STANDARD BASES OVER RINGS

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Dedicated to my Fiancé Muhammad Atiq Jamil

ABSTRACT. The theory of standard bases in polynomial rings with coefficients in a ring R with respect to local orderings is developed. R is a commutative Noetherian ring with 1 and we assume that linear equations are solvable in R.

1. INTRODUCTION

The aim of this paper is to develop the theory of standard bases especially for non-global orderings for polynomial rings with coefficients in a ring. We generalize the concept of Adams and Loustaunau ([1]) and Greuel and Pfister ([2]). In the book of Adams and Loustaunau the concept of Gröbner bases over polynomial rings with coefficients in a ring is developed, i.e, they consider standard bases with respect to global orderings. In the book of Greuel and Pfister the concept of standard bases over polynomial rings with coefficients in a field is developed, i.e, they consider also non-global orderings. We will generalize both concepts to a uniform theory.

Note that the theory of standard bases for ideals developed in this paper can also be also developed for modules without changing the proofs.

First of all we will prove that in the general case the computation of a standard basis with respect to a non-global ordering can be reduced using homogenization to the computation of a Gröbner basis with respect to a suitable global ordering. This is also here a very expensive way to compute a standard basis. Therefore later a more efficient algorithm similar to [2] is presented.

Standard basis computations over the rings \mathbb{Z} and $\mathbb{Z}/\langle m \rangle$ can be performed using the computer algebra system SINGULAR (cf. [4]).

Standard bases are useful in computing elimination of variables, intersection of ideals, quotient of ideals, kernel of the ring map.

This can be done using the method described in [2], all the results remain the same in our case.

2. Basic Definitions

Let R be a Noetherian commutative ring with 1 and $R[x_1, \ldots, x_n]$ the polynomial ring in n variables with coefficients in R. Assume that linear equations are *solvable* in R.

Definition 2.1. Linear equations are solvable in R if the following conditions are satisfied for any $a, a_1, \ldots, a_m \in R$:

(1) there is an algorithm to compute generators for the R-module

$$syz_R(a_1,\ldots,a_m) = \{(b_1,\ldots,b_m) \in R^m \mid a_1b_1 + \ldots + a_mb_m = 0\},\$$

(2) there is an algorithm to determine if $a \in \langle a_1, \ldots, a_m \rangle$,

(3) there is an algorithm to compute $b_1, \ldots, b_m \in R$ such that $a = b_1 a_1 + \ldots + b_m a_m$ if $a \in \langle a_1, \ldots, a_m \rangle$.

We will use the notations from [1] and [2] and repeat them here for the convenience of the reader.

Definition 2.2. A monomial ordering > is a total ordering on the set of monomials $Mon_n = \{x^{\alpha} = x_1^{\alpha_1} \cdot \ldots \cdot x_n^{\alpha_n} | \alpha = (\alpha_1, \ldots, \alpha_n) \in \mathbb{N}^n\}$ in n variables satisfying

$$x^{\alpha} > x^{\beta} \Longrightarrow x^{\gamma} x^{\alpha} > x^{\gamma} x^{\beta}$$

for all $\alpha, \beta, \gamma \in \mathbb{N}^n$. We also say, > is a monomial ordering on $R[x_1, \ldots, x_n]$, if > is a monomial ordering on Mon_n .

Example 2.3. The local lexicographical ordering $>_{ls}$ on Mon_n is defined as follows, $x^{\alpha} > x^{\beta} \iff \exists 1 \le i \le n, \ \alpha_1 = \beta_1, \ldots, \alpha_{i-1} = \beta_{i-1}, \ \alpha_i < \beta_i.$

Example 2.4. Let M be an invertible $(n \times n)$ -matrix with real coefficients and M_1, \ldots, M_n the rows of M. The matrix M defines an ordering > on Mon_n as follows:

 $x^{\alpha} > x^{\beta} \iff \exists 1 \leq i \leq n, \ M_1 \alpha = M_1 \beta, \dots, M_{i-1} \alpha = M_{i-1} \beta, \ M_i \alpha > M_i \beta.$ Every ordering can be defined by a matrix (cf. [2]).

Definition 2.5. Let > be a fixed monomial ordering. Writing $f \in R[x_1, \ldots, x_n]$, $f \neq 0$, in a unique way as a sum of non-zero terms

$$f = a_{\alpha_1} x^{\alpha_1} + a_{\alpha_2} x^{\alpha_2} + \ldots + a_{\alpha_s} x^{\alpha_s}, \qquad x^{\alpha_1} > x^{\alpha_2} > \ldots > x^{\alpha_s},$$

and $a_{\alpha_1}, a_{\alpha_2}, \ldots, a_{\alpha_s} \in R$. We call: $LM(f) := x^{\alpha_1}$, the leading monomial of f, $LE(f) := \alpha_1$, the leading exponent of f, $LT(f) := a_{\alpha_1}x^{\alpha_1}$, the leading term of f, $LC(f) := a_{\alpha_1}$, the leading coefficient of f, ecart(f) := deg(f) - deg(LM(f)). We define the leading monomial and the leading

We define the leading monomial and the leading term of 0 to be 0, and 0 to be smaller than any monomial.

Definition 2.6. Let > be a monomial ordering on Mon_n , > is a called global (resp. local) ordering if $x^{\alpha} > 1$ (resp. $x^{\alpha} < 1$) for all $\alpha \neq (0, ..., 0)$.

Definition 2.7. For any monomial ordering > on Mon_n ,

$$S_{>} := \{ u \in R[x_1, \dots, x_n] \setminus \{0\} | LT(u) = 1 \}$$

is a multiplicatively closed set,

$$R[x_1, \dots, x_n]_{>} := S_{>}^{-1}R[x_1, \dots, x_n] = \{\frac{f}{g} | f, g \in R[x_1, \dots, x_n] \text{ and } g \in S_{>}\}$$

is the localisation of $R[x_1, \ldots, x_n]$ with respect to $S_>$ and we call $R[x_1, \ldots, x_n]_>$ the ring associated to $R[x_1, \ldots, x_n]$ and >.

Definition 2.8. Let > be any monomial ordering. For $f \in R[x_1, \ldots, x_n]_>$, choose $u \in R[x_1, \ldots, x_n]$ such that LT(u) = 1 and $uf \in R[x_1, \ldots, x_n]$. Then

$$LM(f) := LM(uf),$$

$$LC(f) := LC(uf),$$

$$LT(f) := LT(uf),$$

$$LE(f) := LE(uf).$$

Definition 2.9. Let > be any monomial ordering then, for each $G \subset R[x_1, \ldots, x_n]_>$

 $L(G) = \langle \{LT(g) \mid g \in G\} \rangle_{R[x_1, \dots, x_n]}$

is called the leading ideal of G.

Definition 2.10. Let $I < R[x_1, ..., x_n]_>$. (1). A finite set $G \subset R[x_1, ..., x_n]_>$ is called a standard basis of I with respect to >if $G \subset I$, and L(I) = L(G). (2). G is called a strong standard basis¹, of I with respect to >, if $G \subset I$ and for any $f \in I \setminus \{0\}$ there exists $i \in \{1, ..., t\}$ such that $LT(g_i)$ divides LT(f). (3). If > is global, a standard basis is also called a Gröbner basis.

