ON THE RANK OF HIGHER INCLUSION MATRICES

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ABSTRACT. Let $r \geq s \geq 0$ be integers and G be an r-graph. The higher inclusion matrix $M_s^r(G)$ is a $\{0, 1\}$ -matrix with rows indexed by the edges of G and columns indexed by the subsets of V(G) of size s: the entry corresponding to an edge e and a subset S is 1 if $S \subseteq e$ and 0 otherwise. Following a question of Frankl and Tokushige and a result of Keevash, we define the rank-extremal function $\operatorname{rex}(n, t, r, s)$ as the maximum number of edges of an r-graph G having rk $M_s^r(G) \leq {n \choose s} - t$. For t at most linear in n we determine this function as well as the extremal r-graphs. The special case t = 1 answers a question of Keevash.

1. INTRODUCTION

Let $n \ge 1$ and suppose \mathcal{F} is a collection of k-subsets of [n]. For any $p \ge 1$, we can define the lower p-shadow $\partial_l^p \mathcal{F}$ of \mathcal{F} as the set of all (k-p)-subsets of [n] which are contained in at least one element of \mathcal{F} . If p = 1, we may drop the superscript. A fundamental result in extremal combinatorics, the Kruskal-Katona theorem, gives a sharp lower bound for the size of $\partial_l^p \mathcal{F}$. In order to state the theorem, we note that for positive integers m and k there are always unique integers $m_k > m_{k-1} > \ldots > m_j \ge j > 0$ such that $m = \binom{m_k}{k} + \binom{m_{k-1}}{k-1} + \ldots + \binom{m_j}{j}$. We further make the convention that $\binom{x}{y} = 0$ whenever y < 0.

Theorem 1 ([8], [12]). Let $k \ge 1, p \ge 1$ and $m \ge 1$. For every $\mathcal{F} \subseteq {\binom{[n]}{k}}$ with $m = |\mathcal{F}|$ we have

$$|\partial_l^p \mathcal{F}| \ge \binom{m_k}{k-p} + \binom{m_{k-1}}{k-1-p} + \dots + \binom{m_j}{j-p}.$$
(1)

The inequality is best possible for every k, p and $m \leq \binom{n}{k}$. Furthermore, if $m = \binom{m_k}{k}$, then equality holds in (1) if and only if $\mathcal{F} \simeq \binom{[m_k]}{k}$.

Theorem 1 also provides a sharp lower bound for the size of the upper *p*-shadow $\partial_u^p \mathcal{F}$, which is the set of all (k + p)-subsets of [n] which contain at least one element of \mathcal{F} (indeed, $\partial_u^p \mathcal{F} = (\partial_l^p \mathcal{F}^c)^c$, where \mathcal{A}^c is the family obtained by taking the complements of all the sets in \mathcal{A}). For later use, we shall denote

$$K(n,m,k,p) = \min\left\{ |\partial_u^p \mathcal{F}| : \mathcal{F} \subseteq {\binom{[n]}{k}}, \ |\mathcal{F}| = m \right\}.$$

Equivalently, as $|\partial_u^p \mathcal{F}| = |\partial_l^p \mathcal{F}^c|$, K(n, m, k, p) denotes the minimum size of $|\partial_l^p \mathcal{F}|$, taken over all collections $\mathcal{F} \subseteq {\binom{[n]}{n-k}}$ of size m. By the Kruskal-Katona theorem, if $m = {\binom{m_{n-k}}{n-k}} + {\binom{m_{n-k}}{n-k}} = {\binom{m_{n-k}}{n$

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 $\binom{m_{n-k-1}}{n-k-1} + \ldots + \binom{m_j}{j}$ is the unique decomposition of m, then

$$K(n, m, k, p) = \binom{m_{n-k}}{n-k-p} + \binom{m_{n-k-1}}{n-k-1-p} + \dots + \binom{m_j}{j-p}.$$

A weaker but simpler version of the Kruskal-Katona theorem is due to Lovász ([15], Exercise 13.31(b)): if $|\mathcal{F}| = {x \choose k}$, for some real number x, then $|\partial_l^p \mathcal{F}| \ge {x \choose k-p}$, with equality if and only if x is an integer and $\mathcal{F} \simeq {[x] \choose k}$. Several algebraic versions of Lovász's result exist [2], [9]; however, no analogue of the full Kruskal-Katona theorem in linear algebra has so far been obtained.

In this paper we will be mainly concerned with the algebraic generalization due to Keevash [9]. To state it we will first introduce several definitions.

A hypergraph G is an ordered pair of sets (V, E), where $E \subseteq 2^V$, the elements of V are called vertices, and the elements of E are called edges. If $r \ge 1$ and all elements of E have the same size r, we say G is an r-uniform hypergraph, or simply an r-graph. Hence 2-graphs correspond to the usual notion of undirected graphs. We let v(G) denote the number of vertices, or the order of G, and |G| the number of edges, or the size of G. We denote by K_n^r the complete r-graph on n vertices and vertex set [n]. For a subset $E' \subseteq E$ of the edges of an r-graph G = (V, E) we shall denote by G - E the hypergraph with vertex set V and edge set $E \setminus E'$. Furthermore if $F \subseteq {[n] \choose r}$ we shall sometimes identify F with the r-uniform hypergraph ([n], F). Finally, if G = (V, E) is any r-graph we let \overline{G} denote its complement, i.e. the hypergraph $\left(V, {V \choose r} \setminus E\right)$.

Now let G be an r-graph and $s \leq r$. The higher inclusion matrix $M_s^r(G)$ is a $\{0, 1\}$ -matrix with rows indexed by the edges of G and columns indexed by the subsets of V(G) of size s: the entry corresponding to an edge e and a subset S is 1 if $S \subseteq e$ and 0 otherwise. Thus $M_1^2(G)$ is the usual incidence matrix of a graph G.

It is an open problem of Frankl and Tokushige [4] to determine the minimum rank of $M_s^r(G)$ in terms of |G|. A theorem of Gottlieb [7] shows that the matrix $M_s^r(K_n^r)$ has full rank.

Theorem 2 (Gottlieb, [7]). For every $n \ge r \ge s \ge 0$ we have

$$\operatorname{rk} M_s^r(K_n^r) = \min\left\{ \binom{n}{r}, \binom{n}{s} \right\}$$

Here and in the following the rank is considered only over the reals (in fact, we can take any field of characteristic 0). Now the generalization of Lovász's theorem due to Keevash is the following.

Theorem 3 ([9]). For every $r \ge s \ge 0$ there is a number $n_{r,s}$ so that if G is an r-graph with $|G| = \binom{x}{r} \ge n_{r,s}$ then $\operatorname{rk} M_s^r(G) \ge \binom{x}{s}$. Furthermore, if r > s > 0 then equality holds if and only if x is an integer and $G \simeq K_x^r$.

Theorem 3 implies Lovász's result, as the rank of $M_s^r(G)$ is at most the number of non-zero columns, which is the size of the lower (r - s)-shadow of E(G). An important step in the proof of Theorem 3 was provided by the following lemma.

Lemma 4 ([9]). Suppose $n \ge 2r+s$. If $F \subseteq {\binom{[n]}{r}}$ with $|F| < {\binom{r}{s}}^{-1} {\binom{n}{r-s}}$ then $\operatorname{rk} M_s^r(K_n^r - F) = {\binom{n}{s}}$.

Keevash further asked ([10], Question 2 (iii)) whether Lemma 4 remains true under the assumption $|F| < \binom{n-s}{r-s}$, at least for large n. This would be best possible, as removing all edges of K_n^r containing some fixed *s*-set creates a 0-column and hence reduces the rank by at least 1.

