# Coloring 2-intersecting hypergraphs 

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#### Abstract

A hypergraph is 2 -intersecting if any two edges intersect in at least two vertices. Blais, Weinstein and Yoshida asked (as a first step to a more general problem) whether every 2 -intersecting hypergraph has a vertex coloring with a constant number of colors so that each hyperedge has at least $\min \{|e|, 3\}$ colors. We show that there is such a coloring with at most 5 colors (which is best possible).


A proper coloring of a hypergraph is a coloring of its vertices so that no edge is monochromatic, i.e. contains at least two vertices with distinct colors. It is well-known that intersecting hypergraphs without singleton edges have proper colorings with at most three colors. This statement is from the seminal paper of Erdős and Lovász [3]. Recently Blais, Weinstein and Yoshida suggested a generalization in [1]. They consider $t$-intersecting hypergraphs, in which any two edges intersect in at least $t$ vertices and they call a coloring of the vertices $c$-strong if every edge $e$ is colored with at least $\min \{|e|, c\}$ distinct colors. One of the problems they consider is the following.

Problem 1. ([1]) Suppose that $\mathcal{H}$ is a t-intersecting hypergraph. Is there $a(t+1)$ strong vertex coloring of $\mathcal{H}$ where the number of colors is bounded by a function of $t$ ? In particular, is there a $t+1$-strong vertex coloring with at most $2 t+1$ colors? If true, it would be best possible, as the $2 t$-element sets of a $3 t$ element set demonstrate.

Notice that for $t=1$ the answer to Problem 1 is affirmative (for both parts) according to the starting remark but open for $t \geqslant 2$ [1]. Our aim is to give an affirmative answer to both parts of the problem in case of $t=2$. Notice that intersecting hypergraphs do

[^0]not always have 3 -strong colorings with any fixed number of colors: if every edge of a $(k+1)$-chromatic graph is extended by the same new vertex, the resulting intersecting hypergraph has no 3 -strong coloring with $k$ colors. Thus the 2 -intersecting condition is important in the following theorem.

Theorem 2. Every 2-intersecting hypergraph $G$ has a 3-strong coloring with at most five colors.

We learned from a referee that a weaker form of Theorem 2 (with 21 colors instead of 5) is proved recently in [2]. We also prove a lemma that will be used in the proof of Theorem 2 but has independent interest. A hypergraph has property $P_{t}$ for some integer $t \geqslant 2$ if any $i$ edges intersect in at least $t+1-i$ vertices, for all $i, 2 \leqslant i \leqslant t$.

Lemma 3. Suppose that $\mathcal{H}$ is a hypergraph with property $P_{t}$. Then $\mathcal{H}$ has a t-strong coloring with at most $t+1$ colors.

Proof. Let $\mathcal{H}$ be a hypergraph with property $P_{t}$ for $t \geqslant 2$. Select an edge $e$ of $\mathcal{H}$ which is minimal for containment. Let $\mathcal{F}$ be the hypergraph defined on the vertex set of $e$ with edge set $\{h \cap e: h \in E(\mathcal{H})\}$. Color each vertex not in $e$ with color $t+1$. If $t=2$, color the vertices of $e$ arbitrarily using colors 1,2 (or just color 1 if $e$ has just one vertex). If $|e|=t-1$, color vertices of $e$ by $1,2, \ldots, t-1$. Otherwise, since $\mathcal{F}$ has property $P_{t-1}$, we can find by induction a $(t-1)$-strong coloring $C$ on $\mathcal{F}$ with colors $1,2, \ldots, t$. We may suppose that $C$ uses all colors $1,2, \ldots, t$ on $e$, otherwise we may change some repeated colors to the missing colors maintaining the $(t-1)$-strong coloring. Thus $C$ colors $e$ with at least $t$ colors and, since for any other edge $h \in \mathcal{H},|h \cap e| \geqslant t-1, C$ uses at least $t-1$ colors on $h \cap e$ and $h$ also has at least one vertex of color $t+1$. Therefore we have a $t$-strong coloring of $\mathcal{H}$ with $t+1$ colors.

It is worth noting that Lemma 3 does not hold if we require a $t$-strong coloring with at most $t$ colors. Indeed, all $t$-sets of $t+1$ elements have property $P_{t}$ but a $t$-strong coloring must use $t+1$ colors.

Proof of Theorem 2. By the condition, there are no singleton edges. Also, if some edge $e$ has just two vertices, coloring them with colors 1,2 and all other vertices by 3 , we obviously have a 3 -strong coloring. Thus we may assume that every edge has at least three vertices, therefore a 3 -strong coloring on the minimal edges of $G$ is also a 3 -strong coloring on $G$. Thus we may assume that $G$ is an antichain.

If any three edges of $G$ have non-empty intersection, we can apply Lemma 3 and get a 3 -strong coloring with at most 4 colors. Thus, we may suppose that $G$ contains three edges with empty intersection, select them with the smallest possible union, let these edges be $e_{1}, e_{2}, e_{3}$ and set $X=e_{1} \cup e_{2} \cup e_{3}$. A vertex $v \in X$ is called a private part of $e_{i}$ $(i=1,2,3)$ if $v \in e_{i}$ but $v$ is not covered by any of the other two $e_{j}$-s.

