



## ASSEMBLER IN A FORTRAN ENVIRONMENT WITH A NEW DEBUGGING AID

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### Introduction

An Assembly Language course is generally taken as a second course in computer science. Many students, after finishing successfully the first (introductory) course, still find the Assembly Language course extremely difficult and time-consuming. It is felt that the course should be designed in a manner that allows a student to acquire the basic ideas and techniques, while practicing in actual Assembly programming, within a reasonable amount of effort and time to avoid adverse effects on his other studies. We suggest that two important contributions towards the above goal are (1) employing an approach that conceives any Assembly program as a subroutine called by FORTRAN (or another higher level language used in the first course), and (2) providing basic easy-to-use debugging facilities. In this paper we report on the experience obtained while teaching the Assembly course at Rensselaer Polytechnic Institute according to the above considerations.

### Assembly in a FORTRAN Environment

The Assembly Language course at Rensselaer Polytechnic Institute is designed to introduce the student to the machine structure of the IBM 360, the techniques and ideas involved with programming at the Assembly level, the syntax of the IBM Assembler [1], interactions between the users' programs and the Operating System, macro instructions and Input/Output programming. It is felt that the student should be given the opportunity to run his own Assembler program as early as possible. If a student's first program is a main program written in Assembler, then he must be provided with Input/Output instructions, conversion procedures for numerical data types, and possibly some general information about macro instructions, and the referencing of files (data sets) by symbolic names in the Input/Output macro instructions.

Trying to base the first program on non-numeric data has the demerit that most students at this stage are neither motivated nor reasonably experienced with alphanumeric processing.

An easier and more productive approach is to begin the course with the fixed-point arithmetic instructions and base the first computer projects on processing integers. The reading of input and, more important, printing of results are done in a main program that calls the Assembler program as a subroutine, and is written in a higher level language like FORTRAN. This way the student does not get into Input/Output procedures, conversions, and operating system considerations until later in the course, thus concentrating at the initial phase on the basics of register usage, addressing, and the elementary operations. Printing the results by a FORTRAN program makes it easy for the beginning student to label his output with explanatory titles, and to edit his results almost without effort. The only subject that may seem difficult to some students is the linkage conventions for subprograms, but since these conventions are used partly also by "main" programs, it appears helpful to introduce the conventions at this early stage, and explain that, to some extent, a "main" program can also be considered a subroutine called by the operating system. Owing to the fact that the students taking the Assembler course have had experience with FORTRAN in their first computer course, this approach was very instructive and effective.

In [3] the handling of Input/Output at the beginning is done by invoking calls to reading and printing subroutines. The above suggestion of using the FORTRAN environment seems more direct and natural. The approach in [2] employs special instructor-prepared macro instructions for reading data from cards and printing lines. This does not

avoid the necessary conversions, and the student still has to be introduced to them quite early in the course.

As discussed below, running Assembly programs in a FORTRAN environment also helps in providing the students with debugging facilities, whose vitality at the initial stages cannot be overemphasized. In this respect, the approach is self-consistent: it puts students to actual program running at a very early stage and, at the same time, provides a tool without which debugging would be extremely difficult for the beginner.

### Debugging

Perhaps the single most important debugging technique is that of observing intermediate results during program execution. In most high level programming languages this presents no difficulty, just the insertion of output (print) statements at strategic locations in the program. Assembler Language programming, however, presents some rather unique problems. Many Assembler Language programs are intimately related with a particular computer's hardware features and the normal resources used to print diagnostic information may already be in use by the program, thus precluding their use for debugging purposes. On the IBM System/360, in particular, we refer to the registers used for subroutine linkages and addressing purposes.

The basic requirements for an effective debugging system are ease of use and transparency to the program being debugged. By ease of use, it is meant that as few changes as possible be made to the source program and Job Control Language statements. In addition, simplicity, especially for the beginner, should be a major consideration. The novice programmer should not be required to learn the intricate details of the operating system or input/output processing. As far as transparency is concerned, the debugging routines should have a minimal effect on the user's program, and ideally should not require special resource allocations either.

Although OS/360 does provide debugging facilities through the TRACE and SNAP macro instructions, it is felt that their use is difficult for the beginning student. The major problems are their nontransparency and presupposed knowledge of the operating system (JCL) and input/output procedures. Following are some specific examples:

1. The macros may alter General Registers 0, 1, 14, and 15. This is

a case of nontransparency to the program being debugged. The beginning student may be unaware of this fact, and might get confused if his program makes use of any of these registers.

