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2019

Cai, Q., Pratama, M., Alam, S., Ma, C., & Liu, J. (2020). Breakup of directed multipartite networks. *IEEE Transactins on Network Science and Engineering*, 7(3), 947-960.  
doi:10.1109/TNSE.2019.2894142

<https://hdl.handle.net/10356/144369>

<https://doi.org/10.1109/TNSE.2019.2894142>

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# Breakup of Directed Multipartite Networks

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**Abstract**—A complex network in reality often consists of profuse components each of which may suffer from unpredictable perturbations. Because the components of a network could be interdependent, therefore the failures of some components may trigger catastrophes to the whole network. It is thus pivotal to exploit the robustness of complex networks to perturbations. Existing studies on network robustness mainly deal with interdependent or multilayer networks, little work is done to investigate the robustness of multipartite networks which are an indispensable part of complex networks. Here we plumb the robustness of directed multipartite networks. To be specific, we exploit the robustness of bi-directed and unidirectional multipartite networks in face of random node failures. We respectively establish cascading and non-cascading models based on the largest connected component concept for depicting the dynamical processes on bi-directed and unidirectional multipartite networks subject to random node attacks. Based on our developed models, we respectively derive the corresponding percolation theories for mathematically computing the robustness of directed multipartite networks to random node failures. We theoretically unravel the first-order and second-order phase transition phenomena on the robustness of directed multipartite networks. The correctness of our developed theories coincide quite well with simulations on computer-generated multipartite networks.

**Index Terms**—Complex networks, directed multipartite networks, network robustness, percolation, largest connected component

## 1 INTRODUCTION

COMPLEX systems are ubiquitous in our daily life [1]. The form of a complex system ranges from the macroscopic level like the power grid systems [2], to the microscopic level like the metabolic systems [3]. A complex system in real world is usually composed of countless components, which makes it difficult to be controlled [4]. With the advent of network science, the situation of system control has been significantly improved by modeling a complex system as a network composed of vertices and edges where the vertices represent the system components while edges denote the relationships between components [5, 6]. Network modeling has been proved to be an effective instrument not only for system control [7–10] but also for data science [11, 12].

Due to the fact that the components of a complex system may suffer from internal and/or external perturbations which may induce the breakdown of the whole system, an effective method for predicting system stability so as to avoid potential catastrophe is therefore imperative [13–16]. As a consequence, network robustness analysis has emerged and is gaining momentum [17–20].

### 1.1 Previous Work

Network robustness analysis aims to investigate how robust a network is in face of perturbations. To this end,

tremendous efforts have been made along this line. Because the perturbations could occur to vertices and/or edges of a network, existing studies thus can be categorized into three groups, i.e., vertex level [17, 21–23] (Fig. 1(a)), edge level [19, 24–26] (Fig. 1(b)), and mixed level [27–31] (Fig. 1(c)), with their main ideas literally self-explained. From the perspective of perturbation manners, existing studies can be roughly classified as: network robustness to random failure [21] (every vertex is given the same probability to be attacked) and network robustness to target attacks [22, 32–34] (the probability for a vertex to be attacked depends on the importance of the focal vertex, see Fig. 1(d)). If the methods for analyzing network robustness are of concern, existing studies therefore could be archived into two groups: simulations based studies [17, 33, 35] (the curve in Fig. 1(e) is obtained by simulations) and theoretical analysis [19, 23, 36] (the curve in Fig. 1(e) is obtained by theoretical computing).

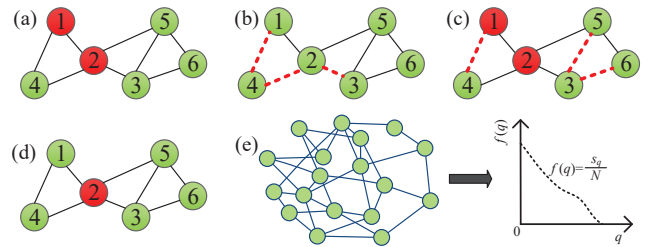


Fig. 1. Graphical examples of existing studies for network robustness analysis. (a) vertex level analysis. (b) edge level analysis. (c) mixed level analysis. (d) robustness to target attacks. (e) simulation based or theoretical analysis.

Due to the fact that a real-world network normally consists of many sub-networks which are interdependent, existing studies therefore can be divided into two branches, i.e., robustness of a single network [18] and robustness of interdependent networks [21, 23, 36, 37]. It has been discovered that a single network could be robust to perturbations

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(the curve in Fig. 1(e) is smooth), while an interdependent network or a network of networks could be extremely vulnerable [38–40] (the curve in Fig. 1(e) exhibits an abrupt jump from a finite value to zero).

Note that the above classification for existing studies on network robustness can be continued, depending on the angle from which the problem is viewed. For instance, in reality it is necessary to repair or recover a damaged network, thus existing studies can be categorized as: network robustness under attacks and network robustness under recoveries [33, 41, 42]. While many studies make efforts to answer the question of how robust a network is in face of perturbations, another research direction explores answers towards questions like what are the optimal structures of a network that are robust to perturbations and what kind of measures can be taken to enhance the robustness of a network [25, 34, 40, 43–45].

## 1.2 Motivation and Contribution

Multipartite networks are an essential part of complex networks [46–48]. The notion of multipartite network is the counterpart of monopartite or unipartite network [1, 2]. Although the works in [17, 26, 27, 29] have explored the robustness of multipartite networks, they are all empirical studies. Putting it another way, they can only tell whether a focal multipartite network is robust or not but cannot tell to what extent a multipartite network can survive perturbations.

To eliminate the deficiency of empirical studies on the robustness of multipartite networks, theoretical analysis on the robustness of bipartite networks which are a special case of multipartite networks have been erected [49–51]. However, on the one hand, the dynamics of bipartite networks as studied in [49–51] require domain-specific knowledge. On the other hand, the corresponding theories are only devised for bipartite networks and therefore are not applicable to multipartite network scenarios.

Note that theoretical methods like those in [21, 23, 36, 37] for analyzing the robustness of interdependent networks are mature, they are monopartite networks oriented and cannot be applied to handle multipartite networks. Though the latest work in [37] puts forward a mathematical method for analyzing the robustness of interdependent directed networks (still monopartite networks oriented), the method is not straightforward because it is developed in a manner similar to the seminal work in [21] where the robustness analysis involves recursive calculations of many transcendental equations.

To circumvent the above mentioned shortcomings, in this paper we present a precise yet direct theoretical method for analyzing the robustness of directed multipartite networks. The main contributions of this paper are threefold:

1) Network models for depicting the dynamical processes of bi-directed and unidirectional multipartite networks subject to vertex perturbations are respectively established. The structural differences between multipartite and monopartite networks render direct technology transfers from monopartite networks infeasible. The dynamic model used in [21] requires one-to-one interdependency, while the interdependency of a multipartite network is in a one-to-many mode.

2) Mathematical methods for calculating the proportions of vertices that eventually survive the perturbations occurred to multipartite networks with arbitrary degree distributions are accordingly developed. Although the theories developed in [23, 37, 52] are capable of handling interdependent networks with one-to-many interdependency, the transcendental equations involved in the calculation process require the variable  $P_i(k)$  which denotes the degree distribution of the vertices in the  $i$ -th network. Note that for an interdependent network, the vertices of the  $i$ -th network are interconnected, i.e.,  $P_i(k) \neq 0$ , while for a multipartite network there is no connections between vertices in the same partite set, i.e.,  $P_i(k) = 0$ . As a consequence, existing models and methods are not amenable to the robustness analysis of multipartite networks.

3) Our proposed theories unravel the first-order and second-order phase transition phenomena on the robustness of directed multipartite networks. Experimental simulations on random multipartite networks with Poisson degree distributions are carried out to validate the correctness of our proposed mathematical methods. The experiments coincide quite well with our theoretical results.

## 1.3 Paper Organization

The remainder of this paper is structured as follows. Section 2 presents the preliminaries including basic network notations, canonical network models for robustness analysis, and robustness evaluation metrics. Section 3 delineates in detail our proposed method for analyzing the robustness of multipartite network subject to random vertex loss. Section 4 validates the correctness of our proposed method through experiments on random multipartite networks with Poisson degree distributions. Section 5 concludes the paper.

## 2 PRELIMINARIES

### 2.1 Network Notation

Given a network denoted by  $G = \{V, E\}$ , where  $V$  and  $E$  respectively represent the sets of vertices and edges. We use  $e_{ij}$  to represent the edge between vertices  $i$  and  $j$ . If the vertex set  $V$  of a network  $G$  is composed of  $L$  different types of vertices, i.e.,  $V = \{v_1, v_2, \dots, v_L\}$ ,  $G$  is said to be a multipartite (or  $L$ -partite) network if the following condition is satisfied:

$$\begin{cases} v_i \cap v_j = \emptyset, & \forall i, j \wedge i \neq j \\ \exists e_{ij}, & \text{iff } i = j - 1 \vee i = j + 1 \end{cases}, \quad (1)$$

where  $v_i$  is called a partite set which contains  $n_i = |v_i|$  vertices. The total number of vertices of  $G$  is thus  $N = \sum n_i$ .

