# Robust Multigrid for Isogeometric Analysis Based on Stable Splittings of Spline Spaces 

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May 25, 2022


#### Abstract

We present a robust and efficient multigrid method for single-patch isogeometric discretizations using tensor product B-splines of maximum smoothness. Our method is based on a stable splitting of the spline space into a large subspace of "interior" splines which satisfy a robust inverse inequality, as well as one or several smaller subspaces which capture the boundary effects responsible for the spectral outliers which occur in Isogeometric Analysis. We then construct a multigrid smoother based on an additive subspace correction approach, applying a different smoother to each of the subspaces. For the interior splines, we use a mass smoother, whereas the remaining components are treated with suitably chosen Kronecker product smoothers or direct solvers.

We prove that the resulting multigrid method exhibits iteration numbers which are robust with respect to the spline degree and the mesh size. Furthermore, it can be efficiently realized for discretizations of problems in arbitrarily high geometric dimension. Some numerical examples illustrate the theoretical results and show that the iteration numbers also scale relatively mildly with the problem dimension.


## 1 Introduction

Isogeometric Analysis (IgA) is a method for the numerical solution of partial differential equations (PDEs) introduced in the seminal paper [18] which has since attracted a sizable research community. Spline spaces, such as spaces spanned by tensor product B-splines or NURBS, are commonly used for geometry representation in industrial CAD systems. The foundational idea in $\operatorname{Ig} A$ is to use such spline spaces both for the representation of the computational domain and for the discretization of the quantities of interest when solving a PDE. The overall goal is to create a tighter integration between geometric design and analysis.

There is a need for efficient solvers for the large, sparse linear systems which arise when applying isogeometric discretizations to boundary value problems. By now, most established solution strategies known from the finite element literature have been applied in one way or another to IgA: among these, direct solvers [2], non-overlapping and overlapping domain decomposition methods [19, 4, 5, 6, and multilevel and multigrid methods [1, 11, 17, 10, 15]. A recent contribution

[^0][20] constructs preconditioners based on fast solvers for Sylvester equations. The above list is certainly not comprehensive.

In IgA, we typically encounter as discretization parameters the mesh size and the spline degree. In the early IgA solver literature, the focus was on translating solvers from the finite element world to IgA with minimal adaptations. As a rule, it was found that such an approach results in methods that work well for low spline degrees, but deteriorate in performance as the degree is increased; often dramatically so. This motivated the search for IgA solvers that are robust not only with respect to the mesh size (which is often easy to achieve), but also with respect to the spline degree.

Within the class of multigrid methods for $\operatorname{IgA}$, advances towards a robust method were made using two approaches. In [9], a careful analysis of the symbol of isogeometric stiffness matrices served as the basis for the construction of multigrid methods. This theoretical approach is somewhat related to the technique known as Local Fourier Analysis (LFA) in the multigrid literature (see, e.g., [22]). It appears that the method presented in [9] is roughly comparable to the one studied in [16], which uses mass matrices as multigrid smoothers, an approach itself motivated by LFA. For both methods, an increase in the number of smoothing steps, roughly linearly with the spline degree, is required in order to maintain robust convergence. They can thus not be considered totally robust and efficient in the strict meaning that we will use in the present work.

A second approach towards robust and efficient multigrid was presented in [15. Based on a robust inverse inequality and approximation error estimate in a large subspace of maximally smooth spline spaces derived in [21, it was shown that mass matrices can be used as robust smoothers in this large subspace. For the remaining, relatively few degrees of freedom, a lowrank correction was constructed. (These degrees of freedom are associated with the boundary of the domain and cannot be captured by LFA, which assumes periodic boundary conditions.) This approach resulted in a provably robust and efficient multigrid method for two-dimensional problems with splines of maximum smoothness. It was however not clear how to extend this approach efficiently to three and higher dimensions.

The present work can be viewed as a continuation of 15]. Based on the theoretical results from [21, we construct a splitting of the tensor product spline space into a large, regular interior part and several smaller spaces which capture boundary effects. The splitting is $L_{2}$-orthogonal and $H^{1}$-stable with respect to both the mesh size and the spline degree. This stability enables us to construct a multigrid smoother based on an additive subspace correction approach, applying a different smoother in each of the subspaces. In the regular interior subspace, we use a mass smoother. In the other subspaces, we construct smoothers which exploit the particular structure of the subspaces while still permitting an efficient application through a Kronecker product representation. In one small subspace associated with the corners of the domain, we apply a direct solver.

Unlike the low-rank correction approach from [15], the subspace correction approach generalizes easily to three-dimensional problems, and indeed to problems of arbitrary space dimension. We show that the method converges robustly with respect to mesh size and spline degree, and that one iteration is asymptotically not more expensive than an application of the stiffness matrix. The result is a quasi-optimal solution method for problems of arbitrary space dimensions.

It appears that the stable splitting of the tensor product spline space presented in Section 3 is an interesting theoretical result in its own right. It may have future applications to other aspects of $\operatorname{IgA}$ beyond the one presented here.

The remainder of the paper is organized as follows. In Section 2 we introduce the needed spline spaces and present an isogeometric model problem. We also present an algorithmic multigrid framework and an abstract convergence result which forms the basis of our later analysis. In Section 3, we derive the main new theoretical result used in our construction: the $L_{2}$-orthogonal and $H^{1}$-stable splitting of the spline space into a large, regular interior part and smaller spaces which capture boundary effects. In Section 4, we use this space splitting to construct a multigrid smoother based on the idea of additive subspace correction and show that it results in a robust solver. In Section 5, we present details on the computational realization of the proposed smoother and show that it permits an efficient implementation in arbitrary space dimensions. In Section 6 , we present numerical experiments which demonstrate the performance of the proposed method in
practice.

## 2 Preliminaries

### 2.1 Spline spaces and B-splines

Consider a subdivision of the interval $(0,1)$ into $m \in \mathbb{N}$ intervals of length $h=1 / m$. We introduce the spline space of degree $p \in \mathbb{N}$ with maximum smoothness,

$$
S:=\left\{u \in C^{p-1}(0,1):\left.u\right|_{((j-1) h, j h)} \in \mathcal{P}^{p} \quad \forall j=1, \ldots, m\right\}
$$

where $C^{p-1}(0,1)$ is the space of all $p-1$ times continuously differentiable functions on $(0,1)$ and $\mathcal{P}^{p}$ is the space of all polynomials of degree at most $p$. We have $n:=\operatorname{dim} S=m+p$. As a basis for $S$, we use the normalized (i.e., satisfying a partition of unity; cf. [8) B-splines with an open knot vector. In higher dimensions $d>1$, we introduce the space of tensor product splines (cf. [8])

$$
S^{d}:=S \otimes \ldots \otimes S
$$

defined over $(0,1)^{d}$ with $\operatorname{dim} S^{d}=n^{d}$ and the corresponding tensor product B-spline basis. For notational convenience, we assume that the same spline space $S$ is used in each of the $d$ coordinate directions. Both our construction and our analysis are however straightforward to generalize to the case where different spline spaces are used in different coordinate directions.

### 2.2 Isogeometric model problem

Let $\Omega=(0,1)^{d}$ with $d \in \mathbb{N}$. As a model problem, we consider a pure Neumann boundary value problem for the $\operatorname{PDE}-\Delta u+u=f$. The variational formulation reads: find $u \in H^{1}(\Omega)$ such that

$$
\begin{equation*}
a(u, v)=\langle f, v\rangle \quad \forall v \in H^{1}(\Omega) \tag{1}
\end{equation*}
$$

where

$$
\begin{equation*}
a(u, v)=\int_{\Omega}(\nabla u \cdot \nabla v+u v) d x \quad \forall u, v \in H^{1}(\Omega) \tag{2}
\end{equation*}
$$

and $f$ is a linear functional on $H^{1}(\Omega)$. We will sometimes refer to the operator $A: H^{1}(\Omega) \rightarrow H^{1}(\Omega)^{\prime}$ given by $A v=a(v, \cdot)$, where $H^{1}(\Omega)^{\prime}$ denotes the continuous dual. Note that $\|v\|_{A}^{2}=a(v, v)=$ $\|v\|_{H^{1}(\Omega)}^{2}$.

Discretizing (1) using tensor product splines, we seek $u_{h} \in S^{d}$ such that

$$
\begin{equation*}
a\left(u_{h}, v_{h}\right)=\left\langle f, v_{h}\right\rangle \quad \forall v_{h} \in S^{d} \tag{3}
\end{equation*}
$$

We are interested in robust and efficient iterative solvers for the discrete problem (3). Here, by "robust" we mean that the number of iterations to solve the problem should stay uniformly bounded with respect to both the mesh size $h$ and the spline degree $p$, and by "efficient" we mean that one iteration of the method should not be asymptotically more expensive than computing the product of the stiffness matrix with a vector. Combined, these properties allow us to solve (3) in quasi-optimal time.