3. Computing Standard Bases By Using Homogenization

Theorem 3.1. Let $f_1, \ldots, f_m \in R[x_1, \ldots, x_n]$ generating the ideal $I < R[x_1, \ldots, x_n]_>$, where > is a monomial ordering given by a matrix M. Let $F_i := f_i^h \in R[t, x_1, \ldots, x_n]$ be the homogenization of f_i and $>_h$ be the monomial ordering given by the matrix $\begin{pmatrix} 1 & 1 & \cdots & 1 \\ 0 & & & \end{pmatrix}$

$$\begin{pmatrix} 0 \\ \vdots \\ 0 \end{pmatrix}.$$

Let $\{G_1, \dots, G_s\}$ be a Gröbner basis, respectively strong Gröbner basis of

¹Strong standard bases do not exist in general. They exist always if R is a principal ideal domain, (cf. theorem 6.4). A strong standard basis is a standard basis.

 $J = \langle F_1, \ldots, F_m \rangle$ with respect to $>_h$. If we denote $g_i = G_i|_{t=1}$ then, $\{g_1, \ldots, g_s\}$ is a standard basis, respectively strong standard basis of the ideal I with respect to >.

Proof. Assume that $\{G_1, \ldots, G_s\}$ is a *Gröbner basis* with respect to $>_h$. Let $f \in I \cap R[x_1, \ldots, x_n]$. Then there exists $u \in S_>$ and $\eta_i \in R[x_1, \ldots, x_n]$ such that

$$u \cdot f = \sum_{i=1}^{m} \eta_i \cdot f_i.$$

Then there exists ρ , $\rho_i \in \mathbb{Z}, \rho$, $\rho_i \ge 0$ such that

$$t^{\rho} \cdot u^h \cdot f^h = \sum_{i=1}^m t^{\rho_i} \cdot \eta_i^h \cdot f_i^h = \sum_{i=1}^m t^{\rho_i} \cdot \eta_i^h \cdot F_i.$$

As $\{G_1, \ldots, G_s\}$ is a Gröbner basis J so there exist $\xi_i \in R[x_1, \ldots, x_n]$ such that

$$LT(t^{\rho} \cdot u^{h} \cdot f^{h}) = \sum_{i=1}^{s} \xi_{i} \cdot LT(G_{i})$$

putting t = 1 we obtain the result.

Now assume that $\{G_1, \ldots, G_s\} \subseteq J$ is a strong Gröbner basis with respect to > and let $f \in I \cap R[x_1, \ldots, x_n]$. We want to show there exists *i* such that $LT(g_i)|LT(f)$ and that $g_1, \ldots, g_s \in IR[x_1, \ldots, x_n]_>$. As $\{G_1, \ldots, G_s\} \subseteq J$

$$G_i = \sum_{j=1}^m \xi_{i,j} \cdot F_j$$

with $\xi_{i,j} \in R[t, x_1, \dots, x_n]$. Put t = 1, we get

$$g_i = \sum_{j=1}^m \xi_{i,j|_{t=1}} \cdot f_j$$

this implies $g_1, \ldots, g_s \in IR[x_1, \ldots, x_n]_>$. Now for f there exists $w \in S_>$ such that

$$w \cdot f = \sum_{j=1}^{m} \eta_j \cdot f_j$$

for suitable $\eta_j \in R[x_1, \ldots, x_n]$. Then there exists $\rho, \rho_i \in \mathbb{Z}, \rho, \rho_i \geq 0$ such that

$$t^{\rho} \cdot w^{h} \cdot f^{h} = \sum t^{\rho_{i}} \cdot \eta_{i}^{h} \cdot f_{j}^{h} = \sum_{j=1}^{m} t^{\rho_{i}} \cdot \eta_{i}^{h} \cdot F_{j}.$$

As $\{G_1, \ldots, G_s\}$ is a strong Gröbner basis of J there exists i such that $LT(G_i)|LT(t^{\rho} \cdot w^h \cdot f^h)$. This implies $LT(g_i)|LT(f)$.

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4. Normal Form

The concept of a normal form with respect to a given system of polynomials is the basis of the theory of standard bases. Normal forms for non-global orderings are different and more complicated than normal forms for global orderings. This is already the case for polynomial rings over a field.

Definition 4.1. Let \mathcal{G} denote the set of all finite lists $G \subset R[x_1, \ldots, x_n]_>$,

$$NF: R[x_1, \ldots, x_n]_> \times \mathcal{G} \longrightarrow R[x_1, \ldots, x_n]_>, \ (f, G) \longmapsto NF(f|G),$$

is called a normal form on $R[x_1, \ldots, x_n]_>$ if for all $G \subset \mathcal{G}$ and for all $f \in R[x_1, \ldots, x_n]_>$, (1). NF(0|G) = 0, (2). $NF(f|G) \neq 0 \Longrightarrow LT(NF(f|G)) \notin L(G)$. (3). If $G = \{g_1, \ldots, g_s\}$, then there exists $u \in S_>$ such that r := uf - NF(f|G)has a standard representation with respect to G, that is,

$$r = \sum_{i=1}^{s} \xi_i \cdot g_i$$

for suitable $\xi_i \in R[x_1, \ldots, x_n]$ and $LM(r) = max_{i=1}^s \{LM(\xi_i)LM(g_i)\}.$

To prove the existence of a normal form we give an algorithm to compute it.

Definition 4.2. Let $h \in R[x_1, \ldots, x_n]$, $T \subseteq R[x_1, \ldots, x_n]$ be finite. If h = 0, let $S(T, h) = \emptyset$. If $h \neq 0$, let $S(T, h) := \{\sum_{g \in T} c_g \cdot x^{\alpha_g} \cdot g | LT(\sum_{g \in T} c_g \cdot x^{\alpha_g} \cdot g) = LT(h), c_g \in R \text{ and } x^{\alpha_g} \cdot LM(c_g \cdot g) = LM(h) \text{ if } c_g \neq 0\}.$

Remark 4.3. The set S(T,h) can be infinite. Algorithm 1 requires to choose an element of S(T,h) which is of minimal ecart. This is achieved by computing a generating system of S(T,h), which is a kind of a syzygy module, so by the assumption on R this can be done.

Algorithm 1. NF(f|G)Let > be any monomial ordering. Input: $f \in R[x_1, ..., x_n], G = \{g_1, ..., g_s\} \subset R[x_1, ..., x_n]$ with $g_i \neq 0 \quad \forall i = 1, ..., s$. Output: $h \in R[x_1, ..., x_n]$ a normal form of f with respect to G. • h := f;• T := G;• while $(S(T, h) \neq \emptyset)$ choose $k \in S(T, h)$ such that ecart(k) is minimal; if(ecart(k) > ecart(h)) $T := T \cup \{h\};$

h := h - k;

• return h;

Proposition 4.4. The algorithm terminates and defines a normal form.

