

Digraphs and variable degeneracy

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Abstract

Let D be a digraph, let $p \geq 1$ be an integer, and let $f : V(D) \rightarrow \mathbb{N}_0^p$ be a vector function with $f = (f_1, f_2, \dots, f_p)$. We say that D has an f -partition if there is a partition (D_1, D_2, \dots, D_p) into induced subdigraphs of D such that for all $i \in [1, p]$, the digraph D_i is weakly f_i -degenerate, that is, in every non-empty subdigraph D' of D_i there is a vertex v such that $\min\{d_{D'}^+(v), d_{D'}^-(v)\} < f_i(v)$. In this paper, we prove that the condition $f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$ is almost sufficient for the existence of an f -partition and give a full characterization of the bad pairs (D, f) . Moreover, we describe a polynomial time algorithm that (under the previous conditions) either verifies that (D, f) is a bad pair or finds an f -partition. Among other applications, this leads to a generalization of Brooks' Theorem as well as the list-version of Brooks' Theorem for digraphs, where a coloring of digraph is a partition of the digraph into acyclic induced subdigraphs. We furthermore obtain a result bounding the s -degenerate chromatic number of a digraph in terms of the maximum of maximum in-degree and maximum out-degree.

AMS Subject Classification: 05C20

Keywords: Digraph coloring, Digraph degeneracy, Brooks' theorem, acyclic digraph

Most of our terminology is defined as in [6] and similar to the papers [4, 5] (see also Section 2). Throughout this paper, for $0 \leq \ell \leq k$, let $[\ell, k] = \{i \in \mathbb{N}_0 \mid \ell \leq i \leq k\}$.

1 Introduction

Even though most people would define a k -coloring of an (undirected) graph G as a function that assigns colors from a color set of cardinality k to the vertices of G such that the same-colored vertices induce edgeless subgraphs of G , this is nothing else than a partition of G into vertex disjoint induced subgraphs (G_1, G_2, \dots, G_k) such that G_i is

*Research supported by the Independent Research Fond Denmark under grant number DFF 7014-00037B

edgeless for all $i \in [1, k]$. Naturally, both ways how to regard a coloring have their benefits. In this paper, we examine a more general approach regarding the latter definition that will allow us to obtain various well-known coloring results. But first of all, we need to clarify what digraph coloring refers to. A **coloring** and **k -coloring** of a digraph D is a function $\varphi : V(D) \rightarrow [1, k]$ such that each color class $\varphi^{-1}(i)$ induces an acyclic subdigraph of D , that is, a subdigraph that does not contain any directed cycle. The **dichromatic number** $\vec{\chi}(D)$ of a digraph D is the least integer k such that D admits a k -coloring. This digraph coloring concept was originally introduced by Neumann-Lara [30] in 1982; however, it took over twenty years until it was rediscovered by Mohar [27] in 2003. Ever since, it has attracted much attention amongst graph theorists (see, e.g., [2, 3, 4, 5, 16, 17, 18, 19, 20, 22, 28, 29]). Although this coloring concept might not seem intuitive at first sight, there are various factors stressing why it is especially reasonable. First of all, the dichromatic number of a bidirected graph and the chromatic number of its underlying (undirected) graph coincide. Consequently, many theorems on digraph coloring are generalizations of theorems on coloring of undirected graphs. Moreover, it has been shown that plenty of well-known theorems in graph coloring indeed have digraph counterparts. For example, Harutyunyan and Mohar [20] proved that there exist digraphs D of maximum total degree Δ and arbitrary large digirth such that $\vec{\chi}(D) \geq \frac{c\Delta}{\log \Delta}$ for some constant c , thereby obtaining the analogue of a famous result of Bollobás [9], respectively Kostochka and Mazurova [23]. Moreover, Andres and Hochstättler [1] obtained the digraph analogue of Chudnovsky, Robertson, Seymour, and Thomas' celebrated Strong Perfect Graph Theorem [12]. In this paper, we will mainly focus on the analogue to Brooks' famous theorem [10], which was discovered by MOHAR in 2010 [29]. Note that, given a digraph D , its maximum out-degree (respectively in-degree) is denoted by $\Delta^+(D)$ (respectively $\Delta^-(D)$).

Theorem 1 (Mohar, 2010) *Let D be a connected digraph. Then, D satisfies $\vec{\chi}(D) \leq \max\{\Delta^-(D), \Delta^+(D)\} + 1$ and equality holds if and only if D is*

- (a) *a directed cycle of length ≥ 2 , or*
- (b) *a bidirected cycle of odd length ≥ 3 , or*
- (c) *a bidirected complete graph.*

◇

In fact, it turns out that it is possible to obtain a choosability version of Brooks' Theorem for digraphs, *i.e.*, a version regarding list-colorings. Given a digraph D , a color set Γ , and a function $L : V(D) \rightarrow 2^\Gamma$ (a so-called **list-assignment**), an **L -coloring** of D is a function $\varphi : V(D) \rightarrow \Gamma$ such that $\varphi(v) \in L(v)$ for all $v \in V(D)$ and $D[\varphi^{-1}(\alpha)]$ contains no directed cycle for each $\alpha \in \Gamma$ (if such a coloring exists, we say that D is **L -colorable**). Harutyunyan and Mohar [19] proved the following, thereby extending a well-known theorem of Erdős, Rubin and Taylor [14] for undirected graphs. Note that a **block** B of a digraph is a maximal connected subdigraph that does not contain a separating vertex, *i.e.*, the underlying graph $G(B)$ of B is a block of the underlying graph $G(D)$ of D . By $\mathcal{B}(D)$ we denote the set of blocks of a digraph D . For $v \in V(D)$, $\mathcal{B}_v(D)$ denotes the set of blocks of D containing v .

Theorem 2 (Harutyunyan and Mohar, 2011) *Let D be a connected digraph, and let L be a list-assignment such that $|L(v)| \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$. Suppose that D is not L -colorable. Then, the following statements hold:*

- (a) D is Eulerian and $|L(v)| = d_D^+(v) = d_D^-(v)$ for all $v \in V(D)$.
- (b) If $B \in \mathcal{B}(D)$, then B is either a directed cycle of length ≥ 2 , a bidirected complete graph, or a bidirected cycle of odd length ≥ 5 .
- (c) For each $B \in \mathcal{B}(D)$ there is a set Γ_B of $\Delta^+(B)$ colors such that for every $v \in V(D)$, the sets Γ_B with $B \in \mathcal{B}_v(D)$ are pairwise disjoint and $L(v) = \bigcup_{B \in \mathcal{B}_v(D)} \Gamma_B$. \diamond

In the present paper, we will obtain a generalization of the two previously mentioned results by examining degenerate digraphs. The concept of digraph degeneracy was introduced by Bokal et al. [7] in 2004. Given a positive integer k , a digraph D is **weakly k -degenerate** if every non-empty subdigraph D' contains a vertex v with $\min\{d_{D'}^+(v), d_{D'}^-(v)\} < k$. As a consequence, a digraph is acyclic if and only if it is weakly 1-degenerate and so a coloring of a digraph coincides with a partition of the digraph into induced subdigraphs which all are weakly 1-degenerate. We shall extend this definition to the case of variable degeneracy, based on the model of Borodin, Kostochka, and Toft [8] for undirected graphs. Let D be a digraph and let $h : V(D) \rightarrow \mathbb{N}_0$ be a function. Then, D is **weakly h -degenerate** if every non-empty subdigraph D' contains a vertex v with $\min\{d_{D'}^+(v), d_{D'}^-(v)\} < h(v)$. Clearly, if $h \equiv k$ is the constant function, then D is weakly h -degenerate if and only if D is weakly k -degenerate.

We will connect the concept of degeneracy with partitions of digraphs. A **partition** and **p -partition** of a digraph D is a sequence (D_1, D_2, \dots, D_p) of pairwise disjoint induced subdigraphs of D such that $V(D) = V(D_1) \cup V(D_2) \cup \dots \cup V(D_p)$. Now let $f : V(D) \rightarrow \mathbb{N}_0^p$ be a vector function. Then we denote by f_i the i -th coordinate of f , i.e. $f = (f_1, f_2, \dots, f_p)$. Then, an **f -partition** of D is a p -partition (D_1, D_2, \dots, D_p) of D such that D_i is weakly f_i -degenerate for $i \in [1, p]$. If D admits an f -partition, then we also say that D is **f -partitionable**. The main aim of this paper is to determine under which degree conditions D admits an f -partition. First of all, let us motivate why this is worthwhile considering. To this end, let D be a digraph and let L be a list-assignment for D . Moreover, let $\Gamma = \bigcup_{v \in V(D)} L(v)$ be the set of all colors appearing in some list. By renaming the colors if necessary we may assume $\Gamma = [1, p]$. Let $f : V(D) \rightarrow \mathbb{N}_0^p$ be the vector function with

$$f_i(v) = \begin{cases} 1 & \text{if } i \in L(v), \text{ and} \\ 0 & \text{if } i \notin L(v) \end{cases}$$

for $i \in [1, p]$ (see also Figure 1).