In $\operatorname{Ig} A$, one introduces a bijective geometry map from $\Omega$ to the actual domain of interest in order to be able to treat more complicated computational domains. Basis functions on the transformed domain are defined by composing the basis functions on the reference domain with the inverse of the geometry map. Furthermore, one is often interested in more general PDEs with varying and possibly matrix-valued coefficients. Discretizations for such more general problems can be preconditioned with a solver for the model problem (3), and the resulting condition number depends only on the geometry map and the coefficient functions, but not on discretization parameters like the mesh size $h$ or the spline degree $p$. This principle has been widely used in the literature on IgA solvers (see, e.g., [9, 15]) and formalized in [20. Therefore, a robust and efficient
solver for the model problem (3) immediately yields robust and efficient solvers for a more general class of problems with "benign" geometry maps and mildly varying coefficients. This justifies the study of solvers for the model problem.

Three different refinement strategies for IgA discretizations were proposed in [18]: $h$-refinement (reducing the mesh size), $p$-refinement (increasing the spline degree), and the so-called $k$-refinement. The latter is unique to IgA and maintains the maximum possible smoothness $C^{p-1}$ for the spline space of degree $p$. Already in [18, the favorable performance of $k$-refinement was observed, and it appears to be the most popular refinement strategy in the wider IgA literature. This motivates the study of solvers for spline spaces with maximum smoothness.

### 2.3 A multigrid method framework

Given a discretization space $V$ and a coarse space $V_{c} \subset V$, we denote by $P: V_{c} \rightarrow V$ the canonical embedding. Let $A: V \rightarrow V^{\prime}$ denote the operator in a (discretized) equation

$$
A u=f
$$

to be solved for $u \in V$. The corresponding coarse-space operator is given by $A_{c}:=P^{\prime} A P$. Furthermore, we assume that we are given a self-adjoint and positive definite smoothing operator $L: V \rightarrow V^{\prime}$.

Given a previous iterate $u^{(k)}$, we let $u^{(k, 0)}:=u^{(k)}$ and perform $\nu \in \mathbb{N}$ smoothing steps given by

$$
u^{(k, j)}:=u^{(k, j-1)}+\tau L^{-1}\left(f-A u^{(k, j-1)}\right), \quad j=1, \ldots, \nu
$$

where $\tau>0$ is a damping parameter. Then, we perform one coarse-grid correction step given by

$$
u^{(k+1)}:=u^{(k, \nu)}+P A_{c}^{-1} P^{\prime}\left(f-A u^{(k, \nu)}\right) .
$$

Together, these updates describe one iteration $u^{(k)} \mapsto u^{(k+1)}$ of a two-grid method. Given an entire sequence of nested spaces $V_{0} \subset \ldots \subset V_{L}=V$, we can replace the exact inversion of $A_{c}$ in the coarse-grid correction step by one or two recursive applications of the two-grid method on the next coarser level $V_{L-1}$, and so on until we reach the coarsest level $V_{0}$, where an exact solver is used. Using one or two recursive iteration steps results in the $V$-cycle or the $W$-cycle multigrid method, respectively.

The following theorem is an abstract convergence result for the two-grid method with the abovementioned smoother. Its proof is given in [15, Theorem 3] and is based on a variant of the standard multigrid theory as developed by Hackbusch [14]. In [15, Theorem 4], it was shown that under the same assumptions also a W-cycle multigrid method converges.

Theorem 1 ([15]). Assume that there are constants $C_{A}$ and $C_{I}$ such that the inverse inequality

$$
\begin{equation*}
\|u\|_{A}^{2} \leq C_{I}\|u\|_{L}^{2} \quad \forall u \in V \tag{4}
\end{equation*}
$$

and the approximation property for the $A$-orthogonal projector $T_{c}: V \rightarrow V_{c}$

$$
\begin{equation*}
\left\|\left(I-T_{c}\right) u\right\|_{L}^{2} \leq C_{A}\|u\|_{A}^{2} \quad \forall u \in V \tag{5}
\end{equation*}
$$

hold. Then the two-grid method converges for any choice of the damping parameter $\tau \in\left(0, C_{I}^{-1}\right]$ and any number of smoothing steps $\nu>\nu_{0}:=\tau^{-1} C_{A}$ with rate $q=\nu_{0} / \nu<1$.

In particular, if $C_{A}$ and $C_{I}$ do not depend on the mesh size $h$ and the spline degree $p$, then the two-grid method converges with a rate $q<1$ which does not depend on $h$ and $p$. In other words, the two-grid method is then robust.

In addition to properties (4) and (5), care must be taken that the smoother can be realized efficiently. In other words, it should be possible to apply the inverse $L^{-1}$ with a computational cost which is roughly comparable to that for applying $A$.

## 3 Stable splittings of spline spaces

Consider first the univariate case, $d=1$, with $\Omega=(0,1)$. In [21], the subspace

$$
S_{0}:=\left\{u \in S: u^{(2 l+1)}(0)=u^{(2 l+1)}(1)=0 \quad \forall l \in \mathbb{N}_{0} \text { with } 2 l+1<p\right\}
$$

of splines with vanishing odd derivatives of order less than $p$ at the boundaries was introduced (denoted in 21 by $\widetilde{S}_{p, h}(\Omega)$ ). It is a large subspace of $S$ in the sense that $\operatorname{dim} S_{0} \geq \operatorname{dim} S-p$.

The subspace $S_{0}$ has the very desirable property of satisfying both a (first-order) approximation property and an inverse inequality, both with constants which are independent of the spline degree $p$. To formulate these results, let $Q_{0}: L_{2}(\Omega) \rightarrow S_{0}$ denote the $L_{2}$-orthogonal projector into $S_{0}$, and let $\Pi_{0}: H^{1}(\Omega) \rightarrow S_{0}$ denote the projector into $S_{0}$ which is orthogonal with respect to the scalar product

$$
(u, v)_{H_{\circ}^{1}(\Omega)}:=(\nabla u, \nabla v)_{L_{2}(\Omega)}+\frac{1}{|\Omega|}\left(\int_{\Omega} u(x) d x\right)\left(\int_{\Omega} v(x) d x\right) .
$$

We abbreviate the $L_{2}(\Omega)$-norm by $\|\cdot\|_{0}$, and the full $H^{1}(\Omega)$-norm and the seminorm by $\|\cdot\|_{1}$ and $|\cdot|_{1}$, respectively. Furthermore, we write $c$ for a generic positive constant which does not depend on the mesh size $h$ or the spline degree $p$.

Theorem 2 ([21, Theorem 6.1]). For any spline degree $p \in \mathbb{N}$, we have the inverse inequality

$$
|u|_{1} \leq 2 \sqrt{3} h^{-1}\|u\|_{0} \quad \forall u \in S_{0}
$$

Theorem 3 ([21, Corollary 5.1], [15, Theorem 14]). For any spline degree $p \in \mathbb{N}$ and any $u \in$ $H^{1}(\Omega)$, we have the approximation error estimates

$$
\left\|\left(I-Q_{0}\right) u\right\|_{0} \leq \sqrt{2} h|u|_{1} \quad \text { and } \quad\left\|\left(I-\Pi_{0}\right) u\right\|_{0} \leq \sqrt{2} h|u|_{1} .
$$

Contrast these properties with the entire spline space $S$, which does satisfy a robust approximation property, but whose inverse inequality deteriorates with increasing degree $p$ ([21]). On the other hand, a smaller space of only "interior" splines, built by discarding the $p$ leftmost and $p$ rightmost B-splines, does satisfy a robust inverse inequality but loses the approximation property.

We remark that the non-robustness of the inverse inequality in $S$ is the root cause of the spectral "outliers" commonly observed when solving eigenvalue problems using $\operatorname{IgA}$ (cf. 3). No such outliers appear in the space $S_{0}$.

### 3.1 A stable splitting in one dimension

Let $S_{1}:=S_{0}^{\perp_{L_{2}}}$ denote the $L_{2}$-orthogonal complement of $S_{0}$ in $S$. Consider the splitting of $S$ into the direct sum

$$
S=S_{0} \oplus S_{1} \quad \longleftrightarrow \quad u=Q_{0} u+\left(I-Q_{0}\right) u
$$

of $S_{0}$ and its complement, illustrated in Fig. 1. Due to orthogonality, we have

$$
\begin{equation*}
\|u\|_{0}^{2}=\left\|Q_{0} u\right\|_{0}^{2}+\left\|\left(I-Q_{0}\right) u\right\|_{0}^{2} \tag{6}
\end{equation*}
$$

Crucially, we can prove that this splitting is stable also in the $H^{1}$-norm. This is a direct result of the space $S_{0}$ satisfying both an approximation property and an inverse inequality.

Theorem 4. For any spline $u \in S$, we have

$$
c^{-1}|u|_{1}^{2} \leq\left|Q_{0} u\right|_{1}^{2}+\left|\left(I-Q_{0}\right) u\right|_{1}^{2} \leq c|u|_{1}^{2}
$$

and the corresponding result for the full $H^{1}$-norm.


Figure 1: Bases for the space $S_{0}$ (left) and its orthogonal complement $S_{1}$ (right) for $p=4$, $h=1 / 20$. Here, $\operatorname{dim} S_{0}=20$ and $\operatorname{dim} S_{1}=4$.