Now let $r > s \ge 1, n \ge r + s$ and $0 \le t \le {n \choose s}$. In view of the above and the question of Frankl and Tokushige [4], we define the *rank-extremal function*

$$\operatorname{rex}(n,t,r,s) := \max\left\{ |G| : G \text{ is an } r \text{-graph on } [n] \text{ and } \operatorname{rk} M_s^r(G) \le \binom{n}{s} - t \right\}$$

This notion is analogous to the notion of Turán number of hypergraphs (see [11] for a survey), where the maximum number of edges is sought when the clique number is bounded. Here, instead, the rank of the higher-inclusion matrix is bounded.

In the current paper we investigate this function for small t and further determine the rankextremal r-graphs. As in the case of t = 1, there is a natural construction which provides a lower bound for given t: fixing t subsets of [n] of size s and removing all edges of K_n^r containing at least one of them yields an r-graph G with $\operatorname{rk} M_s^r(G) \leq {n \choose s} - t$, as $M_s^r(G)$ has at least t zero columns. Minimizing over all choices of the t s-subsets, we see that

$$\operatorname{rex}(n,t,r,s) \ge \binom{n}{r} - K(n,t,s,r-s),$$

where recall that K(n, t, s, r-s) is the smallest size of the upper (r-s)-shadow of a collection of s-subsets of [n] of size t, provided by the Kruskal-Katona theorem.

The content of our main theorem is that equality holds here for t at most linear in n.

2. New Results

We call a collection of s-sets an s-star configuration if their common intersection has size at least s - 1. Clearly for any $t \leq n - s + 1$, there exists a unique (up to isomorphism) s-star configuration $S_{n,t,s}$ of size t on ground set [n]. It is easy to see that $|\partial_u^{r-s}S_{n,t,s}| = \binom{n-s+1}{r-s+1} - \binom{n-s+1-t}{r-s+1}$. Removing all edges of K_n^r containing some element of $S_{n,t,s}$ gives an r-graph G(n,t,r,s) with $\operatorname{rk} M_s^r(G(n,t,r,s)) \leq \binom{n}{s} - t$, as there are at least t zero columns. Then we have

$$\operatorname{rex}(n,t,r,s) \ge \binom{n}{r} - K(n,t,s,r-s) \ge e(G(n,t,r,s)) = \binom{n}{r} - \binom{n-s+1}{r-s+1} + \binom{n-s+1-t}{r-s+1} + \binom{n-s+1-t}{r-s+1}$$

Our main result is the following.

Theorem 5. Let $r > s \ge 1$. Then there exist positive constants $c_0 := c_0(r, s)$ and $n_0 := n_0(r, s)$ such that the following holds. Let $n \ge n_0$ and $1 \le t \le c_0 n$ be integers. Then

$$\operatorname{rex}(n,t,r,s) = \binom{n}{r} - \binom{n-s+1}{r-s+1} + \binom{n-s+1-t}{r-s+1}.$$
(3)

Furthermore, G(n, t, r, s) is the unique rank-extremal r-graph up to isomorphism.

In particular, equality holds everywhere in (2), and so $K(n, t, s, r-s) = \binom{n-s+1}{r-s+1} - \binom{n-s+1-t}{r-s+1}$. Now note that for any $t \leq n-s$ we have

$$t = \binom{n-s}{n-s} + \dots + \binom{n-s-t+1}{n-s-t+1},$$
(4)

and so by unicity, the right-hand side of (1) (for k := n - s, p := r - s and m := t) is exactly $\binom{n-s}{n-r} + \ldots + \binom{n-s-t+1}{n-r-t+1} = \binom{n-s+1}{r-s+1} - \binom{n-s-t+1}{r-s+1}$. Thus we have obtained a linear-algebraic proof of Theorem 1 for the lower (r-s)-shadow of a collection of (n-s)-sets of size at most $c_0 n$.

By taking t := 1 in Theorem 5, we obtain $rex(n, 1, r, s) = {n \choose r} - {n-s \choose r-s}$ which gives a positive answer to Keevash's question.

The main step in the proof of Theorem 5, which might be interesting in its own right, is an extension of Gottlieb's Theorem and concerns the robustness or resilience of the rank. It states that if a limited number of somewhat uniformly distributed r-sets are removed from the complete r-graph then the rank of the s-inclusion matrix does not decrease.

Theorem 6. For all $r \ge s \ge 0$ there exist positive constants $\varepsilon := \varepsilon(r, s), \alpha := \alpha(r, s)$ and $n_1 := n_1(r, s)$ such that the following holds. If $F \subseteq \binom{[n]}{r}$ is a family of at most $\varepsilon n \binom{n-s}{r-s}$ r-sets on $n \ge n_1$ vertices, with the property that every s-set $S \in \binom{[n]}{s}$ is contained in less than $\alpha \binom{n-s}{r-s}$ elements of F, then $\operatorname{rk} M_s^r(K_n^r - F) = \binom{n}{s}$.

Theorem 5 implies that for t up to some small constant times n, we have $\binom{n}{r}$ -rex(n, t, r, s) = K(n, t, s, r - s). One may wonder what is the maximum value t_{max} up to which this equality holds, provided n is large enough. Let us define

$$t_{\max}(n,r,s) := \max\left\{1 \le t \le \binom{n}{s}: \text{ for any } t' \le t, \binom{n}{r} - \operatorname{rex}(n,t',r,s) = K(n,t',s,r-s)\right\}$$

Our next theorem shows that our bound on t in Theorem 5 is best possible up to a constant factor.

Theorem 7. For arbitrary integers $r > s \ge 1$ there exists a positive constant $n'_0 := n'_0(r, s)$ such that for every $n \ge n'_0$, $t_{max}(n, r, s) < n - r - 1$.

The rest of this paper is organized as follows. In Section 3 we prove Theorem 6. In Section 4 we prove the main theorem. Finally, in Section 5 we prove Theorem 7.

Remark. After completion of this work we were informed by Harout Aydinian that Ahlswede, Aydinian and Khachatrian considered before the problem of determining rex(n, t, r, 1) and solved it in [1] for all values of n, t and r. Their proof technique is very different from ours. In [1] it is also mentioned that the function rex(n, 1, r, 1) was first studied by Longstaff [14], with the complete determination of its value being made by Odlyzko [16].

3. Resilience of the rank

The main goal of this section is to prove Theorem 6.

For an arbitrary r-graph H and vertex $x \in V(H)$ we define two derived hypergraphs. We shall denote by H/x the (r-1)-graph on vertex set $V(H) - \{x\}$ with edge set $\{A \setminus \{x\} : A \in E(H), x \in A\}$ and denote by H-x the r-graph on vertex set $V(H) - \{x\}$ with edge set $\{A \in E(H), x \notin A\}$. Observe that for every $x \in V(H)$, |H| = |H/x| + |H-x|.

Let G be any r-graph with vertex set V. For any s-subset $S \subseteq V$ and any r-edge $e \in E(G)$, we shall denote by $c_S(G)$ the column corresponding to S in $M_s^r(G)$, and by $c_S(G)_e$ the entry of $M_s^r(G)$ corresponding to the edge e and the s-subset S. A sequence $\{\alpha_S\}_{S \in \binom{V}{s}}$ of real numbers is called a *dependence sequence* for G if

$$\sum_{S \in \binom{V}{s}} \alpha_S c_S(G)_e = 0, \tag{5}$$

for every $e \in E(G)$. If not all the coefficients α_S are 0, then the columns of $M_s^r(G)$ are linearly dependent and we call the sequence *non-trivial*. For a dependence sequence sequence we construct the associated s-graph G' on vertex set V(G') = V(G) with edge set $E(G') = \{S \in \binom{V}{s} : \alpha_S \neq 0\}$.

For later use we collect here four simple, but useful observations about an arbitrary dependence sequence $\{\alpha_S\}_{S \in \binom{V}{s}}$ and its associated *s*-graph *G'*.

Observation 8. If $V' \subseteq V(G)$ and H := G[V'] is the induced r-graph on V' then $\{\alpha_S\}_{S \in \binom{V'}{s}}$ is a dependence sequence for H.

