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Strategies for Data Reengineering

Henrard, Jean; Hick, Jean-Marc; Hainaut, Jean-Luc

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Strategies for Data Reengineering

Jean Henrard, Jean-Marc Hick,
Philippe Thiran, Jean-Luc Hainaut

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Laboratory of Database Application Engineering

Laboratoire d'Ingénierie des Bases de Données

www.info.fundp.ac.be/libd



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Plan

- Introduction
- Problem statement
- Reengineering strategies
 - 2 dimensions
 - 3 strategies
- Conclusion



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Introduction

- Legacy system =
 - large and old programs build around legacy DBMS
 - vital to the organization
 - significantly resists modifications and changes
 - expensive to maintain
- Solution : migrate to new platform and technologies
 - expensive and complex process
- Incremental strategy is less risky
 - migrate the DB is one of the steps



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Problem statement

- Data reengineering =
deriving a new database from a legacy database and adapting the software components
 - the functionalities of the system do not change
- Three main steps:
 - schema conversion
 - data conversion
 - program modification



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Problem statement

- Schema conversion
 - translation of the legacy schema into equivalent schema in the new technology
 - DBRE + database design
- Data conversion
 - migration of the data instances from the legacy system to the new one
 - depends on the schema conversion



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Problem statement

■ Program modification

- modification of the programs so that they access the new DB instead of the legacy one
- functionalities, programming language, user interface unchanged
- complex process that relies on the schema conversion



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Reengineering strategies

- 2 dimensions
 - *database* dimension (schema migration)
 - *program* dimension (program modification)
- Data conversion is directly dependent on the database dimension

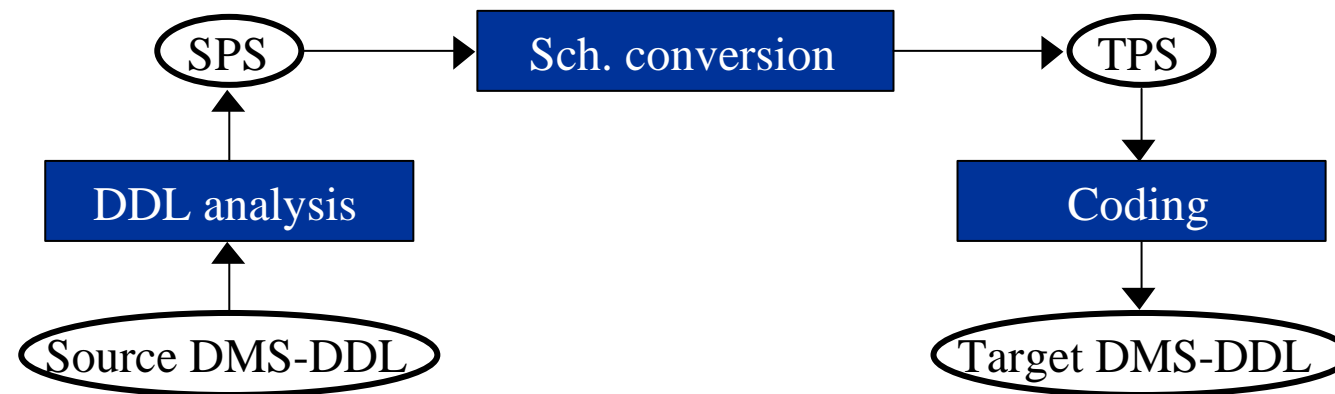


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Database migration strategies

Physical conversion (D1):

- translate to the closer construct into the target DMS (e.g. 1 file \Rightarrow 1 table)
- no semantic interpretation
- cheap but poor quality DB



