax)$ is the same as that of $\operatorname{Tr}_n(bf(x))$ for any nonzero b in \mathbb{F}_{2^n} . Therefore, the Walsh transform of f at (a, b) can be given, similar to (12), as below:

$$W_f(a,b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(ax+bf(x))} = \begin{cases} \pm 2^{\frac{n+d_b}{2}}, & \text{if } \operatorname{Tr}_n(bf(x)+ax) = 0 \text{ for all } x \in V_{\varphi_b}, \\ 0, & \text{otherwise.} \end{cases}$$
(13)

where d_b is the dimension of V_{φ_b} over \mathbb{F}_2 and V_{φ_b} is the kernel of the bilinear mapping of $\varphi_b(x) = \text{Tr}_n(bf(x))$.

In Sections 3 and 4, we will discuss the properties of linear codes defined in (2) and (6) from generalized quadratic functions. The Walsh transform of quadratic functions f(x) in (13) will be heavily used in the discussion.

C. Factorization of low-degree polynomials

The following lemma describes the factorization of a cubic polynomial over \mathbb{F}_{2^m} . If f factors over \mathbb{F}_{2^m} as a product of three linear factors we write f = (1, 1, 1), if f factors as a product of a linear factor and an irreducible quadratic factor we write f = (1, 2) and finally if f is irreducible over \mathbb{F}_{2^m} we write f = (3).

Lemma 2. [28] Let $f(x) = x^3 + ax + b \in \mathbb{F}_{2^m}[x]$ and $b \neq 0$. Let t_1, t_2 denote the solutions of $t^2 + bt + a^3 = 0$. Then the factorizations of f(x) over \mathbb{F}_{2^m} are characterized as follows:

- (1) f = (1, 1, 1) if and only if $\operatorname{Tr}_m\left(\frac{a^3}{b^2}\right) = \operatorname{Tr}_m(1)$, t_1, t_2 cubes in \mathbb{F}_{2^m} (*m* even), $\mathbb{F}_{2^{2m}}$ (*m* odd); (2) f = (1, 2) if and only if $\operatorname{Tr}_m\left(\frac{a^3}{b^2}\right) \neq \operatorname{Tr}_m(1)$;
- (3) f = (3) if and only if $\operatorname{Tr}_m\left(\frac{a^3}{b^2}\right) = \operatorname{Tr}_m(1)$, t_1, t_2 not cubes in \mathbb{F}_{2^m} (m even), $\mathbb{F}_{2^{2m}}$ (m odd).

Lemma 3. [28] Let $f(x) = x^3 + ax + b \in \mathbb{F}_{2^m}[x]$ and $b \neq 0$. Let t be one solution of $t^2 + bt + a^3 = 0$ and ϵ be one solution of $x^3 = t$. Then $r = \epsilon + \frac{a}{\epsilon}$ is a solution of f(x) = 0.

The following lemma characterizes the factorization of a quartic polynomial over \mathbb{F}_{2^n} .

Lemma 4. [21] Let $f(x) = x^4 + a_2x^2 + a_1x + a_0$ with $a_i \in \mathbb{F}_{2^n}$ and $a_0a_1 \neq 0$. Let $f_1(y) = y^3 + a_2y + a_1$ and r_1, r_2, r_3 denote roots of $f_1(y) = 0$ when they exist in \mathbb{F}_{2^n} . Set $w_i = a_0 \frac{r_i^2}{a_1^2}$. Then the factorization of f(x) over \mathbb{F}_{2^n} is characterized as follows:

(1) f = (1, 1, 1, 1) if and only if $f_1 = (1, 1, 1)$ and $\operatorname{Tr}_n(w_1) = \operatorname{Tr}_n(w_2) = \operatorname{Tr}_n(w_3) = 0$; (2) f = (2, 2) if and only if $f_1 = (1, 1, 1)$ and $\operatorname{Tr}_n(w_1) = 0$, $\operatorname{Tr}_n(w_2) = \operatorname{Tr}_n(w_3) = 1$; (3) f = (1, 3) if and only if $f_1 = (3)$; (4) f = (1, 1, 2) if and only if $f_1 = (1, 2)$ and $\operatorname{Tr}_n(w_1) = 0$; (5) f = (4) if and only if $f_1 = (1, 2)$ and $\operatorname{Tr}_n(w_1) = 1$.

3. BINARY LINEAR CODES FROM KNOWN TWO-TO-ONE TRINOMIALS AND QUADRINOMIALS

In this section, we will propose several binary codes with few weights, which are constructed from known two-to-one functions . We first recall some two-to-one functions recently obtained in [22].

Lemma 5. [22] Let $f(x) = x^{\frac{2^{n-1}+2^m-1}{3}} + x^{2^m} + \omega x \in \mathbb{F}_{2^n}[x]$, where n = 2m, m is odd and $\omega \in \mathbb{F}_{2^2} \setminus \mathbb{F}_2$. Then f(x) is two-to-one over \mathbb{F}_{2^n} .

Lemma 6. [22] Let n = 2m + 1. Then the following quadrinomials are all two-to-one over \mathbb{F}_{2^n} :

(1) $f(x) = x^{2^{m+1}+2} + x^{2^{m+1}} + x^2 + x;$ (2) $f(x) = x^{2^{m+1}+2} + x^{2^{m+1}+1} + x^2 + x;$ (3) $f(x) = x^{2^{m+2}+4} + x^{2^{m+1}+2} + x^2 + x;$ (4) $f(x) = x^{2^n - 2^{m+1}+2} + x^{2^{m+1}} + x^2 + x.$

Lemma 7. [22] Let n = 3m. Then the following quadrinomials are two-to-one over \mathbb{F}_{2^n} : (1) $f(x) = x^{2^{2m}+1} + x^{2^{m+1}} + x^{2^m+1} + x$ with $m \not\equiv 1 \pmod{3}$; (2) $f(x) = x^{2^{2m}+2^m} + x^{2^{2m}+1} + x^{2^m+1} + x$.

Below we shall investigate the parameters of the constructed linear codes C_f and $C_{D(f)}$. According to different forms of n in Lemmas 5 - 7, we divide them into three subsections.

A. The case n = 2m

The following binary linear code is derived from the two-to-one polynomial in Lemma 5.

Theorem 8. Let $f(x) = x^{\frac{2^{n-1}+2^m-1}{3}} + x^{2^m} + \omega x \in \mathbb{F}_{2^n}[x]$ with n = 2m, where m > 1 is odd and $\omega \in \mathbb{F}_{2^2} \setminus \mathbb{F}_2$. Define two linear codes C_f and $C_{D(f)}$ as in (2) and (6), respectively. Then, (1) C_f is a $[2^n - 1, 2n]$ binary linear code with weight distribution in Table I.

TABLE IThe weight distribution of the codes \mathcal{C}_f in Theorem 8

Weight	Multiplicity
0	1
$2^{n-1} - 2^m$	$2^{4m-3} + 2^{3m-2} - 2^{2m-3} - 2^{m-2}$
2^{n-1}	$3 \cdot 2^{4m-2} + 2^{2m-2} - 1$
$2^{n-1} + 2^m$	$2^{4m-3} + 2^{m-2} - 2^{3m-2} - 2^{2m-3}$

TABLE II The weight distribution of the codes $\mathcal{C}_{D(f)}$ in Theorem 8

Weight	Multiplicity
0	1
$2^{n-2} - 2^{m-1}$	$2^{n-3} + 2^{m-2}$
2^{n-2}	$3 \cdot 2^{n-2} - 1$
$2^{n-2} + 2^{m-1}$	$2^{n-3} - 2^{m-2}$

(2) $C_{D(f)}$ is a $[2^{n-1}-1,n]$ binary linear code with weight distribution in Table II.

Proof. We first compute the value of $W_f(a, b)$, i.e., the Walsh transform of f, and $W_f(b)$ defined as in (9), for any $a, b \in \mathbb{F}_{2^n}$. It is obvious that $W_f(a, b) = 2^n$ when a = b = 0. Below we consider the cases where $(a, b) \neq (0, 0)$.

By the Euclidean algorithm we have

$$(2^{2m} - 1) \times (2^{m+1} + 5) = (2^{2m-1} + 2^m - 1) \times (2^{m+2} + 2) - 3,$$

which implies $gcd\left(\frac{2^{n-1}+2^{m}-1}{3}, 2^{2m}-1\right) = 1$. Define

$$f_1(x) = f\left(x^{2^{m+2}+2}\right) = x + x^{2^{m+1}+4} + \omega x^{2^{m+2}+2}$$

and

$$Q(x) = ax^{2^{m+2}+2} + bf_1(x) = bx + bx^{2^{m+1}+4} + (b\omega + a)x^{2^{m+2}+2}.$$

Then the Walsh transform

$$W_f(a,b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n\left(ax^{2^{m+2}+2} + bf\left(x^{2^{m+2}+2}\right)\right)} = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(Q(x))}.$$

Note that the bilinear form of $\varphi_{a,b}(x) = \operatorname{Tr}_n(Q(x))$ is given by

$$B_{\varphi_{a,b}}(x,y) = \varphi_{a,b}(x+y) + \varphi_{a,b}(x) + \varphi_{a,b}(y)$$

= $\operatorname{Tr}_n \left(by^4 x^{2^{m+1}} + by^{2^{m+1}} x^4 + (b\omega + a)y^2 x^{2^{m+2}} + (b\omega + a)y^{2^{m+2}} x^2 \right)$
= $\operatorname{Tr}_n \left(L_{a,b}(y) x^{2^{m+2}} \right)$

with

$$L_{a,b}(y) = \Delta^2 y^8 + \Delta^{2^m} y^2$$
, and $\Delta = b^{2^m} \omega^2 + b + a^{2^m}$.

Let $\ker(L_{a,b}) = \{y \in \mathbb{F}_{2^n} : L_{a,b}(y) = 0\}$. From (12), we have

$$W_f(a,b) = 2^n \sum_{y \in \ker(L_{a,b})} (-1)^{\varphi_{a,b}(y)} = \begin{cases} \pm 2^{\frac{n+2}{2}}, & \text{if } \varphi_{a,b}(x) = 0 \text{ for all } x \in \ker(L_{a,b}), \\ 0, & \text{otherwise.} \end{cases}$$

Now we discuss the values of $\varphi_{a,b}(x)$ on the kernel ker $(L_{a,b})$. When $\Delta = 0$, we have $L_{a,b}(y) = 0$ for any $y \in \mathbb{F}_{2^n}$. When $\Delta \neq 0$, by computation we have

$$\ker(L_{a,b}) = \{0, y_0, y_0\omega, y_0\omega^2\}, \text{ where } y_0 = \Delta^{\frac{2^{m-1}-1}{3}}.$$

Moreover,

$$\varphi_{a,b}(y_0) = \operatorname{Tr}_n \left(by_0 + \frac{b}{b^{2^m} \omega^2 + b + a^{2^m}} + \frac{\omega b + a}{b\omega + b^{2^m} + a} \right)$$
$$= \operatorname{Tr}_n \left(by_0 + \frac{b + b^{2^m} \omega^2 + a^{2^m}}{b^{2^m} \omega^2 + b + a^{2^m}} \right) = \operatorname{Tr}_n \left(by_0 \right).$$

Similarly, we have

$$\varphi_{a,b}(\omega y_0) = \operatorname{Tr}_n(\omega b y_0) \text{ and } \varphi_{a,b}(\omega^2 y_0) = \operatorname{Tr}_n(\omega^2 b y_0) = \varphi_{a,b}(y_0) + \varphi_{a,b}(\omega y_0).$$

In the following, we assume a = 0 and show that there exist some b's $\in \mathbb{F}_{2^n}$ such that $\varphi_{0,b}(x) = 0$ for all $x \in \ker(L_{a,b})$ and $\varphi_{0,b}(x) = 1$ for some $x \in \ker(L_{a,b})$, which implies

$$W_f(b) \in \left\{0, \pm 2^{\frac{n+2}{2}}\right\}.$$

It is well known that for any elements $b \in \mathbb{F}_{2^n}$, there exist unique $b_1, b_2 \in \mathbb{F}_{2^m}$ such that $b = b_1 + b_2 \omega$ since m is odd. Plugging $b = b_1 + \omega b_2$ into the expression of y_0 , we get

$$y_0 = \left(b^{2^m}\omega^2 + b\right)^{\frac{2^{m-1}-1}{3}} = \left((b_1 + b_2\omega^2)\omega^2 + b_1 + b_2\omega\right)^{\frac{2^{m-1}-1}{3}} = (b_1\omega)^{\frac{2^{m-1}-1}{3}}$$

Let $g(b) = by_0 = b \left(b^{2^m} \omega^2 + b \right)^{\frac{2^{m-1}-1}{3}}$. Then

$$g(b) = (b_1 + b_2\omega)(b_1\omega)^{\frac{2^{m-1}-1}{3}} = \omega^{\frac{2^{m-1}-1}{3}} \left(b_1^{\frac{2^{m-1}+2}{3}} + b_1^{\frac{2^{m-1}-1}{3}} b_2\omega \right).$$

If there exist two elements $\hat{b} = \hat{b}_1 + \hat{b}_2 \omega$, $\tilde{b} = \tilde{b}_1 + \tilde{b}_2 \omega$ with $\hat{b}_1, \hat{b}_2, \tilde{b}_1, \tilde{b}_2 \in \mathbb{F}_{2^m}$ such that $g(\hat{b}) = g(\tilde{b})$, then we have

$$\hat{b}_1^{\frac{2^{m-1}+2}{3}} = \tilde{b}_1^{\frac{2^{m-1}+2}{3}}$$
 and $\hat{b}_1^{\frac{2^{m-1}-1}{3}} \hat{b}_2 = \tilde{b}_1^{\frac{2^{m-1}-1}{3}} \tilde{b}_2$

Since $gcd\left(\frac{2^{m-1}+2}{3}, 2^m-1\right) = 1$, we have $\hat{b}_1 = \tilde{b}_1$. Then if $\hat{b}_1 = \tilde{b}_1 \neq 0$, $\hat{b}_2 = \tilde{b}_2$; if $\hat{b}_1 = \tilde{b}_1 = 0$, $\hat{b}_2, \tilde{b}_2 \in \mathbb{F}_{2^m}$. Thus for $b = b_1 + \omega b_2$, if $b_1 = 0$, g(b) = 0; if $b_1 \neq 0$, g(b) is bijective. Let $Im(g) = \{g(b) : b \in \mathbb{F}_{2^n}\}$.