Proof. Termination is proved by using homogenization with respect to t: We start with $h := f^h$ and $T^h := \{g^h | g \in G\}$. The while loop looks as follows

• while $(S(T^h, t^{\alpha}h) \neq \emptyset \text{ for some } \alpha)$

choose $k \in S(T^h, t^{\alpha}h)$ such that $\alpha \ge 0$ is minimal; if $(\alpha > 0)$ $T^h := T^h \cup \{h\};$ $h := t^{\alpha}h - k;$ $h := (h|_{t=1})^h;$

Since $R[x_1, \ldots, x_n]_>$ is Noetherian, there exists some positive integer N such that $L(T_v^h)$ becomes stable for $v \ge N$, where T_v^h denotes the set T^h after the v-th turn of the while loop. The next h, therefore, satisfies $LT(h) \in L(T_N^h) = L(T^h)$, whence, $LT(\sum_{g \in T^h} c_g x^{\alpha_g} g)$

= LT(h) for some $\sum_{g \in T^h} c_g x^{\alpha_g} g$ and $\alpha = 0$. That is, T_v^h itself becomes stable for $v \geq N$ and the algorithm continues with fixed T^h . Then it terminates, since $>_h$ is a well ordering on $R[t, x_1, \ldots, x_n]$. For the correctness consider the v-th while loop of Algorithm 1. There we create (with $h_0 := f$)

$$h_v = h_{v-1} - \sum_{g \in T} c_g x^{\alpha_g} g$$

for some $\sum_{g \in T} c_g x^{\alpha_g} g$ such that $LT(\sum_{g \in T} x^{\alpha_g} g) = LT(h_{v-1})$ and $x^{\alpha_g} LM(c_g g) = LM(h_{v-1})$ if $c_g \neq 0$, from the construction of T^h we have

$$\sum_{g \in T} c_g x^{\alpha_g} g = \sum_{i=1}^{s} c_i x^{\alpha_i} g_i + \sum_{j=0}^{v-2} d_j x^{\beta_j} h_j$$

where

$$c_g = \begin{cases} c_i & g = g_i \\ d_j & g = h_j \end{cases}$$
$$\alpha_g = \begin{cases} \alpha_i & g = g_i \\ \beta_j & g = h_j \end{cases}$$

which implies

$$h_{v} = h_{v-1} - \left(\sum_{i=1}^{s} c_{i} x^{\alpha_{i}} g_{i} + \sum_{j=0}^{v-2} d_{j} x^{\beta_{j}} h_{j}\right).$$

Especially for $v \ge 2$ we have $LM(f) > LM(h_{v-1}) = LM(x^{\alpha_i}c_ig_i)$ if $c_i \ne 0$ and $LM(f) > LM(h_{v-1})$

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 $= LM(x^{\beta_j}d_jh_j)$ if $d_j \neq 0$. This implies especially $x^{\beta_j} < 1$. Suppose by induction, that in the first v - 1 steps $(v \ge 1)$ we have constructed standard representations

$$u_j f = \sum_{i=1}^s a_i^{(j)} g_i + h_j, \ u_j \in S_>, \ a_i^{(j)} \in R[x_1, \dots, x_n],$$

with $LM(u_j f - h_j) = max_{i=1}^s \{LM(a_i^{(j)})LM(g_i)\}$ for $0 \le j \le v - 1$, starting with $u_0 = 1, a_i^{(0)} = 0.$

Consider this standard representation for j = v - 1. We replace h_{v-1} by $h_v + (\sum_{i=1}^s c_i x^{\alpha_i} g_i + \sum_{j=0}^{v-2} d_j x^{\beta_j} h_j)$, hence we obtain

$$u_{v-1}f = \sum_{i=1}^{s} a_i^{(v-1)}g_i + h_v + (\sum_{i=1}^{s} c_i x^{\alpha_i}g_i + \sum_{j=0}^{v-2} d_j x^{\beta_j}h_j),$$

where each h_i has a standard representation as above

$$u_{v-1}f = \sum_{i=1}^{s} a_i^{(v-1)}g_i + h_v + (\sum_{i=1}^{s} c_i x^{\alpha_i}g_i + \sum_{j=0}^{v-2} d_j x^{\beta_j}(u_j f - \sum_{i=1}^{s} a_i^{(j)}g_i))$$

$$(u_{v-1} - \sum_{j=0}^{v-2} d_j x^{\beta_j}u_j)f = (\sum_{i=1}^{s} a_i^{(v-1)}g_i + \sum_{i=1}^{s} c_i x^{\alpha_i}g_i - \sum_{j=0}^{v-2} d_j x^{\beta_j} \sum_{i=1}^{s} a_i^{(j)}g_i) + h_v.$$

$$= \sum_{i=1}^{s} (a_i^{(v-1)} + c_i x^{\alpha_i} - \sum_{j=0}^{v-2} d_j x^{\beta_j}a_i^{(j)})g_i + h_v.$$

Let $u_v := (u_{v-1} - \sum_{j=0}^{v-2} d_j x^{\beta_j} u_j)$ and $a_i^{(v)} := a_i^{(v-1)} + c_i x^{\alpha_i} - \sum_{j=0}^{v-2} d_j x^{\beta_j} a_i^{(j)}$. We have to prove that $u_v \in S_>$ and

$$u_v f = \sum_{i=1}^{s} a_i^{(v)} g_i + h_v$$

is a standard representation, i.e,

 $LM(f) = LM(u_v f - h_v) = \max_{i=1}^s \{LM(a_i^{(v)}) LM(g_i)\}.$ Since $x^{\beta_j} < 1$ in case $d_j \neq 0$ it follows $u_v \in S_>$.

Since $LM(f) > LM(x^{\alpha_k}c_kg_k)$ if $c_k \neq 0$ and $LM(f) \ge LM(a_k^{(j)}g_k), x^{\beta_j} < 1$, it follows $LM(a_k^{(v)}g_k) \le LM(f)$. If $LM(a_i^{(v-1)}g_i) = LM(f)$ then with the same argument we obtain $LM(a_i^{(v)}g_i) = LM(a_i^{(v-1)}g_i) = LM(f)$. \Box

Example 4.5. We consider $R = \mathbb{Z}$ and use the local lexicographical ordering ls with x > y in $\mathbb{Z}[x, y]$. Let $f = xy^4 - 12x^2$ then ecart(f) = 0 and let $G = \{f_1, f_2, f_3\}$ where

$$f_1 = -3x + xy, f_2 = y^2 - 2x^2y, f_3 = 6x^2 - x^3y^2$$

then $ecart(f_1) = 1$, $ecart(f_2) = 1$, $ecart(f_3) = 3$. In step 1: $h_0 = xy^4 - 12x^2$, T := G and $S(T, h_0) = \{ky^4 f_1 + (3k+1)xy^2 f_2, k \in \mathbb{Z}\}.$