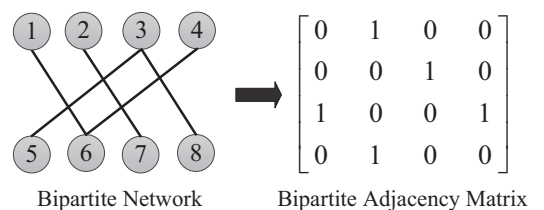


Fig. 2. Graphical illustration of a bipartite network and its bipartite matrix.

The left part of Fig. 2 exhibits an example of a multi-partite network. Because  $L = 2$  for the network shown in Fig. 2,  $G$  is therefore commonly called a bipartite network or a two-mode network [1]. The right part of Fig. 2 shows the bipartite matrix  $\mathbf{B}$  of network  $G$ . The entry  $b_{ij}$  of  $\mathbf{B}$  denotes the interaction between vertices  $i$  and  $j$ . Note that  $\mathbf{B}$  is generally asymmetric.

## 2.2 Network Robustness Analysis

Studies on network robustness aim to investigate the problem of to what extent a network can withstand perturbations occurring to vertices and/or edges. In the literature, two typical network models for depicting the dynamics of networks subject to perturbations are commonly utilized to analyze the robustness of networks. Fig. 3 takes the vertex robustness analysis as an example to delineate the two network models.

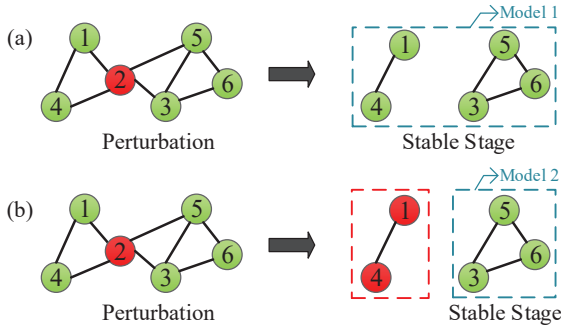


Fig. 3. Schematic illustrations of network models for depicting the dynamics of networks subject to perturbations. Network model 1 is based on the calculation of remaining vertices while network model 2 is based on the calculation of the largest connected component.

In Fig. 3(a), vertex 2 is removed from the network and the removal breaks the network into two clusters, i.e.,  $c_1 = \{1, 4\}$  and  $c_2 = \{3, 5, 6\}$ . The network model based on the calculation of remaining vertices is interested in vertices in clusters  $c_1$  and  $c_2$ . As shown in Fig. 3(b), the network model based on the calculation of the largest connected component (LCC) only concerns cluster  $c_2$  since it contains the largest number of vertices, and as a result cluster  $c_1$  will be disfunctional and removed.

## 2.3 Robustness Evaluation Metric

Fig. 3 exhibits two commonly used network models for network robustness analysis. In the literature, three metrics are accordingly proposed and widely adopted to quantitatively measure the robustness of a network. Other kinds of measures can be found in [53].

- 1) *Node Robustness Index*: Let  $s_q$  be the number of remaining vertices after removing a fraction  $q$  of vertices from a network  $G$ . The node robustness index  $R_n$  is defined as  $R_n = \frac{1}{N} \sum_{q=0}^1 s_q$ . The node robustness index is first used in [22]. The larger the value of  $R_n$ , the more robust the network is. Two successors can be found in [25, 54] where ref. [25] put forward a link robustness index  $R_l$  in a similar manner to  $R_n$  and ref. [54] developed a community

robustness index  $R_c$  which is a combination of  $R_n$  and  $R_l$ .

- 2) *Area Based Robustness Index*: For a focal network  $G$ , if the value of  $q$  is known, then we can easily obtain  $s_q$ . We thus can draw  $q$  and  $s_q$  in a 2-D space for all  $q \in [0, 1]$  and a curve will be yielded. The area based robustness index is then calculated as the area covered by the X-Y axis and the yielded curve. The larger the value of the area, the more robust the network is. The area based index is commonly used in the field of ecology, e.g., refs. [17] and [29] respectively make use of this index to measure the robustness of ecological networks to the loss of species and species community.
- 3) *LCC Based Index*: Given that a fraction  $1 - p$  of vertices are removed from  $G$ . The LCC based index quantifies the robustness of a network as the proportion  $P^\infty$  of vertices contained in the LCC. The robustness of a network in this context can be formulated as  $P^\infty = \frac{N_{LCC}}{N}$ , where  $N_{LCC}$  denotes the number of vertices in the LCC after removing a fraction  $1 - p$  of vertices from  $G$ . This kind of robustness evaluation index is widely used in the fields of network science and physics [36, 37].

## 3 METHODOLOGY

For an  $L$ -partite network, we randomly remove a fraction  $1 - p_i$  of vertices from  $v_i$  for all  $i \in [1, L]$ . Our purpose is to mathematically figure out the proportions of vertices  $P_i^\infty$  in  $v_i$  that are contained in the LCC after perturbations. To do so, we first list all related notations in Table 1.

TABLE 1  
Notations concerning the mathematical analysis of the robustness of an  $L$ -partite network  $G$ .

Variable	Definition
$p_i$	a fraction $1 - p_i$ of vertices in $v_i$ are randomly removed
$P_i^\infty$	fraction of vertices in $v_i$ that are contained in the LCC
$z_{ij}$	probability that a vertex in $v_i$ is not connected to the LCC via a vertex in $v_j$
$P_i(k)$	degree distribution of vertices in $v_i$
$q_i(k)$	excess degree distribution of $P_i(k)$
$P_{ij}(k)$	degree distribution of vertices in $v_i$ which are connected to vertices in $v_j$
$k$	vertex degree
$\langle k_i \rangle$	$\langle k_i \rangle = \sum_{k=0}^{\infty} k P_i(k)$ , mean degree of vertices in $v_i$
$\langle k_{ij} \rangle$	$\langle k_{ij} \rangle = \sum_{k=0}^{\infty} k P_{ij}(k)$ , mean degree of vertices in $v_i$ which are connected to vertices in $v_j$

For an  $L$ -partite network, the vertex perturbation occurred to one partite may affect vertices in other partite sets and eventually cascading failures are likely to occur. In the following, we present our establish dynamic models and theoretical methods for analyzing the robustness of bidirected and unidirectional  $L$ -partite networks subject to random vertex losses.

To enhance the elegance of the theoretical calculations, we make use of the mathematical tool of generating functions [21]. Given a probability distribution function  $P_{ij}(k)$ , its generating function can be defined as

$$G_{ij}(x) = \sum_{k=0}^{\infty} x^k P_{ij}(k). \quad (2)$$

Based on Eq. 2 we define another function which reads

$$H_{ij}(x) = \frac{G'_{ij}(x)}{G'_{ij}(1)} = \frac{\sum_{k=0}^{\infty} kx^{k-1} P_{ij}(k)}{\sum_{k=0}^{\infty} kx^{k-1} P_{ij}(k)|_{x=1}} = \frac{\sum_{k=0}^{\infty} (k+1)x^k P_{ij}(k+1)}{\langle k_{ij} \rangle} \quad (3)$$

### 3.1 Dynamic Model for Unidirectional Multipartite Networks

Models exhibited in Fig. 3 are for single networks, while networks in reality are often interdependent and could be vulnerable to perturbations. Fig. 4 takes a toy interdependent network as an example to introduce the widely adopted network model for analyzing the robustness of interdependent networks.

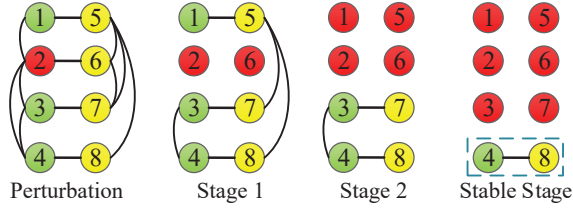


Fig. 4. A schematic illustration of the widely used network model for depicting the dynamics of an interdependent network subject to vertex perturbations.

In Fig. 4, the toy interdependent network consists of two networks (distinguished by different colors) which are one-to-one connected. Originally, vertex 2 from one network is removed. In stage 1, the edges attached to vertex 2 are removed. The same process occurs to vertex 6 since vertices 2 and 6 are interdependent. The removal of vertex 2 fragments the network in the left side into two parts and it is assumed that only the vertices in the LCC will be of interest. As a result, in stage 2 vertices 1 and 5 are removed. The removal of vertex 5 leads to the fragmentation of the network in the right side. This process is continued until no further vertex remove is possible. In the final stage, only vertices 4 and 8 are remained functioning.

As aforementioned, there is no connection between vertices in the same partite set of a multipartite network. Thus the model demonstrated in Fig. 4 is not applicable to multipartite networks.

We establish the dynamic model shown in Fig. 5 to depict the dynamic process of a unidirectional multipartite network subject to vertex perturbations.