Then

$$|\mathrm{Im}(g)| = 2^n - 2^m + 1.$$

According to the property of the trace function, there exist $(2^n - 2^{n-2})$ elements $x \in \mathbb{F}_{2^n}$ such that $\operatorname{Tr}_{n/2}(x) \neq 0$. Since $2^n - 2^m + 1 > 2^n - 2^{n-2}$, there must exist some $b \in \mathbb{F}_{2^n}$ such that $\operatorname{Tr}_{n/2}(g(b)) = 0$ and then

$$\operatorname{Tr}_{n}(by_{0}) = \operatorname{Tr}_{2}\left(\operatorname{Tr}_{n/2}(g(b))\right) = 0$$

and

$$\operatorname{Tr}_{n}(\omega by_{0}) = \operatorname{Tr}_{2}\left(\operatorname{Tr}_{n/2}(\omega g(b))\right) = \operatorname{Tr}_{2}\left(\omega \operatorname{Tr}_{n/2}(g(b))\right) = 0.$$

Clearly, there also exist some $b \in \mathbb{F}_{2^n}$ such that $\operatorname{Tr}_{n/2}(g(b)) = 1$ and then $\operatorname{Tr}_n(\omega by_0) = \operatorname{Tr}_2(\omega) = 1$. Thus $W_f(b) \in \{0, \pm 2^{\frac{n+2}{2}}\}$ and obviously, $W_f(a, b) \in \{0, \pm 2^{\frac{n+2}{2}}\}$ for $(a, b) \neq (0, 0)$.

With the analysis of possible values of $W_f(a, b)$ and $W_f(b)$, we're now ready to determine the parameters of C_f and $C_{D(f)}$ in the following.

(1) For the linear code C_f , since $W_f(a,b) = 0$ if and only if a = b = 0, it follows from (3) that the dimension is 2n. Moreover, for any $a, b \in \mathbb{F}_{2^n}$, $W_f(a,b) \in \left\{0, 2^n, \pm 2^{\frac{n+2}{2}}\right\}$. Let

$$v_1 = -2^{\frac{n+2}{2}}, \quad v_2 = 0, \quad v_3 = 2^{\frac{n+2}{2}}.$$

According to (5), the occurrences of $W_f(a, b) = v_i$'s, i = 1, 2, 3 in the Walsh spectrum of f are

$$\begin{cases} X_1 = 2^{4m-3} + 2^{m-2} - 2^{3m-2} - 2^{2m-3} \\ X_2 = 3 \cdot 2^{4m-2} + 2^{2m-2} - 1 \\ X_3 = 2^{4m-3} + 2^{3m-2} - 2^{2m-3} - 2^{m-2}. \end{cases}$$

Then the desired weight distribution of C_f in Table I follows directly from (4).

(2) For the linear code $C_{D(f)}$, since $W_f(b) = 2^n$ if and only if b = 0, it follows from (7) that the dimension of $C_{D(f)}$ is n. Note that for any $b \in \mathbb{F}_{2^n}$, $W_f(b) \in \left\{0, 2^n, \pm 2^{\frac{n+2}{2}}\right\}$. By (8), the weights of the codewords \mathbf{c}_b in $C_{D(f)}$ satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-2}, 0, 2^{n-2} - 2^{\frac{n-2}{2}}, 2^{n-2} + 2^{\frac{n-2}{2}} \right\}$

Denote

$$w_1 = 2^{n-2} - 2^{\frac{n-2}{2}}, \quad w_2 = 2^{n-2}, \quad w_3 = 2^{n-2} + 2^{\frac{n-2}{2}}.$$

The desired weight distribution of $C_{D(f)}$ in Table II can be easily obtained by solving (10) accordingly. **Example 1.** When m = 3, the code C_f in Theorem 8 is a [63, 12, 24] binary linear code with the weight enumerator

$$1 + 630z^{24} + 3087z^{32} + 378z^{36}$$

TABLE III The weight distribution of the codes $\mathcal{C}_{D(f)}$ in Theorem 9

Weight	Multiplicity
0	1
$2^{n-1} - 2^m$	$2^{4m} + 2^{3m} - 2^{2m-1} - 2^{m-1}$
2^{n-1}	$2^{4m+1} + 2^{2m} - 1$
$2^{n-1} + 2^m$	$2^{4m} + 2^{m-1} - 2^{3m} - 2^{2m-1}$

TABLE IV The weight distribution of the codes $\mathcal{C}_{D(f)}$ in Theorem 9

Weight	Multiplicity
0	1
$2^{n-2} - 2^{m-1}$	$2^{n-2} + 2^{m-1}$
2^{n-2}	$2^{n-1} - 1$
$2^{n-2} + 2^{m-1}$	$2^{n-2} - 2^{m-1}$

Referring to the code table [16], the linear code has the best known parameter.

B. The case n = 2m + 1

From the four classes of two-to-one functions in Lemma 6, this subsection presents five classes of 3weight linear codes, two classes of constant-weight linear codes and one class of at most 5-weight linear codes.

Theorem 9. Let n = 2m + 1 and $f(x) = x^{2^{m+1}+2} + x^{2^{m+1}} + x^2 + x$. Define two linear codes C_f and $C_{D(f)}$ as in (2) and (6), respectively. Then,

(1) C_f is a [2ⁿ - 1, 2n] binary linear code with weight distribution in Table III.
 (2) C_{D(f)} is a [2ⁿ⁻¹ - 1, n] binary linear code with weight distribution in Table IV.

Proof. In a similar manner as in Theorem 8, we will first investigate the value of $W_f(a, b)$ and then discuss the parameters of C_f and $C_{D(f)}$.

It is clear that $W_f(a,b) = 2^n$ when (a,b) = (0,0). For $(a,b) \neq (0,0)$, let $\varphi_b(x) = \text{Tr}_n(bf(x))$. Since f is quadratic, according to (13), we need to compute the dimension of the kernel of the bilinear form of $\varphi_b(x)$. Note that the bilinear form of $\varphi_b(x)$ is given by

$$B_{\varphi_b}(x,y) = \varphi_b(x+y) + \varphi_b(x) + \varphi_b(y) = \operatorname{Tr}_n\left(by^2 x^{2^{m+1}} + by^{2^{m+1}} x^2\right) = \operatorname{Tr}_n\left(L_b(y) x^{2^{m+1}}\right),$$

where $L_b(y) = by^2 + b^{2^m}y$. Clearly, $\ker(L_b) = \{0, b^{2^m-1}\}$. According to (13), we have

$$W_f(a,b) = \begin{cases} \pm 2^{\frac{n+1}{2}}, & \text{if } \operatorname{Tr}_n(bf(x) + ax) = 0 \text{ for all } x \in \{0, b^{2^m - 1}\}\\ 0, & \text{otherwise.} \end{cases}$$

Moreover, for $x = b^{2^m - 1}$,

$$\operatorname{Tr}_{n}(bf(x) + ax) = \operatorname{Tr}_{n} \left(b^{-1} + \left(b + b^{2^{m}} + b^{2^{m+1}} \right) b^{1-2^{m+1}} + ab^{2^{m}-1} \right)$$

=
$$\operatorname{Tr}_{n} \left(b^{-1} + b + b^{1-2^{m}} + b^{2-2^{m+1}} + ab^{2^{m}-1} \right)$$

=
$$\operatorname{Tr}_{n} \left(b^{-1} + b + ab^{2^{m}-1} \right).$$

Obviously, if $(a,b) = (0,b) \neq (0,0)$, $\operatorname{Tr}_n(b^{-1} + b + ab^{2^m - 1}) \in \{0,1\}$ and thus for any $(a,b) \neq (0,0)$,

$$W_f(a,b), W_f(b) \in \left\{0, \pm 2^{\frac{n+1}{2}}\right\}.$$

Next, we consider the parameters of C_f and $C_{D(f)}$, respectively.

(1) For the linear code C_f , since $W_f(a,b) = 2^n$ if and only if (a,b) = (0,0), the dimension is 2n from (3). Moreover, for any $a, b \in \mathbb{F}_{2^n}$, $W_f(a,b) \in \left\{0, 2^n, \pm 2^{\frac{n+1}{2}}\right\}$. Let

$$v_1 = -2^{\frac{n+1}{2}}, \quad v_2 = 0, \quad v_3 = 2^{\frac{n+1}{2}}.$$

Then by computing (5), we can get that the occurrences of $W_f(a, b) = v_i$'s, i = 1, 2, 3 in the Walsh spectrum of f are

$$\begin{cases} X_1 = 2^{4m} + 2^{m-1} - 2^{3m} - 2^{2m-1} \\ X_2 = 2^{4m+1} + 2^{2m} - 1 \\ X_3 = 2^{4m} + 2^{3m} - 2^{2m-1} - 2^{m-1}. \end{cases}$$

Finally, by (4), the desired weight distribution of C_f can be obtained.

(2) For the linear code $\mathcal{C}_{D(f)}$, $W_f(b) = 2^n$ if and only if b = 0, which means that the dimension of $\mathcal{C}_{D(f)}$ is *n* according to (7). Since for any $b \in \mathbb{F}_{2^n}$, $W_f(b) \in \left\{0, 2^n, \pm 2^{\frac{n+1}{2}}\right\}$, by (8), the weights of the codewords \mathbf{c}_b in $\mathcal{C}_{D(f)}$ satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-2}, 0, 2^{n-2} - 2^{\frac{n-3}{2}}, 2^{n-2} + 2^{\frac{n-3}{2}} \right\}.$

In the following, we determine the weight distribution of $C_{D(f)}$. Define

$$w_1 = 2^{n-2} - 2^{\frac{n-3}{2}}, \quad w_2 = 2^{n-2}, \quad w_3 = 2^{n-2} + 2^{\frac{n-3}{2}}.$$

Then solving (10) gives the desired weight distribution.

Example 2. When m = 3, the code C_f in Theorem 9 is a [127, 14, 56] binary linear code with the weight enumerator

$$1 + 4572z^{56} + 8255z^{64} + 3556z^{72}$$

Referring to the code table [16], the linear code is optimal.

Theorem 10. Let n = 2m + 1 and $f(x) = x^{2^n - 2^{m+1} + 2} + x^{2^{m+1}} + x^2 + x$. Define two linear codes C_f and $C_{D(f)}$ as in (2) and (6), respectively. Then,

(1) C_f is a $[2^n - 1, 2n]$ binary linear code with at most five weights.

(2) $C_{D(f)}$ is a $[2^{n-1}-1,n]$ binary linear code with weight distribution in Table IV.