Proof. The proof follows by noting that the submatrix of $M_s^r(G)$ indexed by the rows from $E(H) = \binom{V'}{r} \cap E(G)$ and the columns from $\binom{V}{s} \setminus \binom{V'}{s}$ has only zero entries, and hence when (5) is applied to G and $e \in E(H)$, all terms of the form $\alpha_S c_S(G)_e$ with $S \notin \binom{V'}{s}$ vanish. That is, (5) holds for H and e as well.

An r-subset $R \subseteq V(G')$ is called a 1-*clique* if the induced subgraph G'[R] contains exactly one edge of G'.

Observation 9. No edge of G is a 1-clique.

Proof. Indeed, if $R \in E(G)$ is an *r*-set containing exactly one edge S' of G', then $\alpha_S = 0$ for any $S \in \binom{R}{s} - \{S'\}$. Applying (5) to G and R we obtain $\alpha_{S'} = \alpha_{S'}c_{S'}(G)_R = 0$, so S' is not an edge of G', a contradiction.

Observation 10. If $x \in V(G)$ and E(G'-x) is empty, then the sequence $\beta_{S'} := \alpha_{\{x\}\cup S'}, S' \in \binom{V(G/x)}{s-1}$ is a dependence sequence for G/x.

Proof. Note that the condition of E(G' - x) being empty implies $\alpha_S = 0$, for any s-set S not containing x. Then for any $e \in E(G/x)$ we have

$$\sum_{S' \in \binom{V(G/x)}{s-1}} \beta_{S'} c_{S'}(G/x)_e = \sum_{\substack{S' \in \binom{V(G/x)}{s-1}}} \alpha_{\{x\} \cup S'} c_{\{x\} \cup S'}(G)_{\{x\} \cup e}, \text{ by definition,}$$
$$= \sum_{S \in \binom{V(G)}{s}, S \ni x} \alpha_S c_S(G)_{\{x\} \cup e},$$
$$= \sum_{S \in \binom{V(G)}{s}} \alpha_S c_S(G)_{\{x\} \cup e}, \text{ as } \alpha_S = 0 \text{ when } x \notin S,$$
$$= 0, \text{ by (5) applied to } G \text{ and } \{x\} \cup e.$$

So (5) holds for G/x and $\{\beta_{S'}\}$ as well.

Our last observation extends the notion of a 1-clique. A k-semistar with center x and set of leaves L is an induced subgraph of G' on the (k + 1)-set of vertices $\{x\} \cup L$, such that L is an independent set in G' of size k and at least one of the s-sets $\{x\} \cup S', S' \in \binom{L}{s-1}$, is present in E(G').

Observation 11. For any (r + s - 2)-semistar Z in G', with center x and set of leaves L, at least one of the r-edges $\{x\} \cup R', R' \in \binom{L}{r-1}$, must be in F.

Proof. If s = 1, then V(Z) is a 1-clique, because $\{x\} \in E(G')$ and $V(Z) \setminus \{x\}$ is independent in G'. Hence V(Z) is in F by Observation 9 and the claim holds.

So we may assume for a contradiction that $s \ge 2$ and every edge $\{x\} \cup R', R' \in {L \choose r-1}$, is in G. Consider the (r-1)-graph H induced by G/x on L. Then $H \simeq K_{r+s-2}^{r-1}$ is complete. Hence $M_{s-1}^{r-1}(H)$ has full rank by Gottlieb's Theorem.

On the other hand, we can first apply Observation 8 to G and V(Z) followed by Observation 10 to x and G[V(Z)], and conclude that the sequence defined by $\beta_{S'} := \alpha_{\{x\}\cup S'}, S' \in \binom{L}{s-1}$, is a dependence sequence for H = G[V(Z)]/x. Since Z is a semistar, at least one of the s-sets $\{x\} \cup S', S' \in \binom{L}{s-1}$, must be present in E(G'), hence $\beta_{S'} \neq 0$ for some $S' \in \binom{L}{s-1}$. In conclusion, the dependence sequence is non-trivial and this contradicts the columns of $M_{s-1}^{r-1}(H)$ being linearly independent.

Before we turn to the proof of Theorem 6, we show a lemma ensuring the existence of a large independent set in G' provided the rank drops below $\binom{n}{s}$ upon the deletion of a not so large family F of edges.

Lemma 12. Let $F \subseteq E(K_n^r)$ be such that $\operatorname{rk} M_s^r(K_n^r - F) < \binom{n}{s}$ and $(n - r - s)\left(\binom{n}{r} - |F|\right) \ge |F|\binom{r}{s}\binom{n-r+s}{s}$, then any associated s-graph G' has an independent set of size at least

$$n - \frac{|F|\binom{r}{s}\binom{n-r+s}{s}}{\binom{n}{r} - |F|}.$$

Proof. Let $G := K_n^r - F$. To begin with, we find an edge $R \in E(G)$ intersecting not too many r-sets $e \in F$ in at least r - s elements. We estimate the number of such pairs (R, e) (where $R \in E(G)$, $e \in F$ with $|e \cap R| \ge r - s$) from above by $|F|\binom{r}{s}\binom{n-r+s}{s}$. Indeed, we can choose $e \in F$ in |F| ways, then specify a subset of e of size r - s in $\binom{r}{r-s}$ ways, and extend it to an element $R \in \binom{V(G)}{r} \setminus F$ in at most $\binom{n-r+s}{s}$ ways. By averaging, we find a set $R \in E(G)$ for which there are at most

$$\frac{|F|\binom{r}{s}\binom{n-r+s}{s}}{\binom{n}{r}-|F|}$$

r-sets $e \in F$ with $|e \cap R| \ge r - s$.

We now fix one such edge $R \in E(G)$, an arbitrary non-trivial dependence sequence $\{\alpha_S\}_{S \in \binom{V}{s}}$, and deduce that the associated s-graph G' has a large independent set. For every edge $e \in F$ with $|e \cap R| \ge r-s$, we choose an arbitrary vertex $v_e \in e \setminus R$. This is possible, as e is also an r-subset and $e \ne R$. Now let $A := V(G') \setminus \{v_e : e \in F \text{ and } |e \cap R| \ge r-s\}$. Note that $R \subseteq A$. By the assumption of the lemma, we also have that

$$|A| \ge n - \frac{|F|\binom{r}{s}\binom{n-r+s}{s}}{\binom{n}{r} - |F|} \ge r+s.$$

We claim that A is independent in G'. Suppose for a contradiction that an s-set $Q \subseteq A$ is an edge in G'. Form a set $V' \subseteq A$ on r + s vertices containing R and Q. This is possible, as $|R \cup Q| \leq r + s$ and there are enough vertices in A. By construction, no r-set $e \in F$ is contained in V', as otherwise $|e \cap R| \geq |e| - |V' \setminus R| = r - s$ hence v_e is defined and $v_e \in e \subseteq V' \subseteq A$, a contradiction. Therefore $G = K_n^r - F$ induces on V' a complete r-graph $\binom{V'}{r}$, which by Gottlieb's theorem must have an inclusion matrix $M_s^r(G[V'])$ of full rank. In particular, every dependence sequence for G[V'] must be trivial.

However, by Observation 8, the coefficients $\alpha_T, T \in \binom{V'}{s}$, form a dependence sequence for the *r*-graph G[V'], which is non-trivial, as $Q \subseteq V'$ and $\alpha_Q \neq 0$. This is a contradiction, completing the proof of the lemma.

We are now ready to prove Theorem 6.

Proof of Theorem 6. Note that for r = s or s = 0, we may simply choose $\varepsilon := 1, n_1 := r$ and $\alpha := 1$. Hence we may assume r > s > 0.