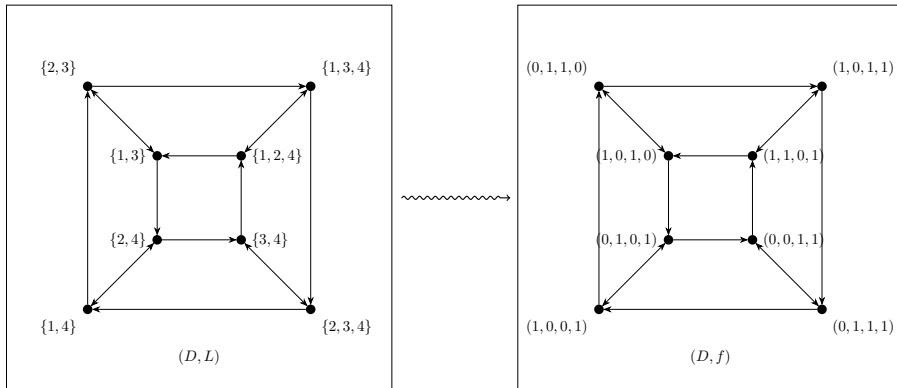


Figure 1: Transforming a list-assignment into a vector function

If D is L -colorable, then it is easy to confirm that the partition (D_1, D_2, \dots, D_p) of D where D_i is the subdigraph of D induced by vertices of color i is indeed an f -partition of D . Conversely, if D has an f -partition (D_1, D_2, \dots, D_p) , then setting $\varphi(v) = i$ if $v \in V(D_i)$ leads to an L -coloring of D . Hence, D is L -colorable if and only if D has an f -partition and so the task of finding an f -partition generalizes the usual list-coloring problem (and, of course, the usual coloring problem, too).

Since deciding whether the dichromatic number is at most two is already NP-hard (see [7]), it is pointless to try to determine whether a digraph D is f -partitionable in general. Instead, we need to find a reasonable condition for the function f that would allow a digraph satisfying this condition to be f -partitionable. Note that in the above transformation, it follows from the definition of f that

$$|L(v)| = f_1(v) + f_2(v) + \dots + f_p(v)$$

for all $v \in V(D)$. Thus, Theorem 2 suggests that

$$f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_D^+(v), d_D^-(v)\}$$

for all $v \in V(D)$ might be the right condition to investigate. Indeed, we will prove that this condition is always sufficient for the existence of an f -partition, unless (D, f) belongs to the following, recursively defined class of configurations. Clearly, D admits an f -partition if and only if each connected component of D has one and, hence, it suffices to examine connected digraphs.

Let D be a connected digraph, let $p \geq 1$, and let $f : V(D) \rightarrow \mathbb{N}_0^p$ be a vector function. We say that (D, f) is a **hard pair** and that D is **f -hard** if one of the following four conditions hold.

(H1) D is a block, D is Eulerian, and there exists an index $j \in [1, p]$ such that

$$f_i(v) = \begin{cases} d_D^+(v) = d_D^-(v) & \text{if } i = j, \text{ and} \\ 0 & \text{otherwise} \end{cases}$$

for all $i \in [1, p]$ and for each $v \in V(D)$. In this case, we say that (D, f) is a hard pair of type **(M)**.

(H2) D is a bidirected complete graph and there are integers $n_1, n_2, \dots, n_p \geq 0$ with at least two n_i different from zero such that $n_1 + n_2 + \dots + n_p = |D| - 1$ and that

$$f(v) = (n_1, n_2, \dots, n_p) \quad \text{for all } v \in V(D).$$

In this case, we say that (D, f) is a hard pair of type **(K)**.

(H3) D is a bidirected cycle of odd length and there are two indices $k \neq \ell$ from the set $[1, p]$ such that

$$f_i(v) = \begin{cases} 1 & \text{if } i \in \{k, \ell\}, \text{ and} \\ 0 & \text{otherwise} \end{cases}$$

for all $i \in [1, p]$ and for each $v \in V(D)$. In this case, we say that (D, f) is a hard pair of type **(C)**.

(H4) There are two hard pairs (D^1, f^1) and (D^2, f^2) with $f^j : V(D^j) \rightarrow \mathbb{N}_0^p$ for $j \in \{1, 2\}$ such that D is obtained from D^1 and D^2 by identifying two vertices $v^1 \in V(D^1)$ and $v^2 \in V(D^2)$ to a new vertex v . Furthermore, for $w \in V(D)$, it holds that

$$f(w) = \begin{cases} f^1(w) & \text{if } w \in V(D^1) \setminus \{v^1\}, \\ f^2(w) & \text{if } w \in V(D^2) \setminus \{v^2\}, \\ f^1(v^1) + f^2(v^2) & \text{if } w = v. \end{cases}$$

In this case we say that (D, f) is obtained from (D^1, f^1) and (D^2, f^2) by **merging** v^1 and v^2 to v .

In order to develop a better feeling of how hard pairs may look like, we refer the reader to Figure 2. The main result of this paper is the following.

Theorem 3 *Let D be a connected digraph, let $p \geq 1$ be an integer, and let $f : V(D) \rightarrow \mathbb{N}_0^p$ be a vector function such that $f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$. Then, D is not f -partitionable if and only if (D, f) is a hard pair. \diamond*

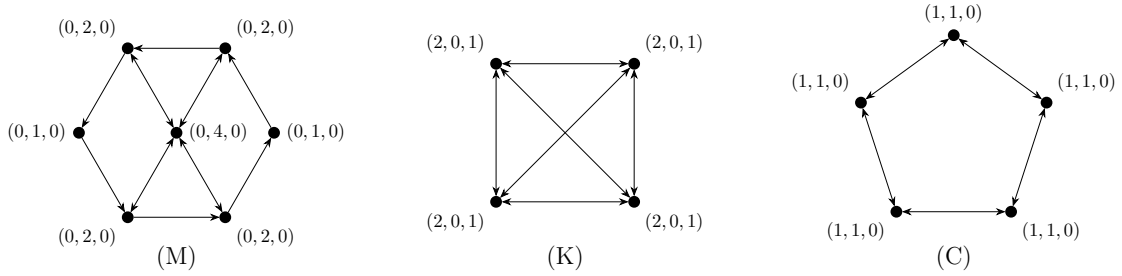


Figure 2: Examples of hard pairs

2 Basic Terminology

Let $D = (V(D), A(D))$ be a digraph, where $V(D)$ is the **set of vertices** of D and $A(D)$ is the **set of arcs** of D . The **order** $|D|$ of D is the size of $V(D)$. The digraphs in this paper do not have loops nor parallel arcs; however, there may be two arcs in opposite directions between two vertices (in this case we say that the arcs are **opposite**). We denote by uv the arc whose **initial vertex** is u and whose **terminal vertex** is v . Two vertices u, v are **adjacent** if at least one of uv and vu belongs to $A(D)$. If u and v are adjacent, we also say that u is a **neighbor** of v and vice versa. If $uv \in A(D)$, then v is called an **out-neighbor** of u and u is called an **in-neighbor** of v . Given a digraph D and a vertex set X , we denote by $D[X]$ the subdigraph of D **induced** by the vertex set X , that is, $V(D[X]) = X$ and $A(D[X]) = \{uv \in A(D) \mid u, v \in X\}$. A digraph D' is said to be an induced subdigraph of D if $D' = D[V(D')]$. As usual, if X is a subset of $V(D)$, we define $D - X = D[V(D) \setminus X]$. If $X = \{v\}$ is a singleton, we use $D - v$ rather than $D - \{v\}$. The **out-degree** of a vertex $v \in V(D)$, denoted $d_D^+(v)$, is the number of arcs whose initial vertex is v . Similarly, the **in-degree** of v , denoted $d_D^-(v)$ is number of arcs whose terminal

vertex is v . A vertex $v \in V(D)$ is **Eulerian** if $d_D^+(v) = d_D^-(v)$. Moreover, the digraph D is **Eulerian** if every vertex of D is Eulerian.

The **underlying** graph $G(D)$ of D is the simple undirected graph with $V(G(D)) = V(D)$ and $\{u, v\} \in E(G(D))$ if and only if at least one of uv and vu belongs to $A(D)$. A **component** of a digraph D is a connected component of $G(D)$ and D is **connected** if it has precisely one component. A **separating vertex** of a connected digraph D is a vertex $v \in V(D)$ such that $D - v$ is not connected. Furthermore, a **block** of D is a maximal subdigraph D' of D such that D' has no separating vertex. By $\mathcal{B}(D)$, we denote the **set of blocks** of D . Moreover, if $v \in V(D)$, we denote by $\mathcal{B}_v(D)$ the set of blocks of D containing v . A **bidirected** graph is a digraph that can be obtained from an undirected (simple) graph G by replacing each edge by two opposite arcs, we denote it by $D(G)$.

3 Proof of the main result

In order to prove Theorem 3, we need two classic results for undirected graphs. The first one is an easy consequence of Menger's Theorem and usually referred to as the Fan Lemma (see, e.g. [13, Corollary 3.3.4]). The second lemma is due to Gallai [15, Satz 1.9]. Recall that a **chord** of a cycle C in an undirected graph G is an edge of G between vertices of C that does not belong to the edges of C .

Lemma 4 (The Fan Lemma) *Let G be a k -connected graph, let $v \in V(G)$, and let $X \subseteq V(G) \setminus \{v\}$ be a set of cardinality at least k . Then, there are k paths from v to vertices of X whose only common vertex is v and whose only intersection with X are the respective end-vertices.* \diamond

Lemma 5 (Gallai, 1963) *If G is a graph in which each even cycle has at least two chords, then every block of G is a complete graph or an odd cycle.* \diamond

The proof of Theorem 3 is divided into two parts. In the first part, we obtain some properties of hard pairs and prove that hard pairs are not f -partitionable. The proof of the next proposition follows easily from the definition of hard pair and can be done via induction on the number of blocks of D .