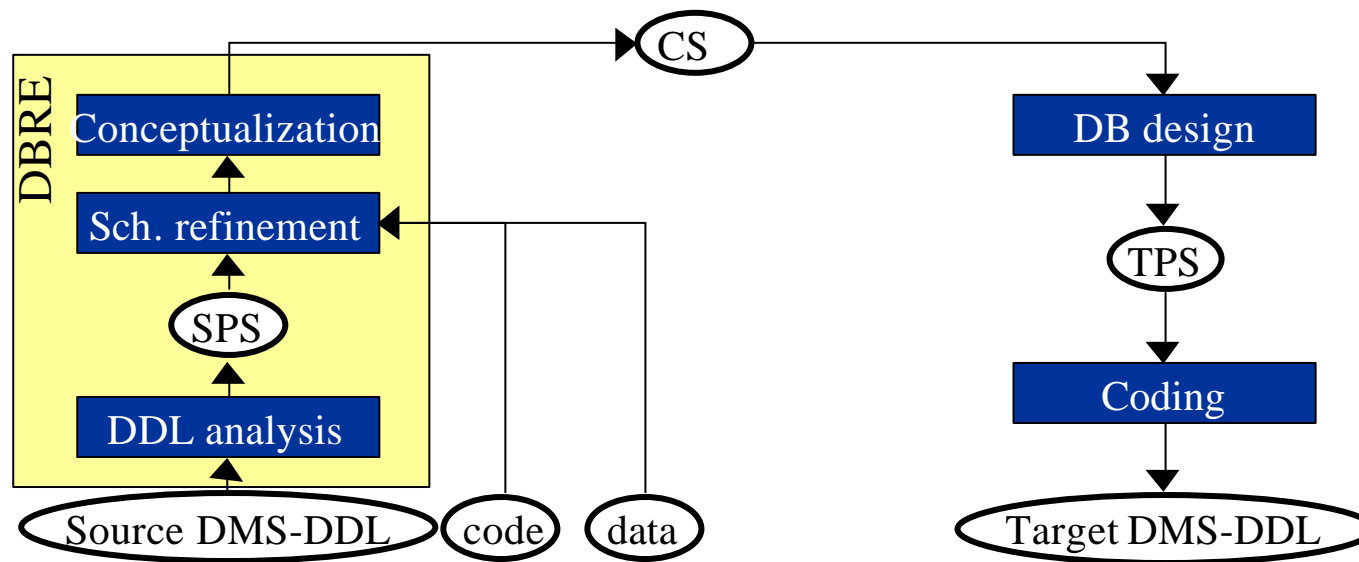


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Database migration strategies

■ *Conceptual conversion (D2)*

- recovering the semantic (conceptual sch) - DBRE
- developing the new DB from the conceptual sch
- good quality and documented DB but expensive





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Database migration strategies

- Schema conversion = schema transformation
- History = chain of transformations
- **Mapping** between the source (SPS) and target (TPS) physical schemas
 - = SPS-to-TPS for physical migration
 - = SPS-to-CS-to-TPS for conceptual migration



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Program modification strategies

■ Wrappers (P1)

- wrappers encapsulate the new database
 - data wrapper =
 - data model conversion
 - semantic conversion
 - functionality simulation
 - “inverse” wrapper: simulate the legacy data interface on the new DB
 - ex: uses COBOL read, write for accessing SQL data
 - SPS -- TPS mapping \Rightarrow automated generation of wrapper
- programs use legacy data access logic
- program logic not changed
- local changes: 1 instruction \Rightarrow x instructions



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Program modification strategies

■ Statement rewriting (P2)

- legacy DMS-DML \Rightarrow target DMS-DML
ex: replace COBOL file access statement by SQL statement
- rewriting the access statements (new DMS-DML)
 - each legacy DML statement must be located and replaced by equivalent statements in the new DML
 - SPT--PTS mapping \Rightarrow automatic program modification
- program logic not changed
- local changes: 1 instruction \Rightarrow x instructions



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Program modification strategies

■ Logic rewriting (P3)

- program rewritten to use the new DMS-DML power
 - explicitly accesses new data
 - takes advantage of the new DML
- logic of the program is changed
 - requires a deep understanding of the program
- global change: x instructions $\Rightarrow y$ instructions)

