Degree powers in graphs with a forbidden even cycle

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Abstract

Let C_l denote the cycle of length l. For $p \ge 2$ and integer $k \ge 1$, we prove that the function

$$\phi(k, p, n) = \max\left\{\sum_{u \in V(G)} d^p(u) : G \text{ is a graph of order } n \text{ containing no } C_{2k+2}\right\}$$

satisfies $\phi(k, p, n) = kn^p (1 + o(1))$. This settles a conjecture of Caro and Yuster. Our proof is based on a new sufficient condition for long paths.

1 Introduction

Our notation and terminology follow [1]; in particular, C_l denotes the cycle of length l.

For $p \ge 2$ and integer $k \ge 1$, Caro and Yuster [3], among other things, studied the function

$$\phi(k, p, n) = \max\left\{\sum_{u \in V(G)} d_G^p(u) : G \text{ is a graph of order } n \text{ without a } C_{2k+2}\right\}$$

and conjectured that

$$\phi(k, p, n) = kn^{p} (1 + o(1)).$$
(1)

The graph $K_k + \overline{K}_{n-k}$, i.e., the join of K_k and \overline{K}_{n-k} , gives $\phi(k, p, n) > k(n-1)^p$, so to prove (1), a matching upper bound is necessary. We give such a bound in Corollary 3 below. Our main tool, stated in Lemma 1, is a new sufficient condition for long paths. This result has other applications as well, for instance, the following spectral bound, proved in [5]:

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Let G be a graph of order n and μ be the largest eigenvalue of its adjacency matrix. If G does not contain C_{2k+2} , then

$$\mu^2 - k\mu \leqslant k \left(n - 1 \right).$$

2 Main results

We write |X| for the cardinality of a finite set X. Let G be a graph, and X and Y be disjoint sets of vertices of G. We write:

- V(G) for the vertex set of G and |G| for |V(G)|;
- $e_G(X)$ for the number of edges induced by X;
- $e_G(X, Y)$ for the number of edges joining vertices in X to vertices in Y;
- G u for the graph obtained by removing the vertex $u \in V(G)$;
- $\Gamma_{G}(u)$ for the set of neighbors of the vertex u and $d_{G}(u)$ for $|\Gamma_{G}(u)|$.

The main result of this note is the following lemma.

Lemma 1 Suppose that $k \ge 1$ and let the vertices of a graph G be partitioned into two sets A and B.

(1) If

$$2e_G(A) + e_G(A, B) > (2k - 2) |A| + k |B|, \qquad (2)$$

then there exists a path of order 2k or 2k + 1 with both ends in A.

(2) If

$$2e_G(A) + e_G(A, B) > (2k - 1)|A| + k|B|, \qquad (3)$$

then there exists a path of order 2k + 1 with both ends in A.

Note that if we choose the set B to be empty, Lemma 1 amounts to a classical result of Erdős and Gallai:

If a graph of order n has more than kn/2 edges, then it contains a path of order k+2.

We postpone the proof of Lemma 1 to Section 3 and turn to two consequences.

Theorem 2 Let G be a graph with n vertices and m edges. If G does not contain a C_{2k+2} , then

$$\sum_{u\in V(G)}d_{G}^{2}\left(u\right)\leqslant 2km+k\left(n-1\right)n.$$

Proof Let u be any vertex of G. Partition the vertices of the graph G - u into the sets $A = \Gamma_G(u)$ and $B = V(G) \setminus (\Gamma_G(u) \cup \{u\})$. Since G contains no C_{2k+2} , the graph G - u does not contain a path of order 2k + 1 with both ends in A. Applying Lemma 1, part (2), we see that

$$2e_{G-u}(A) + e_{G-u}(A, B) \leq (2k - 1) |A| + k |B|,$$

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and therefore,

$$\sum_{v \in \Gamma_G(u)} (d_G(v) - 1) = \sum_{v \in \Gamma_G(u)} d_{G-u}(v) = 2e_{G-u}(A) + e_{G-u}(A, B)$$

$$\leq (2k - 1) |A| + k |B|$$

$$= (2k - 1) d_G(u) + k (n - d_G(u) - 1).$$

Rearranging both sides, we obtain

$$\sum_{v \in \Gamma_G(u)} d_G(v) \leqslant k d_G(u) + k (n-1).$$

Adding these inequalities for all vertices $u \in V(G)$, we find out that

$$\sum_{u \in V(G)} \sum_{v \in \Gamma_G(u)} d_G(v) \leqslant k \sum_{u \in V(G)} d_G(u) + k (n-1) n = 2km + k (n-1) n.$$

