



At execution of $A[f] := r$; on line prior to that labeled *schell*, f has not necessarily been assigned a value. f has a value if, and only if, the Boolean expression $B[k] > 0 \wedge B[k] < B[m]$ is true for at least one of the relevant values of k . In particular when matrix A is set up by $A[i] := i$; for each i the Boolean expression above is false on the first call.

ALGORITHM 202 is the best and fastest algorithm of the exicographic set so far published.

A collected comparison of these algorithms is given in Table I. t_n is the time for complete generation of $n!$ permutations. Times are scaled relative to t_8 for Algorithm 202, which is set at 100. Tests were made on an ICT 1905 computer. The actual time t_8 for Algorithm 202 on this machine was 100 seconds. r_n has the usual definition $r_n = t_n / (n \cdot t_{n-1})$.

TABLE I

Algorithm	t_6	t_7	t_8	r_6	r_7	r_8
87	118	—	—	—	—	—
102	2.1	15.5	135	1.03	1.08	1.1
130	—	—	—	—	—	—
202	1.7	12.4	100	1.00	1.00	1.00

CERTIFICATION OF: ALGORITHM 258 [H] TRANSPORT

[G. Bayer, *Comm. ACM* 8 (June 1965), 381]
ALGORITHM 293 [H]
TRANSPORTATION PROBLEM

[G. Bayer, *Comm. ACM* 9 (Dec. 1966), 869]
LEE S. SIMS (Recd. 21 Feb. 1967 and 17 Mar. 1967)
Kates, Peat, Marwick & Co., Toronto, Ont., Canada

Both of these algorithms were coded in Extended ALGOL 60 and tested on a Burroughs B5500. Three problems were solved correctly, one of them being of medium size (55×167). On this larger problem *transpl* was found to be about twice as fast as *transport*.

In coding and debugging *transpl* three apparent errors were found. In the right-hand column on page 870, after line 27 which is

$i := listu[u]$; $nlvi := nlv[i]$;
a line is missing. This line should read
for $s := (i-1) \times n + 1$ step 1 until $nlvi$ do
Also in the right-hand column, the line

$s4$;
should be inserted ahead of line -12, which begins
comment Step 4. A column j with $b[j]$ has been labeled, $b[j]$
On page 871, in the left-hand column, line -22 which reads
for $s := 1$ step 1 until n do
should read
for $s := l$ step 1 until n do

CERTIFICATION OF ALGORITHM 285 [H]
THE MUTUAL PRIMAL-DUAL METHOD
[Thomas J. Aird, *Comm. ACM* 9 (May 1966), 326]
H. SPÄTH (Recd. 13 Feb. 1967)
Institut für Neutronenphysik und Reaktortechnik,
Kernforschungszentrum, Karlsruhe, Germany

The procedure *Linearprogram* has been translated into FORTRAN II and successfully run on the IBM 7074 Computer. The following corrections had been made (the first two are merely typographical errors).

- P. 328, left column, 1 line after label B3:
reads:
if $A[row[k-1, i], col[k, 0]] > \text{then}$
should read:
if $A[row[k-1, i], col[k, 0]] > 0 \text{ then}$
- P. 328, left column, 1 line after label B4:
reads:
if $A[row[k-1, i], col[k, 0]] > \text{then}$
should read:
if $A[row[k-1, i], col[k, 0]] > 0 \text{ then}$
- P. 328, right column, after the end of the procedure *pickapivot* and before the label *NEXTPIVOT* there must be inserted the statement
 $col[0, 0] := 0$;
Otherwise $col[0, 0]$ has no assigned value when the procedure *subschema* is entered for the first time.

REMARK ON ALGORITHM 301 [S20]
AIRY FUNCTION [Gillian Bond and M.L.V. Pitteway,
Comm. ACM 10 (May 1967), 291]
M.L.V. Pitteway (Recd. 19 May 1967)
Brunel University, ACTON, W.3., England

The initial minus sign has been omitted from the line immediately following the line

end calculation of derivatives;
The statement should read
 $p := - (rtmdx/xi) \times (2 \times A[2] + 4 \times A[4] + 6 \times A[6] + 8 \times A[8] + 10 \times A[10]);$

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