 $\begin{array}{l} \mbox{All elements in } S(T,h_0) \ have \ ecart \ 1 \ and \ we \ choose \ xy^4 - xy^5 + 4x^3y^3 \in S(T,h_0).\\ \mbox{Since } ecart(h_0) < ecart(xy^4 - xy^5 + 4x^3y^3) \ we \ have \ to \ enlarge \ T:\\ T = T \cup \{f_4 := h_0\} \ and\\ h_1 = xy^4 - 12x^2 - (xy^4 - xy^5 + 4x^3y^3) = xy^5 - 12x^2 - 4x^3y^3.\\ \mbox{In step } 2:\\ h_1 = xy^5 - 12x^2 - 4x^3y^3 \ with \ ecart(h_1) = 0,\\ S(T,h_1) = \{ky^5f_1 + lxy^3f_2 + (1 + 3k - l)yf_4, \ k, \ l \in \mathbb{Z}\}.\\ \mbox{We choose } xy^5 - 12x^2y \in S(T,h_1) \ with \ minimal \ ecart \ 0 \ and \ obtain \end{array}$

We choose $xy^3 - 12x^2y \in S(T, h_1)$ with minimal ecarl 0 and obtain $h_2 = xy^5 - 12x^2 - 4x^3y^3 - (xy^5 - 12x^2y) = -12x^2 + 12x^2y - 4x^3y^3$. In step 3: $h_2 = -12x^2 + 12x^2y - 4x^3y^3$ with $ecart(h_2) = 4$, $S(T, h_2) = \{(2k+4)xf_1 + kf_2, k \in \mathbb{Z}\}.$

We choose $-12x^2 + 4x^2y \in S(T, h_2)$ with minimal ecart 1 and obtain $h_3 = -12x^2 + 12x^2y - 4x^3y^3 - (-12x^2 + 4x^2y) = 8x^2y - 4x^3y^3$. In step 4: $h_3 = 8x^2y - 4x^3y^3$ and $LT(h_3) \notin L(T)$, thus $NF(f|G) = 8x^2y - 4x^3y^3$.

Remark 4.6. Assume R has the following property: $c = a_1x_1 + \ldots + a_sx_s$ is solvable in R if and only if there exists j and $x \in R$ such that $c = a_jx$.

Then normal form algorithm is similar to the corresponding normal form algorithm for a polynomial ring over a field, i.e, S(T,h) can be replaced by $S(T,h) = \{g \in T \mid LT(g) \mid LT(h)\}$. In this case each standard basis is a strong standard basis. If R is a discrete valuation ring or $R = \mathbb{Z}/\langle p^n \rangle$, p a prime number, then R has the property above.

5. Computing Standard Bases

Theorem 5.1. Let $I < R[x_1, ..., x_n]$ and let $G = \{g_1, ..., g_t\}$ be a set of non-zero polynomials in I. Then the following are equivalent. (1). L(G) = L(I).

- (2). For any polynomial $f \in R[x_1, \ldots, x_n]_>$, we have $f \in I$ if and only if NF(f|G) = 0.
- (3). For all $f \in I$, $uf = \sum_{i=1}^{t} h_i g_i$ for some polynomials $u, h_1, \ldots, h_t \in R[x_1, \ldots, x_n]$ such that LT(u) = 1 and $LM(f) = max_{i=1}^{t} \{LM(h_i)LM(g_i)\}.$

Proof. (1) \implies (2). We know that if NF(f|G) = 0, then $f \in I$. Conversely assume that $f \in I$. Let r = NF(f|G) and assume $r \neq 0$. Since $G \subset I$ we have $r \in I$. This implies $LT(r) \in L(I) = L(G)$. This is a contradiction to Definition 9. (2) \implies (3). This is obvious from the Definition of normal form.

(3) \implies (1). For $f \in I$ we need to show that $LT(f) \in L(G)$. We have that $uf = \sum_{i=1}^{t} h_i g_i$ such that $LM(f) = max_{i=1}^t \{LM(h_i)LM(g_i)\}$. It is easily seen that $LT(f) = \sum LT(h_i)LT(g_i)$ where the sum is over all *i* satisfying $LM(f) = LM(h_i)LM(g_i)$. Thus $LT(f) \in L(G)$, as desired.

Corollary 5.2. If G is a standard basis of I then I is generated by G.

Proof. Clearly $\langle g_1, \ldots, g_t \rangle \subset I$, since each g_i is in I. For the other inclusion, let $f \in I$. By Theorem 5.1, NF(f|G) = 0 and hence $uf \in \langle g_1, \ldots, g_t \rangle_{R[x_1, \ldots, x_n]}$ for a suitable $u \in S_>$.

Remark 5.3. Let K be a field and $R = K[y_1, \ldots, y_m]$ be the polynomial ring with variables y_1, \ldots, y_m . Let > be a product ordering on $K[y_1, \ldots, y_m, x_1, \ldots, x_n]$ such that the x_1, \ldots, x_n dominate y_1, \ldots, y_m and the restriction of > to R is global. Let $I < K[y_1, \ldots, y_m, x_1, \ldots, x_n]_>$ an ideal and $G = \{f_1, \ldots, f_s\}$ a standard basis of I. Let >₁ be the ordering on $R[x_1, \ldots, x_n]$ considered as polynomial ring with coefficients in R induced by >. Then G is a standard basis of I with respect to >₁.

Proof. Let $f \in I$ then $NF_{>}(f|G) = 0$ ($NF_{>}$ the normal form with respect to >). Analyzing the algorithm for $NF_{>}$ we obtain $a_1, \ldots, a_s \in R$ such that $\sum_{i=1}^s a_i LT_{>_1}(f_i) = LT_{>_1}(f)$. Here $LT_{>_1}$ is the leading term with respect to the ordering $>_1$.

We use the following Definition from [1].

Definition 5.4. Given monomials $x^{\alpha_1}, \ldots, x^{\alpha_s}$ and non-zero elements c_1, \ldots, c_s in R set $L = (c_1 x^{\alpha_1}, \ldots, c_s x^{\alpha_s})$. Then for a given monomial x^{α} , we call a syzygy $h = (h_1, \ldots, h_s) \in syz(L) \subset (R[x_1, \ldots, x_n])^s$ homogeneous of degree x^{α} provided that each h_i is a term and $x^{\alpha_i} LM(h_i) = x^{\alpha}$ for all i such that $h_i \neq 0$.