To complete the proof of the theorem note that the term $d_G(v)$ appears in the left-hand sum exactly $d_G(v)$ times, and so

$$\sum_{u \in V(G)} \sum_{v \in \Gamma_G(u)} d_G(v) = \sum_{v \in V(G)} d_G^2(v) \,.$$

Here is a corollary of Theorem 2 that gives the upper bound for the proof of (1).

Corollary 3 Let G be a graph with n vertices. If G does not contain a C_{2k+2} , then for every $p \ge 2$,

$$\sum_{u \in V(G)} d_G^p\left(u\right) \leqslant kn^p + O\left(n^{p-1/2}\right).$$

Proof Letting m be the number of edges of G, we first deduce an upper bound on m. Theorem 2 and the AM-QM inequality imply that

$$\frac{4m^2}{n} \leqslant \sum_{u \in V(G)} d_G^2(u) \leqslant 2km + k(n-1)n,$$

and so,

$$m \leqslant -kn + n\sqrt{k(n-1) + k^2} < n\sqrt{kn}.$$
(4)

Note that much stronger upper bounds on m are known (e.g., see [2] and [6]), but this one is simple and unconditional.

Now Theorem 2 and inequality (4) imply that

$$\sum_{u \in V(G)} d_G^p(u) < \sum_{u \in V(G)} n^{p-2} d_G^2(u) < kn^p + 2kmn^{p-2} < kn^p + 2(kn)^{3/2} n^{p-2}$$
$$= kn^p + O\left(n^{p-1/2}\right),$$

completing the proof.

Note that we need only part (2) of Lemma 1 to prove Theorem 2 and Corollary 3. However, part (1) of Lemma 1 may have also applications, as shown in [5].

3 Proof of Lemma 1

To simplify the proof of Lemma 1 we state two routine lemmas whose proofs are omitted.

Lemma 4 Let $P = (v_1, \ldots, v_p)$ be a path of maximum order in a connected non-Hamiltonian graph G. Then $p \ge d_G(v_1) + d_G(v_p) + 1$.

Lemma 5 Let $P = (v_1, \ldots, v_p)$ be a path of maximum order in a graph G. Then either v_1 is joined to two consecutive vertices of P or G contains a cycle of order at least $2d_G(v_1)$.

Proof of Lemma 1 For convenience we shall assume that the set B is independent. Also, we shall call a path with both ends in A an A-path.

Claim 6 If G contains an A-path of order p > 2, then G contains an A-path of order p - 2.

Indeed, let (v_1, \ldots, v_p) be an A-path. If $v_2 \in B$, then $v_3 \in A$, and so (v_3, \ldots, v_p) is an A-path of order p-2. If $v_{p-1} \in B$, then $v_{p-2} \in A$, and so (v_1, \ldots, v_{p-2}) is an A-path of order p-2. Finally, if both $v_2 \in A$ and $v_{p-1} \in A$, then (v_2, \ldots, v_{p-1}) is an A-path of order p-2.

The proofs of the two parts of Lemma 1 are very similar, but since they differ in the details, we shall present them separately.

Proof of part (1)

From Claim 6 we easily obtain the following consequence:

Claim 7 If G contains an A-path of order $p \ge 2k$, then G contains an A-path of order 2k or 2k + 1.

This in turn implies

Claim 8 If G contains a cycle C_p for some $p \ge 2k + 1$, then G contains an A-path of order 2k or 2k + 1.