Proof. Since $gcd(2^n - 2^{m+1} + 2, 2^n - 1) = 1$ and

$$(2^n - 2^{m+1} + 2) \times (2^m + 1) \equiv 2^m + 2 \pmod{2^n - 1},$$

we have

$$W_{f}(a,b) = \sum_{x \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n} \left(ax + b \left(x^{2^{n}-2^{m+1}+2} + x^{2^{m+1}} + x^{2} + x \right) \right)} \\ = \sum_{x \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n} \left(ax^{2^{m}+1} + b \left(x^{2^{m}+2} + x^{2^{m+1}+1} + x^{2^{m+1}+2} + x^{2^{m+1}} \right) \right)} \\ = \sum_{x \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n} \left((a^{2}+b+b^{2})x^{2^{m+1}+2} + bx^{2^{m+1}+1} + bx^{2^{m+2}} \right)}.$$

Define

$$Q(x) = (a^{2} + b + b^{2})x^{2^{m+1}+2} + bx^{2^{m+1}+1} + bx^{2^{m}+2}.$$

It is clear that when (a,b) = (0,0), $W_f(a,b) = 0$. For $(a,b) \neq (0,0)$, the bilinear form of $\varphi_{a,b}(x) = \text{Tr}_n(Q(x))$ is given by

$$B_{\varphi_{a,b}}(x,y) = \varphi_{a,b}(x+y) + \varphi_{a,b}(x) + \varphi_{a,b}(y)$$

= $\operatorname{Tr}_n\left((a^2+b+b^2)\left(y^2x^{2^{m+1}}+y^{2^{m+1}}x^2\right) + b\left(yx^{2^{m+1}}+y^{2^{m+1}}x\right) + b\left(y^2x^{2^m}+y^{2^m}x^2\right)\right)$
= $\operatorname{Tr}_n\left(L_{a,b}(y)x^{2^{m+2}}\right),$

where

$$L_{a,b}(y) = b^4 y^8 + \left(b^{2^{m+2}} + b^4 + b^2 + a^4\right) y^4 + \left(b^{2^{m+2}} + b^{2^{m+1}} + b^2 + a^{2^{m+2}}\right) y^2 + b^{2^{m+1}} y^4$$

Let $\ker(L_{a,b}) = \{y \in \mathbb{F}_{2^n} : L_{a,b}(y) = 0\}$. According to (12), we have

$$W_f(a,b) = 2^n \sum_{y \in \ker(L_{a,b})} (-1)^{\varphi_{a,b}(y)} = \begin{cases} \pm 2^{\frac{n+d_{a,b}}{2}}, & \text{if } \varphi_{a,b}(x) = 0 \text{ for all } x \in \ker(L_{a,b}), \\ 0, & \text{otherwise,} \end{cases}$$

where $d_{a,b}$ is the dimension of ker $(L_{a,b})$. From the expression of $L_{a,b}$, it is obvious that $d_{a,b} \leq 3$. Moreover, since n + d must be even and n is odd, $d_{a,b} \in \{1,3\}$. Hence the Walsh transform

$$W_f(a,b) \in \left\{2^n, 0, \pm 2^{\frac{n+1}{2}}, \pm 2^{\frac{n+3}{2}}\right\}.$$

Next, we will show that when a = 0 the Walsh transform $W_f(0, b) = 0$ when the dimension of $L_{a,b}(y) = 0$ equals 3. Namely, there exists some $y_0 \in \ker(L_{a,b})$ such that $\varphi_{a,b}(y_0) = 1$. In this case, we have

$$L_b(y) = L_{0,b}(y) = b^4 y^8 + \left(b^{2^{m+2}} + b^4 + b^2\right) y^4 + \left(b^{2^{m+2}} + b^{2^{m+1}} + b^2\right) y^2 + b^{2^{m+1}} y.$$

Denote $\varphi_b(y) = \varphi_{0,b}(y) = \operatorname{Tr}_n \left(b \left(y^{2^{m+1}+2} + y^{2^{m+1}+1} + y^{2^m+2} + y^{2^m+1} \right) \right)$. Let $z = y^2 + y$. Then we have $L_b = b^4 z^4 + b^2 z^2 + b^{2^{m+2}} z^2 + b^{2^{m+1}} z = (b^2 z^2 + b^{2^{m+1}} z)^2 + b^2 z^2 + b^{2^{m+1}} z$.

If $d_{0,b} = 3$, i.e., the number of solutions of $L_b = 0$ equals 8, then the equation

$$(b^2 z^2 + b^{2^{m+1}} z)^2 + b^2 z^2 + b^{2^{m+1}} z = 0$$
(14)

has 4 solutions in \mathbb{F}_{2^n} since y and y + 1 correspond to one same $z = y^2 + y$. Clearly, from (14), we have $b^2 z^2 + b^{2^{m+1}} z = 0$ or $b^2 z^2 + b^{2^{m+1}} z = 1$. From $b^2 z^2 + b^{2^{m+1}} z = 0$, we get two solutions $z_0 = 0$ and $z_1 = b^{2^{m+1}-2}$ in \mathbb{F}_{2^n} . Similarly, we also obtain two solutions $z = z_2, z_3$ from $b^2 z^2 + b^{2^{m+1}} z = 1$. Thus if $d_b = 3$, $y^2 + y = z_i$ for i = 0, 1, 2, 3 exactly has two solutions in \mathbb{F}_{2^n} . Namely, $\operatorname{Tr}_n(z_i) = 0$ for i = 0, 1, 2, 3. Particularly, $\operatorname{Tr}_n(z_1) = \operatorname{Tr}_n(b^{2^{m+1}-2}) = 0$. Therefore, there exists some element $y_0 \in \mathbb{F}_{2^n}$ such that $b^{2^{m+1}-2} = y_0^2 + y_0$, i.e., $b = \frac{1}{(y_0^2 + y_0)^{2^m+1}}$. In fact, such y_0 belongs to $\ker(L_b)$ and is what we need. Indeed,

$$\begin{aligned} \varphi_b(y_0) &= \operatorname{Tr}_n \left(b \left(y_0^{2^{m+1}+2} + y_0^{2^{m+1}+1} + y_0^{2^m+2} + y_0^{2^m+1} \right) \right) \\ &= \operatorname{Tr}_n \left(\frac{y_0^{2^{m+1}+2} + y_0^{2^{m+1}+1} + y_0^{2^m+2} + y_0^{2^m+1}}{(y_0^2 + y_0)^{2^m+1}} \right) \\ &= \operatorname{Tr}_n(1) = 1. \end{aligned}$$

Hence, $W_f(b) = 0$ when the dimension of ker (L_b) is 3.

Next, we consider the parameters of C_f and $C_{D(f)}$, respectively.

(1) For the linear code C_f , since $W_f(a, b) = 2^n$ if and only if (a, b) = (0, 0), the dimension is 2n from (3). Moreover, the possible Hamming weights of codewords in C_f are given by

$$\mathrm{wt}(\mathbf{c}_{a,b}) \in \left\{0, 2^{n-1}, 2^{n-1} \pm 2^{\frac{n-1}{2}}, 2^{n-1} \pm 2^{\frac{n+1}{2}}\right\}$$

(2) For the linear code $C_{D(f)}$, $W_f(b) = 2^n$ if and only if b = 0, which means that the dimension of $C_{D(f)}$ is *n* according to (7). Since for any $b \in \mathbb{F}_{2^n}$, $W_f(b) \in \left\{2^n, 0, \pm 2^{\frac{n+1}{2}}\right\}$, by (8), the weights of the codewords \mathbf{c}_b in $\mathcal{C}_{D(f)}$ satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-2}, 0, 2^{n-2} - 2^{\frac{n-3}{2}}, 2^{n-2} + 2^{\frac{n-3}{2}} \right\}$.

Similar to the proof in Theorem 9 (2), the desired weight distribution of $C_{D(f)}$ can be obtained by solving (10) accordingly.

TABLE V The weight distribution of the codes C_f in Theorem 11

Weight	Multiplicity
0	1
$2^{n-1} - 2^m$	$2^{4m-1} + 2^{3m-1} - 2^{2m-1} - 2^{m-1}$
2^{n-1}	$2^{4m} + 2^{2m} - 1$
$2^{n-1} + 2^m$	$2^{4m-1} + 2^{m-1} - 2^{3m-1} - 2^{2m-1}$

Theorem 11. Let n = 2m + 1 and $f(x) = x^{2^{m+1}+2} + x^{2^{m+1}+1} + x^2 + x$. Define two linear codes C_f and $C_{D(f)}$ as in (2) and (6), respectively. Then,

C_f is a [2ⁿ − 1, 2n − 1] binary linear code with weight distribution in Table V.
 C_{D(f)} is a [2^{n−1} − 1, n − 1] binary linear code with weight enumerator 1 + (2^{n−1} − 1)z^{2^{n−2}}.

Proof. Firstly, we shall compute the value $W_f(a, b)$. It is clear that when (a, b) = (0, 0), $W_f(a, b) = 2^n$. For $(a, b) \neq (0, 0)$, let $\varphi_b(x) = \text{Tr}_n(bf(x))$. Since f is quadratic, according to (13), we need to determine the dimension of kernel of the bilinear form of $\varphi_b(x)$. Note that the bilinear form of $\varphi_b(x)$ is given by

$$B_{\varphi_b}(x,y) = \varphi_b(x+y) + \varphi_b(x) + \varphi_b(y)$$

= $\operatorname{Tr}_n\left(b\left(y^2x^{2^{m+1}} + y^{2^{m+1}}x^2\right) + b\left(yx^{2^{m+1}} + y^{2^{m+1}}x\right)\right) = \operatorname{Tr}_n\left(L_b(y)x^{2^{m+1}}\right)$

where

$$L_b(y) = \left(b^{2^{m+1}} + b\right) y^2 + \left(b^{2^m} + b\right) y.$$

Since gcd(m, n) = gcd(m + 1, n) = 1, $b^{2^{m+1}} + b = 0$ and $b^{2^m} + b = 0$ both have two solutions b = 0, 1 in \mathbb{F}_{2^n} . Thus $L_b(y) = 0$ if and only if b = 0, 1. Moreover, according to the expression of $W_f(a, b)$, we have

$$W_f(a,1) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n\left(ax + x^{2^{m+1}+2} + x^{2^{m+1}+2}\right)} = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(ax)}$$

equals 2^n if a = 0 and 0 otherwise.

When $b \neq 0, 1$, then it is clear that $\ker(L_b) = \left\{0, \frac{b^{2^m} + b}{b^{2^m+1} + b}\right\}$. It follows from (13) that

$$W_f(a,b) = \begin{cases} \pm 2^{\frac{n+1}{2}}, & \text{if } \operatorname{Tr}_n(bf(x) + ax) = 0 \text{ for all } x \in \left\{0, \frac{b^{2^m} + b}{b^{2^m+1} + b}\right\}, \\ 0, & \text{otherwise.} \end{cases}$$

Let $y_0 = \frac{b^{2^m} + b}{b^{2^m+1} + b} = \left(b^{2^{m+1}} + b\right)^{2^m - 1}$. Note that $y_0^{2^{m+1} + 2} = \frac{1}{b^{2^m+1}}$

$$y_0^{2^{m+1}+2} = \frac{1}{b^{2^{m+1}}+b}, \quad y_0^{2^{m+1}+1} = \frac{1}{b^{2^m}+b}$$

and

$$Tr_n \left(by_0^{2^{m+1}+2} + by_0^{2^{m+1}+1} \right)$$

= $Tr_n \left(\frac{b}{b^{2^{m+1}} + b} + \frac{b}{b^{2^m} + b} \right)$
= $Tr_n \left(\frac{b^{2^m} + b}{b^{2^m} + b} \right) = 1.$

Moreover, denote $b^{2^{m+1}} + b = \Delta$. Then $y_0 = \Delta^{2^m - 1}$, $b^{2^{m+1}} + b^2 = \Delta^{2^{m+1}}$ and $b^2 + b = \Delta + \Delta^{2^{m+1}}$. Thus

$$\operatorname{Tr}_{n}\left(\left(b+b^{2}\right)y_{0}^{2}\right) = \operatorname{Tr}_{n}\left(\frac{\left(\Delta+\Delta^{2^{m+1}}\right)\Delta^{2^{m+1}}}{\Delta^{2}}\right) = \operatorname{Tr}_{n}\left(\Delta^{2^{m+1}-1}+\Delta^{2^{m+2}-2}\right) = 0.$$

Therefore,

$$\operatorname{Tr}_{n}(bf(y_{0})) = \operatorname{Tr}_{n}\left(by_{0}^{2^{m+1}+2} + by_{0}^{2^{m+1}+1} + (b+b^{2})y_{0}^{2}\right) = 1 + 0 = 1$$

and

$$\operatorname{Tr}_n(bf(y_0) + ay_0) = 1 + \operatorname{Tr}_n(ay_0)$$

Clearly, $\operatorname{Tr}_n(bf(y_0) + ay_0) \in \{0, 1\}$ and thus for any $b \neq 0, 1$,

$$W_f(a,b) \in \left\{0, \pm 2^{\frac{n+1}{2}}\right\}.$$

Moreover, since $\operatorname{Tr}_n(bf(y_0)) = 1$ here, $W_f(b) = 0$ for $b \in \mathbb{F}_{2^n} \setminus \{1\}$.