We first set some positive constants, depending on r and s:

$$\varepsilon < \min\left\{\frac{(r-s)!}{2r!}, \frac{1}{4{r \choose s}^2}, \frac{1}{r! 2^{2r-s}}\right\},$$
(6)

$$\delta = 2\varepsilon \binom{r}{s}^2,\tag{7}$$

$$\alpha < \left(\frac{1-2\delta}{2(r-s)}\right)^{r-s},\tag{8}$$

$$n_1 = \left\lceil \frac{4r}{\frac{1}{2} - \delta} \right\rceil. \tag{9}$$

Note that, by (6), $\delta < 1/2$, hence α and n_1 can indeed be chosen to be positive.

Suppose for a contradiction that Theorem 6 is false. Then there exists a family $F \subset E(K_n^r)$ of r-subsets on $n \ge n_1$ vertices, with $|F| < \varepsilon n \binom{n-s}{r-s}$ and every s-set $S \subseteq V(G)$ is contained in less than $\alpha \binom{n-s}{r-s}$ members of F, such that for $G := K_n^r - F$ we have $\operatorname{rk} M_s^r(G) < \binom{n}{s}$. Let $\{\alpha_S\}_{S \in \binom{V}{s}}$ be a non-trivial dependence sequence for G and let G' be the associated s-graph. Note that E(G') is non-empty, as the sequence is non-trivial. Since $s \ge 1$ and $\varepsilon < \frac{(r-s)!}{2r!}$ we have that

$$\binom{n}{r} > 2\varepsilon n \binom{n-s}{r-s} > 2|F|.$$

Then

$$\frac{|F|\binom{r}{s}\binom{n-r+s}{s}}{\binom{n}{r}-|F|} \le 2\varepsilon n\binom{n-s}{r-s}\binom{r}{r-s}\binom{n-r+s}{s}\binom{n}{r}^{-1} \le 2\varepsilon \binom{r}{s}^2 n = \delta n,$$

so by Lemma 12, G' contains an independent set A of maximum size at least $(1-\delta)n$. Define $B := V(G') \setminus A$. Then B is non-empty, because $\{\alpha_S\}_{S \in {V(G)}}$ is non-trivial.

In the following we will reach the desired contradiction by showing that the size of F is larger than assumed.

For $x \in B$ we call $S' \in {A \choose s-1}$ an x-leaf if $\{x\} \cup S' \in E(G')$. Let $d_A(x)$ be the number of x-leaves. We also let $\mu_A(x)$ be the number of vertices in A covered by the x-leaves. Note that for s > 1 we have $d_A(x) \ge \frac{\mu_A(x)}{s-1} \ge 1$, since otherwise $A \cup \{x\}$ would be a larger independent set, contradicting the maximality of A.

First assume that for some $x \in B$ we have $\mu_A(x) \leq n/2$. Fix this x and let S_x be any x-leaf. Note that for s = 1, we have $S_x = \emptyset$.

If T is any (r-s)-subset of A not intersecting any x-leaf then $\{x\} \cup S_x \cup T$ forms a 1-clique, because $S_x \cup T \subseteq A$ is an independent set. By Observation 9, this shows that there are at least $\binom{|A|-\mu_A(x)}{r-s}$ elements of F containing the s-set $\{x\} \cup S_x$. This is a contradiction to the choice of F, because

$$\binom{|A| - \mu_A(x)}{r - s} \ge \left(\frac{|A| - \mu_A(x)}{r - s}\right)^{r - s} \ge \left(\frac{1 - 2\delta}{2(r - s)}\right)^{r - s} n^{r - s} > \alpha n^{r - s} > \alpha \binom{n - s}{r - s}.$$

Here we use the fact that $|A| - \mu_A(x) \ge (\frac{1}{2} - \delta)n_1 \ge 4r \ge r - s$ by (9). Hence for any $x \in B$ we have that $\mu_A(x) > n/2$. In particular, s > 1.

We double-count the triples (x, S, R) where $x \in B$, S is an x-leaf, and $R \in {A \choose r-1}$ with $R \cap S = \emptyset$. Note that for any such triple, $\{x\} \cup S \cup R$ is an (r + s - 2)-semistar with center x, and their number is $\sum_{x \in B} d_A(x) \binom{|A| - s + 1}{r - 1}$. On the other hand, by Observation 11, for any (r + s - 2)-star Z with center x we can fix an r-set R_Z such that $x \in R_Z$ and $R_Z \in \binom{V(Z)}{r} \cap F$. Now, for any such r-set $R^* \in F$ the number of those triples (x, S, R) for which $R_{\{x\}\cup S\cup R} = R^*$ is at most $\binom{|A|-r+1}{s-1}$. Note that given R^* , the center x is determined uniquely, provided $|R^* \cap B| = 1$ (otherwise there is no appropriate triple at all).

These imply the following lower bound for the size of F:

$$\begin{aligned} |F| &\geq \sum_{x \in B} d_A(x) \binom{|A| - s + 1}{r - 1} \binom{|A| - r + 1}{s - 1}^{-1} \\ &\geq \left(\sum_{x \in B} d_A(x)\right) \frac{1}{(r - 1)!} \left(|A| - s - r + 3\right)^{r - 1} \frac{(s - 1)!}{|A|^{s - 1}} \\ &\geq \frac{(s - 1)!}{(r - 1)! 2^{r - 1}} |A|^{r - s} \left(\sum_{x \in B} d_A(x)\right), \text{ as } |A| \geq 2(r + s - 3), \\ &\geq \frac{(s - 1)!}{(r - 1)! 2^{r - 1}} |A|^{r - s} \left(\sum_{x \in B} \frac{\mu_A(x)}{s - 1}\right) \\ &\geq \frac{(s - 2)!}{(r - 1)! 2^{2r - s}} n^{r - s + 1}, \text{ because } \mu_A(x) \geq \frac{n}{2} \text{ and hence } |A| \geq n/2. \\ &> \varepsilon n^{r - s + 1}, \text{ by } (6) \\ &\geq \varepsilon n \binom{n - s}{r - s} \end{aligned}$$

This contradicts the assumption on the size of F and finishes the proof of Theorem 6. \Box

Remark 13. The argument above can be adapted to show that in the case t = 1 of Theorem 5 n_0 can be taken 2^{5r} .

4. Proof of the main result

For any $r \ge s \ge 1$ and $0 \le t \le n$ we define

$$N(n,t,r,s) := \binom{n-s+1}{r-s+1} - \binom{n-s+1-t}{r-s+1}.$$
(10)

Then for any s > 1 it holds that

$$N(n-1,t,r-1,s-1) = N(n,t,r,s).$$
(11)

Further note that

$$N(n,t,r,s) = \sum_{i=1}^{t} \binom{n-s+1-i}{r-s}.$$
(12)

We shall need the following two technical results about the behaviour of $N(\cdot)$.

Lemma 14. Let $r > s \ge 1$ and $\alpha \in (0,1)$. Then there exist positive constants $\gamma_0 < \gamma_1 < 1$ and n_2 such that the following holds. For any $n \ge n_2, t \le \gamma_0 n$ and $p \ge \gamma_1 n$ we have that

$$\alpha N(n, p, r, s) > N(n, t, r, s).$$
(13)

Proof. By definition of $N(\cdot)$ it is enough to prove the following

$$(1-\alpha)\binom{n-s+1}{r-s+1} < \binom{n-s+1-t}{r-s+1} - \binom{n-s+1-p}{r-s+1}.$$

The left hand side is at most

$$(1-\alpha)\frac{n^{r-s+1}}{(r-s+1)!},$$

while the right hand side is at least

$$\frac{1}{(r-s+1)!} \left((n-s+1-t-(r-s))^{r-s+1} - ((1-\gamma_1)n)^{r-s+1} \right).$$

Provided $n_2 \geq \frac{r-1}{\gamma_0}$, this is at least

$$\frac{n^{r-s+1}}{(r-s+1)!} \left((1-2\gamma_0)^{r-s+1} - (1-\gamma_1)^{r-s+1} \right).$$

Hence it is enough if $1 - \alpha < (1 - 2\gamma_0)^{r-s+1} - (1 - \gamma_1)^{r-s+1}$, which clearly has a solution $(\gamma_0, \gamma_1) \in (0, 1)^2$ as desired.