Proof. The left inequality follows from the Cauchy-Schwarz inequality with $c=2$. For the right inequality, we observe that

$$
\left|Q_{0} u\right|_{1} \leq\left|\Pi_{0} u\right|_{1}+\left|\left(\Pi_{0}-Q_{0}\right) u\right|_{1} \leq|u|_{1}+c h^{-1}\left(\left\|\left(I-\Pi_{0}\right) u\right\|_{0}+\left\|\left(I-Q_{0}\right) u\right\|_{0}\right)
$$

because of the triangle inequality, the stability of the $H_{0}^{1}$-projector $\Pi_{0}$ in the $H^{1}$-seminorm and the robust inverse inequality in $S_{0}$ (Theorem 22). With the approximation error estimate (Theorem 3) we obtain $H^{1}$-stability of the $L_{2}$-projector,

$$
\begin{equation*}
\left|Q_{0} u\right|_{1} \leq c|u|_{1} \tag{7}
\end{equation*}
$$

The desired result follows from (7) and

$$
\left|\left(I-Q_{0}\right) u\right|_{1} \leq|u|_{1}+\left|Q_{0} u\right|_{1} \leq(1+c)|u|_{1} .
$$

The result for the full $H^{1}$-norm follows by adding the identity (6).

### 3.2 A stable splitting in two dimensions

The two-dimensional tensor product spline space is given by $S^{2}=S \otimes S$. Since the tensor product distributes over direct sums, we obtain the splitting

$$
S^{2}=\left(S_{0} \otimes S_{0}\right) \oplus\left(S_{0} \otimes S_{1}\right) \oplus\left(S_{1} \otimes S_{0}\right) \oplus\left(S_{1} \otimes S_{1}\right)=S_{00} \oplus S_{01} \oplus S_{10} \oplus S_{11}
$$

with the abbreviations $S_{\alpha_{1}, \alpha_{2}}:=S_{\alpha_{1}} \otimes S_{\alpha_{2}}$ for $\alpha_{j} \in\{0,1\}$. A visualization of this splitting is shown in Fig. 2. Note that the shaded regions do not correspond to the supports of the function spaces; in fact, each of the subspaces has global support. However, the shaded regions roughly correspond to regions where the corresponding functions are "largest", and their areas roughly correspond to the space dimensions. In view of this, it makes sense to think of $S_{00}$ as an "interior" space, of $S_{01}$ and $S_{10}$ as "edge" spaces, and of $S_{11}$ as a "corner" space.


Figure 2: Visualization of the splitting in 2D.
Again, we can prove that the splitting is $H^{1}$-stable. In the following, we let $M: S \rightarrow S^{\prime}$, $K: S \rightarrow S^{\prime}$ denote the operators in the univariate spline space associated with the bilinear forms

$$
\langle M u, v\rangle:=\int_{0}^{1} u(x) v(x) d x, \quad\langle K u, v\rangle:=\int_{0}^{1} u^{\prime}(x) v^{\prime}(x) d x \quad \forall u, v \in S
$$

that is, the one-dimensional mass and stiffness operators, respectively. For any $\left(\alpha_{1}, \alpha_{2}\right) \in\{0,1\}^{2}$, we furthermore introduce the abbreviations

$$
\begin{aligned}
Q_{1} & :=I-Q_{0}: S \rightarrow S_{1}, & Q_{\alpha_{1}, \alpha_{2}} & :=Q_{\alpha_{1}} \otimes Q_{\alpha_{2}}: S^{2} \rightarrow S_{\alpha_{1}, \alpha_{2}} \\
K_{\alpha_{j}} & :=Q_{\alpha_{j}}^{\prime} K Q_{\alpha_{j}}: S_{\alpha_{j}} \rightarrow S_{\alpha_{j}}^{\prime}, & M_{\alpha_{j}} & :=Q_{\alpha_{j}}^{\prime} M Q_{\alpha_{j}}: S_{\alpha_{j}} \rightarrow S_{\alpha_{j}}^{\prime}
\end{aligned}
$$

As tensor products of $L_{2}(0,1)$-orthogonal projectors, the projectors $Q_{\alpha_{1}, \alpha_{2}}$ are $L_{2}(\Omega)$-orthogonal, as one easily verifies. Thus the splitting of $S^{2}$ given above is a direct sum of $L_{2}$-orthogonal subspaces, and we have

$$
\begin{equation*}
\|u\|_{0}^{2}=\sum_{\left(\alpha_{1}, \alpha_{2}\right)}\left\|Q_{\alpha_{1}, \alpha_{2}} u\right\|_{0}^{2} \tag{8}
\end{equation*}
$$

where here and below sums over $\left(\alpha_{1}, \alpha_{2}\right)$ are taken to run over the set $\{0,1\}^{2}$.
Theorem 5. For any tensor product spline $u \in S^{2}$, we have

$$
c^{-1}|u|_{1}^{2} \leq \sum_{\left(\alpha_{1}, \alpha_{2}\right)}\left|Q_{\alpha_{1}, \alpha_{2}} u\right|_{1}^{2} \leq c|u|_{1}^{2}
$$

and the corresponding result for the full $H^{1}$-norm.
Proof. The left inequality follows by the Cauchy-Schwarz inequality. For the right one, fix $\left(\alpha_{1}, \alpha_{2}\right) \in\{0,1\}^{2}$. The $H^{1}$-seminorm can be written using tensor products of one-dimensional operators as

$$
\begin{equation*}
\left|Q_{\alpha_{1}, \alpha_{2}} u\right|_{1}^{2}=\left|Q_{\alpha_{1}, \alpha_{2}} u\right|_{K \otimes M}^{2}+\left|Q_{\alpha_{1}, \alpha_{2}} u\right|_{M \otimes K}^{2} \tag{9}
\end{equation*}
$$

The first term can be rewritten, using the definitions and basic identities for tensor products of operators, as

$$
\left|Q_{\alpha_{1}, \alpha_{2}} u\right|_{K \otimes M}^{2}=\left\langle Q_{\alpha_{1}, \alpha_{2}}^{\prime}(K \otimes M) Q_{\alpha_{1}, \alpha_{2}} u, u\right\rangle=\left\langle\left(K_{\alpha_{1}} \otimes M_{\alpha_{2}}\right) u, u\right\rangle
$$

Due to orthogonality and Theorem 4, we have $M_{0}+M_{1}=M$ and $K_{0}+K_{1} \leq c K$, where all summands are positive semidefinite operators. This implies that we can estimate, in the spectral sense, $K_{\alpha_{1}} \leq c K$ and $M_{\alpha_{2}} \leq M$, and we obtain

$$
\left|Q_{\alpha_{1}, \alpha_{2}} u\right|_{K \otimes M}^{2} \leq c|u|_{K \otimes M}^{2}
$$

Treating the second term in (9) analogously, we obtain

$$
\left|Q_{\alpha_{1}, \alpha_{2}} u\right|_{1}^{2} \leq c\left(|u|_{K \otimes M}^{2}+|u|_{M \otimes K}^{2}\right)=c|u|_{1}^{2}
$$

The right inequality now follows by summing up over all ( $\alpha_{1}, \alpha_{2}$ ). The result for the full $H^{1}$-norm follows by adding the identity (8).

### 3.3 Stable splitting in arbitrary dimensions

For any $d \in \mathbb{N}$, we define multiindices $\alpha \in\{0,1\}^{d}$ and generalize the notations from Section 3.2 in the straightforward way to higher dimensions. We obtain the splitting into the direct sum of $2^{d}$ subspaces

$$
S^{d}=\bigoplus_{\alpha} S_{\alpha}, \quad \text { where } \quad S_{\alpha}=S_{\alpha_{1}} \otimes \ldots \otimes S_{\alpha_{d}}
$$

The $L_{2}$-orthogonal projectors into the subspaces are given by

$$
Q_{\alpha}=Q_{\alpha_{1}} \otimes \ldots \otimes Q_{\alpha_{d}}: S^{d} \rightarrow S_{\alpha}
$$

As in the two-dimensional case, we can prove that this splitting is $H^{1}$-stable.

Theorem 6. For any d-dimensional tensor product spline $u \in S^{d}$, we have

$$
c^{-1}|u|_{1}^{2} \leq \sum_{\alpha=(0, \ldots, 0)}^{(1, \ldots, 1)}\left|Q_{\alpha} u\right|_{1}^{2} \leq c|u|_{1}^{2}
$$

and the corresponding result for the full $H^{1}$-norm.
Proof. Completely analogous to Theorem 5

## 4 Construction of a robust multigrid method

Recall that $S$ was a univariate spline space of degree $p$ and mesh size $h$. Let $S_{c} \subset S$ be the analogous coarse spline space with uniform mesh size $2 h$. For the construction of our two-grid method in $d$ dimensions in accordance with the framework introduced in Section 2.3 , we let

$$
V:=S^{d}, \quad V_{c}:=\left(S_{c}\right)^{d} \subset V
$$

The prolongation $P: V_{c} \rightarrow V$ is the canonical embedding of the coarse tensor product spline space in the fine one. It can be represented as the $d$-fold tensor product of prolongations for the univariate spline spaces, $I: S_{c} \rightarrow S$.

