

PERSISTENTLY ϵ -OPTIMAL STRATEGIES

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Abstract

There are strategies available for nonnegative gambling problems which are not only ϵ -optimal but persist in being conditionally ϵ -optimal along every history.

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1. Introduction.

Let (F, Γ, u) be a gambling problem where, as in [3], F is the set of fortunes or states, Γ is the gambling house, and u the utility function. The gambling problem is nonnegative if u is a nonnegative, possibly unbounded function.

Let $0 < \epsilon < 1$, $f \in F$, and $V(f)$ be the most a gambler with fortune f can achieve ([3], section 3.3). A strategy σ available at f in Γ is (multiplicatively) ϵ -optimal at f if $u(\sigma) \geq (1-\epsilon)V(f)$. Let $p = (f_1, \dots, f_n)$. Then σ is conditionally (multiplicatively) ϵ -optimal given p if the conditional strategy $\sigma[p]$ is multiplicatively ϵ -optimal at f_n . If σ is (multiplicatively) ϵ -optimal at f and is conditionally (multiplicatively) ϵ -optimal given p for every partial history p , then σ is said to be persistently (multiplicatively) ϵ -optimal at f . If u is bounded or, more generally, if u is nonnegative and V is everywhere finite, then persistently ϵ -optimal strategies are always available (Theorem 1 below). If, in addition, the gambling problem is sufficiently measurable, there exist measurable strategies which are persistently ϵ -optimal (Theorem 2).

The notion of persistently (or thoroughly) optimal strategies was introduced in [3, section 3.5]. It is related to the concept of stationarity. For suppose $\bar{\sigma}$ is a stationary family of ϵ -optimal strategies. Then, for every f , $\bar{\sigma}(f)$ is persistently ϵ -optimal since $\bar{\sigma}(f)[f_1, \dots, f_n] = \bar{\sigma}(f_n)$.

Gambles in [3] were taken to be finitely additive probability measures defined on all subsets of F . Here a gamble γ can be regarded as a nonnegative functional with domain the collection \mathcal{H} of nonnegative,

real-valued functions defined on F and satisfying

$$\gamma(u + v) = \gamma u + \gamma v ,$$

$$\gamma(a u) = a \gamma u$$

for $u, v \in \mathcal{H}$ and $a \geq 0$. A gamble in the present sense is, when restricted to the collection of indicator functions, a gamble in the sense of [3]. Thus the gambling problems considered here are slightly more general than those of [3]. However, much of the theory remains unchanged and appropriate definitions and results from [3] will be used without further comment.

2. Preliminaries.

Let $H = F \times F \times \dots$ be the set of histories. A mapping t from H to $\{1, 2, \dots\} \cup \{\infty\}$ is called a stopping time if, whenever h and h' are in H , $t(h) = n < \infty$, and h' agrees with h in the first n coordinates, then $t(h') = n$. A stopping time which assumes only finite values is a stop rule. Let n be a positive integer and $h = (f_1, \dots, f_n, \dots)$ be a history. Recall that $p_n(h) = (f_1, \dots, f_n)$. (See [3] for a detailed explanation of the notation and terminology.) If r is a stopping time, the conditional stopping time given $p_n(h)$ is defined by $t[p_n(h)](h') = t(p_n(h)h') - n$ for $h' \in H$. Notice that $t[p_n(h)]$ is a stopping time or identically zero according as $r(p_n(h)h')$ is greater than or equal to n .

Let (σ, t) be a policy and s be a stop rule. Make the convention that $u(\sigma[p_s(h)], t[p_s(h)])$ is $u(f_t(h))$ when $s(h) \geq t(h)$. Then the formula

$$(1) \quad u(\sigma, t) = \int u(\sigma[p_s], t[p_s]) d\sigma$$

is a special case of formula 3.7.1, [3]. The similar formula

$$(2) \quad u(\sigma) = \int u(\sigma[p_s]) d\sigma$$

follows from Theorem 3.2.1, [3] and an argument by induction on the structure of p_s . The first lemma below generalizes (2) and gives a formula which, in a sense, separates the utility of a strategy into that earned before a given time r and that earned afterwards.