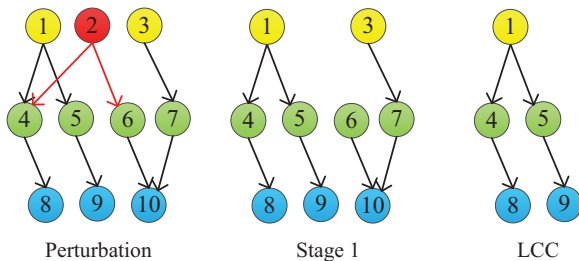


Fig. 5. The network model for depicting the robustness of a unidirectional multipartite network subject to vertex loss.

As shown in Fig. 5, vertex 2 is removed from a unidirectional tripartite network. This removal breaks the network into two parts. Because the network is unidirectional, no cascading failures will occur. Therefore, in the final stage we only consider the vertices in the LCC which is shown in the right panel of Fig. 5.

Note that, since we are considering a unidirectional multipartite network, we should focus on the LCC of the whole network when analyzing its robustness. For example, let us consider the network shown in Fig. 5 as a complex control system where vertices 1, 2, and 3 represent the controllers. In order to ensure the success of the control mission, it is generally required that there exists a LCC which contains as many components as possible.

### 3.2 Dynamic Model for Bi-directed Multipartite Networks

We establish the dynamic model shown in Fig. 6 to depict the dynamic process of a bi-directed multipartite network subject to vertex perturbations.

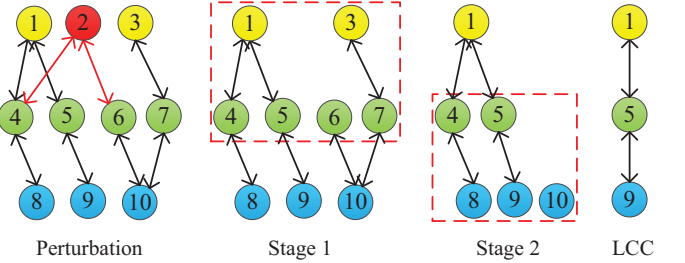


Fig. 6. The network model for depicting the robustness of a bi-directed multipartite network subject to vertex loss.

As shown in Fig. 6, vertex 2 is removed from a bi-directed tripartite network. The removal also fragments the network into two parts. However, we cannot simply compute the robustness of this network in the same manner as shown in Fig. 5 since the focal network is bi-directed and its dynamics are different from that of a unidirectional multipartite network.

After removing vertex 2, the bipartite network (surrounded by red circle in stage 1) breaks into three parts and only vertices 1, 4, and 5 will still be functional. As a consequence, in stage 1 vertices 3, 6, and 7 will be disfunctional. The removal of vertices 3, 6, and 7 further divides the second bipartite network (surrounded by red circle in stage 2) into three groups and only vertices in the LCC of the second bipartite network (because there are two LCCs, we randomly choose the one that contains vertices 5 and 9) will survive. In the final stage, only vertices 1, 5, and 9 are remained in the final LCC.

### 3.3 Theoretical Method for Bipartite Networks

Existing theoretical methods for analyzing the robustness of interdependent networks are largely based on the network model displayed in Fig. 4. As discussed above, those theories are not amenable to multipartite networks and therefore new theories are desired.



Because the simplest form of a multipartite network is the bipartite network, here we first investigate the robustness of bipartite networks to random vertex perturbations.

In order to figure out  $P_1^\infty$  and  $P_2^\infty$  for a bipartite network, we first define  $z_{12}$  as the probability that a vertex  $i \in v_1$  is not connected to the LCC via a vertex  $j \in v_2$ . Analogously, we can define another probability  $z_{21}$ . Suppose the degree of vertex  $i$  is  $k$ , i.e., it has  $k$  neighbors in  $v_2$ . As a consequence, the average probability  $P_r(i \notin \text{LCC})$  that vertex  $i$  does not belong to the LCC is

$$P_r(i \notin \text{LCC}) = \sum_{k=0}^{\infty} P_{12}(k) z_{12}^k = G_{12}(z_{12}). \quad (4)$$

Therefore, the average probability  $P_r(i \in \text{LCC})$  that vertex  $i$  belongs to the LCC is  $P_r(i \in \text{LCC}) = 1 - P_r(i \notin \text{LCC})$ . Since we randomly remove a fraction  $1 - p_1$  and a fraction  $1 - p_2$  of vertices from the bipartite network, the proportion of vertices in  $v_1$  that belong to the LCC, i.e.,  $P_1^\infty$ , thus can be written as

$$P_1^\infty = p_1 P_r(i \in \text{LCC}) = p_1 (1 - G_{12}(z_{12})). \quad (5)$$

Analogously, we can derive the expression of  $P_2^\infty$  which reads

$$P_2^\infty = p_2 (1 - G_{21}(z_{21})). \quad (6)$$

The key to Eqs. 5 and 6 is to figure out the expressions of variables  $z_{12}$  and  $z_{21}$ . Note that the event that a vertex  $i \in v_1$  is not connected to the LCC via a vertex  $j \in v_2$  happens under two independent cases: 1) the vertex in  $v_2$  is removed, which happens with a probability  $1 - p_2$ ; 2) the vertex in  $v_2$  is not removed (this happens with a probability  $p_2$ ) and it is not connected to the LCC via its  $k$  extra neighbors in  $v_1$  (this happens with a probability  $z_{21}^k$ ).

Note that the probability  $q_{21}(k)$  for vertex  $j$  to have  $k$  extra neighbors in  $v_1$  is not  $P_{21}(k+1)$ . Because for a bipartite network there are totally  $n_2 P_{21}(k+1)$  vertices in  $v_2$  each of which has degree  $k+1$ , thus  $q_{21}(k)$  can be calculated as follows

$$q_{21}(k) = \frac{n_2 P_{21}(k+1)(k+1)}{n_2 \langle k_2 \rangle} = \frac{(k+1) P_{21}(k+1)}{\langle k_2 \rangle}. \quad (7)$$

Based on Eq. 7, we can get the expression of  $z_{12}$  as

$$\begin{aligned} z_{12} &= \sum_{k=0}^{\infty} (1 - p_2 + p_2 z_{21}^k) q_{21}(k) \\ &= 1 - p_2 + p_2 H_{21}(z_{21}) \end{aligned} \quad (8)$$

Analogously, we can get the expression of  $z_{21}$  as

$$z_{21} = 1 - p_1 + p_1 H_{12}(z_{12}). \quad (9)$$

By substituting the expression of  $z_{21}$  into that of  $z_{12}$  we can get the following self-consistent equation

$$z_{12} = 1 - p_2 + p_2 H_{21}(1 - p_1 + p_1 H_{12}(z_{12})). \quad (10)$$

A possible non-trivial solution  $z_{12}$  may appear if the two curves  $f_1 = z_{12}$  and  $f_2 = 1 - p_2 +$

$p_2 H_{21}(1 - p_1 + p_1 H_{12}(z_{12}))$  meet with each other tangentially at  $z_{12} = 1$ . Putting it another way, a critical value  $p_c$  occurs when

$$\left. \frac{df_1}{dz_{12}} \right|_{z_{12}=1} = \left. \frac{df_2}{dz_{12}} \right|_{z_{12}=1}. \quad (11)$$

From Eq. 11 we further derive the following relations

$$p_c = p_1 p_2 = \frac{\langle k_1 \rangle \langle k_2 \rangle}{(\langle k_1^2 \rangle - \langle k_1 \rangle^2)(\langle k_2^2 \rangle - \langle k_2 \rangle^2)}. \quad (12)$$

### 3.4 Theoretical Method for Unidirectional Multipartite Networks

For generality and simplicity, hereafter we take a tripartite network as an example to delineate our mathematical derivations for the robustness analysis of multipartite networks. Analogous to the analysis of bipartite networks, we therefore have four variables  $z_{12}$ ,  $z_{21}$ ,  $z_{23}$ , and  $z_{32}$ .

We can notice from Fig. 5 that the LCC may not encompass vertices in  $v_1$  or  $v_3$ . Consequently, it is easy to get the expressions of  $P_1^\infty$  and  $P_3^\infty$  as

$$\begin{cases} P_1^\infty = p_1 \left( 1 - \sum_{k=0}^{\infty} P_{12}(k) z_{12}^k \right) = p_1 (1 - G_{12}(z_{12})) \\ P_3^\infty = p_3 \left( 1 - \sum_{k=0}^{\infty} P_{32}(k) z_{32}^k \right) = p_3 (1 - G_{32}(z_{32})) \end{cases} \quad (13)$$

For a vertex  $j \in v_2$ , if it does not belong to the LCC, then it should not be connected to the LCC via its neighbors in  $v_1$  and  $v_3$ . Thus, the average probability  $P_r(j \notin \text{LCC})$  that vertex  $j$  does not belong to the LCC is

$$\begin{aligned} P_r(j \notin \text{LCC}) &= \sum_{k=0}^{\infty} P_{21}(k) z_{21}^k \cdot \sum_{k=0}^{\infty} P_{23}(k) z_{23}^k \\ &= G_{21}(z_{21}) G_{23}(z_{23}) \end{aligned} \quad (14)$$

As a consequence, we can get the expression of  $P_2^\infty$  as

$$\begin{aligned} P_2^\infty &= p_2 (1 - P_r(j \notin \text{LCC})) \\ &= p_2 (1 - G_{21}(z_{21}) G_{23}(z_{23})) \end{aligned} \quad (15)$$

In the next step we are going to derive the relations between the four variables  $z_{12}$ ,  $z_{21}$ ,  $z_{23}$ , and  $z_{32}$ .