Proposition 6 *Let D be a connected digraph, let $p \geq 1$, and let $f : V(D) \rightarrow \mathbb{N}_0^p$ be a vector function such that D is f -hard. Then, for each $B \in \mathcal{B}(D)$ there is a uniquely determined function $f_B : V(B) \rightarrow \mathbb{N}_0^p$ such that the following statements hold:*

- (a) (B, f_B) is a hard pair of type (M), (K), or (C).
- (b) $f(v) = \sum_{B \in \mathcal{B}_v(D)} f_B(v)$ for all $v \in V(D)$. In particular, $f_B(v) = f(v)$ for all non-separating vertices v of D belonging to B . \diamond

Proposition 7 *Let D be a connected digraph, and let $f : V(D) \rightarrow \mathbb{N}_0^p$ be a vector function with $p \geq 1$. If (D, f) is a hard pair, then the following statements hold:*

- (a) $f_1(v) + f_2(v) + \dots + f_p(v) = d_D^+(v) = d_D^-(v)$ for all $v \in V(D)$. As a consequence, D is Eulerian.
- (b) D is not f -partitionable. \diamond

Proof: Statement (a) follows from an easy induction on the number of blocks of D . Furthermore, if D is a block, then (D, f) is of type (M), (K), or (C) and it is an easy exercise to check that D is indeed not f -partitionable. To show the reader how it may be done, suppose that (D, f) is of type (K), *i.e.*, $D = D(K_n)$ and there are integers $n_1, n_2, \dots, n_p \geq 0$ with at least two n_i different from zero such that $n_1 + n_2 + \dots + n_p = n - 1$ and that $f(v) = (n_1, n_2, \dots, n_p)$ for all $v \in V(D)$. As D is a bidirected complete graph, an induced subdigraph D_i of D is weakly f_i -degenerate if and only if $|D_i| \leq n_i$. Consequently, if there exists an f -partition (D_1, D_2, \dots, D_p) , we have

$$n = |D| = |D_1| + |D_2| + \dots + |D_p| \leq n_1 + n_2 + \dots + n_p = n - 1,$$

which is impossible. This shows that (D, f) is not of type (K).

Now assume that D is not a block and let (D, f) be a minimal counter-example, *i.e.*, (D, f) is a hard pair, D admits an f -partition, and $|D|$ is minimum with respect to the previous two conditions. As D is not a block it follows with (H4) that there are two hard pairs (D^1, f^1) and (D^2, f^2) with $|D^j| < |D|$ for $j \in \{1, 2\}$ such that (D, f) is obtained from (D^1, f^1) and (D^2, f^2) by merging vertices $v^j \in V(D^j)$ to a new vertex v . For the sake of readability, we use v below for $v^1 = v^2 = v$. By the choice of (D, f) , the digraph D^j is not f^j -partitionable for $j \in \{1, 2\}$. Now let (D_1, D_2, \dots, D_p) be an f -partition of D and let $D_i^j = D^j \cap D_i$ for $j \in \{1, 2\}$ and $i \in [1, p]$. By symmetry, we may assume that $v \in V(D_1)$. Then, D_i^j is strictly f_i^j -degenerate for all $i \in [2, p]$ and $j \in \{1, 2\}$ (as $D_i^j \subseteq D_i$ and $f_i^j(w) = f_i(w)$ for all $w \in V(D_i^j)$). As D^j is not f^j -partitionable, it follows that D_1^j is not f_1^j -partitionable for $j \in \{1, 2\}$ and so there are non-empty subdigraphs $\tilde{D}^j \subseteq D_1^j$ with $\min\{d_{\tilde{D}^j}^+(w), d_{\tilde{D}^j}^-(w)\} \geq f_1^j(w)$ for all $w \in V(\tilde{D}^j)$. Let $\tilde{D} = \tilde{D}^1 \cup \tilde{D}^2$. If $v \notin \tilde{D}$, then \tilde{D} is the disjoint union of \tilde{D}^1 and \tilde{D}^2 and we clearly have $\min\{d_{\tilde{D}}^+(w), d_{\tilde{D}}^-(w)\} \geq f_1(w)$ for all $w \in V(\tilde{D})$. If $v \in \tilde{D}$, we obtain that

$$f_1(v) = f_1^1(v) + f_1^2(v) \leq \min\{d_{\tilde{D}^1}^+(v), d_{\tilde{D}^1}^-(v)\} + \min\{d_{\tilde{D}^2}^+(v), d_{\tilde{D}^2}^-(v)\} \leq \min\{d_{\tilde{D}}^+(v), d_{\tilde{D}}^-(v)\}.$$

Consequently, \tilde{D} is a subdigraph of D_1 with $\min\{d_{\tilde{D}}^+(w), d_{\tilde{D}}^-(w)\} \geq f_1(w)$ for all $w \in V(\tilde{D})$ and so D_1 is not weakly f_1 -degenerate, contradicting the assumption that (D_1, D_2, \dots, D_p) is an f -partition of D . ■

Thus, the “if”-direction of Theorem 3 is proved. The hard part, *i.e.*, the “only if”-direction, is covered in the next theorem.

Theorem 8 *Let D be a connected digraph, let $p \geq 1$ be an integer, and let $f : V(D) \rightarrow \mathbb{N}_0^p$ be a vector function such that $f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$. If D is not f -partitionable, then (D, f) is a hard pair.* ◇

Proof: The proof is by reductio ad absurdum. So let (D, f) be a smallest counterexample, that is,

- (1) $f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$,
- (2) D is not f -partitionable,
- (3) (D, f) is not a hard pair, and
- (4) $|D|$ is minimum subject to (1), (2), and (3).

In order to derive a contradiction, we establish a sequence of eight claims.

Claim 1 $D - v$ is f -partitionable for every $v \in V(D)$. \diamond

Proof : Suppose that $D - v$ is not f -partitionable for some $v \in V(D)$. Then there is a component D' of $D - v$ such that D' is not f -partitionable and so (D', f) is a hard pair (by (4)). As D is connected, there is a neighbor u of v in D' . Since (D', f) is a hard pair, Proposition 7(a) implies that $f_1(u) + f_2(u) + \dots + f_p(u) = d_{D'}^+(u) = d_{D'}^-(u)$. As D contains at least one of the arcs vu and uv , we conclude that

$$f_1(u) + f_2(u) + \dots + f_p(u) = d_{D'}^+(u) = d_{D'}^-(u) < \max\{d_D^+(u), d_D^-(u)\},$$

contradicting (1). \square

The next claim is central for the proof of Theorem 8. Casually speaking, it says that not only is D Eulerian and in- and out-degree of each vertex coincides with the sum of its f -values, but also, given a fixed vertex v and an f -partition (D_1, D_2, \dots, D_p) of $D - v$, the component of each $D_i + v$ containing v is Eulerian, too, and the respective degrees coincide with the f_i -values.

Claim 2 Let $v \in V(D)$ be an arbitrary vertex, let (D_1, D_2, \dots, D_p) be an f -partition of $D - v$, and let $i \in [1, p]$. Then, the following hold:

- (a) $d_{D_i+v}^+(v) = d_{D_i+v}^-(v) = f_i(v)$.
- (b) $f_1(v) + f_2(v) + \dots + f_p(v) = d_D^+(v) = d_D^-(v)$. As a consequence, D is Eulerian.
- (c) Let u be a neighbor of v in D_i . Then, the sequence $(D'_1, D'_2, \dots, D'_p)$ with $D'_i = (D_i + v) - u$ and $D'_j = D_j$ for $j \neq i$ is an f -partition of $D - u$.
- (d) The component D' of $D_i + v$ that contains v is Eulerian and $d_{D'}^+(w) = d_{D'}^-(w) = f_i(w)$ for all $w \in V(D')$. \diamond

Proof : Let $i \in [1, p]$ be arbitrary. As D is not f -partitionable, the digraph $D_i + v$ is not weakly f_i -degenerate. Thus, there is a subdigraph D'' of $D_i + v$ such that $\min\{d_{D''}^+(w), d_{D''}^-(w)\} \geq f_i(w)$ for all $w \in V(D'')$. Clearly, D'' must contain v (as D_i is weakly f_i -degenerate) and so

$$f_i(v) \leq \min\{d_{D''}^+(v), d_{D''}^-(v)\} \leq \min\{d_{D_i+v}^+(v), d_{D_i+v}^-(v)\}.$$

Since i was chosen arbitrarily, we conclude that

$$\begin{aligned} \sum_{i \in [1, p]} f_i(v) &\leq \sum_{i \in [1, p]} \min\{d_{D_i+v}^+(v), d_{D_i+v}^-(v)\} \leq \min\{d_D^+(v), d_D^-(v)\} \\ &\leq \max\{d_D^+(v), d_D^-(v)\} \leq \sum_{i \in [1, p]} f_i(v) \end{aligned}$$

and so we have equality everywhere. Thus, (a) and (b) hold.

For the proof of (c) and (d), let D' be the component of $D_i + v$ containing v , and let u be a neighbor of v in $D_i + v$ and therefore in D' (if such a vertex does not exist, then

D' consists only of v , and by (a) $f_i(v) = 0 = d_{D_i+v}^+(v) = d_{D_i+v}^-(v)$; so there is nothing to prove). Again, let D'' be the subdigraph of $D_i + v$ with $\min\{d_{D''}^+(w), d_{D''}^-(w)\} \geq f_i(w)$ for all $w \in V(D'')$. By (a),

$$f_i(v) = d_{D_i+v}^+(v) = d_{D_i+v}^-(v) \geq \max\{d_{D''}^+(v), d_{D''}^-(v)\} \geq f_i(v),$$

implying that the digraph D'' must contain all neighbors of v in D' and, in particular, D'' contains u . Consequently, $(D_i + v) - u$ is weakly f_i -degenerate (as D'' was chosen arbitrarily and so any "bad" subdigraph must contain u) and, hence, $(D'_1, D'_2, \dots, D'_p)$ with $D'_i = (D_i + v) - u$ and $D'_j = D_j$ for $j \neq i$ is an f -partition of $D - u$, which proves (c). Clearly, the component of $D'_i + u$ containing u is still D' and so applying statement (a) to the vertex u and the partition $(D'_1, D'_2, \dots, D'_p)$ leads to $d_{D_i+v}^+(u) = d_{D_i+v}^-(u) = f_i(u)$ (as $D_i + v = D'_i + u$). Note that if z is a neighbor of u in D' we can swap u and z , *i.e.*, regard the f -partition $(D''_1, D''_2, \dots, D''_p)$ with $D''_i = (D'_i + u) - z$ and $D''_j = D'_j = D_j$ for $j \neq i$, and obtain the same conclusion for z . By repeating this procedure, we eventually reach every vertex of D' and so (d) follows. \square