(1) For the linear code C_f , since $W_f(a,b) = 2^n$ if and only if (a,b) = (0,1), (0,0), the dimension is 2n-1 from (3). Moreover, for any $a, b \in \mathbb{F}_{2^n}$, $W_f(a,b) \in \left\{0, 2^n, \pm 2^{\frac{n+1}{2}}\right\}$. Let

$$v_1 = -2^{\frac{n+1}{2}}, \quad v_2 = 0, \quad v_3 = 2^{\frac{n+1}{2}}.$$

Then by computing (5), we can get that the occurrences of $W_f(a, b) = v_i$'s, i = 1, 2, 3 in the Walsh spectrum of f are

$$\begin{cases} X_1 = 2^{4m} + 2^m - 2^{3m} - 2^{2m} \\ X_2 = 2^{4m+1} + 2^{2m+1} - 2 \\ X_3 = 2^{4m} + 2^{3m} - 2^{2m} - 2^m \end{cases}$$

Finally, by (4), the desired weight distribution of C_f can be obtained.

(2) For the linear code $C_{D(f)}$, since there are two b's (0 and 1) such that $W_f(b) = 2^n$ and $(2^n - 2)$ b's such that $W_f(b) = 0$, the dimension of $C_{D(f)}$ equals n - 1 by (7). Moreover, by (8), the Hamming weights of the codewords \mathbf{c}_b in $C_{D(f)}$ satisfy

$$\operatorname{wt}(\mathbf{c}_b) \in \{0, 2^{n-2}\}.$$

The desired weight enumerator thus follows.

TABLE VI THE WEIGHT DISTRIBUTION OF THE CODES C_f in Theorem 13

Weight	Multiplicity
0	1
$2^{n-1} - 2^{2m-1}$	$2^{4m-1} + 2^{3m-1} - 2^{2m-1} - 2^{m-1}$
2^{n-1}	$2^{5m} + 2^{2m} - 2^{4m} - 1$
$2^{n-1} + 2^{2m-1}$	$2^{4m-1} + 2^{m-1} - 2^{3m-1} - 2^{2m-1}$

Theorem 12. Let n = 2m + 1 and $f(x) = x^{2^{m+2}+4} + x^{2^{m+1}+2} + x^2 + x$. Define two linear codes C_f and $C_{D(f)}$ as in (2) and (6), respectively. Then,

C_f is a [2ⁿ − 1, 2n − 1] binary linear code with weight distribution in Table V.
 C_{D(f)} is a [2^{n−1} − 1, n − 1] binary linear code with weight enumerator 1 + (2^{n−1} − 1)z^{2^{n−2}}.

Proof. The proof is very similar with that of Theorem 11 and we omit it here.

Example 3. When m = 3, the codes C_f in Theorems 11 and 12 are two [127, 13, 56] binary linear codes with the same weight enumerator

$$1 + 2268z^{56} + 4159z^{64} + 1764z^{72}$$

Referring to the code table [16], the linear codes C_f have the best known parameter. When m = 3, the codes $C_{D(f)}$ in Theorems 11 and 12 are two [63, 6, 32] binary linear codes with the same weight enumerator

$$1 + 63z^{32}$$

Referring to the code table [16], the linear codes $C_{D(f)}$ are optimal.

C. The case n = 3m

In this subsection, we consider binary linear codes from the first two-to-one polynomial in Lemma 7. The second one will be generalized in Section 4 and the corresponding linear code will be discussed later.

Theorem 13. Let n = 3m with $m \equiv 0 \pmod{3}$ and $f(x) = x^{2^{2m}+1} + x^{2^{m+1}} + x^{2^m+1} + x$. Define two linear codes C_f and $C_{D(f)}$ as in (2) and (6), respectively. Then,

(1) C_f is a [2ⁿ - 1, 5m] binary linear code with weight distribution in Table VI.
(2) C_{D(f)} is a [2ⁿ⁻¹ - 1, n - 1] binary linear code with weight enumerator 1 + (2ⁿ⁻¹ - 1)z^{2ⁿ⁻²}.

Proof. Firstly, we shall compute the value $W_f(a, b)$. Clearly, when (a, b) = (0, 0), $W_f(a, b) = 2^n$. For $(a, b) \neq (0, 0)$, let $\varphi_b(x) = \text{Tr}_n(bf(x))$. If (a, b) = (a, 1), then

$$W_f(a,b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n\left(ax + x^{2^{2m}+1} + x^{2^{m+1}} + x^{2^{m+1}} + x^{2^{m+1}} + x\right)} = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(ax)},$$

which equals 2^n if a = 0 and 0 otherwise. Moreover, if $b \in \mathbb{F}_{2^m} \setminus \mathbb{F}_2$, then

$$W_f(a,b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n\left(\left(a^{2^{m+1}} + b + b^2\right)x^{2^{m+1}}\right)},$$

which equals 2^n if $a^{2^{m+1}} + b + b^2 = 0$ and 0 otherwise. Next, we assume $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$.

Since $\text{Tr}_n(f)$ is a quadratic function, according to (13), we need to determine the dimension of kernel of the bilinear form of $\varphi_b(x)$. Note that the bilinear form of $\varphi_b(x)$ is given by

$$B_{\varphi_b}(x,y) = \varphi_b(x+y) + \varphi_b(x) + \varphi_b(y)$$

= $\operatorname{Tr}_n\left(\left(b+b^{2^m}\right)\left(yx^{2^m}+y^{2^m}x\right)\right) = \operatorname{Tr}_n\left(L_b(y)x^{2^m}\right),$

where $L_b(y) = (b + b^{2^m}) y + (b^{2^m} + b^{2^{2^m}}) y^{2^{2^m}}$. Let $\ker(L_b) = \{y : y \in \mathbb{F}_{2^n} \text{ and } L_b(y) = 0\}$. Since $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$, we have

$$\ker(L_b) = \left\{ \left(b^{2^m} + b^{2^{2m}} \right) \eta : \eta \in \mathbb{F}_{2^m} \right\}.$$

It follows from (13) that

$$W_f(a,b) = \begin{cases} \pm 2^{2m}, & \text{if } \operatorname{Tr}_n(bf(x) + ax) = 0 \text{ for all } x \in \ker(L_b), \\ 0, & \text{otherwise.} \end{cases}$$

For any $x = (b^{2^m} + b^{2^{2m}}) \eta \in \ker(L_b)$ with $\eta \in \mathbb{F}_{2^m}$,

$$\operatorname{Tr}_{n}(bf(x) + ax) = \operatorname{Tr}_{n}\left(U_{a,b}\eta^{2}\right) = \operatorname{Tr}_{m}\left(\operatorname{Tr}_{n/m}(U_{a,b})\eta^{2}\right).$$

where

$$U_{a,b} = a^2 \left(b + b^{2^m}\right)^{2^{m+1}} + \left(b + b^{2^m}\right)^{2^{2m} + 2^m + 1} + \left(b + b^{2^m}\right)^{2^{2m+1}} \left(b^{2^{m+1}} + b\right).$$

Clearly, $\operatorname{Tr}_n(bf(x) + ax) = 0$ if and only if $a^2 = (b + b^{2^m})^{2^{2m} - 2^m + 1} + (b + b^{2^m})^{2^{2m+1} - 2^{m+1}} (b^{2^{m+1}} + b)$. Thus for any $a \in \mathbb{F}_{2^n}, b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$,

$$W_f(a,b) \in \left\{0, \pm 2^{2m}\right\}.$$

As for $W_f(b)$, we need the following claim which will be showed at the end of the proof.

Claim. For any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$, $\operatorname{Tr}_{n/m}(U_b) \neq 0$, where $U_b = U_{0,b}$.

According to the above claim, it is clear that $W_f(b) = 0$ for any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$.

Next, we consider the parameters of C_f and $C_{D(f)}$, respectively.

(1) For the linear code C_f , since $W_f(a, b) = 2^n$ if and only if $b \in \mathbb{F}_{2^m}$ and $a^{2^{m+1}} + b + b^2 = 0$, by (3), the dimension of $K_1 = \{a, b \in \mathbb{F}_{2^n} : W_f(a, b) = 2^n\}$ is m and thus the dimension of C_f is 2n - m = 5m. Moreover, for any $a, b \in \mathbb{F}_{2^n}$, $W_f(a, b) \in \{0, 2^n, \pm 2^{2m}\}$. Let

$$v_1 = -2^{2m}, \quad v_2 = 0, \quad v_3 = 2^{2m}.$$

21

Then by computing (5), we can obtain that the occurrences of $W_f(a, b) = v_i$'s, i = 1, 2, 3 in the Walsh spectrum of f are

$$\begin{cases} X_1 = 2^{5m-1} + 2^{2m-1} - 2^{4m-1} - 2^{3m-1} \\ X_2 = 2^{6m} + 2^{3m} - 2^{5m} - 2^m \\ X_3 = 2^{5m-1} + 2^{4m-1} - 2^{3m-1} - 2^{2m-1}. \end{cases}$$

Finally, by (4), the desired weight distribution of C_f can be obtained.

(2) For the linear code $C_{D(f)}$, since there are two b's (0 and 1) such that $W_f(b) = 2^n$ and $(2^n - 2)$ b's such that $W_f(b) = 0$, by (7), the dimension of $C_{D(f)}$ equals n - 1. Moreover, by (8), we know that the weights of the codewords \mathbf{c}_b in $C_{D(f)}$ satisfy

$$\operatorname{wt}(\mathbf{c}_b) \in \{0, 2^{n-2}\}$$

Furthermore, the weight enumerator is clear.

Finally, we prove the claim, i.e., for any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$, $\operatorname{Tr}_{n/m}(U_b) \neq 0$. By direct computation, we have

$$\operatorname{Tr}_{n/m}(U_b) = b^3 + b^{3 \cdot 2^m} + b^{3 \cdot 2^{2m}} + b^{2^{2m+1} + 2^m} + b^{2^{2m} + 2} + b^{2^{m+1} + 1}.$$
(15)

For any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$, define

$$\begin{cases} b + b^{2^m} + b^{2^{2m}} = \alpha \\ b^{2^m + 1} + b^{2^{2m} + 1} + b^{2^{2m} + 2^m} = \beta \\ b^{2^{2m} + 2^m + 1} = \gamma \end{cases}$$

and $g(x) = x^3 + \alpha x^2 + \beta x + \gamma \in \mathbb{F}_{2^m}[x]$. Then it is clear that $g(x) = (x+b)(x+b^{2^m})(x+b^{2^{2m}})$ and g(x) is irreducible on \mathbb{F}_{2^m} . Let $u = b^{2^{2m+1}+2^m} + b^{2^{2m}+2} + b^{2^{m+1}+1}$ and $v = b^{2^{2m+1}+1} + b^{2^m+2} + b^{2^{m+1}+2^{2m}}$. Then we have

$$u + v = (b + b^{2^{m}}) (b + b^{2^{2m}}) (b^{2^{m}} + b^{2^{2m}})$$

= $(\alpha + b) (\alpha + b^{2^{m}}) (\alpha + b^{2^{2m}}) = g(\alpha) = \alpha\beta + \gamma.$

In addition, from the expanded form of $(b + b^{2^m} + b^{2^{2m}})^3$, we know that

$$b^{3} + b^{3 \cdot 2^{m}} + b^{3 \cdot 2^{2m}} = \alpha^{3} + u + v = \alpha^{3} + \alpha\beta + \gamma.$$

Moreover, since $b, b^{2^m}, b^{2^{2m}}$ are the roots of g(x) in $\mathbb{F}_{2^n}, \frac{1}{b}, \frac{1}{b^{2^m}}, \frac{1}{b^{2^{2m}}}$ are the roots of

$$g'(x) = \frac{1}{\gamma} x^3 g\left(\frac{1}{x}\right) = x^3 + \frac{\beta}{\gamma} x^2 + \frac{\alpha}{\gamma} x + \frac{1}{\gamma}.$$

in \mathbb{F}_{2^n} . Similarly, we have

$$\frac{1}{b^3} + \frac{1}{b^{3 \cdot 2^m}} + \frac{1}{b^{3 \cdot 2^{2m}}} = \frac{\beta^3}{\gamma^3} + \frac{\alpha\beta}{\gamma^2} + \frac{1}{\gamma} = \frac{\beta^3 + \alpha\beta\gamma + \gamma^2}{\gamma^3}.$$

Furthermore,

$$\begin{split} uv &= \left(b^{2^{2m+1}+2^m} + b^{2^{2m}+2} + b^{2^{m+1}+1}\right) \left(b^{2^{2m+1}+1} + b^{2^m+2} + b^{2^{m+1}+2^{2m}}\right) \\ &= \gamma^2 + \gamma \left(b^3 + b^{3\cdot 2^m} + b^{3\cdot 2^{2m}}\right) + b^{3\cdot (2^m+1)} + b^{3\cdot (2^{2m}+2^m)} + b^{3\cdot (2^{2m}+1)} \\ &= \gamma^2 + \gamma \left(\alpha^3 + \alpha\beta + \gamma\right) + \gamma^3 \left(\frac{1}{b^3} + \frac{1}{b^{3\cdot 2^m}} + \frac{1}{b^{3\cdot 2^{2m}}}\right) \\ &= \alpha^3 \gamma + \alpha\beta\gamma + \beta^3 + \alpha\beta\gamma + \gamma^2 \\ &= \alpha^3 \gamma + \beta^3 + \gamma^2. \end{split}$$

