## A Cell-based Model for

# Multi-class Doubly Stochastic Dynamic Traffic Assignment

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#### **ABSTRACT**

This paper proposes a cell-based multi-class dynamic traffic assignment problem that considers the random evolution of traffic states. Travelers are assumed to select routes based on perceived effective travel time, where effective travel time is the sum of mean travel time and safety margin. The proposed problem is formulated as a fixed point problem, which includes a Monte-Carlo-based stochastic cell transmission model to capture the effect of physical queues and the random evolution of traffic states during flow propagation. The fixed point problem is solved by the self-regulated averaging method. Numerical examples are set up to illustrate the properties of the problem and the effectiveness of the solution method. The key findings include the following: i) Reducing perception errors on traffic conditions may not be able to reduce the uncertainty of estimating system performance, ii) Using the selfregulated averaging method can give a much faster rate of convergence in most test cases compared with using the method of successive averages, iii) The combination of the values of the step size parameters highly affects the speed of convergence, iv) A higher demand, a better information quality, or a higher degree of the risk aversion can lead to a higher computation time, v) More driver classes do not necessary results in a longer computation time, and vi) Computation time can be significantly reduced by using small sample sizes in the early stage of solution processes.

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#### 1. INTRODUCTION

Dynamic traffic assignment (DTA) is one of important and hot research topics nowadays because of its wide applications in transport planning and management. DTA consists of two components: *travel choice principle* and *traffic flow component* (Szeto and Lo, 2006).

The travel choice principle describes the travel choice behavior of the trip makers, especially their route choice. The commonly used route choice principles in DTA include: the dynamic user equilibrium principle (e.g., Lo and Szeto, 2002a,b), the dynamic stochastic user equilibrium principle (e.g., Szeto and Lo, 2005), the bounded rationality principle (e.g., Peeta and Zhou, 2006) and the tolerance-based principle (e.g., Szeto and Lo, 2006). These principles can be used in single and multiple user class DTA problems. However, these principles only consider the "mean" travel time and ignore the variability of travel time. In fact, Jackson and Jucker (1982) and Abdel-Aty *et al.* (1997) show that the variability of travel times is one of key factors in route choice. It is therefore important to capture the variability in the travel choice component. Many efforts (e.g., Chen *et al.*, 2002a; Lo *et al.*, 2006; Shao *et al.*, 2006, 2008; Lam *et al.*, 2008; Siu and Lo, 2008; Zhou and Chen, 2008; Connors and Sumalee, 2010; Szeto *et al.*, 2011) focus on incorporating the variability in the route choice component of static trip assignment models. However, to our best knowledge, very limited studies (e.g., Boyce *et al.*, 1999; Liu *et al.*, 2002; Yin *et al.*, 2004) consider travel time variability in DTA.

Traffic flow component depicts how traffic flow propagates within a traffic network and determines the travel time. Early efforts develop the traffic flow component in DTA based on point queue models (e.g., Carey and McCartney, 2004), exit flow models (e.g., Nie and Zhang, 2005a) and dynamic travel time functions (e.g., Nie and Zhang, 2005b). To capture the effect of spatial queues such as queue spillback and junction blockage, physical queue models (e.g., Daganzo, 1994, 1995, 1999; Muñoz *et al.*, 2003; Szeto, 2008) are developed

later. These models are often "macroscopic" and are derived from the hydrodynamic theory. They assume a steady-state speed-density relationship which adopts a number of deterministic parameters (e.g., free-flow speed, jam-density, capacity, etc.), and does not allow fluctuations around the equilibrium (nominal) fundamental flow-density diagram (FD). However, research and empirical studies on the FD have revealed that the FD admits large variations (see Figure 1) due to the variabilities in driving behavior and the characteristics (e.g., acceleration and deceleration abilities) of vehicles, the changing weather conditions, estimation error, and others (Chen *et al.*, 2001; Boel and Mihaylova, 2006; Kim and Zhang, 2008; Chiabaut *et al.*, 2009; Geistefeldt and Brilon, 2009; Ngoduy, 2009). Therefore, macroscopic "stochastic" traffic flow models (e.g., Boel and Mihaylova, 2006; Sumalee *et al.*, 2010) are developed to capture the random evolution of traffic states for freeways with onramps and off-ramps. Nevertheless, these models have not been extended for general network applications.

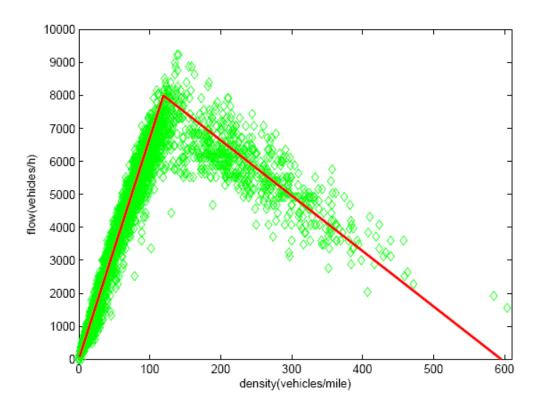


Figure 1 A fundamental flow density diagram of traffic flow with the 24-hour traffic flow data of Interstate 210 West in Los Angeles collected on April 22, 2008.

To address the effect of travel time variability on route choice, the random evolution of traffic states, and the effect of spatial queues simultaneously, this paper proposes a new DTA problem called the multi-class doubly stochastic dynamic user equilibrium problem (MDS-DUE-P). The MDS-DUE-P is to determine the temporal flow pattern in a stochastic traffic network given a fixed demand at each departure time period, where different classes of travelers are assumed to have imperfect information about the network conditions and different attitudes towards risks. The route choice component of the proposed DTA problem is described by the dynamic extension of the travel time budget or effective travel time concept in Lo et al. (2006) and Shao et al. (2006), which considers the standard deviation (SD) of travel time instead of its variance as in Yin et al. (2004). The traffic flow component is depicted by the proposed Monte-Carlo based stochastic cell transmission model (MC-SCTM). Free flow speed, backward wave speed, saturation flow rate, and jam density are modeled as random variables, which can be defined by any distributions. The effect of spatial queues can also be captured.

The MDS-DUE-P is formulated as a fixed point problem and solved by the self-regulated averaging method (SAM) proposed by Liu *et al.* (2006). The SAM works well for solving the deterministic network stochastic user equilibrium problem. However, to our best knowledge, the SAM has not been used to solve stochastic network assignment problems including ours, and hence the performance of the SAM for solving these problems has not been known yet. Hence, numerical examples are set up to the performance the solution method. Examples are also set up to illustrate the properties of the MDS-DUE-P. The main contribution of this paper is to propose a methodology to formulate and solve the problem that considers the random evolution of traffic states in a "macroscopic" traffic flow modeling framework while the effect of variability on route choice behavior is simultaneously taken into account.

The rest of the paper is organized as follows: Section 2 formulates the problem. Section 3

describes the MC-SCTM. Section 4 depicts the SAM. Section 5 presents the numerical studies. Section 6 discusses the implementation issues of the proposed model and the solution method to realistic networks. Finally, Section 7 gives some concluding remarks.

#### 2. PROBLEM FORMULATION

Consider a transportation network with multiple origin-destination (OD) pairs and multiple driver classes. The road segments are represented by a series of cells that have physical length, while the links are there merely to delineate connectivity between cells. Traffic begins at an origin cell (denoted as r) and terminates at a destination cell (denoted as s). The study horizon is discretized into many equal length intervals. The evolution of traffic states is assumed to be random so that capacity and travel time are uncertain.

All drivers in the same class have the same attitude towards the risk of late arrivals and have imperfect information on traffic network conditions such as the mean and SD of the travel time of each route. They are assumed to select routes based on the dynamic extension of the reliability-based stochastic user equilibrium principle (Shao *et al.*, 2006) called the *reliability-based stochastic-dynamic-user-equilibrium* (RSDUE) principle. This RSDUE principle states that for each class of drivers departing at any time, they select routes with the minimum perceived effective travel time at the time of departure.