Lemma 15. Let $r > s \ge 1$ and $\alpha \in (0,1)$. Then there exist positive constants $\sigma < 1$ and n_3 such that the following holds. For any $n \ge n_3$ and $t \le \sigma n$ it holds that

$$N(n,t,r,s) - \alpha \binom{n-s}{r-s} < N(n-1,t,r,s).$$

$$\tag{14}$$

Proof. It is equivalent to show

$$\binom{n-s-t}{r-s} > (1-\alpha)\binom{n-s}{r-s},$$

which holds for $n_3 > \frac{r-1}{\sigma}$ and σ such that $(1-2\sigma)^{r-s} > 1-\alpha$.

We shall frequently use the following inequality, whose simple proof forms part of Lemma 12 in [9].

Lemma 16. Suppose G is an r-graph, x is a vertex of G and $1 \le s \le r-1$. Then

$$\operatorname{rk} M_{s}^{r}(G) \ge \operatorname{rk} M_{s}^{r}(G-x) + \operatorname{rk} M_{s-1}^{r-1}(G/x).$$
 (15)

Note that in Lemma 16, if G/x has no edges, then G - x has the same edge set as G and $\operatorname{rk} M_s^r(G - x) = \operatorname{rk} M_s^r(G)$. We will always assume that for an *r*-graph G with no edges, $\operatorname{rk} M_s^r(G) = 0$.

We will now prove the following equivalent version of Theorem 5.

Theorem 17. For every $r > s \ge 1$, there exist positive constants $c_0 := c_0(r, s)$ and $n_0 := n_0(r, s)$ such that the following holds. For arbitrary integers n and t with $n \ge n_0$ and $1 \le t \le c_0 n$ and r-graph F on n vertices, we have that

- (i) if |F| < N(n, t, r, s), then $\operatorname{rk} M_s^r(\overline{F}) > {n \choose s} t$
- (ii) |F| = N(n, t, r, s) and $\operatorname{rk} M_s^r(\overline{F}) \leq {n \choose s} t$ if and only if F is the upper (r-s)-shadow of an s-star configuration of size t.

Remark. Observe that part (i) implies that if F is the upper (r-s)-shadow of an s-star configuration of size t, then $\operatorname{rk} M_s^r(K_n^r - F) = \binom{n}{s} - t$. Indeed, |F| = N(n, t, r, s), hence deleting any one set from F will make the rank strictly larger than $\binom{n}{s} - t$, but the addition of a single row can not increase the rank by more than one.

Remark. Note that Theorem 17 indeed implies Theorem 5 with the same constants c_0 and n_0 . Let $n \ge n_0, 1 \le t \le c_0 n$ and G be any r-graph on [n] with $\operatorname{rk} M_s^r(G) \le {\binom{n}{s}} - t$. We define $F = \overline{G}$. By part (i) of Theorem 17, we must then have $|F| \ge N(n, t, r, s)$, so $\operatorname{rex}(n, t, r, s) \le {\binom{n}{r}} - N(n, t, r, s)$. Part (ii) shows that $\operatorname{rex}(n, t, r, s) = {\binom{n}{r}} - N(n, t, r, s)$ and the unique family G with $e(G) = \operatorname{rex}(n, t, r, s)$ and $\operatorname{rk} M_s^r(G) \le {\binom{n}{s}} - t$ is the complement of the upper (r-s)-shadow of an s-star configuration of size t.

Proof of Theorem 17. Recall first that we have already seen the "if" part of (ii) in the introduction: The upper (r-s)-shadow of an s-star configuration of size t has size N(n, t, r, s) and for its complement G = G(n, t, r, s) (cf. the discussion before Theorem 5) rk $M_s^r(G) \leq {n \choose s} - t$, since the matrix contains t 0-columns.

For (i) and the "only if" part of (ii) we apply double induction, first on s and then on a parameter measuring how far the hypergraph is from Theorem 6 being applicable. For the second induction we need to prove a somewhat more technical statement requiring a bit of preparation.

Let $\alpha := \alpha(r, s)$ be the constant given by Theorem 6; we may assume that $\alpha < 1$. For every r-graph F' = (V, E), we define $\Lambda(F', s)$ as the smallest integer $\ell \ge 0$, such that there exists a (possibly empty) sequence $v_1, \ldots, v_\ell \in V$ of the vertices, such that for every i = $1, \ldots, \ell v_i$ is of maximum degree in $F_{i-1} = F - \{v_1, \ldots, v_{i-1}\}$ and the maximum degree of $F_\ell = F - \{v_1, \ldots, v_\ell\}$ is at most $\alpha(\frac{|V(F_\ell)| - s}{r - s})$. This function is finite, because the maximum degree of K_{r-1}^r is $0 = \alpha\binom{r-1-s}{r-s}$. By the definition of Λ , if $\Lambda(F', s) > 0$ then there exists a vertex $x \in V(F')$ of maximum degree such that

$$\Lambda(F', s) = \Lambda(F' - x, s) + 1.$$

The rôle of $\Lambda(F', s)$ will become clear in Claim 1 below, where, intuitively, if zero, it will allow us to appeal to Theorem 6 immediately, and otherwise it measures "how far" we are from being able to do so.

Let $r > s \ge 1$ be integers. The base case s = 1 will be similar to the induction step, so we will treat them parallel, but will always clearly distinguish which case we deal with. If $s \ge 2$ we assume that the induction hypothesis (i.e. Theorem 17) holds for all r' > s' with $1 \le s' < s$, in particular that the appropriate constants $n_0(r-1, s-1)$ and $c_0(r-1, s-1)$ exist.

For our technical claim we now define two positive constants n' := n'(r, s) and c' := c'(r, s)for every $r > s \ge 1$. In the definition we will need $\varepsilon = \varepsilon(r, s)$ and $n_1 = n_1(r, s)$, the two other constants (besides α) as asserted by Theorem 6, and $\sigma = \sigma(r, s, \alpha)$ and $n_3 = n_3(r, s, \alpha)$ the constants in Lemma 15. For s = 1, we set

$$n' = \max\{r + s + 1, n_1, n_3\},\$$

$$c' < \min\{\varepsilon, \sigma\}.$$

For s > 1, we set

$$n' = \max\{r + s + 1, n_1, n_3, n_0(r - 1, s - 1) + 1\},\$$

$$c' < \min\{\varepsilon, \sigma, \frac{1}{2}c_0(r - 1, s - 1)\}.$$

Recall that s is now fixed. The following claim is formulated using the constants defined above.

Claim 1. Let $t \ge 0$ be an integer and F be an r-graph on n vertices with $n - \Lambda(F, s) \ge \max\{n', \frac{t}{r'}\}$. Then

- (i) if |F| < N(n, t, r, s), then $\operatorname{rk} M_s^r(\overline{F}) > {n \choose s} t$.
- (ii) if |F| = N(n, t, r, s) and $\operatorname{rk} M_s^r(\overline{F}) \leq {\binom{n}{s}} t$ then F is the upper (r-s)-shadow of an s-star configuration of size t.

Let us discuss shortly how this claim implies the induction statement for s. Recall that $\gamma_0 = \gamma_0(r, s, \alpha) < \gamma_1 = \gamma_1(r, s, \alpha)$ and $n_2 = n_2(r, s, \alpha)$ are the constants given by Lemma 14. We set

$$n_0 = n_0(r, s) = \max\left\{\frac{n'}{1 - \gamma_1}, n_2\right\},\c_0 = c_0(r, s) = \min\{c'(1 - \gamma_1), \gamma_0\}.$$

Let $n \ge n_0, 1 \le t \le c_0 n$ and $F \subseteq {\binom{[n]}{r}}$ be from Theorem 17. Then $n - n' \ge \gamma_1 n$. We can also assume that $\Lambda(F, s) < \gamma_1 n$, as otherwise by the definition of Λ , (12), and by Lemma 14,

$$|F| > \alpha \sum_{i=1}^{\gamma_1 n} \binom{n-s+1-i}{r-s} = \alpha N(n,\gamma_1 n,r,s) > N(n,t,r,s).$$

Therefore $c'(n - \Lambda(F, s)) > c'(1 - \gamma_1)n \ge c_0(r, s)n \ge t$, so the assumptions of Claim 1 are satisfied, hence (i) and (ii) both hold for F.