2. Both macros assume the establishment of a Save Area and its associated address in General Register 13. This is a specialized resource allocation which the beginning student may very easily forget to do. In fact, this is one of the most common errors. Using either the SNAP or TRACE macro without a valid Save Area and address in Register 13 may cause unpredictable results or even an Addressing Exception.
3. Using the macros requires some knowledge of the operating system and input/output procedures, including the OPEN and DCB macros. This topic inevitably causes the most trouble to Assembler students, and it would seem that trying to teach this to the student in the early stages of the course might serve to overpower the student's comprehension. After all it is rather absurd that debugging be more difficult and complex than the original program.
4. In order to use the macros the student has to modify his Control Cards (JCL). We have found from experience that many new students are initially confused by the Control Cards necessary to run their program. Unfortunately the JCL cards must be modified in order to use the SNAP or TRACE macros. It would seem that this only adds to the complexity of the debugging process, and should be minimized if possible.

As a result of these problems with the currently available debugging software, we propose a macro instruction for the OS/360 Assembler Language, in which we have endeavored to simplify the debugging process as much as possible. This macro has proven to be very effective and useful in the course.

### Regdump

Macro REGDUMP has been written in response to a need for a simple debugging aid for Assembler Language programs for new students. The function of REGDUMP is to provide the user with a snapshot of the contents of the 16 General Registers, the 4 Floating Point Registers and the value of the current Condition Code. As explained above, we have found a FORTRAN environment very helpful for teaching a first course in Assembler, and REGDUMP was written with this in mind. REGDUMP can thus be used without modifications to the JCL cards

for the run, since the FORTRAN output routines called by REGDUMP are readily available for use. This, then, eliminates one of the sources of trouble to students.

Attacking the next problem, we have attempted to minimize the interference of REGDUMP with the normal operation of the Assembler program being debugged. As detailed below, it does not alter the contents of any registers or other information that the student is likely to ever use.

A typical example of using REGDUMP by a student would be as follows: the student has decided that it would be helpful to observe the contents of the registers at several points in his program. All he does is to insert at those points a card with 'REGDUMP' on it, and run his job. The program will be compiled and the code generated by REGDUMP will produce a register dump each time a REGDUMP call is encountered during execution.

The output from REGDUMP, illustrated by Fig. 1, provides the user with the following information at execution time:

1. General Registers 0-15 are printed in decimal and hexadecimal formats for convenience in interpreting both arithmetic and bit operations.
2. Floating Point Registers 0-6 are printed in both double precision decimal and hexadecimal formats. This facilitates checkout of floating point operations...
3. The current value of the Condition Code is printed in decimal in order to check conditional branches and follow the flow of logic in the program.
4. The current Program Mask is listed in hexadecimal for monitoring possible Program Interrupts and/or to determine what the exact results of some arithmetic operations will yield.
5. Some identification information (detailed below) of the REGDUMP call is printed to allow determining of which REGDUMP call in the course program is being executed, thus following program execution.
6. The address of the last instruction generated by the macro call is given for convenience in locating the macro coding when studying Loader Maps, dumps, or other object listings. This address aids the new student by giving him an address which lies between two consecutive source

instructions thereby specifying precisely when the dump occurred.

All output is printed in an easy to read format which is blocked off from the program output by a border thus facilitating differentiation between diagnostic and program output.

In its final form, REGDUMP was chosen to be a macro call. The advantage of using this approach over that of standard subroutine linkage is that the beginning student need not concern himself with all of the linkage details. In REGDUMP, all housekeeping and linkage to auxiliary subroutines is done internally. Thus it is possible to use REGDUMP without the knowledge of linkage conventions.

In order to avoid the pitfall of a student forgetting to establish a Save Area and its address in General Register 13, REGDUMP creates its own internal Save Area. This way, REGDUMP performs normally and the student might detect the trouble by checking the REGDUMP output that contains the contents of register 13. Unreliable results and a possible Addressing Exception are thus avoided.

One of the functions of REGDUMP is to preserve the current status of the program. This is done by copying all pertinent data to a temporary storage area and later restoring this information. In REGDUMP, this data must be saved before the general registers can be set up for subroutine linkage. For this purpose, an inline storage area is generated as part of the REGDUMP expansion.

To minimize the amount of inline code, it was decided to perform the output functions by calling a subroutine external to the macro. In our implementation, the external subroutine was written in FORTRAN because, as explained earlier, REGDUMP was used in a FORTRAN environment. Furthermore, it is much easier for the course instructor to control the output format in FORTRAN.