Since the event that a vertex in  $v_2$  is not connected to the LCC via a vertex in  $v_1$  happens under two cases which are very similar to that of a bipartite network, we therefore can formulate  $z_{21}$  and  $z_{23}$  as

$$\begin{cases} z_{21} = 1 - p_1 + p_1 H_{12}(z_{12}) \\ z_{23} = 1 - p_3 + p_3 H_{32}(z_{32}) \end{cases} \quad (16)$$

Now let us consider the event that a vertex  $i \in v_1$  is not connected to the LCC via a vertex  $j \in v_2$ , i.e., the probability  $z_{12}$ . This event happens under two cases: 1)  $j$  is removed; 2)  $j$  is not removed. Case 1 happens with a probability  $1 - p_2$ . Now the key is to work out the probability for case 2.

Because a vertex  $j \in v_2$  could have neighbors in  $v_1$  and  $v_3$ , if  $j$  does not belong to the LCC, then it must not be connected to the LCC via its neighbors in  $v_1$  and  $v_3$ . Note that the event for  $j$  to have  $k$  neighbors in  $v_1$  and the event

for  $j$  to have  $k$  neighbors in  $v_3$  are independent. Therefore, the probability  $P_r(j \notin \text{LCC})$  in this situation becomes

$$P_r(j \notin \text{LCC}) = \sum_{k=0}^{\infty} q_{21}(k) z_{21}^k \cdot \sum_{k=0}^{\infty} q_{23}(k) z_{23}^k = H_{21}(z_{21}) H_{23}(z_{23}). \quad (17)$$

As a result, the probability for case 2 to happen is  $p_2 H_{21}(z_{21}) H_{23}(z_{23})$ . With all these, the expression for  $z_{12}$  becomes

$$z_{12} = 1 - p_2 + p_2 H_{21}(z_{21}) H_{23}(z_{23}). \quad (18)$$

Note that the event that a vertex  $l \in v_3$  is not connected to the LCC via a vertex  $j \in v_2$  happens in the same way as that of vertex  $i$  does, we can easily derive the expression of  $z_{32}$ , which has the same form as Eq. 18.

### 3.5 Theoretical Method for Bi-directed Multipartite Networks

By comparing Figs. 5 and 6 we can notice that the main difference between the two network models lies in the calculation of the LCC. For a bi-directed tripartite network, the derivation process of  $P_1^\infty$  and  $P_3^\infty$  is the same as that of a unidirectional tripartite network. It is easy to prove that  $P_1^\infty$  and  $P_3^\infty$  have the same form as Eq. 13, while  $z_{21}$  and  $z_{23}$  have the same form as Eq. 16. The most difficult part lies in the calculations of  $P_2^\infty$ ,  $z_{12}$ , and  $z_{32}$ .

In order to figure out  $P_2^\infty$ ,  $z_{12}$ , and  $z_{32}$ , we first figure out the probability that a vertex  $j$  in  $v_2$  belongs to the LCC. As mentioned above, vertex  $j$  may have neighbors in  $v_1$  and  $v_3$ . If  $j \in \text{LCC}$ , then at least one neighbor vertex  $i$  in  $v_1$  should connect  $j$  to the LCC, and this event happens with a probability  $1 - z_{21}^k$ . Because the LCC contains vertices from  $v_1$ ,  $v_2$ , and  $v_3$ , thus there should exist at least one neighbor vertex  $l$  in  $v_3$  which also connects  $j$  to the LCC, and this event happens with a probability  $1 - z_{23}^k$ . As a consequence, the probability that a vertex  $j$  in  $v_2$  belongs to the LCC is  $(1 - z_{21}^k)(1 - z_{23}^k)$ .

Bear in mind that probabilities for  $j$  to have  $k$  neighbors in  $v_1$  and  $k$  neighbor in  $v_3$  are respectively  $P_{21}(k)$  and  $P_{23}(k)$ . Note that the probabilities are not  $q_{21}(k)$  and  $q_{23}(k)$  since vertex  $j$  is randomly picked but not arrived from vertex  $i$ . Therefore, we can formulate  $P_2^\infty$  as

$$\begin{aligned} P_2^\infty &= p_2 \sum_{k=0}^{\infty} P_{21}(k) (1 - z_{21}^k) \sum_{k=0}^{\infty} P_{23}(k) (1 - z_{23}^k) \\ &= p_2 \left( 1 - \sum_{k=0}^{\infty} P_{21}(k) z_{21}^k \right) \left( 1 - \sum_{k=0}^{\infty} P_{23}(k) z_{23}^k \right) \\ &= p_2 (1 - G_{21}(z_{21})) (1 - G_{23}(z_{23})) \end{aligned} \quad (19)$$

Analogously, we can figure out  $z_{12}$  and  $z_{32}$  as

$$\begin{cases} z_{12} &= 1 - p_2 \left( 1 - \sum_{k=0}^{\infty} q_{21}(k) z_{21}^k \right) \left( 1 - \sum_{k=0}^{\infty} q_{23}(k) z_{23}^k \right) \\ &= 1 - p_2 (1 - H_{21}(z_{21})) (1 - H_{23}(z_{23})) \\ z_{32} &= z_{12} \end{cases} \quad (20)$$

In summary, the robustness analysis for a bi-directed tripartite network with arbitrary degree distributions can be written as

$$\begin{cases} z_{12} &= z_{32} = 1 - p_2 (1 - H_{21}(z_{21})) (1 - H_{23}(z_{23})) \\ z_{21} &= 1 - p_1 + p_1 H_{12}(z_{12}) \\ z_{23} &= 1 - p_3 + p_3 H_{32}(z_{32}) \end{cases}, \quad (21)$$

$$\begin{cases} P_1^\infty &= p_1 (1 - G_{12}(z_{12})) \\ P_2^\infty &= p_2 (1 - G_{21}(z_{21})) (1 - G_{23}(z_{23})) \\ P_3^\infty &= p_3 (1 - G_{32}(z_{32})) \end{cases}. \quad (22)$$

## 4 RESULTS

### 4.1 Random Networks

In order to validate the correctness of our proposed method for analyzing the robustness of directed multipartite networks subject to random vertex perturbations, here we generate random multipartite networks.

Let us define a probability vector  $\mathbf{R} = (r_1, r_2, \dots, r_{L-1})$ . Given an empty  $L$ -partite network, we connect two arbitrary vertices with one comes from  $v_i$  and the other one comes from  $v_{i+1}$  with a probability  $r_i$ . Then the degree distribution  $P_{ij}(k)$  becomes

$$\begin{aligned} P_{ij}(k) &= \binom{n_j}{k} r_j^k (1 - r_j)^{n_j - k} \\ &\approx e^{-\langle k_{ij} \rangle} \frac{\langle k_{ij} \rangle^k}{k!}, \end{aligned} \quad (23)$$

where  $j = i - 1$  and  $\langle k_{ij} \rangle = n_j r_j$ . Analogously, we can get the degree distribution  $P_{il}(k)$  which read

$$\begin{aligned} P_{il}(k) &= \binom{n_l}{k} r_l^k (1 - r_l)^{n_l - k} \\ &\approx e^{-\langle k_{il} \rangle} \frac{\langle k_{il} \rangle^k}{k!}, \end{aligned} \quad (24)$$

where  $l = i + 1$  and  $\langle k_{il} \rangle = n_l r_l$ .

Eqs. 23 and 24 indicate that the generated network follows the Poisson degree distributions. The main reason for only generating networks with Poisson distributions is that a Poisson distribution  $P_{ij}(k)$  has good mathematical properties. To be specific, for a Poisson distribution  $P_{ij}(k) = e^{-\lambda} \frac{\lambda^k}{k!}$ , we have the following equations:

$$G_{ij}(x) = \sum_{k=0}^{\infty} x^k P_{ij}(k) = \sum_{k=0}^{\infty} x^k e^{-\lambda} \frac{\lambda^k}{k!} = e^{\lambda(x-1)}, \quad (25)$$

$$H_{ij}(x) = \frac{G'_{ij}(x)}{G'_{ij}(1)} = e^{\lambda(x-1)} = G_{ij}(x), \quad (26)$$

$$\begin{aligned} \langle k_{ij}^2 \rangle &= \sum_{k=0}^{\infty} k^2 P_{ij}(k) = \sum_{k=0}^{\infty} k^2 e^{-\lambda} \frac{\lambda^k}{k!} \\ &= \lambda e^{-\lambda} \left( \sum_{k=0}^{\infty} (k-1) \frac{\lambda^{k-1}}{(k-1)!} + \sum_{k=0}^{\infty} \frac{\lambda^{k-1}}{(k-1)!} \right) \\ &= \lambda e^{-\lambda} (\lambda e^{\lambda} + e^{\lambda}) \\ &= \langle k_{ij} \rangle^2 + \langle k_{ij} \rangle \end{aligned} \quad (27)$$

where  $\langle k_{ij}^2 \rangle$  is the second moment of  $P_{ij}(k)$ .