Indeed, let $C = (v_1, \ldots, v_p, v_1)$ be a cycle of order $p \ge 2k+1$. The assertion is obvious if C is entirely in A, so let assume that C contains a vertex of B, say $v_1 \in B$. Then $v_2 \in A$ and $v_p \in A$; hence, (v_2, \ldots, v_p) is an A-path of order at least 2k. In view of Claim 7, this completes the proof of Claim 8.

To complete the proof of part (1) we shall use induction on the order of G. First we show that condition (2) implies that $|G| \ge 2k$. Indeed, assume that $|G| \le 2k-1$. We have

$$|A|^{2} - |A| + |A| |B| \ge 2e_{G}(A) + e_{G}(A, B) > (2k - 2) |A| + k |B|$$

and so,

$$|G|(|A| - k) = (|A| + |B|)(|A| - k) > (k - 1)|A|.$$

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Hence, we find that

$$(2k-1)(|A|-k) > (k-1)|A|$$

and so, |A| > 2k - 1, a contradiction with $|A| \leq |G|$.

The conclusion of Lemma 1, part (1) follows when $|G| \leq 2k - 1$ since then the hypothesis is false. Assume now that $|G| \geq 2k$ and that the Lemma holds for graphs with fewer vertices than G. It is easy to see that this assumption implies the assertion if G is disconnected. Indeed, let G_1, \ldots, G_s be the components of G. Assuming that G has no A-path of order 2k + 1, the inductive assumption implies that each component G_i satisfies

$$2e_{G_i}(A_i) + e_{G_i}(A_i, B_i) \leq (2k - 2) |A_i| + k |B_i|, \qquad (5)$$

where

$$A_i = A \cap V(G_i)$$
 and $B_i = B \cap V(G_i)$

Summing (5) for i = 1, ..., s, we obtain a contradiction to (2).

Thus, to the end of the proof we shall assume that G is connected. Also, we can assume that G is non-Hamiltonian. Indeed, in view of Claim 8, this is obvious when |G| > 2k. If |G| = 2k and G is Hamiltonian, then no two consecutive vertices along the Hamiltonian cycle belong to A, and since B is independent, we have |B| = |A| = k. Then

$$k(2k-1) \ge 2e_G(A) + e_G(A,B) > (2k-2)|A| + k|B| = k(2k-1),$$

contradicting (2). Thus, we shall assume that G is non-Hamiltonian.

The induction step is completed if there is a vertex $u \in B$ such that $d_G(u) \leq k$. Indeed the sets A and $B' = B \setminus \{u\}$ partition the vertices of G - u and also

$$2e_{G-u}(A) + e_{G-u}(A, B) = 2e_G(A) + e_G(A, B) - d_G(u) > (2k-2)|A| + k|B| - k$$

= (2k-2)|A| + k|B'|;

hence G - u contains an A-path of order at least 2k, completing the proof. Thus, to the end of the proof we shall assume that

(a) $d_G(u) \ge k+1$ for every vertex $u \in B$.

For every vertex $u \in A$, write $d'_G(u)$ for its neighbors in A and $d''_G(u)$ for its neighbors in B. The induction step can be completed if there is a vertex $u \in A$ such that $2d'_G(u) + d''_G(u) \leq 2k - 2$. Indeed, if u is such a vertex, note that the sets $A' = A \setminus \{u\}$ and Bpartition the vertices of G - u and also

$$2e_{G-u}(A) + e_{G-u}(A, B) = 2e_G(A) + e_G(A, B) - 2d'_G(u) - d''_G(u)$$

> $(2k-2)|A| + k|B| - 2k + 2$
= $(2k-2)|A'| + k|B|;$

hence G - u contains an A-path of order at least 2k, completing the proof. Hence we have $2d'_G(u) + d''_G(u) \ge 2k - 1$, and so $d_G(u) \ge k$. Thus, to the end of the proof, we shall assume that:

(b) $d_G(u) \ge k$ for every vertex $u \in A$.

Select now a path $P = (v_1, \ldots, v_p)$ of maximum length in G. To complete the induction step we shall consider three cases: (i) $v_1 \in B$, $v_p \in B$; (ii) $v_1 \in B$, $v_p \in A$, and (iii) $v_1 \in A$, $v_p \in A$.