Now we go back to the expression of $\operatorname{Tr}_{n/m}(U_b)$, i.e., (15). If $\operatorname{Tr}_{n/m}(U_b) = 0$, we have

$$u = b^3 + b^{3 \cdot 2^m} + b^{3 \cdot 2^{2m}} = \alpha^3 + \alpha\beta + \gamma$$

and then

$$v = \alpha\beta + \gamma + u = \alpha^3.$$

Thus $uv = \alpha^6 + \alpha^4 \beta + \alpha^3 \gamma = \alpha^3 \gamma + \beta^3 + \gamma^2$, namely,

$$\alpha^6 + \alpha^4 \beta + \beta^3 + \gamma^2 = 0. \tag{16}$$

Next, we show that under (16), g(x) = 0 has three solutions in \mathbb{F}_{2^m} , which is in contradiction with the irreduciblity of g(x). Firstly, using $x + \alpha$ to replace x in g(x) = 0 and simplifying it, we obtain

$$x^{3} + (\alpha^{2} + \beta)x + \alpha\beta + \gamma = 0.$$
(17)

Moreover,

$$\operatorname{Tr}_m\left(\frac{(\alpha^2+\beta)^3}{(\alpha\beta+\gamma)^2}\right) = \operatorname{Tr}_m\left(\frac{\alpha^6+\beta^3+\alpha^4\beta+\alpha^2\beta^2}{(\alpha\beta+\gamma)^2}\right) = \operatorname{Tr}_m(1),$$

where the last equality is derived from (16). Furthermore, it is easy to get that the equation

$$t^2 + (\alpha\beta + \gamma)t + (\alpha^2 + \beta)^3 = 0$$

has a solution $t_1 = (\alpha\beta + \gamma)\omega$, where $\omega^3 = 1$. Since $m \equiv 0 \pmod{3}$, ω is a cube in \mathbb{F}_{2^m} (*m* even), $\mathbb{F}_{2^{2m}}$ (*m* odd). In addition, $\alpha\beta + \gamma = \sqrt{(\alpha^2 + \beta)^3}$ is also a cube in \mathbb{F}_{2^m} . Thus according to Lemma 2, (17) has three solutions in \mathbb{F}_{2^m} , which is a contradiction, and thus $\operatorname{Tr}_{n/m}(U_b) \neq 0$.

4. BINARY LINEAR CODES FROM NEW TWO-TO-ONE POLYNOMIALS

In this section, we construct two new classes of two-to-one functions, of which the first one is a generalization of (2) in Lemma 7. Then we also obtain some binary linear codes C_f and $C_{D(f)}$ from these two new two-to-one functions.

Theorem 14. Let n = km with k, m odd and $f(x) = \text{Tr}_{n/m}(x^{2^m+1}) + x$. Then f(x) is two-to-one over \mathbb{F}_{2^n} .

Proof. According to the definition of two-to-one functions, it suffices to prove that for any $a \in \mathbb{F}_{2^n}$, $|f^{-1}(a)| \in \{0, 2\}$. Namely, for any $a \in \mathbb{F}_{2^n}$, f(x + a) + f(a) = 0 has exactly two solutions in \mathbb{F}_{2^n} . By simplifying the equation, we have

$$\operatorname{Tr}_{n/m}\left(x^{2^{m}+1} + ax^{2^{m}} + a^{2^{m}}x\right) = x.$$
(18)

This implies $x \in \mathbb{F}_{2^m}$, and then (18) becomes $x^2 + x = 0$, which has exactly two solutions x = 0, 1.

Theorem 15. Let n = 3m with m odd and $f(x) = x^{2^{2m+1}+1} + x^{2^{m+1}+1} + x^4 + x^3$. Then f(x) is two-to-one over \mathbb{F}_{2^n} .

Proof. It suffices to prove that for any $a \in \mathbb{F}_{2^n}$, the equation f(x+a) + f(a) = 0,

$$x^{2^{2m+1}+1} + ax^{2^{2m+1}} + x^{2^{m+1}+1} + ax^{2^{m+1}} + x^4 + x^3 + ax^2 + \left(a^{2^{2m+1}} + a^{2^{m+1}} + a^2\right)x = 0,$$
 (19)

has exactly two solutions in \mathbb{F}_{2^n} . In fact, since x = 0 is clearly a solution of (19), we shall only show that (19) has at most two solutions in \mathbb{F}_{2^n} .

Let $y = x^{2^{m}}$, $z = y^{2^{m}}$, $b = a^{2^{m}}$ and $c = b^{2^{m}}$. Then (19) becomes

$$xz^{2} + xy^{2} + x^{3} + x^{4} + a(x^{2} + y^{2} + z^{2}) + (a^{2} + b^{2} + c^{2})x = 0.$$
 (20)

Raising (20) to the 2^m -th power and the 2^{2m} -th power, we get

$$yx^{2} + yz^{2} + y^{3} + y^{4} + b(x^{2} + y^{2} + z^{2}) + (a^{2} + b^{2} + c^{2})y = 0$$
(21)

and

$$zy^{2} + zx^{2} + z^{3} + z^{4} + c(x^{2} + y^{2} + z^{2}) + (a^{2} + b^{2} + c^{2})z = 0,$$
(22)

respectively. Let t = x + y + z and s = a + b + c. Computing the summation of (20), (21) and (22), we obtain

$$t^4 + t^3 + st^2 + s^2t = 0.$$

Thus t = 0 or $t^3 + t^2 + st + s^2 = 0$.

If t = 0, plugging it into (20), we have $x^4 + sx = 0$ and thus x = 0 or $x^3 = s$. It is clear that x = 0 is a solution of (19). If $x^3 = s = a + b + c \in \mathbb{F}_{2^m}$, then $x = s^{\frac{1}{3}} \in \mathbb{F}_{2^m}$ and $y = z = x \in \mathbb{F}_{2^m}$. Thus x = x + y + z = t = 0.

If $t^3 + t^2 + st + s^2 = 0$, using $(t_1 + 1)$ to replace t, we get

$$t_1^3 + (s+1)t_1 + s^2 + s = 0.$$
(23)

Since

$$\operatorname{Tr}_{m}\left(\frac{(s+1)^{3}}{(s^{2}+s)^{2}}\right)$$
$$= \operatorname{Tr}_{m}\left(\frac{1}{s}+\frac{1}{s^{2}}\right) = 0 \neq \operatorname{Tr}_{m}(1),$$

(23) has exactly one solution in \mathbb{F}_{2^m} according to Lemma 2. Moreover, we can get the expression of the unique solution by Lemma 3. For the equation $u^2 + (s^2 + s)u + (s + 1)^3 = 0$, we have

$$\left(\frac{u}{s^2+s}\right)^2 + \frac{u}{s^2+s} = \frac{1}{s} + \frac{1}{s^2}$$

and thus u = s + 1 is a solution and $\epsilon = (s + 1)^{\frac{1}{3}}$ is a solution of $x^3 = u$ since $gcd(3, 2^m - 1) = 1$. Furthermore, $r = \epsilon + \frac{a}{\epsilon} = (s + 1)^{\frac{1}{3}} + (s + 1)^{\frac{2}{3}}$ is a solution of (23) and thus

$$\bar{t} = r + 1 = \epsilon + \epsilon^2 + 1,$$

where $\epsilon = (s+1)^{\frac{1}{3}}$, is the unique solution of $t^3 + t^2 + st + s^2 = 0$. Namely, x + y + z equals a constant. Plugging $x + y + z = \overline{t}$ into (20), we get

$$x^4 + (s^2 + \bar{t}^2)x + a\bar{t}^2 = 0.$$
(24)

Next, using Lemma 4, we will prove that the above equation has two solutions in \mathbb{F}_{2^n} . However, we will also show that the two solutions can not satisfy $x + y + z = \overline{t}$ at the same time and thus (20) has at most one solution in this case. Together with the zero solution, (20) has at most two solutions in \mathbb{F}_{2^n} and thus f(x) is two-to-one.

Recall that $\epsilon^3 = s + 1$ and $\bar{t} = \epsilon + \epsilon^2 + 1$. Since $s^2 + \bar{t}^2 = \epsilon^6 + \epsilon^4 + \epsilon^2$, if $s^2 + \bar{t}^2 = 0$, then $\epsilon = 0$ clearly ($\epsilon^2 + \epsilon + 1 \neq 0$ due to *m* odd). Moreover, $s = \epsilon^3 + 1 = 1$ and t = 1. Thus if s = 1 and t = 1, (24) becomes $x^4 = a$, which has exactly one solution in \mathbb{F}_{2^n} . In the following, we assume that $s^2 + \bar{t}^2 \neq 0$. Let $f_1(r) = r^3 + (s^2 + \bar{t}^2)$. Then it is clear that $f_1 = (1, 2)$, which means that f_1 can factor as a product of a linear factor and an irreducible quadratic factor. Moreover,

$$r_1 = (s^2 + \bar{t}^2)^{\frac{1}{3}} = (s^2 + \epsilon^2 + \epsilon^4 + 1)^{\frac{1}{3}} = (\epsilon^2 + \epsilon^4 + \epsilon^6)^{\frac{1}{3}}$$

TABLE VIITHE WEIGHT DISTRIBUTION OF THE CODES C_f IN THEOREM 16

Weight	Multiplicity
0	1
$2^{n-1} - 2^{\frac{n+m-1}{2}}$	$2^{n-1} + 2^{\frac{n+m}{2}-1} - 2^{n-m-1} - 2^{\frac{n-m}{2}-1}$
2^{n-1}	$2^{n+m} + 2^{n-m} - 2^n - 1$
$2^{n-1} + 2^{\frac{n+m-1}{2}}$	$2^{\frac{n-m}{2}-1} + 2^{n-1} - 2^{\frac{n+m}{2}-1} - 2^{n-m-1}$

TABLE VIII THE WEIGHT DISTRIBUTION OF THE CODES $\mathcal{C}_{D(f)}$ in Theorem 16

Weight	Multiplicity
0	1
$2^{n-2} - 2^{\frac{n+m-4}{2}}$	$2^{n-m-1} + 2^{\frac{n-m-2}{2}}$
2^{n-2}	$2^n - 2^{n-m} - 1$
$2^{n-2} + 2^{\frac{n+m-4}{2}}$	$2^{n-m-1} - 2^{\frac{n-m-2}{2}}$

is the unique solution of $f_1(r) = 0$. Set $w_1 = a\bar{t}^2 \frac{r_1^2}{(s^2 + \bar{t}^2)^2}$. In addition,

$$\begin{aligned} \mathrm{Tr}_{n}(w_{1}) &= \mathrm{Tr}_{n}\left(\frac{a(\epsilon+\epsilon^{2}+1)^{2}}{(\epsilon^{2}+\epsilon^{4}+\epsilon^{6})^{\frac{4}{3}}}\right) \\ &= \mathrm{Tr}_{m}\left(\mathrm{Tr}_{n/m}\left(\frac{a\bar{t}^{2}}{(\epsilon^{2}+\epsilon^{4}+\epsilon^{6})^{\frac{4}{3}}}\right)\right) \\ &= \mathrm{Tr}_{m}\left(\frac{s\bar{t}^{2}}{(\epsilon^{2}+\epsilon^{4}+\epsilon^{6})^{\frac{4}{3}}}\right) \\ &= \mathrm{Tr}_{m}\left(\frac{(\epsilon^{3}+1)(\epsilon+\epsilon^{2}+1)^{2}}{(\epsilon^{2}+\epsilon^{4}+\epsilon^{6})^{\frac{4}{3}}}\right) \\ &= \mathrm{Tr}_{m}\left(\frac{(\epsilon^{2}+\epsilon+1)^{\frac{4}{3}}}{\epsilon^{\frac{8}{3}}}+\frac{(\epsilon^{2}+\epsilon+1)^{\frac{1}{3}}}{\epsilon^{\frac{2}{3}}}\right) = 0. \end{aligned}$$

Thus according to Lemma 4, (24) has exactly two solutions in \mathbb{F}_{2^n} , denoted by x_1, x_2 . Next, we show that the two solutions can not satisfy $x + y + z = \overline{t}$ at the same time. Clearly, there exist some $\alpha, \beta \in \mathbb{F}_{2^n}$ such that (24) becomes

$$(x^{2} + \alpha x + \beta) (x^{2} + \alpha x + \alpha^{2} + \beta) = 0$$

and by comparing the coefficient of x, we know that $\alpha^3 = (s^2 + \bar{t}^2) \neq 0$. In addition, by the Vieta theorem, $x_1 + x_2 = \alpha \neq 0$. Thus the two solutions can not satisfy $x + y + z = \bar{t}$ at the same time.