The *perceived effective travel time* at the time of departure is defined as the sum of the modified effective route travel time at that departure time and the associated perception error:

$$U_{p,i}^{rs}(t) = V_{p,i}^{rs}(t) + \varepsilon_{p,i}^{rs}(t), \qquad (1)$$

where  $U_{p,i}^{rs}(t)$ ,  $V_{p,i}^{rs}(t)$ , and  $\varepsilon_{p,i}^{rs}(t)$  denote the perceived effective travel time, the modified effective travel time, and the perception error on route p between OD pair rs for class-i drivers departing at time t, respectively. There are many alternatives to model  $\varepsilon_{p,i}^{rs}(t)$ , namely the Multinomial Logit model, the C-Logit model, the Path Size Logit model, the

Probit model and others (see Ben-Akiva and Bierlaire, 1994). The C-logit model is selected for this framework due to its tractability, simplicity, and the ability to handle path correlation.

The *modified effective route travel time* in (1) is the sum of the commonality factor and actual effective route travel time:

$$V_{ni}^{rs}(t) = CF_n^{rs} + \hat{\eta}_{ni}^{rs}(t), \qquad (2)$$

where  $CF_p^{rs}$  is the commonality factor of route p, and  $\hat{\eta}_{p,i}^{rs}(t)$  denotes the actual effective travel time on route p between OD pair rs for class-i drivers departing at time t.

The *commonality factor* in (2) is used to capture the degree of similarity between paths as in Cascetta *et al.* (1996). It is expressed as follows:

$$CF_p^{rs} = \lambda \cdot \ln \sum_{h \in R_{rs}} \left( \frac{L_{hp}}{L_h^{1/2} \cdot L_p^{1/2}} \right)^{\psi},$$
 (3)

where  $L_{hp}$  is the "length" of links or free flow travel time on links common to both paths h and k.  $L_h$  and  $L_p$  are, respectively, the free flow travel times on paths h and p belonging to the same OD pair.  $\lambda$  and  $\psi$  are parameters.  $R_{rs}$  is the set of paths connecting OD pair rs.

The actual effective route travel time in (2) is the sum of the mean route travel time and the safety margin:

$$\hat{\eta}_{p,i}^{rs}\left(t\right) = \eta_{p}^{rs}\left(t\right) + s_{p,i}^{rs}\left(t\right),\tag{4}$$

where  $\eta_p^{rs}(t)$  is the mean travel time of all drivers on route p between OD pair rs departing at t.  $s_{p,i}^{rs}(t)$  denotes the safety margin of route  $p \in R_{rs}$  for class-i drivers departing at time t, which is the additional time reserved to counter the travel time variation to avoid late arrivals.

The safety margin  $s_{p,i}^{rs}(t)$  in (4) is a linear function of the SD,  $\sigma_p^{rs}(t)$ , of the travel time on route  $p \in R_{rs}$  as in Lo *et al.* (2006) and Shao *et al.* (2006):

$$S_{n,i}^{rs}(t) = Z_i \sigma_n^{rs}(t), \tag{5}$$

where  $Z_i$  is the parameter representing the degree of the risk aversion of class-i drivers. The larger is the value of  $Z_i$ , the greater is the degree of the risk aversion of class-i drivers. In particular,  $Z_i > 0$  if class-i drivers are risk-averse and tend to depart earlier to avoid late arrivals;  $Z_i = 0$  if class-i drivers are risk-neutral and ignore  $\sigma_p^{rs}(t)$  when selecting routes. Of course,  $Z_i$  can also be related to the importance of the trip, i.e., by trip purpose (higher  $Z_i$  for more important trip and vice versa). The value of  $Z_i$  can be calibrated using the survey data. The details of the survey setting and calibration procedure can be found in Jackson and Jucker (1982) but the SD is considered instead of the variance in this paper.

The motivation of (5) is that a higher degree of risk aversion implies a larger safety margin, and a traveler reserves more additional time for a more uncertain route, where the uncertainty is reflected by the SD of the route travel time.

The parameter  $Z_i$  in (5) is related to the probability  $\rho_i$  that the actual travel time is not greater than the actual effective route travel time:

$$P\{C_p^{rs}(t) \le \hat{\eta}_{p,i}^{rs}(t) = \eta_p^{rs}(t) + Z_i \sigma_p^{rs}(t)\} = \rho_i, \tag{6}$$

where  $C_p^{rs}(t)$  is the actual travel time on route p between OD pair rs when the traveler departs at time t. The probability  $\rho_i$  can be regarded as the within budget time reliability, where the budget is defined by the actual effective travel time. Then, a higher value of  $Z_i$  implies that the passenger is willing to have a higher probability of arriving on time or equivalently the actual travel time not being greater than the actual effective travel time.

 $\eta_p^{rs}(t)$  in (4) and  $\sigma_p^{rs}$  in (5) are functions of all the route flows  $\mathbf{f} = \left[f_{p,i}^{rs}(t)\right]$ :

$$\begin{bmatrix} \eta_p^{rs}(t) \\ \sigma_p^{rs}(t) \end{bmatrix} = \mathbf{\Phi}(\mathbf{f}), \tag{7}$$

where  $\Phi$  represents a unique mapping yielding the means and SDs of route travel times from

route flows. This mapping is defined by the stochastic traffic flow model and the travel time determination procedure. In this study,  $\Phi$  is modeled by the proposed MC-SCTM, which will be described in the next section.

Assume the *perception errors* in (1) follow the identical and independent Gumbel distributions with mean zero and identical SD. The RSDUE route flow pattern can then be formulated as:

$$f_{p,i}^{rs}(t) = W_{p,i}^{rs}(t) \cdot q_i^{rs}(t), \forall rs, p, t, i, \text{ and}$$
(8)

$$w_{p,i}^{rs}(t) = \frac{\exp(-\theta_i \cdot (\hat{\eta}_{p,i}^{rs}(t) - CF_p^{rs}))}{\sum_{k} \exp(-\theta_i \cdot (\hat{\eta}_{k,i}^{rs}(t) - CF_k^{rs}))}, \forall rs, p, t, i,$$
(9)

where  $f_{p,i}^{rs}(t)$  and  $w_{p,i}^{rs}(t)$  are, respectively, the flow and proportion of class-i drivers on route p between OD pair rs departing at time t;  $q_i^{rs}(t)$  is the demand of class-i drivers between OD pair rs at time t, which is assumed to be fixed;  $\theta_i$  is the dispersion parameter representing the effective travel time perception variation of class-i drivers. A higher value of  $\theta_i$  represents a better information quality received by the drivers. When  $\theta_i$  tends to infinity, the flow pattern tends to be the reliability-based user equilibrium flow pattern. If  $Z_i = 0$ , the flow pattern becomes the dynamic stochastic user equilibrium flow pattern. If  $\theta_i$  tends to infinity and  $Z_i = 0$ , the flow pattern tends to be the dynamic user equilibrium flow pattern.

The MDS-SDUE-P is to find **f** to satisfy (3)-(9). Since  $w_{p,i}^{rs}(t)$  is a function of the path flow  $f_{p,i}^{rs}(t)$ , this problem can be formulated as a fixed point problem:

$$\mathbf{f} = \mathbf{Y}(\mathbf{f}) = [w_{p,i}^{rs}(t) \cdot q_i^{rs}(t)], \tag{10}$$

where  $w_{p,i}^{rs}(t)$  is defined by (3)-(5), (7), and (9). As (10) can be reformulated as a variational inequality (VI), the existence and uniqueness of solutions for the proposed problem can be analyzed by using the existence and uniqueness theorems for VIs. If the means and SDs of

route travel times are continuous of  $\mathbf{f}$ , a solution exists to the proposed problem. The uniqueness of the solution further requires  $\mathbf{F}(\mathbf{f}) = \left[ f_{p,i}^{rs}(t) - w_{p,i}^{rs}(t) \cdot q_i^{rs}(t) \right]$  to be strongly monotone, which is almost unlikely in general cases.