Case (i): $v_1 \in B, v_p \in B$

In view of assumption (a) we have $d_G(v_1) + d_G(v_p) \ge 2k + 2$, and Lemma 4 implies that $p \ge 2k+3$. We see that (v_2, \ldots, v_{p-1}) is an A-path of order at least 2k+1, completing the proof by Claim 7.

Case (ii): $v_1 \in B, v_p \in A$

In view of assumptions (a) and (b) we have $d_G(v_1) + d_G(v_p) \ge 2k + 1$, and Lemma 4 implies that $p \ge 2k + 2$, and so, (v_2, \ldots, v_p) is an A-path of order at least 2k + 1. This completes the proof by Claim 7.

Case (iii): $v_1 \in A, v_p \in A$

In view of assumption (b) we have $d_G(v_1) + d_G(v_p) \ge 2k$, and Lemma 4 implies that $p \ge 2k + 1$. Since (v_1, \ldots, v_p) is an A-path of order at least 2k + 1, by Claim 7, the proof of part (A) of Lemma 1 is completed.

Proof of part (2)

From Claim 6 we easily obtain the following consequence:

Claim 9 If G contains an A-path of odd order $p \ge 2k+1$, then G contains an A-path of order exactly 2k + 1.

From Claim 9 we deduce another consequence:

Claim 10 If G contains a cycle C_p for some $p \ge 2k + 1$, then G contains an A-path of order exactly 2k + 1.

Indeed, let $C = (v_1, \ldots, v_p, v_1)$ be a cycle of order $p \ge 2k + 1$. If p is odd, then some two consecutive vertices of C belong to A, say the vertices v_1 and v_2 . Then (v_2, \ldots, v_p, v_1) is an A-path of odd order $p \ge 2k + 1$, and by Claim 9 the assertion follows. If p is even, then $p \ge 2k + 2$. The assertion is obvious if C is entirely in A, so let assume that Ccontains a vertex of B, say $v_1 \in B$. Then $v_2 \in A$ and $v_p \in A$; hence (v_2, \ldots, v_p) is an A-path of odd order at least 2k + 1, completing the proof of Claim 10.

To complete the proof of Lemma 1 we shall use induction on the order of G. First we show that condition (3) implies that $|G| \ge 2k+1$. Indeed, assume that $|G| \le 2k$. We have

$$|A|^{2} - |A| + |A| |B| \ge 2e_{G} (A) + e_{G} (A, B) > (2k - 1) |A| + k |B|$$

and so,

$$|G|(|A| - k) = (|A| + |B|)(|A| - k) > k |A|.$$

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Hence, we find that 2k(|A|-k) > k|A|, and |A| > 2k, contradicting that $|A| \leq |G|$.

The conclusion of Lemma 1, part (2) follows when $|G| \leq 2k$ since then the hypothesis is false. Assume now that $|G| \geq 2k + 1$ and that the assertion holds for graphs with fewer vertices than G. As in part (1), it is easy to see that this assumption implies the assertion if G is disconnected, so to the end of the proof we shall assume that G is connected. Also, in view of Claim 10 and $|G| \geq 2k + 1$, we shall assume that G is non-Hamiltonian.

The induction step is completed if there is a vertex $u \in B$ such that $d_G(u) \leq k$. Indeed the sets A and $B' = B \setminus \{u\}$ partition the vertices of G - u and also

$$2e_{G-u}(A) + e_{G-u}(A, B) = 2e_G(A) + e_G(A, B) - d_G(u)$$

> (2k - 1) |A| + k |B| - k
= (2k - 1) |A| + k |B'|;

hence G - u contains an A-path of order 2k + 1, completing the proof. Thus, to the end of the proof we shall assume that:

(a) $d_G(u) \ge k+1$ for every vertex $u \in B$.