For simplicity, assume henceforth that u is nonnegative and all strategies have finite utility.

Lemma 1. Let σ be a strategy and r be a stopping time. Then

$$(3) \quad u(\sigma) = \lim_{t \rightarrow \infty} \sup u_r(\sigma, t)$$

$$\text{where} \quad u_r(\sigma, t) = \int_{[t \leq r]} u(f_t) d\sigma + \int_{[t > r]} u(\sigma[p_r]) d\sigma$$

and the $\lim \sup$ is taken over the directed set of stop rules.

Proof: By definition ([3], section 3.2)

$$(4) \quad u(\sigma) = \lim_{t \rightarrow \infty} \sup u(\sigma, t) .$$

For t a stop rule, apply (1) with $s = t \wedge r$ to get

$$\begin{aligned} (5) \quad u(\sigma, t) &= \int u(\sigma[p_{t \wedge r}], t[p_{t \wedge r}]) d\sigma \\ &= \int_{[t \leq r]} u(f_t) d\sigma + \int_{[t > r]} u(\sigma[p_r], t[p_r]) d\sigma . \end{aligned}$$

Let $\epsilon > 0$.

Claim: There is a stop rule t_0 such that, for every stop rule t ,

$$\int_{[t > r, t_0 \leq r]} u(\sigma[p_r], t[p_r]) d\sigma < \epsilon .$$

Suppose the claim is false. Then, for every stop rule t , there is a stop rule s such that

$$\int_{[s > r, t \leq r]} u(\sigma[p_r], s[p_r]) d\sigma \geq \epsilon .$$

Define $t_1 = t$ if $t > r$ or $s \leq t$,
 $= s$ if $t \leq r$ and $s > t$.

Then t_1 is a stop rule, $t_1 \geq t$ and

$$\begin{aligned} \int_{[t_1 > r]} u(\sigma[p_r], t_1[p_r]) d\sigma &= \int_{[t > r]} u(\sigma[p_r], t[p_r]) d\sigma + \int_{[t \leq r, s > r]} u(\sigma[p_r], s[p_r]) d\sigma \\ &\geq \int_{[t > r]} u(\sigma[p_r], t[p_r]) d\sigma + \epsilon . \end{aligned}$$

Similarly, there exist stop rules t_2, t_3, \dots such that $t_n \leq t_{n+1}$ and

$$\int_{[t_{n+1} > r]} u(\sigma[p_r], t_{n+1}[p_r]) d\sigma \geq \int_{[t_n > r]} u(\sigma[p_r], t_n[p_r]) d\sigma + \epsilon .$$

Thus, for $n = 1, 2, \dots$, $t_n \geq t$ and

$$u(\sigma, t_n) \geq \int_{[t_n > r]} u(\sigma[p_r], t_n[p_r]) d\sigma \geq n \epsilon .$$

Hence, $u(\sigma) = \infty$, a contradiction which establishes the claim.

Next, define a stop rule t' thus: If $t_0(h) \leq r(h)$, set $t'(h) = t_0(h)$. If $t_0(h) > r(h)$, choose a stop rule t_h such that $t_h \geq t_0[p_r(h)]$ and, for every stop rule $s \geq t_h$,

$$(6) \quad u(\sigma[p_r(h)], s) \leq u(\sigma[p_r(h)]) + \epsilon .$$

Then set $t'(h) = r(h) + t_h(f_{r(h)+1}, f_{r(h)+2}, \dots)$ so that $t'[p_r(h)] = t_h$.

Notice that $t' \geq t_0$.