Claim 3 D is a block. \diamond

Proof : Suppose, to the contrary, that D is the union of two induced subdigraphs D^1 and D^2 with $V(D^1) \cap V(D^2) = \{v\}$ and $|D^j| < |D|$ for $j \in \{1, 2\}$. We will model two functions f^1 and f^2 such that (D^1, f^1) and (D^2, f^2) are hard pairs and (D, f) is obtained from the two hard pairs via the merging operation, thereby giving us the desired contradiction. To this end, let (D_1, D_2, \dots, D_p) be an f -partition of $D - v$ and, for $i \in [1, p]$ and $j \in \{1, 2\}$, let $D_i^j = D_i \cap D^j$. Clearly, the digraphs D_i^1 and D_i^2 are disjoint for all $i \in [1, p]$. Now let $i \in [1, p]$ and let D' be the component of $D_i + v$ containing v . Then,

$$D' = (D' \cap (D_i^1 + v)) \cup (D' \cap (D_i^2 + v)).$$

Note that $(D' \cap (D_i^1 + v))$ and $(D' \cap (D_i^2 + v))$ have only v in common and that there are no arcs between vertices of D_i^1 and D_i^2 . By Claim 2(d), we have $d_{D'}^+(w) = d_{D'}^-(w) = f_i(w)$ for all $w \in V(D')$. Thus, it follows from the above observation that

$$f_i(w) = d_{D'}^+(w) = d_{D'}^-(w) = d_{(D' \cap (D_i^j + v))}^+(w) = d_{(D' \cap (D_i^j + v))}^-(w)$$

for $j \in \{1, 2\}$ and all $w \in V(D_i^j)$. Consequently, all vertices from $D' \cap (D_i^j + v)$ besides v are Eulerian in $D' \cap (D_i^j + v)$. Since in each digraph, the sum of out-degrees over all vertices equals the sum of in-degrees over all vertices, we conclude that vertex v is also Eulerian in $D' \cap (D_i^j + v)$ and, therefore, in $D_i^j + v$ for $j \in \{1, 2\}$, *i.e.*

$$d_{D_i^j+v}^+(v) = d_{(D' \cap (D_i^j + v))}^+(v) = d_{(D' \cap (D_i^j + v))}^-(v) = d_{D_i^j+v}^-(v). \quad (3.1)$$

This gives us a nice way to define the functions f^1 and f^2 . For $i \in [1, p]$ and $j \in \{1, 2\}$ let

$$f_i^j(w) = \begin{cases} f_i(w) & \text{if } w \in V(D^j) \setminus \{v\}, \text{ and} \\ d_{D_i^j}^+(v) & \text{if } w = v. \end{cases} \quad (3.2)$$

Then, for $j \in \{1, 2\}$ and all $w \in V(D^j) \setminus \{v\}$ we have

$$\begin{aligned} f_1^j(w) + f_2^j(w) + \dots + f_p^j(w) &= f_1(w) + f_2(w) + \dots + f_p(w) \\ &= \max\{d_D^+(w), d_D^-(w)\} = \max\{d_{D^j}^+(w), d_{D^j}^-(w)\}. \end{aligned} \quad (3.3)$$

Moreover, since v is Eulerian in $D_i^j + v$ we obtain

$$d_{D^j+v}^+(v) = \sum_{i \in [1, p]} d_{D_i^j+v}^+(v) = \sum_{i \in [1, p]} d_{D_i^j+v}^-(v) = d_{D^j+v}^-(v), \quad (3.4)$$

and by the choice of $f^j(v)$, it follows that

$$f_1^j(v) + f_2^j(v) + \dots + f_p^j(v) = d_{D^j}^+(v) = d_{D^j}^-(v). \quad (3.5)$$

As a consequence, both (D^1, f^1) and (D^2, f^2) fulfil the requirements of Theorem 8.

First assume that for some $j \in \{1, 2\}$, the digraph D^j is f^j -partitionable. By symmetry, we may assume $j = 1$. Then, D^1 admits an f^1 -partition $(D'_1, D'_2, \dots, D'_p)$ and, by symmetry, $v \in V(D'_1)$. Let $(D_1^*, D_2^*, \dots, D_p^*)$ be a partition of D with $D_1^* = D'_1 \cup (D_1^2 + v)$ and $D_i^* = D'_i \cup D_i^2$ for $i \in [2, p]$. As D'_i and D_i^2 are disjoint and as $(D_1^2, D_2^2, \dots, D_p^2)$ is an f -partition of $D^2 - v$, D_i^* is weakly f_i -degenerate for $i \in [2, p]$. We claim that D_1^* is weakly f_1 -degenerate. To this end, let \tilde{D} be a non-empty subdigraph of D_1^* . If \tilde{D} is a subdigraph of D_1^2 , then \tilde{D} is weakly f_1 -degenerate (as D_1^2 is weakly f_1 -degenerate). So assume that $\tilde{D}^1 = \tilde{D} \cap D^1$ is non-empty. Since \tilde{D}^1 is a subdigraph of D'_1 and therefore weakly f_1^1 -degenerate, there is a vertex $w \in V(\tilde{D}^1)$ with $\min\{d_{\tilde{D}^1}^+(w), d_{\tilde{D}^1}^-(w)\} < f_1^1(w)$. If $w \neq v$, then

$$\min\{d_{\tilde{D}}^+(w), d_{\tilde{D}}^-(w)\} = \min\{d_{\tilde{D}^1}^+(w), d_{\tilde{D}^1}^-(w)\} < f_1^1(w) = f_1(w), \quad (3.6)$$

and we are done. It remains to consider the case that $w = v$. Since $d_{D_1^2+v}^+(v) = d_{D_1^2+v}^-(v) = f_1^2(v)$ (by (3.1) and (3.2)), this implies that

$$\min\{d_{\tilde{D}}^+(v), d_{\tilde{D}}^-(v)\} \leq \min\{d_{\tilde{D}^1}^+(v), d_{\tilde{D}^1}^-(v)\} + d_{D_1^2}^+(v) < f_1^1(v) + f_1^2(v) = f(v).$$

Consequently, D_1^* is weakly f_1 -degenerate, as claimed, and so $(D_1^*, D_2^*, \dots, D_p^*)$ is an f -partition of D , which is impossible.

Thus, D^j is not f^j -partitionable for $j \in \{1, 2\}$ and, hence by the minimality of D , both (D^1, f^2) and (D^2, f^2) are hard pairs and (D, f) is obtained from (D^1, f^2) and (D^2, f^2) via the merging operation. As a consequence, (D, f) is a hard pair, contradicting (3). This contradiction completes the proof of the claim. \square

By the above claim, D is a block. Hence it remains to show that (D, f) is a hard pair of type (M), (K), or (C), giving us the desired contradiction. The next claim eliminates the pairs of type (M).

Claim 4 *For every $v \in V(D)$ and each f -partition (D_1, D_2, \dots, D_p) of $D - v$, there are two indices $i \neq j$ from $[1, p]$ such that D_i and D_j are non-empty.* \diamond

Proof : Suppose that there is a vertex $v \in V(D)$ and a partition (D_1, D_2, \dots, D_p) of $D - v$ such that exactly one part of the partition is non-empty, say D_1 . Then $D_1 + v = D$ and so it follows from Claim 2(d) and the fact that D is connected that $d_D^+(w) = d_D^-(w) = f_1(w)$ for all $w \in V(D)$. Thus it follows from Claim 2(b) that $f_j(w) = 0$ for all $j \neq 1$ from $[1, p]$ and for all $w \in V(D)$ and (D, f) is a hard pair of type (M), contradicting (3). \square

Actually, Claim 2(c) provides us with a powerful tool that we shall use in the following. Let $v \in V(D)$ be an arbitrary vertex and let (D_1, D_2, \dots, D_p) be an f -partition of $D - v$. Moreover, let $u \in V(D)$ be a neighbor of v . Then, $u \in V(D_i)$ for some $i \in [1, p]$ and Claim 2(c) implies that replacing D_i with $(D_i + v) - u$ leads to an f -partition of $D - u$ in which v is contained in what was previously D_i . Thus, we can swap v and any neighbor u of v and obtain a new f -partition of $D - u$. In order to make this observation a bit more graphical, we introduce the following terms. Given a vertex $v \in V(D)$, we call an f -partition (D_1, D_2, \dots, D_p) of $D - v$ an **f -coloring** φ of $D - v$. Moreover, for $w \in V(D) \setminus \{v\}$, we set $\varphi(w) = i$ if $w \in V(D_i)$ and say that w has **color** i . Finally, we say that the vertex v is **uncolored**. The following Claim is just a reformulation of Claim 2(a) and Claim 4 to fit the new terminology.

Claim 5 *Let $v \in V(D)$ be an arbitrary vertex and let φ be an f -coloring of $D - v$. Then, the following statements hold:*

- (a) *At least two color-classes of φ are non-empty.*
- (b) $d_{D[\varphi^{-1}(\{i\})+v]}^+(v) = d_{D[\varphi^{-1}(\{i\})+v]}^-(v) = f_i(v)$ for all $i \in [1, p]$. \diamond

Using this terminology, the above described method of swapping two vertices is nothing else than assigning to v the color of some neighbor u of v and uncoloring u . We will call this process **shifting** the color from u to v . Note that this leads to an f -coloring of $D - u$. The original idea of shifting goes back to GALLAI [15].

Now let C be a cycle in $G(D)$, let $v \in V(C)$ be an arbitrary vertex, and let φ be an f -coloring of $D - v$. Moreover, let u and w be the vertices such that uv and vw are in $E(C)$. Then, we can shift the color from u to v . Afterwards, we shift the color from the other neighbor of u on C to u . Continuing like this, we can shift the color of each vertex of C , one after another, clockwise, until eventually we shift the color from v to w . This gives us a new f -coloring φ' of $D - v$ (see Figure 3). In particular, $\varphi'(w) = \varphi(u)$.