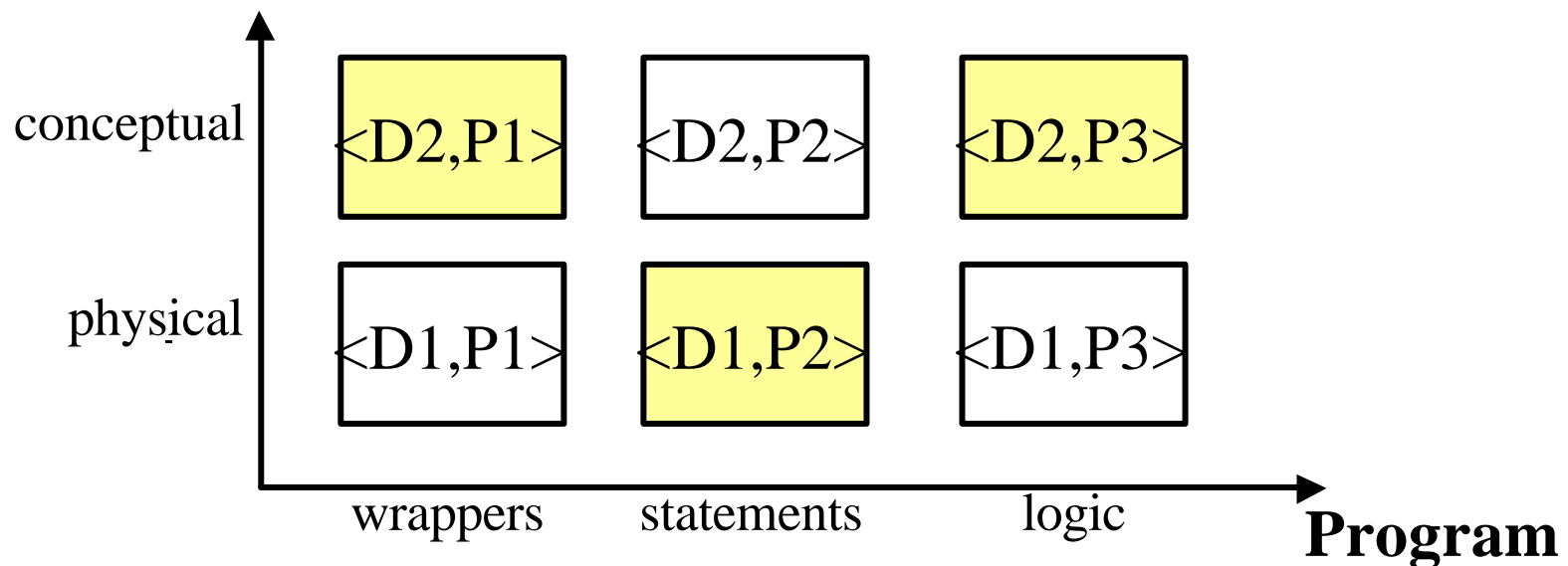


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Reengineering strategies (summary)

■ Six strategies

Database (schema)



■ <D1,P3> useless



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Wrapper strategy <D2,P1>

- modification of the legacy code is minimal,
new DB well structured and optimized w.r.t. the new DMS
- good solution for complete migration:
first the DB and later the programs... Illustration

READ PRODUCT

```
KEY IS PROD-CODE  
INVALID KEY  
    GO TO ERR-123.
```

DELETE PRODUCT

```
END-DELETE.
```



```
CALL WR-ORD-MNGMT  
    USING "READKEY", "PRODUCT",  
        "PROD-CODE",  
        PRODUCT, WR-STATE.  
IF STATUS OF WR-STATE NOT= 0  
    GO TO ERR-123.
```

```
CALL WR-ORD-MNGMT  
    USING "DELETE ", "PRODUCT",  
        " ", PRODUCT, WR-STATE.
```




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Statement rewriting strategy <D1,P2>

■ Illustration

```
MOVE CUS-CODE TO ORD-CUSTOMER.  
START ORDER KEY >= ORD-CUSTOMER.  
MOVE 0 TO END-FILE.  
PERFORM READ-ORD UNTIL END-FILE = 1.  
READ-ORD SECTION.  
BEG-ORD.  
READ ORDER NEXT  
  
AT END MOVE 1 TO END-FILE  
GO TO EXIT-ORD.  
  
<<processing current ORD record>>  
EXIT-ORD.  
EXIT.
```



```
EXEC SQL declare cursor ORD_GE_K1 for  
select ORD_CODE,ORD_CUSTOMER,ORD_DETAIL  
from ORDER where ORD_CODE >= :ORD-CODE  
order by ORD_CODE END-EXEC.
```