Proof of Claim 1. If t = 0, then (i) is vacuously true, while (ii) holds because the \emptyset is the upper shadow of itself.

For $t \ge 1$ we prove Claim 1 by induction on $\Lambda = \Lambda(F, s)$. Let F be an r-graph on n vertices with $n - \Lambda(F, s) \ge \max\{n', \frac{t}{c'}\}$ and $|F| \le N(n, t, r, s)$. Set $G := \overline{F}$.

We first consider the base case $\Lambda(F,s) = 0$. Then the maximum degree of F is at most $\alpha\binom{n-s}{r-s}$. So every s-subset of V(G) is contained in at most $\alpha\binom{n-s}{r-s}$ elements of F. Also note that (12) implies

$$|F| \le N(n,t,r,s) \le t \binom{n-s}{r-s} \le c' n \binom{n-s}{r-s} \le \varepsilon n \binom{n-s}{r-s}.$$

As $n \ge n' \ge n_1$, the conditions of Theorem 6 are satisfied. Hence the matrix $M_s^r(G)$ has rank $\binom{n}{s}$, which proves both parts (i) and (ii) because $t \ge 1$.

Let us assume now that $\Lambda(F, s) \ge 1$ and the claim holds for hypergraphs F' with $\Lambda(F', s) < \Lambda(F, s)$.

Let $x \in [n]$ be a vertex of maximum degree in F, such that $\Lambda(F - x, s) = \Lambda(F, s) - 1$. Hence

$$|V(F-x)| - \Lambda(F-x,s) = n - \Lambda(F,s) \ge \max\left\{n', \frac{t}{c'}\right\}.$$

so Claim 1 can be applied to F - x and any integer $t', 0 \le t' \le t$ by the induction on Λ .

When s > 1 we can also apply the induction of Theorem 17 for s - 1 with r - 1 > s - 1and F/x, as $|V(F/x)| \ge n' - 1 \ge n_0(r - 1, s - 1)$ and for $n \ge 2$ by the definition of c' we have

$$t \le c'n < c_0(r-1,s-1)|V(F/x)| = c_0(r-1,s-1)(n-1)$$

We shall now distinguish two cases according to how large $\deg_F(x)$ is.

Case 1: $\binom{n-s}{r-s} \leq \deg_F(x) \leq N(n,t,r,s).$

We first deal with the case s = 1. Then $\deg_F(x) = \binom{n-1}{r-1}$, so G/x is just the empty graph, and hence $\operatorname{rk} M_0^{r-1}(G/x) = 0$. To estimate the rank of the matrix of $\operatorname{rk} M_1^r(G-x)$ we use Claim 1 for F-x and t-1. If |F| < N(n,t,r,1), then

$$|F - x| = |F| - \binom{n-1}{r-1} < N(n, t, r, 1) - \binom{n-1}{r-1} = N(n-1, t-1, r, 1),$$

and by part (i) $\operatorname{rk} M_1^r(G-x) > (n-1) - (t-1) = n-t$. Lemma 16 then implies $\operatorname{rk} M_1^r(G) = \operatorname{rk} M_1^r(G-x) > n-t$. For part (ii) let |F| = N(n,t,r,1) and $\operatorname{rk} M_1^r(G) \le n-t$. Then |F-x| = N(n-1,t-1,r,1) and $\operatorname{rk} M_1^r(G-x) \le n-t$. Hence by part (ii) of Claim 1, F-x is the upper (r-1)-shadow of a 1-star configuration \mathcal{S} of size t-1. Because all r-subsets of V(G) containing x are edges of F, F is the upper (r-1)-shadow of a 1-star configuration of size t (namely $\mathcal{S} \cup \{\{x\}\})$. This finishes the proof of Case 1 when s = 1.

Now assume s > 1. Let $p, 1 \le p \le t$ be the largest integer such that

$$|F/x| = \deg_F(x) \ge \sum_{i=1}^p \binom{n-s+1-i}{r-s} = N(n,p,r,s) = N(n-1,p,r-1,s-1).$$

Whenever p < t, we can apply part (i) of Theorem 17 to s - 1 < r - 1, F/x and p + 1. Since |F/x| < N(n-1, p+1, r-1, s-1) by the definition of p, we conclude that

$$\operatorname{rk} M_{s-1}^{r-1}(G/x) \ge \binom{n-1}{s-1} - p.$$
(16)

Whenever |F - x| < N(n - 1, t - p, r, s) we can apply part (i) of Claim 1 to F - x and t - p and conclude that

$$\operatorname{rk} M_s^r(G-x) > \binom{n-1}{s} - (t-p).$$
(17)

If we assume |F| < N(n, t, r, s) we certainly have p < t, as well as

$$|F - x| = |F| - |F/x| < N(n, t, r, s) - N(n, p, r, s) = \sum_{i=p+1}^{t} \binom{n-s+1-i}{r-s}$$
$$\leq \sum_{i=1}^{t-p} \binom{n-s-i}{r-s} = N(n-1, t-p, r, s),$$
(18)

so both (16) and (17) apply. Hence by Lemma 16 we obtain $\operatorname{rk} M_s^r(G) > \binom{n}{s} - t$ and the proof of our part (i) is complete.

For part (*ii*) we assume that $\operatorname{rk} M_s^r(G) \leq {\binom{n}{s}} - t$. Then p = t or |F - x| = N(n-1, t-p, r, s), otherwise both (16) and (17) apply, contradicting our rank assumption.

Assume first that |F - x| = N(n - 1, t - p, r, s). In this case we must also have equality everywhere in (18): |F/x| = N(n, p, r, s) = N(n - 1, p, r - 1, s - 1) and p = 1. We will arrive at a contradiction.

Applying Theorem 17 to s - 1 < r - 1, F/x and p, we get that $\operatorname{rk} M_{s-1}^{r-1}(G/x) \ge {\binom{n-1}{s-1}} - p$ and applying Claim 1 to F - x and t - p we have $\operatorname{rk} M_s^r(G - x) \ge {\binom{n-1}{s}} - (t - p)$. This means that, in order to avoid contradiction with $\operatorname{rk} M_s^r(G) \le {\binom{n}{s}} - t$, we must have equalities in both rank-inequalities. Then by part (*ii*) of our induction F/x is the upper (r - s)-shadow of an (s - 1)-star configuration of size p and F - x is the upper (r - s)-shadow of an s-star configuration of size (t - p). Now we use that p = 1: there must be a vertex $z \in V(G) \setminus \{x\}$ which occurs in all members of F/x. Since F - x is an upper shadow, the degree of z in F - xis at least 1. Hence z has degree at least $|F/x| + 1 = {\binom{n-s}{r-s}} + 1$ in F. As p = 1, the maximum degree in F should only be $N(n, 1, r, s) = {\binom{n-s}{r-s}}$, a contradiction.

Assume now p = t. Then $\deg_F(x) = N(n, t, r, s) = |F|$, so every edge of F contains x. Hence $G - x \simeq K_{n-1}^r$, so by Gottlieb's Theorem we have

$$\operatorname{rk} M_s^r(G-x) = \binom{n-1}{s}$$

as n' > r + s.