Theorem 5.5. Let $G = \{g_1, \ldots, g_t\}$ be a set of non-zero polynomials in $R[x_1, \ldots, x_n]$. Let \mathcal{B} be a homogeneous generating set for $syz(LT(g_1), \ldots, LT(g_t))$. Then G is a standard basis for the ideal $\langle g_1, \ldots, g_t \rangle R[x_1, \ldots, x_n]_>$ if and only if for all $(h_1, \ldots, h_t) \in \mathcal{B}$, we have

$$NF(\sum_{i=1}^{t} h_i g_i | G) = 0.$$

Proof. If G is a standard basis, then by Theorem 5.1,

$$NF(\sum_{i=1}^{t} h_i g_i | G) = 0.$$

Conversely, let $g \in \langle g_1, \ldots, g_t \rangle R[x_1, \ldots, x_n]_>$, then there exists $u \in S_>$ and $u_i \in R[x_1, \ldots, x_n]$ such that

$$ug = \sum_{i=1}^{t} u_i g_i.$$

Choose a representation as in above equation with $x^{\alpha} = max_{i=1}^{t}(LM(u_i)LM(g_i))$ minimal. Since by Theorem 5.1, we need to show that $LM(g) = x^{\alpha}$. We assume $LM(g) < x^{\alpha}$ and show that we can obtain an equation for g with a smaller value for x^{α} . Let $S = \{i \in \{1, \ldots, t\} | LM(u_i)LM(g_i) = x^{\alpha}\}$. Then

$$\sum_{i \in S} LT(u_i)LT(g_i) = 0.$$

Let $\mathbf{h} = \sum_{i \in S} LT(u_i)\mathbf{e}_i$ (where $\mathbf{e}_1 = (1, \ldots, 0), \ldots, \mathbf{e}_t = (0, \ldots, 1)$ is a generating set for $R[x_1, \ldots, x_n]^t$). Then $\mathbf{h} \in syz(LT(g_1), \ldots, LT(g_t))$ and \mathbf{h} is homogeneous of degree x^{α} . Now let $\mathcal{B} = {\mathbf{h}_1, \ldots, \mathbf{h}_l}$, with $\mathbf{h}_j = (h_{1,j}, \ldots, h_{t,j})$ then $\mathbf{h} = \sum_{j=1}^l a_j \mathbf{h}_j$. Since \mathbf{h} is a homogeneous syzygy, we may assume that the a_j are terms such that $LM(a_j)LM(h_{i,j})LM(g_i) = x^{\alpha}$ for all i, j such that $a_jh_{i,j} \neq 0$. By hypothesis, for each $j, NF(\sum_{i=1}^t h_{i,j}g_i|G) = 0$. Thus by Theorem 5.1, for each $j = 1, \ldots, l$ there exist $w_j \in S_>$ and $v_{i,j} \in R[x_1, \ldots, x_n]$ such that

$$w_j \sum_{i \in S} h_{i,j} g_i = \sum_{i=1}^{l} v_{i,j} g_i,$$

and

$$max_{i=1}^{t}LM(v_{i,j}g_{i}) = LM(\sum_{i=1}^{t}h_{i,j}g_{i}) < max_{i=1}^{t}LM(h_{i,j})LM(g_{i}).$$

The latter strict inequality is because $\sum_{i=1}^{t} h_{i,j} LT(g_i) = 0$. We may assume $w = w_j$ for all j. Thus,

$$wug = \sum_{i=1}^{l} wu_i g_i$$
$$= \sum_{i\in S} wLT(u_i)g_i + \sum_{i\in S} (wu_i - wLT(u_i))g_i + \sum_{i\notin S} wu_i g_i$$
$$= \sum_{j=1}^{l} \sum_{i\in S} wa_j h_{i,j} g_i + terms lower than x^{\alpha}$$
$$= \sum_{i=1}^{l} \sum_{i=1}^{t} a_j v_{i,j} g_i + terms lower than x^{\alpha}.$$

We have $max_{i,j}LM(a_j)LM(v_{i,j})LM(g_i) < max_{i,j}LM(a_j)LM(h_{i,j})LM(g_i) = x^{\alpha}$. We have a representation of g as a linear combination of the g_i such that the maximum of the leading monomials of any summand is less than x^{α} . Thus the theorem is proved.

As a consequence of Theorem 5.5, we obtain that the following algorithm computes a standard basis for a given ideal I in $R[x_1, \ldots, x_n]_>$.

Algorithm 2. Standardbasis(G)*Input*: $F = \{f_1, \ldots, f_r\} \subset R[x_1, \ldots, x_n]$ with $f_i \neq 0$ $(1 \le i \le r)$, **Output**: G a Standard basis for $\langle f_1, \ldots, f_r \rangle R[x_1, \ldots, x_n]_>$. • G := F: • P := a finite homogeneous generating set (considered as ordered set) for $syz(\{LT(f_i)\}_{1 \le i \le r});$ • $while(P \neq \emptyset)$ Let $G = \{f_1, ..., f_k\}$ choose $(s_1,\ldots,s_k) \in P$ $P := P \setminus \{(s_1, \ldots, s_k)\};$ $h := NF(\sum_{i=1}^k s_i f_i | G);$ $if(h \neq 0);$ $f_{k+1} := h;$ $G := \{f_1, \ldots, f_{k+1}\};$ H:=a finite homogeneous generating set for $syz({LT(f_i)}_{1 \le i \le k+1});$ $P := (P \times \{0\}) \cup \{h = (h_1, \dots, h_{k+1}) \in H | h_{k+1} \neq 0\};$

• return G;

Using ideas of M.Möller ([3]) Adams and Loustaunau propose a more efficient algorithm for computing Gröbner bases. This applies also in our situation with the same proof.

We use the following Definition and Theorem from [1].

Definition 5.6. Let $x^{\alpha_1}, \ldots, x^{\alpha_s}$ be a set of monomials. For any subset $J \subseteq \{1, \ldots, s\}$, set $x^{\gamma_J} = lcm(x^{\alpha_j}|j \in J)$. We say that J is saturated with respect to $x^{\alpha_1}, \ldots, x^{\alpha_s}$ provided that for all $j \in \{1, \ldots, s\}$ if x^{α_j} divides x^{γ_J} , then $j \in J$. For any subset $J \in \{1, \ldots, s\}$ we call the saturation of J the set J_1 consisting of all $j \in \{1, \ldots, s\}$ such that x^{α_j} divides x^{γ_J} . (Note that $x^{\gamma_J} = x^{\gamma_{J_1}}$.)

Theorem 5.7. (cf. [1], page 214) Given monomials $x^{\alpha_1}, \ldots, x^{\alpha_s}$ and non-zero elements c_1, \ldots, c_s in R. For each set $J \subseteq \{1, \ldots, s\}$, which is saturated with respect to $x^{\alpha_1}, \ldots, x^{\alpha_s}$, let $\mathcal{B}_J = \{b_{1,J}, \ldots, b_{\nu_J,J}\}$ be a set of generators of the R-module of syzygies $syz_R(c_j|j \in J)$. (Note that each of the vectors $b_{\mu,J}$ is in the R-module $R^{|J|}$, where |J| denotes the cardinality of J). For each such $b_{\nu,J}$, denote its jth coordinate, for $j \in J$, by $b_{j\nu,J}$ and $x^{\gamma_J} = lcm(x^{\alpha_j}|j \in J)$. Set

$$s_{\nu,J} = \sum_{j \in J} b_{j\nu,J} \frac{x^{\gamma_J}}{x^{\alpha_j}} e_j$$

(Note that each of the vectors $s_{\nu,J}$ is in $R[x_1, \ldots, x_n]^s$). Then the vectors $s_{\nu,J}$, for J running over all such saturated subsets of $\{1, \ldots, s\}$, and, $1 \leq \nu \leq \nu_J$, forms a homogeneous generating set for the syzygy module $syz(c_1x^{\alpha_1}, \ldots, c_sx^{\alpha_s})$.