### 3. THE MONTE-CARLO-BASED STOCHASTIC CELL TRANSMISSION MODEL

The proposed Monte-Carlo-based stochastic cell transmission model (MC-SCTM) is developed from the cell transmission model (CTM) proposed by Daganzo (1994). The CTM is a convergent numerical approximation scheme to the hydrodynamic model of Lighthill and Whitham (1955) and Richards (1956). Daganzo's CTM (1994) is originally developed for a simple highway based on a trapezoidal flow-density relation, which includes the triangular one as a special case. However, the CTM can be extended to consider general networks (Daganzo, 1995a), general flow-density relations (Daganzo, 1995b), priority vehicles and special lanes (Daganzo, 1997 and Daganzo *et al.*, 1997) and lane-changing behavior (Laval and Daganzo, 2006). In terms of formulation representation, the CTM can use either cell occupancy (e.g., Daganzo, 1995a,b) or cell density (Daganzo, 1999; Muñoz *et al.*, 2003) as the state variable, where the cell occupancy is defined as the number of vehicles in a cell, and the cell density is defined as the average number of vehicles in a cell per unit distance.

The CTM can capture traffic dynamics such as shockwaves, queues formation, queue dissipation, and dynamic traffic interactions across multiple links (Lo and Szeto, 2002a, b). However, the CTM requires dividing a link into many cells with equal lengths, leading to a high computation burden. To reduce computation time, various modifications to the CTM have been proposed. These modifications allow variable cell lengths (Ziliaskopoulos and Lee, 1997; Daganzo, 1999; Muñoz *et al.*, 2003; Ishak *et al.*, 2006; Szeto, 2008), and hence require few cells and can reduce the computation burden. Nevertheless, their developments are based on deterministic flow-density relations. The random evolution of traffic states due to

randomness in driving behavior, the changing weather conditions, etc. is not captured.

The stochastic cell transmission model (SCTM), proposed by Sumalee *et al.* (2010) considers the random evolution of traffic states. This model extends the CTM by defining parameters governing sending and receiving functions explicitly as random variables as well as by specifying the dynamics of the basic model parameters of the FD including free flow speed, backward wave speed, saturation flow rate, and jam density in each cell. The SCTM employs a triangular flow-density relation and uses densities as state variables instead of cell occupancies. Moreover, the SCTM allows variable cell lengths. The SCTM is shown to be able to give a good estimate on the mean actual traffic flow on freeways (Sumalee *et al.*, 2010). However, the SCTM does not consider route choice and the First-In-First-Out (FIFO) rule at diverges. Hence, the SCTM cannot be used to model traffic propagation in general networks. In addition, the correlations between the model parameters between different cells are not captured.

The proposed MC-SCTM considers route choice, FIFO, and the correlations between the model parameters of trapezoidal flow-density relationships of various cells at the expense of the increase in computation time. The MC-SCTM allows the parameters to be modeled by any distribution with correlation between the model parameters to fit the observed data better. The MC-SCTM relies on the Monte Carlo simulation (MCS) to generate traffic states for network loading and collect travel time statistics from each network loading. The network loading is done by the proposed modified CTM (M-CTM) which is extended from the existing network version of the CTM (Daganzo, 1995) that use occupancies as state variables. The M-CTM allows cells having unequal lengths regardless of actual free flow speed, and is more suitable for large networks with random evolution of traffic states than the existing network version of the CTM.

In the M-CTM, a highway is discretized into many cells and the study horizon is

discretized into many equal length intervals. The length of each cell is equal to the product of the maximum free flow speed for that cell and the length of each time interval. Then, the flow propagation equation of the M-CTM for a simple highway is the same as that of the CTM:

$$n_{j}(\omega+1) = n_{j}(\omega) - y_{j}(\omega) + y_{j-1}(\omega), \tag{11}$$

where the subscript j refers to cell j, and the subscript j-1 represents the cell upstream of j. The variables  $n_j(\omega)$  and  $y_j(\omega)$  denote the number of vehicles in cell j and the actual number of vehicles entering cell j at time  $\omega$ , respectively. Basically, (11) states that the number of vehicles in cell j at time  $\omega + 1$  is equal to the number of vehicles in that cell at time  $\omega$  plus those entering minus those leaving.

As in the CTM, the M-CTM uses the Godunov scheme (see Godunov, 1959) to set the actual number of vehicles entering cell j at time  $\omega$  to be the minimum of the maximum possible number of vehicles leaving the upstream cell j-1 at that time,  $S_{j-1}(\omega)$ , and the maximum possible number of vehicles entering cell j at that time,  $R_j(\omega)$ :

$$y_{j}(\omega) = \min \left\{ S_{j-1}(\omega), R_{j}(\omega) \right\}, \tag{12}$$

To account for non-uniform cell lengths, the M-CTM calculates  $S_{j-1}(\omega)$  and  $R_{j}(\omega)$  based on the following:

$$S_{j}(\omega) = \min \left\{ Q_{j}(\omega), n_{j}(\omega), \frac{N_{j}(\omega)}{d_{j}} \right\} \text{ and}$$
 (13)

$$R_{j}(\omega) = \min \left\{ Q_{j}(\omega), \frac{w_{j}(\omega)}{v_{j}(\omega)} \left[ N_{j}(\omega) - n_{j}(\omega) \right], \frac{N_{j}(\omega)}{d_{j}} \right\}, \tag{14}$$

where  $Q_j(\omega)$  denotes the inflow capacity of cell j or the maximum number of vehicles that can flow into cell j during time  $\omega$ .  $n_j(\omega)$  and  $N_j(\omega)$  are, respectively, the occupancy of cell j and the maximum number of vehicles that can be present in cell j (or the holding

capacity of cell j) at time  $\omega$ . The latter equals the product of the number of lanes, the jam density, and the length of the cell.  $w_j(\omega)$  and  $v_j(\omega)$  are, respectively, the backward wave speed and the free flow speed of cell j during time  $\omega$ .  $d_j$  is the length of cell j. Equations (13)-(14) assume that any cell j is made up of n "imaginary" standard cells connected in series, where n is a positive integer and a standard cell is the shortest cell in the network. According to (13), the sending flow of cell j is constrained by the inflow capacity, the cell occupancy, and the holding capacity of the last "imaginary" standard cell of cell j. Equation (14) states that the available capacity of cell j is the minimum of the inflow capacity of cell j, the product of the vacant space  $\left[N_j(\omega)-n_j(\omega)\right]$  and the factor  $w_j(\omega)/v_j(\omega)$  that accounts for the effect of shockwave on the vacant space in cell j, and the holding capacity of the first "imaginary" standard cell embedded in cell j.

In a general network with many OD pairs, the general principle of (11)-(14) is still applicable to the vehicle movements between cells. However, to maintain the intended routes of vehicles and the FIFO property, the number of vehicles in each cell is disaggregated by route (p) and by waiting time  $(\tau, \text{ modeled as a discrete time index})$ . The route variable is used to direct vehicles at merges and diverges. By tracking  $\tau$ , the FIFO property is maintained by ensuring that earlier arrivals (with a larger  $\tau$ ) will leave sooner.

For a simple highway, the conservation of traffic flow leads to the following

$$n_{j,p,\tau+1}(\omega+1) = n_{j,p,\tau}(\omega) - y_{j,p,\tau}(\omega), \text{ and}$$
 (15)

$$n_{j+1,p,1}(\omega+1) = \sum_{\tau} y_{j,p,\tau}(\omega), \tag{16}$$

where  $n_{j,p,\tau}(\omega)$  denotes the number of vehicles in cell j at time  $\omega$  on route p and enters the cell at time  $\omega - \tau + 1$ , and  $y_{j,p,\tau}(\omega)$  denotes as the number of vehicles (on route p and that has been waiting for  $\tau$  time intervals) leaving cell j at time  $\omega$ . Equation (15) states that the

remaining number of vehicles in cell j at time  $\omega + 1$  (determined by the difference between the number of vehicles in cell j at time  $\omega$  minus those leaving) increases their waiting time from  $\tau$  to  $\tau + 1$ . Equation (16) states that the vehicles enters cell j + 1 at  $\tau = 1$ .