Suppose t is a stop rule and $t \geq t'$. Then

$$u(\sigma, t) = \int_{[t \leq r]} u(f_t) d\sigma + \int_{[t_0 > r]} u(\sigma[p_r], t[p_r]) d\sigma + \int_{[t > r, t_0 \leq r]} u(\sigma[p_r], t[p_r]) d\sigma$$

(by (5))

$$\leq \int_{[t \leq r]} u(f_t) d\sigma + \int_{[t_0 > r]} u(\sigma[p_r]) d\sigma + 2 \epsilon$$

(by (6) and the Claim)

$$\begin{aligned}
&\leq \int_{[t \leq r]} u(f_t) d\sigma + \int_{[t > r]} u(\sigma[p_r]) d\sigma + 2\epsilon \\
&= u_r(\sigma, t) + 2\epsilon.
\end{aligned}$$

Hence, $u(\sigma) \leq \limsup_{t \rightarrow \infty} u_r(\sigma, t)$.

To prove the opposite inequality, again let $\epsilon > 0$ and s be a stop rule. Define a stop rule t as follows: If $s(h) \leq r(h)$, let $t(h) = s(h)$. If $s(h) > r(h)$, choose t_h to be a stop rule at least as large as $s[p_r(h)]$ and such that

$$(7) \quad u(\sigma[p_r(h)], t_h) \geq u(\sigma[p_r(h)] - \epsilon).$$

Set $t(h) = r(h) + t_h(f_{r(h)+1}, f_{r(h)+2}, \dots)$ so that $t[p_r(h)] = t_h$.

Then $t \geq s$ and

$$\begin{aligned}
u(\sigma, t) &\geq u_r(\sigma, t) - \epsilon \quad (\text{by (5) and (7)}) \\
&= u_r(\sigma, s) - \epsilon.
\end{aligned}$$

Hence, $u(\sigma) \geq \limsup_{t \rightarrow \infty} u_r(\sigma, t)$. \square

Let r be a stopping time. The strategies σ and σ' agree up to time $r-$ if $\sigma_0 = \sigma'_0$ and, for every $h \in H$ and $0 < n < r(h)$, $\sigma_n(p_n(h)) = \sigma'_n(p_n(h))$. Here is a formulation of the nearly obvious fact that if σ and σ' agree up to some time and if, given the past up to that time, σ' does conditionally nearly as well as σ , then unconditionally σ' does nearly as well as σ .

Lemma 2. If r is a stopping time, σ and σ' are strategies which agree up to time $r-$ and $\epsilon > 0$, then each of the following conditions

implies its successor:

- (i) σ' is available in Γ and $u(\sigma[p_r(h)]) \geq (1-\epsilon)V(f_r(h))$ whenever $r(h) < \infty$;
- (ii) $u(\sigma[p_r(h)]) \geq (1-\epsilon)u(\sigma'[p_r(h)])$ whenever $r(h) < \infty$;
- (iii) $u(\sigma) \geq (1-\epsilon)u(\sigma')$.

Proof: To see that (i) implies (ii), observe that

$u(\sigma'[p_r(h)]) \leq V(f_r(h))$ when $r(h) < \infty$. Use Lemma 1 for an easy proof that (ii) implies (iii). \square

The next lemma states that a good strategy earns relatively little income after it becomes conditionally bad. Moreover, any strategy which agrees with a good strategy until it becomes conditionally bad is itself a fairly good strategy. A related fact (Lemma 3.1 in [6]) is that a sufficiently good strategy is unlikely to ever become conditionally bad. First some notation: If σ is a strategy and $0 < \alpha < 1$, let $r(\sigma, \alpha)$ be the first time (if any) when σ is not conditionally α -optimal. That is,

$$(8) \quad r(\sigma, \alpha)(h) = \inf\{n: u(\sigma[p_n(h)]) < (1-\alpha)V(f_n)\}.$$

The infimum of the empty set is taken to be infinite.