...

```
EXEC SQL declare cursor ORD_GE_K2 for  
select ORD_CODE,ORD_CUSTOMER,ORD_DETAIL  
from ORDER where ORD_CUSTOMER >= :ORD-CUSTOMER  
ORDER BY ORD_CUSTOMER END-EXEC.
```

...

```
MOVE CUS-CODE TO ORD-CUSTOMER.
```

```
EXEC SQL open ORD_GE_K2 END-EXEC.
```

```
MOVE "ORD_GE_K2" to ORD-SEQ.
```

```
IF ORD-SEQ = "ORD_GE_K1"
```

```
EXEC SQL fetch ORD_GE_K1 into :ORD-CODE,  
:ORD-CUSTOMER,:ORD-DETAIL END-EXEC
```

```
ELSE IF ORD-SEQ = "ORD_GE_K2"
```

```
EXEC SQL fetch ORD_GE_K2 into :ORD-CODE,  
:ORD-CUSTOMER,:ORD-DETAIL END-EXEC
```

```
ELSE IF ...
```

```
END-IF.
```

```
IF SQLCODE NOT = 0
```

```
MOVE 1 TO END-FILE GO TO EXIT-ORD.
```

```
<<processing current ORD record>>
```



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Statement rewriting strategy <D1,P2>

- modification of the legacy code is minimal, DB not restructured, mimics the legacy DB
- Quick and dirty solution....



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Logic rewriting strategy <D2,P3>

■ Illustration

DISP-ORD.

READ ORDER KEY IS ORD-CODE

INVALID KEY

GO TO ERR-ORD-NOT-FOUND.

PERFORM DISP-ORD-CUS-NAME.

...

DISP-ORD-CUS-NAME.

MOVE ORD-CUSTOMER TO CUS-CODE

READ CUSTOMER

INVALID KEY

DISPLAY "ERROR: UNKOWN CUST"

NOT INVALID KEY

DISPLAY "ORD-CODE: "

ORD-CODE NAME.



DISP-ORD.

EXEC SQL

SELECT O.CODE, C.NAME

INTO :ORD-CODE, :NAME

FROM ORDER O, CUSTOMER C

WHERE O.CUS_CODE = C.CODE

AND O.CODE = :ORD-CODE

END-EXEC.

IF SQLCODE = 0

DISPLAY "ORD-CODE: »

ORD-CODE NAME

ELSE

GO TO ERR-ORD-NOT-FOUND.



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Logic rewriting strategy <D2,P3>

- program is rewritten (long, difficult, risky)
new DB well structured and optimized w.r.t. the new DMS
programs optimized w.r.t. the new DMS
- good solution if no program migration planned, only the DB is migrated



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Conclusion

| Strategy | Database migration | Program conversion | Quality |
|---------------|-----------------------------|---|--|
| D2, P1 | complete DBRE, expensive | cheap, fully automated, wrapper semi- automatically generated | good quality DB, the programs unchanged (call to the wrapper) |
| D1, P2 | cheap, fully automated | cheap, fully automated | poor quality DB, the programs unchanged (call the new DML) |
| D2, P3 | complete DBRE, expensive | very expensive, requires a deep understanding of the programs | good quality DB, programs semi- renovated |



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Conclusion

| Strategy | Performance | Maintenance | Evolution |
|---------------|---|---|---|
| D2, P1 | Poor: legacy logic, mismatch, emulation | like the legacy system, but the semantics of the DB is known and data access simulated by the wrapper | easier, the new functions can directly access to the new DB |
| D1, P2 | Poor: legacy logic, mismatch | like the legacy system, the semantics of the DB is not recovered but data access are simulated by the new DML | difficult, the DB simulates the legacy one |
| D2, P3 | Good: new logic, matching | easier, the semantics of the DB is known | easier, the new functions can directly access to the new DB |