The following result states that a robust approximation error estimate holds for the Galerkin projector to the coarse spline space. It was proved for $d=1$ and $d=2$ in [15]. We extend the proof to arbitrary dimensions in the Appendix.
Lemma 7. The $A$-orthogonal projector $T_{c}: S^{d} \rightarrow\left(S_{c}\right)^{d}$ satisfies the approximation error estimate

$$
\left\|\left(I-T_{c}\right) u\right\|_{L_{2}(\Omega)} \leq c h\|u\|_{A} \quad \forall u \in S^{d}
$$

with a constant $c$ which is independent of $h$ and $p$ (but may depend on $d$ ).
In the following subsections, we construct a smoother for the two-grid method on these nested spline spaces which leads to a robust and efficient iterative method.

### 4.1 A multigrid smoother based on subspace correction

In each of the $2^{d}$ subspaces $S_{\alpha} \subset S^{d}, \alpha \in\{0,1\}^{d}$, defined in Section 3.3. we prescribe a local, symmetric and positive definite smoothing operator $L_{\alpha}: S_{\alpha} \rightarrow S_{\alpha}^{\prime}$. The overall smoothing operator is then given by the additive subspace operator

$$
\begin{equation*}
L:=\sum_{\alpha} Q_{\alpha}^{\prime} L_{\alpha} Q_{\alpha}: S^{d} \rightarrow S^{d^{\prime}} \tag{10}
\end{equation*}
$$

from $S^{d}$ to its dual $S^{d^{\prime}}$, and its inverse has the form

$$
L^{-1}=\sum_{\alpha} L_{\alpha}^{-1} Q_{\alpha}^{\prime}: S^{d^{\prime}} \rightarrow S^{d}
$$

The assumptions of Theorem 1 for $L$, and thus the convergence of the two-grid method with such a smoother, can be guaranteed under simple assumptions on the subspace operators $L_{\alpha}$, as the following two lemmas show. The stability of the space splitting is crucial to both proofs. Although we do not explicitly use any results from the literature on subspace correction methods, we rely heavily on the ideas developed therein; cf., e.g., [23, 13].
Lemma 8. Assume that for every $\alpha \in\{0,1\}^{d}$, we have

$$
\begin{equation*}
\left\langle A v_{\alpha}, v_{\alpha}\right\rangle \leq c\left\langle L_{\alpha} v_{\alpha}, v_{\alpha}\right\rangle \quad \forall v_{\alpha} \in S_{\alpha} \tag{11}
\end{equation*}
$$

Then the subspace correction smoother satisfies

$$
\langle A v, v\rangle \leq c\langle L v, v\rangle \quad \forall v \in S^{d}
$$

Proof. Due to Theorem 6 and (11), we have

$$
\langle A v, v\rangle \leq c \sum_{\alpha}\left\langle A Q_{\alpha} v, Q_{\alpha} v\right\rangle \leq c \sum_{\alpha}\left\langle L_{\alpha} Q_{\alpha} v, Q_{\alpha} v\right\rangle=c\langle L v, v\rangle .
$$

Lemma 9. Assume that for every $\alpha \in\{0,1\}^{d}$, we have

$$
\begin{equation*}
\left\langle L_{\alpha} v_{\alpha}, v_{\alpha}\right\rangle \leq c\left\langle\left(A+h^{-2} M^{d}\right) v_{\alpha}, v_{\alpha}\right\rangle \quad \forall v_{\alpha} \in S_{\alpha} \tag{12}
\end{equation*}
$$

where $M^{d}: S^{d} \rightarrow S^{d^{\prime}}$ is the mass operator in the tensor product spline space. Then the subspace correction smoother satisfies

$$
\left\|\left(I-T_{c}\right) v\right\|_{L} \leq c\|v\|_{A} \quad \forall v \in S^{d}
$$

Proof. From $\sqrt{12}$, Theorem 6 and $L_{2}$-orthogonality, we obtain

$$
\langle L v, v\rangle \leq c \sum_{\alpha}\left\langle\left(A+h^{-2} M^{d}\right) Q_{\alpha} v, Q_{\alpha} v\right\rangle \leq c\left\langle\left(A+h^{-2} M^{d}\right) v, v\right\rangle
$$

Thus, it follows

$$
\left\|\left(I-T_{c}\right) v\right\|_{L}^{2} \leq c\left\|\left(I-T_{c}\right) v\right\|_{A}^{2}+c h^{-2}\left\|\left(I-T_{c}\right) v\right\|_{M^{d}}^{2} \leq c\|v\|_{A}^{2}
$$

where we used the stability of the coarse-grid projector and the coarse-grid approximation property Lemma 7.

### 4.2 Choice of the local smoothing operators

We now construct suitable local operators $L_{\alpha}$ which satisfy the assumptions of Lemma 8 and Lemma 9. In the two-dimensional case, the operator associated with the bilinear form (2) admits the representation

$$
A=K \otimes M+M \otimes K+M \otimes M
$$

in terms of the stiffness and mass operators for the univariate case. Restricting $A$ to a subspace $S_{\alpha}=S_{\alpha_{1}, \alpha_{2}}$, we obtain

$$
A_{\alpha}:=Q_{\alpha}^{\prime} A Q_{\alpha}=K_{\alpha_{1}} \otimes M_{\alpha_{2}}+M_{\alpha_{1}} \otimes K_{\alpha_{2}}+M_{\alpha_{1}} \otimes M_{\alpha_{2}}
$$

The inverse inequality in $S_{0}$ (Theorem 2 ) allows us to estimate

$$
K_{0} \leq \sigma M_{0}
$$

where $\sigma=12 h^{-2}$. We obtain subspace smoothers $L_{\alpha}$ by replacing $K_{0}$ by $\sigma M_{0}$,

$$
\begin{array}{ll}
A_{00} \leq(1+2 \sigma) M_{0} \otimes M_{0} & =: L_{00} \\
A_{01} \leq M_{0} \otimes\left((1+\sigma) M_{1}+K_{1}\right) & =: L_{01} \\
A_{10} \leq\left((1+\sigma) M_{1}+K_{1}\right) \otimes M_{0} & =: L_{10} \\
A_{11}=M_{1} \otimes M_{1}+K_{1} \otimes M_{1}+M_{1} \otimes K_{1}=: L_{11}
\end{array}
$$

where (11), the assumption of Lemma 8, holds by construction. It is easy to see that each $L_{\alpha}$ can be spectrally bounded from above by a constant times the matrix $Q_{\alpha}^{\prime}\left(A+h^{-2} M \otimes M\right) Q_{\alpha}$, which proves the assumption (12) of Lemma 9 . Using the statements of these two lemmas, Theorem 1 implies the two-grid convergence.

The same approach generalizes to higher dimensions, and we illustrate this in the threedimensional setting. Here, we have

$$
A=K \otimes M \otimes M+M \otimes K \otimes M+M \otimes M \otimes K+M \otimes M \otimes M
$$

Again, we define $A_{\alpha}$ as above and obtain the operators $L_{\alpha}$ by replacing $K_{0}$ by $\sigma M_{0}$,

$$
\begin{array}{lr}
A_{000} \leq(1+3 \sigma) M_{0} \otimes M_{0} \otimes M_{0} & =: L_{000}, \\
A_{001} \leq M_{0} \otimes M_{0} \otimes\left((1+2 \sigma) M_{1}+K_{1}\right) & =: L_{001}, \\
A_{010} \leq M_{0} \otimes\left((1+2 \sigma) M_{1}+K_{1}\right) \otimes M_{0} & =: L_{010}, \\
A_{100} \leq\left((1+2 \sigma) M_{1}+K_{1}\right) \otimes M_{0} \otimes M_{0} & =: L_{100}, \\
A_{011} \leq M_{0} \otimes\left((1+\sigma) M_{1} \otimes M_{1}+K_{1} \otimes M_{1}+M_{1} \otimes K_{1}\right) & =: L_{011}, \\
A_{110} \leq\left((1+\sigma) M_{1} \otimes M_{1}+K_{1} \otimes M_{1}+M_{1} \otimes K_{1}\right) \otimes M_{0} & =: L_{110}, \\
A_{101} \leq K_{1} \otimes M_{0} \otimes M_{1}+(1+\sigma) M_{1} \otimes M_{0} \otimes M_{1}+M_{1} \otimes M_{0} \otimes K_{1} & =: L_{101}, \\
A_{111}=M_{1} \otimes M_{1} \otimes M_{1}+K_{1} \otimes M_{1} \otimes M_{1}+M_{1} \otimes K_{1} \otimes M_{1}+M_{1} \otimes M_{1} \otimes K_{1}=: L_{111} .
\end{array}
$$

We point out that, whereas $L_{011}$ and $L_{110}$ permit a tensor product factorization, the operator $L_{101}$ cannot directly be factorized due to the ordering of the involved spaces. However, the tensor product space $S_{101}$ is isomorphic to $S_{011}$ by a simple swapping of the order of the involved tensor products. We exploit this in Section 5.3 below by a simple renumbering of the degrees of freedom in order to obtain an efficient method for inverting $L_{101}$.