Lemma 3. Let α and β be numbers in $(0,1)$. Suppose σ is available at f in Γ and $u(\sigma) \geq (1-\beta)V(f)$. If $r = r(\sigma, \alpha)$, then, for every stop rule t ,

$$(9) \quad \int_{[t > r]} V(f_r) d\sigma \leq \beta \alpha^{-1} V(f).$$

Furthermore, if σ' is available at f in Γ and if σ' and σ agree up to time $r-$, then

$$(10) \quad u(\sigma') \geq (1-\epsilon)V(f)$$

where $\epsilon = \beta(1 + \alpha^{-1})$.

Proof: Let t be a stop rule and set $v_t = \int_{[t>r]} V(f_r) d\sigma$. Then

$$\begin{aligned} (1-\beta)V(f) &\leq u(\sigma) \\ &= \int u(\sigma[p_{t \wedge r}]) d\sigma \quad (\text{by (2)}) \\ &= \int_{[t \leq r]} u(\sigma[p_t]) d\sigma + \int_{[t > r]} u(\sigma[p_r]) d\sigma \\ &\leq \int_{[t \leq r]} V(f_t) d\sigma + (1-\alpha) \int_{[t > r]} V(f_r) d\sigma \\ &= \int V(f_{t \wedge r}) d\sigma - \alpha v_t \\ &\leq V(f) - \alpha v_t \quad (\text{Corollary 3.3.4, [3]}) . \end{aligned}$$

Inequality (9) is now clear.

Now let s be a stop rule and $\epsilon' > 0$. By Lemma 1 there is a stop rule t such that $t \geq s$ and

$$\begin{aligned} (11) \quad (1-\beta)V(f) - \epsilon' &\leq \int_{[t \leq r]} u(f_t) d\sigma + \int_{[t > r]} u(\sigma[p_r]) d\sigma \\ &\leq \int_{[t \leq r]} u(f_t) d\sigma + v_t . \end{aligned}$$

Hence,

$$\begin{aligned} \int_{[t \leq r]} u(f_t) d\sigma' + \int_{[t > r]} u(\sigma'[p_r]) d\sigma' &\geq \int_{[t \leq r]} u(f_t) d\sigma' \\ &= \int_{[t \leq r]} u(f_t) d\sigma \end{aligned}$$

(since σ and σ' agree up to r -)

$$\geq (1-\beta)V(f) - v_t - \epsilon'$$

(by (11))

$$\geq (1-\epsilon)V(f) - \epsilon' .$$

(by (9))

Inequality (10) now follows from another application of Lemma 1. \square

3. The existence proof.

Let (F, Γ, u) be a nonnegative gambling problem. A family of strategies is a mapping from F to the set of strategies. Let $\bar{\rho}$ be a family of strategies and $0 < \alpha < 1$. Define another family of strategies $\bar{\sigma} = \psi(\bar{\rho}, \alpha)$ as follows: For every $f' \in F$, let

$$(12) \quad r(f') = r(\bar{\rho}(f'), \alpha) \quad (\text{see formula (8)}) .$$

Fix $f \in F$. Set $s(1) = r(f)$ and, for $n = 1, 2, \dots$, let $s(n+1)$ be the composition of $s(n)$ with the family $r(\cdot)$ (p. 22, [3]); that is, for $h = (f_1, f_2, \dots) \in H$, let

$$(13) \quad \begin{aligned} s(n+1)(h) &= s(n)(h) + r(f_{s(n)(h)})(f_{s(n)(h)+1}, f_{s(n)(h)+2}, \dots) \\ &\quad \text{if } s(n)(h) < \infty , \\ &= \infty \quad \text{if } s(n)(h) = \infty . \end{aligned}$$

Define $\bar{\sigma}(f)$ to be that strategy which agrees with $\bar{\rho}(f)$ up to time $r(f)$ - and is such that, for $n \geq 1$ and $h \in H$, $\sigma[p_{s(n)}(h)]$ agrees with $\bar{\rho}(f_{s(n)}(h))$ up to time $r(f_{s(n)}(h))$ -. The family $\bar{\sigma} = \psi(\bar{\rho}, \alpha)$ is now defined.