B. Binary linear codes from these new two-to-one functions

Theorem 16. Let n = km with k, m odd and $f(x) = \text{Tr}_{n/m}(x^{2^m+1}) + x$. Define two linear codes C_f and $C_{D(f)}$ as in (2) and (6), respectively. Then,

(1) C_f is a $[2^n - 1, n + m]$ binary linear code with weight distribution in Table VII.

(2) $C_{D(f)}$ is a $[2^{n-1}-1,n]$ binary linear code with weight distribution in Table VIII

Proof. Firstly, we shall determine the value $W_f(a, b)$. It is clear that when $(a, b) = (0, 0), W_f(a, b) = 2^n$. Note that $\operatorname{Tr}_n(b\operatorname{Tr}_{n/m}(x^{2^m+1})) = \operatorname{Tr}_n(\operatorname{Tr}_{n/m}(b)x^{2^m+1})$. Thus, if $\operatorname{Tr}_{n/m}(b) = 0$, then

$$W_f(a,b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n((a+b)x)}$$

which equals 2^n if a = b and 0 otherwise.

In the following, we consider the case that $\operatorname{Tr}_{n/m}(b) \neq 0$. Let $\varphi_b(x) = \operatorname{Tr}_n(bf(x))$. Since it is quadratic, according to (13), we need to determine the dimension of kernel of the bilinear form of $\varphi_b(x)$. Note that the bilinear form of $\varphi_b(x)$ is given by

$$B_{\varphi_b}(x,y) = \varphi_b(x+y) + \varphi_b(x) + \varphi_b(y)$$

= $\operatorname{Tr}_n\left(\operatorname{Tr}_{n/m}(b)\left(yx^{2^m} + y^{2^m}x\right)\right) = \operatorname{Tr}_n\left(L_b(y)x^{2^m}\right),$

where $L_b(y) = \operatorname{Tr}_{n/m}(b)(y+y^{2^{2m}})$. Then $\ker(L_b) = \mathbb{F}_{2^m}$. It follows from (13) that

$$W_f(a,b) = \begin{cases} \pm 2^{\frac{n+m}{2}}, & \text{if } \operatorname{Tr}_n(bf(x) + ax) = 0 \text{ for all } x \in \mathbb{F}_{2^m}, \\ 0, & \text{otherwise.} \end{cases}$$

Moreover, for $x \in \mathbb{F}_{2^m}$,

$$\operatorname{Tr}_n(bf(x) + ax) = \operatorname{Tr}_n((a^2 + b + b^2)x^2) = \operatorname{Tr}_m\left(\operatorname{Tr}_{n/m}(a^2 + b + b^2)x^2\right),$$

always equals 0 if and only if $\operatorname{Tr}_{n/m}(a^2+b+b^2) = 0$. Thus for any a, b satisfying $\operatorname{Tr}_{n/m}(a^2+b^2+b) \neq 0$, $W_f(a,b) \in \left\{0, \pm 2^{\frac{n+m}{2}}\right\}$. Similarly, for any b with $\operatorname{Tr}_{n/m}(b) \notin \{0,1\}$, $W_f(b) \in \left\{0, \pm 2^{\frac{n+m}{2}}\right\}$.

Next, we consider the parameters of C_f and $C_{D(f)}$, respectively.

(1) For the linear code C_f , since $W_f(a,b) = 2^n$ if and only if a = b with $\operatorname{Tr}_{n/m}(b) = 0$, by (3), the dimension of $K_1 = \{a, b \in \mathbb{F}_{2^n} : W_f(a,b) = 2^n\}$ is n-m and thus the dimension of C_f is 2n - (n-m) = n+m. Moreover, for any $a, b \in \mathbb{F}_{2^n}$, $W_f(a,b) \in \{0, 2^n, \pm 2^{\frac{n+m}{2}}\}$. Let

$$v_1 = -2^{\frac{n+m}{2}}, \quad v_2 = 0, \quad v_3 = 2^{\frac{n+m}{2}}.$$

Then by computing (5), we can obtain that the occurrences of $W_f(a, b) = v_i$'s, i = 1, 2, 3 in the Walsh

TABLE IXTHE WEIGHT DISTRIBUTION OF THE CODES $\mathcal{C}_{D(f)}$ IN THEOREM 17

Weight	Multiplicity
0	1
$2^{n-2} - 2^{\frac{n+2m-3}{2}}$	$2^{m-2} + 2^{\frac{m-3}{2}}$
2^{n-2}	$2^n - 2^{m-1} - 1$
$2^{n-2} + 2^{\frac{n+2m-3}{2}}$	$2^{m-2} - 2^{\frac{m-3}{2}}$

spectrum of f are

$$X_{1} = 2^{\frac{3n-3m}{2}-1} + 2^{2n-m-1} - 2^{\frac{3n-m}{2}-1} - 2^{2n-2m-1}$$
$$X_{2} = 2^{2n} + 2^{2n-2m} - 2^{2n-m} - 2^{n-m}$$
$$X_{3} = 2^{2n-m-1} + 2^{\frac{3n-m}{2}-1} - 2^{2n-2m-1} - 2^{\frac{3n-3m}{2}-1}.$$

Finally, by (4), the desired weight distribution of C_f can be obtained.

(2) For the linear code $C_{D(f)}$, since $W_f(b) = 2^n$ if and only if b = 0 and for $b \in \mathbb{F}_{2^n}$, by (7), the dimension of $C_{D(f)}$ equals n. Moreover, we have

$$W_f(b) \in \left\{0, 2^n, \pm 2^{\frac{n+m}{2}}\right\}$$

and by (8), the weights of the codewords \mathbf{c}_b in $\mathcal{C}_{D(f)}$ satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-2}, 0, 2^{n-2} - 2^{\frac{n+m-4}{2}}, 2^{n-2} + 2^{\frac{n+m-4}{2}} \right\}.$

In the following, we determine the weight distribution of $C_{D(f)}$. Define

$$w_1 = 2^{n-2} - 2^{\frac{n+m-4}{2}}, \quad w_2 = 2^{n-2}, \quad w_3 = 2^{n-2} + 2^{\frac{n+m-4}{2}}$$

Then solving (10) gives the desired weight distribution.

Theorem 17. Let n = 3m with m odd and $f(x) = x^{2^{2m+1}+1} + x^{2^{m+1}+1} + x^4 + x^3$. Define two linear codes C_f and $C_{D(f)}$ as in (2) and (6), respectively. Then,

(1) C_f is a $[2^n - 1, 2n]$ binary linear code with 5 weights. Moreover, the weights of the codewords \mathbf{c}_b in C_f satisfy

$$wt(\mathbf{c}_b) \in \left\{ 2^{n-1}, 0, 2^{n-1} - 2^{\frac{n+2m-1}{2}}, 2^{n-1} + 2^{\frac{n+2m-1}{2}}, 2^{n-1} - 2^{\frac{n+m-2}{2}}, 2^{n-1} + 2^{\frac{n+m-2}{2}} \right\}.$$

(2) $C_{D(f)}$ is a $[2^{n-1}-1,n]$ binary linear code with weight distribution in Table IX.

Proof. First of all, we shall determine the value $W_f(a, b)$. It is clear that when (a, b) = (0, 0), $W_f(a, b) = 2^n$. Let $\varphi_b(x) = \text{Tr}_n(bf(x))$. Since f is quadratic, according to (13), we need to determine the dimension of

$$B_{\varphi_b}(x,y) = \varphi_b(x+y) + \varphi_b(x) + \varphi_b(y)$$

= $\operatorname{Tr}_n\left(b(x^{2^{2m+1}}y + xy^{2^{2m+1}} + x^{2^{m+1}}y + xy^{2^{m+1}} + x^2y + xy^2)\right) = \operatorname{Tr}_n\left(L_b(y)x^2\right),$

where

$$L_b(y) = b^2 y^{2^{2m+2}} + b^{2^{2m}} y^{2^{2m}} + b^2 y^{2^{m+2}} + b^{2^m} y^{2^m} + b^2 y^4 + by.$$

Let the dimension of ker (L_b) be d_b . It follows from (13) that

$$W_f(a,b) = \begin{cases} \pm 2^{\frac{n+d_b}{2}}, & \text{if } \operatorname{Tr}_n(bf(x) + ax) = 0 \text{ for all } x \in \mathbb{F}_{2^m}, \\ 0, & \text{otherwise.} \end{cases}$$

Next, we consider the equation $L_b(y) = 0$, i.e.,

$$b^2 \operatorname{Tr}_{n/m}(y)^4 = \operatorname{Tr}_{n/m}(by)$$

Since $\operatorname{Tr}_{n/m}(y), \operatorname{Tr}_{n/m}(by) \in \mathbb{F}_{2^m}$, we have $b \in \mathbb{F}_{2^m}^*$ or $\operatorname{Tr}_{n/m}(y) = \operatorname{Tr}_{n/m}(by) = 0$.

Case 1: If $b \in \mathbb{F}_{2^m}^*$, then the equation $L_b(y) = 0$ becomes $\operatorname{Tr}_{n/m}(y) = 0$ or $\sqrt[3]{b^{-1}}$. Thus in this case, the number of solutions of $L_b(y)$ is 2^{2m+1} . Namely, $d_b = 2m + 1$. In the following, we show that there exist some $b \in \mathbb{F}_{2^m}^*$ such that the restriction of $\operatorname{Tr}_n(b(f(x)))$ on $\ker(L_b) = \left\{ y \in \mathbb{F}_{2^n} : \operatorname{Tr}_{n/m}(y) = 0 \text{ or } \sqrt[3]{b^{-1}} \right\}$ is the all-zero mapping or not, i.e., $W_f(b) \in \left\{ 0, \pm 2^{\frac{n+2m+1}{2}} \right\}$ for $b \in \mathbb{F}_{2^m}^*$. On one hand, if $\operatorname{Tr}_{n/m}(y) = 0$,

$$\operatorname{Tr}_n(bf(y)) = \operatorname{Tr}_n\left(b(y\operatorname{Tr}_{n/m}(y)^2 + y^4)\right) = \operatorname{Tr}_n(by^4) = \operatorname{Tr}_m\left(b\operatorname{Tr}_{n/m}(y^4)\right) = 0$$

On the other hand, if $\operatorname{Tr}_{n/m}(y) = \sqrt[3]{b^{-1}}$,

$$\begin{aligned} \operatorname{Tr}_{n}(bf(y)) &= \operatorname{Tr}_{n}\left(b(y\operatorname{Tr}_{n/m}(y)^{2} + y^{4})\right) \\ &= \operatorname{Tr}_{m}\left(\operatorname{Tr}_{n/m}\left(by\operatorname{Tr}_{n/m}(y)^{2}\right) + \operatorname{Tr}_{n/m}(by^{4})\right) \\ &= \operatorname{Tr}_{m}\left(b\operatorname{Tr}_{n/m}(y)^{3} + b\operatorname{Tr}_{n/m}(y)^{4}\right) \\ &= \operatorname{Tr}_{m}(1 + \operatorname{Tr}_{n/m}(y)) = 1 + \operatorname{Tr}_{m}\left(\sqrt[3]{b^{-1}}\right). \end{aligned}$$

Then $\varphi_b(y) = 0$ if and only if $\operatorname{Tr}_m\left(\sqrt[3]{b^{-1}}\right) = 1$ and thus the restriction of $\operatorname{Tr}_n(b(f(x)))$ on $\ker(L_b)$ is the all-zero mapping if and only if $\operatorname{Tr}_m\left(\sqrt[3]{b^{-1}}\right) = 1$. Therefore $W_f(b) \in \left\{0, \pm 2^{\frac{n+2m+1}{2}}\right\}$ and then clearly $W_f(a, b) \in \left\{0, \pm 2^{\frac{n+2m+1}{2}}\right\}$ in this case. **Case 2:** If $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$, then

$$\ker(L_b) = \left\{ y : y \in \mathbb{F}_{2^n} \text{ and } \operatorname{Tr}_{n/m}(y) = \operatorname{Tr}_{n/m}(by) = 0 \right\}.$$

For any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$, define

$$\begin{cases} b + b^{2^m} + b^{2^{2m}} = \alpha \\ b^{2^m + 1} + b^{2^{2m} + 1} + b^{2^{2m} + 2^m} = \beta \\ b^{2^{2m} + 2^m + 1} = \gamma \end{cases}$$

and $g(x) = x^3 + \alpha x^2 + \beta x + \gamma \in \mathbb{F}_{2^m}[x]$. Then it is clear that $g(x) = (x+b)(x+b^{2^m})(x+b^{2^{2m}})$ and g(x) is irreducible on \mathbb{F}_{2^m} . Since $g(x+\alpha) = x^3 + (\alpha^2 + \beta)x + \alpha\beta + \gamma$ is also irreducible, we have

$$\alpha^2 + \beta \neq 0$$
 and $\alpha\beta + \gamma \neq 0$.