Example 5.8. We consider $R = \mathbb{Z}$ and let $c_1 x^{\alpha_1} = 3xy^2$, $c_2 x^{\alpha_2} = 7xyz$, $c_3 x^{\alpha_3} = 2y^2z^2$. The saturated subsets of $\{1, 2, 3\}$ are $\{1\}, \{2\}, \{3\}, \{1, 2\}$ and $\{1, 2, 3\}$. Since \mathbb{Z} is an integral domain, the singletons $\{1\}, \{2\}, \{3\}$ do not give rise to any non-zero syzygy.

For $J = \{1, 2\}$ we need to solve in $R = \mathbb{Z}$ the equation $3b_1 + 7b_2 = 0$. The module of all solutions is generated by (7, -3). Since $x^{\gamma_J} = xy^2 z$, the corresponding syzygy is $s_{\nu,J} = 7\frac{xy^2 z}{xy^2} e_1 + 3\frac{xy^2 z}{xyz} e_2 = (-7z, 3y)$.

Now for $J = \{1, 2, 3\}$ we need to solve $3b_1 + 7b_2 + 2b_3 = 0$. The module of all solutions is generated by (-4, 2, -1) and (-7, 3, 0). Then with $x^{\gamma_J} = xy^2z^2$ we obtain the syzygies are

$$\begin{aligned} s_{\nu,J} &= -4 \frac{xy^2 z^2}{xy^2} \, \mathbf{e}_1 + 2 \frac{xy^2 z^2}{xyz} \, \mathbf{e}_2 \, - \frac{xy^2 z^2}{y^2 z^2} \, \mathbf{e}_3 = (-4z^2, 2yz, -x). \\ s_{\nu,J} &= -7 \frac{xy^2 z^2}{xy^2} \, \mathbf{e}_1 \, + \, 3 \frac{xy^2 z^2}{xyz} \, \mathbf{e}_2 \, = (-7z^2, 3yz, 0). \\ So \ we \ obtain \ that \end{aligned}$$

$$syz(3xy^2, 7xyz, 2y^2z^2) = \langle (-7z, 3y, 0), (-4z^2, 2yz, -x), (-7z^2, 3yz, 0) \rangle$$

The theorem is the basis of the following modified standard basis algorithm.

Algorithm 3. Standardbasis(G) Input: $F = \{f_1, \ldots, f_s\} \subseteq R[x_1, \ldots, x_n]$ with $f_i \neq 0$ ($1 \le i \le s$), Output: G a standard basis for $\langle f_1, \ldots, f_s \rangle R[x_1, \ldots, x_n]_>$.

- G := F;
- $\sigma := 1;$
- m := s;
- $while(\sigma \leq m)$

Compute $S = \{$ subsets of $\{1, ..., \sigma\}$, saturated with respect to $LM(f_1), ..., LM(f_s)$, which contain $\sigma \}$; for $(J \in S)$ $x^{\gamma} := lcm(LM(f_j)|j \in J)$; Compute a generating set $\{b_{i,J}, i = 1, ..., \mu_J\}$ for $\langle LC(f_j)|j \in J, j \neq \sigma \rangle_R : \langle LC(f_\sigma) \rangle_R$ for $(i := 1, ..., \mu_J)$ Compute $b_j \in R, j \in J, j \neq \sigma$ such that $\sum_{j \in J, j \neq \sigma} b_j LC(f_j) + b_{i,J}LC(f_\sigma) = 0$

$$\begin{split} r := NF(\sum_{j \in J, j \neq \sigma} b_j \frac{x^{\gamma}}{LM(f_j)} f_j + b_{i,J} \frac{x^{\gamma}}{LM(f_\sigma)} f_\sigma | G); \\ if(r \neq 0) \\ f_{m+1} &:= r; \\ G := G \cup \{f_{m+1}\}; \\ m &:= m+1; \\ \sigma := \sigma + 1; \end{split}$$

• return
$$G$$
;

Example 5.9. We consider $R = \mathbb{Q}[x, y]$ with the local lexicographical ordering ls in R[z]. Let $G = \{f_1, f_2\}$ and $I = \langle G \rangle$ where

$$f_1 = y - x^3 z^2$$
 and $f_2 = x^2 y - z$

Then $ecart(f_1) = 2$ and $ecart(f_2) = 1$. In step 1 ($\sigma = 1$): $S = \{\{1\}\} = saturated subsets of \{1\} containing 1.$ Since R is a domain and f_1 , f_2 are irreducible we have no non trivial syzygy. In step 2 ($\sigma = 2$): $S = \{\{1,2\}\} = saturated subsets of \{1,2\} containing 2.$ $J = \{1, 2\}.$ $x^{\gamma} = lcm(1,1) = 1$ A generating set for $\langle y \rangle : \langle x^2 y \rangle$ is $\{1\}$. The solution of $yb_1 + x^2y = 0$ is $b_1 = -x^2$. $h = -x^{2}(y - x^{3}z^{2}) + 1(x^{2}y - z) = -z + x^{5}z^{2}$ which is reduced with respect to $G = \{f_1, f_2\}.$ $f_3 := h$, and $G = \{f_1, f_2, f_3\}.$ In step 3 ($\sigma = 3$): $S = \{\{1, 2, 3\}\} = saturated subsets of \{1, 2, 3\} containing 3.$ $J = \{1, 2, 3\}.$ $x^{\gamma} = lcm(1, 1, z) = z.$ A generating set for $\langle y, x^2y \rangle : \langle -1 \rangle$ is $\{y\}$. The solution of $yb_1 + x^2yb_2 + x^2y = 0$ is $b_1 = 1, b_2 = 0$. $h = z(y - x^{3}z^{2}) + y(-z + x^{5}z^{2}) = x^{5}yz^{2} - x^{3}z^{3}.$ NF(h|G) = 0.So $G = \{f_1, f_2, f_3\}$ is a standard basis.

Remark 5.10. Let R be a local ring of the type $K[y_1, \ldots, y_m]_{\langle y_1, \ldots, y_m \rangle}/I$, where $I < K[x_1, \ldots, x_n]_{\langle x_1, \ldots, x_n \rangle}$. Then we can compute Hilbert-Samuel function of ideals in $R[x_1, \ldots, x_n]_>$ by using the same method described in [2]. Carlo Traverso [5], used the Hilbert-Samuel function to speed up the Buchberger Algorithm. The same method can be applied in our case too.

6. Standard Bases Over Principal Ideal Domains

If R is a principal ideal domain (for short PID) then there is a standard basis algorithm similar to the corresponding algorithm for a polynomial ring over a field (cf. [2], page 54) with the following notion of the s-polynomial.

Definition 6.1. Let $f, g \in R[x_1, \dots, x_n] \setminus \{0\}$. $spoly(f,g) := \frac{lcm(LT(f), LT(g))}{LT(f)}f - \frac{lcm(LT(f), LT(g))}{LT(g)}g.$

This is a consequence of [1], Proposition 4.5.3 and Theorem 5.5.