It is clear that the rule of replacing $K_{0}$ by $\sigma M_{0}$ in each operator $A_{\alpha}$ to obtain $L_{\alpha}$ extends directly to arbitrary dimension $d$. By the same arguments as above, we see that the resulting subspace correction smoother satisfies the assumptions of Lemma 8 and Lemma 9 . Thus Theorem 1 shows that the resulting two-grid method converges robustly with respect to $h$ and $p$. We summarize this in the following theorem.

Theorem 10. For any $d \in \mathbb{N}$, there exist choices for $\tau$ and $\nu$, independent of $h$ and $p$, such that the two-grid method in $S^{d}$ with the smoother induced by the subspace operators $L_{\alpha}$ as constructed above converges with a rate $q<1$ which does not depend on the grid size $h$ or the spline degree $p$.

The robust convergence of the W-cycle multigrid method follows using standard arguments, cf. 14, 15].

## 5 Computational realization

In Section 4, we have proposed a smoother and shown that it leads to a robust two-grid method. In this section, we provide details on the realization of the method and show that it permits an efficient implementation.

### 5.1 Computation of a basis for $S_{0}$ and $S_{1}$

In order to be able to work with the space $S_{0}$ and its orthogonal complement, we require bases for them. The aim of this subsection is to provide an algorithm for computing such bases as linear combinations of B-splines.

Recall that the univariate spline space $S$ with $m$ knot spans of width $h=1 / m$, degree $p$ and maximum smoothness $C^{p-1}$ has dimension $n=m+p$. Let

$$
\mathcal{B}:=\left\{\varphi_{1}, \ldots, \varphi_{n}\right\}
$$

denote the normalized (i.e., satisfying a partition of unity, cf. [8) B-spline basis of $S$. We have $\operatorname{supp} \varphi_{j}=[(j-p-1) h, j h] \cap[0,1]$. All interior B-splines

$$
\mathcal{B}^{I}:=\left\{\varphi_{p+1}, \ldots, \varphi_{n-p}\right\}
$$

vanish with all their derivatives up to the $p-1$ st at the boundaries of the interval $[0,1]$ and therefore lie in $S_{0}$. (Here and in the following we assume that $p+1 \leq m$ such that $\mathcal{B}^{I}$ is nonempty.)

It remains to find linear combinations of the first and last $p$ B-splines which complete $\mathcal{B}^{I}$ to a basis of $S_{0}$. Recall that $u \in S$ lies in $S_{0}$ iff

$$
u^{(2 l+1)}(0)=u^{(2 l+1)}(1)=0 \quad \forall l \in \mathbb{N}_{0} \text { with } 2 l+1<p
$$

Consider first the left boundary. We need to satisfy $k:=\lfloor p / 2\rfloor$ conditions on the derivatives of the splines. Let

$$
\tilde{D}=\left(h^{2 i-1} \varphi_{j}^{(2 i-1)}(0)\right)_{i=1, \ldots, k, j=1, \ldots, p} \in \mathbb{R}^{k \times p}
$$

denote the matrix of the relevant B-spline derivatives at 0 , scaled with a suitable power of $h$ in order to avoid numerical instabilities. We pad $\tilde{D}$ with $p-k$ zero rows to obtain a square matrix $D \in \mathbb{R}^{p \times p}$. Computing the singular value decomposition (SVD), we obtain

$$
D=U \Sigma V^{\top}
$$

with $U, V \in \mathbb{R}^{p \times p}$ orthogonal and $\Sigma \in \mathbb{R}^{p \times p}$ being the diagonal matrix of singular values in descending order. By construction, $\Sigma$ contains $k$ nonzero and $p-k$ zero singular values. Therefore, the rightmost $p-k$ columns of $V$ span the kernel of $D$, and the linear combinations

$$
\mathcal{B}_{0}^{L}:=\left\{\sum_{i=1}^{p} V_{i, j} \varphi_{i}: j=p-k+1, \ldots, p\right\}
$$

lie in $S_{0}$. By the analogous procedure at the right boundary, we compute a set $\mathcal{B}_{0}^{R}$ of $p-k$ linear combinations of the last $p \mathrm{~B}$-splines. Then, the functions in the set

$$
\mathcal{B}_{0}:=\mathcal{B}_{0}^{L} \cup B^{I} \cup \mathcal{B}_{0}^{R}
$$

are by construction linearly independent and lie in $S_{0}$. Since $n_{0}:=\left|\mathcal{B}_{0}\right|=n-2 k=\operatorname{dim} S_{0}$, we have

$$
\operatorname{span} \mathcal{B}_{0}=S_{0}
$$

In practice, we collect the coefficients in a sparse block diagonal matrix

$$
P_{0}=\left[\begin{array}{ccc}
V^{L}[:, p-k+1: p] & & \\
& I_{n-2 p} & \\
& & V^{R}[:, p-k+1: p]
\end{array}\right] \in \mathbb{R}^{n \times n_{0}}
$$

where $V^{L}[:, p-k+1: p] \in \mathbb{R}^{p \times(p-k)}$ denotes the last $p-k$ columns of the matrix $V$ computed for the left boundary, analogously $V^{R}$ that for the right boundary, and $I_{d}$ is the $d \times d$ identity matrix. Then clearly, splines in $S_{0}$ can be uniquely represented in terms of the B-spline basis as

$$
u \in S_{0} \quad \Longleftrightarrow \quad \exists \underline{u} \in \mathbb{R}^{n_{0}}: u=\sum_{j=1}^{n}\left(P_{0} \underline{u}\right)_{j} \varphi_{j}
$$

Since the SVD produces an orthonormal basis, collecting the remaining columns of $V^{L}$ and $V^{R}$ in a second sparse block matrix

$$
P_{\perp}=\left[\begin{array}{cc}
V^{L}[:, 1: k] & 0 \\
0 & 0 \\
0 & V^{R}[:, 1: k]
\end{array}\right] \in \mathbb{R}^{n \times 2 k}
$$

satisfies $P_{0}^{\top} P_{\perp}=0$. In fact, the columns of the concatenation $\left[\begin{array}{ll}P_{0} & P_{\perp}\end{array}\right]$ form an orthonormal basis of $\mathbb{R}^{n}$. Let

$$
P_{1}:=\underline{M}^{-1} P_{\perp} \in \mathbb{R}^{n \times 2 k}
$$

where $\underline{M}$ denotes the $\mathcal{B}$-mass matrix. Note that $P_{1}$ is no longer sparse. Furthermore, let $\underline{u} \in \mathbb{R}^{n_{0}}$ and $\underline{v} \in \mathbb{R}^{2 k}$ with associated splines

$$
u=\sum_{j=1}^{n}\left(P_{0} \underline{u}\right)_{j} \varphi_{j}, \quad v=\sum_{j=1}^{n}\left(P_{1} \underline{v}\right)_{j} \varphi_{j} .
$$

By construction, $u \in S_{0}$. We have

$$
\langle u, v\rangle_{L_{2}(\Omega)}=\left\langle\underline{M} P_{0} \underline{u}, \underline{M}^{-1} P_{\perp} \underline{v}\right\rangle=\underline{u}^{\top} P_{0}^{\top} P_{\perp} \underline{v}=0 .
$$

Since this holds for all $u \in S_{0}, v$ lies in the $L_{2}$-orthogonal complement of $S_{0}$. All in all, we have constructed basis representations or "prolongation matrices"

$$
P_{0} \in \mathbb{R}^{n \times(n-2 k)}, \quad P_{1}=\underline{M}^{-1} P_{\perp} \in \mathbb{R}^{n \times 2 k}
$$

for $S_{0}$ and its $L_{2}$-orthogonal complement $S_{1}$, respectively.
For $d>1$, we let $\alpha \in\{0,1\}^{d}$ and introduce the Kronecker products

$$
P_{\alpha}:=P_{\alpha_{1}} \otimes \ldots \otimes P_{\alpha_{d}} \in \mathbb{R}^{n^{d} \times n_{\alpha}}
$$

where $n_{\alpha}=\operatorname{dim} S_{\alpha}$, which represent bases for the spaces $S_{\alpha}$ in terms of the coefficients of linear combinations of the tensor product B-spline basis $\mathcal{B}^{\otimes d}$.