If the minimum waiting time  $a_j(\omega)$  in cell j at time  $\omega$  is known, then  $y_{j,p,\tau}(\omega)$  in (15) and (16) can be determined by:

$$y_{j,p,\tau}(\omega) = \begin{cases} n_{j,p,\tau}(\omega) & \text{if } \tau > |a_{j}(\omega)|^{+}, \\ (\tau - a_{j}(\omega))n_{j,p,\tau}(\omega) & \text{if } \tau = |a_{j}(\omega)|^{+}, \\ 0 & \text{if } \tau < |a_{j}(\omega)|^{+}, \end{cases}$$

$$(17)$$

where  $\left|a_{j}\left(\omega\right)\right|^{+}$  denotes the smallest integer equal to or greater than  $a_{j}\left(\omega\right)$ . Equation (17) states that if the waiting time  $\tau$  of any vehicle is longer than  $\left|a_{j}\left(\omega\right)\right|^{+}$ , the vehicle advances to the next cell; none of the vehicles leave cell j if their waiting time is less than  $a_{j}\left(\omega\right)$ ; and lastly, for all the vehicles in cell j with waiting time  $\tau=\left|a_{j}\left(\omega\right)\right|^{+}$ , only  $\left(\tau-a_{j}\left(\omega\right)\right)n_{j,p,\tau}\left(\omega\right)$  of them leave the cell.

To avoid vehicles traveling at a speed higher than the free flow speed, according to the Courant-Friedrichs-Levy condition (Courant *et al.*, 1967), the minimum waiting time  $a_j(\omega)$ , or equivalently the minimum number of intervals required by a vehicle to stay in cell j, must be at least equal to the minimum number of intervals required to leave cell j,  $t_j(\omega)$ :

$$a_j(\omega) \ge t_j(\omega)$$
. (18)

The latter depends on the length and actual free flow speed of the cell as follows:

$$t_{j}(\omega) = \frac{d_{j}}{d_{0}} \cdot \frac{v_{\max,j}}{v_{j}(\omega)},\tag{19}$$

where  $v_j(\omega)$  and  $v_{\max,j}$ , respectively, denote the actual and maximum free flow speeds of cell

j with  $v_j(\omega) \le v_{\max,j}$ , and  $v_j(\omega) = v_{\max,j}$  when there is no uncertainty for  $v_j(\omega)$ .  $d_j$  and  $d_0$ , respectively, denote the lengths of cell j and the shortest cell. They represent the distance traveled by a vehicle at the maximum free flow speed during one time interval. Equation (19) states that  $t_j(\omega)$  is directly proportional to the length of cell j and inversely proportional to the actual free flow speed  $v_j(\omega)$ . In the extreme case, when  $d_j = d_0$  and  $v_j(\omega) = v_{\max,j}$ ,  $\forall j, \omega$ , the M-CTM is reduced to the CTM.

The M-CTM also has the following relationship:

$$n_{j}(\omega) = \sum_{\tau \ge t_{j}(\omega)} \sum_{p} n_{j,p,\tau}(\omega)$$
, and (20)

$$y_{j}(\omega) = \sum_{\tau} \sum_{p} y_{j,p,\tau}(\omega). \tag{21}$$

Equations (12)-(21) can be used to determine the traffic propagation for a simple highway over time. Suppose the disaggregate occupancies  $n_{j,p,\tau}(\omega), \forall j,p,\tau$  and the necessary boundary conditions are known. Then, we can obtain  $n_j(\omega), \forall j$  by (20). After that, by solving the system of equations (12)-(14), (17), (19), (21), we can determine  $a_j(\omega)$  and  $y_{j,p,\tau}(\omega)$  simultaneously, which are used to compute  $n_{j,p,\tau+1}(\omega+1)$  and  $n_{j+1,p,1}(\omega+1), \forall j,p$  through (15) and (16), respectively. The new disaggregate occupancies  $n_{j,p,\tau}(\omega+1), \forall j,p,\tau$  can then be used to obtain the disaggregate occupancies for next time interval. This procedure is repeated in the ascending order of  $\omega$ .

Compared with Daganzo (1995), the new feature of the M-CTM is described by (13), (14), (19), and (20). For general networks, the M-CTM still adopts the merge and diverge concepts proposed in Daganzo (1995) but the maximum possible numbers of vehicles leaving and entering a cell on each branch are determined by (13) and (14), respectively, and the disaggregate cell occupancy on each branch is determined by (19) and (20).

The M-CTM discussed above is to determine the travel time given the path inflow  $\mathbf{f}^k$  for one realization in the MC-SCTM. To obtain the means and SDs of route travel times, many realizations are generated by the MCS. To reduce computation time, the MC-SCTM incorporates the common random number MCS method and recursive formulas for determining the means and SDs of route travel times. The algorithmic steps of the MC-SCTM are as follows:

- 1.) Select the maximum number of samples for each iteration  $M_{\rm max}$ . Set the sample number m=0, the estimates of the means and SDs of route travel times  $\bar{\eta}_{p,t,m}^{rs}=0, \; \bar{\sigma}_{p,t,m}^{rs}=0, \forall rs,p,t$ .
- 2.) Set m=m+1. Generate an observation  $\mu^m$  for free flow speed, backward wave speed, jam density, and saturation flow for each cell at each time interval from the predefined distributions.
- 3.) Load the path inflows  $\mathbf{f}^{k}$  into the cell network using the M-CTM.
- 4.) Deduce the cumulative inflow and outflow curves. Calculate the travel time  $\eta_{p,t}^{rs}(\mathbf{f}^k, \mathbf{\mu}^m), \forall rs, p, t$  using the procedure in Tong and Wong (2000), Lo and Szeto (2002) or Long *et al.* (2010).
- 5.) Update the mean route travel time and the corresponding SD based on

$$\overline{\eta}_{p,t,m}^{rs} = \overline{\eta}_{p,t,m-1}^{rs} + \frac{\eta_{p,t}^{rs} \left(\mathbf{f}^k, \mathbf{\omega}^m\right) - \overline{\eta}_{p,t,m-1}^{rs}}{m}, \forall rs, p, t, \text{ and}$$
(22)

$$\overline{\sigma}_{p,t,m}^{rs} = \sqrt{\left(1 - \frac{1}{m-1}\right) \left[\overline{\sigma}_{p,t,m-1}^{rs}\right]^2 + m\left(\overline{\eta}_{p,t,m}^{rs} - \overline{\eta}_{p,t,m-1}^{rs}\right)^2}, \forall rs, p, t . \tag{23}$$

6.) If  $m < M_{\text{max}}$ , go to 2.). Otherwise, set  $\eta_p^{rs}(t) = \overline{\eta}_{p,t,M_{\text{max}}}^{rs}$ , and  $\sigma_p^{rs}(t) = \overline{\sigma}_{p,t,M_{\text{max}}}^{rs}$ .

To sum up, given the path inflow  $\mathbf{f}^k$ , the MC-SCTM can calculate the occupancy of each cell together with the means and SDs of route travel times.

#### 4. THE SELF-REGULATED AVERAGING SOLUTION METHOD

The method of successive averages (MSA) was extensively used to solve the fixed point problem during the past decades. This method relies on a predetermined step size for guaranteeing convergence. However, the convergence speed of the MSA is slow due to two reasons: 1) The predetermined step size is too large at some iterations such that the next solution is even farther away from the optimal solution than the current one (i.e., overshooting occurs), and 2) When the current solution is close to the optimal solution, the predetermined step size is too small and the convergence process is slowed down.