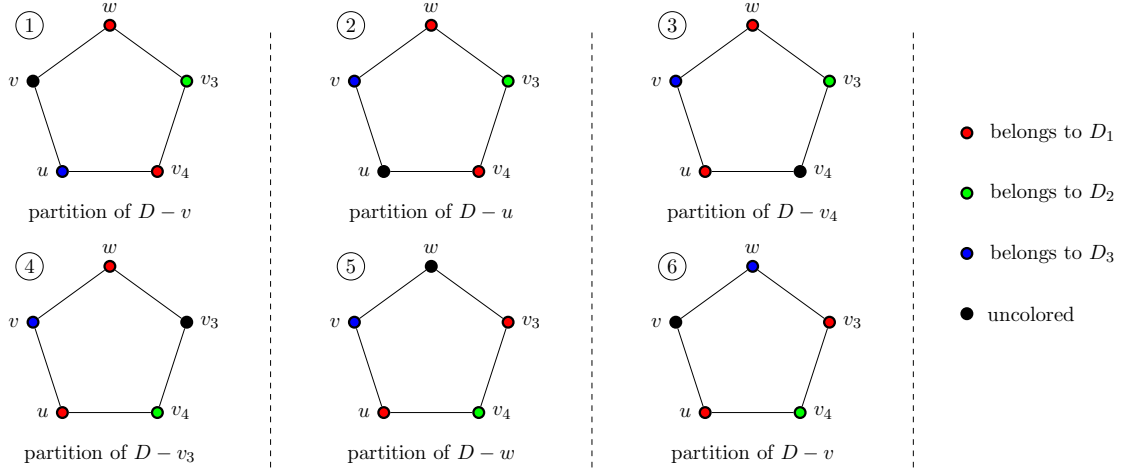


Figure 3: Clockwise shifting of colors around a cycle in $G(D)$.

Similarly, starting from φ with shifting the color from w to v , we can shift the color of each vertex counter-clockwise on the cycle and obtain a third f -coloring of $D - v$.

By repeated clockwise, respectively counter-clockwise shifting, it is easy to see that the following claim is true.

Claim 6 *Let C be a cycle in $G(D)$, let $v \in V(C)$ be an arbitrary vertex, and let φ be an f -coloring of $D - v$. Moreover, let u and w be the neighbours of v on C , that is, $\{uv, vw\} \subseteq E(C)$. Then, for any pair v_1, v_2 of vertices distinct from v with $v_1v_2 \in E(C)$ and for each $i \in \{1, 2\}$ there is an f -coloring φ^* of $D - v$ such that $\varphi^*(u) = \varphi(v_i)$ and $\varphi^*(w) = \varphi(v_{3-i})$.* \diamond

Using Claims 5 and 6, we are able to prove the next claim.

Claim 7 *Let C be a cycle in $G(D)$ and let $v \in V(C)$ be a vertex of C that is not contained in a chord of C in $G(D)$. Then, C is an odd cycle and there is an f -coloring φ^* of $D - v$ and indices $k \neq \ell$ from $[1, p]$ such that the vertices of $C - v$ are colored alternately with k and ℓ .* \diamond

Proof : Let u and w be the vertices with $\{uv, vw\} \subseteq E(C)$. As v is not contained in a chord, u and w are the only neighbors of v from $V(C)$ in $G(D)$. We first claim that there is an f -coloring φ^* of $D - v$ with $\varphi^*(u) \neq \varphi^*(w)$. To this end, let φ be an f -coloring of $D - v$ and assume that $\varphi(u) = \varphi(w)$. From Claim 5(a) we know that at least two color classes of φ are non-empty. Thus, in $D - v$ there is a vertex z with $\varphi(z) \neq \varphi(u)$. As C is a cycle contained in the block $G(D)$, we have $|G(D)| \geq 3$ and so $G(D)$ is 2-connected. Then, it follows from Lemma 4 that in $G(D)$ there are two paths P, P' from z into the set $\{u, v, w\}$ whose only common vertex is z and whose internal vertices are not in $\{u, v, w\}$. Note that the union $P \cup P'$ together with either one or both of the edges $\{uv, vw\}$ forms a cycle C' in $G(D)$, which contains v . Since $\varphi(z) \neq \varphi(u) = \varphi(w)$, C' contains two consecutive vertices of different colors, say v_1 and v_2 . If C' contains both u and w , it follows from Claim 6 that there is a coloring φ^* of $D - v$ with $\varphi^*(u) = \varphi(v_1) \neq \varphi(v_2) = \varphi^*(w)$, and we are done. In order to complete the first part of the claim, it remains to consider the case that C' contains only one of u, w , say u . By symmetry, we may assume $\varphi(v_1) \neq \varphi(w)$. Then, it follows from Claim 6 that there is a coloring φ^* of $D - v$ with $\varphi^*(u) = \varphi(v_1) \neq \varphi(w) = \varphi^*(w)$ and so we are done.

Now let $\varphi^*(u) = k$ and $\varphi^*(w) = \ell$. We claim that the vertices of $C - v$ are colored alternately with k and ℓ . Otherwise, there are vertices u_1, u_2 distinct from v with $u_1u_2 \in E(C)$ and $\{\varphi^*(u_1), \varphi^*(u_2)\} \neq \{k, \ell\}$. By symmetry, we may assume that $\ell \notin \{\varphi^*(u_1), \varphi^*(u_2)\}$. Then, by Claim 6, there is a coloring φ' of $D - v$ with $\varphi'(u) = \varphi^*(u_1)$ and $\varphi'(w) = \varphi^*(u_2)$. As $\ell \notin \{\varphi^*(u_1), \varphi^*(u_2)\}$, and as u and w are the only neighbors of v from $V(C)$ in $G(D)$, we obtain that either

$$d_{D[(\varphi')^{-1}(\{\ell\})+v]}^+(v) \neq d_{D[(\varphi^*)^{-1}(\{\ell\})+v]}^+(v) = f_\ell(v)$$

or

$$d_{D[(\varphi')^{-1}(\{\ell\})+v]}^-(v) \neq d_{D[(\varphi^*)^{-1}(\{\ell\})+v]}^-(v) = f_\ell(v),$$

in contradiction to Claim 5(b). This proves the claim that the vertices of $C - v$ are colored alternately with k and ℓ . As $\varphi^*(u) \neq \varphi^*(w)$, this implies that C is an odd cycle and so the proof of the claim is complete. \square

As a consequence of the above claim, we obtain that every even cycle in $G(D)$ has at least two chords. Thus, it follows from Lemma 5 that every block of $G(D)$ is an odd cycle

or a complete graph. As D and therefore $G(D)$ itself is a block, we conclude that $G(D)$ is either a cycle of odd length or a complete graph. To complete the proof of Theorem 3, we show that both cases are impossible.

Claim 8 $G(D)$ is not an odd cycle. \diamond

Proof : For otherwise, let $v \in V(D)$ be an arbitrary vertex. As $C = G(D)$ is an odd cycle, C obviously has no chords and so it follows from Claim 7 that there is a coloring φ of $D - v$ and indices $k \neq \ell$ from $[1, p]$ so that the vertices of $C - v$ are colored alternately with k and ℓ . Then, it follows from Claim 5(b) that

$$f_k(v) = f_\ell(v) = 1 \text{ and } f_i(v) = 0 \text{ for } i \in [1, p] \setminus \{k, \ell\} \quad (3.7)$$

and $\{uv, vu, vw, wv\} \subseteq A(D)$ where u and w are the neighbors of v in C .

By shifting the color from the neighbor u of v to v we obtain (3.7) for u instead of v . Repeating this argument proves that (D, f) is a hard pair of type (C), a contradiction. \square

It remains to consider the case that $G(D)$ is a complete graph. Claim 5 implies that $|D| \geq 3$. First we claim that D is bidirected. For otherwise, there are two vertices $u, v \in V(D)$ with $uv \in A(D)$ but $vu \notin A(D)$. Let φ be an f -coloring of $D - v$. Then, by Claim 5(a), there is a vertex $w \in V(D) \setminus \{v\}$ with $\varphi(w) \neq \varphi(u)$. As $G(D[\{u, v, w\}])$ is a triangle, it follows from Claim 6 that swapping the colors of u and w results in another f -coloring of $D - v$. Then, by Claim (5)(b), we obtain that $wv \in A(D)$ but $vw \notin A(D)$. By shifting the color from w to v we conclude from Claim 5(b) that $uw \in A(D)$ but $wu \notin A(D)$. Similarly, by instead shifting the color from u to w we conclude from Claim (5)(b) that $wu \in A(D)$, but $uw \notin A(D)$, which clearly is impossible. This proves the claim that D is bidirected and so D is a bidirected complete graph.