### 5.2 Implementation of the subspace correction smoother

For any $\alpha \in\{0,1\}^{d}$, the matrices $P_{\alpha}$ as defined in Section 5.1 describe a basis for $S_{\alpha}$. Let $\underline{L}_{\alpha} \in \mathbb{R}^{n_{\alpha} \times n_{\alpha}}$ be the (symmetric and positive definite) matrix representation of $L_{\alpha}: S_{\alpha} \rightarrow S_{\alpha}^{\prime}$ (as defined in Section 4.2) with respect to that basis. Then the matrix representation of

$$
L^{-1}=\sum_{\alpha} L_{\alpha}^{-1} Q_{\alpha}^{\prime}=\sum_{\alpha} I_{S_{\alpha} \rightarrow S^{d}} L_{\alpha}^{-1} I_{S^{d^{\prime}} \rightarrow S_{\alpha}^{\prime}} Q_{\alpha}^{\prime} I_{S^{d^{\prime}} \rightarrow S_{\alpha}^{\prime}}
$$

is given by

$$
\begin{equation*}
\underline{L}^{-1}=\sum_{\alpha} P_{\alpha} \underline{L}_{\alpha}^{-1} P_{\alpha}^{\top} \underline{M} P_{\alpha} \underline{M}_{\alpha}^{-1} P_{\alpha}^{\top}=\sum_{\alpha} P_{\alpha} \underline{L}_{\alpha}^{-1} P_{\alpha}^{\top} \tag{13}
\end{equation*}
$$

where we used that the matrix representation of the embedding $I_{S_{\alpha} \rightarrow S^{d}}$ is $P_{\alpha}$ and the matrix representation of the $L_{2}$-projector $Q_{\alpha}$ is

$$
\underline{M}_{\alpha}^{-1} P_{\alpha}^{\top} \underline{M}, \quad \text { where } \quad \underline{M}_{\alpha}=P_{\alpha}^{\top} \underline{M} P_{\alpha}
$$

Hence (13) can be used to implement the subspace correction smoother using only the prolongation matrices $P_{\alpha}$ and a fast method for applying $\underline{L}_{\alpha}^{-1}$. It is never necessary to explicitly apply the $L_{2^{-}}$ projectors $Q_{\alpha}$. Furthermore, due to the use of additive subspace correction, the residual needs to be computed only once, and the individual subspace corrections may be done in parallel.

### 5.3 Inversion of the subspace operators

The final required algorithmic component is a fast method for applying the inverse of the local smoothing matrices $\underline{L}_{\alpha} \in \mathbb{R}^{n_{\alpha} \times n_{\alpha}}$. We illustrate this in the three-dimensional setting as described in Section 4.2, but the principles are the same regardless of dimension. A detailed discussion of the computational costs for arbitrary dimension is given in Section 5.4.

Interior space and face spaces. The interior space $S_{000}$ and the face spaces $S_{001}, S_{010}, S_{100}$ contain the complement space $S_{1}$ as a factor space at most once, and thus the matrices associated with their smoothing operators can be represented as Kronecker products of three one-dimensional discretization matrices, e.g.,

$$
\underline{L}_{000}=(1+3 \sigma) \underline{M}_{0} \otimes \underline{M}_{0} \otimes \underline{M}_{0}, \quad \underline{L}_{001}=\underline{M}_{0} \otimes \underline{M}_{0} \otimes\left((1+2 \sigma) \underline{M}_{1}+\underline{K}_{1}\right) .
$$

Here the symmetric matrices $\underline{M}_{\beta}, \underline{K}_{\beta} \in \mathbb{R}^{\operatorname{dim} S_{\beta} \times \operatorname{dim} S_{\beta}}, \beta \in\{0,1\}$, are the matrix representations of $M_{\beta}$ and $K_{\beta}$, respectively, with respect to the bases described by $P_{\beta}$ as computed in Section 5.1 above. For $\beta=0, \underline{M}_{\beta}$ and $\underline{K}_{\beta}$ have dimension $\mathcal{O}(n)$ and bandwidth $\mathcal{O}(p)$, whereas for $\beta=1$ they have dimension $\mathcal{O}(p)$ and are dense.

Since the Kronecker product can be inverted componentwise, we obtain, e.g.,

$$
\underline{L}_{001}^{-1}=\underline{M}_{0}^{-1} \otimes \underline{M}_{0}^{-1} \otimes\left((1+2 \sigma) \underline{M}_{1}+\underline{K}_{1}\right)^{-1} .
$$

Instead of computing this (dense) inverse explicitly, we employ the algorithm described by de Boor [7] for computing the application of a Kronecker product of matrices to a vector, given only routines for applying the individual Kronecker factors. For the latter, we use Cholesky factorization.

Edge spaces. The spaces $S_{011}, S_{110}, S_{101}$ contain the complement space $S_{1}$ as a factor twice. In $S_{011}$, the matrix to be inverted has the form

$$
\underline{L}_{011}=\underline{M}_{0} \otimes\left((1+\sigma) \underline{M}_{1} \otimes \underline{M}_{1}+\underline{K}_{1} \otimes \underline{M}_{1}+\underline{M}_{1} \otimes \underline{K}_{1}\right)
$$

It again has Kronecker product structure and can be inverted using the algorithm described in the previous case. The same holds for $S_{110}$.

In the case of the space $S_{101}$, the associated matrix

$$
\underline{L}_{101}=\underline{K}_{1} \otimes \underline{M}_{0} \otimes \underline{M}_{1}+(1+\sigma) \underline{M}_{1} \otimes \underline{M}_{0} \otimes \underline{M}_{1}+\underline{M}_{1} \otimes \underline{M}_{0} \otimes \underline{K}_{1}
$$

does not permit a Kronecker product factorization due to the order of the involved spaces. However, by a simple renumbering of the degrees of freedom, $S_{101}$ can be identified with $S_{011}$, and then $\underline{L}_{011}^{-1}$ can be applied as above.

Alternatively, the matrix $\underline{L}_{101}$ could be directly computed and inverted in its entirety using Cholesky factorization. This would exceed asymptotically (for $p \rightarrow \infty$ ) the computational costs derived in the following subsection, however this slowdown appears to be negligible in practice. For $d>3$, this shortcut seems no longer viable.

Corner space. The space $S_{111}$ is the tensor product of the three complement spaces and has dimension $\operatorname{dim}\left(S_{1}\right)^{3} \leq p^{3}$. The associated matrix

$$
\underline{L}_{111}=M_{1} \otimes M_{1} \otimes M_{1}+K_{1} \otimes M_{1} \otimes M_{1}+M_{1} \otimes K_{1} \otimes M_{1}+M_{1} \otimes M_{1} \otimes K_{1}
$$

is dense and is inverted by means of its Cholesky factorization.

### 5.4 Computational costs

We now study the computational complexity for applying the subspace correction smoother in the general $d$-dimensional setting. In our analysis, we ignore multiplicative constants which depend only on $d$. Repeatedly, we make use of the fact that the Cholesky factorization of a symmetric matrix of dimension $N$ and bandwidth $q$ can be computed in $\mathcal{O}\left(N q^{2}\right)$ operations, and its inverse can then be applied in $\mathcal{O}(N q)$ operations. If the matrix is not banded but dense, the factorization and inversion require $\mathcal{O}\left(N^{3}\right)$ and $\mathcal{O}\left(N^{2}\right)$ operations, respectively (cf. [12]).

By the renumbering of degrees of freedom described in Section 5.3, we can always rearrange the factor spaces such that we only need to consider spaces of the form

$$
\underbrace{S_{0} \otimes \ldots \otimes S_{0}}_{k \text { times }} \otimes \underbrace{S_{1} \otimes \ldots \otimes S_{1}}_{d-k \text { times }}
$$

The smoothing matrices to be inverted, constructed as in Section 5.3, have the form

$$
\underline{L}_{\{k, d-k\}}:=\underbrace{M_{0} \otimes \ldots \otimes \underline{M}_{0}}_{k \text { times }} \otimes \underline{X}_{d-k}
$$

where $\underline{X}_{j} \in \mathbb{R}^{\left(\operatorname{dim} S_{1}\right)^{j} \times\left(\operatorname{dim} S_{1}\right)^{j}}$ is a dense, symmetric matrix. Recall that $\operatorname{dim} S_{1} \leq p$.

Setup costs. The computation of the basis for $S_{0}$ and its $L_{2}$-orthogonal complement as described in Section 5.1 requires computing the SVD of two matrices of dimension $\mathcal{O}(p)$ as well as $\mathcal{O}(p)$ applications of the inverse of $\underline{M}$, which has dimension $n=m+p$ and bandwidth $\mathcal{O}(p)$, where $m$ is the number of subintervals. The costs for this step are thus $\mathcal{O}\left(p^{3}+n p^{2}\right)=\mathcal{O}\left(p^{3}+m p^{2}\right)$.

The one-dimensional mass matrix in $S_{0}, \underline{M}_{0}$, has dimension $\mathcal{O}(m)$ and bandwidth $\mathcal{O}(p)$ and thus requires $\mathcal{O}\left(m p^{2}\right)$ operations to factorize.

The matrices $\underline{X}_{j}, j=1, \ldots, d$, are dense and therefore require $\mathcal{O}\left(p^{3 j}\right)$ operations to factorize.
The overall setup costs are therefore $\mathcal{O}\left(m p^{2}+p^{3 d}\right)$.
Application costs. After factorization, the cost for applying the inverse $\underline{M}_{0}^{-1}$ is $\mathcal{O}(m p)$, and for $\underline{X}_{j}^{-1}$, it is $\mathcal{O}\left(p^{2 j}\right)$. To apply $\underline{L}_{\{k, d-k\}}^{-1}$ using the Kronecker product algorithm from [7], we need to perform $m^{d-1}$ applications of each of the $k$ factors $\underline{M}_{0}^{-1}$ and $m^{k}$ applications of $\underline{X}_{j}^{-1}$. Thus, the cost is $\mathcal{O}\left(k m^{d} p+m^{k} p^{2(d-k)}\right)$.