A self-regulated averaging method (SAM) was proposed by Liu *et al.* (2009) to deal with the slow convergence problem, and was used to solve the deterministic-network stochastic user equilibrium (SUE) problem. Nevertheless, to our best knowledge, this method has not been used to solve stochastic network assignment problems yet. In this paper, we adopt this method to solve our proposed stochastic problem, which relies on the Euclidean distances  $\|\mathbf{f}^k - \mathbf{Y}(\mathbf{f}^k)\|$  in the two consecutive iterations to select an appropriate step size. If  $\|\mathbf{f}^k - \mathbf{Y}(\mathbf{f}^k)\| \le \|\mathbf{f}^{k-1} - \mathbf{Y}(\mathbf{f}^{k-1})\|$ , a step size larger than that of the MSA at iteration k will be chosen to speed up the convergence process. Otherwise, a step size smaller than that of the MSA at iteration k will be chosen to reduce the chance of overshooting.

Due to the sampling errors of MCS, the error estimate for each (major) iteration may not be exact. Therefore, instead of just using the Euclidean distance at the current iteration k, we use a simple moving average of the last three Euclidean distances as the convergence measure. The detailed algorithm steps are as follows:

1.) Initialization: Set the number of major iterations k=1 and the initial step size  $\beta^0=1$ . Define parameters such as the maximum number of iterations  $k_{\max}$ , the convergence tolerance  $\mathcal{E}$ , the step size parameters  $\Gamma$  and  $\gamma$ . Set the initial route flows  $\mathbf{f}^0$ .

- 2.) Updating the mapping function: Based on  $\mathbf{f}^k$ , determine the means and SDs of route travel times by the MC-SCTM. Calculate  $\mathbf{Y}(\mathbf{f}^k)$ .
- 3.) Checking convergence: Report results and stop if i)  $k > k_{\text{max}}$ , or ii)  $\frac{1}{3} \sum_{i=k-2}^{k} \left\| \mathbf{f}^{j} \mathbf{Y}^{j}(\mathbf{f}^{j}) \right\| \leq \varepsilon, \ k \geq 3, \text{ where}$

$$\|\mathbf{f}^{k} - \mathbf{Y}(\mathbf{f}^{k})\| = \sum_{rs} \sum_{p} \sum_{i} \sum_{t} \left[ \sqrt{(f_{p,i}^{rs}(t) - w_{p,i}^{rs}(t) \cdot q_{i}^{rs}(t))^{2}} \right].$$
 (24)

4.) Step size determination: The self-regulated step size  $a^k = 1/\beta^k$ , and  $\beta^k$  is updated by:

$$\boldsymbol{\beta}^{k} = \begin{cases} \boldsymbol{\beta}^{k-1} + \boldsymbol{\Gamma}, \boldsymbol{\Gamma} > 1, & \text{if } \|\mathbf{f}^{k} - \mathbf{Y}(\mathbf{f}^{k})\| \ge \|\mathbf{f}^{k-1} - \mathbf{Y}(\mathbf{f}^{k-1})\|, \\ \boldsymbol{\beta}^{k-1} + \boldsymbol{\gamma}, \boldsymbol{\gamma} < 1, & \text{if } \|\mathbf{f}^{k} - \mathbf{Y}(\mathbf{f}^{k})\| < \|\mathbf{f}^{k-1} - \mathbf{Y}(\mathbf{f}^{k-1})\|. \end{cases}$$
(25)

5.) Updating the flow pattern:  $\mathbf{f}^{k+1} = \mathbf{f}^k + a^k \cdot (\mathbf{Y}^k(\mathbf{f}^k) - \mathbf{f}^k)$ , k = k+1 and go to 2.) Note that when  $\Gamma = \gamma = 1$ , the SAM becomes the MSA.

#### 5. NUMERICAL STUDIES

Eight examples are set up to illustrate the properties of the problem and the performance of the solution method. These examples are summarized in Table 1. For examples 1-5, 7 and 8, a 10-cell network in Figure 2 is used. This network has one OD pair where cell 1 is the origin cell and cell 10 is the destination cell. The basic settings for these examples are as follows:

- Cell lengths: a) cells 3, 4, and 6: 5 miles; b) cell 7: 3 miles, and c) other cells: 1 mile;
- The number of lanes: a) cells 5 and 8: 2; b) cells 3, 4, 6, and 7: 3; c) other cells: 6;
- Convergence tolerance:  $\varepsilon = 0.01$ ;
- Length of a discretized time interval: 1 minute;
- Sample size of the MCS: 1000;
- Route choice parameters:  $Z=1.65, \ \theta_1=0.05; \ \lambda=\psi=1;$

- Demand: a) level: 1800 veh/hr/lane; b) loading period: [1,15];
- Parameters for the FD of each cell: symmetric triangular distributions with
  - a) Free flow speed: The maximum is 60 mph, and the minimum is 50 mph;
  - b) Backward wave speed: The maximum is -10 mph, and the minimum is -20 mph;
  - c) Jam density: The maximum is 240 veh/mile, and the minimum is 160 veh/mile;
  - d) Saturation flow: The maximum is 1800 veh/hr/lane, and the minimum is 1620 veh/hr/lane, except cells 6 8, where the minimum is 1080 veh/hr/lane.

With these settings, the lower route is shorter than the upper one, but has a larger SD and a lower mean saturation flow.

For example 6, the Mong Kok district (see Figure 3) is considered and modeled as shown in Figure 4. The network has 4 OD pairs and 33 cells. The shortest and longest cells are 0.02 mile and 0.08 mile long, respectively. Each time interval is 5 seconds long. The rest of parameter settings are similar to those for the small network, except the following:

- Cell lengths: a) cells 5, 12, and 21: 0.08 mile; b) cells 6, 7, 8, 11, 22, and 23: 0.06 mile,
   and c) cells 15, 16, 18, 19, 20, 29, and 30: 0.02 mile; d) other cells: 0.04 mile;
- The numbers of lanes: a) cells 1 to 9: 3; b) cells 10, 11, 13, 14, 24, 26, 27, 28, 32, and 33: 2; c) other cells: 1;
- Demand levels: OD pair 1-9: 960 veh/hr/lane; OD pair 10-9: 1080 veh/hr/lane, and OD pairs 1-33 and 10-33: 720 veh/hr/lane;
- Parameters for the FD of each cell: symmetric triangular distributions with
  - a) Free flow speed: The maximum is 14.4 mph, and the minimum is 8.64 mph;
  - b) Backward wave speed: The maximum is -6 mph, and the minimum is -8 mph.

The nominal fundamental diagrams for the two networks are shown in figure 5.

All examples were run on a computer with an Intel (R) Core(TM) 2 Duo T9400 2.53GHz CPU and a 3GB RAM. Unless stated otherwise, each example is run for 20 times with

different seeds, and the average results are reported. Initially, the demand in each period is evenly assigned to each path at the first iteration.

Example	Network	No. of	Illustration	Setting/ Remarks
		classes N		
1	10-cell	1	Effects of $\Gamma$ and $\gamma$ on	Different combinations of
	network		computation time	$\Gamma$ and $\gamma$
2	10-cell	1	Effect of $\theta$ on computation	$\theta = 0.05, 0.1, 0.15, 0.2,$
	network		time	0.25
3	10-cell	1	Effect of demand on	Demand = 600, 900, 1200,
	network		computation time	1500, 1800 veh/hr/lane
4	10-cell	1	Effect of Z on computation	Z = 0, 0.825, 1.65
	network		time	
5	10-cell	varied	Effect of the number of	N = 1, 2,  and  3
	network		classes on computation time	
6	Mong	1	Computation performances	Change the sample size of
	Kok		of constant and varying	MCS from 100 to 1000 at
	Network		sample size strategies	different convergence levels
7	10-cell	1	Effect of the degree of risk	See below
	network		aversion and $\theta$ on route	
			choice	
8	10-cell	2	Effect of information quality	See below
	network		on the mean and SD of total	
			system travel time	

**Table 1 Example setting** 

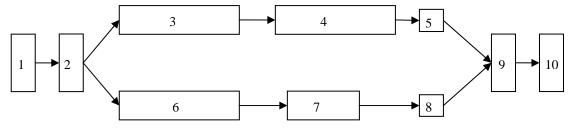


Figure 2 A ten-cell network for examples 1-5, 7 and 8

# Example 1

This example is set up to illustrate the effect of the different combinations of  $\Gamma$  and  $\gamma$  on computation time and compare the performance between the SAM and the MSA. In this example,  $\Gamma$  is in the range from 1.0 to 2.0, while  $\gamma$  is in the range from 0.2 to 1.0, thereby capturing the range recommended by Liu *et al.* (2006). The increments for these parameters

are set to 0.2, and hence 30 combinations can be obtained as shown in Table 2. In particular, when  $\Gamma=1.0$  and  $\gamma=1.0$ , the test case represents the case using the MSA.