Note that the generated multipartite networks for testing purpose are undirected. This is because that the directions of the edges have already been taken into account when developing the network models and mathematical methods for the robustness analysis. While calculating the robustness of a directed multipartite network, we only need to know the degree distributions for vertices in each partite set.

## 4.2 Robustness of bipartite networks

For a bipartite network with Poisson degree distributions, we can simplify the robustness analysis functions into the following forms

$$\begin{cases} z_{12} = 1 - p_2 + p_2 e^{\langle k_2 \rangle (z_{21} - 1)} \\ z_{21} = 1 - p_1 + p_1 e^{\langle k_1 \rangle (z_{12} - 1)} \end{cases}, \quad (28)$$

$$\begin{cases} P_1^\infty = p_1 (1 - e^{\langle k_1 \rangle (z_{12} - 1)}) \\ P_2^\infty = p_2 (1 - e^{\langle k_2 \rangle (z_{21} - 1)}) \end{cases}. \quad (29)$$

We can notice that as long as parameters  $n_1$ ,  $n_2$ , and  $r$  are given, then we can directly solve Eqs. 28 and 29 so as to investigate the robustness of the focal bipartite network.

Now let us consider bipartite networks with the following configurations:  $N = n_1 + n_2 = 8 \times 10^4$ ,  $n_1 = \alpha N$ , and  $r = C/N$ , where  $\alpha \in (0, 1)$  and  $C$  is a constant. We test the robustness of two sets of bipartite networks:

- 1)  $\alpha = \frac{1}{8}$  and  $C = \{4, 5, 6, 7\}$ ;
- 2)  $\alpha = \frac{3}{8}$  and  $C = \{4, 5, 6, 7\}$ .

For simplicity we set  $p_1 = p_2 = p$ , i.e., for each network we randomly remove a fraction  $1 - p$  of vertices from the whole network.

Fig. 7 displays the robustness results on bipartite networks with Poisson degree distributions. The size for each network is controlled by  $\alpha$  and the degree is controlled by  $C$ . The simulation results are averaged over 1000 independent trials while numerical results are obtained by solving Eqs. 28 and 29 by substituting the corresponding values of  $\langle k_1 \rangle$  and  $\langle k_2 \rangle$ . We can clearly see from Fig. 7 that our theoretical results coincide quite well with the simulation results.

For a bipartite network with Poisson degree distributions, the critical value as given in Eq. 12 now becomes

$$\begin{aligned} p_c^2 &= \frac{\langle k_1 \rangle \langle k_2 \rangle}{(\langle k_1^2 \rangle - \langle k_1 \rangle^2)(\langle k_2^2 \rangle - \langle k_2 \rangle^2)} \\ &= \frac{\langle k_1 \rangle \langle k_2 \rangle}{\alpha(1 - \alpha)C^2} \\ \Rightarrow p_c &= \frac{1}{C\sqrt{\alpha(1 - \alpha)}} \end{aligned} \quad (30)$$

Based on Eq. 30 we have

- 1) when  $\alpha = \frac{1}{8}$  and  $C = \{4, 5, 6, 7\}$ , the critical values are  $p_c = \{0.7559, 0.6047, 0.5040, 0.4320\}$ ;
- 2) when  $\alpha = \frac{3}{8}$  and  $C = \{4, 5, 6, 7\}$ , the critical values are  $p_c = \{0.5773, 0.4619, 0.3850, 0.3299\}$ ,

which are in accordance with the simulation results shown in Fig. 7.

Fig. 8 exhibits the simulations on the robustness of two bipartite networks in a more general way. We respectively remove a fraction  $1 - p_1$  and a fraction  $1 - p_2$  of vertices

from  $v_1$  and  $v_2$  of each bipartite network.  $P_1^\infty$  and  $P_2^\infty$  are respectively shown with respect to different settings of  $p_1$  and  $p_2$ . We can clearly see from Fig. 8 that the robustness curves are smooth and the turning points of  $p_1$  and  $p_2$  are small in values, which indicates that bipartite networks are extremely robustness to random vertex perturbations.

## 4.3 Robustness of Unidirectional multipartite networks

For a unidirectional tripartite network with Poisson degree distributions, we can simplify Eqs. 16, 18, 13 and 15 into the following forms

$$\begin{cases} z_{21} = 1 - p_1 + p_1 e^{\langle k_{12} \rangle (z_{12} - 1)} \\ z_{23} = 1 - p_3 + p_3 e^{\langle k_{32} \rangle (z_{32} - 1)} \\ z_{12} = 1 - p_2 + p_2 e^{\langle k_{23} \rangle (z_{23} - 1)} e^{\langle k_{21} \rangle (z_{21} - 1)} \\ z_{32} = z_{12} \end{cases}, \quad (31)$$

$$\begin{cases} P_1^\infty = p_1 (1 - e^{\langle k_{12} \rangle (z_{12} - 1)}) \\ P_2^\infty = p_2 (1 - e^{\langle k_{23} \rangle (z_{23} - 1)} e^{\langle k_{21} \rangle (z_{21} - 1)}) \\ P_3^\infty = p_3 (1 - e^{\langle k_{32} \rangle (z_{32} - 1)}) \end{cases}. \quad (32)$$

For simplicity we set  $p_1 = p_2 = p_3 = p$ . In the experiments we generate three tripartite networks with their parameter configurations respectively given as

- 1)  $(n_1, n_2, n_3) = (3, 1, 2) \times 10^4$ ,  $\mathbf{R} = (\frac{C_1}{n_1 + n_2}, \frac{C_2}{n_3 + n_2})$ ,  $C_1 = 8, C_2 = 6$ ;
- 2)  $(n_1, n_2, n_3) = (1, 4, 1) \times 10^4$ ,  $\mathbf{R} = (\frac{C_1}{n_1 + n_2}, \frac{C_2}{n_3 + n_2})$ ,  $C_1 = 12, C_2 = 14$ ;
- 3)  $(n_1, n_2, n_3) = (2, 2, 2) \times 10^4$ ,  $\mathbf{R} = (\frac{C_1}{n_1 + n_2}, \frac{C_2}{n_3 + n_2})$ ,  $C_1 = 4, C_2 = 3$ .

Fig. 9 displays the robustness of three tripartite networks with each network subject to random vertices remove from every partite set. The theoretical results  $P^\infty$  as shown in the last subfigure of Fig. 9 are simply obtained by solving  $P^\infty = \sum P_i^\infty$ . Fig. 9 clearly validates the correctness of our proposed method.

By combining Eqs. 11 and 31 we can get

$$\begin{aligned} 1 &= p_2 \frac{d}{dz_{12}} (e^{\langle k_{23} \rangle (z_{23} - 1)} e^{\langle k_{21} \rangle (z_{21} - 1)}) \Big|_{z_{12}=1} \\ \Rightarrow 1 &= p_2 (p_3 \langle k_{23} \rangle \langle k_{32} \rangle + p_1 \langle k_{21} \rangle \langle k_{12} \rangle) \end{aligned} \quad (33)$$

Since we set  $p_1 = p_2 = p_3 = p$  in our experiments, from the above equation we can figure out the critical value of  $p_c$  which reads

$$p_c = \sqrt{\frac{1}{\langle k_{23} \rangle \langle k_{32} \rangle + \langle k_{21} \rangle \langle k_{12} \rangle}}. \quad (34)$$

For the three tested tripartite networks, we have

- 1)  $\langle k_{12} \rangle = 2, \langle k_{21} \rangle = 6, \langle k_{23} \rangle = 4, \langle k_{32} \rangle = 2$ ;
- 2)  $\langle k_{12} \rangle = 9.6, \langle k_{21} \rangle = 2.4, \langle k_{23} \rangle = 2.8, \langle k_{32} \rangle = 11.2$ ;
- 3)  $\langle k_{12} \rangle = 2, \langle k_{21} \rangle = 2, \langle k_{23} \rangle = 1.5, \langle k_{32} \rangle = 1.5$ .

By substituting the mean degrees into Eq. 34 we respectively get the critical values for the three networks which are  $p_c = 0.224, p_c = 0.136$ , and  $p_c = 0.400$ . We can observe from Fig. 9 that our theoretical results are the same as the simulations.