For every vertex $u \in A$, write $d'_G(u)$ for its neighbors in A and $d''_G(u)$ for its neighbors in B. The induction step can be completed if there is a vertex $u \in A$ such that $2d'_G(u) + d''_G(u) \leq 2k - 1$. Indeed, if u is such a vertex, note that the sets $A' = A \setminus \{u\}$ and Bpartition the vertices of G - u and also

$$2e_{G-u}(A) + e_{G-u}(A, B) = 2e_G(A) + e_G(A, B) - 2d'_G(u) - d''_G(u)$$

> (2k - 1) |A| + k |B| - 2k + 1
= (2k - 1) |A'| + k |B|;

hence G - u contains an A-path of order 2k + 1, completing the proof. Thus, to the end of the proof, we shall assume that:

(b) $d_G(u) \ge k$ for every vertex $u \in A$ and if u has neighbors in B, then $d_G(u) \ge k+1$.

Select now a path $P = (v_1, \ldots, v_p)$ of maximum length in G. To complete the induction step we shall consider three cases: (i) $v_1 \in B$, $v_p \in B$; (ii) $v_1 \in B$, $v_p \in A$, and (iii) $v_1 \in A$, $v_p \in A$.

Case (i): $v_1 \in B, v_p \in B$

In view of assumption (b) we have $d_G(v_1) + d_G(v_p) \ge 2k + 2$, and Lemma 4 implies that $p \ge 2k+3$. If p is odd, we see that (v_2, \ldots, v_{p-1}) is an A-path of order at least 2k+1, and by Claim 9, the proof is completed.

Suppose now that p is even. Applying Lemma 5, we see that either G has a cycle of order at least $2d_G(v_1) \ge 2k+2$, or v_1 is joined to v_i and v_{i+1} for some $i \in \{2, \ldots, p-2\}$. In the first case we complete the proof by Claim 10; in the second case we see that the sequence

$$(v_2, v_3, \ldots, v_i, v_1, v_{i+1}, v_{i+2}, \ldots, v_{p-1})$$

is an A-path of order p-1. Since p-1 is odd and $p-1 \ge 2k+3$, the proof is completed by Claim 9.

Case (ii): $v_1 \in B, v_p \in A$

In view of assumptions (a) and (b) we have $d_G(v_1) + d_G(v_p) \ge 2k + 1$, and Lemma 4 implies that $p \ge 2k + 2$. If p is even, we see that (v_2, \ldots, v_{p-1}) is an A-path of order at least 2k + 1, and by Claim 9, the proof is completed.

Suppose now that p is odd. Applying Lemma 5, we see that either G has a cycle of order at least $2d_G(v_1) \ge 2k+2$, or v_1 is joined to v_i and v_{i+1} for some $i \in \{2, \ldots, p-1\}$. In the first case we complete the proof by Claim 10; in the second case we see that the sequence

$$(v_2, v_3, \ldots, v_i, v_1, v_{i+1}, v_{i+2}, \ldots, v_p)$$

is an A-path of order p. Since p is odd and $p \ge 2k+2$, the proof is completed by Claim 9.

Case (iii): $v_1 \in A, v_p \in A$

In view of assumption (b) we have $d_G(v_1) + d_G(v_p) \ge 2k$, and Lemma 4 implies that $p \ge 2k + 1$. If p is odd, the proof is completed by Claim 9.

Suppose now that p is even, and therefore, $p \ge 2k + 2$. If $v_2 \in A$, then the sequence (v_2, \ldots, v_p) is an A-path of odd order $p - 1 \ge 2k + 1$, completing the proof by Claim 9. If $v_2 \in B$, we see that v_1 has a neighbor in B, and so, $d_G(v_1) \ge k + 1$.

Applying Lemma 5, we see that either G has a cycle of order at least $2d_G(v_1) \ge 2k+2$, or v_1 is joined to v_i and v_{i+1} for some $i \in \{2, \ldots, p-2\}$. In the first case we complete the proof by Claim 10. In the second case we shall exhibit an A-path of order p-1. Indeed, if i = 2, let

$$Q = (v_1, v_3, v_4, \ldots, v_p),$$

and if $i \ge 3$, let

$$Q = (v_3, \ldots, v_i, v_1, v_{i+1}, v_{i+2}, \ldots, v_p).$$

In either case Q is an A-path of order p-1. Since p-1 is odd and $p-1 \ge 2k+1$, the proof is completed by Claim 9.

This completes the proof of Lemma 1.

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