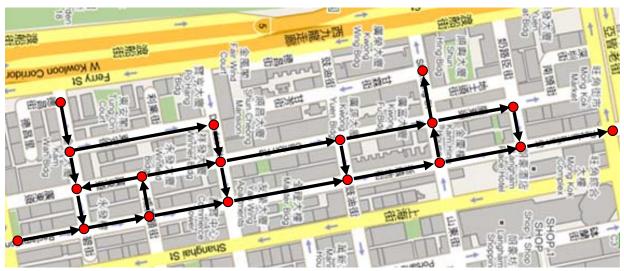


Figure 3 The Mong Kok district in Hong Kong

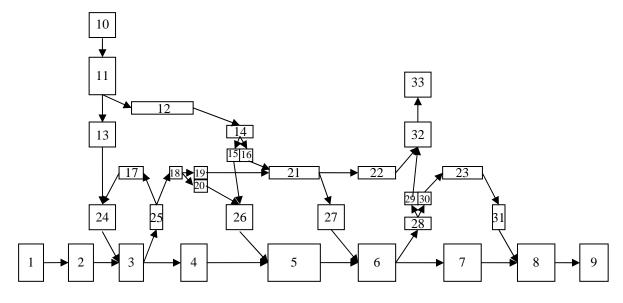


Figure 4 The Mong Kok network for example 6

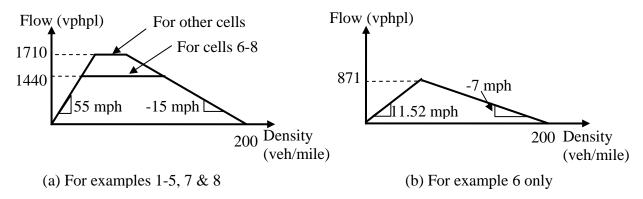


Figure 5 The nominal flow-density diagrams for the example networks

Table 2 The effect of step size on number of iterations required

		Γ							
		1	1.2	1.4	1.6	1.8	2		
γ	1	176.2	212.5	249.3	282.0	321.0	355.5		
	0.8	86.4	102.9	117.8	135.3	151.2	170.7		
	0.6	44.4	51.9	59.6	67.9	75.6	82.8		
	0.4	25.0	28.7	32.2	36.1	40.6	44.2		
	0.2	15.7	17.7	19.4	21.3	23.2	25.2		

The numbers reported in table 2 are the average numbers of iterations required to obtain optimal solutions. From this table, we can observe that a higher value of  $\gamma$  results in a larger number of iterations required on average and hence a longer computation time. This means that a large step size is preferable for this problem. While it may not be always true to use a very large step size for all the problems, this example clearly shows that the combination of the step size parameters highly affects the speed of convergence. The ratio of the average number of iterations of the worst combination to that of the best combination is larger than 20. More importantly, this example clearly demonstrates the benefit of using the SAM over the MSA to solve stochastic problems, which has not been examined before. The computation time of the best combination (i.e.,  $\Gamma$ =1.0,  $\gamma$ =0.2.) is even less than one-tenth of that of the MSA. This is because for most of the time, the MSA used too small step sizes along the descent direction, whereas the SAM can regulate the step size to speed up the convergence process. Moreover, for all the test cases shown, all 24 out of 29 cases can obtain optimal solutions quicker than the MSA, as these 24 cases use  $\gamma$  <1.0 which allows increasing step sizes for speeding up the solution process when necessary. For the remaining 5 cases, ( $\Gamma > 1$  and  $\gamma = 1.0$ ),  $\Gamma > 1$  causes the initial overshooting (the new solution is farther away from the optimal solution than the initial solution) but there is no speeding-up strategy. Therefore, these 5 cases require a longer computation time than the case using the MSA. In particular, the larger is the value of  $\Gamma$ , the worse is the performance of the SAM, as the overshooting problem is more serious.

### Example 2-Example 4

Examples 2-4 aim to illustrate the effect of parameters on computation time. For each example, to avoid being smeared by other factors, we only vary one parameter value while keeping other values fixed as those used in the basic setting. The best combination of  $\Gamma$  and  $\gamma$  (i.e.,  $\Gamma$ =1.0,  $\gamma$ =0.2.) is used to solve for solutions. The results are presented in Figures 6-8. Figure 6 show that information quality has significant effect on computation time as reflected by the average number of iterations required to get optimal solutions. According to Figure 6, when  $\theta$  increases from 0.05 to 0.25, (representing the case when the information quality increases), the average number of iterations required increases nearly double. While the exact increase of average computation time depends on the specific problem, the increasing trend here agrees with the observation obtained by solving the SUE problem by the MSA under various values of  $\theta$  (see Sheffi, 1985), where the SAM is a variant of the MSA. When  $\theta$  is small, effective travel time is less important and the information quality plays a more important role in determining the flow pattern. In the limiting case, when  $\theta$ =0, effective travel time is not used for determining the flow pattern.

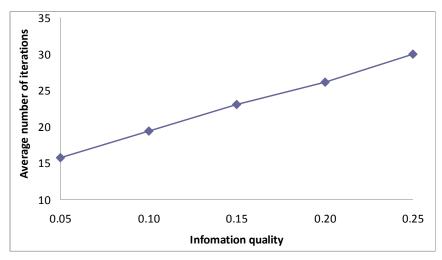


Figure 6 Effects of information quality on computation time

When the demand level increases from 600 veh/hr/lane to 1800 veh/hr/lane, the network condition changes from free-flow to congested, and the average number of iterations

increases nearly four times (see Figure 7). This observation is again consistent with the one obtained by solving the SUE problem by the MSA. When the demand is low, the change in effective travel time in successive iterations is small, and hence the convergence speed is fast.

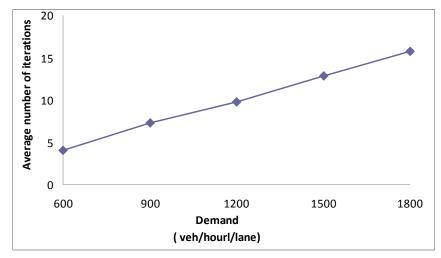


Figure 7 The Effect of demand on computation time

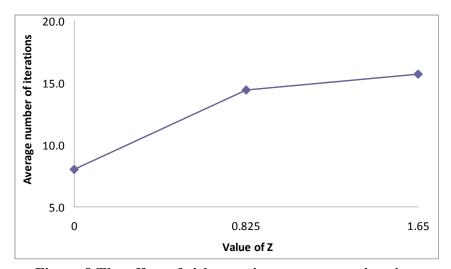


Figure 8 The effect of risk aversion on computation time

Figure 8 demonstrates the effect of degree of the risk aversion of drivers on computation time. Basically, a higher value of Z results in a higher number of iterations required, as a higher Z implies a larger safety margin and a larger effective travel cost. When Z is small, the change in effective travel cost in successive iterations is small, and hence the average computation time is small. However, the relationship between the average computation time and Z is not linear, and the increase in Z is marginal when Z is greater than 0.825. This phenomenon deserves a deeper analysis. We leave this to future studies.