In addition, for any fixed $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$, it is well known that for any $y \in \mathbb{F}_{2^n}$, there exist unique $y_0, y_1, y_2 \in \mathbb{F}_{2^m}$ such that $y = y_0 + y_1 b + y_2 b^2$. Then

$$\operatorname{Tr}_{n/m}(y) = y_0 + y_1 \alpha + y_2 \alpha^2 = 0$$

and

$$\begin{aligned} \operatorname{Tr}_{n/m}(by) &= \operatorname{Tr}_{n/m} \left(y_0 b + y_1 b^2 + y_2 b^3 \right) \\ &= \operatorname{Tr}_{n/m} \left((y_1 + y_2 \alpha) b^2 + (y_0 + y_2 \beta) b + y_2 \gamma \right) \\ &= \alpha y_0 + \alpha^2 y_1 + \left(\gamma + \alpha \beta + \alpha^3 \right) y_2 = 0. \end{aligned}$$

Plugging $y_0 = y_1 \alpha + y_2 \alpha^2$ into the above equation and simplifying it, we obtain $(\gamma + \alpha \beta)y_2 = 0$ and then $y_2 = 0$ since $\gamma + \alpha \beta \neq 0$. Thus

$$\ker(L_b) = \{(\alpha + b)\eta : \eta \in \mathbb{F}_{2^m}\}.$$

Clearly, in this case, the dimension of $\ker(L_b)$ is m. Moreover, for $x \in \ker(L_b)$,

$$Tr_n(bf(x) + ax) = Tr_n (b(xTr_{n/m}(x)^2 + x^4) + ax) = Tr_n((b + a^4)x^4)$$

= Tr_n ((b + a^4)(\alpha + b)^4\eta^4)
= Tr_m (Tr_{n/m} ((b + a^4)(\alpha + b)^4) \eta^4).

Obviously, if $a^4 = b$, the restriction of $\operatorname{Tr}_n(b(f(x)) + ax)$ on $\ker(L_b)$ is the all-zero mapping and thus $W_f(a,b) = \pm 2^{\frac{n+m}{2}}$.

Moreover, if a = 0, $\operatorname{Tr}_n(bf(x)) = \operatorname{Tr}_m\left(\operatorname{Tr}_{n/m}\left(b(\alpha+b)^4\right)\eta^4\right) = \operatorname{Tr}_m\left(U_b\eta^4\right)$, where

$$U_b = \operatorname{Tr}_{n/m}(b(\alpha^4 + b^4)) = \alpha^5 + b^5 + b^{5 \cdot 2^m} + b^{5 \cdot 2^{2m}}.$$

In the following, we will show that $U_b \neq 0$ for any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$. If there exist some $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$ such that $U_b = 0$, then by simplifying it, we get

$$(b+b^{2^m})(b+b^{2^{2^m}})(b^{2^m}+b^{2^{2^m}})(\alpha^2+\beta) = 0,$$

which is impossible since $b \notin \mathbb{F}_{2^m}$ and $\alpha^2 + \beta \neq 0$. Thus for any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_{2^m}$, $U_b \neq 0$ and then the restriction of $\operatorname{Tr}_n(b(f(x)))$ on $\ker(L_b)$ can not be the all-zero mapping. Thus $W_f(b) = 0$.

In conclusion, for any $a, b \in \mathbb{F}_{2^n}$,

$$W_f(a,b) \in \left\{0, 2^n, \pm 2^{\frac{n+2m+1}{2}}, \pm 2^{\frac{n+m}{2}}\right\}.$$

However, for any $b \in \mathbb{F}_{2^n}$,

$$W_f(b) \in \left\{0, 2^n, \pm 2^{\frac{n+2m+1}{2}}\right\}.$$

Next, we consider the parameters of C_f and $C_{D(f)}$, respectively.

(1) For the linear code C_f , since $W_f(a,b) = 2^n$ if and only if (a,b) = (0,0), by (3), the dimension of C_f is 2n. Moreover, since for any $a, b \in \mathbb{F}_{2^n}$, $W_f(a,b) \in \left\{0, 2^n, \pm 2^{\frac{n+2m+1}{2}}, \pm 2^{\frac{n+m}{2}}\right\}$, by (4), the weights of the codewords \mathbf{c}_b in C_f satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-1}, 0, 2^{n-1} - 2^{\frac{n+2m-1}{2}}, 2^{n-1} + 2^{\frac{n+2m-1}{2}}, 2^{n-1} - 2^{\frac{n+m-2}{2}}, 2^{n-1} + 2^{\frac{n+m-2}{2}} \right\}.$

(2) For the linear code $C_{D(f)}$, $W_f(b) = 2^n$ if and only if b = 0, which means that the dimension of $C_{D(f)}$ is *n* according to (7). Since for any $b \in \mathbb{F}_{2^n}$, $W_f(b) \in \left\{0, 2^n, \pm 2^{\frac{n+2m+1}{2}}\right\}$, by (8), the weights of the codewords \mathbf{c}_b in $C_{D(f)}$ satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-2}, 0, 2^{n-2} - 2^{\frac{n+2m-3}{2}}, 2^{n-2} + 2^{\frac{n+2m-3}{2}} \right\}$

In the following, we determine the weight distribution of $C_{D(f)}$. Define

$$w_1 = 2^{n-2} - 2^{\frac{n+2m-3}{2}}, \quad w_2 = 2^{n-2}, \quad w_3 = 2^{n-2} + 2^{\frac{n+2m-3}{2}}.$$

Then solving (10) gives the desired weight distribution.

Remark 18. In Theorems 11 - 13, the linear codes $C_{D(f)}$ has the same parameters as the shortened Hadamard codes, which is locally decodable code that provides a way to recover parts of the original message with high probability, while only looking at a small fraction of the received word. This property gives rise to applications in the computational complexity theory and in the CDMA communication system. The dual codes of $C_{D(f)}$ are the binary Hamming codes with parameters $[2^{n-1} - 1, 2^{n-1} - n, 3]$.

Remark 19. In [23] there are several other classes of two-to-one quadratic polynomials. The experiment results show that we can obtain 3-weight or 5-weight binary linear codes as well from generalized quadratic polynomials. Due to the similarities of the parameters of those codes and the proofs, we choose some representatives of them that are more difficult and omitted the others in this paper. In addition, the linear code C_f in Theorem 10 appears to be a 3-weight code by numerical results. Nevertheless, we didn't manage to prove it by the techniques used in this paper. We cordially invite interested readers to determine the weight distribution of the linear codes C_f in Theorem 10 and Theorem 17.

Problem 20. Determine the weight distribution of the linear codes C_f in Theorem 10 and Theorem 17.

According to the experiment results, we also have the following conjecture.

Conjecture 21. Let n = 2m + 1 and $f(x) = x^{3 \cdot 2^{m+1}} + x^{2^{m+2}+1} + x^{2^{m+1}+1} + x$. Then f(x) is two-to-one over \mathbb{F}_{2^n} . Moreover, when $m \ge 4$, the linear code $\mathcal{C}_{D(f)}$ has the parameters $\left[2^{n-1} - 1, n, 2^{n-1} - 2^{\frac{n-1}{2}}\right]$ and the weights of the codewords \mathbf{c}_b in $\mathcal{C}_{D(f)}$ satisfy

wt(**c**_b)
$$\in \left\{ 2^{n-2}, 0, 2^{n-2} - 2^{\frac{n-3}{2}}, 2^{n-2} + 2^{\frac{n-3}{2}}, 2^{n-2} - 2^{\frac{n-1}{2}}, 2^{n-2} + 2^{\frac{n-1}{2}} \right\}.$$

If possible, determine the weight distribution of the linear code $C_{D(f)}$.

5. BINARY LINEAR CODES FROM $(x^{2^t} + x)^e$

It is clear that the function $(x^{2^t} + x)^e$ with gcd(t, n) = 1 and $gcd(e, 2^n - 1)$ is a two-to-one function from \mathbb{F}_{2^n} to itself. In this section, we construct binary linear codes from two-to-one functions in this form.

Recall that given a two-to-one function f, the parameters of the linear codes $C_{D(f)}$ in (6) depend on the investigation of the value $W_f(b)$. We first present an interesting relation on $W_f(b)$ for $f(x) = (x^{2^t} + x)^e$ and the Walsh transform of $\operatorname{Tr}_n(x^e)$. Actually, we consider a general form $f(x) = P(\psi(x))$, where P is a permutation polynomial over \mathbb{F}_{2^n} and $\psi(x)$ is two-to-one with $\operatorname{Im}(\psi) = \{y \in \mathbb{F}_{2^n} : \operatorname{Tr}_n(y) = 0\}$.

Proposition 22. Let $f(x) = P(\psi(x))$, where P is a permutation polynomial over \mathbb{F}_{2^n} and $\psi(x)$ is two-to-one with $\operatorname{Im}(\psi) = \{y \in \mathbb{F}_{2^n} : \operatorname{Tr}_n(y) = 0\}$. Then for any $b \in \mathbb{F}_{2^n}^*$,

$$W_f(b) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(bP(y)+y)}$$

Proof. Let

$$T_0 = \{ y \in \mathbb{F}_{2^n} : \operatorname{Tr}_n(y) = 0 \}.$$

Take an element $a \in \mathbb{F}_{2^n}$ with $\operatorname{Tr}_n(a) = 1$. Then

$$T_1 := \{ y \in \mathbb{F}_{2^n} : \operatorname{Tr}_n(y) = 1 \} = \{ a + y : y \in T_0 \}.$$

For any $b \in \mathbb{F}_{2^n}^*$, the fact $\sum_{y \in \mathbb{F}_{2^n}} (-1)^{\operatorname{Tr}_n(bP(y))} = 0$ implies

$$\sum_{y \in T_0} (-1)^{\operatorname{Tr}_n(bP(y))} = -\sum_{y \in T_1} (-1)^{\operatorname{Tr}_n(bP(y))} = \sum_{y \in T_0} (-1)^{\operatorname{Tr}_n(bP(y+a)+a)}.$$

TABLE X KNOWN ALMOST BENT POWER FUNCTIONS x^e over \mathbb{F}_{2^n} , n odd

Functions	e	Conditions	References
Gold	$2^{i} + 1$	gcd(i,n) = 1	[15, 26]
Kasami	$2^{2i} - 2^i + 1$	gcd(i,n) = 1	[20]
Welch	$2^m + 3$	n = 2m + 1	[3, 18]
Niho-1	$2^m + 2^{\frac{m}{2}} - 1$	n = 2m + 1, m even	[18]
Niho-2	$2^m + 2^{\frac{3m+1}{2}} - 1$	n=2m+1, m odd	[18]

Thus,

$$W_{f}(b) = \sum_{x \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n}(bP(\psi(x)))} = 2 \sum_{y \in T_{0}} (-1)^{\operatorname{Tr}_{n}(bP(y))}$$

$$= \sum_{y \in T_{0}} (-1)^{\operatorname{Tr}_{n}(bP(y))} + \sum_{y \in T_{0}} (-1)^{\operatorname{Tr}_{n}(bP(y+a)+a)}$$

$$= \sum_{y \in T_{0}} (-1)^{\operatorname{Tr}_{n}(bP(y)+y)} + \sum_{y \in T_{0}} (-1)^{\operatorname{Tr}_{n}(bP(y+a)+y+a)}$$

$$= \sum_{y \in T_{0}} (-1)^{\operatorname{Tr}_{n}(bP(y)+y)} + \sum_{y \in T_{1}} (-1)^{\operatorname{Tr}_{n}(bP(y)+y)}$$

$$= \sum_{y \in \mathbb{F}_{2^{n}}} (-1)^{\operatorname{Tr}_{n}(bP(y)+y)}.$$

Remark 23. It is well known that given an integer e with $gcd(e, 2^n - 1) = 1$, the calculations of the weight distribution of C_f with $f(x) = x^e$, the Walsh spectrum of x^e , the cross-correlation distribution of m-sequences and their e-decimated sequences are equivalent. The relation has provided a great amount of interesting results which originated from cryptography, coding theory and sequence design. Proposition 22 exhibits a similar relation, which indicates the equivalence between the computation of the weight distribution of $C_{D(f)}$ for $f(x) = (x^{2^t} + x)^e$ and the Walsh spectrum of $Tr_n(x^e)$. In other words, any power function x^e , $gcd(e, 2^n - 1) = 1$, with t-valued Walsh spectrum can be employed to construct linear codes $C_{D(f)}$ with t nonzero weights.