The inline area begins with the storage area, skipped over by a branch instruction, and followed by a sequence of instructions that stores the relevant information into the storage area, calls the printing subroutine, restores the registers and condition code, and resumes normal execution. The storage area is organized as follows:

<u>WORDS</u>	<u>BYTES</u>	<u>CONTENTS</u>
1-3	0-11	Linkage Data to FORTRAN output subroutine
	8	current program mask and condition code
4-19	12-75	general registers (0-15) when REGDUMP is called
20-27	76-107	floating point registers (0-6)
28-45	108-179	conventional save area for FORTRAN output subroutine
46	180-183	identification information

Since the inline storage area requires over 180 bytes, provisions were made to generate it only once in each control section. Therefore the macro creates a compilation-timetable of the control sections in which it has been called. In the present implementation, there is provision for up to 10 control sections in which the macro may be called. If more than 10 control sections are used (which is rare), the macro will still work, but the full storage area will be generated for each REGDUMP call from the 11th control section onward.

As discussed earlier, just using 'REGDUMP' in the operation-code field is sufficient to invoke the debugging routine. Even with such a simple call, each dump of information will be uniquely identified by an internally generated sequence number. Optionally, the user may specify his own identification number by entering it in the operand field. Other than misspelling 'REGDUMP', it is nearly impossible to commit an error in trying to call REGDUMP. However, should the user supply an invalid identification number, the macro will be sure to print a message to this effect and it will generate a number of its own and specify it to the user. If no identification number was supplied in the call, again the macro will note this and assign a unique integer for an ID and print it for the user's information. Examples of what REGDUMP calls might look like, with the macro reaction shown on the following line appear in Table 1.

To summarize, REGDUMP is absolutely transparent to the student's program in that it keeps unchanged the contents of all registers and the condition code. It provides clear identification information so that the student can easily distinguish between different REGDUMP calls within the same program.

Excellent legibility was made possible by blocking off the dump output from any program output that may occur between dumps. Thus by reading the output, the student may actually follow the dynamic flow of his program.

#### Concluding Remarks

The approach described above of teaching Assembler in a FORTRAN environment together with the REGDUMP macro has proved to be very effective and successful. A significant improvement has been noticed relative to previous years.

A natural addition to REGDUMP would be to enable dumping of storage areas. This was not done in the present implementation in order to make the calling sequence as simple as possible. This extension can be added in an obvious manner.

In order to implement REGDUMP, one should have the FORTRAN printing subroutine in a library consulted by the Loader. Such a standard library exists in almost every installation, and if not, can be appended through the use of a single JCL card.

Figure 2 contains a source listing of REGDUMP and Figure 3 shows the FORTRAN output subroutine.

#### References

1. IBM Corporation, OS Assembler Language, GC28-6514.
2. W. G. Rudd, Assembly Language Programming and the IBM 360 and 370 Computers, Prentice-Hall Inc., Englewood Cliffs (1976).
3. G. W. Struble, Assembler Language Programming: the IBM System/360 and 370, 2nd ed., Addison Wesley, Reading (1975).

- a. REGDUMP  
\*, REGDUMP ID = 0005 (no id given)  
(id assigned by REGDUMP)
- b. REGDUMP 25 (id of 25 specified)  
(no message given)
- c. REGDUMP XYZ (invalid id given)  
\*, ILLEGAL REGDUMP ID. 0008 USED. (legal id assigned)

Table 1.