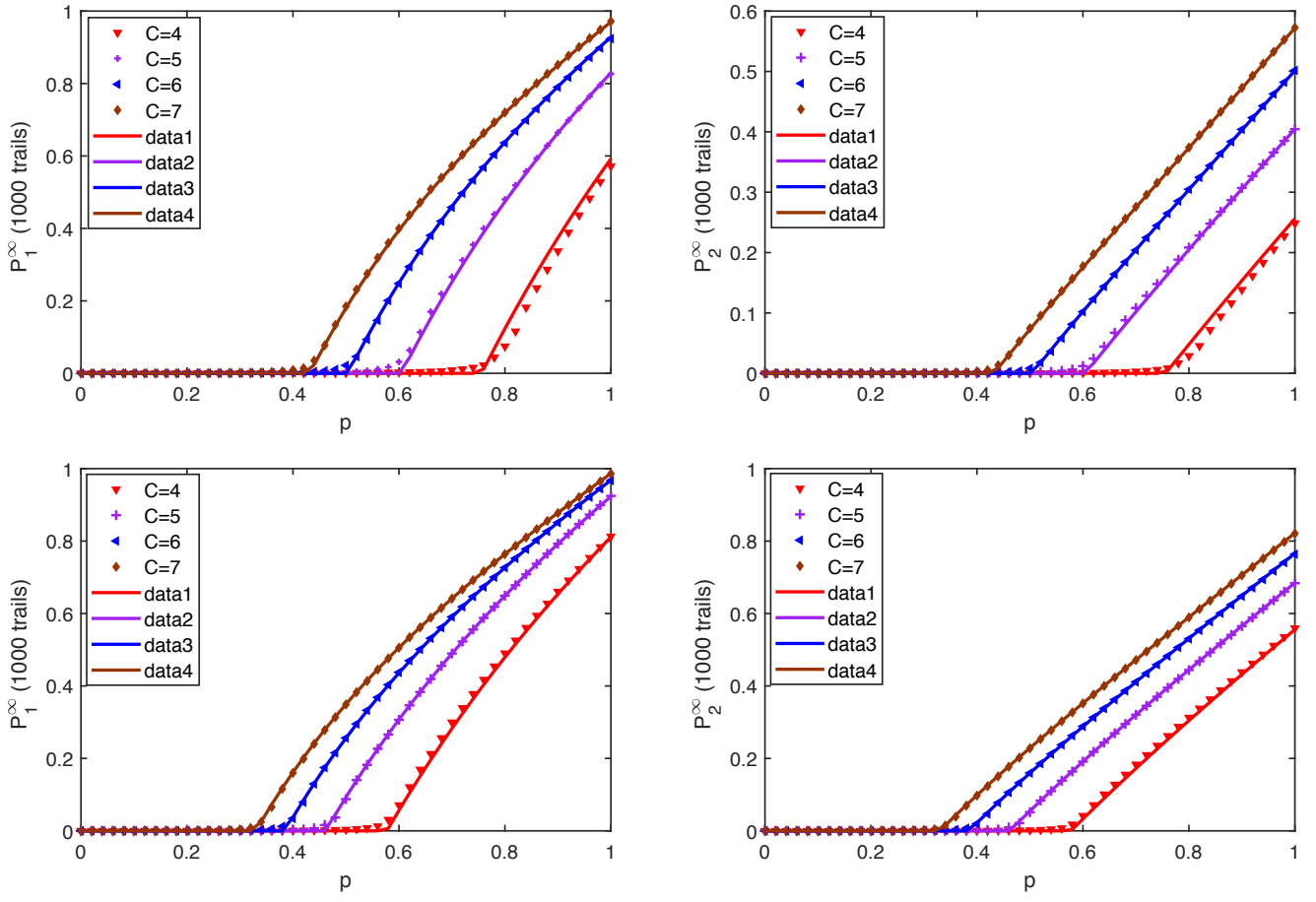


Fig. 7. Theoretical (denoted by lines) and simulation (denoted by symbols) results on the robustness of bipartite networks with Poisson degree distributions. A fraction  $1 - p$  of vertices are randomly removed from each network. We set  $\alpha = 1/8$  and  $\alpha = 2/8$  respectively for networks in the first row and the second row.

#### 4.4 Robustness of Bi-directed multipartite networks

For a bi-directed tripartite network with Poisson degree distributions, we can simplify Eqs. 19 and 20 into the following forms

$$\begin{cases} z_{21} = 1 - p_1 + p_1 e^{\langle k_{12} \rangle (z_{12} - 1)} \\ z_{23} = 1 - p_3 + p_3 e^{\langle k_{32} \rangle (z_{32} - 1)} \\ z_{12} = 1 - (p_2 - p_2 e^{\langle k_{21} \rangle (z_{21} - 1)}) (1 - e^{\langle k_{23} \rangle (z_{23} - 1)}) \\ z_{32} = z_{12} \end{cases} \quad (35)$$

$$\begin{cases} P_1^\infty = p_1 (1 - e^{\langle k_{12} \rangle (z_{12} - 1)}) \\ P_2^\infty = p_2 (1 - e^{\langle k_{21} \rangle (z_{21} - 1)}) (1 - e^{\langle k_{23} \rangle (z_{23} - 1)}) \\ P_3^\infty = p_3 (1 - e^{\langle k_{32} \rangle (z_{32} - 1)}) \end{cases} \quad (36)$$

In our experiments, we set  $p_1 = p_2 = p_3 = p$ . We still carry out experiments and theoretical analysis on the previously generated three tripartite networks. For the third network, we slightly change the constant  $C_2$  to be  $C_2 = 7$ .

Fig. 10 displays the robustness of bi-directed tripartite networks with Poisson degree distributions. Each network is subject to random removal of a fraction  $1 - p$  of vertices from the whole network. Fig. 10 once again proves the correctness of our proposed method for analyzing the robustness of directed multipartite networks.

It can be seen from Fig. 10 that there exist abrupt jumps for the robustness curves, i.e., at the critical point  $p_c$ ,  $P_i^\infty$  suddenly falls from a finite value to zero. This phenomenon indicates that bi-directed multipartite networks are less robust than unidirectional multipartite networks to random vertex perturbations.

Note that the critical value  $p_c$  cannot be obtained by solving Eq. 11. The reason can be discovered from Fig. 11 in which we draw the two curves  $f_1(z_{12}) = z_{12}$  and  $f_2(z_{12})$ . The expression of  $f_2(z_{12})$  is the right term of the expression of  $z_{12}$  in Eq. 36 where  $z_{21}$  and  $z_{23}$  are functions of  $z_{12}$ .

It can be seen from Fig. 11 that there always exists a trivial solution of  $z_{12} = 1$  to the function  $f_2(z_{12}) = z_{12}$ . When  $p$  increases from 0 to 1, there exist a critical value of  $p$  which makes function  $f_2(z_{12}) = z_{12}$  have the first non-trivial solution. Note that  $f_2(z_{12}) = z_{12}$  is a transcendental function, we use numerical analysis to figure out critical value of  $p$ .

## 5 CONCLUDING REMARKS

To explore the robustness of complex networks to perturbations is of pivotal significance toward system control. While existing studies are primarily developed for monopartite networks, in this paper we proposed a theoretical method to investigate the robustness of multipartite networks to