Finally, we prove that (D, f) is a hard pair of type (K), leading to the desired contradiction. To this end, let $v \in V(D)$ be an arbitrary vertex, let (D_1, D_2, \dots, D_p) be an f -partition of $D - v$, and let $n_i = |D_i|$. As D is a bidirected complete graph, it then follows from Claim 2(a) that $f_i(v) = n_i$ for all $i \in [1, p]$. Moreover, if $u \in V(D) \setminus \{v\}$, say $u \in V(D_1)$ (by symmetry), then the sequence $(D'_1, D'_2, \dots, D'_p)$ with $D'_1 = (D_1 + v) - u$ and $D'_j = D_j$ for $j \in [2, p]$ is an f -partition of $D - u$ (by Claim 2(c)) with $|D'_i| = |D_i|$ for all $i \in [1, p]$ and applying Claim 2(a) to u leads to $f_i(u) = n_i$ for all $i \in [1, p]$. Note that

$$n_1 + n_2 + \dots + n_p = |D_1| + |D_2| + \dots + |D_p| = |V(D)| - 1$$

and so (D, f) is a hard pair of type (K), contradicting (3). This contradiction completes the proof of Theorem 3. \blacksquare

Since a bidirected graph D is weakly h -degenerate if and only if its underlying graph $G(D)$ is strictly h -degenerate (*i.e.* in each non-empty subgraph G' of G there is a vertex v with $d_{G'}(v) < h(v)$), the restriction of Theorem 3 to bidirected graphs gives us also a result regarding strict degeneracy of undirected graphs. This result for undirected graphs was originally proven by Borodin, Kostochka, and Toft in the paper [8]. The proof structure of Claim 3 is similar to that of the first part of their proof. However, apart from this, our proof leads to another proof for the undirected case that uses completely different methods than the original proof. A major benefit of our proof is that it—contrary to the original proof in [8]—generalizes easily to the case of directed multigraphs respectively

multigraphs. Of course, the definition of weak degeneracy and of f -partitions also works if we allow multiple arcs between vertices going in the same direction. Also, the definition of hard pair needs only to be generalized slightly. To this end, we need the following term. Let G be a simple graph and let $t \geq 1$ be an integer. Then we denote by tG the graph that results from G by replacing each edge by t parallel edges. Now let \mathcal{D} be a connected multidigraph, and let $f : V(\mathcal{D}) \rightarrow \mathbb{N}_0^p$ be a vector function. We say that (\mathcal{D}, f) is a **hard pair**, if either

- (\mathcal{D}, f) is of type (M), or
- \mathcal{D} is a bidirected tK_n and there are integers $n_1, n_2, \dots, n_p \geq 0$ with at least two n_i different from zero such that $n_1 + n_2 + \dots + n_p = n - 1$ and

$$f(v) = (tn_1, tn_2, \dots, tn_p) \text{ for all } v \in V(\mathcal{D}), \text{ or}$$

- \mathcal{D} is a bidirected tC_n with n odd and there are two indices $k \neq \ell$ from the set $[1, p]$ such that for all $v \in V(\mathcal{D})$, we have $f_i(v) = t$ if $i \in \{k, \ell\}$ and $f_i(v) = 0$ otherwise, or
- (\mathcal{D}, f) is obtained from two hard pairs via the merging operation.

For multidigraphs, we obtain the following theorem.

Theorem 9 *Let \mathcal{D} be a connected multidigraph, let $p \geq 1$, and let $f : V(\mathcal{D}) \rightarrow \mathbb{N}_0^p$ be a vector function such that $f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_{\mathcal{D}}^+(v), d_{\mathcal{D}}^-(v)\}$ for all $v \in V(\mathcal{D})$. Then, \mathcal{D} is not f -partitionable if and only if (\mathcal{D}, f) is a hard pair. \diamond*

By inspecting their proofs, it is easy to check that Claims 1-7 also hold for multidigraphs. Only the remaining part of the proof needs to be changed slightly, but still can easily be done adapting the methods described there. Therefore, we abstain from giving an extra proof.

4 Applications of Theorem 3

4.1 Brooks' Theorem for list-colorings of digraphs

As mentioned in the introduction, Theorem 3 implies Harutyunyan and Mohar's Theorem 2 [19]. Let us recall the theorem for the reader's convenience.

Theorem 2 (Harutyunyan and Mohar, 2011) *Let D be a connected digraph, and let L be a list-assignment such that $|L(v)| \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$. Suppose that D is not L -colorable. Then, the following statements hold:*

- D is Eulerian and $|L(v)| = \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$.*
- If $B \in \mathcal{B}(D)$, then B is a directed cycle of length ≥ 2 , or B is a bidirected complete graph, or B is a bidirected cycle of odd length ≥ 5 .*
- For each $B \in \mathcal{B}(D)$ there is a set Γ_B of $\Delta^+(B)$ colors such that for every $v \in V(D)$, the sets Γ_B with $B \in \mathcal{B}_v(D)$ are pairwise disjoint and $L(v) = \bigcup_{B \in \mathcal{B}_v(D)} \Gamma_B$. \diamond*

Proof of Theorem 2: Let D and L be as described in the theorem. By using the method from the introduction, we transform the list-coloring problem to that of finding an f -partition: Let $\Gamma = \bigcup_{v \in V(D)} L(v)$. By renaming the colors if necessary, we may assume $\Gamma = [1, p]$. Now let $f : V(D) \rightarrow \mathbb{N}_0^p$ be the vector function with

$$f_i(v) = \begin{cases} 1 & \text{if } i \in L(v), \text{ and} \\ 0 & \text{if } i \notin L(v) \end{cases}$$

for $i \in [1, p]$. By the definition of f , we have $f_1(v) + f_2(v) + \dots + f_p(v) = |L(v)| \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$ and D is not f -partitionable. Thus, it follows from Theorem 3 that (D, f) is a hard pair. Then, statement (a) follows from Proposition 7(a). Moreover, Proposition 6(a) implies that for each block $B \in \mathcal{B}$, (B, f_B) is of type (M), (K), or (C), where f_B is the function as defined in Proposition 6. Note that if (B, f_B) is of type (M), then from the definition of f it follows that the only non-zero coordinate of f_B is constant 1 and, hence, $d_B^+(v) = d_B^-(v) = 1$ for all $v \in V(B)$. Consequently, B is a directed cycle. Thus, (b) holds true. Statement (c) follows easily from Proposition 6(a) and (b). \blacksquare

4.2 The s -degenerate dichromatic number

In [16], Golowich introduces a generalization of the dichromatic number as follows. Let $s \geq 1$ be an integer. Then, the **s -degenerate dichromatic number** of a digraph D , denoted by $\vec{\chi}_s(D)$, is the least integer p such that D admits a p -partition into weakly s -degenerate subdigraphs. Clearly, χ_1 corresponds to the dichromatic number. Note that the s -degenerate dichromatic number is the digraph counterpart to the **point partition number** of an undirected graph, which was introduced by Lick and White [25] and which is also known as the **s -chromatic number**. The point partition number χ_s of an undirected graph G is usually defined as the minimum number p such that G admits a p -partition into s -degenerate subgraphs. By setting $f_i \equiv s$ in Theorem 3, we easily obtain the following theorem, which was proven for undirected graphs by Mitchem [26].

Theorem 10 *Let D be a connected digraph and let $m = \max_{v \in V(D)} \{d_D^+(v), d_D^-(v)\}$. Moreover, let $s \geq 1$ be an integer and let $p = \lceil \frac{m}{s} \rceil$. If D is neither a bidirected odd cycle, a bidirected complete graph, nor an Eulerian digraph in which every vertex has in- and out-degree s , then $\chi_s(D) \leq p$.* \diamond

To conclude the paper, let us demonstrate how to obtain a list-version of Theorem 10. Given a digraph D and a list-assignment L of D , we say that D is **(L, s) -colorable** if D admits an L -coloring in which every color class induces a weakly s -degenerate subdigraph.

Theorem 11 *Let D be a connected digraph and let $m = \max_{v \in V(D)} \{d_D^+(v), d_D^-(v)\}$. Moreover, let $s \geq 1$ be an integer and let L be a list-assignment with $|L(v)| \geq \frac{m}{s}$ for all $v \in V(D)$. Then, D is not (L, s) -colorable if and only if the following two conditions are fulfilled:*

- (a) *D is a bidirected complete graph with $|D| - 1 \equiv 0 \pmod{s}$, or D is a bidirected odd cycle and $s = 1$, or D is an Eulerian digraph in which every vertex has in-degree and out-degree s .*
- (b) *There is a color set Γ such that $L(v) = \Gamma$ for all $v \in V(D)$ and $|\Gamma| = m/s$.* \diamond

Proof: If $m = 0$, then D consists of just one vertex and the statement is evident. So assume $m \geq 1$. Let $\Gamma = \bigcup_{v \in V(D)} L(v)$. By renaming the colors if necessary, we can assume that $\Gamma = [1, p]$. Now we define a vector function $f : V(D) \rightarrow \mathbb{N}_0^p$ as follows. Let

$$f_i(v) = \begin{cases} s & \text{if } i \in L(v), \\ 0 & \text{otherwise} \end{cases}$$

for $i \in [1, p]$. Then, D is not (L, s) -colorable if and only if D is not f -partitionable. By definition, we have $f_1(v) + f_2(v) + \dots + f_p(v) = s|L(v)| \geq m$ for all $v \in V(D)$ and so D is not (L, s) -colorable if and only if (D, f) is a hard pair. (by Theorem 3)

First assume that D and L satisfy (a) and (b). Then it is easy to check that (D, f) is indeed a hard pair and so D is not (L, s) -colorable. Now assume that D is not (L, s) -colorable. Then (D, f) is a hard pair and from Proposition 7(a) it follows that $m \leq s|L(v)| = d_D^+(v) = d_D^-(v)$ for all $v \in V(D)$ and so D is Eulerian and each vertex v satisfies $d_D^+(v) = d_D^-(v) = m$. If $m = s$, then $|L(v)| = 1$ for all $v \in V(D)$ and D is not (L, s) -colorable if and only if there is a color $\alpha \in \Gamma$ with $L(v) = \{\alpha\}$ for all $v \in V(D)$, and we are done. So assume $m > s$. Then, $|L(v)| \geq 2$ and so for each vertex v there are two indices $i \neq j$ such that $f_i(v) \geq 1$ and $f_j(v) \geq 1$. Consequently, by Proposition 6, no end-block of D is a mono-block. Since every vertex of D has in-degree and out-degree m , this implies that D itself is a block and so D is either a bidirected complete graph or a bidirected odd cycle. In the first case, (D, f) is of type (K) and, since every coordinate of f is either s or zero we easily conclude from the definition of hard pair of type (K) that $|D| - 1 \equiv 0 \pmod{s}$. If (D, f) is of type (C), we conclude that $s = 1$. Moreover, in both cases the function f is constant and so the list-assignment L is constant, too, and $|L(v)| = m/s$ for all $v \in V(D)$. This completes the proof. ■

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5 Appendix: Constructing an f -partition

The problem of deciding whether the dichromatic number of a digraph is at most two is already NP-complete [7]. Since this problem is the same as deciding whether there is an f -partition for $f \equiv (1, 1)$, we cannot hope to find an efficient algorithm for deciding, for a given pair (D, f) , whether D admits an f -partition. The good news, however, is that it is possible to deduce a polynomial time algorithm from our proof that, given a pair (D, f) with $f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$, either verifies that (D, f) is a hard pair or returns an f -partition of D . The algorithm uses the following six subroutines.