### Example 5

Table 3 compares the means and SDs of the computation times of the three cases (i.e., N = 1, N = 2, and N = 3) over 20 runs. (Note that the average level of information quality and the total demand of all the cases are the same in order to have a fair computation). The results of the t-test on the difference in their means are shown in Table 4. The t-test shows that the computation time for the case of N = 1 is smaller than that of N = 2 and that of N = 3 at the significant level of 5 %, but the computation time for the case of N = 2 is greater than that of N = 3 at the same significant level. These results indicate that increasing more variables in the problem does not necessarily increase computation time.

Table 3 The computation times in seconds for various number of user classes

	N=1	N=2	N=3
	$(\theta_1 = 0.1,$	$(\theta_1 = 0.05, \theta_2 = 0.15,$	$(\theta_1 = 0.05, \ \theta_2 = 0.1, \ \theta_3 = 0.15,$
	$q_1^{1,10}(t)=1800$ )	$q_1^{1,10}(t) = q_2^{1,10}(t) = 900$ )	$q_1^{1,10}(t) = q_2^{1,10}(t) = q_3^{1,10}(t) = 600$
Mean	58.17	91.12	80.44
SD	2.07	12.07	11.18

Table 4 t-tests on the difference in mean computation times

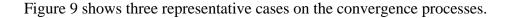
Comparisons	Difference in mean	<i>t</i> -value	Probability
Between the cases of N=1 and N=2	-32.95	-12.03	0.000
Between the cases of N=1 and N=3	-22.76	-8.76	0.000
Between the cases of N=2 and N=3	10.68	2.90	0.003

#### Example 6

This example aims to illustrate the computation performance of varying sample sizes during the solution process. The Mong Kok network is used. Drivers are assumed to select paths without cycles. Three strategies are considered:

- 1. Strategy a: A constant sample size of 1000 is used throughout the computation process.
- 2. Strategy b: A sample size of 100 is used initially. A sample is 1000 is used when the error  $\|\mathbf{f}^k \mathbf{Y}(\mathbf{f}^k)\|$  is less than 1.

3. Strategy c: Similar to strategy b except switching occurs at the error of 0.1.



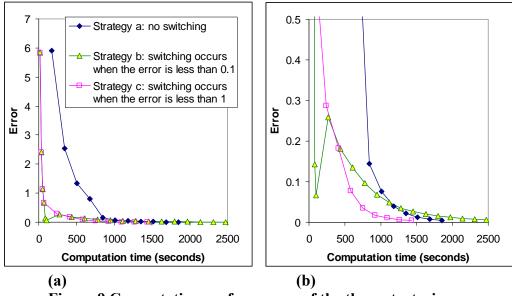


Figure 9 Computation performances of the three strategies

While the three strategies can decrease the error quickly in very early stages (figure 9a), the three strategies show markedly different convergence behavior. As shown in Figure 9a, the error of strategy a is monotonically decreasing in a smooth manner. However, the errors of strategies b and c are not because of the occurrence of switching (see figure 9b). In the worse case, the switching in strategy c causes a temporary increase in error due to the realization of more accurate means and SDs of route travel times using a large sample. This prolongs the convergence process compared with the no-switching strategy (Strategy a).

The average performances of 20 runs for the three strategies are shown in Table 5 and the t-test results on their differences in means at the significant level of 0.05 are reported in Table 6. This table shows that their means are quite difference at this significant level. Strategy b is better than strategy a by 29% in terms of their mean performances, which is in turn better than strategy c by 30%. The implication is that while switching at earlier stages can reduce computation time, switching at later stages can increase computation time. A careful selection of the parameter for triggering switching is necessary to optimize the speed of convergence.

Nevertheless, in general, switching at earlier stages of the solution process can reduce unnecessary sampling, as we only need a rough descent direction anyway. However, at later stages, we need to have accurate mean and SD to find out an exact descent direction and an optimal solution. Switching too late can prolong the convergence process.

Table 5 The computation times in seconds for various computation strategies

	Strategy a	Strategy b	Strategy c
Mean	1930.83	1500.01	2503.99
SD	83.74	269.03	332.31

Table 6 t-tests on the difference in mean computation times

	Difference in mean	<i>t</i> -value	Probability
Comparison between strategies a and b	430.82	6.84	0.000
Comparison between strategies a and c	-573.16	-7.48	0.000
Comparison between strategies b and c	-1003.99	-10.50	0.000

## Example 7

To illustrate the effect of the degree of risk aversion and information quality on route choice, four cases are set up: (1) Z=0.00,  $\theta=0.05$  (risk-neutral drivers with poor information quality), (2) Z=0.00,  $\theta=0.10$  (risk-neutral drivers with better information quality), (3) Z=1.65,  $\theta=0.05$  (risk-averse drivers with poor information quality), and (4) Z=1.65,  $\theta=0.10$  (risk-averse drivers with better information quality).

The optimal result of each case is depicted in Figure 10.

In Figure 10, we can see that the four lines are distinct implying that different route choice assumptions give significantly different equilibrium route flow proportions. For example, for the 1<sup>st</sup> departure time interval, the difference in route flow proportions between cases 1 and 4 is about 8%. This observation further implies that the route choice assumption can greatly affect the results of performance evaluation because most performance measures, e.g., total system travel time (TSTT), highly depend on the flow proportions.

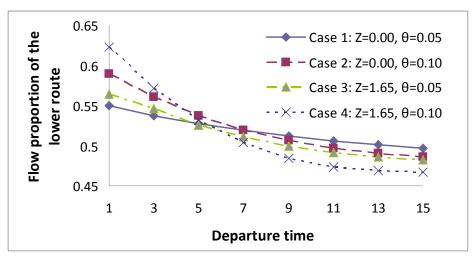


Figure 10 Proportions of drivers on the lower route

We can also observe that the best route changes with the departure time. When the route flow proportions are above 0.5, the route is the best one. According to figure 10, for all cases, the lower route (with a shorter mean travel time but a higher SD) is the best route only during early periods. For cases 1 and 2 (i.e., risk-neutral drivers), the lower route is better in terms of mean travel time in early periods. However, since the mean saturation flow on the lower route is less than that of the upper route, the rate of increase of the mean travel time on the lower route is faster than that of the upper one. Eventually, the mean travel time of the upper route becomes better than that of the lower route. For cases 3 and 4 (i.e., risk-averse drivers), similar to the case with risk-neutral drivers, the lower route is better initially in terms of effective travel time, and then the mean and SD of travel time on this route become worse than those on the upper route in later time periods. In particular, the increase of the SD of the travel time on this route is because the queue formation occurs on this route and the uncertainty of queue length accumulates over time.

The best route also varies from case to case. Consider the 9<sup>th</sup> departure time interval, the proportion is greater (less) than 0.5 for case 1 (case 3), which means that the lower (upper) route is better for case 1 (case 3). The implication is that when only mean travel time is considered, the lower route is better as the mean is lower. However, when the effective travel

time is considered, the upper route is better due to its lower travel time variability. Similar findings can also be observed when comparing between cases 2 and 4. For the effect of perception errors, a higher  $\theta$  value results in more travelers selecting the better route. For example, consider the 1<sup>st</sup> departure time interval. More drivers are on the better route in case 2 (4) than in case 1 (3).

### Example 8

The setting of this example is similar to the basic setting except that there are two classes of drivers: risk-neutral drivers (i.e.,  $Z_1=0$ ) and risk-averse drivers (i.e.,  $Z_2=1.65$ ). Four scenarios are considered as shown in table 7. For each scenario, two levels of information quality are considered: low (i.e.,  $\theta_1=\theta_2=0.05$ ) and high (i.e.,  $\theta_1=\theta_2=1.0$ ). The objective is to examine the effect of information quality on the mean and SD of TSTT obtained by the MCS after the optimal flow pattern is determined. The procedure is similar to that for determining the mean and SD of route travel time but the route travel times in the formulas are replaced by TSTT.