We apply Theorem 17 for s-1 with r-1, F/x, and t. By (11) |F/x| = N(n-1, t, r-1, s-1), so we conclude that

$$\operatorname{rk} M_{s-1}^{r-1}(G/x) \ge \binom{n-1}{s-1} - t,$$
(19)

with equality only if |F/x| = N(n-1,t,r-1,s-1) = N(n,t,r,s) and F/x is the upper (r-s)-shadow of an (s-1)-star configuration S of size t. By Lemma 16,

$$\operatorname{rk} M_s^r(G) \ge \operatorname{rk} M_s^r(G-x) + \operatorname{rk} M_{s-1}^{r-1}(G/x) \ge \binom{n}{s} - t$$

and for equality to hold we must have equality in (19), that is |F| = |F/x| = N(n, t, r, s) and F is the upper (r-s)-shadow of an s-star configuration of size t (x appended to each member of S). This finishes the proof in Case 1.

Case 2:
$$\alpha \binom{n-s}{r-s} \leq \deg_F(x) < \binom{n-s}{r-s}$$
.

In this case we show that if $|F| \leq N(n,t,r,s)$, then $\operatorname{rk} M_s^r(G) > {n \choose s} - t$, concluding the proof of Claim 1.

We first show that $\operatorname{rk} M_{s-1}^{r-1}(G/x) = \binom{n-1}{s-1}$.

If s = 1, we have $\operatorname{rk} M_0^{r-1}(G/x) = 1$, since $\deg_F(x) < \binom{n-1}{r-1}$.

If s > 1 we apply part (i) of Theorem 17 to r - 1, s - 1, F/x and t = 1. Since $|F/x| = \deg_F(x) < N(n-1, 1, r-1, s-1)$ we conclude that $\operatorname{rk} M_{s-1}^{r-1}(G/x) = \binom{n-1}{s-1}$.

For rk $M_s^r(G-x)$ we apply Claim 1 to F-x and t. Since $n' \ge n_3$ and $t \le c'n < \sigma n$, we may apply Lemma 15 to obtain $|F-x| \le N(n,t,r,s) - \alpha \binom{n-s}{r-s} < N(n-1,t,r,s)$. Hence we conclude that rk $M_s^r(G-x) > \binom{n-1}{s} - t$.

Consequently by Lemma 16 we obtain $\operatorname{rk} M_s^r(G) > \binom{n}{s} - t$.

By the remarks following the statement of Claim 1, this implies Theorem 17.

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5. TIGHTNESS OF THE MAIN RESULT

Fix r and s and suppose $n \ge n'_0(r,s)$ is large enough. Recall the definition of $t_{\max} := t_{\max}(n,r,s)$. It is the maximum value of $t \le \binom{n}{s}$ such that for any $t' \le t$, $\binom{n}{r} - \operatorname{rex}(n,t',r,s)$ equals K(n,t',s,r-s), the lower bound given by the Kruskal-Katona theorem for the upper (r-s)-shadow of a collection of s-subsets of [n] of size t.

We will show that $t_{\text{max}} < n - r - 1$. The first step is provided by the following lemma.

Lemma 18. Let $r > s \ge 1$. For any $n \ge r+2$ there exists an r-graph R := R(n,r,s) on n vertices with $\operatorname{rk} M_s^r(R) \le \binom{n}{s} - (n-r-1)$, but $\binom{n}{r} - |R| < N(n,n-(r+1),r,s)$.

Proof. The proof of this lemma is partially based on a construction of Keevash from [9].

We prove the lemma by induction on s.

First assume s = 1.

If r = 2, we define $R(4, 2, 1) := C_4$. Then R(4, 2, 1) verifies the conditions of the lemma.

If r > 2, we define R := R(r+2, r, 1) as the complete r-graph on the set [r+2], from which all the edges spanned by the vertices $\{1, \ldots, r+1\}$ were removed. Then $\alpha_i := 1, 1 \le i \le r+1$, and $\alpha_{r+2} := -(r-1)$, is a dependence sequence for R, showing that $\operatorname{rk} M_1^r(R) < r+2$. As $\binom{r+2}{r} - |R| = \binom{r+1}{r} < \binom{r+1}{r-1} = N(r+2, 1, r, 1)$ by (10), the graph R also verifies the conditions of the lemma.

For any n > r + 2, we let R(n, r, 1) have the same edges as R(r + 2, r, 1), but extend the vertex set to [n]. Then R(n, r, 1) has rank of the associated inclusion matrix at most r + 1, and

$$\binom{n}{r} - |R(n,r,1)| < \binom{n}{r} - \binom{r+1}{r} = N(n,n-(r+1),r,1).$$

Hence the claim holds for s = 1.

Now assume s > 1 and the claim holds for any pair (s', r') with s' + r' < s + r.

We define the edge-set of R(n, r, s) as the collection $\binom{[n-1]}{r}$, to which we add all sets $\{n\} \cup e$, where e is an edge in R(n-1, r-1, s-1).

Note that

$$\operatorname{rk} M_{s}^{r}(R(n,r,s)) \leq {\binom{n-1}{s}} + \operatorname{rk} M_{s-1}^{r-1}(R(n-1,r-1,s-1)).$$
(20)

Indeed, from the columns corresponding to the s-sets containing n we can pick at most $\operatorname{rk} M_{s-1}^{r-1}(R(n-1,r-1,s-1))$ linearly independent ones, to which we may further add at most $\binom{n-1}{s}$ columns, to form a linearly independent collection of size $\operatorname{rk} M_s^r(R(n,r,s))$. Therefore,

$$\binom{n}{s} - \operatorname{rk} M_s^r(R(n,r,s)) \ge \binom{n-1}{s-1} - \operatorname{rk} M_{s-1}^{r-1}(R(n-1,r-1,s-1)) \ge n - (r+1).$$

Furthermore, by construction,

$$\binom{n}{r} - |R(n,r,s)| = \binom{n-1}{r-1} - |R(n-1,r-1,s-1)| < N(n-1,n-(r+1),r-1,s-1),$$

where the last inequality follows by the induction hypothesis. The proof now follows from (11).

We are now in position to prove Theorem 7.

Proof of Theorem 7. Let $r > s \ge 1$. We set $n'_0(r,s) := r+2$ and let $n \ge n'_0(r,s)$.

We first claim that for t = n - r - 1, the lower bound provided by the Kruskal-Katona theorem is still equal to N(n,t,r,s). As $n-r-1 \leq n-s$, by (4) we have $t = \binom{n-s}{n-s} + \dots + \binom{n-s-t+1}{n-s-t+1}$. In the same manner as in the Introduction we obtain that the lower bound provided by the Kruskal-Katona theorem for the upper (r-s)-shadow of a collection of *s*-subsets of [n] of size *t* equals $\binom{n-s+1}{r-s+1} - \binom{n-s+1-t}{r-s+1} = N(n,t,r,s)$, hence the claim. However, by Lemma 18, for $n \ge r+2$ and t = n-r-1, $\binom{n}{r} - \operatorname{rex}(n,t,r,s)$ is no longer

equal to N(n, t, r, s). Therefore $t_{\max} < n - r - 1$, as desired.

6. Concluding Remarks

We close our discussion with several remarks and questions.

6.1. Rank-extremal function for graphs and other ranges of t. In the particular case of s = 1 and r = 2 the rank-extremal function as well as the rank-extremal graphs can be completely determined for all values of $t \leq n$. Indeed, since the rank of a graph G is the sum of the ranks of its components, and a component C of G has full rank |V(C)| if it is not bipartite and has rank |V(C)| - 1 otherwise, one obtains the following. For any graph G on n vertices, $\operatorname{rk} M_1^2(G) = n - b(G)$, where b(G) is the number of bipartite components of G. With this characterization at hand, it is not difficult to compute rex(n, t, 2, 1) and the rankextremal graphs: if G is a rank-extremal graph with n vertices and $\operatorname{rk} M_1^2(G) = n - t$, then G consists of t bipartite components and of possibly one non-bipartite complete subgraph. All one has to do now is to choose the sizes of components such that the number of edges is maximized. We omit the straightforward details, but describe briefly rex(n, t, 2, 1) for $n \ge 6$. Let G be maximal such that $\operatorname{rk} M_1^2(G) \leq n-t$. Remove all isolated vertices from G to form a new graph H. The structure of H is now easy to describe. If t < n - 4 then $H = K_{n-t}$. If t = n - 4 then H is either $K_{2,3}$ or K_4 , while if t = n - 3, H must be C_4 . Finally, for t = n - 2, H is either P_2 or two disjoint edges, for t = n - 1, H is just K_2 and for t = n, H is the empty graph. In particular, Theorem 5 does not hold for t = n - r - 2, as there are two extremal graphs.