Example 6.2. We consider $R = \mathbb{Z}$ with the local lexicographical ordering ls with x > y in $\mathbb{Z}[x, y]$. Let $I = \langle f_1, f_2 \rangle$ where $f_1 = -3y + xy$ and $f_2 = y^2 - 2x$. Initialization: $G = \{f_1, f_2\}, \ \mathcal{G} = \{\{f_1, f_2\}\}.$ In step 1: $LT(f_1) = -3y, LT(f_2) = y^2,$ $c = lcm(c_1, c_2) = lcm(3, 1) = 3, \ x^{\gamma} = lcm(LM(f_1), LM(f_2)) = lcm(y, y^2) = y^2$ $h := spoly(f_1, f_2) = \frac{3y^2}{-3y}(-3y + xy) - \frac{3y^2}{y^2}(y^2 - 2x) = 6x - xy^2$ which is reduced with respect to G. So $G \cup \{f_3 := h\}$ and $\mathcal{G} = \{\{f_1, f_3\}, \{f_2, f_3\}\}.$ In step 2: $LT(f_1) = -3y, LT(f_3) = 6x,$ $c = lcm(c_1, c_3) = lcm(3, 6) = 6, x^{\gamma} = lcm(LM(f_1), LM(f_3)) = lcm(y, x) = xy$ $spoly(f_1, f_3) = \frac{6xy}{-3y}(-3y + xy) - \frac{6xy}{6x}(6x - xy^2) = xy^3 - 2x^2y = xy \cdot f_2$ $NF(xy^3 - 2x^2y|G) = 0.$ In step 3: $LT(f_2) = y^2, LT(f_3) = 6x,$ $c = lcm(c_2, c_3) = lcm(1, 6) = 6, \ x^{\gamma} = lcm(LM(f_2), LM(f_3)) = lcm(y^2, x) = xy^2$ $spoly(f_2, f_3) = \frac{6xy^2}{y^2}(y^2 - 2x) - \frac{6xy^2}{6x}(6x - xy^2) = xy^4 - 12x^2$

 $NF(xy^4 - 12x^2|G) = 0$. Since $xy^4 - 12x^2 = xy^2f_2 - 12x^2 + 2x^2y^2 = xy^2f_2 + 2xf_3$. So $G = \{f_1, f_2, f_3\}$ is a standard basis.

Example 6.3. We consider $R = \mathbb{Z}$ with the local degree lexicographical ordering ds (cf.[2]) with x > y > z in $\mathbb{Z}[x, y, z]$. Let $I = \langle f_1, f_2, f_3, f_4 \rangle$ where $f_1 = 15x^2 + 28y^2z^6$ and $f_2 = 3x^2y + 7yz^5$ $f_3 = 4xy^2 - 5xz^{10}$, $f_4 = -28y^3 + 35yz^{11}$. Similar to Example 6.2 we have to compute the normal form of the spoly's of all pairs (f_i, f_j) .

$$\begin{split} Nf(spoly(f_1, f_2)|G) &= 35yz^5 - 28y^3z^6 =: f_5, \ G = G \cup \{f_5\} \\ Nf(spoly(f_2, f_3)|G) &= 28y^2z^5 + 15x^2z^{10} =: f_6, \ G = G \cup \{f_6\} \\ Nf(spoly(f_3, f_5)|G) &= 35xz^{15} =: f_7, \ G = G \cup \{f_7\}. \\ The normal form of all the other spoly's is zero. \\ The standard basis is \ G &= \{f_1, f_2, f_3, f_4, f_5, f_6, f_7\}. \ G \ is \ a \ standard \ basis \ of \ I \ in \\ \mathbb{Z}[x, y, z]_{\langle x, y, z \rangle}. \\ Note \ that \ f_1 - zf_5 &= 15x^2 - 15x^2yz^{10} = 15x^2(1 - yz^{10}). \ Therefore \ 15x^2 \in I. \\ Similarly \ one \ can \ see \ that \ 28y^3, 28y^2z^5, 35yz^5 \ and \ 35xz^{15} \in I. \ This \ implies \ that \\ \{15x^2, 3x^2y + 7yz^5, 4xy^2 - 5xz^{10}, 28y^3, 28y^2z^5, 35yz^5, 35xz^{15}\} \ is \ a \ standard \ basis \ of \ I. \end{split}$$

Theorem 6.4. (cf. [1] Theorem 4.5.9 page 251) Let R be a PID, and I be an ideal of $R[x_1, \ldots, x_n]_>$. Assume that $\{f_1, \ldots, f_n\}$ is a standard basis for I. Let $LT(f_i) = c_i x^{\alpha_i}$, for a saturated subset J of $\{1, \ldots, s\}$, let $c_J = gcd(c_j|j \in J)$ and write $c_J = \sum_{j \in J} a_j c_j$ (any such representation will do). Also, let $x^{\alpha_J} = lcm(x^{\alpha_j}|j \in J)$. Then the set

$$\{f_J = \sum_{j \in J} a_j \frac{x^{\alpha_J}}{x^{\alpha_j}} f_j | J \text{ is a saturated subset of } \{1, \dots, s\}\}$$

is a strong standard basis for I. In particular, every non-zero ideal in $R[x_1, \ldots, x_n]_>$ has a strong standard basis.

7. Standard Bases In The Formal Power Series Rings

Finally we want to apply our results to compute standard bases in the formal power series ring $R[[x_1, \ldots, x_n]]$ with coefficients in a ring R. Let > be a local degree ordering, i.e, > is a local ordering and $x^{\alpha} > x^{\beta}$ implies $deg(x^{\alpha}) \leq deg(x^{\beta})$. A non-zero element $f \in R[[x_1, \ldots, x_n]]$ can be written as $\sum_{v=0}^{\infty} a_v x^{\alpha_v}$, $a_v \in R$, $a_0 \neq 0$ and $x^{\alpha_v} > x^{\alpha_{v+1}}$ for all v. As in Definition 2.5, we define LM(f), LE(f), LT(f), LC(f) and tail(f). As in Definition 2.11, we define a standard basis (respectively a strong standard basis) of an ideal $I \subseteq R[[x_1, \ldots, x_n]]$.

Proposition 7.1. Let $I < R[x_1, ..., x_n]$ and G is a standard basis (respectively a strong standard basis) of I with respect to >, where > is a local degree orderig. Then G is a standard basis (respectively a strong standard basis) of $IR[[x_1, ..., x_n]]$.

Proof. Let $\{g_1, \ldots, g_s\}$ be a standard basis of I and $\overline{g} = \sum_{i=1}^s \overline{a_i}g_i \in IR[[x_1, \ldots, x_n]], \overline{g} \neq 0$. Let c be an integer such that $LM(\overline{g}) \notin \langle x_1, \ldots, x_n \rangle^c$. Choose $a_i \in R[x_1, \ldots, x_n]$ such that $\overline{a_i} - a_i \in \langle x_1, \ldots, x_n \rangle^c$. Let $g = \sum_{i=1}^s a_i g_i$. Then $g \in I$ and $\overline{g} - g \in \langle x_1, \ldots, x_n \rangle^c$. This implies $LT(\overline{g}) = LT(g)$. If G is a strong standard basis for I then there exists i such that $LT(g_i)|LT(g) = LT(\overline{g})$, i.e. G is also a

strong standard basis of $IR[[x_1, \ldots, x_n]]$. Similarly it follows that a standard basis of I is a standard basis of $IR[[x_1, \ldots, x_n]]$.