Let $0 < \beta < 1$. Let $S(\beta)$ be the collection of families of strategies $\bar{\rho}$ such that, for every f , $\bar{\rho}(f)$ is available at f in Γ and $u(\bar{\rho}(f)) \geq (1-\beta)V(f)$. Define $P(\alpha, \beta)$ to be the set of families of strategies $\bar{\sigma}$ such that $\bar{\sigma} = \psi(\bar{\rho}, \alpha)$ for some $\bar{\rho}$ in $S(\beta)$. Intuitively, to construct an element of $P(\alpha, \beta)$, start with a family $\bar{\rho}$ of β -optimal strategies; use the strategy which $\bar{\rho}$ specifies at the initial fortune until it becomes conditionally less than α -optimal; then switch

to the strategy specified by $\bar{\rho}$ at the current fortune and use it until it becomes conditionally less than α -optimal and so forth.

Notice that if V is everywhere finite then $S(\beta)$ and $P(\alpha, \beta)$ are not empty.

Theorem 1. Let (F, Γ, u) be a nonnegative gambling problem with V everywhere finite. For every $f \in F$ and $0 < \epsilon < 1$, there is available at f in Γ a strategy σ which is persistently (multiplicatively) ϵ -optimal. Indeed if

$$(14) \quad 0 < \alpha < 1, 0 < \beta < 1, (1 - \beta(1 + \alpha^{-1}))(1 - \alpha) \geq 1 - \epsilon,$$

and $\bar{\sigma} \in P(\alpha, \beta)$, then $\bar{\sigma}(f)$ is persistently (multiplicatively) ϵ -optimal at f .

Proof: Choose α, β to satisfy (14). Let $\bar{\rho} \in S(\beta)$ and $\bar{\sigma} = \psi(\bar{\rho}, \alpha) \in P(\alpha, \beta)$. Set $\sigma = \bar{\sigma}(f)$. Let $r(f)$ and $s(n)$ be as in the definition of $\bar{\sigma}$. Set $\epsilon' = \beta(1 + \alpha^{-1})$. Since $u(\bar{\rho}(f)) \geq (1 - \beta)V(f)$ and σ agrees with $\bar{\rho}(f)$ up to time $r(f)-$, Lemma 3 applies to show $u(\sigma) \geq (1 - \epsilon')V(f) \geq (1 - \epsilon)V(f)$. Similarly, for each positive integer m and $h \in H$, it follows from Lemma 3 that, whenever $s(m)(h) < \infty$,

$$(15) \quad \begin{aligned} u(\sigma[p_{s(m)}(h)]) &\geq (1 - \epsilon')V(f_{s(m)}(h)) \\ &\geq (1 - \epsilon)V(f_{s(m)}(h)). \end{aligned}$$

Consider next a partial history $p = (f_1, \dots, f_n)$ which, for every $k = 1, 2, \dots$, is not of the form $p_{s(k)}(h)$. Let m be the least positive integer such that $n < s(m)(ph')$ for some h' (and, hence, every h'). If $m \geq 2$, let $k = s(m-1)(f_1, \dots, f_n, \dots)$ so that $f_k = f_{s(m-1)}(f_1, \dots, f_n, \dots)$. If $m = 1$, let $k = 0$ and $f_k = f$. Set $\sigma' = \bar{\rho}(f_k)[(f_{k+1}, \dots, f_n)]$ and

$r_0 = r(f_k)[(f_{k+1}, \dots, f_n)]$. Use the definitions of $r(f_k)$ and $s(m)$ (formulas (12) and (13)) to see that

$$(16) \quad u(\sigma') \geq (1 - \alpha)V(f_n) .$$

Check that $\sigma[p]$ and σ' agree up to time r_0 . Also, for $h' \in H$,

$$\begin{aligned} u(\sigma[p][p_{r_0}(h')]) &= u(\sigma[p_{s(m)}(ph')]) \\ &\geq (1 - \epsilon')V(f_{s(m)}(ph')) \quad (\text{by (15)}) \\ &= (1 - \epsilon')V(f_{r_0}(h')) , \end{aligned}$$

whenever $r_0(h') < \infty$. Apply Lemma 2 to the strategies $\sigma[p]$ and σ' to conclude

$$(17) \quad u(\sigma[p]) \geq (1 - \epsilon')u(\sigma') .$$

By (16) and (17),

$$\begin{aligned} u(\sigma[p]) &\geq (1 - \epsilon')(1 - \alpha)V(f_n) \\ &\geq (1 - \epsilon)V(f_n) . \end{aligned}$$