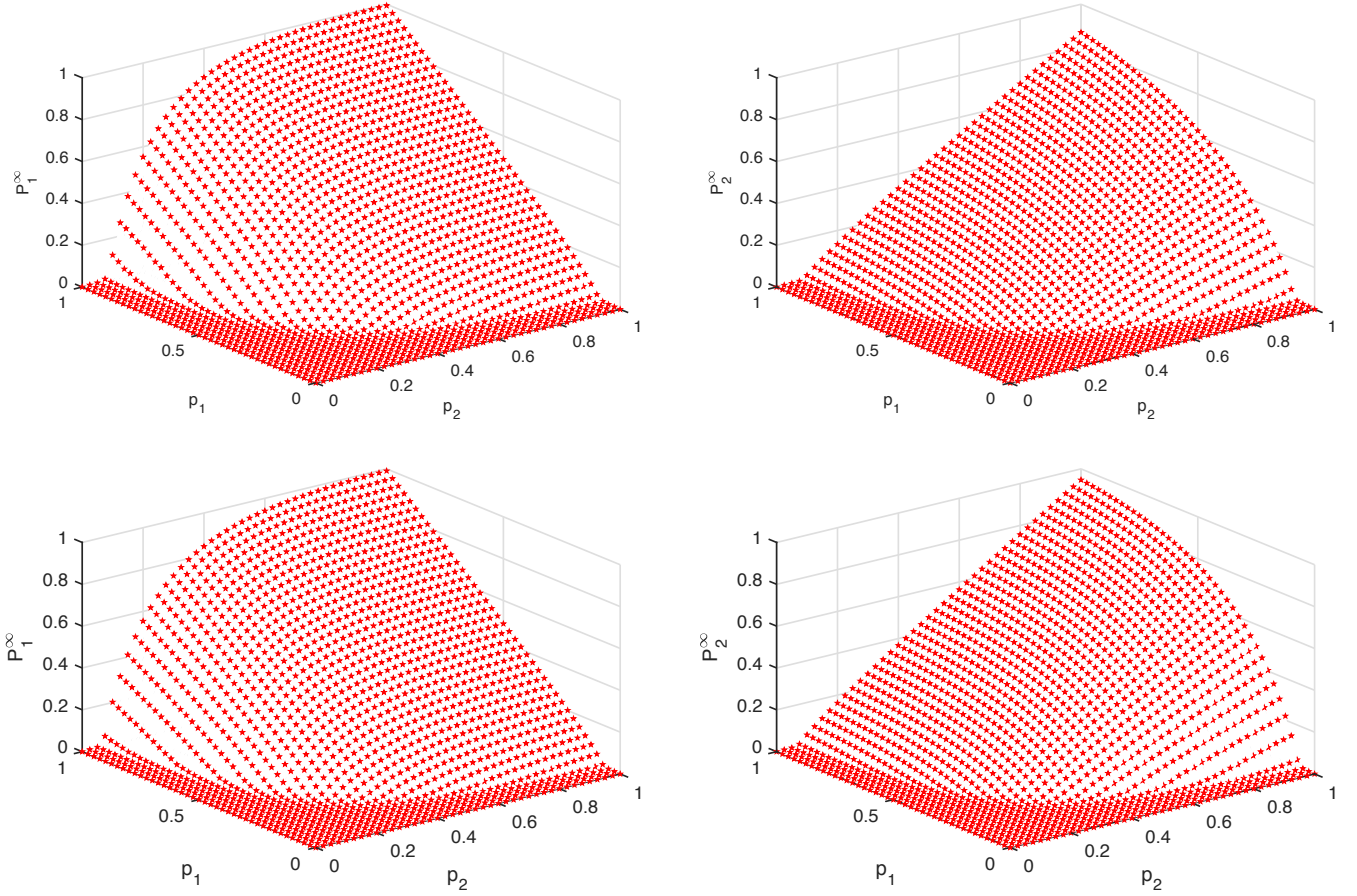


Fig. 8. Simulation results on the robustness of two bipartite network with Poisson degree distributions. We respectively set  $\alpha = 1/8$ ,  $C = 8$  and  $\alpha = 3/8$ ,  $C = 8$  for the two bipartite networks. A fraction  $1 - p_1$  and a fraction  $1 - p_2$  of vertices are randomly removed from  $v_1$  and  $v_2$ , respectively.

random vertex perturbations. The correctness of our proposed methods have been verified through experiments on multipartite networks with Poisson degree distributions. Although we only mathematically investigated the robustness of bipartite and tripartite networks, it is easy to extend our proposed method to  $L$ -partite networks with  $L > 3$ .

Existing studies on network robustness indicate that a single network could be robust to random perturbations (exhibit second order phase transition) while an interdependent network could be vulnerable to random perturbations (exhibit first order phase transition). To some extent, a multipartite network could be regarded as a special case of an interdependent network. However, in this study we discovered that the robustness of bi-directed multipartite networks show first order phase transition while unidirectional multipartite networks display second order phase transition.

Note that it has long been reported that many real-world networks exhibit fat-tail degree distributions [1, 55], i.e., power law distributions. These networks are generally called scale-free (SF) networks. However, in the experiments multipartite networks with power law distributions are not tested. For one thing, our proposed methods can theoretically analyze the robustness of multipartite networks with arbitrary degree distributions. For another thing, SF net-

works in real-world are rare according to the latest research [56]. Meanwhile, although many efforts have been done to generate single networks with arbitrary degree distributions [2, 57–66], generating multipartite networks with power law degree distributions is non-trivial and still demands tremendous efforts. Below are some useful discussions for generating multipartite SF networks.

For generality, let us consider generating a bipartite SF network  $\mathbf{BP}_{m \times n}$  with  $m$  and  $n$  respectively denote the number of vertices in  $v_1$  and  $v_2$ . The degree distributions of the vertices in  $v_1$  and  $v_2$  can be respectively formulated as

$$\begin{aligned} P_1(d) &= C_1 \cdot d^{-\lambda_1} \\ P_2(k) &= C_2 \cdot k^{-\lambda_2} \end{aligned} \quad (37)$$

where  $C_1$  and  $C_2$  are two constants and  $\lambda_1$  and  $\lambda_2$  are the exponents of the power law distributions.

In order to generate a bipartite SF network, one needs to do the inverse transform sampling, i.e., to sample the degree sequences  $\mathbf{D} = (d_1, d_2, \dots, d_m)$  and  $\mathbf{K} = (k_1, k_2, \dots, k_n)$  respectively from  $P_1(d)$  and  $P_2(k)$ , where  $d_i \in [d_{\min}, n]$  and  $k_i \in [k_{\min}, m]$ . Once  $\mathbf{D}$  and  $\mathbf{K}$  are obtained, then one can utilize the Configuration Model or its improved variants to graphically realize  $\mathbf{D}$  and  $\mathbf{K}$  and the bipartite network is thus generated.

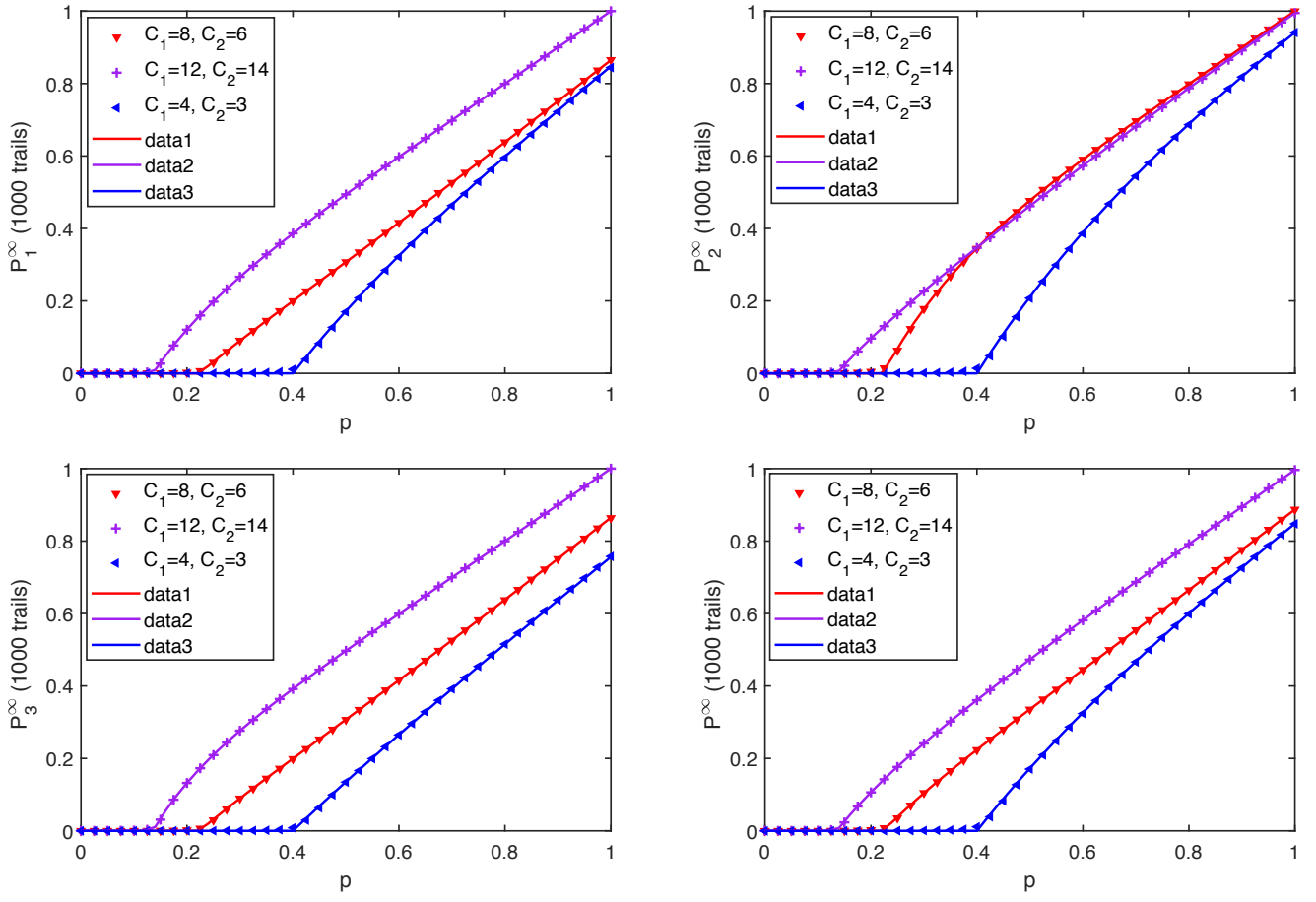


Fig. 9. Theoretical (denoted by lines) and simulation (denoted by symbols) results on the robustness of unidirectional tripartite networks with Poisson degree distributions. A fraction  $1 - p$  of vertices are randomly removed from each network.