We color the vertices in $X$ as follows. The private parts of $e_{1}, e_{2}, e_{3}$ (if they exist) are colored with $1,2,3$ respectively. Notice that each intersection has at least two vertices, color $e_{1} \cap e_{3}$ with colors 1,3 so that color 1 is used only once, color $e_{1} \cap e_{2}$ with colors 2 , 4 so that color 2 is used only once. Vertices in $e_{2} \cap e_{3}$ are all colored with color 5 .

The coloring outside $X$ varies according to the number of private parts of $e_{i}$-s.
Case 1. Each $e_{i}$ has private parts, $i=1,2,3$.
Here we color vertices not covered by $X$ one by one with 1 or 2 by the following greedy type algorithm: if an uncolored vertex $w \notin X$ completes an edge $f$ such that all vertices of $f-\{w\}$ are colored with colors 2,3 only (both present otherwise $\left|f \cap e_{1}\right| \leqslant 1$ or $\left|f \cap e_{2}\right| \leqslant 1$ ) then color $w$ with color 1 , otherwise color it with color 2 . We claim that a 3 -strong coloring is obtained.

Suppose there is an edge $f_{i j}$ with colors $i, j$ only, $1 \leqslant i<j \leqslant 5$. Edges $f_{12}, f_{14}, f_{24}$ would intersect $e_{3}$ in at most one vertex, edge $f_{25}$ would intersect $e_{1}$ in at most one vertex and $f_{13}$ would not intersect $e_{2}$ at all. Edges $f_{35}, f_{45}$ would form a proper subset of $e_{3}, e_{2}$, respectively, contradicting the antichain property.

Edge $f_{34}$ cannot exist because the triple $f_{34}, e_{2}, e_{3}$ has no intersection and $Y=f_{34} \cup$ $e_{2} \cup e_{3}$ is a proper subset of $X$ because $e_{1}$ has a private vertex. Thus we get a contradiction with the definition of $e_{1}, e_{2}, e_{3}$. The same argument can be applied to exclude $f_{15}, f_{23} \subset X$ (with $Y=f_{15} \cup e_{1} \cup e_{2}, Y=f_{23} \cup e_{2} \cup e_{3}$ and using that $e_{3}, e_{1}$ have private vertices).

Thus the only possibility is that there is an edge $f_{15}$ or $f_{23}$ with some vertex $w \notin X$. However, no such $f_{15}$ exists since $w \notin X$ is colored with 1 only if there exists edge $f$ of $G$ such that $f-\{w\}$ is colored with colors 2,3 only thus $\left|f \cap f_{15}\right|=1$ contradiction. Moreover, no such $f_{23}$ can exist either, because its vertex in $V-X$ colored last got color 1 according to the rule governing Case 1.

Case 2. Two of $e_{1}, e_{2}, e_{3}$ have private parts, by suitable relabeling we may suppose that the private part of $e_{2}$ is empty.

In this case vertices not covered by $X$ are colored with color 2 and claim that we have a 3 -strong coloring. The nonexistence of $f_{12}, f_{13}, f_{14}, f_{24}, f_{25}$ follow as in Case 1 and here $f_{23}$ can be excluded the same way since $\left|f_{23} \cap e_{2}\right| \leqslant 1$. The exclusion of $f_{34}, f_{35}, f_{45}$ and $f_{15} \subset X$ is also exactly the same as in Case 1 . Thus here we have to exclude only the existence of an edge $f_{15}$ containing some vertices $w \notin X$. However, this cannot happen since here every vertex outside $X$ is colored with color 2 .
Case 3. Exactly one of $e_{1}, e_{2}, e_{3}$ has a private part, by suitable relabeling we may suppose that it is $e_{2}$.

Here all vertices not covered by $X$ are colored with 1 . Edges $f_{12}, f_{13}, f_{14}, f_{15}, f_{24}, f_{25}$ are all excluded since there is some $e_{i}$ intersecting them in at most one vertex. The edges $f_{34}, f_{35}, f_{45}$ are excluded since they are proper subsets of some $e_{i}$. The only possible edge is $f_{23}$ but in this case we can replace the triple $e_{1}, e_{2}, e_{3}$ by the non-intersecting triple $f_{23}, e_{2}, e_{3}$ which has the same union but they have two private parts: the vertices of color 4 in $e_{2}$ and the vertex of color 1 in $e_{3}$. This reduces Case 3 to Case 2.

Case 4. None of the edges $e_{1}, e_{2}, e_{3}$ have private parts.
Vertices uncovered by $X$ are colored with 1 . Here $f_{12}, f_{13}, f_{14}, f_{15}, f_{23}, f_{24}, f_{25}$ are all excluded since there is some $e_{i}$ intersecting them in at most one vertex. The other three edges $f_{34}, f_{35}, f_{45}$ are excluded since they are proper subsets of some $e_{i}$.

In all cases we found a 3 -strong coloring with at most five colors.

## References

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