The proof of the theorem is now complete. \square

Even if optimal strategies are available at every fortune, persistently 0-optimal strategies need not be as can be seen from the following example.

Example. Let $F = \{0, 1, \dots\}$; $u(1) = 1$, $u(n) = 0$ if $n \neq 1$;
 $\Gamma(n) = \{\delta(n)\}$ for $n \leq 1$, $\Gamma(n) = \{\delta(n), (1-n^{-1})\delta(1) + n^{-1}\delta(0), \gamma\}$
for $n \geq 2$ where γ is a diffuse gamble on F .

Corollary 1. Let (F, Γ, u) be a gambling problem with u bounded.

For each $f \in F$ and $\epsilon > 0$, there is a strategy σ available at f in Γ such that $u(\sigma) \geq V(f) - \epsilon$ and $u(\sigma[p]) \geq V(f_n) - \epsilon$ for every $p = (f_1, \dots, f_n)$.

Proof: There is no real loss of generality in assuming u is nonnegative.

Since V is bounded, the conclusion follows easily from the theorem. \square

An example of Blackwell in [1] can easily be modified to show that for unbounded, even nonnegative u , there need not exist strategies which are persistently (additively) ϵ -optimal, that is, strategies which satisfy the conclusion of Corollary 1.

4. Measurable gambling problems.

Let X be a separable metric space. Call X analytic if there is a continuous function from the set of irrationals in the unit interval onto X . (See Kuratowski [4] or Blackwell, Freedman, and Orkin [2] for a discussion of analytic sets.) Let $\mathcal{B}(X)$ denote the sigma-field of Borel subsets of X . A nonnegative, real-valued function g defined on X is semi-analytic if $\{x: g(x) > a\}$ is analytic for all nonnegative a . (See [2] and [4].) Denote by $\mathcal{P}(X)$ the set of countably additive probability measures defined on $\mathcal{B}(X)$. Equip $\mathcal{P}(X)$ with the weak-star topology. Then $\mathcal{P}(X)$ is an analytic set if X is. (Lemma (25) in [2]). A function g from an analytic set X into an analytic set Y is universally measurable or measurable for short if, for every $S \in \mathcal{B}(Y)$ and $p \in \mathcal{P}(X)$, $g^{-1}(S)$ is in the completion of $\mathcal{B}(X)$ under p .

A gambling problem (F, Γ, u) is nonnegative analytic if F is an analytic set, u is semi-analytic, and the set $\{(f, \gamma): \gamma \in \Gamma(f)\}$ is an analytic subset of $F \times \mathcal{P}(F)$. (Here each gamble γ is identified with its restriction to $\mathcal{B}(F)$ and is assumed to be countably additive on $\mathcal{B}(F)$.) This definition of analytic gambling problems was inspired by Blackwell, Freedman, and Orkin [2]. Analytic gambling problems include the Borel measurable gambling problems defined by Strauch [5].