REGDUMP 25 AT 1E61B6; CC IS 2, PGM MASK IS 2, REGS ARE:									
HEX	00000114	001E607E	00000015	40F43260	001E5700	00800000	001E5780	46000014	
DEC	276	1990782	00000015	1083741808	2053888	8388638	2054064	1174405140	
HEX	00000114	001E5700	001E5700	001E6000	421E6016	001E6010	421E6450	001E6010	
DEC	51760	2053688	2054136	1990656	1109286934	1991118	1109288016	1990672	
HEX	41000053	043C8165	3847405F	741D0000	4100E695	00000000	397137A5	00000000	
DEC	0.138155508601630 02	0.9965178760553920-08	0.9965178760553920-08	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	
REGDUMP 2 AT 1E6208; CC IS 2, PGM MASK IS 2, REGS ARE:									
HEX	00000114	001E607E	00000015	40F43260	001E5700	00800000	001E5780	46000014	
DEC	276	1990782	00000015	1083741808	2053888	8388638	2054064	1174405140	
HEX	00000114	001E5700	001E5700	001E6000	421E6016	001E6010	421E6450	001E6010	
DEC	51760	2053688	2054136	1990656	1109286934	1991118	1109288016	1990672	
HEX	41000053	043C8165	3847405F	741D0000	4100E695	00000000	397137A5	00000000	
DEC	0.138155508601630 02	0.9965178760553920-08	0.9965178760553920-08	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	
REGDUMP 25 AT 1E61B6; CC IS 2, PGM MASK IS 2, REGS ARE:									
HEX	00000114	001E607E	00000015	40F43260	001E5700	00800000	001E5780	46000014	
DEC	276	1990782	00000015	1083741808	2053888	8388638	2054064	1174405140	
HEX	00000114	001E5700	001E5700	001E6000	421E6016	001E6010	421E6450	001E6010	
DEC	51760	2053688	2054136	1990656	1109286934	1991118	1109288016	1990672	
HEX	41000053	043C8165	3847405F	741D0000	4100E695	00000000	397137A5	00000000	
DEC	0.138155508601630 02	0.9965178760553920-08	0.9965178760553920-08	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	
REGDUMP 2 AT 1E6208; CC IS 2, PGM MASK IS 2, REGS ARE:									
HEX	00000114	001E607E	00000015	40F43260	001E5700	00800000	001E5780	46000014	
DEC	276	1990782	00000015	1083741808	2053888	8388638	2054064	1174405140	
HEX	00000114	001E5700	001E5700	001E6000	421E6016	001E6010	421E6450	001E6010	
DEC	51760	2053688	2054136	1990656	1109286934	1991118	1109288016	1990672	
HEX	41000053	043C8165	3847405F	741D0000	4100E695	00000000	397137A5	00000000	
DEC	0.138155508601630 02	0.9965178760553920-08	0.9965178760553920-08	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	0.138155508601630 02	