The inverse of $\underline{L}_{\{k, d-k\}}$ needs to be applied $\binom{d}{k}$ times since that is the number of multiindices $\alpha \in\{0,1\}^{d}$ which permute to $(0, \ldots, 0,1, \ldots, 1)$ with exactly $k$ leading zeros. The binomial coefficient satisfies $\binom{d}{k}=\mathcal{O}\left(2^{d} / \sqrt{d}\right)$ and in particular can be bounded from above by a constant which depends only on $d$. The overall cost for one application of the subspace correction smoother is then

$$
\sum_{k=0}^{d}\binom{d}{k} \mathcal{O}\left(k m^{d} p+m^{k} p^{2(d-k)}\right)=\mathcal{O}\left(m^{d} p+\max _{k=0, \ldots, d} m^{k} p^{2(d-k)}\right)=\mathcal{O}\left(m^{d} p+p^{2 d}\right)
$$

Overall costs. For $d \geq 2$, we have $m p^{2} \leq m^{2}+p^{4} \leq m^{d}+p^{2 d}$. Therefore, the overall costs for setting up and applying the smoother are bounded by

$$
\mathcal{O}\left(m^{d} p+p^{3 d}\right)
$$

Assuming $p^{2} \lesssim m$, the overall costs are asymptotically not more expensive than one application of the stiffness matrix, which has complexity $\mathcal{O}\left(n^{d} p^{d}\right)=\mathcal{O}\left(m^{d} p^{d}+p^{2 d}\right)$.

In a multigrid setting, assuming $\mathcal{O}(\log \mathfrak{m})$ levels with $m=\mathfrak{m}, \frac{\mathfrak{m}}{2}, \frac{\mathfrak{m}}{4}, \frac{\mathfrak{m}}{8} \ldots$ intervals per dimension, one obtains for $d \geq 2$ by summing up the overall costs of

$$
\mathcal{O}\left(\mathfrak{m}^{d} p+(\log \mathfrak{m}) p^{3 d}\right) \quad \text { and } \quad \mathcal{O}\left(\mathfrak{m}^{d} p+\mathfrak{m} p^{2 d}+(\log \mathfrak{m}) p^{3 d}\right)
$$

for smoothing in the V-cycle and the W-cycle, respectively. The full complexity including the costs for the exact coarse-grid solver and the intergrid transfers is asymptotically the same. Under mild assumptions on the relation between $p$ and $\mathfrak{m}$, again the overall effort is asymptotically not higher than that for one application of the stiffness matrix.

## 6 Numerical experiments

### 6.1 Experiments for the model problem

We solve the problem (1), i.e.,

$$
-\Delta u+u=f \quad \text { in } \Omega=(0,1)^{d}, \quad \partial_{n} u=0 \quad \text { on } \partial \Omega
$$

for $d=1,2,3$ with the right-hand side

$$
\begin{equation*}
f(x)=d \pi^{2} \prod_{j=1}^{d} \sin \left(\pi\left(x_{j}+\frac{1}{2}\right)\right) \tag{14}
\end{equation*}
$$

We perform a (tensor product) B-spline discretization using equidistant knot spans and maximumcontinuity splines for varying spline degrees $p$. We refer to the coarse discretization with only
one single interval as level $\ell=0$ and perform uniform, dyadic refinement to obtain the finer discretization levels $\ell$ with $2^{\ell d}$ elements and $h_{\ell}=2^{-\ell}$.

We set up a V-cycle multigrid method as described in Section 2.3 and using on each level the proposed smoother (10) as constructed in Section 4. We always use one pre- and one postsmoothing step with $\tau=1$. The parameter $\sigma$ was chosen as $\frac{1}{0.09} h^{-2}$ in $1 \mathrm{D}, \frac{1}{0.18} h^{-2}$ in 2D, and $\frac{1}{0.19} h^{-2}$ in 3D. In each test, the coarsest grid was chosen in such a way that the spaces $S_{0}$ on each higher level are non-empty, i.e., such that the smoother is well-defined. We perform tests both using the V-cycle multigrid method and a conjugate gradient solver preconditioned with one V-cycle. The iteration numbers required to reduce the $\ell^{2}$-norm of the initial residual by a factor of $10^{-8}$ for the $1 \mathrm{D}, 2 \mathrm{D}$ and 3D problem are given in Tables 133 , respectively.

Table 1: Iteration numbers: unit interval (1D)

|  | $\ell \backslash p$ | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 |
| :--- | :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| V-cycle | 9 | 33 | 34 | 34 | 33 | 33 | 33 | 32 | 31 | 31 | 31 | 28 | 28 | 29 |
|  | 8 | 33 | 34 | 34 | 32 | 33 | 33 | 31 | 30 | 30 | 31 | 28 | 28 | 27 |
|  | 7 | 33 | 34 | 34 | 32 | 33 | 33 | 31 | 28 | 30 | 29 | 28 | 25 | 26 |
| PCG | 9 | 13 | 13 | 13 | 13 | 13 | 13 | 13 | 13 | 12 | 12 | 12 | 12 | 12 |
|  | 8 | 13 | 13 | 13 | 13 | 13 | 13 | 12 | 12 | 12 | 12 | 12 | 12 | 11 |
|  | 7 | 13 | 13 | 13 | 13 | 13 | 12 | 12 | 12 | 12 | 11 | 11 | 11 | 11 |

Table 2: Iteration numbers: unit square (2D).

|  | $\ell \backslash p$ | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| ---: | :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| V-cycle | 8 | 38 | 39 | 39 | 39 | 38 | 38 | 37 | 37 | 36 |
|  | 7 | 38 | 39 | 39 | 38 | 38 | 37 | 36 | 36 | 34 |
|  | 6 | 38 | 38 | 38 | 37 | 37 | 35 | 34 | 34 | 32 |
|  | 5 | 36 | 37 | 34 | 34 | 32 | 30 | 28 | 26 | 24 |
|  | 8 | 14 | 14 | 14 | 14 | 14 | 14 | 14 | 14 | 13 |
|  | 7 | 14 | 14 | 14 | 14 | 14 | 14 | 14 | 13 | 13 |
|  | 6 | 14 | 14 | 14 | 14 | 14 | 13 | 13 | 13 | 12 |
|  | 5 | 14 | 14 | 13 | 13 | 13 | 12 | 11 | 11 | 10 |

Table 3: Iteration numbers: unit cube (3D).

|  | $\ell \backslash p$ | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :--- | ---: | ---: | ---: | ---: | ---: | ---: |
| V-cycle | 6 | 46 | 44 | 43 | 43 | 42 | 41 |
|  | 5 | 44 | 43 | 42 | 39 | 38 | 35 |
|  | 4 | 39 | 36 | 32 | 29 | 25 | 23 |
|  | 3 | 30 | 42 | 18 | 22 | 12 | 17 |
|  | 6 | 17 | 16 | 15 | 15 | 15 | 15 |
|  | 5 | 17 | 16 | 15 | 15 | 14 | 13 |
|  | 4 | 14 | 16 | 13 | 14 | 11 | 12 |
|  | 3 | 12 | 13 | 9 | 10 | 7 | 8 |

The method was implemented in $\mathrm{C}++$ based on the $\mathrm{G}+\mathrm{SMO}$ library ${ }^{1}$ which is developed in the framework of the National Research Network "Geometry + Simulation" at Johannes Kepler University, Linz.

[^1]We observe that the iteration numbers are robust with respect to both the discretization level $\ell$ (and thus $h$ ) and the spline degree $p$. They do increase with the space dimension $d$, but this dependence, which we have not fully analyzed, appears to be relatively mild. In particular, the 2D iteration numbers are significantly lower than those obtained using the boundary-corrected mass smoother in 15.

### 6.2 Experiments for non-trivial computational domains

We perform experiments with varying, matrix-valued diffusion coefficients on the non-trivial geometries shown in Fig. 3. The geometry map for the quarter annulus in the two-dimensional example is described exactly with NURBS, that for the three-dimensional object with B-splines. On these objects, we solve

$$
-\operatorname{div}(A(x) \nabla u(x))=f(x) \quad \text { in } \Omega
$$

with Dirichlet boundary conditions $g(x)$ on $\Gamma_{D}$ as indicated in Fig. 3 and homogeneous Neumann boundary conditions on the remaining part of the boundary. Furthermore, $f$ is given by (14) and the diffusion coefficient is given by

$$
A^{(2 \mathrm{D})}(x)=\left(\begin{array}{cc}
1+x_{1}^{2} & -x_{1} x_{2} \\
-x_{1} x_{2} & 1+x_{2}^{2}
\end{array}\right), \quad A^{(3 \mathrm{D})}(x)=\left(\begin{array}{ccc}
1+x_{1}^{2} & -\frac{1}{3} x_{1} x_{2} & -\frac{1}{3} x_{1} x_{3} \\
-\frac{1}{3} x_{1} x_{2} & 1+x_{2}^{2} & -\frac{1}{3} x_{2} x_{3} \\
-\frac{1}{3} x_{1} x_{3} & -\frac{1}{3} x_{2} x_{3} & 1+x_{3}^{2}
\end{array}\right)
$$




Figure 3: Computational domains for 2D and 3D example.
Table 4 gives the iteration numbers for a conjugate gradient method, preconditioned with one V-cycle of the proposed multigrid solver, where the multigrid solver was set up as solver for the model problem $-\Delta u+u=f$ on the parameter domain.