Note that for a bipartite network,  $D$  and  $K$  should satisfy the following condition

$$\sum_{i=1}^m d_i \equiv \sum_{i=1}^n k_i \equiv |E|. \quad (38)$$

Eq. 38 has an equivalent form which reads

$$m \cdot \langle k_1 \rangle = n \cdot \langle k_2 \rangle. \quad (39)$$

Recall the definition for the first moment  $\langle k \rangle$ , along with Eq. 39 we can further derive

$$\begin{aligned} m \cdot \sum_{d=d_{min}}^{\infty} d \cdot P_1(d) &= n \cdot \sum_{k=k_{min}}^{\infty} k \cdot P_2(k) \\ \Rightarrow m \cdot C_1 \cdot \int_{d_{min}}^{\infty} d^{-\lambda_1+1} dd &= n \cdot C_2 \cdot \int_{k_{min}}^{\infty} k^{-\lambda_2+1} dk \end{aligned} \quad (40)$$

Assume that  $\lambda_1 - 1 \geq 1$  and  $\lambda_2 - 1 \geq 1$ , the above equation thus can be simplified as

$$\frac{mC_1}{\lambda_1 - 2} d_{min}^{2-\lambda_1} = \frac{nC_2}{\lambda_2 - 2} k_{min}^{2-\lambda_2}. \quad (41)$$

Note that we always have  $\sum P_1(d) = \sum P_2(k) = 1$ . To be specific, in the limit of  $m, n \rightarrow \infty$ , we have

$$\begin{aligned} \sum P_1(d) &= C_1 \sum_{d=d_{min}}^n d^{-\lambda_1} \simeq C_1 \int_{d=d_{min}}^{\infty} d^{-\lambda_1} dd \\ &= \frac{C_1}{\lambda_1 - 1} d_{min}^{-(\lambda_1-1)} \end{aligned} \quad (42)$$

$$\begin{aligned} \sum P_2(k) &= C_2 \sum_{k=k_{min}}^m k^{-\lambda_2} \simeq C_2 \int_{k=k_{min}}^{\infty} k^{-\lambda_2} dk \\ &= \frac{C_2}{\lambda_2 - 1} k_{min}^{-(\lambda_2-1)} \end{aligned} \quad (43)$$

From the above equation we can easily figure out the constants as

$$\begin{aligned} C_1 &= (\lambda_1 - 1) d_{min}^{\lambda_1-1} \\ C_2 &= (\lambda_2 - 1) k_{min}^{\lambda_2-1} \end{aligned} \quad (44)$$

Putting the above equation back into Eq. 41 we get the relationships between the exponents  $\lambda_1$  and  $\lambda_2$  as follows

$$m \cdot \frac{\lambda_1 - 1}{\lambda_1 - 2} \cdot d_{min} = n \cdot \frac{\lambda_2 - 1}{\lambda_2 - 2} \cdot k_{min}. \quad (45)$$

Therefore, when generating a bipartite SF network, the following deterrents should be taken into account:

- 1) Exponents  $\lambda_1$  and  $\lambda_2$  should obey the constraint shown in Eq. 45. Note that it is easy to generating a

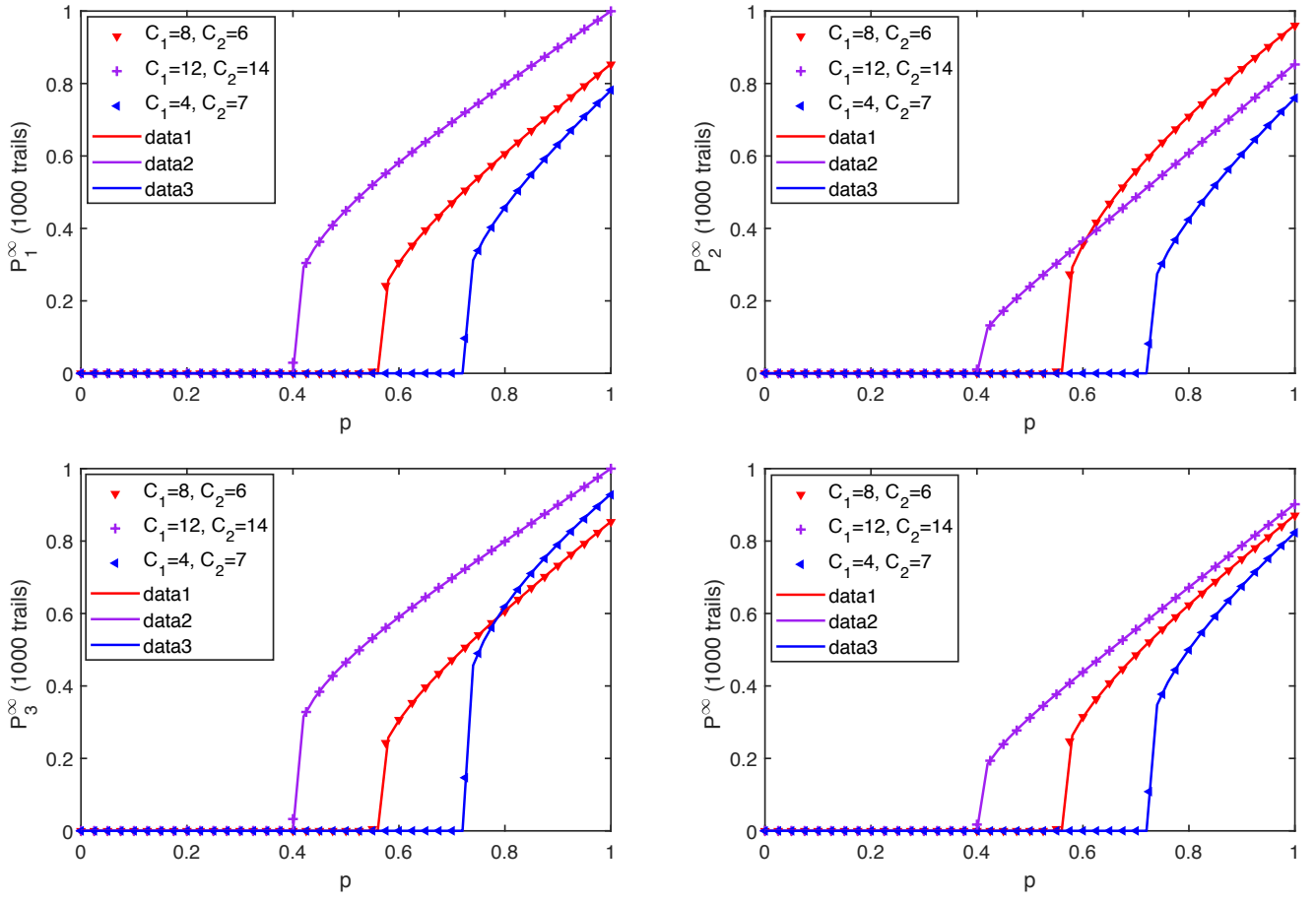


Fig. 10. Theoretical (denoted by lines) and simulation (denoted by symbols) results on the robustness of bi-directed tripartite networks with Poisson degree distributions. A fraction  $1 - p$  of vertices are randomly removed from each network.

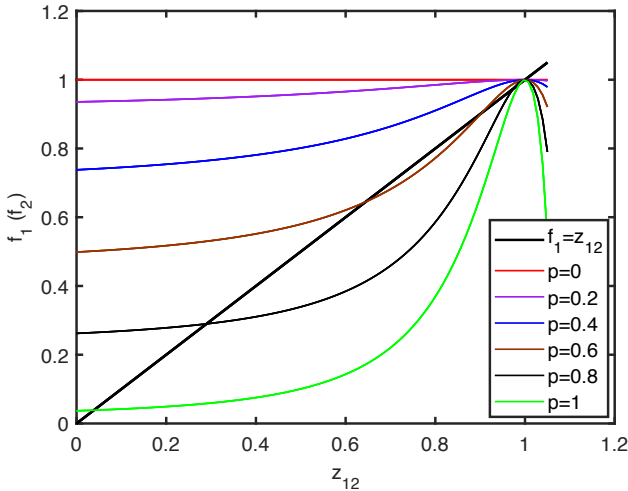


Fig. 11. Function plot of  $f_1(z_{12})$  and  $f_2(z_{12})$ . The parameter settings are  $z_{12} = [0, 1]$ ,  $p = \{0, 0.2, 0.4, 0.6, 0.8, 1\}$ ,  $\langle k_{12} \rangle = 2$ ,  $\langle k_{21} \rangle = 6$ ,  $\langle k_{23} \rangle = 4$ ,  $\langle k_{32} \rangle = 2$ ;

SF-SF interdependent network, since exponents  $\lambda_1$  and  $\lambda_2$  of a SF-SF interdependent network can be set to arbitrary values.

- 2) Degree sequences  $\mathbf{D}$  and  $\mathbf{K}$  should satisfy the

condition shown in Eq. 38. On the one hand, it is hard to sample two degree sequences  $\mathbf{D}$  and  $\mathbf{K}$  which exactly follow the distributions  $P_1(k) \sim \lambda_1^k$  and  $P_2(k) \sim \lambda_2^k$ . On the other hand, even if  $\mathbf{D}$  and  $\mathbf{K}$  strictly follow  $P_1(k)$  and  $P_2(k)$ , there is no guarantee that the condition shown in Eq. 38 will be satisfied. As a consequence, modifications like increasing or decreasing the values of some  $d_i \in \mathbf{D}$  and  $k_j \in \mathbf{K}$  have to be made, while which  $d_i$  and  $k_j$  should be chosen and modified is yet an open question.

- 3) Graphicality condition checking. Even though  $\mathbf{D}$  and  $\mathbf{K}$  simultaneously satisfy Eqs. 45 and 38, one still has to check the graphicality condition as shown in [63, 65] to see whether  $\mathbf{D}$  and  $\mathbf{K}$  can be graphically realized or not.
- 4) Unbiased graph sampling. Note that the number of networks that can graphically realize sequences  $\mathbf{D}$  and  $\mathbf{K}$  could grow exponentially as  $N$  increases [62, 64]. How to efficiently sample a small amount of networks which are less correlated with each other is an open issue.

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