Algorithm 1 Greedy Coloring

Input: (D, f, v^*) such that D is connected, $f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$ and $f_1(v^*) + f_2(v^*) + \dots + f_p(v^*) > \min\{d_D^+(v^*), d_D^-(v^*)\}$.

Output: An f -partition (D_1, D_2, \dots, D_p) of D

```

1:  $n := |D|$ ,  $v_n := v^*$ 
2: for  $i = n - 1, n - 2, \dots, 1$  do
3:   let  $v_i$  be a vertex such that  $v_i$  has a neighbor from  $\{v_{i+1}, v_{i+2}, \dots, v_n\}$ 
4: end for
5:  $D_j := \emptyset$  for  $j \in [1, p]$ 
6: for  $i = 1, 2, \dots, n$  do
7:    $D_j := D_j \cup \{v_i\}$  where  $j$  is the minimum integer with
8:    $f_j(v_i) > \min\{d_{D_j+v_i}^+(v_i), d_{D_j+v_i}^-(v_i)\}$ .
9:
10: end for
11: return  $(D_1, D_2, \dots, D_p)$ .

```

Algorithm 2 Greedy Coloring of $D - v$

Input: (D, f, v) where D is connected and $f_1(v) + f_2(v) + \dots + f_p(v) = \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$

Output: An f -partition (D_1, D_2, \dots, D_p) of $D - v$

- 1: Let K_1, \dots, K_q be the connected components of $D - v$.
 - 2: **for** $j = 1, 2, \dots, q$ **do**
 - 3: find a vertex $v_{K_j}^*$ with $f_1(v_{K_j}^*) + f_2(v_{K_j}^*) + \dots + f_p(v_{K_j}^*) > \min\{d_K^+(v_{K_j}^*), d_K^-(v_{K_j}^*)\}$
 - 4: apply Algorithm 1 to $(K, f, v_{K_j}^*)$ to obtain an f -partition $(D_{K_j,1}, D_{K_j,2}, \dots, D_{K_j,p})$ of K_j .
 - 5: **end for**
 - 6: **for** $i = 1, 2, \dots, p$ **do**
 - 7: $D_i := \bigcup_{j \in [1,q]} D_{K_j,i}$
 - 8: **end for**
 - 9: **return** (D_1, D_2, \dots, D_p) .
-

The correctness of the first two algorithms follows from Claims 1 and 2. In particular, in step 7 of Algorithm 1, we always find such an index j since

$$\begin{aligned} \sum_{k \in [1,p]} \min\{d_{D_k+v_i}^+(v_i), d_{D_k+v_i}^-(v_i)\} &\leq \min\{d_{D[\{v_1, v_2, \dots, v_i\}]}^+(v_i), d_{D[\{v_1, v_2, \dots, v_i\}]}^-(v_i)\} \\ &< \max\{d_D^+(v_i), d_D^-(v_i)\} \leq f_1(v_i) + f_2(v_i) + \dots + f_p(v_i). \end{aligned}$$

Here the last in-equality follows from the fact that v_i has a neighbor in the set $\{v_{i+1}, \dots, v_n\}$.

Algorithm 3 describes the shifting procedure that we introduced after Claim 4. Given a vertex v and a neighbor w of v , we first check if we can add v to the f -partition in order to obtain the desired f -partition of D . If this is not the case, then it follows from Claim 2(c) that we can "swap" v and w and obtain an f -partition of $D - w$. The idea of the Main Algorithm 7 is to iteratively split end-blocks B together with the appropriate block-function f_B (as in Proposition 6) from D and try to partition those separately. If we find an f_B -partition of some block B , we describe a procedure how to extend this partition to an f -partition of the whole digraph.

Algorithm 4 deals with the case that the underlying graph $C = G(B)$ of a block B is a cycle and (B, f_B) is not a hard pair. Then, it follows from the proof of Claim 8 that by repeated application of the shifting procedure we will eventually obtain an f_B -partition. As (B, f_B) is not a hard pair, there are two consecutive vertices v, w on C with $f_B(v) \neq f_B(w)$. By beginning the shifting with the vertex v , we ensure that the algorithm will terminate after at most $\text{length}(C) + 1$ iterations since then the condition in Algorithm 3 line 2 will be violated at the latest. This is due to the fact that if u, u' are the neighbors of v , respectively w on C distinct from $\{v, w\}$ and (D_1, D_2, \dots, D_p) is an f_B -partition of $D - v$, then u and u' need to belong to different parts of the partition (by Claim 4(b) applied to (D_1, D_2, \dots, D_p) and to the f -partition of $D - w$ that results from (D_1, D_2, \dots, D_p) by swapping v and w). Thus, after $\text{length}(C) + 1$ iterations, v is again uncolored but the neighbors of v will belong to other partition parts than previous. Hence, Algorithm 4 runs in polynomial time.

Algorithm 3 Shifting Procedure

Input: $(D, f, v, w, (D_1, D_2, \dots, D_p))$ where

- D is a Eulerian block
- v is a vertex of D
- w is a neighbor of v in D
- $f_1(u) + f_2(u) + \dots + f_p(u) = d_D^+(u) = d_D^-(u)$ for all $u \in V(D)$
- (D_1, D_2, \dots, D_p) f -partition of $D - v$

Output: Either an f -partition of D , or an f -partition of $D - w$.

```
1: for  $i = 1, 2, \dots, p$  do
2:   if  $\min\{d_{D_i+v}^+(v), d_{D_i+v}^-(v)\} < f_i(v)$  then
3:      $D_i := D_i + v$ 
4:     return the  $f$ -partition  $(D_1, D_2, \dots, D_p)$  of  $D$ 
5:   end if
6: end for
7: let  $i$  be the index with  $w \in V(D_i)$ 
8:  $D_i := D_i + v - w$ 
9: return the  $f$ -partition  $(D_1, D_2, \dots, D_p)$  of  $D - w$ .
```

Algorithm 4 f -partition of a cycle

```
1: procedure CYCLEPARTITION( $D, f$ )
2:   let  $v, w$  be two adjacent vertices of  $D$  with  $f(v) \neq f(w)$ 
3:   apply Algorithm 2 to  $(D, f, v)$ 
4:   while true do
5:     apply Algorithm 3 to  $(D, f, v, w, (D_1, D_2, \dots, D_p))$ 
6:     if Algorithm 3 returns an  $f$ -partition  $(D'_1, D'_2, \dots, D'_p)$  of  $D$  then
7:       return the  $f$ -partition  $(D'_1, D'_2, \dots, D'_p)$  of  $D$ 
8:     else
9:       Algorithm 3 returns an  $f$ -partition  $(D'_1, D'_2, \dots, D'_p)$  of  $D - w$ 
10:       $u :=$  neighbor of  $w$  in  $D$  distinct from  $v$ ,  $v := w$ ,  $w := u$ ,
       $(D_1, D_2, \dots, D_p) := (D'_1, D'_2, \dots, D'_p)$ 
11:    end if
12:  end while
13: end procedure
```

Algorithm 5 describes a procedure that settles the case that the underlying graph of a block B is a complete graph. Here, we first need to check if B is bidirected; otherwise, we describe a method how to obtain an f_B -partition of the non-bidirected block (lines 1 till 24). That the procedure works follows from the proof's part after Claim 8. If B is bidirected, we argue as in the same proof section to get an f -partition.

Algorithm 5 f -partition of a complete digraph

```

1: procedure COMPLETEDIGRAPHPARTITION( $D, f$ )
2:   if there is an arc  $vu \in A(D)$  with  $uv \notin A(D)$  then
3:     apply Algorithm 2 to  $(D, f, v)$  to obtain  $f$ -partition  $(D_1, D_2, \dots, D_p)$  of  $D - v$ .
4:     if  $\min\{d_{D_j+v}^+(v), d_{D_j+v}^-(v)\} < f_j(v_i)$  for some  $j \in [1, p]$  then
5:        $D_j := D_j + v_i$ 
6:       return  $f$ -partition  $(D_1, D_2, \dots, D_p)$ .
7:     else if all but one  $D_i$  are empty then
8:       find vertex  $w \in V(D)$  with  $f_j(w) > 0$  for some  $j \neq i$  (if possible,  $w = v$ ).
9:        $D_i := (D_i + v) - w$ ,  $D_j := D_j + w$ 
10:      return  $(D_1, D_2, \dots, D_p)$   $f$ -partition of  $D$ 
11:   else
12:     let  $i$  be the index with  $u \in V(D_i)$ . Find  $w$  with  $w \in V(D_j)$  and  $i \neq j$ .
13:      $v^* := v$ ,  $C = G(D)[\{u, v, w\}]$ 
14:     while true do
15:       apply Algorithm 3 to  $(D, f, C, v^*, (D_1, D_2, \dots, D_p))$ 
16:       if Algorithm 3 returns  $f$ -partition  $(D'_1, D'_2, \dots, D'_p)$  of  $D$  then
17:          $(D_1, D_2, \dots, D_p) := (D'_1, D'_2, \dots, D'_p)$ 
18:         return  $(D_1, D_2, \dots, D_p)$  is  $f$ -partition of  $D$ 
19:       else
20:         Algorithm 3 returns  $f$ -partition  $(D'_1, D'_2, \dots, D'_p)$  of  $D - w^*$ 
21:         where  $w^*$  is the right neighbor of  $v^*$  on  $C$ 
22:          $v^* := w^*$ ,  $(D_1, D_2, \dots, D_p) := (D'_1, D'_2, \dots, D'_p)$ 
23:       end if
24:     end while
25:   end if
26:   else  $D$  is a bidirected complete graph. Let  $v \in V(D)$  be arbitrary and apply
27:     Algorithm 2 to obtain an  $f$ -partition  $(D_1, D_2, \dots, D_p)$  of  $D - v$ .
28:     let  $n_i = |D_i|$  for  $i = 1, \dots, p$ .
29:     if  $f_i(v) > n_i$  for some  $i \in [1, p]$  then
30:        $D_i := D_i + v$ 
31:       return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ 
32:     else if there is a vertex  $w$  with  $w \in V(D_i)$  ( $i \in [1, p]$ ) and  $f_i(w) \geq n_i + 2$  then
33:        $D_i := D_i + v$ 
34:       return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ 
35:     else
36:       find a vertex  $w$  and indices  $i \neq j$  with  $w \in V(D_i)$  and  $f_j(w) > n_j$ 
37:        $D_i := (D_i + v) - w$ ,  $D_j := D_j + w$ 
38:       return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ 
39:     end if
40:   end if
41: end procedure