Recently Li and Zeng in [24] surveyed the exponents e that allow for 3-valued, 4-valued, 5-valued Walsh spectra of x^e . All the exponents e listed in [24] with $gcd(e, 2^n - 1) = 1$ can be employed to generate binary linear codes $C_{D(f)}$ with few weights.

For simplicity, we only provide the result from almost bent functions over \mathbb{F}_{2^n} with n odd, which has three-valued Walsh spectrum $\left\{0, \pm 2^{\frac{n+1}{2}}\right\}$ [6]. The known almost bent exponents e is listed in Table X. From Proposition 22, we have the following theorem on the linear codes $\mathcal{C}_{D(f)}$ defined as in (6).

Theorem 24. Let n = 2m + 1, $f(x) = (x^{2^t} + x)^e$ with gcd(t, n) = 1 and e being one of the almost bent

TABLE XITHE WEIGHT DISTRIBUTION OF THE CODES $\mathcal{C}_{D(f)}$ IN THEOREM 24

Weight	Multiplicity
0	1
$2^{n-2} - 2^{m-1}$	$2^{n-2} + 2^{m-1}$
2^{n-2}	$2^{n-1} - 1$
$2^{n-2} + 2^{m-1}$	$2^{n-2} - 2^{m-1}$

exponents in Table X. Let $C_{D(f)}$ is defined as in (6). Then $C_{D(f)}$ is a $[2^{n-1}-1,n]$ binary linear code with weight distribution in Table XI.

Proof. From Proposition 22 and the almost bent property of x^e , we know that for $b \in \mathbb{F}_{2^n}^*$,

$$W_f(b) \in \left\{0, \pm 2^{\frac{n+1}{2}}\right\}.$$

Moreover, it is clear that $W_f(b) = 2^n$ if and only if b = 0, which means that the dimension of $C_{D(f)}$ is n according to (7). Furthermore, by (8), the weights of the codewords \mathbf{c}_b in $C_{D(f)}$ satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-2}, 0, 2^{n-2} - 2^{\frac{n-3}{2}}, 2^{n-2} + 2^{\frac{n-3}{2}} \right\}.$

Finally, define

$$w_1 = 2^{n-2} - 2^{\frac{n-3}{2}}, \quad w_2 = 2^{n-2}, \quad w_3 = 2^{n-2} + 2^{\frac{n-3}{2}}.$$

Then solving (10) gives the desired weight distribution.

As for the linear codes C_f defined as in (2), where $f = (x^2 + x)^e$ for examples, it seems hard to compute the Walsh transform $W_f(a, b)$ for any $a, b \in \mathbb{F}_{2^n}$. However, for the Gold function, we manage to determine its possible values.

Theorem 25. Let n = 2m + 1 and *i* be a positive integer with gcd(i, n) = 1. Let $f(x) = (x^{2^t} + x)^{2^{t+1}}$ with gcd(t, n) = 1. Define the linear code C_f as in (2). Then, C_f is a $[2^n - 1, 2n]$ binary code with five weights. Moreover, the weights of the codewords c_b in C_f satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-1}, 0, 2^{n-1} - 2^{\frac{n-1}{2}}, 2^{n-1} + 2^{\frac{n-1}{2}}, 2^{n-1} - 2^{\frac{n+1}{2}}, 2^{n-1} + 2^{\frac{n+1}{2}} \right\}.$

Proof. First of all, we shall compute the value $W_f(a, b)$. It is clear that when (a, b) = (0, 0), $W_f(a, b) = 2^n$. For $(a, b) \neq (0, 0)$, let $\varphi_b(x) = \text{Tr}_n(b(f(x)))$. Since f is quadratic, from (13), we need to compute the dimension of kernel of the bilinear form of $\varphi_b(x)$. Note that the bilinear form of $\varphi_b(x)$ is given by

$$B_{\varphi_b}(x,y) = \varphi_b(x+y) + \varphi_b(x) + \varphi_b(y)$$

= $\operatorname{Tr}_n\left(b\left(y^{2^t}x^{2^{i+t}} + y^{2^{i+t}}x^{2^t} + yx^{2^{i+t}} + y^{2^{i+t}}x + y^{2^t}x^{2^i} + y^{2^i}x^{2^t} + yx^{2^i} + y^{2^i}x\right)\right)$
= $\operatorname{Tr}_n\left(L_b(y)x^{2^{i+t}}\right),$

where

$$L_b(y) = \left(b^{2^t}y^{2^{t+2t}} + (b+b^{2^t})y^{2^{t+t}} + by^{2^t}\right)^{2^t} + b^{2^t}y^{2^{2t}} + (b+b^{2^t})y^{2^t} + by$$

Let $\ker(L_b) = \{y : y \in \mathbb{F}_{2^n} \text{ and } L_b(y) = 0\}$. Next we determine the dimension of $\ker(L_b)$. Let $y^{2^t} + y = z$. Then

$$L_b = b^{2^{i+t}} z^{2^{2^{i+t}}} + b^{2^i} z^{2^{2^i}} + b^{2^t} z^{2^t} + bz = 0,$$

which means $b^{2^i}z^{2^{2^i}} + bz = 0$ or 1. From $b^{2^i}z^{2^{2^i}} + bz = 0$, since gcd(i, n) = 1, we have $z = b^{-\frac{1}{2^{i+1}}}$. Thus $L_b(y) = 0$ has at most eight solutions in \mathbb{F}_{2^n} , namely, the dimension of $ker(L_b)$ is at most 3. Moreover, since $W_f(a,b) \in \left\{0, \pm 2^{\frac{n+d_b}{2}}\right\}$ and n is odd, where d_b is the dimension of $ker(L_b)$, $d_b = 1$ or 3. Therefore, for any $a, b \in \mathbb{F}_{2^n}$,

$$W_f(a,b) = \left\{0, 2^n, \pm 2^{\frac{n+1}{2}}, \pm 2^{\frac{n+3}{2}}\right\}.$$

Next, we consider the parameters of C_f . Since $W_f(a,b) = 2^n$ if and only if (a,b) = (0,0), by (3), the dimension of C_f is 2n. Moreover, since for any $a, b \in \mathbb{F}_{2^n}$, $W_f(a,b) \in \left\{0, 2^n, \pm 2^{\frac{n+1}{2}}, \pm 2^{\frac{n+3}{2}}\right\}$, by (4), the weights of the codewords \mathbf{c}_b in C_f satisfy

wt(
$$\mathbf{c}_b$$
) $\in \left\{ 2^{n-1}, 0, 2^{n-1} - 2^{\frac{n-1}{2}}, 2^{n-1} + 2^{\frac{n-1}{2}}, 2^{n-1} - 2^{\frac{n+1}{2}}, 2^{n-1} + 2^{\frac{n+1}{2}} \right\}.$

Moreover, according to the experiment results, we have the following conjecture.

Conjecture 26. Let n = 2m+1 and e be the almost bent exponents as given in Table X. Let $f(x) = (x^{2^t}+x)^e$ with gcd(t,n) = 1. Define the linear codes C_f as in (2). Then, the parameters of the linear codes C_f are the same as that in Theorem 25. If possible, determine the weight distribution of the linear codes C_f .

6. CONCLUSION

In this paper, we employed some known and new two-to-one functions in two generic constructions of binary linear codes. By investigating the Walsh transform of relevant quadratic functions, we obtained the possible Hamming weights of the codewords in constructed linear codes. The two-to-one functions with few-valued Walsh transforms are particularly studied. As a result, a large number of new binary codes with few weights are presented. Moreover, the weight distributions of the codes with one nonzero weight and with three nonzero weights are determined.

REFERENCES

- [1] AR Calderbank and JM Goethals. Three-weight codes and association schemes. *Philips J. Res*, 39(4-5):143–152, 1984.
- [2] Robert Calderbank and William M Kantor. The geometry of two-weight codes. Bulletin of the London Mathematical Society, 18(2):97–122, 1986.

- [3] Anne Canteaut, Pascale Charpin, and Hans Dobbertin. Binary m-sequences with three-valued crosscorrelation: a proof of Welch's conjecture. *IEEE Transactions on Information Theory*, 46(1):4–8, 2000.
- [4] Claude Carlet, Pascale Charpin, and Victor Zinoviev. Codes, bent functions and permutations suitable for DES-like cryptosystems. *Designs, Codes and Cryptography*, 15(2):125–156, 1998.
- [5] Claude Carlet, Cunsheng Ding, and Jin Yuan. Linear codes from perfect nonlinear mappings and their secret sharing schemes. *IEEE Transactions on Information Theory*, 51(6):2089–2102, 2005.
- [6] Florent Chabaud and Serge Vaudenay. Links between differential and linear cryptanalysis. In *Workshop* on the Theory and Application of of Cryptographic Techniques, pages 356–365. Springer, 1994.
- [7] P. Delsarte. On subfield subcodes of modified Reed-Solomon codes (Corresp.). *IEEE Transactions on Information Theory*, 21(5):575–576, 1975.
- [8] Cunsheng Ding. Linear codes from some 2-designs. *IEEE Transactions on information theory*, 61(6):3265–3275, 2015.
- [9] Cunsheng Ding. A construction of binary linear codes from boolean functions. *Discrete Mathematics*, 339(9):2288 2303, 2016.
- [10] Cunsheng Ding. A construction of binary linear codes from Boolean functions. *Discrete mathematics*, 339(9):2288–2303, 2016.
- [11] Cunsheng Ding and Harald Niederreiter. Cyclotomic linear codes of order 3. IEEE Transactions on information theory, 53(6):2274–2277, 2007.
- [12] Cunsheng Ding and Xuesong Wang. A coding theory construction of new systematic authentication codes. *Theoretical computer science*, 330(1):81–99, 2005.
- [13] Kelan Ding and Cunsheng Ding. A class of two-weight and three-weight codes and their applications in secret sharing. *IEEE Transactions on Information Theory*, 61(11):5835–5842, 2015.
- [14] Hans Dobbertin, Patrick Felke, Tor Helleseth, and Petri Rosendahl. Niho type cross-correlation functions via Dickson polynomials and Kloosterman sums. *IEEE Transactions on Information Theory*, 52(2):613–627, 2006.
- [15] Robert Gold. Maximal recursive sequences with 3-valued recursive cross-correlation functions (Corresp.). *IEEE transactions on Information Theory*, 14(1):154–156, 1968.
- [16] Markus Grassl. Bounds on the minimum distance of linear codes and quantum codes. Online available at http://www.codetables.de, 2007. Accessed on 2020-06-16.
- [17] Ziling Heng, Qin Yue, and Chengju Li. Three classes of linear codes with two or three weights. *Discrete Mathematics*, 339(11):2832–2847, 2016.
- [18] Henk DL Hollmann and Qing Xiang. A proof of the Welch and Niho conjectures on cross-correlations of binary *m*-sequences. *Finite Fields and Their Applications*, 7(2):253–286, 2001.
- [19] W Cary Huffman and Vera Pless. *Fundamentals of error-correcting codes*. Cambridge university press, 2010.
- [20] Tadao Kasami. The weight enumerators for several classes of subcodes of the 2nd order binary Reed-

Muller codes. Information and Control, 18(4):369-394, 1971.

- [21] Philip A Leonard and Kenneth S Williams. Quartics over $GF(2^n)$. Proceedings of the American Mathematical Society, pages 347–350, 1972.
- [22] Kangquan Li, Sihem Mesnager, and Longjiang Qu. Further study of 2-to-1 mappings over \mathbb{F}_{2^n} . *arXiv:1910.06654*, 2019.
- [23] Nian Li and Sihem Mesnager. Recent results and problems on constructions of linear codes from cryptographic functions. *Cryptography and Communications*, 2020.
- [24] Nian Li and Xiangyong Zeng. A survey on the applications of niho exponents. *Cryptography and Communications*, 11(3):509–548, may 2018.
- [25] Sihem Mesnager and Longjiang Qu. On two-to-one mappings over finite fields. *IEEE Transactions on Information Theory*, 65(12):7884–7895, 2019.
- [26] Kaisa Nyberg. Differentially uniform mappings for cryptography. In Workshop on the Theory and Application of of Cryptographic Techniques, pages 55–64. Springer, 1993.
- [27] Chunming Tang, Nian Li, Yanfeng Qi, Zhengchun Zhou, and Tor Helleseth. Linear codes with two or three weights from weakly regular bent functions. *IEEE Transactions on Information Theory*, 62(3):1166–1176, 2016.
- [28] Kenneth S Williams. Note on cubics over $GF(2^n)$ and $GF(3^n)$. Journal of Number Theory, 7(4):361–365, 1975.
- [29] Jin Yuan and Cunsheng Ding. Secret sharing schemes from three classes of linear codes. *IEEE Transactions on Information Theory*, 52(1):206–212, 2005.
- [30] Zhengchun Zhou, Nian Li, Cuiling Fan, and Tor Helleseth. Linear codes with two or three weights from quadratic bent functions. *Designs, Codes and Cryptography*, 81(2):283–295, 2016.