Figure 1. Sample output from REGDUMP (numbers outside boxes are program output).

```

1      MACRC
2  GNAM  REGDUMP  EID
3      **
4      ** REGDUMP PRINTS A SNAP OF THE GENERAL PURPOSE AND FLOATING
5      ** POINT REGS DURING EXECUTION. IT ALSO GIVES THE CCNDITION CODE,
6      ** PROGRAM MASK, ADDR OF THE LAST INSTR. OF THE MACRC EXPANSION,
7      ** AND AN ID NUMBER. REGDUMP DOES NOT ALTER ANY OF THIS INFORMATION.
8      ** REGDUMP CALLS FORTRAN SUBROUTINE REGCMP FOR OUTPUT
9      **
10     ** REGDUMP WRITTEN BY ADIR PRIDOR AND RAY PAYLAK - 11/14/76
11     **
12     ** CSECT LIST OF CSECT NAMES. ONLY ONE STORAGE AREA FOR EACH CSECT.
13     **
14     **      GBLC  EGS(10)
15     **
16     **      GPC=0PM, CURRENT NUMBER OF NAMES IN EGS.
17     **
18     **      GULA  EPC,0PM
19     **      LCLA  EP,6IDMCD
20     **
21     **      EA WILL BE NAME OF STORAGE AREA.
22     **
23     **      LCLC  EA,EN
24     **      AIF  (I'EID EQ 'C'),CONT1  TO CHOOSE ID IF OMITTED
25     **      AIF  (I'EID NE 'N'),MESS  CHECK IF ID VALID
26     **      AIF  (EID LT 0),MESS
27     **      SETA  EID-((EID/1000)*1000)  REDUCE MCD 1000
28     **      AGO  .CONT
29     **      MNOTE *,*ILLEGAL REGDUMP ID. ESYNDX USED.*
30     **      AGO  .CONT2
31     **      MNOTE *,*REGDUMP ID = ESYNDX*
32     **      .CONT2  ANCP
33     **      EIDMCD  SETA  ESYNDX  CHOOSE ID
34     **      .CONT  ANCP
35     **
36     **      EPM=10 IS LENGTH OF CSECT LIST ( EGS ).
37     **
38     **      SETA  10
39     **      SETA  1
40     **
41     **      LCLC  FOR CSECT NAME IN TABLE.
42     **
43     **      .CHKCS  AIF  (EP GT EPC),NEWCS
44     **      AIF  (I'ESYSECT' EG 'EGS(EP)'),NODEF1
45     **      SETA  EP+1
46     **      AGO  .CHKCS
47     **      .NODEF1  ANCP
48     **      SETC  'CS','EP','PREG'  CREATE AREA NAME
49     **      SETC  'ENAM'  ATTACH ENAM TO B INSTRUCTION
50     **      AGO  .NODEF
51     **
52     **      IF EGS NOT FULL, INSERT CSECT NAME.
53     **
54     **      .NEWCS  AIF  (EPC LT EPM),INSCS
55     **      MNOTE *,*SEPERATE AREA ASSIGNED FOR THIS REGDUMP CALL.*
56     **
57     **      SETC  'PREG','ESYNDX'  CREATE AREA NAME
58     **      AGO  .DEF
59     **      .INSCS  ANCP
60     **      SETA  EPC+1  ADVANCE COUNTER
61     **      SETC  'ESYSECT'  ADD NAME TO LIST
62     **      SETC  'CS','EPC','PREG'  CREATE AREA NAME
63     **      .DEF  ANCP
64     **      CNCP  0,B  SAVE STORAGE AREA
65     **      ENAM  EA+1B4  SKIP AREA
66     **      DC  A(*+B)  ADDRESS PARAM SENT TO REGDMP
67     **      CC  V(REGDMP)
68     **      DS  44F
69     **      .NODEF  ANCP
70     **
71     **      SAVE REGISTERS AND STATUS, GC TO OUTPUT ROUTINE, AND RESTORE DATA
72     **
73     **      EN  STM  0,15,EA+12  SAVE AND PASS GENERAL REGISTERS
74     **      LA  1,EA
75     **      STD  0,76(1)  PASS FLOATING POINT REGISTERS
76     **      STD  2,84(1)
77     **      STD  4,92(1)
78     **      STD  6,100(1)
79     **      LA  13,108(1)  ESTABLISH SAVE AREA
80     **      L  15,4(1)
81     **      LA  2,EICVOL  PASS IC
82     **      ST  2,18C(1)
83     **      BALR  2,0  TAKE CC, MASK, PRG COUNTER
84     **      ST  2,8(1)
85     **      BALN  14,15  GO TO OUTPUT SUBROUTINE
86     **      LD  0,76(1)  RESTORE FLOATING POINT REGISTERS
87     **      LD  2,84(1)
88     **      LD  4,92(1)
89     **      LD  6,100(1)
90     **      SPM  2  RESTORE CC
91     **      LM  0,15,12(1)  RESTORE GENERAL REGISTERS
92     **      MEND

```

Figure 2. Source listing of REGDUMP.

```

0001      SUBROUTINE REGCMP(M)
C
C      PRINTING SUBROUTINE USED BY REGDUMP MACRO.
C
0002      DATA K24 /21600000/
0003      DIMENSION M(44),MM(8)
0004      DOUBLE PRECISION F(4)
0005      EQUIVALENCE (F,MM)
0006      INTEGER ADDR,STATUS,CC
C
0007      DO 1 I=1,8
0008      1   MM(I) = M(I+17)
C
0009      STATUS = M(1)/K24
0010      ADDR = M(1)-STATUS*K24+24
0011      ICC = STATUS/16
0012      MASK = STATUS-16*ICC
0013      CC = ICC-4
C
C      M(44) = REGDUMP ID NUMBER
C      M(2) --> M(17) ARE GENERAL REGS 0-15
C      M(18) --> M(25) ARE FLOATING POINT REGS
C
0014      WRITE(6,100)M(44),ADDR,CC,MASK,(M(I),I=2,9),(M(I),I=2,9),
X      (M(I),I=10,17),(M(I),I=10,17),(M(I),I=18,25),F
0015      100  FORMAT(
C      A      1X,132(' ')
C      B      /' 1 REGDUMP',14,' AT ',Z6,'; CC 15',12,'; PGW MASK 15 ',
C      C      Z1,'; REGS ARE:',73X,' ')
C      D      2(/' 1',130X,' ')
C      E      /' 1',15X,'HEX',8Z14,' ')
C      F      /' 1',15X,'DEC',8I14,' ')
C      G      /' 1',130X,' ')
C      I      /' 1',15X,4(Z19,Z9),3X,' ')
C      J      /' 1',12X,4C28,16,' ')
C      K      /1H ,132(' ')
C      L      )
C
0016      RETURN
0017      END

```

Figure 3. Source listing of the FORTRAN output subroutine.