A strategy σ is measurable if, for $n = 1, 2, \dots$, the mappings $(f_1, \dots, f_n) \rightarrow \sigma_n(f_1, \dots, f_n)$ are measurable from F^n to $\mathcal{P}(F)$. A measurable strategy σ determines a probability measure $p(\sigma)$ on the Borel subsets of the product space $H = F \times F \times \dots$ as follows: the $p(\sigma)$ -marginal distribution of f_1 is σ_0 and, for every (f_1, \dots, f_n)

the $p(\sigma)$ -distribution of f_{n+1} given (f_1, \dots, f_n) is $\sigma_n(f_1, \dots, f_n)$. For simplicity, denote $p(\sigma)$ by σ below. A measurable family of strategies is a mapping $\bar{\sigma}$ which assigns to each $f \in F$ a measurable strategy $\bar{\sigma}(f)$ in such a way that, for every $n = 0, 1, \dots$, the function $(f, f_1, \dots, f_n) \rightarrow \bar{\sigma}(f)_n(f_1, \dots, f_n)$ is measurable from F^{n+1} to $\mathcal{P}(F)$. The family is available if $\bar{\sigma}(f)$ is available at f for every f .

Lemma 4. If $\bar{\sigma}$ is a measurable family of strategies and u is a bounded Borel function from F to the reals, then the mapping $\varphi: f \rightarrow u(\bar{\sigma}(f))$ is measurable from F to the reals and the mappings $\varphi_n: (f, h) \rightarrow u(\bar{\sigma}(f)[p_n(h)])$ are measurable from $F \times H$ to the reals for $n = 1, 2, \dots$.

Proof: If g is a nonnegative Borel function from H to the real line and if g depends only on a finite number of coordinates, then, by Corollary (41a) of [2], the mapping

$$\varphi_g: f \rightarrow \int g \, d\bar{\sigma}(f)$$

is measurable. The collection \mathcal{Q} of functions g such that φ_g is measurable includes, in particular, the indicator functions of Borel cylinder sets. Since \mathcal{Q} is closed under linear combinations and increasing limits of nonnegative functions, conclude that \mathcal{Q} contains all nonnegative Borel functions on H . By Theorem 3.2 of [8],

$$u(\bar{\sigma}(f)) = \int u^* \, d\bar{\sigma}(f) \text{ where } u^*(f_1, f_2, \dots) = \lim_{n \rightarrow \infty} \sup u(f_n). \text{ Thus}$$

$\varphi = \varphi_{u^*}$ is measurable. The proof that the φ_n are measurable is similar. \square

Theorem 2. Let (F, Γ, u) be a nonnegative analytic gambling problem and assume u is a bounded, Borel measurable function. Suppose that, for every $\beta \in (0, 1)$, there is a measurable family of (multiplicatively)

β -optimal strategies available. Then, for every $\epsilon \in (0,1)$, there is available a measurable family of strategies which are persistently (multiplicatively) ϵ -optimal.

Proof: First notice that V is measurable. To see this, let $\bar{\sigma}_n$ be a measurable family of n^{-1} -optimal strategies for $n = 1, 2, \dots$. Then $V(f) = \sup u(\bar{\sigma}_n(f))$ and, by Lemma 4, V is measurable.

To prove the assertion of the theorem, it suffices, by Theorem 1, to show that, given $\alpha, \beta \in (0, 1)$, there is a measurable family $\bar{\sigma}$ in $P(\alpha, \beta)$. By assumption, there is a measurable family $\bar{\rho} \in S(\beta)$. Let $\bar{\sigma} = \psi(\bar{\rho}, \alpha)$. By definition, $\bar{\sigma} \in P(\alpha, \beta)$. In checking that $\bar{\sigma}$ is measurable, the only real difficulty lies in showing that the map $(f, h) \rightarrow r(f)(h)$ is measurable from $F \times H$ to the real line. (See formulas (12) and (8).) This in turn follows easily from the measurability of V together with Lemma 4. \square

Perhaps, the conclusion of Theorem 2 holds for every nonnegative analytic gambling problem which has V everywhere finite. However, up to the present only two special cases have been treated. It was shown in [4] that good measurable strategies are available for leavable, Borel measurable problems with a bounded utility function. In fact, good measurable families of strategies are always available for leavable, nonnegative analytic problems and so Theorem 2 applies. The same is true for nonleavable problems in which u is the indicator of a single fortune (see [7] for the Borel case). Of course, if F is countable, every strategy is measurable and Theorem 1 yields the existence of persistently good measurable strategies.

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