The case of n = 4, s = 1, r = 2 and C_4 also shows that Theorem 5 may fail to hold for small values of n, even under the additional assumption that $n \ge r + s$. Hence the constant n_0 in the theorem statement is necessary. It would be interesting to determine the best constant c_0 in Theorem 5. Our argument only shows it is at most 1.

Using Theorem 3, we can determine rex(n, t, r, s) asymptotically. Indeed, for any large enough real $x \in [0, n]$, we have that $\binom{x}{r} \ge \operatorname{rex}(n, \binom{n}{s} - \binom{x}{s}, r, s) > \binom{x-1}{r}$. The upper bound follows from the definition and Theorem 3, while the lower bound follows by considering the rgraph G on [n] and edge set $\binom{[\lfloor x \rfloor]}{r}$. Consequently the rank-extremal function is asymptotically like $x^r/r!$. However, the error of this approximation becomes very large when x approaches n, exactly the case covered by our Theorem 5. In general, the exact values of the rank-extremal function as well as the structure of the rank-extremal hypergraphs are still an open problem.

6.2. Constructive proof of Theorem 5 for s = 1. The proof of Theorem 5 is non-constructive in the sense that we only show that for an r-graph G with more than rex(n, t, r, s) edges there exist some $\binom{n}{s} - t + 1$ linearly independent rows in the inclusion matrix $M_s^r(G)$. In the case t = s = 1 we can show that $M_s^r(G)$ has full (column) rank by finding a particular subgraph of G with n edges, whose rows are linearly independent, as follows.

Let H' be an (r-1)-tight Hamilton cycle on [n'] with $n' \equiv r-1 \mod r$ and $n' \ge n-r+1$ (this tight Hamilton cycle has its edges being any r consecutive elements from [n'] modulo n'). The r-graph H on the vertex set [n] is obtained by adding to H'(n-n') new vertices $n'+1, \ldots, n$ and adding new disjoint edges $e_{n'+1}, \ldots, e_n$ such that e_i contains vertex i and (r-1) further consecutive vertices from the cycle H'. It is not difficult to see that $M_s^r(H)$ has full rank n (the inclusion matrix of the (r-1)-tight Hamilton cycle H' on n' vertices has full rank and adding each of the edges e_i increases the rank each time by 1). It follows then from the work in [6] on the extremal number of Hamilton cycles that any r-graph G on n vertices with $\operatorname{rex}(n, 1, r, 1) + 1 = \binom{n}{r} - \binom{n-1}{r-1} + 1$ edges must contain a copy of H. This can be seen as the generalization of the graph argument described above: a graph

This can be seen as the generalization of the graph argument described above: a graph with $\binom{n-1}{2} + 1$ edges and *n* vertices $(n \ge 5)$ is clearly connected and not bipartite, yielding that its inclusion matrix has full column rank. It would be interesting to investigate what structures will replace the bipartite components for *r*-graphs when t > 1 and to look for minimal structures whose inclusion matrices have full rank in the case r > s > 1.

6.3. Random inclusion matrices. Lots of research in random graphs dealt recently with questions of transferring results from extremal combinatorics to a probabilistic setting. Given a random r-graph G(n,p) which is the product probability space $\{0,1\}^{\binom{n}{r}}$ where $\Pr(G) = p^{|G|}(1-p)^{\binom{n}{r}-|G|}$. The basic question one might ask is the following: what is the threshold for the property that rk $M_s^r(G(n,p)) = \binom{n}{s}$?

In the graph case (r = 2, s = 1), the inclusion matrix $M_1^2(G(n, p))$ gets full column rank, exactly at the very moment when G(n, p) becomes connected (since it contains a connected spanning subgraph with an odd cycle and with n edges - and its inclusion matrix already has full rank). Further, the threshold p is sharp and equals $\frac{\ln n}{n}$. For $r > s \ge 1$, rk $M_s^r(G(n, p))$ stays clearly below $\binom{n}{s}$ if there is a zero column. It is further easy to see that the threshold p for the property that $M_s^r(G(n, p))$ has no zero columns is precisely at $\frac{s(r-s)!\ln n}{n^{r-s}}$, which is a standard first moment-second moment calculation. This gives us a lower bound on the threshold for $M_s^r(G(n, p))$ to possess full rank.

On the other hand, it is not hard to prove the upper bound on $p = O(\frac{\ln^2 n}{n^{r-s}})$ which is only slightly apart from the lower bound. Here is a rough sketch. First, from Theorem 3 and some straightforward estimates for binomial coefficients one can obtain the following lemma.

Lemma 19. If G is an r-uniform hypergraph with n vertices and at least $\binom{n}{r} - tn^{r-s}$ edges, then

$$\operatorname{rk} M_s^r(G) \ge \binom{n}{s} - c \cdot t,$$

where c > 0 depends on r only.

Next, suppose that we are given G with n vertices and of the rank $\binom{n}{s} - ct - 1$. Then, by Lemma 19, we obtain $\Pr\left[\operatorname{rk} M_s^r(G \cup G(n,p)) < \binom{n}{s} - c \cdot t\right] \leq (1-p)^{tn^{r-s}} = \exp(-ptn^{r-s})$. Thus, starting with the empty hypergraph on n vertices, and performing at most $\binom{n}{s}$ times multi-round exposure, we obtain

$$\Pr\left[\operatorname{rk} M_s^r(\cup_i G(n, p_i)) < \binom{n}{s}\right] \le \sum_{i=1}^{\binom{n}{s}} \exp(-p_i i n^{r-s}).$$
(21)

Therefore, it is enough to choose $p_i = \frac{C \log n}{in^{r-s}}$ for C sufficiently large, so that the sum in (21) is o(1). Moreover, since $\sum_{i=1}^{\binom{n}{s}} p_i = O(\frac{\log^2 n}{n^{r-s}})$, the upper bound on the threshold for $M_s^r(G(n,p))$ to possess full rank follows.

Finally, it should be mentioned, that Friedgut's theorem [5] implies that the property $\operatorname{rk} M_s^r(G(n,p)) = \binom{n}{s}$ has a sharp threshold, since the phase transition happens between $\frac{\ln n}{n^{r-s}}$ and $\frac{\ln^2 n}{n^{r-s}}$ and thus p does not have a rational exponent. It would be interesting to determine the precise threshold for the property above. In fact, we believe it should match the lower bound.

6.4. Local resilience of the rank. The study of graph resilience was initiated by Sudakov and Vu in [17]. In view of Theorem 6, one can formulate a similar question about the rank of K_n^r .

Let $r \ge s \ge 0$. If $F \subseteq {\binom{[n]}{r}}$ and $S \in {\binom{[n]}{s}}$ we define $\deg_F(S)$ as the number of elements of F containing S. Then one has the following problem.

Problem 1. For n sufficiently large, determine the maximum value m(n, r, s) such that for any $F \subseteq {\binom{[n]}{r}}$ with $\deg_F(S) \leq m(n, r, s), \forall S \in {\binom{[n]}{s}}$, we have $\operatorname{rk} M_s^r(K_n^r - F) = {\binom{n}{s}}$.

Theorem 6 shows that for $r \ge 2s - 1$ we have $m(n, r, s) \ge cn^{r+1-2s}$, for some c > 0.

6.5. Fields of finite characteristic. It turns out that the inclusion matrix of the complete hypergraph does not necessarily possess the full rank. In the case of GF(2) the rank was computed by Linial and Rothschild [13] and in the case of GF(p) by Wilson [18], see also Frankl [3].

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