```

Algorithm 6 uses the two previously described procedures in order to find an f_B -partition of any block B where (B, f_B) is not a hard pair. We first check if (B, f_B) is a hard pair (which can be done efficiently). If this is not the case, we determine if the underlying graph is a cycle or a complete graph and apply the respective procedures to get an f_B -partition. If the block's underlying graph is neither a complete graph nor a cycle, we argue as in Claim 7. Note that Cranston and Rabern describe in [11] a constructive proof of Gallai's Lemma 5. In particular, they provide an efficient method how to find an even cycle that has at most one chord if the block is neither a complete graph nor an odd cycle. Moreover, in the worst case we need to shift a lot here but we still stay polynomial in the size of the block.

Algorithm 7, which is the main algorithm, first tries to color the vertices greedily by using Algorithm 1. If this is not possible, it immediately follows that $f_1(v) + f_2(v) + \dots + f_p(v) = d_D^+(v) = d_D^-(v)$ for all $v \in V(D)$ and, hence, D is Eulerian. Note that this implies that every block is Eulerian, too (see Claim 3). We then compute a block-decomposition of D (which can be done efficiently, see [21]). Afterwards, we take an end-block B_1 of D as well as the unique separating vertex v_1 of D contained in B_1 and color the vertices of $D - v_1$ greedily using Algorithm 2. If we cannot add v_1 to any of the partition parts, then it follows from Claim 3 that we can split $f(v_1)$ into $f_{B_1}(v_1)$ and an updated $f(v_1) := f(v_1) - f_{B_1}(v_1)$ as in lines 14-22, so that both (B_1, f_{B_1}) as well as the remaining part $D \setminus (V(B_1) \setminus \{v_1\})$ together with the new function f fulfill the degree-condition. Then we try to find an f_{B_1} -partition of B_1 , if this is not possible we record that (B_1, f_{B_1}) is a hard pair. Afterwards, we continue splitting end-blocks from D until we can either confirm that every block together with its corresponding function is a hard pair or we find an f_{B_i} -partition of some block B_i . Then, it follows from the proof of Claim 3 that we can extend this partition to an f -partition of the current D . Now we can greedily color the vertices of $B_{i-1} \setminus \{v_{i-1}\}$ and combine this partition with the partition of the previous D in order to get an f -partition of $D + B_{i-1}$. By repeating this procedure until B_1 gets added to D , we eventually obtain the desired f -partition.

Algorithm 6 Finding an f -partition of a block

Input: (D, f) where

- D is a Eulerian block, and
- $f_1(v) + f_2(v) + \dots + f_p(v) = d_D^+(v) = d_D^-(v)$ for all $v \in V(D)$

Output: Either that (D, f) is a hard pair, or an f -partition (D_1, D_2, \dots, D_p) of D

```
1: if  $(D, f)$  is a hard pair then
2:   return  $(D, f)$  is a hard pair
3: else if  $G(D)$  is a cycle then
4:    $(D_1, D_2, \dots, D_p) = \text{CYCLEPARTITION}(D, f)$ 
5:   return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ 
6: else if  $G(D)$  is a complete graph then
7:    $(D_1, D_2, \dots, D_p) = \text{COMPLETEDIGRAPHPARTITION}(D, f)$ 
8:   return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ 
9: else
10:  find an even cycle  $C$  of  $D(G)$  and vertex  $v \in V(C)$  that is not contained in a chord
11:  apply Algorithm 2 to obtain an  $f$ -partition  $(D_1, D_2, \dots, D_p)$  of  $D - v$ 
12:  if  $\min\{d_{D_j+v}^+(v), d_{D_j+v}^-(v)\} < f_j(v_i)$  for some  $j \in [1, p]$  then
13:     $D_j := D_j + v_i$ 
14:    return  $f$ -partition  $(D_1, D_2, \dots, D_p)$ .
15:  end if
16:  if all but one  $D_i$  are empty then
17:    find a vertex  $w$  with  $f_j(w) > 0$  for some  $j \neq i$ 
18:     $D_i := D_i - w$ ,  $D_j := (\{w\}, \emptyset)$ .
19:  end if
20:  let  $u, w$  be the neighbors of  $v$  on  $C$ 
21:  if  $u, w$  are contained in the same  $D_i$  then
22:    find vertex  $z$  with  $z \in V(D_j)$  and  $j \neq i$ 
23:    find cycle  $C'$  in  $G(D)$  containing  $z$  and the edge  $vw$ 
    and choose a cyclic ordering of the vertices of  $C'$ 
24:  end if
25:   $v^* := v$ 
26:  while  $u, w$  are contained in the same  $D_i$  or  $v \neq v^*$  do
27:    apply Algorithm 3 to  $(D, f, v^*, C', (D_1, \dots, D_p))$ 
28:    if Algorithm 3 returns  $f$ -partition  $(D_1, D_2, \dots, D_p)$  of  $D$  then
29:      return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ 
30:    else  $v^* :=$  right neighbor of  $v^*$  on  $C'$ 
31:    end if
32:  end while
33:  choose a cyclic ordering of the vertices of  $C$ 
34:  while true do
35:    apply Algorithm 3 to  $(D, f, v^*, C, (D_1, \dots, D_p))$ 
36:    if Algorithm 3 returns  $f$ -partition  $(D_1, D_2, \dots, D_p)$  of  $D$  then
37:      return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ 
38:    else  $v^* :=$  right neighbor of  $v^*$  on  $C$ 
39:    end if
40:  end while
41: end if
```

Algorithm 7 Main Algorithm

Input: (D, f) where

- D is a connected digraph
- $f_1(v) + f_2(v) + \dots + f_p(v) \geq \max\{d_D^+(v), d_D^-(v)\}$ for all $v \in V(D)$

Output: Either (D, f) is a hard pair, or f -partition (D_1, D_2, \dots, D_p) of D

```
1: if there is a vertex  $v^*$  with  $f_1(v^*) + f_2(v^*) + \dots + f_p(v^*) > \min\{d_D^+(v^*), d_D^-(v^*)\}$  then
2:   apply Algorithm 1 to  $(D, f, v^*)$  to obtain an  $f$ -partition  $(D_1, D_2, \dots, D_p)$  of  $D$ 
3:   return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ 
4: end if
5: Compute the block-decomposition of  $D$ 
6:  $i := 1$ 
7: while  $D \neq \emptyset$  do
8:   let  $B_i$  be an end-block of  $D$ , let  $v_i$  be the separating vertex contained in  $B_i$ .
9:   apply Algorithm 2 to  $(D, v_i)$  to obtain an  $f$ -partition of  $D - v_i$ .
10:  if  $\min\{d_{D_j+v_i}^+(v_i), d_{D_j+v_i}^-(v_i)\} < f_j(v_i)$  for some  $j \in [1, p]$  then
11:     $D_j := D_j + v_i$ 
12:    go to step 27
13:  else
14:     $f_{B_i}(v_i) := (f_{B_{i,1}}(v_i), f_{B_{i,2}}(v_i), \dots, f_{B_{i,p}}(v_i))$  with
15:     $f_{B_{i,j}}(v_i) = d_{B_i \cap (D_j+v_i)}^+(v_i)$  for  $j \in [1, p]$ 
16:     $f_{B_i}(w) = f(w)$  for  $w \in V(B_i) \setminus \{v_i\}$ .
17:    apply Algorithm 6 to  $(B_i, f_{B_i})$ 
18:    if Algorithm 6 returns  $f_{B_i}$ -partition of  $B_i$  then
19:       $D_i := D[(V(D_i) \cap (V(D) \setminus V(B))) \cup V(D'_i)]$  for  $i \in [1, p]$ 
20:      go to step 27
21:    else
22:       $D := D - (V(B) \setminus \{v_i\})$ ,  $f(v_i) := f(v_i) - f_{B_i}(v_i)$ , update block-decomposition
    of  $D$ ,  $i := i + 1$ 
23:    end if
24:  end if
25: end while
26: return  $D$  is  $f$ -hard and results from merging  $(B_1, f_{B_1}), (B_2, f_{B_2}), \dots, (B_{i-1}, f_{B_{i-1}})$ 
27: while  $i > 1$  do
28:   apply Algorithm 2 on  $(B_{i-1}, f_{B_{i-1}}, v_{i-1})$ , get an  $f_{B_{i-1}}$ -partition  $(D'_1, D'_2, \dots, D'_p)$ 
    of  $B_{i-1} - v_{i-1}$ 
29:    $D_i := D[V(D_i) \cup V(D'_i)]$  for  $i \in [1, p]$ 
30:    $i := i - 1$ 
31: end while
32: return  $(D_1, D_2, \dots, D_p)$  is an  $f$ -partition of  $D$ .
```