Table 7 The effect of information quality and the degree of risk aversion on the mean and SD of TSTT

	Demand	Proportion Proportion		$\theta = 0.05$		θ=1.0	
	(veh/hr/lane)	of risk- neutral	of risk- aversion	Mean (min)	SD (min)	Mean (min)	SD (min)
Scenario 1	1800	10%	90%	79970	7034	87145	7903
Scenario 2	1800	90%	10%	79925	6595	86001	8120
Scenario 3	1200	10%	90%	44615	3186	49159	3258
Scenario 4	1200	90%	10%	44570	3329	48477	3635

From table 7, we can observe that when the information quality is improved, both the mean and SD of TSTT increase. This observation implies that it may not be cost-effective to improve the traffic information provision system where the original objective is to provide

better information for travelers to avoid passing through the congested area and reduce the uncertainty of trip travel time due to unpredictable congestion. While the mean is increased by around 9 % in all scenarios, the SD is increased by 2-23%, depending on scenarios. The extent of increase in SD depends on both the driver composition and the demand level. When there are more risk-neutral drivers as in scenarios 2 and 4, improving information quality results a higher increase in the SD of TSTT, because more risk-neutral drivers will select the route with a lower mean travel time but a higher travel time variability. On the other hand, when the demand level is higher (i.e., Scenarios 1 and 2), the percentage increase of the SD of TSTT is also higher, since a more congested network leads to more uncertainty in trip time.

### 6. IMPLEMENTATION ISSUES TO REALISTIC NETWORKS

When implementing the proposed fixed-point model (including the submodel, the MC-SCTM) and the solution method to realistic networks, the following issues should be considered. First, the proposed cell-based modeling requires the following two conditions to be satisfied simultaneously: 1) The cell length must equal the maximum free flow speed times the length of a discretized time interval, and 2) The lengths of longer cells must be multiples of the length of the shortest cell. These two conditions will restrict the numbers of discretized time intervals and cells used, and the lengths of each discretized time interval and cell. Normally, short time intervals and short cells can satisfy these conditions. However, these can increase the computation time as the number of cells and time intervals increase.

Second, one cell for a corridor should be considered to reduce the number of cells used. However, if there are lane additions or lane drops within a link, more cells should be considered for modeling a link for increasing the modeling accuracy at the expense of increasing computation time. For modeling junctions, sometimes dummy cells are added to apply the merge and diverge models proposed in Daganzo (1995). More general intersection

models (e.g., Ishak, et al., 2006) can be used to avoid using dummy cells.

Third, the path set is assumed to be given in the SAM. In a realistic network, not many paths between an OD pair will be used by travelers. The used path set indeed is not large. This path set can be generated by the k-shortest path algorithms using a realistic number of k. According to Cascetta *et al.*, (1997), 4 to 7 paths are enough to obtain a reasonable solution. Alternatively, this path set can be obtained from personal interviews and hence constitutes a set of likely used paths (Chen *et al.*, 2002b).

Fourth, the proposed MC-SCTM requires the MCS that can lead to long computation times for large networks. One of the numerical studies has tested the idea of using smaller sample sizes for early stages of solution processes and larger ones at the later stages. The results are quiet promising in terms of saving computation time. Although the optimal parameter for triggering switching varies case by case, this strategy can be used for reducing the computation time to solve large problems by selecting a nearly optimal switching triggering value. Another strategy is to extend the analytical SCTM proposed by Sumalee et al. (2010) for general networks, and use the analytical SCTM to approximately solve the problem in the early or some stages of the solution process. One more strategy is to use antithetic sampling and Latin hypercube sampling techniques, and single point approximation techniques to reduce the sample sizes and hence the computation time while maintaining an acceptable level of solution quality. According to Sharma et al. (2010), using Antithetic Sampling can reduce the sample size 10 times to solve in their problem about the Network of Fort Area, Mumbai while reducing the computation time to 20% of the original time required. Using single point approximations can further lead to a 99% computation time saving. We expect the computation time can be reduced significantly by introducing these techniques in our solution method, but the exact reduction magnitudes may be slightly different. Due to space limitation, we did not include their comparison in this paper but they are definitely

deserved to be further studied and compared comprehensively.

Fifthly, the proposed model requires a fairly accurate time-dependent OD matrix. If this matrix is not accurate enough, one can consider extending the model to consider departure time choice as well, as the simultaneous route and departure time choice model only requires an OD matrix like the one for the static traffic assignment. The cost is to estimate the desirable arrival time and desirable arrival time interval of travelers. For the peak hour traffic, this information is not difficult to obtain from surveys. Alternatively, one can consider time-dependent OD matrix with a larger discretized time interval and the demand for each larger interval is assumed to be equally split into several smaller time intervals for network loading.

### 7. CONCLUDING REMARKS

This paper proposes the MDS-DUE-P considering the random evolution of traffic states, the effect of travel time variation on route choice, and the effect of spatial queues. The proposed problem consists of two components: traffic flow and route choice components. The traffic flow component extends the SCTM proposed in Sumalee *et al.* (2010) to consider general networks. The route choice component is described by the dynamic extension of the effective travel time concept. The proposed problem is formulated as a fixed point problem and solved by the SAM. Numerical examples are set up to illustrate the properties of the problem and performance of the solution method. The key findings and implications are as follows:

- 1. Reducing perception errors on traffic conditions may not be able to reduce the uncertainty of estimating system performance. This has a strong implication on providing better but free traffic information to travelers, as providing better information requires investment but the investment may not result in reducing the occurrence of the worse case performance.
- 2. Using the SAM can get solutions much quicker compared to using the MSA in most cases.

Moreover, the combination of the values of the step size parameters highly affects the speed of convergence. This implies that when the combination of the parameters is wrongly chosen, the rate of convergence can be worse than that of the MSA.

- 3. A higher demand, a better information quality, or a higher degree of the risk aversion can lead to a higher computation time.
- 4. More driver classes do not necessary results in a longer computation time.
- 5. Computation time can be greatly reduced by using small sample sizes in the early stage of solution processes.

A discussion on the implementation issues of the proposed model and the SAM to realistic networks is also provided.

This study opens many research directions. Firstly, to use the proposed model for real applications, we need to calibrate the parameters and validate the models. This can be one future research direction. Secondly, the effects of physical queues have been captured in the proposed model. Yet, their impacts on travel time uncertainty have not been investigated. Indeed, it is interesting and important to investigate if there is a propagation of travel time uncertainty from one route to another, as the travel time uncertainty affects the route choice. This investigation should be carried out in the future. Thirdly, this paper does not consider departure time choice and stochastic travel demand. Including these two components into the model framework should be an interesting research direction. Fourthly, multiple-pipe and variational theories as well as the modeling theories for lane-changing behaviour and moving bottlenecks (Daganzo and Laval, 2005; Daganzo, 2006; Leclercq, 2007) have not been advanced to deal with large-scale network assignment applications which requires modeling the traffic interaction at junctions and reasonably short computation time. Advancing these theories can be another future research direction. Fifthly, antithetic sampling and Latin hypercube sampling techniques and single point approximation techniques have not been

incorporated in the solution method. A comprehensive evaluation and comparison of using these techniques for solving the proposed problem can be done in the future. Finally, the proposed model can be applied to transport planning and management by extending it to a bilevel model and solved the resultant model by heuristics such as tabu search (e.g., Fan and Machemehl, 2008; Mohan Rao and Shyju, 2010), genetic algorithms (Adeli and Kumar, 1995a,b; Sarma and Adeli, 2000a,b, 2001, 2002; Kim and Adeli, 2001; Teklu *et al.*, 2007; Jiang and Adeli, 2008; Cheng and Yen, 2009; Kang *et al.*, 2009; Ng *et al.*, 2009; Zeferino, *et al.*, 2009; Lee and Wei, 2010; Al-Bazi and Dawood, 2010), and ant colony heuristics (e.g., Vitins and Axhausen, 2009).

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