8. PROCEDURES

Let $I < R[x_1, \ldots, x_n]$ be an ideal and $G = \{g_1, \ldots, g_m\}$ be a standard basis of I. Then we can compute a strong standard basis of I using the SINGULAR-procedures below.

```
LIB"poly.lib";
```

```
//computes the set of all subsets of \{1, \ldots, n\}
proc powerSet(int n)
ł
    list L,K,S;
    int i;
    if(n==0)
    ł
         L[1]=L;
         return(L);
    }
       if(n==1)
       {
         L[1]=L;
         L[2]=list(1);
         return(L);
       }
       S=powerSet(n-1);
       int r=size(S);
       S[r+1]=list(n);
       for(i=2;i<=r;i++)</pre>
       {
         K=S[i];
         K[size(K)+1]=n;
         S[size(S)+1]=K;
       }
       return(S);
}
proc satt(ideal I)
                             //computes the saturated subsets of \{1, \ldots, size(I)\}
                             //w.r.t {LM(f) \mid f \in I}
ł
    int j;
```

```
list K,J;
    I=lead(I);
    J=powerSet(size(I));
    for(j=2; j<=size(J); j++)</pre>
    {
         if(saturat(I,J[j]))
        {
            K[size(K)+1]=J[j];
        }
    }
    return(K);
}
proc specialGCD2(int a, int n)
ł
    int x=a mod n;
    if(x==0)return(list(0,1,n));
    list L=specialGCD2(n,x);
    return(list(L[2],L[1]-(a-x)*L[2]/n,L[3]));
}
proc specialGCD(list L)
                              //computes the gcd over \mathbb{Z}
ł
    int i;
    for(i=1;i<=size(L);i++)L[i]=int(L[i]);</pre>
    if(size(L)==1)return(L);
    if(size(L)==2)return(specialGCD2(L[1],L[2]));
    bigint p=L[size(L)];
    L=delete(L,size(L));
    list T=specialGCD(L);
    list S=specialGCD2(T[size(T)],p);
    for(i=1;i<=size(T)-1;i++)</pre>
    ł
         T[i]=T[i]*S[1];
    }
    p=T[size(T)];
    T[size(T)]=S[2];
    T[size(T)+1]=S[3];
    return(T);
}
```

```
// computes the coefficients a_j such that
proc coeffJ(list C,list J)
                                  // gcd(c_j|j \in J) = \sum_{j \in J} a_j c_j
{
    if(size(J)==1)return(list(1));
    int n=size(J);
    int i;
    list L,M;
    for(i=1;i<=n;i++)</pre>
    {
        L[size(L)+1]=C[J[i]];
    L=specialGCD(L);
    L=delete(L,size(L));
    return(L);
}
proc lcmJ(ideal X,list J) // computes lcm(f_j | f_j \in X, j \in J)
{
    poly p=X[J[1]];
    int i;
    for(i=2;i<=size(J);i++)</pre>
    ł
          p=lcmS(p,X[J[i]]);
    }
    return(p);
}
proc maxZ(intvec a,intvec b)
                                       // computes the maximum of two
                                 //integer vectors
ł
    int i,j;
    intvec c;
    for(i=1;i<=size(a);i++)</pre>
    {
         if(a[i]>=b[i])
         ł
             c[i]=a[i];
         }
         else
         {
             c[i]=b[i];
```

```
}
    }
     return(c);
}
                                       // computes the LCM of two monomials
proc lcmS(poly p,poly q)
    intvec a=leadexp(p);
    intvec b=leadexp(q);
    intvec c=maxZ(a,b);
     int i,j;
     poly s=1;
    for(i=1;i<=nvars(basering);i++)</pre>
    ł
         s=s*var(i)^c[i];
    }
     return(s);
}
proc leadTerm(ideal I)
                                // give a list containing two lists, first
                            //contains leading coefficients and second contains
    int i;
                                  //leading monomials
    list L;
    ideal J;
    I=lead(I);
    for(i=1;i<=size(I);i++)</pre>
    ł
         L[size(L)+1]=leadcoef(I[i]);
         J[size(J)+1]=leadmonom(I[i]);
    }
     return(list(L,J));
}
                             // delete from an ideal f_i such that
proc clean(ideal G)
                     // there exists f_j \in I with LM(f_j) divides LM(f_i)
{
    int i,j;
    while(i<=size(G)-1)</pre>
    {
         i++;
         j=i;
         while(j < size(G))</pre>
```

```
{
             j++;
            if((G[i]!=0)&&(G[j]!=0))
            {
                if (lead(G[j])/lead(G[i])!=0)
                {
                       G[j]=0;
                       G=simplify(G,2);
                       j--;
                }
                   if(G[j]!=0)
                   {
                       if (lead(G[i])/lead(G[j])!=0)
                       {
                             G[i]=0;
                             G=simplify(G,2);
                             i--;
                             break;
                       }
                  }
            }
        }
    }
        return(simplify(G,2));
}
proc strongSB(ideal I)
                               // computes a strong standard basis
                        // for an ideal if standard basis is given
{
    def R=basering;
    list rl=ringlist(R);
    rl[1]=0;
    def S=ring(rl);
    setring S;
    ideal I=imap(R,I);
    list L=leadTerm(I);
    list C=L[1];
    ideal X=L[2];
    list J=satt(X);
    int i,j;
    ideal G;
```

```
list P;
poly f,q;
for(i=1;i<=size(J);i++)
{
    P=coeffJ(C,J[i]);
    q=lcmJ(X,J[i]);
    f=0;
    for(j=1;j<=size(J[i]);j++)
    {
       f=f+q/X[J[i][j]]*I[J[i][j]]*P[j];
    }
    G[size(G)+1]=f;
}
setring R;
ideal G=imap(S,G);
return(clean(G));
```

}

```
Example 8.1. Consider Example 6.3, we now compute strong standard basis for 

I = \{15x^2, 3x^2y + 7yz^5, 4xy^2 - 5xz^{10}, 28y^3, 28y^2z^5, 35yz^5, 35xz^{15}\}.

> ring R = integer, (x, y, z), ds;

> ideal I = 15x2, 3x2y + 7yz5, 4xy2 - 5xz10, 28y3, 28y2z5, 35yz5, 35xz15;

> strongSB(I);

[1] = 15x2

[2] = 3x2y + 7yz5

[3] = 4xy2 - 5xz10

[4] = x2y2 - 7y2z5 - 5x2z10

[5] = 28y3

[6] = 35yz5

[7] = x2yz5 + 84yz10

[8] = 7y2z5

[9] = 35xz15

[10] = 5x2z15
```

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