Obviously, the condition number of the preconditioned system depends only on the geometry transformation, the diffusion coefficient and on the contraction number of the multigrid method (as a solver for the model problem on the parameter domain). All of these quantities are independent of the grid size and the polynomial degree $p$. This is reflected in the numerical results, which are robust in those two parameters.

## Appendix

The aim of this section is to prove Lemma 7, an approximation result for the coarse spline space Galerkin projector in $d$ dimensions. It was shown in [15] for $d=1$ and $d=2$, and here we extend it to arbitrary dimensions by induction.

Before we give the proof, we need an auxiliary lemma which is a variant of the Aubin-Nitsche duality argument in a finite-dimensional Hilbert space $V$. By the choice of a suitable basis, we

Table 4: Iteration numbers for the nontrivial 2D (top) and 3D (bottom) domains as shown in Fig. 3

| $\ell \backslash p$ | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 8 | 53 | 55 | 56 | 56 | 55 | 55 | 55 | 54 | 54 |
| 7 | 52 | 53 | 54 | 53 | 53 | 52 | 51 | 50 | 51 |
| 6 | 47 | 50 | 50 | 48 | 48 | 48 | 46 | 46 | 45 |
| 5 | 43 | 45 | 45 | 44 | 44 | 41 | 41 | 40 | 41 |


| $\ell \backslash p$ | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: |
| 5 | 87 | 90 | 91 | 90 | 89 | 90 |
| 4 | 73 | 76 | 76 | 79 | 81 | 83 |
| 3 | 55 | 61 | 66 | 67 | 72 | 75 |

can identify $V$ with $\mathbb{R}^{n}$, and operators $A$ on $V$ with matrices. We use this matrix representation implicitly in the following, and operations like $A^{1 / 2}$ and $A^{\top}$ are to be understood in the matrix sense.

Lemma 11. Let $A$ and $M$ be self-adjoint and positive definite linear operators on $V, T: V \rightarrow$ $W \subset V$ an $A$-orthogonal projector, and $\theta>0$. Then, the statements

$$
\begin{equation*}
\|T u\|_{M} \leq \theta\|u\|_{A} \quad \forall u \in V \quad \text { and } \quad\|T u\|_{A} \leq \theta\|u\|_{A M^{-1} A} \quad \forall u \in V \tag{15}
\end{equation*}
$$

are equivalent.
Proof. We first observe that the statements in are equivalent to

$$
\begin{equation*}
\left\|M^{1 / 2} T A^{-1 / 2}\right\| \leq \theta \quad \text { and } \quad\left\|A^{1 / 2} T A^{-1} M^{1 / 2}\right\| \leq \theta \tag{16}
\end{equation*}
$$

respectively. Since $T$ is self-adjoint in the scalar product $(\cdot, \cdot)_{A}, A T=T^{\top} A$ and further

$$
\begin{equation*}
T A^{-1}=A^{-1} T^{\top} \tag{17}
\end{equation*}
$$

hold. Using (17) as well as the self-adjointness of $M$ and $A$, we obtain

$$
\left\|M^{1 / 2} T A^{-1 / 2}\right\|=\left\|M^{1 / 2} A^{-1} T^{\top} A^{1 / 2}\right\|=\left\|\left(A^{1 / 2} T A^{-1} M^{1 / 2}\right)^{\top}\right\|=\left\|A^{1 / 2} T A^{-1} M^{1 / 2}\right\|
$$

This proves that the two statements in (16) and, consequently, those in (15) are equivalent.
Proof of Lemma 7. Within this proof, we denote the dimensions explicitly and use a recursive representation,

$$
\begin{array}{ll}
M_{1}:=M, & A_{1}:=K+M \\
M_{d}:=M_{d-1} \otimes M, & A_{d}:=A_{d-1} \otimes M+M_{d-1} \otimes K
\end{array}
$$

Furthermore we let $T_{d}$ denote the $A_{d}$-orthogonal projector into $\left(S_{c}\right)^{d}$.
In [15], the desired result was proved for $d=1$, namely,

$$
\begin{equation*}
\left\|\left(I-T_{1}\right) u\right\|_{M_{1}} \leq c h\|u\|_{A_{1}} \quad \forall u \in S \tag{18}
\end{equation*}
$$

By Lemma 11, this is equivalent to

$$
\begin{equation*}
\left\|\left(I-T_{1}\right) u\right\|_{A_{1}} \leq c h\|u\|_{A_{1} M_{1}^{-1} A_{1}} \quad \forall u \in S \tag{19}
\end{equation*}
$$

Stability of the $A_{1}$-orthogonal projector means that

$$
\begin{equation*}
\left\|\left(I-T_{1}\right) u\right\|_{A_{1}} \leq\|u\|_{A_{1}} \quad \forall u \in S \tag{20}
\end{equation*}
$$

We now show the desired result using induction. Assume that we have already shown

$$
\begin{equation*}
\left\|\left(I-T_{d-1}\right) u\right\|_{M_{d-1}} \leq c h\|u\|_{A_{d-1}} \quad \forall u \in S^{d-1} \tag{21}
\end{equation*}
$$

for some $d>1$. Using Lemma 11, this implies

$$
\begin{equation*}
\left\|\left(I-T_{d-1}\right) u\right\|_{A_{d-1}} \leq c h\|u\|_{A_{d-1} M_{d-1}^{-1} A_{d-1}} \quad \forall u \in S^{d-1} \tag{22}
\end{equation*}
$$

Stability of the $A_{d-1}$-orthogonal projector means that

$$
\begin{equation*}
\left\|\left(I-T_{d-1}\right) u\right\|_{A_{d-1}} \leq\|u\|_{A_{d-1}} \quad \forall u \in S^{d-1} \tag{23}
\end{equation*}
$$

Using equations (18) and the fact that the operator norm of a tensor product is the product of the individual operator norms, we obtain for all $u \in S^{d}$

$$
\begin{aligned}
& \left\|\left(I-T_{d-1}\right) \otimes\left(I-T_{1}\right) u\right\|_{A_{d-1} \otimes M_{1}+M_{d-1} \otimes A_{1}} \leq c h\|u\|_{A_{d-1} \otimes A_{1}}, \\
& \left\|\left(I-T_{d-1}\right) \otimes I u\right\|_{A_{d-1} \otimes M_{1}+M_{d-1} \otimes A_{1}} \leq c h\|u\|_{A_{d-1} M_{d-1}^{-1} A_{d-1} \otimes M_{1}+A_{d-1} \otimes A_{1}}, \\
& \left\|I \otimes\left(I-T_{1}\right) u\right\|_{A_{d-1} \otimes M_{1}+M_{d-1} \otimes A_{1}} \leq c h\|u\|_{A_{d-1} \otimes A_{1}+M_{d-1} \otimes A_{1} M_{1}^{-1} A_{1}} .
\end{aligned}
$$

Since $I-T_{d-1} \otimes T_{1}=\left(I-T_{d-1}\right) \otimes I+I \otimes\left(I-T_{1}\right)-\left(I-T_{d-1}\right) \otimes\left(I-T_{1}\right)$, this implies using the triangle inequality

$$
\left\|\left(I-T_{d-1} \otimes T_{1}\right) u\right\|_{A_{d-1} \otimes M_{1}+M_{d-1} \otimes A_{1}} \leq c h\|u\|_{A_{d-1} M_{d-1}^{-1} A_{d-1} \otimes M_{1}+A_{d-1} \otimes A_{1}+M_{d-1} \otimes A_{1} M_{1}^{-1} A_{1}} .
$$

As the norm on the left-hand side is bounded from below by $\|\cdot\|_{A_{d}}$ and the norm on the right-hand side is bounded from above by $c\|\cdot\|_{A_{d} M_{d}^{-1} A_{d}}$, we further obtain

$$
\left\|\left(I-T_{d-1} \otimes T_{1}\right) u\right\|_{A_{d}} \leq c h\|u\|_{A_{d} M_{d}^{-1} A_{d}} \quad \forall u \in S^{d}
$$

Both $T_{d-1} \otimes T_{1}$ and $T_{d}$ are projectors into $\left(S_{c}\right)^{d}$. Since the latter projector produces the best approximation in the $A_{d}$-norm, we have

$$
\left\|\left(I-T_{d}\right) u\right\|_{A_{d}} \leq c h\|u\|_{A_{d} M_{d}^{-1} A_{d}} \quad \forall u \in S^{d}
$$

which, by Lemma 11, is equivalent to the desired result

$$
\left\|\left(I-T_{d}\right) u\right\|_{M_{d}} \leq c h\|u\|_{A_{d}} \quad \forall u \in S^{d}
$$

## Acknowledgments

We gratefully acknowledge the discussions with Ludmil Zikatanov (Penn State University) which were instrumental in developing some of the ideas underlying this work.

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[^0]:    *The first author was supported by the National Research Network "Geometry + Simulation" (NFN S117, 2012-2016), funded by the Austrian Science Fund (FWF)

[^1]:    ${ }^{1}$ http://www.gs.jku.at/gismo

